# Humuhumunukunukuapua'a UFMG

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#### 1 Estruturas

#### 1.1 BIT

```
// BIT de soma 1-based, v 0-based
// Para mudar o valor da posicao p para x,
// faca: poe(x - query(p, p), p)
// l_bound(x) retorna o menor p tal que
// query(1, p+1) > x (0 based!)
//
// Complexidades:
// build - O(n)
// poe - O(\log(n))
// query - O(log(n))
// l_bound - O(log(n))
// d432a4
1a8 int n;
7f4 int bit[MAX];
b69 int v[MAX];
0a8 void build() {
b91
        bit[0] = 0;
        for (int i = 1; i <= n; i++) bit[i] = v[i - 1];</pre>
33 c
```

```
for (int i = 1; i <= n; i++) {
78a
edf
            int j = i + (i \& -i);
            if (j <= n) bit[j] += bit[i];</pre>
b8a
 cbb
cbb }
    // soma x na posicao p
235 void poe(int x, int p) {
        for (; p <= n; p += p & -p) bit[p] += x;
9 c7
cbb }
    // soma [1, p]
Obf int pref(int p) {
7 c 9
       int ret = 0;
        for (; p; p -= p & -p) ret += bit[p];
805
         return ret:
edf
cbb }
    // soma [a, b]
4ea int query(int a, int b) {
70c
        return pref(b) - pref(a - 1);
cbb }
e4a int l_bound(ll x) {
1ba
        int p = 0;
        for (int i = MAX2; i+1; i--) if (p + (1<<i) <= n
676
             and bit[p + (1 << i)] <= x) x -= bit[p += (1 << i)];
729
74e
         return p;
cbb }
1.2 BIT 2D
// BIT de soma, update incrementa posicao
// Tem que construir com um vetor com todos os pontos
// que vc quer um dia atualizar (os pontos q vc vai chamar
    update)
//
// Complexidades:
// construir - O(n log(n))
// update e query - O(\log^2(n))
// 6a760a
```

```
a6b template < class T = int> struct bit2d {
acf
        vector<T> X;
a84
        vector < vector < T >> Y, t;
709
        int ub(vector<T>& v, T x) {
dde
            return upper_bound(v.begin(), v.end(), x) -
   v.begin();
        }
cbb
        bit2d(vector<pair<T, T>> v) {
5cb
            for (auto [x, y] : v) X.push_back(x);
2e1
            sort(X.begin(), X.end());
fd4
            X.erase(unique(X.begin(), X.end()), X.end());
1ee
            t.resize(X.size() + 1);
d56
d12
            Y.resize(t.size());
3d0
            sort(v.begin(), v.end(), [](auto a, auto b) {
                 return a.second < b.second; });</pre>
43 d
            for (auto [x, y] : v) for (int i = ub(X, x); i < v)
961
   t.size(); i += i&-i)
                if (!Y[i].size() or Y[i].back() != y)
b75
   Y[i].push_back(y);
            for (int i = 0; i < t.size(); i++)</pre>
7 c 7
   t[i].resize(Y[i].size() + 1);
cbb
       }
e78
        void update(T x, T y, T v) {
            for (int i = ub(X, x); i < t.size(); i += i&-i)</pre>
2a9
                 for (int j = ub(Y[i], y); j < t[i].size(); j</pre>
cd2
   += j\&-j) t[i][j] += v;
cbb
        T query(T x, T y) {
5d2
966
            T ans = 0:
c54
            for (int i = ub(X, x); i; i = i\&-i)
                 for (int j = ub(Y[i], y); j; j -= j\&-j) ans
4fb
   += t[i][i];
ba7
            return ans;
cbb
46 d
        T \text{ query}(T x1, T y1, T x2, T y2)  {
            return query (x2, y2) -query (x2, y1-1) -query (x1-1, y2)
fcf
```

```
y2)+query(x1-1, y1-1);
cbb }
214 };
```

#### 1.3 BIT com update em range

```
// Operacoes 0-based
// query(1, r) retorna a soma de v[1..r]
// update(1, r, x) soma x em v[1..r]
//
// Complexidades:
// build - O(n)
// query - O(log(n))
// update - O(log(n))
// f91737
e04 namespace bit {
3ba
        11 bit[2][MAX+2];
1 a8
        int n;
61c
        void build(int n2, int* v) {
1 e3
            n = n2;
535
            for (int i = 1; i <= n; i++)
                bit[1][min(n+1, i+(i&-i))] += bit[1][i] +=
edd
   v[i-1];
cbb
637
        11 get(int x, int i) {
b73
            11 \text{ ret} = 0;
            for (; i; i -= i&-i) ret += bit[x][i];
360
edf
            return ret;
cbb
20 c
        void add(int x, int i, ll val) {
503
            for (; i <= n; i += i&-i) bit[x][i] += val;
cbb
162
        11 get2(int p) {
            return get(0, p) * p + get(1, p);
с7с
cbb
02a
        11 query(int 1, int r) {
ff5
            return get2(r+1) - get2(1);
cbb
089
        void update(int 1, int r, 11 x) {
e5f
            add(0, 1+1, x), add(0, r+2, -x);
```

```
f58
            add(1, 1+1, -x*1), add(1, r+2, x*(r+1));
cbb
        }
214 };
1.4 DSU
// Une dois conjuntos e acha a qual conjunto um elemento
   pertence por seu id
// find e unite: O(a(n)) \sim = O(1) amortizado
// 8e197e
8d3 struct dsu {
825
        vector < int > id, sz;
        dsu(int n) : id(n), sz(n, 1) { iota(id.begin(), }
b33
   id.end(), 0); }
        int find(int a) { return a == id[a] ? a : id[a] =
0cf
   find(id[a]); }
440
        void unite(int a, int b) {
            a = find(a), b = find(b);
605
           if (a == b) return;
d54
           if (sz[a] < sz[b]) swap(a, b);
956
            sz[a] += sz[b], id[b] = a;
6d0
cbb
        }
214 };
    // DSU de bipartido
```

// Une dois vertices e acha a qual componente um vertice

// Informa se a componente de um vertice e bipartida

 $dsu(int n) : id(n), sz(n, 1), bip(n, 1), c(n) {$ 

//

pertence

// 118050

8d3 struct dsu {

6f7

5b4

// find e unite:  $O(\log(n))$ 

vector<int> id, sz, bip, c;

```
cbb
       int find(int a) { return a == id[a] ? a :
   find(id[a]); }
        int color(int a) { return a == id[a] ? c[a] : c[a] ^
   color(id[a]); }
440
        void unite(int a, int b) {
            bool change = color(a) == color(b);
263
            a = find(a), b = find(b);
605
a 89
            if (a == b) {
                if (change) bip[a] = 0;
4 ed
505
                return;
            }
cbb
d41
956
            if (sz[a] < sz[b]) swap(a, b);
            if (change) c[b] = 1;
efe
2cd
            sz[a] += sz[b], id[b] = a, bip[a] &= bip[b];
cbb
214 };
    // DSU Persistente
    // Persistencia parcial, ou seja, tem que ir
    // incrementando o 't' no une
    // find e unite: O(\log(n))
    // 6c63a4
8d3 struct dsu {
33c
        vector < int > id, sz, ti;
733
        dsu(int n) : id(n), sz(n, 1), ti(n, -INF) {
db8
            iota(id.begin(), id.end(), 0);
cbb
        int find(int a, int t) {
5e6
            if (id[a] == a or ti[a] > t) return a;
6ba
ea5
            return find(id[a], t);
cbb
        }
```

iota(id.begin(), id.end(), 0);

db8

```
fa0
        void unite(int a, int b, int t) {
            a = find(a, t), b = find(b, t);
84f
            if (a == b) return;
d54
           if (sz[a] < sz[b]) swap(a, b);
956
            sz[a] += sz[b], id[b] = a, ti[b] = t;
35 d
cbb
        }
214 };
    // DSU com rollback
    // checkpoint(): salva o estado atual de todas as
       variaveis
    // rollback(): retorna para o valor das variaveis para
    // o ultimo checkpoint
    // Sempre que uma variavel muda de valor, adiciona na
       stack
    // find e unite: O(\log(n))
    // checkpoint: O(1)
    // rollback: O(m) em que m e o numero de vezes que alguma
   // variavel mudou de valor desde o ultimo checkpoint
    // c6e923
8d3 struct dsu {
825
        vector<int> id, sz;
27 c
        stack < stack < pair < int & , int >>> st;
98 d
        dsu(int n) : id(n), sz(n, 1) {
            iota(id.begin(), id.end(), 0), st.emplace();
1cc
cbb
d41
        void save(int &x) { st.top().emplace(x, x); }
bdf
        void checkpoint() { st.emplace(); }
30 d
d41
5cf
        void rollback() {
            while(st.top().size()) {
ba9
                auto [end, val] = st.top().top();
6bf
   st.top().pop();
149
                end = val;
```

```
cbb
            }
25a
            st.pop();
cbb
        int find(int a) { return a == id[a] ? a :
ef0
   find(id[a]); }
440
        void unite(int a, int b) {
            a = find(a), b = find(b);
605
d54
            if (a == b) return;
956
            if (sz[a] < sz[b]) swap(a, b);
803
            save(sz[a]), save(id[b]);
6d0
            sz[a] += sz[b], id[b] = a;
cbb
214 };
1.5 Li-Chao Tree
// Adiciona retas (ax+b), e computa o minimo entre as retas
// em um dado 'x'
// Cuidado com overflow!
// Se tiver overflow, tenta comprimir o 'x' ou usar
// convex hull trick
//
// O(log(MA-MI)), O(n) de memoria
// 59ba68
5b0 template < 11 MI = 11(-1e9), 11 MA = 11(1e9) > struct
   lichao {
        struct line {
b3a
12d
            ll a, b:
cef
            array<int, 2> ch;
            line(ll a_{-} = 0, ll b_{-} = LINF):
fdf
423
                a(a_{-}), b(b_{-}), ch(\{-1, -1\}) \{\}
888
            11 operator ()(11 x) { return a*x + b; }
214
        };
17b
        vector < line > ln;
        int ch(int p, int d) {
df8
```

if (ln[p].ch[d] == -1) {

ln.emplace\_back();

ln[p].ch[d] = ln.size();

e85

9af

cdc

```
cbb
ef2
            return ln[p].ch[d];
cbb
021
        lichao() { ln.emplace_back(); }
c33
        void add(line s, ll l=MI, ll r=MA, int p=0) {
3e3
            11 m = (1+r)/2;
            bool L = s(1) < ln[p](1);
911
            bool M = s(m) < ln[p](m);
d37
            bool R = s(r) < ln[p](r);
03b
           if (M) swap(ln[p], s), swap(ln[p].ch, s.ch);
825
           if (s.b == LINF) return;
cac
           if (L != M) add(s, 1, m-1, ch(p, 0));
f6d
            else if (R != M) add(s, m+1, r, ch(p, 1));
898
cbb
092
        11 query(int x, 11 1=MI, 11 r=MA, int p=0) {
            11 m = (1+r)/2, ret = ln[p](x);
11b
            if (ret == LINF) return ret;
9db
           if (x < m) return min(ret, query(x, 1, m-1,</pre>
529
   ch(p, 0));
81a
            return min(ret, query(x, m+1, r, ch(p, 1));
cbb
214 };
```

#### 1.6 MergeSort Tree

```
// Se for construida sobre um array:
        count(i, j, a, b) retorna quantos
//
        elementos de v[i..j] pertencem a [a, b]
//
        report(i, j, a, b) retorna os indices dos
//
//
        elementos de v[i..j] que pertencem a [a, b]
        retorna o vetor ordenado
// Se for construida sobre pontos (x, y):
        count(x1, x2, y1, x2) retorna quantos pontos
//
//
        pertencem ao retangulo (x1, y1), (x2, y2)
//
        report(x1, x2, y1, y2) retorna os indices dos pontos
   que
        pertencem ao retangulo (x1, y1), (x2, y2)
//
//
        retorna os pontos ordenados lexicograficamente
        (assume x1 \le x2, y1 \le y2)
// kth(y1, y2, k) retorna o indice do ponto com k-esimo menor
```

```
// x dentre os pontos que possuem y em [y1, y2] (0 based)
// Se quiser usar para achar k-esimo valor em range,
   construir
// com ms_tree t(v, true), e chamar kth(l, r, k)
//
// Usa O(n log(n)) de memoria
// Complexidades:
// construir - O(n log(n))
// count - O(\log(n))
// report - O(log(n) + k) para k indices retornados
// kth - O(log(n))
// 1cef03
c6c template <typename T = int> struct ms_tree {
6f7
        vector < tuple < T, T, int >> v;
1 a8
        int n;
5ee
        vector < vector < tuple < T, T, int >>> t; // {y, idx, left}
        vector <T> vy;
6ae
78 c
        ms_tree(vector<pair<T, T>>& vv) : n(vv.size()),
   t(4*n), vy(n) {
e80
            for (int i = 0; i < n; i++)
   v.push_back({vv[i].first, vv[i].second, i});
            sort(v.begin(), v.end());
fca
224
            build(1, 0, n-1);
           for (int i = 0; i < n; i++) vy[i] =
   get <0 > (t[1][i+1]);
cbb
        }
        ms_tree(vector<T>& vv, bool inv = false) { // inv:
dac
   inverte indice e valor
            vector < pair < T, T >> v2;
8 e 8
           for (int i = 0; i < vv.size(); i++)</pre>
e1e
                inv ? v2.push_back({vv[i], i}) :
196
   v2.push_back({i, vv[i]});
            *this = ms_tree(v2);
cca
cbb
2 c 6
        void build(int p, int l, int r) {
            t[p].push_back({get<0>(v[1]), get<0>(v[r]), 0});
1 d2
   // {min x, max x, 0}
            if (1 == r) return t[p].push_back({get<1>(v[1]),
5 c8
   get <2 > (v[1]), 0});
```

```
int m = (1+r)/2;
ee4
bd9
             build(2*p, 1, m), build(2*p+1, m+1, r);
32d
             int L = 0, R = 0;
             while (t[p].size() \le r-l+1) {
a03
68 e
                 int left = get <2>(t[p].back());
                 if (L > m-1 \text{ or } (R+m+1 \le r \text{ and } t[2*p+1][1+R])
4aa
   < t[2*p][1+L])) {
                      t[p].push_back(t[2*p+1][1 + R++]);
8cf
                      get < 2 > (t[p].back()) = left;
da0
5e2
                      continue;
cbb
249
                 t[p].push_back(t[2*p][1 + L++]);
339
                 get < 2 > (t[p].back()) = left+1;
             }
cbb
cbb
        }
        int get_l(T y) { return lower_bound(vy.begin(),
dd3
   vy.end(), y) - vy.begin(); }
        int get_r(T y) { return upper_bound(vy.begin(),
   vy.end(), y) - vy.begin(); }
f62
        int count(T x1, T x2, T y1, T y2) {
             function < int(int, int, int) > dfs = [&](int p,
902
   int 1, int r) {
7c6
                 if (1 == r \text{ or } x2 < get < 0 > (t[p][0]) \text{ or }
   get<1>(t[p][0]) < x1) return 0;
                 if (x1 \le get < 0 > (t[p][0]) and
2bb
   get<1>(t[p][0]) <= x2) return r-1;
                 int nl = get<2>(t[p][1]), nr =
784
   get <2>(t[p][r]);
eb6
                 return dfs(2*p, nl, nr) + dfs(2*p+1, l-nl,
   r-nr);
             };
214
7cb
             return dfs(1, get_l(y1), get_r(y2));
cbb
002
        vector<int> report(T x1, T x2, T y1, T y2) {
4b8
             vector<int> ret;
            function < void(int, int, int) > dfs = [&](int p,
85 e
   int 1, int r) {
882
                 if (1 == r \text{ or } x2 < get < 0 > (t[p][0]) \text{ or }
   get<1>(t[p][0]) < x1) return;</pre>
```

```
8da
                if (x1 \le get < 0 > (t[p][0]) and
   get<1>(t[p][0]) <= x2) {
                    for (int i = 1; i < r; i++)
e00
   ret.push_back(get<1>(t[p][i+1]));
505
                     return:
cbb
784
                int nl = get <2>(t[p][1]), nr =
   get <2 > (t[p][r]);
194
                dfs(2*p, nl, nr), dfs(2*p+1, l-nl, r-nr);
214
            };
8ad
            dfs(1, get_l(y1), get_r(y2));
edf
            return ret;
cbb
985
        int kth(T y1, T y2, int k) {
902
            function < int(int, int, int) > dfs = [&](int p,
   int 1, int r) {
150
                if (k >= r-1) {
                    k -= r-1:
941
                    return -1;
daa
cbb
8da
                if (r-l == 1) return get <1 > (t[p][1+1]);
784
                int nl = get <2>(t[p][1]), nr =
   get <2 > (t[p][r]);
072
                int left = dfs(2*p, nl, nr);
3 b 6
                if (left != -1) return left;
04d
                return dfs(2*p+1, l-nl, r-nr);
214
            };
7cb
            return dfs(1, get_l(y1), get_r(y2));
cbb
214 };
1.7 Min queue - deque
// Tudo O(1) amortizado
// c13c57
```

```
// Tudo 0(1) amortizado
// c13c57

1dc template < class T > struct minqueue {
2d8     deque < pair < T, int >> q;

3fc     void push(T x) {
56e         int ct = 1;
953         while (q.size() and x < q.front().first)</pre>
```

#### 1.8 Min queue - stack

```
// Tudo O(1) amortizado
// fe0cad
557 template < class T> struct minstack {
81f
        stack <pair <T, T>> s;
3fc
        void push(T x) {
            if (!s.size()) s.push({x, x});
12b
9d9
            else s.emplace(x, std::min(s.top().second, x));
cbb
4f0
        T top() { return s.top().first; }
94a
        } () gog T
            T ans = s.top().first;
1f2
2eb
            s.pop();
            return ans;
ba7
cbb
614
        int size() { return s.size(); }
        T min() { return s.top().second; }
13b
214 };
1dc template < class T> struct minqueue {
        minstack <T> s1, s2;
cdc
        void push(T x) { s1.push(x); }
7cd
c96
        void move() {
d4d
            if (s2.size()) return;
d92
            while (s1.size()) {
                T x = s1.pop();
7ae
489
                s2.push(x);
cbb
            }
```

```
cbb
787
        T front() { return move(), s2.top(); }
        T pop() { return move(), s2.pop(); }
23a
        int size() { return s1.size()+s2.size(); }
7 f 3
19 c
        T min() {
cd6
            if (!s1.size()) return s2.min();
            else if (!s2.size()) return s1.min();
58e
31d
            return std::min(s1.min(), s2.min());
cbb
214 };
```

#### 1.9 Order Statistic Set

```
// Funciona do C++11 pra cima
774 #include <ext/pb_ds/assoc_container.hpp>
30f #include <ext/pb_ds/tree_policy.hpp>
0d7 using namespace __gnu_pbds;
4fc template <class T>
        using ord_set = tree<T, null_type, less<T>,
def
   rb_tree_tag,
3a1
        tree_order_statistics_node_update>;
d41
    // para declarar:
b36 ord_set <int> s;
    // coisas do set normal funcionam:
e6f for (auto i : s) cout << i << endl;</pre>
738 cout << s.size() << endl;
    // k-esimo maior elemento O(log|s|):
    // k=0: menor elemento
e46 cout << *s.find_by_order(k) << endl;
    // quantos sao menores do que k O(\log |s|):
df7 cout << s.order_of_key(k) << endl;</pre>
    // Para fazer um multiset, tem que
    // usar ord_set<pair<int, int>> com o
    // segundo parametro sendo algo para diferenciar
    // os ementos iguais.
    // s.order_of_key({k, -INF}) vai retornar o
    // numero de elementos < k
```

#### 1.10 Range color

```
// update(1, r, c) colore o range [1, r] com a cor c,
// e retorna os ranges que foram coloridos {1, r, cor}
// query(i) returna a cor da posicao i
//
// Complexidades (para q operacoes):
// update - O(log(q)) amortizado
// query - 0(log(q))
// 9e9cab
df6 template < typename T> struct color {
        set < tuple < int , int , T >> se;
f0c
071
        vector<tuple<int, int, T>> update(int 1, int r, T
   val) {
9c4
            auto it = se.upper_bound({r, INF, val});
            if (it != se.begin() and get<1>(*prev(it)) > r) {
753
e91
                auto [L, R, V] = *--it;
                se.erase(it);
3f0
                se.emplace(L, r, V), se.emplace(r+1, R, V);
bfd
cbb
d9e
            it = se.lower_bound({1, -INF, val});
            if (it != se.begin() and get<1>(*prev(it)) >= 1)
516
   {
e91
                auto [L, R, V] = *--it;
                se.erase(it);
3f0
                se.emplace(L, 1-1, V), it = se.emplace(1, R,
75a
   V).first;
cbb
d7b
            vector<tuple<int, int, T>> ret;
7a1
            for (; it != se.end() and get<0>(*it) <= r; it =
   se.erase(it))
                ret.push_back(*it);
8c0
            se.emplace(1, r, val);
b4a
edf
            return ret;
cbb
ff9
        T query(int i) {
            auto it = se.upper_bound({i, INF, T()});
c31
            if (it == se.begin() or get<1>(*--it) < i)</pre>
8e7
   return -1; // nao tem
53 d
            return get <2 > (*it);
```

```
cbb
214 };
1.11 RMQ \langle O(n), O(1) \rangle - min queue
// O(n) pra buildar, query O(1)
// Se tiver varios minimos, retorna
// o de menor indice
// bab412
1a5 template < typename T> struct rmq {
517
        vector <T> v;
        int n; static const int b = 30;
fcc
70e
        vector < int > mask, t;
        int op(int x, int y) { return v[x] \leftarrow v[y] ? x : y; }
183
        int msb(int x) { return
ee1
   __builtin_clz(1)-__builtin_clz(x); }
        int small(int r, int sz = b) { return
   r-msb(mask[r]&((1<<sz)-1));}
        rmq() {}
6ad
        rmq(const vector < T > \& v_) : v(v_), n(v.size()),
43c
   mask(n), t(n) {
2e5
            for (int i = 0, at = 0; i < n; mask[i++] = at |=
   1) {
a61
                at = (at << 1) &((1 << b) -1);
c00
                while (at and op(i-msb(at&-at), i) == i) at
   ^= at&-at;
cbb
            for (int i = 0; i < n/b; i++) t[i] =
ea4
   small(b*i+b-1):
            for (int j = 1; (1<<j) <= n/b; j++) for (int i =
39 d
   0; i+(1<<j) <= n/b; i++)
                t[n/b*j+i] = op(t[n/b*(j-1)+i],
ba5
   t[n/b*(j-1)+i+(1<<(j-1))]);
cbb
e34
        int index_query(int 1, int r) {
            if (r-l+1 <= b) return small(r, r-l+1);</pre>
27b
e80
            int x = 1/b+1, y = r/b-1;
fd3
            if (x > y) return op(small(1+b-1), small(r));
a4e
            int j = msb(y-x+1);
            int ans = op(small(1+b-1), op(t[n/b*j+x],
```

ea3

```
t[n/b*j+y-(1<<j)+1]));
be6
            return op(ans, small(r));
cbb
        T query(int 1, int r) { return v[index_query(1, r)];
093
   }
214 };
1.12 SegTreap
// Muda uma posicao do plano, e faz query de operacao
// associativa e comutativa em retangulo
// Mudar ZERO e op
// Esparso nas duas coordenadas, inicialmente eh tudo ZERO
//
// Para query com distancia de manhattan <= d, faca
// nx = x+y, ny = x-y
// Update em (nx, ny), query em ((nx-d, ny-d), (nx+d, ny+d))
//
// Valores no X tem que ser de O ateh NX
// Para q operacoes, usa O(q log(NX)) de memoria, e as
// operacoes custa O(log(q) log(NX))
// 75f2d0
55b const int ZERO = INF;
560 const int op(int 1, int r) { return min(1, r); }
878 mt19937 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
aa1 template < typename T> struct treap {
3c9
        struct node {
b19
            node *1, *r;
ee1
            int p;
            pair<11, 11> idx; // {y, x}
850
36 d
            T val, mi;
bc2
            node(ll x, ll y, T val_) : l(NULL), r(NULL),
   p(rng()),
                idx(pair(y, x)), val(val_), mi(val) {}
1b5
01e
            void update() {
                mi = val;
d6e
                if (1) mi = op(mi, 1->mi);
182
```

if (r) mi = op(mi, r->mi);

b68

```
cbb
             }
214
        };
        node* root;
bb7
84b
        treap() { root = NULL; }
cec
        \simtreap() {
609
             vector < node *> q = {root};
402
             while (q.size()) {
                 node* x = q.back(); q.pop_back();
e5d
ee9
                 if (!x) continue;
1 c7
                 q.push_back(x->1), q.push_back(x->r);
                 delete x;
bf0
             }
cbb
cbb
225
        treap(treap&& t) : treap() { swap(root, t.root); }
         void join(node* 1, node* r, node*& i) { // assume
bcf
   que 1 < r
986
             if (!l or !r) return void(i = 1 ? 1 : r);
            if (1->p > r->p) join(1->r, r, 1->r), i = 1;
80e
fa0
             else join(l, r\rightarrow l, r\rightarrow l), i = r;
bda
             i->update();
cbb
c82
        void split(node* i, node*& 1, node*& r, pair<11, 11>
   idx) {
26a
             if (!i) return void(r = 1 = NULL);
13 c
             if (i-)idx < idx) split(i-)r, i-)r, r, idx), l =
   i;
d26
             else split(i \rightarrow l, l, i \rightarrow l, idx), r = i;
             i->update();
bda
cbb
d3b
        void update(ll x, ll y, T v) {
             node *L, *M, *R;
df9
8 b 2
             split(root, M, R, pair(y, x+1)), split(M, L, M,
   pair(y, x));
1 e 4
             if (M) M->val = M->mi = v;
9 e 5
             else M = new node(x, y, v);
             join(L, M, M), join(M, R, root);
69 d
cbb
91b
        T query(ll ly, ll ry) {
df9
             node *L, *M, *R;
```

```
1 c 0
            split(root, M, R, pair(ry, LINF)), split(M, L,
   M, pair(ly, 0));
            T \text{ ret} = M ? M -> mi : ZERO;
0f7
69 d
            join(L, M, M), join(M, R, root);
            return ret;
edf
cbb
        }
214 };
46a template < typename T > struct segtreap {
        vector<treap<T>> seg;
c4f
6e7
        vector < int > ch[2];
e4e
        11 NX;
        segtreap(ll NX_) : seg(1), NX(NX_) {
253
   ch[0].push_back(-1), ch[1].push_back(-1); }
a71
        int get_ch(int i, int d){
            if (ch[d][i] == -1) {
e51
2d6
                ch[d][i] = seg.size();
                 seg.emplace back();
23 e
                ch[0].push_back(-1), ch[1].push_back(-1);
842
cbb
968
            return ch[d][i];
cbb
        }
10 c
        T query(11 lx, 11 rx, 11 ly, 11 ry, int p, 11 1, 11
  r) {
            if (rx < 1 or r < 1x) return ZERO;
003
            if (lx <= 1 and r <= rx) return seg[p].query(ly,</pre>
fOf
   ry);
            11 m = 1 + (r-1)/2;
e6a
354
            return op(query(lx, rx, ly, ry, get_ch(p, 0), l,
   m).
060
                     query(lx, rx, ly, ry, get_ch(p, 1), m+1,
   r)):
cbb
        T query(ll lx, ll rx, ll ly, ll ry) { return
   query(lx, rx, ly, ry, 0, 0, NX); }
        void update(ll x, ll y, T val, int p, ll l, ll r) {
249
73с
            if (l == r) return seg[p].update(x, y, val);
```

## 1.13 SegTree

```
// Recursiva com Lazy Propagation
// Query: soma do range [a, b]
// Update: soma x em cada elemento do range [a, b]
// Pode usar a seguinte funcao para indexar os nohs:
// f(1, r) = (1+r) | (1!=r), usando 2N de memoria
//
// Complexidades:
// build - O(n)
// query - 0(log(n))
// update - O(log(n))
// Oafec1
aa4 namespace seg {
005
        11 \text{ seg}[4*MAX], lazy[4*MAX];
052
        <u>int</u> n, *v;
        11 build(int p=1, int l=0, int r=n-1) {
d22
3 c7
            lazy[p] = 0;
6cd
            if (1 == r) return seg[p] = v[1]:
ee4
            int m = (1+r)/2:
193
            return seg[p] = build(2*p, 1, m) + build(2*p+1,
   m+1, r);
cbb
        void build(int n2, int* v2) {
0 d8
            n = n2, v = v2;
680
6f2
            build();
cbb
ceb
        void prop(int p, int 1, int r) {
            seg[p] += lazy[p]*(r-l+1);
cdf
2 c 9
            if (1 != r) lazy[2*p] += lazy[p], lazy[2*p+1] +=
```

```
lazy[p];
3c7
            lazv[p] = 0;
cbb
2c3
        11 query(int a, int b, int p=1, int l=0, int r=n-1) {
6b9
            prop(p, 1, r);
            if (a <= 1 and r <= b) return seg[p];</pre>
527
            if (b < 1 or r < a) return 0;
786
ee4
            int m = (1+r)/2;
b1f
            return query (a, b, 2*p, 1, m) + query(a, b,
   2*p+1, m+1, r);
cbb
        ll update(int a, int b, int x, int p=1, int l=0, int
cfb
   r=n-1) {
6b9
            prop(p, 1, r);
9a3
            if (a \le 1 \text{ and } r \le b) {
b94
                 lazy[p] += x;
6b9
                 prop(p, 1, r);
                 return seg[p];
534
            }
cbb
e9f
            if (b < l or r < a) return seg[p];</pre>
            int m = (1+r)/2;
ee4
fdb
            return seg[p] = update(a, b, x, 2*p, 1, m) +
7fd
                 update(a, b, x, 2*p+1, m+1, r);
cbb
        }
214 };
    // Se tiver uma seg de max, da pra descobrir em O(\log(n))
    // o primeiro e ultimo elemento >= val numa range:
    // primeira posicao >= val em [a, b] (ou -1 se nao tem)
    // 68c3e5
119 int get_left(int a, int b, int val, int p=1, int l=0,
   int r=n-1) {
6b9
        prop(p, 1, r);
        if (b < l \text{ or } r < a \text{ or seg}[p] < val) return -1;
f38
205
        if (r == 1) return 1;
        int m = (1+r)/2;
ee4
753
        int x = get_left(a, b, val, 2*p, l, m);
50e
        if (x != -1) return x;
сЗс
        return get_left(a, b, val, 2*p+1, m+1, r);
cbb }
```

```
// ultima posicao >= val em [a, b] (ou -1 se nao tem)
    // 1b71df
992 int get_right(int a, int b, int val, int p=1, int l=0,
   int r=n-1) {
6 b 9
        prop(p, 1, r);
f38
        if (b < l or r < a or seg[p] < val) return -1;</pre>
205
        if (r == 1) return 1;
        int m = (1+r)/2;
ee4
        int x = get_right(a, b, val, 2*p+1, m+1, r);
1 b 1
50e
        if (x != -1) return x;
6a7
        return get_right(a, b, val, 2*p, 1, m);
cbb }
    // Se tiver uma seg de soma sobre um array nao negativo
       v, da pra
    // descobrir em O(\log(n)) o maior j tal que
       v[i] + v[i+1] + ... + v[j-1] < val
    // 2b8ea7
6a9 int lower_bound(int i, ll& val, int p, int l, int r) {
6 b 9
        prop(p, 1, r);
6e8
        if (r < i) return n;
        if (i <= l and seg[p] < val) {</pre>
b5d
bff
            val -= seg[p];
041
            return n;
cbb
Зсе
        if (1 == r) return 1;
ee4
        int m = (1+r)/2;
514
        int x = lower_bound(i, val, 2*p, 1, m);
ee0
        if (x != n) return x;
8 b 9
        return lower_bound(i, val, 2*p+1, m+1, r);
cbb }
1.14 SegTree 2D Iterativa
// Consultas 0-based
// Um valor inicial em (x, y) deve ser colocado em
```

```
seg[x+n][y+n]
// Query: soma do retangulo ((x1, y1), (x2, y2))
// Update: muda o valor da posicao (x, y) para val
// Nao pergunte como que essa coisa funciona
//
```

```
// Para query com distancia de manhattan <= d, faca
// nx = x+y, ny = x-y
// Update em (nx, ny), query em ((nx-d, ny-d), (nx+d, ny+d))
// Se for de min/max, pode tirar os if's da 'query', e fazer
// sempre as 4 operacoes. Fica mais rapido
// Complexidades:
// build - O(n^2)
// \text{ query - O(log^2(n))}
// update - O(log^2(n))
// 67b9e5
731 int seg[2*MAX][2*MAX], n;
0a8 void build() {
919
         for (int x = 2*n; x; x--) for (int y = 2*n; y; y--) {
             if (x < n) seg[x][y] = seg[2*x][y] +
   seg[2*x+1][y];
             if (y < n) seg[x][y] = seg[x][2*y] +
   seg[x][2*y+1];
       }
cbb
cbb }
251 int query(int x1, int y1, int x2, int y2) {
827
         int ret = 0, v3 = v1 + n, v4 = v2 + n;
         for (x1 += n, x2 += n; x1 <= x2; ++x1 /= 2, --x2 /=
83e
   2)
             for (y1 = y3, y2 = y4; y1 \le y2; ++y1 /= 2, --y2
0f2
   /= 2) {
                 if (x1\%2 == 1 \text{ and } y1\%2 == 1) \text{ ret } +=
554
   seg[x1][v1];
                 if (x1\%2 == 1 \text{ and } y2\%2 == 0) \text{ ret } +=
6b0
   seg[x1][y2];
                 if (x2\%2 == 0 \text{ and } y1\%2 == 1) \text{ ret } +=
c01
   seg[x2][v1];
5d4
                  if (x2\%2 == 0 \text{ and } y2\%2 == 0) \text{ ret } +=
   seg[x2][y2];
            }
cbb
d41
edf
         return ret;
cbb }
```

```
767 void update(int x, int y, int val) {
66a
        int y2 = y += n;
192
        for (x += n; x; x /= 2, y = y2) {
970
            if (x >= n) seg[x][y] = val;
ba9
            else seg[x][y] = seg[2*x][y] + seg[2*x+1][y];
d41
3 b 1
            while (y /= 2) seg[x][y] = seg[x][2*y] +
   seg[x][2*y+1];
cbb
       }
cbb }
```

#### 1.15 SegTree Beats

```
// \text{ query}(a, b) - \{\{\min(v[a..b]), \max(v[a..b])\}, \sup(v[a..b])\}
// updatemin(a, b, x) faz com que v[i] \leftarrow min(v[i], x),
// para i em [a, b]
// updatemax faz o mesmo com max, e updatesum soma x
// em todo mundo do intervalo [a, b]
//
// Complexidades:
// build - O(n)
// query - O(log(n))
// update - O(\log^2(n)) amortizado
// (se nao usar updatesum, fica log(n) amortizado)
// 41672b
7c6 #define f first
Oab #define s second
f39 namespace beats {
3 c 9
        struct node {
526
            int tam;
            ll sum, lazy; // lazy pra soma
125
4f3
            ll mi1, mi2, mi; // mi = #mi1
            ll ma1, ma2, ma; // ma = #ma1
c61
426
             node(11 x = 0) {
ba6
                 sum = mi1 = ma1 = x;
                 mi2 = LINF, ma2 = -LINF;
b29
62 c
                 mi = ma = tam = 1;
c60
                lazv = 0;
```

```
}
cbb
            node(const node& 1, const node& r) {
770
a95
                 sum = 1.sum + r.sum, tam = 1.tam + r.tam;
c60
                lazv = 0;
797
                if (1.mi1 > r.mi1) {
230
                     mi1 = r.mi1, mi = r.mi;
                     mi2 = min(1.mi1, r.mi2);
ea2
                } else if (1.mi1 < r.mi1) {</pre>
dcd
e34
                     mi1 = l.mi1, mi = l.mi;
4b3
                     mi2 = min(r.mi1, 1.mi2);
9d9
                } else {
a39
                     mi1 = l.mi1, mi = l.mi+r.mi;
83 d
                     mi2 = min(1.mi2, r.mi2);
                }
cbb
cd0
                if (1.ma1 < r.ma1) {
6a0
                     ma1 = r.ma1, ma = r.ma;
96 d
                     ma2 = max(1.ma1, r.ma2);
                } else if (1.ma1 > r.ma1) {
5f0
                     ma1 = l.ma1, ma = l.ma;
ae0
                     ma2 = max(r.ma1, l.ma2);
2ca
                } else {
9d9
db2
                     ma1 = 1.ma1, ma = 1.ma+r.ma;
c05
                     ma2 = max(1.ma2, r.ma2);
                }
cbb
            }
cbb
            void setmin(ll x) {
4b4
55e
                if (x >= ma1) return;
                sum += (x - ma1)*ma;
463
be5
                if (mi1 == ma1) mi1 = x;
                if (mi2 == ma1) mi2 = x;
0a0
b81
                ma1 = x:
cbb
6cb
            void setmax(ll x) {
                if (x <= mi1) return;</pre>
e25
7e8
                sum += (x - mi1)*mi;
0bb
                if (ma1 == mi1) ma1 = x;
                if (ma2 == mi1) ma2 = x:
c32
1ff
                mi1 = x:
            }
cbb
            void setsum(ll x) {
4cf
fe8
                mi1 += x, mi2 += x, ma1 += x, ma2 += x;
620
                 sum += x*tam:
```

```
c46
                lazv += x:
cbb
            }
214
        };
62b
        node seg[4*MAX];
052
        int n, *v;
93b
        node build(int p=1, int l=0, int r=n-1) {
            if (1 == r) return seg[p] = {v[1]};
d84
ee4
            int m = (1+r)/2;
3 d6
            return seg[p] = \{build(2*p, 1, m), build(2*p+1,
   m+1, r);
       }
cbb
0 d8
        void build(int n2, int* v2) {
            n = n2, v = v2;
680
6f2
            build():
cbb
        }
ceb
        void prop(int p, int 1, int r) {
            if (1 == r) return;
8 ce
abd
            for (int k = 0; k < 2; k++) {
d07
                if (seg[p].lazy)
   seg[2*p+k].setsum(seg[p].lazy);
843
                seg[2*p+k].setmin(seg[p].ma1);
f79
                 seg[2*p+k].setmax(seg[p].mi1);
cbb
            }
431
            seg[p].lazv = 0;
cbb
055
        pair <pair <11, 11>, 11> query (int a, int b, int p=1,
   int 1=0, int r=n-1) {
e07
            if (b < 1 or r < a) return {{LINF, -LINF}, 0};</pre>
            if (a \le 1 \text{ and } r \le b) \text{ return } \{\{seg[p].mi1,
   seg[p].ma1}, seg[p].sum};
6 b 9
            prop(p, 1, r);
            int m = (1+r)/2:
ee4
            auto L = query(a, b, 2*p, 1, m), R = query(a, b,
   2*p+1, m+1, r):
96d
            return {{min(L.f.f, R.f.f), max(L.f.s, R.f.s)},
   L.s+R.s}:
cbb
2 c8
        node updatemin(int a, int b, ll x, int p=1, int l=0,
   int r=n-1) {
744
            if (b < l or r < a or seg[p].ma1 <= x) return
```

```
seg[p];
309
             if (a \le 1 \text{ and } r \le b \text{ and } seg[p].ma2 < x) {
                 seg[p].setmin(x);
ccd
534
                 return seg[p];
             }
cbb
6b9
             prop(p, l, r);
             int m = (1+r)/2;
ee4
96a
             return seg[p] = \{updatemin(a, b, x, 2*p, 1, m),
                               updatemin(a, b, x, 2*p+1, m+1,
faf
   r)};
cbb
044
        node updatemax(int a, int b, 11 x, int p=1, int 1=0,
   int r=n-1) {
b59
             if (b < l or r < a or seg[p].mi1 >= x) return
   seg[p];
a9e
             if (a \le 1 \text{ and } r \le b \text{ and } seg[p].mi2 > x) {
                 seg[p].setmax(x);
e8a
534
                 return seg[p];
cbb
6b9
             prop(p, 1, r);
ee4
             int m = (1+r)/2;
ee3
             return seg[p] = \{updatemax(a, b, x, 2*p, 1, m),
                               updatemax(a, b, x, 2*p+1, m+1,
bd2
   r)};
        }
cbb
        node updatesum(int a, int b, ll x, int p=1, int l=0,
aee
   int r=n-1) {
e9f
             if (b < l or r < a) return seg[p];</pre>
             if (a \le 1 \text{ and } r \le b) {
9a3
                 seg[p].setsum(x);
8f4
534
                 return seg[p];
cbb
6b9
             prop(p, 1, r);
             int m = (1+r)/2:
ee4
7b6
             return seg[p] = \{updatesum(a, b, x, 2*p, 1, m),
                               updatesum(a, b, x, 2*p+1, m+1,
ddb
   r)};
cbb
214 };
```

#### 1.16 SegTree Colorida

```
// Cada posicao tem um valor e uma cor
// O construtor receve um vector de {valor, cor}
// e o numero de cores (as cores devem estar em [0, c-1])
// query(c, a, b) retorna a soma dos valores
// de todo mundo em [a, b] que tem cor c
// update(c, a, b, x) soma x em todo mundo em
// [a, b] que tem cor c
// paint(c1, c2, a, b) faz com que todo mundo
// em [a, b] que tem cor c1 passe a ter cor c2
//
// Complexidades:
// construir - O(n \log(n)) espaco e tempo
// query - O(log(n))
// update - O(log(n))
// paint - O(log(n)) amortizado
// 2938e8
04f struct seg_color {
3 c 9
        struct node {
b19
            node *1, *r;
0f9
            int cnt;
9ca
            ll val, lazy;
277
            node() : l(NULL), r(NULL), cnt(0), val(0),
   lazy(0) {}
01e
            void update() {
d0a
                cnt = 0, val = 0;
bc4
                for (auto i : {1, r}) if (i) {
c89
                    i->prop();
281
                     cnt += i->cnt, val += i->val;
                }
cbb
            }
cbb
a9c
            void prop() {
2 dd
                if (!lazy) return;
3f7
                val += lazy*(ll)cnt;
                for (auto i : {1, r}) if (i) i->lazy += lazy;
b64
c60
                lazy = 0;
cbb
            }
214
        };
1 a8
        int n;
9b0
        vector < node * > seg;
```

```
6e0
        seg_color(vector<pair<int, int>>& v, int c) :
   n(v.size()), seg(c, NULL) {
830
            for (int i = 0; i < n; i++)
                seg[v[i].second] = insert(seg[v[i].second],
9b7
   i, v[i].first, 0, n-1);
cbb
3 c 7
        \simseg color() {
dde
            queue < node *> q;
3a6
            for (auto i : seg) q.push(i);
402
            while (q.size()) {
20b
                auto i = q.front(); q.pop();
                if (!i) continue;
dab
7 c 7
                q.push(i->1), q.push(i->r);
5ce
                delete i;
            }
cbb
cbb
        }
        node* insert(node* at, int idx, int val, int l, int
40b
r) {
            if (!at) at = new node();
1a4
            if (1 == r) return at->cnt = 1, at->val = val,
232
   at;
ee4
            int m = (1+r)/2;
            if (idx \le m) at->1 = insert(at->1, idx, val, 1,
137
  m);
3e6
            else at->r = insert(at->r, idx, val, m+1, r);
cff
            return at->update(), at;
cbb
870
        11 query(node* at, int a, int b, int 1, int r) {
61b
            if (!at or b < 1 or r < a) return 0;
d9f
            at->prop();
cb2
           if (a <= l and r <= b) return at->val;
           int m = (1+r)/2;
ee4
            return query(at->1, a, b, 1, m) + query(at->r,
4c4
   a, b, m+1, r);
cbb
e54
        11 query(int c, int a, int b) { return query(seg[c],
   a, b, 0, n-1); }
        void update(node* at, int a, int b, int x, int 1,
91c
   int r) {
            if (!at or b < l or r < a) return;
fba
d9f
            at->prop();
```

```
9a3
            if (a \le 1 \text{ and } r \le b)
e9a
                 at -> lazv += x;
cb2
                 return void(at->prop());
cbb
ee4
            int m = (1+r)/2;
0 b 0
            update(at->1, a, b, x, 1, m), update(at->r, a,
   b, x, m+1, r);
7b4
            at->update();
cbb
        void update(int c, int a, int b, int x) {
a40
   update(seg[c], a, b, x, 0, n-1); }
        void paint(node*& from, node*& to, int a, int b, int
   1, int r) {
            if (to == from or !from or b < 1 or r < a)
10f
   return;
e85
            from ->prop();
889
            if (to) to->prop();
9a3
            if (a \le 1 \text{ and } r \le b) {
24d
                 if (!to) {
38f
                     to = from;
140
                     from = NULL;
505
                     return;
cbb
ee4
                int m = (1+r)/2;
1 cb
                 paint(from->1, to->1, a, b, 1, m),
   paint(from -> r, to -> r, a, b, m+1, r);
72d
                to->update();
270
                 delete from;
140
                from = NULL;
505
                 return;
cbb
            }
019
            if (!to) to = new node();
ee4
            int m = (1+r)/2;
1 cb
            paint(from->1, to->1, a, b, 1, m),
   paint(from->r, to->r, a, b, m+1, r);
45a
            from->update(), to->update();
cbb
471
        void paint(int c1, int c2, int a, int b) {
   paint(seg[c1], seg[c2], a, b, 0, n-1); }
214 };
```

