

WEEK:2-3

SETS – RELATIONS-FUNCTIONS



Outlines

- ☐ Sets and Relations
- ☐ Functions



SETS

Mathematical objects are often seen as collections of other objects: $(x, f(x)), y = 3x$

- a square is a collection of points in a plane; $\{...\}$ (...)
- a function is a collection of pairs linking arguments with values.

These collections are called **sets**. Set theory is fundamental part of Mathematics and Mathematics forms the basis of modern software engineering.

The elements in a set are not ordered Examples:

e.g $\{a, b, c\}$ is the same as $\{c, b, a\}$

The elements are not repeated

e.g. $\{a, a, b\}$ is the same as $\{a, b\}$

- the passwords that may be generated using eight lower-case letters
- the collection of programs written in C++ that halt if run for a sufficient time on a computer with unlimited storage

Examples: $\mathbb{N} == \{0, 1, 2, 3, \dots\}$ *set of natural numbers*

$\mathbb{N}_1 == \{1, 2, 3, \dots\}$



Examples:

i) The set of even integers is given as

$$\{x: \mathbb{Z} \mid \exists k: \mathbb{Z} \cdot x = 2k\}$$

1. Questions: Write the set of natural numbers which when divided by 7 leave a remainder of 4



MEMBERSHIP

Suppose $x \in X$ is a predicate, then the predicate is ε

- True if x is in the set X
- False if x is not in the set X

Note: $\forall x: \mathbb{Z} \cdot x > 5 \Rightarrow x \in \mathbb{N}$ means $x: \mathbb{Z}$ declares a new variable of type \mathbb{Z}
 $x \in \mathbb{N}$ is a predicate which is true or false depending on the value previously declared x

- The *cardinality* (#) of a finite set is the *number of its elements*
- Examples:
 - $\# \{\text{red, yellow, blue}\} = 3$
 - $\# \{1, 23\} = 2$
 - $\# \mathbb{Z} = ?$

CARDINALITY



CARTESIAN PRODUCT

- Given two sets A and B , the cartesian **product** of A and B , usually denoted $A \times B$, is the set of all possible pairs (a, b) where $a \in A$ and $b \in B$

$$A \times B \equiv \{ (a, b) \mid a \in A, b \in B \}$$

- Example: PrimaryColor \times Boolean:
 $(\text{red}, \text{true})$, $(\text{red}, \text{false})$,
 $(\text{blue}, \text{true})$, $(\text{blue}, \text{false})$,
 $(\text{yellow}, \text{true})$, $(\text{yellow}, \text{false})$

In a formal description of a software system, we may wish to associate objects of different kinds: names; numbers; various forms of composite data. We may also wish to associate two or more objects of the same kind, respecting order and multiplicity, in that case Z notation is used



POWER SETS

- The **power set** of set S (denoted $Pow(S)$ or $\mathbb{P} S$) is the set of all subsets of S , i.e.,

$$Pow(S) \equiv \{e \mid e \subseteq S\} \text{ or } \underline{\mathbb{P} S = \{s \mid s \subseteq S\}}$$

Note: If S has k elements, $\mathbb{P} S$ has 2^k elements

- Example:
 - $Pow(\{a,b,c\}) = \{\emptyset, \{a\}, \{b\}, \{c\}, \{a,b\}, \{a,c\}, \{b,c\}, \{a,b,c\}\}$

Note: for any S , $\emptyset \subseteq S$ and thus $\emptyset \in Pow(S)$



COMPUTER REPRESENTATION OF SETS

The set M is represented in a computer as a string of binary digits $b_1, b_2, b_3, \dots, b_n$ where n is the cardinality of the universal set \mathbf{U} . The bits b_i (where i ranges over the values $1, 2, \dots, n$) are determined according to the rule:

$b_i = 1$ if i th element of \mathbf{U} is in M .
 $b_i = 0$ if i th element of \mathbf{U} is not in M .

For example, if $\mathbf{U} = \{1, 2, \dots, 10\}$, then the representation of $M = \{1, 2, 5, 8\}$ is given by the bit string 1100100100 where this is given by looking at each element of \mathbf{U} in turn and writing down 1 if it is in M and 0 otherwise.

Similarly, the bit string 0100101100 represents the set $M = \{2, 5, 7, 8\}$, and this is determined by writing down the corresponding element in \mathbf{U} that corresponds to a 1 in the bit string.

