

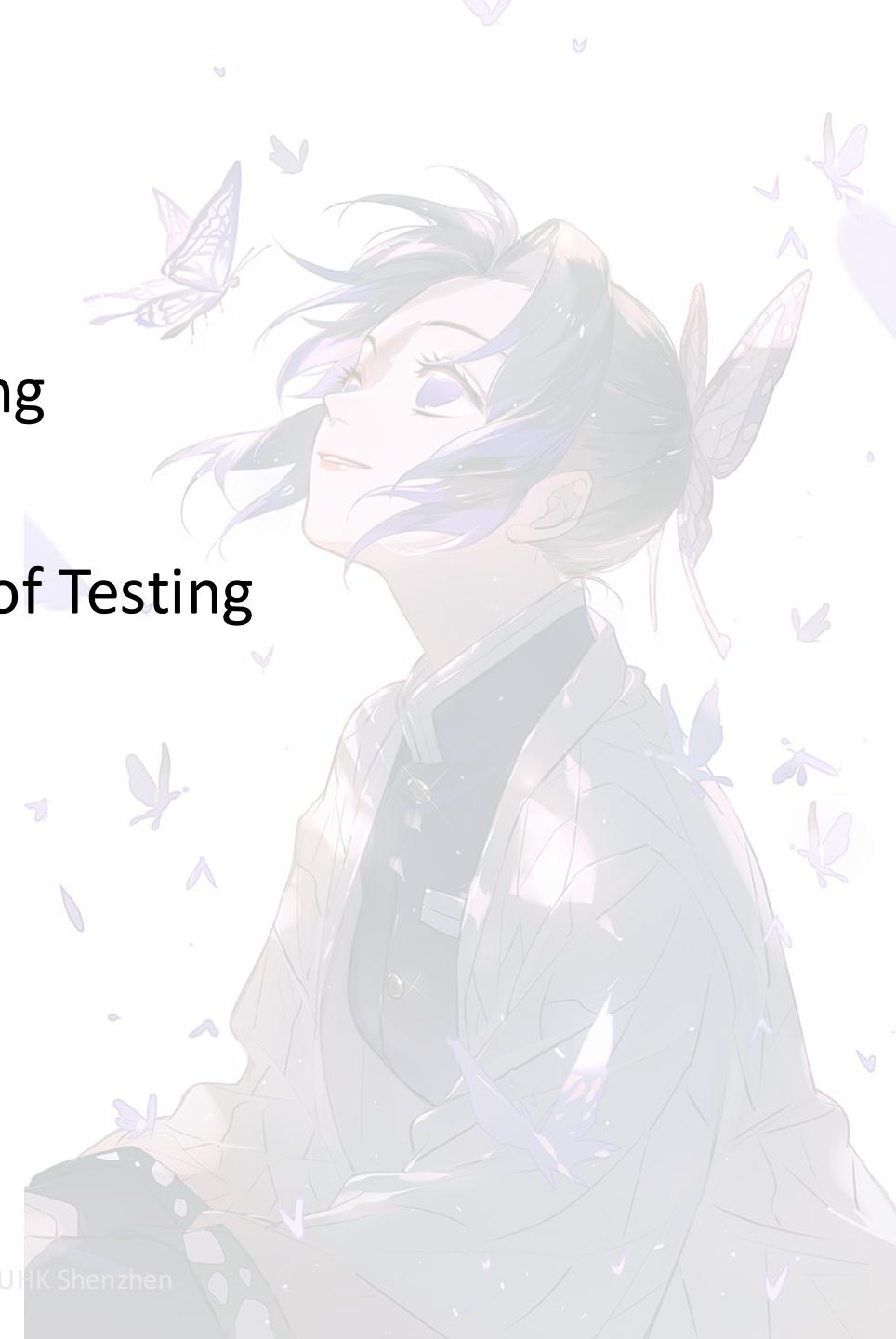
Software Engineering

Testing Basics

CUHK Shenzhen

Pinjia He

1. What is Software Testing
2. Why Software Testing
3. Stages and Categories of Testing
4. Coverage Metrics



What is Software Testing

Software testing is the process of evaluating and verifying that a software product or application does what it is supposed to do.

— IBM

<https://www.ibm.com/topics/software-testing>



说人话

Find as many **bugs** as possible.

The Patriot Accident

- The Patriot missile air defense system tracks and intercepts incoming missiles
- On February 25, 1991, a Patriot system ignored an incoming Scud missile
- 28 soldiers died; 98 were injured



History of Software Testing

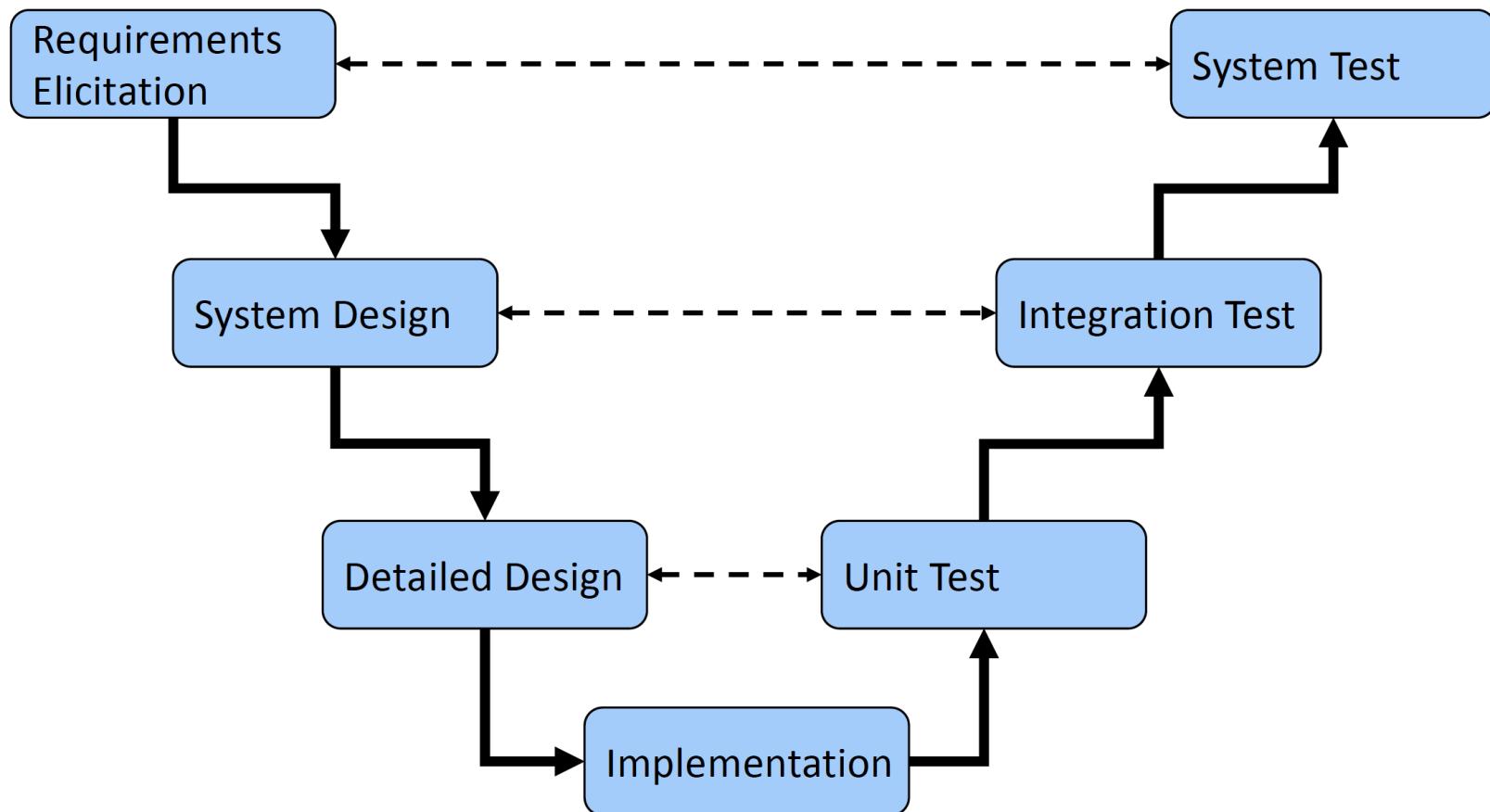
- Debugging era
- Demonstration era
- Destruction era
- Evaluation era
- Prevention era

Testing and Analysis

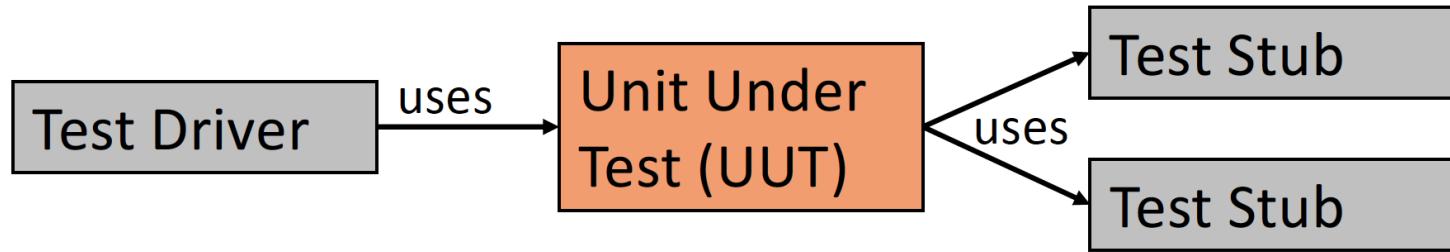
	manual	automatic
empirical	manual testing	automated testing
analytical	inspection	program analysis

<https://andrewbegel.com/info461/readings/verification.html>

Testing Stages



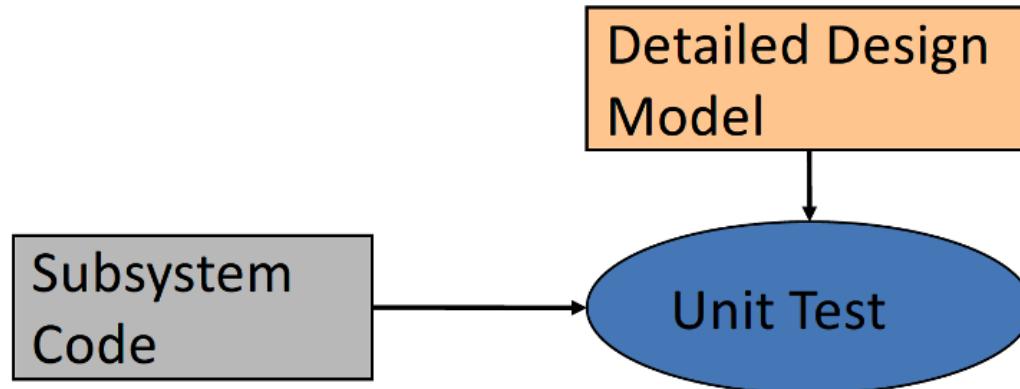
Creation of Test Harness



- Test driver
 - ✓ Apply tests to UUT, including setup & clean-up
- Test stub
 - ✓ *Partial, temporary implementation of a component used by UUT*
 - ✓ *Mock a missing component back fake data*

Unit Testing

- Testing individual units



- **Goal:** Confirm units are correct & meet intended functionalities

Unit Testing Example

```
class SavingsAccount {  
    ...  
  
    public void deposit( int amount ) { ... }  
    public void withdraw( int amount ) { ... }  
    public int getBalance( ) { ... }  
}
```

```
@Test  
public void withdrawTest( ) {  
    SavingsAccount target = new SavingsAccount();  
    target.deposit( 300 );  
    int amount = 100;  
    target.withdraw( amount );  
    Assert.assertTrue( target.getBalance( ) == 200 );  
}
```

Implement
test driver

Create
test data

Create
test oracle

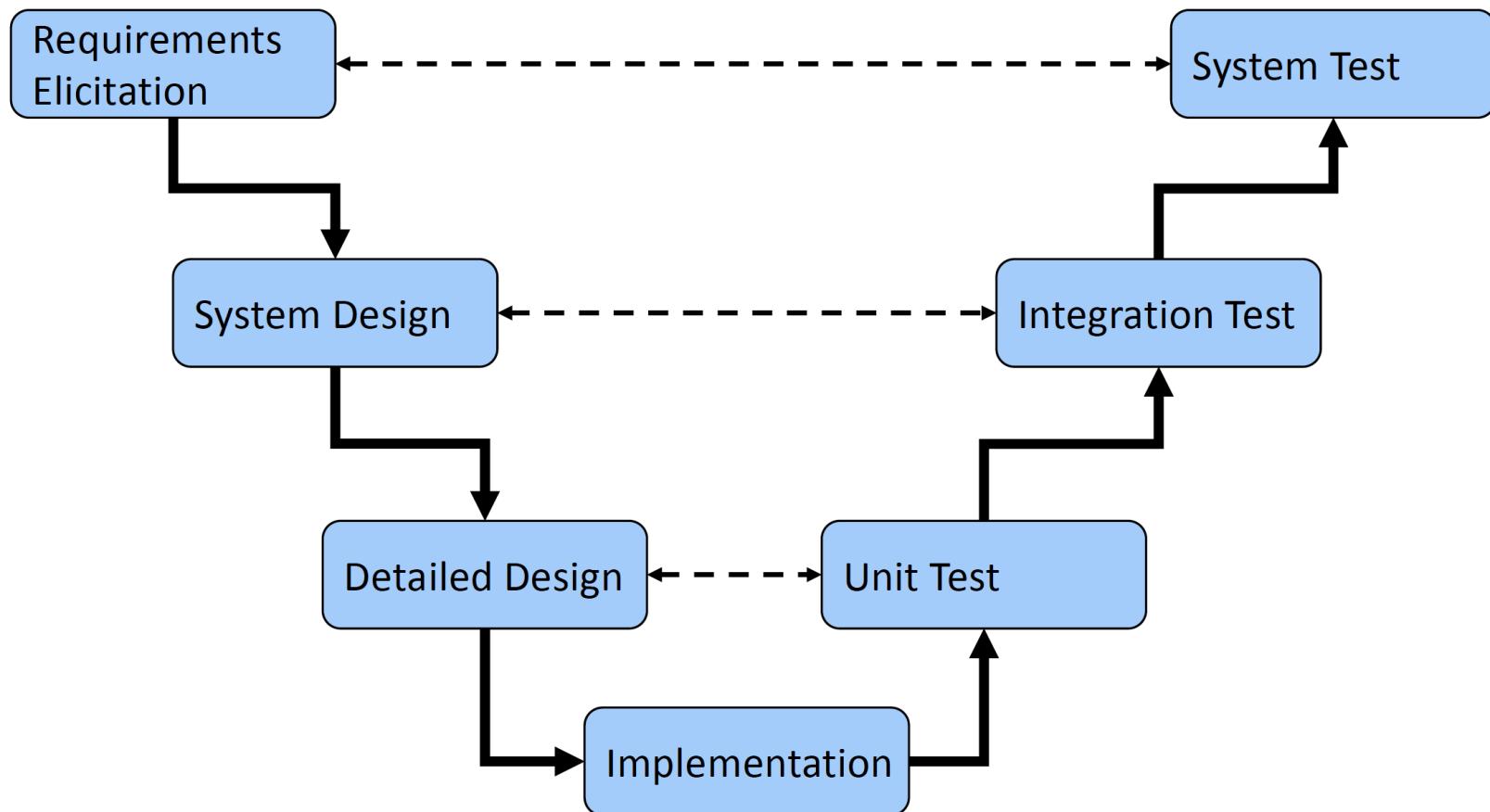
Test Execution

- Execute the tests

Eight Rules of Testing [M. Fowler]

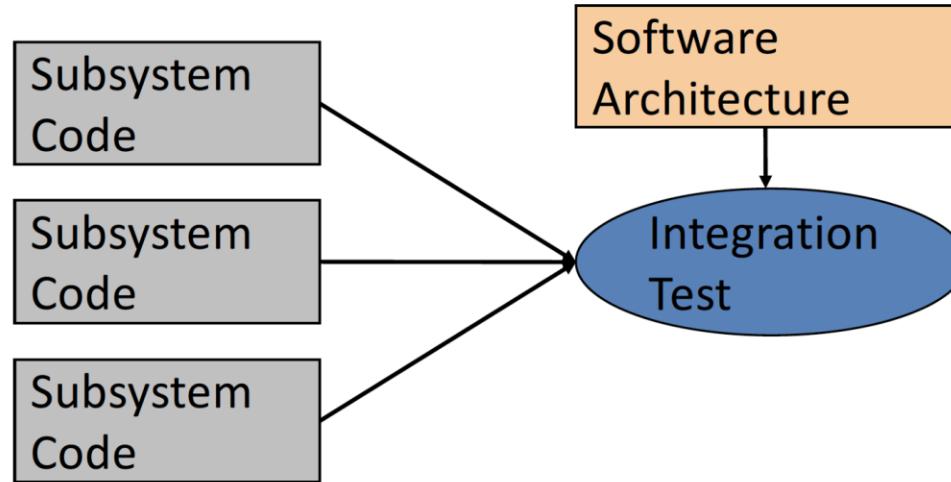
1. Make sure all tests are **fully automatic** and check their own results
2. A **test suite** is a powerful bug detector that reduces the time to find bugs
3. **Run** tests **frequently** – every test at least once a day
4. When you get a bug report, start by writing a **unit test to expose the bug**
5. Better to write and run incomplete tests than not run complete tests
6. Concentrate your tests on **boundary conditions**
7. Don't forget to test raised **exceptions** when things are expected to go wrong
8. Do not let the fear that testing can't catch all bugs stop you from writing tests that will catch most bugs

Testing Stages



Integration Testing

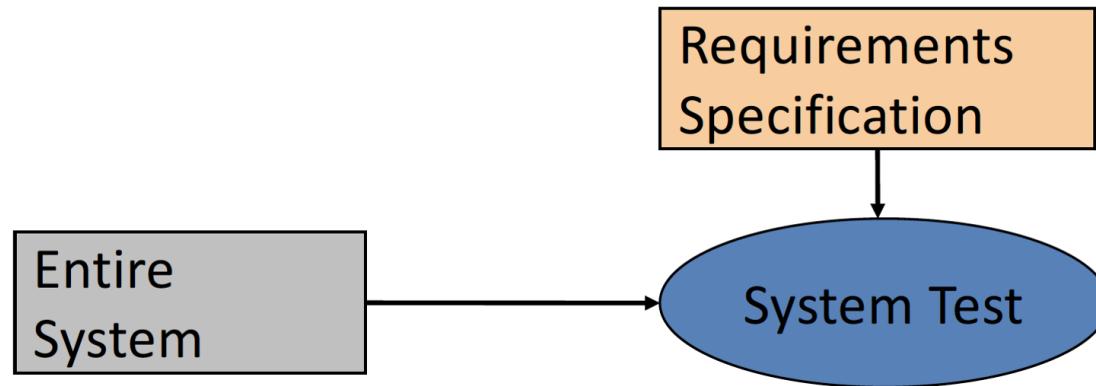
- Test groups of subsystem & eventually the entire system



- Goal: Test interfaces between subsystems

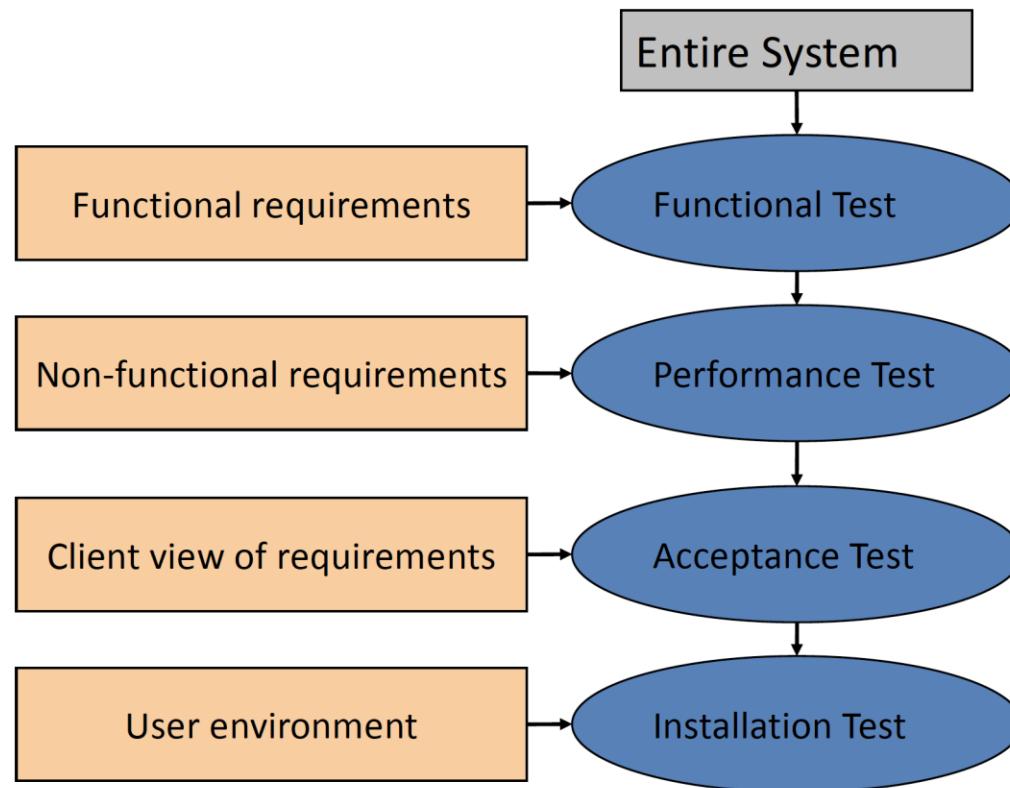
System Testing

- Test the entire system



- Goal: Determine if the system meets the requirements, both functional and non-functional

System Testing Stages



Functional Testing

- Goal: Test functionality of system
 - System is treated as a black box
- Test cases are design from requirements
 - Based on use cases
 - Alternative source: user manual
- Test cases describe
 - Input data
 - Flow of events
 - Results to check

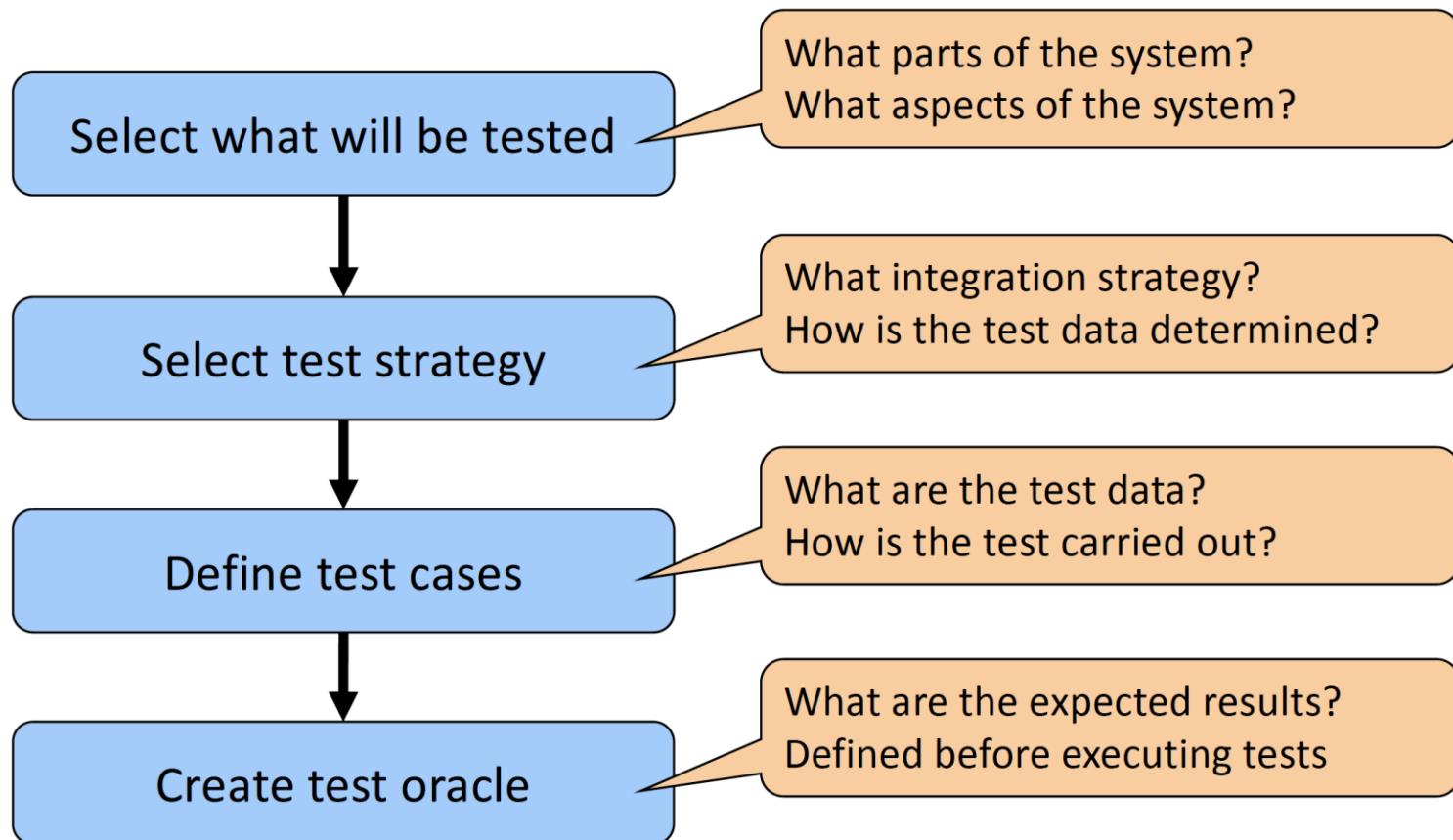
Acceptance Testing

- **Goal:** Demonstrate system meets customer requirements
- Performed by the client, not by the developer
- **Alpha test**
 - Client uses the software at the developer's site
 - Software used in controlled setting --- developers ready to fix bugs
- **Beta test**
 - Conducted at client's site (developer not present)
 - Software gets a realistic workout in target environments

Independent Testing

- Hard for programmers to accept they made a mistake
 - Plus a vested interest in not finding mistakes
 - Often stick to the data that makes the program work
- Designing and programming are constructive tasks
 - Testers must seek to break the software
- Testing is done best by independent testers

Testing Steps



Testing Strategies

- **Exhaustive testing:** Check UUT for all possible inputs
 - Not feasible, even for trivial programs
- **Random testing:** Select test data randomly & uniformly
- **Functional testing:** Use [requirements](#) to decide test cases
- **Structural testing:** Use [design knowledge](#) about system structure, algorithms, data structures to determine test cases that exercise a large portion of the code

