Improving Automated Feedback Building a Rule Feedback Generator

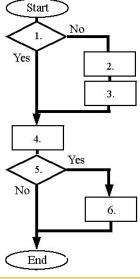
Eric Bouwers

September 27, 2007

(Generic) Outline

- 1 The problem
- Our solution
- 3 Evaluation
- **4** Conclusion

Procedural skills



- 1. Are the denominators the same? If yes, skip to step 4.
- 2. Find the lowest common denominator.
- 3. Convert both fractions to the common denominator.
- 4. Add the numerators, and place the sum over the common denominator.
- 5. Can the answer be simplified? (If no, you are finished.)
- 6. Simplify the answer.

Universiteit Utrecht

Feedback





Feedback in Education

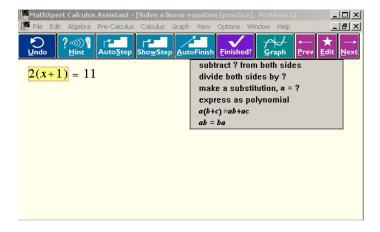


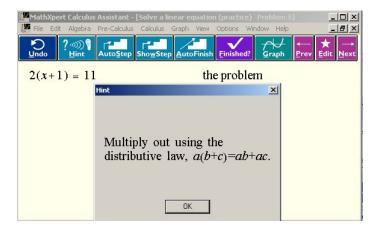


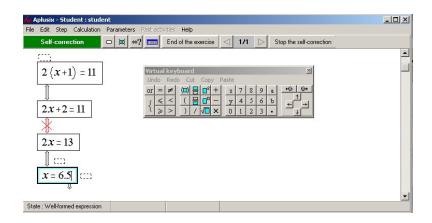
Feedback in Education

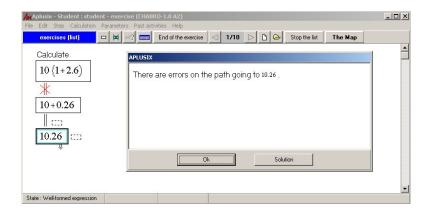


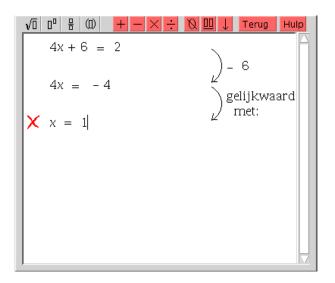


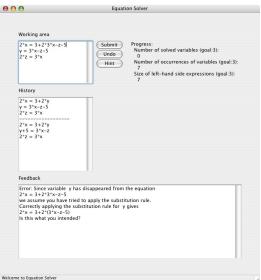


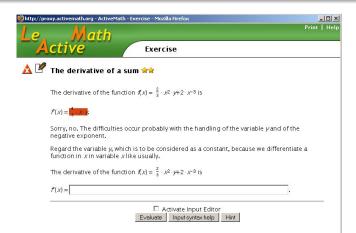


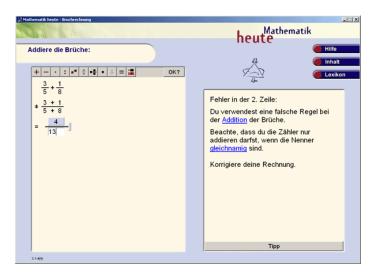












Synopsis

- Feedback is important in learning
- Personal feedback is not feasible without (computerized) help
- Current solutions are either:
 - Only Correct/Incorrect solutions
 - Result of extensive research

Goal

Research question

How can we make the generation of high-quality, domain-specific feedback easier?

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How can we make the generation of high-quality, domain-specific feedback easier?

Basic idea

Create a generic framework which separates the knowledge of rule-feedback generation from knowledge about the domain.

