

Efficient Coinductives through State-Machine Corecursors

Optimising the terminal F-coalgebra in Set
for computational behaviour in Lean4

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Polynomials (background)



- A head H ,
- child families c_h in $h : H$.
- Morally, a collection of constructors H , and data at each constructor c_h .
- For now we will not consider “recursive” (inductive) data structures, rather I will have this be a variable we take the fix point in.
- Examples:
 - ▶ natural numbers $\text{NatF}(X) := 1 + X$,
 - ▶ lists $\text{ListF}(A, X) := 1 + A * X$,
 - ▶ streams $\text{StreamF}(A, X) := A * X$.

$$(h : H) \times c_h \rightarrow \alpha \equiv \sum_{h:H} \alpha^{c_h}$$

See (nLab authors, 2026) and (Neel Krishnaswami, 2026)

Coinductives (background)



- Unbounded tree of data.
- Greatest cofixpoint of a polynomial.
- Morally, the process of putting the word “lazy” in front of a functor,
 - on lists it sends them to lazy lists.
- Have a corecursion principle rather than an recursion principle,
 - instead of `rec : (f a → a) → fix f → a`,
 - we have `corec : (a → f a) → a → cofix f`.
- On our examples and their correspondencens,
 - `cofix NatF`, natural numbers union infinity,
 - `cofix ListF(A)`, lists of `As` with arbitrary length union streams of `A`,
 - `cofix StreamF(A)`, infinite streams of `As`.

See (Avigad et al., 2019)

Performance of coinductives (core)



- Previous work has demonstrated implementation of coinductives.
- Computational behaviour was not a concern.

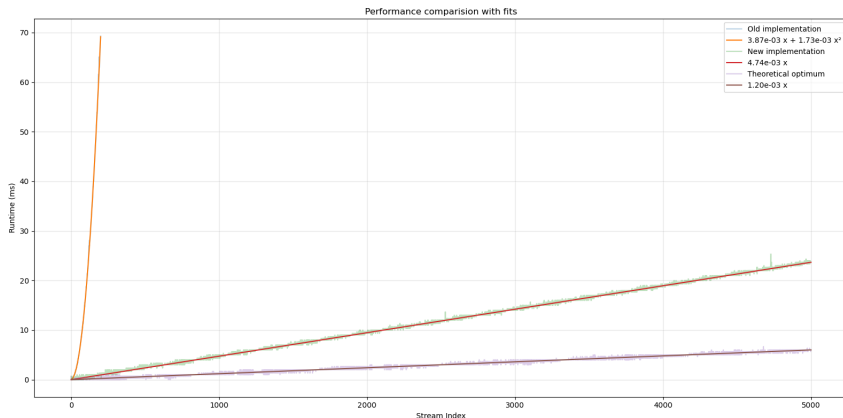


Figure 1: Performance of the old, new, and a perfect implementation

ABI Type (extention)



- Given $\alpha : \text{Type}_{\mathcal{U}}$, $\beta : \text{Type}_{\mathcal{V}}$ and $\text{eq} : \alpha \simeq \beta$.
- Univalence gives uniqueness.
- Adds a form of “weak-univalence”.
- Allows for a zero cost universe transliteration.
- Symetric form of **Shrink**.

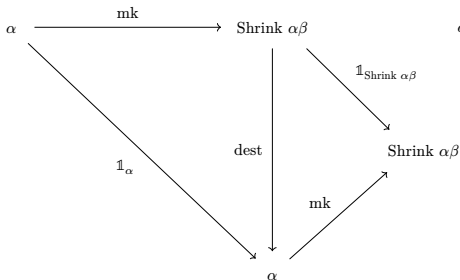


Figure 2: Operations on Shrink

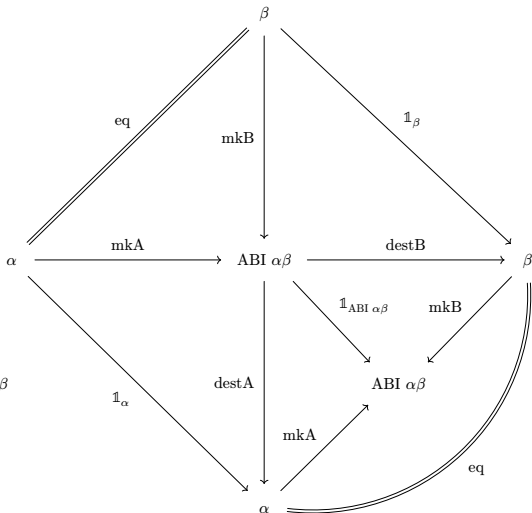


Figure 3: Operations on ABI

Interaction Trees (extention)



