# *Ada*<sup>-</sup>: A Simple Ada Programming Language

# Programming Assignment 2

Syntactic and Semantic Definitions

Due Date: 1:20PM, Tuesday, May 25, 2021

Your assignment is to write an LALR(1) parser for the  $Ada^-$  language. You will have to write the grammar and create a parser using **yacc**. Furthermore, you will do some simple checking of semantic correctness. Code generation will be performed in the third phase of the project.

## 1 Assignment

You first need to write your symbol table, which should be able to perform the following tasks:

- Push a symbol table when entering a scope and pop it when exiting the scope.
- Insert entries for variables, constants, and procedure declarations.
- Lookup entries in the symbol table.

You then must create an LALR(1) grammar using **yacc**. You need to write the grammar following the syntactic and semantic definitions in the following sections. Once the LALR(1) grammar is defined, you can then execute **yacc** to produce a C program called "**y.tab.c**", which contains the parsing function **yyparse**(). You must supply a main function to invoke **yyparse**(). The parsing function **yyparse**() calls **yylex**(). You will have to revise your scanner function **yylex**().

#### 1.1 What to Submit

You should submit the following items:

- revised version of your lex scanner
- a file describing what changes you have to make to your scanner
- your yacc parser

Note: comments must be added to describe statements in your program

- Makefile
- · test programs

## 1.2 Implementation Notes

Since **yyparse()** wants tokens to be returned back to it from the scanner. You should modify the definitions of **token**, **tokenInteger**, **tokenString**. For example, the definition of **token** should be revised to:

```
#define token(t) {LIST; printf("<\%s>\n","t"); return(t);}
```

## 2 Syntactic Definitions

#### 2.1 Constant and Variable Declarations

There are two types of constants and variables in a program:

- global constants and variables declared inside the program
- local constants and variables declared inside procedures and blocks

#### **Data Types and Declarations**

The predefined data types are integer, string, boolean, and float.

#### 2.1.1 Constants

A constant declaration has the form:

```
identifier : constant <: type > := constant\_exp;
```

where the item in the < > pair is optional, and the type of the declared constant must be inferred based on the constant expression on the right-hand side. Note that constants cannot be reassigned or this code would cause an error. For example,

```
s: constant := "Hey There";
i: constant := -25;
f: constant : float := 3.14;
b: constant := true;
```

#### 2.1.2 Variables

A variable declaration has the form:

```
identifier <: type >< := constant\_exp >;
```

where *type* is one of the predefined data types. When both the type attribute declaration, i.e : *type* and initialization are omitted from variable declarations, the default data type is **int**. For example,

```
s: string;
i := 10;
d: float;
b: boolean = false;
```

#### **Arrays**

Arrays declaration has the form:

```
identifier : type [ num ];
```

For example,

### 2.2 Program Units

The two program units are the program and procedures.

## 2.2.1 Program

A program has the form:

where the item in the <> pair is optional.

#### 2.2.2 Procedures

Procedure declaration has the following form:

```
procedure identifier < ( formal arguments ) > < return type > block end identifier;
```

where *block* is a block statement (see Section 2.3.2), **return** *type* is optional, and *type* can be one of the predefined types. The formal arguments are declared in the following form:

```
identifier: type <; identifier: type; ...; identifier: type>
```

Parentheses are not required when no arguments are declared. No procedures may be declared inside a procedure. For example,

```
program Example
  -- constant and variable declaration
  declare
   a: integer := 5;
    c: integer;
  -- function declaration
  procedure add (a: integer; b: integer) return integer
   begin
      return a+b;
    end;
  end add;
  -- main block
  begin
    c := add(a, 10);
    println c;
  end
end Example
```

Note that procedures with no retuen type can not be used in expressions, and procedures with retuen values are called functions and can be used in expressions.

#### 2.3 Statements

There are several distinct types of statements in  $Ada^-$ .

### **2.3.1** simple

```
The simple statement has the form:
```

```
identifier := expression;

or
    identifier[integer_expression] := expression;

or
    print <(> expression <)>;

or
    println <(> expression <)>;

or
    read identifier;

or
    return;
```

## expressions

Arithmetic expressions are written in infix notation, using the following operators with the precedence:

```
(1) – (unary)
```

- (2) \* /
- (3) + -
- (4) < <= = => > /=

return expression;

- (5) not
- (6) and
- (7) or

Associativity is the left. Valid components of an expression include literal constants, variable names, function invocations, and array reference of the form

```
A [ integer_expression ]
```

## function invocation

A function invocation has the following form:

```
identifier < ( comma-separated expressions ) >
```

#### 2.3.2 block

A block is a collection of statements enclosed by **begin** and **end** with an optional **declare** section. The simple statement has the form:

```
< declare
zero or more variable and constant declarations>
begin
<one or more statements>
end;
```

### 2.3.3 conditional

The conditional statement may appear in two forms:

```
if boolean_expr then
a block or simple statement
else
a block or simple statement
end if;
or

if boolean_expr then
a block or simple statement
end if;
```

## 2.3.4 loop

or

The loop statement has two forms:

```
while boolean_expr loop
a block or simple statement
end loop;

for ( identifier in num . . num ) loop
a block or simple statement
end loop;
```

## 2.3.5 procedure invocation

A procedure has no return value. It has the following form:

```
identifier < ( semicomma-separated expressions ) > ;
```

## **3 Semantic Definition**

The semantics of the constructs are the same as the corresponding Pascal and C constructs, with the following exceptions and notes:

- The parameter passing mechanism for procedures in call-by-value.
- Scope rules are similar to C.
- The identifier after the **end** of program or procedure declaration must be the same identifiers as the name given at the beginning of the declaration.
- Types of the left-hand-side identifier and the right-hand-side expression of every assignment must be matched.
- The types of formal parameters must match the types of the actual parameters.

# 4 yacc Template (yacctemplate.y)

```
int Opt_P = 1;
응 }
/* tokens */
%token SEMICOLON
응응
program:
             identifier semi
             Trace("Reducing to program\n");
semi:
             SEMICOLON
             Trace("Reducing to semi\n");
응응
#include "lex.yy.c"
yyerror(msg)
char *msg;
  fprintf(stderr, "%s\n", msg);
}
main()
  yyparse();
```