Testing Strategies

Functional testing

- Goal: Cover all requirements
- Black-box test
- Suitable for all test stages

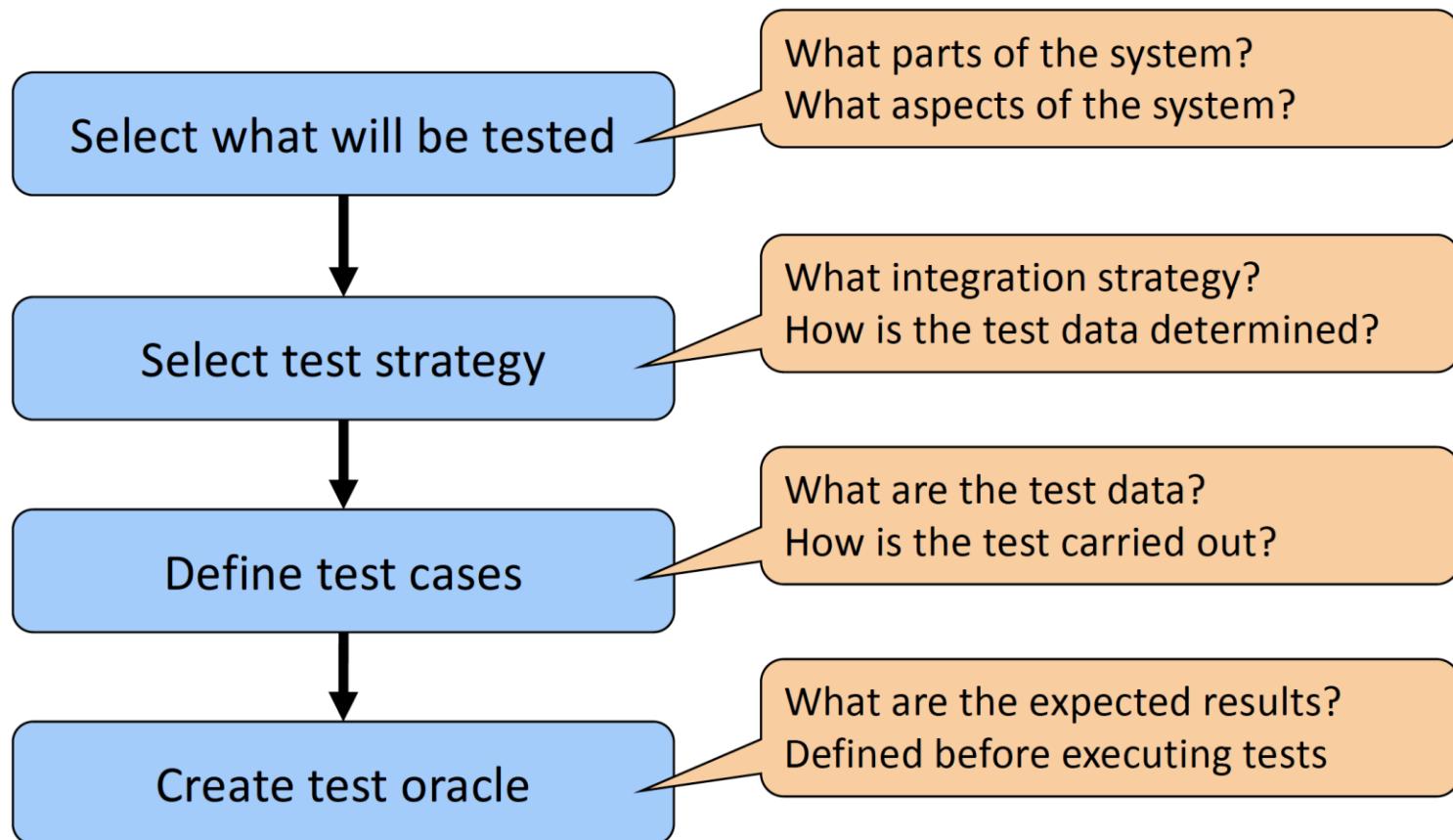
Structural testing

- Goal: Cover all the code
- White-box test
- Suitable for unit testing

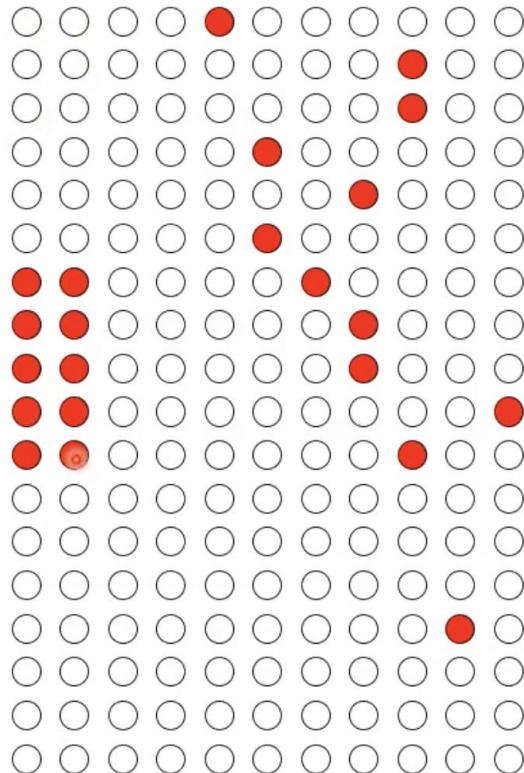
Random testing

- Goal: Cover corner cases
- Black-box test
- Suitable for all test stages

Testing Steps



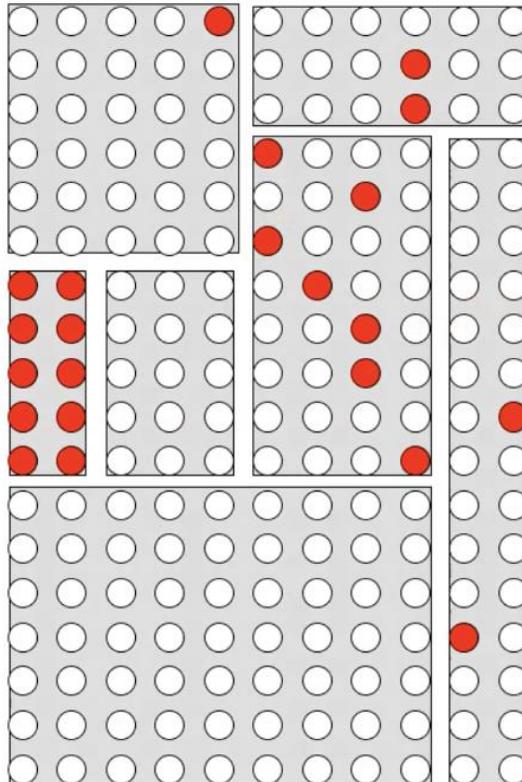
Finding representative inputs



● Failure

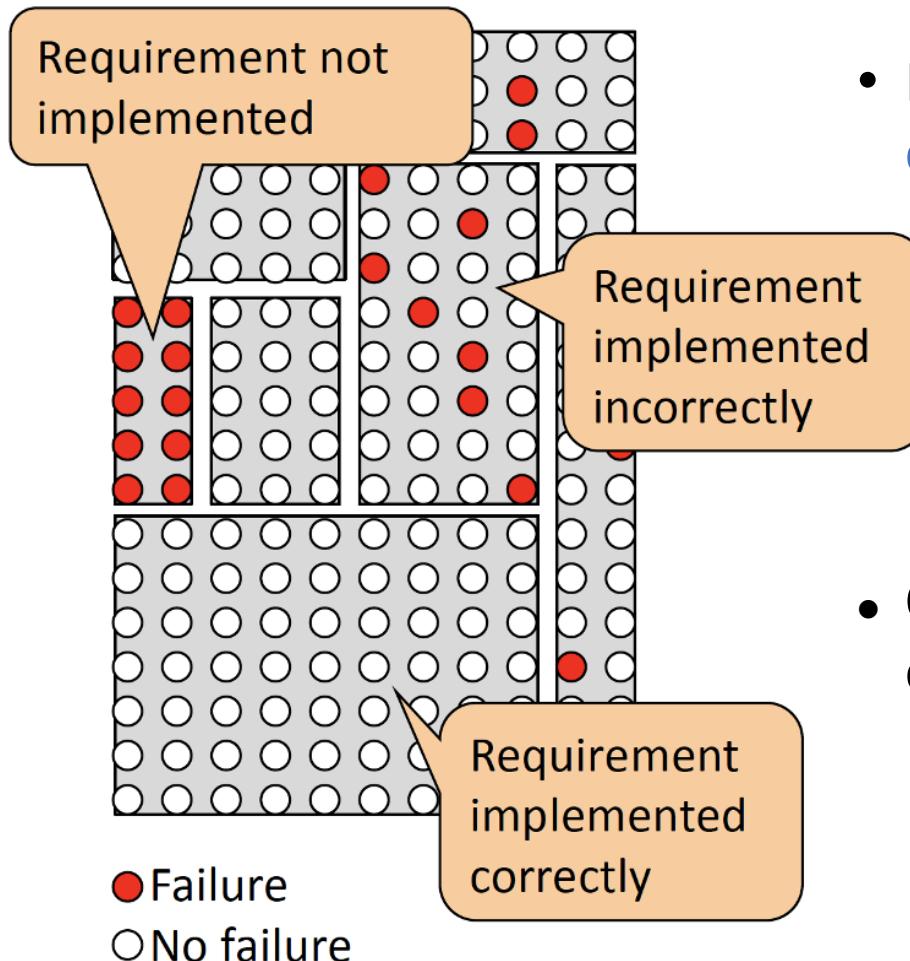
○ No failure

Finding representative inputs



- Divide inputs into **equivalence classes**
 - Each possible input belongs to one equivalence class
 - Goal: some classes have higher failure density
- Choose tests for each equivalence class

Finding representative inputs



- Divide inputs into equivalence classes
 - Each possible input belongs to one equivalence class
 - Goal: some classes have higher failure density
- Choose tests for each equivalence class

Selecting representative values

- Once partitioned the inputs, need to select **concrete values** for the tests for each **equivalence class**
- Input from a range of valid values
 - Below, within, and above the range
- Input from a discrete set of values
 - Valid and invalid discrete values
 - Instances of each subclass

Boundary Testing

Given an integer x ,
determine the
absolute value of x

x	
Valid	
	all values

```
int abs( int x ) {  
    if( 0 <= x ) return x;  
    return -x;  
}
```

Negative result for
 $x == \text{Integer.MIN_VALUE}$

`Integer.MIN_VALUE: -2147483648`
`MAX_VALUE: 2147483647`

- Many errors occur at boundaries of the input domain
 - Overflows
 - Comparisons ('<' instead of '<='), etc.)
 - Missing emptiness checks (e.g., collections)
 - Wrong number of iterations

Combinatorial Testing

- Many input values: equiv. classes + boundary testing
- Testing all combinations leads to combinatorial explosion
- Reduce test cases to make effort feasible
 - Random selection
 - Semantic constraints
 - Combinatorial selection

Semantic Constraints

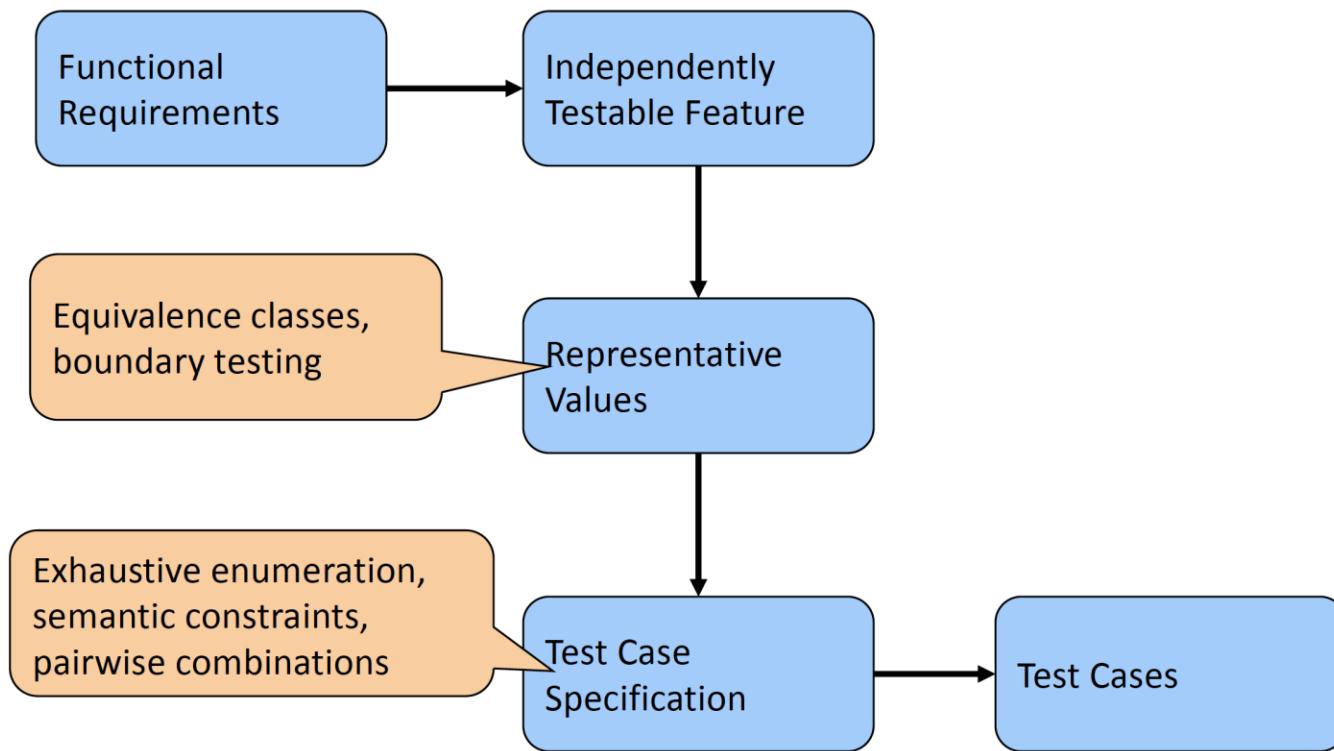
Given three values, a , b , c , compute all solutions of the equation $ax^2 + bx + c = 0$;
report an error if all three values are zero

	Two solutions	One solution	No solution
Linear equation		$a = 0$ and $b \neq 0$	$a = 0$, $b = 0$, and $c \neq 0$
(Truly) quadratic equation	$a \neq 0$ and $b^2 - 4ac > 0$	$a \neq 0$ and $b^2 - 4ac = 0$	$a \neq 0$ and $b^2 - 4ac < 0$
Invalid input	$a = 0$, $b = 0$, $c = 0$		

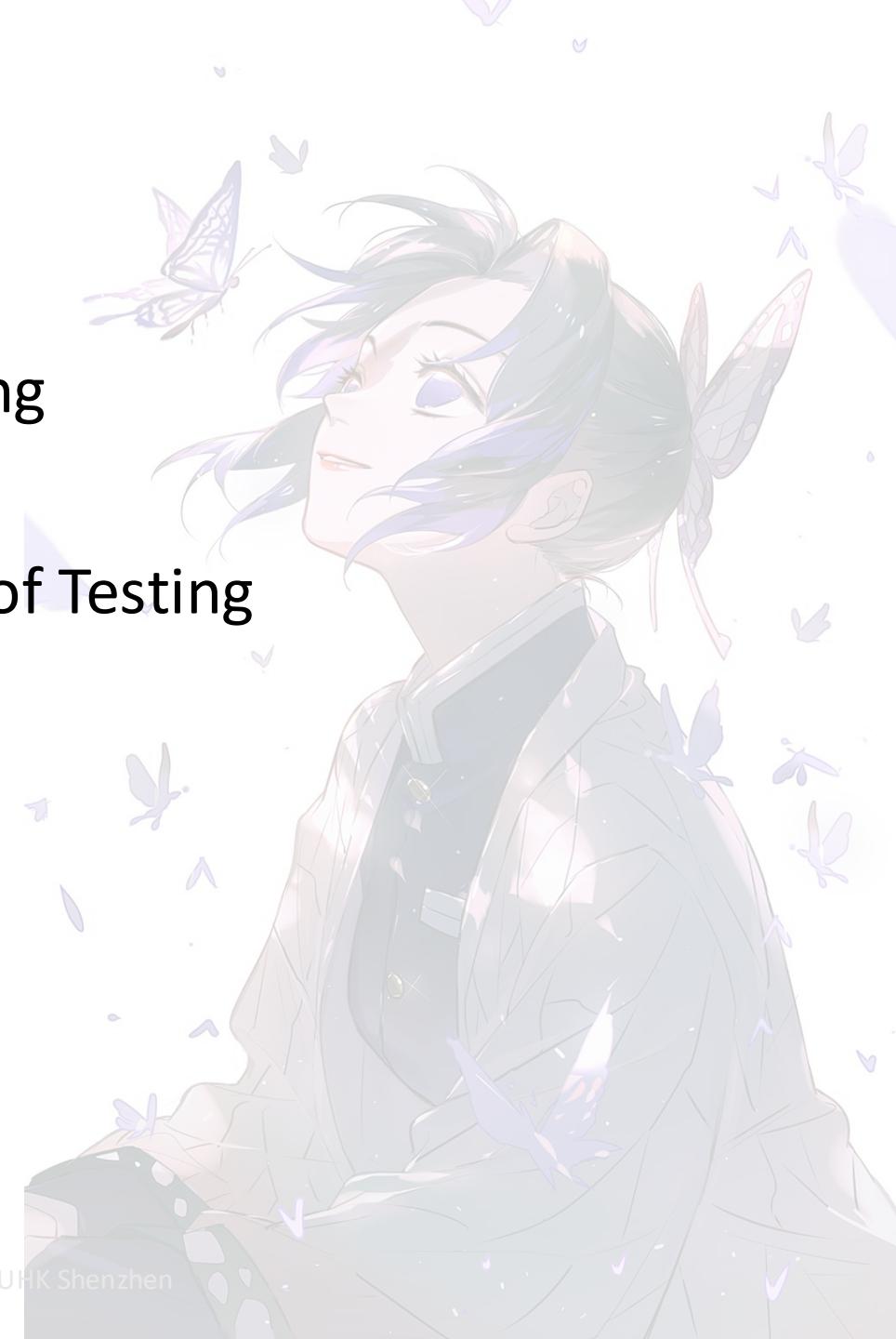
Pairwise Combinatorial Testing

- Focus on possible combinations of pairs of inputs
- Example: Consider a method with 4 Boolean parameters
 - Combinatorial testing requires $2^4 = 16$ test cases
 - Pairwise-combinations testing requires 5 test cases:
TTTT, TFFF, FTFF, FFTF, FFFT
- Can be generalized to k-tuples (k-way testing)
- For n parameters with d values per parameter, the number of test cases grow logarithmically in n and quadratic in d
- Results hold for large n & d, and for all k in k-way testing

Functional Testing: Summary

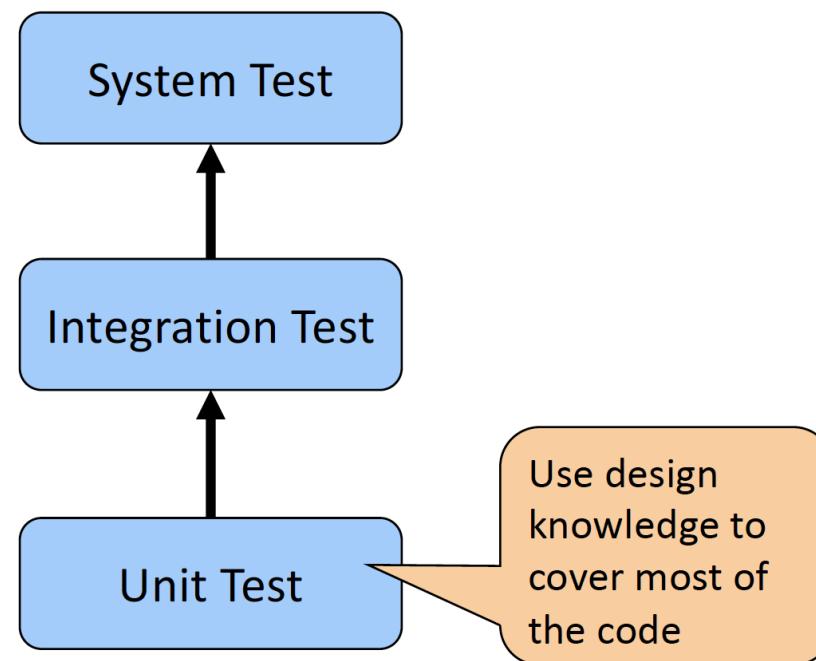


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Structural Testing

- White-box test a unit to cover a large portion of its code



Basic Blocks

- **Basic block:** a sequence of statements having
 - **One entry point:** no code within is a jump target
 - **One exit point:** only the last instruction exits the block
 - If the first instruction is executed, the rest are also executed

Basic Blocks: Example

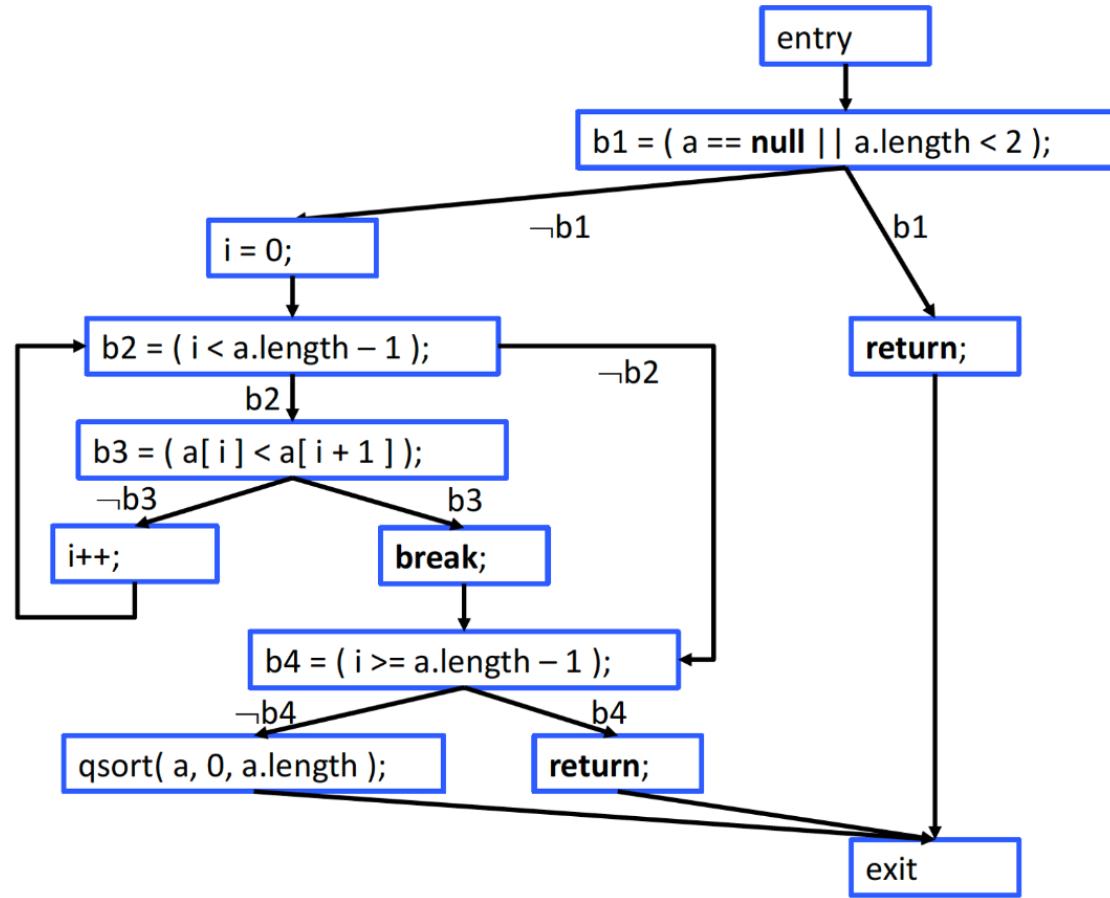
```
public void sort( int[ ] a ) {  
    if( a == null || a.length < 2 )  
        return;  
  
    int i;  
    for( i = 0; i < a.length - 1; i++ ) {  
        if( a[ i ] < a[ i + 1 ] )  
            break;  
    }  
    if( i >= a.length - 1 )  
        return;  
    qsort( a, 0, a.length );  
}
```

Intraprocedural Control Flow Graphs

- An **intraprocedural CFG** of a proc. p is a graph (N, E) where
 - Nodes
 - Basic blocks in p
 - Designated entry and exit blocks
 - Edges
 - **(a, b)**: an edge from a to b if there is direct control flow from a to b
 - **(entry, a)**: a is the first basic block of p
 - **(b , exit)**: for each b ending with a (possibly implicit) return statement

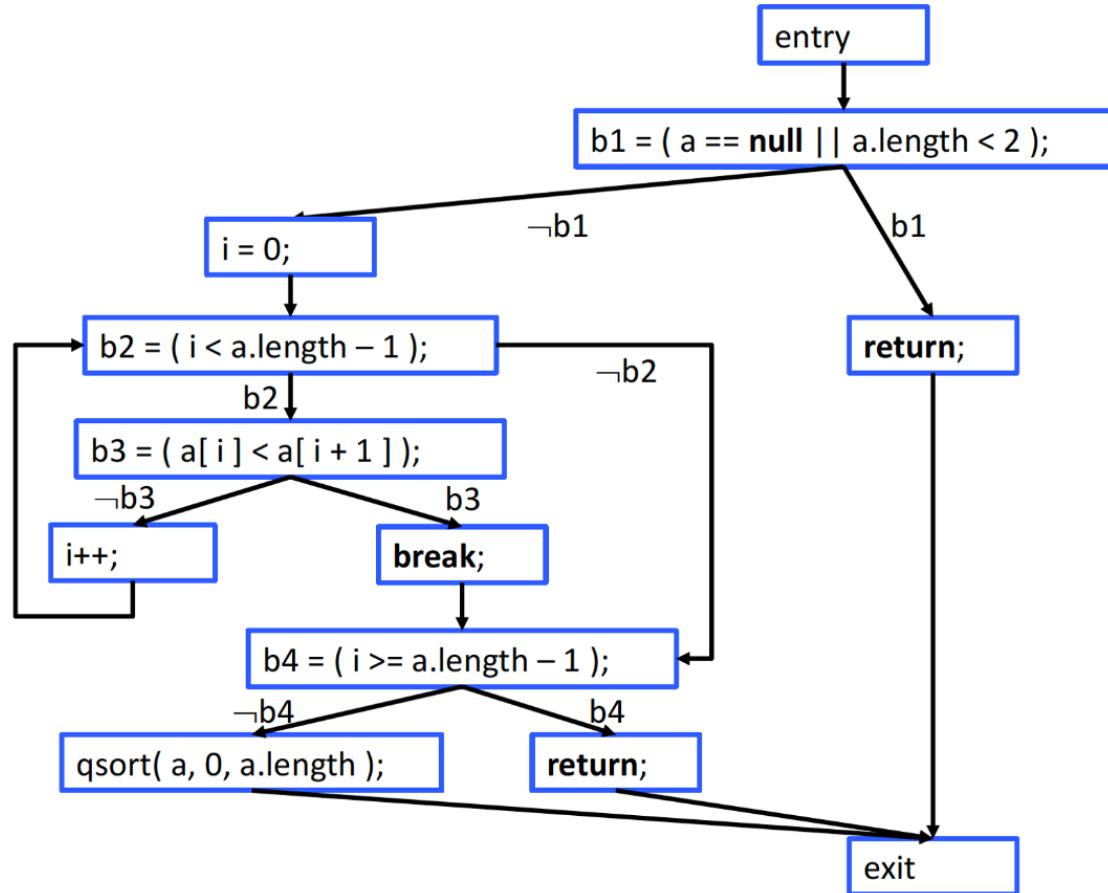
Control Flow Graphs: Example

```
public void sort( int[ ] a ) {
    if( a == null || a.length < 2 )
        return;
    int i;
    for( i = 0; i < a.length - 1; i++ ) {
        if( a[ i ] < a[ i + 1 ] )
            break;
    }
    if( i >= a.length - 1 )
        return;
    qsort( a, 0, a.length );
}
```



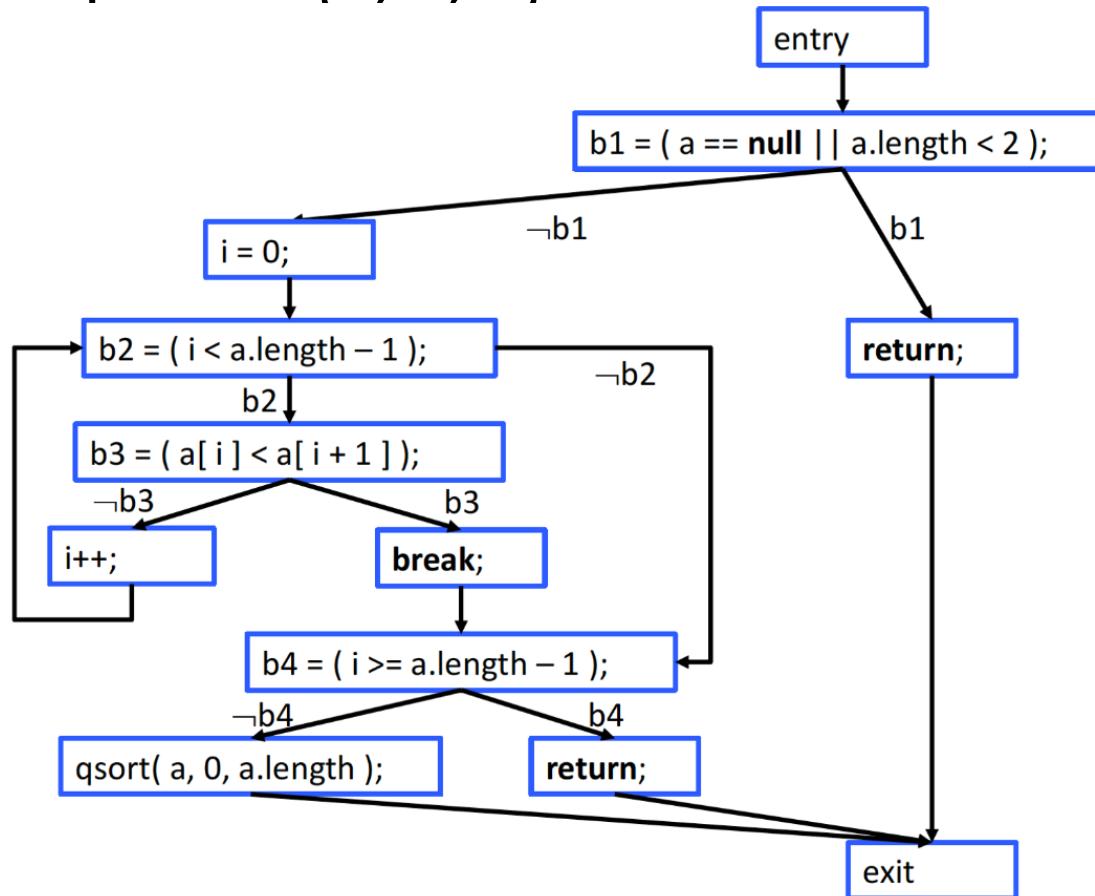
Test Coverage

- The CFG serves as an **adequacy criterion** for tests
- The more parts executed, the higher chance to find bugs
- “parts” can be
 - Nodes
 - Edges
 - Paths
 - etc.



