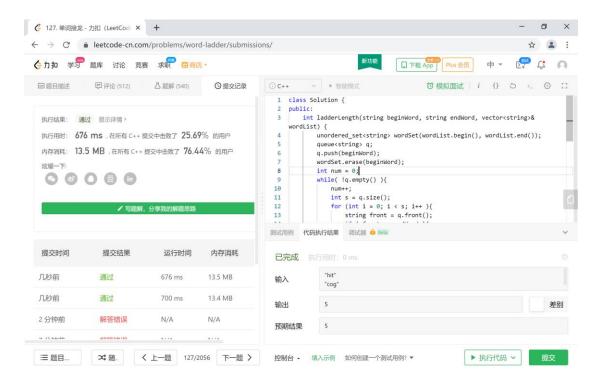
1. 算法思路:深度优先搜索,对于陆地上的点,搜索其相邻的陆地节点,并置 0 复杂度分析: *Q(MM)*,其中 M*M*和 N*N*分别为行数和列数

```
class Solution {
    void dfs(vector<vector<char>> & grid, int r, int c ){
         grid[r][c] = '0';
         if (r - 1 >= 0 \&\& grid[r - 1][c] == '1'){
              dfs(grid, r - 1, c);
         }
         if (r + 1 < grid.size() && grid[r + 1][c] == '1'){</pre>
              dfs(grid, r + 1, c);
         }
         if (c - 1 >= 0 && grid[r][c - 1] == '1'){
              dfs(grid, r, c - 1);
         }
         if (c + 1 < grid[0].size() && grid[r][c + 1] == '1'){</pre>
              dfs(grid, r, c + 1);
         }
     }
public:
     int numIslands(vector<vector<char>>& grid) {
         int num = 0;
         for (int r = 0; r < grid.size(); ++r) {</pre>
              for (int c = 0; c < grid[0].size(); ++c) {</pre>
                   if (grid[r][c] == '1') {
                        ++num;
                        dfs(grid, r, c);
                   }
              }
         }
         return num;
    }
};
← → C • leetcode-cn.com/problems/number-of-islands/submissions/
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                                                 □下载App Plus会员 中 ▼ E+ C+ C
 ◆ 力力 学习 题库 讨论 竞赛 求职 □商店。
回 题目描述 □ 评论 (1.2k) △ 题解 (1.6k) ○ 提交记录
                                                     び 模拟面试 | i (} 5 > □ ① []
                                         执行结果: 通过 显示详情 >
  执行用时: 20 ms , 在所有 C++ 提交中击败了 49.94% 的用户
  内存消耗: 9.2 MB , 在所有 C++ 提交中击败了 85.79% 的用户
                                         (r + 1 < grid.size() && grid[r + 1][c] == '1'){
dfs(grid, r + 1, c);
  00000
                                       if (c - 1 >= 0 && grid[r][c - 1] == '1'){
    dfs(grid, r, c - 1);
                                        几秒前
                        9.2 MB
                                                           ▶ 执行代码 ∨ 提交
 ※ 题目… ★ 随. 〈 上一题 200/2056 下一题 〉 控制台 → 贡献 i
```

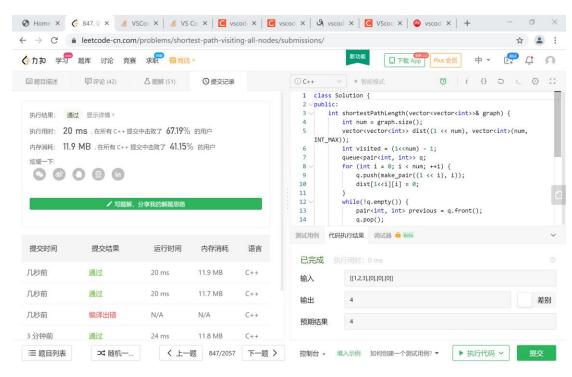
2. 算法思路: 广度优先搜索, 使用队列存储 复杂度分析:  $O(N^2 \times C)$ , 其中 N 为字符串个数, C 为字符串长度

```
class Solution {
public:
    int ladderLength(string beginWord, string endWord, vector<string>& wordList
) {
        unordered_set<string> wordSet(wordList.begin(), wordList.end());
        queue<string> q;
        q.push(beginWord);
        wordSet.erase(beginWord);
        int num = 0;
        while( !q.empty() ){
            num++;
            int s = q.size();
            for (int i = 0; i < s; i++ ){
                string front = q.front();
                if ( front == endWord ){
                    return num;
                }
                q.pop();
                for ( auto it = wordSet.begin(); it != wordSet.end(); ){
                    string word = *it;
                    int d = 0;
                    for (int j = 0; j < word.length(); j++ ){</pre>
                        if ( word[j] != front[j] ){
                             d++;
                        }
                        if (d > 1){
                             break;
                         }
                    }
                    it++;
                    if ( d <= 1 ){</pre>
                        q.push(word);
                        wordSet.erase(word);
                    }
                }
            }
        }
        return 0;
   }
};
```