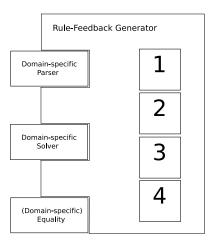
- Should always be able to produce basic feedback

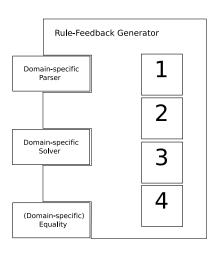
- •

- Should always be able to produce basic feedback
- More detailed/complete input leads to better feedback
- 0
- •

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- Can be instantiated on different domains (with little effort)

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- More detailed/complete input leads to better feedback
- Can be instantiated on different domains (with little effort)
- Adaptable to a single class-room (or student!)



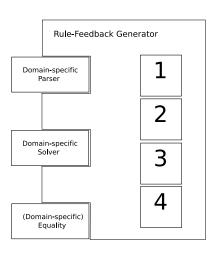


Phase 1:

Input: CT + PT

Output:

 ${\sf Correct/Incorrect\ message}$

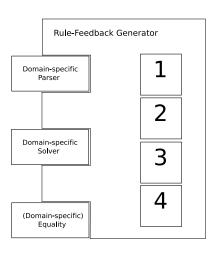


Phase 2:

Input: CT + PT + Student rule

Output:

 ${\sf Correct/Incorrect\ message-tuple}$

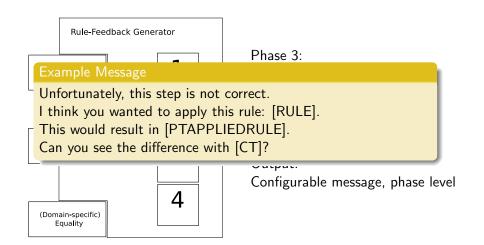


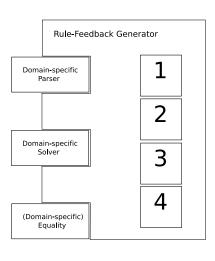
Phase 3:

Input: CT + PT + Allowed rules

Output:

Configurable message, phase level





Phase 4:

Input:

 $\mathsf{CT} + \mathsf{PT}$

+ Allowed rules + Buggy rules

Output:

Configurable message, rule level

Input: PT, CT

Output: Correct or Incorrect

Algorithm: Solve both terms and check whether their results are

semantically equal.

Input: PT, CT

Output: Correct or Incorrect

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semantically equal.

Implementation

firstPhase :: RFG a => RFGSolver a -> RFGEqual a -> a -> a -> String

firstPhase solve equal pt ct =

let resultCt = solve ct

resultPt = solve pt

in if equal resultPt resultCt

then getConfig Correct

else getConfig Incorrect

Input: PT, CT

Output: Correct or Incorrect

Algorithm: Solve both terms and check whether their results are

semantically equal.

Implementation

resultPt = solve pt

in if equal resultPt resultCt
 then getConfig Correct

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Implementation

```
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firstPhase solve equal pt ct =

let resultCt = solve ct
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Input: PT, CT and allowed rules

Output: Configurable message

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Output: Configurable message

Algorithm

Given two terms, calculate the rewrite rule δ between these terms.

Determine which allowed rule is closest to δ .

Apply rules from the set to the PT if possible and recurse.

The result is the rule which is a closest match.

Input: PT, CT and allowed rules

Output: Configurable message

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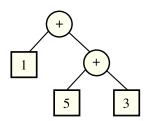
Calculating rewrite rules

$$PT = 1 + 5 + 3$$

 $CT = 3 + 6$

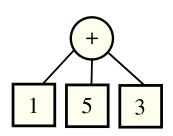
$$PT = 1 + 5 + 3$$

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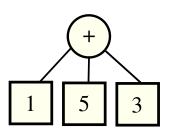
$$PT = 1 + 5 + 3$$

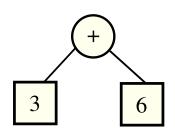
 $CT = 3 + 6$



$$PT = 1 + 5 + 3$$

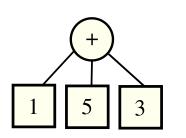
 $CT = 3 + 6$

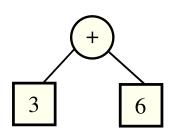




$$PT = 1 + 5 + 3$$

 $CT = 3 + 6$





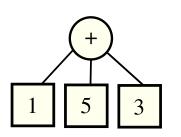
Non-Equal

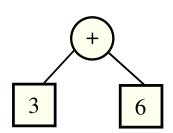
 $([1,5,3],[3,6,_])$



$$PT = 1 + 5 + 3$$

 $CT = 3 + 6$





Equal, Associative

([1,5,3],[3,6,_]), ([1,5,3],[_,3,6])

$$PT = 1 + 5 + 3$$

 $CT = 3 + 6$