DATA STRUCTURES

- Objects from discrete mathematics can model *data structures*.
 - Tuples (records)
 - Relations (tables, linked data structures)
 - Functions (lookup tables, trees and lists)
 - Sequences (lists, arrays)

In a formal specification, it is often necessary to describe relationships between objects:

- this record is stored under that key;
- this input channel is connected to that output channel;
- this action takes priority over that one.

These relationships, and others like them, can be described using simple mathematical objects called **relations**

RELATIONS

- *Relations* can be sets of tuples. They can resemble *tables* or *databases*.
- In Z this can be expressed

```
Employee:  $\mathbb{P}$  EMPLOYEE  
-----  
Employee = {  
  (0019, frank, admin),  
  (0308, philip, research),  
  (7408, aki, research),  
  ...  
}
```

ID	NAME	DEPT
0019	Frank	Admin
0308	Philip	Research
7408	Aki	Research
...

DOMAIN AND RANGE

A relation may contain a great deal of information; often, we require only a small part. To enable us to extract the information that we need the concept of **domain** and **range** functions is used

R is a relation of type $X \leftrightarrow Y$, then the domain of R is the set of elements in X related to something in Y :

$$\text{dom } R = \{x : X; y : Y \mid x \rightarrow y \sqcap R \circ x\}$$

The range of R is the set of elements of Y to which some element of X is related:

$$\text{ran } R = \{x : X; y : Y \mid x \rightarrow y \sqcap R \circ y\}$$

Example: The of people that drive is the domain of *drives*:

$$\text{dom } \textit{drives} . \{helen, Kate, jim, John\}$$

The set of cars that are driven is the range:

$$\text{ran } \textit{drives} . \{Volvo, Benz, Toyota\}$$



BINARY RELATIONS

- *Binary relations* are sets of pairs.
- $\mathbb{P}(\text{NAME} \times \text{PHONE})$ OR
- $\text{NAME} \leftrightarrow \text{PHONE}$
- Binary relations can model lookup tables
- Binary relations are many-to-many relations

NAME	PHONE
Aki	4019
Philip	4107
Doug	4107
Doug	4136
Philip	0113
Frank	0110
Frank	6190
...	...

BINARY RELATION

- A **binary relation** R between A and B is an element of $Pow(A \times B)$, i.e., $R \subseteq A \times B$
- Examples:
 - Parent : Person \times Person
 - Parent == {(John, Jerry), (John, Sam)}
 - Square : $\mathbb{Z} \times \mathbb{N}$
 - Square == {(1,1), (-1,1), (-2,4)}
 - ClassGrades : Person \times {A, B, C, D, F}
 - ClassGrades == {(Todd,A), (Jane,B)}



TERNARY RELATION

- A **ternary relation** R between A , B and C is an element of $Pow(A \times B \times C)$
- Example:
 - FavoriteDrink : Person \times Drink \times Price
 - FavoriteDrink == {(John, Miller, \$2), (Ted, Heineken, \$4), (Steve, Miller, \$2)}
- **N-ary relations** with $n > 3$ are defined analogously (n is the **arity** of the relation)



