# CSCE 222-501 Discrete Structures for Computing Fall 2014 – Hyunyoung Lee

### Problem Set 10

**Due dates:** Electronic submission of hw10.tex and hw10.pdf files of this homework is due on **Monday 12/8/2014 before 23:59** on csnet.cs.tamu.edu. A signed paper copy of the pdf file is due on **Tuesday 12/9/2014** at the beginning of class.

Name: Eric E. Gonzalez

Resources. Discrete Mathematics and its Applications 7th Ed.(Rosen)

On my honor, as an Aggie, I have neither given nor received any unauthorized aid on any portion of the academic work included in this assignment. Furthermore, I have disclosed all resources (people, books, web sites, etc.) that have been used to prepare this homework.

In this problem set, you will earn total 100 + 20 (extra credit) points. Each problem is worth ten points unless otherwise specified.

**Problem 1.** Chapter 13.1, Exercise 4, page 856

#### Solution.

- (a)  $S \to 1S$
- $\rightarrow 11S$
- $\rightarrow 111S$
- $\rightarrow 11100A$
- $\rightarrow 111000$
- 111000 belongs to language generated by G.
- (b) Any sentence of the language generated by G ends with 0
- 11001 does not belong to the language generated by G.
- (c)  $L(G) = 1^m 0^n / m \ge 0, n \ge 3$

Problem 2. (20 points) Chapter 13.1, Exercise 6, page 856

# Solution.

- (a)  $\{abbb\}$
- (b)  $\{aba, aa\}$
- (c)  $\{abb, abab\}$
- (d)  $\{a2^n | n \ge 2\} \cup \{b^n | n \ge 1\}$
- (e)  $\{a^n b^{n+m} a^m | m > 0, n > 0\}$

Problem 3. (15 points) Chapter 13,1, Exercise 14, page 856

## Solution.

(a) Let  $L = \{10, 01, 101\}$  and G = (V, T, S, P)

Grammar G will generate language L if

 $V=\{0,1,S\}, T=\{0,1\},$  S is the starting symbol, and the productions are  $S\to 10, S\to 01, S\to 101$ 

(b) Let L be a bit string that starts with 00 and ends with one or more 1s and G = (V, T, S, P)

Grammar G will generate language L if

 $V = \{0, 1, S, A, B\}, T = \{0, 1\}, S$  is the starting symbol, and the productions are

 $S \rightarrow 00AB, A \rightarrow AA, A \rightarrow 0, A \rightarrow 1, B \rightarrow BB, B \rightarrow 1$ 

(c) Let L be a bit string consisting of an even number of 1s followed by a final 0 and G=(V,T,S,P)

Grammar G will generate language L if

 $V = \{0,1,S,A,B,C\}, T = \{0,1\},$  S is the starting symbol, and the productions are

$$S \to A0, A \to \lambda, A \to BBC, A \to BCB, A \to CBB, C \to CC, B \to CB, B \to BC, B \to 1, C \to 0$$

(d) Let L be a bit string that has neither two consecutive 0s nor two consecutive 1s and G=(V,T,S,P)

Grammar G will generate language L if

 $V = \{0, 1, S, A, B\}, T = \{0, 1\},$  S is the starting symbol, and the productions are

$$S\to A, A\to AA, A\to A0, A\to 01, A\to \lambda, S\to B, B\to BB, B\to B1, B\to 10, B\to \lambda$$

Problem 4. (15 points) Chapter 13,1, Exercise 18, page 856

### Solution.

(a) Let  $L = \{01^{2n}/n \ge 0\}$  and G = (V, T, S, P)

Grammar G will generate language L if

 $V=\{0,1,S,A\}, T=\{0,1\},$  S is the starting symbol, and the productions are  $S\to 0A, A\to \lambda, A\to AA, A\to 11$ 

(b) Let  $L = \{0^n 1^{2n} / n \ge 0\}$  and G = (V, T, S, P)

Grammar G will generate language L if

 $V=\{1,0,S\}, T=\{0,1\},$  S is the starting symbol, and the productions are  $S\to 0S11, S\to \lambda$ 