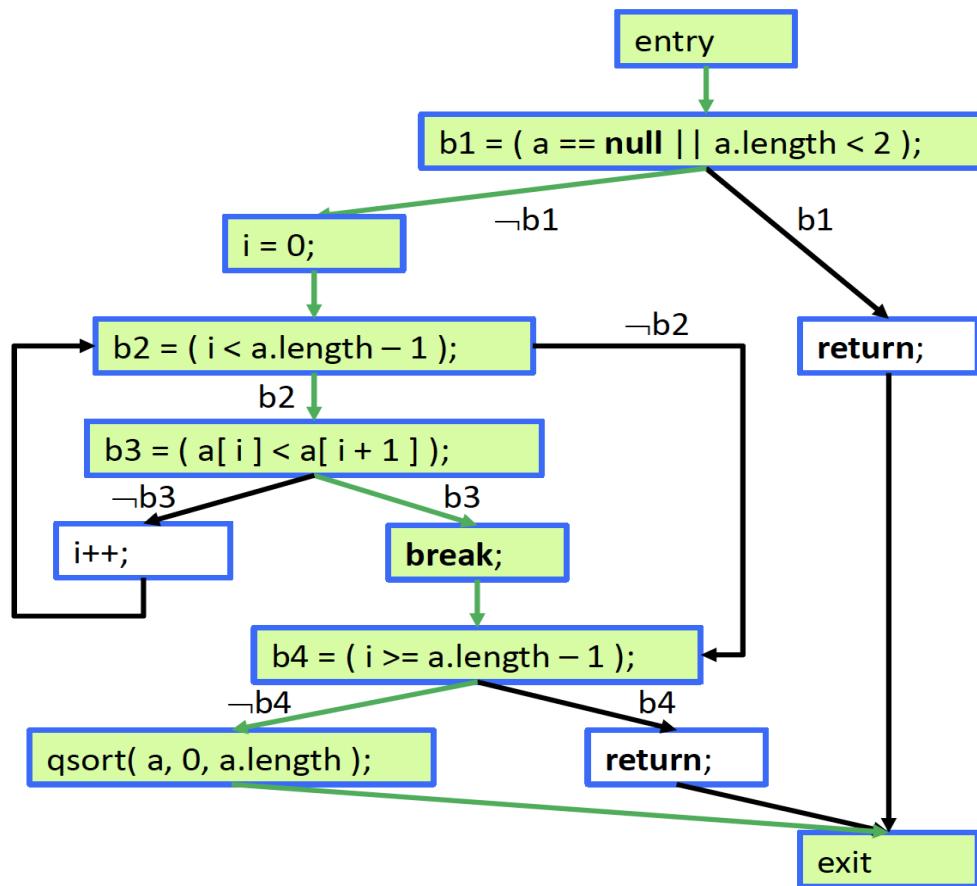
Test Coverage

- Consider input $a = \{3, 7, 5\}$



Test Coverage

- Consider input $a = \{3, 7, 5\}$



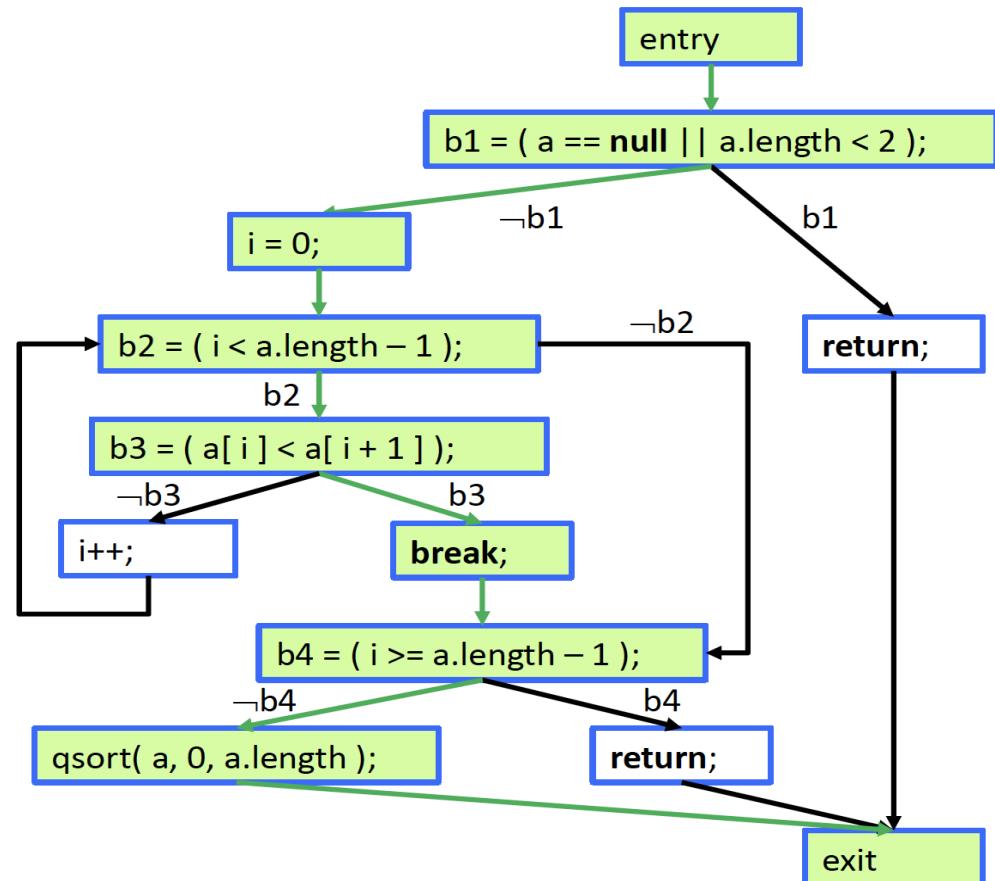
Statement Coverage

- Judge a test suite's quality by how much the CFG is covered
- Idea: Can detect a bug in a statement only by executing it
 - Can also be defined on basic blocks

$$\text{Statement Coverage} = \frac{\text{Number of executed statements}}{\text{Total number of statements}}$$

Statement Coverage: Example

- Consider input $a = \{3, 7, 5\}$
- It executes 7/10 basic blocks
- Statement coverage: 70%
- How to achieve 100% statement coverage?

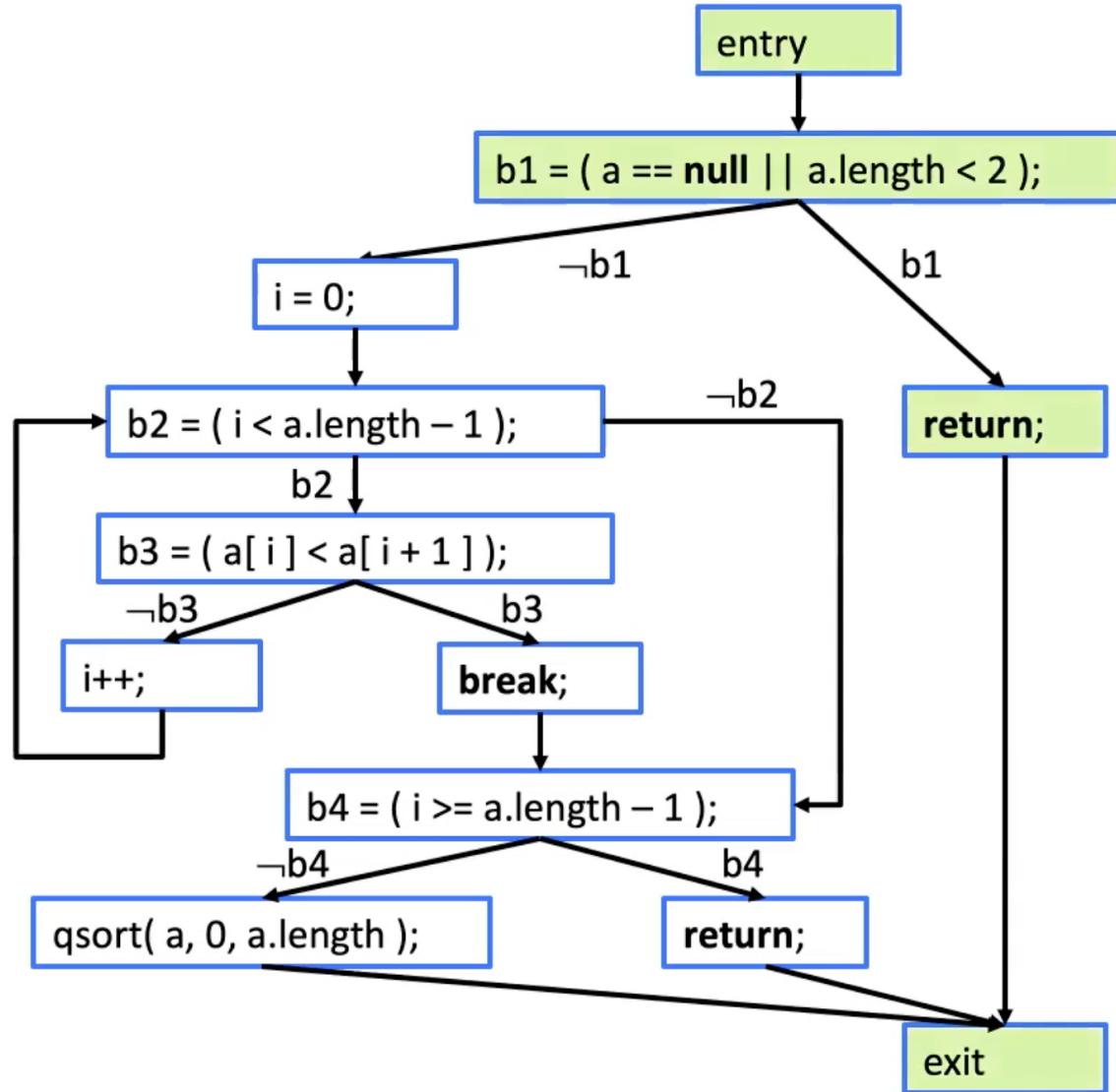


Statement Coverage: Example (cont.)

- We can achieve 100% statement coverage with three test cases

- $a = \{ 1 \}$
- $a = \{ 5, 7 \}$
- $a = \{ 7, 5 \}$

- The last detects the bug

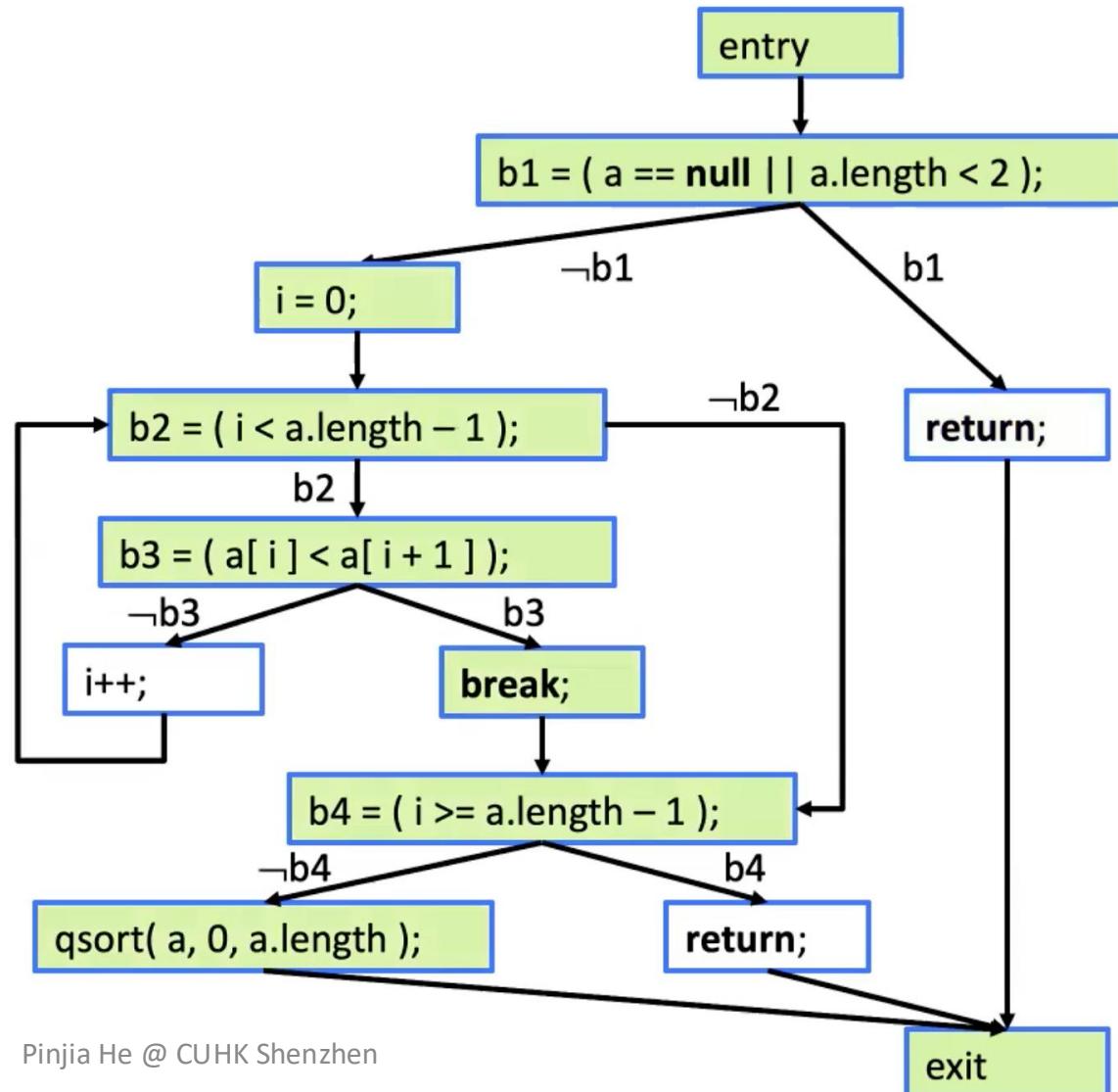


Statement Coverage: Example (cont.)

- We can achieve 100% statement coverage with three test cases

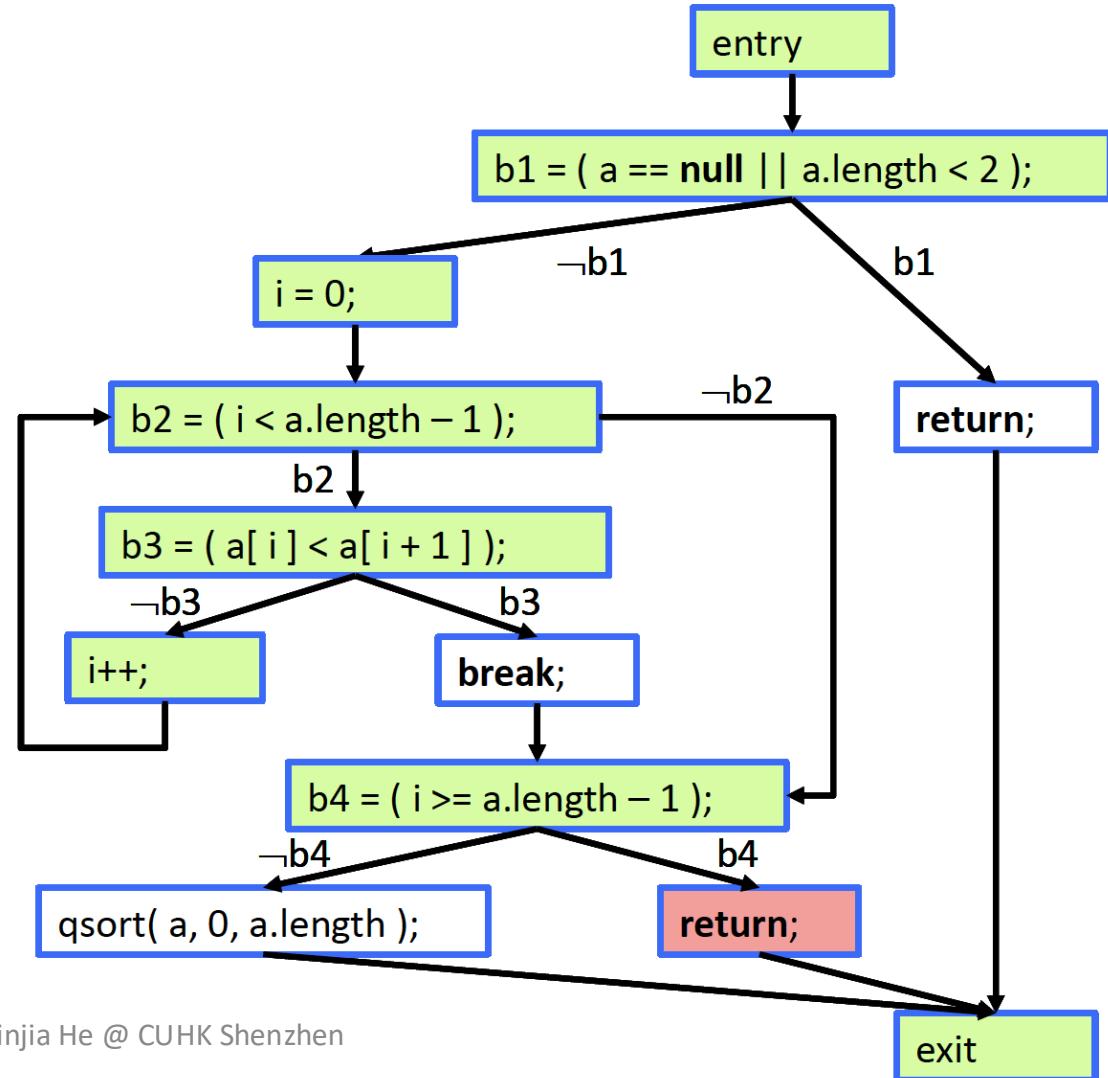
- $a = \{ 1 \}$
- $a = \{ 5, 7 \}$
- $a = \{ 7, 5 \}$

- The last detects the bug



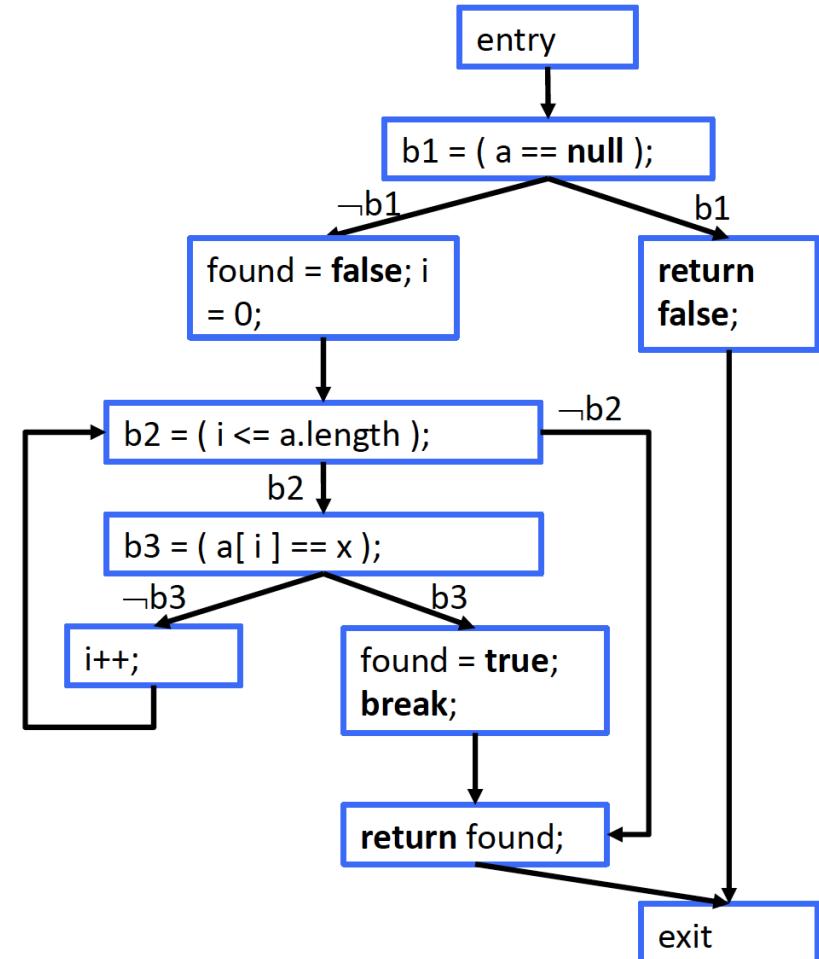
Statement Coverage: Example (cont.)

- We can achieve 100% statement coverage with three test cases
 - $a = \{ 1 \}$
 - $a = \{ 5, 7 \}$
 - $a = \{ 7, 5 \}$
- The last detects the bug



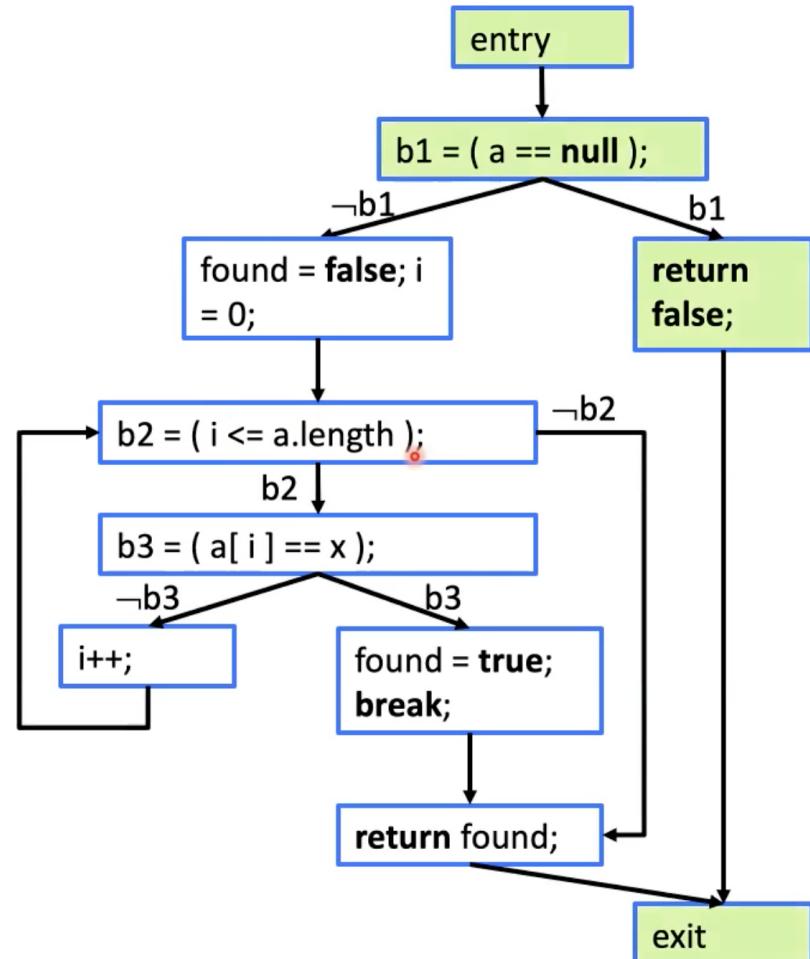
Statement Coverage: Discussion

```
boolean contains( int[ ] a, int x ) {  
    if( a == null ) return false;  
  
    boolean found = false;  
  
    for( int i = 0; i <= a.length; i++ ) {  
  
        if( a[ i ] == x ) {  
            found = true;  
            break;  
        }  
    }  
  
    return found;  
}
```



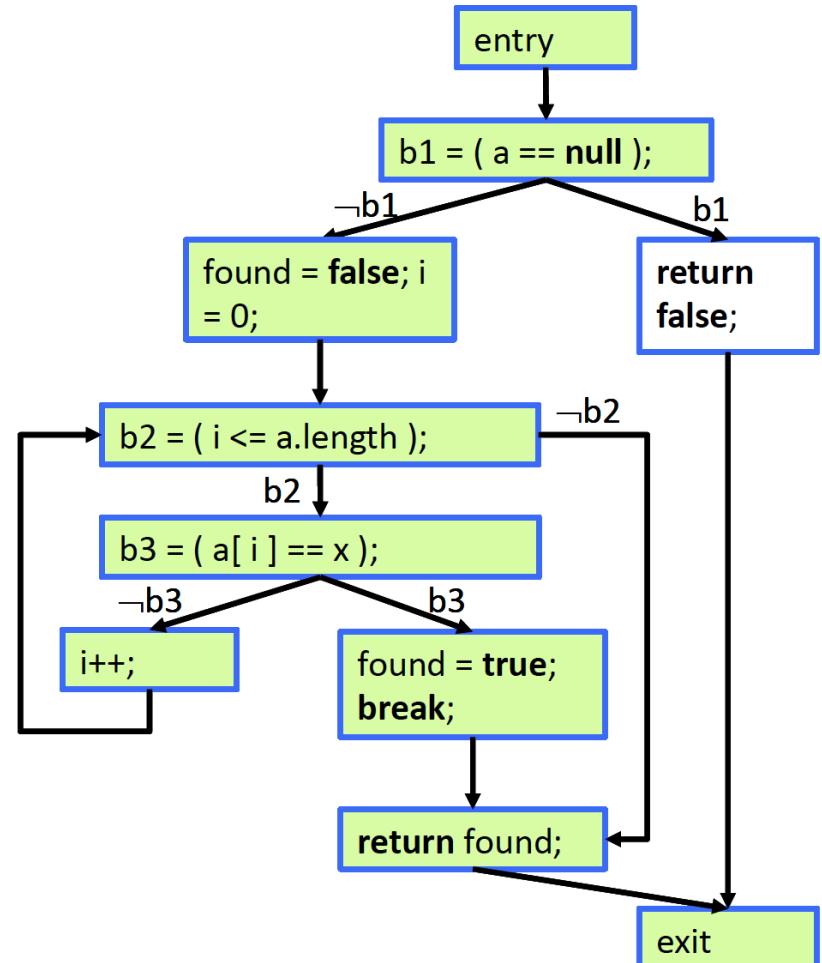
Statement Coverage: Discussion

- Achieve 100% statement coverage with 2 tests
 - $a = \text{null}$
 - $a = \{ 1, 2 \}, x = 2$



Statement Coverage: Discussion

- Achieve 100% statement coverage with 2 tests
 - $a = \text{null}$
 - $a = \{ 1, 2 \}, x = 2$
- The tests do not detect the bug
- More thorough testing is needed



Branch Coverage

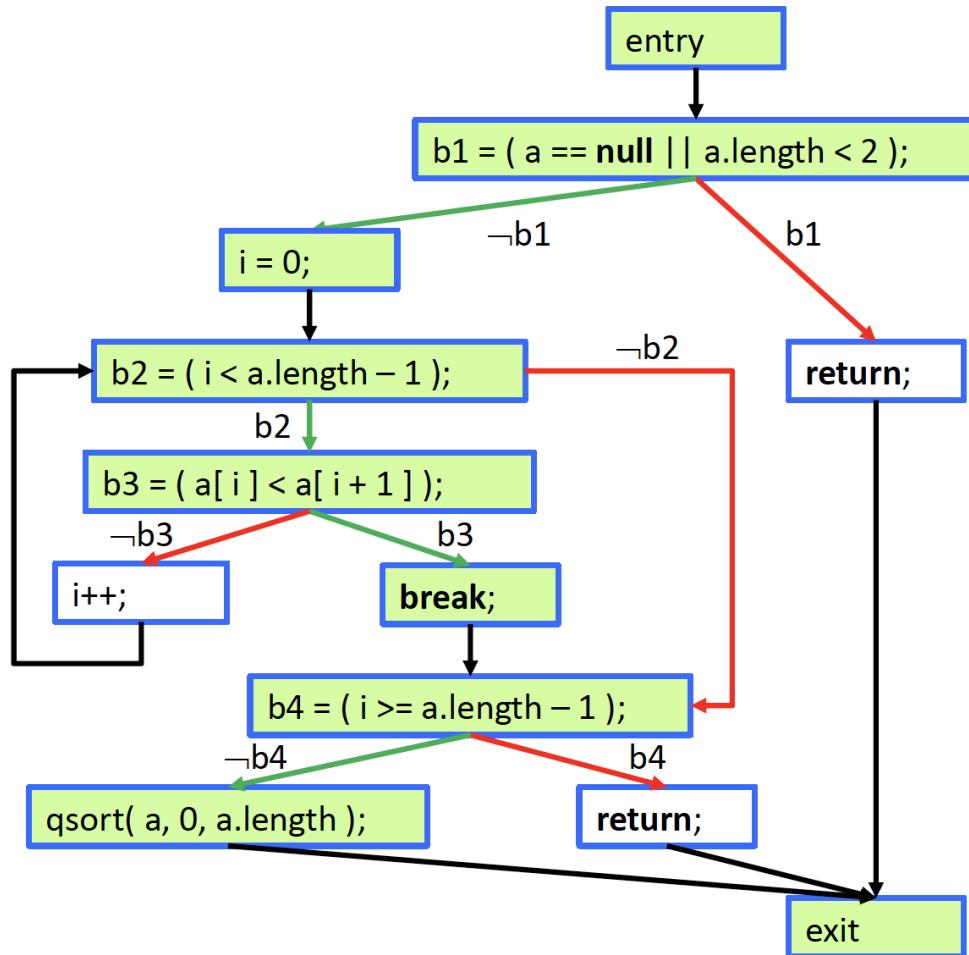
- **Idea:** test all possible branches in the control flow
- Edge (m,n) is a branch iff there is edge (m, n') with $n \neq n'$

$$\text{Branch Coverage} = \frac{\text{Number of executed branches}}{\text{Total number of branches}}$$

- Branch coverage is 100% if code has no branches
- Most widely-used adequacy criterion in industry
- Branch coverage is more thorough than statement coverage
 - Q: 100% branch coverage \Rightarrow 100% statement coverage ?
 - Q: “ $>=n\%$ branch cov.” \Rightarrow “ $>=n\%$ statement cov.” ?