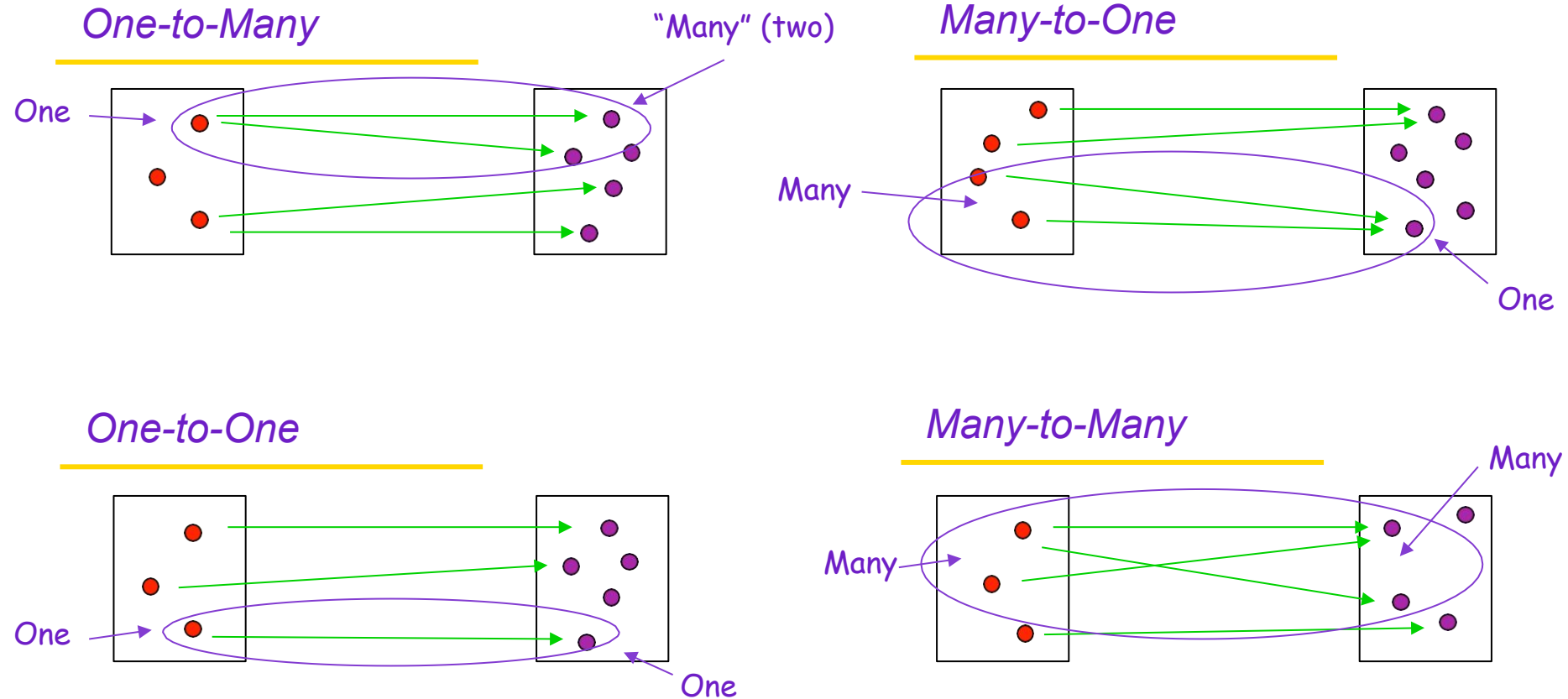
PAIRS

- *Pairs* are tuples with just two components.
(aki, 4117)
- The *maplet* arrow provides alternate syntax without parentheses.

aki \mapsto 4117

- The *projection* operators ***first*** and ***second*** extract the components of a pair.
- `first(aki,4117) = aki`
- `second(aki, 4117) = 4117`

COMMON RELATION STRUCTURES



RELATIONAL CALCULUS

- **Restriction operators** can model database queries.
- **Domain restriction** selects pairs based on their first component.

$\{ \text{doug}, \text{philip} \} \triangleleft \text{phone} =$

$\{ \text{philip} \mapsto 4107,$
 $\text{doug} \mapsto 4107,$
 $\text{doug} \mapsto 4136,$
 $\text{philip} \mapsto 0113 \}$

Range restriction selects according to the second element.

$\text{phone} \triangleright (4000 \dots 4999) = \{$

\dots
 $\text{aki} \mapsto 4019,$
 $\text{philip} \mapsto 4107,$
 $\text{doug} \mapsto 4107,$
 $\text{doug} \mapsto 4136,$
 \dots
 $\}$

OPERATOR SYMBOLS

Operator symbols are defined systematically.

$X \triangleleft R$	Domain restriction
$X \triangleleft\!\!\!\triangleleft R$	Domain anti-restriction
$R \triangleright Y$	Range restriction
$R \triangleright\!\!\!\triangleright Y$	Range anti-restriction

Pictorial Domain & Range restriction operators can also be combined

$\{ \text{doug}, \text{philip} \} \triangleleft \text{phone} \triangleright (4000 \dots 4999) =$

$\{ \text{philip} \mapsto 4107,$
 $\text{doug} \mapsto 4107,$
 $\text{doug} \mapsto 4136 \}$

SEQUENCES

It is sometimes necessary to record the order in which objects are arranged:

For example, *data may be indexed by an ordered collection of keys; messages may be stored in order of arrival; tasks may be performed in order of importance*

Definition and Notation

A **sequence** is an ordered collection of objects. If there are no objects in the collection, the sequence is the *empty sequence*, and is denoted as ‘ $< >$ ’.

