

Introduction to Algorithms

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GUROBI
OPTIMIZATION

The World's Fastest Solver

About the Speaker



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Agenda

1. Linear programming (LP)
 - a. Define LP
 - b. LP examples
 - c. Algorithms (high-level)
2. Mixed-integer programming (MIP)
 - a. Define MIP
 - b. MIP example
3. Basic techniques for solving MIP
 - a. Bounding the optimal solution
 - b. Bound-and-bound
 - c. Improvements (heuristics, cuts, etc.)
 - d. Termination criteria
4. Log file interpretation
5. Outlook: Tuning

Introducing LP

What is LP?

$$\begin{array}{llll} \min_{x \in \mathbb{R}^n} & f(x) & & \\ \text{subject to} & h_1(x) & \leq & 0 \\ & h_2(x) & \leq & 0 \\ & \vdots & & \\ & h_m(x) & \leq & 0 \end{array}$$

- This is a Linear Programming (LP) instance if:
 - The objective function $f(x)$ is linear
 - e.g., $f(x) = 3x_1 + 2x_2$
 - The functions $h_1(x), \dots, h_m(x)$ are linear
 - e.g., $h_1(x) = x_1 - 17x_2 + 10$, etc.

Quiz time!

$$\begin{array}{llllllll} \min_{x \in \mathbb{R}^2} & x_1 & - & x_2 & & & & \\ \text{subject to} & x_1 & + & 17x_2 & - & 10 & \leq & 0 \\ & 2x_1 & + & 3x_2 & - & 5 & \leq & 0 \\ & -x_1 & & & & & \leq & 0 \\ & & & & & -x_2 & \leq & 0 \end{array}$$



Quiz time!

$$\begin{array}{llllllll} \min_{x \in \mathbb{R}^2} & x_1^2 & - & x_2 & & & & \\ \text{subject to} & x_1 & + & 17x_2 & - & 10 & \leq & 0 \\ & 2x_1 & + & 3x_2 & - & 5 & \leq & 0 \\ & -x_1 & & & & & \leq & 0 \\ & & & & & -x_2 & \leq & 0 \end{array}$$



Quiz time!

$$\begin{array}{ll} \min_{x \in \mathbb{R}^2} & |x_1 - x_2| \\ \text{subject to} & x_1 + 17x_2 - 10 \leq 0 \\ & 2x_1 + 3x_2 - 5 \leq 0 \\ & -x_1 \leq 0 \\ & -x_2 \leq 0 \end{array}$$

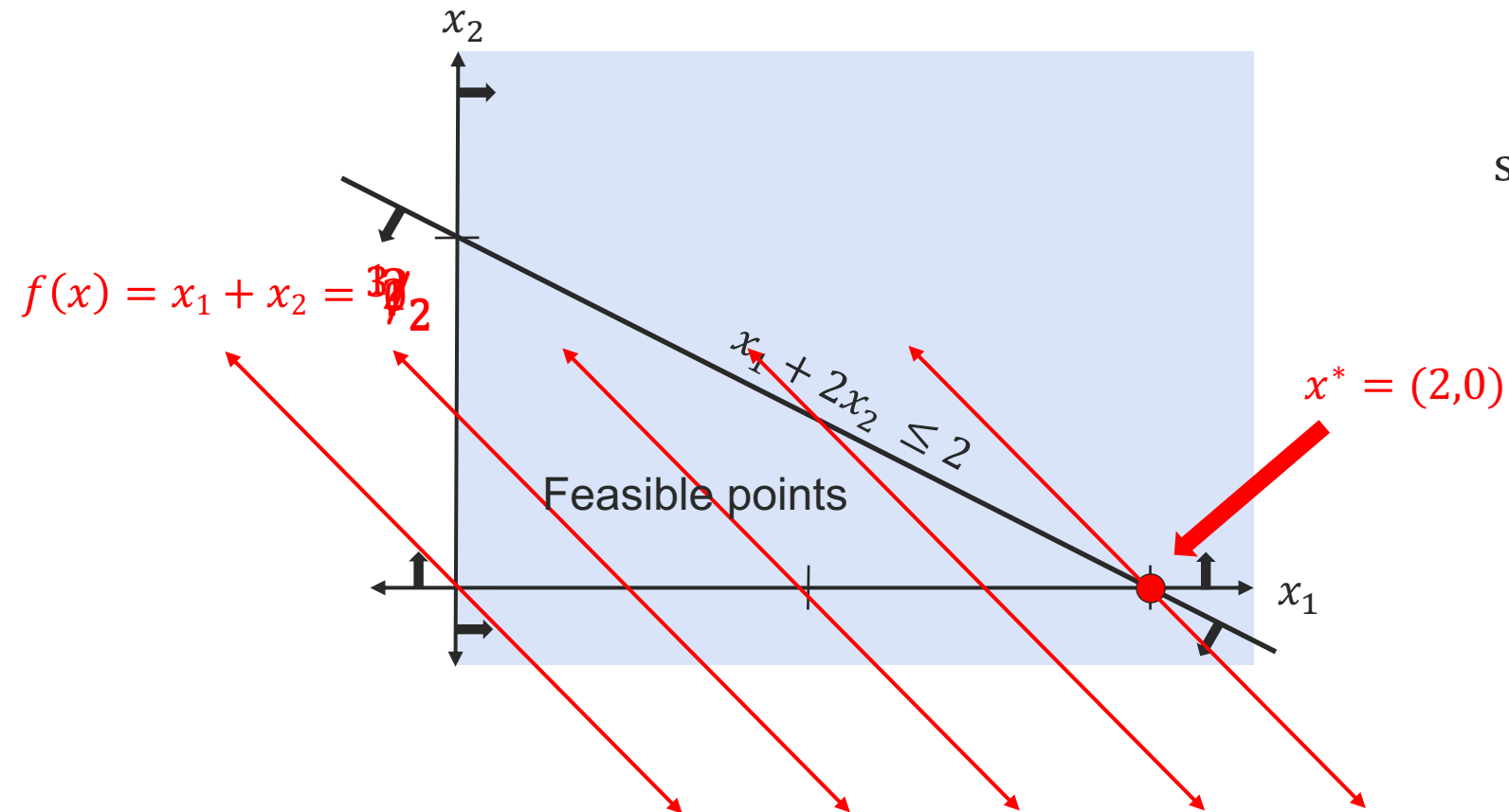


Quiz time!

$$\begin{array}{ll} \max_{x \in \mathbb{R}^2} & x_1 - x_2 \\ \text{subject to} & x_1 + 17x_2 \leq 10 \\ & 2x_1 + 3x_2 - 5 \geq 0 \\ & 4x_1 + x_2 - 9 = 0 \\ & -x_1 \leq 0 \\ & -x_2 \leq 0 \end{array} \quad \Leftrightarrow \quad \begin{array}{l} -\min_{x \in \mathbb{R}^2} -(x_1 - x_2) \\ x_1 + 17x_2 - 10 \leq 0 \\ -2x_1 - 3x_2 + 5 \leq 0 \\ \begin{cases} 4x_1 + x_2 - 9 \leq 0, \\ -4x_1 - x_2 + 9 \leq 0, \end{cases} \end{array}$$



LP by picture

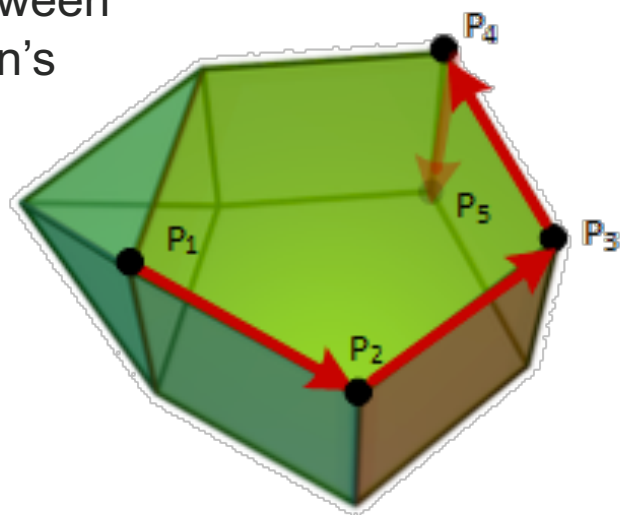


$$\begin{array}{llll} \max_{x \in \mathbb{R}^2} & x_1 & + & x_2 \\ \text{subject to} & x_1 & + & 2x_2 \leq 2 \\ & x_1 & & \geq 0 \\ & & & x_2 \geq 0 \end{array}$$

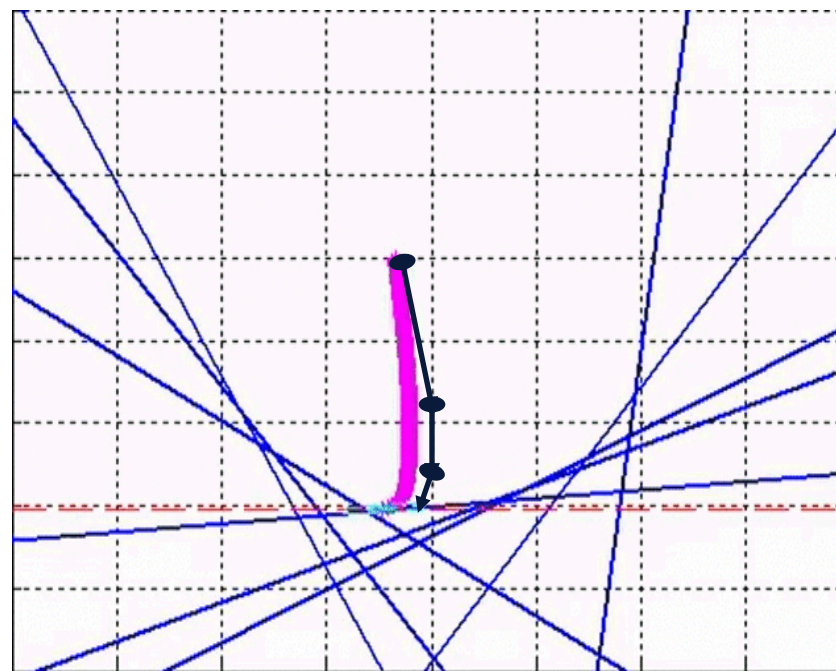
Algorithmic LP ideas

$$\begin{array}{ll} \min_{x \in \mathbb{R}^2} & c^T x \\ \text{subject to} & Ax \leq b \end{array}$$

Traverses
edges between
polyhedron's
vertices



Simplex



Barrier

Generates
sequence
of points in
interior of
polyhedron

Introducing MIP

What is MILP?

$$\begin{array}{llll} \min_{x \in \mathbb{R}^n} & f(x) & & \\ \text{subject to} & h_1(x) & \leq & 0 \\ & h_2(x) & \leq & 0 \\ & \vdots & & \\ & h_m(x) & \leq & 0 \\ & \text{Some } x_i & \text{are integer} & \end{array}$$

- This is a Mixed-Integer Linear Programming (MILP) instance if:
 - The objective function $f(x)$ is linear
 - e.g., $f(x) = 3x_1 + 2x_2$
 - The functions $h_1(x), \dots, h_m(x)$ are linear
 - e.g., $h_1(x) = x_1 - 17x_2 + 10$, etc.
- MILP = LP + integrality constraints!