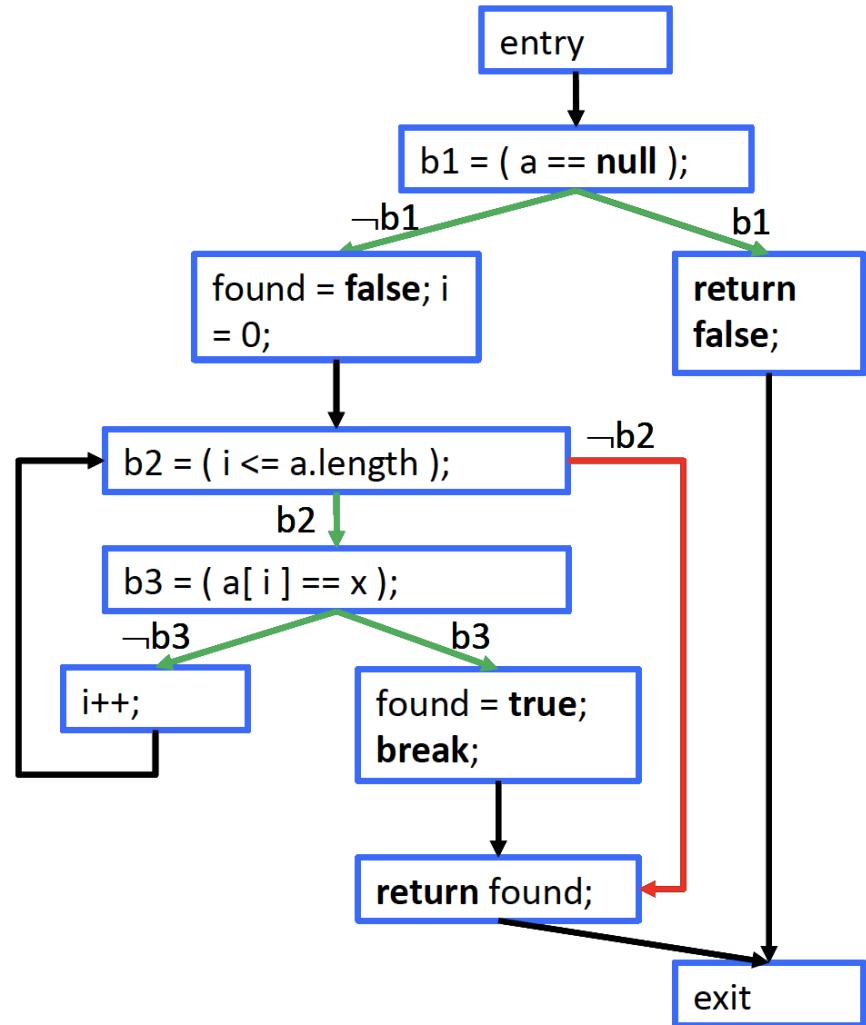
Branch Coverage: Example 1

- Consider input $a = \{3, 7, 5\}$
- It executes 4/8 branches
- Branch coverage 50%
- What tests to include for 100% branch coverage



Branch Coverage: Example 2

- The two tests:
 - $a = \text{null}$
 - $a = \{1,2\}$, $x = 2$execute 5/6 branches
- Branch coverage: 83%
- How to achieve 100% branch coverage?

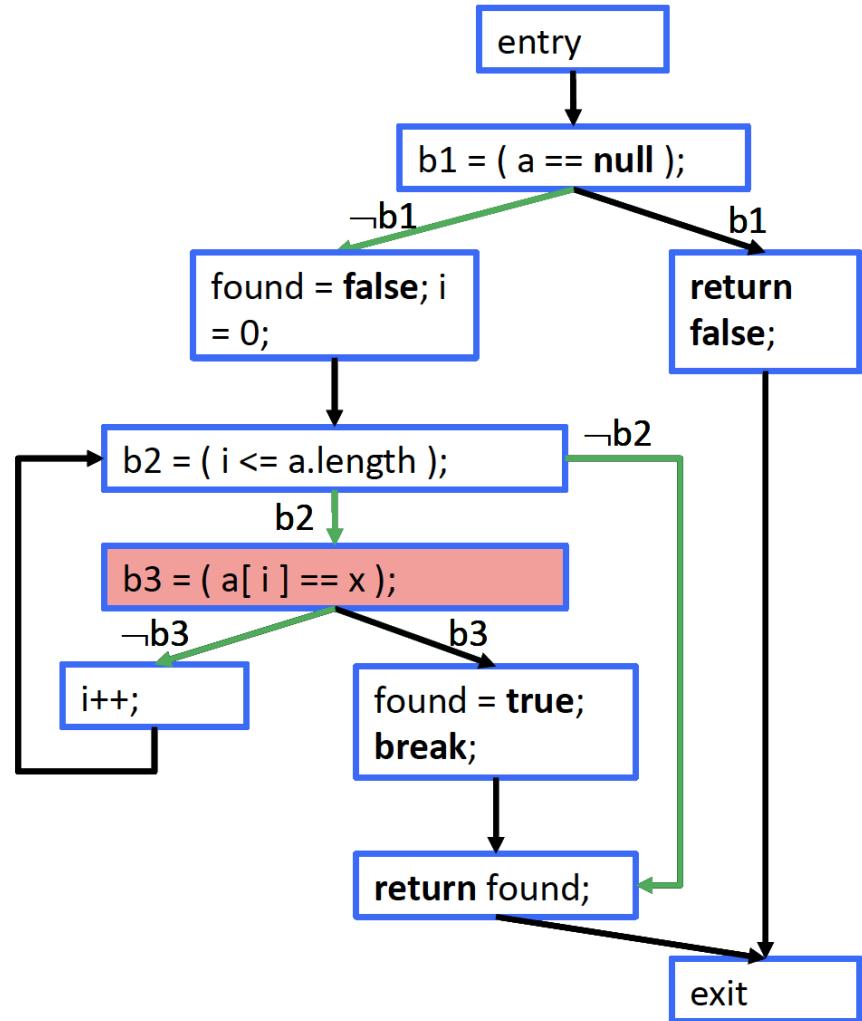


Branch Coverage: Example 2

- To achieve 100% branch coverage requires a test that runs the loop to the end

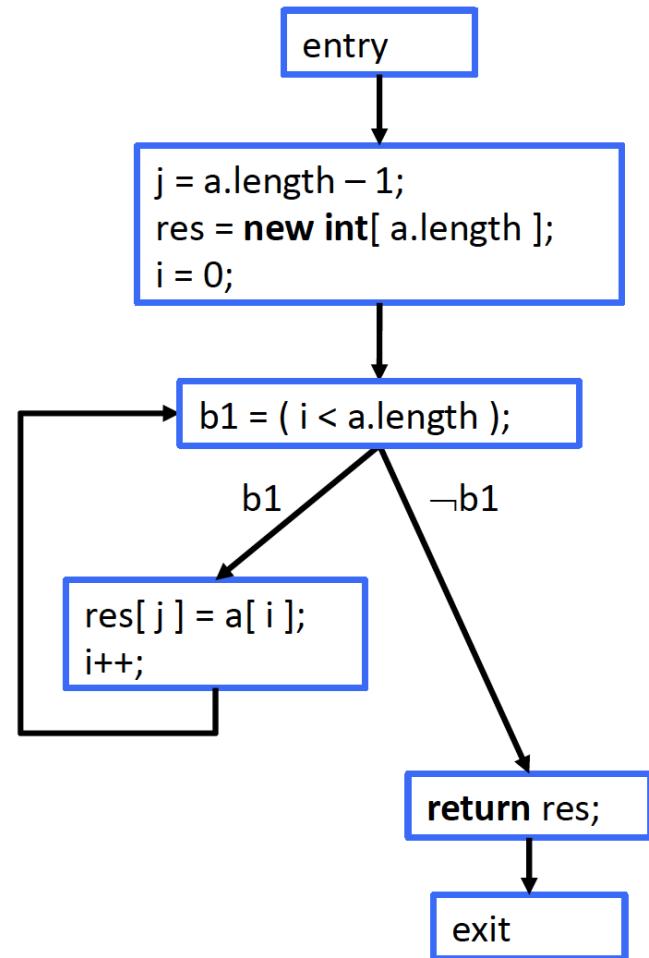
- $a = \text{null}$
- $a = \{1\}, x = 1$
- $a = \{1\}, x = 3$

- The last one detects the bug



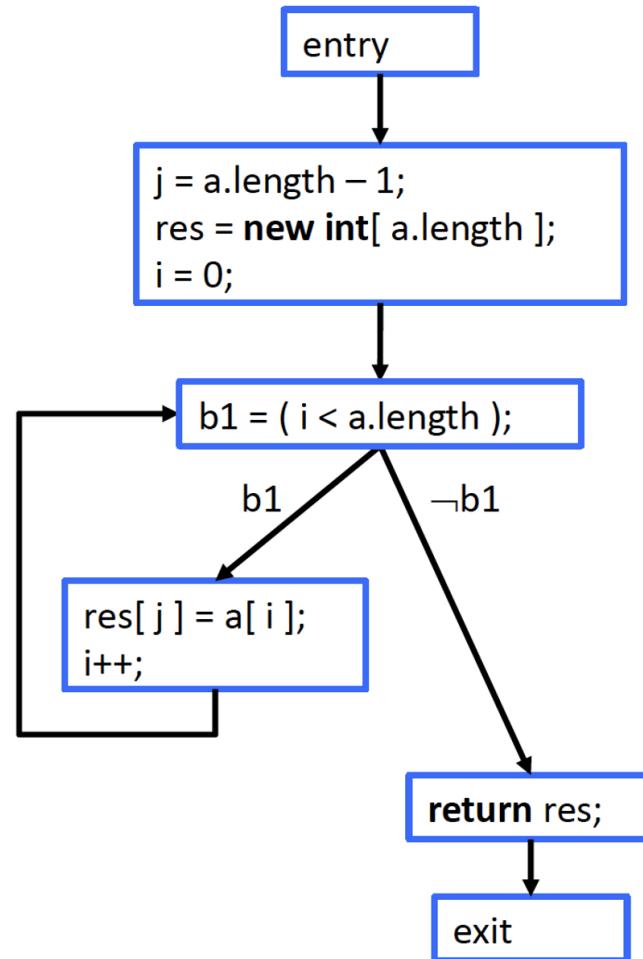
Branch Coverage: Discussion

```
int[ ] reverse( int[ ] a ) {  
    int j = a.length - 1;  
    int[ ] res = new int[ a.length ];  
    for( int i = 0; i < a.length; i++ ) {  
        res[ j ] = a[ i ];  
    }  
    return res;  
}
```



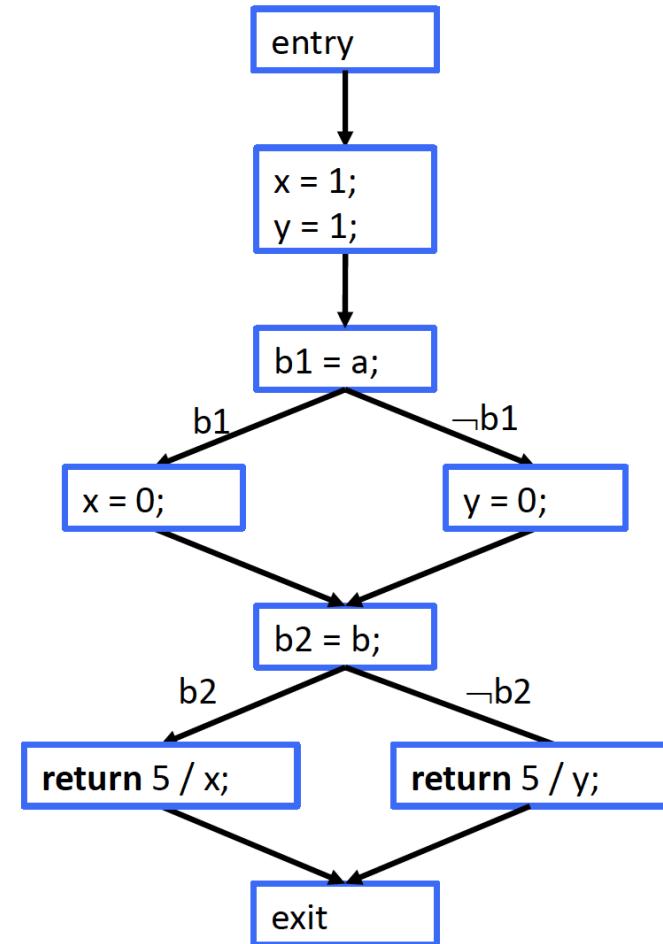
Branch Coverage: Discussion

- Can obtain 100% branch coverage with 1 test
 - $a = \{1\}$
- The test doesn't detect the bug
- More tests are needed



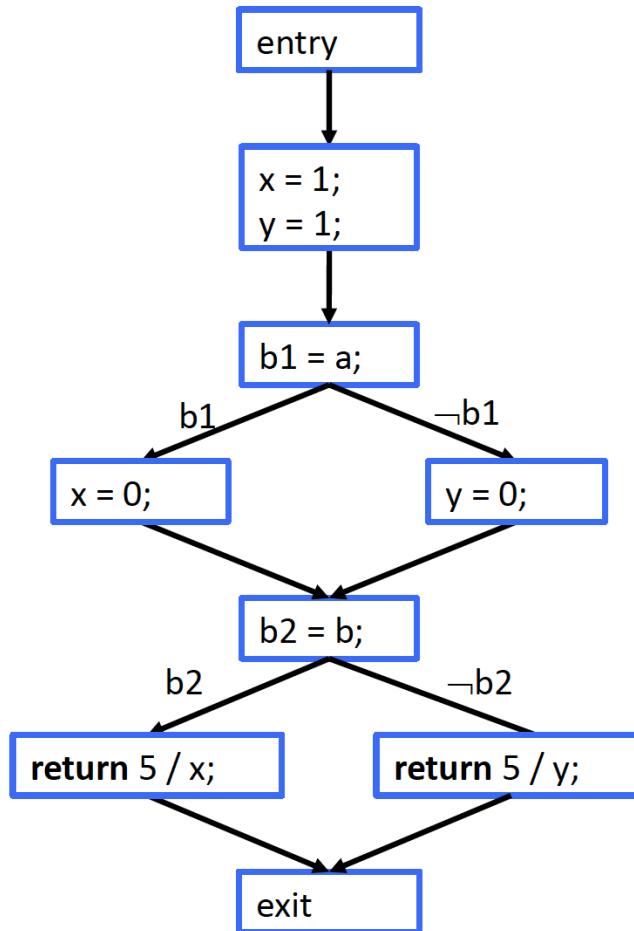
Branch Coverage: Discussion

```
int foo( boolean a, boolean b ) {  
    int x = 1;  
    int y = 1;  
    if( a )  
        x = 0;  
    else  
        y = 0;  
    if( b )  
        return 5 / x;  
    else  
        return 5 / y;  
}
```



Branch Coverage: Discussion

- We can achieve 100% branch coverage with two tests
 - `a = true, b = false`
 - `a = false, b = true`
- The tests do not detect the bug
- More thorough testing is needed



Path Coverage

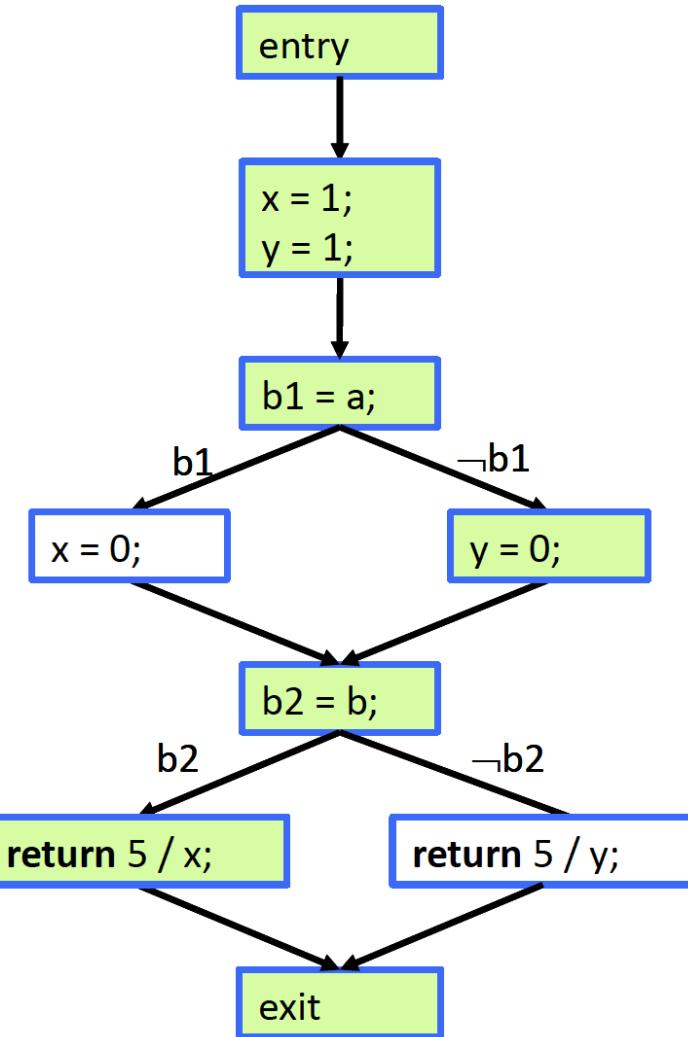
- **Idea:** test all possible paths through the CFG
- A path is a sequence of nodes $n_1, \dots n_k$ such that
 - $n_1 = \text{entry}$
 - $n_k = \text{exit}$
 - There is an edge (n_i, n_{i+1}) in the CFG

$$\text{Path Coverage} = \frac{\text{Number of executed paths}}{\text{Total number of paths}}$$

- Path coverage is more thorough than statement and branch coverage
 - 100% path coverage => 100% complete statement and branch coverage
 - But, “ $\geq n\%$ path coverage” does not generally imply “ $\geq n\%$ statement coverage” or “ $\geq n\%$ branch coverage”
- Complete path coverage is not feasible for input-dependent loops
 - Unbounded number of paths

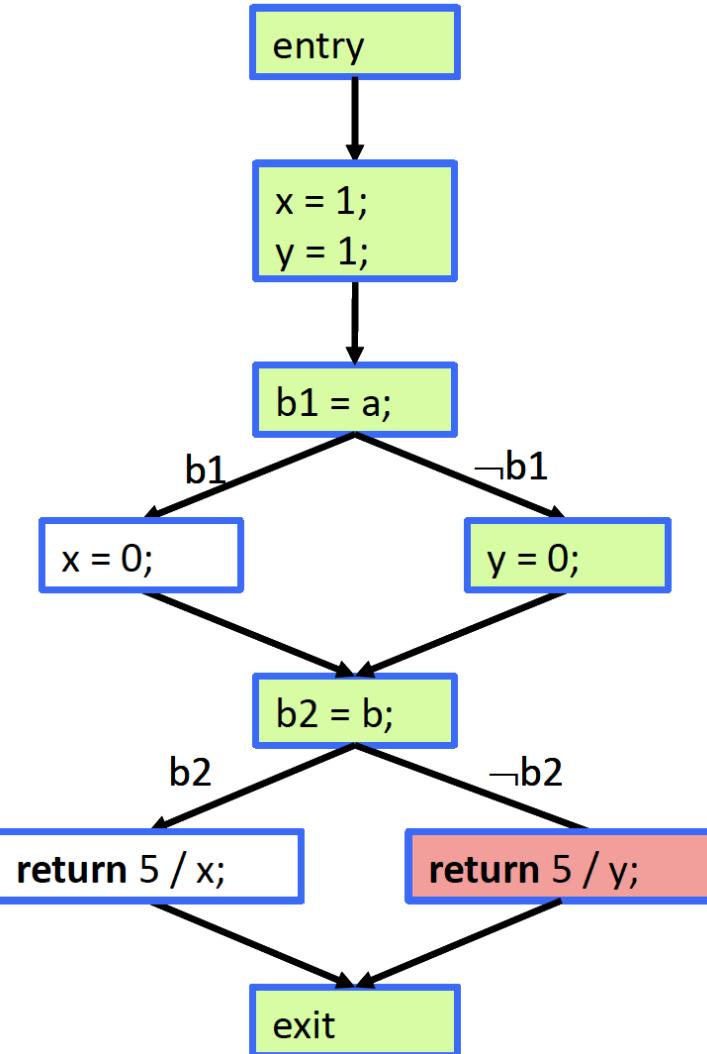
Path Coverage: Example 1

- The two tests
 - $a = \text{true}, b = \text{false}$
 - $a = \text{false}, b = \text{true}$execute 2/4 paths
- Path coverage: 50%



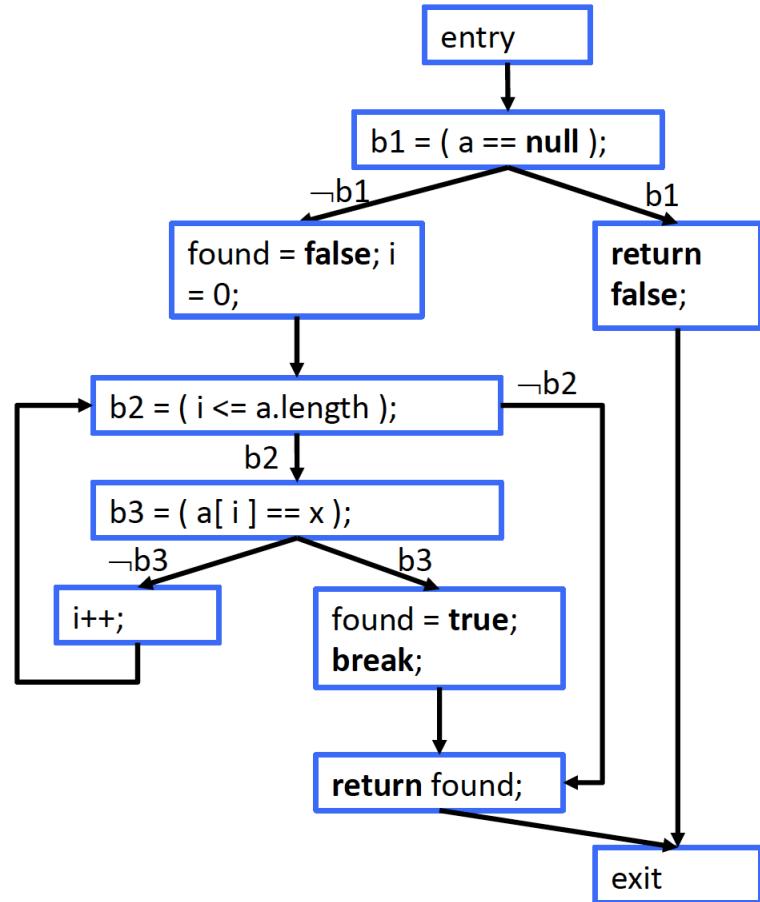
Path Coverage: Example 1

- We can achieve 100% path coverage with four tests
 - a = **true**, b = **false**
 - a = **false**, b = **true**
 - a = **true**, b = **true**
 - a = **false**, b = **false**
- The two new tests detect the bug



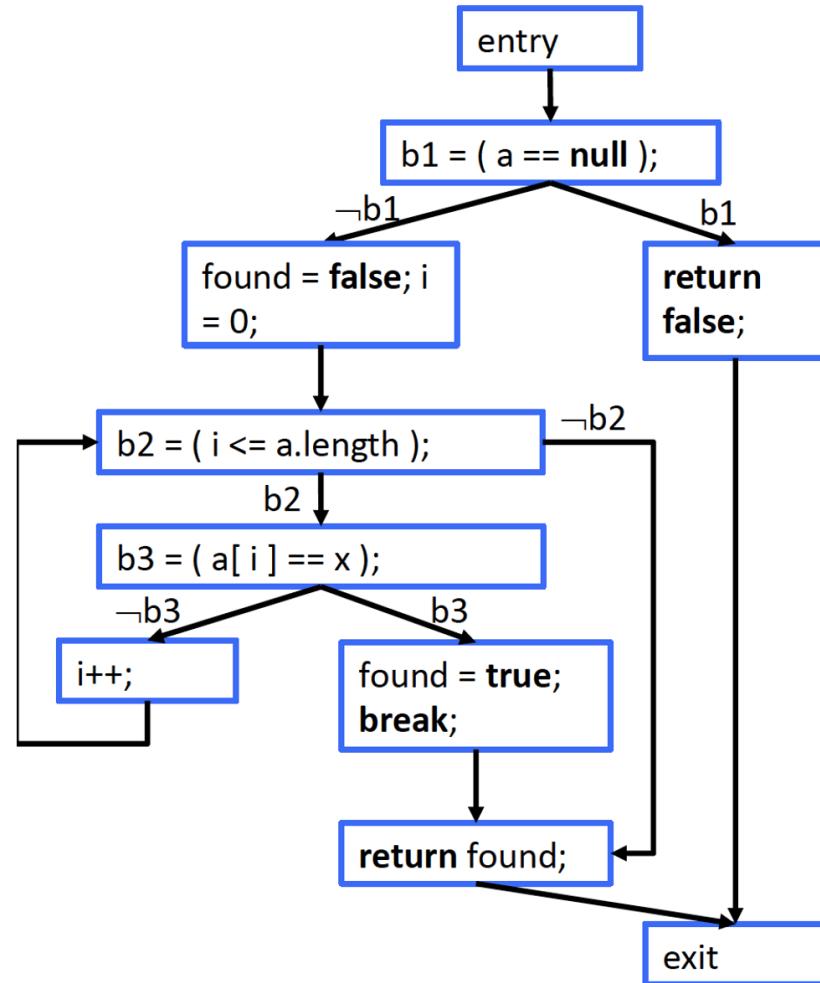
Path Coverage: Example 2

```
boolean contains( int[ ] a, int x ) {  
    if( a == null ) return false;  
  
    boolean found = false;  
  
    for( int i = 0; i <= a.length; i++ ) {  
        if( a[ i ] == x ) {  
            found = true;  
            break;  
        }  
    }  
    return found;  
}
```



Path Coverage: Example 2

- Number of loop iterations is unknown statically
- Arbitrarily many tests are needed for complete path coverage



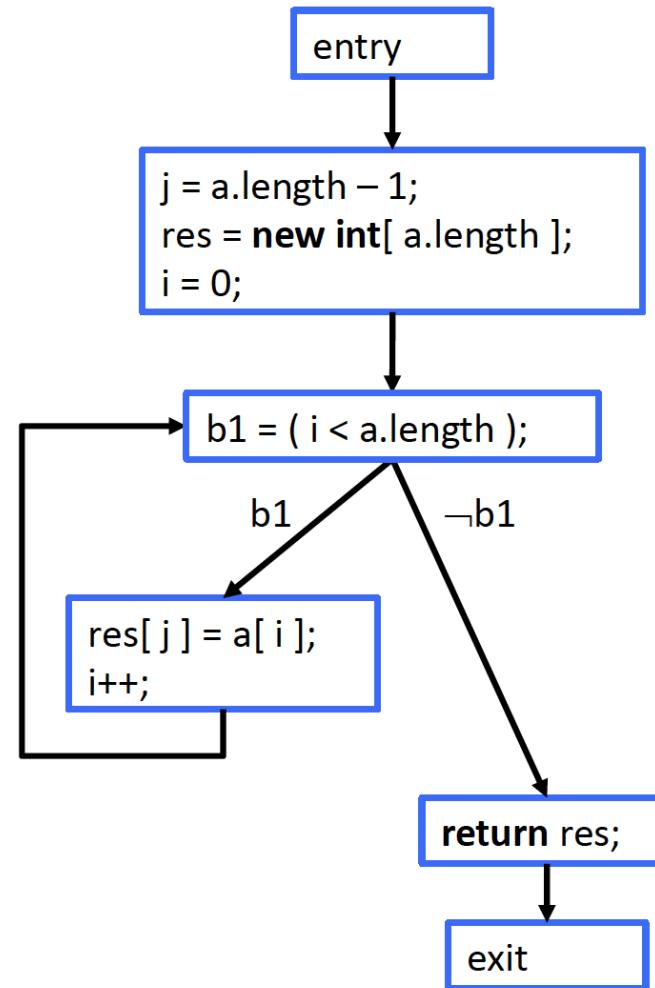
Loop Coverage

- Idea: for each loop, test 0,1,>=2 (consecutive) iterations

$$\text{Loop Coverage} = \frac{\text{Number of executed loops with 0, 1, and more than 1 iterations}}{\text{Total number of loops} * 3}$$

Loop Coverage: Example

- We can achieve 100% loop coverage with 3 tests
 - $a = \{\}$
 - $a = \{1\}$
 - $a = \{1,2\}$
- The last detects the bug