```
Equal, Associative and Commutative

([1,5,3],[3,6,-]), ([1,5,3],[-,3,6]), ([1,5,3],[6,3,-]),

([1,5,3],[-,6,3]), ([1,3,5],[3,6,-]), ([1,3,5],[-,3,6]),

([1,3,5],[6,3,-]), ([1,3,5],[-,6,3]), ([3,1,5],[3,6,-]),

1 ([3,1,5],[-,3,6]), ([3,1,5],[6,3,-]), ([3,1,5],[-,6,3]),

([3,5,1],[3,6,-]), ([3,5,1],[-,3,6]), ([3,5,1],[6,3,-]),

([3,5,1],[-,6,3]), ([5,3,1],[3,6,-]), ([5,3,1],[-,3,6]),

([5,3,1],[6,3,-]), ([5,3,1],[-,6,3]), ([5,1,3],[3,6,-]),

([5,1,3],[-,3,6]), ([5,1,3],[6,3,-]), ([5,1,3],[-,6,3])
```

6

$$PT = 1 + 5 + 3$$

 $CT = 3 + 6$

```
Equal, Associative and Commutative

([1,5,3],[3,6,_]), ([1,5,3],[-,3,6]), ([1,5,3],[6,3,_]),

([1,5,3],[-,6,3]), ([1,3,5],[3,6,_]), ([1,3,5],[-,3,6]),

([1,3,5],[6,3,_]), ([1,3,5],[-,6,3]), ([3,1,5],[-,6,3]),

([3,5,1],[3,6,_]), ([3,5,1],[-,3,6]), ([3,5,1],[6,3,_]),

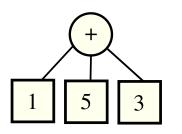
([3,5,1],[-,6,3]), ([5,3,1],[3,6,_]), ([5,3,1],[-,3,6]),

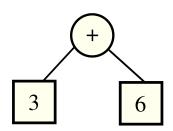
([5,3,1],[6,3,_]), ([5,3,1],[-,6,3]), ([5,1,3],[3,6,_]),