MIP example: furniture manufacturing



	Stool (x)	Chair (y)	Available
Profit	\$10	\$11	↑↑↑

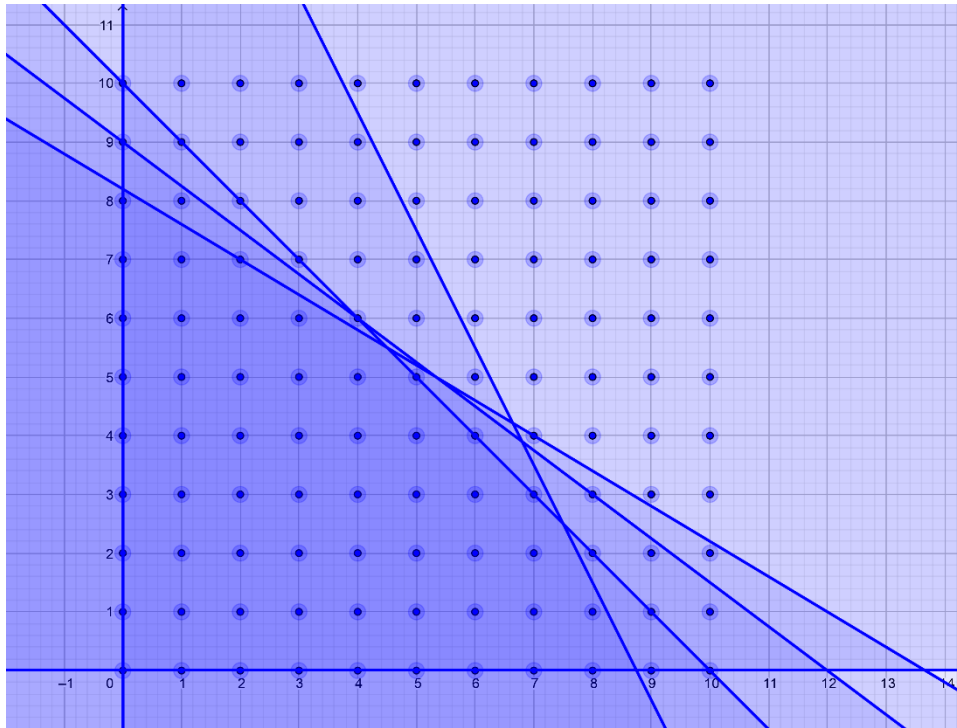
Variables:

- x : number of stools we make
- y : number of chairs we make

Solution: $x^* = 4.5, y^* = 5.5$

$$\begin{aligned} &\Leftrightarrow \max_{x,y} && 10x &+& 11y \\ &\Leftrightarrow \text{subject to} && x &+& y &\leq 10 \\ &\Leftrightarrow && 3x &+& 4y &\leq 36 \\ &\Leftrightarrow && 3x &+& 5y &\leq 41 \\ &\Leftrightarrow && 4x &+& 2y &\leq 35 \\ &&& \text{integer} && x, y &\geq 0 \end{aligned}$$

MIP by picture







$$\begin{array}{llllll} \max_{x,y} & 10x & + & 11y & & \\ \text{subject to} & x & + & y & \leq & 10 \\ & 3x & + & 4y & \leq & 36 \\ & 3x & + & 5y & \leq & 41 \\ & 4x & + & 2y & \leq & 35 \\ & \text{integer } x, y & \geq & 0 & & \end{array}$$

Why not enumerate (try) all the integer feasible points?

Why not enumerate all combinations of 0-1 values?

- Try all 2^n points on Summit supercomputer at ORNL
 - 200 petaflops (2×10^{17})

n	0 – 1 points to check	Time	
10	1024	0 seconds	
60	1.15×10^{18}	6 seconds	
70	1.18×10^{21}	1.6 hours	
120	1.33×10^{36}	210 billion years	

Solving MIP

Bounds and MIP Gap

Let $f(x, y) := 10x + 11y$

$$\begin{aligned} z^* &:= \max_{x,y} && \overbrace{10x + 11y} \\ \text{s. t.} &&& x + y \leq 10 \\ &&& 3x + 4y \leq 36 \\ &&& 3x + 5y \leq 41 \\ &&& 4x + 2y \leq 35 \\ &&& \text{integer } x, y \geq 0 \end{aligned}$$

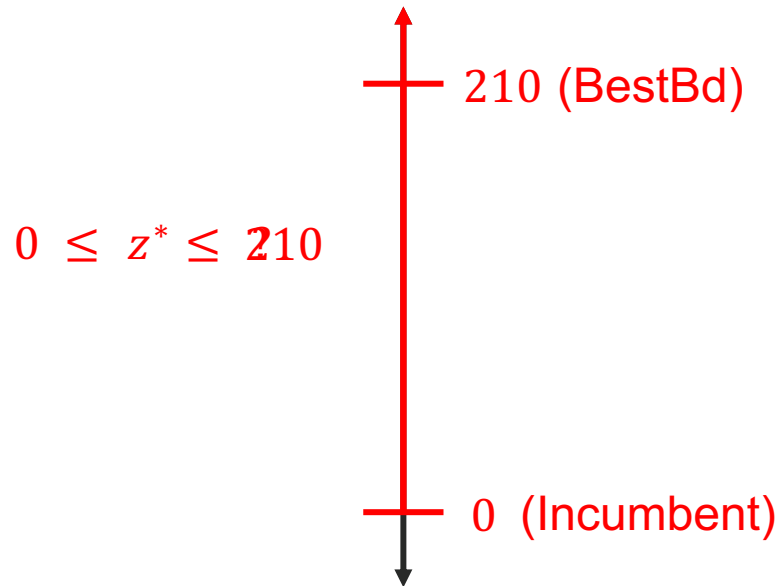
$$0 \leq z^* = ?$$

0 (Incumbent)

Can we try something simple?

- $x = 0, y = 0$ is a feasible solution!
- Then z^* must *at least as large as* $f(0,0) = 0$
- Currently, $(0,0)$ is our *incumbent* solution

Bounds and MIP Gap



What about upper bounds?

- Use linear constraints to get an upper bound on the optimal objective z^*
- All solutions have objective $z^* \leq 210$
 - This is currently our *best (upper) bound*

$$z^* := \max_{x,y} \\ \text{s. t.}$$

$$10x + 11y$$

$$x + y \leq 10$$

$$3x + 4y \leq 36$$

$$3x + 5y \leq 41$$

$$4x + 2y \leq 35$$

$$\text{integer } x, y \geq 0$$

$$x + y \leq 10$$



$$x \leq 10 \text{ and } y \leq 10$$

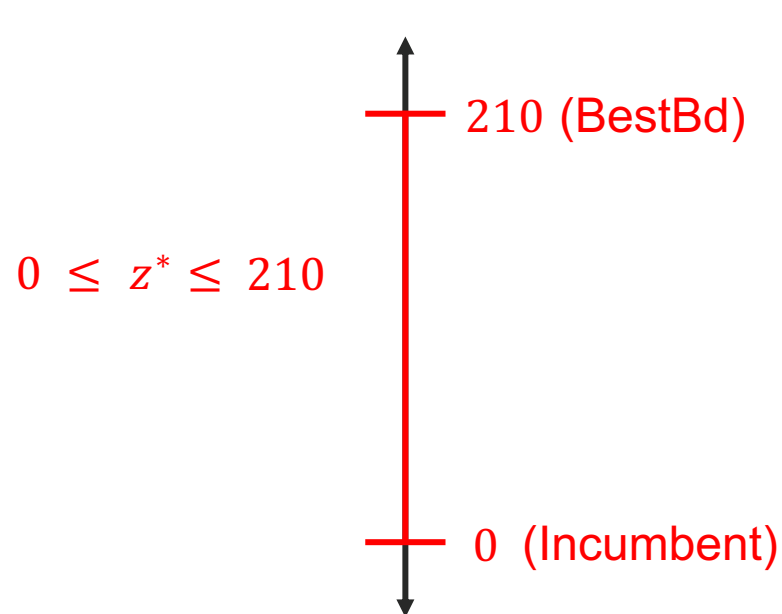


$$10x \leq 100 \text{ and } 11y \leq 110$$



$$f(x,y) = 10x + 11y \leq 210$$

Bounds and MIP Gap

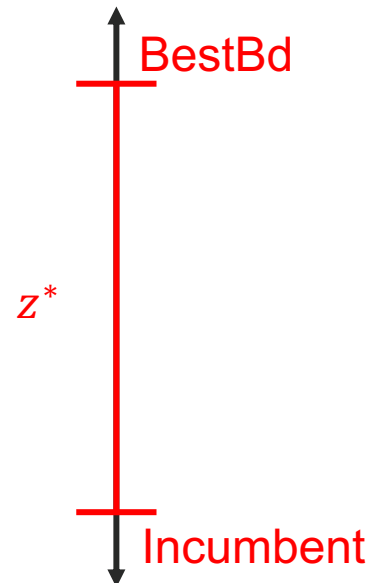


$$\begin{aligned} z^* := \max_{x,y} \quad & 10x + 11y \\ \text{s. t.} \quad & x + y \leq 10 \\ & 3x + 4y \leq 36 \\ & 3x + 5y \leq 41 \\ & 4x + 2y \leq 35 \\ & \text{integer } x, y \geq 0 \end{aligned}$$

How good is the incumbent?