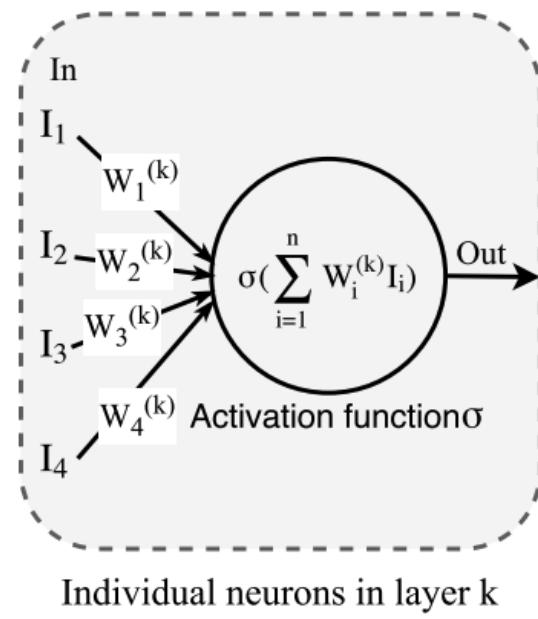
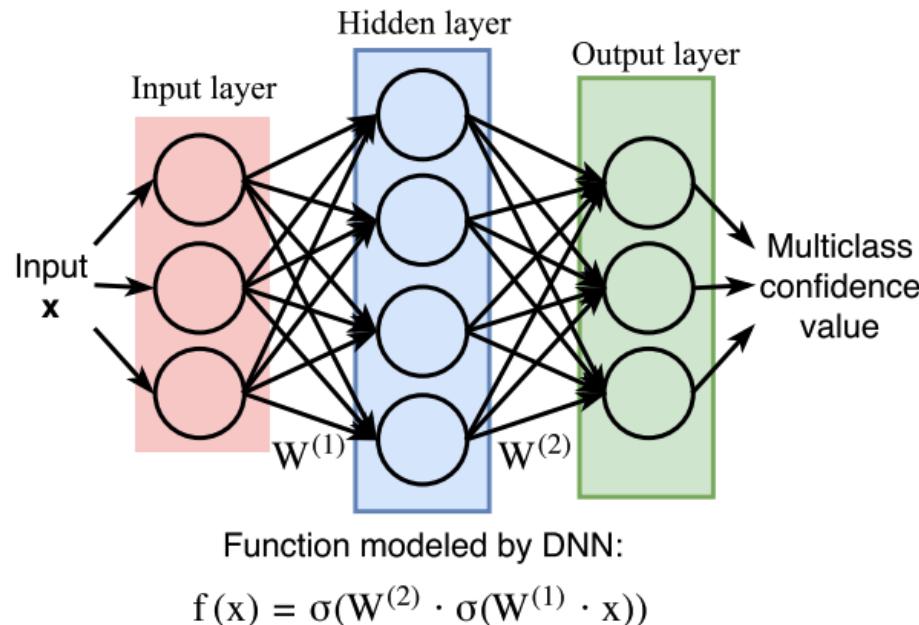
Many Others

- Condition coverage
- Multiple condition coverage
- Decision coverage
- Condition/decision coverage
- Modified condition/decision coverage

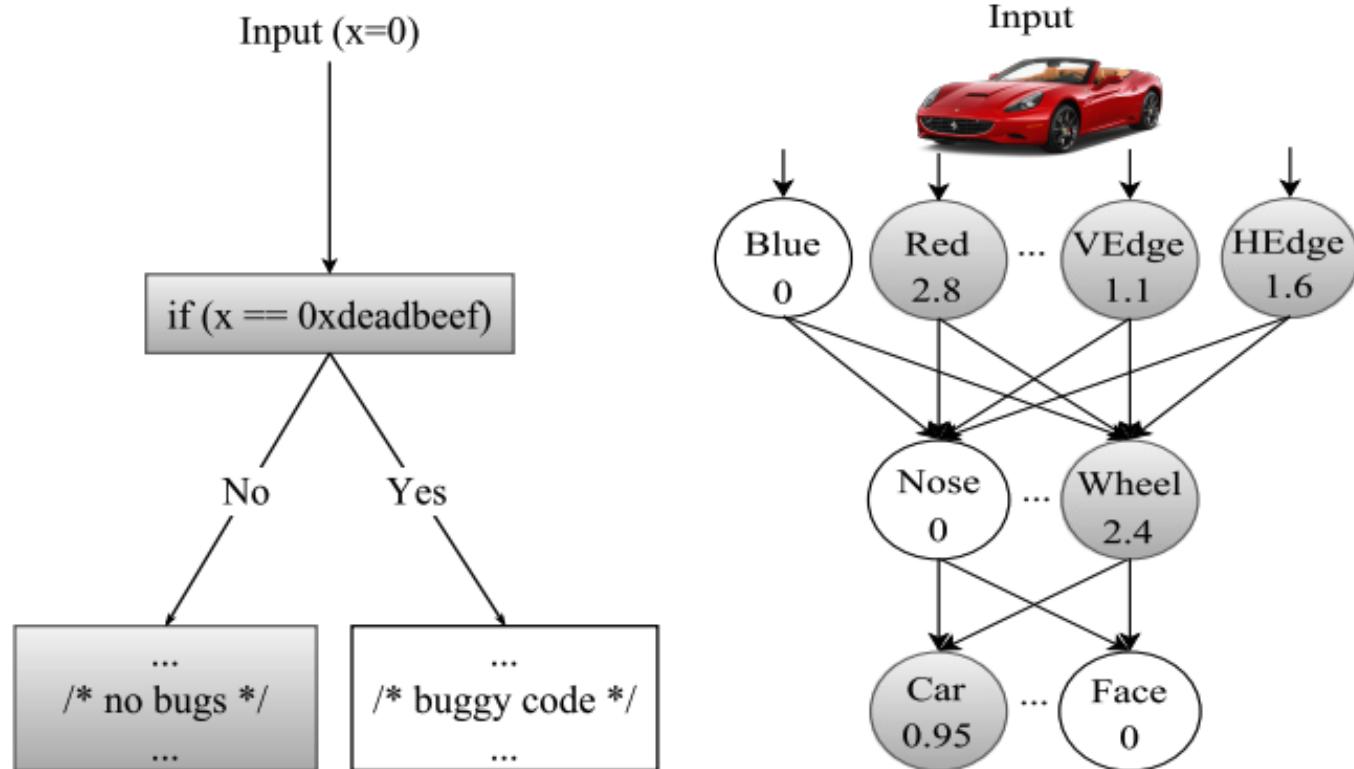
...

Coverage Metrics for DL

- How about deep neural networks?



Coverage Metrics for DL



Neuron Coverage

- How many neurons were activated

$$NCov(T, \mathbf{x}) = \frac{|\{n | \forall x \in T, out(n, \mathbf{x}) > t\}|}{|N|}$$

- N: all neurons
- T: test inputs
- out: function returns the output value of neuron
- t: threshold

Many others

- K-multisection neuron coverage
- Neuron boundary coverage
- Strong neuron activation coverage
- Top-k neuron patterns

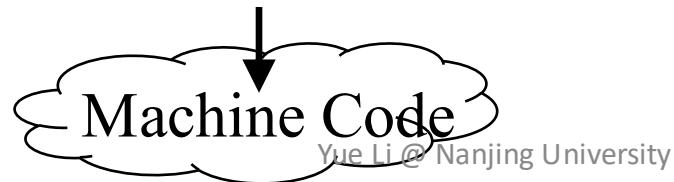
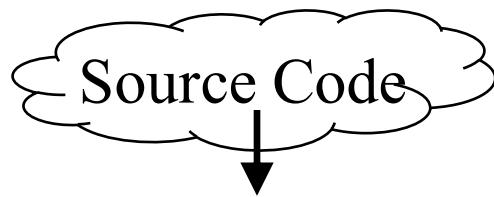
...

Useful?

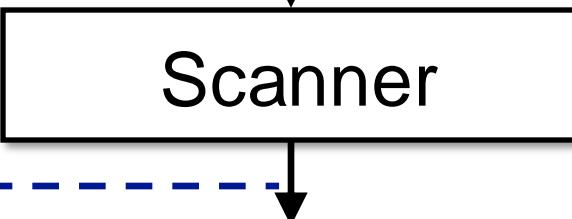
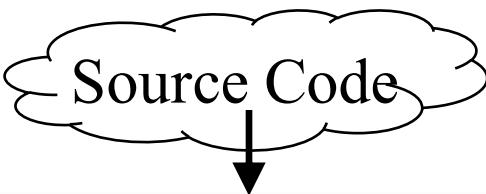
- [ESEC/FSE20] Is neuron coverage a meaningful measure for testing deep neural networks?
- [ESEC/FSE20] Correlations between deep neural network model coverage criteria and model quality

...

Compiler



Compiler

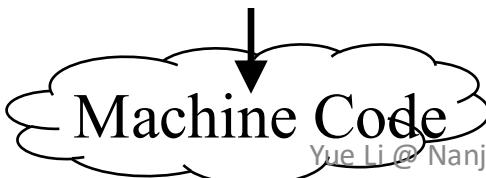


Tokens

Lexical Analysis

You ↗ goouojd

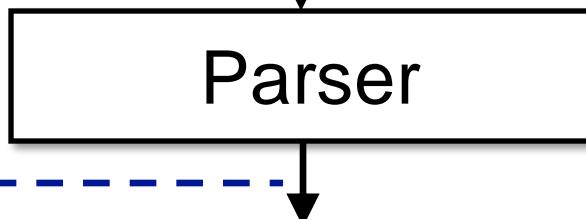
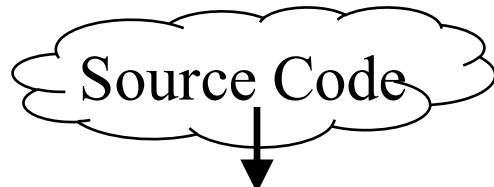
Regular
Expression



Compiler

Tokens

AST



Lexical Analysis

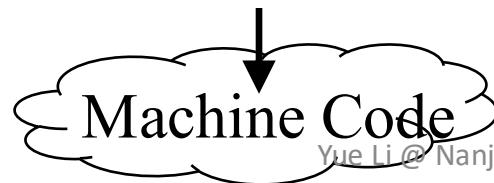
You ↗ goouojd

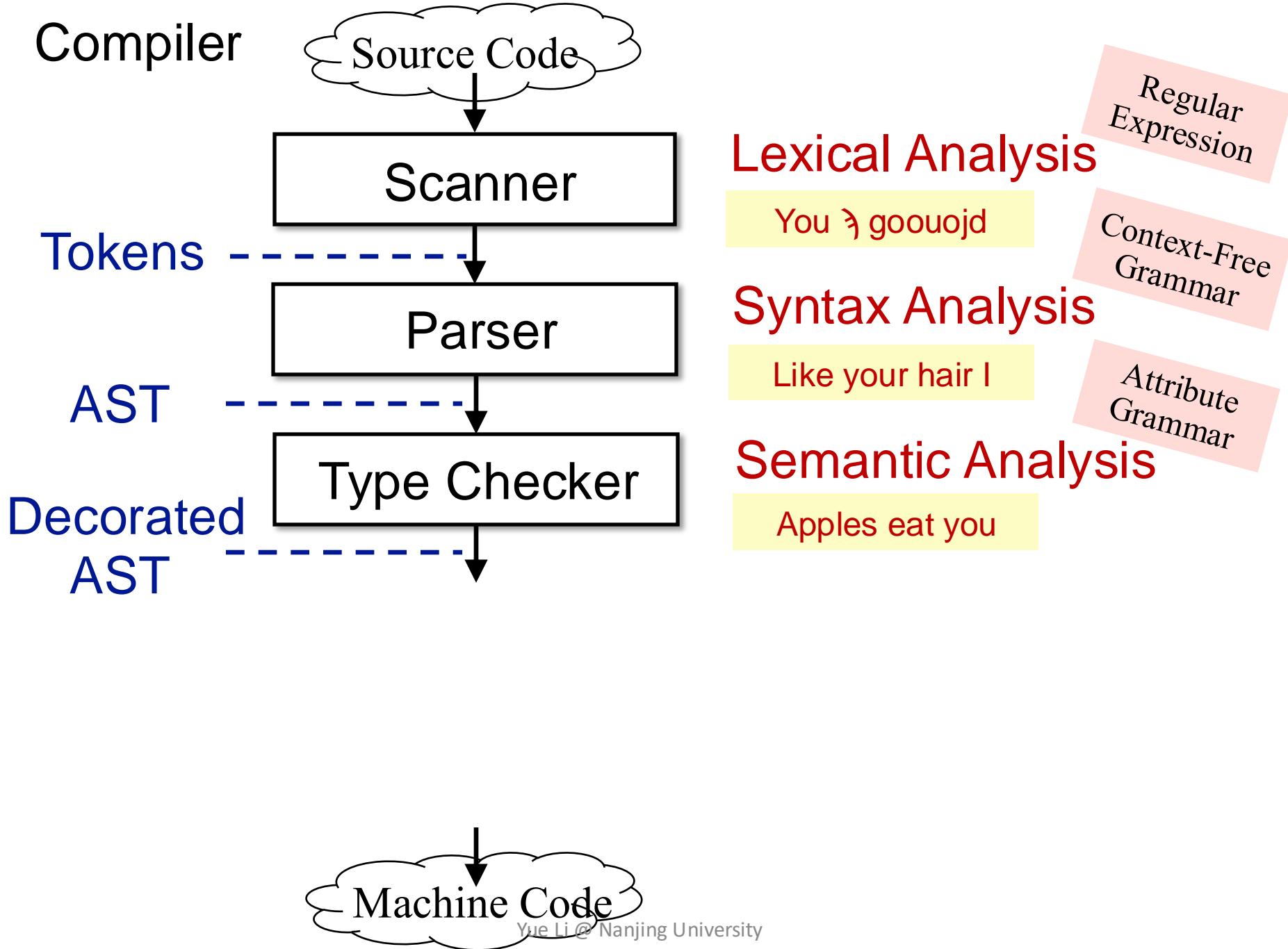
Syntax Analysis

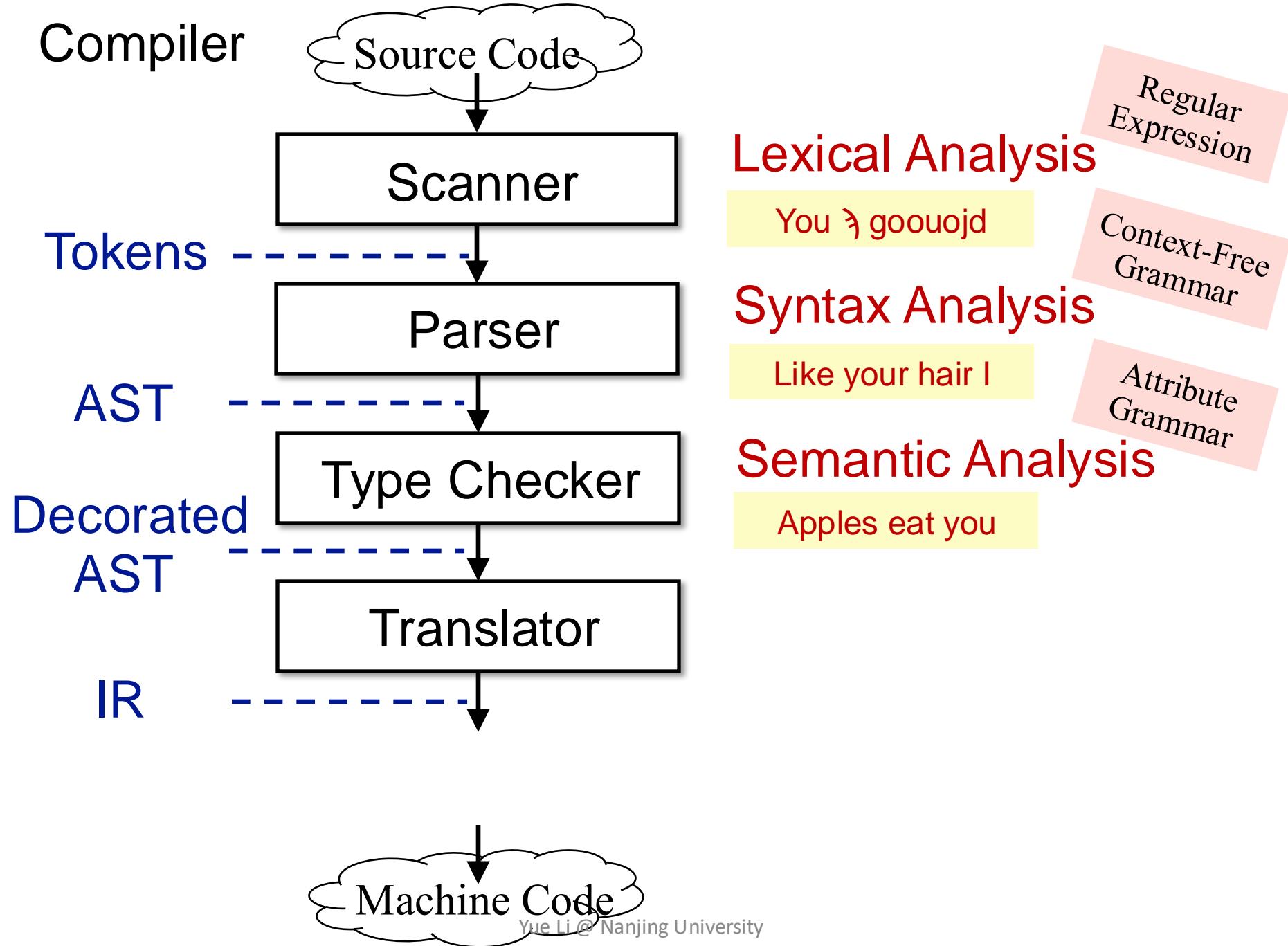
Like your hair I

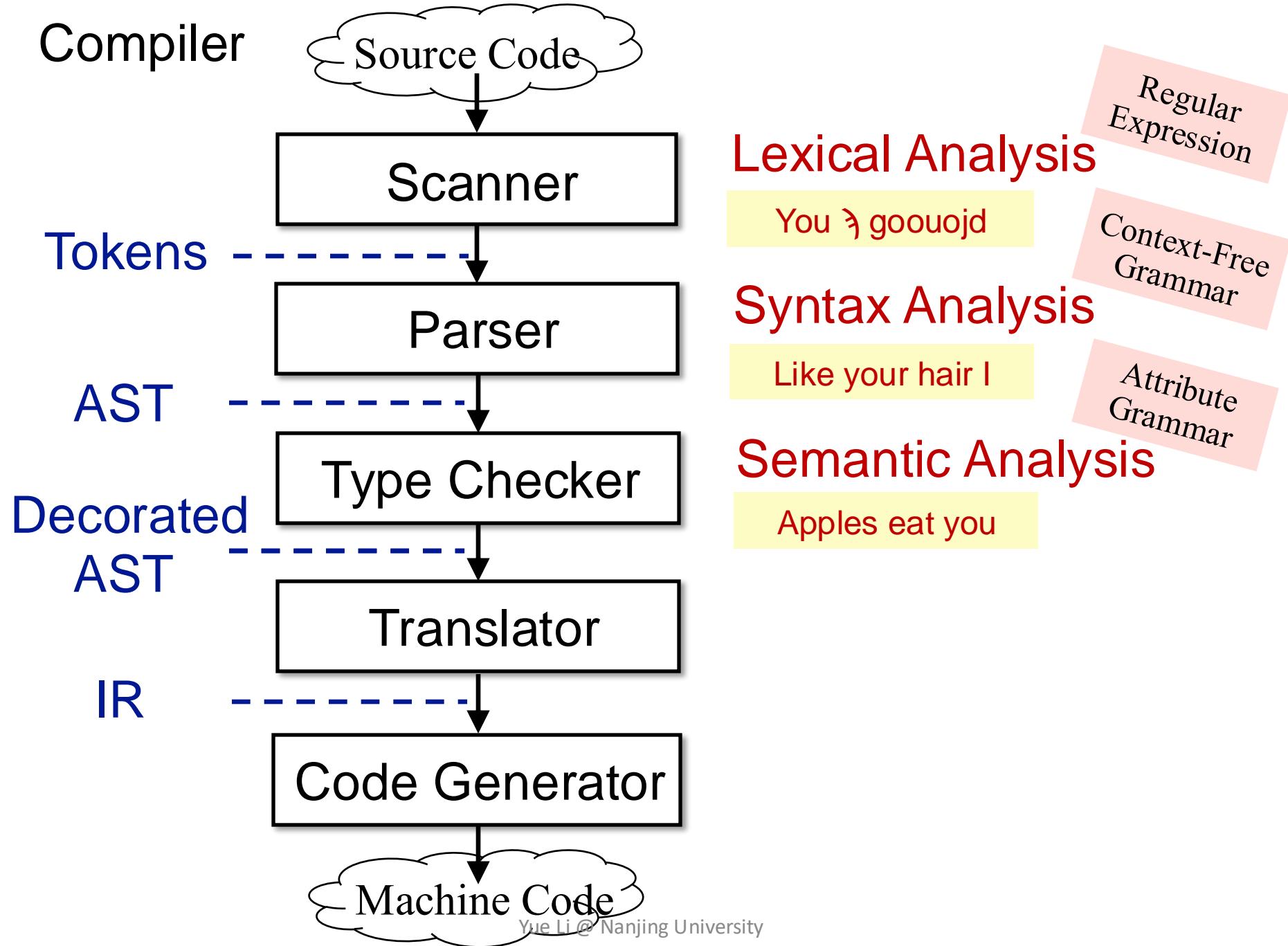
Regular
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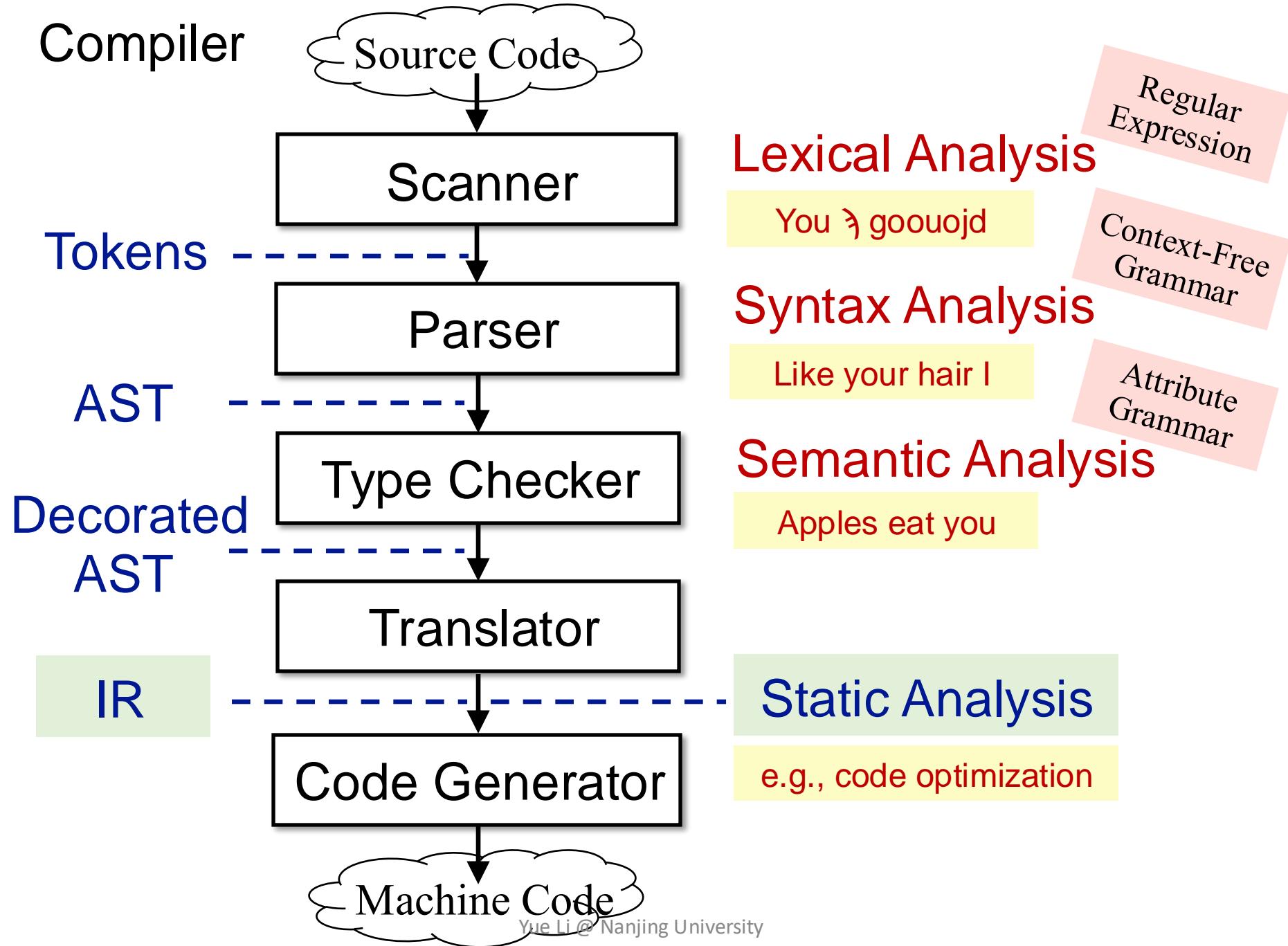
Context-Free
Grammar



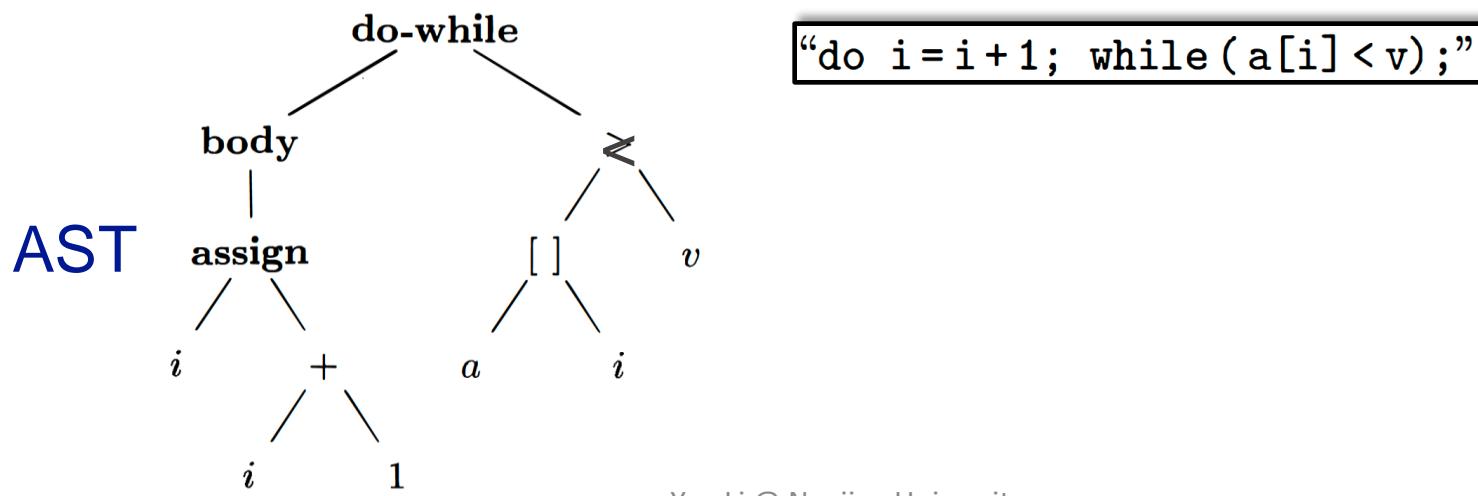






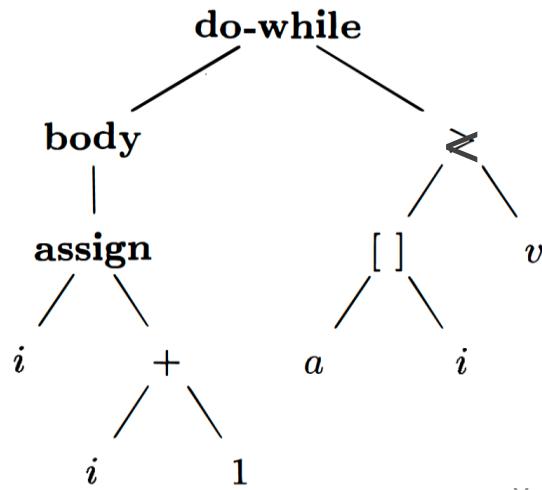


AST vs. IR



AST vs. IR

AST



“do `i = i + 1;` while (`a[i] < v`) ;”

IR

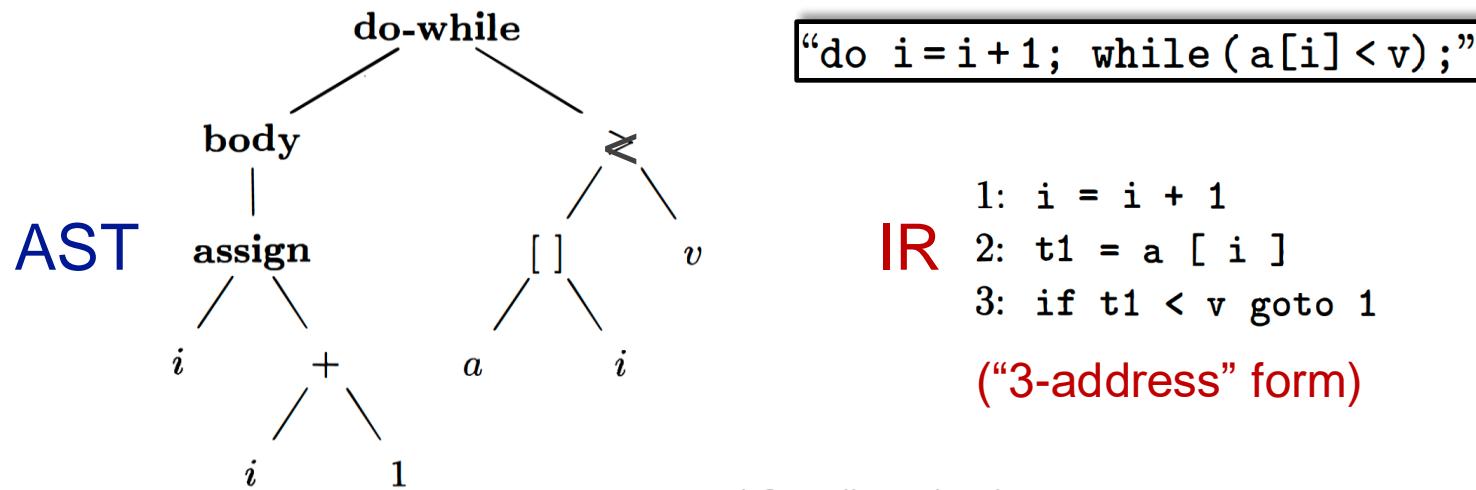
1: `i = i + 1`
2: `t1 = a [i]`
3: `if t1 < v goto 1`

