```
1. 算法思路: 递归, 回溯法, 枚举
    复杂度分析: \frac{1}{n+1}\binom{2n}{n}, n 为括号的对数
    class Solution {
         public:
               vector<string> generateParenthesis(int n) {
                    string s;
                    pair<int, int> left(n, n);
                    return search(s, left);
              }
              vector<string> search(string s, pair<int, int> & left){
                    vector<string> v, v1, v2;
                    if ( left.first == 0 && left.second == 0 ){
                         v.push_back(s);
                         return v;
                    if ( left.first > 0 ){
                         left.first--;
                         v1 = search(s+'(', left);
                         left.first++;
                    }
                    if ( left.second > left.first ){
                         left.second--;
                         v2 = search(s+')', left);
                         left.second++;
                    set_union(v1.begin(), v1.end(), v2.begin(), v2.end(),
                              std::back_inserter(v));
                    return v;
              }
    };
    ③ hw x | ■ SFS x | ③ h16 x | ③ 扫描 x | ② 用描 x | ③ Alg x | ⑤ (5) + x | ⑥ (5) + x | ⑥ (5) + x | Ⅰ 中止 x | + ● □ □ X
     ← → C · leetcode.com/problems/generate-parentheses/submissions/
     	riangle Description 	riangle Solution 	riangle Discuss (999+) 	riangle Submissions 	riangle C++ 	riangle 	riangle Autocomplete 	riangle (	riangle) 	riangle 	riangle 	riangle 	riangle 	riangle Autocomplete 	riangle (	riangle) 	riangle
                                                Runtime: 12\, ms, faster than 15.05\% of C++ online submissions for Generate
      Memory Usage: 26.2\ MB, less than 6.72\% of C++ online submissions for
                                                   vector<string> search(string s, pair<int, int> &
                                               Testcase Run Code Result Debugger 🔓
      Letter Combinations of a Phone Number Valid Parentheses
      Show off your acceptance: f y in
                                               Time Submitted Status
                        Runtime Memory Language
                                               Expected ["((()))","(()()","(()()","()(())","()(()"]
      Use Example Testcases
```

 X Pick One
 < Prev</td>
 22/1901
 Next >
 ► Run Code ^
 Submit

```
2. 算法思路: 递归, 回溯法, 枚举
   复杂度分析: 0(4<sup>n</sup>), N是 digits 字符串的长度
   class Solution {
       public:
           vector<string> letterCombinations(string digits) {
               if (digits == ""){
                   vector<string> v;
                   return v;
               }
               string s;
               return search(s, digits, 0);
           }
           vector<string> search(string s, string & digits, int i){
               if ( i>=digits.length() ){
                   vector<string> v;
                   v.push_back(s);
                   return v;
               }
               vector<string> v, w, x, y, z;
               char c = 'a' + (digits[i]-'2')*3;
               switch( digits[i] ){
                   case '2':
                   case '3':
                   case '4':
                   case '5':
                   case '6':
                       x = search(s+c, digits, i+1);
                       y = search(s+(char)(c+1), digits, i+1);
                       z = search(s+(char)(c+2), digits, i+1);
                       break;
                   case '7':
                       x = search(s+'p', digits, i+1);
                       y = search(s+'q', digits, i+1);
                       z = search(s+'r', digits, i+1);
                       w = search(s+'s', digits, i+1);
                       break;
                   case '8':
                       x = search(s+'t', digits, i+1);
                       y = search(s+'u', digits, i+1);
                       z = search(s+'v', digits, i+1);
                       break;
                   default:
                       x = search(s+'w', digits, i+1);
```

```
y = search(s+'x', digits, i+1);
                          z = search(s+'y', digits, i+1);
                          w = search(s+'z', digits, i+1);
                }
                set_union(x.begin(), x.end(), y.begin(), y.end(),
                          std::back_inserter(v));
                set_union(z.begin(), z.end(), w.begin(), w.end(),
                          std::back_inserter(v));
                return v;
          }
};
 \leftarrow \rightarrow C \blacksquare leetcode.com/problems/letter-combinations-of-a-phone-number/submissions/
 class Solution {
public:
 Success Details >
                                                                 ltc:
vector<string> letterCombinations(string digits) {
   if (digits == ""){
      vector<string> v;
      return v;
}
  Runtime: 4 ms, faster than 16.88% of C++ online submissions for Letter
                                                                   string s;
return search(s, digits, 0);
 Memory Usage: 12.5 MB, less than 5.52% of C++ online submissions for
                                                         Testcase Run Code Result Debugger 🔒
 Letter Combinations of a Phone Number.
 Next challenges:
                                                         Accepted
                                                                  Runtime: 0 ms
  Combination Sum Binary Watch
                                                         Your input
                                                                   "23"
 Show off your acceptance:
                                                                    ["ad","ae","af","bd","be","bf","cd","ce","c
                                                         Output
                                                                                                     Diff
                                                                    ["ad","ae","af","bd","be","bf","cd","ce","cf"]
                                                         Expected
   Time Submitted
                 Status
                                      Memory
                                                         Console -
                                                               ► Run Code ^
```