- Measure the quality by the *(relative) MIP gap*: $\frac{|\text{BestBd} - \text{Incumbent}|}{|\text{Incumbent}|} \times 100\%$
- Current gap is $\frac{210-0}{0} \times 100\%$ (undefined or $+\infty$)
 - This is a very poor gap!

Bounds and MIP Gap

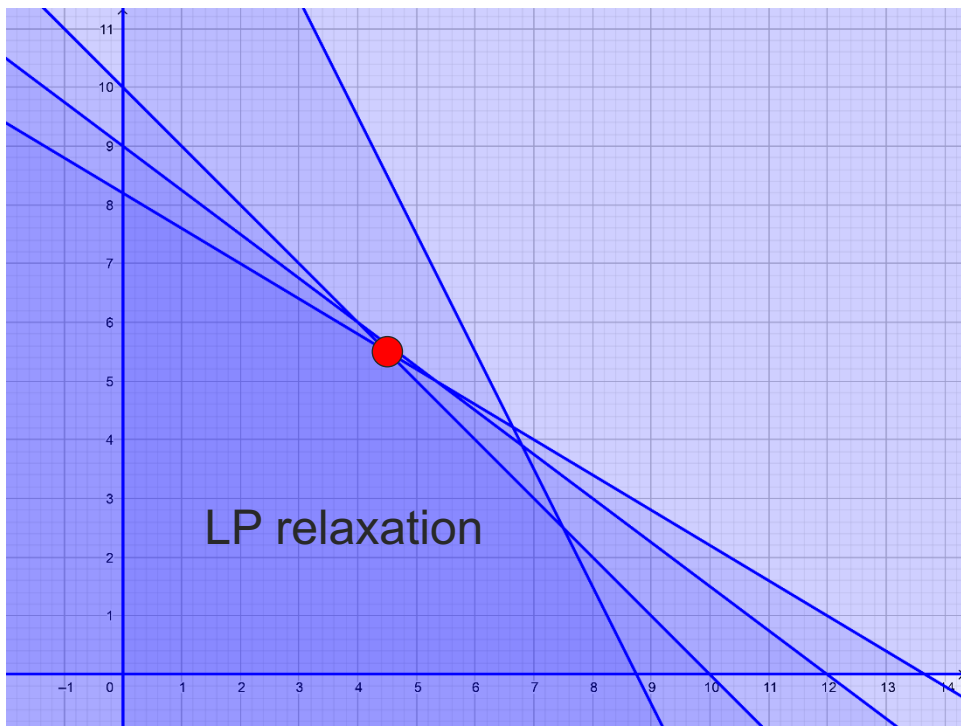


$$\begin{aligned} z^* := \max_{x,y} \quad & 10x + 11y \\ \text{s. t.} \quad & x + y \leq 10 \\ & 3x + 4y \leq 36 \\ & 3x + 5y \leq 41 \\ & 4x + 2y \leq 35 \\ & \text{integer } x, y \geq 0 \end{aligned}$$

Can we do better?

- A MIPGap of 0 means the incumbent solution is optimal
 - This happens when the bound and the incumbent have equal value
- Our goal: try to improve the bound and/or incumbent

Upper bound from relaxation

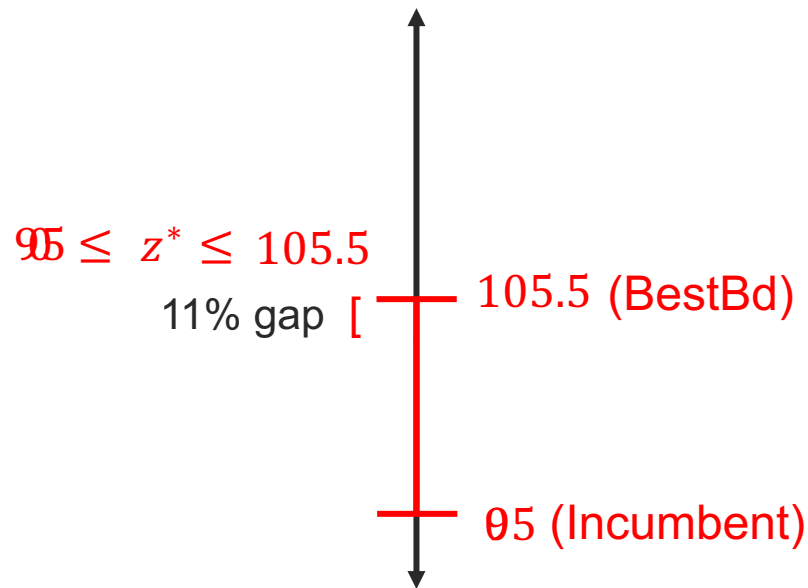


$$\begin{aligned} z^{\text{LP}} := \max_{x,y} \quad & 10x + 11y \\ \text{s. t.} \quad & x + y \leq 10 \\ & 3x + 4y \leq 36 \\ & 3x + 5y \leq 41 \\ & 4x + 2y \leq 35 \\ & \text{integer } x, y \geq 0 \end{aligned}$$

Can we improve our upper bound?

- The integrality constraints are hard...let's temporarily ignore them
- *LP (root) relaxation*: relax (remove) integrality restrictions
- Optimal solution is (4.5, 5.5)
 $\Rightarrow z^{\text{LP}} = f(4.5, 5.5) = 105.5$ is an upper bound on z^* !

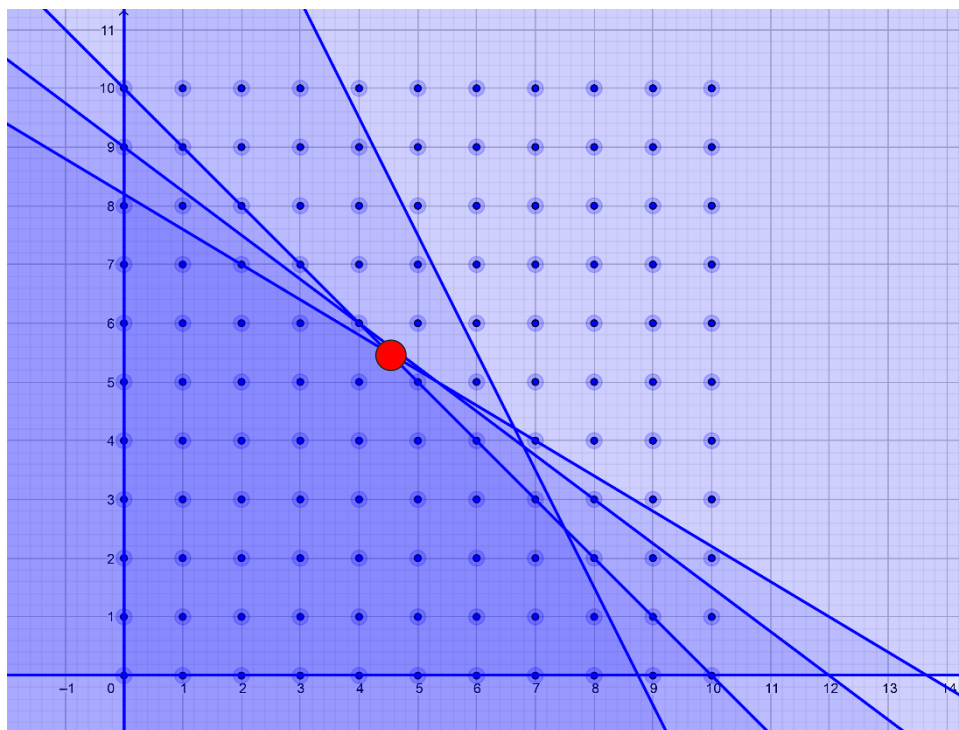
Lower bound from heuristics



$$\begin{aligned}
 z^* &:= \max_{x,y} && 10x &+& 11y \\
 \text{s. t.} &&& x &+& y &\leq 10 \\
 &&& 3x &+& 4y &\leq 36 \\
 &&& 3x &+& 5y &\leq 41 \\
 &&& 4x &+& 2y &\leq 35 \\
 &&& \text{integer } x, y &\geq 0
 \end{aligned}$$

Can we improve our incumbent?

- Let's "guess" at a good solution
 - This is commonly called a "heuristic"
- Try: round down the relaxed solution (4.5, 5.5):
 $\Rightarrow ([4.5], [5.5]) = (4, 5)$ is feasible!
- $f(4, 5) = 95$...we have improved our incumbent!



$$\begin{aligned} z^* &:= \max_{x,y} && 10x &+& 11y \\ \text{s. t.} &&& x &+& y &\leq 10 \\ &&& 3x &+& 4y &\leq 36 \\ &&& 3x &+& 5y &\leq 41 \\ &&& 4x &+& 2y &\leq 35 \\ &&& \text{integer } x, y &\geq 0 \end{aligned}$$

LP optimum:
 $x = 4.5, y = 5.5$

Fractional solutions are bad!

- Is there an inequality for our MIP that “cuts off” the LP optimum $(4.5, 5.5)$?
- *Cuts*: inequalities that tighten our relaxation of the MIP feasible region
- Let’s try to derive a cut that is violated by $(4.5, 5.5)$!

