Evolution, extension, modularity





Recap

- Grammar -> Parser -> Parse Tree -> AST
- Name resolution: recover referential structure
- Checking: find errors not captured by syntax
- Semantics: interpret or compile
- Transformations: at the core of SLE
- Today: after "release" 0.1 -> language evolution
 - and how to engineer to support it

Software evolution

Lehman's laws (some of them)



Manny Lehman

- Continuing change: software must adapt to changes in the environment
- Increasing complexity: maintenance increases complexity unless measures to mitigate are taken.
- Continuing growth: software has to accumulate more features to keep users satisfied.
- Declining quality: without rigorous counter measures, the quality of a system will decline.

Software evolution

different kinds of maintenance

- Corrective: fixing bugs
- Preventive: anticipating things, like code quality problems, or changes in libraries etc.
- Adaptive: respond to changes in environment, security problems in dependencies, deprecation of frameworks, new versions etc.
- Perfective: improvements to, e.g., user-friendliness, performance, ...

Language evolution: possible changes

- Changes to the language design:
 - new constructs (or deprecation), additional checks, change in semantics, ...
- Changes to the implementation:
 - performance improvements, different back-end, new parser algorithm, ...
- Changes w.r.t. code quality:
 - refactoring, rearchitecting, documentation
- Some changes affect users, others may not.

A brief and incomplete history of Java

Java	Year Language Features		Library Features	Other			
1.0	1996	[base language]					
1.1	1997	Inner classes					
1.2	1998						
1.3	2000		Hotspot JVM				
1.4	2002	Assertions					
5	2004	Generics, for-each					
6	2006						
7	2007	Strings in switch, Try-with-resources					
8	2014	Lambdas, default methods					
9	2017	Modules, private methods in interfaces					
10	2018	Var keyword					
11	2018						
12-16	Minor things and previews						
17	2021	Sealed classes, Pattern matching instanced					
18-20	Minor things and previews						
21	2023	Pattern matching in switch	Generational ZGC				

Evolution is inevitable

how to facilitate effective maintenance?

- Separation of concerns: "one concern, one module" (ideally)
 - e.g., separate parsing from type checking
- Modular decomposition: isolate and encapsulate (information hiding)
 - hide implementation details behind an interface
- Layering: built higher-level components on top of foundational components
 - e.g., intermediate languages
- Reuse: "code we don't write, is code we don't have to maintain"

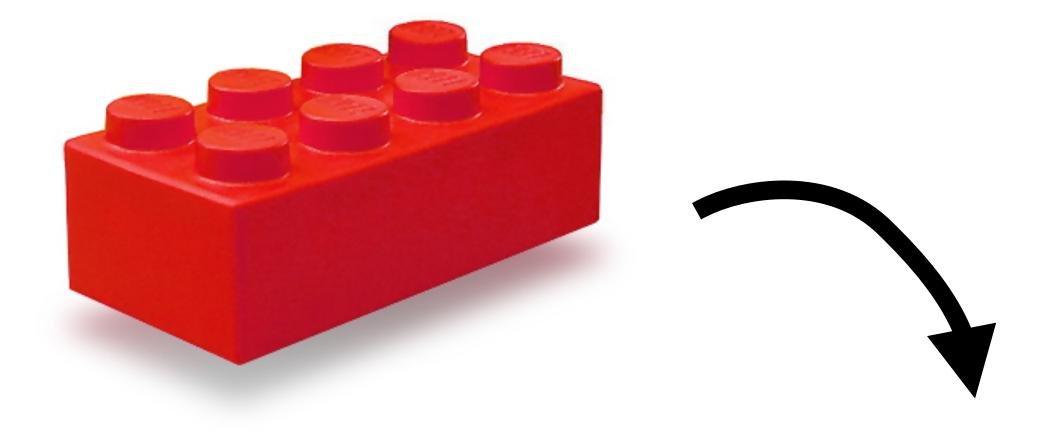
var x = 1 + 2 * 3;