(“3-address” form)

AST vs. IR

- high-level and closed to grammar structure
- usually language dependent
- suitable for fast type checking
- lack of control flow information

- low-level and closed to machine code
- usually language independent
- compact and uniform
- contains control flow information
- usually considered as the basis for static analysis



Intermediate Representation (IR)

- 3-Address Code (3AC)

There is at most one operator on the right side of an instruction.

$$c = a + b + 3 \Rightarrow \begin{array}{l} t1 = a + b \\ c = t1 + 3 \end{array}$$

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- Name: a, b, c
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Address can be one of the following:

- Name: a, b, c
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- Compiler-generated temporary: t1

Each type of instructions has its own 3AC form

Some Common 3AC Forms

- $x = y \ bop\ z$
- $x = uop\ y$
- $x = y$
- goto L
- if x goto L
- if x rop y goto L

x, y, z: addresses

bop: binary arithmetic or logical operation

uop: unary operation (minus, negation, casting)

L: a label to represent a program location

rop: relational operator (>, <, ==, >=, <=, etc.)

goto L: unconditional jump

if ... goto L: conditional jump

Some Common 3AC Forms

- $x = y \ bop\ z$
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uop: unary operation (minus, negation, casting)
- $x = y$
L: a label to represent a program location
- goto *L*
rop: relational operator ($>$, $<$, $==$, \geq , \leq , etc.)
goto *L*: unconditional jump
- if *x* goto *L*
if ... goto *L*: conditional jump
- if *x* *rop* *y* goto *L*

Let's see some more real-world complicated forms

Soot and Its IR: Jimple

- Soot

Most popular static analysis framework for Java

<https://github.com/Sable/soot>

<https://github.com/Sable/soot/wiki/Tutorials>

Soot's IR is Jimple: typed 3-address code

```
package nju.sa.examples;
public class DoWhileLoop3AC {
    public static void main(String[] args) {
        int[] arr = new int[10];
        int i = 0;
        do {
            i = i + 1;
        } while (arr[i] < 10);
    }
}
```

Java Src



Do-While Loop

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package nju.sa.examples;
public class DoWhileLoop3AC {
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        int[] arr = new int[10];
        int i = 0;
        do {
            i = i + 1;
        } while (arr[i] < 10);
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}
```

Java Src

Do-While Loop

```
public static void main(java.lang.String[])
{
    java.lang.String[] r0;
    int[] r1;
    int $i0, i1;

    r0 := @parameter0: java.lang.String[];
    r1 = newarray (int)[10];
    i1 = 0;

    label1:
    i1 = i1 + 1;
    $i0 = r1[i1];
    if $i0 < 10 goto label1;
}

return;
}
```

Yue Li @ Nanjing University

3AC (jimple)

```
package nju.sa.examples;
public class MethodCall3AC {

    String foo(String para1, String para2) {
        return para1 + " " + para2;
    }

    public static void main(String[] args) {
        MethodCall3AC mc = new MethodCall3AC();
        String result = mc.foo("hello", "world");
    }
}
```

Java Src



```

java.lang.String foo(java.lang.String, java.lang.String)
{
    nju.sa.examples.MethodCall3AC r0;
    java.lang.String r1, r2, $r7;
    java.lang.StringBuilder $r3, $r4, $r5, $r6;

    r0 := @this: nju.sa.examples.MethodCall3AC;

    r1 := @parameter0: java.lang.String;

    r2 := @parameter1: java.lang.String;

    $r3 = new java.lang.StringBuilder;

    specialinvoke $r3.<java.lang.StringBuilder: void <init>()>();

    $r4 = virtualinvoke $r3.<java.lang.StringBuilder: java.lang.StringBuilder append(java.lang.String)>(r1);

    $r5 = virtualinvoke $r4.<java.lang.StringBuilder: java.lang.StringBuilder append(java.lang.String)>(" ");

    $r6 = virtualinvoke $r5.<java.lang.StringBuilder: java.lang.StringBuilder append(java.lang.String)>(r2);

    $r7 = virtualinvoke $r6.<java.lang.StringBuilder: java.lang.String toString()>();

    return $r7;
}

```

```

package nju.sa.examples;
public class MethodCall3AC {

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Java Src



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    }
}
```

Java Src



```
public static void main(java.lang.String[])
{
    java.lang.String[] r0;
    nju.sa.examples.MethodCall3AC $r3;

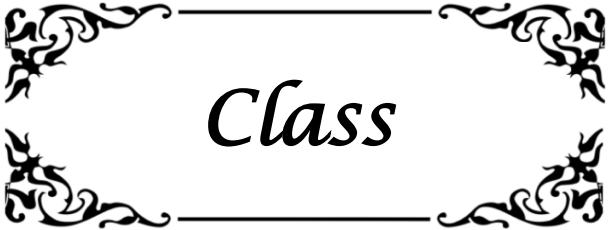
    r0 := @parameter0: java.lang.String[];

    $r3 = new nju.sa.examples.MethodCall3AC;

    specialinvoke $r3.<nju.sa.examples.MethodCall3AC: void <init>()>();

    virtualinvoke $r3.<nju.sa.examples.MethodCall3AC:
        java.lang.String foo(java.lang.String,java.lang.String)>("hello", "world");

    return;
```



Class

```
package nju.sa.examples;
public class Class3AC {

    public static final double pi = 3.14;
    public static void main(String[] args) {

    }
}
```

Java Src

```
public class nju.sa.examples.Class3AC extends java.lang.Object
{
    public static final double pi;

    public void <init>()
    {
        nju.sa.examples.Class3AC r0;

        r0 := @this: nju.sa.examples.Class3AC;

        specialinvoke r0.<java.lang.Object: void <init>()>();

        return;
    }

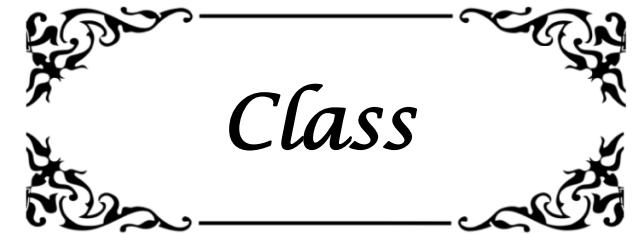
    public static void main(java.lang.String[])
    {
        java.lang.String[] r0;

        r0 := @parameter0: java.lang.String[];

        return;
    }

    public static void <clinit>()
    {
        <nju.sa.examples.Class3AC: double pi> = 3.14;

        return;
    }
}
```



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package nju.sa.examples;
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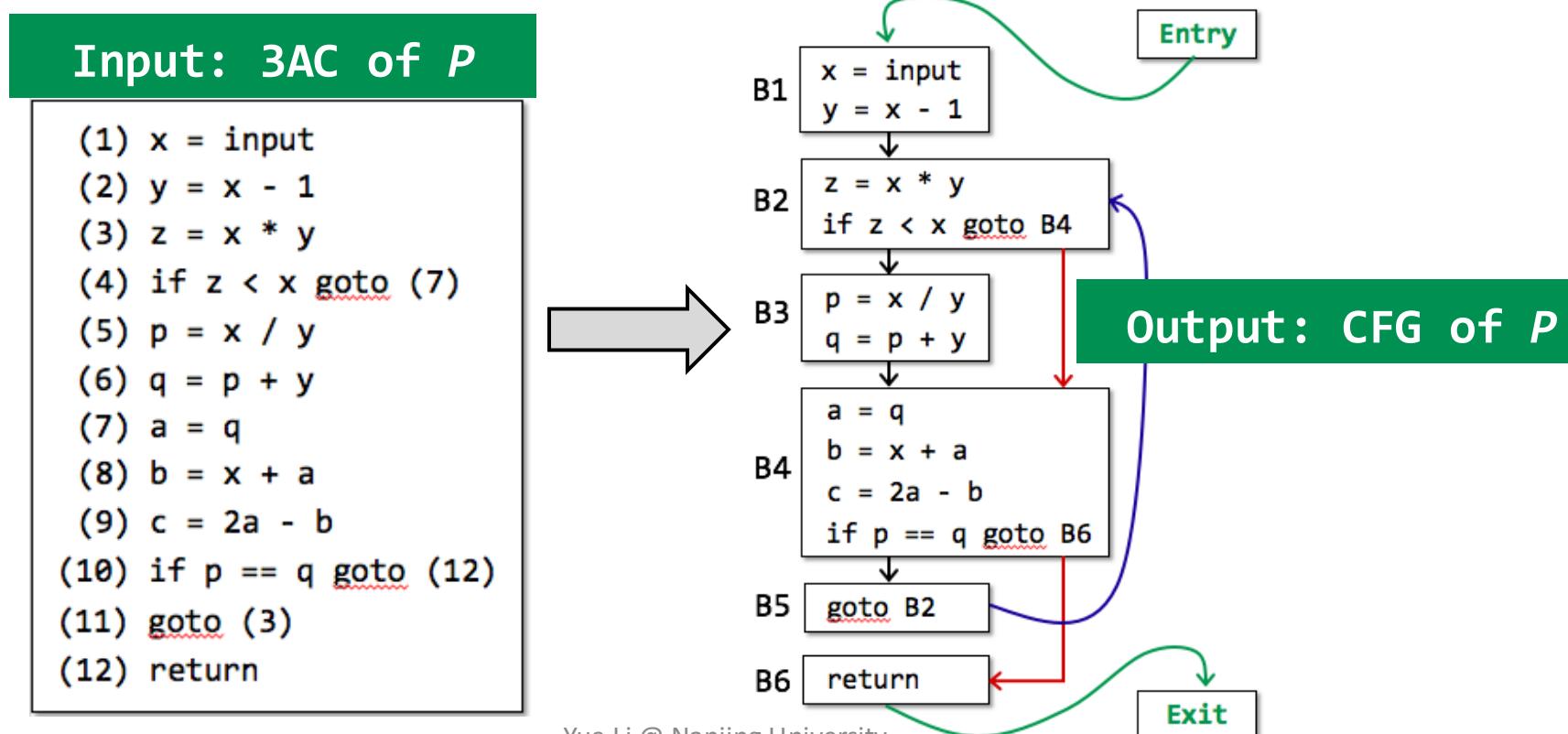
Java Src

Control Flow Analysis

- Usually refer to building Control Flow Graph (CFG)

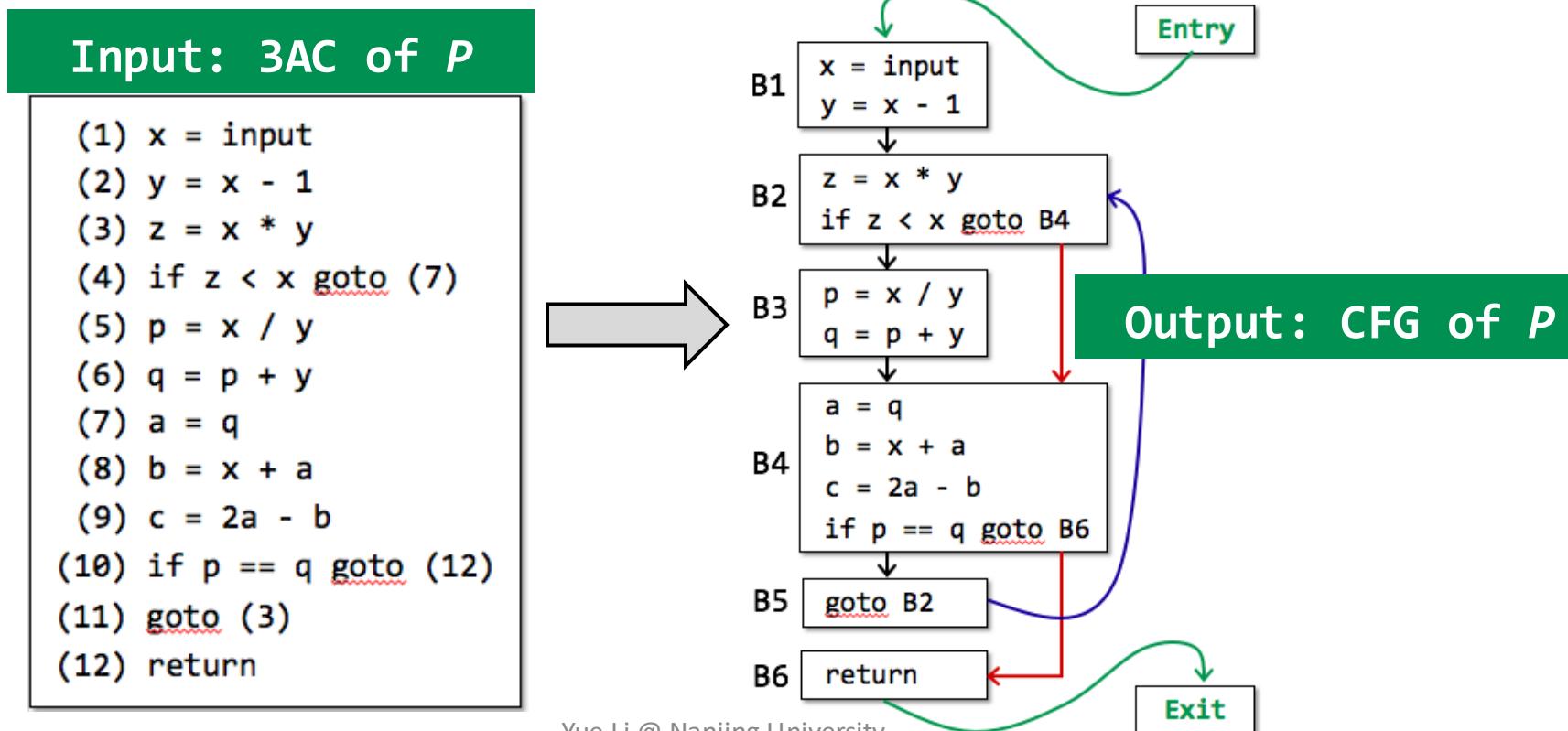
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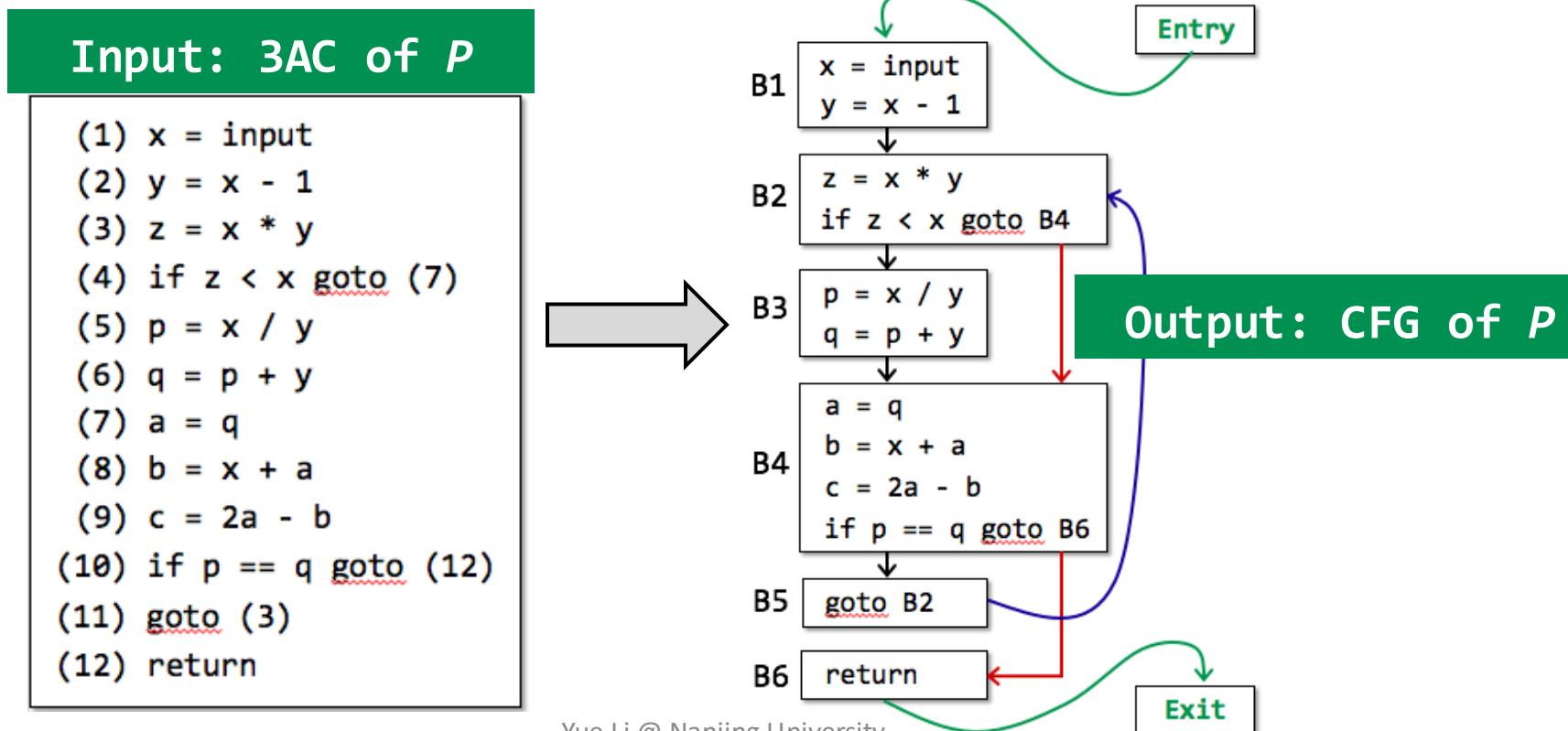
Control Flow Analysis

- Usually refer to building Control Flow Graph (CFG)
- CFG serves as the basic structure for static analysis



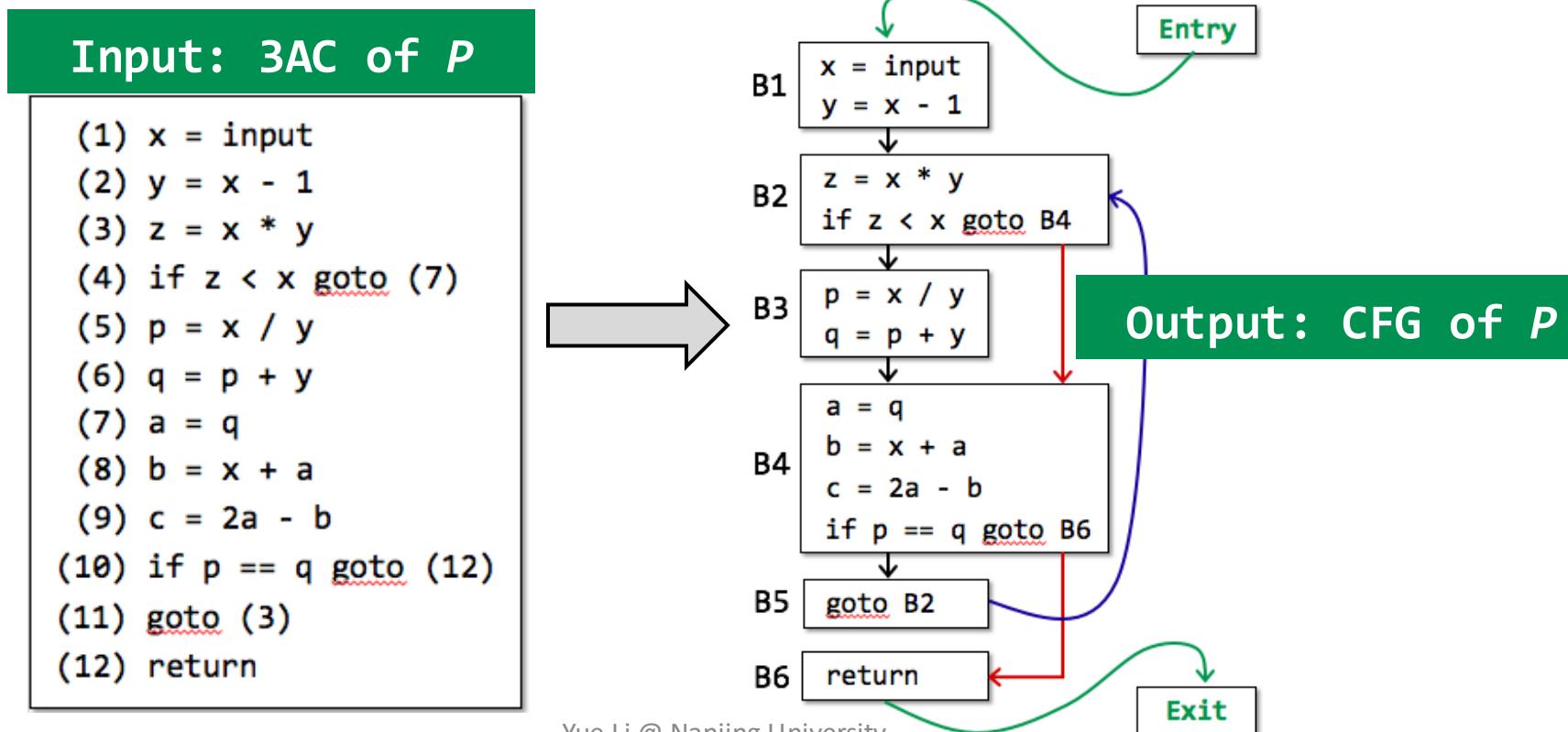
Control Flow Analysis

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- The node in CFG can be an individual 3-address instruction, or (usually) a Basic Block (BB)



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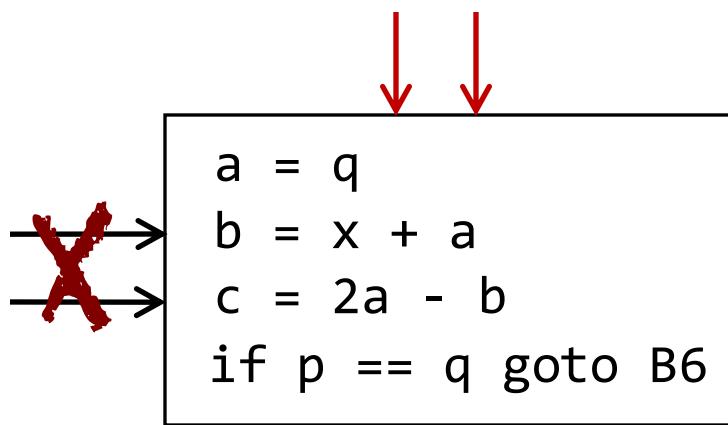
Basic Blocks (BB)

- Basic blocks (BB) are maximal sequences of consecutive three-address instructions with the properties that

```
a = q  
b = x + a  
c = 2a - b  
if p == q goto B6
```

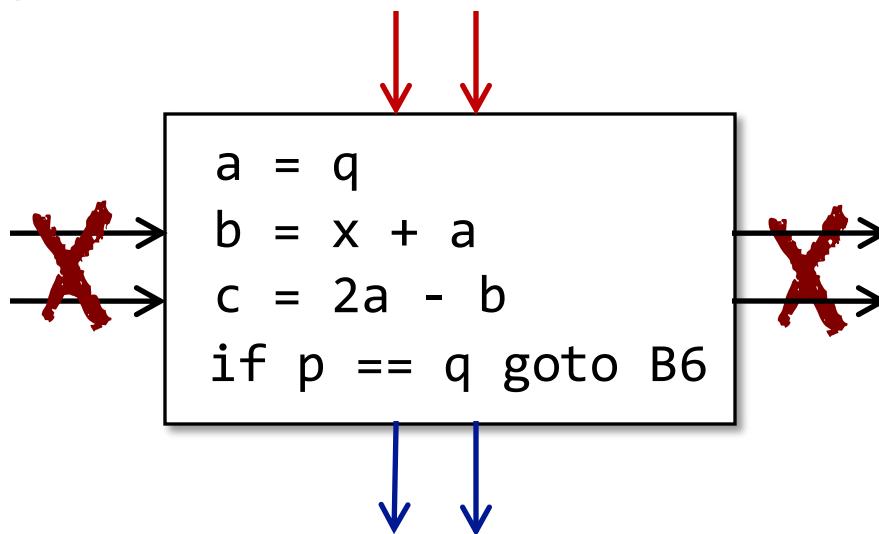
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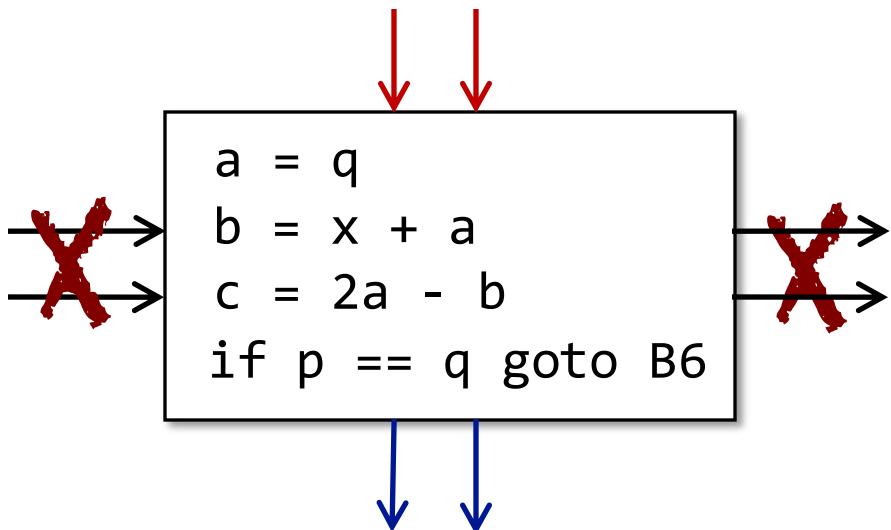
Basic Blocks (BB)

- Basic blocks (BB) are maximal sequences of consecutive three-address instructions with the properties that
 - It can be entered only at the beginning, i.e., *the first instruction in the block*
 - It can be exited only at the end, i.e., *the last instruction in the block*



```
(1) x = input  
(2) y = x - 1  
(3) z = x * y  
(4) if z < x goto (7)  
(5) p = x / y  
(6) q = p + y  
(7) a = q  
(8) b = x + a  
(9) c = 2a - b  
(10) if p == q goto (12)  
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```

Now try to design the algorithm to build BBs by yourself!