Cuts

LP optimum:
 $x = 4.5, y = 5.5$

$$\begin{array}{ll}
 \max & 10x + 11y \\
 \text{s. t.} & x + y \leq 10 \\
 & 3x + 4y \leq 36 \\
 & 3x + 5y \leq 41 \\
 & 4x + 2y \leq 35 \\
 & x, y, s_1, s_2, s_3, s_4 \text{ integer} \\
 & s_1, s_2, s_3, s_4, x, y \geq 0
 \end{array}$$

Satisfied
without
slack at
 $(4.5, 5.5)$

1. Aggregate

$$-3x - 3y - 3s_1 = -30$$

and

$$3x + 5y + s_3 = 41$$

\Downarrow

$$2y - 3s_1 + s_3 = 11$$

\Downarrow

$$y - \frac{3}{2}s_1 + \frac{1}{2}s_3 = \frac{11}{2}$$

2. Weaken

$$y - \frac{3}{2}s_1 + \frac{1}{2}s_3 = \frac{11}{2}$$

\Downarrow

$$y - \left\lfloor -2\left(-\frac{3}{2}\right) \right\rfloor s_1 + \left\lfloor \frac{1}{2} \right\rfloor s_3 \leq \left\lfloor \frac{11}{2} \right\rfloor$$

\Downarrow

$$2x + 3y \leq 25$$

3. Strengthen

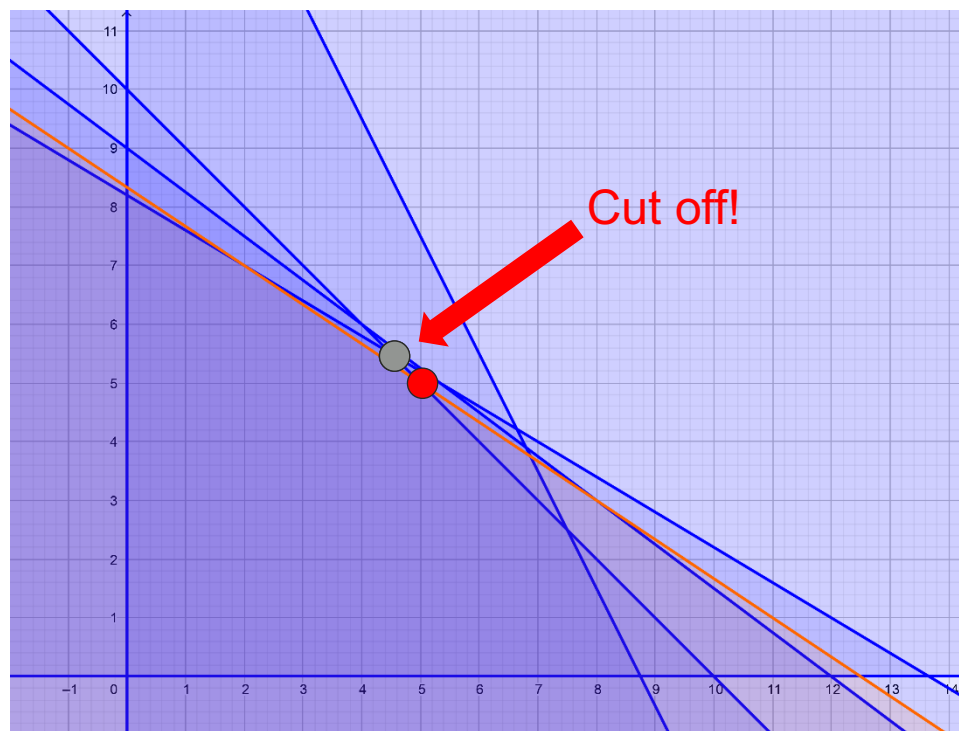
$$y - 2s_1 \leq \frac{11}{2}$$

\Downarrow

$$y - 2s_1 \leq \left\lfloor \frac{11}{2} \right\rfloor = 5$$

**New
inequality!**

Cuts



LP optimum:

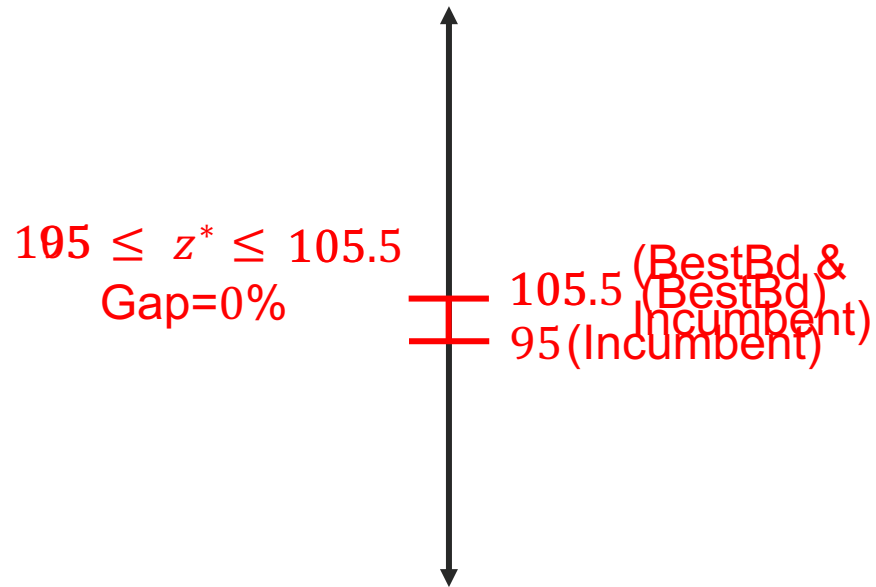
$$x = 4.5, y = 5.5$$

$$\begin{aligned} z^* &:= \max_{x,y} && 10x &+& 11y \\ \text{s. t.} &&& x &+& y &\leq 10 \\ &&& 3x &+& 4y &\leq 36 \\ &&& 3x &+& 5y &\leq 41 \\ &&& 4x &+& 2y &\leq 35 \\ &&& 3x &+& 2y &\leq 25 \\ &&& \text{integer } x, y &\geq 0 \end{aligned}$$

Is this new inequality any good?

- It cuts off the previous optimal solution!
- Solving the new, tightened relaxation yields a new upper bound
- New upper bound on z^* is $10(5) + 11(5) = 105$

Improving bounds

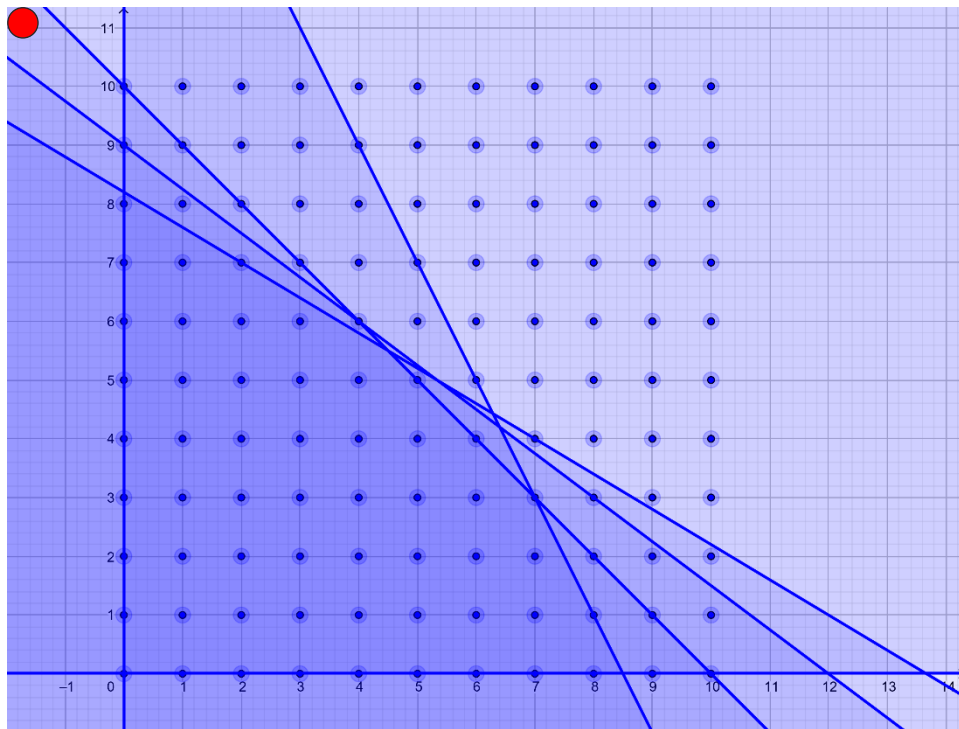


$$\begin{aligned} z^* &:= \max_{x,y} && 10x &+& 11y \\ \text{s. t.} &&& x &+& y &\leq 10 \\ &&& 3x &+& 4y &\leq 36 \\ &&& 3x &+& 5y &\leq 41 \\ &&& 4x &+& 2y &\leq 35 \\ &&& \text{integer } x, y &\geq 0 \end{aligned}$$

What can we say about z^* ?

- New upper bound: $z^* \leq 105$
- But the optimal solution to the relaxation is (5, 5)
 - This is feasible to the original problem!
 - 105 is a new incumbent objective value
- **We have solved the problem: $z^* = 105$**

Presolve

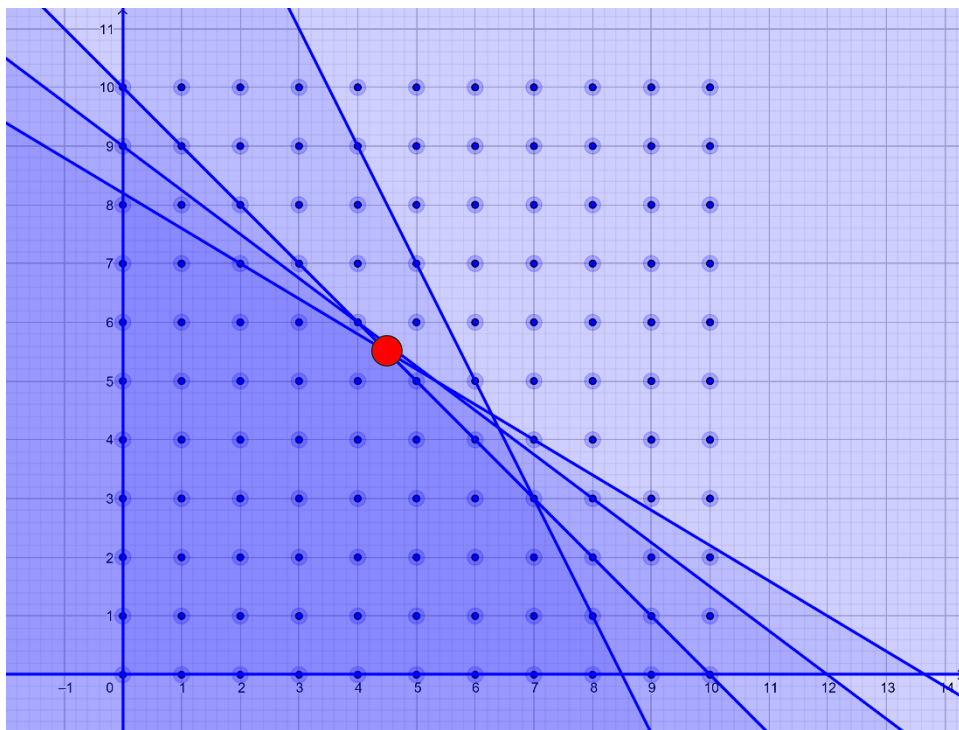


$$\begin{aligned} z^* &:= \max_{x,y} && 10x &+& 11y \\ &\text{s. t.} && x &+& y &\leq 10 \\ &&& 3x &+& 4y &\leq 36 \\ &&& 3x &+& 5y &\leq 41 \\ &&& 4x &+& 2y &\leq 34 \\ &&& \text{integer } x, y &\geq 0 \end{aligned}$$

Can we tighten the formulation just by looking at it?