([5,1,3],[-,3,6]), ([5,1,3],[6,3,_]), ([5,1,3],[-,6,3])
```

$$PT = 1 + 5 + 3$$

 $CT = 3 + 6$



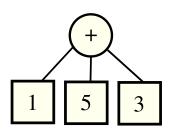


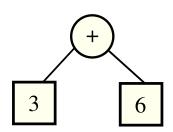
Result:

([1,5],[6])

$$PT = 1 + 5 + 3$$

 $CT = 3 + 6$



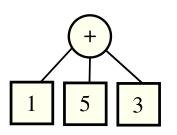


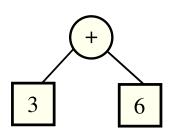
Result

$$\delta = 1 + 5 \Rightarrow +6$$

$$PT = 1 + 5 + 3$$

 $CT = 3 + 6$

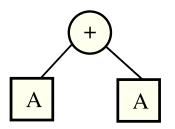


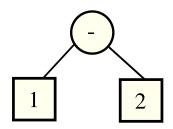


Result.

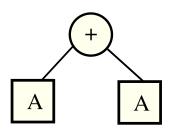
$$\delta = 1 + 5 \Rightarrow 6$$

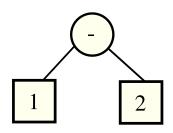
Distance between rules:





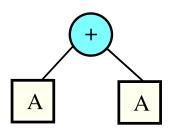
Distance between terms:

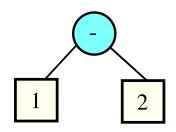




Environment $= \{\}$ Distance = 0

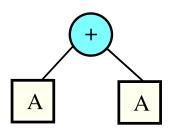
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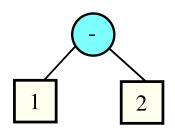




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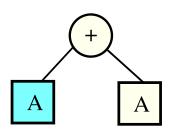
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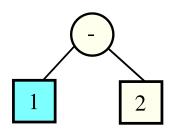




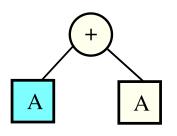
Environment $= \{\}$ Distance = 2

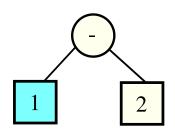
Distance between terms:





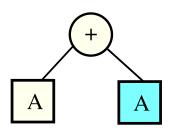
Environment = {}
Distance = 2

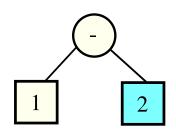




Environment =
$$\{(A, 1)\}$$

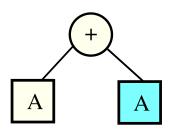
Distance = 2

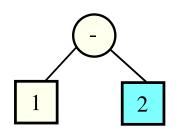




Environment =
$$\{(A, 1)\}$$

Distance = 2





Environment =
$$\{(A, 1)\}$$

Distance = 4

$$\mathsf{PT} = (\mathsf{a} \lor \mathsf{b}) \to \mathsf{c}$$
$$\mathsf{CT} = (\neg \mathsf{a} \land \mathsf{b}) \lor \mathsf{c}$$

```
Rules = {
    ImpElimination :
    A \rightarrow B \Rightarrow \neg(A) \lor B
, MorganOr :
    \neg(A \lor B) \Rightarrow \neg(A) \land \neg(B)
}
```

$$\mathsf{PT} = (a \lor b) \to c$$
 $\mathsf{CT} = (\neg a \land b) \lor c$

$$\label{eq:Rules} \begin{split} & \text{Rules} = \{ \\ & \text{ImpElimination}: \\ & A \to B \Rightarrow \neg(A) \lor B \\ & \text{, MorganOr}: \\ & \neg(A \lor B) \Rightarrow \neg(A) \land \neg(B) \\ \} \end{split}$$

Calculated Rule

$$\delta = (a \lor b) \to c \Rightarrow (\neg a \land b) \lor c$$

$$\mathsf{PT} = (a \lor b) \to c$$
 $\mathsf{CT} = (\neg a \land b) \lor c$

Rules = {
 ImpElimination :
$$A \rightarrow B \Rightarrow \neg(A) \lor B$$
 , MorganOr :
$$\neg(A \lor B) \Rightarrow \neg(A) \land \neg(B)$$
 }

Calculated Rule

$$\delta = (a \lor b) \to c \Rightarrow (\neg a \land b) \lor c$$

Distances

$$dist(\delta, ImpElimination) = 9$$

 $dist(\delta, MorganOr) = 15$

$$PT = (a \lor b) \to c$$

$$CT = (\neg a \land b) \lor c$$

Rules = {
 ImpElimination :

$$A \rightarrow B \Rightarrow \neg(A) \lor B$$
, MorganOr :
 $\neg(A \lor B) \Rightarrow \neg(A) \land \neg(B)$
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$$PT' = \neg(a \lor b) \lor c$$
$$CT = (\neg a \land b) \lor c$$

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$$\mathsf{PT'} = \neg(a \lor b) \lor c$$

$$\mathsf{CT} = (\neg a \land b) \lor c$$

Rules = {
 ImpElimination :

$$A \rightarrow B \Rightarrow \neg(A) \lor B$$
, MorganOr :
 $\neg(A \lor B) \Rightarrow \neg(A) \land \neg(B)$
}

Calculated Rule

$$\delta = \neg(a \lor b) \Rightarrow (\neg a \land b)$$

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Rules = {
 ImpElimination :

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Calculated Rule

$$\delta = \neg(a \lor b) \Rightarrow (\neg a \land b)$$

Distances

$$dist(\delta, ImpElimination) = 9$$

 $dist(\delta, MorganOr) = 4$

$$\mathsf{PT'} = \neg(a \lor b) \lor c$$

$$\mathsf{CT} = (\neg a \land b) \lor c$$

Rules = {
 ImpElimination :

$$A \rightarrow B \Rightarrow \neg(A) \lor B$$

 , MorganOr :
 $\neg(A \lor B) \Rightarrow \neg(A) \land \neg(B)$ }

Calculated Rule

$$\delta = \neg(a \lor b) \Rightarrow (\neg a \land b)$$

Distances

$$dist(\delta, ImpElimination) = 9$$

 $dist(\delta, MorganOr) = 4$

$$PT'' = (\neg a \land \neg b) \lor c$$
$$CT = (\neg a \land b) \lor c$$

