

ButterRoti ICPC Team Notebook (2017-18)

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1 Misc

1.1 Build

```

1 {
2 "cmd": ["g++ -std=c++14 -g -Wall '${file}' &&
1     timeout 15s '${file_path}/./a.out' <'${file_path
1     }/input.txt'>'${file_path}/output.txt'"],
2 "shell":true
4 }

```

1.2 Snippet

```

1 #include <bits/stdc++.h>
2
3 using namespace std;
4
5 #define SYNC std::ios::sync_with_stdio(false);
6     cin.tie(NULL);
7
6 template<typename T> using V = vector<T>;
7 template<typename T, typename V> using P = pair<T,
8     V>;
9
14 using LL = long long;
14 using ll = LL;
15 using LD = long double;
16 using ld = long double;
17
18 #define fi first
19 #define ff first
17 #define se second
18 #define ss second
19 #define endl '\n'
20 #define FOR0(i,n) for(int i=0, _##i=(n); i<_##i;
21     ++i)
21 #define FOR(i,l,r) for(int i=(l), _##i=(r); i<_##i
22     ; ++i)
23 #define FORD(i,l,r) for(int i=(r), _##i=(l); --i>=
24     _##i; )
25 #define dzx cerr << "here";
26 #define her cerr << "HERE\n"
27
26 const int MOD = (int)1e9 + 7, inf = 1e9;
27 const ll INF = 1e18;
28

```

```

29 int32_t main() {SYNC;
30
31     return 0;
32 }

```

1.3 Stack Size Increase

```

1  #include <sys/resource.h>
2
3  int main() {
4      rlimit R;
5      getrlimit(RLIMIT_STACK, &R);
6      R.rlim_cur = R.rlim_max;
7      setrlimit(RLIMIT_STACK, &R);
8  }

```

1.4 Variadic Multiplication and Addition

```

1
2  const int MOD = (int)1e9 + 7;
3
4  int add() { return 0; }
5
6  template<typename... T> int add(int a, T... arg) {
7      int b = add(arg...);
8      return (a + b >= MOD ? a + b - MOD : a + b);
9  }
10
11 int multiply() { return 1; }
12
13 template<typename... Args> int multiply(int a,
14     Args... arg) {
15     return (a * 1LL * multiply(arg...)) % MOD;
16 }

```

2 Combinatorial optimization

2.1 Lowest Common Ancestor

```

1 // 0-based vertex indexing. memset to -1
2 int log(int t) {
3     int res = 1;
4     for(; 1 << res <= t; res++);
5     return res;
6 }

```

```

7 int lca(int u, int v) {
8     if(h[u] < h[v]) swap(u, v);
9     int L = log(h[u]);
10    for(int i = L - 1; i >= 0; i--) {
11        if(par[u][i] + 1 && h[u] - (1 << i) >= h[v])
12            u = par[u][i];
13    }
14    if(v == u) return u;
15    for(int i = L - 1; i >= 0; i--) {
16        if(par[u][i] + 1 && par[u][i] != par[v][i]) {
17            u = par[u][i]; v = par[v][i];
18        }
19    }
20    return par[u][0];
21 }

```

2.2 Heavy-Light Decomposition

```

1 V<V<int>> > g;
2 int N;
3 V<int> cpar, id, depth, parent;
4 V<int> chain;
5
6 int dfs(int c, int p) {
7     parent[c] = p;
8     depth[c] = depth[p] + 1;
9     int sz = 1;
10    auto it = find(g[c].begin(), g[c].end(), p);
11    if(it != g[c].end())
12        g[c].erase(it);
13    if(g[c].empty())
14        return 1;
15    int mx = 0;
16    for(auto &i: g[c]) {
17        int cur = dfs(i, c);
18        sz += cur;
19        if(cur > mx)
20            mx = cur, swap(i, g[c][0]);
21    }
22    return sz;
23 }
24
25 void form_chains(int c, int cp) {
26     cpar[c] = cp;
27     id[c] = (int)chain.size();
28     chain.push_back(c);
29 }

```

```

29 for(int i=0;i<(int)g[c].size();i++){
30     if(i)
31         form_chains(g[c][i], g[c][i]);
32     else
33         form_chains(g[c][i], cp);
34 }
35 }
36
37 void update(int u, int v){
38     while(u!=v){
39         if(cpar[v] == cpar[u]){
40             if(depth[v] < depth[u])
41                 swap(v,u);
42             supdate(0, N - 1, 1, id[u]+1, id[v]);
43             v = u;
44         }
45         else{
46             if(depth[cpar[v]] < depth[cpar[u]])
47                 swap(v,u);
48             supdate(0, N - 1, 1, id[cpar[v]], id[v]);
49             v = parent[cpar[v]];
50         }
51     }
52 }
53
54 void preprocess(int r) {
55     depth.resize(N);
56     depth[r] = 0;
57     cpar.resize(N);
58     parent.resize(N);
59     chain.clear();
60     chain.reserve(N);
61     id.resize(N);
62     dfs(r, r);
63     form_chains(r, r);
64 }

```

2.3 Auxiliary Tree

```

1
2 //std::vector<int> a contains vertices to form the
  aux t
3 sort(ALL(a), [](const int & a, const int & b) ->
  bool{
4     return st[a] < st[b];
5 });

```

```

6 set<int> s(a);
7 for(int i = 0, k = (int)a.size(); i + 1 < k; i++){
8     int v = lca(a[i], a[i + 1]);
9     if(s.find(v) == s.end())
10         a.push_back(v);
11     s.insert(v);
12 }
13
14 sort(ALL(a), [](const int & a, const int & b) ->
  bool{
15     return st[a] < st[b];
16 });
17
18 stack<int> S;
19 S.push(a[0]);
20
21 auto anc = [](int & a, int & b) -> bool{
22     return st[b] >= st[a] && en[b] <= en[a];
23 };
24
25 for(int i = 1; i < (int)a.size(); i++){
26     while(!anc(S.top(), a[i])) S.pop();
27     G[S.top()].pp(a[i]);
28     G[a[i]].pp(S.top());
29     S.push(a[i]);
30 }
31 //G is the Aux tree

```

2.4 Articulation Point and Bridges

```

1 #include <bits/stdc++.h>
2
3 using namespace std;
4 const int N = 50;
5 int dis[N], low[N], par[N], AP[N], vis[N], tits;
6 void update(int u, int i, int child) {
7     //For Cut Vertices
8     if(par[u] != -1 && low[i] >= dis[u]) AP[u] =
  true;
9     if(par[u] == -1 && child > 1) AP[u] = true;
10
11     //For Finding Cut Bridge
12     if(low[i] > dis[u]){
13         //articulation bridge found.
14     }
15 }

```

```

16 void dfs(int u) {
17     vis[u] = true;
18     low[u] = dis[u] = (++tits); int child = 0;
19     for(int i : g[u]) {
20         if(!vis[i]) {
21             child++;
22             par[i] = u;
23             dfs(i);
24             low[u] = min(low[u], low[i]);
25             update(u, i, child);
26         }
27         else if(i != par[u]) {
28             low[u] = min(low[u], dis[i]);
29         }
30     }
31 }

```

2.5 Biconnected Components

```

1  #include <bits/stdc++.h>
2  using namespace std;
3  const int N = (int)2e5 + 10;
4
5  vector<vector<int>> tree, g;
6  bool isBridge[N << 2], vis[N];
7  int Time, arr[N], U[N], V[N], cmpno, comp[N];
8  vector<int> temp; //temp stores component values
9
10 int adj(int u, int e) {
11     return (u == U[e] ? V[e] : U[e]);
12 }
13
14 int find_bridge(int u, int edge) {
15     vis[u] = true;
16     arr[u] = Time++;
17     int x = arr[u];
18
19     for(auto & i : g[u]) {
20         int v = adj(u, i);
21         if(!vis[v]) {
22             x = min(x, find_bridge(v, i));
23         }
24         else if(i != edge) {
25             x = min(x, arr[v]);
26         }
27     }

```

```

28
29     if(x == arr[u] && edge != -1) {
30         isBridge[edge] = true;
31     }
32     return x;
33 }
34
35 void dfs1(int u) {
36     int current = cmpno;
37     queue<int> q;
38     q.push(u);
39     vis[u] = 1;
40     temp.push_back(current);
41
42     while(!q.empty()) {
43         int v = q.front();
44         q.pop();
45         comp[v] = current;
46
47         for(auto & i : g[v]) {
48             int w = adj(v, i);
49             if(vis[w]) continue;
50             if(isBridge[i]) {
51                 cmpno++;
52                 tree[current].push_back(cmpno);
53                 tree[cmpno].push_back(current);
54                 dfs1(w);
55             }
56             else {
57                 q.push(w);
58                 vis[w] = 1;
59             }
60         }
61     }
62 }
63
64 int main() {
65     int n, m;
66     cin >> n >> m;
67     g.resize(n + 2); tree.resize(n + 2);
68
69     for(int i = 0; i < m; i++) {
70         cin >> U[i] >> V[i];
71         g[U[i]].push_back(i);
72         g[V[i]].push_back(i);
73     }
74

```

```

75 cmpno = Time = 0;
76 memset(vis, false, sizeof vis);
77
78 for(int i = 0; i < n; i ++){
79     if(!vis[i]){
80         find_bridge(i, -1);
81     }
82 }
83
84 memset(vis, false, sizeof vis);
85 cmpno = 0;
86
87 for(int i = 0; i < n; i ++){
88     if(!vis[i]){
89         temp.clear();
90         cmpno++;
91         dfs1(i);
92     }
93 }
94 }