For example, the expression $<a, b, c>$ denotes the sequence containing objects a , b , and c , in that order

Concatenation: This is a means of composing sequences in which two sequences are combined in such a way that the elements of one follow the elements of the other, and order is maintained. For example: $<a,b,c>$ and $<d,e>$ becomes $<a,b,c>^{\wedge}<d,e> = <a,b,c,d,e>$



Application: The ticket office in a railway station has a choice of two counters at which tickets may be purchased. There are two queues of people, one at each counter; these may be modelled as sequences:

queue_a = <john, joy, ashley> and **queue_b** = <kim, max, thomas>

John and Kim are at the head of their respective queues, but just as Kim is about to be served the ticket machine at Counter *b* breaks down, and the people waiting there join the end of other queue. Order is maintained, so the result is given by $queue_a \wedge queue_b$, the sequence is given by

<john, joy, Ashley, kim, max, thomas>

A queue of six people forms at Counter *a*

A sequence contains information about a **collection of elements** and the *order* in which they occur. It may be that not all of this information concerns us: we may restrict our attention to elements from a given set using the filter operator: if *s* is a sequence, then $s \mid A$ is the largest subsequence of *s* containing *only those objects that are elements of A*

$\langle a, b, c, d, e, d, c, b, a \rangle \mid \{a, d\} = \langle a, d, d, a \rangle$

The order and multiplicity of elements is preserved



APPLICATION EXAMPLE

Example:

In a train station, there is a destination board displaying a list of trains, arranged in order of departure: This may be modelled as a sequence of pairs, each recording a time and a destination.

$trains = \langle (10:15, London), (10:38, Edinburgh), (10:40, London), (11:15, Birmingham), (11:20, Reading), (11:40, London) \rangle$

Suppose John is interested only in those trains that are going to London; he would be satisfied with the filtered sequence

$$trains \mid \{ t : Time \mid (t, london) \}$$

that is, $\langle (10:15, london), (10:40, london), (11:40, london) \rangle$

It may be that we need to refer to the first element of a sequence, or to the part of the sequence that follows the first element; these are called the **head** and **tail**, respectively.

$head \langle a, b, c, d, e \rangle = a$

$tail \langle a, b, c, d, e \rangle = \langle b, c, d, e \rangle$

Note: the head of a sequence is an element, while the tail is another sequence. If s is any non-empty sequence, then

$S = \langle head\ s \rangle ^ tail\ s$



FUNCTIONS

- *Functions* are binary relations where each element in the domain appears just once. Each domain element is a *unique key*.
- A function cannot be a many-to-many or even one-to-many relation

phonef: NAME \mapsto PHONE

phonef = {
...
aki \mapsto 4019,
philip \mapsto 4107,
doug \mapsto 4107,
frank \mapsto 6190,
...
}

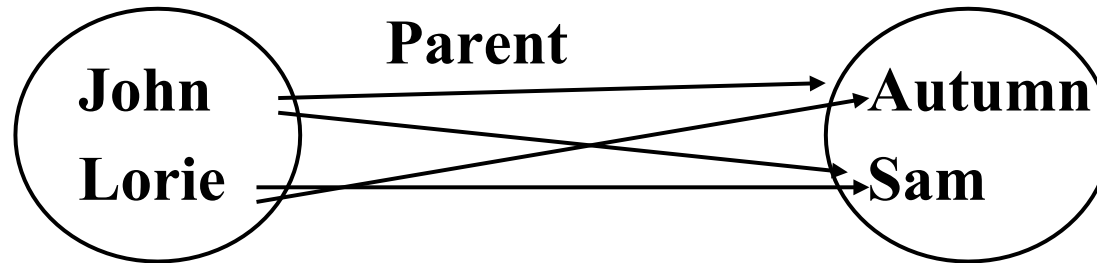
Function application is a special case of relational image. It associates a domain element with its unique range element.

