## Homework 5

Topic: NP and NP-Completeness

### 1. (10 Points) Interval Scheduling. Consider the decision version of the problem:

Given a collection of intervals in the timeline, and a bound k, does the collection contain a subset of non-overlapping intervals of size at least k?

For each statement below, state whether it is "True", "False", or "Unknown", and briefly justify your answer.

- (a) (2 Points) Interval Scheduling is in P.
- (b) (2 Points) Interval Scheduling is in NP.
- (c) (2 Points) Interval Scheduling is NP-complete.
- (d) (2 Points) Interval Scheduling  $\leq_p$  Vertex Cover.
- (e) (2 Points) Independent Set  $\leq_p$  Interval Scheduling.

### 2. (35 Points) Vertex Cover on Trees.

For a graph G = (V, E) with |V| = n nodes and |E| = m edges, a *vertex cover* is defined as a subset of nodes  $S \subset V$  s.t. that all edges e are covered (*i.e.*, each edge is touched by at least one node in S). Let us restrict our attention specifically to graphs that are trees.

- (a) (15 Points) Design a greedy algorithm that finds the minimum size vertex cover on a tree; or argue that such an algorithm does not exist.
- (b) (15 Points) Design a dynamic programming algorithm that finds a minimum size vertex cover on a tree in O(m+n) time. (Note: justify why the running time of your algorithm is indeed O(m+n).)
- (c) (5 Points) Is the polynomial time algorithm in (b) in contradiction with the fact that the vertex cover is a well-known NP-complete problem? Please explain.

#### 3. (30 Points) The Path Selection Problem.

Consider the path selection problem stated below:

Given a directed graph G = (V, E), a set of paths  $P_1, P_2, ...P_c$ , and an integer k > 0, is it possible to select at least k out of the c paths so that no two of the selected paths share any nodes?

- (a) (20 points) Prove that the path selection problem is NP-complete.
- (b) (10 points) Is this a decision or an optimization problem? In which of the six broad categories of NP-complete problems does this problem belong?

# 4. (25 Points) The Traveling Salesman Problem.

Let's consider the optimization version of the traveling salesman problem (TSP): We are given a graph G(V, E) with n = |V| nodes. Each edge of the graph has an associated distance  $d_{ij}$ . The goal is to find a tour, i.e., a cycle that passes through every node exactly once, with minimum total distance.

- (a) (6 Points) State the decision version of the TSP problem. Show that the decision version of TSP is in NP.
- (b) (6 Points) How would you solve the TSP problem using a brute force approach? What would be the running time of that approach?
- (c) (13 Points) Develop a dynamic programming approach that solves this problem and report the running time of your algorithm.