var Check Result **Execute**

Detour: the expression problem

Extensibility

Multiplication feature



Language with addition



Variants

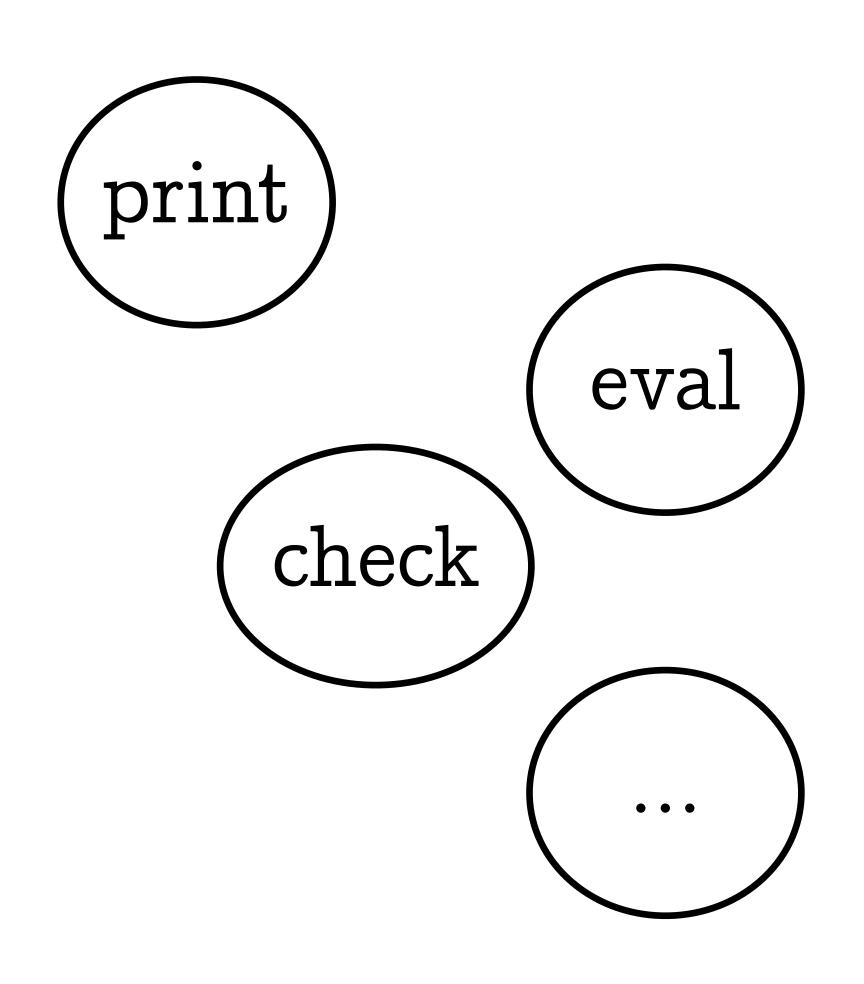
Add

Mul

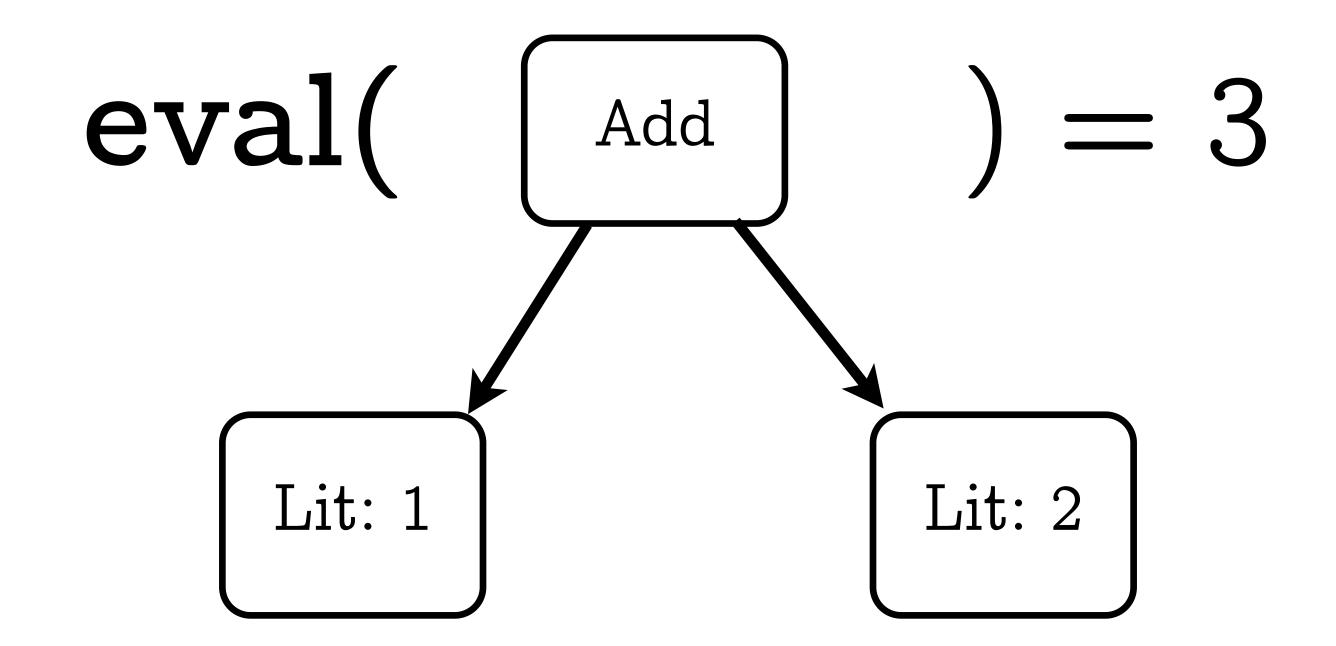
Lit

• • •

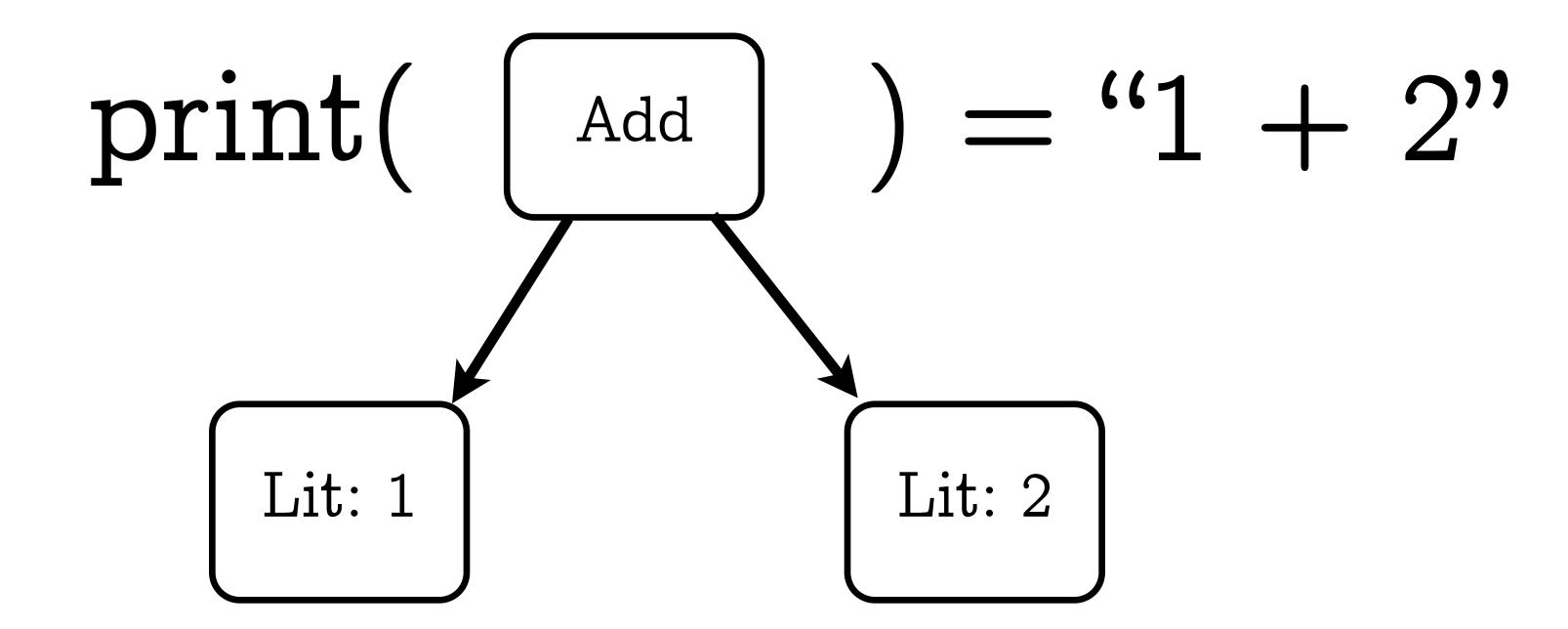
Operations



Evaluate operation



Print operation