```
Rules = { ImpElimination : A \rightarrow B \Rightarrow \neg(A) \lor B , MorganOr : \neg(A \lor B) \Rightarrow \neg(A) \land \neg(B) }
```

$$PT'' = (\neg a \land \neg b) \lor c$$
$$CT = (\neg a \land b) \lor c$$

$$\label{eq:Rules} \begin{split} & \text{Rules} = \{ \\ & \text{ImpElimination}: \\ & A \rightarrow B \Rightarrow \neg(A) \vee B \\ & \text{, MorganOr}: \\ & \neg(A \vee B) \Rightarrow \neg(A) \wedge \neg(B) \\ \} \end{split}$$

Calculated Rule

$$\delta = \neg b \Rightarrow b$$

$$\mathsf{PT}" = (\neg a \land \neg b) \lor c$$

$$\mathsf{CT} = (\neg a \land b) \lor c$$

$$\label{eq:Rules} \begin{split} & \text{Rules} = \{ \\ & \text{ImpElimination}: \\ & A \to B \Rightarrow \neg(A) \lor B \\ & \text{, MorganOr}: \\ & \neg(A \lor B) \Rightarrow \neg(A) \land \neg(B) \\ \} \end{split}$$

Calculated Rule

$$\delta = \neg b \Rightarrow b$$

Distances

$$dist(\delta, ImpElimination) = 9$$

 $dist(\delta, MorganOr) = 10$

Fourth phase

M. Hennecke

Online Diagnose in intelligenten mathematischen Lehr-Lern-Systemen.

PhD thesis, Hildesheim University, 1999.

Fourth phase

M Hennecke

Online Diagnose in intelligenten mathematischen Lehr-Lern-Systemen.

PhD thesis, Hildesheim University, 1999.

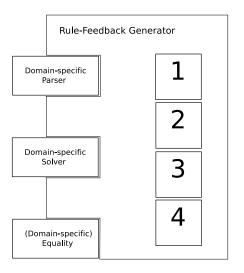
Example

$$\frac{A}{B} + \frac{C}{D} \Rightarrow \frac{A+C}{B+D}$$

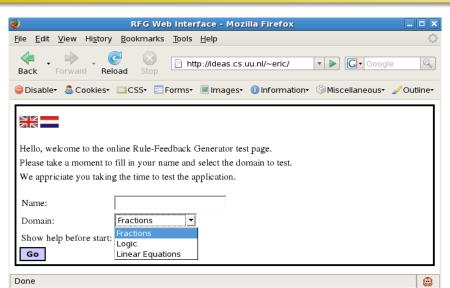
: AddError, You can only add fractions with equal denominators

: 0.3

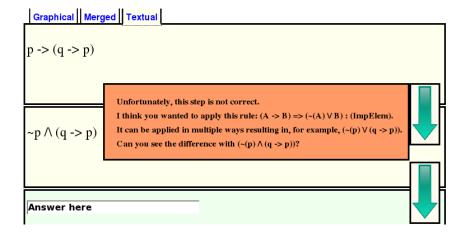
Overview



Receiving feedback



Feedback message



Wrong Guesses