```

2.6 2-SAT

```

1  /*
2   * Make sure to give the size of n atleast a
3   * larger than original (n + 100).
4   * In the truth_table(), u, v are 1 based indexed.
5   * Truth value of the nodes is calculated in the
6   * satisfiable() function i.e. val[] vector
7   * st[] -> stack
8   * comp[] -> component number of every node
9   */
10
11 class sat_2{
12 public:
13     int n, m, tag;
14     vector<vector<int>> g, grev;
15     vector<bool> val;
16     vector<int> st;
17     vector<int> comp;
18
19     sat_2() {}
20     sat_2(int n) : n(n), m(2 * n), tag(0), g(m +
        1), grev(m + 1), val(n + 1) {}
21
22     void add_edge(int u, int v) { //u or v

```

```

21     auto make_edge = [&](int a, int b) {
22         if(a < 0) a = n - a;
23         if(b < 0) b = n - b;
24         g[a].push_back(b);
25         grev[b].push_back(a);
26     };
27
28     make_edge(-u, v);
29     make_edge(-v, u);
30 }
31
32 void truth_table(int u, int v, vector<int> t)
33 {
34     for(int i = 0; i < 2; i ++){
35         for(int j = 0; j < 2; j ++){
36             if(!t[i * 2 + j]){
37                 add_edge((2 * (i ^ 1) - 1) * u, (2 * (j
38                     ^ 1) - 1) * v);
39             }
40         }
41     }
42 }
43
44 void dfs(int u, vector<vector<int>> & G, bool
45     first) {
46     comp[u] = tag;
47     for(int & i : G[u]) if(comp[i] == -1)
48         dfs(i, G, first);
49     if(first) st.push_back(u);
50 }
51
52 bool satisfiable() {
53     tag = 0; comp.assign(m + 1, -1);
54     for(int i = 1; i <= m; i ++){
55         if(comp[i] == -1)
56             dfs(i, g, true);
57     } reverse(ALL(st));
58
59     tag = 0; comp.assign(m + 1, -1);
60     for(int & i : st){
61         if(comp[i] != -1) continue;
62         tag++;
63         dfs(i, grev, false);
64     }
65
66     for(int i = 1; i <= n; i ++){
67         if(comp[i] == comp[i + n]) return false;
68         val[i] = comp[i] > comp[i + n];
69     }

```

```

64         return true;
65     }
66 };
67 };

```

2.7 Dinic's Max Flow

```

1  // from stanford notebook
2  struct edge {
3      int u, v;
4      ll c, f;
5      edge() { }
6      edge(int _u, int _v, ll _c, ll _f = 0): u(_u), v
          (_v), c(_c), f(_f) { }
7  };
8  int n;
9  vector<edge> edges;
10 vector<vector<int>> > g;
11 vector<int> d, pt;
12
13 void addEdge(int u, int v, ll c, ll f = 0) {
14     g[u].emplace_back(edges.size());
15     edges.emplace_back(edge(u, v, c, f));
16     g[v].emplace_back(edges.size());
17     edges.emplace_back(edge(v, u, 0, 0));
18 }
19 bool bfs(int s, int t) {
20     queue<int> q({s});
21     d.assign(n+1, n+2);
22     d[s] = 0;
23     while(!q.empty()) {
24         int u = q.front(); q.pop();
25         if (u == t) break;
26         for(int k : g[u]) {
27             edge &e = edges[k];
28             if(e.f < e.c && d[e.v] > d[e.u] + 1){
29                 d[e.v] = d[e.u] + 1;
30                 q.push(e.v);
31             }
32         }
33     }
34     return d[t] < n+2;
35 }
36
37 ll dfs(int u, int t, ll flow = -1) {
38     if(u == t || !flow) return flow;

```

```

39     for(int &i = pt[u]; i < (int)(g[u].size()); i++)
40     {
41         edge &e = edges[g[u][i]], &oe=edges[g[u][i]
42             ^1];
43         if(d[e.v] == d[e.u] + 1) {
44             ll amt = e.c - e.f;
45             if (flow != -1 && amt > flow) amt = flow;
46             if(ll pushed = dfs(e.v, t, amt)) {
47                 e.f += pushed;
48                 oe.f -= pushed;
49                 return pushed;
50             }
51         }
52     }
53     return 0;
54 }
55
56 ll flow(int s, int t) {
57     ll ans = 0;
58     while(bfs(s, t)) {
59         pt.assign(n+1, 0);
60         while(ll val = dfs(s, t)) ans += val;
61     }
62     return ans;
63 }

```

2.8 Min Cost Max Flow

```

1  class CostFlowGraph{
2  public:
3      struct Edge{
4          int v, f, c;
5          Edge() {}
6          Edge(int v, int f, int c): v(v), f(f), c(c) {}
7      };
8      V<V<int>> > g;
9      V<Edge> e;
10     V<int> pot;
11     int n;
12     int flow;
13     int cost;
14     CostFlowGraph(int sz) {
15         n=sz;
16         g.resize(n);
17         pot.assign(n, 0);
18         flow=0;

```

```

19     cost=0;
20 }
21 void addEdge(int u,int v,int cap,int c) {
22     g[u].pb((int)e.size());
23     e.pb(Edge(v, cap, c));
24     g[v].pb((int)e.size());
25     e.pb(Edge(u, 0, -c));
26 }
27 void assignPots(int s) {
28     priority_queue<pii, V<pii>, greater<pii>> q;
29     V<int> npot(n, inf);
30     q.push({s, 0});
31     while(!q.empty()) {
32         auto cur=q.top(); q.pop();
33         if(npot[cur.fi] <= cur.se)
34             continue;
35         npot[cur.fi]=cur.se;
36         for(auto i:g[cur.fi]) if(e[i].f>0) {
37             int cst=pot[cur.fi]-pot[e[i].v]+e[i].c;
38             q.push({e[i].v, cst+cur.se});
39         }
40     }
41     for(int i=0; i<n; i++) if(npot[i] != inf) {
42         pot[i] += npot[i];
43     }
44 }
45 void negativeEdges(int s) {
46     pot.assign(n, inf);
47     pot[s]=0;
48     for(int j=0; j<n; j++)
49         for(int i=0; i<(int)e.size(); i++) if(e[i].f
50             >0 && pot[e[i^1].v] != inf) {
51             pot[e[i].v]=min(pot[e[i].v], pot[e[i^1].v]+
52                 e[i].c);
53         }
54 }
55 int augment(int s, int t, int fl, V<bool> &v) {
56     if(s==t)
57         return fl;
58     v[s]=1;
59     for(auto i:g[s]) if(!v[e[i].v] && e[i].f>0 &&
60         (pot[s]-pot[e[i].v]+e[i].c)==0) {
61         int cf=augment(e[i].v, t, min(fl, e[i].f), v);
62         if(cf!=0) {
63             e[i].f-=cf;
64             e[i^1].f+=cf;

```

```

62         return cf;
63     }
64 }
65 return 0;
66 }
67 void mcf(int s, int t, bool neg=0) {
68     int cur=0;
69     V<bool> vis;
70     if(neg)
71         negativeEdges(s);
72     do{
73         vis.assign(n, 0);
74         flow+=cur;
75         cost+=(pot[t]-pot[s]);
76         assignPots(s);
77         cur=augment(s, t, inf, vis);
78     }while(cur);
79 }
80 };

```

2.9 Global Min Cut

```

1 // Adj mat. Stoer-Wagner min cut algorithm.
2 // Running time: O(|V|^3)
3 typedef vector<int> VI;
4 typedef vector<VI> VVI;
5
6 const int INF = 1000000000;
7
8 pair<int, VI> GetMinCut(VVI &weights) {
9     int N = weights.size();
10    VI used(N), cut, best_cut;
11    int best_weight = -1;
12
13    for (int phase = N-1; phase >= 0; phase--) {
14        VI w = weights[0];
15        VI added = used;
16        int prev, last = 0;
17        for (int i = 0; i < phase; i++) {
18            prev = last;
19            last = -1;
20            for (int j = 1; j < N; j++)
21                if (!added[j] && (last == -1 || w[j] > w[last]))
22                    last = j;
23            if (i == phase-1) {

```

```

23 for (int j = 0; j < N; j++) weights[prev][j] +=
    weights[last][j];
24 for (int j = 0; j < N; j++) weights[j][prev] =
    weights[prev][j];
25 used[last] = true;
26 cut.push_back(last);
27 if (best_weight == -1 || w[last] < best_weight)
    {
28     best_cut = cut;
29     best_weight = w[last];
30 }
31     } else {
32     for (int j = 0; j < N; j++)
33         w[j] += weights[last][j];
34     added[last] = true;
35     }
36 }
37 }
38 return make_pair(best_weight, best_cut);
39 }
40
41 int main() {
42     int N;
43     cin >> N;
44     for(int i = 0; i < N; i++) {
45         int n, m;
46         cin >> n >> m;
47         VVI weights(n, VI(n));
48         for (int j = 0; j < m; j++) {
49             int a, b, c;
50             cin >> a >> b >> c;
51             weights[a-1][b-1] = weights[b-1][a-1] = c;
52         }
53         pair<int, VI> res = GetMinCut(weights);
54         cout << "Case #" << i+1 << ": " << res.first
55             << endl;
56     }

```

2.10 Bipartite Matching

```

1 // maximum cardinality bipartite matching using
  augmenting paths.
2 // assumes that first n elements of graph
  adjacency list belong to the left vertex set.