- Free monad with non-termination as an effect.

```
1 coinductive ITree
2   (E : Type u → Type u) (R : Type v) :=
3   | tau : ITree E R → ITree E R
4   | ret : R → ITree E R
5   | vis {A : Type u} (e : E A) (k : A → ITree E R)
```



Lean

- ITerm E forms a monad,
 - With a method trigger : $EV \rightarrow \text{ITree } E \ V$,
 - and iter : $(A \rightarrow \text{ITree } E(A + B)) \rightarrow A \rightarrow \text{ITree } E \ B$.
- And an equivalence relation $A \approx B$,
 - satisfying $A \approx B \leftrightarrow \text{tau } A \approx B$.

Project statistics



```
> nix-shell -p cloc --run "cloc . --exclude-dir=.lake"
  100 text files.
   92 unique files.
   14 files ignored.
```

github.com/AlDanial/cloc v 2.06 T=0.08 s (1110.8 files/s, 124656.0 lines/s)

Language	files	blank	comment	code
Lean	59	1229	252	6621
Typst	20	332	122	1099
SVG	1	0	0	171
JSON	3	0	0	125
TeX	1	4	0	76
YAML	3	8	7	69
Markdown	2	27	0	68
Python	2	22	18	57
Nix	1	2	0	15
SUM:	92	1624	399	8301

(excluding mathlib PRs)

Bibliography



- Avigad, J., Carneiro, M., & Hudon, S. (2019). Data Types as Quotients of Polynomial Functors. In J. Harrison, J. O’Leary, & A. Tolmach (Eds.), *10th International Conference on Interactive Theorem Proving (ITP 2019): Vol. 141. 10th International Conference on Interactive Theorem Proving (ITP 2019)*. <https://doi.org/10.4230/LIPIcs.ITP.2019.6>
- Chappe, N., He, P., Henrio, L., Zakowski, Y., & Zdancewic, S. (2023). Choice Trees: Representing Nondeterministic, Recursive, and Impure Programs in Coq. *Proceedings of the ACM on Programming Languages*, 7(POPL), 1770–1800. <https://doi.org/10.1145/3571254>
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- Zakowski, Y., Beck, C., Yoon, I., Zaichuk, I., Zaliva, V., & Zdancewic, S. (2021). Modular, compositional, and executable formal semantics for LLVM IR. *Proc. ACM Program. Lang.*, 5(ICFP). <https://doi.org/10.1145/3473572>

Further work

Futomorphic productivity (extention)



- A futomorphism is a morphism $(c \rightarrow f (\text{Free } f \ c)) \rightarrow c \rightarrow \text{cofix } f$,
- a terminating function in the futomorphism is exactly a productive one,
- futomorphisms have a unfold lemma which is very hard to prove.

See (Yang & Wu, 2022)

LLVM semantics (extention)



- Tobias Grosser's project VeIR wants formalized semantics,
- for this they want to use an ITree interpreter,
- semantics are detailed in (Zakowski et al., 2021).

CTree (extention)



- Another coinductive datastructure,
- has visible and silent non-determinism,

```
1 coinductive
2   (E B : Type u → Type u) (R : Type v) :=
3   | ret : R → ITree E R
4   | step : ITree E R → ITree E R
5   | guard : ITree E R → ITree E R
6   | vis {A : Type u} (e : E A) (k : A → ITree E R)
7   | br {A : Type u} (e : B A) (k : A → ITree E R)
8   | stuck
```



Lean

See (Chappe et al., 2023)