

CSCE 222-501 Discrete Structures for Computing
Fall 2014 – Hyunyoung Lee

Problem Set 6

Due dates: Electronic submission of hw6.tex and hw6.pdf files of this homework is due on **10/27/2014 before 23:59** on csnet.cs.tamu.edu. Please do not archive or compress the files. A signed paper copy of the pdf file is due on **10/28/2014** at the beginning of class.

Name: Eric E. Gonzalez

Resources. Discrete Mathematics and its Applications 7th Ed.(Rosen)

On my honor, as an Aggie, I have neither given nor received any unauthorized aid on any portion of the academic work included in this assignment. Furthermore, I have disclosed all resources (people, books, web sites, etc.) that have been used to prepare this homework.

Signature: _____

Problem 1. Section 5.3, Exercise 4, page 357

Solution.

$$f(0) = f(1) = 1$$

$$(a) f(n+1) = f(n) - f(n-1)$$

$$f(2) = f(1) - f(0)$$

$$= 1 - 1$$

$$= 0$$

$$f(3) = f(2) - f(1)$$

$$= 0 - 1$$

$$= -1$$

$$f(4) = f(3) - f(2)$$

$$= -1 - 0$$

$$= -1$$

$$f(5) = f(4) - f(3)$$

$$= -1 - (-1)$$

$$= 0$$

$$(b) f(n+1) = f(n)f(n-1)$$

$$f(2) = f(1)f(0)$$

$$= 1 * 1$$

$$= 1$$

$$f(3) = f(2)f(1)$$

$$= 1 * 1$$

$$= 1$$

$$f(4) = f(3)f(2)$$

$$= 1 * 1$$

$$\begin{aligned}
&= 1 \\
f(5) &= f(4)f(3) \\
&= 1 * 1 \\
&= 1
\end{aligned}$$

$$\begin{aligned}
\text{(c) } f(n+1) &= f(n)^2 + f(n-1)^3 \\
f(2) &= f(1)^2 + f(0)^3 \\
&= 1^2 + 1^3 \\
&= 1 + 1 \\
&= 2 \\
f(3) &= f(2)^2 + f(1)^3 \\
&= 2^2 + 1^3 \\
&= 4 + 1 \\
&= 5 \\
f(4) &= f(3)^2 + f(2)^3 \\
&= 5^2 + 2^3 \\
&= 25 + 8 \\
&= 33 \\
f(5) &= f(4)^2 + f(3)^3 \\
&= 33^2 + 5^3 \\
&= 1089 + 125 \\
&= 1214
\end{aligned}$$

$$\begin{aligned}
\text{(d)} \\
f(n+1) &= f(n)/f(n-1) \\
f(2) &= f(1)/f(0) \\
&= 1/1 = 1 \\
f(3) &= f(2)/f(1) \\
&= 1/1 = 1 \\
f(4) &= f(3)/f(2) \\
&= 1/1 = 1 \\
f(5) &= f(4)/f(3) \\
&= 1/1 = 1
\end{aligned}$$

Problem 2. Section 5.3, Exercise 6, page 357

Solution.

(a) The function is well-defined

$$f(n) = -f(n-1) \text{ for } n \geq 1.$$

$$\text{Basis step: } f(0) = 1 = (-1)^0$$

$$\text{Induction step: } f(n) = -f(n-1)$$

$$= -(-1)^{n-1}$$

$$= (-1)^n$$

Therefore, claim proven by induction.

(b) The function is well-defined

$$2^m \text{ if } n = 3k,$$

$$0 \text{ if } n = 3k + 1,$$

and 2^{m+1} if $n = 3k + 2$

for some integer k .

Basis step:

$$f(0) = 1 = 2^0$$

$$f(1) = 0$$

$$f(2) = 2 = 2^{0+1}$$

Induction step: $f(n) = 2f(n-3)$ for $n \geq 3$

Therefore, $f(n) = 2^k$, or, $= 0$, or, $= 2^{k+1}$

The claim is proven by strong induction.

c) Function is invalid because no base case can be reached.

d) Function is not well-defined because $f(1)$ and $2f(1-1)$ equal different values.

e) The function is defined by $f(n) = 2^{\lfloor n/2 \rfloor + 1}$

Basis Step:

$$f(0) = 2 = 2^1 = 2^{\lfloor 0/2 \rfloor + 1}$$

$$f(1) = f(0) = 2 = 2^1 = 2^{\lfloor 1/2 \rfloor + 1}$$

Induction Step:

$$f(n) = 2f(n-2)$$

$$= 2 * 2^{\lfloor (n-2)/2 \rfloor + 1}$$

$$= 2 * 2^{\lfloor n/2 \rfloor} = 2^{\lfloor n/2 \rfloor + 1}$$

Thus, the claim is proven by induction.

Problem 3. Section 5.3, Exercise 8, page 358

Solution.

a) $a_1 = 2$ and $a_n = 4 + a_{n-1}$ for $n \geq 2$.

b) $a_1 = 0$, $a_2 = 2$, and $a_n = a_{n-2}$ for $n \geq 3$.

c) $a_1 = 2$ and $a_n = a_{n-1} + 2n$ for $n \geq 2$.

d) $a_1 = 1$ and $a_n = a_{n-1} + 2n - 1$ for $n \geq 2$.

Problem 4. Section 5.3, Exercise 12, page 358

Solution.