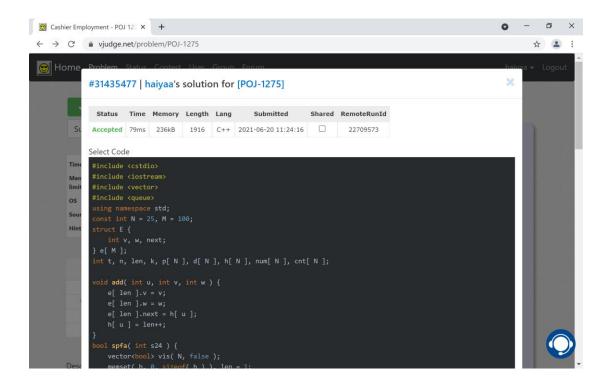
3. 算法思路:用 SPFA 解差分约束

```
复杂度分析: O(|V|*|E|), V 为构建的图中的节点数, E 为构建的图中的边数
#include <cstdio>
#include <iostream>
#include <vector>
#include <queue>
using namespace std;
const int N = 25, M = 100;
struct E {
   int v, w, next;
} e[ M ];
int t, n, len, k, p[ N ], d[ N ], h[ N ], num[ N ], cnt[ N ];

void add( int u, int v, int w ) {
   e[ len ].v = v;
```

```
e[len].w = w;
    e[ len ].next = h[ u ];
    h[ u ] = len++;
}
bool spfa( int s24 ) {
    vector<bool> vis( N, false );
    memset( h, 0, sizeof( h ) ), len = 1;
    memset( d, -0x3f, sizeof( d ) );
    memset( cnt, 0, sizeof( cnt ) );
    d[0] = 0;
    for ( int i = 1; i <= 24; i++ ) {
        add( i - 1, i, 0 ), add( i, i - 1, -num[ i ] );
       if ( i >= 8 ) {
            add( i - 8, i, p[ i ] );
        }
        else {
            add( i + 16, i, -s24 + p[ i ] );
        }
    }
    add( 0, 24, s24 ), add( 24, 0, -s24 );
    queue<int> q;
    q.push( 0 );
    while ( !q.empty() ) {
       int u = q.front();
        q.pop();
        vis[ u ] = false;
        for ( int j = h[ u ]; j; j = e[ j ].next ) {
            int v = e[ j ].v;
            int w = d[u] + e[j].w;
           if (w > d[v]) {
               d[v] = w;
                cnt[ v ] = cnt[ u ] + 1;
                if ( cnt[ v ] >= 25 ) return true;
                if ( !vis[v] ) q.push(v), vis[v] = true;
            }
        }
    }
    return false;
}
int main() {
    cin >> t;
    while ( t-- ) {
        memset( num, 0, sizeof( num ) );
        for ( int i = 1; i <= 24; i++ ){
```

```
cin >> p[ i ];
        }
        cin >> n;
        for ( int i = 1; i <= n; i++ ) {
            cin >> k;
            num[ ++k ]++;
        bool ok = false;
        for ( int i = 1; i <= n; i++ ) {
            if (!spfa(i)) {
                 cout << d[ 24 ] << endl;</pre>
                 ok = true;
                 break;
            }
        }
        if (!ok){
            cout << "No Solution\n";</pre>
        }
    }
}
```



4. P问题: 所有能在多项式时间内解决的搜索问题叫做 P问题。

NP 问题: 所有能在多项式时间内验证一个解是否正确的搜索问题叫做 NP 问题。

NPC 问题: 一个搜索问题是搜索问题, 当且仅当所有的搜索问题都能归约到该问题。

证明一个问题 Q 是 NP 难问题的方法: 将另一个 NP 难问题归约到 Q