phonef $\{ \{ \text{doug} \} \} = 4107$

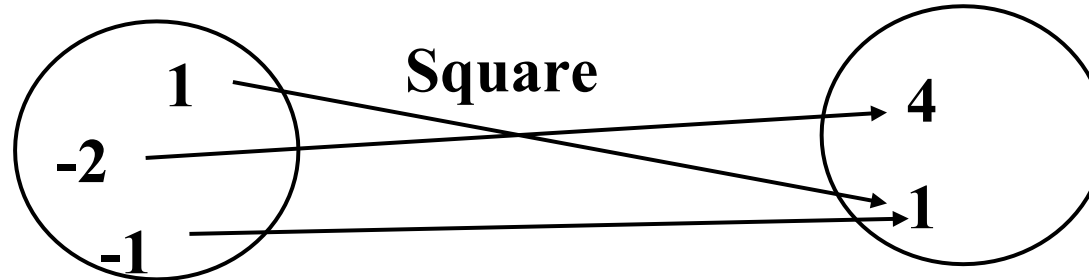
phonef (doug) = 4107

phonef doug = 4107

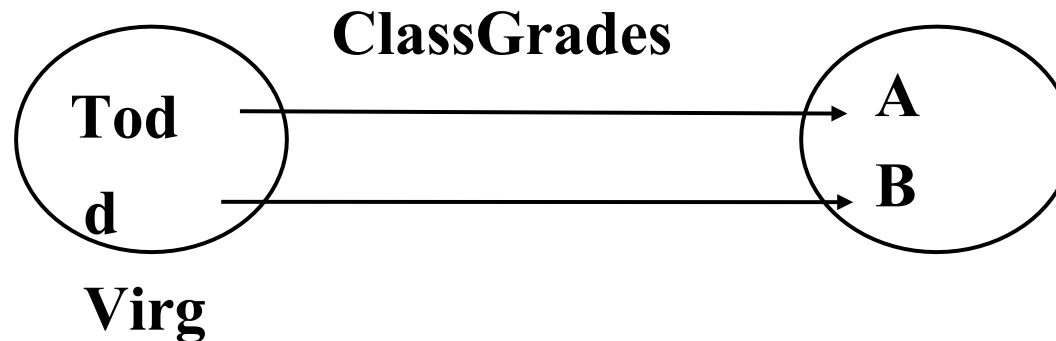
RELATIONS VS. FUNCTIONS



Many-to-many



Many-to-one



One-to-one



SPECIAL KINDS OF FUNCTIONS

- Consider a function f from S to T
- f is *total* if defined for all values of S
- f is *partial* if defined for some values of S
- Examples
 - Squares : $\mathbb{Z} \rightarrow \mathbb{N}$, Squares = $\{(-1,1), (2,4)\}$
 - Abs = $\{(x,y) : \mathbb{Z} \times \mathbb{N} \mid$
 $(x < 0 \text{ and } y = -x) \text{ or } (x \geq 0 \text{ and } y = x)\}$



SPECIAL FUNCTIONS ON SEQUENCES

Concatenation

$$\langle a, b \rangle \wedge \langle b, a, c \rangle = \langle a, b, b, a, c \rangle$$

Head

$$\left| \begin{array}{l} \text{Head} : \text{seq}_1 A \rightarrow A \\ \hline \forall s : \text{seq}_1 A \bullet \text{head}(s) = s(1) \end{array} \right|$$

$$\text{head}\langle c, b, \mid b \rangle = c$$

Tail

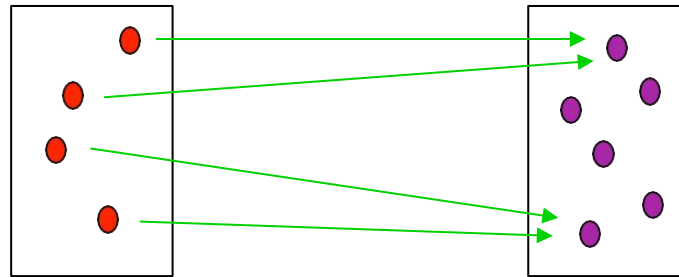
$$\left| \begin{array}{l} \text{Tail} : \text{seq}_1 A \rightarrow A \\ \hline \forall s : \text{seq}_1 A \bullet \langle \text{head}(s) \wedge \text{tail}(s) \rangle = s \end{array} \right|$$

$$\text{tail}\langle c, \textcolor{red}{b}, \textcolor{red}{b} \rangle = \langle \textcolor{red}{b}, \textcolor{red}{b} \rangle$$