```

```

3 int n;
4 vector<vector<int>> graph;
5 vector<int> match, vis;
6
7 int augment(int l) {
8     if(vis[l]) return 0;
9     vis[l] = 1;
10    for(auto r: graph[l]) {
11        if(match[r]==-1 || augment(match[r])) {
12            match[r]=l; return 1;
13        }
14    }
15    return 0;
16 }
17
18 int matching() {
19     int ans = 0;
20     for(int l = 0; l < n; l++) {
21         vis.assign(n, 0);
22         ans += augment(l);
23     }
24     return ans;
25 }

```

2.11 Hopcraft-Karp

```

1 #define MAX 100001
2 #define NIL 0
3 #define INF (1<<28)
4
5 vector< int > G[MAX];
6 int n, m, match[MAX], dist[MAX];
7 // n: number of nodes on left side, nodes are
  numbered 1 to n
8 // m: number of nodes on right side, nodes are
  numbered n+1 to n+m
9
10 bool bfs() {
11     int i, u, v, len;
12     queue< int > Q;
13     for(i=1; i<=n; i++) {
14         if(match[i]==NIL) {
15             dist[i] = 0;
16             Q.push(i);
17         }
18         else dist[i] = INF;

```



```

19 }
20 dist[NIL] = INF;
21 while(!Q.empty()) {
22     u = Q.front(); Q.pop();
23     if(u!=NIL) {
24         len = G[u].size();
25         for(i=0; i<len; i++) {
26             v = G[u][i];
27             if(dist[match[v]]==INF) {
28                 dist[match[v]] = dist[u] + 1;
29                 Q.push(match[v]);
30             }
31         }
32     }
33 }
34 return (dist[NIL]!=INF);
35 }
36
37 bool dfs(int u) {
38     int i, v, len;
39     if(u!=NIL) {
40         len = G[u].size();
41         for(i=0; i<len; i++) {
42             v = G[u][i];
43             if(dist[match[v]]==dist[u]+1) {
44                 if(dfs(match[v])) {
45                     match[v] = u;
46                     match[u] = v;
47                     return true;
48                 }
49             }
50         }
51         dist[u] = INF;
52         return false;
53     }
54     return true;
55 }
56
57 int hopcroft_karp() {
58     int matching = 0, i;
59     // match[] is assumed NIL for all vertex in G
60     while(bfs())
61         for(i=1; i<=n; i++)
62             if(match[i]==NIL && dfs(i))
63                 matching++;
64     return matching;

```

```

65 }

```

2.12 Hungarian

```

1 // Min cost BPM via shortest augmenting paths
2 // O(n^3).Solves 1000x1000 in ~1s
3 // cost[i][j] = cost for pairing left node i with
4 // right node j
5 // Lmate[i] = index of right node that left node
6 // i pairs with
7 // Rmate[j] = index of left node that right node
8 // j pairs with
9 // The values in cost[i][j] may be +/- . To
10 // perform
11 // maximization, negate cost[][].
12 typedef vector<double> VD;
13 typedef vector<VD> VVD;
14 typedef vector<int> VI;
15
16 double MinCostMatching(const VVD &cost, VI &Lmate,
17                         VI &Rmate) {
18     int n = int(cost.size());
19
20     // construct dual feasible solution
21     VD u(n);
22     VD v(n);
23     for (int i = 0; i < n; i++) {
24         u[i] = cost[i][0];
25         for (int j = 1; j < n; j++) u[i] = min(u[i],
26             cost[i][j]);
27     }
28     for (int j = 0; j < n; j++) {
29         v[j] = cost[0][j] - u[0];
30         for (int i = 1; i < n; i++) v[j] = min(v[j],
31             cost[i][j] - u[i]);
32     }
33
34     // construct primal solution satisfying
35     // complementary slackness
36     Lmate = VI(n, -1);
37     Rmate = VI(n, -1);
38     int mated = 0;
39     for (int i = 0; i < n; i++) {
40         for (int j = 0; j < n; j++) {
41             if (Rmate[j] != -1) continue;

```

```

34     if (fabs(cost[i][j] - u[i] - v[j]) < 1e-10)
35     {
36         Lmate[i] = j;
37         Rmate[j] = i;
38         mated++;
39         break;
40     }
41 }
42
43 VD dist(n);
44 VI dad(n);
45 VI seen(n);
46
47 // repeat until primal solution is feasible
48 while (mated < n) {
49     // find an unmatched left node
50     int s = 0;
51     while (Lmate[s] != -1) s++;
52
53     // initialize Dijkstra
54     fill(dad.begin(), dad.end(), -1);
55     fill(seen.begin(), seen.end(), 0);
56     for (int k = 0; k < n; k++)
57         dist[k] = cost[s][k] - u[s] - v[k];
58
59     int j = 0;
60     while (true) {
61         // find closest
62         j = -1;
63         for (int k = 0; k < n; k++) {
64             if (seen[k]) continue;
65             if (j == -1 || dist[k] < dist[j]) j = k;
66         }
67         seen[j] = 1;
68
69         // termination condition
70         if (Rmate[j] == -1) break;
71
72         // relax neighbors
73         const int i = Rmate[j];
74         for (int k = 0; k < n; k++) {
75             if (seen[k]) continue;
76             const double new_dist = dist[j] + cost[i][k] -
77                 u[i] - v[k];

```

```

79     if (dist[k] > new_dist) {
80         dist[k] = new_dist;
81         dad[k] = j;
82     }
83 }
84
85 // update dual variables
86 for (int k = 0; k < n; k++) {
87     if (k == j || !seen[k]) continue;
88     const int i = Rmate[k];
89     v[k] += dist[k] - dist[j];
90     u[i] -= dist[k] - dist[j];
91 }
92 u[s] += dist[j];
93
94 // augment along path
95 while (dad[j] >= 0) {
96     const int d = dad[j];
97     Rmate[j] = Rmate[d];
98     Lmate[Rmate[j]] = j;
99     j = d;
100 }
101 Rmate[j] = s;
102 Lmate[s] = j;
103
104 mated++;
105 }
106
107 double value = 0;
108 for (int i = 0; i < n; i++)
109     value += cost[i][Lmate[i]];
110
111 return value;
112 }

```

2.13 Link

```

1 struct Node { // Splay tree. Root's pp contains
2     tree's parent.
3     Node *p = 0, *pp = 0, *c[2];
4     bool flip = 0;
5     Node() { c[0] = c[1] = 0; fix(); }
6     void fix() {
7         if (c[0]) c[0]->p = this;
8         if (c[1]) c[1]->p = this;
9         // (+ update sum of subtree elements etc.

```

```

        if wanted)
    }
    void push_flip() {
        if (!flip) return;
        flip = 0; swap(c[0], c[1]);
        if (c[0]) c[0]->flip ^= 1;
        if (c[1]) c[1]->flip ^= 1;
    }
    int up() { return p ? p->c[1] == this : -1; }
    void rot(int i, int b) {
        int h = i ^ b;
        Node *x = c[i], *y = b == 2 ? x : x->c[h],
            *z = b ? y : x;
        if ((y->p = p)) p->c[up()] = y;
        c[i] = z->c[i ^ 1];
        if (b < 2) {
            x->c[h] = y->c[h ^ 1];
            z->c[h ^ 1] = b ? x : this;
        }
        y->c[i ^ 1] = b ? this : x;
        fix(); x->fix(); y->fix();
        if (p) p->fix();
        swap(pp, y->pp);
    }
    void splay() { /// Splay this up to the root.
        Always finishes without flip set.
        for (push_flip(); p; ) {
            if (p->p) p->p->push_flip();
            p->push_flip(); push_flip();
            int c1 = up(), c2 = p->up();
            if (c2 == -1) p->rot(c1, 2);
            else p->p->rot(c2, c1 != c2);
        }
    }
    Node* first() { /// Return the min element of
        the subtree rooted at this, splayed to the
        top.
        push_flip();
        return c[0] ? c[0]->first() : (splay(),
            this);
    }
};

struct LinkCut {
    vector<Node> node;
    LinkCut(int N) : node(N) {}

```

```

    }
    void link(int u, int v) { // add an edge (u, v)
        assert(!connected(u, v));
        make_root(&node[u]);
        node[u].pp = &node[v];
    }
    void cut(int u, int v) { // remove an edge (u,
        v)
        Node *x = &node[u], *top = &node[v];
        make_root(top); x->splay();
        assert(top == (x->pp ? x->c[0]));
        if (x->pp) x->pp = 0;
        else {
            x->c[0] = top->p = 0;
            x->fix();
        }
    }
    bool connected(int u, int v) { // are u, v in
        the same tree?
        Node* nu = access(&node[u])->first();
        return nu == access(&node[v])->first();
    }
    /// Move u to root of represented tree.
    void make_root(Node* u) {
        access(u);
        u->splay();
        if (u->c[0]) {
            u->c[0]->p = 0;
            u->c[0]->flip ^= 1;
            u->c[0]->pp = u;
            u->c[0] = 0;
            u->fix();
        }
    }
    /// Move u to root aux tree. Return the root
    of the root aux tree.
    Node* access(Node* u) {
        u->splay();
        while (Node* pp = u->pp) {
            pp->splay(); u->pp = 0;
            if (pp->c[1]) {
                pp->c[1]->p = 0; pp->c[1]->pp = pp

```

```

92         ; }
93         pp->c[1] = u; pp->fix(); u = pp;
94     }
95     return u;
96 }

```

3 Data Structures

3.1 Implicit Treap

```

1 //1-based with lazy-updates, range sum query
2 struct node {
3     int val, sum, lazy, prior, size;
4     node *l, *r;
5 };
6 const int N = 2e5;
7 node pool[N]; int poolptr=0;
8 typedef node* pnode;
9 int sz(pnode t) { return t?t->size:0; }
10 void upd_sz(pnode t) { if(t) t->size = sz(t->l) +
    1 + sz(t->r); }
11 void lazy(pnode t) {
12     if(!t || !t->lazy) return;
13     t->val+=t->lazy;
14     t->sum+=t->lazy*sz(t);
15     if(t->l)t->l->lazy+=t->lazy;
16     if(t->r)t->r->lazy+=t->lazy;
17     t->lazy = 0;
18 }
19 void reset(pnode t) {
20     if(t) t->sum=t->val;
21 }
22 void combine(pnode& t, pnode l, pnode r) {
23     if(!l || !r) return void(t=l?l:r);
24     t->sum = l->sum + r->sum;
25 }
26 void operation(pnode t) {
27     if(!t) return;
28     reset(t);
29     lazy(t->l); lazy(t->r);
30     combine(t,t->l,t); combine(t,t,t->r);
31 }
32 void split(pnode t, pnode& l, pnode& r, int pos,
    int add = 0) {