#### 1.17 SegTree Esparsa - Lazy

```
// Query: soma do range [a, b]
// Update: flipa os valores de [a, b]
// O MAX tem q ser Q log N para Q updates
//
// Complexidades:
// build - 0(1)
// query - O(log(n))
// update - O(\log(n))
// dc37e6
aa4 namespace seg {
6de
        int seg[MAX], lazy[MAX], R[MAX], L[MAX], ptr;
e9a
        int get_l(int i){
3db
            if (L[i] == 0) L[i] = ptr++;
a96
            return L[i];
cbb
943
        int get_r(int i){
71b
            if (R[i] == 0) R[i] = ptr++;
283
            return R[i];
cbb
        }
e71
        void build() { ptr = 2; }
ceb
        void prop(int p, int 1, int r) {
b77
            if (!lazy[p]) return;
            seg[p] = r-l+1 - seg[p];
76 c
213
            if (1 != r) lazy[get_l(p)]^=lazy[p],
   lazy[get_r(p)]^=lazy[p];
3c7
            lazv[p] = 0;
        }
cbb
158
        int query(int a, int b, int p=1, int l=0, int r=N-1)
  {
6b9
            prop(p, 1, r);
786
            if (b < l or r < a) return 0;
            if (a <= 1 and r <= b) return seg[p];</pre>
527
           int m = (1+r)/2;
ee4
            return query(a, b, get_l(p), l, m)+query(a, b,
818
   get_r(p), m+1, r);
```

```
cbb
       }
51f
        int update(int a, int b, int p=1, int l=0, int
   r = N - 1) {
6 b 9
            prop(p, 1, r);
e9f
            if (b < l or r < a) return seg[p];
            if (a \le 1 \text{ and } r \le b)
9 a 3
ab6
                lazy[p] ^= 1;
6 b 9
                prop(p, 1, r);
534
                return seg[p];
cbb
            }
ee4
            int m = (1+r)/2;
            return seg[p] = update(a, b, get_l(p), l,
   m)+update(a, b, get_r(p), m+1, r);
cbb
214 };
1.18 SegTree Esparsa - O(q) memoria
```

```
// Query: min do range [a, b]
// Update: troca o valor de uma posicao
// Usa O(q) de memoria para q updates
//
// Complexidades:
// query - O(log(n))
// update - O(log(n))
// 072a21
13d template < typename T> struct seg {
3 c 9
        struct node {
d53
            node* ch[2]:
970
            char d;
ca0
            T v;
c4e
            T mi;
            node(int d_, T v_, T val) : d(d_), v(v_) {
d4e
                 ch[0] = ch[1] = NULL;
e71
d6e
                mi = val;
cbb
b32
            node(node* x) : d(x->d), v(x->v), mi(x->mi) {
c99
                 ch[0] = x -> ch[0], ch[1] = x -> ch[1];
```

```
cbb
01e
            void update() {
909
                mi = numeric limits <T>::max();
151
                for (int i = 0; i < 2; i++) if (ch[i])
                     mi = min(mi, ch[i]->mi);
b5a
cbb
214
        };
bb7
        node* root;
9 c 5
        char n;
        seg() : root(NULL), n(0) {}
ba7
512
        \simseg() {
4c0
            std::vector<node*> q = {root};
402
            while (q.size()) {
                node* x = q.back(); q.pop_back();
e5d
                if (!x) continue;
ee9
                q.push_back(x->ch[0]), q.push_back(x->ch[1]);
73f
                delete x;
bf0
            }
cbb
cbb
        }
        char msb(T v, char l, char r) { // msb in range (l,
1a6
  r]
8e4
            for (char i = r; i > l; i--) if (v > i \& 1) return
  i;
            return -1;
daa
cbb
        void cut(node* at, T v, char i) {
430
677
            char d = msb(v ^a at -> v, at -> d, i);
23b
            if (d == -1) return; // no need to split
            node* nxt = new node(at);
ebf
            at -> ch[v>>d&1] = NULL;
d43
            at - ch[!(v > d&1)] = nxt:
34f
150
            at->d=d:
cbb
        }
6e5
        node* update(node* at, T idx, T val, char i) {
            if (!at) return new node(-1, idx, val);
c8c
d67
            cut(at, idx, i);
            if (at -> d == -1) { // leaf }
1a2
792
                at->mi = val;
```

```
ce6
                return at:
cbb
b29
            bool dir = idx>>at->d&1;
c8f
            at->ch[dir] = update(at->ch[dir], idx, val,
   at - > d - 1);
7 b 4
            at->update();
ce6
            return at;
cbb
85 c
        void update(T idx, T val) {
8f4
            while (idx >> n) n++;
61e
            root = update(root, idx, val, n-1);
cbb
9 d8
        T query(node* at, T a, T b, T l, T r, char i) {
            if (!at or b < l or r < a) return
   numeric limits <T>::max():
fd3
            if (a <= l and r <= b) return at->mi;
            T m = 1 + (r-1)/2:
841
            if (at->d < i) {</pre>
c85
                if ((at->v>>i&1) == 0) return query(at, a,
c59
   b, 1, m, i-1);
ca4
                else return query(at, a, b, m+1, r, i-1);
cbb
373
            return min(query(at->ch[0], a, b, 1, m, i-1),
   query(at->ch[1], a, b, m+1, r, i-1));
cbb
6f6
        T query(T 1, T r) { return query(root, 1, r, 0,
   (1 << n) -1, n-1);
214 };
1.19 SegTree Iterativa
// Consultas 0-based
// Valores iniciais devem estar em (seg[n], ..., seg[2*n-1])
// Query: soma do range [a, b]
// Update: muda o valor da posicao p para x
//
// Complexidades:
// build - O(n)
```

// query - O(log(n))

// 779519

// update -  $O(\log(n))$ 

```
6a4 int seg[2 * MAX];
1a8 int n;
0a8 void build() {
        for (int i = n - 1; i; i--) seg[i] = seg[2*i] +
   seg[2*i+1];
cbb }
4ea int query(int a, int b) {
7c9
        int ret = 0:
728
        for(a += n, b += n; a <= b; ++a /= 2, --b /= 2) {
            if (a % 2 == 1) ret += seg[a];
4ea
           if (b \% 2 == 0) ret += seg[b];
244
cbb
edf
        return ret;
cbb }
ff3 void update(int p, int x) {
        seg[p += n] = x;
37 d
        while (p /= 2) seg[p] = seg[2*p] + seg[2*p+1];
c8c
cbb }
```

#### 1.20 SegTree Iterativa com Lazy Propagation

```
// Query: soma do range [a, b]
// Update: soma x em cada elemento do range [a, b]
// Para mudar, mudar as funcoes junta, poe e query
// LOG = ceil(log2(MAX))
//
// Complexidades:
// build - O(n)
// query - O(log(n))
// update - O(\log(n))
// 6dc475
aa4 namespace seg {
6db
        11 \text{ seg}[2*MAX], lazy[2*MAX];
1a8
        int n;
9b3
        ll junta(ll a, ll b) {
534
            return a+b;
```

```
cbb
        }
d41
        // soma x na posicao p de tamanho tam
1 b 4
        void poe(int p, ll x, int tam, bool prop=1) {
517
            seg[p] += x*tam;
6ae
            if (prop and p < n) lazy[p] += x;</pre>
cbb
d41
        // atualiza todos os pais da folha p
b1e
        void sobe(int p) {
            for (int tam = 2; p /= 2; tam *= 2) {
d5a
4 ca
                seg[p] = junta(seg[2*p], seg[2*p+1]);
388
                poe(p, lazy[p], tam, 0);
            }
cbb
        }
cbb
        // propaga o caminho da raiz ate a folha p
d41
        void prop(int p) {
a0a
076
            int tam = 1 << (LOG-1);</pre>
            for (int s = LOG; s; s--, tam /= 2) {
0a8
4b1
                int i = p \gg s;
27 c
                if (lazy[i]) {
860
                     poe(2*i, lazy[i], tam);
                     poe(2*i+1, lazy[i], tam);
e38
                     lazy[i] = 0;
b97
cbb
cbb
            }
cbb
        }
61c
        void build(int n2, int* v) {
1 e3
            n = n2:
95f
            for (int i = 0; i < n; i++) seg[n+i] = v[i];
c 4 1
            for (int i = n-1; i; i--) seg[i] =
   junta(seg[2*i], seg[2*i+1]);
            for (int i = 0; i < 2*n; i++) lazy[i] = 0;
f4c
cbb
4f3
        11 query(int a, int b) {
b73
           11 \text{ ret} = 0;
            for (prop(a+=n), prop(b+=n); a <= b; ++a/=2,
b48
   --b/=2) {
                if (a%2 == 1) ret = junta(ret, seg[a]);
a8e
```

```
c58
                if (b%2 == 0) ret = junta(ret, seg[b]);
cbb
edf
            return ret;
cbb
        }
a28
        void update(int a, int b, int x) {
            int a2 = a += n, b2 = b += n, tam = 1;
c2d
            for (; a <= b; ++a/=2, --b/=2, tam *= 2) {
Off
                if (a\%2 == 1) poe(a, x, tam);
32a
                if (b\%2 == 0) poe(b, x, tam);
9da
cbb
0 f 7
            sobe(a2), sobe(b2);
cbb
214 };
```

#### 1.21 SegTree PA

```
// Segtree de PA
// update_set(1, r, A, R) seta [1, r] para PA(A, R),
// update_add soma PA(A, R) em [1, r]
// query(1, r) retorna a soma de [1, r]
//
// PA(A, R) eh a PA: [A+R, A+2R, A+3R, ...]
//
// Complexidades:
// construir - O(n)
// update_set, update_add, query - O(log(n))
// bc4746
dc7 struct seg_pa {
350
        struct Data {
8f5
            ll sum:
662
            ll set_a, set_r, add_a, add_r;
9b7
            Data() : sum(0), set_a(LINF), set_r(0),
   add_a(0), add_r(0) {}
214
        };
16a
        vector < Data > seg;
1a8
        int n;
        seg_pa(int n_) {
d45
e95
            n = n_{\cdot};
fc3
            seg = vector < Data > (4*n);
```

```
cbb
        }
        void prop(int p, int 1, int r) {
ceb
            int tam = r-l+1;
d5a
            11 &sum = seg[p].sum, &set_a = seg[p].set_a,
c3f
   &set_r = seg[p].set_r
a1b
                &add_a = seg[p].add_a, &add_r = seg[p].add_r;
c02
            if (set a != LINF) {
660
                set_a += add_a, set_r += add_r;
                sum = set_a*tam + set_r*tam*(tam+1)/2;
06e
579
                if (1 != r) {
                    int m = (1+r)/2;
ee4
886
                    seg[2*p].set_a = set_a;
358
                    seg[2*p].set_r = set_r;
                    seg[2*p].add_a = seg[2*p].add_r = 0;
ed6
                    seg[2*p+1].set_a = set_a + set_r *
f0c
   (m-1+1);
471
                    seg[2*p+1].set_r = set_r;
d48
                    seg[2*p+1].add_a = seg[2*p+1].add_r = 0;
cbb
823
                set_a = LINF, set_r = 0;
953
                add_a = add_r = 0;
            } else if (add a or add r) {
105
18b
                sum += add_a*tam + add_r*tam*(tam+1)/2;
579
                if (1 != r) {
                    int m = (1+r)/2;
ee4
ff0
                    seg[2*p].add_a += add_a;
ec0
                    seg[2*p].add_r += add_r;
                    seg[2*p+1].add_a += add_a + add_r *
06c
   (m-1+1);
a6d
                    seg[2*p+1].add_r += add_r;
cbb
953
                add_a = add_r = 0;
            }
cbb
        }
cbb
0 b 7
        int inter(pair<int, int> a, pair<int, int> b) {
```

```
98 c
            if (a.first > b.first) swap(a, b);
eef
            return max(0, min(a.second, b.second) - b.first
 + 1);
cbb
be1
        11 set(int a, int b, ll aa, ll rr, int p, int l, int
 r) {
6b9
            prop(p, 1, r);
            if (b < 1 or r < a) return seg[p].sum;</pre>
457
            if (a <= 1 and r <= b) {</pre>
9a3
91c
                seg[p].set_a = aa;
                seg[p].set_r = rr;
774
6b9
                prop(p, 1, r);
254
                return seg[p].sum;
            }
cbb
ee4
            int m = (1+r)/2;
963
            int tam_l = inter({1, m}, {a, b});
c34
            return seg[p].sum = set(a, b, aa, rr, 2*p, 1, m)
                set(a, b, aa + rr * tam_1, rr, 2*p+1, m+1,
365
  r);
cbb
        }
f55
        void update_set(int 1, int r, 11 aa, 11 rr) {
6f7
            set(1, r, aa, rr, 1, 0, n-1);
cbb
5f6
        ll add(int a, int b, ll aa, ll rr, int p, int l, int
  r) {
6b9
            prop(p, 1, r);
457
            if (b < l or r < a) return seg[p].sum;</pre>
            if (a \le 1 \text{ and } r \le b)
9a3
359
                seg[p].add_a += aa;
                seg[p].add_r += rr;
1ee
6b9
                prop(p, 1, r);
254
                return seg[p].sum;
            }
cbb
ee4
            int m = (1+r)/2:
963
            int tam_l = inter({1, m}, {a, b});
586
            return seg[p].sum = add(a, b, aa, rr, 2*p, 1, m)
695
                add(a, b, aa + rr * tam_l, rr, 2*p+1, m+1,
   r);
cbb
848
        void update_add(int 1, int r, 11 aa, 11 rr) {
```

```
afa
            add(1, r, aa, rr, 1, 0, n-1);
cbb
f45
        ll query(int a, int b, int p, int l, int r) {
6 b 9
            prop(p, 1, r);
786
            if (b < l or r < a) return 0;
e9a
            if (a <= 1 and r <= b) return seg[p].sum;
ee4
            int m = (1+r)/2;
b1f
            return query(a, b, 2*p, 1, m) + query(a, b,
   2*p+1, m+1, r);
cbb
        ll query(int 1, int r) { return query(1, r, 1, 0,
214 };
```

#### 1.22 SegTree Persistente

```
// SegTree de soma, update de somar numa posicao
//
// query(a, b, t) retorna a query de [a, b] na versao t
// update(a, x, t) faz um update v[a]+=x a partir da
// versao de t, criando uma nova versao e retornando seu id
// Por default, faz o update a partir da ultima versao
//
// build - O(n)
// query - O(log(n))
// update - O(log(n))
// 50ab73
54a const int MAX = 1e5+10, UPD = 1e5+10, LOG = 18;
6de const int MAXS = 2*MAX+UPD*LOG;
f6e namespace perseg {
bd6
        11 seg[MAXS];
        int rt[UPD], L[MAXS], R[MAXS], cnt, t;
f4e
052
        int n, *v;
3 c 4
        11 build(int p, int 1, int r) {
            if (1 == r) return seg[p] = v[1];
6cd
855
            L[p] = cnt++, R[p] = cnt++;
ee4
            int m = (1+r)/2;
275
            return seg[p] = build(L[p], 1, m) + build(R[p],
   m+1, r);
```

```
}
cbb
0d8
        void build(int n2, int* v2) {
680
           n = n2, v = v2;
856
           rt[0] = cnt++;
           build(0, 0, n-1);
c50
cbb
       }
f45
       11 query(int a, int b, int p, int l, int r) {
786
           if (b < l or r < a) return 0;
527
           if (a <= 1 and r <= b) return seg[p];</pre>
           int m = (1+r)/2;
ee4
           return query(a, b, L[p], 1, m) + query(a, b,
1ed
   R[p], m+1, r);
       }
cbb
182
        ll query(int a, int b, int tt) {
           return query(a, b, rt[tt], 0, n-1);
c13
cbb
       11 update(int a, int x, int lp, int p, int l, int r)
bb3
  {
           if (l == r) return seg[p] = seg[lp]+x;
747
           int m = (1+r)/2;
ee4
ab8
           if (a <= m)
b48
               return seg[p] = update(a, x, L[lp],
  return seg[p] = seg[L[p]=L[lp]] + update(a, x,
8a9
   R[lp], R[p] = cnt++, m+1, r);
cbb
6f6
       int update(int a, int x, int tt=t) {
ab3
           update(a, x, rt[tt], rt[++t]=cnt++, 0, n-1);
e0d
           return t;
cbb
       }
214 };
      Sparse Table
// Resolve RMQ
// MAX2 = log(MAX)
```

```
// Resolve RMQ
// MAX2 = log(MAX)
//
// Complexidades:
// build - O(n log(n))
// query - O(1)
// 7aa4c9
```

```
cca namespace sparse {
710
        int m[MAX2][MAX], n;
61c
        void build(int n2, int* v) {
1 e3
            n = n2;
78e
            for (int i = 0; i < n; i++) m[0][i] = v[i];
            for (int j = 1; (1<<j) <= n; j++) for (int i =
   0; i+(1 << j) <= n; i++)
                m[j][i] = min(m[j-1][i],
5d5
   m[j-1][i+(1<<(j-1))]);
cbb
4ea
        int query(int a, int b) {
            int j = __builtin_clz(1) - __builtin_clz(b-a+1);
ee5
            return min(m[j][a], m[j][b-(1<<j)+1]);</pre>
dc3
cbb
cbb }
```

## 1.24 Sparse Table Disjunta

```
// Resolve qualquer operacao associativa
// MAX2 = log(MAX)
//
// Complexidades:
// build - O(n log(n))
// query - O(1)
// fd81ae
cca namespace sparse {
9bf
        int m[MAX2][2*MAX], n, v[2*MAX];
5f7
        int op(int a, int b) { return min(a, b); }
0 d8
        void build(int n2, int* v2) {
1 e3
            n = n2:
df4
            for (int i = 0; i < n; i++) v[i] = v2[i];
a84
            while (n&(n-1)) n++;
3 d2
            for (int j = 0; (1<<j) < n; j++) {
                int len = 1<<j;</pre>
1 c 0
                for (int c = len; c < n; c += 2*len) {
d9b
                     m[j][c] = v[c], m[j][c-1] = v[c-1];
332
668
                     for (int i = c+1; i < c+len; i++)</pre>
   m[j][i] = op(m[j][i-1], v[i]);
                     for (int i = c-2; i >= c-len; i--)
432
   m[j][i] = op(v[i], m[j][i+1]);
cbb
```

```
}
                                                                   891
                                                                           \simsplaytree() {
cbb
                                                                   609
                                                                                vector < node *> q = {root};
cbb
                                                                                while (q.size()) {
9e3
        int query(int 1, int r) {
                                                                   402
f13
            if (1 == r) return v[1];
                                                                   e5d
                                                                                    node* x = q.back(); q.pop_back();
            int j = __builtin_clz(1) - __builtin_clz(l^r);
                                                                   ee9
                                                                                    if (!x) continue;
e6d
                                                                                    q.push_back(x->ch[0]), q.push_back(x->ch[1]);
d67
            return op(m[j][1], m[j][r]);
                                                                   73f
                                                                   bf0
cbb
        }
                                                                                    delete x;
cbb }
                                                                               }
                                                                   cbb
                                                                           }
                                                                   cbb
1.25 Splay Tree
                                                                   94f
                                                                            void rotate(node* x) { // x vai ficar em cima
                                                                   d9b
                                                                                node *p = x->p, *pp = p->p;
// SEMPRE QUE DESCER NA ARVORE, DAR SPLAY NO
                                                                                if (pp) pp->ch[pp->ch[1] == p] = x;
                                                                   ecf
// NODE MAIS PROFUNDO VISITADO
                                                                   286
                                                                                bool d = p \rightarrow ch[0] == x;
// Todas as operacoes sao O(\log(n)) amortizado
                                                                                p - > ch[!d] = x - > ch[d], x - > ch[d] = p;
                                                                   d63
// Se quiser colocar mais informação no node,
                                                                                if (p->ch[!d]) p->ch[!d]->p = p;
                                                                   bad
// mudar em 'update'
                                                                               x->p = pp, p->p = x;
                                                                   fc2
// 4ff2b3
                                                                                p->update(), x->update();
                                                                   1 ea
                                                                   cbb
538 template < typename T > struct splaytree {
                                                                   3fa
                                                                           node* splay(node* x) {
3c9
        struct node {
                                                                                if (!x) return x;
                                                                   a39
183
            node *ch[2], *p;
                                                                   4ea
                                                                                root = x;
e4d
            int sz;
                                                                   3cf
                                                                                while (x->p) {
f48
            T val;
                                                                   d9b
                                                                                    node *p = x->p, *pp = p->p;
            node(T v) {
da0
                                                                                    if (!pp) return rotate(x), x; // zig
                                                                   359
696
                 ch[0] = ch[1] = p = NULL;
                                                                                    if ((pp->ch[0] == p)^(p->ch[0] == x))
                                                                   e3c
a26
                 sz = 1:
                                                                   a2b
                                                                                        rotate(x), rotate(x); // zigzag
250
                 val = v;
                                                                   4b2
                                                                                    else rotate(p), rotate(x); // zigzig
cbb
                                                                   cbb
                                                                                }
01e
            void update() {
                                                                   ea5
                                                                                return x;
a26
                 sz = 1;
                                                                   cbb
                 for (int i = 0; i < 2; i++) if (ch[i]) {</pre>
с7с
                                                                   313
                                                                           node* insert(T v, bool lb=0) {
d5f
                     sz += ch[i]->sz:
                                                                                if (!root) return lb ? NULL : root = new node(v);
                                                                   b64
cbb
                 }
                                                                                node *x = root, *last = NULL;;
                                                                   002
            }
cbb
                                                                   31e
                                                                                while (1) {
214
        };
                                                                   5 d.7
                                                                                    bool d = x -> val < v;
                                                                   0fd
                                                                                    if (!d) last = x;
bb7
        node* root;
```

fbc

214

cbf

cbb

}

splaytree() { root = NULL; }

splaytree(const splaytree& t) {

throw logic\_error("Nao copiar a splaytree!");

c2e

c16

4e6

dea

055

if (x->val == v) break;

if (lb) break;

else {

if (x->ch[d]) x = x->ch[d];

x - ch[d] = new node(v);

```
99 c
                     x -> ch[d] -> p = x;
30e
                     x = x -> ch[d];
c2b
                     break:
                }
cbb
            }
cbb
0b6
            splay(x);
61 c
            return lb ? splay(last) : x;
cbb
сОс
        int size() { return root ? root->sz : 0; }
        int count(T v) { return insert(v, 1) and root->val
2ca
   == v: 
        node* lower_bound(T v) { return insert(v, 1); }
111
26 b
        void erase(T v) {
            if (!count(v)) return;
446
            node *x = root, *1 = x -> ch[0];
bce
268
            if (!1) {
8 b 1
                root = x - > ch[1];
32e
                if (root) root->p = NULL;
8f3
                return delete x;
            }
cbb
5e7
            root = 1, 1->p = NULL;
902
            while (1->ch[1]) 1 = 1->ch[1];
bab
            splay(1);
            1 - ch[1] = x - ch[1];
f0e
            if (1->ch[1]) 1->ch[1]->p = 1;
7d9
bf0
            delete x;
62a
            1->update();
cbb
24a
        int order_of_key(T v) {
62b
            if (!lower_bound(v)) return root ? root->sz : 0;
            return root->ch[0] ? root->ch[0]->sz : 0;
1cc
cbb
        node* find_by_order(int k) {
db6
            if (k >= size()) return NULL;
084
52f
            node* x = root;
31e
            while (1) {
                if (x->ch[0] \text{ and } x->ch[0]->sz >= k+1) x =
20 f
   x - > ch[0]:
4e6
                 else {
a1c
                     if (x->ch[0]) k -= x->ch[0]->sz;
1dc
                     if (!k) return splay(x);
eb8
                     k--, x = x->ch[1];
```

#### 1.26 Splay Tree Implicita

```
// vector da NASA
// Um pouco mais rapido q a treap
// O construtor a partir do vector
// eh linear, todas as outras operacoes
// custam O(log(n)) amortizado
// a3575a
081 template < typename T> struct splay {
3 c 9
        struct node {
183
            node *ch[2], *p;
e4d
            int sz;
875
            T val, sub, lazy;
aa6
            bool rev;
da0
            node(T v) {
                ch[0] = ch[1] = p = NULL;
696
a26
                sz = 1;
1 e 4
                sub = val = v;
c60
                lazy = 0;
b67
                rev = false:
cbb
a9c
            void prop() {
0ec
                if (lazy) {
                     val += lazy, sub += lazy*sz;
924
                     if (ch[0]) ch[0]->lazy += lazy;
091
1 a8
                     if (ch[1]) ch[1]->lazy += lazy;
cbb
1 bb
                if (rev) {
80a
                     swap(ch[0], ch[1]);
628
                     if (ch[0]) ch[0]->rev ^= 1;
adc
                     if (ch[1]) ch[1] -> rev ^= 1;
```

```
cbb
                                                                             int size(node* x) { return x ? x->sz : 0; }
a32
                 lazv = 0, rev = 0;
                                                                    73с
            }
                                                                             void rotate(node* x) { // x vai ficar em cima
cbb
                                                                    94f
01e
            void update() {
                                                                    d9b
                                                                                 node *p = x->p, *pp = p->p;
0 c 3
                 sz = 1, sub = val;
                                                                                 if (pp) pp -> ch[pp -> ch[1] == p] = x;
                                                                    ecf
с7с
                 for (int i = 0; i < 2; i++) if (ch[i]) {
                                                                    286
                                                                                 bool d = p -> ch[0] == x;
                     ch[i]->prop();
                                                                    d63
                                                                                 p - ch[!d] = x - ch[d], x - ch[d] = p;
05f
                                                                                 if (p->ch[!d]) p->ch[!d]->p = p;
                     sz += ch[i] -> sz;
d5f
                                                                    bad
                                                                                 x->p = pp, p->p = x;
                     sub += ch[i]->sub;
4a1
                                                                    fc2
                }
                                                                    1 ea
                                                                                 p->update(), x->update();
cbb
cbb
                                                                    cbb
214
        };
                                                                    6a0
                                                                             node* splaya(node* x) {
                                                                    a39
                                                                                 if (!x) return x;
                                                                                 root = x, x->update();
                                                                    be6
bb7
        node* root;
                                                                    3 cf
                                                                                 while (x->p) {
                                                                                     node *p = x->p, *pp = p->p;
5d9
        splay() { root = NULL; }
                                                                    d9b
9b1
        splay(node* x) {
                                                                    359
                                                                                     if (!pp) return rotate(x), x; // zig
                                                                                     if ((pp->ch[0] == p)^(p->ch[0] == x))
4ea
            root = x;
                                                                    e3c
32e
            if (root) root->p = NULL;
                                                                    a2b
                                                                                          rotate(x), rotate(x); // zigzag
                                                                    4b2
cbb
                                                                                     else rotate(p), rotate(x); // zigzig
1b7
        splay(vector < T > v) { // O(n)}
                                                                    cbb
                                                                                 }
950
            root = NULL;
                                                                    ea5
                                                                                 return x;
806
            for (T i : v) {
                                                                    cbb
                 node* x = new node(i);
                                                                    a7f
2a0
                                                                             node* find(int v) {
                                                                    a2e
bd1
                x -  ch[0] = root;
                                                                                 if (!root) return NULL;
                                                                    52f
37a
                if (root) root->p = x;
                                                                                 node *x = root;
4ea
                 root = x;
                                                                    6cd
                                                                                 int kev = 0;
                                                                    31e
                                                                                 while (1) {
a0a
                 root->update();
            }
                                                                    857
                                                                                     x - > prop();
cbb
cbb
                                                                    ba1
                                                                                     bool d = key + size(x->ch[0]) < v;
a9e
        splay(const splay& t) {
                                                                    877
                                                                                     if (\text{key} + \text{size}(x->\text{ch}[0]) != v \text{ and } x->\text{ch}[d]) {
                                                                                          if (d) key += size(x->ch[0])+1;
e62
            throw logic_error("Nao copiar a splay!");
                                                                    15e
        }
                                                                    30e
                                                                                         x = x -> ch[d];
cbb
        \simsplay() {
                                                                    9af
                                                                                     } else break:
5ab
609
            vector<node*> q = {root};
                                                                    cbb
                                                                                 }
402
            while (q.size()) {
                                                                    152
                                                                                 return splaya(x);
                 node* x = q.back(); q.pop_back();
e5d
                                                                    cbb
                                                                             int size() { return root ? root->sz : 0; }
                 if (!x) continue;
                                                                    сОс
ee9
                 q.push_back(x->ch[0]), q.push_back(x->ch[1]);
                                                                             void join(splay<T>& 1) { // assume que 1 < *this</pre>
73f
                                                                    c26
                                                                    690
                                                                                 if (!size()) swap(root, 1.root);
bf0
                 delete x;
            }
                                                                    579
                                                                                 if (!size() or !l.size()) return;
cbb
cbb
        }
                                                                    bee
                                                                                 node* x = 1.root;
```

```
31e
            while (1) {
857
                x->prop();
34d
                if (!x->ch[1]) break;
bd8
                 x = x -> ch[1];
            }
cbb
147
            1.splaya(x), root->prop(), root->update();
42b
            x -> ch[1] = root, x -> ch[1] -> p = x;
            root = l.root, l.root = NULL;
0aa
a0a
            root->update();
cbb
        }
5ed
        node* split(int v) { // retorna os elementos < v</pre>
398
            if (v <= 0) return NULL;</pre>
            if (v >= size()) {
060
                node* ret = root;
f87
950
                 root = NULL;
8c9
                 ret ->update();
edf
                 return ret;
            }
cbb
            find(v);
adc
            node*1 = root -> ch[0];
a59
4df
            root -> ch [0] = NULL;
5a3
            if (1) 1->p = NULL;
a0a
            root ->update();
792
            return 1;
        }
cbb
511
        T& operator [](int i) {
9d4
            find(i);
ae0
            return root->val;
cbb
231
        void push_back(T v) \{ // 0(1) \}
a 0 1
            node* r = new node(v);
            r -> ch[0] = root;
0de
b11
            if (root) root->p = r;
            root = r, root->update();
b13
cbb
b7a
        T query(int 1, int r) {
95f
            splay <T> M(split(r+1));
5ff
            splay <T> L(M.split(1));
            T ans = M.root->sub;
d1c
49 с
            M.join(L), join(M);
ba7
            return ans;
cbb
        }
```

```
41f
        void update(int 1, int r, T s) {
95f
             splav <T> M(split(r+1));
5ff
             splay <T> L(M.split(1));
996
            M.root->lazy += s;
49 c
            M. join(L), join(M);
cbb
8 c 1
        void reverse(int 1, int r) {
95f
             splay<T> M(split(r+1));
            splay<T> L(M.split(1));
5ff
945
            M.root->rev ^= 1;
            M.join(L), join(M);
49 c
cbb
2fb
        void erase(int 1, int r) {
95f
             splay<T> M(split(r+1));
5ff
            splay<T> L(M.split(1));
dcc
            join(L);
cbb
        }
214 };
```

#### 1.27 Split-Merge Set

```
// Representa um conjunto de inteiros nao negativos
// Todas as operacoes custam O(log(N)),
// em que N = maior elemento do set,
// exceto o merge, que custa O(log(N)) amortizado
// Usa O(\min(N, n \log(N))) de memoria, sendo 'n' o
// numero de elementos distintos no set
// 2d2d8a
2dc template < typename T, bool MULTI = false, typename
   SIZE T=int> struct sms {
3 c 9
        struct node {
b19
            node *1, *r;
15f
            SIZE_T cnt;
658
            node() : l(NULL), r(NULL), cnt(0) {}
01e
            void update() {
a01
                cnt = 0;
d8a
                if (1) cnt += 1->cnt;
e49
                if (r) cnt += r->cnt;
cbb
            }
214
        };
```

```
bb7
        node* root:
fd0
        T N;
f34
        sms() : root(NULL), N(0) {}
83b
        sms(T v) : sms() { while (v >= N) N = 2*N+1; }
        sms(const sms& t) : root(NULL), N(t.N) {
5e1
            for (SIZE T i = 0; i < t.size(); i++) {</pre>
3af
a0f
                T at = t[i];
                SIZE_T qt = t.count(at);
e6d
                insert(at, qt);
a43
f42
                i += qt-1;
            }
cbb
        }
cbb
        sms(initializer_list<T> v) : sms() { for (T i : v)
   insert(i): }
        \simsms() {
2dd
609
            vector<node*> q = {root};
            while (q.size()) {
402
                node* x = q.back(); q.pop_back();
e5d
                if (!x) continue;
ee9
                q.push_back(x->1), q.push_back(x->r);
1 c 7
bf0
                delete x;
            }
cbb
cbb
        }
fdc
        friend void swap(sms& a, sms& b) {
49e
            swap(a.root, b.root), swap(a.N, b.N);
cbb
83 e
        sms& operator =(const sms& v) {
768
            sms tmp = v;
            swap(tmp, *this);
420
357
            return *this;
cbb
d06
        SIZE_T size() const { return root ? root->cnt : 0; }
        SIZE_T count(node* x) const { return x ? x->cnt : 0;
17 f
  }
75a
        void clear() {
0a0
            sms tmp;
            swap(*this, tmp);
4ac
cbb
a06
        void expand(T v) {
bc3
            for (; N < v; N = 2*N+1) if (root) {
```

```
63 c
                node* nroot = new node();
956
                nroot ->1 = root;
897
                root = nroot;
a0a
                root ->update();
cbb
            }
cbb
        }
b14
        node* insert(node* at, T idx, SIZE_T qt, T 1, T r) {
1a4
            if (!at) at = new node();
893
            if (1 == r) {
435
                at->cnt += qt;
beb
                if (!MULTI) at->cnt = 1;
ce6
                return at;
            }
cbb
841
            T m = 1 + (r-1)/2;
            if (idx \le m) at->1 = insert(at->1, idx, qt, 1,
a 02
   m):
8 d 9
            else at->r = insert(at->r, idx, qt, m+1, r);
            return at->update(), at;
cff
cbb
cf7
        void insert(T v, SIZE T gt=1) { // insere 'qt'
   ocorrencias de 'v'
            if (qt <= 0) return erase(v, -qt);</pre>
882
            assert(v >= 0);
72b
f52
            expand(v);
            root = insert(root, v, qt, 0, N);
5e9
cbb
       }
        node* erase(node* at, T idx, SIZE_T qt, T 1, T r) {
f06
28 c
            if (!at) return at:
54b
            if (1 == r) at->cnt = at->cnt < qt ? 0 : at->cnt
 - qt;
            else {
4e6
                T m = 1 + (r-1)/2:
841
281
                if (idx \le m) at->1 = erase(at->1, idx, qt,
  1, m);
ba1
                else at->r = erase(at->r, idx, qt, m+1, r);
7 b 4
                at->update();
cbb
135
            if (!at->cnt) delete at, at = NULL;
ce6
            return at;
        }
cbb
```

```
void erase(T v, SIZE_T qt=1) { // remove 'qt'
43 d
   ocorrencias de 'v'
9 c 3
            if (v < 0 \text{ or } v > N \text{ or } !qt) \text{ return};
9dc
            if (qt < 0) insert(v, -qt);</pre>
            root = erase(root, v, qt, 0, N);
b1d
cbb
        void erase all(T v) { // remove todos os 'v'
8d6
            if (v < 0 \text{ or } v > N) return;
347
            root = erase(root, v,
9f2
   numeric_limits <SIZE_T >:: max(), 0, N);
cbb
        }
        SIZE_T count(node* at, T a, T b, T 1, T r) const {
0fe
            if (!at or b < l or r < a) return 0;</pre>
61b
            if (a <= 1 and r <= b) return at->cnt:
0fe
841
            T m = 1 + (r-1)/2:
84a
            return count(at->1, a, b, 1, m) + count(at->r,
   a, b, m+1, r);
cbb
        SIZE_T count(T v) const { return count(root, v, v,
0a9
   0, N); }
        SIZE_T order_of_key(T v) { return count(root, 0,
ffc
   v-1, 0, N); }
        SIZE_T lower_bound(T v) { return order_of_key(v); }
df2
        const T operator [](SIZE_T i) const { // i-esimo
e68
   menor elemento
809
            assert(i >= 0 and i < size());
            node* at = root;
c43
4a5
            T 1 = 0, r = N;
40 c
            while (1 < r) {
                 T m = 1 + (r-1)/2;
841
5c2
                if (count(at->1) > i) at = at->1, r = m;
4e6
                 else {
b4a
                     i -= count(at->1):
ded
                     at = at -> r; l = m+1;
                }
cbb
cbb
            }
792
             return 1;
cbb
        }
78 c
        node* merge(node* 1, node* r) {
```

```
347
            if (!l or !r) return 1 ? 1 : r:
504
            if (!l->l \text{ and } !l->r) \{ // \text{ folha} \}
599
                if (MULTI) 1->cnt += r->cnt;
55d
                 delete r;
792
                return 1;
cbb
            1->1 = merge(1->1, r->1), 1->r = merge(1->r,
f58
   r->r);
            1->update(), delete r;
f4f
792
            return 1;
cbb
f59
        void merge(sms& s) { // mergeia dois sets
068
             if (N > s.N) swap(*this, s);
785
             expand(s.N);
938
            root = merge(root, s.root);
            s.root = NULL;
ee2
cbb
        }
        node* split(node*& x, SIZE_T k) {
dc6
7ca
            if (k <= 0 or !x) return NULL;</pre>
6 d 0
             node* ret = new node();
386
            if (!x->1 \text{ and } !x->r) x->cnt -= k, ret->cnt += k;
4e6
             else {
85e
                if (k \le count(x->1)) ret->1 = split(x->1,
   k);
4e6
                 else {
06f
                     ret->r = split(x->r, k - count(x->1));
                     swap(x->1, ret->1);
cfd
cbb
674
                 ret->update(), x->update();
cbb
d5b
            if (!x->cnt) delete x, x = NULL;
edf
            return ret;
cbb
02b
        void split(SIZE_T k, sms& s) { // pega os 'k' menores
e63
             s.clear():
6e5
             s.root = split(root, min(k, size()));
еЗс
             s.N = N;
cbb
        // pega os menores que 'k'
d41
        void split_val(T k, sms& s) { split(order_of_key(k),
131
   s); }
```

#### 1.28 Split-Merge Set - Lazy

```
// Representa um conjunto de inteiros nao negativos
// Todas as operacoes custam O(log(N)),
// em que N = maior elemento do set,
// exceto o merge e o insert_range, que custa O(log(N))
   amortizado
// Usa O(\min(N, n \log(N))) de memoria, sendo 'n' o
// numero de elementos distintos no set
// 3828d0
fb1 template < typename T> struct sms {
         struct node {
3c9
             node *1, *r;
b19
0f9
            int cnt;
393
             bool flip;
0fa
             node(): 1(NULL), r(NULL), cnt(0), flip(0) {}
            void update() {
01e
a01
                 cnt = 0;
d8a
                 if (1) cnt += 1->cnt;
e49
                 if (r) cnt += r->cnt;
             }
cbb
214
        };
         void prop(node* x, int size) {
aee
             if (!x or !x->flip) return;
bb3
            x - > flip = 0;
f2c
            x \rightarrow cnt = size - x \rightarrow cnt;
fec
23f
            if (size > 1) {
641
                 if (!x->1) x->1 = new node();
756
                 if (!x->r) x->r = new node();
                 x - > 1 - > flip ^= 1;
ddd
Off
                 x \rightarrow r \rightarrow flip = 1;
             }
cbb
cbb
        }
bb7
        node* root;
fd0
        T N;
f34
         sms() : root(NULL), N(0) {}
```

```
83b
        sms(T v) : sms() { while (v >= N) N = 2*N+1; }
        sms(sms& t) : root(NULL), N(t.N) {
bdd
dc5
            for (int i = 0; i < t.size(); i++) insert(t[i]);</pre>
cbb
        sms(initializer_list<T> v) : sms() { for (T i : v)
a96
   insert(i); }
b2a
        void destroy(node* r) {
685
            vector < node *> q = {r};
402
            while (q.size()) {
e5d
                node* x = q.back(); q.pop_back();
ee9
                if (!x) continue;
1 c7
                q.push_back(x->1), q.push_back(x->r);
bf0
                delete x;
            }
cbb
cbb
b58
        \simsms() { destroy(root); }
        friend void swap(sms& a, sms& b) {
fdc
49e
            swap(a.root, b.root), swap(a.N, b.N);
cbb
83e
        sms& operator =(const sms& v) {
768
            sms tmp = v;
420
            swap(tmp, *this);
357
            return *this;
cbb
ff8
        int count(node* x, T size) {
a66
            if (!x) return 0;
793
            prop(x, size);
            return x->cnt;
ead
cbb
4fe
        int size() { return count(root, N+1); }
75a
        void clear() {
0a0
            sms tmp;
            swap(*this, tmp);
4ac
cbb
a06
        void expand(T v) {
            for (; N < v; N = 2*N+1) if (root) {
bc3
edf
                prop(root, N+1);
63 c
                node* nroot = new node();
956
                nroot ->1 = root;
897
                root = nroot;
a0a
                root->update();
```

```
cbb
cbb
        }
        node* insert(node* at, T idx, T 1, T r) {
fde
1a4
            if (!at) at = new node();
5ae
            else prop(at, r-l+1);
893
            if (1 == r) {
44b
                at->cnt = 1;
ce6
                return at;
            }
cbb
841
            T m = 1 + (r-1)/2:
95 a
            if (idx \le m) at -> l = insert(at -> l, idx, l, m);
            else at->r = insert(at->r, idx, m+1, r);
018
            return at->update(), at;
cff
cbb
        void insert(T v) {
c27
72b
            assert(v >= 0):
            expand(v);
f52
            root = insert(root, v, 0, N);
7f2
cbb
        }
393
        node* erase(node* at, T idx, T 1, T r) {
28 c
            if (!at) return at;
553
            prop(at, r-l+1);
            if (1 == r) at->cnt = 0;
4be
            else {
4e6
841
                T m = 1 + (r-1)/2;
d2d
                if (idx \le m) at->1 = erase(at->1, idx, 1,
   m);
                else at->r = erase(at->r, idx, m+1, r);
f3c
7 b 4
                at->update();
cbb
ce6
            return at;
cbb
26b
        void erase(T v) {
347
            if (v < 0 \text{ or } v > N) return:
980
            root = erase(root, v, 0, N);
        }
cbb
        int count(node* at, T a, T b, T l, T r) {
b4f
            if (!at or b < l or r < a) return 0;</pre>
61b
553
            prop(at, r-l+1);
```

```
0fe
            if (a <= 1 and r <= b) return at->cnt:
            T m = 1 + (r-1)/2;
841
84a
            return count(at->1, a, b, 1, m) + count(at->r,
   a, b, m+1, r);
cbb
b36
        int count(T v) { return count(root, v, v, 0, N); }
        int order_of_key(T v) { return count(root, 0, v-1,
eb0
   0, N); }
        int lower_bound(T v) { return order_of_key(v); }
fb8
         const T operator [](int i) { // i-esimo menor
   elemento
809
            assert(i >= 0 and i < size());</pre>
c43
             node* at = root:
             T 1 = 0, r = N:
4 a 5
40 c
             while (1 < r) {
553
                 prop(at, r-l+1);
841
                 T m = 1 + (r-1)/2;
                 if (count(at->1, m-1+1) > i) at = at->1, r =
4 e 7
   m;
4e6
                 else {
e6c
                     i \rightarrow count(at \rightarrow l, r-m);
ded
                     at = at -> r; l = m+1;
cbb
                 }
            }
cbb
792
             return 1:
cbb
        node* merge(node* a, node* b, T tam) {
63d
c48
             if (!a or !b) return a ? a : b;
10e
             prop(a, tam), prop(b, tam);
             if (b\rightarrow cnt == tam) swap(a, b);
abd
bb3
             if (tam == 1 or a->cnt == tam) {
a9e
                 destroy(b);
3f5
                 return a:
cbb
c 14
             a - > 1 = merge(a - > 1, b - > 1, tam > > 1), a - > r =
   merge(a->r, b->r, tam>>1);
496
            a->update(), delete b;
3f5
             return a;
cbb
f59
        void merge(sms& s) { // mergeia dois sets
```

```
068
            if (N > s.N) swap(*this, s);
785
            expand(s.N);
            root = merge(root, s.root, N+1);
707
            s.root = NULL;
ee2
cbb
        }
f76
        node* split(node*& x, int k, T tam) {
            if (k <= 0 or !x) return NULL;</pre>
7ca
            prop(x, tam);
e3b
6d0
            node* ret = new node();
37b
            if (tam == 1) x -> cnt = 0, ret -> cnt = 1;
4e6
                if (k \le count(x->1, tam>>1)) ret->1 =
b20
   split(x->1, k, tam>>1);
                 else {
4e6
5d8
                     ret -> r = split(x->r, k - count(x->l,
   tam >> 1), tam >> 1);
                     swap(x->1, ret->1);
cfd
cbb
674
                 ret->update(), x->update();
            }
cbb
edf
            return ret;
cbb
049
        void split(int k, sms& s) { // pega os 'k' menores
            s.clear();
e63
            s.root = split(root, min(k, size()), N+1);
eb6
e3c
            s.N = N;
cbb
d41
        // pega os menores que 'k'
131
        void split_val(T k, sms& s) { split(order_of_key(k),
   s); }
        void flip(node*& at, T a, T b, T 1, T r) {
ecf
            if (!at) at = new node();
1a4
5ae
            else prop(at, r-l+1);
9a3
            if (a \le 1 \text{ and } r \le b) {
747
                at->flip ^= 1;
553
                prop(at, r-1+1);
505
                return;
cbb
сс9
            if (r < a or b < 1) return;
            T m = 1 + (r-1)/2;
841
```

```
2a1
            flip(at->1, a, b, 1, m), flip(at->r, a, b, m+1,
   r);
7b4
            at->update();
cbb
        void flip(T 1, T r) \{ // \text{ flipa os valores em } [1, r] 
1ee
63e
             assert(1 >= 0 \text{ and } 1 <= r);
34b
             expand(r);
de7
            flip(root, 1, r, 0, N);
cbb
d41
        // complemento considerando que o universo eh [0,
 liml
042
        void complement(T lim) {
             assert(lim >= 0);
2 e 9
95c
            if (lim > N) expand(lim);
11a
            flip(root, 0, lim, 0, N);
0a0
             sms tmp;
180
             split_val(lim+1, tmp);
4ac
             swap(*this, tmp);
cbb
0eb
        void insert_range(T 1, T r) { // insere todo os
   valores em [1, r]
0a0
             sms tmp;
5fa
            tmp.flip(l, r);
7 f 7
            merge(tmp);
cbb
       }
214 };
```