- *Tighten*: reduce the feasible region of relaxation without removing any feasible points
- $4x$ is even, and so is $2y$...
- This doesn't change optimal solution relaxation

Branch-and-bound

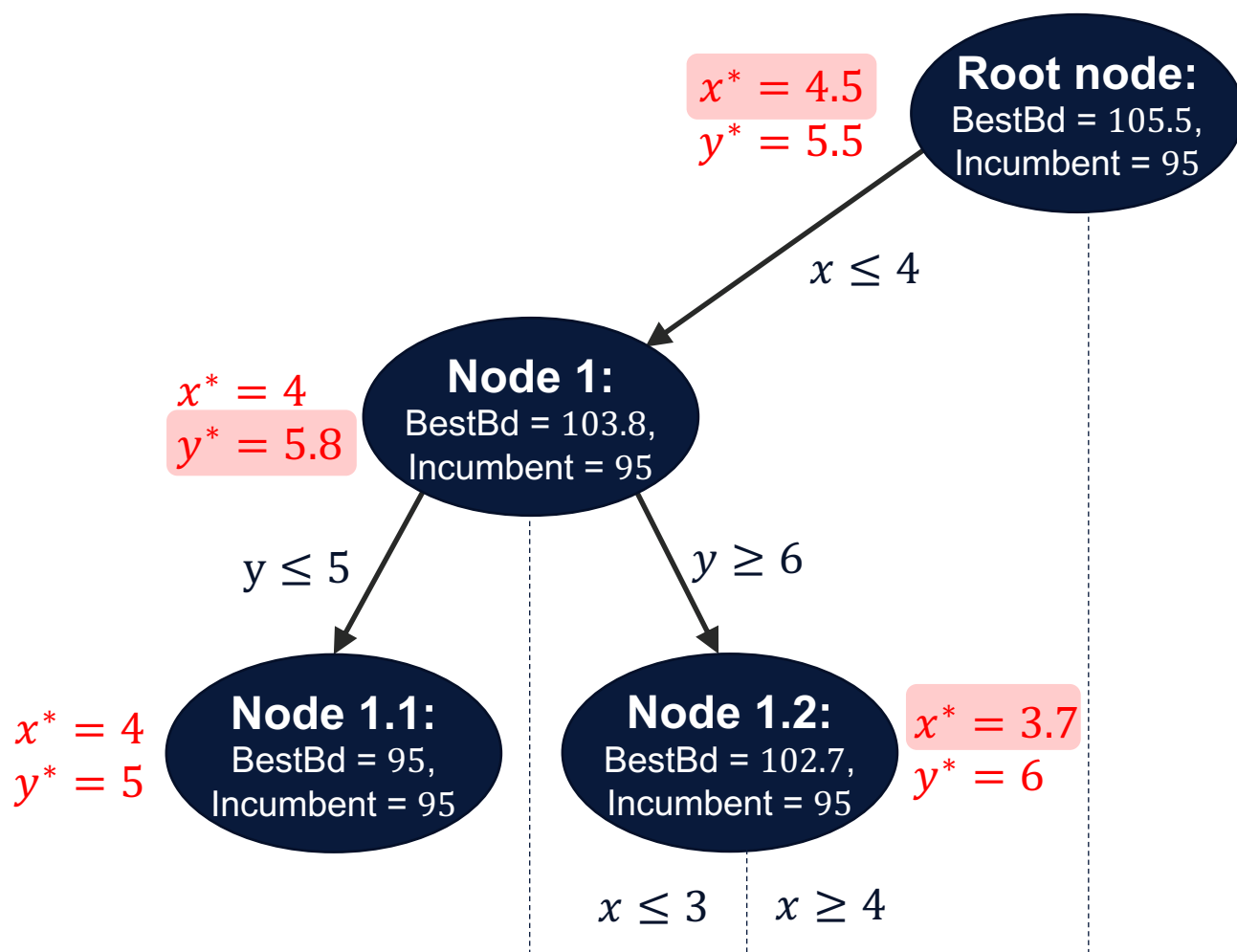


$$\begin{aligned} z^* &:= \max_{x,y} && 10x &+& 11y \\ &\text{s. t.} && x &+& y &\leq 10 \\ &&& 3x &+& 4y &\leq 36 \\ &&& 3x &+& 5y &\leq 41 \\ &&& 4x &+& 2y &\leq 34 \\ &&& \text{integer } x, y &\geq 0 \end{aligned}$$

Can we combine these ideas with brute-force enumeration?

- Yes! For example, consider the relaxation solution $(x, y) = (4.5, 5.5)$
 - x has to be integer. So is $x \leq 4$ or $x \geq 5$?
 - Let's investigate!

Branch-and-bound tree

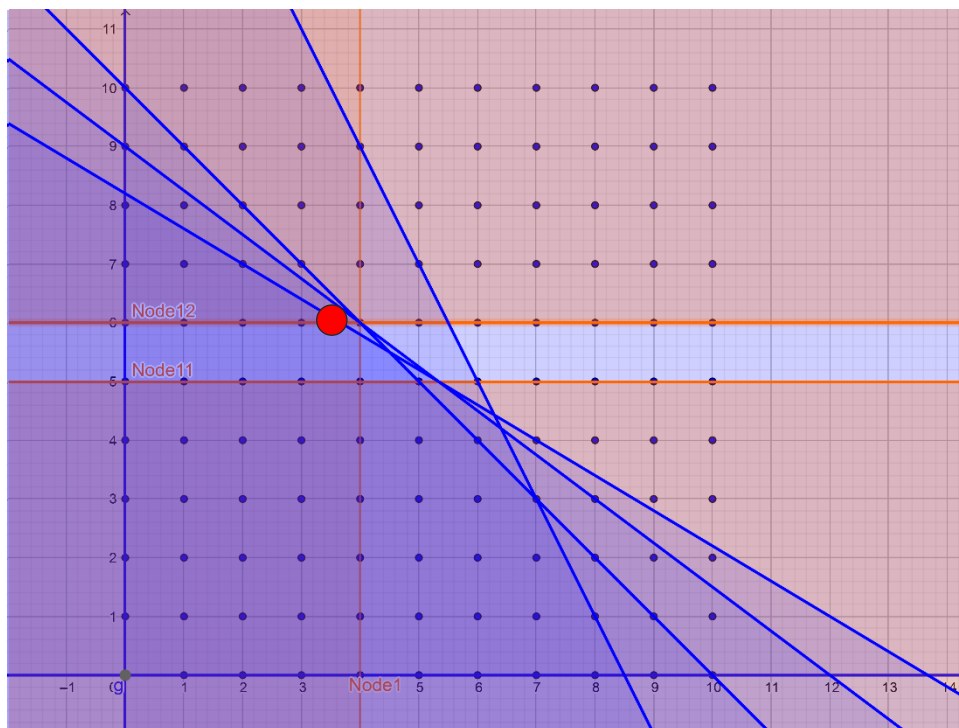


$$z^* := \max_{x,y} \quad 10x + 11y$$
$$\text{s. t.} \quad \begin{array}{rclcl} x & + & y & \leq & 10 \\ 3x & + & 4y & \leq & 36 \\ 3x & + & 5y & \leq & 41 \\ 4x & + & 2y & \leq & 34 \\ & & \text{integer } x, y & \geq & 0 \end{array}$$

$x \leq 4$

$y \geq 5$

Branch-and-bound tree



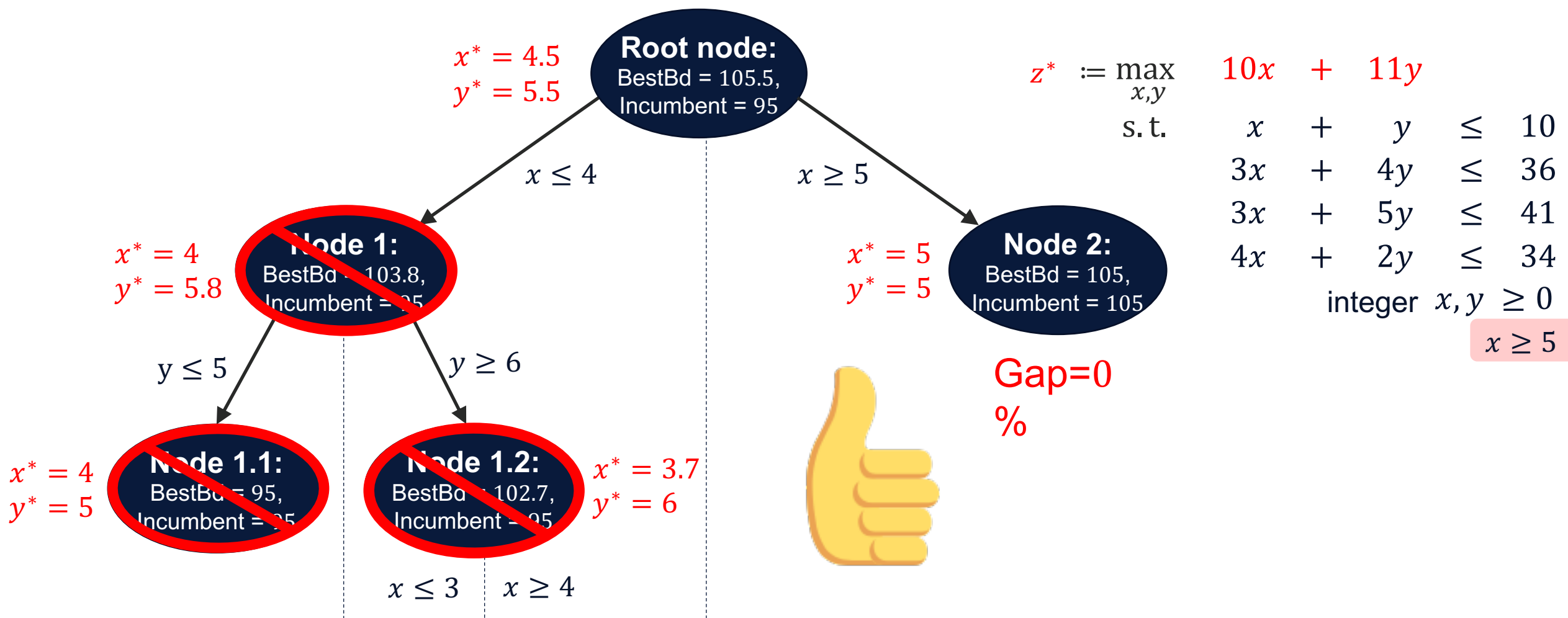
Geometrically, what does this look like?