```

```

33     if(!t) return void(l=r=NULL);
34     lazy(t); int curr_pos = add + sz(t->l);
35     if(curr_pos<pos) split(t->r,t->r,r,pos,
        curr_pos+1),l=t;
36     else split(t->l,l,t->r,pos,add),r=t;
37     upd_sz(t); operation(t);
38 }
39 void merge(pnode& t, pnode l, pnode r) {
40     lazy(l); lazy(r);
41     if(!l || !r) t = l?l:r;
42     else if(l->prior > r->prior) merge(l->r,l->r,r,
        ),t=l;
43     else merge(r->l, l, r->l), t=r;
44     upd_sz(t); operation(t);
45 }
46 pnode init(int val) {
47     pnode ret = &(pool[poolptr++]);
48     ret->prior = rand(); ret->size = 1;
49     ret->val= val; ret->sum = val; ret->lazy = 0;
50     return ret;
51 }
52 int query(pnode t, int l, int r) {
53     pnode L,mid,R;
54     split(t, L, mid, l-1); split(mid, t, R, r-1);
55     int ans = t->sum;
56     merge(mid, L, t); merge(t, mid, R);
57     return ans;
58 }
59 void upd(pnode t, int l, int r, int val) {
60     pnode L, mid, R;
61     split(t, L, mid, l-1); split(mid, t, R, r-1);
62     t->lazy += val;
63     merge(mid, L, t); merge(t, mid, R);
64 }
65 void insert(pnode& t, int val, int pos) {
66     pnode l;
67     split(t,l,t,pos-1); merge(l,l,init(val));
        merge(t,l,t);
68 }

```

3.2 Segment Tree

```

1 // This code solves problem Help Ashu on
    hackerearth
2 // Iterative segment tree supporting non

```

```

    commutative combiner function
3 // The combiner function and identity of the
   combiner function are taken as constructor
   arguments
4 // Assign the initial input into t[size] to t[2*
   size-1] then call build
5 // Memory 2*size*sizeof(T)
6 // Time complexity O(log(size))
7 #include <bits/stdc++.h>
8 using namespace std;
9 /* Equinox */
10 template<typename T>
11 class SegTree{
12 public:
13     vector<T> t;
14     T identity;
15     T (*combine)(T,T);
16     int size;
17     SegTree(T (*op)(T,T), T e, int n) {
18         combine=op;
19         identity=e;
20         t.assign(2*n,e);
21         size=n;
22     }
23     void build() {for(int i=size-1;i>0;i--)t[i]=
        combine(t[i<<1],t[i<<1|1]);}
24     T query(int l,int r){
25         T lt=identity;
26         T rt=identity;
27         for(l+=size,r+=size;l<=r;r>=1,l>=1){
28             if(l&1) lt=combine(lt,t[l++]);
29             if(!(r&1)) rt=combine(t[r--],rt);
30         }
31         return combine(lt,rt);
32     }
33     void update(int p,T v){for(t[p+=size]=v;p>=1;)t
        [p]=combine(t[p<<1],t[p<<1|1]);}
34 };
35 int32_t main(){
36     int n;
37     cin>>n;
38     SegTree<int> tree([](int a,int b){return a+b
        ;},0,n);
39     for(int i=0;i<n;i++){
40         int a;

```

```

41         cin>>a;
42         tree.t[i+n]=a&1;
43     }
44     tree.build();
45     int q;
46     cin>>q;
47     while(q--){
48         int c,x,y;
49         cin>>c>>x>>y;
50         switch(c){
51             case 0:
52                 tree.update(x-1,y&1);
53             break;
54             case 1:
55                 cout<<(y-x+1)-tree.query(x-1,y-1)<<"\n";
56             break;
57             case 2:
58                 cout<<tree.query(x-1,y-1)<<"\n";
59         }
60     }
61     return 0;
62 }

```

3.3 Lazy Propagation

```

1 // This code solves problem LITE on spoj
2 // Iterative segment tree with lazy propagation
   supporting non commutative combiner functions
3 // The combiner function and identity of the
   combiner function are taken as constructor
   arguments
4 // Also the function for application of lazy nodes
   onto tree nodes is taken as parameter along
   with Zero of lazy node
5 // Assign the initial input into t[size] to t[2*
   size-1] then call build
6 // Memory 2*size*sizeof(T)+2*size*sizeof(L)
7 // Time complexity O(log(size))
8 #include <bits/stdc++.h>
9 using namespace std;
10 /* Equinox */
11 template<typename T,typename L>
12 class SegTree{
13 public:
14     vector<T> t;
15     vector<T> lz;

```

```

16 T identity;
17 L zero;
18 T (*combine) (T,T);
19 void (*apply) (T&,L&,L&,int k);
20 int size;
21 int height;
22 SegTree(T (*op) (T,T), T e, void (*pro) (T&,L&,L&,
    int k), L z, int n) {
23     combine=op;
24     apply=pro;
25     identity=e;
26     zero=z;
27     t.assign(2*n,e);
28     lz.assign(2*n,z);
29     size=n;
30     height = sizeof(int)*8-__builtin_clz(n);
31 }
32 void build() {for(int i=size-1;i>0;i--)t[i]=
    combine(t[i<<1],t[i<<1|1]);}
33 void push(int p) {
34     for(int s=height;s>0;s--) {
35         int i=p>>s;
36         apply(t[i<<1],lz[i<<1],lz[i],1<<(s-1));
37         apply(t[i<<1|1],lz[i<<1|1],lz[i],1<<(s-1));
38         lz[i]=zero;
39     }
40 }
41 void reassign(int p) {
42     for(p>>=1;p>0;p>>=1)
43         if(lz[p]==zero)
44             t[p]=combine(t[p<<1],t[p<<1|1]);
45 }
46 T query(int l,int r) {
47     push(l+=size);
48     push(r+=size);
49     T lt=identity;
50     T rt=identity;
51     for(;l<=r;r>>=1,l>>=1) {
52         if(l&1) lt=combine(lt,t[l++]);
53         if(!(r&1)) rt=combine(t[r--],rt);
54     }
55     return combine(lt,rt);
56 }
57 void update(int p,T v) {push(p+=size);for(t[p]=v;
    p>>=1;)t[p]=combine(t[p<<1],t[p<<1|1]);}

```

```

58 void update(int l,int r,L v) {
59     push(l+=size);
60     push(r+=size);
61     int k=1;
62     int l0=l,r0=r;
63     for(;l<=r;r>>=1,l>>=1,k<<=1) {
64         if(l&1) apply(t[l],lz[l],v,k),l++;
65         if(!(r&1)) apply(t[r],lz[r],v,k),r--;
66     }
67     reassign(l0);
68     reassign(r0);
69 }
70 };
71 int32_t main() {
72     int n,m;
73     cin>>n>>m;
74     SegTree<int,int> s([] (int a, int b) {return a + b
        ;},0,[] (int &v,int &l,int &u,int k) {if(u)v=k-
        v;l^=u;},0,n);
75     while(m--){
76         int c;
77         cin>>c;
78         if(!c){
79             int l,r;
80             cin>>l>>r;
81             s.update(l-1,r-1,1);
82         }
83         else{
84             int l,r;
85             cin>>l>>r;
86             cout<<s.query(l-1,r-1)<<"\n";
87         }
88     }
89     return 0;
90 }

```

4 Math

4.1 Extended Euclid

```

1 #include <bits/stdc++.h>
2
3 using namespace std;
4 using LL = long long;
5

```

```

6 template<typename T> T gcd(T a , T b){return (a ?
    gcd(b % a , a): b);} //supposing a is small and
    b is large.
7 template<typename T> pair<T,T> extend_euclid(T a,
    T b){ //supposing a is small and b is large.
8     pair<T,T> a_one = {1, 0} , b_one = {0 , 1};
9     // b_one is just the second last step's
        coefficient, a_one is the last step's
        coefficient
10    if(!b)return a_one;
11    while(a){
12        /* We first start from writing
13        b = 0(a) + 1(b), for which it's b_one
14        a = 1(a) + 0(b), for which it's a_one
15        b = b % a + (b / a)*a, then
16        */
17        T q = b / a; T r = b % a;
18        T dx = b_one.first - q*a_one.first;
19        T dy = b_one.second - q*a_one.second;
20        b = a; a = r;
21        b_one = a_one;
22        a_one = {dx , dy};
23    }
24    return b_one;
25 }
26
27 int main(){
28     LL a, m; cin >> a >> m;
29     auto ans = extend_euclid(a, m);
30     LL x = (ans.first + m)%m; //Inverse Modulo (m) $
        ax=1 mod(m) and gcd(a,m) == 1
31     cout << (ans.first + m) % m << endl;
32     return 0;
33 }

```

4.2 Fast Fourier Transform

```

1 const long double PI=acos(-1.0);
2 typedef long long ll;
3 typedef long double ld;
4 typedef vector<ll> VL;
5 int bits(int x){
6     int r=0;
7     while(x){
8         r++;
9         x>>=1;

```

```

10     }
11     return r;
12 }
13 int reverseBits(int x,int b){
14     int r=0;
15     for(int i=0;i<b;i++){
16         r<<=1;
17         r|=(x&1);
18         x>>=1;
19     }
20     return r;
21 }
22 class Complex{
23 public:
24     ld r,i;
25     Complex() {r=0.0;i=0.0;}
26     Complex(ld a,ld b) {r=a;i=b;}
27 };
28 Complex operator*(Complex a,Complex b){
29     return Complex(a.r*b.r-a.i*b.i,a.r*b.i+a.i*b.r);
30 }
31 Complex operator-(Complex a,Complex b){
32     return Complex(a.r-b.r,a.i-b.i);
33 }
34 Complex operator+(Complex a,Complex b){
35     return Complex(a.r+b.r,a.i+b.i);
36 }
37 Complex operator/(Complex a,ld b){
38     return Complex(a.r/b,a.i/b);
39 }
40 Complex EXP(ld theta){
41     return Complex(cos(theta),sin(theta));
42 }
43
44 typedef vector<Complex> VC;
45
46 void FFT(VC& A,int inv){
47     int l=A.size();
48     int b=bits(l)-1;
49     VC a(A);
50     for(int i=0;i<l;i++){
51         A[reverseBits(i,b)]=a[i];
52     }
53     for(int i=1;i<=b;i++){
54         int m=(1<<i);
55         int n=m>>1;