```
EqElem : A \leftrightarrow B \rightarrow (A \land B) \lor (\neg(A) \land \neg(B))
```

 $ImpElem \quad : \ A \to B \qquad \quad \to \quad \neg(A) \ \lor \ B$

 $\mathsf{MorganA} \quad : \ \neg(\mathsf{A} \, \wedge \, \mathsf{B}) \qquad \rightarrow \quad \neg(\mathsf{A}) \, \vee \, \neg(\mathsf{B})$

 $\mathsf{NotNot} \quad : \, \neg \neg \mathsf{A} \qquad \longrightarrow \quad \mathsf{A}$

DistAO : $A \land (B \lor C) \rightarrow (A \land B) \lor (A \land C)$

 $\mathsf{Idem}\mathsf{A} \quad : \; \mathsf{A} \wedge \mathsf{A} \qquad \to \; \mathsf{A}$

Wrong Guesses

EqElem: $A \leftrightarrow B$ \rightarrow $(A \land B) \lor (\neg(A) \land \neg(B))$ ImpElem: $A \rightarrow B$ \rightarrow $\neg(A) \lor B$ MorganA: $\neg(A \land B)$ \rightarrow $\neg(A) \lor \neg(B)$ NotNot: $\neg\neg A$ \rightarrow ADistAO: $A \land (B \lor C)$ \rightarrow $(A \land B) \lor (A \land C)$

IdemA : $A \wedge A \rightarrow A$

PT	CT	Expected rule	Actual output
$a \rightarrow (b \land c)$	a∨(b∧c)	ImpElem	NotNot
¬¬a↔b	$(\neg a \land b) \lor (\neg \neg \neg a \lor \neg b)$	EqElem	MorganA
a∨c⇔b	$(a \lor c \land b) \lor \neg (a \lor c) \land \neg b$	EqElem	IdemA
$a \leftrightarrow b \land (c \lor d)$	$(a \leftrightarrow b \land c) \lor (a \leftrightarrow b \land d)$	EqElem	DistAO

Definition of Rules

Multiplication : $(A/B) * (C/D) \Rightarrow (A*C) / (B*D)$ Division : $(A/B) / (C/D) \Rightarrow (A/B) * (D/C)$

SameMixed : A|B/B \Rightarrow A+1

AddFractions : $(A/B) + (C/B) \Rightarrow (A+C)/B$ SubFractions : $(A/B) - (C/B) \Rightarrow (A-C)/B$

Definition of Rules

Multiplication : $(A/B) * (C/D) \Rightarrow (A*C) / (B*D)$ Division : $(A/B) / (C/D) \Rightarrow (A/B) * (D/C)$

SameMixed : A|B/B \Rightarrow A+1

AddFractions : $(A/B) + (C/B) \Rightarrow (A+C) / B$ SubFractions : $(A/B) - (C/B) \Rightarrow (A-C) / B$

SolveMin : $A - B \Rightarrow C$ where C := solve(A - B)SolveAdd : $A + B \Rightarrow C$ where C := solve(A + B)SolveMul : $A * B \Rightarrow C$ where C := solve(A * B)

Conclusion

- Partial/Configurable approach is possible
- Adding an extra domain takes little effort
- Combining techniques leads to surprisingly good results
- Already useful, both for research as well as practical

Resources

Thesis page:

http://www.cs.uu.nl/wiki/Students/EricBouwersThesisPage

Prototype:

http://ideas.cs.uu.nl/~eric/

Contact me:

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