#### 1.29 SQRT Tree

```
// RMQ em O(log log n) com O(n log log n) pra buildar
// Funciona com qualquer operacao associativa
// Tao rapido quanto a sparse table, mas usa menos memoria
// (log log (1e9) < 5, entao a query eh praticamente O(1))
//
// build - O(n log log n)
// query - O(log log n)
// 8ff986

97a namespace sqrtTree {
052    int n, *v;
ec7    int pref[4][MAX], sulf[4][MAX], getl[4][MAX],
    entre[4][MAX], sz[4];</pre>
```

```
5f7
        int op(int a, int b) { return min(a, b); }
c72
        inline int getblk(int p, int i) { return
   (i-getl[p][i])/sz[p]; }
        void build(int p, int l, int r) {
2c6
bc8
            if (1+1 >= r) return;
            for (int i = 1; i \le r; i++) getl[p][i] = 1;
368
            for (int L = 1; L <= r; L += sz[p]) {</pre>
f16
                int R = min(L+sz[p]-1, r);
191
                pref[p][L] = v[L], sulf[p][R] = v[R];
89 c
                for (int i = L+1; i <= R; i++) pref[p][i] =</pre>
59f
   op(pref[p][i-1], v[i]);
                for (int i = R-1; i >= L; i--) sulf[p][i] =
d9a
   op(v[i], sulf[p][i+1]);
221
                build(p+1, L, R);
cbb
695
            for (int i = 0; i <= sz[p]; i++) {
                int at = entre[p][1+i*sz[p]+i] =
ca5
   sulf[p][l+i*sz[p]];
                for (int j = i+1; j <= sz[p]; j++)
759
   entre[p][1+i*sz[p]+j] = at =
                        op(at, sulf[p][1+j*sz[p]]);
23a
cbb
cbb
        }
        void build(int n2, int* v2) {
0d8
            n = n2, v = v2;
680
            for (int p = 0; p < 4; p++) sz[p] = n2 =
44 c
   sqrt(n2);
            build(0, 0, n-1);
c50
cbb
9e3
        int query(int 1, int r) {
            if (1+1 >= r) return 1 == r ? v[1] : op(v[1],
792
   v[r]);
            int p = 0;
1ba
            while (getblk(p, 1) == getblk(p, r)) p++;
4ba
            int ans = sulf[p][1], a = getblk(p, 1)+1, b =
9e4
   getblk(p, r)-1;
8bf
            if (a \le b) ans = op(ans,
   entre[p][getl[p][1]+a*sz[p]+b]);
            return op(ans, pref[p][r]);
dea
cbb
cbb }
```

#### 1.30 Treap

```
// Todas as operacoes custam
// O(log(n)) com alta probabilidade, exceto meld
// meld custa O(log^2 n) amortizado com alta prob.,
// e permite unir duas treaps sem restricao adicional
// Na pratica, esse meld tem constante muito boa e
// o pior caso eh meio estranho de acontecer
// bd93e2
878 mt19937 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
aa1 template < typename T > struct treap {
        struct node {
3 c 9
b19
            node *1, *r;
284
            int p, sz;
36d
            T val, mi;
4 c7
            node(T \ v) : l(NULL), r(NULL), p(rng()), sz(1),
   val(v), mi(v) {}
01e
            void update() {
a26
                sz = 1;
d6e
                mi = val;
bd7
                if (1) sz += 1->sz, mi = min(mi, 1->mi);
a54
                if (r) sz += r->sz, mi = min(mi, r->mi);
cbb
            }
214
        };
        node* root;
bb7
84b
        treap() { root = NULL; }
2 d8
        treap(const treap& t) {
            throw logic_error("Nao copiar a treap!");
465
cbb
        \simtreap() {
cec
            vector < node *> q = {root};
609
402
            while (q.size()) {
                node* x = q.back(); q.pop_back();
e5d
ee9
                if (!x) continue;
                q.push_back(x->1), q.push_back(x->r);
1 c7
bf0
                delete x;
cbb
            }
```

```
}
cbb
73с
         int size(node* x) { return x ? x->sz : 0; }
b2b
         int size() { return size(root); }
         void join(node* 1, node* r, node*& i) { // assume
bcf
   que 1 < r
986
             if (!l or !r) return void(i = 1 ? 1 : r);
80e
             if (1->p > r->p) join(1->r, r, 1->r), i = 1;
fa0
             else join(l, r\rightarrow l, r\rightarrow l), i = r;
bda
             i->update();
cbb
         void split(node* i, node*& 1, node*& r, T v) {
есе
             if (!i) return void(r = l = NULL);
26a
             if (i\rightarrow val < v) split(i\rightarrow r, i\rightarrow r, r, v), l = i;
f05
807
             else split(i \rightarrow 1, l, i \rightarrow 1, v), r = i;
bda
             i->update();
         }
cbb
3fc
         void split_leg(node* i, node*& 1, node*& r, T v) {
             if (!i) return void(r = 1 = NULL);
26a
             if (i\rightarrow val \leftarrow v) split_leq(i\rightarrow r, i\rightarrow r, r, v), l
181
   = i;
58f
             else split_leq(i \rightarrow l, l, i \rightarrow l, v), r = i;
             i->update();
bda
cbb
e13
         int count(node* i, T v) {
6b4
             if (!i) return 0;
352
             if (i->val == v) return 1;
8d0
             if (v < i->val) return count(i->1, v);
             return count(i->r, v);
4d0
cbb
         void index_split(node* i, node*& 1, node*& r, int v,
26 d
   int key = 0) {
             if (!i) return void(r = 1 = NULL);
26a
             if (key + size(i->1) < v) index_split(i->r,
c10
   i - r, r, v, key+size(i - r)+1), l = i;
             else index_split(i \rightarrow 1, l, i \rightarrow l, v, key), r = i;
e5a
bda
             i->update();
         }
cbb
a1f
         int count(T v) {
e06
             return count(root, v);
cbb
c27
         void insert(T v) {
```

```
980
            if (count(v)) return;
031
            node *L, *R;
d42
            split(root, L, R, v);
585
            node* at = new node(v);
59f
            join(L, at, L);
a28
            join(L, R, root);
cbb
26b
        void erase(T v) {
df9
            node *L, *M, *R;
b6b
            split_leq(root, M, R, v), split(M, L, M, v);
f17
            if (M) delete M;
f38
            M = NULL;
a28
            join(L, R, root);
cbb
e77
        void meld(treap& t) { // segmented merge
4a6
            node *L = root, *R = t.root;
950
            root = NULL:
6 b 1
            while (L or R) {
                if (!L or (L and R and L->mi > R->mi))
fe2
   std::swap(L, R);
5e1
                if (!R) join(root, L, root), L = NULL;
3 c 9
                 else if (L->mi == R->mi) {
a76
                     node* LL;
439
                     split(L, LL, L, R->mi+1);
359
                     delete LL;
9 d 9
                } else {
a76
                     node* LL;
537
                     split(L, LL, L, R->mi);
dbb
                     join(root, LL, root);
                }
cbb
            }
cbb
689
            t.root = NULL;
cbb
214 };
```

#### 1.31 Treap Implicita

```
// Todas as operacoes custam
// O(log(n)) com alta probabilidade
// 63ba4d
878 mt19937 rng((int)
```

```
chrono::steady_clock::now().time_since_epoch().count());
aa1 template < typename T> struct treap {
3 c 9
        struct node {
b19
            node *1, *r;
284
            int p, sz;
875
            T val, sub, lazy;
aa6
            bool rev;
            node(T v) : l(NULL), r(NULL), p(rng()), sz(1),
8dc
   val(v), sub(v), lazy(0), rev(0) {}
            void prop() {
a9c
0ec
                if (lazy) {
924
                     val += lazy, sub += lazy*sz;
                     if (1) 1->lazy += lazy;
b87
d3b
                     if (r) r->lazy += lazy;
                }
cbb
1bb
                if (rev) {
e4f
                     swap(1, r);
                     if (1) 1->rev ^= 1;
dc8
                     if (r) r->rev ^= 1;
f2f
                }
cbb
a32
                 lazv = 0, rev = 0;
            }
cbb
            void update() {
01e
0 c 3
                 sz = 1, sub = val;
a09
                if (1) 1->prop(), sz += 1->sz, sub += 1->sub;
095
                if (r) r\rightarrow prop(), sz += r\rightarrow sz, sub += r\rightarrow sub;
cbb
214
        };
        node* root;
bb7
84b
        treap() { root = NULL; }
2d8
        treap(const treap& t) {
465
            throw logic_error("Nao copiar a treap!");
cbb
        }
        \simtreap() {
cec
609
            vector < node *> q = {root};
            while (q.size()) {
402
                node* x = q.back(); q.pop_back();
e5d
ee9
                if (!x) continue;
1 c 7
                 q.push_back(x->1), q.push_back(x->r);
                                                                    df9
```

```
bf0
                 delete x:
cbb
            }
        }
cbb
        int size(node* x) { return x ? x->sz : 0; }
73c
b2b
        int size() { return size(root); }
        void join(node* 1, node* r, node*& i) { // assume
bcf
   que 1 < r
986
            if (!l or !r) return void(i = 1 ? l : r);
            1->prop(), r->prop();
161
            if (1->p > r->p) join(1->r, r, 1->r), i = 1;
80e
fa0
            else join(l, r \rightarrow l, r \rightarrow l), i = r;
            i->update();
bda
cbb
a20
        void split(node* i, node*& 1, node*& r, int v, int
   kev = 0) {
26a
            if (!i) return void(r = 1 = NULL);
            i->prop();
c89
            if (key + size(i->1) < v) split(i->r, i->r, r,
5 bd
   v, kev + size(i - > 1) + 1), l = i;
            else split(i \rightarrow 1, l, i \rightarrow 1, v, key), r = i;
219
bda
            i->update();
cbb
231
        void push_back(T v) {
2e0
             node* i = new node(v);
7ab
             join(root, i, root);
cbb
b7a
        T query(int 1, int r) {
df9
            node *L, *M, *R;
dca
             split(root, M, R, r+1), split(M, L, M, 1);
d43
            T ans = M->sub;
69 d
             join(L, M, M), join(M, R, root);
ba7
            return ans;
cbb
41f
        void update(int 1, int r, T s) {
df9
             node *L, *M, *R;
dca
             split(root, M, R, r+1), split(M, L, M, 1);
8f6
            M - > lazy += s;
69 d
             join(L, M, M), join(M, R, root);
cbb
8 c 1
        void reverse(int 1, int r) {
             node *L, *M, *R;
```

# 1.32 Treap Persistent Implicita

```
// Todas as operacoes custam
// O(log(n)) com alta probabilidade
// fb8013
6cf mt19937_64 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
3c9 struct node {
b19
        node *1, *r;
f14
        ll sz, val, sub;
        node(ll v) : l(NULL), r(NULL), sz(1), val(v), sub(v)
304
   {}
        node(node* x) : l(x->l), r(x->r), sz(x->sz),
c12
   val(x->val), sub(x->sub) {}
        void update() {
01e
0 c 3
            sz = 1, sub = val;
77 e
            if (1) sz += 1->sz, sub += 1->sub;
d6e
            if (r) sz += r->sz, sub += r->sub;
124
            sub %= MOD;
        }
cbb
214 };
bc9 ll size(node* x) { return x ? x->sz : 0; }
761 void update(node* x) { if (x) x->update(); }
828 node* copy(node* x) { return x ? new node(x) : NULL; }
b02 node* join(node* 1, node* r) {
        if (!l or !r) return 1 ? copy(1) : copy(r);
e1f
48b
        node* ret;
        if (rng() % (size(l) + size(r)) < size(l)) {</pre>
49 f
7eb
            ret = copy(1);
            ret -> r = join(ret -> r, r);
cc1
9d9
        } else {
            ret = copy(r);
4 c 5
```

```
551
            ret->1 = join(1, ret->1);
cbb
74f
        return update(ret), ret;
cbb }
723 void split(node* x, node*& 1, node*& r, 11 v, 11 key =
    0) {
421
        if (!x) return void(l = r = NULL);
        if (key + size(x->1) < v) {
b4b
72f
            1 = copy(x);
             split(1->r, 1->r, r, v, key+size(1->1)+1);
d70
9d9
        } else {
            r = copy(x);
303
417
             split(r->1, l, r->l, v, key);
cbb
da2
         update(1), update(r);
cbb }
f9e vector < node *> treap;
139 void init(const vector<ll>& v) {
bbd
         treap = {NULL};
        for (auto i : v) treap[0] = join(treap[0], new
    node(i));
cbb }
1.33
      Wavelet Tree
// Usa O(sigma + n log(sigma)) de memoria,
// onde sigma = MAXN - MINN
// Depois do build, o v fica ordenado
// count(i, j, x, y) retorna o numero de elementos de
// v[i, j) que pertencem a [x, y]
// kth(i, j, k) retorna o elemento que estaria
// na poscicao k-1 de v[i, j), se ele fosse ordenado
// sum(i, j, x, y) retorna a soma dos elementos de
// v[i, j) que pertencem a [x, y]
// sumk(i, j, k) retorna a soma dos k-esimos menores
// elementos de v[i, j) (sum(i, j, 1) retorna o menor)
//
// Complexidades:
// build - O(n log(sigma))
```

```
// count - O(log(sigma))
// kth - 0(log(sigma))
// sum - 0(log(sigma))
// sumk - O(log(sigma))
// 782344
597 int n, v[MAX];
578 vector <int> esq[4*(MAXN-MINN)], pref[4*(MAXN-MINN)];
f8d void build(int b = 0, int e = n, int p = 1, int l =
   MINN, int r = MAXN) {
        int m = (1+r)/2; esq[p].push_back(0);
   pref[p].push_back(0);
        for (int i = b; i < e; i++) {</pre>
f2f
            esq[p].push_back(esq[p].back()+(v[i]<=m));</pre>
6b9
            pref[p].push_back(pref[p].back()+v[i]);
26f
cbb
        }
        if (1 == r) return;
8ce
        int m2 = stable_partition(v+b, v+e, [=](int
3a7
   i) {return i <= m;}) - v;
        build(b, m2, 2*p, 1, m), build(m2, e, 2*p+1, m+1, r);
347
cbb }
540 int count(int i, int j, int x, int y, int p = 1, int l =
   MINN, int r = MAXN) {
       if (y < 1 \text{ or } r < x) \text{ return } 0;
2ad
       if (x \le 1 \text{ and } r \le y) \text{ return } j-i;
4db
        int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
ddc
        return count(ei, ej, x, y, 2*p, 1, m)+count(i-ei,
0a5
   j-ej, x, y, 2*p+1, m+1, r);
cbb }
f62 int kth(int i, int j, int k, int p=1, int l = MINN, int
   r = MAXN) {
        if (1 == r) return 1:
Зсе
        int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
ddc
        if (k <= ej-ei) return kth(ei, ej, k, 2*p, 1, m);</pre>
585
28b
        return kth(i-ei, j-ej, k-(ej-ei), 2*p+1, m+1, r);
cbb }
f2c int sum(int i, int j, int x, int y, int p = 1, int l =
   MINN, int r = MAXN) {
```

```
if (y < 1 or r < x) return 0;
2ad
       if (x <= 1 and r <= y) return pref[p][j]-pref[p][i];</pre>
2a9
       int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
ddc
        return sum(ei, ej, x, y, 2*p, 1, m) + sum(i-ei,
43b
   j-ej, x, y, 2*p+1, m+1, r);
cbb }
b84 int sumk(int i, int j, int k, int p = 1, int l = MINN,
   int r = MAXN) {
       if (1 == r) return l*k;
8a1
      int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
50 c
       if (k <= ej-ei) return sumk(ei, ej, k, 2*p, 1, m);</pre>
        return pref[2*p][ej]-pref[2*p][ei]+sumk(i-ei, j-ej,
   k-(ej-ei), 2*p+1, m+1, r);
cbb }
```

## 2 Grafos

### 2.1 AGM Direcionada

```
// Fala o menor custo para selecionar arestas tal que
// o vertice 'r' alcance todos
// Se nao tem como, retorna LINF
//
// O(m log(n))
// dc345b
3c9 struct node {
f31
        pair<11, int> val;
4e4
        ll lazv:
       node *1, *r;
b19
        node() {}
f93
        node(pair<int, int> v) : val(v), lazy(0), l(NULL),
c53
   r(NULL) {}
a9c
        void prop() {
768
            val.first += lazy;
b87
            if (1) 1->lazy += lazy;
d3b
            if (r) r->lazy += lazy;
c60
            lazv = 0;
```

```
}
cbb
214 };
de5 void merge(node*& a, node* b) {
        if (!a) swap(a, b);
c11
        if (!b) return;
802
626
        a->prop(), b->prop();
        if (a->val > b->val) swap(a, b);
d04
        merge(rand()%2 ? a->1 : a->r, b);
4b0
cbb }
d01 pair<ll, int> pop(node*& R) {
e8f
        R->prop();
22 e
        auto ret = R->val;
af0
        node* tmp = R;
3f3
        merge(R->1, R->r);
6c9
        R = R -> 1:
3 e 4
        if (R) R->lazy -= ret.first;
7c3
        delete tmp;
edf
        return ret;
cbb }
6f6 void apaga(node* R) { if (R) apaga(R->1), apaga(R->r),
   delete R; }
f13 ll dmst(int n, int r, vector <pair <pair <int, int >, int >> &
   ar) {
        vector < int > p(n); iota(p.begin(), p.end(), 0);
94e
        function <int(int) > find = [&](int k) { return
a23
   p[k] == k?k:p[k] = find(p[k]); };
2d7
        vector < node *> h(n);
        for (auto e : ar) merge(h[e.first.second], new
56f
   node({e.second, e.first.first}));
        vector < int > pai(n, -1), path(n);
fd1
66 e
        pai[r] = r;
        11 \text{ ans} = 0;
04b
603
        for (int i = 0; i < n; i++) { // vai conectando todo
   mundo
2a3
            int u = i, at = 0;
             while (pai[u] == -1) {
cae
                if (!h[u]) { // nao tem
daa
947
                     for (auto i : h) apaga(i);
77 c
                     return LINF;
cbb
                }
```

```
167
                path[at++] = u, pai[u] = i;
                auto [mi, v] = pop(h[u]);
55e
64c
                ans += mi;
                if (pai[u = find(v)] == i) { // ciclo
5e2
86f
                    while (find(v = path[--at]) != u)
621
                         merge(h[u], h[v]), h[v] = NULL,
   p[find(v)] = u;
57a
                    pai[u] = -1;
cbb
cbb
            }
cbb
        for (auto i : h) apaga(i);
947
ba7
        return ans;
cbb }
```

## 2.2 Bellman-Ford

```
// Calcula a menor distancia
// entre a e todos os vertices e
// detecta ciclo negativo
// Retorna 1 se ha ciclo negativo
// Nao precisa representar o grafo,
// soh armazenar as arestas
//
// O(nm)
// 03059ъ
14e int n, m;
248 int d[MAX];
e93 vector <pair <int, int>> ar; // vetor de arestas
9e2 vector < int > w:
                                // peso das arestas
6be bool bellman_ford(int a) {
        for (int i = 0; i < n; i++) d[i] = INF;</pre>
8ec
        d[a] = 0;
8a8
4e3
        for (int i = 0; i <= n; i++)
            for (int j = 0; j < m; j++) {
891
6 e 4
                if (d[ar[j].second] > d[ar[j].first] + w[j])
   {
705
                    if (i == n) return 1;
```

### 2.3 Block-Cut Tree

```
// Cria a block-cut tree, uma arvore com os blocos
// e os pontos de articulação
// Blocos sao componentes 2-vertice-conexos maximais
// Uma 2-coloracao da arvore eh tal que uma cor sao
// os blocos, e a outra cor sao os pontos de art.
// Funciona para grafo nao conexo
// art[i] responde o numero de novas componentes conexas
// criadas apos a remocao de i do grafo g
// Se art[i] >= 1, i eh ponto de articulação
//
// Para todo i <= blocks.size()</pre>
// blocks[i] eh uma componente 2-vertce-conexa maximal
// edgblocks[i] sao as arestas do bloco i
// tree[i] eh um vertice da arvore que corresponde ao bloco i
//
// pos[i] responde a qual vertice da arvore vertice i
   pertence
// Arvore tem no maximo 2n vertices
// O(n+m)
// 056fa2
d10 struct block_cut_tree {
        vector < vector < int >> g, blocks, tree;
d8e
43b
        vector < vector < pair < int , int >>> edgblocks;
4ce
        stack < int > s;
6c0
        stack<pair<int, int>> s2;
2bb
        vector<int> id, art, pos;
d41
763
        block_cut_tree(vector < vector < int >> g_) : g(g_) {
af1
            int n = g.size();
```

```
37a
            id.resize(n, -1), art.resize(n), pos.resize(n);
6f2
            build();
cbb
        int dfs(int i, int& t, int p = -1) {
df6
cf0
            int lo = id[i] = t++;
18e
            s.push(i);
d41
827
            if (p != -1) s2.emplace(i, p);
            for (int j : g[i]) if (j != p and id[j] != -1)
   s2.emplace(i, j);
d 4 1
            for (int j : g[i]) if (j != p) {
сас
                if (id[j] == -1) {
9 a 3
121
                    int val = dfs(j, t, i);
0 c3
                     lo = min(lo, val);
                    if (val >= id[i]) {
588
66a
                         art[i]++;
483
                         blocks.emplace_back(1, i);
                         while (blocks.back().back() != j)
110
138
   blocks.back().push_back(s.top()), s.pop();
128
                         edgblocks.emplace_back(1, s2.top()),
   s2.pop();
47e
                         while (edgblocks.back().back() !=
   pair(j, i))
bce
   edgblocks.back().push_back(s2.top()), s2.pop();
cbb
d 4 1
                     // if (val > id[i]) aresta i-j eh ponte
cbb
328
                else lo = min(lo, id[j]);
cbb
            }
d41
3bd
            if (p == -1 and art[i]) art[i]--;
253
            return lo;
cbb
0a8
        void build() {
6 bb
            int t = 0;
```

```
for (int i = 0; i < g.size(); i++) if (id[i] ==
abf
   -1) dfs(i, t, -1);
d41
56 c
            tree.resize(blocks.size());
            for (int i = 0; i < g.size(); i++) if (art[i])</pre>
f7d
965
                pos[i] = tree.size(), tree.emplace back();
973
           for (int i = 0; i < blocks.size(); i++) for (int</pre>
   j : blocks[i]) {
                if (!art[j]) pos[j] = i;
403
                else tree[i].push_back(pos[j]),
101
   tree[pos[j]].push_back(i);
cbb
        }
cbb
214 };
```

## 2.4 Blossom - matching maximo em grafo geral

```
// O(n^3)
// Se for bipartido, nao precisa da funcao
// 'contract', e roda em O(nm)
// 4426a4
042 vector <int> g[MAX];
128 int match[MAX]; // match[i] = com quem i esta matchzado
1f1 int n, pai[MAX], base[MAX], vis[MAX];
26a queue < int > q;
107 void contract(int u, int v, bool first = 1) {
165
        static vector <bool> bloss:
fbe
        static int 1:
418
        if (first) {
            bloss = vector < bool > (n, 0);
a47
            vector < bool > teve(n, 0);
042
            int k = u; l = v;
ddf
31e
            while (1) {
                teve[k = base[k]] = 1;
297
                if (match[k] == -1) break;
116
                k = pai[match[k]];
dfa
            }
cbb
            while (!teve[l = base[l]]) l = pai[match[l]];
d31
```

```
cbb
        while (base[u] != 1) {
2e9
            bloss[base[u]] = bloss[base[match[u]]] = 1;
e29
8fa
            pai[u] = v;
0 b 0
            v = match[u];
a51
            u = pai[match[u]];
cbb
71c
        if (!first) return;
        contract(v, u, 0);
95e
        for (int i = 0; i < n; i++) if (bloss[base[i]]) {</pre>
6ee
594
            base[i] = 1;
ca7
            if (!vis[i]) q.push(i);
29a
            vis[i] = 1:
        }
cbb
cbb }
f10 int getpath(int s) {
        for (int i = 0; i < n; i++) base[i] = i, pai[i] =</pre>
88f
   -1, vis[i] = 0;
        vis[s] = 1; q = queue < int > (); q.push(s);
ded
402
        while (q.size()) {
            int u = q.front(); q.pop();
be1
           for (int i : g[u]) {
bdc
7a2
                if (base[i] == base[u] or match[u] == i)
   continue;
                if (i == s or (match[i] != -1 and
e35
   pai[match[i]] != -1))
4f2
                     contract(u, i);
e2e
                else if (pai[i] == -1) {
545
                     pai[i] = u;
                    if (match[i] == -1) return i;
f6a
                    i = match[i];
818
29 d
                    vis[i] = 1; q.push(i);
                }
cbb
cbb
            }
cbb
daa
        return -1;
cbb }
83f int blossom() {
1a4
        int ans = 0;
315
        memset(match, -1, sizeof(match));
```

```
2e3
        for (int i = 0: i < n: i++) if (match[i] == -1)
            for (int j : g[i]) if (match[j] == -1) {
f76
1bc
                match[i] = j;
                match[j] = i;
f1d
0df
                ans++;
c2b
                break:
cbb
        for (int i = 0; i < n; i++) if (match[i] == -1) {</pre>
da8
7e3
            int j = getpath(i);
5f2
            if (j == -1) continue;
0df
           ans++;
3a0
            while (j != -1) {
                int p = pai[j], pp = match[p];
ef0
               match[p] = j;
348
               match[i] = p;
fe9
55 d
                j = pp;
            }
cbb
        }
cbb
ba7
        return ans;
cbb }
```

## 2.5 Centro de arvore

```
// Retorna o diametro e o(s) centro(s) da arvore
// Uma arvore tem sempre um ou dois centros e estes estao no
   meio do diametro
//
// O(n)
// cladeb
042 vector <int> g[MAX];
df1 int d[MAX], par[MAX];
544 pair <int, vector <int>> center() {
a95
        int f, df;
        function < void(int) > dfs = [&] (int v) {
36 d
d47
           if (d[v] > df) f = v, df = d[v];
           for (int u : g[v]) if (u != par[v])
e68
                d[u] = d[v] + 1, par[u] = v, dfs(u);
1a5
214
        };
d41
1 b 0
        f = df = par[0] = -1, d[0] = 0;
```

```
41e
        dfs(0):
c2d
        int root = f;
        f = df = par[root] = -1, d[root] = 0;
0f6
14e
        dfs(root);
d41
761
        vector < int > c;
87e
        while (f != -1) {
999
            if (d[f] == df/2 \text{ or } d[f] == (df+1)/2)
   c.push_back(f);
19 c
           f = par[f];
cbb
d41
00f
        return {df, c};
cbb }
2.6 Centroid
```

```
// Computa os 2 centroids da arvore
//
// O(n)
// e16075
97a int n, subsize[MAX];
042 vector < int > g[MAX];
98f void dfs(int k, int p=-1) {
bd2
        subsize[k] = 1:
        for (int i : g[k]) if (i != p) {
6 e 5
801
            dfs(i, k);
2 e 3
             subsize[k] += subsize[i];
cbb
cbb }
2e8 int centroid(int k, int p=-1, int size=-1) {
        if (size == -1) size = subsize[k];
e73
        for (int i : g[k]) if (i != p) if (subsize[i] >
8df
   size/2)
bab
            return centroid(i, k, size);
839
        return k;
cbb }
f20 pair <int, int > centroids(int k=0) {
```

```
051     dfs(k);
909     int i = centroid(k), i2 = i;
8dd     for (int j : g[i]) if (2*subsize[j] == subsize[k])
     i2 = j;
0cb     return {i, i2};
cbb }
```

# 2.7 Centroid decomposition

```
// decomp(0, k) computa numero de caminhos com 'k' arestas
// Mudar depois do comentario
//
// O(n log(n))
// fe2541
042 vector <int> g[MAX];
ba8 int sz[MAX], rem[MAX];
747 void dfs(vector<int>& path, int i, int l=-1, int d=0) {
        path.push_back(d);
547
        for (int j : g[i]) if (j != l and !rem[j]) dfs(path,
   j, i, d+1);
cbb }
071 int dfs_sz(int i, int l=-1) {
02c
        sz[i] = 1:
        for (int j : g[i]) if (j != l and !rem[j]) sz[i] +=
e5c
   dfs_sz(j, i);
       return sz[i];
191
cbb }
85a int centroid(int i, int 1, int size) {
        for (int j : g[i]) if (j != l and !rem[j] and sz[j]
994
   > size / 2)
735
            return centroid(j, i, size);
d9a
        return i;
cbb }
d79 ll decomp(int i, int k) {
        int c = centroid(i, i, dfs_sz(i));
106
        rem[c] = 1;
a67
```

```
d41
                                                           // gasta O(n) aqui - dfs sem ir pros caras removidos
04b
                                                           11 \text{ ans} = 0;
                                                           vector < int > cnt(sz[i]);
020
878
                                                            cnt[0] = 1;
                                                          for (int j : g[c]) if (!rem[j]) {
0a8
5b4
                                                                                         vector < int > path;
baf
                                                                                          dfs(path, j);
1 a 1
                                                                                         for (int d : path) if (0 \le k-d-1 \text{ and } k-d-1 \le k-d-1 \text{ and } k
                        sz[i])
 285
                                                                                                                       ans += cnt[k-d-1];
                                                                                        for (int d : path) cnt[d+1]++;
 e8b
 cbb
1 c 1
                                                           for (int j : g[c]) if (!rem[j]) ans += decomp(j, k);
3 f 1
                                                            rem[c] = 0:
ba7
                                                            return ans;
cbb }
```

### 2.8 Centroid Tree

```
// Constroi a centroid tree
// p[i] eh o pai de i na centroid-tree
// dist[i][k] = distancia na arvore original entre i
// e o k-esimo ancestral na arvore da centroid
// O(n log(n)) de tempo e memoria
// a0e7c7
845 vector <int> g[MAX], dist[MAX];
cle int sz[MAX], rem[MAX], p[MAX];
071 int dfs_sz(int i, int l=-1) {
02c
        sz[i] = 1:
        for (int j : g[i]) if (j != l and !rem[j]) sz[i] +=
e5c
   dfs_sz(j, i);
        return sz[i];
191
cbb }
85a int centroid(int i, int 1, int size) {
        for (int j : g[i]) if (j != l and !rem[j] and sz[j]
   > size / 2)
735
           return centroid(j, i, size);
```

```
d9a
        return i:
cbb }
324 void dfs_dist(int i, int 1, int d=0) {
        dist[i].push_back(d);
541
5a1
        for (int j : g[i]) if (j != l and !rem[j])
            dfs_dist(j, i, d+1);
82a
cbb }
27e void decomp(int i, int l = -1) {
        int c = centroid(i, i, dfs_sz(i));
106
1 b 9
        rem[c] = 1, p[c] = 1;
        dfs_dist(c, c);
534
        for (int j : g[c]) if (!rem[j]) decomp(j, c);
a2a
cbb }
76c void build(int n) {
        for (int i = 0; i < n; i++) rem[i] = 0.</pre>
   dist[i].clear();
        decomp(0);
867
96b
        for (int i = 0; i < n; i++) reverse(dist[i].begin(),</pre>
   dist[i].end());
cbb }
```

# 2.9 Dijkstra

```
// encontra menor distancia de x
// para todos os vertices
// se ao final do algoritmo d[i] = LINF,
// entao x nao alcanca i
//
// O(m log(n))
// 695ac4
eff ll d[MAX];
c0d vector<pair<int, int>> g[MAX]; // {vizinho, peso}
1a8 int n;
d41
abc void dijkstra(int v) {
        for (int i = 0; i < n; i++) d[i] = LINF;</pre>
22 c
        d[v] = 0;
a7f
```

```
88 c
        priority_queue <pair <11, int>> pq;
b32
        pq.emplace(0, v);
d41
265
        while (pq.size()) {
a25
            auto [ndist, u] = pq.top(); pq.pop();
953
            if (-ndist > d[u]) continue;
d41
            for (auto [idx, w] : g[u]) if (d[idx] > d[u] +
cda
   w) {
331
                d[idx] = d[u] + w;
a84
                pq.emplace(-d[idx], idx);
cbb
        }
cbb
cbb }
```

### 2.10 Dinitz

```
// O(min(m * max_flow, n^2 m))
// Grafo com capacidades 1: O(\min(m \text{ sqrt}(m), m * n^2(2/3)))
// Todo vertice tem grau de entrada ou saida 1: O(m sqrt(n))
// 67ce89
472 struct dinitz {
        const bool scaling = false; // com scaling -> 0(nm
   log(MAXCAP)),
206
        int lim;
                                      // com constante alta
670
        struct edge {
358
            int to, cap, rev, flow;
7f9
            bool res;
d36
            edge(int to_, int cap_, int rev_, bool res_)
a94
                : to(to_), cap(cap_), rev(rev_), flow(0),
   res(res_) {}
214
        }:
002
        vector < vector < edge >> g;
        vector < int > lev, beg;
216
a71
        11 F;
190
        dinitz(int n) : g(n), F(0) {}
        void add(int a, int b, int c) {
087
            g[a].emplace_back(b, c, g[b].size(), false);
bae
            g[b].emplace_back(a, 0, g[a].size()-1, true);
4 c 6
```

```
cbb
123
        bool bfs(int s, int t) {
90f
            lev = vector \langle int \rangle (g.size(), -1); lev[s] = 0;
64 c
            beg = vector < int > (g.size(), 0);
8b2
            queue < int > q; q.push(s);
402
            while (q.size()) {
be1
                 int u = q.front(); q.pop();
bd9
                 for (auto& i : g[u]) {
                     if (lev[i.to] != -1 or (i.flow ==
dbc
   i.cap)) continue;
                     if (scaling and i.cap - i.flow < lim)</pre>
   continue;
                     lev[i.to] = lev[u] + 1;
185
8ca
                     q.push(i.to);
                 }
cbb
cbb
            return lev[t] != -1;
0de
cbb
        int dfs(int v, int s, int f = INF) {
dfb
            if (!f or v == s) return f;
50b
            for (int& i = beg[v]; i < g[v].size(); i++) {</pre>
88f
027
                 auto& e = g[v][i];
206
                if (lev[e.to] != lev[v] + 1) continue;
                 int foi = dfs(e.to, s, min(f, e.cap -
ee0
   e.flow));
749
                if (!foi) continue;
                 e.flow += foi, g[e.to][e.rev].flow -= foi;
3c5
45 c
                 return foi;
            }
cbb
            return 0;
bb3
cbb
ff6
        11 max_flow(int s, int t) {
            for (lim = scaling ? (1 << 30) : 1; lim; lim /= 2)
a86
                 while (bfs(s, t)) while (int ff = dfs(s, t))
9d1
   F += ff:
            return F:
4ff
        }
cbb
214 };
    // Recupera as arestas do corte s-t
    // d23977
dbd vector<pair<int, int>> get_cut(dinitz& g, int s, int t) {
```

```
f07
        g.max_flow(s, t);
68 c
        vector < pair < int , int >> cut;
1 b0
        vector < int > vis(g.g.size(), 0), st = \{s\};
321
        vis[s] = 1;
        while (st.size()) {
3 c 6
            int u = st.back(); st.pop_back();
b17
322
            for (auto e : g.g[u]) if (!vis[e.to] and e.flow
   < e.cap)
c17
                 vis[e.to] = 1, st.push_back(e.to);
cbb
481
        for (int i = 0; i < g.g.size(); i++) for (auto e :</pre>
   g.g[i])
9d2
             if (vis[i] and !vis[e.to] and !e.res)
   cut.emplace_back(i, e.to);
        return cut:
d1b
cbb }
```

## 2.11 Dominator Tree - Kawakami

```
// Se vira pra usar ai
//
// build - O(n)
// dominates - O(1)
// c80920
1a8 int n;
bbf namespace d_tree {
042
        vector < int > g[MAX];
d41
        // The dominator tree
b39
        vector < int > tree[MAX]:
        int dfs_l[MAX], dfs_r[MAX];
5af
        // Auxiliary data
d41
        vector < int > rg[MAX], bucket[MAX];
a2e
3ef
        int idom[MAX], sdom[MAX], prv[MAX], pre[MAX];
44b
        int ancestor[MAX], label[MAX];
563
        vector < int > preorder;
76a
        void dfs(int v) {
6a1
            static int t = 0;
```

```
db6
            pre[v] = ++t;
            sdom[v] = label[v] = v;
767
a3d
            preorder.push_back(v);
            for (int nxt: g[v]) {
80b
                if (sdom[nxt] == -1) {
56 c
                     prv[nxt] = v;
eed
900
                     dfs(nxt);
cbb
                rg[nxt].push_back(v);
2b5
cbb
cbb
62 e
        int eval(int v) {
            if (ancestor[v] == -1) return v;
c93
            if (ancestor[v]] == -1) return label[v];
a75
            int u = eval(ancestor[v]):
f33
            if (pre[sdom[u]] < pre[sdom[label[v]]]) label[v]</pre>
b49
   = u;
            ancestor[v] = ancestor[u];
66e
            return label[v];
c24
cbb
4b2
        void dfs2(int v) {
6a1
            static int t = 0;
330
            dfs l[v] = t++;
5e0
            for (int nxt: tree[v]) dfs2(nxt);
8e2
            dfs_r[v] = t++;
cbb
        }
c2c
        void build(int s) {
            for (int i = 0; i < n; i++) {</pre>
603
e6f
                 sdom[i] = pre[i] = ancestor[i] = -1;
2e1
                rg[i].clear();
                tree[i].clear();
50a
666
                bucket[i].clear();
            }
cbb
772
            preorder.clear():
c6c
            dfs(s):
12b
            if (preorder.size() == 1) return;
            for (int i = int(preorder.size()) - 1; i >= 1;
3 c 7
   i--) {
                int w = preorder[i];
6c6
                for (int v: rg[w]) {
a52
                    int u = eval(v);
5c1
                    if (pre[sdom[u]] < pre[sdom[w]]) sdom[w]</pre>
a 17
```

```
= sdom[u]:
cbb
680
                bucket[sdom[w]].push_back(w);
                ancestor[w] = prv[w];
ea7
b99
                for (int v: bucket[prv[w]]) {
5 c 1
                     int u = eval(v);
977
                     idom[v] = (u == v) ? sdom[v] : u;
cbb
2cc
                bucket[prv[w]].clear();
cbb
d0c
            for (int i = 1; i < preorder.size(); i++) {</pre>
6c6
                int w = preorder[i];
                 if (idom[w] != sdom[w]) idom[w] =
14b
   idom[idom[w]]:
32f
                 tree[idom[w]].push_back(w);
cbb
8ac
            idom[s] = sdom[s] = -1;
1 b6
            dfs2(s);
        }
cbb
d41
        // Whether every path from s to v passes through u
490
        bool dominates(int u, int v) {
            if (pre[v] == -1) return 1; // vacuously true
c75
            return dfs_l[u] <= dfs_l[v] && dfs_r[v] <=</pre>
2ea
   dfs_r[u];
cbb
       }
214 };
```

## 2.12 Euler Path / Euler Cycle

```
// Para declarar: 'euler < true > E(n); ' se quiser
// direcionado e com 'n' vertices
// As funcoes retornam um par com um booleano
// indicando se possui o cycle/path que voce pediu,
// e um vector de {vertice, id da aresta para chegar no vertice}
// Se for get_path, na primeira posicao o id vai ser -1
// get_path(src) tenta achar um caminho ou ciclo euleriano
// comecando no vertice 'src'.
// Se achar um ciclo, o primeiro e ultimo vertice serao 'src'.
// Se for um P3, um possiveo retorno seria [0, 1, 2, 0]
```

```
// get_cycle() acha um ciclo euleriano se o grafo for
   euleriano.
// Se for um P3, um possivel retorno seria [0, 1, 2]
// (vertie inicial nao repete)
//
// O(n+m)
// 7113df
63f template <bool directed=false > struct euler {
1a8
        int n;
4c0
        vector < vector < pair < int , int >>> g;
        vector<int> used:
d63
        euler(int n_) : n(n_), g(n) {}
30f
50 f
        void add(int a, int b) {
            int at = used.size();
4cd
c51
            used.push_back(0);
            g[a].emplace_back(b, at);
74e
            if (!directed) g[b].emplace_back(a, at);
fab
cbb
d41 #warning chamar para o src certo!
        pair < bool, vector < pair < int, int >>> get_path(int src)
eed
   {
            if (!used.size()) return {true, {}};
baf
            vector<int> beg(n, 0);
b25
            for (int& i : used) i = 0;
4ec
            // {{vertice, anterior}, label}
d41
            vector<pair<int, int>, int>> ret, st =
363
   {{src, -1}, -1}};
            while (st.size()) {
3 c 6
                int at = st.back().first.first;
8ff
002
                int& it = beg[at];
                while (it < g[at].size() and
8a1
   used[g[at][it].second]) it++;
8e4
                if (it == g[at].size()) {
                    if (ret.size() and
9dd
   ret.back().first.second != at)
                         return {false, {}};
b82
420
                     ret.push_back(st.back()), st.pop_back();
9d9
                } else {
                     st.push_back({{g[at][it].first, at},
daa
   g[at][it].second});
```

```
used[g[at][it].second] = 1;
eb8
cbb
cbb
            }
             if (ret.size() != used.size()+1) return {false,
a 19
   {}};
f77
             vector < pair < int , int >> ans;
fdf
            for (auto i : ret)
   ans.emplace_back(i.first.first, i.second);
459
             reverse(ans.begin(), ans.end());
997
             return {true, ans};
cbb
9 b 6
        pair < bool, vector < pair < int, int >>> get_cycle() {
             if (!used.size()) return {true, {}};
baf
            int src = 0;
ad1
34b
            while (!g[src].size()) src++;
            auto ans = get_path(src);
687
33 c
            if (!ans.first or ans.second[0].first !=
   ans.second.back().first)
b82
                 return {false, {}};
350
             ans.second[0].second = ans.second.back().second;
8 b 8
             ans.second.pop_back();
ba7
             return ans;
cbb
214 };
```

## 2.13 Euler Tour Tree

```
// Mantem uma floresta enraizada dinamicamente
// e permite queries/updates em sub-arvore
//
// Chamar ETT E(n, v), passando n = numero de vertices
// e v = vector com os valores de cada vertice (se for vazio,
// constroi tudo com 0
//
// link(v, u) cria uma aresta de v pra u, de forma que u se torna
// o pai de v (eh preciso que v seja raiz anteriormente)
// cut(v) corta a resta de v para o pai
// query(v) retorna a soma dos valores da sub-arvore de v
// update(v, val) soma val em todos os vertices da sub-arvore de v
// update_v(v, val) muda o valor do vertice v para val
```

```
// is_in_subtree(v, u) responde se o vertice u esta na
   sub-arvore de v
// Tudo O(log(n)) com alta probabilidade
// c97d63
878 mt19937 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
9f9 template < typename T> struct ETT {
d41
        // treap
3 c 9
        struct node {
ed1
            node *1, *r, *p;
fa4
            int pr, sz;
875
            T val, sub, lazy;
53e
            int id;
ffd
            bool f; // se eh o 'first'
            int qt_f; // numero de firsts na subarvore
5ef
            node(int id_, T v, bool f_ = 0) : 1(NULL),
7 a 8
   r(NULL), p(NULL), pr(rng()),
62b
                 sz(1), val(v), sub(v), lazy(), id(id_),
   f(f_), qt_f(f_) {}
a9c
            void prop() {
d09
                if (lazy != T()) {
021
                     if (f) val += lazy;
971
                     sub += lazv*sz;
b87
                     if (1) 1->lazy += lazy;
d3b
                     if (r) r \rightarrow lazy += lazy;
cbb
                }
bfd
                 lazy = T();
            }
cbb
01e
            void update() {
                 sz = 1, sub = val, qt_f = f;
8da
                if (1) 1 - > prop(), sz += 1 - > sz, sub +=
171
   1->sub, qt_f += 1->qt_f;
                if (r) r\rightarrow prop(), sz += r\rightarrow sz, sub +=
117
   r->sub, qt_f += r->qt_f;
cbb
214
        };
bb7
        node* root;
```

```
73c
        int size(node* x) { return x ? x->sz : 0; }
bcf
        void join(node* 1, node* r, node*& i) { // assume
  que 1 < r
986
             if (!1 or !r) return void(i = 1 ? 1 : r);
            1->prop(), r->prop();
161
ff5
            if (1->pr > r->pr) join(1->r, r, 1->r), 1->r->p
= i = 1;
982
             else join(1, r \rightarrow 1, r \rightarrow 1), r \rightarrow 1 \rightarrow p = i = r;
bda
             i->update();
cbb
a20
        void split(node* i, node*& 1, node*& r, int v, int
26a
             if (!i) return void(r = 1 = NULL);
c89
            i->prop();
d9e
            if (key + size(i->1) < v) {
448
                 split(i->r, i->r, r, v, key+size(i->l)+1), l
   = i;
a21
                 if (r) r - > p = NULL;
6 e 8
                 if (i->r) i->r->p = i;
9 d 9
             } else {
98d
                 split(i\rightarrow 1, 1, i\rightarrow 1, v, key), r = i;
5 a 3
                 if (1) 1->p = NULL;
899
                 if (i->1) i->1->p = i;
cbb
bda
            i->update();
cbb
ac7
        int get_idx(node* i) {
6cf
             int ret = size(i->1);
482
             for (; i->p; i = i->p) {
fbf
                 node* pai = i->p;
8a6
                 if (i != pai->l) ret += size(pai->l) + 1;
cbb
             }
edf
             return ret;
cbb
        node* get_min(node* i) {
048
             if (!i) return NULL;
433
f8e
             return i->1 ? get_min(i->1) : i;
cbb
        node* get_max(node* i) {
f03
433
             if (!i) return NULL;
424
             return i->r ? get_max(i->r) : i;
cbb
```

```
d41
        // fim da treap
4fb
        vector<node*> first, last;
        ETT(int n, vector<T> v = {}) : root(NULL), first(n),
f82
   last(n) {
            if (!v.size()) v = vector < T > (n);
с5е
            for (int i = 0; i < n; i++) {</pre>
603
                first[i] = last[i] = new node(i, v[i], 1);
a00
469
                join(root, first[i], root);
            }
cbb
        }
cbb
83f
        ETT(const ETT& t) { throw logic_error("Nao copiar a
   ETT!"); }
c09
        \simETT() {
609
            vector<node*> q = {root};
402
            while (q.size()) {
                node* x = q.back(); q.pop_back();
e5d
                if (!x) continue;
ee9
                q.push_back(x->1), q.push_back(x->r);
1c7
bf0
                delete x;
cbb
            }
cbb
        }
153
        pair<int, int> get_range(int i) {
            return {get_idx(first[i]), get_idx(last[i])};
670
cbb
7af
        void link(int v, int u) { // 'v' tem que ser raiz
890
            auto [lv, rv] = get_range(v);
f13
            int ru = get_idx(last[u]);
d41
4b4
            node* V;
            node *L, *M, *R;
df9
            split(root, M, R, rv+1), split(M, L, M, lv);
117
f1e
            V = M:
a28
            join(L, R, root);
e66
            split(root, L, R, ru+1);
            join(L, V, L);
367
            join(L, last[u] = new node(u, T() /* elemento
7 e 8
   neutro */), L);
a28
            join(L, R, root);
```

```
cbb
        }
4e6
        void cut(int v) {
892
            auto [1, r] = get_range(v);
            node *L, *M, *R;
df9
dca
            split(root, M, R, r+1), split(M, L, M, 1);
            node *LL = get max(L), *RR = get min(R);
de6
            if (LL and RR and LL->id == RR->id) { // remove
710
   duplicata
                 if (last[RR->id] == RR) last[RR->id] = LL;
e8b
992
                  node *A. *B:
6 b 3
                  split(R, A, B, 1);
10 c
                  delete A;
9 d5
                  R = B;
cbb
            }
a28
            join(L, R, root);
a0d
            join(root, M, root);
cbb
808
        T query(int v) {
            auto [1, r] = get_range(v);
892
df9
            node *L, *M, *R;
dca
            split(root, M, R, r+1), split(M, L, M, 1);
d43
            T ans = M->sub;
69 d
            join(L, M, M), join(M, R, root);
ba7
            return ans;
cbb
93b
        void update(int v, T val) { // soma val em todo
   mundo da subarvore
892
            auto [1, r] = get_range(v);
df9
            node *L, *M, *R;
            split(root, M, R, r+1), split(M, L, M, 1);
dca
            M->lazy += val;
409
69 d
            join(L, M, M), join(M, R, root);
cbb
129
        void update_v(int v, T val) { // muda o valor de v
   pra val
ac1
            int l = get_idx(first[v]);
df9
            node *L, *M, *R;
            split(root, M, R, l+1), split(M, L, M, l);
d0c
            M \rightarrow val = M \rightarrow sub = val;
25e
69 d
            join(L, M, M), join(M, R, root);
cbb
        }
```

```
934
        bool is_in_subtree(int v, int u) { // se u ta na
   subtree de v
890
            auto [lv, rv] = get_range(v);
6ec
            auto [lu, ru] = get_range(u);
732
            return lv <= lu and ru <= rv;</pre>
cbb
        }
355
        void print(node* i) {
eae
            if (!i) return;
            print(i->1);
a1e
743
            cout << i->id+1 << " ";
f 15
            print(i->r);
cbb
        }
        void print() { print(root); cout << endl; }</pre>
065
214 };
```