```

```

56 Complex wn=EXP((ld)inv*(ld)2.0*PI/(ld)m);
57 for(int j=0; j<l; j+=m) {
58     Complex w(1.0,0.0);
59     for(int k=j; k<j+n; k++) {
60         Complex t1=A[k]+w*A[k+n];
61         Complex t2=A[k]-w*A[k+n];
62         A[k]=t1;
63         A[k+n]=t2;
64         w=w*wn;
65     }
66 }
67 }
68 if(inv==--1) {
69     for(auto &i:A) {
70         i=i/(ld)l;
71     }
72 }
73 }
74
75 VL Convolution(VL & a, VL & b) {
76     int tot_size = (int)a.size() + (int)b.size();
77     int bit = bits(tot_size);
78     int l = 1 << bit;
79     VC A, B, C;
80     A.reserve(l); B.reserve(l); C.reserve(l);
81     for(int i = 0; i < l; i++) {
82         if(i < (int)a.size()) A.pb({(ld)a[i], 0.0});
83         else A.pb({0.0, 0.0});
84         if(i < (int)b.size()) B.pb({(ld)b[i], 0.0});
85         else B.pb({0.0, 0.0});
86     }
87     FFT(A, 1);
88     FFT(B, 1);
89     for(int i = 0; i < l; i++) {
90         C.pb(A[i] * B[i]);
91     }
92     FFT(C, -1);
93     VL c;
94     for(auto & i : C) {
95         c.pb(round(i.r));
96     }
97     return c;
98 }

```

```

1  const ll mod = 163577857, root = 8*23; // =
   998244353
2
3
4  ll modpow(ll x, ll y)
5  {
6      ll ret=1;
7      for (; y; y >>= 1, x = x*x%mod)
8          if (y&1) ret = x*ret%mod;
9      return ret;
10 }
11
12 ll modinv(ll a)
13 {
14     return modpow(a, mod-2);
15 }
16
17 // https://arxiv.org/pdf/1708.01873.pdf
18 void BitReverse(vl& vec, int n) {
19     int L = 31 - __builtin_clz(n);
20     int rev = 0;
21     rep(i, 0, n)
22     {
23         if (i < rev) swap(vec[i], vec[rev]);
24         int tail = i ^ (i + 1), shift = 32 -
25             __builtin_clz(tail);
26         rev ^= tail << (L - shift);
27     }
28
29
30 // For p < 2^30 there is also e . g . (5 << 25, 3)
31 // , (7 << 26, 3) ,
32 // (479 << 21, 3) and (483 << 21, 5) . The last
33 // two are > 10^9.
34 void ntt(vl& x, vl roots) {
35     int n = x.size();
36     if (n == 1) return;
37     BitReverse(x, n);
38
39     int L = 31 - __builtin_clz(n);
40     int mask = (1<<L)-1;
41
42     for (int s=1; s<=L; ++s) {
43         int m = 1<<s;
44         int dx = n/m;
45         for (int k = 0; k<n; k+=m) {

```



```

44     int index = 0;
45     for (int j = 0; j < m/2; ++j)
46     {
47         ll t = roots[index] * (x[k+j+m/2] %
48             mod);
49         ll u = x[k+j];
50         x[k+j] = (u+t);
51         x[k+j+m/2] = (u-t);
52         index = (index+dx) & mask;
53     }
54 }
55
56 for (int i=0; i<n; ++i)
57     x[i] %= mod;
58 }
59
60 // Remember roots to speed up fft
61 vl precalced_roots[30][2];
62
63 void ntt(vl& x, bool inv = false) {
64     int L = sz(x);
65     int s = 31 - __builtin_clz(L);
66
67     if (precalced_roots[s][inv].size() == 0)
68     {
69         ll e = modpow(root, (mod-1) / L);
70         if (inv)
71             e = modpow(e, mod-2);
72         vl roots(L, 1);
73         rep(i, 1, L)
74         {
75             roots[i] = roots[i-1] * e % mod;
76         }
77         precalced_roots[s][inv] = roots;
78     }
79     ntt(x, precalced_roots[s][inv]);
80 }
81
82 vl conv(vl a, vl b) {
83     int s = sz(a) + sz(b) - 1; if (s <= 0) return {};
84     int L = s > 1 ? 32 - __builtin_clz(s - 1) : 0,
85         n = 1 << L;
86     if (s <= 100 || sz(a) < 100 || sz(b) < 100) { // (
87         factor 10 optimization for jaj, j b j =

```

```

10)
86     vl c(s);
87     rep(i, 0, sz(a)) rep(j, 0, sz(b))
88     {
89         c[i+j] = (c[i+j] + a[i] * b[j]) %
90             mod;
91     }
92     return c;
93 }
94 a.resize(n); ntt(a);
95 b.resize(n); ntt(b);
96
97 ll d = modpow(n, mod-2);
98 rep(i, 0, n) a[i] = a[i] * b[i] % mod * d % mod; // c[i] = a
99 [i] * b[i] % mod * d % mod;
100 ntt(a, true);
101 a.resize(s);
102 return a;

```

4.4 Large Factorial

```

1 ll fmod(ll x, ll md, ll p) {
2     V<ll> pre(md);
3     pre[0] = 1;
4     for (ll i = 1; i < md; i++) {
5         if (i % p != 0)
6             pre[i] = (pre[i-1] * i) % md;
7         else
8             pre[i] = pre[i-1];
9     }
10    ll r = 1;
11    while (x) {
12        ll cy = x / md;
13        r = (r * modexp(pre[md-1], cy, md)) % md;
14        r = (r * pre[x % md]) % md;
15        x /= p;
16    }
17    return r;
18 }

```

4.5 Large Modulo Multiplication

```

1 // Finds (a*b)%m when either can be as big as
  10^18

```

```

2 #define ll long long
3 #define ld long double
4 ll mulmod(ll a, ll b, ll m){
5     a%=m;b%=m;
6     ll q = (ll)((ld)a*(ld)b) / (ld)m);
7     ll r = a*b - q*m;
8     if (r > m) r %= m;
9     if (r < 0) r += m;
10    return r;
11 }

```

4.6 Segmented Sieve

```

1 // Segmented Seive
2 // N=sqrt(b)
3 // Time complexity: O(N.log(B-A))
4 #define A 10000000000000LL
5 #define B 10000001000000LL
6 bitset<B-A> p;
7 void seive() {
8     p.set();
9     for(ll i=2; i*i<=B; i++) {
10         for(ll j=((A+i-1)/i)*i; j<=B; j+=i) {
11             p.reset(j-A);
12         }
13     }
14 }

```

4.7 Miller Rabin

```

1 V<int> A{2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31,
2     37, 41};
3 bool Miller(long long p) {
4     if(p < 2) {
5         return false;
6     }
7     if(p != 2 && p % 2 == 0) {
8         return false;
9     }
10    long long s = p - 1;
11    while(s % 2 == 0) {
12        s /= 2;
13    }
14    for(auto & a : A) {

```

```

15        long long temp = s;
16        long long mod = power(a, temp, p);
17        while(temp != p - 1 && mod != 1 && mod != p-1)
18            {
19                mod = mulmod(mod, mod, p);
20                temp *= 2;
21            }
22        if(mod != p - 1 && temp % 2 == 0) {
23            return false;
24        }
25        return true;
26    }

```

4.8 Random Number Generator

```

1 random_device rd;
2 mt19937 gen(rd()); // mersenne twister (only for
3     32 bit unsigned numbers)
4 uniform_int_distribution<int> dis(1, 10000); //
5     dis(L, R) uniformly generates [L, R] (inclusive)
6
7 /* For 64 bit numbers */
8 mt19937_64 gen2(rd());
9 uniform_int_distribution<LL> dis2((int)1e9 + 7, (
10     int)1e10);
11 cout << dis2(gen) << endl << dis2(gen) << endl;

```

5 Strings

5.1 Aho Corasick

```

1 const int N = 500*5005;
2 map<char, int> nxt[N], go[N];
3 int par[N], occ[N], sz = 1, link[N];
4 char parc[N];
5 void add(string& s, int i) {
6     int cur = 1;
7     for(char c : s) {
8         if(!nxt[cur][c]) {
9             sz++;
10            parc[sz]=c, par[sz]=cur, nxt[cur][c]=sz,
11            cur=sz;

```

```

12     else cur=nxt[cur][c];
13 }
14 occ[cur]++;
15 }
16 int GO(int p, char c);
17 int getlink(int p) {
18     if(!link[p]) {
19         if(p==1 || par[p]==1) link[p]=1;
20         else {
21             link[p]=GO(getlink(par[p]),parc[p]);
22             occ[p] += occ[link[p]];
23         }
24     }
25     return link[p];
26 }
27 int GO(int p, char c) {
28     auto it = nxt[p].find(c);
29     if(it == nxt[p].end()) {
30         auto it = go[p].find(c);
31         if (it==go[p].end())
32             return (go[p][c]= p==1 ? 1: GO(getlink
33                 (p),c));
34         else return it->ss;
35     } else return it->ss;
36 }

```

5.2 Suffix Array

```

1 vector<int> suffix_array(string s) {
2     s += "$";
3     int n = s.size();
4     const int alphabet = 256;
5     vector<int> p(n), c(n), cnt(max(alphabet, n),
6         0);
7     for (int i = 0; i < n; i++)
8         cnt[s[i]]++;
9     for (int i = 1; i < alphabet; i++)
10         cnt[i] += cnt[i-1];
11     for (int i = 0; i < n; i++)
12         p[--cnt[s[i]]] = i;
13     c[p[0]] = 0;
14     int classes = 1;
15     for (int i = 1; i < n; i++) {
16         if (s[p[i]] != s[p[i-1]])
17             classes++;
18     }
19 }

