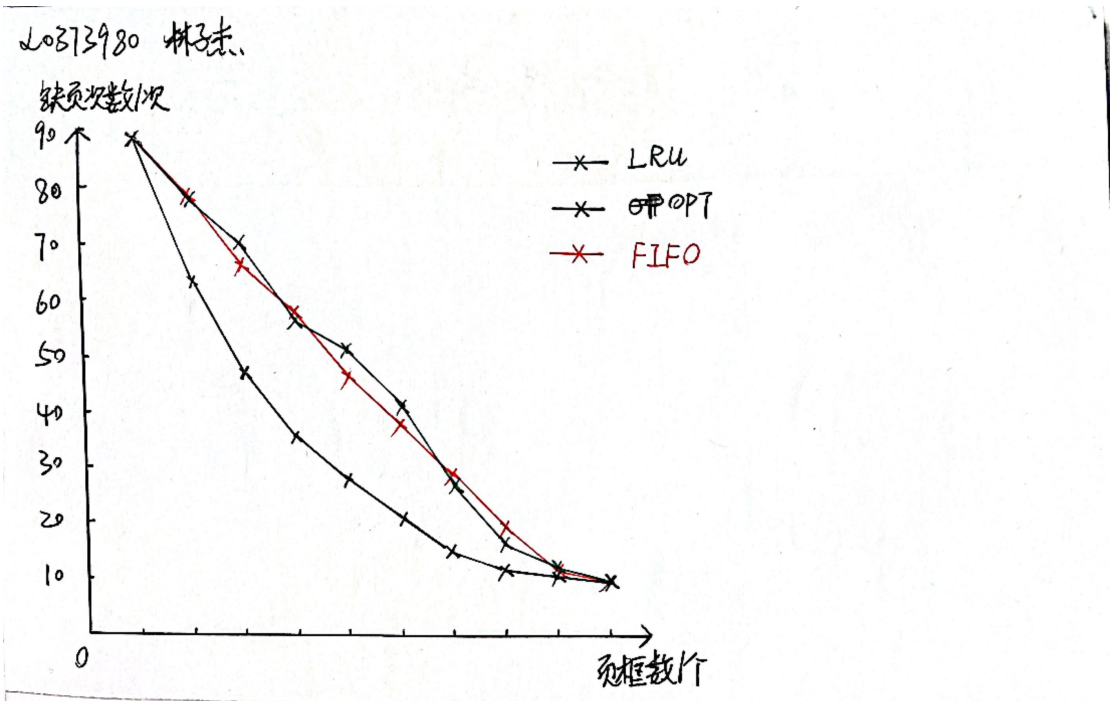
```
#include <stdio.h>

int visit[100] = {0, 9, 8, 4, 4, 3, 6, 5, 1, 5, 0, 2, 1, 1, 1, 1, 8, 8, 5,
                  3, 9, 8, 9, 9, 6, 1, 8, 4, 6, 4, 3, 7, 1, 3, 2, 9, 8, 6, 2, 9, 2, 7, 2, 7, 8, 4, 2, 3, 0, 1,
9, 4,
                  7, 1, 5, 9, 1, 7, 3, 4, 3, 7, 1, 0, 3, 5, 9, 9, 4, 9, 6, 1, 7, 5, 9, 4, 9, 7, 3, 6, 7, 7, 4,
5, 3, 5,
                  3, 1, 5, 6, 1,
                  1, 9, 6, 6, 4, 0, 9, 4, 3};

int FIFO(int frames) {
    int miss_times = 0, flag, i, j, max_cnt, max_pos;
    int tlb[2][100];
    for (i = 0; i < 100; i++) {
        tlb[0][i] = -1;
        tlb[1][i] = -1;
    }
    for (i = 0; i < 100; i++) {
        for (j = 0, flag = 0; j < frames; j++) {
            if (visit[i] == tlb[0][j]) {
                flag = 1;
                break;
            }
        }
        if (flag == 0) {
            for (j = 0, max_cnt = -1, max_pos = 0; j < frames; j++) {
                if (tlb[1][j] == -1) {
                    max_pos = j;
                    break;
                }
                if (tlb[1][j] > max_cnt) {
                    max_cnt = tlb[1][j];
                    max_pos = j;
                }
            }
            tlb[0][max_pos] = visit[i];
            tlb[1][max_pos] = 0;
            miss_times++;
        }
    }
    return miss_times;
}

int LRU(int frames) {
```

```
    int miss_times = 0, flag, i, j, max_cnt, max_pos;
    int tlb[2][100];
    for (i = 0; i < 100; i++) {
        tlb[0][i] = -1;
        tlb[1][i] = -1;
    }
    for (i = 0; i < 100; i++) {
        for (j = 0, flag = 0; j < frames; j++) {
            if (visit[i] == tlb[0][j]) {
                flag = 1;
                tlb[1][j] = 0;
                break;
            }
        }
        if (flag == 0) {
            for (j = 0, max_cnt = -1, max_pos = 0; j < frames; j++) {
                if (tlb[1][j] == -1) {
                    max_pos = j;
                    break;
                }
                if (tlb[1][j] > max_cnt) {
                    max_cnt = tlb[1][j];
                    max_pos = j;
                }
            }
            tlb[0][max_pos] = visit[i];
            tlb[1][max_pos] = 0;
            miss_times++;
        }
        for (j = 0; j < frames; j++) {
            if (tlb[1][j] != -1) {
                tlb[1][j]++;
            }
        }
    }
    return miss_times;
}

int main() {
    for (int i = 1; i <= 10; i++) {
        printf("LRU : Miss %d times when frame is %d\n", LRU(i), i);
    }
    for (int i = 1; i <= 10; i++) {
        printf("FIFO : Miss %d times when frame is %d\n", FIFO(i), i);
```

```
    }
    return 0;
}
```

OPT 算法不知为何我未能实现，故而借鉴了他人的算法，就不当做作业提交了。



2.解:
从0x80000000开始映射4MB的页表空间
则该地址对应的页框号为 $0x80000000/2^{12}=0x200$，该页的起始地址为 $0x80000000+4k\cdot 0x200=0x80200000$

3.解:
(1) 整个空间有 $2^{32}B=4GB$ 字节，一页有 $2^{12}B=4KB$ 字节

(2)
A: 0x0，访问一级页表，得到 0x0，页面无效，应产生缺页中断

B: 0x00803004，访问一级页表中的 0x002，得到 0x5001，且页面有效，访问 0x5000 对应的二级页表，由虚拟地址，访问其 0x003，得到 0x20001，有效，访问 0x20000 的页表项，得到物理地址 0x20000，页内偏移为4，取出 0x326001

C: 类似，0x00402001，访问 0x001，得到 0x1001，有效，且预留的，由虚拟地址，查询页目录的 0x2，得到 0x5001，有效，组合 0x5000 与 0x1，得数据 0x000(一个Byte)

(3) 应访问虚拟地址 0x00C01028

一级: 0000000011 ⇒ 取出 0x20000

二级: 0000000001 ⇒ 取出 0x326000

页偏移 000000101000 ⇒ 取 0x326028