Variants

00 languages





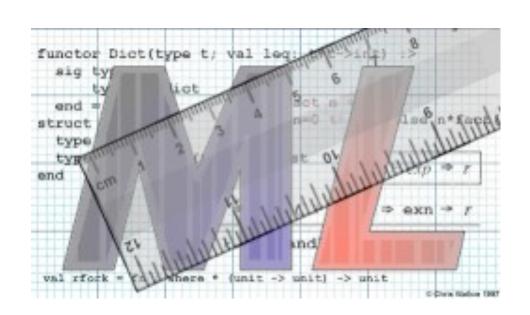




Operations

 $FP\ languages$









The Expression Problem

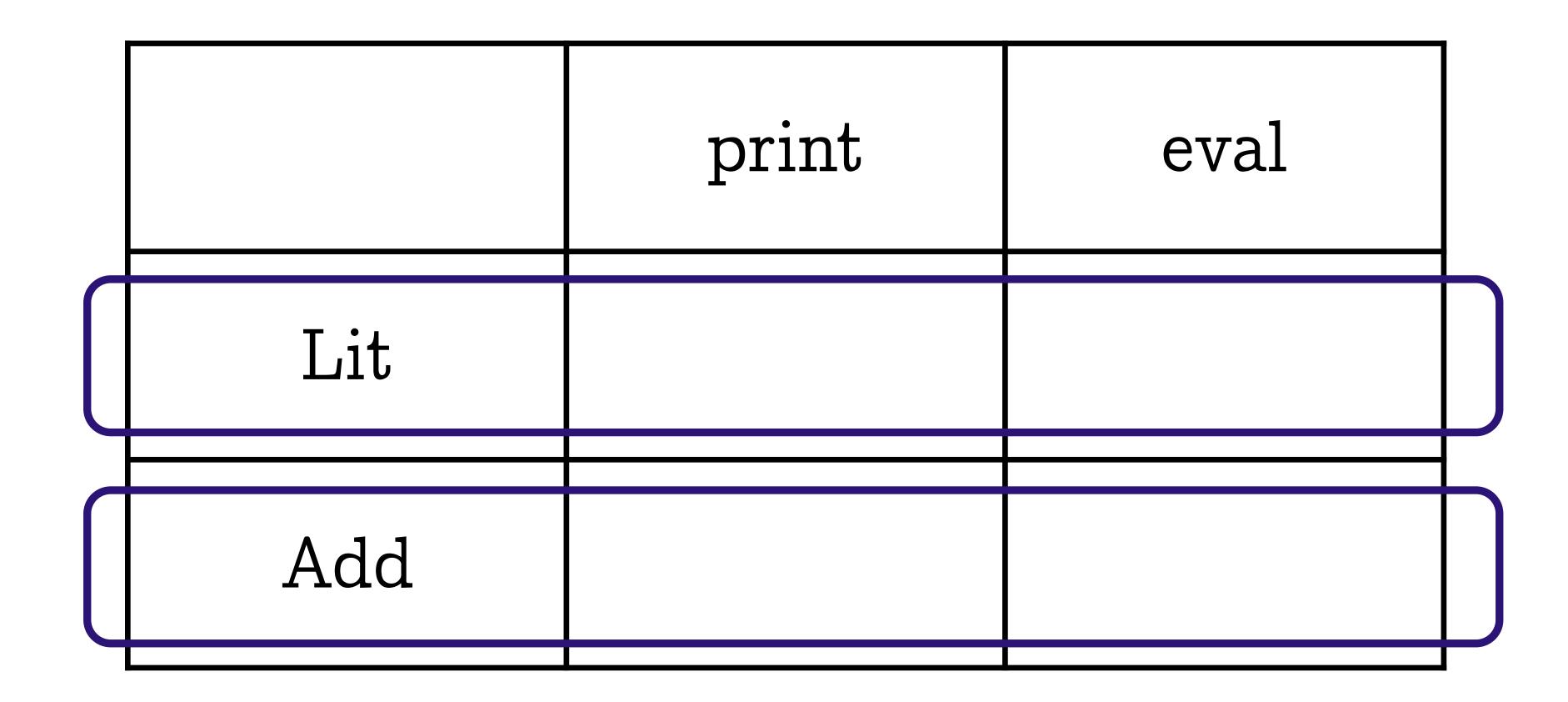
Philip Wadler, 12 November 1998

The Expression Problem is a new name for an old problem. The goal is to define a datatype by cases, where one can add new cases to the datatype and new functions over the datatype, without recompiling existing code, and while retaining static type safety (e.g., no casts). For the concrete example, we take expressions as the data type, begin with one case (constants) and one function (evaluators), then add one more construct (plus) and one more function (conversion to a string).