```

```

17     c[p[i]] = classes - 1;
18 }
19 vector<int> pn(n), cn(n);
20 for (int h = 0; (1 << h) < n; ++h) {
21     for (int i = 0; i < n; i++) {
22         pn[i] = p[i] - (1 << h);
23         if (pn[i] < 0)
24             pn[i] += n;
25     }
26     fill(cnt.begin(), cnt.begin() + classes,
27         0);
28     for (int i = 0; i < n; i++)
29         cnt[c[pn[i]]]++;
30     for (int i = 1; i < classes; i++)
31         cnt[i] += cnt[i-1];
32     for (int i = n-1; i >= 0; i--)
33         p[--cnt[c[pn[i]]]] = pn[i];
34     cn[p[0]] = 0;
35     classes = 1;
36     for (int i = 1; i < n; i++) {
37         pair<int, int> cur = {c[p[i]], c[(p[i]
38             + (1 << h)) % n]};
39         pair<int, int> prev = {c[p[i-1]], c[(p
40             [i-1] + (1 << h)) % n]};
41         if (cur != prev)
42             ++classes;
43         cn[p[i]] = classes - 1;
44     }
45     c.swap(cn);
46 }
47 p.erase(p.begin()); // remove "$" suffix
48 return p;
49 }
50 vector<int> lcp_array(string s, vector<int> sa) {
51     int n=s.size(), k=0;
52     vector<int> lcp(n,0), r(n,0);
53     for(int i=0; i<n; i++) r[sa[i]]=i;
54     for(int i=0; i<n; i++, k?k--:0) {
55         if(r[i]==n-1) {k=0; continue;}
56         int j=sa[r[i]+1];
57         while(i+k<n && j+k<n && s[i+k]==s[j+k]) k
58             ++;
59         lcp[r[i]]=k;
60     }
61     return lcp;
62 }

```

59 }

5.3 Suffix Tree

```

1  const int N=1000000,    // maximum possible number
    of nodes in suffix tree
2  INF=1000000000; // infinity constant
3  string a;              // input string for which the
    suffix tree is being built
4  int t[N][26],          // array of transitions (state,
    letter)
5  l[N],                  // left...
6  r[N],                  // ...and right boundaries of the
    substring of a which correspond to incoming
    edge
7  p[N],                  // parent of the node
8  s[N],                  // suffix link
9  tv,                    // the node of the current suffix (if
    we're mid-edge, the lower node of the edge)
10 tp,                    // position in the string which
    corresponds to the position on the edge (
    between l[tv] and r[tv], inclusive)
11 ts,                    // the number of nodes
12 la;                    // the current character in the string
13
14 void ukkadd(int c) { // add character s to the
    tree
15     suff;;             // we'll return here after each
    transition to the suffix (and will add
    character again)
16     if (r[tv]<tp) { // check whether we're still
    within the boundaries of the current edge
    // if we're not, find the next edge. If it
    doesn't exist, create a leaf and add
    it to the tree
17         if (t[tv][c]==-1) {t[tv][c]=ts;l[ts]=la;p[
    ts++]=tv;tv=s[tv];tp=r[tv]+1;goto suff
    ;}
18         tv=t[tv][c];tp=l[tv];
19     } // otherwise just proceed to the next edge
20     if (tp==-1 || c==a[tp]-'a')
21         tp++; // if the letter on the edge equal c
    , go down that edge
22     else {
23         // otherwise split the edge in two with
24

```

```

    middle in node ts
25     l[ts]=l[tv];r[ts]=tp-1;p[ts]=p[tv];t[ts][a
    [tp]-'a']=tv;
26     // add leaf ts+1. It corresponds to
    transition through c.
27     t[ts][c]=ts+1;l[ts+1]=la;p[ts+1]=ts;
28     // update info for the current node -
    remember to mark ts as parent of tv
29     l[tv]=tp;p[tv]=ts;t[p[ts]][a[l[ts]]-'a']=
    ts;ts+=2;
30     // prepare for descent
31     // tp will mark where are we in the
    current suffix
32     tv=s[p[ts-2]];tp=l[ts-2];
33     // while the current suffix is not over,
    descend
34     while (tp<=r[ts-2]) {tv=t[tv][a[tp]-'a'];
    tp+=r[tv]-l[tv]+1;}
35     // if we're in a node, add a suffix link
    to it, otherwise add the link to ts
36     // (we'll create ts on next iteration).
37     if (tp==r[ts-2]+1) s[ts-2]=tv; else s[ts
    -2]=ts;
38     // add tp to the new edge and return to
    add letter to suffix
39     tp=r[tv]-(tp-r[ts-2])+2;goto suff;
40 }
41 }
42
43 void Build() {
44     ts=2;
45     tv=0;
46     tp=0;
47     fill(r,r+N,(int)a.size()-1);
48     // initialize data for the root of the tree
49     s[0]=1;
50     l[0]=-1;
51     r[0]=-1;
52     l[1]=-1;
53     r[1]=-1;
54     memset (t, -1, sizeof t);
55     fill(t[1],t[1]+26,0);
56     // add the text to the tree, letter by letter
57     for (la=0; la<(int)a.size(); ++la)
58         ukkadd (a[la]-'a');

```

5.4 Z Algorithm

```

1 vector<int> z_compute(string &S_) {
2     int zn = S_.size();
3     vector<int> z((int)S_.size() + 10);
4     int l = 0, r = 0;
5     z[0] = 0;
6     for (int i = 1; i < zn; i++) {
7         if (i <= r)
8             z[i] = min(r - i + 1, z[i - l]);
9         while (i + z[i] < zn && S_[z[i]] == S_[i + z[i]
10             ])
11             ++z[i];
12         if (i + z[i] - 1 > r)
13             l = i, r = i + z[i] - 1;
14     }
15     return z;

```

6 Geometry

6.1 Geometry Library

```

1
2 /*Returns the orientation of Point C wrt line from
   B to A
   * It returns :-
   * -1 if C lies to left
   * +1 if C lies to the right of the line
   * 0 if C lies on the line
   */
3
4
5
6
7
8
9 int ccw(Point a, Point b, Point c) {
10     int ans = (a - c) ^ (b - c);
11     return ans < 0 ? -1 : ans > 0;
12 }
13
14 /* 0 means outside, 1 means looselyinside the
   polygon(include on the edges of the polygon)*/
15
16 /* To change it strictly inside
   * change the type of this function to int
   * 0 means on the edge / point
   * +1 means strictly inside
17
18
19

```

```

20 * -1 means strictly outside
21 * winding number = 0 means outside
22 * winding number != 0 means inside
23 */
24
25 bool is_inside(auto & p, auto & pt) {
26     int n = (int)p.size();
27     int cnt = 0;
28
29     for(int i = 0; i < n; i++) {
30         if(p[i] == pt) return true;
31         int j = (i + 1) % n;
32         if(p[i].y == pt.y && p[j].y == pt.y) {
33             if(pt.x >= min(p[i].x, p[j].x) && pt.x <=
34                 max(p[i].x, p[j].x))
35                 return true;
36         } else {
37             bool below = p[i].y < pt.y;
38             if(below != (p[j].y < pt.y)) {
39                 auto orientation = ccw(p[i], p[j], pt);
40                 if(!orientation) return true;
41                 if(below == (orientation > 0)) cnt +=
42                     below ? 1 : -1;
43             }
44         }
45     }
46     return (cnt != 0);

```

6.2 Convex Hull

```

1 vector<Point> half_hull(vector<Point> &pts, int t) {
2     vector<Point> hull;
3     hull.pb(pts[0]);
4     hull.pb(pts[1]);
5     for(int i=2; i<n; i++) {
6         while((int)hull.size()>1) {
7             Point p1=hull[(int)hull.size()-2];
8             Point p2=hull.back();
9             if(((p1-pts[i])*(p2-pts[i]))*t)>=0) {
10                 hull.pop_back();
11             }
12             else
13                 break;
14         }

```

```

15     hull.pb(pts[i]);
16 }
17 return move(hull);
18 }
19 vector<Point> convex_hull(vector<Point> &pts) {
20     sort(pts.begin(), pts.end(), [](Point &a, Point
21         &b) {
22         if(a.x==b.x)
23             return a.y<b.y;
24             return a.x<b.x;

```

```

24     });
25     vector<Point> uh, lh;
26     uh=half_hull(pts, 1);
27     lh=half_hull(pts, -1);
28     lh.pop_back();
29     reverse(lh.begin(), lh.end());
30     uh.insert(uh.end(), lh.begin(), lh.end());
31     return move(uh);
32 }