$$\begin{aligned} z^* &:= \max_{x,y} && 10x &+& 11y \\ &\text{s. t.} && x &+& y &\leq 10 \\ &&& 3x &+& 4y &\leq 36 \\ &&& 3x &+& 5y &\leq 41 \\ &&& 4x &+& 2y &\leq 34 \\ &&& \text{integer } x, y &\geq 0 \end{aligned}$$

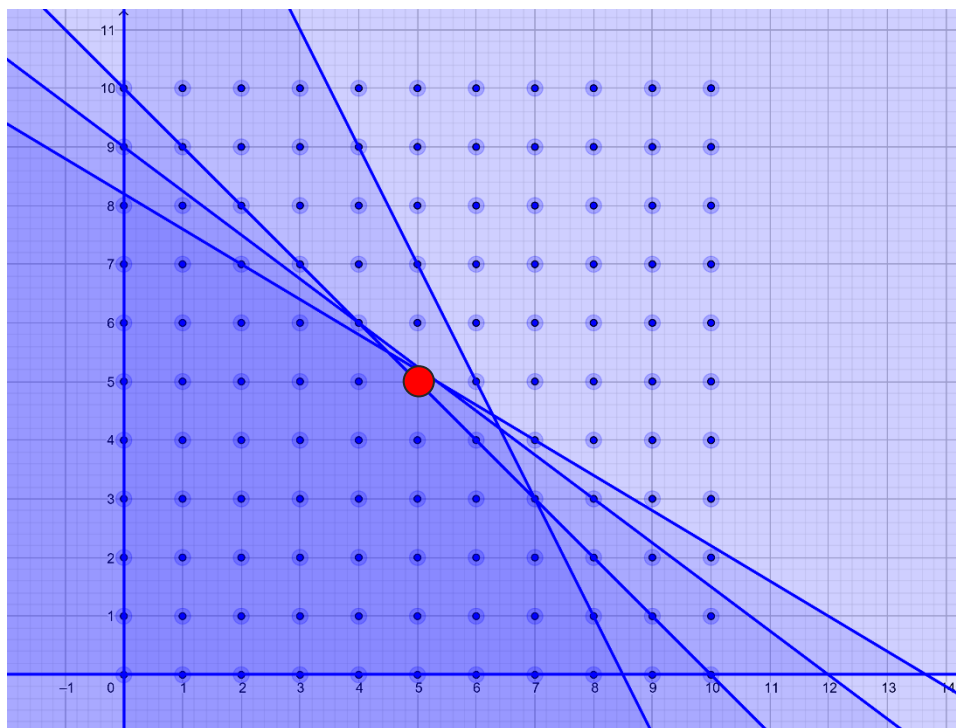
$x \leq 4$

$y \geq 5$

Branch-and-bound tree



Branch-and-bound



$$\begin{aligned} z^* &:= \max_{x,y} && 10x &+& 11y \\ \text{s. t.} &&& x &+& y &\leq 10 \\ &&& 3x &+& 4y &\leq 36 \\ &&& 3x &+& 5y &\leq 41 \\ &&& 4x &+& 2y &\leq 34 \\ &&& \text{integer } x, y &\geq 0 \end{aligned}$$

Looks like a lot of work!

- For the computer – not for you!
- Gurobi does presolve, cuts, heuristics, branch-and-bound and more “auto-magically”!

Termination criteria

- **MIP gap**
 - E.g., terminate if *relative MIP gap* $\leq 1\%$ or *absolute MIP gap* ≤ 10 ,
- **Time spent solving**
 - E.g., terminate after *3600 seconds* have elapsed
- **Number of nodes explored**
 - E.g., terminate after solver visits *100,000* nodes
- **Number of feasible solutions found**
- ...and more!

See the Termination section of the **Parameters** page in the Reference Manual

Gurobi logs

Recognize these ideas in the Gurobi logs

Log file structure

LP

1. User input + model statistics
2. Presolve
3. Simplex or Barrier iterations
4. Termination statistics

MILP

1. User input + model statistics
2. Presolve
3. Root node relaxation
4. Branch-and-bound tree exploration
5. Termination statistics

Simplex log

Iteration	Objective	Primal Inf.	Dual Inf.	Time
0	1.7748600e+04	6.627132e+03	0.000000e+00	0s
9643	1.1574611e+07	1.418653e+03	0.000000e+00	5s
14440	1.1607748e+07	4.793500e+00	0.000000e+00	10s
15213	1.1266396e+07	0.000000e+00	0.000000e+00	11s

Solved in 15213 iterations and 10.86 seconds
Optimal objective 1.126639605e+07

Barrier log

Barrier statistics:

AA' NZ : 2.836e+03
Factor NZ : 3.551e+03 (roughly 40 MBytes of memory)
Factor Ops : 1.739e+05 (less than 1 second per iteration)
Threads : 4

Iter	Objective		Residual		Compl	Time
	Primal	Dual	Primal	Dual		
0	1.30273209e+06	0.00000000e+00	5.90e+02	0.00e+00	7.32e+00	12s
1	1.04326180e+05	-5.84079103e+02	4.84e+01	1.69e+00	5.95e-01	12s
2	9.46325157e+03	-4.40392705e+02	2.92e+00	1.35e+00	5.46e-02	12s
3	3.66683689e+03	9.27381244e+02	1.94e-01	5.35e-01	1.41e-02	12s
4	3.37449982e+03	1.79938013e+03	1.29e-01	2.41e-01	7.64e-03	12s
5	3.13244138e+03	1.90266941e+03	8.89e-02	2.07e-01	6.00e-03	12s
6	2.71282610e+03	2.11401255e+03	3.20e-02	1.15e-01	2.96e-03	12s
7	2.48856811e+03	2.18107490e+03	1.06e-02	7.26e-02	1.56e-03	12s
8	2.35427593e+03	2.21183615e+03	3.20e-03	4.52e-02	7.36e-04	12s
9	2.30239737e+03	2.22464753e+03	1.53e-03	2.38e-02	4.03e-04	12s
10	2.25547118e+03	2.23096162e+03	3.00e-04	1.40e-02	1.30e-04	12s
11	2.24052450e+03	2.23917612e+03	4.10e-06	6.33e-04	7.20e-06	12s
12	2.23967243e+03	2.23966346e+03	2.01e-08	5.01e-06	4.82e-08	12s
13	2.23966667e+03	2.23966666e+03	1.11e-10	1.14e-13	4.81e-11	13s

Barrier solved model in 13 iterations and 12.51 seconds
Optimal objective 2.23966667e+03

Barrier crossover log

Barrier solved model in 13 iterations and 12.51 seconds
Optimal objective 2.23966667e+03

Root crossover log...

40 DPushes remaining with DInf 0.0000000e+00	13s
0 DPushes remaining with DInf 7.8159701e-14	13s
1176 PPushes remaining with PInf 0.0000000e+00	13s
0 PPushes remaining with PInf 0.0000000e+00	13s
Push phase complete: Pinf 0.0000000e+00, Dinf 1.2079227e-13	13s

Root simplex log...

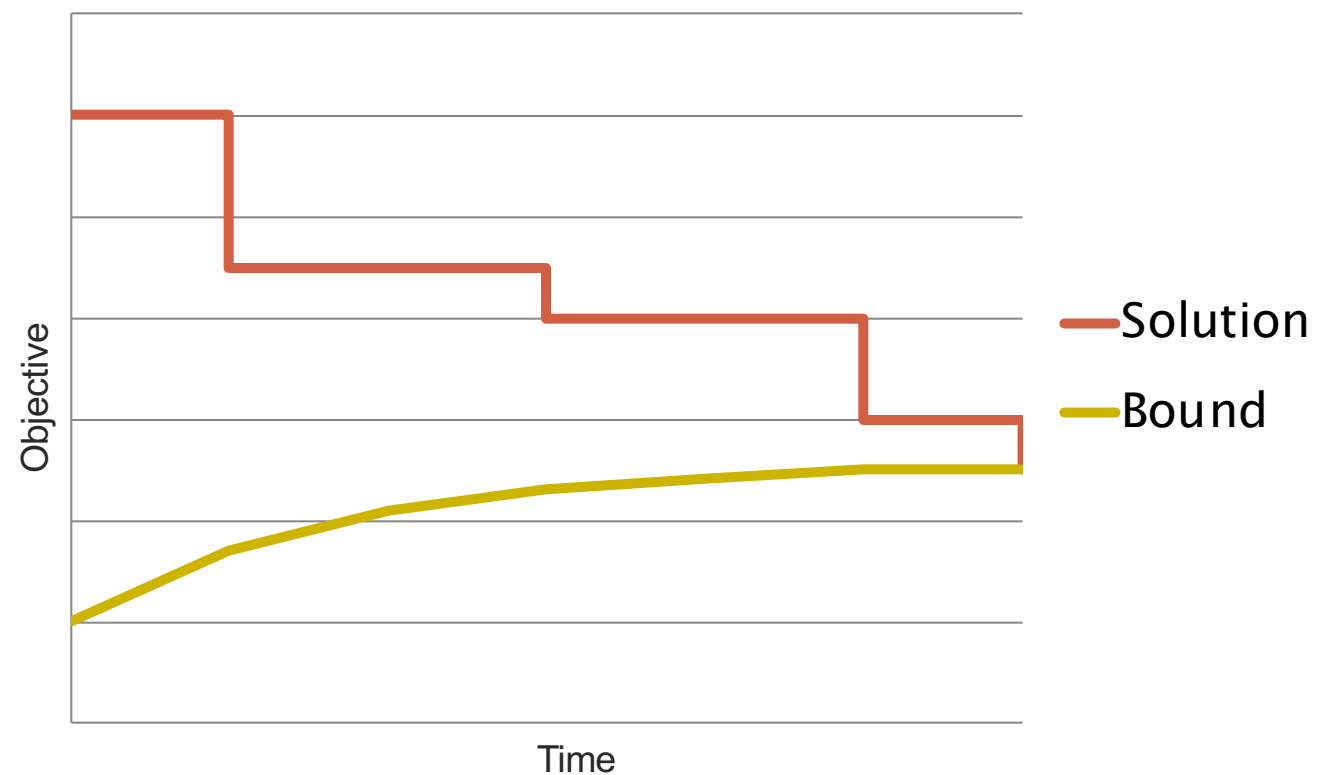
Iteration	Objective	Primal Inf.	Dual Inf.	Time
1219	2.2396667e+03	0.0000000e+00	0.0000000e+00	13s
1219	2.2396667e+03	0.0000000e+00	0.0000000e+00	13s