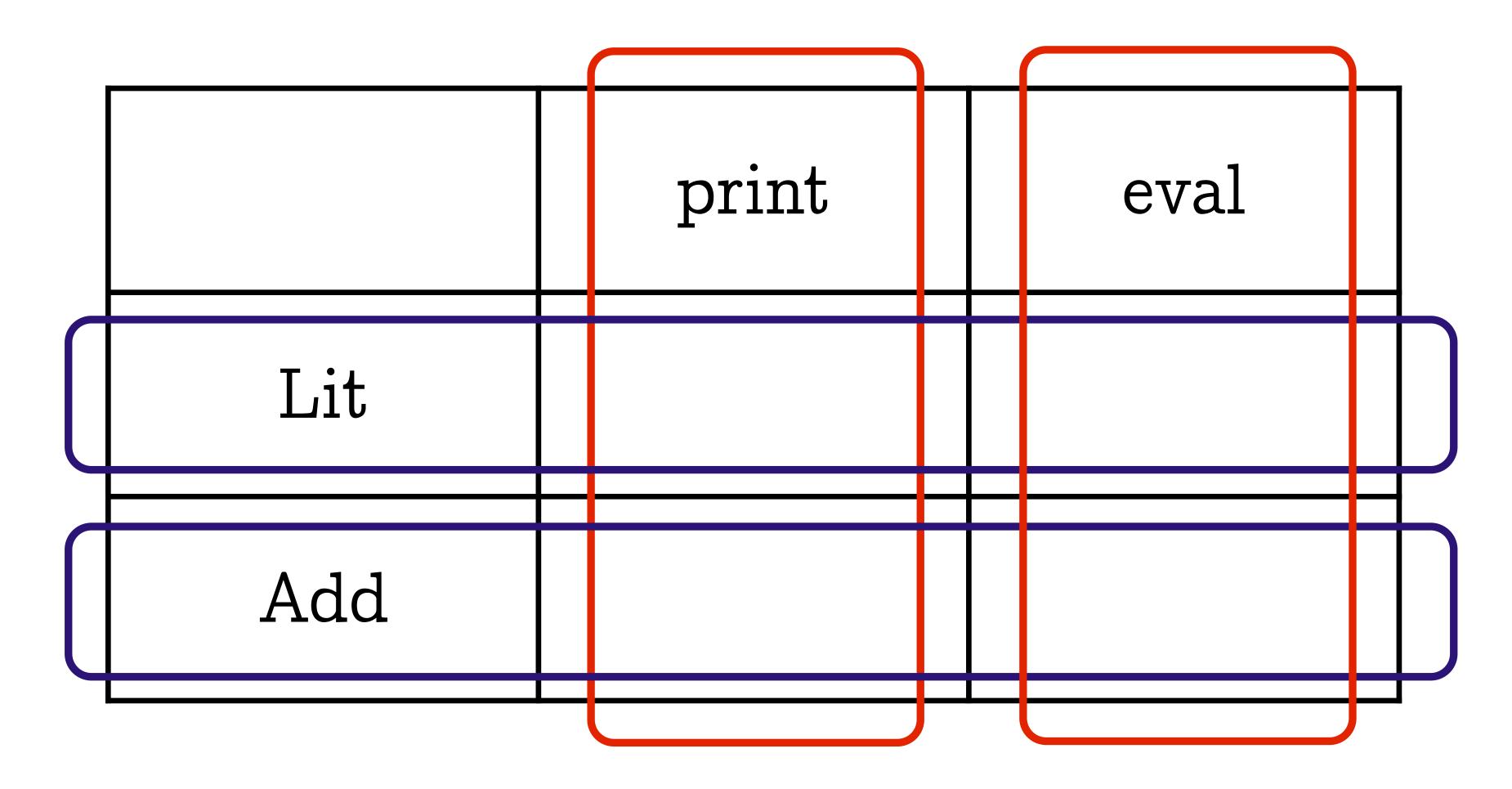
Decomposition

	print	eval
Lit		
Add		

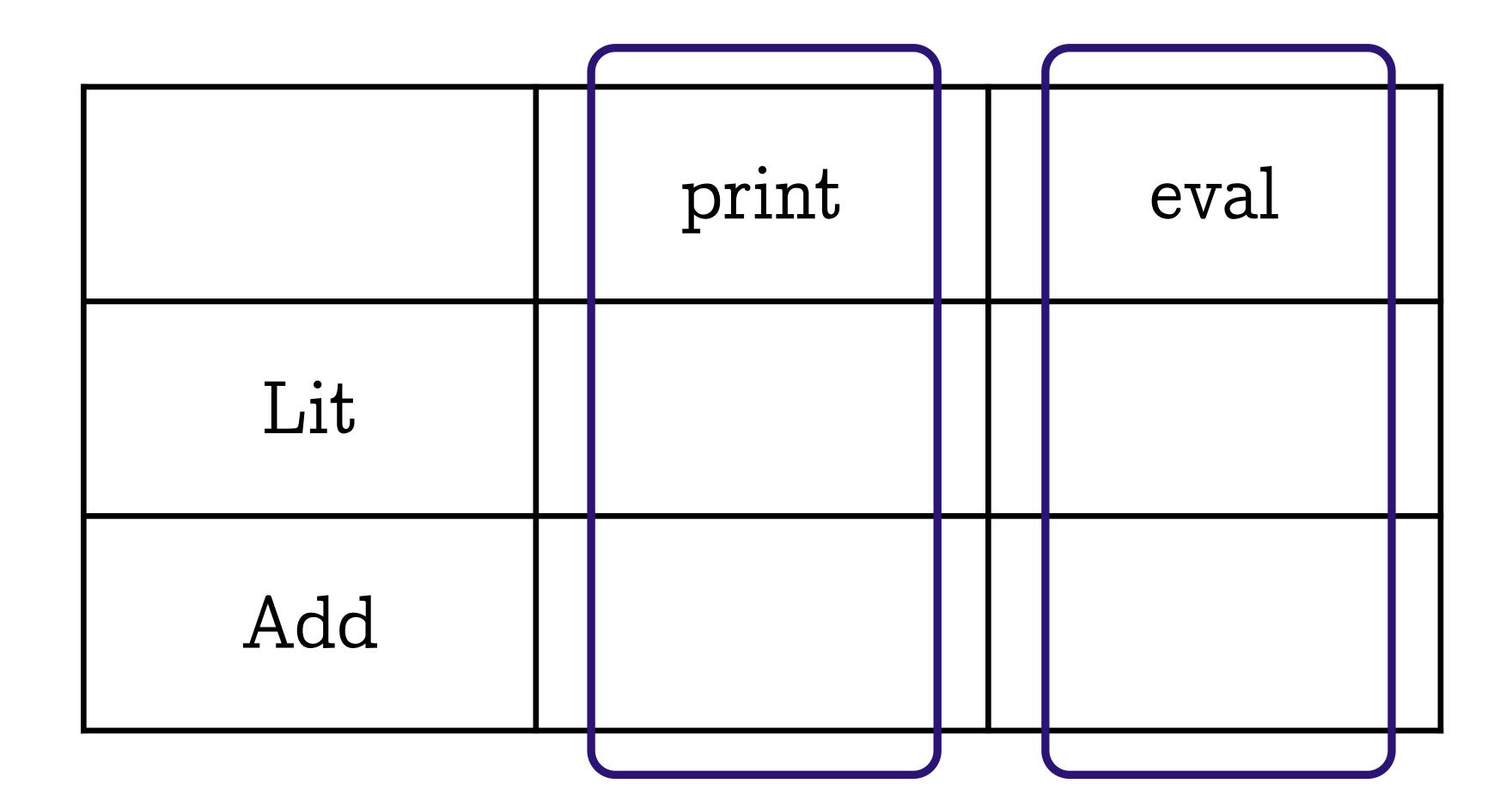
Object-oriented



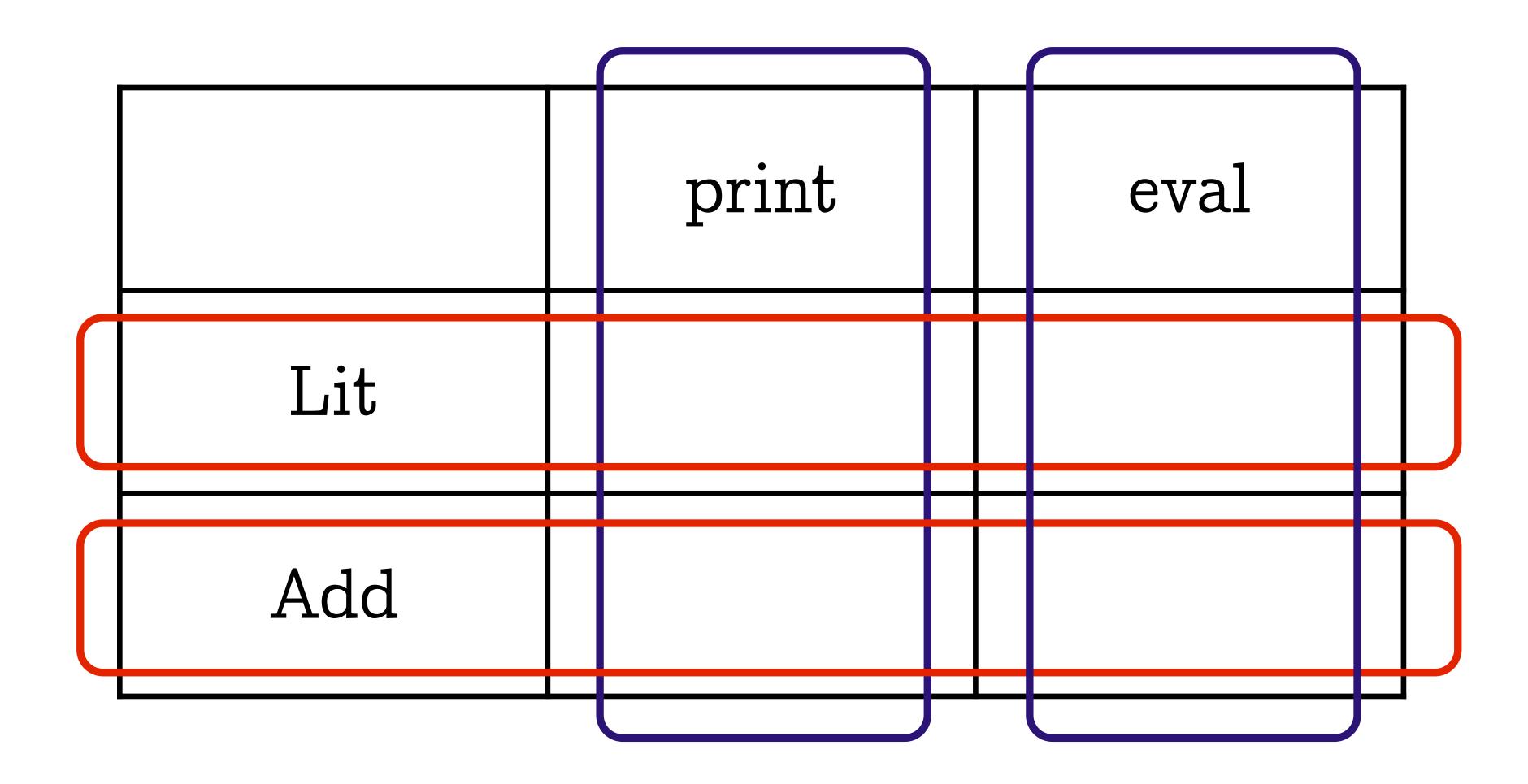
Object-oriented

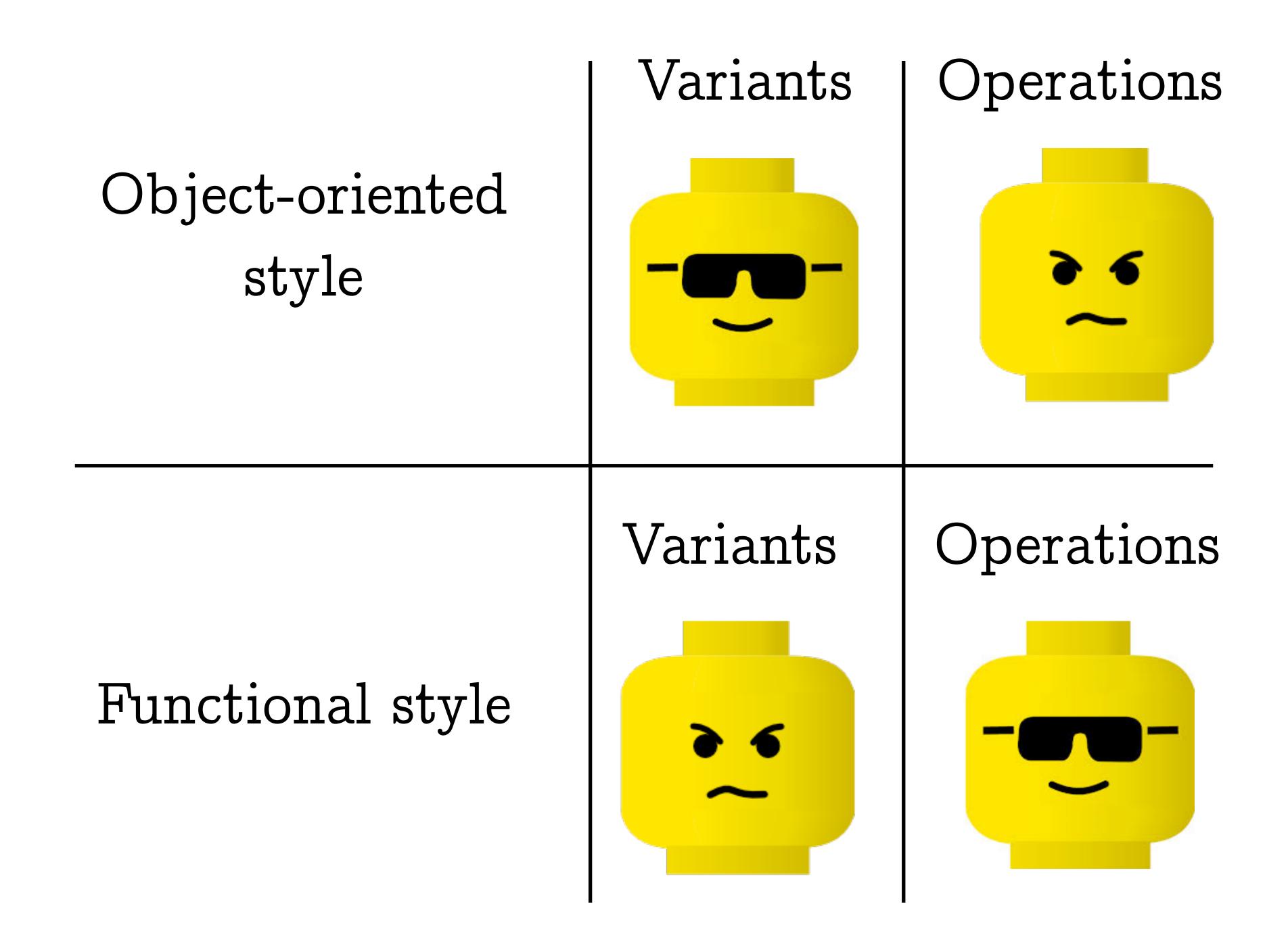


Functional



Functional





Solving the expression problem

	print	eval
Lit		
Add		

Extensibility for the Masses Practical Extensibility with Object Algebras

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Rascal modularity in action

Oberon-0 Language Levels

- L1: Basic control flow statements, constant, type and variable declarations, assignments and expressions.
- L2: Extension with FOR-statements and CASE-statements.
- L3: Definition of (nested) procedures and procedure call statements.
- L4: Support for arrays and records, including subscript and field selection expressions and assignment statements.

	T ₁ Syntax	T ₂ Bind	Т ₃ Снеск	T ₄ Desugar	T ₅ Compile	T ₆ Eval	T' _{7,8} Normalize	T ₇ ToJava	T ₈ ToJVM	T ₉ CFLow
$\overline{L_1}$	244	162	146	_	60	211	_	_	_	121
L_2	80	44	26	39	_	_	_	_	_	71