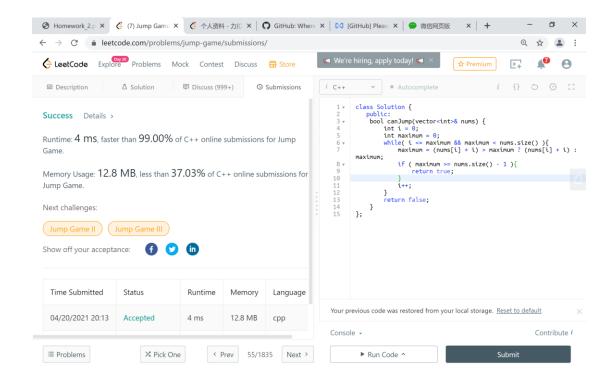
3. 算法思路: 广度优先搜索, 使用 bitmask 复杂度分析:  $O(N \times 2^n)$ , 其中 N 为节点数量。

```
class Solution {
public:
    int shortestPathLength(vector<vector<int>>& graph) {
        int num = graph.size();
        vector<vector<int>> dist((1 << num), vector<int>(num, INT_MAX));
        int visited = (1<<num) - 1;</pre>
        queue<pair<int, int>> q;
        for (int i = 0; i < num; ++i) {</pre>
            q.push(make_pair((1 << i), i));</pre>
            dist[1<<i][i] = 0;
        }
        while(!q.empty()) {
            pair<int, int> previous = q.front();
            q.pop();
            int d = dist[previous.first][previous.second];
            if (previous.first == visited){
                 return d;
            for (auto & vertex : graph[previous.second]) {
                 int v_new = previous.first |(1 << vertex);</pre>
                 if (dist[v_new][vertex] > d + 1) {
                     dist[v_new][vertex] = d + 1;
                     q.push(make_pair(v_new, vertex));
```



4. 算法思路: 贪心算法,每一步记录最大能到达的位置 复杂度分析: O(N)

```
class Solution {
  public:
  bool canJump(vector<int>& nums) {
    int i = 0;
    int maximum = 0;
    while (i <= maximum && maximum < nums.size()) {
        maximum = (nums[i] + i) > maximum ? (nums[i] + i) : maximum;
        if (maximum >= nums.size() - 1) {
            return true;
        }
        i++;
    }
    return false;
}
```



5. 算法思路: 贪心算法, 在每一步记录剩余油量的值

```
复杂度: O(N)
class Solution {
   public:
    int canCompleteCircuit(vector<int>& gas, vector<int>& cost) {
        int start = 0;
        int current = 0;
        int tank = 0;
        bool flag = false;
        while (start - current != 1 && current - start != gas.size() - 1) {
            tank = tank + gas[current] - cost[current]; //gas in tank when arr
iving (i+1)th station
             current = (current + 1) % gas.size();
            if (current == 0) {
                 flag = true;
             }
             if (tank < 0) {</pre>
                 start = current;
                 tank = 0;
                 if (start == 0 || (flag && start >= current)) {
                     return -1;
                 }
             }
        if (tank + gas[current] - cost[current] < 0) {</pre>
```

```
return -1;
}
return start;
}
```

