

1 February

RECURSION

Meaning \rightarrow ~~Calling~~ breaking a large problem into same problem of smaller size

Imagine I want to find $n!$

$$n! = n \times (n-1) \times (n-2) \dots \times 1$$

$$5! = 5 \times 4 \times 3 \times 2 \times 1$$

If we use loops for this that is called the Iterative approach

Code using loops \rightarrow

Iterative Algorithm.

```
int factorial(int n) {  
    int i, ans;  
    if (n == 0)  
        return 1;  
    for (int i = 1; i <= n; i++)  
    {  
        ans *= i;  
    }  
    return ans;  
}
```

$0! = 1$

Recursion \rightarrow base Case
 \rightarrow Recursive Case

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Recursive Way \rightarrow

We can write $n! \rightarrow n \times (n-1)!$

Code \rightarrow

```
int factorial(int n)
```

```
{
    if (n == 0) {
```

```
        return 1; // base case
```

```
    else
```

```
        return (n * factorial(n-1));
```

Explanation

5!

$n \times \text{factorial}(n-1)$

\downarrow

$5 \times \text{factorial}(4)$

Main:

`int n = factorial(3)`

`factorial(3)`

`n = 3`

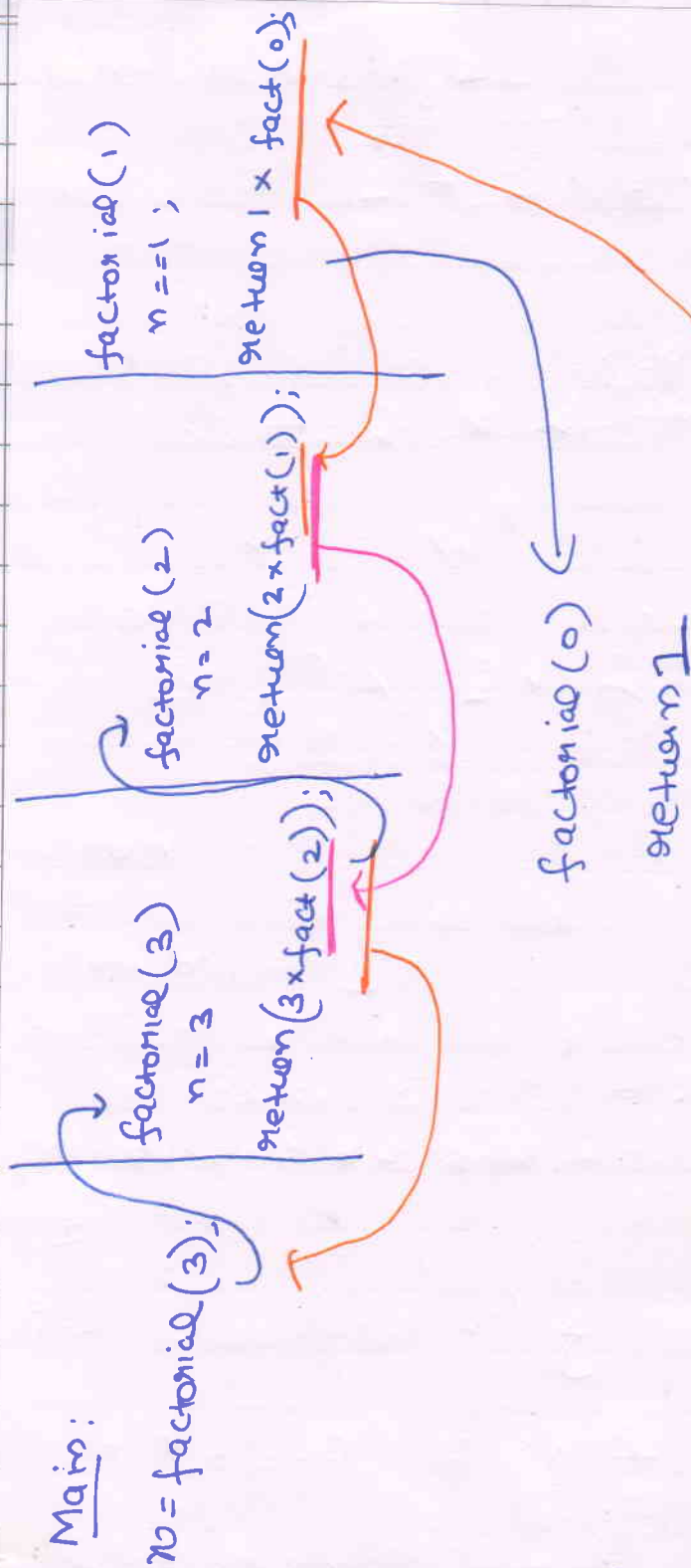
`return 3 * factorial(2)`

`factorial(2)`

`n = 2`

`return 2 * factorial(1)`

`factorial(1)`
`n = 1`



First we want fact 3 we stored $3 \times \text{fact}(2)$, $\text{fact}(2)$ will go to last one and when $f(0)$ comes we returned 1 the value from last goes to second last hence recursion works like winding and unwinding of spring

Questions on Recursion

- Q1) print nth Fibonacci number \rightarrow Applied
- Q2) print power of 2 $\rightarrow 2^n \rightarrow$ babbar