# 2.14 Floyd-Warshall

```
// encontra o menor caminho entre todo
// par de vertices e detecta ciclo negativo
// returna 1 sse ha ciclo negativo
// d[i][i] deve ser 0
// para i != j, d[i][j] deve ser w se ha uma aresta
// (i, j) de peso w, INF caso contrario
//
// O(n^3)
// ea05be
1a8 int n;
ae5 int d[MAX][MAX];
73c bool floyd_warshall() {
e22
        for (int k = 0; k < n; k++)
        for (int i = 0; i < n; i++)</pre>
830
        for (int j = 0; j < n; j++)
f90
            d[i][j] = min(d[i][j], d[i][k] + d[k][j]);
0ab
830
        for (int i = 0; i < n; i++)
753
            if (d[i][i] < 0) return 1;
bb3
        return 0;
cbb }
```

# 2.15 Functional Graph

```
// rt[i] fala o ID da raiz associada ao vertice i
// d[i] fala a profundidade (0 sse ta no ciclo)
// pos[i] fala a posicao de i no array que eh a concat. dos
   ciclos
// build(f, val) recebe a funcao f e o custo de ir de
// i para f[i] (por default, val = f)
// f_k(i, k) fala onde i vai parar se seguir k arestas
// path(i, k) fala o custo (soma) seguir k arestas a partir
   de i
// Se quiser outra operacao, da pra alterar facil o codigo
// Codigo um pouco louco, tenho que admitir
//
// build - O(n)
// f_k - O(log(min(n, k)))
// path - O(log(min(n, k)))
// 51fabe
6ef namespace func_graph {
1 a8
        int n;
ce2
        int f[MAX], vis[MAX], d[MAX];
f82
        int p[MAX], pp[MAX], rt[MAX], pos[MAX];
ebd
        int sz[MAX], comp;
        vector < vector < int >> ciclo;
6a9
405
        11 val[MAX], jmp[MAX], seg[2*MAX];
97 c
        11 op(ll a, ll b) { return a+b; }; // mudar a
   operacao aqui
        void dfs(int i, int t = 2) {
27b
9 c 9
            vis[i] = t:
f09
            if (vis[f[i]] >= 2) \{ // comeca ciclo - f[i] eh
   o rep.
                d[i] = 0, rt[i] = comp;
e0a
74c
                sz[comp] = t - vis[f[i]] + 1;
                p[i] = pp[i] = i, jmp[i] = val[i];
97b
                ciclo.emplace_back();
15c
                ciclo.back().push_back(i);
bfb
9d9
            } else {
c16
                if (!vis[f[i]]) dfs(f[i], t+1);
8 c 0
                rt[i] = rt[f[i]];
                if (sz[comp]+1) { // to no ciclo
195
```

```
d[i] = 0:
d0f
                    p[i] = pp[i] = i, jmp[i] = val[i];
97b
bfb
                    ciclo.back().push_back(i);
                } else { // nao to no ciclo
9d9
                    d[i] = d[f[i]]+1, p[i] = f[i];
00 d
511
                    pp[i] = 2*d[pp[f[i]]] ==
  d[pp[pp[f[i]]]+d[f[i]] ? pp[pp[f[i]]] : f[i];
                    jmp[i] = pp[i] == f[i] ? val[i] :
   op(val[i], op(jmp[f[i]], jmp[pp[f[i]]]));
cbb
            }
           if (f[ciclo[rt[i]][0]] == i) comp++; // fim do
e4a
   ciclo
            vis[i] = 1;
29a
cbb
1da
        void build(vector<int> f_, vector<int> val_ = {}) {
           n = f_size(), comp = 0;
bcb
           if (!val_.size()) val_ = f_;
527
           for (int i = 0; i < n; i++)
830
                f[i] = f[i], val[i] = val[i], vis[i] = 0,
998
   sz[i] = -1;
e74
           ciclo.clear();
158
           for (int i = 0; i < n; i++) if (!vis[i]) dfs(i);
           int t = 0;
6bb
           for (auto& c : ciclo) {
daa
336
                reverse(c.begin(), c.end());
ea5
                for (int j : c) {
                    pos[i] = t;
85b
948
                    seg[n+t] = val[j];
c82
                    t++;
                }
cbb
            }
cbb
            for (int i = n-1; i; i--) seg[i] = op(seg[2*i],
dc1
   seg[2*i+1]);
       }
cbb
283
        int f_k(int i, ll k) {
1 b 1
            while (d[i] and k) {
               int big = d[i] - d[pp[i]];
77b
ded
               if (big \le k) k = big, i = pp[i];
                else k--, i = p[i];
584
```

```
cbb
77e
           if (!k) return i;
           return ciclo[rt[i]][(pos[i] -
   pos[ciclo[rt[i]][0]] + k) % sz[rt[i]]];
cbb
047
        ll path(int i, ll k) {
            auto query = [&](int 1, int r) {
3cf
3 e 4
                11 q = 0;
                for (1 += n, r += n; 1 <= r; ++1/=2, --r/=2)
47a
27 e
                    if (1\%2 == 1) q = op(q, seg[1]);
1f2
                    if (r\%2 == 0) q = op(q, seg[r]);
                }
cbb
bef
                return q;
214
            };
b73
            11 \text{ ret} = 0:
1 b 1
            while (d[i] and k) {
77b
                int big = d[i] - d[pp[i]];
                if (big <= k) k -= big, ret = op(ret,</pre>
327
   jmp[i]), i = pp[i];
f9e
                else k--, ret = op(ret, val[i]), i = p[i];
cbb
            }
e3c
           if (!k) return ret;
           int first = pos[ciclo[rt[i]][0]], last =
   pos[ciclo[rt[i]].back()];
d41
           // k/sz[rt[i]] voltas completas
430
           if (k/sz[rt[i]]) ret = op(ret, k/sz[rt[i]] *
   query(first, last));
9af
            k %= sz[rt[i]];
e3c
            if (!k) return ret;
           int l = pos[i], r = first + (pos[i] - first + k
8ea
- 1) % sz[rt[i]]:
982
           if (1 <= r) return op(ret, query(1, r));</pre>
687
            return op(ret, op(query(l, last), query(first,
  r)));
      }
cbb
cbb }
```

## 2.16 Heavy-Light Decomposition - aresta

```
// SegTree de soma
// query / update de soma das arestas
// Complexidades:
// build - O(n)
// query path - 0(\log^2 (n))
// update path - O(\log^2 (n))
// query_subtree - O(log(n))
// update_subtree - O(log(n))
556 namespace seg { ... }
    // 599946
826 namespace hld {
c0d
        vector<pair<int, int> > g[MAX];
e65
        int pos[MAX], sz[MAX];
7c0
        int sobe[MAX], pai[MAX];
        int h[MAX], v[MAX], t;
096
        void build hld(int k, int p = -1, int f = 1) {
0ce
            v[pos[k] = t++] = sobe[k]; sz[k] = 1;
180
            for (auto& i : g[k]) if (i.first != p) {
418
                auto [u, w] = i;
dd2
                sobe[u] = w; pai[u] = k;
a76
                h[u] = (i == g[k][0] ? h[k] : u);
0 c 1
                build hld(u, k, f); sz[k] += sz[u];
da7
865
                if (sz[u] > sz[g[k][0].first] or
   g[k][0].first == p)
9a3
                     swap(i, g[k][0]);
cbb
            if (p*f == -1) build_hld(h[k] = k, -1, t = 0);
667
cbb
1f8
        void build(int root = 0) {
a34
            t = 0:
295
            build_hld(root);
            seg::build(t, v);
c83
cbb
        11 query_path(int a, int b) {
3fc
2d5
            if (a == b) return 0;
            if (pos[a] < pos[b]) swap(a, b);</pre>
aa1
```

```
29b
            if (h[a] == h[b]) return seg::query(pos[b]+1,
   pos[a]);
fca
            return seg::query(pos[h[a]], pos[a]) +
   query_path(pai[h[a]], b);
cbb
920
        void update_path(int a, int b, int x) {
            if (a == b) return;
d54
            if (pos[a] < pos[b]) swap(a, b);
aa1
            if (h[a] == h[b]) return
881
   (void)seg::update(pos[b]+1, pos[a], x);
701
            seg::update(pos[h[a]], pos[a], x);
   update_path(pai[h[a]], b, x);
cbb
d0a
        11 query_subtree(int a) {
b9f
            if (sz[a] == 1) return 0;
2f6
            return seg::query(pos[a]+1, pos[a]+sz[a]-1);
cbb
acc
        void update_subtree(int a, int x) {
            if (sz[a] == 1) return;
a5a
            seg::update(pos[a]+1, pos[a]+sz[a]-1, x);
9cd
cbb
7be
        int lca(int a, int b) {
            if (pos[a] < pos[b]) swap(a, b);</pre>
aa1
ca5
            return h[a] == h[b] ? b : lca(pai[h[a]], b);
cbb
cbb }
     Heavy-Light Decomposition - vertice
```

```
// SegTree de soma
// query / update de soma dos vertices
//
// Complexidades:
// build - O(n)
// query_path - O(log^2 (n))
// update_path - O(log^2(n))
// query_subtree - O(log(n))
// update subtree - O(\log(n))
556 namespace seg { ... }
```

```
// de3d84
826 namespace hld {
042
        vector < int > g[MAX];
e65
        int pos[MAX], sz[MAX];
        int peso[MAX], pai[MAX];
bd4
096
        int h[MAX], v[MAX], t;
0ce
        void build_hld(int k, int p = -1, int f = 1) {
            v[pos[k] = t++] = peso[k]; sz[k] = 1;
b18
b94
            for (auto& i : g[k]) if (i != p) {
78 d
                pai[i] = k;
26 e
                h[i] = (i == g[k][0] ? h[k] : i);
                build_hld(i, k, f); sz[k] += sz[i];
193
                if (sz[i] > sz[g[k][0]] or g[k][0] == p)
cd1
   swap(i, g[k][0]);
           }
cbb
            if (p*f == -1) build_hld(h[k] = k, -1, t = 0);
667
cbb
        void build(int root = 0) {
1f8
a34
            t = 0;
295
            build hld(root);
c83
            seg::build(t, v);
cbb
3fc
        11 query_path(int a, int b) {
aa1
            if (pos[a] < pos[b]) swap(a, b);
4bf
           if (h[a] == h[b]) return seg::query(pos[b],
   pos[a]);
            return seg::query(pos[h[a]], pos[a]) +
fca
   query_path(pai[h[a]], b);
cbb
        void update_path(int a, int b, int x) {
920
            if (pos[a] < pos[b]) swap(a, b);
aa1
198
            if (h[a] == h[b]) return
   (void)seg::update(pos[b], pos[a], x);
            seg::update(pos[h[a]], pos[a], x);
701
  update_path(pai[h[a]], b, x);
cbb
d0a
        11 query_subtree(int a) {
            return seg::query(pos[a], pos[a]+sz[a]-1);
b3e
```

## 2.18 Heavy-Light Decomposition sem Update

```
// query de min do caminho
//
// Complexidades:
// build - O(n)
// query_path - O(log(n))
// ee6991
826 namespace hld {
        vector < pair < int , int > > g[MAX];
c0d
e65
        int pos[MAX], sz[MAX];
7 c0
        int sobe[MAX], pai[MAX];
096
        int h[MAX], v[MAX], t;
ea2
        int men[MAX], seg[2*MAX];
        void build_hld(int k, int p = -1, int f = 1) {
0ce
180
            v[pos[k] = t++] = sobe[k]; sz[k] = 1;
            for (auto& i : g[k]) if (i.first != p) {
418
                sobe[i.first] = i.second; pai[i.first] = k;
1f5
6fa
                h[i.first] = (i == g[k][0] ? h[k] : i.first);
                men[i.first] = (i == g[k][0] ? min(men[k],
87b
   i.second) : i.second);
                build_hld(i.first, k, f); sz[k] +=
4 b 2
   sz[i.first];
bc3
                if (sz[i.first] > sz[g[k][0].first] or
   g[k][0].first == p)
                    swap(i, g[k][0]);
9 a 3
cbb
           if (p*f == -1) build_hld(h[k] = k, -1, t = 0);
667
cbb
```

```
1f8
        void build(int root = 0) {
            t = 0;
a34
295
            build hld(root);
           for (int i = 0; i < t; i++) seg[i+t] = v[i];</pre>
3ae
            for (int i = t-1; i; i--) seg[i] = min(seg[2*i],
8db
   seg[2*i+1]);
       }
cbb
f04
        int query_path(int a, int b) {
           if (a == b) return INF;
490
           if (pos[a] < pos[b]) swap(a, b);
aa1
98f
           if (h[a] != h[b]) return min(men[a],
   query_path(pai[h[a]], b));
           int ans = INF, x = pos[b]+1+t, y = pos[a]+t;
46 b
           for (; x \le y; ++x/=2, --y/=2) ans = min({ans,
   seg[x], seg[y]);
            return ans:
ba7
cbb
214 };
```

## 2.19 Isomorfismo de arvores

```
// thash() retorna o hash da arvore (usando centroids como
   vertices especiais).
// Duas arvores sao isomorfas sse seu hash eh o mesmo
//
// O(|V|.log(|V|))
// 8fb6bb
91f map < vector < int >, int > mphash;
df6 struct tree {
1a8
        int n:
789
        vector < vector < int >> g;
347
        vector<int> sz, cs;
        tree(int n_) : n(n_), g(n_), sz(n_) {}
1b5
76b
        void dfs_centroid(int v, int p) {
588
            sz[v] = 1;
fa7
            bool cent = true;
            for (int u : g[v]) if (u != p) {
18 e
```

```
365
                dfs centroid(u, v), sz[v] += sz[u]:
e90
                if(sz[u] > n/2) cent = false;
cbb
1f6
            if (cent and n - sz[v] <= n/2) cs.push_back(v);</pre>
cbb
784
        int fhash(int v, int p) {
544
            vector < int > h;
            for (int u : g[v]) if (u != p)
332
   h.push_back(fhash(u, v));
1 c9
            sort(h.begin(), h.end());
            if (!mphash.count(h)) mphash[h] = mphash.size();
3ac
bbc
            return mphash[h];
cbb
38f
        11 thash() {
23a
            cs.clear():
3a5
            dfs_centroid(0, -1);
16d
            if (cs.size() == 1) return fhash(cs[0], -1);
772
            11 h1 = fhash(cs[0], cs[1]), h2 = fhash(cs[1],
   cs[0]);
fae
            return (min(h1, h2) << 30) + max(h1, h2);
cbb
214 };
2.20 Kosaraju
```

```
// O(n + m)
// a4f310
1a8 int n;
042 vector <int > g[MAX];
58d vector <int > gi[MAX]; // grafo invertido
c5a int vis[MAX]:
ee6 stack<int> S:
a52 int comp[MAX]; // componente conexo de cada vertice
1ca void dfs(int k) {
        vis[k] = 1;
59a
        for (int i = 0; i < (int) g[k].size(); i++)</pre>
54f
            if (!vis[g[k][i]]) dfs(g[k][i]);
8 d5
        S.push(k);
58f
cbb }
```

```
436 void scc(int k, int c) {
        vis[k] = 1;
59a
52 c
        comp[k] = c;
        for (int i = 0; i < (int) gi[k].size(); i++)</pre>
ff0
bf6
            if (!vis[gi[k][i]]) scc(gi[k][i], c);
cbb }
db8 void kosaraju() {
        for (int i = 0; i < n; i++) vis[i] = 0;
991
        for (int i = 0; i < n; i++) if (!vis[i]) dfs(i);</pre>
158
        for (int i = 0; i < n; i++) vis[i] = 0;</pre>
991
        while (S.size()) {
d32
            int u = S.top();
70b
7de
            S.pop();
f43
            if (!vis[u]) scc(u, u);
        }
cbb
cbb }
```

### 2.21 Kruskal

```
// Gera e retorna uma AGM e seu custo total a partir do
   vetor de arestas (edg)
// do grafo
//
// O(m log(m) + m a(m))
// 864875
1b9 vector<tuple<int, int, int>> edg; // {peso,[x,y]}
    // DSU em O(a(n))
4a6 void dsu_build();
d78 int find(int a);
369 void unite(int a, int b);
c67 pair<ll, vector<tuple<int, int, int>>> kruskal(int n) {
8d2
        dsu build(n);
e31
        sort(edg.begin(), edg.end());
d41
854
        11 cost = 0;
979
        vector<tuple<int, int, int>> mst;
```

```
fea
        for (auto [w,x,y] : edg) if (find(x) != find(y)) {
            mst.emplace_back(w, x, y);
9de
45f
            cost += w;
05a
            unite(x,y);
cbb
5df
        return {cost, mst};
cbb }
2.22 Kuhn
// Computa matching maximo em grafo bipartido
// 'n' e 'm' sao quantos vertices tem em cada particao
// chamar add(i, j) para add aresta entre o cara i
// da particao A, e o cara j da particao B
// (entao i < n, j < m)
// Para recuperar o matching, basta olhar 'ma' e 'mb'
// 'recover' recupera o min vertex cover como um par de
// {caras da particao A, caras da particao B}
//
// O(|V| * |E|)
// Na pratica, parece rodar tao rapido quanto o Dinic
878 mt19937 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
    // b0dda3
6c6 struct kuhn {
14e
        int n, m;
789
        vector < vector < int >> g;
        vector < int > vis, ma, mb;
d3f
        kuhn(int n_, int m_) : n(n_), m(m_), g(n),
40e
8af
            vis(n+m), ma(n, -1), mb(m, -1) {}
        void add(int a, int b) { g[a].push_back(b); }
ba6
caf
        bool dfs(int i) {
29a
            vis[i] = 1;
            for (int j : g[i]) if (!vis[n+j]) {
29b
8 c 9
                vis[n+j] = 1;
                if (mb[j] == -1 or dfs(mb[j])) {
2cf
```

ma[i] = j, mb[j] = i;

bfe

```
8a6
                     return true:
                }
cbb
cbb
d1f
             return false;
cbb
bf7
        int matching() {
1ae
            int ret = 0, aum = 1;
5a8
            for (auto& i : g) shuffle(i.begin(), i.end(),
   rng);
392
             while (aum) {
618
                 for (int j = 0; j < m; j++) vis[n+j] = 0;
c5d
                for (int i = 0; i < n; i++)</pre>
830
                     if (ma[i] == -1 \text{ and } dfs(i)) \text{ ret++, aum} =
0.1 f
  1:
cbb
            }
edf
             return ret;
cbb
214 };
    // 55fb67
ebf pair < vector < int > , vector < int > > recover(kuhn & K) {
e80
        K.matching();
        int n = K.n, m = K.m;
50 c
        for (int i = 0; i < n+m; i++) K.vis[i] = 0;
9d0
        for (int i = 0; i < n; i++) if (K.ma[i] == -1)
bde
   K.dfs(i);
        vector < int > ca, cb;
8ad
        for (int i = 0; i < n; i++) if (!K.vis[i])</pre>
576
   ca.push_back(i);
        for (int i = 0; i < m; i++) if (K.vis[n+i])
f24
   cb.push_back(i);
        return {ca, cb};
aad
cbb }
2.23 LCA com binary lifting
// Assume que um vertice eh ancestral dele mesmo, ou seja,
// se a eh ancestral de b, lca(a, b) = a
// MAX2 = ceil(log(MAX))
//
// Complexidades:
```

```
// build - O(n log(n))
// lca - O(log(n))
677 vector < vector < int > > g(MAX);
41c int n, p;
e75 int pai[MAX2][MAX];
999 int in[MAX], out[MAX];
1ca void dfs(int k) {
fdf
        in[k] = p++;
        for (int i = 0; i < (int) g[k].size(); i++)</pre>
54f
9 b 7
            if (in[g[k][i]] == -1) {
                pai[0][g[k][i]] = k;
ba6
c38
                dfs(g[k][i]);
cbb
            }
26f
        out[k] = p++;
cbb }
c11 void build(int raiz) {
        for (int i = 0; i < n; i++) pai[0][i] = i;</pre>
a67
c63
        p = 0, memset(in, -1, sizeof in);
ecb
        dfs(raiz);
d41
        // pd dos pais
        for (int k = 1; k < MAX2; k++) for (int i = 0; i < MAX2
511
   n; i++)
            pai[k][i] = pai[k - 1][pai[k - 1][i]];
d38
cbb }
00f bool anc(int a, int b) { // se a eh ancestral de b
bfe
        return in[a] <= in[b] and out[a] >= out[b];
cbb }
7be int lca(int a, int b) {
86d
        if (anc(a, b)) return a;
        if (anc(b, a)) return b;
e52
        // sobe a
d41
        for (int k = MAX2 - 1; k >= 0; k--)
f70
acf
            if (!anc(pai[k][a], b)) a = pai[k][a];
847
        return pai[0][a];
```

```
cbb }
    // Alternativamente:
    // 'binary lifting' gastando O(n) de memoria
    // Da pra add folhas e fazer queries online
    // 3 vezes o tempo do binary lifting normal
    // build - O(n)
    // kth, lca, dist - O(log(n))
9c6 int d[MAX], p[MAX], pp[MAX];
d40 void set_root(int i) { p[i] = pp[i] = i, d[i] = 0; }
e9d void add_leaf(int i, int u) {
       p[i] = u, d[i] = d[u]+1;
e0b
       pp[i] = 2*d[pp[u]] == d[pp[pp[u]]]+d[u] ? pp[pp[u]]
b15
  : u:
cbb }
c37 int kth(int i, int k) {
4e3
        int dd = max(0, d[i]-k);
        while (d[i] > dd) i = d[pp[i]] >= dd? pp[i] : p[i];
935
d9a
        return i;
cbb }
7be int lca(int a, int b) {
a69
        if (d[a] < d[b]) swap(a, b);
        while (d[a] > d[b]) a = d[pp[a]] >= d[b] ? pp[a] :
6cd
   p[a];
984
        while (a != b) {
            if (pp[a] != pp[b]) a = pp[a], b = pp[b];
932
            else a = p[a], b = p[b];
e7c
        }
cbb
3f5
        return a;
cbb }
4fe int dist(int a, int b) { return d[a]+d[b]-2*d[lca(a,b)];
  }
042 vector < int > g[MAX];
```

```
3ab void build(int i, int pai=-1) {
5cf     if (pai == -1) set_root(i);
15f     for (int j : g[i]) if (j != pai) {
d31         add_leaf(j, i);
b21         build(j, i);
cbb }
```

## 2.24 LCA com HLD

```
// Assume que um vertice eh ancestral dele mesmo, ou seja,
// se a eh ancestral de b, lca(a, b) = a
// Para buildar pasta chamar build(root)
// anc(a, b) responde se 'a' eh ancestral de 'b'
//
// Complexidades:
// build - O(n)
// lca - O(\log(n))
// anc - O(1)
// fb22c1
042 vector < int > g[MAX];
713 int pos[MAX], h[MAX], sz[MAX];
ff1 int pai[MAX], t;
8bf void build(int k, int p = -1, int f = 1) {
        pos[k] = t++; sz[k] = 1;
bce
        for (int& i : g[k]) if (i != p) {
e26
78d
            pai[i] = k;
            h[i] = (i == g[k][0] ? h[k] : i);
26e
            build(i, k, f); sz[k] += sz[i];
cb8
d41
            if (sz[i] > sz[g[k][0]] or g[k][0] == p) swap(i,
cd1
   g[k][0]);
cbb
        if (p*f == -1) t = 0, h[k] = k, build (k, -1, 0);
3da
cbb }
7be int lca(int a, int b) {
        if (pos[a] < pos[b]) swap(a, b);</pre>
aa1
        return h[a] == h[b] ? b : lca(pai[h[a]], b);
ca5
cbb }
```

```
00f bool anc(int a, int b) {
db5     return pos[a] <= pos[b] and pos[b] <= pos[a]+sz[a]-1;
cbb }</pre>
```

## 2.25 LCA com RMQ

```
// Assume que um vertice eh ancestral dele mesmo, ou seja,
// se a eh ancestral de b, lca(a, b) = a
// dist(a, b) retorna a distancia entre a e b
// Complexidades:
// build - O(n)
// lca - 0(1)
// dist - 0(1)
// 22cde8 - rmg + lca
// 0214e8
1a5 template < typename T> struct rmq {
517
        vector<T> v;
fcc
        int n; static const int b = 30;
70e
        vector < int > mask, t;
        int op(int x, int y) { return v[x] < v[y] ? x : y; }</pre>
18 e
ee1
        int msb(int x) { return
   __builtin_clz(1)-__builtin_clz(x); }
        rmq() {}
6ad
        rmq(const vector < T > \& v_) : v(v_), n(v.size()),
43 c
   mask(n), t(n) {
            for (int i = 0, at = 0; i < n; mask[i++] = at |=
2e5
   1) {
                at = (at << 1) &((1 << b) -1);
a61
                while (at and op(i, i-msb(at&-at)) == i) at
76a
   ^= at&-at:
cbb
243
            for (int i = 0; i < n/b; i++) t[i] =
   b*i+b-1-msb(mask[b*i+b-1]);
            for (int j = 1; (1<<j) <= n/b; j++) for (int i =
   0; i+(1<<j) <= n/b; i++)
                t[n/b*j+i] = op(t[n/b*(j-1)+i],
   t[n/b*(j-1)+i+(1<<(j-1))]);
cbb
```

```
c92
        int small(int r, int sz = b) { return
   r-msb(mask[r]&((1<<sz)-1));}
b7a
        T query(int 1, int r) {
            if (r-l+1 <= b) return small(r, r-l+1);</pre>
27b
7bf
            int ans = op(small(1+b-1), small(r));
e80
            int x = 1/b+1, y = r/b-1;
e25
            if (x <= y) {
                int j = msb(y-x+1);
a4e
                ans = op(ans, op(t[n/b*j+x],
   t[n/b*j+y-(1<<j)+1]));
cbb
ba7
            return ans;
        }
214 };
    // 645120
065 namespace lca {
        vector < int > g[MAX];
042
        int v[2*MAX], pos[MAX], dep[2*MAX];
8ec
8 bd
        int t;
2de
        rmq<int> RMQ;
        void dfs(int i, int d = 0, int p = -1) {
4cf
            v[t] = i, pos[i] = t, dep[t++] = d;
c97
            for (int j : g[i]) if (j != p) {
cac
8ec
                dfs(i, d+1, i);
                v[t] = i, dep[t++] = d;
cf2
cbb
            }
cbb
789
        void build(int n, int root) {
a34
            t = 0;
14e
            dfs(root):
3f4
            RMQ = rmq < int > (vector < int > (dep, dep + 2*n - 1));
cbb
7be
        int lca(int a, int b) {
ab7
            a = pos[a], b = pos[b];
            return v[RMQ.query(min(a, b), max(a, b))];
9 c 0
cbb
b5d
        int dist(int a, int b) {
            return dep[pos[a]] + dep[pos[b]] -
670
   2*dep[pos[lca(a, b)]];
cbb
```

```
cbb }
```

## 2.26 Line Tree

```
// Reduz min-query em arvore para RMQ
// Se o grafo nao for uma arvore, as queries
// sao sobre a arvore geradora maxima
// Queries de minimo
// build - O(n log(n))
// query - O(log(n))
// b1f418
1a8 int n;
3ae namespace linetree {
f37
        int id[MAX], seg[2*MAX], pos[MAX];
43f
        vector < int > v[MAX], val[MAX];
430
        vector<pair<int, pair<int, int> > ar;
        void add(int a, int b, int p) { ar.push_back({p, {a,
dc6
   b}}); }
      void build() {
0a8
b09
           sort(ar.rbegin(), ar.rend());
           for (int i = 0; i < n; i++) id[i] = i, v[i] =
   {i}, val[i].clear();
         for (auto i : ar) {
8bb
               int a = id[i.second.first], b =
c.9.1
   id[i.second.second];
               if (a == b) continue;
f6f
c58
               if (v[a].size() < v[b].size()) swap(a, b);
                for (auto j : v[b]) id[j] = a,
fb8
   v[a].push_back(j);
482
                val[a].push_back(i.first);
                for (auto j : val[b]) val[a].push_back(j);
78b
                v[b].clear(), val[b].clear();
e39
cbb
           }
8e8
            vector<int> vv;
           for (int i = 0; i < n; i++) for (int j = 0; j <
2ce
   v[i].size(); j++) {
               pos[v[i][j]] = vv.size();
e52
               if (j + 1 < v[i].size())</pre>
941
```

```
vv.push_back(val[i][j]);
1 cb
                else vv.push_back(0);
cbb
           for (int i = n; i < 2*n; i++) seg[i] = vv[i-n];</pre>
bb4
           for (int i = n-1; i; i--) seg[i] = min(seg[2*i],
   seg[2*i+1]);
cbb
4ea
        int query(int a, int b) {
            if (id[a] != id[b]) return 0; // nao estao
   conectados
            a = pos[a], b = pos[b];
ab7
           if (a > b) swap(a, b);
d11
199
           b - -;
           int ans = INF;
38a
          for (a += n, b += n; a <= b; ++a/=2, --b/=2) ans
   = min({ans, seg[a], seg[b]});
            return ans:
ba7
       }
cbb
214 };
2.27 Link-cut Tree
// Link-cut tree padrao
//
// Todas as operacoes sao O(\log(n)) amortizado
// e4e663
1ef namespace lct {
3 c 9
        struct node {
19f
            int p, ch[2];
062
            node() \{ p = ch[0] = ch[1] = -1; \}
214
        }:
        node t[MAX];
5f3
971
        bool is root(int x) {
            return t[x].p == -1 or (t[t[x].p].ch[0] != x and
657
   t[t[x].p].ch[1] != x);
cbb
        void rotate(int x) {
ed6
          int p = t[x].p, pp = t[p].p;
497
```

if (!is\_root(p)) t[pp].ch[t[pp].ch[1] == p] = x;

fc4

```
251
            bool d = t[p].ch[0] == x;
            t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
461
            if (t[p].ch[!d]+1) t[t[p].ch[!d]].p = p;
a76
8fa
            t[x].p = pp, t[p].p = x;
cbb
        }
07с
        void splay(int x) {
18 c
            while (!is root(x)) {
                int p = t[x].p, pp = t[p].p;
497
                if (!is_root(p)) rotate((t[pp].ch[0] ==
0 c 5
   p)^(t[p].ch[0] == x) ? x : p);
64f
                rotate(x);
            }
cbb
        }
cbb
        int access(int v) {
f 16
0eb
            int last = -1:
01a
            for (int w = v; w+1; last = w, splay(v), w =
   t[v].p)
                splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
024
3d3
            return last;
cbb
e89
        int find_root(int v) {
5e3
            access(v);
3de
            while (t[v].ch[0]+1) v = t[v].ch[0];
f05
            return splay(v), v;
cbb
        void link(int v, int w) { // v deve ser raiz
142
5e3
            access(v);
10 d
            t[v].p = w;
cbb
        void cut(int v) { // remove aresta de v pro pai
4e6
            access(v);
5e3
264
            t[v].ch[0] = t[t[v].ch[0]].p = -1;
cbb
        int lca(int v, int w) {
bbb
948
            return access(v), access(w);
cbb
cbb }
      Link-cut Tree - aresta
2.28
// Valores nas arestas
```

```
// rootify(v) torna v a raiz de sua arvore
```

```
// querv(v, w) retorna a soma do caminho v--w
// update(v, w, x) soma x nas arestas do caminho v--w
// Todas as operacoes sao O(\log(n)) amortizado
// 9ce48f
1ef namespace lct {
        struct node {
3 c 9
19f
            int p, ch[2];
810
            ll val, sub;
aa6
            bool rev;
04a
            int sz, ar;
4e4
            ll lazy;
f93
            node() {}
7 a 8
            node(int v, int ar_) :
            p(-1), val(v), sub(v), rev(0), sz(ar_), ar(ar_),
546
   lazy(0) {
                ch[0] = ch[1] = -1;
b07
            }
cbb
214
        };
c53
        node t[2*MAX]; // MAXN + MAXQ
99e
        map < pair < int , int > , int > aresta;
e4d
        int sz;
95a
        void prop(int x) {
dc1
            if (t[x].lazy) {
25e
                if (t[x].ar) t[x].val += t[x].lazy;
2ab
                t[x].sub += t[x].lazy*t[x].sz;
                if (t[x].ch[0]+1) t[t[x].ch[0]].lazy +=
edc
   t[x].lazy;
942
                if (t[x].ch[1]+1) t[t[x].ch[1]].lazy +=
   t[x].lazy;
cbb
            if (t[x].rev) {
aa2
                swap(t[x].ch[0], t[x].ch[1]);
f95
379
                if (t[x].ch[0]+1) t[t[x].ch[0]].rev ^= 1;
                if (t[x].ch[1]+1) t[t[x].ch[1]].rev ^= 1;
c3d
cbb
230
            t[x].lazy = 0, t[x].rev = 0;
cbb
564
        void update(int x) {
```

```
1a3
            t[x].sz = t[x].ar, t[x].sub = t[x].val:
                                                                 9f0
                                                                              while (t[v].ch[0]+1) v = t[v].ch[0], prop(v);
            for (int i = 0; i < 2; i++) if (t[x].ch[i]+1) {
                                                                 637
8ca
                                                                              return splay(v);
                prop(t[x].ch[i]);
621
                                                                 cbb
                t[x].sz += t[t[x].ch[i]].sz;
c4f
                                                                 82f
                                                                         bool conn(int v, int w) {
                t[x].sub += t[t[x].ch[i]].sub;
269
                                                                 2cf
                                                                              access(v), access(w);
                                                                              return v == w ? true : t[v].p != -1;
cbb
            }
                                                                 b9b
                                                                 cbb
cbb
971
                                                                 277
        bool is root(int x) {
                                                                          void rootify(int v) {
            return t[x].p == -1 or (t[t[x].p].ch[0] != x and
657
                                                                 5 e 3
                                                                              access(v);
   t[t[x].p].ch[1] != x);
                                                                 a02
                                                                              t[v].rev ^= 1;
cbb
                                                                 cbb
        void rotate(int x) {
                                                                 971
                                                                         11 query(int v, int w) {
ed6
            int p = t[x].p, pp = t[p].p;
                                                                              rootify(w), access(v);
497
                                                                 b54
            if (!is_root(p)) t[pp].ch[t[pp].ch[1] == p] = x;
                                                                 249
                                                                              return t[v].sub;
fc4
            bool d = t[p].ch[0] == x;
251
                                                                 cbb
            t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
                                                                          void update(int v, int w, int x) {
461
                                                                 3fa
a76
           if (t[p].ch[!d]+1) t[t[p].ch[!d]].p = p;
                                                                 b54
                                                                              rootify(w), access(v);
8fa
            t[x].p = pp, t[p].p = x;
                                                                 12c
                                                                              t[v].lazy += x;
            update(p), update(x);
444
                                                                 cbb
                                                                 204
cbb
                                                                          void link (int v, int w) {
238
                                                                 821
                                                                              rootify(w);
        int splay(int x) {
18 c
            while (!is root(x)) {
                                                                 389
                                                                              t[w].p = v;
                int p = t[x].p, pp = t[p].p;
497
                                                                 cbb
77b
                if (!is_root(p)) prop(pp);
                                                                 6 b 8
                                                                          void link(int v, int w, int x) { // v--w com peso x
                prop(p), prop(x);
                                                                 379
                                                                              int id = MAX + sz++;
be5
                if (!is_root(p)) rotate((t[pp].ch[0] ==
                                                                              aresta[make_pair(v, w)] = id;
0 c 5
                                                                 110
  p)^{(t[p].ch[0] == x)} ? x : p);
                                                                 a88
                                                                              make tree(id, x, 1);
64f
                rotate(x);
                                                                 c88
                                                                              link_(v, id), link_(id, w);
            }
                                                                 cbb
cbb
aab
            return prop(x), x;
                                                                 e63
                                                                          void cut_(int v, int w) {
cbb
                                                                 b54
                                                                              rootify(w), access(v);
        int access(int v) {
                                                                 264
                                                                              t[v].ch[0] = t[t[v].ch[0]].p = -1;
f16
0eb
            int last = -1;
                                                                 cbb
            for (int w = v; w+1; update(last = w), splay(v),
                                                                 031
                                                                          void cut(int v, int w) {
d9f
  w = t[v].p
                                                                 b0f
                                                                              int id = aresta[make_pair(v, w)];
024
                splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
                                                                 a4a
                                                                              cut_(v, id), cut_(id, w);
3d3
            return last;
                                                                 cbb
cbb
                                                                 bbb
                                                                         int lca(int v, int w) {
        void make_tree(int v, int w=0, int ar=0) { t[v] =
                                                                 5 e 3
                                                                              access(v);
9f1
   node(w, ar); }
                                                                 a8b
                                                                              return access(w);
        int find_root(int v) {
                                                                 cbb
e89
            access(v), prop(v);
                                                                 cbb }
13f
```

### 2.29 Link-cut Tree - vertice

```
// Valores nos vertices
// make tree(v, w) cria uma nova arvore com um
// vertice soh com valor 'w'
// rootify(v) torna v a raiz de sua arvore
// query(v, w) retorna a soma do caminho v--w
// update(v, w, x) soma x nos vertices do caminho v--w
// Todas as operacoes sao O(\log(n)) amortizado
// f9f489
1ef namespace lct {
3c9
        struct node {
19f
            int p, ch[2];
810
            ll val, sub;
           bool rev;
aa6
e4d
           int sz;
4e4
           ll lazv;
f93
            node() {}
            node(int v) : p(-1), val(v), sub(v), rev(0),
aa0
   sz(1), lazv(0) {
                ch[0] = ch[1] = -1;
b07
cbb
            }
214
       };
5f3
       node t[MAX];
95a
        void prop(int x) {
dc1
            if (t[x].lazy) {
9f7
                t[x].val += t[x].lazy, t[x].sub +=
   t[x].lazy*t[x].sz;
                if (t[x].ch[0]+1) t[t[x].ch[0]].lazy +=
edc
   t[x].lazy;
                if (t[x].ch[1]+1) t[t[x].ch[1]].lazy +=
942
   t[x].lazv;
cbb
            if (t[x].rev) {
aa2
                swap(t[x].ch[0], t[x].ch[1]);
f95
379
                if (t[x].ch[0]+1) t[t[x].ch[0]].rev ^= 1;
c3d
                if (t[x].ch[1]+1) t[t[x].ch[1]].rev ^= 1;
cbb
            }
```

```
230
            t[x].lazv = 0, t[x].rev = 0:
cbb
564
        void update(int x) {
            t[x].sz = 1, t[x].sub = t[x].val;
ec2
8ca
            for (int i = 0; i < 2; i++) if (t[x].ch[i]+1) {
621
                prop(t[x].ch[i]);
c4f
                t[x].sz += t[t[x].ch[i]].sz;
                t[x].sub += t[t[x].ch[i]].sub;
269
cbb
            }
cbb
971
        bool is_root(int x) {
657
            return t[x].p == -1 or (t[t[x].p].ch[0] != x and
   t[t[x].p].ch[1] != x);
cbb
ed6
        void rotate(int x) {
            int p = t[x].p, pp = t[p].p;
497
fc4
            if (!is_root(p)) t[pp].ch[t[pp].ch[1] == p] = x;
251
            bool d = t[p].ch[0] == x;
461
            t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
a76
            if (t[p].ch[!d]+1) t[t[p].ch[!d]].p = p;
            t[x].p = pp, t[p].p = x;
8fa
444
            update(p), update(x);
cbb
238
        int splay(int x) {
18 c
            while (!is_root(x)) {
                int p = t[x].p, pp = t[p].p;
497
77b
                if (!is_root(p)) prop(pp);
be5
                prop(p), prop(x);
0 c5
                if (!is_root(p)) rotate((t[pp].ch[0] ==
   p)^(t[p].ch[0] == x) ? x : p);
64f
                rotate(x);
cbb
aab
            return prop(x), x;
cbb
f16
        int access(int v) {
0eb
            int last = -1:
            for (int w = v; w+1; update(last = w), splay(v),
d9f
   w = t[v].p
024
                splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
3 d3
            return last;
cbb
f 17
        void make_tree(int v, int w) { t[v] = node(w); }
```

```
e89
        int find root(int v) {
13f
            access(v), prop(v);
            while (t[v].ch[0]+1) v = t[v].ch[0], prop(v);
9f0
637
            return splay(v);
cbb
        }
f94
        bool connected(int v, int w) {
            access(v), access(w);
2cf
            return v == w ? true : t[v].p != -1;
b9b
cbb
        }
277
        void rootify(int v) {
            access(v);
5e3
a02
            t[v].rev ^= 1;
        }
cbb
971
        11 query(int v, int w) {
            rootify(w), access(v);
b54
249
            return t[v].sub;
        }
cbb
        void update(int v, int w, int x) {
3fa
            rootify(w), access(v);
b54
            t[v].lazy += x;
12 c
cbb
        }
142
        void link(int v, int w) {
821
            rootify(w);
389
            t[w].p = v;
cbb
031
        void cut(int v, int w) {
b54
            rootify(w), access(v);
            t[v].ch[0] = t[t[v].ch[0]].p = -1;
264
cbb
        int lca(int v, int w) {
bbb
            access(v);
5e3
a8b
            return access(w);
        }
cbb
cbb }
```

## 2.30 Max flow com lower bound nas arestas

```
// add(a, b, 1, r):
// adiciona aresta de a pra b, onde precisa passar f de
   fluxo, l <= f <= r
// add(a, b, c):
// adiciona aresta de a pra b com capacidade c</pre>
```

```
//
// Mesma complexidade do Dinic
// 5f2379
919 struct lb_max_flow : dinic {
5 се
        vector < int > d;
        lb max flow(int n): dinic(n + 2), d(n, 0) {}
331
        void add(int a, int b, int 1, int r) {
b12
c97
            d[a] -= 1;
f1b
            d[b] += 1;
017
            dinic::add(a, b, r - 1);
cbb
087
        void add(int a, int b, int c) {
107
            dinic::add(a, b, c);
cbb
7a1
        bool has_circulation() {
50 c
            int n = d.size();
854
            11 cost = 0;
            for (int i = 0; i < n; i++) {</pre>
603
                if (d[i] > 0) {
c69
f56
                     cost += d[i];
d06
                     dinic::add(n, i, d[i]);
                } else if (d[i] < 0) {</pre>
9 c 7
                     dinic::add(i, n+1, -d[i]);
76b
cbb
                }
cbb
            }
283
            return (dinic::max_flow(n, n+1) == cost);
cbb
7bd
        bool has_flow(int src, int snk) {
65d
            dinic::add(snk, src, INF);
e40
            return has_circulation();
cbb
4eb
        11 max_flow(int src, int snk) {
ee8
            if (!has_flow(src, snk)) return -1;
            dinic::F = 0:
ea5
626
            return dinic::max_flow(src, snk);
cbb
214 };
```

