

## 1 Asymptotic Analysis

### 1.1 Big O

#### 1.11 Explanation

**Big-O notation** is a mathematical concept to describe the efficiency of algorithms in computer science, specially to describe the time-complexity and space-complexity. Big-O notation allows people to discuss in the abstract that how the performance of an algorithm changes when the size grows. For example:

**Time-complexity:** it refers to the relationship between the execution time of algorithms and input size  $n$ .

- $O(1)$ : constant time. Whatever the input size, the execution time is always the same.
- $O(n)$ : linear time. When the input size increases, the execution time increases linearly as well.
- $O(n^2)$ : quadratic time. When the input size increases, the execution time increases quadratically.
- $O(\log n)$ : logarithmic time. When the input size increases, the execution time increases logarithmically.

**Space-complexity:** it refers to the relationship between the required memory space and the size of the inputs by an algorithm during execution. The pattern is similar to the time complexity.

**Big O** is an upper bound notation that describes the maximum resource consumption of an algorithm in the worst case. This is particularly important for measuring the performance of an algorithm under the most unfavorable conditions.

#### 1.12 Exercises

Analyze the time complexity of the Merge Sort algorithm for the following array:

[3,1,4,1,5,9,2,6]

#### 1.13 Solution

Merge sort is a divide-and-conquer sorting algorithm. It works by recursively splitting the array into two halves, sorting each half, and then merging the sorted halves together. The specific step is as follows:

**Divide:**

- Split the array into two halves that [3,1,4,1] and [5,9,2,6].
- Further split each subarray until it only has one element, such as [3],[1],[4],[1],[5],[9],[2],[6].

**Conquer:**

- Merge adjacent two subarrays and sort them like [1,3],[1,4],[5,9],[2,6].
- Keep merging adjacent two sorted subarrays, [1,1,3,4],[2,5,6,9].

**Combine:**

- Merge the two halves, [1,1,2,3,4,5,6,9].

**Time Complexity: is  $O(n \log n)$**

**Divide Step:**

- The array will be split into two subarrays recursively, and each split takes  $O(1)$ . Due to the array is divided  $\log n$  time, so the total time is  $O(n \log n)$ .

**Merge Step:**

- Each step merges  $n$  elements of two subarrays, the merging process for each recursion level is linear, so it's  $O(n)$ .

**Total Time Complexity:**

- Given there are  $O(\log n)$  levels need to be split, and each level takes  $O(n)$  to merge, so the total time complexity of Merge Sort is  $O(n \log n)$ .

**Space Complexity: is  $O(n)$**

- Merge Sort need additional array to store the auxiliary arrays for merging, so for an array of size  $n$ , the additional space requirement is  $O(n)$ .

## 1.2 Big Omega

### 1.21 Explanation

**Big- $\Omega$  notation is used to describe the time complexity or space complexity of an algorithm in the best case. It indicates how much time or space that the algorithm takes at least for all possible inputs. Big- $\Omega$  could help to understand the best performance under the best conditions.**

**Time-complexity: it refers to the relationship between the execution time of algorithms and input size  $n$ .**

- $\Omega(1)$ : constant time. Whatever the input size, the execution time is always the same.
- $\Omega(n)$ : linear time. When the input size increases, the execution time increases linearly as well.
- $\Omega(n^2)$ : quadratic time. When the input size increases, the execution time increases quadratically.
- $\Omega(\log n)$ : logarithmic time. When the input size increases, the execution time increases logarithmically.

**Space-complexity: it refers to the relationship between the required memory space and the size of the inputs by an algorithm during execution. The pattern is similar to the time complexity.**

**Big- $\Omega$  is often used to describe the lower bound of an algorithm, indicating that the algorithm can not be faster than  $\Omega$ . While the worst case is more practical, the best case can also serve as an additional analysis of the algorithm.**

## 1.22 Exercises

Use binary search to find the target element 13 in the following sorted array and analyze the time complexity:

[1,3,5,7,9,11,13]

## 1.23 Solution

Binary search is an efficient search algorithm for finding a target element in an ordered array. It divides the search interval repeatedly into two halves to narrow the possible range of the target element, until the target element be found or the search interval is empty. The specific step is as follows:

**Initial array and start at the middle element of the array**

- The search interval is [1,3,5,7,9,11,13], and the target element is 13

**Fisrt comparison, compare current element with target number, if smaller than target element, continue to search the left half, if bigger than target element, continue to search right half, if equal to target, target element found.**

- The middle element is 7, and the target element is 13, the target element should be in the right half[9,11,13].

**Second comparison**

- The middle element is 11, and the target element is 13, the target element should be in the right half [13].

**Third comparison**

- The middle element is 13, and the target element is 13, the target element found in three comparisons.

**Time Complexity:**

**Best Case:  $\Omega(1)$**

- The best case for Binary Search is when the target element [13] is found as the middle element of the array on the first comparison. This means that only one comparison is required to find the target element [13]. So the time complexity is  $\Omega(1)$ .

**Worst Case:  $O(\log n)$**

- The worst case is when the target element [13] is not present in the array or is located at the far front or end of the array. In such cases, the array needs to be repeatedly divided until the interval size becomes 1, which means  $\log n$  comparisons are required. So the time complexity is  $O(\log n)$ .

**Average Case:  $O(\log n)$**

- For a randomly selected target element in the array, the average number of comparisons required to find the element is proportional to  $\log n$ . This is because the array is halved with each comparison, thus exponentially reducing the search space. So the time complexity is  $O(\log n)$ .