```

7 Formulas

Catalan	$C_0 = 1, C_n = \frac{1}{n+1} \binom{2n}{n} = \sum_{i=0}^{n-1} C_i C_{n-i-1} = \frac{4n-2}{n+1} C_{n-1}$	
Stirling 1st kind	$\begin{bmatrix} 0 \\ 0 \end{bmatrix} = 1, \begin{bmatrix} n \\ 0 \end{bmatrix} = \begin{bmatrix} 0 \\ n \end{bmatrix} = 0, \begin{bmatrix} n \\ k \end{bmatrix} = (n-1) \begin{bmatrix} n-1 \\ k \end{bmatrix} + \begin{bmatrix} n-1 \\ k-1 \end{bmatrix}$	#perms of n objs with exactly k cycles
Stirling 2nd kind	$\left\{ \begin{matrix} n \\ 1 \end{matrix} \right\} = \left\{ \begin{matrix} n \\ n \end{matrix} \right\} = 1, \left\{ \begin{matrix} n \\ k \end{matrix} \right\} = k \left\{ \begin{matrix} n-1 \\ k \end{matrix} \right\} + \left\{ \begin{matrix} n-1 \\ k-1 \end{matrix} \right\}$	#ways to partition n objs into k nonempty sets
Euler	$\langle \begin{matrix} n \\ 0 \end{matrix} \rangle = \langle \begin{matrix} n \\ n-1 \end{matrix} \rangle = 1, \langle \begin{matrix} n \\ k \end{matrix} \rangle = (k+1) \langle \begin{matrix} n-1 \\ k \end{matrix} \rangle + (n-k) \langle \begin{matrix} n-1 \\ k-1 \end{matrix} \rangle$	#perms of n objs with exactly k ascents
Euler 2nd Order	$\llbracket \begin{matrix} n \\ k \end{matrix} \rrbracket = (k+1) \llbracket \begin{matrix} n-1 \\ k \end{matrix} \rrbracket + (2n-k-1) \llbracket \begin{matrix} n-1 \\ k-1 \end{matrix} \rrbracket$	#perms of $1, 1, 2, 2, \dots, n, n$ with exactly k ascents
Bell	$B_1 = 1, B_n = \sum_{k=0}^{n-1} B_k \binom{n-1}{k} = \sum_{k=0}^n \left\{ \begin{matrix} n \\ k \end{matrix} \right\}$	#partitions of $1..n$ (Stirling 2nd, no limit on k)

#labeled rooted trees	n^{n-1}
#labeled unrooted trees	n^{n-2}
#forests of k rooted trees	$\frac{k}{n} \binom{n}{k} n^{n-k}$
$\sum_{i=1}^n i^2 = n(n+1)(2n+1)/6$	$\sum_{i=1}^n i^3 = n^2(n+1)^2/4$
$!n = n \times !(n-1) + (-1)^n$	$!n = (n-1)(!(n-1) + !(n-2))$
$\sum_{i=1}^n \binom{n}{i} F_i = F_{2n}$	$\sum_i \binom{n-i}{i} = F_{n+1}$
$a \equiv b \pmod{x, y} \Rightarrow a \equiv b \pmod{\text{lcm}(x, y)}$	$\sum_{d n} \phi(d) = n$
$ac \equiv bc \pmod{m} \Rightarrow a \equiv b \pmod{\frac{m}{\text{gcd}(c, m)}}$	$(\sum_{d n} \sigma_0(d))^2 = \sum_{d n} \sigma_0(d)^3$
$p \text{ prime} \Leftrightarrow (p-1)! \equiv -1 \pmod{p}$	$\text{gcd}(n^a - 1, n^b - 1) = n^{\text{gcd}(a, b)} - 1$
$\sigma_x(n) = \prod_{i=0}^r \frac{p_i^{(a_i+1)x} - 1}{p_i^x - 1}$	$\sigma_0(n) = \prod_{i=0}^r (a_i + 1)$
$\sum_{k=0}^m (-1)^k \binom{n}{k} = (-1)^m \binom{n-1}{m}$	
$2^{\omega(n)} = O(\sqrt{n})$	$\sum_{i=1}^n 2^{\omega(i)} = O(n \log n)$
$d = v_i t + \frac{1}{2} a t^2$	$v_f^2 = v_i^2 + 2ad$
$v_f = v_i + at$	$d = \frac{v_i + v_f}{2} t$

7.1 The Twelfold Way

Putting n balls into k boxes.

Balls	same	distinct	same	distinct	
Boxes	same	same	distinct	distinct	Remarks
-	$p_k(n)$	$\sum_{i=0}^k \left\{ \begin{matrix} n \\ i \end{matrix} \right\}$	$\binom{n+k-1}{k-1}$	k^n	$p_k(n)$: #partitions of n into $\leq k$ positive parts
size ≥ 1	$p(n, k)$	$\left\{ \begin{matrix} n \\ k \end{matrix} \right\}$	$\binom{n-1}{k-1}$	$k! \left\{ \begin{matrix} n \\ k \end{matrix} \right\}$	$p(n, k)$: #partitions of n into k positive parts
size ≤ 1	$[n \leq k]$	$[n \leq k]$	$\binom{k}{n}$	$n! \binom{k}{n}$	$[cond]$: 1 if $cond = true$, else 0

7.2 DP Optimizations

- **1D1D** $dp(x) = \min_{i=1}^{x-1} \{dp(i) + w[i, x]\}$ and follows quadrangle inequality $w[i, j] + w[i+1, j+1] \leq w[i, j+1] + w[i+1, j]$. Then arg array is non-decreasing. Construct arg array using binary search.

7.3 Some primes

- 7 digits - 2171159, 9368299, 1874351, 9873623, 3934741, 3932941, 4753739, 1251703, 8324893, 5610793
- 8 digits - 59707699, 84765091, 64216913, 36853373, 91814719, 29647939, 99082553, 68007601, 35386633, 91221883
- 9 digits - 267222157, 248334941, 853519241, 879700489, 529560481, 160736231, 308615471, 722344243, 546428819, 528094447
- 12 digits - 744903658181, 805685255317, 901677551977, 645778995493, 951016942451, 743768119319, 463374658853, 390290791217, 730300933471
- 16 digits - 6934008823912991, 6133523110774669, 4707120596051539, 5856250400014373, 5824952666729017, 5619411481414127, 6239941242022171, 3765554534448349, 3773976086888701, 6077904809921143

• **Legendre symbol:** $\left(\frac{a}{b}\right) = a^{(b-1)/2} \pmod{b}$, b odd prime.

• **Heron's formula:** A triangle with side lengths a, b, c has area $\sqrt{s(s-a)(s-b)(s-c)}$ where $s = \frac{a+b+c}{2}$.

• **Pick's theorem:** A polygon on an integer grid strictly containing i lattice points and having b lattice points on the boundary has area $i + \frac{b}{2} - 1$. (Nothing similar in higher dimensions)

• **Euler characteristic:** A finite, connected, planar graph is drawn in the plane without any edge intersections where v denotes $|V|$, e denotes $|E|$ and f denotes the number of faces, then $v - e + f = 2$

• **Baby Step Giant Step:** Given a cyclic group \mathcal{G} of order n , a generator α of the group and a group element β , find x such that $\alpha^x = \beta$

Algorithm:

- Write x as $x = im + j$, where $m = \lceil \sqrt{n} \rceil$ and $0 \leq i < m$ and $0 \leq j < m$.
- Hence, we have $\beta(\alpha^{-m})^i = \alpha^j$.
- $\forall j$ where $0 \leq j < m$: calculate α^j and add them to `std::unordered_map<int, int>`
- $\forall i$ where $0 \leq i < m$: check if $\beta(\alpha^{-m})^i$ exists in the

`std::unordered_map<int, int>`
or not

• **Euler's totient:** The number of integers less than n that are coprime to n are $n \prod_{p|n} \left(1 - \frac{1}{p}\right)$ where each p is a distinct prime factor of n .

Calculation of $\phi(n)$ $\forall n$ where $2 \leq n < 10^6$

- In the regular sieve initialize $\phi(i) = i \forall i$.
- As soon as a prime i is found, update $\phi(j) = \phi(j) - \phi(j)/i$

• **Gauss Generalization and Wilson's theorem:** Let p be an odd prime and α be a positive integer, then in $\mathbb{Z}/(n)$

$$\prod_{k=1}^{\phi(n)} = \begin{cases} 0 & n = 1, \\ -1 & n = 4, p^\alpha, 2p^\alpha, \\ 1 & \text{otherwise} \end{cases}$$

• **Chinese Remainder Theorem:** Given pairwise coprime positive integers n_1, n_2, \dots, n_k and arbitrary integers a_1, a_2, \dots, a_k , the system of simultaneous congruences such that

$$\begin{aligned} x &\equiv a_1 \pmod{n_1} \\ x &\equiv a_2 \pmod{n_2} \\ &\vdots \\ x &\equiv a_k \pmod{n_k} \end{aligned}$$

has a solution, and the solution is unique modulo $N = n_1 n_2 \cdots n_k$. To construct the solution, do the following

1. Compute $N = n_1 \times n_2 \cdots \times n_k$.
2. For each $i = 1, 2, \dots, k$, compute

$$y_i = \frac{N}{n_i} = n_1 n_2 \cdots n_i n_{i+1} \cdots n_k.$$

3. For each $i = 1, 2, \dots, k$, compute $z_i \equiv y_i^{-1} \pmod{n_i}$ using Euclid's extended algorithm
4. The integer $x = \sum_{i=1}^k a_i y_i z_i$ is a solution to the system of the congruences and $x \pmod{N}$ is the unique solution modulo N .

• **Shoelace Formula for Area of Simple Polygon:** Polygon represented by $(x_0, y_0), \dots, (x_{n-1}, y_{n-1})$, then it's area \mathcal{A} is

$$\mathcal{A} = \frac{1}{2} \left| \sum_{i=0}^{n-1} x_i (y_{i+1} - y_{i-1}) \right|$$

where $(i+1) \equiv (i+1) \pmod{n}$
where $(i-1) \equiv (i-1+n) \pmod{n}$

• **Line Intersection Formula:** Given 2 lines

$$\begin{cases} A_1 x + B_1 y + C_1 = 0, \\ A_2 x + B_2 y + C_2 = 0 \end{cases}$$

We find their intersection using Cramer's rule where **Note the minus signs in front of them**

$$x = -\frac{\begin{vmatrix} C_1 & B_1 \\ C_2 & B_2 \end{vmatrix}}{\begin{vmatrix} A_1 & B_1 \\ A_2 & B_2 \end{vmatrix}}, y = -\frac{\begin{vmatrix} A_1 & C_1 \\ A_2 & C_2 \end{vmatrix}}{\begin{vmatrix} A_1 & B_1 \\ A_2 & B_2 \end{vmatrix}},$$

- **Circle-Line Intersection:** Intersection of a circle and a line given by

$$\begin{cases} x^2 + y^2 = r^2 \\ Ax + By + C = 0 \end{cases}$$

If the circle is centered at point (x_c, y_c) , transform the coordinate system using

$$\begin{aligned} x &= X + x_c \\ y &= Y + y_c \end{aligned}$$

Calculate the point closest to origin (x_0, y_0) . It's distance from origin is $d_0 = \frac{|C|}{\sqrt{A^2 + B^2}}$, therefore Point (x_0, y_0) ,

$$\begin{aligned} x_0 &= \frac{-AC}{A^2 + B^2} \\ y_0 &= \frac{-BC}{A^2 + B^2} \end{aligned}$$

If $d_0 < r$, then there are 2 intersections. If $d_0 = r$, then there is only one intersection. If $d_0 > r$, no intersection. Calculate $d = \sqrt{r^2 - \frac{C^2}{A^2 + B^2}}$ and $m = \sqrt{\frac{d^2}{A^2 + B^2}}$. The two points of intersections (a_x, a_y) and (b_x, b_y) are (if $d_0 < r$)

$$\begin{aligned} a_x &= x_0 + B \cdot m, a_y = y_0 - A \cdot m \\ b_x &= x_0 - B \cdot m, b_y = y_0 + A \cdot m \end{aligned}$$

If $d_0 = r$, then (x_0, y_0) is the intersection point which is tangent to the surface.