L_3	73	117	39	_	106	167	_	_	_	36
L_4	90	67	53	_	48	117	743	78	116	_
Total	487	390	264	39	214	495	743	78	116	228

```
🔥 AddLit.rsc U 🗡
                                                      🔥 Mul.rsc U 🔀
sle-master-course > src > expr > 🐴 AddLit.rsc > { } expr
                                                      sle-master-course > src > expr > 🐴 Mul.rsc > ...
                                                              Import in new Rascal terminal
        Import in new Rascal terminal
                                                              module expr::Mul
        module expr::AddLit
        extend lang::std::Layout;
                                                              extend expr::AddLit;
        syntax Expr
                                                              syntax Expr
                                                                 = left Expr "*" Expr;
          = Lit
           bracket "(" Expr ")"
            left Expr "+" Expr
  10
        lexical Lit = [0-9]+;
  11
  12
  13
                                                                                                                 ξη □ ·
🔼 EvalAddLit.rsc U 🔀
                                                      🚵 EvalMul.rsc U 🗡
sle-master-course > src > expr > 🐴 EvalAddLit.rsc > {}
                                                      sle-master-course > src > expr > 🐴 EvalMul.rsc > ...
        Import in new Rascal terminal
                                                              Import in new Rascal terminal
        module expr::EvalAddLit
                                                              module expr::EvalMul
        import expr::AddLit;
                                                              import expr::Mul;
        import String;
                                                              extend expr::EvalAddLit;
        int eval((Expr)`<Lit n>`) = toInt("<n>");
                                                              int eval((Expr)`<Expr lhs> * <Expr rhs>`)
                                                                  = eval(lhs) * eval(rhs);
        int eval((Expr)`(<Expr e>)`) = eval(e);
  10
        int eval((Expr)`<Expr lhs> + <Expr rhs>`)
                                                              test bool testMul() = eval((Expr)\2 * 3\) == 6;
 11
            = eval(lhs) + eval(rhs);
                                                        11
  12
                                                              test bool testAddMul()
                                                        12
  13
        test bool testAdd() = eval((Expr)`1 + 2`)
                                                                  = eval((Expr)^1 + (2 * 3)) == 7;
                                                        13
```

In the lab

- Pick one of the provided language extensions and implement them
- Think about:
 - which operations (check, eval, etc.) are affected
 - how to modularize? (no copy paste, no reimplementation, reuse as much as possible)
 - can you organize the code so that existing ("base QL") code remains the same (minimize modifications to original code)
- In other words: try to extend rather than change your code.