**Space Complexity:**

#### **Iterative: $O(1)$**

- In the iterative implementation of Binary Search, no additional space is used other than a few variables to track the left, right, and middle indices. The space requirement is independent of the input size. So the space complexity is  $O(1)$ .

#### **Recursive: $O(\log n)$**

- In the recursive implementation of Binary Search, each recursive call adds a new frame to the call stack. The depth of the recursion is  $\log n$  because the array is halved at each step. So the space complexity is  $O(\log n)$ .

## **1.3 Big Theta**

### **1.31 Explanation**

**Big- $\theta$**  is used to describe the average time or space complexity of an algorithm. It represents the exact growth rating of an algorithm in all possible cases. That means it's both the upper and lower bounds of an algorithm's complexity. Big- $\theta$  common time and space complexities:

**Time-Complexity:** it refers to the relationship between the execution time of algorithms and input size  $n$ .

- $\theta(1)$ : constant time. Whatever the input size, the execution time is always the same.
- $\theta(n)$ : linear time. When the input size increases, the execution time increases linearly as well.
- $\theta(n^2)$ : quadratic time. When the input size increases, the execution time increases quadratically.
- $\theta(\log n)$ : logarithmic time. When the input size increases, the execution time increases logarithmically.

**Space-Complexity:** it refers to the relationship between the required memory space and the size of the inputs by an algorithm during execution. The pattern is similar to the time complexity

If an algorithm's time complexity is  $\theta(n)$ , it means that the execution time is proportional to  $n$  in both the best and worst case.

### **1.32 Exercises**

Use Bubble Sort to sort the following array and show the state of the array after each step:

[5,2,9,1,5,6]

### **1.33 Solution**

Bubble sort is a simple and inefficient sorting algorithm. It works by repeatedly iterating over an array, comparing adjacent elements and swapping them if they are in the wrong order. The largest element "bubbles" into its correct position at the end of the array. In short, after each swap, the largest unsorted element is moved to its correct position.

**The steps of Bubble Sort are as follows:**

- Start from the first element of an array and compare adjacent elements. If the previous one is larger than the next one, swap them.
- After each move, the last element of the unsorted part is in its correct position.
- Repeat the above step until each element is in its correct position and the array is sorted.

**Solution:**

**Initial Array:**

[5,2,9,1,5,6]

**first round comparison and swap:**

- Compare 5 and 2,  $5 > 2$ , then swap:

[2,5,9,1,5,6]

- Compare 5 and 9,  $5 < 9$ , no need swap:
- Compare 9 and 1,  $9 > 1$ , then swap:

[2,5,1,9,5,6]

- Compare 9 and 5,  $9 > 5$ , then swap:

[2,5,1,5,9,6]

- Compare 9 and 6,  $9 > 6$ , then swap:

[2,5,1,5,6,9]

**Second round comparison and swap:**

- Compare 2 and 5,  $2 < 5$ , no need swap:
- Compare 5 and 1,  $5 > 1$ , then swap:

[2,1,5,5,6,9]

- Compare 5 and 5,  $5 = 5$ , then swap:

[2,1,5,5,6,9]

- Compare 5 and 6,  $5 < 6$ , no need swap:

[2,1,5,5,6,9]

**Third round comparison and swap:**

- Compare 2 and 1,  $2 > 1$ , then swap:

[1,2,5,5,6,9]

- Compare 2 and 5,  $2 < 5$ , no need swap:
- No need for swaps for the remaining elements.

**The final sorted array is:**

[1,2,5,5,6,9]

**Time Complexity is  $\theta(n^2)$**

**Best Case:  $O(n)$** 

- The array is already sorted such as [1,2,5,5,6,9]. Bubble Sort only needs to pass through the array once without any swaps. So the best case time complexity is  $O(n)$ .

**Worst Case:  $O(n^2)$** 

- The array is in reverse order such as [9,6,5,5,2,1]. In this case, each move requires the maximum number of swaps. So the worst case time complexity is  $O(n^2)$ .

**Average Case:  $\theta(n^2)$** 

- Even though in this array [5,2,9,1,5,6], the total number of comparisons was smaller than the worst case, it still follows a quadratic growth pattern, because the quadratic nature of the algorithm remains dominant over all passes. The cumulative sum of operations will always be of the order  $\theta(n^2)$ .

**Space Complexity is  $\theta(1)$** 

- Bubble Sort is an in-place sorting algorithm, meaning that it does not require additional memory for extra variables. So the space complexity is  $\theta(1)$ .

## 2 Divide and Conquer

### 2.1 Technique Definition

#### 2.11 Explanation

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#### 2.12 Exercises

#### 2.13 Solution

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## 2.2 Divide And Conquer Multiplication

### 2.21 Explanation

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### 2.22 Exercises

### 2.23 Solution

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## 2.3 Divide And Conquer Algorithmic Design

### 2.31 Explanation

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### 2.32 Exercises

### 2.33 Solution

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## 3 Recurrences and Master Theorem

### 3.1 Recurrences

#### 3.11 Explanation

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### 3.12 Exercises

### 3.13 Solution

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## 3.2 Master Theorem

### 3.21 Explanation

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### 3.22 Exercises

### 3.23 Solution

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## 3.3 Using Master Theorem To Analyze A Recursive Algorithm

### 3.31 Explanation

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### 3.32 Exercises

### 3.33 Solution

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