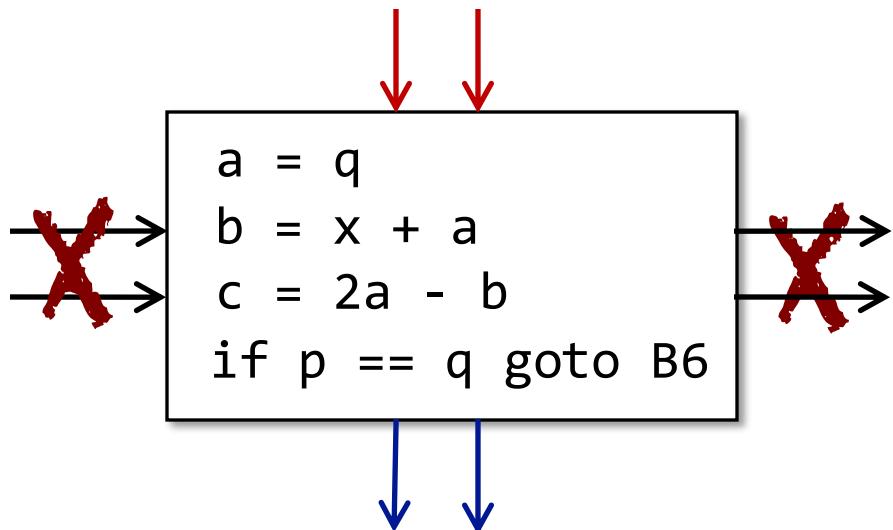


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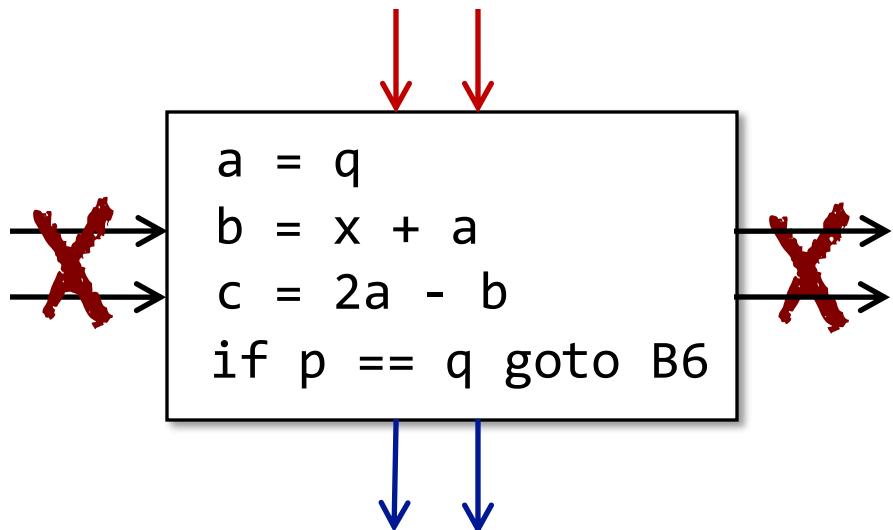


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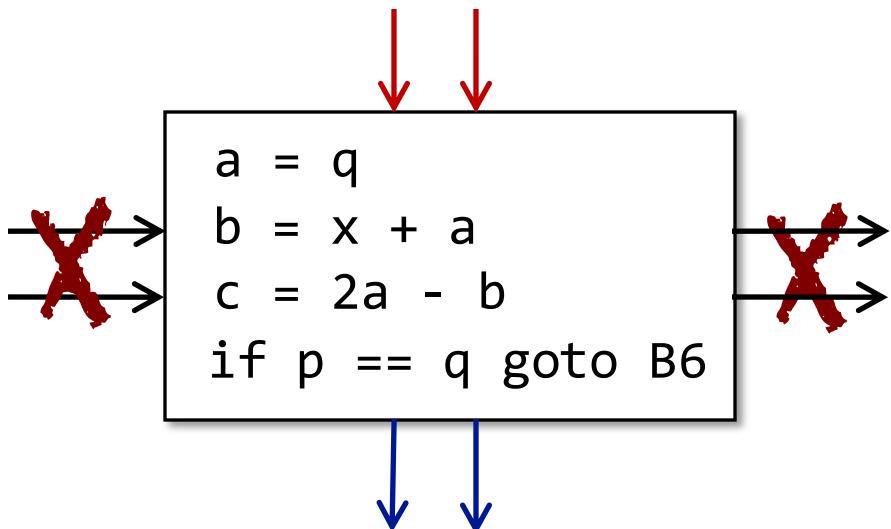
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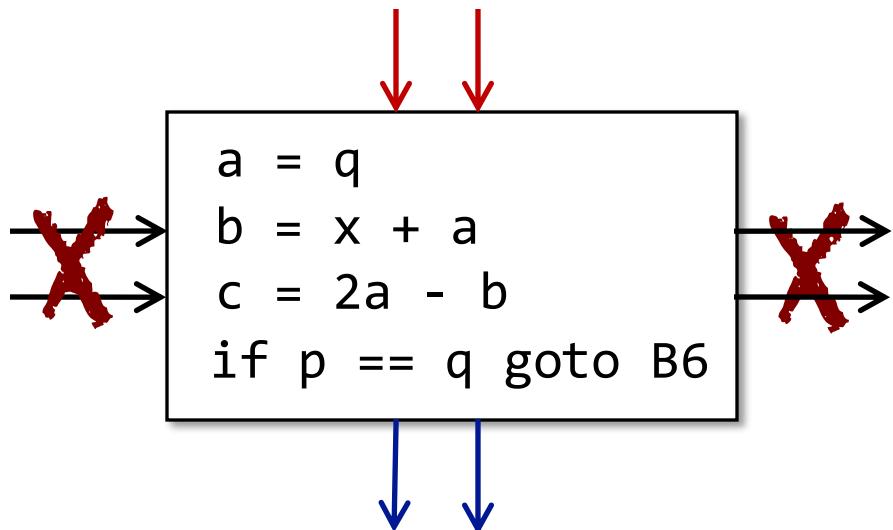


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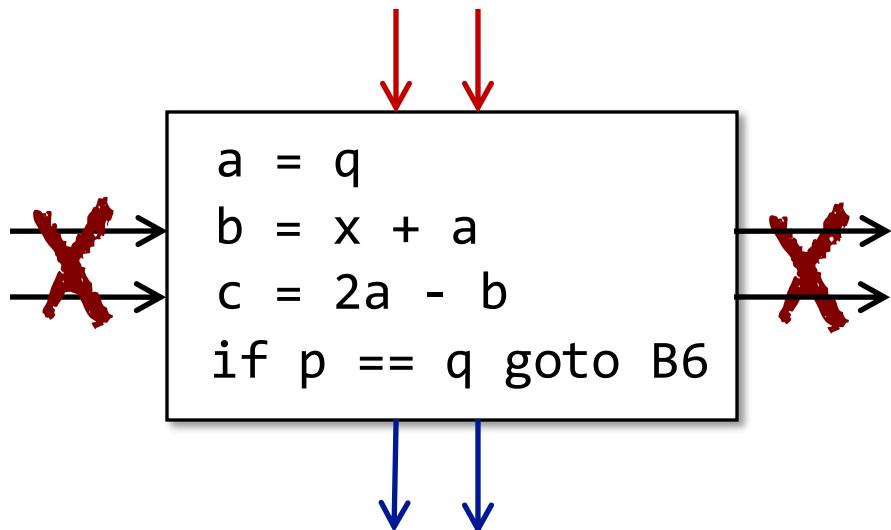


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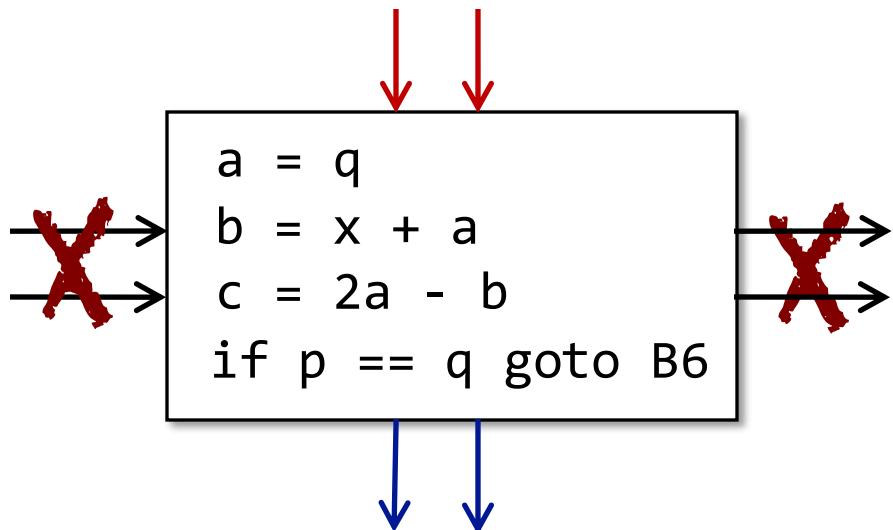


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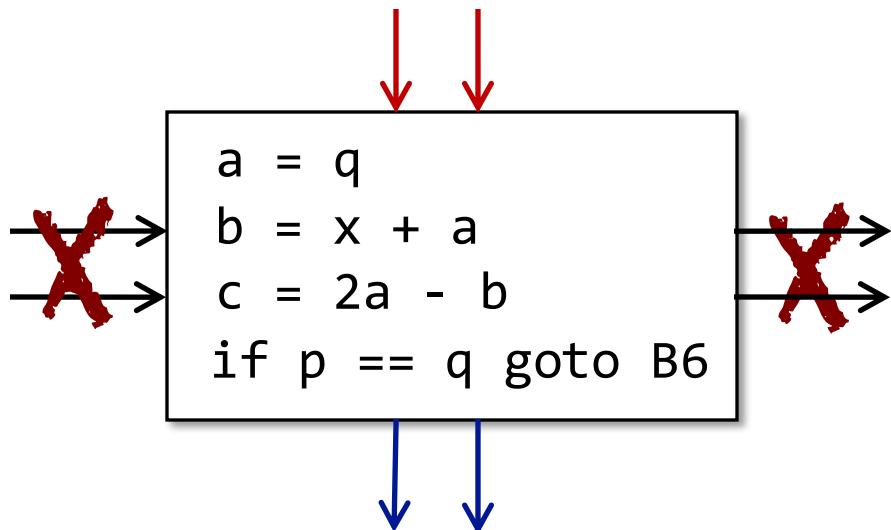


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How to build Basic Blocks?

INPUT: A sequence of three-address instructions of P

OUTPUT: A list of basic blocks of P

METHOD: (1) Determine the leaders in P

- The first instruction in P is a leader
- Any target instruction of a conditional or unconditional jump is a leader
- Any instruction that immediately follows a conditional or unconditional jump is a leader

(2) Build BBs for P

- A BB consists of a leader and all its subsequent instructions until the next leader

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(5) p = x / y  
(6) q = p + y  
(7) a = q  
(8) b = x + a  
(9) c = 2a - b  
(10) if p == q goto (12)  
(11) goto (3)  
(12) return
```

Output: BBs of P

(1) Determine the leaders in P

- (1)
- Any target instruction of a conditional or unconditional jump is a leader

Input: BAC of P

```
(1) x = input  
(2) y = x - 1  
(3) z = x * y  
(4) if z < x goto (7)  
(5) p = x / y  
(6) q = p + y  
(7) a = q  
(8) b = x + a  
(9) c = 2a - b  
(10) if p == q goto (12)  
(11) goto (3)  
(12) return
```

Output: BBs of P

(1) Determine the leaders in P

- (1)
- (3), (7), (12)

Input: BAC of P

```
(1) x = input  
(2) y = x - 1  
(3) z = x * y  
(4) if z < x goto (7)  
(5) p = x / y  
(6) q = p + y  
(7) a = q  
(8) b = x + a  
(9) c = 2a - b  
(10) if p == q goto (12)  
(11) goto (3)  
(12) return
```

Output: BBs of P

(1) Determine the leaders in P

- (1)
- (3), (7), (12)
- Any instruction that immediately follows a conditional or unconditional jump is a leader

Input: BAC of P

```
(1) x = input  
(2) y = x - 1  
(3) z = x * y  
(4) if z < x goto (7)  
(5) p = x / y  
(6) q = p + y  
(7) a = q  
(8) b = x + a  
(9) c = 2a - b  
(10) if p == q goto (12)  
(11) goto (3)  
(12) return
```

Output: BBs of P

(1) Determine the leaders in P

- (1)
- (3), (7), (12)
- Any instruction that immediately follows a conditional or unconditional jump is a leader

Input: BAC of P

```
(1) x = input  
(2) y = x - 1  
(3) z = x * y  
(4) if z < x goto (7)  
(5) p = x / y  
(6) q = p + y  
(7) a = q  
(8) b = x + a  
(9) c = 2a - b  
(10) if p == q goto (12)  
(11) goto (3)  
(12) return
```

Output: BBs of P

(1) Determine the leaders in P

- (1)
- (3), (7), (12)
- (5), (11), (12)

Input: BAC of P

```
(1) x = input  
(2) y = x - 1  
(3) z = x * y  
(4) if z < x goto (7)  
(5) p = x / y  
(6) q = p + y  
(7) a = q  
(8) b = x + a  
(9) c = 2a - b  
(10) if p == q goto (12)  
(11) goto (3)  
(12) return
```

Output: BBs of P

(1) Determine the leaders in P

- (1) Leaders: (1), (3), (5), (7), (11), (12)
- (3), (7), (12)
- (5), (11), (12)

Input: 3AC of P

```
(1) x = input  
(2) y = x - 1  
(3) z = x * y  
(4) if z < x goto (7)  
(5) p = x / y  
(6) q = p + y  
(7) a = q  
(8) b = x + a  
(9) c = 2a - b  
(10) if p == q goto (12)  
(11) goto (3)  
(12) return
```

Output: BBs of P

(1) Determine the leaders in P

- (1) Leaders: (1), (3), (5), (7), (11), (12)
- (3), (7), (12)
- (5), (11), (12)

(2) Build BBs for P

- A BB consists of a leader and all its subsequent instructions until the next leader

Input: 3AC of P

```
(1) x = input  
(2) y = x - 1  
(3) z = x * y  
(4) if z < x goto (7)  
(5) p = x / y  
(6) q = p + y  
(7) a = q  
(8) b = x + a  
(9) c = 2a - b  
(10) if p == q goto (12)  
(11) goto (3)  
(12) return
```

Output: BBs of P

(1) Determine the leaders in P

- (1) Leaders: (1), (3), (5), (7), (11), (12)
- (3), (7), (12)
- (5), (11), (12)

(2) Build BBs for P

- A BB consists of a leader and all its subsequent instructions until the next leader
- B1 {(1)}
- B2 {(3)}
- B3 {(5)}
- B4 {(7)}
- B5 {(11)}
- B6 {(12)}

Input: 3AC of P

```
(1) x = input  
(2) y = x - 1  
(3) z = x * y  
(4) if z < x goto (7)  
(5) p = x / y  
(6) q = p + y  
(7) a = q  
(8) b = x + a  
(9) c = 2a - b  
(10) if p == q goto (12)  
(11) goto (3)  
(12) return
```

Output: BBs of P

(1) Determine the leaders in P

- (1) Leaders: (1), (3), (5), (7), (11), (12)
- (3), (7), (12)
- (5), (11), (12)

(2) Build BBs for P

- A BB consists of a leader and all its subsequent instructions until the next leader
- B1 {(1), (2)}
- B2 {(3), (4)}
- B3 {(5), (6)}
- B4 {(7), (8), (9), (10)}
- B5 {(11)}
- B6 {(12)}

Input: BAC of P

```
(1) x = input  
(2) y = x - 1  
(3) z = x * y  
(4) if z < x goto (7)  
(5) p = x / y  
(6) q = p + y  
(7) a = q  
(8) b = x + a  
(9) c = 2a - b  
(10) if p == q goto (12)  
(11) goto (3)  
(12) return
```

Output: BBs of P

- B1 $\{(1), (2)\}$
- B2 $\{(3), (4)\}$
- B3 $\{(5), (6)\}$
- B4 $\{(7), (8), (9), (10)\}$
- B5 $\{(11)\}$
- B6 $\{(12)\}$

Input: BAC of P

```
(1) x = input  
(2) y = x - 1  
(3) z = x * y  
(4) if z < x goto (7)  
(5) p = x / y  
(6) q = p + y  
(7) a = q  
(8) b = x + a  
(9) c = 2a - b  
(10) if p == q goto (12)  
(11) goto (3)  
(12) return
```

Output: BBs of P

B1

```
(1) x = input  
(2) y = x - 1
```

B2

```
(3) z = x * y  
(4) if z < x goto (7)
```

B3

```
(5) p = x / y  
(6) q = p + y
```

B4

```
(7) a = q  
(8) b = x + a  
(9) c = 2a - b  
(10) if p == q goto (12)
```

B5

```
(11) goto (3)
```

B6

```
(12) return
```

Input: 3AC of P

```
(1) x = input  
(2) y = x - 1  
(3) z = x * y  
(4) if z < x goto (7)  
(5) p = x / y  
(6) q = p + y  
(7) a = q  
(8) b = x + a  
(9) c = 2a - b  
(10) if p == q goto (12)  
(11) goto (3)  
(12) return
```

How to build CFG
on top of BBs?

Output: BBs of P

B1

```
(1) x = input  
(2) y = x - 1
```

B2

```
(3) z = x * y  
(4) if z < x goto (7)
```

B3

```
(5) p = x / y  
(6) q = p + y
```

B4

```
(7) a = q  
(8) b = x + a  
(9) c = 2a - b  
(10) if p == q goto (12)
```

B5

```
(11) goto (3)
```

B6

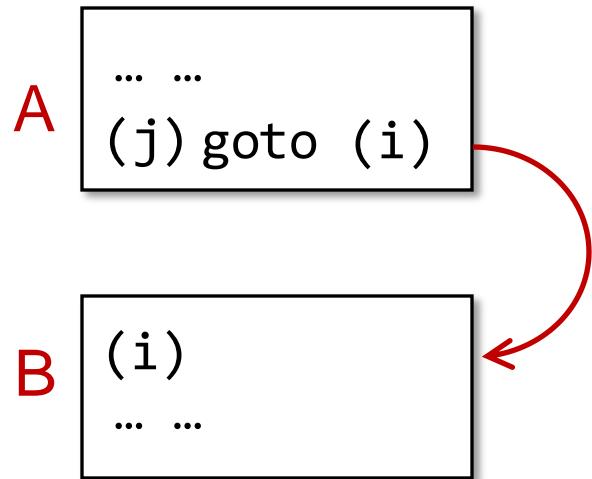
```
(12) return
```

Control Flow Graph (CFG)

- The nodes of CFG are basic blocks
- There is an edge from block A to block B if and only if

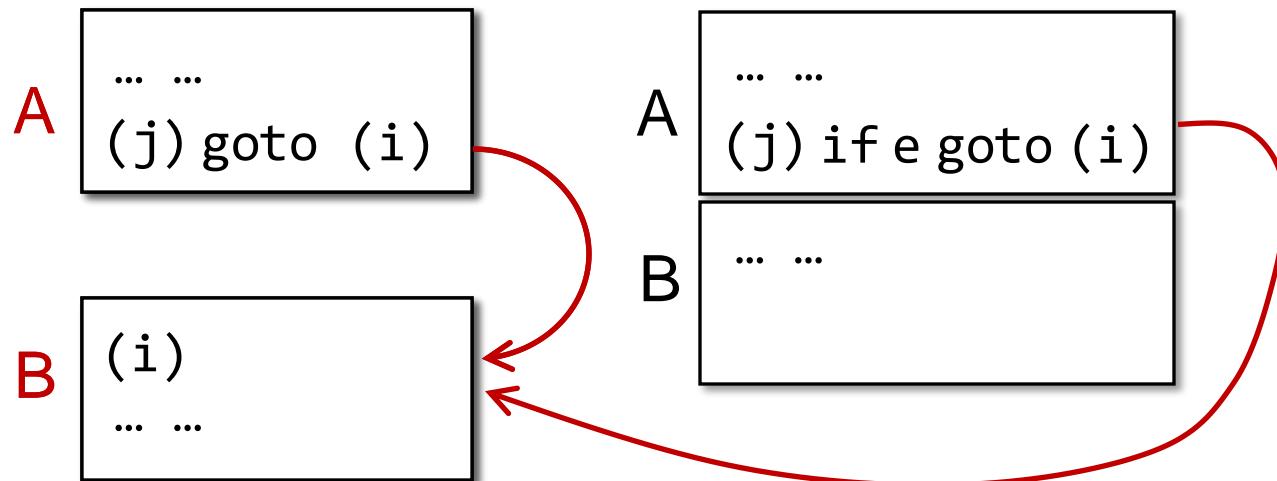
Control Flow Graph (CFG)

- The nodes of CFG are basic blocks
- There is an edge from block A to block B if and only if



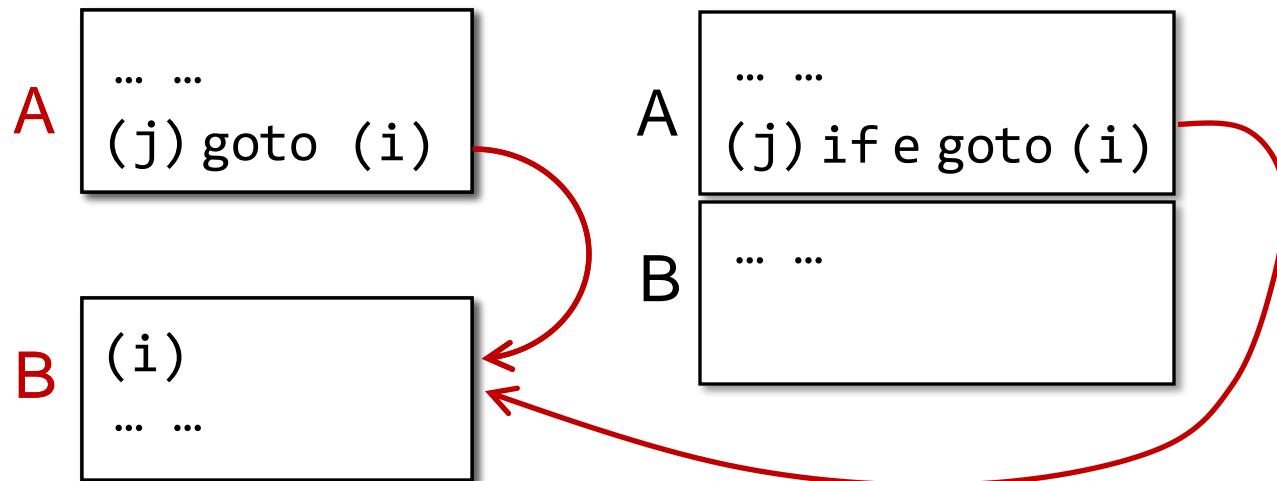
Control Flow Graph (CFG)

- The nodes of CFG are basic blocks
- There is an edge from block A to block B if and only if



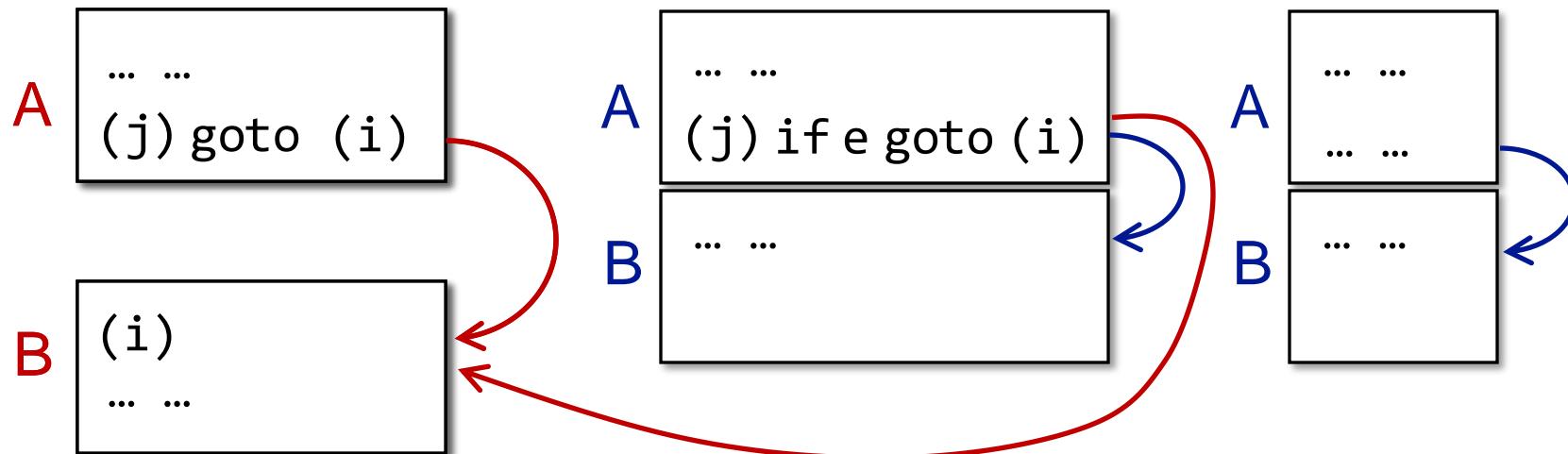
Control Flow Graph (CFG)

- The nodes of CFG are basic blocks
- There is an edge from block A to block B if and only if
 - There is a conditional or unconditional jump from the end of A to the beginning of B



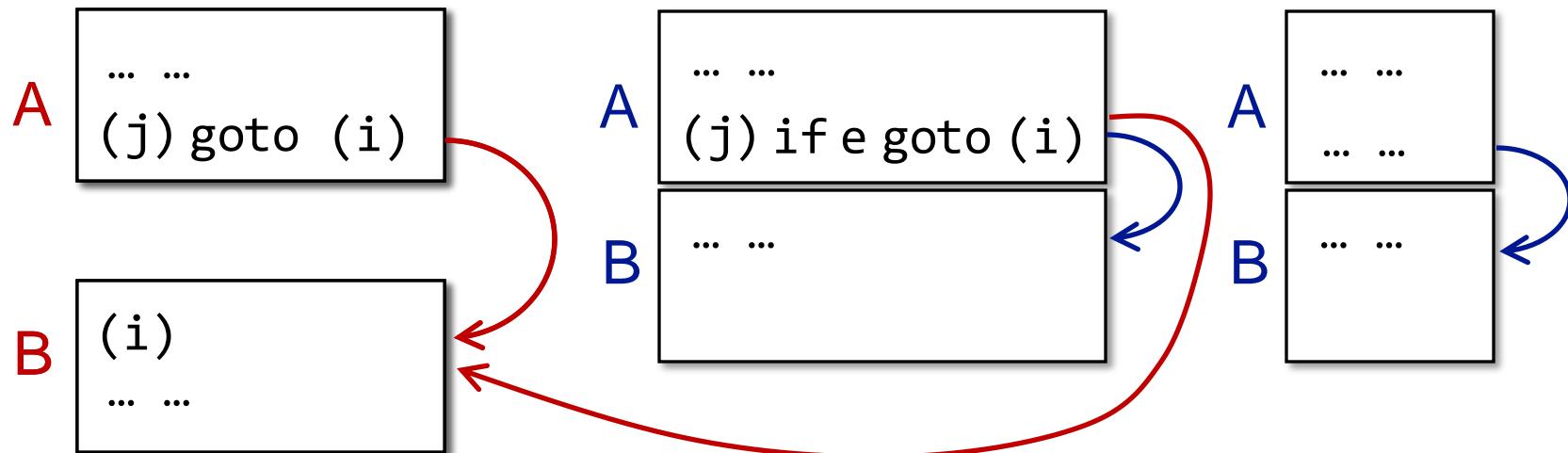
Control Flow Graph (CFG)

- The nodes of CFG are basic blocks
- There is an edge from block A to block B if and only if
 - There is a conditional or unconditional jump from the end of A to the beginning of B



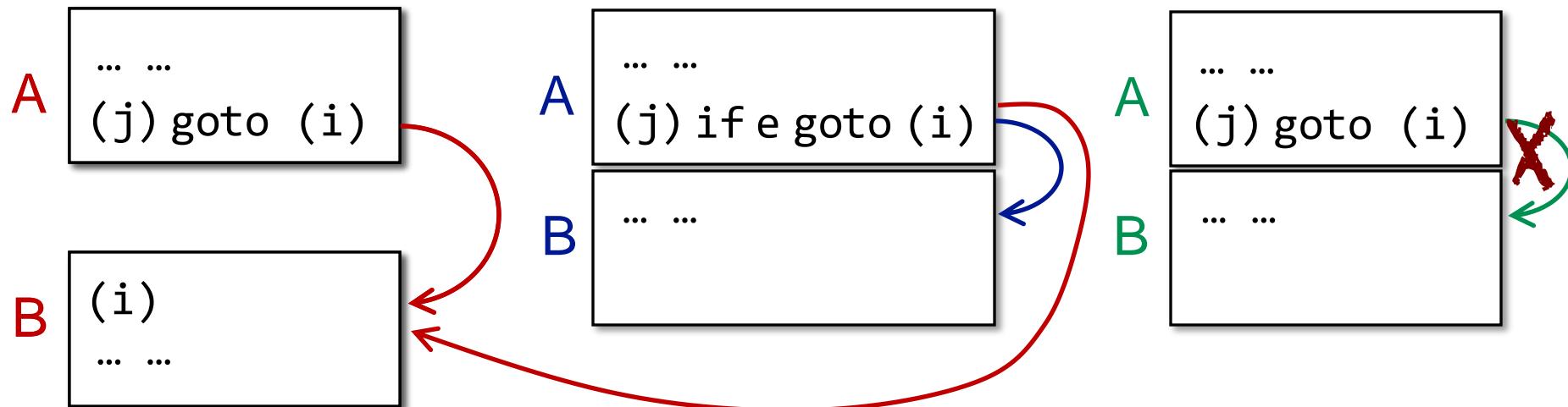
Control Flow Graph (CFG)

- The nodes of CFG are basic blocks
- There is an edge from block A to block B if and only if
 - There is a conditional or unconditional jump from the end of A to the beginning of B
 - B immediately follows A in the original order of instructions



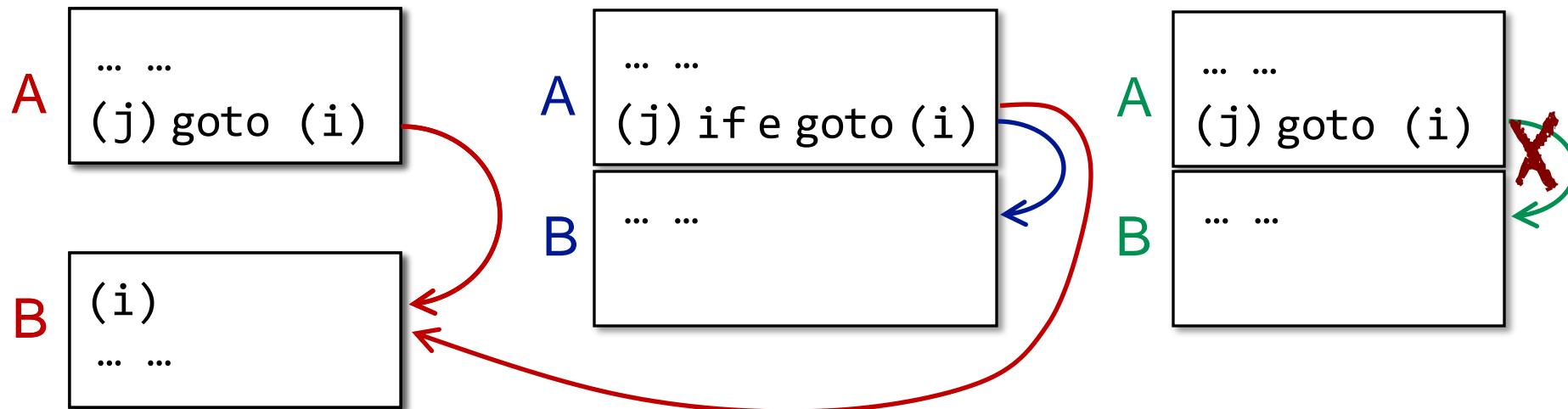
Control Flow Graph (CFG)

- The nodes of CFG are basic blocks
- There is an edge from block A to block B if and only if
 - There is a conditional or unconditional jump from the end of A to the beginning of B
 - B immediately follows A in the original order of instructions



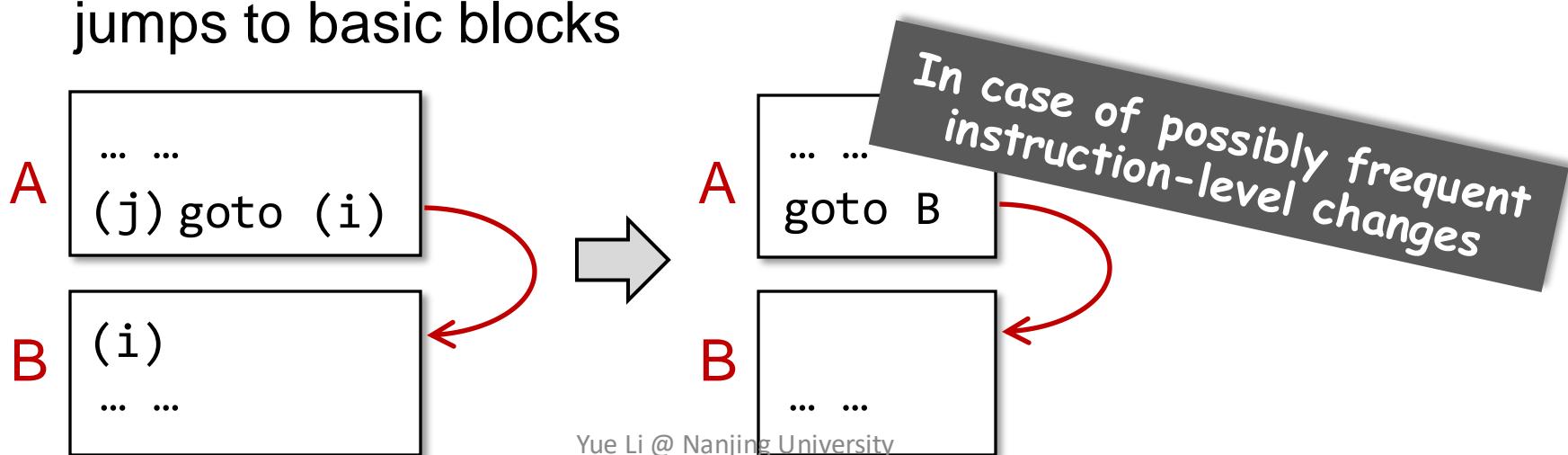
Control Flow Graph (CFG)

- The nodes of CFG are basic blocks
- There is an edge from block A to block B if and only if
 - There is a conditional or unconditional jump from the end of A to the beginning of B
 - B immediately follows A in the original order of instructions and A does not end in an unconditional jump



Control Flow Graph (CFG)

- The nodes of CFG are basic blocks
- There is an edge from block A to block B if and only if
 - There is a conditional or unconditional jump from the end of A to the beginning of B
 - B immediately follows A in the original order of instructions and A does not end in an unconditional jump
- It is normal to replace the jumps to instruction labels by jumps to basic blocks



B1

(1) $x = \text{input}$
(2) $y = x - 1$

B2

(3) $z = x * y$
(4) if $z < x$ goto (7)

B3

(5) $p = x / y$
(6) $q = p + y$

B4

(7) $a = q$
(8) $b = x + a$
(9) $c = 2a - b$
(10) if $p == q$ goto (12)

B5

(11) goto (3)

B6

(12) return

B1

$x = \text{input}$
 $y = x - 1$

B2

$z = x * y$
if $z < x$ goto B4

B3

$p = x / y$
 $q = p + y$

B4

$a = q$
 $b = x + a$
 $c = 2a - b$
if $p == q$ goto B6

B5

goto B2

B6

return

Add edges in CFG

B1

