



Benodigdhede vir hierdie vraestel/Requirements for this paper:

Multikeusekaarte/  
Multi-choice cards:

☐

Nie-programmeerbare sakrekenaar/  
Non-programmable calculator:

☒

Grafiekpapier/  
Graph paper:

☐

Draagbare Rekenaar/  
Laptop:

☐

Oopboek-eksamen/  
Open book examination?

NEE/  
NO

EKSAMEN/TOETS  
EXAMINATION/TEST:

Semestertoets /  
Semester Test

KWALIFIKASIE/  
QUALIFICATION:

Honns. B.Sc., M.Sc.

MODULEKODE/  
MODULE CODE:

ITRI626

TYDSDUUR/  
DURATION: 90  
minute/minutes

MODULEBESKRYWING/  
MODULE DESCRIPTION:

Kunsmatige Intelligensie / Artificial  
Intelligence

MAKS/  
MAX: 62

EKSAMINATOR(E)/  
EXAMINER(S):

DR. J. V. DU TOIT

DATUM/  
DATE: 06/09/2016

TYD/TIME: 07:30

MODERATOR:

DR. J. V. DU TOIT

Vraag 1 (Logiese Agente) / Question 1 (Logical Agents)

1.1 Gee drie definisies vir logiese gevolgtrekking ( $\alpha \models \beta$ ).

Give three definitions of logical entailment ( $\alpha \models \beta$ ).

[3 x 3 = 9]

Answer:

- For any sentences  $\alpha$  and  $\beta$ ,  $\alpha \models \beta$  if and only if  $M(\alpha) \subseteq M(\beta)$  [3 marks].
- For any sentences  $\alpha$  and  $\beta$ ,  $\alpha \models \beta$  if and only if the sentence  $(\alpha \Rightarrow \beta)$  is valid [3 marks].
- For any sentences  $\alpha$  and  $\beta$ ,  $\alpha \models \beta$  if and only if  $(\alpha \wedge \neg\beta)$  is unsatisfiable [3 marks].

1.2 Teken die volgende waarheidstabel oor en voltooi dit.

Draw the following truth table and complete it.

[5]

P	Q	$\neg P$	$P \wedge Q$	$P \vee Q$	$P \Rightarrow Q$	$P \Leftrightarrow Q$
T	T	?	?	?	?	?
T	F	?	?	?	?	?
F	T	?	?	?	?	?
F	F	?	?	?	?	?

Answer:

P	Q	$\neg P$	$P \wedge Q$	$P \vee Q$	$P \Rightarrow Q$	$P \Leftrightarrow Q$
T	T	F	T	T	T	T
T	F	F	F	T	F	F
F	T	T	F	T	T	F
F	F	T	F	F	T	T

Each column counts 1 mark.

1.3 Voltooi die volgende tabel met logiese ekwivalensies.

Complete the following table with logical equivalences.

[5 x 2 = 10]

$(\alpha \Rightarrow \beta)$	$\equiv$	1.3 (a)
$(\alpha \Rightarrow \beta)$	$\equiv$	1.3 (b)
$(\alpha \Leftrightarrow \beta)$	$\equiv$	1.3 (c)
1.3 (d)	$\equiv$	$(\neg\alpha \wedge \neg\beta)$
1.3 (e)	$\equiv$	$((\alpha \wedge \beta) \vee (\alpha \wedge \gamma))$

Answer:

$(\alpha \Rightarrow \beta)$	$\equiv$	$(\neg\beta \Rightarrow \neg\alpha)$ [1.3 (a)]	2 marks
$(\alpha \Rightarrow \beta)$	$\equiv$	$(\neg\alpha \vee \beta)$ [1.3 (b)]	2 marks
$(\alpha \Leftrightarrow \beta)$	$\equiv$	$((\alpha \Rightarrow \beta) \wedge (\beta \Rightarrow \alpha))$ [1.3 (c)]	2 marks
$\neg(\alpha \vee \beta)$ [1.3 (d)]	$\equiv$	$(\neg\alpha \wedge \neg\beta)$	2 marks
$(\alpha \wedge (\beta \vee \gamma))$ [1.3 (e)]	$\equiv$	$((\alpha \wedge \beta) \vee (\alpha \wedge \gamma))$	2 marks

- 1.4 Bewys dat  $P \Rightarrow P \vee Q$  altyd waar is deur van 'n bewys en ekwivalensies gebruik te maak. Wys jou redenasies en stappe volledig.

Show that  $P \Rightarrow P \vee Q$  is always true by using a proof and equivalences. Show all your reasoning and steps clearly. [10]

Answer:

$$P \Rightarrow P \vee Q$$

$$\equiv \neg P \vee (P \vee Q) \quad \text{Implication elimination}$$

$$\equiv (\neg P \vee P) \vee Q \quad \text{Association}$$

$$\equiv (P \vee \neg P) \vee Q \quad \text{Commutative}$$

$$\equiv (\text{True} \vee Q) \quad \text{Negation}$$

$$\equiv (Q \vee \text{True}) \quad \text{Commutative}$$

$$\equiv \text{True}$$

The left side counts 5 marks and the right side counts 5 marks. The exact steps and number of them can differ.

- 1.5 Toon aan dat  $((P \wedge Q) \vee R) \wedge (R \Rightarrow S) \models (P \vee S)$  deur van die Resolusie algoritme gebruik te maak. Wys jou redenasies en stappe volledig.

Show that  $((P \wedge Q) \vee R) \wedge (R \Rightarrow S) \models (P \vee S)$  by using the Resolution algorithm. Show all your reasoning and steps clearly. [28]

$$\text{Let } KB = ((P \wedge Q) \vee R) \wedge (R \Rightarrow S) \text{ and } \alpha = (P \vee S)$$

$\therefore$  Proof that  $KB \wedge \neg\alpha$  is unsatisfiable.

Convert  $KB \wedge \neg\alpha$  to conjunctive normal form (CNF):

$$((P \wedge Q) \vee R) \wedge (R \Rightarrow S) \wedge \neg(P \vee S)$$

$$\equiv ((P \wedge Q) \vee R) \wedge (R \Rightarrow S) \wedge (\neg P \wedge \neg S) \quad \text{De Morgan}$$

$$\equiv ((P \wedge Q) \vee R) \wedge (\neg R \vee S) \wedge (\neg P) \wedge (\neg S) \quad \text{Implication elimination}$$

$$\equiv (R \vee (P \wedge Q)) \wedge (\neg R \vee S) \wedge (\neg P) \wedge (\neg S) \quad \text{Association}$$

$$\equiv (R \vee P) \wedge (R \vee Q) \wedge (\neg R \vee S) \wedge (\neg P) \wedge (\neg S) \quad \text{Association}$$

2 mark for each correct clause produced = 10 marks. 5 marks for the conversion and 5 marks for the explanations on the right = 10 marks.

Number the clauses:

- 1)  $(R \vee P)$
- 2)  $(R \vee Q)$
- 3)  $(\neg R \vee S)$
- 4)  $(\neg P)$
- 5)  $(\neg S)$

1 mark.

Perform resolution:

Resolution between (5) and (3) gives:

- 6)  $(\neg R)$

Resolution between (6) and (1) gives:

- 7)  $(P)$

Resolution between (7) and (4) gives:

- 8)  $\square$

6 marks for the application of the resolution rule.

$\therefore$  Since the empty clause was found,  $KB \models \alpha$

$\therefore ((P \wedge Q) \vee R) \wedge (R \Rightarrow S) \models (P \vee S)$

1 mark for the final answer.

TOTAAL/TOTAL: 62