- Q3) print factorial of a number \rightarrow Applied
- Q4) print Counting 1 \rightarrow n and n \rightarrow 1
- Q5) Sum of N natural number using Recursion \rightarrow babbar
Sheet

Solutions

Answer 1) Nth fibonacci number

```
int main() {
```

```
    int x;
```

```
    x = fibonacci(5);
```

```
}
```

Main function

```
int fibonacci(int n) {
```

```
    if (n == 0) return 0;
```

```
    if (n == 1) return 1;
```

else

```
    return (fibonacci(n-1) + fibonacci(n-2));
```

Q Print Powers of 2

$$2^5 \rightarrow 2^4 \times 2^3 \times 2^2 \times 2^1 \times 2^0$$



base case / termination case

~~$2^5 = 2 \times 2 \times 2 \times 2 \times 2$~~

$$2^5 = 2 \times 2^4$$

$2^n \rightarrow 2 \times 2^{(n-1)}$

→ Recursive logic

Code →

```
int main () {
```

```
    x = Power (5) ; → 25
```

```
    cout << x << endl;
```

```
}
```

```
int power (int x) {
```

```
    if (x == 0) return 1 → (20 → 1)
```

```
    return (2 * Power (x-1));
```


2 Print counting from 1 — 10 & 10 — 1
with the help of Recursion.

```
#include <iostream>
```

```
int main () {
```

```
    int start = 1
```

```
    int end = 10
```

```
    Count (end);
```

```
    }
```

```
int Count (int end) {
```

```
    if (end == 0) return;    => 10 - 1
```

```
    cout << end << " ";
```

```
    end --;
```

```
    Count (end);
```

```
}
```

```
int Count (int end) {
```

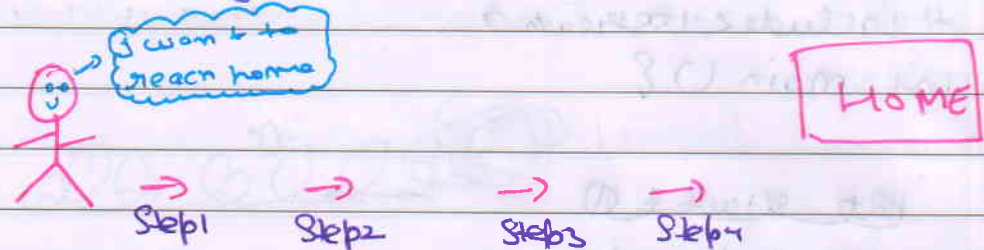
```
    Count (-end);
```

```
    cout << end << " ";
```

```
}
```

Lecture - 2

Recursion kya bolti hai?