Root relaxation: objective 2.239667e+03, 1219 iterations, 0.43 seconds

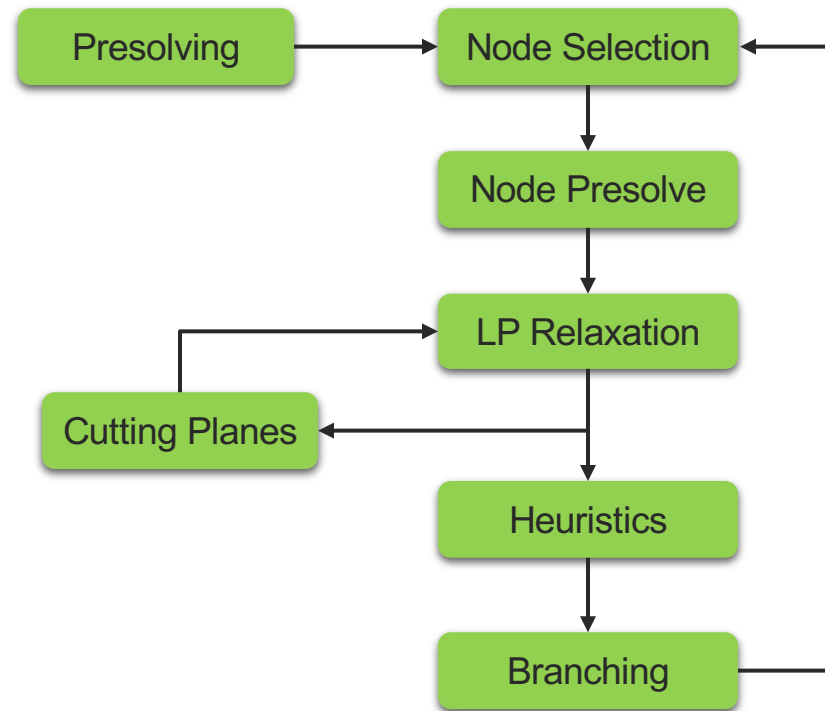
MIP tree exploration

Nodes			Current Node			Objective Bounds			Work	
Expl	Unexpl		Obj	Depth	IntInf	Incumbent	BestBd	Gap	It/Node	Time
	0	0	8.0000e+08	0	72	3.1334e+09	8.0000e+08	74.5%	–	0s
H	0	0				2.400019e+09	8.0000e+08	66.7%	–	0s
	0	0	8.0000e+08	0	72	2.4000e+09	8.0000e+08	66.7%	–	0s
H	0	0				2.066683e+09	8.0000e+08	61.3%	–	0s
	0	0	8.0000e+08	0	72	2.0667e+09	8.0000e+08	61.3%	–	0s
	0	0	8.0000e+08	0	72	2.0667e+09	8.0000e+08	61.3%	–	0s
	0	2	8.0000e+08	0	72	2.0667e+09	8.0000e+08	61.3%	–	0s
H	796	711				2.050016e+09	8.0000e+08	61.0%	4.4	0s
H	796	675				2.050016e+09	8.0000e+08	61.0%	4.4	0s
*	3353	2003		109		2.000021e+09	8.0000e+08	60.0%	4.5	0s
H	5097	3281				2.000019e+09	8.0000e+08	60.0%	4.1	0s
H	5571	3544				1.950018e+09	8.0000e+08	59.0%	4.0	0s
*	6228	4017		82		1.940019e+09	8.0000e+08	58.8%	3.9	1s
*	6347	4038		97		1.925019e+09	8.0000e+08	58.4%	3.9	1s
H	6422	3912				1.833347e+09	8.0000e+08	56.4%	3.9	1s
*	7999	4981		104		1.816683e+09	8.0000e+08	56.0%	3.8	1s
H	9958	6572				1.800016e+09	8.0000e+08	55.6%	3.8	1s

Solving a MIP model



Branch-and-cut



Branch-and-cut

Presolving

Cutting Planes

```
Gurobi Optimizer version 6.0.0 (linux64)
Copyright (c) 2014, Gurobi Optimization, Inc.

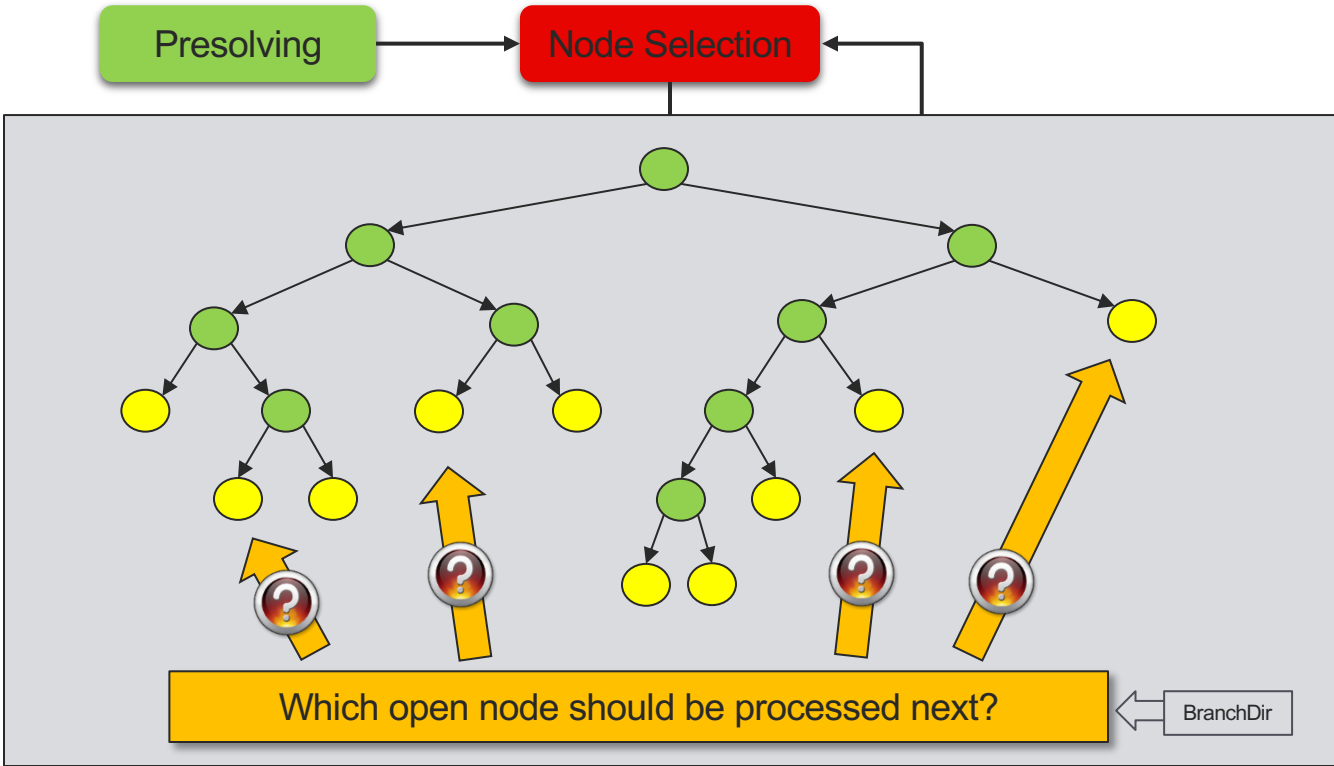
Read MPS format model from file /models/mip/roll3000.mps.bz2
Reading time = 0.03 seconds
roll3000: 2295 rows, 1166 columns, 29386 nonzeros
Optimize a model with 2295 rows, 1166 columns and 29386 nonzeros
Coefficient statistics:
  Matrix range    [2e-01, 3e+02]
  Objective range [1e+00, 1e+00]
  Bounds range    [1e+00, 1e+09]
  RHS range       [6e-01, 1e+03]
Presolve removed 1308 rows and 311 columns
Presolve time: 0.08s
Presolved: 987 rows, 855 columns, 19346 nonzeros
Variable types: 211 continuous, 644 integer (545 binary)

Root relaxation: objective 1.112003e+04, 1063 iterations, 0.03 seconds

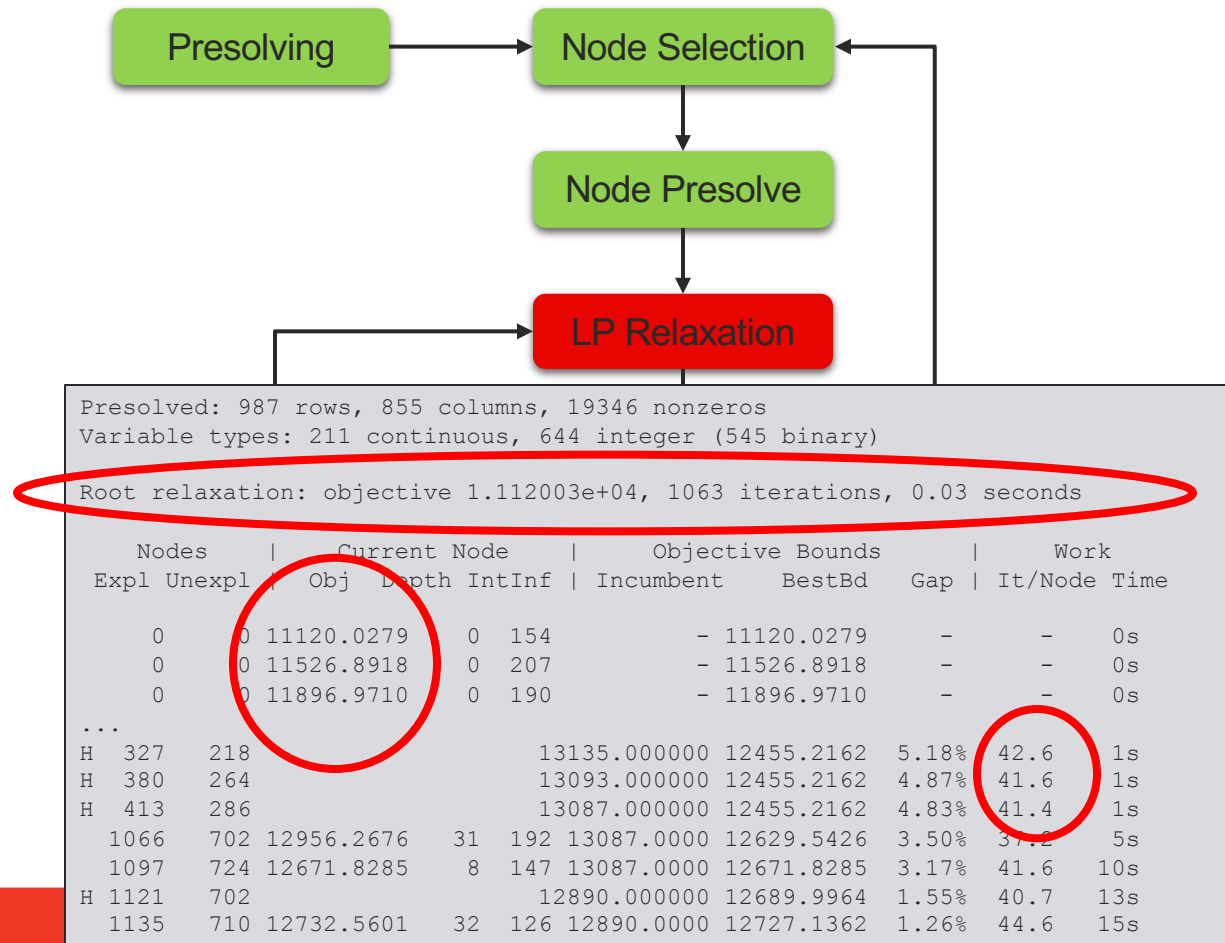
   Nodes      |   Current Node   |   Objective Bounds   |   Work
 Expl Unexpl |  Obj  Depth IntInf | Incumbent    BestBd  Gap   | It/Node Time
-----
    0     0 11120.0279    0  154      - 11120.0279    -     -    0s
    0     0 11526.8918    0  207      - 11526.8918    -     -    0s
    0     0 11896.9710    0  190      - 11896.9710    -     -    0s
```