### 2.31 MinCostMaxFlow

```
// min cost flow(s, t, f) computa o par (fluxo, custo)
// com max(fluxo) <= f que tenha min(custo)</pre>
// min_cost_flow(s, t) -> Fluxo maximo de custo minimo de s
// Se for um dag, da pra substituir o SPFA por uma DP pra nao
// pagar O(nm) no comeco
// Se nao tiver aresta com custo negativo, nao precisa do
//
// O(nm + f * m log n)
// 697b4c
123 template < typename T> struct mcmf {
        struct edge {
670
b75
            int to, rev, flow, cap; // para, id da reversa,
  fluxo, capacidade
7f9
           bool res; // se eh reversa
            T cost; // custo da unidade de fluxo
635
892
           edge(): to(0), rev(0), flow(0), cap(0),
   cost(0), res(false) {}
1d7
            edge(int to_, int rev_, int flow_, int cap_, T
   cost , bool res )
               : to(to_), rev(rev_), flow(flow_),
   cap(cap_), res(res_), cost(cost_) {}
214
        }:
002
        vector < vector < edge >> g;
168
        vector<int> par_idx, par;
f1e
        T inf:
a03
        vector<T> dist:
        mcmf(int n) : g(n), par_idx(n), par(n),
b22
   inf(numeric_limits <T>::max()/3) {}
        void add(int u, int v, int w, T cost) { // de u pra
91 c
   v com cap w e custo cost
            edge a = edge(v, g[v].size(), 0, w, cost, false);
2fc
            edge b = edge(u, g[u].size(), 0, 0, -cost, true);
234
            g[u].push_back(a);
b24
```

```
g[v].push_back(b);
c12
cbb
        vector <T> spfa(int s) { // nao precisa se nao tiver
8bc
   custo negativo
871
             deque < int > q;
            vector <bool> is_inside(g.size(), 0);
3 d 1
            dist = vector <T > (g.size(), inf);
577
a93
             dist[s] = 0;
a30
             q.push_back(s);
ecb
            is_inside[s] = true;
14d
             while (!q.empty()) {
b1e
                 int v = q.front();
ced
                 q.pop_front();
48 d
                 is_inside[v] = false;
76e
                 for (int i = 0; i < g[v].size(); i++) {</pre>
                     auto [to, rev, flow, cap, res, cost] =
9d4
   g[v][i];
e61
                     if (flow < cap and dist[v] + cost <</pre>
   dist[to]) {
943
                          dist[to] = dist[v] + cost;
ed6
                         if (is_inside[to]) continue;
020
                         if (!q.empty() and dist[to] >
   dist[q.front()]) q.push_back(to);
b33
                         else q.push_front(to);
b52
                         is_inside[to] = true;
cbb
                     }
                 }
cbb
cbb
            }
8 d7
             return dist;
cbb
2a2
        bool dijkstra(int s, int t, vector <T>& pot) {
489
             priority_queue <pair <T, int>, vector <pair <T,</pre>
   int>>, greater<>> q;
577
             dist = vector <T>(g.size(), inf);
a93
             dist[s] = 0;
115
             q.emplace(0, s);
402
             while (q.size()) {
```

```
91b
                 auto [d, v] = q.top();
833
                 q.pop();
                 if (dist[v] < d) continue;</pre>
68b
                 for (int i = 0; i < g[v].size(); i++) {</pre>
76 e
                      auto [to, rev, flow, cap, res, cost] =
9d4
   g[v][i];
                      cost += pot[v] - pot[to];
e8c
                     if (flow < cap and dist[v] + cost <</pre>
e61
   dist[to]) {
                          dist[to] = dist[v] + cost;
943
441
                          q.emplace(dist[to], to);
88b
                          par_idx[to] = i, par[to] = v;
                     }
cbb
                 }
cbb
cbb
1 d 4
             return dist[t] < inf;</pre>
        }
cbb
        pair < int , T > min_cost_flow(int s, int t, int flow =
3d2
   INF) {
3dd
             vector<T> pot(g.size(), 0);
9e4
             pot = spfa(s); // mudar algoritmo de caminho
   minimo aqui
d22
            int f = 0;
            T ret = 0;
ce8
4a0
             while (f < flow and dijkstra(s, t, pot)) {</pre>
bda
                 for (int i = 0; i < g.size(); i++)</pre>
                     if (dist[i] < inf) pot[i] += dist[i];</pre>
d2a
71b
                 int mn_flow = flow - f, u = t;
                 while (u != s){
045
90f
                      mn_flow = min(mn_flow,
07 d
                          g[par[u]][par_idx[u]].cap -
   g[par[u]][par_idx[u]].flow);
3d1
                      u = par[u];
                 }
cbb
1f2
                 ret += pot[t] * mn_flow;
476
                 u = t;
045
                 while (u != s) {
```

```
e09
                     g[par[u]][par_idx[u]].flow += mn_flow;
                     g[u][g[par[u]][par_idx[u]].rev].flow -=
d98
   mn flow;
                     u = par[u];
3 d 1
                }
cbb
04d
                f += mn flow;
cbb
15b
            return make_pair(f, ret);
cbb
d41
        // Opcional: retorna as arestas originais por onde
   passa flow = cap
182
        vector < pair < int , int >> recover() {
24a
            vector <pair <int, int>> used;
2a4
            for (int i = 0; i < g.size(); i++) for (edge e:
   g[i])
                if (e.flow == e.cap && !e.res)
587
   used.push back({i, e.to});
f6b
            return used;
cbb
       }
214 };
2.32 Prufer code
// Traduz de lista de arestas para prufer code
// e vice-versa
// Os vertices tem label de 0 a n-1
// Todo array com n-2 posicoes e valores de
// O a n-1 sao prufer codes validos
//
// O(n)
// d3b324
47d vector <int > to_prufer(vector <pair <int, int >> tree) {
1fa
        int n = tree.size()+1;
2cf
        vector < int > d(n, 0);
4aa
        vector < vector < int >> g(n);
        for (auto [a, b] : tree) d[a]++, d[b]++,
f87
f60
            g[a].push_back(b), g[b].push_back(a);
```

vector < int > pai(n, -1);

c5a

```
260
        queue < int > q; q.push(n-1);
        while (q.size()) {
402
            int u = q.front(); q.pop();
be1
            for (int v : g[u]) if (v != pai[u])
34 c
9c9
                 pai[v] = u, q.push(v);
cbb
        }
399
        int idx, x;
        idx = x = find(d.begin(), d.end(), 1) - d.begin();
897
4b8
        vector<int> ret;
        for (int i = 0; i < n-2; i++) {</pre>
b28
            int y = pai[x];
d4b
e81
            ret.push_back(y);
           if (--d[y] == 1 \text{ and } y < idx) x = y;
666
            else idx = x = find(d.begin()+idx+1, d.end(), 1)
367
   - d.begin();
cbb
edf
        return ret;
cbb }
    // 765413
4d8 vector<pair<int, int>> from_prufer(vector<int> p) {
455
        int n = p.size()+2;
126
        vector < int > d(n, 1);
650
        for (int i : p) d[i]++;
        p.push_back(n-1);
85b
399
        int idx, x;
        idx = x = find(d.begin(), d.end(), 1) - d.begin();
897
1df
        vector<pair<int, int>> ret;
        for (int y : p) {
b06
            ret.push_back({x, y});
dab
666
            if (--d[y] == 1 \text{ and } y < idx) x = y;
            else idx = x = find(d.begin()+idx+1, d.end(), 1)
367
   - d.begin();
        }
cbb
        return ret:
edf
cbb }
      Sack (DSU em arvores)
// Responde queries de todas as sub-arvores
// offline
//
```

```
// O(n log(n))
// bb361f
6bf int sz[MAX], cor[MAX], cnt[MAX];
042 vector < int > g[MAX];
d41
6df void build(int k, int d=0) {
        sz[k] = 1;
e8f
01a
        for (auto& i : g[k]) {
            build(i, d+1); sz[k] += sz[i];
30f
            if (sz[i] > sz[g[k][0]]) swap(i, g[k][0]);
925
cbb
cbb }
d41
74f void compute(int k, int x, bool dont=1) {
de9
        cnt[cor[k]] += x;
828
        for (int i = dont; i < g[k].size(); i++)</pre>
b5c
            compute(g[k][i], x, 0);
cbb }
d41
dc4 void solve(int k, bool keep=0) {
32a
        for (int i = int(g[k].size())-1; i >= 0; i--)
b4c
            solve(g[k][i], !i);
4a0
        compute(k, 1);
d41
        // agora cnt[i] tem quantas vezes a cor
d41
d41
        // i aparece na sub-arvore do k
830
        if (!keep) compute(k, -1, 0);
cbb }
2.34 Tarjan para SCC
```

```
// O(n + m)
// 573bfa
042 vector < int > g[MAX];
4ce stack<int> s;
a42 int vis[MAX], comp[MAX];
3fd int id[MAX];
    // se quiser comprimir ciclo ou achar ponte em grafo nao
```

```
46e
                                                                          vector < int > ret(n,-1), vis(n,0);
       direcionado,
    // colocar um if na dfs para nao voltar pro pai da DFS
                                                                  d41
                                                                  f51
                                                                          int pos = n-1, dag = 1;
                                                                          function < void(int) > dfs = [&](int v) {
f32 int dfs(int i, int& t) {
                                                                  36d
        int lo = id[i] = t++;
                                                                              vis[v] = 1;
cf0
                                                                  cca
18 e
        s.push(i);
                                                                  440
                                                                              for (auto u : g[v]) {
0 c 2
        vis[i] = 2;
                                                                                  if (vis[u] == 1) dag = 0;
                                                                  152
                                                                                  else if (!vis[u]) dfs(u);
                                                                  532
        for (int j : g[i]) {
48 e
                                                                  cbb
                                                                              }
            if (!vis[j]) lo = min(lo, dfs(j, t));
                                                                  d44
                                                                              ret[pos--] = v, vis[v] = 2;
740
994
            else if (vis[j] == 2) lo = min(lo, id[j]);
                                                                  214
                                                                          }:
        }
cbb
                                                                  158
                                                                          for (int i = 0; i < n; i++) if (!vis[i]) dfs(i);
        // aresta de i pro pai eh uma ponte (no caso nao
d41
   direcionado)
                                                                  d8f
                                                                          if (!dag) ret.clear();
3de
        if (lo == id[i]) while (1) {
                                                                  edf
                                                                          return ret:
3c3
           int u = s.top(); s.pop();
                                                                  cbb }
           vis[u] = 1, comp[u] = i;
9 c 5
           if (u == i) break;
2ef
                                                                  2.36 Vertex cover
cbb
        }
                                                                  // Encontra o tamanho do vertex cover minimo
253
        return lo;
                                                                  // Da pra alterar facil pra achar os vertices
cbb }
                                                                  // Parece rodar com < 2 s pra N = 90
                                                                  //
f93 void tarjan(int n) {
                                                                  // O(n * 1.38^n)
        int t = 0;
6bb
                                                                  // 9c5024
        for (int i = 0; i < n; i++) vis[i] = 0;
991
                                                                 76a namespace cover {
        for (int i = 0; i < n; i++) if (!vis[i]) dfs(i, t);
3be
                                                                  5a4
                                                                          const int MAX = 96;
cbb }
                                                                  042
                                                                          vector < int > g[MAX];
                                                                  823
                                                                          bitset < MAX > bs [MAX]:
       Topological Sort
2.35
                                                                  1 a8
                                                                          int n:
// Retorna uma ordenacaoo topologica de g
                                                                  697
                                                                          void add(int i, int j) {
// Se g nao for DAG retorna um vetor vazio
                                                                  bd0
                                                                              if (i == j) return;
                                                                              n = max({n, i+1, j+1});
//
                                                                  78 c
                                                                              bs[i][j] = bs[j][i] = 1;
// O(n + m)
                                                                  200
// bdc95e
                                                                          }
                                                                  cbb
042 vector < int > g[MAX];
                                                                          int rec(bitset < MAX > m) {
                                                                  6 c 0
                                                                  1 a 4
                                                                              int ans = 0;
                                                                              for (int x = 0; x < n; x++) if (m[x]) {
b6a vector<int> topo_sort(int n) {
                                                                  25b
```

```
002
                 bitset < MAX > comp;
                 function < void(int) > dfs = [&](int i) {
4bf
b96
                     comp[i] = 1, m[i] = 0;
                     for (int j : g[i]) if (m[j]) dfs(j);
0 c 3
214
                 };
963
                 dfs(x);
d34
                 int ma, deg = -1, cyc = 1;
417
                 for (int i = 0; i < n; i++) if (comp[i]) {
                     int d = (bs[i]&comp).count();
d0b
18a
                     if (d \le 1) cyc = 0;
                     if (d > deg) deg = d, ma = i;
c1f
                 }
cbb
                 if (deg <= 2) { // caminho ou ciclo</pre>
269
340
                     ans += (comp.count() + cyc) / 2;
5e2
                     continue:
                 }
cbb
                 comp[ma] = 0;
3f9
                // ou ta no cover, ou nao ta no cover
d41
1 dd
                 ans += min(1 + rec(comp), deg + rec(comp & \sim
   bs[ma]));
cbb
             return ans;
ba7
cbb
f5c
        int solve() {
3c5
            bitset < MAX > m;
603
            for (int i = 0; i < n; i++) {
939
                 m[i] = 1;
                 for (int j = 0; j < n; j++)
f90
741
                     if (bs[i][j]) g[i].push_back(j);
cbb
4f9
            return rec(m);
        }
cbb
cbb }
      Virtual Tree
2.37
```

```
// Comprime uma arvore dado um conjunto S de vertices, de
  forma que
// o conjunto de vertices da arvore comprimida contenha S e
  seja
```

```
// minimal e fechado sobre a operacao de LCA
// Se |S| = k, a arvore comprimida tem menos que 2k vertices
// As arestas de virt possuem a distancia do vertice ate o
   vizinho
// Retorna a raiz da virtual tree
//
// lca::pos deve ser a ordem de visitacao no dfs
// voce pode usar o LCAcomHLD, por exemplo
//
// O(k log(k))
// 42d990
b36 vector <pair <int, int>> virt[MAX];
d41 #warning lembrar de buildar o LCA antes
c14 int build_virt(vector<int> v) {
        auto cmp = [&](int i, int j) { return lca::pos[i] <</pre>
   lca::pos[i]; };
074
        sort(v.begin(), v.end(), cmp);
        for (int i = v.size()-1; i; i--)
e85
   v.push_back(lca::lca(v[i], v[i-1]));
074
        sort(v.begin(), v.end(), cmp);
d76
        v.erase(unique(v.begin(), v.end()), v.end());
        for (int i = 0; i < v.size(); i++)</pre>
37 c
   virt[v[i]].clear();
197
        for (int i = 1; i < v.size(); i++)</pre>
   virt[lca::lca(v[i-1], v[i])].clear();
ad7
        for (int i = 1; i < v.size(); i++) {</pre>
            int parent = lca::lca(v[i-1], v[i]);
51b
290
            int d = lca::dist(parent, v[i]);
d41 #warning soh to colocando aresta descendo
4 d 0
            virt[parent].emplace_back(v[i], d);
cbb
832
        return v[0]:
cbb }
```

# 3 Problemas

# 3.1 Algoritmo Hungaro

```
// Resolve o problema de assignment (matriz n x n)
// Colocar os valores da matriz em 'a' (pode < 0)</pre>
// assignment() retorna um par com o valor do
// assignment minimo, e a coluna escolhida por cada linha
//
// O(n^3)
// 64c53e
a6a template < typename T > struct hungarian {
1a8
        int n;
a08
        vector < vector < T >> a;
        vector <T> u, v;
f36
5ff
        vector < int > p, way;
f1e
        T inf;
c3f
        hungarian(int n_{-}): n(n_{-}), u(n+1), v(n+1), p(n+1),
   way(n+1) {
             a = vector < vector < T > (n, vector < T > (n));
b2f
             inf = numeric_limits <T>::max();
1f3
cbb
d67
        pair <T, vector <int>> assignment() {
78a
             for (int i = 1; i <= n; i++) {</pre>
8c9
                 p[0] = i;
625
                 int j0 = 0;
                 vector < T > minv(n+1, inf);
ce7
241
                 vector < int > used(n+1, 0);
016
                 do {
472
                      used[j0] = true;
                     int i0 = p[j0], j1 = -1;
d24
7e5
                     T delta = inf;
                     for (int j = 1; j \le n; j++) if
9ac
   (!used[j]) {
7bf
                          T cur = a[i0-1][j-1] - u[i0] - v[j];
                          if (cur < minv[j]) minv[j] = cur,</pre>
9f2
   way[j] = j0;
                          if (minv[j] < delta) delta =</pre>
821
   minv[j], j1 = j;
cbb
                      for (int j = 0; j \le n; j++)
f63
2c5
                          if (used[j]) u[p[j]] += delta, v[j]
   -= delta;
                          else minv[j] -= delta;
6ec
```

```
6 d 4
                      j0 = j1;
233
                 } while (p[j0] != 0);
016
                 do {
4 c 5
                      int j1 = way[j0];
0 d7
                      p[j0] = p[j1];
6 d 4
                      j0 = j1;
ca1
                 } while (j0);
cbb
306
             vector < int > ans(n);
             for (int j = 1; j <= n; j++) ans[p[j]-1] = j-1;</pre>
6 db
da3
             return make_pair(-v[0], ans);
cbb
214 };
```

# 3.2 Algoritmo MO - queries em caminhos de arvore

```
// Problema que resolve: https://www.spoj.com/problems/COT2/
//
// Complexidade sendo c = O(update) e SQ = sqrt(n):
// O((n + q) * sqrt(n) * c)
// 395329
1bc const int MAX = 40010, SQ = 400;
042 vector < int > g[MAX];
c54 namespace LCA { ... }
249 int in[MAX], out[MAX], vtx[2 * MAX];
81b bool on [MAX];
4c3 int dif, freq[MAX];
9e2 vector < int > w;
d9a void dfs(int v, int p, int &t) {
        vtx[t] = v, in[v] = t++;
659
18e
        for (int u : g[v]) if (u != p) {
c53
            dfs(u, v, t);
cbb
        vtx[t] = v, out[v] = t++;
217
cbb }
```

```
e5f void update(int p) { // faca alteracoes aqui
        int v = vtx[p];
bbc
0ec
        if (not on[v]) { // insere vtx v
            dif += (freq[w[v]] == 0);
31 c
b20
            freq[w[v]]++;
cbb
        }
4e6
        else { // retira o vertice v
            dif -= (freq[w[v]] == 1);
0a9
fd3
            freq[w[v]]--;
        }
cbb
        on[v] = not on[v];
73e
cbb }
a3a vector < tuple < int , int >> build_queries (const
   vector < pair < int , int > > & q) {
ea6
        LCA::build(0):
        vector<tuple<int, int, int>> ret;
f77
        for (auto [1, r] : q){
aa9
            if (in[r] < in[l]) swap(l, r);</pre>
d24
            int p = LCA::lca(1, r);
6f9
            int init = (p == 1) ? in[1] : out[1];
826
07a
            ret.emplace_back(init, in[r], in[p]);
cbb
edf
        return ret;
cbb }
f31 vector<int> mo_tree(const vector<pair<int, int>>& vq){
6bb
        int t = 0;
        dfs(0, -1, t);
dab
        auto q = build_queries(vq);
af1
d41
        vector<int> ord(q.size());
f48
        iota(ord.begin(), ord.end(), 0);
be8
d01
        sort(ord.begin(), ord.end(), [&] (int 1, int r) {
            int b1 = get<0>(q[1]) / SQ, br = <math>get<0>(q[r]) /
d8d
   SQ;
            if (bl != br) return bl < br;</pre>
596
            else if (bl % 2 == 1) return get<1>(q[1]) <</pre>
158
   get<1>(q[r]);
f1d
            else return get<1>(q[1]) > get<1>(q[r]);
сОс
        }):
```

```
80e
        memset(freq, 0, sizeof freq);
bf6
        dif = 0;
d41
ff2
        vector < int > ret(q.size());
3d9
        int 1 = 0, r = -1;
        for (int i : ord) {
8 b 0
            auto [ql, qr, qp] = q[i];
3 c7
            while (r < qr) update(++r);</pre>
af7
d6b
             while (1 > q1) update(--1);
            while (1 < q1) update(1++);</pre>
951
6a1
            while (r > qr) update(r--);
3 d8
            if (qp < 1 \text{ or } qp > r)  { // se LCA estah entre as
   pontas
74b
                 update(qp);
2 e 1
                 ret[i] = dif;
74b
                 update(qp);
            }
cbb
0fe
             else ret[i] = dif;
cbb
edf
        return ret;
cbb }
     Angle Range Intersection
// Computa intersecao de angulos
// Os angulos (arcos) precisam ter comprimeiro < pi
// (caso contrario a intersecao eh estranha)
//
// Tudo 0(1)
// 5e1c85
32a struct angle_range {
        static constexpr ld ALL = 1e9, NIL = -1e9;
75e
395
        ld 1, r;
```

angle\_range() : 1(ALL), r(ALL) {}

void fix(ld& theta) {

angle\_range(ld l\_, ld r\_) :  $l(l_), r(r_) \{ fix(l), \}$ 

if (theta == ALL or theta == NIL) return;

c77

894

4ee

da7

fix(r);

```
323
            if (theta > 2*pi) theta -= 2*pi;
868
            if (theta < 0) theta += 2*pi;
cbb
2ee
        bool empty() { return l == NIL; }
931
        bool contains(ld q) {
40f
            fix(q);
4d7
            if (1 == ALL) return true;
            if (1 == NIL) return false;
fec
6a6
           if (1 < r) return 1 < q and q < r;
075
            return q > 1 or q < r;
cbb
9c7
        friend angle_range operator &(angle_range p,
   angle_range q) {
743
            if (p.1 == ALL or q.1 == NIL) return q;
20f
            if (q.l == ALL or p.l == NIL) return p;
7 d 5
            if (p.l > p.r \text{ and } q.l > q.r) \text{ return } \{\max(p.l, q.r)\}
   q.1) , min(p.r, q.r)};
            if (q.1 > q.r) swap(p.1, q.1), swap(p.r, q.r);
aa6
            if (p.1 > p.r) {
8d8
249
                if (q.r > p.1) return \{max(q.1, p.1), q.r\};
                else if (q.1 < p.r) return {q.1, min(q.r,</pre>
6f7
   p.r)};
270
                return {NIL, NIL};
cbb
            if (max(p.1, q.1) > min(p.r, q.r)) return {NIL,
5a8
NIL };
bcb
            return {max(p.1, q.1), min(p.r, q.r)};
cbb
214 };
```

# 3.4 Area da Uniao de Retangulos

```
// O(n log(n))
// bea565
aa4 namespace seg {
6b3
        pair < int , ll > seg[4*MAX];
b1b
        11 lazy[4*MAX], *v;
1a8
        int n;
d41
e01
        pair<int, ll> merge(pair<int, ll> l, pair<int, ll>
  r){
```

```
719
            if (1.second == r.second) return
   {l.first+r.first, l.second};
             else if (1.second < r.second) return 1;</pre>
53b
             else return r;
aa0
        }
cbb
d41
6fc
        pair < int, ll > build(int p=1, int l=0, int r=n-1) {
3 c7
            lazv[p] = 0;
           if (1 == r) return seg[p] = {1, v[1]};
bf8
ee4
           int m = (1+r)/2;
            return seg[p] = merge(build(2*p, 1, m),
   build(2*p+1, m+1, r);
        }
cbb
        void build(int n2, l1* v2) {
d9e
             n = n2, v = v2;
680
6f2
             build();
cbb
ceb
        void prop(int p, int 1, int r) {
208
             seg[p].second += lazy[p];
2 c 9
            if (1 != r) lazy[2*p] += lazy[p], lazy[2*p+1] +=
   lazy[p];
3 c7
            lazy[p] = 0;
cbb
693
        pair<int, 11> query(int a, int b, int p=1, int 1=0,
int r=n-1) {
6 b 9
             prop(p, 1, r);
527
            if (a \le 1 \text{ and } r \le b) \text{ return seg}[p];
            if (b < 1 \text{ or } r < a) \text{ return } \{0, LINF\};
9 b 7
ee4
           int m = (1+r)/2;
           return merge(query(a, b, 2*p, 1, m), query(a, b,
   2*p+1, m+1, r);
cbb
07c
        pair <int, ll> update(int a, int b, int x, int p=1,
   int 1=0, int r=n-1) {
6 b 9
             prop(p, 1, r);
9 a 3
            if (a \le 1 \text{ and } r \le b) {
                 lazy[p] += x;
b94
6 b 9
                 prop(p, 1, r);
534
                 return seg[p];
cbb
e9f
            if (b < l or r < a) return seg[p];</pre>
ee4
            int m = (1+r)/2;
```

```
086
            return seg[p] = merge(update(a, b, x, 2*p, 1, m),
579
                    update(a, b, x, 2*p+1, m+1, r));
cbb
        }
214 };
d41
eb5 ll seg_vec[MAX];
8be 1l area_sq(vector<pair<int, int>, pair<int, int>>>
   &sq){
        vector<pair<int, int>, pair<int, int>>> up;
28 c
60a
        for (auto it : sq){
619
            int x1, y1, x2, y2;
            tie(x1, y1) = it.first;
ae0
68 e
            tie(x2, y2) = it.second;
            up.push_back({{x1+1, 1}, {y1, y2}});
80f
            up.push_back(\{\{x2+1, -1\}, \{y1, y2\}\}\});
aee
        }
cbb
092
        sort(up.begin(), up.end());
049
        memset(seg_vec, 0, sizeof seg_vec);
6fe
        11 H MAX = MAX;
156
        seg::build(H_MAX-1, seg_vec);
7ba
        auto it = up.begin();
04b
        11 \text{ ans} = 0;
f14
        while (it != up.end()){
07f
            ll L = (*it).first.first;
            while (it != up.end() && (*it).first.first == L){
718
127
                int x, inc, y1, y2;
d35
                tie(x, inc) = it->first;
                tie(y1, y2) = it->second;
d3d
                seg::update(y1+1, y2, inc);
5d1
40 d
                it++;
cbb
852
            if (it == up.end()) break;
            11 R = (*it).first.first;
d8a
d41
f59
            11 W = R-L;
efd
            auto jt = seg::query(0, H_MAX-1);
91a
            11 H = H_MAX - 1;
            if (jt.second == 0) H -= jt.first;
e8a
8df
            ans += W*H;
cbb
ba7
        return ans;
```

```
cbb }
```

#### 3.5 Area Maxima de Histograma

```
// Assume que todas as barras tem largura 1,
// e altura dada no vetor v
//
// O(n)
// e43846
15e ll area(vector<int> v) {
b73
        ll ret = 0;
4ce
        stack<int> s;
d41
        // valores iniciais pra dar tudo certo
447
        v.insert(v.begin(), -1);
d56
        v.insert(v.end(), -1);
        s.push(0);
1f8
        for(int i = 0; i < (int) v.size(); i++) {</pre>
0be
            while (v[s.top()] > v[i]) {
78e
265
                ll h = v[s.top()]; s.pop();
                ret = max(ret, h * (i - s.top() - 1));
de1
cbb
18e
            s.push(i);
cbb
d41
edf
        return ret;
cbb }
```

#### 3.6 Binomial modular

```
// Computa C(n, k) mod m em O(m + log(m) log(n))
// = O(rapido)
// ed4344

97c ll divi[MAX];

398 ll expo(ll a, ll b, ll m) {
1c1    if (!b) return 1;
399    ll ans = expo(a*a%m, b/2, m);
751    if (b%2) ans *= a;
```

```
2e9
        return ans%m:
cbb }
f0a ll inv(ll a, ll b){
        return 1 < a ? b - inv(b\%a,a)*b/a : 1;
bca
cbb }
153 template < typename T > tuple < T, T, T > ext_gcd(T a, T b) {
        if (!a) return {b, 0, 1};
3bd
        auto [g, x, y] = ext_gcd(b\%a, a);
550
        return \{g, y - b/a*x, x\};
c59
cbb }
bfe template < typename T = 11> struct crt {
627
        Ta, m;
        crt(): a(0), m(1) {}
5f3
        crt(T a_, T m_) : a(a_), m(m_) {}
7eb
        crt operator * (crt C) {
911
            auto [g, x, y] = ext_gcd(m, C.m);
238
           if ((a - C.a) \% g) a = -1;
dc0
           if (a == -1 or C.a == -1) return crt(-1, 0);
4f9
d09
            T lcm = m/g*C.m;
            T ans = a + (x*(C.a-a)/g \% (C.m/g))*m;
eb2
            return crt((ans % lcm + lcm) % lcm, lcm);
d8d
cbb
        }
214 };
6f2 pair<11, 11> divide_show(11 n, int p, int k, int pak) {
4f7
        if (n == 0) return \{0, 1\};
        11 blocos = n/pak, falta = n%pak;
d02
        ll periodo = divi[pak], resto = divi[falta];
2ce
616
        11 r = expo(periodo, blocos, pak)*resto%pak;
        auto rec = divide_show(n/p, p, k, pak);
445
a51
        11 y = n/p + rec.first;
        r = r*rec.second % pak;
bb9
        return {y, r};
90f
cbb }
6e6 ll solve_pak(ll n, ll x, int p, int k, int pak) {
```

```
d34
        divi[0] = 1:
        for (int i = 1; i <= pak; i++) {</pre>
f2b
            divi[i] = divi[i-1];
901
            if (i%p) divi[i] = divi[i] * i % pak;
840
        }
cbb
        auto dn = divide_show(n, p, k, pak), dx =
   divide_show(x, p, k, pak),
162
             dnx = divide\_show(n-x, p, k, pak);
768
        11 y = dn.first-dx.first-dnx.first, r =
b64
            (dn.second*inv(dx.second,
   pak)%pak)*inv(dnx.second, pak)%pak;
        return expo(p, y, pak) * r % pak;
035
cbb }
9dd ll solve(ll n, ll x, int mod) {
490
        vector < pair < int , int >> f;
        int mod2 = mod;
c3b
        for (int i = 2; i*i \le mod2; i++) if (mod2\%i = = 0) {
7 b 4
            int c = 0;
aff
75b
            while (mod2\%i==0) mod2 /= i, c++;
            f.push_back({i, c});
2a1
cbb
Off
        if (mod2 > 1) f.push_back({mod2, 1});
        crt ans(0, 1);
e96
        for (int i = 0; i < f.size(); i++) {</pre>
a13
702
            int pak = 1;
7 e 4
            for (int j = 0; j < f[i].second; j++) pak *=
   f[i].first;
304
            ans = ans * crt(solve_pak(n, x, f[i].first,
   f[i].second, pak), pak);
cbb
5fb
        return ans.a;
cbb }
3.7 Closest pair of points
```

```
// O(nlogn)
// f90265
915 pair <pt, pt > closest_pair_of_points(vector <pt > v) {
3d2
        int n = v.size();
```

```
sort(v.begin(), v.end());
fca
        for (int i = 1; i < n; i++) if (v[i] == v[i-1])
31 c
   return {v[i-1], v[i]};
        auto cmp_y = [&](const pt &1, const pt &r) {
c20
            if (1.y != r.y) return 1.y < r.y;</pre>
b53
920
             return l.x < r.x;</pre>
214
        };
62e
        set < pt, decltype(cmp_y) > s(cmp_y);
3d9
        int 1 = 0, r = -1;
6a2
        11 d2_min = numeric_limits < ll >:: max();
4d5
        pt pl, pr;
        const int magic = 5;
bd1
        while (r+1 < n) {
a55
7 f 1
             auto it = s.insert(v[++r]).first;
c92
            int cnt = magic/2;
773
             while (cnt-- and it != s.begin()) it--;
a 0 1
            cnt = 0:
             while (cnt++ < magic and it != s.end()) {</pre>
d68
                 if (!((*it) == v[r])) {
f19
67 e
                     11 d2 = dist2(*it, v[r]);
74e
                     if (d2 min > d2) {
229
                          d2_min = d2;
                          pl = *it;
841
4f2
                          pr = v[r];
                     }
cbb
cbb
                 }
40 d
                 it++;
cbb
            while (1 < r \text{ and } sq(v[1].x-v[r].x) > d2_min)
eb0
   s.erase(v[1++]);
cbb
c74
        return {pl, pr};
cbb }
```

#### 3.8 Coloração de Grafo de Intervalo

```
// Colore os intervalos com o numero minimo
// de cores de tal forma que dois intervalos
// que se interceptam tem cores diferentes
// As cores vao de 1 ate n
//
// O(n log(n))
```

```
// 83a32d
615 vector <int > coloring (vector <pair <int , int >>& v) {
        int n = v.size();
3d2
        vector<pair<int, pair<int, int>>> ev;
c08
603
        for (int i = 0; i < n; i++) {</pre>
            ev.push_back({v[i].first, {1, i}});
150
            ev.push_back({v[i].second, {0, i}});
cda
cbb
49e
        sort(ev.begin(), ev.end());
        vector < int > ans(n), avl(n);
360
265
        for (int i = 0; i < n; i++) avl.push_back(n-i);
        for (auto i : ev) {
4bf
            if (i.second.first == 1) {
che
021
                ans[i.second.second] = avl.back();
a 0 0
                avl.pop_back();
296
            } else avl.push_back(ans[i.second.second]);
cbb
ba7
        return ans;
cbb }
3.9 Conectividade Dinamica
// Offline com Divide and Conquer e
// DSU com rollback
// O(n log^2(n))
// 043d93
8f2 typedef pair <int, int > T;
1cd namespace data {
553
        int n, ans;
573
        int p[MAX], sz[MAX];
        stack<int> S;
ee6
        void build(int n2) {
e5b
1 e3
            n = n2;
8a6
            for (int i = 0; i < n; i++) p[i] = i, sz[i] = 1;
0b2
            ans = n;
cbb
        int find(int k) {
1 b 1
006
            while (p[k] != k) k = p[k];
```

```
839
            return k:
cbb
        }
072
        void add(T x) {
700
            int a = x.first, b = x.second;
605
            a = find(a), b = find(b);
843
            if (a == b) return S.push(-1);
            ans - - ;
e7d
3c6
            if (sz[a] > sz[b]) swap(a, b);
4c2
            S.push(a);
            sz[b] += sz[a];
582
84b
            p[a] = b;
        }
cbb
5eb
        int query() {
            return ans;
ba7
cbb
5cf
        void rollback() {
465
            int u = S.top(); S.pop();
            if (u == -1) return;
61 c
270
            sz[p[u]] -= sz[u];
            p[u] = u;
546
0df
            ans++;
cbb
        }
214 };
357 int ponta[MAX]; // outra ponta do intervalo ou -1 se for
   querv
4f0 int ans[MAX], n, q;
487 T qu[MAX];
47b void solve(int l = 0, int r = q-1) {
0 b 1
        if (1 >= r) {
8 c 0
            ans[1] = data::query(); // agora a estrutura ta
   certa
505
            return;
cbb
962
        int m = (1+r)/2, qnt = 1;
        for (int i = m+1; i <= r; i++) if (ponta[i]+1 and</pre>
fc7
   ponta[i] < 1)
37 d
            data::add(qu[i]), qnt++;
221
        solve(1, m);
        while (--qnt) data::rollback();
593
        for (int i = 1; i <= m; i++) if (ponta[i]+1 and</pre>
a2c
```

#### 3.10 Conectividade Dinamica 2

```
// Offline com link-cut trees
// O(n log(n))
// d38e4e
1ef namespace lct {
3c9
        struct node {
19f
            int p, ch[2];
a2a
            int val, sub;
aa6
           bool rev;
f93
            node() {}
            node(int v) : p(-1), val(v), sub(v), rev(0) 
   ch[0] = ch[1] = -1; }
214
       };
d41
c53
        node t[2*MAX]; // MAXN + MAXQ
99e
        map < pair < int , int > , int > aresta;
e4d
        int sz;
d41
95a
        void prop(int x) {
aa2
            if (t[x].rev) {
f95
                swap(t[x].ch[0], t[x].ch[1]);
379
                if (t[x].ch[0]+1) t[t[x].ch[0]].rev ^= 1;
c3d
                if (t[x].ch[1]+1) t[t[x].ch[1]].rev ^= 1;
cbb
693
            t[x].rev = 0;
cbb
564
        void update(int x) {
            t[x].sub = t[x].val;
e8d
            for (int i = 0; i < 2; i++) if (t[x].ch[i]+1) {
8ca
621
                prop(t[x].ch[i]);
78 d
                t[x].sub = min(t[x].sub, t[t[x].ch[i]].sub);
cbb
            }
cbb
971
        bool is root(int x) {
```

```
657
            return t[x].p == -1 or (t[t[x].p].ch[0] != x and
                                                                  cbb
   t[t[x].p].ch[1] != x);
                                                                  204
                                                                          void link (int v, int w) {
                                                                  821
cbb
                                                                              rootify(w);
        void rotate(int x) {
                                                                  389
                                                                              t[w].p = v;
ed6
497
            int p = t[x].p, pp = t[p].p;
                                                                  cbb
            if (!is_root(p)) t[pp].ch[t[pp].ch[1] == p] = x;
fc4
                                                                  6 b 8
                                                                          void link(int v, int w, int x) { // v--w com peso x
251
            bool d = t[p].ch[0] == x;
                                                                  379
                                                                               int id = MAX + sz++;
            t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
                                                                              aresta[make_pair(v, w)] = id;
461
                                                                  110
            if (t[p].ch[!d]+1) t[t[p].ch[!d]].p = p;
a76
                                                                  ab6
                                                                              make_tree(id, x);
            t[x].p = pp, t[p].p = x;
8fa
                                                                  c88
                                                                              link_(v, id), link_(id, w);
            update(p), update(x);
444
                                                                  cbb
cbb
        }
                                                                  e63
                                                                          void cut_(int v, int w) {
238
                                                                              rootify(w), access(v);
        int splay(int x) {
                                                                  b54
            while (!is_root(x)) {
                                                                  264
                                                                              t[v].ch[0] = t[t[v].ch[0]].p = -1;
18 c
                int p = t[x].p, pp = t[p].p;
497
                                                                  cbb
                if (!is_root(p)) prop(pp);
77b
                                                                  031
                                                                          void cut(int v, int w) {
                                                                              int id = aresta[make_pair(v, w)];
be5
                prop(p), prop(x);
                                                                  b0f
                if (!is_root(p)) rotate((t[pp].ch[0] ==
                                                                  a4a
                                                                               cut_(v, id), cut_(id, w);
0c5
  p)^{(t[p].ch[0]} == x) ? x : p);
                                                                  cbb
                rotate(x);
                                                                  cbb }
64f
            }
cbb
aab
            return prop(x), x;
                                                                  893 void dyn_conn() {
cbb
                                                                  c5f
                                                                          int n, q; cin >> n >> q;
f16
        int access(int v) {
                                                                  d6e
                                                                          vector <int > p(2*q, -1); // outra ponta do intervalo
            int last = -1;
                                                                  b4f
                                                                          for (int i = 0; i < n; i++) lct::make_tree(i);</pre>
0eb
d9f
            for (int w = v; w+1; update(last = w), splay(v),
                                                                          vector < pair < int , int >> qu(q);
                                                                  fbf
   w = t[v].p)
                                                                  139
                                                                          map<pair<int, int>, int> m;
                splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
                                                                          for (int i = 0; i < q; i++) {</pre>
024
                                                                  abf
3d3
            return last;
                                                                  3 c2
                                                                              char c; cin >> c;
cbb
        }
                                                                  ef6
                                                                              if (c == '?') continue;
952
        void make_tree(int v, int w=INF) { t[v] = node(w); }
                                                                  602
                                                                              int a, b; cin >> a >> b; a--, b--;
        bool conn(int v, int w) {
                                                                              if (a > b) swap(a, b);
82f
                                                                  d11
2cf
            access(v), access(w);
                                                                  8a1
                                                                              qu[i] = \{a, b\};
                                                                  8 d7
                                                                              if (c == '+') {
b9b
            return v == w ? true : t[v].p != -1;
cbb
                                                                  94b
                                                                                  p[i] = i+q, p[i+q] = i;
277
        void rootify(int v) {
                                                                  906
                                                                                   m[make_pair(a, b)] = i;
                                                                  9 d 9
                                                                              } else {
5e3
            access(v);
a02
            t[v].rev ^= 1;
                                                                  412
                                                                                   int j = m[make_pair(a, b)];
                                                                  ac2
                                                                                  p[i] = j, p[j] = i;
cbb
        int query(int v, int w) {
                                                                  cbb
                                                                              }
a1d
            rootify(w), access(v);
                                                                  cbb
b54
            return t[v].sub;
249
                                                                  447
                                                                          int ans = n;
```

```
for (int i = 0; i < q; i++) {</pre>
abf
             if (p[i] == -1) {
87 d
886
                 cout << ans << endl; // numero de comp</pre>
   conexos
5e2
                 continue;
cbb
69 d
            int a = qu[i].first, b = qu[i].second;
c4d
            if (p[i] > i) { // +
ac5
                 if (lct::conn(a, b)) {
                     int mi = lct::query(a, b);
18 f
                     if (p[i] < mi) {</pre>
993
                         p[p[i]] = p[i];
dd3
5e2
                          continue;
cbb
6f7
                     lct::cut(qu[p[mi]].first,
   qu[p[mi]].second), ans++;
                     p[mi] = mi;
6ea
cbb
                 lct::link(a, b, p[i]), ans--;
d1d
            } else if (p[i] != i) lct::cut(a, b), ans++; // -
cb5
cbb
cbb }
```

#### Conj. Indep. Maximo com Peso em Grafo de 3.11Intervalo

```
// Retorna os indices ordenados dos intervalos selecionados
// Se tiver empate, retorna o que minimiza o comprimento
   total
//
// O(n log(n))
// c4dbe2
31e vector<int> ind_set(vector<tuple<int, int, int>>& v) {
b27
        vector<tuple<int, int, int>> w;
f14
        for (int i = 0; i < v.size(); i++) {</pre>
            w.push_back(tuple(get<0>(v[i]), 0, i));
e85
6f0
            w.push_back(tuple(get<1>(v[i]), 1, i));
cbb
d1d
        sort(w.begin(), w.end());
```

```
844
        vector < int > nxt(v.size());
c22
        vector <pair < ll, int >> dp(v.size());
0eb
        int last = -1;
        for (auto [fim, t, i] : w) {
723
25a
             if (t == 0) {
4ca
                 nxt[i] = last;
5e2
                 continue:
cbb
78b
             dp[i] = \{0, 0\};
            if (last != -1) dp[i] = max(dp[i], dp[last]);
cb8
            pair < 11, int > pega = {get < 2 > (v[i]),
911
   -(get<1>(v[i]) - get<0>(v[i]) + 1);
            if (nxt[i] != -1) pega.first +=
5d3
   dp[nxt[i]].first, pega.second += dp[nxt[i]].second;
            if (pega > dp[i]) dp[i] = pega;
b08
7 cb
             else nxt[i] = last;
381
            last = i:
cbb
977
        pair < ll, int > ans = \{0, 0\};
        int idx = -1;
919
ceb
        for (int i = 0; i < v.size(); i++) if (dp[i] > ans)
   ans = dp[i], idx = i;
4b8
        vector < int > ret;
        while (idx != -1) {
fdd
d69
            if (get < 2 > (v[idx]) > 0 and
a05
                 (nxt[idx] == -1 \text{ or } get<1>(v[nxt[idx]]) <
   get <0 > (v[idx]))) ret.push_back(idx);
e4f
             idx = nxt[idx];
cbb
        sort(ret.begin(), ret.end());
0ea
edf
        return ret;
cbb }
```

#### Distancia maxima entre dois pontos

```
// \max_{dist2(v)} - O(n \log(n))
// max dist manhattan - O(n)
// Quadrado da Distancia Euclidiana (precisa copiar
   convex_hull, ccw e pt)
// bdace4
859 ll max dist2(vector<pt> v) {
```

```
221
        v = convex hull(v):
        if (v.size() <= 2) return dist2(v[0], v[1%v.size()]);</pre>
a14
04b
        11 \text{ ans} = 0;
        int n = v.size(), j = 0;
323
603
        for (int i = 0; i < n; i++) {</pre>
            while (!ccw(v[(i+1)%n]-v[i], pt(0, 0),
057
   v[(j+1)%n]-v[j])) j = (j+1)%n;
            ans = max(\{ans, dist2(v[i], v[j]),
   dist2(v[(i+1)%n], v[j])});
cbb
ba7
        return ans;
cbb }
    // Distancia de Manhattan
c51 template < typename T > T max_dist_manhattan (vector < pair < T,
   T >> v)
        T min_sum, max_sum, min_dif, max_dif;
8eb
        min_sum = max_sum = v[0].first + v[0].second;
4f5
        min_dif = max_dif = v[0].first - v[0].second;
271
        for (auto [x, y] : v) {
c25
            min_sum = min(min_sum, x+y);
1cb
683
            max_sum = max(max_sum, x+y);
782
            min_dif = min(min_dif, x-y);
            max_dif = max(max_dif, x-y);
af7
cbb
        return max(max_sum - min_sum, max_dif - min_dif);
9f0
cbb }
      Distinct Range Query
3.13
```

```
// build - O(n (log n + log(sigma)))
// query - O(log(sigma))
// 5c7aa1

789 namespace perseg { };

53d int qt[MAX];
edc void build(vector<int>& v) {
3d2    int n = v.size();
16b    perseg::build(n);
```

```
663
        map < int , int > last;
05е
        int at = 0;
        for (int i = 0; i < n; i++) {</pre>
603
            if (last.count(v[i])) {
817
                 perseg::update(last[v[i]], -1);
a58
69a
                 at++;
            }
cbb
4f2
            perseg::update(i, 1);
            qt[i] = ++at;
460
efe
            last[v[i]] = i;
cbb
cbb }
9e3 int query(int 1, int r) {
        return perseg::query(1, r, qt[r]);
cbb }
```