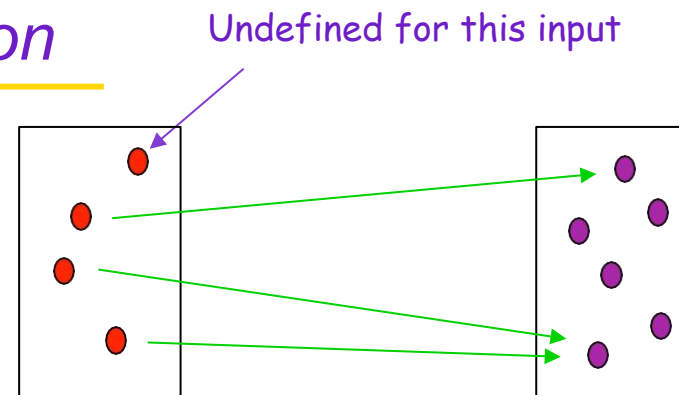


FUNCTION STRUCTURES

Total Function



Partial Function



Note: the empty relation is a partial function



TYPES OF FUNCTIONS

A function $f: S \rightarrow T$ is

Surjective Function

- **Surjective** (onto) if every element of the domain is mapped to some element of the range. some domain elements may be mapped to more than one range elements. (Total Injections)

Injective Function

- *one-to-one* (*injective*) if no image element is associated with multiple domain elements (Partial injections)

Bijjective Function

- **Bijjective** (one-to-one and onto) iff it is both injective and surjective. (Equivalently, every element of the domain is mapped to exactly one element of the range.) A bijective function is a bijection (one-to-one correspondence), and is reversible.

SPECIFYING FUNCTIONS

1. Using a Look-up- Table

If a function $f : A \rightarrow B$ is finite (and not too large) we can specify the function explicitly by listing all the pairs (a, b) in the subset $A \times B$ where $f(a) = b$

e.g

address : PassportNo \rightarrow Address

PassportNo	Address
A001017	77 Sunset Strip
...	...
...	...
G707165	19 Mail Street
...	...



2. Giving an Algorithm

A function $f : A \rightarrow B$ is specified by an algorithm(i.e a program) such that given any element a in the domain of f , the element $f(a)$ can be computed using the algorithm

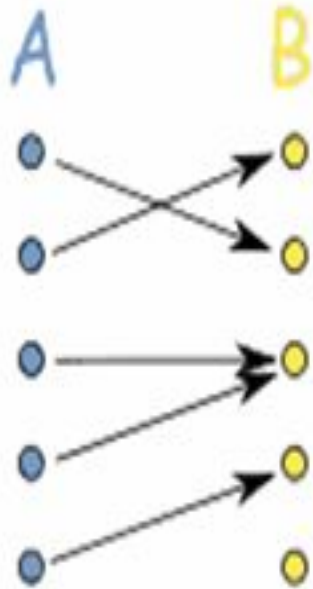
```
input  $n : \mathbb{N}$ 
var  $x, y$ : integer;
begin
 $x := n$ ;  $y := 0$ ;
while  $x \neq 0$  do
  begin
     $x := x - 1$ ;  $y := y + 2$ 
  end;
write( $y$ )
end.
```

This algorithm computes the function. But how can we prove this?

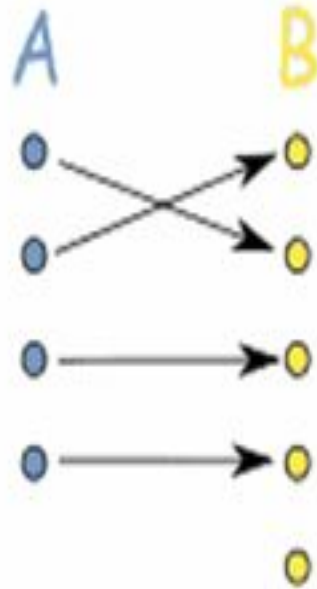
$$\left| \begin{array}{l} \text{double} : \mathbb{N} \rightarrow \mathbb{N} \\ \hline \forall n : \mathbb{N} \bullet \text{double}(n) = 2n \end{array} \right.$$