```
x = input  
y = x - 1
```

B2

```
z = x * y  
if z < x goto B4
```

B3

```
p = x / y  
q = p + y
```

B4

```
a = q  
b = x + a  
c = 2a - b  
if p == q goto B6
```

B5

```
goto B2
```

B6

```
return
```

Add edges in CFG

B1

```
x = input  
y = x - 1
```

B2

```
z = x * y  
if z < x goto B4
```

B3

```
p = x / y  
q = p + y
```

B4

```
a = q  
b = x + a  
c = 2a - b  
if p == q goto B6
```

B5

```
goto B2
```

B6

```
return
```

There is a **conditional** or **unconditional** jump from the end of A to the beginning of B

Add edges in CFG

```
B1
x = input
y = x - 1
```

```
B2
z = x * y
if z < x goto B4
```

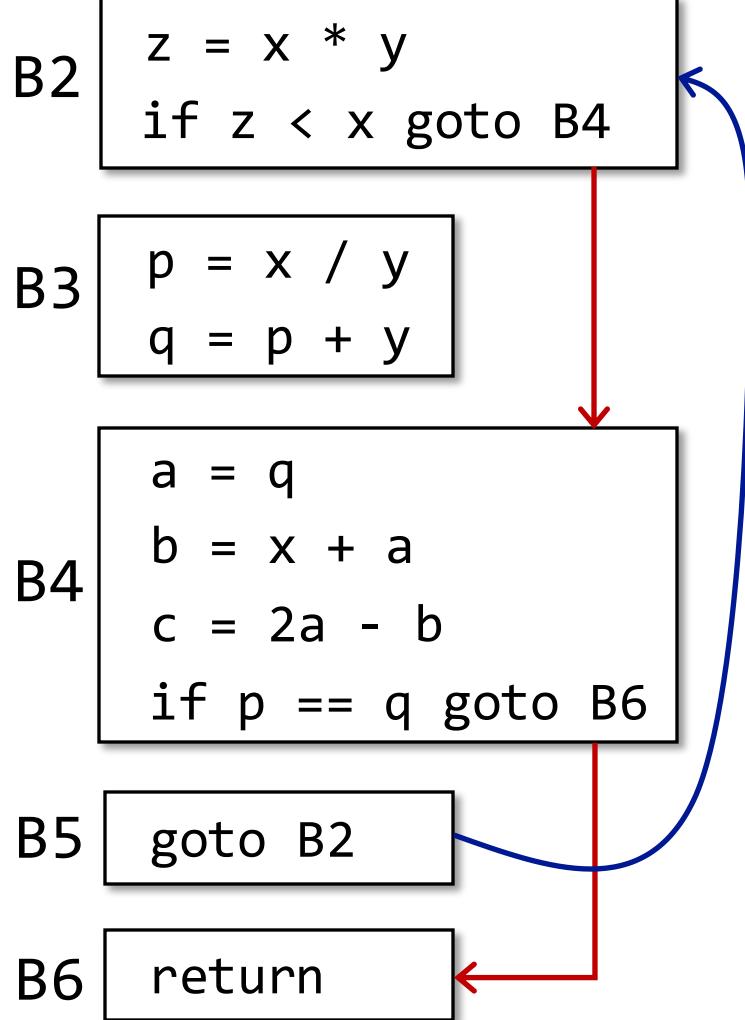
```
B3
p = x / y
q = p + y
```

```
B4
a = q
b = x + a
c = 2a - b
if p == q goto B6
```

```
B5
goto B2
```

```
B6
return
```

There is a **conditional** or **unconditional** jump from the end of **A** to the beginning of **B**



Add edges in CFG

B1
x = input
y = x - 1

B2
z = x * y
if z < x goto B4

B3
p = x / y
q = p + y

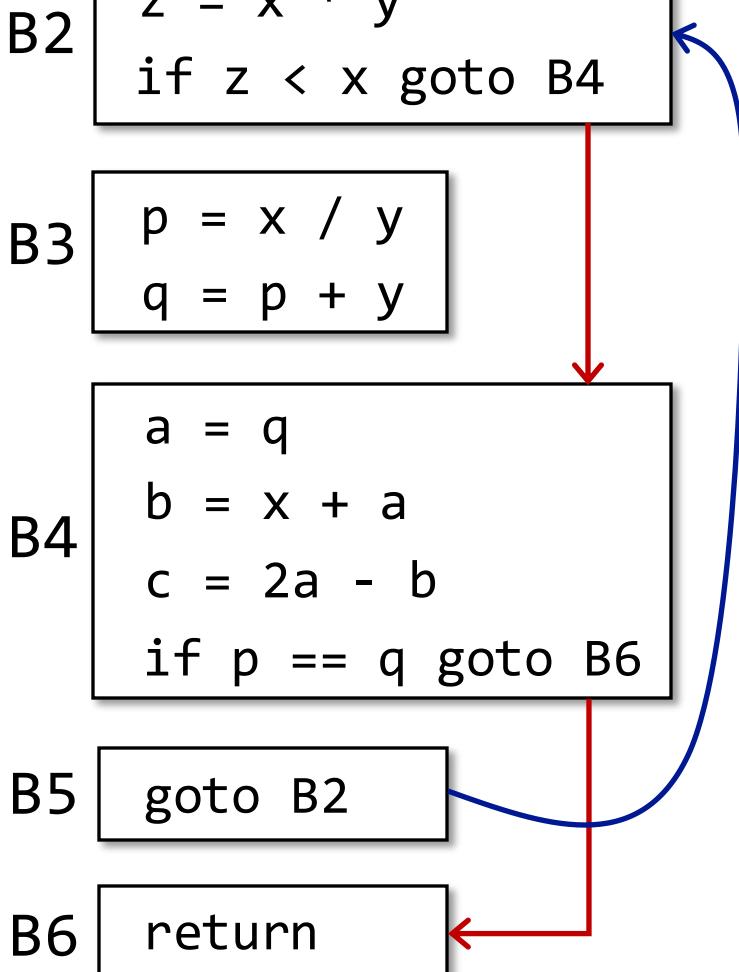
B4
a = q
b = x + a
c = 2a - b
if p == q goto B6

B5
goto B2

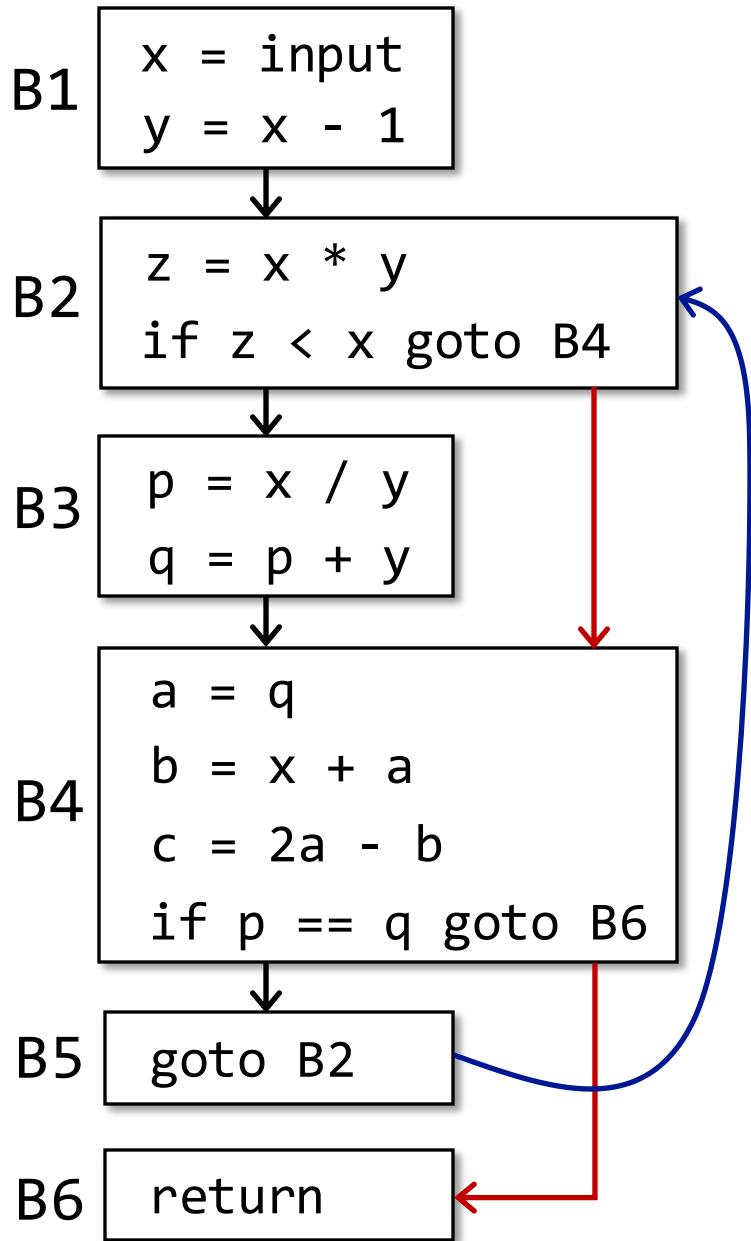
B6
return

There is a **conditional** or **unconditional** jump from the end of **A** to the beginning of **B**

B immediately follows **A** in the original order of instructions and **A** does not end in an unconditional jump



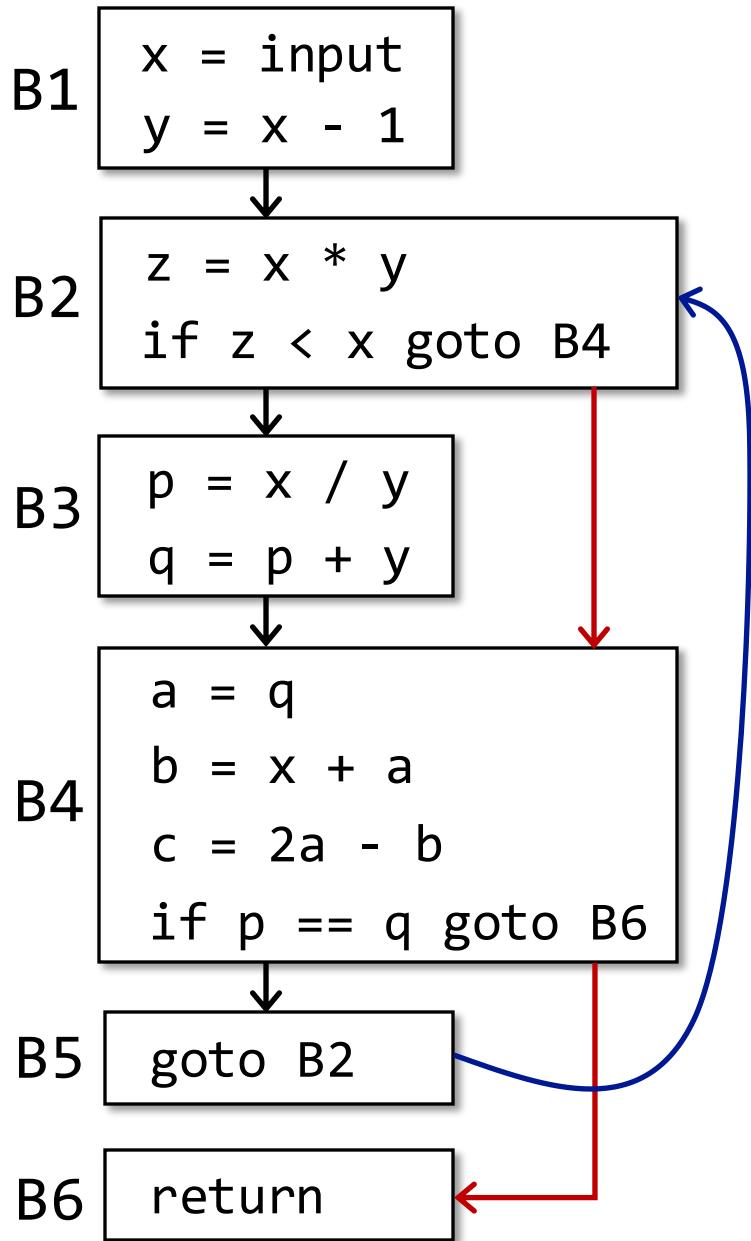
Add edges in CFG



There is a **conditional** or **unconditional** jump from the end of **A** to the beginning of **B**

B immediately follows **A** in the original order of instructions and **A** does not end in an unconditional jump

Add edges in CFG

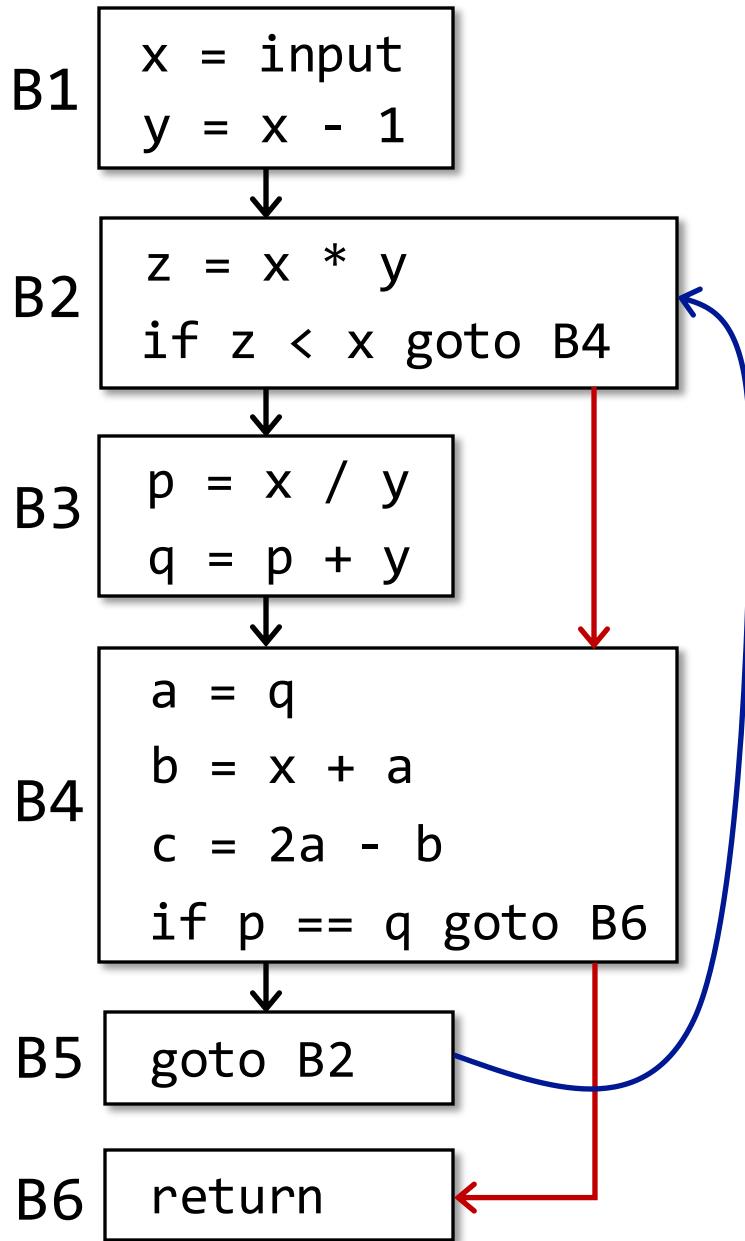


There is a **conditional** or **unconditional** jump from the end of **A** to the beginning of **B**

B immediately follows **A** in the original order of instructions and **A** does not end in an unconditional jump

We say that **A** is a **predecessor** of **B**, and **B** is a **successor** of **A**

Add edges in CFG



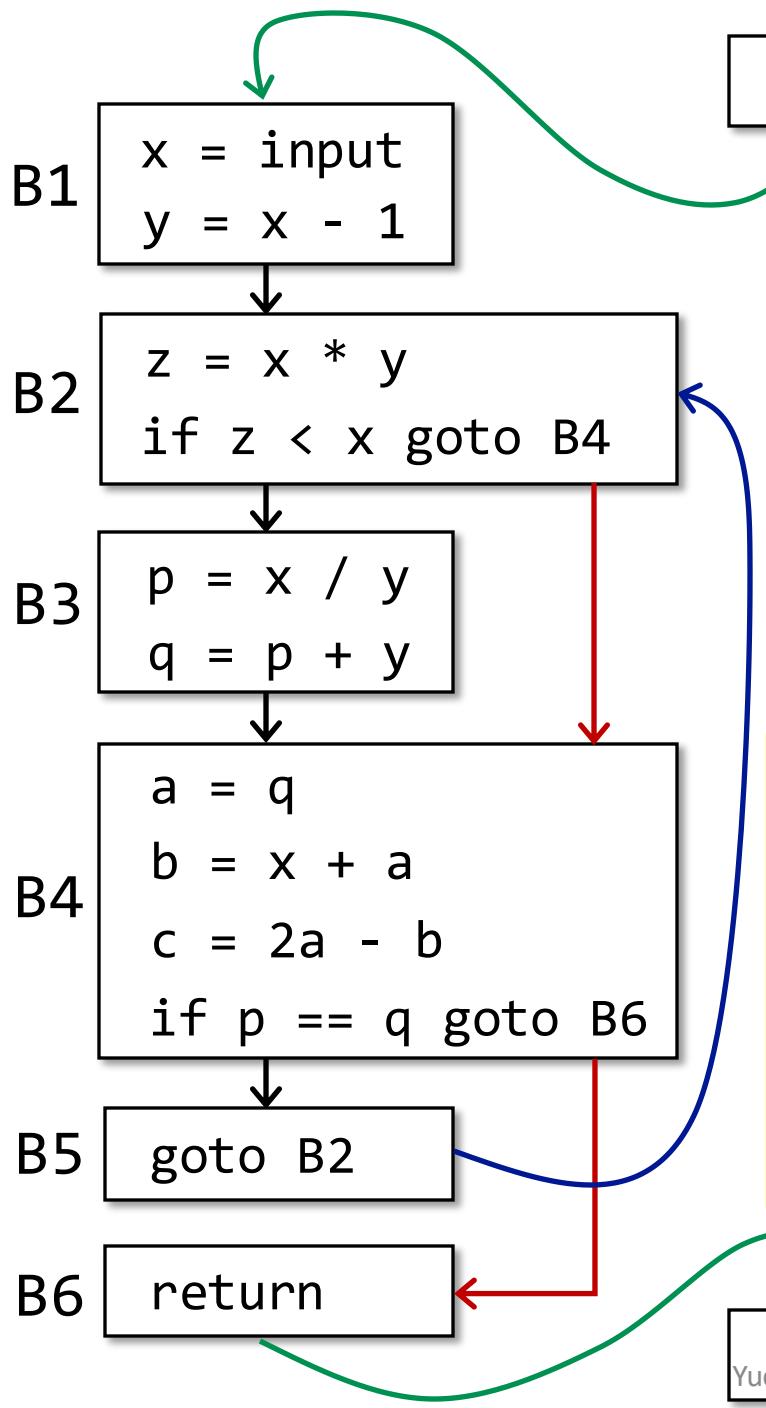
There is a **conditional** or **unconditional jump** from the end of A to the beginning of B

B immediately follows A in the original order of instructions and A does not end in an unconditional jump

We say that A is a **predecessor** of B, and B is a **successor** of A

Usually we add two nodes, **Entry** and **Exit**.

- They do not correspond to executable IR
- An edge from Entry to the BB containing the first instruction of IR
- An edge to Exit from any BB containing an instruction that could be the last instruction of IR



Entry

Add edges in CFG

There is a **conditional** or **unconditional jump** from the end of **A** to the beginning of **B**

B immediately follows **A** in the original order of instructions and **A** does not end in an unconditional jump

We say that **A** is a **predecessor** of **B**, and **B** is a **successor** of **A**

Usually we add two nodes, **Entry** and **Exit**.

- They do not correspond to executable IR
- An edge from Entry to the BB containing the first instruction of IR
- An edge to Exit from any BB containing an instruction that could be the last instruction of IR

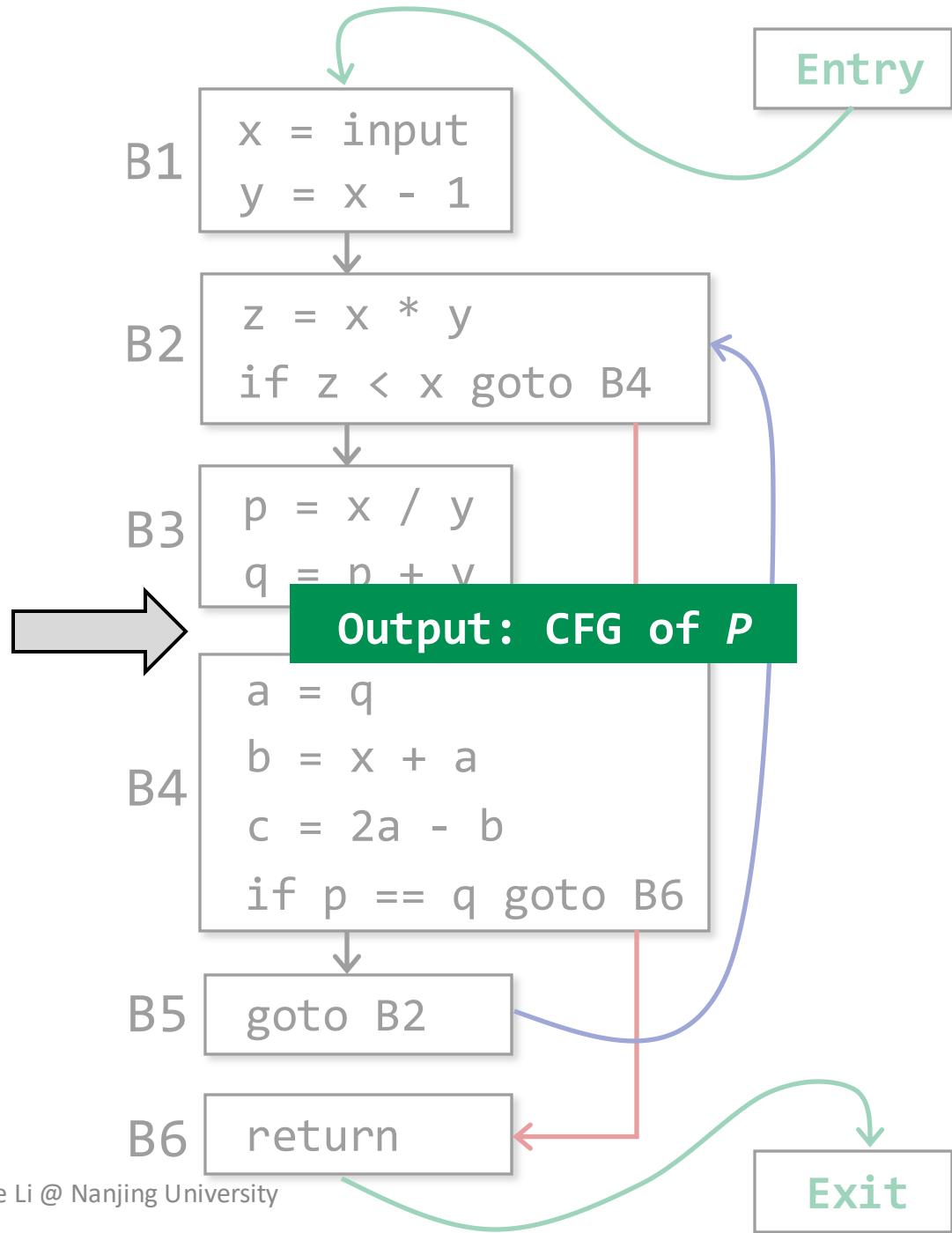
Exit

```

(1) x = input
(2) y = x - 1
(3) z = x * y
(4) if z < x goto (7)
(5) p = x / y
(6) q = p + v
(7) a = q
(8) b = x + a
(9) c = 2a - b
(10) if p == q goto (12)
(11) goto (3)
(12) return

```

Input: 3AC of P



Recommended Reading

The screenshot shows a user profile with a cartoon character icon, a banner for '甜品专家 Lv6 年度大会员' (Sweet Treat Expert Lv6 Annual Conference Member), and the text '真的以为，我能逼出红尘。' (Really think I can逼出 red dust.). The profile includes statistics: 关注数 61, 粉丝数 9779, 获赞数 1万, 播放数 25.7万, 阅读数 1.4万. Below the profile, a navigation bar has links for '主页' (Home), '动态' (Activity), '投稿 28' (Upload 28), '合集和列表 1' (Collection and List 1), '订阅' (Subscribe), and a search bar. A button '已关注' (Followed) is present. The main content area displays a list of 16 video lectures for the '南京大学《软件分析》课程2020' (Nanjing University Software Analysis Course 2020). The lectures are arranged in two rows of eight. Each lecture card includes the title, duration, view count, and upload date. The first row includes a '播放全部' (Play All) button. The second row includes sorting options '默认排序' (Default Sort) and '升序排序' (Ascending Sort).

课程标题	时长	观看次数	上传日期
南京大学《软件分析》课程01 (Introduction)	01:57:19	7万	2020-2-26
南京大学《软件分析》课程02 (Intermediate)	01:56:36	2.4万	2020-3-5
南京大学《软件分析》课程03 (Data Flow Analysis I)	01:59:07	2.3万	2020-3-12
南京大学《软件分析》课程04 (Data Flow Analysis II)	01:52:37	1.6万	2020-3-19
南京大学《软件分析》课程05 (Data Flow Analysis -)	01:43:57	1.3万	2020-3-26
南京大学《软件分析》课程06 (Data Flow Analysis -)	02:01:49	1.3万	2020-4-1
南京大学《软件分析》课程07 (Interprocedural Analysis)	01:46:51	1.2万	2020-4-15
南京大学《软件分析》课程08 (Pointer Analysis)	01:47:37	1.3万	2020-4-22
Our Pointer Analysis Algorithms	01:47:37		
Pointer Analysis in the Presence of Method Invocations	01:54:51		
Static Program Analysis Pointer Analysis Context-Sensitive I	01:48:03		
How to Implement Context-Sensitive Pointer Analysis	01:53:17		
南京大学《软件分析》课程09 (Pointer Analysis -)	02:03:28		
南京大学《软件分析》课程10 (Pointer Analysis -)			
南京大学《软件分析》课程11 (Pointer Analysis -)			
南京大学《软件分析》课程12 (Pointer Analysis -)			