#### 3.14 Distinct Range Query com Update

```
// build - O(n log(n))
// query - O(log^2(n))
// update - 0(log^2(n))
// 2306f3
774 #include <ext/pb_ds/assoc_container.hpp>
30f #include <ext/pb_ds/tree_policy.hpp>
0d7 using namespace __gnu_pbds;
4fc template <class T>
        using ord_set = tree<T, null_type, less<T>,
def
   rb_tree_tag,
3a1
        tree_order_statistics_node_update>;
042 int v[MAX], n, nxt[MAX], prv[MAX];
f60 map < int , set < int > > ocor;
e04 namespace bit {
686
        ord_set<pair<int, int>> bit[MAX];
0a8
        void build() {
3 e 1
             for (int i = 1; i <= n; i++)</pre>
   bit[i].insert({nxt[i-1], i-1});
78a
            for (int i = 1; i <= n; i++) {</pre>
```

```
int i = i + (i\&-i):
edf
                if (j <= n) for (auto x : bit[i])</pre>
d03
   bit[j].insert(x);
            }
cbb
cbb
        int pref(int p, int x) {
d3f
            int ret = 0;
7 c 9
            for (; p; p -= p&-p) ret +=
bbf
   bit[p].order_of_key({x, -INF});
edf
            return ret;
cbb
        int query(int 1, int r, int x) {
d50
            return pref(r+1, x) - pref(1, x);
e55
cbb
ff3
        void update(int p, int x) {
f17
            int p2 = p;
5ed
            for (p++; p \le n; p += p\&-p) {
                bit[p].erase({nxt[p2], p2});
ca8
                bit[p].insert({x, p2});
f6b
cbb
        }
cbb
cbb }
0a8 void build() {
        for (int i = 0; i < n; i++) nxt[i] = INF;</pre>
383
        for (int i = 0; i < n; i++) prv[i] = -INF;
7b3
d07
        vector<pair<int, int>> t;
        for (int i = 0; i < n; i++) t.push_back({v[i], i});</pre>
348
3fd
        sort(t.begin(), t.end());
603
        for (int i = 0; i < n; i++) {
b40
            if (i and t[i].first == t[i-1].first)
                prv[t[i].second] = t[i-1].second;
565
a8b
            if (i+1 < n and t[i].first == t[i+1].first)</pre>
                 nxt[t[i].second] = t[i+1].second:
12f
cbb
        }
        for (int i = 0; i < n; i++) ocor[v[i]].insert(i);</pre>
a23
        bit::build();
1 d7
cbb }
aae void muda(int p, int x) {
```

```
f92
        bit::update(p, x);
        nxt[p] = x;
c3d
cbb }
4ea int query(int a, int b) {
a0a
        return b-a+1 - bit::query(a, b, b+1);
cbb }
ff3 void update(int p, int x) { // mudar valor na pos. p
   para x
c0b
        if (prv[p] > -INF) muda(prv[p], nxt[p]);
        if (nxt[p] < INF) prv[nxt[p]] = prv[p];</pre>
        ocor[v[p]].erase(p);
5bf
        if (!ocor[x].size()) {
4 b 4
            muda(p, INF);
19 d
            prv[p] = -INF;
8 d 4
        } else if (*ocor[x].rbegin() < p) {</pre>
a69
            int i = *ocor[x].rbegin();
5 b 5
f64
            prv[p] = i;
            muda(p, INF);
19 d
            muda(i, p);
5f2
9d9
        } else {
d46
            int i = *ocor[x].lower_bound(p);
33f
            if (prv[i] > -INF) {
                muda(prv[i], p);
f 17
                prv[p] = prv[i];
8f9
            } else prv[p] = -INF;
94f
523
            prv[i] = p;
597
            muda(p, i);
cbb
c96
        v[p] = x; ocor[x].insert(p);
cbb }
```

#### 3.15 Dominator Points

```
// Se um ponto A tem ambas as coordenadas >= B, dizemos
// que A domina B
// is_dominated(p) fala se existe algum ponto no conjunto
// que domina p
// insert(p) insere p no conjunto
// (se p for dominado por alguem, nao vai inserir)
```

```
// o multiset 'quina' guarda informacao sobre os pontos
// nao dominados por um elemento do conjunto que nao dominam
// outro ponto nao dominado por um elemento do conjunto
// No caso, armazena os valores de x+y esses pontos
//
// Complexidades:
// is dominated - O(\log(n))
// insert - O(log(n)) amortizado
// query - 0(1)
// 09ffdc
e2a struct dominator_points {
        set<pair<int, int>> se;
baf
4dd
        multiset < int > quina;
a85
        bool is_dominated(pair<int, int> p) {
            auto it = se.lower_bound(p);
80f
            if (it == se.end()) return 0;
633
            return it->second >= p.second;
ab4
cbb
99b
        void mid(pair<int, int> a, pair<int, int> b, bool
   rem) {
29a
            pair < int , int > m = {a.first+1, b.second+1};
            int val = m.first + m.second;
b19
            if (!rem) quina.insert(val);
638
731
            else quina.erase(quina.find(val));
cbb
7 c 4
        bool insert(pair<int, int> p) {
            if (is dominated(p)) return 0;
fb4
80f
            auto it = se.lower_bound(p);
            if (it != se.begin() and it != se.end())
ca9
                mid(*prev(it), *it, 1);
d4a
            while (it != se.begin()) {
1fa
                it--:
049
23 c
                if (it->second > p.second) break;
b86
                if (it != se.begin()) mid(*prev(it), *it, 1);
                it = se.erase(it);
316
            }
cbb
            it = se.insert(p).first;
433
            if (it != se.begin()) mid(*prev(it), *it, 0);
69 e
            if (next(it) != se.end()) mid(*it, *next(it), 0);
96 d
6a5
            return 1:
```

```
cbb  }
5eb    int query() {
956        if (!quina.size()) return INF;
add        return *quina.begin();
cbb  }
214 };
```

#### 3.16 DP de Dominacao 3D

```
// Computa para todo ponto i,
// dp[i] = 1 + max_{j dominado por i} dp[j]
// em que ser dominado eh ter as 3 coordenadas menores
// Da pra adaptar facil para outras dps
//
// O(n log<sup>2</sup> n), O(n) de memoria
// 7c8896
c53 void lis2d(vector < vector < tuple < int, int, int >>> & v,
   vector < int > & dp, int 1, int r) {
        if (1 == r) {
893
56f
             for (int i = 0; i < v[1].size(); i++) {</pre>
8 b 5
                 int ii = get < 2 > (v[1][i]);
1 ce
                 dp[ii] = max(dp[ii], 1);
cbb
             }
505
             return:
cbb
ee4
        int m = (1+r)/2;
62b
        lis2d(v, dp, 1, m);
325
        vector < tuple < int , int , int >> vv[2];
d44
        vector < int > Z:
871
        for (int i = 1; i \le r; i++) for (auto it : v[i]) {
2ef
             vv[i > m].push_back(it);
             Z.push_back(get<1>(it));
042
cbb
        sort(vv[0].begin(), vv[0].end());
e9f
        sort(vv[1].begin(), vv[1].end());
9 b 5
0 d 1
        sort(Z.begin(), Z.end());
573
        auto get_z = [&](int z) { return
   lower_bound(Z.begin(), Z.end(), z) - Z.begin(); };
         vector < int > bit(Z.size());
c51
```

```
181
        int i = 0:
        for (auto [y, z, id] : vv[1]) {
e9a
             while (i < vv[0].size() and get<0>(vv[0][i]) <</pre>
6bd
   y) {
                 auto [y2, z2, id2] = vv[0][i++];
397
                 for (int p = get z(z_2)+1; p <= Z.size(); p
ea0
   += p\&-p
300
                     bit[p-1] = max(bit[p-1], dp[id2]);
cbb
d3b
            int q = 0;
             for (int p = get_z(z); p; p -= p\&-p) q = max(q,
   bit[p-1]);
            dp[id] = max(dp[id], q + 1);
614
cbb
c25
        lis2d(v, dp, m+1, r);
cbb }
4de vector < int > solve (vector < tuple < int , int , int >> v) {
3d2
        int n = v.size();
        vector<tuple<int, int, int, int>> vv;
cd4
        for (int i = 0; i < n; i++) {</pre>
603
             auto [x, y, z] = v[i];
9be
5bb
             vv.emplace_back(x, y, z, i);
cbb
bd3
        sort(vv.begin(), vv.end());
e11
        vector < vector < tuple < int , int , int >>> V;
603
        for (int i = 0; i < n; i++) {</pre>
            int j = i;
a5b
808
            V.emplace_back();
c 0 1
             while (j < n \text{ and } get < 0 > (vv[j]) == get < 0 > (vv[i]))
                 auto [x, y, z, id] = vv[j++];
ba6
                 V.back().emplace_back(y, z, id);
cbb
            }
cbb
452
             i = j-1;
cbb
388
        vector<int> dp(n);
839
        lis2d(V, dp, 0, V.size()-1);
        return dp;
898
cbb }
```

#### 3.17 Gray Code

```
// Gera uma permutacao de 0 a 2^n-1, de forma que
// duas posicoes adjacentes diferem em exatamente 1 bit
//
// 0(2^n)
// 840df4
df6 vector < int > gray code(int n) {
73f
        vector < int > ret(1 << n);</pre>
        for (int i = 0; i < (1 << n); i++) ret[i] = i^{(i>>1)};
f29
edf
        return ret;
cbb }
3.18 Half-plane intersection
// Cada half-plane eh identificado por uma reta e a regiao
   ccw a ela
//
// O(n log n)
// f56e1c
f4f vector <pt> hp_intersection(vector <line> &v) {
9bc
        deque <pt> dq = {{INF, INF}, {-INF, INF}, {-INF,
   -INF}, {INF, -INF}};
d41 #warning considerar trocar por compare_angle
        sort(v.begin(), v.end(), [&](line r, line s) {
de3
   return angle(r.q-r.p) < angle(s.q-s.p); });</pre>
5e9
        for(int i = 0; i < v.size() and dq.size() > 1; i++) {
c69
            pt p1 = dq.front(), p2 = dq.back();
            while (dq.size() and !ccw(v[i].p, v[i].q,
6c6
   dq.back()))
47b
                p1 = dq.back(), dq.pop_back();
            while (dq.size() and !ccw(v[i].p, v[i].q,
0a2
   dq.front()))
                p2 = dq.front(), dq.pop_front();
7cf
            if (!dq.size()) break;
4d9
            if (p1 == dq.front() and p2 == dq.back())
```

continue:

```
c9b
            dq.push_back(inter(v[i], line(dq.back(), p1)));
65 c
            dq.push_front(inter(v[i], line(dq.front(), p2)));
            if (dq.size() > 1 and dq.back() == dq.front())
fdd
   dq.pop_back();
cbb
        return vector <pt>(dq.begin(), dq.end());
b2b
cbb }
3.19 Heap Sort
// O(n log n)
// 385e91
f18 void down(vector<int>& v, int n, int i) {
        while ((i = 2*i+1) < n) {
e1f
            if (i+1 < n and v[i] < v[i+1]) i++;</pre>
583
b27
            if (v[i] < v[(i-1)/2]) break;
322
            swap(v[i], v[(i-1)/2]);
cbb
cbb }
eb6 void heap_sort(vector<int>& v) {
3d2
        int n = v.size();
61 d
        for (int i = n/2-1; i \ge 0; i--) down(v, n, i);
        for (int i = n-1; i > 0; i--)
917
            swap(v[0], v[i]), down(v, i, 0);
37 f
cbb }
      Inversion Count
3.20
// Computa o numero de inversoes para transformar
// l em r (se nao tem como, retorna -1)
// O(n log(n))
// eef01f
37b template < typename T > 1l inv_count(vector < T > 1, vector < T >
```

 $r = \{\}\}$ 

bb6

796

1bc

if (!r.size()) {

sort(r.begin(), r.end());

r = 1;

```
cbb
874
        int n = 1.size();
8 c 0
        vector < int > v(n), bit(n);
        vector < pair < T, int >> w;
4e9
61c
        for (int i = 0; i < n; i++) w.push_back({r[i], i+1});
d1d
        sort(w.begin(), w.end());
603
        for (int i = 0; i < n; i++) {</pre>
             auto it = lower_bound(w.begin(), w.end(),
bf3
   make_pair(l[i], 0));
             if (it == w.end() or it->first != l[i]) return
1bf
   -1: // nao da
962
             v[i] = it->second;
            it->second = -1;
cbb
04b
        11 \text{ ans} = 0:
45b
        for (int i = n-1; i \ge 0; i--) {
             for (int j = v[i]-1; j; j -= j\&-j) ans += bit[j];
2 d 9
             for (int j = v[i]; j < n; j += j&-j) bit[j]++;
3 a 1
cbb
ba7
        return ans;
cbb }
```

### 3.21 LIS - Longest Increasing Subsequence

```
// Calcula e retorna uma LIS
//
// O(n.log(n))
// 4749e8
121 template < typename T> vector < T> lis(vector < T>& v) {
1fa
        int n = v.size(), m = -1;
f0c
        vector <T> d(n+1, INF);
aec
        vector < int > l(n);
        d[0] = -INF;
007
603
        for (int i = 0; i < n; i++) {
d41
            // Para non-decreasing use upper_bound()
            int t = lower_bound(d.begin(), d.end(), v[i]) -
4 f d
   d.begin();
            d[t] = v[i], l[i] = t, m = max(m, t);
3ad
cbb
```

```
4ff
        int p = n;
5a9
        vector < T > ret;
        while (p--) if (1[p] == m) {
cdf
            ret.push_back(v[p]);
883
76b
            m --;
cbb
969
        reverse(ret.begin(),ret.end());
edf
        return ret;
cbb }
```

#### 3.22 LIS2 - Longest Increasing Subsequence

```
// Calcula o tamanho da LIS
// O(n log(n))
// 402def
84b template < typename T > int lis(vector < T > &v) {
        vector < T > ans;
2da
        for (T t : v){
5e0
           // Para non-decreasing use upper_bound()
d41
           auto it = lower_bound(ans.begin(), ans.end(), t);
fe6
           if (it == ans.end()) ans.push_back(t);
d7f
            else *it = t;
b94
cbb
        return ans.size();
1eb
cbb }
```

#### 3.23 Mininum Enclosing Circle

```
// O(n) com alta probabilidade
// b0a6ba

22c const double EPS = 1e-12;
878 mt19937 rng((int)
    chrono::steady_clock::now().time_since_epoch().count());

b2a struct pt {
662    double x, y;
```

```
be7
       pt(double x_ = 0, double y_ = 0) : x(x_), y(y_) {}
       pt operator + (const pt& p) const { return pt(x+p.x,
7af
b23
        pt operator - (const pt& p) const { return pt(x-p.x,
  y-p.y); }
254
       pt operator * (double c) const { return pt(x*c,
   v*c); }
       pt operator / (double c) const { return pt(x/c,
   v/c): }
214 };
2f9 double dot(pt p, pt q) { return p.x*q.x+p.y*q.y; }
dd5 double cross(pt p, pt q) { return p.x*q.y-p.y*q.x; }
e7c double dist(pt p, pt q) { return sqrt(dot(p-q, p-q)); }
3f4 pt center(pt p, pt q, pt r) {
       pt a = p-r, b = q-r;
5d9
       pt c = pt(dot(a, p+r)/2, dot(b, q+r)/2);
e84
       return pt(cross(c, pt(a.y, b.y)), cross(pt(a.x,
   b.x), c)) / cross(a, b);
cbb }
aa8 struct circle {
f41
        pt cen;
c12
        double r;
       circle(pt cen_, double r_) : cen(cen_), r(r_) {}
898
       circle(pt a, pt b, pt c) {
83c
13d
            cen = center(a, b, c);
1 f 1
           r = dist(cen, a);
cbb
        bool inside(pt p) { return dist(p, cen) < r+EPS; }</pre>
cd5
214 }:
806 circle minCirc(vector <pt> v) {
        shuffle(v.begin(), v.end(), rng);
f 2.1
        circle ret = circle(pt(0, 0), 0);
ae0
        for (int i = 0; i < v.size(); i++) if</pre>
618
   (!ret.inside(v[i])) {
          ret = circle(v[i], 0);
16a
           for (int j = 0; j < i; j++) if
f 1 1
   (!ret.inside(v[i])) {
                ret = circle((v[i]+v[j])/2, dist(v[i],
881
```

#### 3.24 Minkowski Sum

```
// Computa A+B = \{a+b : a \setminus in A, b \setminus in B\}, em que
// A e B sao poligonos convexos
// A+B eh um poligono convexo com no max |A|+|B| pontos
// O(|A|+|B|)
// d7cca8
539 vector<pt> minkowski(vector<pt> p, vector<pt> q) {
        auto fix = [](vector < pt > & P) {
051
515
             rotate(P.begin(), min_element(P.begin(),
   P.end()), P.end());
018
            P.push_back(P[0]), P.push_back(P[1]);
214
        };
889
        fix(p), fix(q);
8af
        vector<pt> ret;
692
        int i = 0, j = 0;
        while (i < p.size()-2 \text{ or } j < q.size()-2) {
2ee
            ret.push_back(p[i] + q[j]);
898
            auto c = ((p[i+1] - p[i]) ^ (q[j+1] - q[j]));
732
ebc
            if (c \ge 0) i = min<int>(i+1, p.size()-2);
81e
            if (c \le 0) j = min \le int > (j+1, q.size()-2);
        }
cbb
edf
        return ret;
cbb }
    // 2f5dd2
c3e ld dist_convex(vector<pt> p, vector<pt> q) {
        for (pt& i : p) i = i * -1;
dc2
44 c
        auto s = minkowski(p, q);
        if (inpol(s, pt(0, 0))) return 0;
95 d
6a5
        return 1;
```

```
921
        ld ans = DINF;
        for (int i = 0; i < s.size(); i++) ans = min(ans,</pre>
073
f04
                disttoseg(pt(0, 0), line(s[(i+1)%s.size()],
   s[i])));
ba7
        return ans;
cbb }
3.25 MO - DSU
// Dado uma lista de arestas de um grafo, responde
// para cada query(1, r), quantos componentes conexos
// o grafo tem se soh considerar as arestas 1, 1+1, ..., r
// Da pra adaptar pra usar MO com qualquer estrutura
   rollbackavel
//
// O(m \ sqrt(q) \log(n))
// f98540
8d3 struct dsu {
553
        int n, ans;
2 e 3
        vector < int > p, sz;
ee6
        stack < int > S;
4b8
        dsu(int n_{-}) : n(n_{-}), ans(n), p(n), sz(n) {
8a6
            for (int i = 0; i < n; i++) p[i] = i, sz[i] = 1;
cbb
1 b 1
        int find(int k) {
006
            while (p[k] != k) k = p[k];
839
            return k;
cbb
553
        void add(pair<int, int> x) {
700
            int a = x.first, b = x.second;
605
            a = find(a), b = find(b);
843
            if (a == b) return S.push(-1);
e7d
3 c 6
            if (sz[a] > sz[b]) swap(a, b);
4c2
            S.push(a);
            sz[b] += sz[a];
582
84b
            p[a] = b;
cbb
        int query() { return ans; }
35 c
5cf
        void rollback() {
```

```
465
            int u = S.top(); S.pop();
            if (u == -1) return;
61 c
            sz[p[u]] -= sz[u];
270
546
            p[u] = u;
0df
            ans++;
cbb
        }
214 };
1a8 int n;
e93 vector<pair<int, int>> ar; // vetor com as arestas
617 vector<int> MO(vector<pair<int, int>> &q) {
        int SQ = ar.size() / sqrt(q.size()) + 1;
d4d
        int m = q.size();
c23
        vector<int> ord(m):
3f8
be8
        iota(ord.begin(), ord.end(), 0);
        sort(ord.begin(), ord.end(), [&](int 1, int r) {
d01
            if (q[1].first / SQ != q[r].first / SQ) return
9c9
   q[1].first < q[r].first;</pre>
            return q[1].second < q[r].second;</pre>
a66
сОс
        });
435
        vector < int > ret(m);
        for (int i = 0; i < m; i++) {</pre>
dd5
            dsu D(n);
176
            int fim = q[ord[i]].first/SQ*SQ + SQ - 1;
ae9
e25
            int last r = fim;
ebc
            int j = i-1;
            while (j+1 < m and q[ord[j+1]].first / SQ ==</pre>
00 c
   q[ord[i]].first / SQ) {
                auto [1, r] = q[ord[++j]];
a0e
                if (1 / SQ == r / SQ) {
acc
                     dsu D2(n):
ce9
495
                     for (int k = 1; k <= r; k++)
   D2.add(ar[k]);
                     ret[ord[j]] = D2.query();
fdf
                     continue;
5e2
                }
cbb
                while (last r < r) D.add(ar[++last r]);
59b
2cf
                for (int k = 1; k \le fim; k++) D.add(ar[k]);
```

#### 3.26 Mo - numero de distintos em range

```
// Para ter o bound abaixo, escolher
// SQ = n / sqrt(q)
//
// O(n * sqrt(q))
// e94f60
0d2 const int MAX = 1e5+10;
6ff const int SQ = sqrt(MAX);
b69 int v[MAX];
b65 int ans, freq[MAX];
9da inline void insert(int p) {
ae0
        int o = v[p];
591
        freq[o]++;
992
        ans += (freq[o] == 1);
cbb }
a25 inline void erase(int p) {
ae0
        int o = v[p];
7ee
        ans -= (freq[o] == 1);
        freq[o]--;
ba2
cbb }
e51 inline ll hilbert(int x, int y) {
        static int N = 1 << (__builtin_clz(0) -</pre>
   __builtin_clz(MAX));
100
        int rx, ry, s;
b72
        11 d = 0;
43b
        for (s = N/2; s > 0; s /= 2) {
```

```
c95
            rx = (x \& s) > 0, rv = (v \& s) > 0;
            d += s * ll(s) * ((3 * rx) ^ ry);
e3e
            if (ry == 0) {
d2e
               if (rx == 1) x = N-1 - x, y = N-1 - y;
5aa
                 swap(x, y);
9dd
cbb
            }
cbb
        }
be2
        return d;
cbb }
bac #define HILBERT true
617 vector<int> MO(vector<pair<int, int>> &q) {
c3b
        ans = 0;
       int m = q.size();
c23
        vector < int > ord(m);
3f8
        iota(ord.begin(), ord.end(), 0);
be8
6a6 #if HILBERT
        vector < ll> h(m);
8 c 4
       for (int i = 0; i < m; i++) h[i] =
   hilbert(q[i].first, q[i].second);
        sort(ord.begin(), ord.end(), [&](int 1, int r) {
   return h[1] < h[r]; });</pre>
8c1 #else
        sort(ord.begin(), ord.end(), [&](int 1, int r) {
d01
            if (q[1].first / SQ != q[r].first / SQ) return
9c9
   q[1].first < q[r].first;</pre>
            if ((q[1].first / SQ) % 2) return q[1].second >
0db
            return q[1].second < q[r].second;</pre>
a66
сОс
        }):
f2e #endif
435
        vector < int > ret(m);
3d9
        int 1 = 0, r = -1;
8b0
        for (int i : ord) {
6c6
            int q1, qr;
            tie(ql, qr) = q[i];
4f5
026
            while (r < qr) insert(++r);</pre>
232
            while (1 > q1) insert(--1);
            while (1 < q1) erase(1++);</pre>
75 e
            while (r > qr) erase(r--);
fe8
            ret[i] = ans;
381
```

```
cbb }
edf return ret;
cbb }
```

#### 3.27 Palindromic Factorization

```
// Precisa da eertree
// Computa o numero de formas de particionar cada
// prefixo da string em strings palindromicas
//
// O(n log n), considerando alfabeto O(1)
// 9e6e22
070 struct eertree { ... };
0e7 ll factorization(string s) {
        int n = s.size(), sz = 2;
b19
        eertree PT(n);
580
        vector < int > diff(n+2), slink(n+2), sans(n+2),
   dp(n+1);
0ec
        dp[0] = 1;
78a
        for (int i = 1; i <= n; i++) {
c58
            PT.add(s[i-1]);
           if (PT.size()+2 > sz) {
a7c
6 c 4
                diff[sz] = PT.len[sz] - PT.len[PT.link[sz]];
241
                if (diff[sz] == diff[PT.link[sz]])
d6f
                    slink[sz] = slink[PT.link[sz]];
f53
                else slink[sz] = PT.link[sz];
eb9
                sz++;
cbb
            for (int v = PT.last; PT.len[v] > 0; v =
911
   slink[v]) {
                sans[v] = dp[i - (PT.len[slink[v]] +
297
   diff[v])]:
                if (diff[v] == diff[PT.link[v]])
85d
                    sans[v] = (sans[v] + sans[PT.link[v]]) %
f20
   MOD;
                dp[i] = (dp[i] + sans[v]) % MOD;
071
           }
cbb
cbb
5f0
        return dp[n];
cbb }
```

#### 3.28 Parsing de Expressao

```
// Operacoes associativas a esquerda por default
// Para mudar isso, colocar em r assoc
// Operacoes com maior prioridade sao feitas primeiro
//
// 68921b
cc1 bool blank(char c) {
        return c == ' ';
cbb }
8e4 bool is_unary(char c) {
        return c == '+' or c == '-';
cbb }
76d bool is_op(char c) {
       if (is_unary(c)) return true;
010
        return c == '*' or c == '/' or c == '+' or c == '-';
31 c
cbb }
fa3 bool r_assoc(char op) {
d41
        // operator unario - deve ser assoc. a direita
cf0
        return op < 0;</pre>
cbb }
79d int priority(char op) {
       // operator unario - deve ter precedencia maior
d41
103
        if (op < 0) return INF;</pre>
727
        if (op == '*' or op == '/') return 2;
439
        if (op == '+' or op == '-') return 1;
        return -1;
daa
cbb }
c15 void process_op(stack<int>& st, stack<int>& op) {
88 c
        char o = op.top(); op.pop();
91c
        if (o < 0) {
4e6
           o *= -1;
           int 1 = st.top(); st.pop();
1e2
Off
           if (o == '+') st.push(1);
7e9
           if (o == '-') st.push(-1);
```

```
9 d 9
        } else {
14c
            int r = st.top(); st.pop();
            int 1 = st.top(); st.pop();
1 e2
           if (o == '*') st.push(1 * r);
1 e 4
f55
            if (o == ',') st.push(1 / r);
605
           if (o == '+') st.push(1 + r);
c40
            if (o == '-') st.push(1 - r);
cbb
cbb }
439 int eval(string& s) {
212
        stack<int> st, op;
d0c
        bool un = true;
1cf
        for (int i = 0; i < s.size(); i++) {</pre>
68d
            if (blank(s[i])) continue;
139
            if (s[i] == '(') {
367
                op.push('(');
99d
                un = true;
            } else if (s[i] == ')') {
130
                while (op.top() != '(') process_op(st, op);
709
75e
                op.pop();
                un = false;
ce2
146
            } else if (is_op(s[i])) {
4 d0
                char o = s[i];
37 c
                if (un and is_unary(o)) o *= -1;
ae3
                while (op.size() and (
cd6
                             (!r_assoc(o) and
   priority(op.top()) >= priority(o)) or
c41
                             (r_assoc(o) and
   priority(op.top()) > priority(o))))
c47
                    process_op(st, op);
c00
                op.push(o);
99d
                un = true;
9 d 9
            } else {
da8
                int val = 0:
c2b
                while (i < s.size() and isalnum(s[i]))</pre>
8a3
                     val = val * 10 + s[i++] - '0';
169
                i--;
25 d
                st.push(val);
ce2
                un = false;
cbb
            }
```

```
cbb }
7f6     while (op.size()) process_op(st, op);
123     return st.top();
cbb }
```

#### 3.29 RMQ com Divide and Conquer

```
// Responde todas as queries em
// O(n log(n))
// 5a6ebd
f74 typedef pair <pair <int, int>, int> iii;
7c6 #define f first
Oab #define s second
87d int n, q, v[MAX];
e3f iii qu[MAX];
aeb int ans[MAX], pref[MAX], sulf[MAX];
0e3 void solve(int l=0, int r=n-1, int ql=0, int qr=q-1) {
        if (1 > r or q1 > qr) return;
8a3
ee4
        int m = (1+r)/2;
        int qL = partition(qu+ql, qu+qr+1, [=](iii x){return
1 b 1
   x.f.s < m; \}) - qu;
        int qR = partition(qu+qL, qu+qr+1, [=](iii x){return
eb0
   x.f.f <= m;}) - qu;
3cd
        pref[m] = sulf[m] = v[m];
        for (int i = m-1; i >= 1; i--) pref[i] = min(v[i],
9f9
   pref[i+1]);
        for (int i = m+1; i <= r; i++) sulf[i] = min(v[i],</pre>
   sulf[i-1]);
b2a
        for (int i = qL; i < qR; i++)
            ans[qu[i].s] = min(pref[qu[i].f.f],
f3a
   sulf[qu[i].f.s]);
        solve(l, m-1, ql, qL-1), solve(m+1, r, qR, qr);
364
cbb }
```

#### 3.30 Segment Intersection

```
// Verifica, dado n segmentos, se existe algum par de
   segmentos
// que se intersecta
//
// O(n log n)
// 3957d8
6e0 bool operator < (const line& a, const line& b) { //
   comparador pro sweepline
        if (a.p == b.p) return ccw(a.p, a.q, b.q);
191
        if (!eq(a.p.x, a.q.x) and (eq(b.p.x, b.q.x) or
   a.p.x+eps < b.p.x)
            return ccw(a.p, a.q, b.p);
780
        return ccw(a.p, b.q, b.p);
dc0
cbb }
8e2 bool has_intersection(vector<line> v) {
        auto intersects = [&](pair<line, int> a, pair<line,</pre>
576
   int> b) {
a08
            return interseg(a.first, b.first);
214
e1b
        vector<pair<pt, pair<int, int>>> w;
        for (int i = 0; i < v.size(); i++) {</pre>
f 14
            if (v[i].q < v[i].p) swap(v[i].p, v[i].q);</pre>
876
e1d
            w.push_back({v[i].p, {0, i}});
            w.push_back({v[i].q, {1, i}});
034
cbb
d1d
        sort(w.begin(), w.end());
        set < pair < line , int >> se;
7f2
e58
        for (auto i : w) {
            line at = v[i.second.second];
bfd
292
            if (i.second.first == 0) {
145
                auto nxt = se.lower_bound({at,
   i.second.second});
                if (nxt != se.end() and intersects(*nxt,
d1e
   {at, i.second.second})) return 1;
                if (nxt != se.begin() and
257
   intersects(*(--nxt), {at, i.second.second})) return 1;
                se.insert({at, i.second.second});
78 f
9d9
            } else {
```

```
884
                auto nxt = se.upper_bound({at,
   i.second.second)), cur = nxt, prev = --cur;
                if (nxt != se.end() and prev != se.begin()
b64
                     and intersects (*nxt, *(--prev))) return
4fb
   1;
                se.erase(cur);
cca
            }
cbb
        }
cbb
        return 0;
bb3
cbb }
```

### 3.31 Sequencia de de Brujin

```
// Se passar sem o terceiro parametro, gera um vetor com
   valores
// em [0, k) de tamanho k^n de forma que todos os subarrays
   ciclicos
// de tamanho n ocorrem exatamente uma vez
// Se passar com um limite lim, gera o menor vetor com
   valores
// em [0, k) que possui lim subarrays de tamanho n distintos
// (assume que lim <= k^n)</pre>
//
// Linear no tamanho da resposta
// 19720c
860 vector<int> de_brujin(int n, int k, int lim = INF) {
        if (k == 1) return vector < int > (lim == INF ? 1 : n,
b55
   0);
5f6
        vector < int > 1 = \{0\}, ret; // 1 eh lyndon word
667
        while (true) {
c86
            if (1.size() == 0) {
1b9
                if (lim == INF) break;
                1.push_back(0);
daf
cbb
686
            if (n % 1.size() == 0) for (int i : 1) {
728
                ret.push_back(i);
c99
                if (ret.size() == n+lim-1) return ret;
cbb
630
            int p = 1.size();
905
            while (l.size() < n) l.push_back(l[l.size()%p]);</pre>
            while (l.size() and l.back() == k-1)
e7f
```

```
1.pop_back();
88a          if (l.size()) l.back()++;
cbb     }
edf     return ret;
cbb }
```

#### 3.32 Shortest Addition Chain

```
// Computa o menor numero de adicoes para construir
// cada valor, comecando com 1 (e podendo salvar variaveis)
// Retorna um par com a dp e o pai na arvore
// A arvore eh tao que o taminho da raiz (1) ate x
// contem os valores que devem ser criados para gerar x
// A profundidade de x na arvore eh dp[x]
// DP funciona para ateh 300, mas a arvore soh funciona
// para ateh 148
//
// 84fcff
// recuperacao certa soh ateh 148 (erra para 149, 233, 298)
3de pair < vector < int > , vector < int > > addition_chain() {
16f
        int MAX = 301;
875
        vector < int > dp(MAX), p(MAX);
1ab
        for (int n = 2; n < MAX; n++) {
7 c 0
            pair < int , int > val = {INF , -1};
            for (int i = 1; i < n; i++) for (int j = i; j; j
212
   = p[j]
94a
                 if (j == n-i) val = min(val, pair(dp[i]+1,
   i));
eb3
            tie(dp[n], p[n]) = val;
            if (n == 9) p[n] = 8;
efe
            if (n == 149 \text{ or } n == 233) \text{ dp}[n] --;
ba1
cbb
717
        return {dp, p};
cbb }
```

#### 3.33 Simple Polygon

```
// Verifica se um poligono com n pontos eh simples // // O(n log n)
```

```
6e0 bool operator < (const line& a, const line& b) { //
   comparador pro sweepline
        if (a.p == b.p) return ccw(a.p, a.q, b.q);
191
231
        if (!eq(a.p.x, a.q.x) and (eq(b.p.x, b.q.x) or
   a.p.x+eps < b.p.x)
780
            return ccw(a.p, a.q, b.p);
        return ccw(a.p, b.q, b.p);
dc0
cbb }
6f3 bool simple(vector<pt> v) {
        auto intersects = [&](pair<line, int> a, pair<line,</pre>
   int> b) {
e72
            if ((a.second+1)%v.size() == b.second or
80e
                 (b.second+1) %v.size() == a.second) return
   false:
            return interseg(a.first, b.first);
a 0.8
214
        };
        vector<line> seg;
41a
        vector<pair<pt, pair<int, int>>> w;
e1b
        for (int i = 0; i < v.size(); i++) {</pre>
f14
            pt at = v[i], nxt = v[(i+1)%v.size()];
0a8
828
            if (nxt < at) swap(at, nxt);</pre>
937
            seg.push_back(line(at, nxt));
            w.push_back({at, {0, i}});
f7e
69 c
            w.push_back({nxt, {1, i}});
d41
            // casos degenerados estranhos
            if (isinseg(v[(i+2)%v.size()], line(at, nxt)))
ae8
   return 0:
            if (isinseg(v[(i+v.size()-1)%v.size()], line(at,
b 88
   nxt))) return 0;
cbb
        sort(w.begin(), w.end());
d1d
7f2
        set < pair < line, int >> se;
e58
        for (auto i : w) {
            line at = seg[i.second.second];
ff8
292
            if (i.second.first == 0) {
                auto nxt = se.lower_bound({at,
145
   i.second.second});
7 c 4
                if (nxt != se.end() and intersects(*nxt,
   {at, i.second.second})) return 0;
```

// c724a4

```
b34
                if (nxt != se.begin() and
   intersects(*(--nxt), {at, i.second.second})) return 0;
                se.insert({at, i.second.second});
78f
9 d 9
            } else {
884
                auto nxt = se.upper_bound({at,
   i.second.second}), cur = nxt, prev = --cur;
                if (nxt != se.end() and prev != se.begin()
b64
                    and intersects(*nxt, *(--prev))) return
403
   0;
                se.erase(cur);
cca
cbb
            }
cbb
6a5
        return 1;
cbb }
3.34 Sweep Direction
// Passa por todas as ordenacoes dos pontos definitas por
   "direcoes"
// Assume que nao existem pontos coincidentes
//
// O(n^2 \log n)
// 6bb68d
4b8 void sweep_direction(vector<pt> v) {
3 d2
        int n = v.size();
163
        sort(v.begin(), v.end(), [](pt a, pt b) {
```

if (a.x != b.x) return a.x < b.x;

for (int i = 0; i < n; i++) for (int j = i+1; j < n;

sort(swapp.begin(), swapp.end(), [&](auto a, auto b)

pt A = rotate90(v[a.first] - v[a.second]);

pt B = rotate90(v[b.first] - v[b.second]);

if (quad(A) == quad(B) and !sarea2(pt(0, 0), A,

swapp.push\_back({i, j}), swapp.push\_back({j, i});

return a.y > b.y;

iota(at.begin(), at.end(), 0);

vector < pair < int , int >> swapp;

vector < int > at(n):

3a5

572

сОс

b89

516

b79

25 e

95f

269

134

247

615

{

j++)

}):

```
B)) return a < b;
224
            return compare_angle(A, B);
сОс
        });
        for (auto par : swapp) {
4e6
            assert(abs(at[par.first] - at[par.second]) == 1);
e24
a96
            int l = min(at[par.first], at[par.second]),
                r = n-1 - max(at[par.first], at[par.second]);
0d3
           // l e r sao quantos caras tem de cada lado do
d41
   par de pontos
            // (cada par eh visitado duas vezes)
d41
            swap(v[at[par.first]], v[at[par.second]]);
9cf
            swap(at[par.first], at[par.second]);
1c0
        }
cbb
cbb }
```

#### 3.35 Triangulação de Delaunay

```
// Computa a triangulacao de Delaunay, o dual
// do diagrama de Voronoi (a menos de casos degenerados)
// Retorna um grafo indexado pelos indices dos pontos, e as
   arestas
// sao as arestas da triangulacao
// As arestas partindo de um vertice ja vem ordenadas por
   angulo,
// ou seja, se o vertice v nao esta no convex hull, (v, v_i,
   v \{i+1\})
// eh um triangulo da triangulacao, em que v_i eh o i-esimo
   vizinho
// Usa o alg d&c, precisa representar MAX_COOR^4, por isso
   __int128
// pra aguentar valores ateh 1e9
//
// Propriedades:
// 1 - 0 grafo tem no max 3n-6 arestas
// 2 - Para todo triangulo, a circunf. que passa pelos 3
   pontos
      nao contem estritamente nenhum ponto
// 3 - A MST euclidiana eh subgrafo desse grafo
// 4 - Cada ponto eh vizinho do ponto mais proximo dele
//
// O(n log n)
// 83ebab
```

```
2ad typedef struct QuadEdge* Q;
ba5 struct QuadEdge {
53e
       int id;
114
        pt o;
41e
        Q rot, nxt;
3 e 5
       bool used;
3fc
        QuadEdge(int id_ = -1, pt o_ = pt(INF, INF)) :
            id(id_), o(o_), rot(nullptr), nxt(nullptr),
4ba
   used(false) {}
        Q rev() const { return rot->rot; }
00f
        Q next() const { return nxt; }
сЗс
        Q prev() const { return rot->next()->rot; }
188
       pt dest() const { return rev()->o; }
0 d4
214 };
91b Q edge(pt from, pt to, int id_from, int id_to) {
        Q e1 = new QuadEdge(id_from, from);
сбе
61b
        Q e2 = new QuadEdge(id_to, to);
8f6
       Q e3 = new QuadEdge;
5ca
        Q e4 = new QuadEdge;
        tie(e1->rot, e2->rot, e3->rot, e4->rot) = \{e3, e4,
e69
   e2, e1};
f22
        e4, e3};
1ad
        return e1;
cbb }
d8d void splice(Q a, Q b) {
a6f
        swap(a->nxt->rot->nxt, b->nxt->rot->nxt);
        swap(a->nxt, b->nxt);
da4
cbb }
167 void del_edge(Q& e, Q ne) { // delete e and assign e <-
   ne
cc0
        splice(e, e->prev());
        splice(e->rev(), e->rev()->prev());
eec
        delete e->rev()->rot, delete e->rev();
7ea
524
        delete e->rot; delete e;
        e = ne;
6 b 2
```

```
cbb }
d08 Q conn(Q a, Q b) {
        Q = edge(a->dest(), b->o, a->rev()->id, b->id);
cc5
        splice(e, a->rev()->prev());
f2b
d37
        splice(e->rev(), b);
6bf
        return e;
cbb }
d64 bool in_c(pt a, pt b, pt c, pt p) { // p ta na circunf.
        _{-}int128 p2 = p*p, A = a*a - p2, B = b*b - p2, C =
268
   c*c - p2;
        return sarea2(p, a, b) * C + sarea2(p, b, c) * A +
   sarea2(p, c, a) * B > 0;
cbb }
540 pair < Q, Q > build_tr(vector < pt > & p, int 1, int r) {
        if (r-1+1 <= 3) {
0.9d
            Q = edge(p[1], p[1+1], 1, 1+1), b =
2eb
   edge(p[l+1], p[r], l+1, r);
           if (r-1+1 == 2) return {a, a->rev()};
912
0ec
            splice(a->rev(), b);
           ll ar = sarea2(p[1], p[1+1], p[r]);
сЗс
            Q c = ar ? conn(b, a) : 0;
1af
021
           if (ar >= 0) return \{a, b -> rev()\};
9db
           return {c->rev(), c};
cbb
        }
       int m = (1+r)/2;
ee4
328
        auto [la, ra] = build_tr(p, 1, m);
        auto [lb, rb] = build_tr(p, m+1, r);
b93
        while (true) {
667
            if (ccw(lb->o, ra->o, ra->dest())) ra =
b99
   ra->rev()->prev();
458
            else if (ccw(lb->o, ra->o, lb->dest())) lb =
   lb->rev()->next():
f97
            else break;
cbb
        Q b = conn(lb -> rev(), ra);
ca5
        auto valid = [&](Q e) { return ccw(e->dest(),
713
   b->dest(), b->o); };
        if (ra->o == la->o) la = b->rev();
ee1
```

```
63f
        if (1b->0 == rb->0) rb = b;
667
        while (true) {
71e
             Q L = b - > rev() - > next();
            if (valid(L)) while (in_c(b->dest(), b->o,
d11
   L->dest(), L->next()->dest()))
1 c 0
                 del_edge(L, L->next());
c76
             Q R = b - > prev();
            if (valid(R)) while (in_c(b->dest(), b->o,
2b0
   R->dest(), R->prev()->dest()))
541
                 del_edge(R, R->prev());
a3a
            if (!valid(L) and !valid(R)) break;
ccd
            if (!valid(L) or (valid(R) and in_c(L->dest(),
   L \rightarrow 0, R \rightarrow 0, R \rightarrow dest())))
                b = conn(R, b->rev());
36 c
            else b = conn(b->rev(), L->rev());
666
cbb
a2b
        return {la, rb};
cbb }
b58 vector < vector < int >> delaunay (vector < pt > v) {
3 d2
        int n = v.size();
397
        auto tmp = v;
        vector < int > idx(n);
135
295
        iota(idx.begin(), idx.end(), 0);
        sort(idx.begin(), idx.end(), [&](int 1, int r) {
fe9
   return v[1] < v[r]; });</pre>
        for (int i = 0; i < n; i++) v[i] = tmp[idx[i]];</pre>
5 d8
        assert(unique(v.begin(), v.end()) == v.end());
780
4aa
        vector < vector < int >> g(n);
4ec
        bool col = true;
        for (int i = 2; i < n; i++) if (sarea2(v[i], v[i-1],
   v[i-2])) col = false;
        if (col) {
bf5
            for (int i = 1; i < n; i++)
aa4
                 g[idx[i-1]].push_back(idx[i]),
839
   g[idx[i]].push_back(idx[i-1]);
96b
             return g;
cbb
d36
        Q = build_tr(v, 0, n-1).first;
        vector <Q> edg = {e};
113
        for (int i = 0; i < edg.size(); e = edg[i++]) {</pre>
5 d 1
             for (Q at = e; !at->used; at = at->next()) {
3ed
```

#### 3.36 Triangulos em Grafos

```
// get_triangles(i) encontra todos os triangulos ijk no grafo
// Custo nas arestas
// retorna {custo do triangulo, {j, k}}
// O(m sqrt(m) log(n)) se chamar para todos os vertices
// fladbc
c0d vector<pair<int, int>> g[MAX]; // {para, peso}
d41 #warning o 'g' deve estar ordenado
9a5 vector<pair<int, pair<int, int>>> get_triangles(int i) {
771
        vector<pair<int, pair<int, int>>> tri;
b23
        for (pair<int, int> j : g[i]) {
2b3
            int a = i, b = j.first;
            if (g[a].size() > g[b].size()) swap(a, b);
6dd
           for (pair < int, int > c : g[a]) if (c.first != b
eb0
   and c.first > j.first) {
                auto it = lower_bound(g[b].begin(),
525
   g[b].end(), make_pair(c.first, -INF));
               if (it == g[b].end() or it->first !=
f55
   c.first) continue;
                tri.push_back({j.second+c.second+it->second,
0aa
   {a == i ? b : a, c.first}});
cbb
cbb
f5e
        return tri;
cbb }
```

#### 4 Matematica

#### 4.1 2-SAT

```
// solve() retorna um par, o first fala se eh possivel
// atribuir, o second fala se cada variavel eh verdadeira
//
// O(|V|+|E|) = O(\#variaveis + \#restricoes)
// ef6b3b
138 struct sat {
e6c
         int n, tot;
789
         vector < vector < int >> g;
0 \, ca
         vector < int > vis, comp, id, ans;
4ce
         stack<int> s;
141
         sat() {}
         sat(int n_{-}) : n(n_{-}), tot(n), g(2*n) {}
172
f32
         int dfs(int i, int& t) {
             int lo = id[i] = t++;
cf0
efc
             s.push(i), vis[i] = 2;
             for (int j : g[i]) {
48e
740
                 if (!vis[j]) lo = min(lo, dfs(j, t));
                 else if (vis[j] == 2) lo = min(lo, id[j]);
994
cbb
             }
3de
             if (lo == id[i]) while (1) {
3 c 3
                 int u = s.top(); s.pop();
                 vis[u] = 1, comp[u] = i;
9 c 5
                 if ((u>>1) < n \text{ and } ans[u>>1] == -1)
91d
    ans [u >> 1] = \sim u \& 1:
                 if (u == i) break;
2ef
cbb
253
             return lo;
cbb
74a
         void add_impl(int x, int y) { // x -> y = !x ou y
26a
             x = x >= 0 ? 2*x : -2*x-1;
2 b8
             y = y >= 0 ? 2*y : -2*y-1;
             g[x].push_back(y);
a1e
             g[y^1].push_back(x^1);
1 e 2
```

```
cbb
        void add_cl(int x, int y) { // x ou y
e85
0b5
             add_impl(\sim x, y);
cbb
        void add_xor(int x, int y) { // x xor y
487
0b7
             add_cl(x, y), add_cl(\simx, \simy);
cbb
978
        void add_eq(int x, int y) { // x = y
c86
             add_xor(\simx, y);
cbb
b10
        void add_true(int x) { // x = T
18b
             add_impl(\sim x, x);
        }
cbb
        void at_most_one(vector<int> v) { // no max um
d14
   verdadeiro
54 d
            g.resize(2*(tot+v.size()));
             for (int i = 0; i < v.size(); i++) {</pre>
f 14
                 add_impl(tot+i, \simv[i]);
8c9
                 if (i) {
a8f
                      add_impl(tot+i, tot+i-1);
b6a
                      add_impl(v[i], tot+i-1);
3d3
                 }
cbb
cbb
258
             tot += v.size();
cbb
        }
a8e
        pair < bool, vector < int >> solve() {
             ans = vector < int > (n, -1);
27b
6bb
             int t = 0:
            vis = comp = id = vector \langle int \rangle (2*tot, 0);
0de
            for (int i = 0; i < 2*tot; i++) if (!vis[i])
53 c
   dfs(i, t);
            for (int i = 0; i < tot; i++)</pre>
f88
                 if (comp[2*i] == comp[2*i+1]) return {false,
4c9
   {}};
997
             return {true, ans};
cbb
214 };
```