A to to Recursion

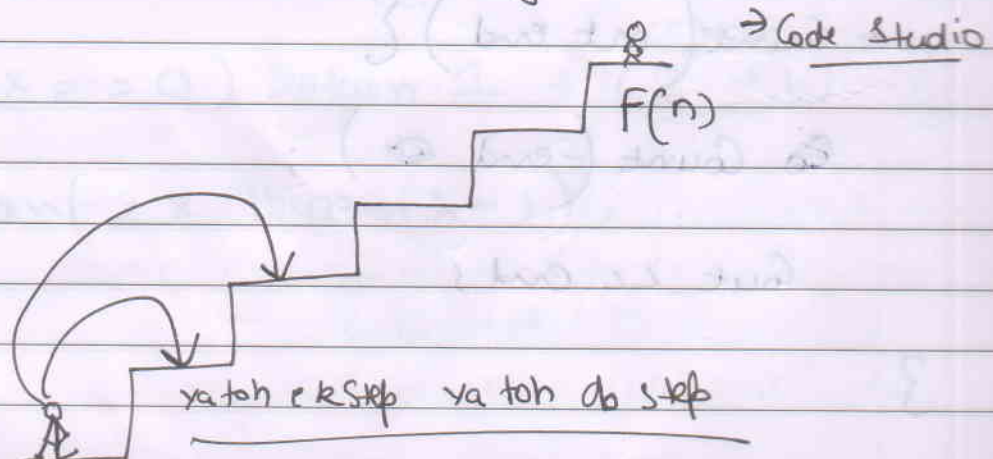
↳ Bhai ek step lena sikh le baaki recursion kardegi.

→ 1 step le lo

→ Recursion use karlo

→ Ruckna kab hai → base case → Jab ghar pounce jaye

Question → Climb Stairs → Count ways to reach n stairs



We need to think for last stair

↳ Hum last stair per Kese aye hain

ya toh 1 chaloge
Man K
(n-1)

ya to 2 chaloge Man K
(n-2)

Recursive Relation $\rightarrow F(n) = F(n-1) + F(n-2)$

Base Case \rightarrow if (n < 0) } \rightarrow Zero K niche
return 0; } to koi stair hi
nai hai

if (n == 1) } \rightarrow Vo already 0 stair per hai
return 1; } So he ~~has~~ has one
way ie vahi per jump
laga ke
oth seedi se oth seedi
per aana.
 \downarrow
Vapis dubara
Kud K aaya

Code \rightarrow int CountStairs (int Stairs)

if (Stairs < 0) return 0; } base Case
if (Stairs == 0) return 1;

Recursive Relation \rightarrow

return (CountStair (n-1) + CountStair (n-2));

8 Feb, 2022

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Recursion - 3

Ques → Check whether ARRAY is sorted or not - ?

Ques → given

| | | | | |
|---|---|---|---|---|
| 5 | 4 | 3 | 2 | 1 |
|---|---|---|---|---|

↑ ↑
We will check for first 2 if true then
check for next two

if Not sorted → Return false!

Code → (Rough)

// base case →

if ~~arr~~ $((\text{Size} == 0) \parallel (\text{Size} == 1))$
Return true;

if $(\text{arr}[0] > \text{arr}[1])$ return false

else

Call function $(\text{arr}+1, \text{Size}-1)$

Next
arr

↓
with less
Size

classmate

Key point

Code →

#include <iostream>

using namespace std;

bool isSorted (int arr[], int size()) {

// base case

if (size == 0 || size == 1) return true;

if (arr[0] > arr[1]) return false;

else
 return isSorted (arr+1, size-1);

Linear Search Using Recursion

Code →

bool linear_search (int arr[], int size, int target)

if (size == 0) return false;

if (arr[0] == target) return true;

else

return linear_search (arr+1, size-1, target);

1 up things to note

Binary Search using Recursion

Note → ARRAY MUST BE SORTED WHILE Binary Search

```

int binarysearch( int start, int end, arrayint arr[], int key)
{
    if (arr.size() == 0) return 0;
    if (start > end) return false;
    if (start < end)
    {
        int mid = (start + end) / 2;

        if (arr[mid] == key)
        {
            return mid + 1;
        }
        else if (arr[mid] > key)
        {
            binarysearch(start, mid - 1, arr);
        }
        else if (arr[mid] < key)
        {
            binarysearch(mid + 1, end, arr);
        }
    }
}
  
```