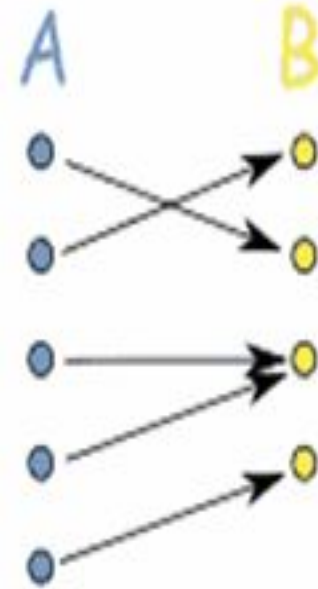
ILLUSTRATIONS



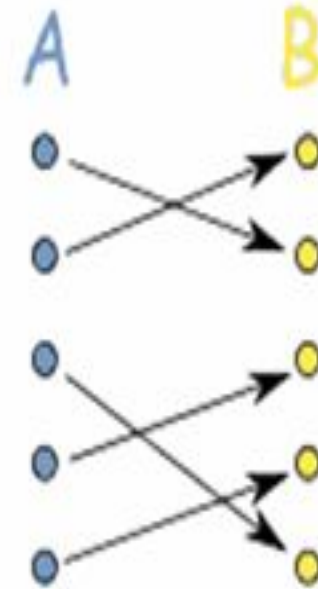
General
Function



Injective
Not surjective



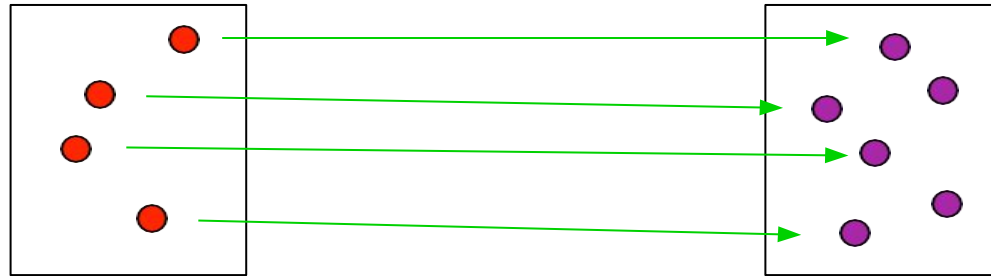
Surjective
Not injective



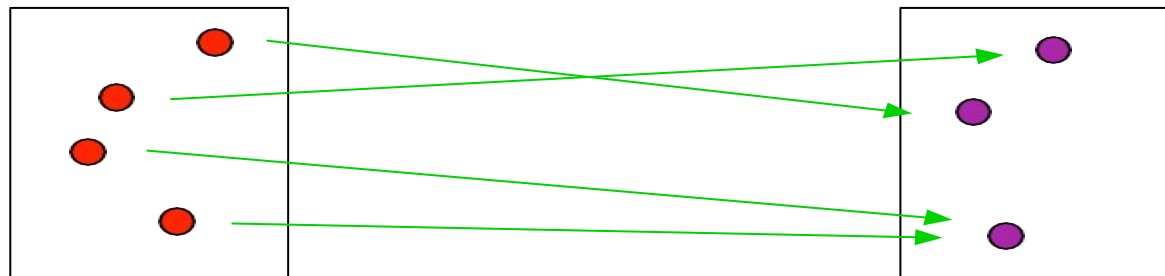
Bijective
(injective and
surjective)

TYPES OF FUNCTION

Injective(one-to-one Function)



Surjective(onto)Function



BINARY RELATIONS AND FUNCTIONS

$X \leftrightarrow Y$	Binary relations: many-to-many
$X \rightharpoonup Y$	Partial functions: many-to-one Some function applications undefined
$X \rightarrow Y$	Total functions: All function applications defined
$X \succcurlyeq Y$	Partial injections: one-to-one Inverse is also a function
$X \succcurlyeq Y$	Total injections
$X \succcurlyeq Y$	Bijections: cover entire range (onto)

All have type $\mathbb{P}(X \times Y)$

CLASS WORK

- What kind of function/relation is Abs?
 - $\text{Abs} = \{(x,y) : \mathbb{Z} \times \mathbb{N} \mid (x < 0 \text{ and } y = -x) \text{ or } (x \geq 0 \text{ and } y = x)\}$
- How about Squares?
 - $\text{Squares} : \mathbb{Z} \times \mathbb{N}, \quad \text{Squares} = \{(-1,1),(2,4)\}$



EXERCISES