Basis Step: $P(1) = f_1^2 = f_1 f_1 = 1 * 1 = 1$

Induction Step:

$$P(k+1) = f_1^2 + f_2^2 + \dots + f_k^2 + f_{k+1}^2 = f_k f_k + 1 + f_{k+1}^2$$

$$= f_k f_{k+1} + f_{k+1} f_{k+1}$$

$$= f_{k+1} [f_k + f_{k+1}]$$

$$= f_{k+1} f_{k+2}$$

Therefore, true by induction.

Problem 5. Section 5.3, Exercise 14, page 358

Solution.

Basis Step: $P(1) = f_{1+1} f_{1-1} - f_1^2 = f_2 f_0 - f_1 f_1 = 1 * 0 - 1 * 1 = (-1)^1$

Induction Step:

$$P(k+1) = f_{(k+1)+1} f_{(k+1)-1} - f_{k+1}^2$$

$$= f_{k+2} f_k - f_{k+1}^2$$

$$= f_{k+2} f_k - f_{k+1} f_{k+1}$$

$$\begin{aligned}
&= (f_{k+1} + f_k)f_k - f_{k+1}fk + 1 \\
&= f_{k+1}f_k + f_k^2 - f_{k+1}[f_k + f_{k-1}] \\
&= f_{k+1}f_k + f_k^2 - f_{k+1}f_k - f_{k+1}f_{k-1} \\
&= f_k^2 - f_{k+1}f_{k-1} \\
&= -(-1)^k \\
&= -1(-1)^k \\
&= (-1)^{k+1}
\end{aligned}$$

Therefore, proven by induction.

Problem 6. Section 5.3, Exercise 16, page 358

Solution.

Basis Step: $f_0 - f_1 + f_2 = 0 - 1 + 1 = 1 - 1 = f_1 - 1$

Induction Step:

$$\begin{aligned}
&\sum_{k=0}^{2n+2} (-1)^k f_k \\
&= f_{2n-1} - 1 - f_{2n+1} + f_{2n+2} \\
&= f_{2n-1} + f_{2n} - 1 \\
&= f_{2n+1} - 1 \\
&= f_{2(n+1)-1} - 1
\end{aligned}$$

Therefore, claim proven by induction.

Problem 7. Section 5.3, Exercise 20, page 358

Solution.

max:

$$\begin{aligned}
\max(a_i, a_{i+1}) &= a_i \text{ if } a_i \geq a_{i+1} \text{ for } i = 0, 1, 2, \dots, n-1 \\
\max(a_1, a_2, \dots, a_n) &= \max(\max(a_1, a_2, \dots, a_{n-1}), a_n)
\end{aligned}$$

min:

$$\begin{aligned}
\min(a_i, a_{i+1}) &= a_i \text{ if } a_i \leq a_{i+1} \text{ for } i = 0, 1, 2, \dots, n-1 \\
\min(a_1, a_2, \dots, a_n) &= \min(\min(a_1, a_2, \dots, a_{n-1}), a_n)
\end{aligned}$$

Problem 8. Section 5.3, Exercise 26, page 358

Solution.

(a)

1st Application: (2, 3), (3, 2)

2nd Application: (4, 6), (5, 5), (6, 4)

3rd Application: (6, 9), (7, 8), (8, 7), (9, 6)

4th Application: (8, 12), (9, 11), (10, 10), (11, 9), (12, 8)

5th Application: (10, 15), (11, 14), (12, 13), (13, 12), (14, 11), (15, 11)

(b)

Basis Step:

$5|a + b$ with (0, 0)

$5|0 + 0$, so basis holds

Inductive Hypothesis: Assume $5|a + b$

Recursive Step:

True by strong induction. Cause I ran out of time.

(c)

Basis Step:

$5|a + b$ when $a=0$ and $b=0$

$5|0$, so basis holds

Inductive Hypothesis: Assume $5|a + b$ for $(a, b) \in S$

Recursive Step:

Show that $(a+2, b+3)$ and $(a+3, b+2)$ are divisible by 5

$5 \nmid a+b$ is true by I.H.

$5 \nmid a+2+b+3 = 5 \nmid a+b+5$

$5 \nmid 5$

Thus, we can prove the claim by structural induction.

Problem 9. Section 5.3, Exercise 36, page 359

Solution.

Let $w_1 = ab$

Let $w_2 = cd$

$w_1 w_2 = abcd$

$(w_1 w_2)^R = dcba$

Recursive Step:

$w_1^R = ba$

$w_2^R = dc$

$w_2^R w_1^R = dcba$

$= (w_1 w_2)^R$

Problem 10. Section 5.3, Exercise 44, page 359

Solution.

Let T be a binary tree with a single root node.

$l(T) = 1, i(T) = 0$

$l(T) = l(T_1) + l(T_2)$

$= (i(T_1) + 1) + (i(T_2) + 1)$

As such, for any binary tree T , $l(T) = i(T) + 1$

The claim is proven by structural induction.

Checklist:

- ☐ Did you add your name?
- ☐ Did you disclose all resources that you have used?
(This includes all people, books, websites, etc. that you have consulted)
- ☐ Did you sign that you followed the Aggie honor code?
- ☐ Did you solve all problems?
- ☐ Did you submit (a) your latex source file and (b) the resulting pdf file of your homework on csnet?
- ☐ Did you submit (c) a signed hardcopy of the pdf file in class?