(c) Let  $L = \{0^n 1^m 0^n / m \ge 0, n \ge 0\}$ 

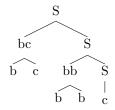
Grammar G will generate language L if

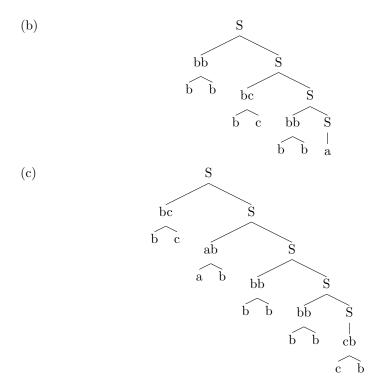
 $V=\{0,1,S,A\}, T=\{0,1\},$  S is the starting symbol, and the productions are  $S\to 0S0, S\to \lambda, S\to A, A\to AA, A\to 1$ 

**Problem 5.** Chapter 13.1, Exercise 24, page 857

## Solution.

(a)





**Problem 6.** (5 points) Chapter 13.2, Exercise 2 a), page 864 Solution.

	1	f	g		
	Inp	out	Input		
State	0	1	0	1	
$s_0$	$s_1$	$s_2$	1	0	
$s_1$	$s_0$	$s_3$	1	0	
$s_2$	$s_3$	$s_0$	0	0	
$s_3$	$s_1$	$s_2$	1	1	

**Problem 7.** (5 points) Chapter 13.2, Exercise 4 a), page 864 Solution.

Input	1	0	0	0	1	-
State	$s_0$	$s_2$	$s_3$	$s_1$	$s_0$	$s_2$
Output	0	0	1	1	0	_

Problem 8. Chapter 13.2, Exercise 18, page 865 (explain your FSM)

# Solution.

The machine will need 5 states:  $s_1,\ s_2,\ s_3,\ s_4,\ s_5,$  aside from the starting

state  $s_0$ , to track the last five digits input. The numbers of the state indicate the number of consecutive 1s that have been seen. The machine will output 1 if the input meets the condition, otherwise it will output 0.

**Problem 9.** Chapter 13.3, Exercise 8 a), b), e) and f), page 875

### Solution.

(a) Let 
$$A = \{v\}$$
 and  $A^2 = \{vv\}$   
  $A \notin A^2$ 

Proven by counterexample that the statement is false.

(b) Let  $A = \emptyset$ . The hypothesis is true, as an empty set squared results in the empty set. However, the conclusion is false, as  $\lambda \notin A$ 

Proven by counterexample that the statement is false.

(e) If A is not the empty set, the left hand side of the expression will be  $A^i$ , i > 0. The right hand side will contain  $A^0$ , which is the empty string  $\lambda$ , so the equality does not hold.

Proven by counterexample that the statement is false.

(f) Let 
$$A = \{1, \lambda\}$$
 and  $A^2 = \{1, 11, \lambda\}$ 

Hence, 
$$|A^2| = 3$$
 and  $|A|^2 = 4$ 

Proving that the equality is not true for  $A = \{1, \lambda\}, n=2$ 

Thus, the equation is false.

Problem 10. (5 points) Chapter 13.3, Exercise 10 b), d) and f), page 875

## Solution.

(b) 
$$\{0\} * \{10\}\{1\} *$$

The set allows consecutive 0s only in the beginning of the string and only a single 0 after the first 1 occurs.

01001 is not present in  $\{0\} * \{10\}\{1\} *$ .

The string 01001 is equivalent to  $\{010\}\{01\}$ , which cane be formed by taking 010 from the first set and 01 from the second set of the given string.

01001 is present in the given set.

(f) 
$$\{01\} * \{01\} *$$

in the set, we cannot have the continuation of two 0s, so 01001 is not present. 01001 is not present in the given set.

**Problem 11.** (5 points) Chapter 13.3, Exercise 16, page 876

**Solution.** 
$$\{\lambda\} \cup \{1\}\{0,1\} * \cup \{0\}\{1\} * \{0\}\{0,1\} *$$

Problem 12. (5 points) Chapter 13.3, Exercise 18, page 876

Solution.  $\{\lambda\} \cup \{0\}\{1\}*$ 

Problem 13. (5 points) Chapter 13.3, Exercise 28, page 876

**Solution.** Have four states, with only  $q_3$  final. For i = 0, 1, 2, 3, transition from si to itself on input 1 and to si+1 on input 0. Both transitions from  $q_3$  are to itself.

# Checklist:

□ Did you add your name?
□ Did you disclose all resources that you have used?
(This includes all people, books, websites, etc. that you have consulted)
□ Did you sign that you followed the Aggie honor code?
□ Did you solve all problems?
□ Did you submit (a) your latex source file and (b) the resulting pdf file of your homework on csnet?
□ Did you submit (c) a signed hardcopy of the pdf file in class?