## 4.2 Algoritmo de Euclides estendido

```
// Acha x e y tal que ax + by = mdc(a, b) (nao eh unico)
```

```
// Assume a, b >= 0
//
// O(log(min(a, b)))
// 35411d
2be tuple < 11, 11, 11 > ext gcd (11 a, 11 b) {
        if (!a) return {b, 0, 1};
        auto [g, x, y] = ext_gcd(b\%a, a);
550
        return \{g, y - b/a*x, x\};
c59
cbb }
4.3 Avaliação de Interpolação
// Dado 'n' pontos (i, y[i]), i \in [0, n),
// avalia o polinomio de grau n-1 que passa
// por esses pontos em 'x'
// Tudo modular, precisa do mint
//
// O(n)
// 4fe929
ee8 mint evaluate_interpolation(int x, vector<mint> y) {
        int n = v.size();
80e
d41
184
        vector \langle mint \rangle sulf (n+1, 1), fat (n, 1), if at (n);
        for (int i = n-1; i >= 0; i--) sulf[i] = sulf[i+1] *
6fa
   (x - i):
29b
        for (int i = 1; i < n; i++) fat[i] = fat[i-1] * i;</pre>
        ifat[n-1] = 1/fat[n-1];
0 da
        for (int i = n-2; i >= 0; i--) ifat[i] = ifat[i+1] *
3 db
   (i + 1):
        mint pref = 1, ans = 0;
ca1
        for (int i = 0; i < n; pref *= (x - i++)) {</pre>
5ea
            mint num = pref * sulf[i+1];
42f
b4e
            mint den = ifat[i] * ifat[n-1 - i];
            if ((n-1 - i)\%2) den *= -1;
0bd
            ans += y[i] * num * den;
03f
cbb
ba7
        return ans;
```

cbb }

#### 4.4 Berlekamp-Massey

```
// guess_kth(s, k) chuta o k-esimo (0-based) termo
// de uma recorrencia linear que gera s
// Para uma rec. lin. de ordem x, se passar 2x termos
// vai gerar a certa
// Usar aritmetica modular
// O(n<sup>2</sup> log k), em que n = |s|
// 8644e3
b7c template < typename T> T evaluate (vector < T> c, vector < T>
   s, 11 k) {
        int n = c.size();
ff2
9ee
        assert(c.size() <= s.size());
d09
        auto mul = [&](const vector<T> &a, const vector<T>
   &b) {
            vector<T> ret(a.size() + b.size() - 1);
564
d75
            for (int i = 0; i < a.size(); i++) for (int j =
   0; j < b.size(); j++)
                 ret[i+j] += a[i] * b[j];
cff
83 d
            for (int i = ret.size()-1; i \ge n; i--) for (int
   j = n-1; j >= 0; j--)
112
                ret[i-j-1] += ret[i] * c[j];
            ret.resize(min<int>(ret.size(), n));
16 d
edf
            return ret;
214
        };
        vector < T > a = n == 1 ? vector < T > ({c[0]}) :
1a6
   vector < T > ({0, 1}), x = {1};
95 f
        while (k) {
7f1
            if (k\&1) x = mul(x, a);
b28
            a = mul(a, a), k >>= 1;
cbb
dd6
        x.resize(n);
        T ret = 0;
ce8
e72
        for (int i = 0; i < n; i++) ret += x[i] * s[i];
edf
        return ret;
```

```
cbb }
192 template < typename T > vector < T >
   berlekamp_massey(vector <T> s) {
        int n = s.size(), l = 0, m = 1;
ce8
222
        vector < T > b(n), c(n);
        T ld = b[0] = c[0] = 1;
46e
620
        for (int i = 0; i < n; i++, m++) {
793
            T d = s[i]:
            for (int j = 1; j <= 1; j++) d += c[j] * s[i-j];</pre>
ab6
5f0
            if (d == 0) continue;
8 b 4
            vector <T> temp = c;
369
            T coef = d / ld;
            for (int j = m; j < n; j++) c[j] -= coef *
ba6
   b[j-m];
88f
            if (2 * 1 \le i) 1 = i + 1 - 1, b = temp, 1d = d,
   m = 0:
cbb
90c
        c.resize(l + 1);
        c.erase(c.begin());
844
0dc
        for (T\& x : c) x = -x;
807
        return c;
cbb }
2cf template < typename T > T guess_kth(const vector < T > & s, 11
   k) {
ссЗ
        auto c = berlekamp_massey(s);
96a
        return evaluate(c, s, k);
cbb }
     Binomial Distribution
// binom(n, k, p) retorna a probabilidade de k sucessos
// numa binomial(n, p)
// 00d38f
361 double logfact[MAX];
9e4 void calc() {
7 a 0
        logfact[0] = 0;
        for (int i = 1; i < MAX; i++) logfact[i] =</pre>
   logfact[i-1] + log(i);
```

```
94c double binom(int n, int k, double p) {
271    return exp(logfact[n] - logfact[k] - logfact[n-k] +
    k * log(p) + (n-k) * log(1 - p));
cbb }
```

#### 4.6 Convolucao de GCD / LCM

```
// O(n log(n))
// multiple_transform(a)[i] = \sum_d a[d * i]
bbe template < typename T> void multiple_transform(vector < T>&
   v, bool inv = false) {
        vector < int > I(v.size()-1);
64a
        iota(I.begin(), I.end(), 1);
847
        if (inv) reverse(I.begin(), I.end());
674
        for (int i : I) for (int j = 2; i*j < v.size(); j++)</pre>
dad
            v[i] += (inv ? -1 : 1) * v[i*j];
a8a
cbb }
    // gcd convolution(a, b)[k] = \sum {gcd(i, j) = k} a i *
       b_j
    // 984f53
fe2 template < typename T > vector < T > gcd_convolution(vector < T >
   a, vector <T> b) {
        multiple_transform(a), multiple_transform(b);
bdf
        for (int i = 0; i < a.size(); i++) a[i] *= b[i];</pre>
799
        multiple_transform(a, true);
dea
3f5
        return a:
cbb }
    // divisor_transform(a)[i] = \sum_{d|i} a[i/d]
    // aa74e5
be7 template < typename T > void divisor_transform (vector < T > &
   v, bool inv = false) {
        vector < int > I(v.size()-1);
64a
847
        iota(I.begin(), I.end(), 1);
        if (!inv) reverse(I.begin(), I.end());
5ea
        for (int i : I) for (int j = 2; i*j < v.size(); j++)</pre>
dad
14f
            v[i*j] += (inv ? -1 : 1) * v[i];
```

#### 4.7 Deteccao de ciclo - Tortoise and Hare

```
// Linear no tanto que tem que andar pra ciclar,
// O(1) de memoria
// Retorna um par com o tanto que tem que andar
// do f0 ate o inicio do ciclo e o tam do ciclo
// 899f20
58d pair<11, 11> find_cycle() {
273
        11 \text{ tort} = f(f0);
b2b
        ll hare = f(f(f0));
b1b
        11 t = 0;
683
        while (tort != hare) {
b4d
             tort = f(tort):
4b2
            hare = f(f(hare));
c82
             t++;
cbb
0 e8
        11 \text{ st} = 0:
909
        tort = f0:
683
        while (tort != hare) {
             tort = f(tort);
b4d
1 a 2
             hare = f(hare);
397
             st++;
cbb
73d
        11 len = 1;
3 cd
        hare = f(tort);
        while (tort != hare) {
683
1a2
             hare = f(hare);
```

```
040 len++;
cbb }
ebd return {st, len};
cbb }
```

```
// Gera o conjunto n/i, pra todo i, em O(sqrt(n))
// copiei do github do tfg50

79c for(int l = 1, r; l <= n; l = r + 1) {
746     r = n / (n / 1);
d41     // n / i has the same value for l <= i <= r
cbb }</pre>
```

#### 4.9 Eliminacao Gaussiana

```
// Resolve sistema linear
// Retornar um par com o numero de solucoes
// e alguma solucao, caso exista
//
// O(n^2 * m)
// 1d10b5
67a template < typename T>
728 pair <int, vector <T>> gauss(vector <vector <T>> a,
   vector <T> b) {
        const double eps = 1e-6;
6ca
        int n = a.size(), m = a[0].size();
f92
        for (int i = 0; i < n; i++) a[i].push_back(b[i]);</pre>
2f0
3cb
        vector<int> where(m, -1);
        for (int col = 0, row = 0; col < m and row < n;</pre>
237
   col++) {
f05
            int sel = row;
b95
            for (int i=row; i < n; ++i)</pre>
                 if (abs(a[i][col]) > abs(a[sel][col])) sel =
e55
  i;
            if (abs(a[sel][col]) < eps) continue;</pre>
2 c 4
            for (int i = col; i <= m; i++)</pre>
1ae
                 swap(a[sel][i], a[row][i]);
dd2
```

```
2 c3
            where [col] = row;
0c0
            for (int i = 0; i < n; i++) if (i != row) {
96c
                T c = a[i][col] / a[row][col];
                for (int j = col; j <= m; j++)</pre>
d5c
c8f
                    a[i][i] -= a[row][i] * c;
            }
cbb
b70
            row++;
cbb
b1d
        vector <T> ans(m, 0);
e1a
        for (int i = 0; i < m; i++) if (where[i] != -1)</pre>
            ans[i] = a[where[i]][m] / a[where[i]][i];
12a
603
        for (int i = 0; i < n; i++) {</pre>
501
            T sum = 0:
            for (int j = 0; j < m; j++)
a 75
5 a 9
                sum += ans[j] * a[i][j];
            if (abs(sum - a[i][m]) > eps)
b1f
                return pair(0, vector<T>());
6 cd
        }
chh
12e
        for (int i = 0; i < m; i++) if (where[i] == -1)
018
            return pair(INF, ans);
280
        return pair(1, ans);
cbb }
4.10 Eliminacao Gaussiana Z2
// D eh dimensao do espaco vetorial
// add(v) - adiciona o vetor v na base (retorna se ele jah
   pertencia ao span da base)
// coord(v) - retorna as coordenadas (c) de v na base atual
   (basis^T.c = v)
// recover(v) - retorna as coordenadas de v nos vetores na
   ordem em que foram inseridos
// coord(v).first e recover(v).first - se v pertence ao span
//
// Complexidade:
// add, coord, recover: O(D^2 / 64)
// d0a4b3
```

2a3 template <int D> struct Gauss\_z2 {

```
3 c 1
        bitset < D > basis[D], keep[D];
b16
        int rk, in;
        vector<int> id;
482
d41
37f
        Gauss_z2 () : rk(0), in(-1), id(D, -1) {};
04e
        bool add(bitset < D > v) {
42 c
            in++;
fb0
            bitset < D > k;
            for (int i = D - 1; i >= 0; i--) if (v[i]) {
659
                 if (basis[i][i]) v ^= basis[i], k ^= keep[i];
189
4e6
                     k[i] = true, id[i] = in, keep[i] = k;
ea6
                     basis[i] = v, rk++;
6ce
8a6
                     return true;
                }
cbb
cbb
            }
d1f
            return false;
cbb
        pair < bool, bitset < D >> coord(bitset < D > v) {
0f6
            bitset <D> c;
944
659
            for (int i = D - 1; i >= 0; i--) if (v[i]) {
                if (basis[i][i]) v ^= basis[i], c[i] = true;
a39
                 else return {false, bitset <D>()};
8af
            }
cbb
            return {true, c};
5db
cbb
        pair < bool, vector < int >> recover(bitset < D > v) {
330
22 e
            auto [span, bc] = coord(v);
af8
            if (not span) return {false, {}};
f79
            bitset < D > aux;
            for (int i = D - 1; i \ge 0; i - -) if (bc[i]) aux
5a0
   ^= keep[i];
            vector<int> oc:
ea9
            for (int i = D - 1; i >= 0; i--) if (aux[i])
ef2
   oc.push back(id[i]):
            return {true, oc};
001
cbb
214 };
```

#### 4.11 Equação Diofantina Linear

```
// Encontra o numero de soluções de a*x + b*y = c,
// em que x \in [lx, rx] e v \in [lv, rv]
// Usar o comentario para recuperar as solucoes
// (note que o b ao final eh b/gcd(a, b))
// Cuidado com overflow! Tem que caber o quadrado dos valores
//
// O(\log(\min(a, b)))
// 2e8259
c5e template < typename T > tuple < 11, T, T > ext_gcd(11 a, 11 b)
3 bd
        if (!a) return {b, 0, 1};
        auto [g, x, y] = ext_gcd < T > (b\%a, a);
c4b
        return \{g, y - b/a*x, x\};
c59
cbb }
   // numero de solucoes de a*[lx, rx] + b*[ly, ry] = c
14c template < typename T = 11> // usar __int128 se for ate
2a4 ll diophantine(ll a, ll b, ll c, ll lx, ll rx, ll ly, ll
   ry) {
c80
     if (lx > rx \text{ or } ly > ry) \text{ return } 0;
       if (a == 0 and b == 0) return c ? 0 :
   (rx-lx+1)*(ry-ly+1);
        auto [g, x, y] = ext_gcd <T > (abs(a), abs(b));
8ce
       if (c % g != 0) return 0;
9 c3
       if (a == 0) return (rx-lx+1)*(ly <= c/b and c/b <=
249
   rv);
       if (b == 0) return (ry-ly+1)*(lx <= c/a and c/a <=
4ce
   rx):
      x *= a/abs(a) * c/g, y *= b/abs(b) * c/g, a /= g, b
fb1
   /=g;
        auto shift = [\&](T qt) \{ x += qt*b, y -= qt*a; \};
b20
        auto test = [&](T& k, ll mi, ll ma, ll coef, int t) {
efa
866
            shift((mi - k)*t / coef);
79 d
            if (k < mi) shift(coef > 0 ? t : -t);
74d
            if (k > ma) return pair T, T > (rx+2, rx+1);
41f
            T x1 = x;
633
            shift((ma - k)*t / coef);
            if (k > ma) shift(coef > 0 ? -t : t);
c5b
4a9
            return pair < T, T > (x1, x);
```

```
214
        }:
639
        auto [11, r1] = test(x, 1x, rx, b, 1);
       auto [12, r2] = test(y, ly, ry, a, -1);
38 e
       if (12 > r2) swap(12, r2);
c43
       T l = max(11, 12), r = min(r1, r2);
50a
       if (1 > r) return 0;
339
       11 k = (r-1) / abs(b) + 1;
42f
        return k; // solucoes: x = 1 + [0, k)*|b|
839
cbb }
```

#### 4.12 Exponenciacao rapida

```
// (x^y mod m) em O(log(y))
03c ll pow(ll x, ll y, ll m) { // iterativo
       ll ret = 1;
c85
1b8
        while (y) {
895
           if (y & 1) ret = (ret * x) % m;
23b
           y >>= 1;
         x = (x * x) % m;
cc5
cbb
edf
        return ret;
cbb }
03c ll pow(ll x, ll y, ll m) { // recursivo
      if (!v) return 1;
13a
       ll ans = pow(x*x\%m, y/2, m);
426
       return y%2 ? x*ans%m : ans;
88 d
cbb }
```

#### 4.13 Fast Walsh Hadamard Transform

```
// FWHT<'!'>(f) eh SOS DP
// FWHT<'&'>(f) eh soma de superset DP
// Se chamar com ^, usar tamanho potencia de 2!!
//
// O(n log(n))
// 50e84f
```

```
382 template < char op, class T > vector < T > FWHT (vector < T > f,
   bool inv = false) {
b75
       int n = f.size();
       for (int k = 0; (n-1) >> k; k++) for (int i = 0; i <
   n; i++) if (i>>k&1) {
29e
           int j = i^{(1 << k)};
           if (op == '^') f[j] += f[i], f[i] = f[j] -
627
   2*f[i];
a38
          if (op == ', ') f[i] += (inv ? -1 : 1) * f[j];
           if (op == '&') f[j] += (inv ? -1 : 1) * f[i];
93c
cbb
578
       if (op == ', and inv) for (auto& i : f) i /= n;
abe
        return f;
cbb }
4.14 FFT
// Chamar convolution com vector < complex < double >> para FFT
// Precisa do mint para NTT
//
// O(n log(n))
// Para FFT
// de56b9
488 void get_roots(bool f, int n, vector<complex<double>>&
   roots) {
f26
        const static double PI = acosl(-1);
71a
       for (int i = 0; i < n/2; i++) {
           double alpha = i*((2*PI)/n);
b1e
           if (f) alpha = -alpha;
1a1
069
            roots[i] = {cos(alpha), sin(alpha)};
cbb
        }
cbb }
   // Para NTT
   // 91cd08
9f7 template <int p>
97b void get_roots(bool f, int n, vector<mod_int<p>>& roots)
1 e 6
        mod_int r;
de9
    int ord;
       if (p == 998244353) {
57a
```

```
9b6
            r = 102292:
            ord = (1 << 23);
81b
        } else if (p == 754974721) {
1 c c
43a
            r = 739831874;
            ord = (1 << 24);
f0a
b60
        } else if (p == 167772161) {
            r = 243;
a2a
            ord = (1 << 25);
033
        } else assert(false);
6e0
547
        if (f) r = r^(p - 1 - ord/n);
ee2
        else r = r^(ord/n);
        roots[0] = 1:
be4
        for (int i = 1; i < n/2; i++) roots[i] =
   roots[i-1]∗r:
cbb }
    // d5c432
8a2 template < typename T> void fft(vector < T> &a, bool f, int
   N, vector < int > & rev) {
        for (int i = 0; i < N; i++) if (i < rev[i])</pre>
bc7
   swap(a[i], a[rev[i]]);
12b
        int 1, r, m;
        vector<T> roots(N);
cb4
        for (int n = 2; n <= N; n *= 2) {
192
            get roots(f, n, roots);
0f4
5dc
            for (int pos = 0; pos < N; pos += n) {</pre>
                1 = pos+0, r = pos+n/2, m = 0;
432
a88
                 while (m < n/2) {
297
                     auto t = roots[m]*a[r];
254
                     a[r] = a[1] - t;
                     a[1] = a[1] + t;
b8f
925
                    l++: r++: m++:
                }
cbb
cbb
            }
cbb
        }
235
        if (f) {
            auto invN = T(1)/T(N);
1c5
            for (int i = 0; i < N; i++) a[i] = a[i]*invN;
557
cbb
cbb }
```

```
bf5 template < typename T > vector < T > convolution (vector < T > &a.
   vector <T> &b) {
        vector <T> l(a.begin(), a.end());
279
        vector <T> r(b.begin(), b.end());
f41
7 c6
        int ln = l.size(), rn = r.size();
287
        int N = ln+rn-1;
        int n = 1, log n = 0;
f03
        while (n <= N) { n <<= 1; log_n++; }</pre>
ac4
        vector < int > rev(n);
808
bae
        for (int i = 0; i < n; ++i) {</pre>
434
            rev[i] = 0;
920
            for (int j = 0; j < log_n; ++j)
836
                if (i & (1 << j)) rev[i] |= 1 << (log_n-1-j);
cbb
143
        assert(N <= n):
fa4
        l.resize(n):
7 e4
        r.resize(n):
        fft(1, false, n, rev);
56e
        fft(r, false, n, rev);
fcf
        for (int i = 0; i < n; i++) l[i] *= r[i];
917
88b
        fft(l, true, n, rev);
5e1
        l.resize(N);
792
        return 1:
cbb }
    // NTT
    // 3bf256
6c8 template <int p, typename T> vector <mod_int <p>>>
   ntt(vector<T>& a, vector<T>& b) {
        vector < mod_int < p >> A(a.begin(), a.end()),
d52
   B(b.begin(), b.end());
        return convolution(A, B);
d29
cbb }
    // Convolucao de inteiro
    // Precisa do CRT
    //
    // Tabela de valores:
    // [0,1]
              - <int, 1>
    // [-1e5, 1e5] - <11, 2>
    // [-1e9, 1e9] - < int128, 3>
```

```
//
    // 053a7d
b3c template < typename T, int mods >
eec vector<T> int_convolution(vector<int>& a, vector<int>&
   b) {
        static const int M1 = 998244353, M2 = 754974721, M3
   = 167772161;
bf5
        auto c1 = ntt < M1 > (a, b);
        auto c2 = (mods >= 2 ? ntt < M2 > (a, b) :
221
   vector < mod_int < M2 >>());
        auto c3 = (mods >= 3 ? ntt < M3 > (a, b) :
   vector < mod_int < M3 >>());
        vector < T > ans:
2da
5 c 5
        for (int i = 0; i < c1.size(); i++) {</pre>
            crt <T> at(c1[i].v, M1);
c09
            if (mods \ge 2) at = at * crt<T>(c2[i].v, M2);
316
            if (mods >= 3) at = at * crt<T>(c3[i].v, M3);
987
            ans.push back(at.a);
b2b
            if (at.a > at.m/2) ans.back() -= at.m;
26 d
cbb
        }
ba7
        return ans;
cbb }
      Integração Numerica - Metodo de Simpson 3/8
// Integra f no intervalo [a, b], erro cresce proporcional a
   (b - a)^5
676 const int N = 3*100; // multiplo de 3
287 ld integrate(ld a, ld b, function < ld(ld) > f) {
        ld s = 0, h = (b - a)/N;
b4d
067
        for (int i = 1; i < N; i++) s += f(a + i*h)*(i%3 ?
        return (f(a) + s + f(b))*3*h/8;
0da
cbb }
      Inverso Modular
4.16
```

// Computa o inverso de a modulo b

```
// Se b eh primo, basta fazer
// a^{(b-2)}
f0a ll inv(ll a, ll b) {
        return a > 1 ? b - inv(b\%a, a)*b/a : 1;
cbb }
    // computa o inverso modular de 1..MAX-1 modulo um primo
a88 ll inv[MAX]:
0f2 inv[1] = 1;
Ofa for (int i = 2; i < MAX; i++) inv[i] = MOD -</pre>
   MOD/i*inv[MOD%i]%MOD;
4.17 Karatsuba
// Os pragmas podem ajudar
// Para n \sim 2e5, roda em < 1 s
//
// O(n^1.58)
// 8065d6
//#pragma GCC optimize("Ofast")
//#pragma GCC target ("avx,avx2")
77a template < typename T > void kar(T* a, T* b, int n, T* r,
   T* tmp) {
d4c
        if (n <= 64) {
510
            for (int i = 0; i < n; i++) for (int j = 0; j < 0
   n; j++)
212
                r[i+j] += a[i] * b[j];
505
             return;
cbb
194
        int mid = n/2:
        T * atmp = tmp, * btmp = tmp + mid, *E = tmp + n;
2d7
        memset(E, 0, sizeof(E[0])*n);
4f1
        for (int i = 0; i < mid; i++) {</pre>
c65
             atmp[i] = a[i] + a[i+mid];
c72
            btmp[i] = b[i] + b[i+mid];
4b9
cbb
38a
        kar(atmp, btmp, mid, E, tmp+2*n);
        kar(a, b, mid, r, tmp+2*n);
b1e
229
        kar(a+mid, b+mid, mid, r+n, tmp+2*n);
```

for (int i = 0; i < mid; i++) {</pre>

c65

```
735
            T \text{ temp} = r[i+mid]:
            r[i+mid] += E[i] - r[i] - r[i+2*mid];
de7
            r[i+2*mid] += E[i+mid] - temp - r[i+3*mid];
f1e
cbb
cbb }
e38 template < typename T > vector < T > karatsuba (vector < T > a,
   vector <T> b) {
ba3
        int n = max(a.size(), b.size());
        while (n&(n-1)) n++;
a84
ca9
        a.resize(n), b.resize(n);
        vector\langle T \rangle ret(2*n), tmp(4*n);
ae0
        kar(&a[0], &b[0], n, &ret[0], &tmp[0]);
644
        return ret:
edf
cbb }
```

#### 4.18 Logaritmo Discreto

```
// Resolve logaritmo discreto com o algoritmo baby step
   giant step
// Encontra o menor x tal que a^x = b (mod m)
// Se nao tem, retorna -1
//
// O(sqrt(m) * log(sqrt(m))
// 739fa8
d41
da8 int dlog(int b, int a, int m) {
        if (a == 0) return b ? -1 : 1; // caso nao definido
9f8
d41
        a \% = m, b \% = m;
a6e
a10
        int k = 1, shift = 0;
31e
        while (1) {
6e3
           int g = gcd(a, m);
d47
            if (g == 1) break;
d41
9bc
            if (b == k) return shift;
642
           if (b % g) return -1;
            b \neq g, m \neq g, shift++;
c36
9ab
            k = (11) k * a / g % m;
cbb
d41
af7
        int sq = sqrt(m)+1, giant = 1;
```

```
975
         for (int i = 0; i < sq; i++) giant = (11) giant * a
   % m;
d41
0 b 5
         vector < pair < int , int >> baby;
 33f
         for (int i = 0, cur = b; i <= sq; i++) {
 496
             baby.emplace back(cur, i);
             cur = (11) cur * a % m;
 16c
 cbb
 eb4
         sort(baby.begin(), baby.end());
d41
9 c 9
         for (int j = 1, cur = k; j \le sq; j++) {
             cur = (11) cur * giant % m;
78b
             auto it = lower_bound(baby.begin(), baby.end(),
    pair(cur, INF));
d26
            if (it != baby.begin() and (--it)->first == cur)
ac3
                 return sq * j - it->second + shift;
 cbb
         }
d41
 daa
         return -1;
 cbb }
      Miller-Rabin
4.19
// Testa se n eh primo, n <= 3 * 10^18
//
// O(log(n)), considerando multiplicação
// e exponenciacao constantes
// 4ebecc
d8b ll mul(ll a, ll b, ll m) {
e7a
        11 \text{ ret} = a*b - 11((long double)1/m*a*b+0.5)*m;
074
         return ret < 0 ? ret+m : ret;</pre>
cbb }
03c ll pow(ll x, ll y, ll m) {
        if (!v) return 1;
13a
        ll ans = pow(mul(x, x, m), y/2, m);
dbc
        return y%2 ? mul(x, ans, m) : ans;
7fa
cbb }
1a2 bool prime(ll n) {
l aa
       if (n < 2) return 0;
```

```
237
        if (n <= 3) return 1:
        if (n \% 2 == 0) return 0;
9de
        ll r = \__builtin\_ctzll(n - 1), d = n >> r;
f6a
d41
        // com esses primos, o teste funciona garantido para
   n <= 2^64
        // funciona para n \leq 3*10^24 com os primos ate 41
d41
771
        for (int a: {2, 325, 9375, 28178, 450775, 9780504,
   795265022}) {
            11 x = pow(a, d, n);
da0
709
            if (x == 1 \text{ or } x == n - 1 \text{ or a } \% n == 0) continue;
d41
4a2
            for (int j = 0; j < r - 1; j++) {
                x = mul(x, x, n);
10 f
                if (x == n - 1) break;
df0
cbb
e1b
            if (x != n - 1) return 0;
cbb
6a5
        return 1;
cbb }
```

#### 4.20 Pollard's Rho Alg

```
// Usa o algoritmo de deteccao de ciclo de Floyd
// com uma otimizacao na qual o gcd eh acumulado
// A fatoracao nao sai necessariamente ordenada
// O algoritmo rho encontra um fator de n,
// e funciona muito bem quando n possui um fator pequeno
//
// Complexidades (considerando mul constante):
// rho - esperado O(n^{(1/4)}) no pior caso
// fact - esperado menos que O(n^{(1/4)} \log(n)) no pior caso
// ъ00653
d8b ll mul(ll a, ll b, ll m) {
        11 \text{ ret} = a*b - 11((long double)1/m*a*b+0.5)*m;
e7a
074
        return ret < 0 ? ret+m : ret;</pre>
cbb }
03c 11 pow(11 x, 11 y, 11 m) {
13a
       if (!y) return 1;
dbc
        ll ans = pow(mul(x, x, m), y/2, m);
```

```
7fa
        return y%2 ? mul(x, ans, m) : ans;
cbb }
1a2 bool prime(ll n) {
        if (n < 2) return 0;
1aa
237
       if (n <= 3) return 1;
       if (n % 2 == 0) return 0;
9de
f6a
        ll r = \_builtin\_ctzll(n - 1), d = n >> r;
771
        for (int a: {2, 325, 9375, 28178, 450775, 9780504,
   795265022}) {
            11 x = pow(a, d, n);
da0
709
            if (x == 1 or x == n - 1 or a % n == 0) continue;
d41
            for (int j = 0; j < r - 1; j++) {
4a2
10f
                x = mul(x, x, n);
                if (x == n - 1) break;
df0
cbb
e1b
            if (x != n - 1) return 0;
cbb
6a5
        return 1;
cbb }
9cf ll rho(ll n) {
       if (n == 1 or prime(n)) return n;
0f9
        auto f = [n](11 x) \{return mul(x, x, n) + 1;\};
f7c
8a5
        11 x = 0, y = 0, t = 30, prd = 2, x0 = 1, q;
        while (t % 40 != 0 or gcd(prd, n) == 1) {
533
            if (x==y) x = ++x0, y = f(x);
8a0
            q = mul(prd, abs(x-y), n);
e13
            if (q != 0) prd = q;
21f
450
            x = f(x), y = f(f(y)), t++;
cbb
002
        return gcd(prd, n);
cbb }
5b7 vector<ll> fact(ll n) {
        if (n == 1) return {};
1 b 9
        if (prime(n)) return {n};
0ec
        11 d = rho(n);
0ed
1 de
        vector \langle 11 \rangle 1 = fact(d), r = fact(n / d);
```

# 4.21 Produto de dois long long mod m

#### 4.22 Simplex

```
// Maximiza c^T x s.t. Ax <= b, x >= 0
//
// O(2^n), porem executa em O(n^3) no caso medio
// 3a08e5
395 const double eps = 1e-7;
493 namespace Simplex {
69 c
        vector < vector < double >> T;
14e
       int n, m;
43 e
       vector<int> X, Y;
        void pivot(int x, int y) {
c51
8e6
            swap(X[y], Y[x-1]);
            for (int i = 0; i \le m; i++) if (i != y) T[x][i]
d03
   /= T[x][y];
           T[x][y] = 1/T[x][y];
33 c
           for (int i = 0; i \le n; i++) if (i != x and
38b
   abs(T[i][y]) > eps) {
               for (int j = 0; j <= m; j++) if (j != y)
774
  T[i][j] -= T[i][y] * T[x][j];
3d8
               T[i][y] = -T[i][y] * T[x][y];
            }
cbb
        }
cbb
```

```
d41
        // Retorna o par (valor maximo, vetor solucao)
        pair < double , vector < double >> simplex(
6f8
e9d
                 vector < vector < double >> A, vector < double >> b,
   vector < double > c) {
            n = b.size(), m = c.size();
5 bb
002
            T = vector(n + 1, vector < double > (m + 1));
            X = vector < int > (m);
2d9
0 c2
            Y = vector < int > (n);
            for (int i = 0; i < m; i++) X[i] = i;</pre>
115
            for (int i = 0; i < n; i++) Y[i] = i+m;</pre>
51f
            for (int i = 0; i < m; i++) T[0][i] = -c[i];</pre>
5 b 5
603
            for (int i = 0; i < n; i++) {</pre>
                for (int j = 0; j < m; j++) T[i+1][j] =</pre>
ba6
   A[i][i];
                T[i+1][m] = b[i];
eca
cbb
667
            while (true) {
714
                int x = -1, y = -1;
2db
                double mn = -eps;
                for (int i = 1; i <= n; i++) if (T[i][m] <
   mn) mn = T[i][m], x = i;
              if (x < 0) break;
af2
               for (int i = 0; i < m; i++) if (T[x][i] <
   -eps) { y = i; break; }
4a6
                if (y < 0) return {-1e18, {}}; // sem</pre>
 solucao para Ax <= b
7fb
                pivot(x, y);
cbb
667
            while (true) {
714
                int x = -1, y = -1;
2db
                double mn = -eps;
562
                for (int i = 0; i < m; i++) if (T[0][i] <
   mn) mn = T[0][i], y = i;
9b0
               if (y < 0) break;
034
                mn = 1e200:
                for (int i = 1; i <= n; i++) if (T[i][y] >
5af
   eps and T[i][m] / T[i][y] < mn)
                  mn = T[i][m] / T[i][y], x = i;
48f
53b
               if (x < 0) return {1e18, {}}; // c^T x eh
   ilimitado
```

```
pivot(x, y);
7fb
cbb
290
             vector < double > r(m);
            for(int i = 0; i < n; i++) if (Y[i] < m) r[Y[i]]</pre>
32f
   = T[i+1][m];
             return {T[0][m], r};
e59
cbb
        }
cbb }
```

#### Teorema Chines do Resto

```
// Combina equacoes modulares lineares: x = a (mod m)
// O m final eh o lcm dos m's, e a resposta eh unica mod o
   1cm
// Os m nao precisam ser coprimos
// Se nao tiver solucao, o 'a' vai ser -1
// 7cd7b3
153 template < typename T > tuple < T, T, T > ext_gcd(T a, T b) {
3bd
        if (!a) return {b, 0, 1};
        auto [g, x, y] = ext_gcd(b%a, a);
550
        return \{g, y - b/a*x, x\};
c59
cbb }
bfe template < typename T = 11> struct crt {
627
        Ta, m;
5f3
        crt(): a(0), m(1) {}
7eb
        crt(T a_, T m_) : a(a_), m(m_) {}
        crt operator * (crt C) {
911
238
            auto [g, x, y] = ext_gcd(m, C.m);
            if ((a - C.a) % g) a = -1;
dc0
4f9
            if (a == -1 or C.a == -1) return crt(-1, 0);
d09
            T lcm = m/g*C.m;
            T ans = a + (x*(C.a-a)/g \% (C.m/g))*m;
eb2
            return crt((ans % lcm + lcm) % lcm, lcm);
d8d
cbb
        }
214 };
```

#### 4.24Totiente

```
// O(sart(n))
// faeca3
a7e int tot(int n){
0f6
        int ret = n;
505
        for (int i = 2; i*i <= n; i++) if (n % i == 0) {
b0c
            while (n \% i == 0) n /= i;
125
            ret -= ret / i;
cbb
af4
        if (n > 1) ret -= ret / n;
edf
        return ret;
cbb }
      Variações do crivo de Eratosthenes
```

```
// "O" crivo
//
// Encontra maior divisor primo
// Um numero eh primo sse divi[x] == x
// fact fatora um numero <= lim
// A fatoracao sai ordenada
//
// \text{ crivo } - O(n \log(\log(n)))
// fact - O(log(n))
f12 int divi[MAX];
fb9 void crivo(int lim) {
f53
        for (int i = 1; i <= lim; i++) divi[i] = 1;
        for (int i = 2; i <= lim; i++) if (divi[i] == 1)
d46
            for (int j = i; j <= lim; j += i) divi[j] = i;</pre>
018
cbb }
470 void fact(vector<int>& v, int n) {
        if (n != divi[n]) fact(v, n/divi[n]);
ac8
ab4
        v.push_back(divi[n]);
cbb }
    // Crivo linear
```

```
//
    // Mesma coisa que o de cima, mas tambem
    // calcula a lista de primos
    //
    // O(n)
f12 int divi[MAX];
fd3 vector<int> primes;
fb9 void crivo(int lim) {
d5a
        divi[1] = 1:
        for (int i = 2; i <= lim; i++) {
f70
            if (divi[i] == 0) divi[i] = i,
   primes.push_back(i);
           for (int j : primes) {
3ba
                if (j > divi[i] or i*j > lim) break;
522
00b
               divi[i*i] = i;
            }
cbb
        }
cbb
cbb }
    // Crivo de divisores
    // Encontra numero de divisores
    // ou soma dos divisores
    //
    // O(n log(n))
f12 int divi[MAX];
fb9 void crivo(int lim) {
f53
        for (int i = 1; i <= lim; i++) divi[i] = 1;
        for (int i = 2; i <= lim; i++)</pre>
424
594
            for (int j = i; j <= lim; j += i) {</pre>
d41
                // para numero de divisores
                divi[j]++;
9e0
                // para soma dos divisores
d41
                divi[j] += i;
278
cbb
cbb }
```

```
// Crivo de totiente
    // Encontra o valor da funcao
    // totiente de Euler
    // O(n log(log(n)))
5f4 int tot[MAX];
fb9 void crivo(int lim) {
a27
        for (int i = 1; i <= lim; i++) {</pre>
bc9
            tot[i] += i;
          for (int j = 2*i; j <= lim; j += i)
feb
                tot[j] -= tot[i];
837
       }
cbb
cbb }
    // Crivo de função de mobius
    // O(n log(log(n)))
4e1 char meb[MAX];
fb9 void crivo(int lim) {
        for (int i = 2; i <= lim; i++) meb[i] = 2;</pre>
        meb[1] = 1;
ace
842
        for (int i = 2; i <= lim; i++) if (meb[i] == 2)
8 d8
            for (int j = i; j <= lim; j += i) if (meb[j]) {
686
                if (meb[j] == 2) meb[j] = 1;
ae1
                meb[j] *= j/i\%i ? -1 : 0;
cbb
            }
cbb }
    // Crivo linear de funcao multiplicativa
    // Computa f(i) para todo 1 <= i <= n, sendo f
   // uma funcao multiplicativa (se gcd(a,b) = 1,
   // entao f(a*b) = f(a)*f(b))
   // f_prime tem que computar f de um primo, e
   // add_prime tem que computar f(p^{(k+1)}) dado f(p^k) e p
   // Se quiser computar f(p^k) dado p e k, usar os
       comentarios
```

```
// O(n)
fd3 vector<int> primes;
623 int f[MAX], pot[MAX];
    //int expo[MAX];
5c4 void sieve(int lim) {
        // Funcoes para soma dos divisores:
d41
        auto f_prime = [](int p) { return p+1; };
fc9
31 c
        auto add_prime = [](int fpak, int p) { return
   fpak*p+1; };
        //auto f_pak = [](int p, int k) {};
d41
02d
        f[1] = 1:
        for (int i = 2; i <= lim; i++) {</pre>
f70
e6b
            if (!pot[i]) {
                primes.push_back(i);
e74
                f[i] = f_prime(i), pot[i] = i;
f05
                //\exp[i] = 1;
d41
            }
cbb
3b9
            for (int p : primes) {
b9f
                if (i*p > lim) break;
569
                if (i\%p == 0) {
                    f[i*p] = f[i / pot[i]] *
b97
   add_prime(f[pot[i]], p);
d41
                    // se for descomentar, tirar a linha de
   cima tambem
                    //f[i*p] = f[i / pot[i]] * f_pak(p,
d41
   expo[i]+1);
                    //\exp[i*p] = \exp[i]+1;
d41
51f
                    pot[i*p] = pot[i] * p;
c2b
                    break;
                } else {
9d9
                    f[i*p] = f[i] * f[p];
9ef
638
                    pot[i*p] = p;
d41
                    //\exp[i*p] = 1;
                }
cbb
            }
cbb
cbb
        }
cbb }
```

#### 5 DP

## 5.1 Convex Hull Trick (Rafael)

```
// adds tem que serem feitos em ordem de slope
// queries tem que ser feitas em ordem de x
//
// linear
// 30323e
4b5 struct CHT {
942
        int it;
ac1
        vector<ll> a, b;
45 e
        CHT():it(0){}
0 bb
        ll eval(int i, ll x){
93d
             return a[i]*x + b[i];
cbb
63a
        bool useless(){
a20
            int sz = a.size();
35f
            int r = sz-1, m = sz-2, l = sz-3;
d71
            return (b[1] - b[r])*(a[m] - a[1]) <
413
                 (b[1] - b[m])*(a[r] - a[1]);
cbb
bf4
        void add(ll A, ll B){
7f5
            a.push_back(A); b.push_back(B);
565
            while (!a.empty()){
233
                 if ((a.size() < 3) || !useless()) break;</pre>
                 a.erase(a.end() - 2);
ecb
568
                b.erase(b.end() - 2);
cbb
            }
cbb
81b
        ll get(ll x){
d27
            it = min(it, int(a.size()) - 1);
46a
             while (it+1 < a.size()){</pre>
3 c 4
                 if (eval(it+1, x) > eval(it, x)) it++;
f97
                 else break;
cbb
420
            return eval(it, x);
cbb
214 };
```

#### 5.2 Convex Hull Trick Dinamico

```
// para double, use LINF = 1/.0, div(a, b) = a/b
// update(x) atualiza o ponto de intersecao da reta x
// overlap(x) verifica se a reta x sobrepoe a proxima
// add(a, b) adiciona reta da forma ax + b
// query(x) computa maximo de ax + b para entre as retas
//
// O(log(n)) amortizado por insercao
// O(log(n)) por query
// 978376
72c struct Line {
        mutable ll a, b, p;
073
        bool operator < (const Line& o) const { return a <</pre>
8e3
   o.a; }
abf
        bool operator<(ll x) const { return p < x; }</pre>
214 };
326 struct dynamic_hull : multiset <Line, less <>> {
33a
        ll div(ll a, ll b) {
             return a / b - ((a ^ b) < 0 and a % b);
a20
cbb
d41
bbb
        void update(iterator x) {
b2a
             if (next(x) == end()) x -> p = LINF;
772
             else if (x->a == next(x)->a) x->p = x->b >=
   next(x)->b ? LINF : -LINF;
             else x \rightarrow p = div(next(x) \rightarrow b - x \rightarrow b, x \rightarrow a -
424
   next(x) -> a);
cbb
        }
        bool overlap(iterator x) {
71c
f18
             update(x);
             if (next(x) == end()) return 0;
cfa
             if (x->a == next(x)->a) return x->b>=
a4a
   next(x)->b;
             return x - p >= next(x) - p;
d40
cbb
d41
176
        void add(ll a, ll b) {
1 c 7
             auto x = insert({a, b, 0});
```

```
4ab
            while (overlap(x)) erase(next(x)), update(x);
dbc
            if (x != begin() and !overlap(prev(x))) x =
   prev(x), update(x);
0fc
            while (x != begin() and overlap(prev(x)))
4d2
                x = prev(x), erase(next(x)), update(x);
cbb
        }
d41
4ad
        11 query(ll x) {
229
            assert(!empty());
7 d 1
            auto 1 = *lower_bound(x);
aba
            return 1.a * x + 1.b;
cbb
214 }:
5.3 Divide and Conquer DP
// Particiona o array em k subarrays
// minimizando o somatorio das queries
//
// O(k n log n), assumindo quer query(1, r) eh O(1)
// 4efe6b
547 ll dp[MAX][2];
94b void solve(int k, int l, int r, int lk, int rk) {
de6
        if (1 > r) return;
109
        int m = (1+r)/2, p = -1;
d2b
        auto& ans = dp[m][k&1] = LINF;
6e2
        for (int i = max(m, lk); i \le rk; i++) {
            int at = dp[i+1][\sim k\&1] + query(m, i);
324
57 d
            if (at < ans) ans = at, p = i;
cbb
        solve(k, l, m-1, lk, p), solve(k, m+1, r, p, rk);
1ee
cbb }
cf1 ll DC(int n, int k) {
        dp[n][0] = dp[n][1] = 0;
321
f27
        for (int i = 0; i < n; i++) dp[i][0] = LINF;
b76
        for (int i = 1; i \le k; i++) solve(i, 0, n-i, 0,
   n-i);
        return dp[0][k&1];
8e7
cbb }
```