Branching

Branch-and-cut



Branch-and-cut



Branch-and-cut

Presolving

Cutting Planes

```
Presolved: 987 rows, 855 columns, 19346 nonzeros
Variable types: 211 continuous, 644 integer (545 binary)

Root relaxation: objective 1.112003e+04, 1063 iterations, 0.03 seconds

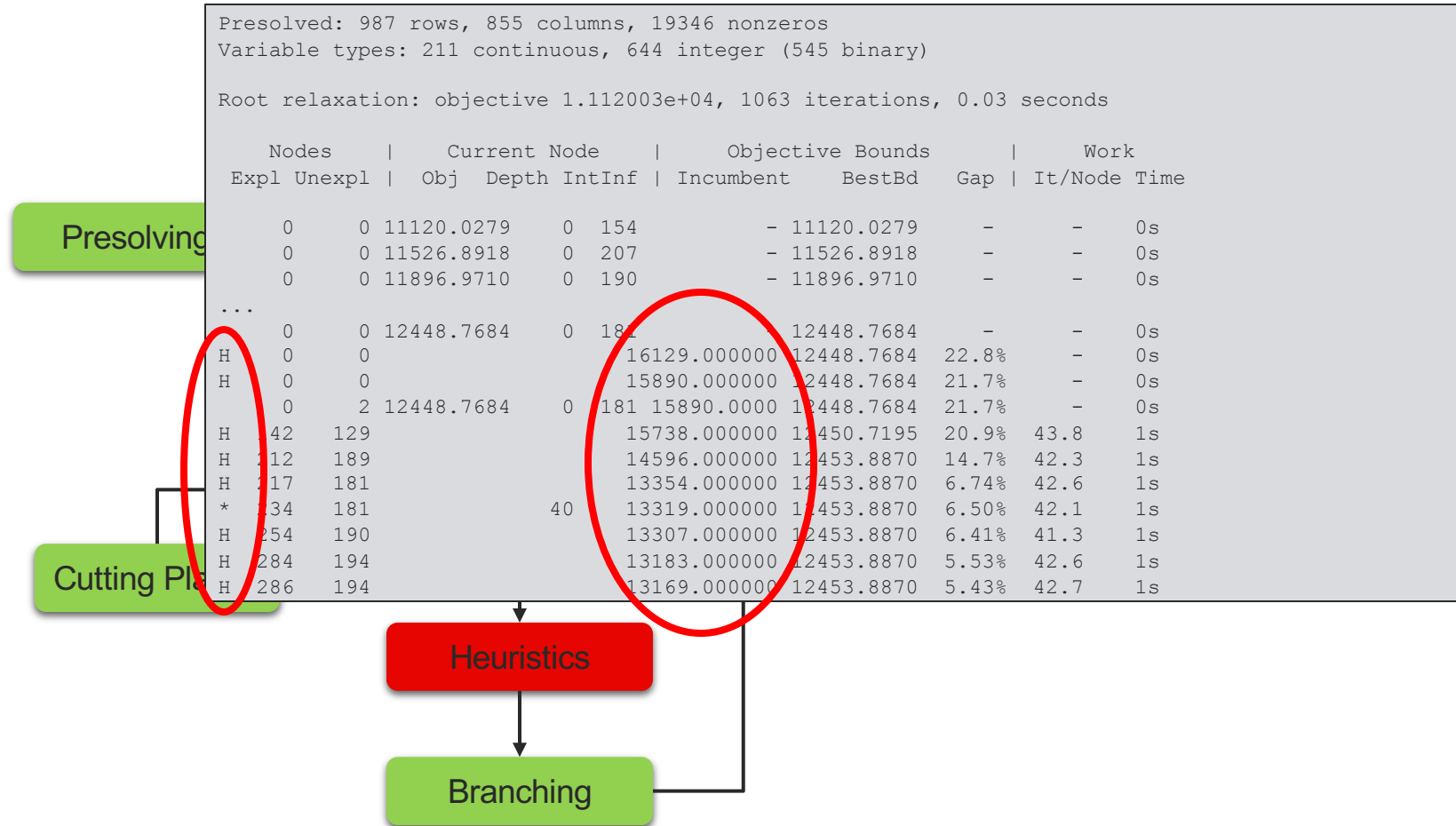
  Nodes      |      Current Node      |      Objective Bounds      |      Work
 Expl Unexpl |  Obj  Depth IntInf | Incumbent    BestBd   Gap | It/Node Time

    0     0 11120.0279    0 154      - 11120.0279    -     -    -    0s
    0     0 11526.8918    0 207      - 11526.8918    -     -    -    0s
    0     0 11896.9710    0 190      - 11896.9710    -     -    -    0s
    0     0 12151.4022    0 190      - 12151.4022    -     -    -    0s
    0     0 12278.3391    0 208      - 12278.3391    -     -    -    0s
...
5489    634 12885.3652   52 143 12890.0000 12829.0134  0.47%  54.5   25s

Cutting planes:
  Learned: 4
  Gomory: 46
  Cover: 39
  Implied bound: 8
  Clique: 2
  MIR: 112
  Flow cover: 27
  GUB cover: 11
  Zero half: 91

Explored 6808 nodes (357915 simplex iterations) in 27.17 seconds
Thread count was 4 (of 8 available processors)
```

Branch-and-cut



Branch-and-cut

Presolving

```
Presolved: 987 rows, 855 columns, 19346 nonzeros
Variable types: 211 continuous, 644 integer (545 binary)

Root relaxation: objective 1.112003e+04, 1063 iterations, 0.03 seconds
```

Nodes		Current Node			Objective Bounds			Work	
Expl	Unexpl	Obj	Depth	IntInf	Incumbent	BestBd	Gap	It/Node	Time
0	0	11120.0279	0	154	-	11120.0279	-	-	0s
0	0	11526.8918	0	207	-	11526.8918	-	-	0s
0	0	11896.9710	0	190	-	11896.9710	-	-	0s
...									
H	0	0			15890.000000	12448.7684	21.7%	-	0s
0	2	12448.7684	0	181	15890.0000	12448.7684	21.7%	-	0s
...									
1066	702	12956.2676	31	192	13087.0000	12629.5426	3.50%	37.2	5s
1097	724	12671.8285	8	147	13087.0000	12671.8285	3.17%	41.6	10s
1135	710	12732.5601	32	126	12890.0000	12727.1362	1.26%	44.6	15s
3416	887	12839.9880	46	136	12890.0000	12780.7059	0.85%	49.7	20s
5485	634	12885.3652	52	143	12890.0000	12829.0134	0.47%	54.5	25s

Cutting P

Branching

Outlook: Algorithmic Performance

- Solving optimization problems is never fast enough
- **Great solver performance is the sum of many parts**
 - Model construction speed
 - Selection and configuration of the underlying algorithms
 - Hardware
 - ...and a lot more
- **The default settings of Gurobi are carefully selected to perform best on a large variety of model instance**
- **However, there is usually some tuning potential left, e.g.**
 - The model formulation itself
 - More efficient use of the API
 - Parameter changes
 - ...and a lot more

Thank you!



GUROBI
OPTIMIZATION

The World's Fastest Solver

Your Next Steps

- Try Gurobi 9.0 Now!
 - www.gurobi.com/free-trial
 - Get a 30-day commercial trial license of Gurobi at academic and research licenses are free!
- For questions about Gurobi pricing, please contact sales@gurobi.com or sales@gurobi.de
- **Need help leveraging Optimization for your business?**
 - Contact us at info@gurobi.com
- A recording of this webinar, including the slides, will be available in roughly one week.

This Week!

- **Wednesday, June 24th**
 - Compute Server and Cloud Demo
 - Introduction to Algorithms
 - Advanced Algorithms
 - Ask the Experts: R&D Team
- **Thursday, June 25th**
 - Customer Case Study: KPMG
 - Customer Case Study: NFL
 - Ask the Experts: Fireside chat with Michael North, VP of Planning & Scheduling at NFL