- **Intersection of Circle and Circle:** Intersection of two circles whose equations are given as follows

$$\begin{cases} x^2 + y^2 = r_1^2 \\ (x - x_2)^2 + (y - y_2)^2 = r_2^2 \end{cases}$$

Subtract these two equations to get the equation of line given as

$$\begin{aligned} Ax + By + C &= 0 \\ A &= -2x_2 \\ B &= -2y_2 \\ C &= x_2^2 + y_2^2 + r_1^2 - r_2^2 \end{aligned}$$

Now, solve this problem for intersection of a line and circle. Now, handle the degenerate case when $x_2 = y_2 = 0$ and equation of line is $C = r_1^2 - r_2^2 = 0$. If the radii of the circles are same, then there are infinitely many intersections, if they differ, then there are no intersections.

- **Konig's theorem:** In any bipartite graph $G = (L \cup R, E)$, the number of edges in a maximum matching is equal to the number of vertices in a minimum vertex cover. Let U be the set of unmatched vertices in L , and Z be the set of vertices that are either in U or are connected to U by an alternating path. Then $K = (L \setminus Z) \cup (R \cap Z)$ is the minimum vertex cover.
- **Dilworth's Theorem:** There exists an antichain A , and a partition of the order into a family P of chains, such that the number of chains in the partition equals the cardinality of A .
- **Mirsky's Theorem:** A poset of height h can be partitioned into h antichains.
- The number of vertices of a graph is equal to its minimum vertex cover number plus the size of a maximum independent set.

- A minimum Steiner tree for n vertices requires at most $n - 2$ additional Steiner vertices.

• **Lagrange polynomial** through points $(x_0, y_0), \dots, (x_k, y_k)$ is $L(x) = \sum_{j=0}^k y_j \prod_{\substack{0 \leq m \leq k \\ m \neq j}} \frac{x - x_m}{x_j - x_m}$

- **Hook length formula:** If λ is a Young diagram and $h_\lambda(i, j)$ is the hook-length of cell (i, j) , then then the number of Young tableaux $d_\lambda = n! / \prod h_\lambda(i, j)$.

• **Moebius inversion formula:** If $f(n) = \sum_{d|n} g(d)$, then $g(n) = \sum_{d|n} \mu(d) f(n/d)$. If $f(n) = \sum_{m=1}^n g(\lfloor n/m \rfloor)$, then $g(n) = \sum_{m=1}^n \mu(m) f(\lfloor \frac{n}{m} \rfloor)$. $\sum_{d|n} \mu(d) = [n = 1]$. $\sum_{i=1}^n \sum_{j=1}^n [gcd(i, j) = 1] = \sum_{d=1}^n \mu(d) \lfloor \frac{n}{d} \rfloor^2$

- **Farey Sequence:** F_n Sequence of reduced fractions with denominators $\leq n$. For neighbors $\frac{a}{b}$ and $\frac{c}{d}$, $bc - ad = 1$.

- #primitive pythagorean triples with hypotenuse $< n$ approx $n/(2\pi)$.

- **Frobenius Number:** largest number which can't be expressed as a linear combination of numbers a_1, \dots, a_n with non-negative coefficients. $g(a_1, a_2) = a_1 a_2 - a_1 - a_2$, $N(a_1, a_2) = (a_1 - 1)(a_2 - 1)/2$. $g(d \cdot a_1, d \cdot a_2, a_3) = d \cdot g(a_1, a_2, a_3) + a_3(d - 1)$. An integer $x > (\max_i a_i)^2$ can be expressed in such a way iff. $x \mid \gcd(a_1, \dots, a_n)$.

7.4 Markov Chains

A Markov Chain can be represented as a weighted directed graph of states, where the weight of an edge represents the probability of transitioning over that edge in one timestep. Let $P^{(m)} = (p_{ij}^{(m)})$ be the probability matrix of transitioning from state i to state j in m timesteps,

and note that $P^{(1)}$ is the adjacency matrix of the graph. **Chapman-Kolmogorov:** $p_{ij}^{(m+n)} = \sum_k p_{ik}^{(m)} p_{kj}^{(n)}$. It follows that $P^{(m+n)} = P^{(m)} P^{(n)}$ and $P^{(m)} = P^m$. If $p^{(0)}$ is the initial probability distribution (a vector), then $p^{(0)} P^{(m)}$ is the probability distribution after m timesteps.

The return times of a state i is $R_i = \{m \mid p_{ii}^{(m)} > 0\}$, and i is *aperiodic* if $\gcd(R_i) = 1$. A MC is aperiodic if any of its vertices is aperiodic. A MC is *irreducible* if the corresponding graph is strongly connected.

A distribution π is stationary if $\pi P = \pi$. If MC is irreducible then $\pi_i = 1/\mathbb{E}[T_i]$, where T_i is the expected time between two visits at i . π_j/π_i is the expected number of visits at j in between two consecutive visits at i . A MC is *ergodic* if $\lim_{m \rightarrow \infty} p^{(0)} P^m = \pi$. A MC is ergodic iff. it is irreducible and aperiodic.

A MC for a random walk in an undirected weighted graph (unweighted graph can be made weighted by adding 1-weights) has $p_{uv} = w_{uv}/\sum_x w_{ux}$. If the graph is connected, then $\pi_u = \sum_x w_{ux}/\sum_v \sum_x w_{vx}$. Such a random walk is aperiodic iff. the graph is not bipartite.

An *absorbing* MC is of the form $P = \begin{pmatrix} Q & R \\ 0 & I_r \end{pmatrix}$. Let $N = \sum_{m=0}^{\infty} Q^m = (I_t - Q)^{-1}$. Then, if starting in state i , the expected number of steps till absorption is the i -th entry in $N1$. If starting in state i , the probability of being absorbed in state j is the (i, j) -th entry of NR .

Many problems on MC can be formulated in terms of a system of recurrence relations, and then solved using Gaussian elimination.

7.5 Burnside's Lemma

Let G be a finite group that acts on a set X . For each g in G let X^g denote the set of elements in

X that are fixed by g . Then the number of orbits

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$$

$$Z(S_n) = \frac{1}{n} \sum_{l=1}^n a_l Z(S_{n-l})$$

7.6 Bezout's identity

If (x, y) is any solution to $ax + by = d$ (e.g. found by the Extended Euclidean Algorithm), then all solutions are given by

$$\left(x + k \frac{b}{\gcd(a, b)}, y - k \frac{a}{\gcd(a, b)} \right)$$

7.7 Misc

7.7.1 Determinants and PM

$$\det(A) = \sum_{\sigma \in S_n} \text{sgn}(\sigma) \prod_{i=1}^n a_{i, \sigma(i)}$$

$$\text{perm}(A) = \sum_{\sigma \in S_n} \prod_{i=1}^n a_{i, \sigma(i)}$$

$$\begin{aligned} pf(A) &= \frac{1}{2^n n!} \sum_{\sigma \in S_{2n}} \text{sgn}(\sigma) \prod_{i=1}^n a_{\sigma(2i-1), \sigma(2i)} \\ &= \sum_{M \in \text{PM}(n)} \text{sgn}(M) \prod_{(i,j) \in M} a_{i,j} \end{aligned}$$

7.7.2 BEST Theorem

Count directed Eulerian cycles. Number of OST given by Kirchoff's Theorem (remove r/c with root) $\# \text{OST}(G, r) \cdot \prod_v (d_v - 1)!$

7.7.3 Primitive Roots

Only exists when n is $2, 4, p^k, 2p^k$, where p odd prime. Assume n prime. Number of primitive roots $\phi(\phi(n))$. Let g be primitive root. All primitive roots are of the form g^k where $k, \phi(p)$ are coprime.

k -roots: $g^{i \cdot \phi(n)/k}$ for $0 \leq i < k$

How to find a primitive root? To test that a is a primitive root of p you need to do the following. First, let $s = \phi(p)$ where $\phi()$ is [the Euler's totient function][1]. If p is prime, then $s = p - 1$. Then you need to determine all the prime factors of s : p_1, \dots, p_k . Finally, calculate $a^{s/p_i} \bmod p$ for all $i = 1 \dots k$, and if you find 1 among residuals then it is NOT a primitive root, otherwise it is.

So, basically you need to calculate and check k numbers where k is the number of different prime factors in $\phi(p)$.

7.7.4 Sum of primes

For any multiplicative f :

$$S(n, p) = S(n, p-1) - f(p) \cdot (S(n/p, p-1) - S(p-1, p-1))$$

7.7.5 Floor

$$\begin{aligned} \lfloor \lfloor x/y \rfloor / z \rfloor &= \lfloor x / (yz) \rfloor \\ x \% y &= x - y \lfloor x/y \rfloor \end{aligned}$$