# 模:

模(Modulus)是一個建立在除法之上的二元運算·考慮整數 A、B、C·我們聲稱:

A mod B = C 當且僅當 存在整數 n 使得 A=nB+C 且 0≦C<B

在模運算之上我們定義等價關係"同餘":

A≡B (mod C) 當且僅當 A mod C = B mod C

#### 基本性質:

- 1.  $A \equiv B, B \equiv C \rightarrow A \equiv C$
- 2.  $A \equiv B \rightarrow B \equiv A$
- 3. A≡A
- 4.  $A \equiv B, C \equiv D \rightarrow A + C \equiv B + D$
- 5.  $A \equiv B, C \equiv D \rightarrow A-C \equiv B-D$
- 6.  $A \equiv B, C \equiv D \rightarrow AC \equiv BD$
- 7.  $nA \equiv nB \pmod{nm} \rightarrow A \equiv B \pmod{m}$
- 8.  $nA \equiv nB \pmod{m}$ , n is relative prime with  $m \rightarrow A \equiv B \pmod{m}$
- 9. Consider P a polynomial,  $A \equiv B \rightarrow P(A) \equiv P(B)$

#### 定理:

費馬小定理: (Fermat's Little Theorem)

 $a^p \equiv a \pmod{p}$ , for all integers a and all primes p

#### 定理二:

 $gcd(a, b) \times lcm(a, b) = a \times b$ , for all integers a, b

# 中國剩餘定理: (Chinese Remainder Theorem)

```
Consider Sequence <A_n>, < m_n>, < M_n> where M_k=(\prod_{i=1}^n m_i)/m_k. Assume that x\equiv A_k\pmod{m_k}, for all k in [1,n]: Then x is a solution if and only if the following equivalence holds. x\equiv \sum_{k=1}^n A_k M_k t_k\pmod{\prod_{k=1}^n m_k}, where t_k M_k\equiv 1\pmod{m_k} for all k
```

# 輾轉相除法:(Euclidean Algorithm)

```
template < typename type >
type GCD(type left, type right) {
    if(!left && !right) throw logic_error("Return value does not exist.\n");
    left = abs(left);
    right = abs(right);
    return left?GCD(right%left,left):right;
}
```

# 最小公倍數:透過定理二求出

```
template<typename type>
type LCM(type left, type right){
    if(!left && !right) return 0;
    return abs(left*right/GCD(left,right));
}
```

# 模逆元:

對所有整數 A. 我們稱 B 為 A 在模 m 下之模逆元·當且僅當 AB≡1 (mod m). 此地·A 在模 m 下之模逆元存在當且僅當 A、m 互質。

#### 模逆元的同餘性:

考慮 $A \times r \equiv 1 \pmod{M}$  · 則所有 A 在模 M 下的模逆元構成如下集合:  $\{r+nM \mid n \in Z\}, \text{ where } Z \text{ stands for the set of all integers }$ 

考慮整數  $A \times B$  以及方程式 Ax+By=GCD(A,B) , 其解空間如下:

 $\{(x,y)|x=\alpha+A\times n,y=\beta-B\times n\}, where \ (\alpha,\beta) \ is \ of \ one \ solution$ 並且,我們可以藉由擴展歐基里德演算法算出其中一組解。

# 擴展歐基里德算法: (Extended Euclidean Algorithm)

```
template<typename type>
pair<type,type> extGCD(type left, type right){
    if(!left && !right) throw logic_error("Return value does not exist.\n");
    pair<type,type>lCor(1,0), rCor(0,1);
    if(left<0) lCor.first = -1;
    if(right<0) rCor.second = -1;
    left = abs(left);
    right = abs(right);
    while(left){
         swap(lCor,rCor);
         swap(left,right);
         lCor.first -= left/right*rCor.first;
         lCor.second -= left/right*rCor.second;
         left %= right;
    }
    return rCor;
}
最小正模逆元之求取:
template<typename type>
type modulusInverse(type input, type mod){
    type output = extGCD(input,mod).first;
    if(output<0) output = output%mod+mod;</pre>
    return output;
}
```