## 5.4 Longest Common Subsequence

```
// Computa a LCS entre dois arrays usando
// o algoritmo de Hirschberg para recuperar
// O(n*m), O(n+m) de memoria
// 337bb3
eaf int lcs_s[MAX], lcs_t[MAX];
a6d int dp[2][MAX];
    // dp[0][j] = max lcs(s[li...ri], t[li, li+j])
d12 void dp_top(int li, int ri, int lj, int rj) {
        memset(dp[0], 0, (r_{j-1}_{j+1})*sizeof(dp[0][0]));
d13
753
        for (int i = li; i <= ri; i++) {</pre>
9aa
            for (int j = rj; j >= lj; j--)
                dp[0][j - 1j] = max(dp[0][j - 1j],
83b
                (lcs_s[i] == lcs_t[j]) + (j > lj ? dp[0][j-1]
741
   - lj] : 0));
            for (int j = lj+1; j <= rj; j++)
04c
939
                dp[0][j-1j] = max(dp[0][j-1j], dp[0][j-1
   -li]);
cbb
      }
cbb }
    // dp[1][j] = max lcs(s[li...ri], t[lj+j, rj])
ca0 void dp_bottom(int li, int ri, int lj, int rj) {
        memset(dp[1], 0, (rj-lj+1)*sizeof(dp[1][0]));
0 d d
        for (int i = ri; i >= li; i--) {
3a2
            for (int j = lj; j <= rj; j++)</pre>
49с
dbb
                dp[1][i - li] = max(dp[1][i - li],
                (lcs_s[i] == lcs_t[j]) + (j < rj ?
4da
   dp[1][j+1 - 1j] : 0));
           for (int j = rj-1; j >= lj; j--)
6ca
769
                dp[1][i - li] = max(dp[1][i - li], dp[1][i+1]
   - 1j]);
cbb
cbb }
93c void solve(vector<int>& ans, int li, int ri, int lj, int
   rj) {
2ad if (li == ri){
```

```
49с
            for (int j = lj; j <= rj; j++)</pre>
                if (lcs_s[li] == lcs_t[j]){
f5b
a66
                     ans.push_back(lcs_t[j]);
c2b
                     break:
                }
cbb
505
            return:
cbb
        if (lj == rj){
534
753
            for (int i = li; i <= ri; i++){</pre>
                if (lcs_s[i] == lcs_t[lj]){
88f
531
                     ans.push_back(lcs_s[i]);
c2b
                     break:
                }
cbb
            }
cbb
505
            return;
cbb
a57
        int mi = (li+ri)/2;
        dp_top(li, mi, lj, rj), dp_bottom(mi+1, ri, lj, rj);
ade
d7a
        int j = 0, mx = -1;
aee
        for (int j = lj-1; j <= rj; j++) {
da8
           int val = 0;
2bb
            if (j >= 1j) val += dp[0][j - 1j];
            if (j < rj) val += dp[1][j+1 - lj];
b9e
ba8
            if (val >= mx) mx = val, j_ = j;
cbb
6f1
        if (mx == -1) return;
        solve(ans, li, mi, lj, j_), solve(ans, mi+1, ri,
   j_+1, rj);
cbb }
058 vector<int> lcs(const vector<int>& s, const vector<int>&
   t) {
953
        for (int i = 0; i < s.size(); i++) lcs_s[i] = s[i];
        for (int i = 0; i < t.size(); i++) lcs_t[i] = t[i];</pre>
577
dab
        vector < int > ans:
        solve(ans, 0, s.size()-1, 0, t.size()-1);
599
ba7
        return ans;
cbb }
```

#### 5.5 Mochila

```
// Resolve mochila, recuperando a resposta
// O(n * cap), O(n + cap) de memoria
// 400885
add int v[MAX], w[MAX]; // valor e peso
582 int dp[2][MAX_CAP];
    // DP usando os itens [1, r], com capacidade = cap
0d6 void get_dp(int x, int 1, int r, int cap) {
        memset(dp[x], 0, (cap+1)*sizeof(dp[x][0]));
       for (int i = 1; i \le r; i++) for (int j = cap; j >=
574
   0; j--)
           if (j - w[i] >= 0) dp[x][j] = max(dp[x][j], v[i]
3a9
   + dp[x][i - w[i]]);
cbb }
5ab void solve(vector<int>& ans, int 1, int r, int cap) {
        if (1 == r) {
893
            if (w[1] <= cap) ans.push_back(1);</pre>
9ff
505
            return;
cbb
        int m = (1+r)/2;
ee4
        get_dp(0, 1, m, cap), get_dp(1, m+1, r, cap);
283
        int left_cap = -1, opt = -INF;
056
        for (int j = 0; j \le cap; j++)
c94
2f2
           if (int at = dp[0][j] + dp[1][cap - j]; at > opt)
                opt = at, left_cap = j;
91d
        solve(ans, 1, m, left_cap), solve(ans, m+1, r, cap -
da3
  left_cap);
cbb }
0d7 vector<int> knapsack(int n, int cap) {
        vector < int > ans;
dab
        solve(ans, 0, n-1, cap);
1e0
ba7
        return ans;
cbb }
```

### 5.6 SOS DP

```
// O(n 2^n)
// soma de sub-conjunto
e03 vector<ll> sos_dp(vector<ll> f) {
      int N = __builtin_ctz(f.size());
        assert((1<<N) == f.size());
        for (int i = 0; i < N; i++) for (int mask = 0; mask
   < (1 << N); mask++)
796
            if (mask>>i&1) f[mask] += f[mask^(1<<i)];</pre>
abe
        return f:
cbb }
    // soma de super-conjunto
e03 vector<ll> sos_dp(vector<ll> f) {
        int N = __builtin_ctz(f.size());
        assert((1<<N) == f.size());
e59
        for (int i = 0; i < N; i++) for (int mask = 0; mask
5 a 5
   < (1 << N); mask++)
            if (\sim mask >> i \&1) f[mask] += f[mask^(1<<i)];
a3c
abe
        return f;
cbb }
```

# 6 Strings

## 6.1 Aho-corasick

```
// query retorna o somatorio do numero de matches de
// todas as stringuinhas na stringona
//
// insert - O(|s| log(SIGMA))
// build - O(N), onde N = somatorio dos tamanhos das strings
// query - O(|s|)
// a30d6e

ea1 namespace aho {
807     map<char, int> to[MAX];
c87     int link[MAX], idx, term[MAX], exit[MAX], sobe[MAX];
```

```
bfc
        void insert(string& s) {
05е
            int at = 0;
            for (char c : s) {
b4f
b68
                auto it = to[at].find(c);
                if (it == to[at].end()) at = to[at][c] =
1 c 9
   ++idx;
361
                else at = it->second;
cbb
            term[at]++, sobe[at]++;
142
cbb
d41 #warning nao esquece de chamar build() depois de inserir
        void build() {
26a
            queue < int > q;
537
            q.push(0);
           link[0] = exit[0] = -1:
dff
402
            while (q.size()) {
379
                int i = q.front(); q.pop();
3 c 4
                for (auto [c, j] : to[i]) {
                    int 1 = link[i];
5da
                    while (1 != -1 and !to[1].count(c)) 1 =
102
   link[1];
7a5
                    link[i] = 1 == -1 ? 0 : to[1][c];
3ab
                     exit[j] = term[link[j]] ? link[j] :
   exit[link[j]];
                    if (exit[j]+1) sobe[j] += sobe[exit[j]];
6f2
113
                    q.push(j);
cbb
                }
            }
cbb
cbb
bc0
        int query(string& s) {
86 d
            int at = 0, ans = 0;
b4f
            for (char c : s){
                while (at != -1 and !to[at].count(c)) at =
1ca
   link[at]:
5b9
                at = at == -1 ? 0 : to[at][c];
2b1
                ans += sobe[at]:
cbb
ba7
            return ans;
        }
cbb
cbb }
```

# 6.2 Algoritmo Z

```
// z[i] = lcp(s, s[i..n))
//
// Complexidades:
// z - O(|s|)
// \text{ match - } O(|s| + |p|)
// 74a9e1
a19 vector <int> get_z(string s) {
163
        int n = s.size();
        vector < int > z(n, 0);
2b1
fae
       int 1 = 0, r = 0;
6f5
       for (int i = 1; i < n; i++) {
0af
            if (i \le r) z[i] = min(r - i + 1, z[i - 1]);
457
            while (i + z[i] < n \text{ and } s[z[i]] == s[i + z[i]])
   z[i]++;
            if (i + z[i] - 1 > r) l = i, r = i + z[i] - 1;
65e
cbb
070
        return z;
cbb }
```

## 6.3 Automato de Sufixo

```
// Automato que aceita os sufixos de uma string
// Todas as funcoes sao lineares
// c37a72
16e namespace sam {
c1a
        int cur, sz, len[2*MAX], link[2*MAX], acc[2*MAX];
       int nxt[2*MAX][26];
0 b8
        void add(int c) {
e6a
17a
           int at = cur;
           len[sz] = len[cur]+1, cur = sz++;
9a6
500
            while (at != -1 and !nxt[at][c]) nxt[at][c] =
   cur, at = link[at];
7ea
          if (at == -1) { link[cur] = 0; return; }
654
           int q = nxt[at][c];
```

```
fd9
            if (len[q] == len[at]+1) { link[cur] = q;
   return; }
31f
           int qq = sz++;
            len[qq] = len[at]+1, link[qq] = link[q];
2c3
            for (int i = 0; i < 26; i++) nxt[qq][i] =
9a9
   nxt[q][i];
            while (at != -1 and nxt[at][c] == q) nxt[at][c]
e76
   = qq, at = link[at];
            link[cur] = link[q] = qq;
8d8
cbb
94e
        void build(string& s) {
889
            cur = 0, sz = 0, len[0] = 0, link[0] = -1, sz++;
            for (auto i : s) add(i-'a');
9fe
17a
            int at = cur:
            while (at) acc[at] = 1, at = link[at]:
121
cbb
        }
        // coisas que da pra fazer:
d41
        11 distinct_substrings() {
28 c
04b
            11 \text{ ans} = 0;
           for (int i = 1; i < sz; i++) ans += len[i] -
a1e
   len[link[i]];
ba7
            return ans;
cbb
        string longest_common_substring(string& S, string&
a6c
  T) {
419
            build(S);
            int at = 0, 1 = 0, ans = 0, pos = -1;
111
           for (int i = 0; i < T.size(); i++) {</pre>
d59
                while (at and !nxt[at][T[i]-'a']) at =
f2c
   link[at], 1 = len[at];
efa
                if (nxt[at][T[i]-'a']) at =
   nxt[at][T[i]-'a'], 1++;
                else at = 0.1 = 0:
749
                if (1 > ans) ans = 1, pos = i;
a1a
cbb
20f
            return T.substr(pos-ans+1, ans);
cbb
        11 dp [2*MAX];
46 e
455
        11 paths(int i) {
2a8
            auto& x = dp[i];
dee
            if (x) return x;
```

```
483
            x = 1;
71c
            for (int j = 0; j < 26; j++) if (nxt[i][j]) x +=
   paths(nxt[i][j]);
            return x;
ea5
cbb
105
        void kth_substring(int k, int at=0) { // k=1 : menor
   substring lexicog.
9d2
            for (int i = 0; i < 26; i++) if (k and
   nxt[at][i]) {
                if (paths(nxt[at][i]) >= k) {
d58
d02
                    cout << char('a'+i);</pre>
c43
                    kth_substring(k-1, nxt[at][i]);
505
                    return:
                }
cbb
5f4
                k -= paths(nxt[at][i]);
cbb
            }
cbb
        }
214 };
6.4 eertree
// Constroi a eertree, caractere a caractere
// Inicializar com a quantidade de caracteres maxima
// size() retorna a quantidade de substrings pal. distintas
// depois de chamar propagate(), cada substring palindromica
// ocorre qt[i] vezes. O propagate() retorna o numero de
// substrings pal. com repeticao
//
// O(n) amortizado, considerando alfabeto O(1)
// a2e693
8eb struct eertree {
7 c.c
        vector < vector < int >> t;
42e
        int n, last, sz;
745
        vector < int > s, len, link, qt;
d36
        eertree(int N) {
            t = vector(N+2, vector(26, int()));
ec8
cee
            s = len = link = qt = vector < int > (N+2);
cd1
            s[0] = -1:
            link[0] = 1, len[0] = 0, link[1] = 1, len[1] =
288
```

-1;

```
688
            sz = 2, last = 0, n = 1;
cbb
        void add(char c) {
244
692
            s[n++] = c -= 'a';
            while (s[n-len[last]-2] != c) last = link[last];
34f
289
            if (!t[last][c]) {
dab
                int prev = link[last];
                while (s[n-len[prev]-2] != c) prev =
553
   link[prev];
                link[sz] = t[prev][c];
fb2
3f5
                len[sz] = len[last]+2;
                t[last][c] = sz++;
1f8
cbb
344
            qt[last = t[last][c]]++;
cbb
        int size() { return sz-2; }
f 17
        11 propagate() {
2af
            11 \text{ ret = 0};
b73
            for (int i = n; i > 1; i--) {
ebb
fd3
                qt[link[i]] += qt[i];
db5
                ret += qt[i];
            }
cbb
edf
            return ret;
        }
cbb
214 };
6.5 KMP
// mathcing(s, t) retorna os indices das ocorrencias
// de s em t
// autKMP constroi o automato do KMP
// Complexidades:
// pi - O(n)
// match - O(n + m)
// construir o automato - O(|sigma|*n)
// n = |padrao| e m = |texto|
// f50359
ea8 template < typename T> vector < int> pi(T s) {
```

vector < int > p(s.size());

019

```
725
         for (int i = 1, j = 0; i < s.size(); i++) {</pre>
             while (j \text{ and } s[j] != s[i]) j = p[j-1];
a51
            if (s[j] == s[i]) j++;
973
f8c
             p[i] = j;
cbb
74e
         return p;
cbb }
    // c82524
c10 template < typename T > vector < int > matching(T& s, T& t) {
         vector < int > p = pi(s), match;
a1b
        for (int i = 0, j = 0; i < t.size(); i++) {</pre>
             while (j and s[j] != t[i]) j = p[j-1];
6be
            if (s[j] == t[i]) j++;
c4d
             if (j == s.size()) match.push_back(i-j+1), j =
   p[j-1];
       }
cbb
ed8
         return match;
cbb }
    // 79bd9e
a2d struct KMPaut : vector < vector < int >> {
47 c
         KMPaut(){}
6 c 7
        KMPaut (string& s) : vector < vector < int >> (26,
   vector < int > (s.size()+1)) {
503
             vector < int > p = pi(s);
04b
             auto& aut = *this;
4fa
             aut[s[0]-'a'][0] = 1;
            for (char c = 0; c < 26; c++)
19a
5 d3
                 for (int i = 1; i <= s.size(); i++)</pre>
                     aut[c][i] = s[i] - a' == c ? i+1 :
   aut[c][p[i-1]];
       }
cbb
214 };
6.6 Manacher
// manacher recebe um vetor de T e retorna o vetor com
   tamanho dos palindromos
// ret[2*i] = tamanho do maior palindromo centrado em i
// ret[2*i+1] = tamanho maior palindromo centrado em i e i+1
//
```

```
// Complexidades:
// manacher - O(n)
// palindrome - <0(n), 0(1)>
// pal_end - O(n)
// ebb184
28a template < typename T > vector < int > manacher (const T& s) {
        int l = 0, r = -1, n = s.size();
        vector < int > d1(n), d2(n);
fc9
603
        for (int i = 0; i < n; i++) {</pre>
            int k = i > r ? 1 : min(d1[l+r-i], r-i);
821
             while (i+k < n \&\& i-k >= 0 \&\& s[i+k] == s[i-k])
61a
   k++;
            d1[i] = k--;
61e
9f6
            if (i+k > r) l = i-k, r = i+k;
cbb
        1 = 0, r = -1;
e03
        for (int i = 0; i < n; i++) {</pre>
603
            int k = i > r ? 0 : min(d2[1+r-i+1], r-i+1); k++;
a64
             while (i+k \le n \&\& i-k >= 0 \&\& s[i+k-1] ==
2c6
   s[i-k]) k++;
            d2[i] = --k;
eaa
26 d
            if (i+k-1 > r) l = i-k, r = i+k-1;
cbb
c41
        vector \langle int \rangle ret(2*n-1);
        for (int i = 0; i < n; i++) ret[2*i] = 2*d1[i]-1;
e6b
e1d
        for (int i = 0; i < n-1; i++) ret[2*i+1] = 2*d2[i+1];
edf
        return ret;
cbb }
    // 60c6f5
    // verifica se a string s[i..j] eh palindromo
cac template < typename T> struct palindrome {
f97
        vector < int > man:
        palindrome(const T& s) : man(manacher(s)) {}
b2d
        bool query(int i, int j) {
9d7
             return man[i+j] >= j-i+1;
bad
cbb
        }
214 };
    // 8bd4d5
```

```
// tamanho do maior palindromo que termina em cada
       posicao
7cb template < typename T > vector < int > pal_end(const T& s) {
        vector < int > ret(s.size());
        palindrome <T> p(s);
fde
d51
        ret[0] = 1;
88e
        for (int i = 1; i < s.size(); i++) {</pre>
a32
            ret[i] = min(ret[i-1]+2, i+1);
6ea
            while (!p.query(i-ret[i]+1, i)) ret[i]--;
cbb
edf
        return ret;
cbb }
     Min/max suffix/cyclic shift
// Computa o indice do menor/maior sufixo/cyclic shift
// da string, lexicograficamente
//
// O(n)
// af0367
016 template < typename T > int max_suffix(T s, bool mi =
   false) {
476
        s.push_back(*min_element(s.begin(), s.end())-1);
1 a 4
        int ans = 0;
88e
        for (int i = 1; i < s.size(); i++) {</pre>
            int j = 0;
eec
708
            while (ans+j < i and s[i+j] == s[ans+j]) j++;
7a2
            if (s[i+j] > s[ans+j]) {
                if (!mi or i != s.size()-2) ans = i;
b52
c05
            } else if (j) i += j-1;
cbb
ba7
        return ans;
cbb }
a1a template < typename T > int min_suffix(T s) {
76b
        for (auto& i : s) i *= -1;
09d
        s.push_back(*max_element(s.begin(), s.end())+1);
925
        return max_suffix(s, true);
cbb }
```

97c template < typename T > int max\_cyclic\_shift(T s) {

```
int n = s.size();
for (int i = 0; i < n; i++) s.push_back(s[i]);
return max_suffix(s);
bb }

8a template < typename T > int min_cyclic_shift(T s) {
for (auto& i : s) i *= -1;
return max_cyclic_shift(s);
bb }
```

# 6.8 String Hashing

```
// Complexidades:
// construtor - O(|s|)
// operator() - 0(1)
878 mt19937 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
d41
463 int uniform(int 1, int r) {
a7f
        uniform_int_distribution <int> uid(1, r);
f54
        return uid(rng);
cbb }
d41
9e0 template <int MOD> struct str_hash { // 116fcb
c63
        static int P;
        vector<ll> h, p;
dcf
        str_hash(string s) : h(s.size()), p(s.size()) {
ea8
            p[0] = 1, h[0] = s[0];
7a2
            for (int i = 1; i < s.size(); i++)</pre>
ad7
                 p[i] = p[i - 1] * P \% MOD, h[i] = (h[i - 1] * P +
84 c
   s[i])%MOD;
cbb
        11 operator()(int 1, int r) { // retorna hash
af7
            ll\ hash = h[r] - (l\ ?\ h[l\ -\ 1]*p[r\ -\ l\ +\ 1]%MOD
749
   : 0);
             return hash < 0 ? hash + MOD : hash;</pre>
dfd
cbb
214 };
217 template <int MOD > int str_hash < MOD > :: P = uniform (256,
   MOD - 1); // 1 > |sigma|
```

# 6.9 String Hashing - modulo 2<sup>61</sup> - 1

```
// Quase duas vezes mais lento
//
// Complexidades:
// build - O(|s|)
// operator() - 0(1)
//
// d3c0f0
9d0 const 11 MOD = (111 < < 61) - 1;
e38 ll mulmod(ll a, ll b) {
        const static ll LOWER = (111<<30) - 1, GET31 =</pre>
   (111 << 31) - 1;
        11 11 = a\&LOWER, h1 = a>>30, 12 = b\&LOWER, h2 =
410
   b>>30;
d54
        11 m = 11*h2 + 12*h1, h = h1*h2;
        ll ans = 11*12 + (h>>1) + ((h&1)<<60) + (m>>31) +
784
   ((m\&GET31) << 30) + 1;
        ans = (ans\&MOD) + (ans >> 61), ans = (ans\&MOD) +
1 \, dd
   (ans >> 61);
        return ans - 1;
c0f
cbb }
798 mt19937 64
   rng(chrono::steady_clock::now().time_since_epoch().count());
f89 ll uniform(ll l, ll r) {
969
        uniform_int_distribution < ll> uid(l, r);
f54
        return uid(rng);
cbb }
d41
d7d struct str_hash {
c20
        static 11 P;
dcf
        vector<11> h, p;
ea8
        str_hash(string s) : h(s.size()), p(s.size()) {
            p[0] = 1, h[0] = s[0];
7a2
ad7
            for (int i = 1; i < s.size(); i++)</pre>
                p[i] = mulmod(p[i - 1], P), h[i] =
632
   (\text{mulmod}(h[i-1], P) + s[i])\%MOD;
cbb
        11 operator()(int 1, int r) { // retorna hash
af7
```

```
s[1...r]
            ll hash = h[r] - (l ? mulmod(h[l - 1], p[r - l +
538
                                                                   481 vector <int > kasai (string s, vector <int > sa) {
   1]) : 0);
                                                                   232
                                                                           int n = s.size(), k = 0;
                                                                   408
                                                                           vector < int > ra(n), lcp(n);
            return hash < 0 ? hash + MOD : hash;</pre>
dfd
                                                                           for (int i = 0; i < n; i++) ra[sa[i]] = i;</pre>
cbb
        }
                                                                   676
214 };
6c5 ll str_hash::P = uniform(256, MOD - 1); // 1 > | sigma |
                                                                   740
                                                                           for (int i = 0; i < n; i++, k -= !!k) {
                                                                   199
                                                                               if (ra[i] == n-1) { k = 0; continue; }
                                                                   1 de
                                                                               int j = sa[ra[i]+1];
6.10 Suffix Array - O(n log n)
                                                                   891
                                                                               while (i+k < n \text{ and } j+k < n \text{ and } s[i+k] == s[j+k])
                                                                      k++;
// kasai recebe o suffix array e calcula lcp[i],
                                                                               lcp[ra[i]] = k;
                                                                   866
// o lcp entre s[sa[i],...,n-1] e s[sa[i+1],..,n-1]
                                                                           }
                                                                   cbb
                                                                   5ed
                                                                           return lcp;
// Complexidades:
                                                                   cbb }
// suffix_array - O(n log(n))
// kasai - O(n)
                                                                   6.11 Suffix Array - O(n)
// d3a6ce
                                                                   // Rapidao
733 vector<int> suffix_array(string s) {
        s += "$";
                                                                   // Computa o suffix array em 'sa', o rank em 'rnk'
b38
043
        int n = s.size(), N = max(n, 260);
                                                                   // e o lcp em 'lcp'
        vector<int> sa(n), ra(n);
                                                                   // query(i, j) retorna o LCP entre s[i..n-1] e s[j..n-1]
2f3
        for(int i = 0; i < n; i++) sa[i] = i, ra[i] = s[i];</pre>
29b
                                                                   //
                                                                   // Complexidades
        for(int k = 0; k < n; k ? k *= 2 : k++) {
0a2
                                                                   // O(n) para construir
5ce
            vector < int > nsa(sa), nra(n), cnt(N);
                                                                   // query - O(1)
            for(int i = 0; i < n; i++) nsa[i] =
                                                                   // bab412
fae
   (nsa[i]-k+n)%n, cnt[ra[i]]++;
                                                                   1a5 template < typename T> struct rmg {
            for(int i = 1; i < N; i++) cnt[i] += cnt[i-1];</pre>
4c4
                                                                   517
                                                                           vector <T> v;
368
            for(int i = n-1; i+1; i--) sa[--cnt[ra[nsa[i]]]]
                                                                   fcc
                                                                           int n; static const int b = 30;
   = nsa[i]:
                                                                   70e
                                                                           vector < int > mask, t;
            for(int i = 1, r = 0; i < n; i++) nra[sa[i]] = r
                                                                   183
                                                                           int op(int x, int y) { return v[x] \le v[y] ? x : y; }
28f
   += ra[sa[i]] !=
                                                                           int msb(int x) { return
                                                                   ee1
                 ra[sa[i-1]] or ra[(sa[i]+k)%n] !=
                                                                      __builtin_clz(1)-__builtin_clz(x); }
f86
                                                                           int small(int r, int sz = b) { return
   ra[(sa[i-1]+k)%n];
26b
            ra = nra;
                                                                      r-msb(mask[r]&((1<<sz)-1)); }
            if (ra[sa[n-1]] == n-1) break;
                                                                           rmq() {}
d5e
                                                                   6ad
                                                                           rmq(const vector<T>& v_) : v(v_), n(v.size()),
cbb
                                                                      mask(n), t(n) {
057
        return vector < int > (sa.begin()+1, sa.end());
                                                                              for (int i = 0, at = 0; i < n; mask[i++] = at |=</pre>
cbb }
                                                                   2e5
```

```
1) {
                at = (at << 1) &((1 << b) -1);
a61
                while (at and op(i-msb(at&-at), i) == i) at
c00
   ^= at&-at;
cbb
           for (int i = 0; i < n/b; i++) t[i] =
   small(b*i+b-1);
           for (int j = 1; (1<<j) <= n/b; j++) for (int i = 1)
   0; i+(1<< j) <= n/b; i++)
                t[n/b*j+i] = op(t[n/b*(j-1)+i],
ba5
  t[n/b*(j-1)+i+(1<<(j-1))]);
cbb
e34
        int index_query(int 1, int r) {
            if (r-l+1 \le b) return small(r, r-l+1);
27b
           int x = 1/b+1, y = r/b-1;
e80
           if (x > y) return op(small(1+b-1), small(r));
fd3
           int j = msb(y-x+1);
a4e
           int ans = op(small(1+b-1), op(t[n/b*j+x],
ea3
t[n/b*j+y-(1<<j)+1]));
            return op(ans, small(r));
be6
cbb
093
        T query(int 1, int r) { return v[index_query(1, r)];
 }
214 };
9d7 struct suffix_array {
ac0
        string s;
1a8
        int n;
        vector < int > sa, cnt, rnk, lcp;
5b4
2de
       rmq < int > RMQ;
d6e
        bool cmp(int a1, int b1, int a2, int b2, int a3=0,
   int b3=0) {
            return a1 != b1 ? a1 < b1 : (a2 != b2 ? a2 < b2
91d
  : a3 < b3):
cbb
4a4
        template < typename T > void radix(int* fr, int* to, T*
  r, int N, int k) {
           cnt = vector < int > (k+1, 0);
c17
            for (int i = 0; i < N; i++) cnt[r[fr[i]]]++;
bac
           for (int i = 1; i \le k; i++) cnt[i] += cnt[i-1];
703
            for (int i = N-1; i+1; i--) to [--cnt[r[fr[i]]]]
000
```

```
= fr[i]:
cbb
d66
        void rec(vector<int>& v, int k) {
a76
            auto &tmp = rnk, &m0 = lcp;
3a9
            int N = v.size()-3, sz = (N+2)/3, sz2 = sz+N/3;
7f8
            vector < int > R(sz2+3);
74f
            for (int i = 1, j = 0; j < sz2; i += i%3) R[j++]</pre>
  = i;
            radix(&R[0], &tmp[0], &v[0]+2, sz2, k);
b30
            radix(&tmp[0], &R[0], &v[0]+1, sz2, k);
207
5 f 1
            radix(&R[0], &tmp[0], &v[0]+0, sz2, k);
af5
            int dif = 0;
ed9
            int 10 = -1, 11 = -1, 12 = -1:
            for (int i = 0: i < sz2: i++) {</pre>
d81
                if (v[tmp[i]] != 10 or v[tmp[i]+1] != 11 or
8de
   v[tmp[i]+2] != 12)
                    10 = v[tmp[i]], 11 = v[tmp[i]+1], 12 =
b43
   v[tmp[i]+2], dif++;
199
                if (tmp[i]%3 == 1) R[tmp[i]/3] = dif;
1f5
                else R[tmp[i]/3+sz] = dif;
cbb
            }
47 f
            if (dif < sz2) {</pre>
                rec(R, dif);
146
746
                for (int i = 0; i < sz2; i++) R[sa[i]] = i+1;
            } else for (int i = 0; i < sz2; i++) sa[R[i]-1]
8 b 7
  = i;
            for (int i = 0, j = 0; j < sz2; i++) if (sa[i] <
6f4
   sz) tmp[j++] = 3*sa[i];
            radix(&tmp[0], &m0[0], &v[0], sz, k);
7 ce
74d
            for (int i = 0: i < sz2: i++)
                sa[i] = sa[i] < sz ? 3*sa[i]+1 :
с9е
   3*(sa[i]-sz)+2;
332
            int at = sz2+sz-1, p = sz-1, p2 = sz2-1;
1 c9
            while (p >= 0 \text{ and } p2 >= 0) {
                if ((sa[p2]%3==1 and cmp(v[m0[p]],
3 b 3
   v[sa[p2]], R[m0[p]/3],
0ce
                    R[sa[p2]/3+sz]) or (sa[p2]\%3==2 and
```

```
cmp(v[m0[p]], v[sa[p2]],
                     v[m0[p]+1], v[sa[p2]+1], R[m0[p]/3+sz],
af6
   R[sa[p2]/3+1]))
                     sa[at--] = sa[p2--];
300
                 else sa[at--] = m0[p--];
cb0
cbb
f2b
            while (p >= 0) sa[at--] = m0[p--];
            if (N\%3==1) for (int i = 0; i < N; i++) sa[i] =
eb6
   sa[i+1];
cbb
        }
938
        suffix_array(const string& s_) : s(s_), n(s.size()),
   sa(n+3),
                 cnt(n+1), rnk(n), lcp(n-1) {
e62
9fe
            vector < int > v(n+3);
f9b
            for (int i = 0; i < n; i++) v[i] = i;
            radix(&v[0], &rnk[0], &s[0], n, 256);
eba
e6d
            int dif = 1;
830
            for (int i = 0; i < n; i++)
                 v[rnk[i]] = dif += (i and s[rnk[i]] !=
419
   s[rnk[i-1]]);
7cf
            if (n >= 2) rec(v, dif);
fb9
            sa.resize(n);
            for (int i = 0; i < n; i++) rnk[sa[i]] = i;</pre>
76f
892
            for (int i = 0, k = 0; i < n; i++, k -= !!k) {
668
                 if (rnk[i] == n-1) {
5a4
                     k = 0;
5e2
                     continue;
                }
cbb
39a
                 int j = sa[rnk[i]+1];
891
                 while (i+k < n \text{ and } j+k < n \text{ and } s[i+k] ==
   s[j+k]) k++;
                 lcp[rnk[i]] = k;
825
cbb
9ff
            RMQ = rmq < int > (lcp);
cbb
d41
        // hash ateh aqui (sem o RMQ): 1ff700
588
        int query(int i, int j) {
d97
            if (i == j) return n-i;
223
            i = rnk[i], j = rnk[j];
```

```
c3a
             return RMQ.query(min(i, j), max(i, j)-1);
cbb
71c
        pair < int, int > next(int L, int R, int i, char c) {
024
             int 1 = L, r = R+1;
40 c
            while (1 < r) {
                int m = (1+r)/2;
ee4
                if (i+sa[m] >= n or s[i+sa[m]] < c) l = m+1;</pre>
e7e
ef3
                 else r = m;
cbb
            }
575
            if (1 == R+1 \text{ or } s[i+sa[1]] > c) \text{ return } \{-1, -1\};
eb7
            L = 1:
9e2
            1 = L, r = R+1;
40 c
            while (1 < r) {
ee4
                int m = (1+r)/2;
1a1
                if (i+sa[m] >= n \text{ or } s[i+sa[m]] <= c) l = m+1;
ef3
                 else r = m:
cbb
            }
56a
            R = 1-1;
e13
            return {L, R};
cbb
d41
        // quantas vezes 't' ocorre em 's' - O(|t| \log n)
66d
        int count_substr(string& t) {
b2b
            int L = 0, R = n-1;
            for (int i = 0; i < t.size(); i++) {</pre>
c9d
de0
                 tie(L, R) = next(L, R, i, t[i]);
4fc
                if (L == -1) return 0;
cbb
fbf
            return R-L+1;
cbb
        }
d41
        // exemplo de f que resolve o problema
d41
   https://codeforces.com/edu/course/2/lesson/2/5/practice/contest
57e
        ll f(ll k) { return k*(k+1)/2; }
e68
        ll dfs(int L, int R, int p) { // dfs na suffix tree
   chamado em pre ordem
            int ext = L != R ? RMQ.query(L, R-1) : n - sa[L];
c54
d41
            // Tem 'ext - p' substrings diferentes que
   ocorrem 'R-L+1' vezes
```

```
d41
            // O LCP de todas elas eh 'ext'
f80
            ll ans = (ext-p)*f(R-L+1);
           // L eh terminal, e folha sse L == R
d41
            if (sa[L]+ext == n) L++;
63 c
           /* se for um SA de varias strings separadas como
   s#t$u&, usar no lugar do if de cima
                (separadores < 'a', diferentes e inclusive
548
   no final)
            while (L <= R && (sa[L]+ext == n \mid | s[sa[L]+ext]
   < 'a')) {
f49
              L++;
           } */
792
add
            while (L <= R) {
                int idx = L != R ? RMQ.index_query(L, R-1) :
5a8
   -1;
                if (idx == -1 \text{ or } lcp[idx] != ext) idx = R;
5ef
478
                ans += dfs(L, idx, ext);
28 d
                L = idx+1;
            }
cbb
ba7
            return ans;
cbb
        }
d41
       // sum over substrings: computa, para toda substring
   t distinta de s,
       // \sum f(# ocorrencias de t em s) - O (n)
d41
        ll sos() { return dfs(0, n-1, 0); }
ca8
214 };
6.12 Suffix Array Dinamico
// Mantem o suffix array, lcp e rank de uma string,
// premitindo push_front e pop_front
// O operador [i] return um par com sa[i] e lcp[i]
// lcp[i] tem o lcp entre sa[i] e sa[i-1] (lcp[0] = 0)
// Complexidades:
// Construir sobre uma string de tamanho n: O(n log n)
// push front e pop front: O(log n) amortizado
```

```
// 4c2a2e
2fe struct dyn_sa {
3 c9
        struct node {
1 d4
             int sa, lcp;
ed1
             node *1, *r, *p;
f0d
            int sz, mi;
             node(int sa_, int lcp_, node* p_) : sa(sa_),
17b
   lcp(lcp_),
543
                 1(NULL), r(NULL), p(p_), sz(1), mi(lcp) {}
01e
             void update() {
58f
                 sz = 1, mi = lcp;
                 if (1) sz += 1->sz, mi = min(mi, 1->mi);
bd7
                 if (r) sz += r \rightarrow sz, mi = min(mi, r \rightarrow mi);
a54
cbb
            }
214
        }:
bb7
        node* root;
        vector<ll> tag; // tag of a suffix (reversed id)
295
        string s; // reversed
ac0
cf4
        dyn_sa() : root(NULL) {}
e45
        dyn_sa(string s_) : dyn_sa() {
ae4
             reverse(s_.begin(), s_.end());
519
             for (char c : s_) push_front(c);
cbb
a86
        \simdyn_sa() {
             vector < node *> q = {root};
609
402
             while (q.size()) {
                 node* x = q.back(); q.pop_back();
e5d
ee9
                 if (!x) continue;
1 c7
                 q.push_back(x->1), q.push_back(x->r);
                 delete x;
bf0
            }
cbb
cbb
        }
        int size(node* x) { return x ? x->sz : 0; }
73c
08e
        int mirror(int i) { return s.size()-1 - i; }
        bool cmp(int i, int j) {
580
            if (s[i] != s[j]) return s[i] < s[j];</pre>
a29
            if (i == 0 \text{ or } j == 0) \text{ return } i < j;
5b4
             return tag[i-1] < tag[j-1];</pre>
988
```

```
cbb
        void fix_path(node* x) { while (x) x->update(), x =
919
   x->p; }
        void flatten(vector < node * > & v , node * x) {
245
             if (!x) return;
8 c 8
e96
            flatten(v, x->1);
            v.push_back(x);
2a2
            flatten(v, x->r);
42 d
cbb
        }
964
        void build(vector<node*>& v, node*& x, node* p, int
   L, int R, 11 1, 11 r) {
            if (L > R) return void(x = NULL);
0.4 c
331
            int M = (L+R)/2;
3e3
            11 m = (1+r)/2;
7e5
            x = v[M]:
63e
            x - p = p;
            tag[x->sa] = m;
bb3
             build(v, x \rightarrow 1, x, L, M-1, 1, m-1), build(v,
ae0
   x - > r, x, M + 1, R, m + 1, r);
            x->update();
ca8
cbb
82f
        void fix(node*& x, node* p, ll l, ll r) {
             if (3*max(size(x->1), size(x->r)) \le 2*size(x))
7f0
   return x->update();
             vector < node *> v;
3d1
Осс
            flatten(v, x);
ea9
             build(v, x, p, 0, v.size()-1, 1, r);
cbb
b19
        node* next(node* x) {
728
            if (x->r) {
a91
                 x = x - > r;
347
                 while (x->1) x = x->1;
ea5
                 return x;
cbb
402
             while (x->p \text{ and } x->p->r == x) x = x->p;
137
             return x->p;
        }
cbb
b68
        node* prev(node* x) {
            if (x->1) {
e41
a26
                 x = x \rightarrow 1;
93 c
                 while (x->r) x = x->r;
ea5
                 return x;
```

```
cbb
6a1
            while (x->p \text{ and } x->p->l == x) x = x->p;
137
            return x->p;
cbb
4f7
        int get lcp(node* x, node* y) {
75a
            if (!x or !y) return 0; // change defaut value
   here
            if (s[x->sa] != s[y->sa]) return 0;
e51
843
            if (x->sa == 0 \text{ or } y->sa == 0) return 1;
4 d 0
            return 1 + query(mirror(x->sa-1),
   mirror(y->sa-1));
        }
cbb
        void add_suf(node*& x, node* p, int id, ll l, ll r) {
ad6
91e
            if (!x) {
8 e 3
                x = new node(id, 0, p);
8 e 2
                node *prv = prev(x), *nxt = next(x);
                int lcp_cur = get_lcp(prv, x), lcp_nxt =
65d
   get_lcp(x, nxt);
                if (nxt) nxt->lcp = lcp nxt, fix path(nxt);
ca3
                x->lcp = lcp_cur;
71f
7 b 4
                tag[id] = (1+r)/2;
ca8
                x->update();
505
                return;
cbb
            }
4a3
            if (cmp(id, x->sa)) add_suf(x->1, x, id, 1,
   tag[x->sa]-1);
c3a
            else add_suf(x->r, x, id, tag[x->sa]+1, r);
            fix(x, p, l, r);
3 db
cbb
ec2
        void push_front(char c) {
cc7
            s += c;
493
            tag.push_back(-1);
05e
            add_suf(root, NULL, s.size() - 1, 0, 1e18);
cbb
        }
7f3
        void rem_suf(node*& x, int id) {
6cf
            if (x->sa != id) {
                if (tag[id] < tag[x->sa]) return
864
   rem suf(x->1, id);
e6f
                return rem_suf(x->r, id);
cbb
            }
```

```
2cf
             node* nxt = next(x):
            if (nxt) nxt -> lcp = min(nxt -> lcp, x -> lcp),
09b
   fix_path(nxt);
b20
             node *p = x->p, *tmp = x;
f3f
            if (!x->1 \text{ or } !x->r) {
2fd
                 x = x - > 1 ? x - > 1 : x - > r;
753
                 if (x) x->p = p;
9d9
            } else {
                 for (tmp = x->1, p = x; tmp->r; tmp =
7 f 7
   tmp -> r) p = tmp;
                 x->sa = tmp->sa, x->lcp = tmp->lcp;
f2a
482
                 if (tmp -> 1) tmp -> 1 -> p = p;
                 if (p->1 == tmp) p->1 = tmp->1;
14 c
a94
                 else p -> r = tmp -> 1;
cbb
            fix_path(p);
b5e
7 c 3
             delete tmp;
cbb
15b
        void pop_front() {
abe
             if (!s.size()) return;
342
             s.pop_back();
436
             rem_suf(root, s.size());
c6e
             tag.pop_back();
cbb
        }
d41
530
        int query(node* x, ll l, ll r, ll a, ll b) {
            if (!x \text{ or } tag[x->sa] == -1 \text{ or } r < a \text{ or } b < 1)
e51
   return s.size();
            if (a <= 1 and r <= b) return x->mi;
ef5
            int ans = s.size();
8eb
            if (a \le tag[x->sa]  and tag[x->sa] \le b) ans =
e1f
   min(ans, x->lcp);
             ans = min(ans, query(x->1, 1, tag[x->sa]-1, a,
d99
   b));
             ans = min(ans, query(x->r, tag[x->sa]+1, r, a,
261
   b));
ba7
             return ans;
cbb
        int query(int i, int j) { // lcp(s[i..], s[j..])
588
209
             if (i == j) return s.size() - i;
29 e
             ll a = tag[mirror(i)], b = tag[mirror(j)];
```

```
710
            int ret = query(root, 0, 1e18, min(a, b)+1.
   max(a, b));
edf
            return ret;
cbb
        // optional: get rank[i], sa[i] and lcp[i]
d41
044
        int rank(int i) {
396
            i = mirror(i);
52f
            node* x = root;
7 c9
            int ret = 0;
f4c
            while (x) {
33e
                if (tag[x->sa] < tag[i]) {</pre>
f9d
                    ret += size(x->1)+1;
a 9 1
                    x = x - > r;
                } else x = x - > 1;
eb5
cbb
            }
edf
            return ret;
cbb
649
        pair<int, int> operator[](int i) {
52f
            node* x = root;
31e
            while (1) {
d4d
                if (i < size(x->1)) x = x->1;
4e6
                else {
85f
                    i = size(x->1);
e03
                    if (!i) return {mirror(x->sa), x->lcp};
040
                    i--, x = x->r;
cbb
cbb
            }
cbb
214 };
6.13 Trie
// trie T() constroi uma trie para o alfabeto das letras
   minusculas
// trie T(tamanho do alfabeto, menor caracter) também pode
   ser usado
//
// T.insert(s) - 0(|s|*sigma)
// T.erase(s) - O(|s|)
// T.find(s) retorna a posicao, 0 se nao achar - O(|s|)
// T.count_pref(s) numero de strings que possuem s como
   prefixo - O(|s|)
```

```
//
// Nao funciona para string vazia
// 979609
ab5 struct trie {
e1a
        vector < vector < int >> to;
450
        vector<int> end, pref;
        int sigma; char norm;
af0
        trie(int sigma_=26, char norm_='a') : sigma(sigma_),
bb1
   norm(norm_) {
            to = {vector < int > (sigma)};
58a
86e
            end = \{0\}, pref = \{0\};
        }
cbb
64e
        void insert(string s) {
c67
            int x = 0:
7e7
            for(auto c : s) {
800
                 int &nxt = to[x][c-norm];
                 if(!nxt) {
dd7
0aa
                     nxt = to.size();
                     to.push_back(vector < int > (sigma));
526
                     end.push_back(0), pref.push_back(0);
770
cbb
                 }
827
                 x = nxt, pref[x]++;
cbb
e4e
             end[x]++;
cbb
6b2
        void erase(string s) {
c67
            int x = 0;
            for(char c : s) {
b4f
800
                 int & nxt = to[x][c-norm];
10 c
                 x = nxt, pref[x] - -;
d8e
                 if(!pref[x]) nxt = 0;
            }
cbb
             end[x]--:
bf0
cbb
        int find(string s) {
aee
            int x = 0;
c67
7e7
            for(auto c : s) {
                 x = to[x][c-norm];
2ec
a66
                 if(!x) return 0;
cbb
ea5
            return x;
```

```
cbb }
839    int count_pref(string s) {
e2f        return pref[find(s)];
cbb }
214 };
```

## 7 Primitivas

#### 7.1 Aritmetica Modular

```
// O mod tem q ser primo
// 5a6efb
429 template <int p> struct mod_int {
        ll pow(ll b, ll e) {
02c
a63
           if (e == 0) return 1;
630
           11 r = pow(b*b\%p, e/2);
475
           if (e\%2 == 1) r = (r*b)\%p;
4c1
            return r;
cbb
ae3
        11 inv(11 b) { return pow(b, p-2); }
4d7
        using m = mod_int;
d93
        int v;
fe0
        mod_int() : v(0) {}
e12
        mod_int(ll v_) {
            if (v_ >= p or v_ <= -p) v_ %= p;
019
bc6
            if (v_{-} < 0) v_{-} += p;
2e7
            v = v_{-};
cbb
74d
        m& operator+=(const m &a) {
2fd
            v += a.v;
ba5
            if (v >= p) v -= p;
357
            return *this:
cbb
        m& operator -= (const m &a) {
eff
8 b 4
            v -= a.v;
            if (v < 0) v += p;
cc8
357
            return *this;
cbb
```

```
4c4
        m& operator*=(const m &a) {
            v = v * ll(a.v) % p;
8a5
357
            return *this;
cbb
3f9
        m& operator/=(const m &a) {
            v = v * inv(a.v) % p;
5d6
357
            return *this;
cbb
        m operator-() { return m(-v); }
d65
        m& operator^=(ll e) {
b3e
06 d
            if (e < 0){
6e2
                v = inv(v);
0.0 c
                e = -e:
cbb
ebf
            v = pow(v, e\%(p-1));
357
            return *this;
cbb
        }
423
        bool operator == (const m &a) { return v == a.v; }
69 f
        bool operator!=(const m &a) { return v != a.v; }
1c6
        friend istream & operator >> (istream & in, m& a) {
d1c
            11 val; in >> val;
            a = m(val);
d48
091
            return in;
cbb
        friend ostream &operator << (ostream &out, m a) {</pre>
44f
5a0
            return out << a.v;</pre>
cbb
399
        friend m operator+(m a, m b) { return a+=b; }
f9e
        friend m operator-(m a, m b) { return a-=b; }
        friend m operator*(m a, m b) { return a*=b; }
9c1
        friend m operator/(m a, m b) { return a/=b; }
51b
08f
        friend m operator^(m a, ll e) { return a^=e; }
214 };
055 typedef mod_int <(int)1e9+7> mint;
7.2 Big Integer
// Complexidades: (para n digitos)
// Soma, subtracao, comparacao - O(n)
// Multiplicacao - O(n log(n))
```

```
// Divisao, resto - O(n^2)
864 struct bint {
669
        static const int BASE = 1e9;
990
        vector < int > v;
3bd
        bool neg;
609
        bint() : neg(0) \{ \}
        bint(int val) : bint() { *this = val; }
d53
        bint(long long val) : bint() { *this = val; }
e8f
a0f
        void trim() {
            while (v.size() and v.back() == 0) v.pop_back();
f42
            if (!v.size()) neg = 0;
df8
cbb
d41
        // converter de/para string | cin/cout
294
        bint(const char* s) : bint() {
   from_string(string(s)); }
        bint(const string& s) : bint() { from_string(s); }
548
        void from_string(const string& s) {
4ab
            v.clear(), neg = 0;
0a6
            int ini = 0;
d72
            while (ini < s.size() and (s[ini] == '-' or</pre>
8e2
   s[ini] == '+' or s[ini] == '0'))
                if (s[ini++] == '-') neg = 1;
71d
           for (int i = s.size()-1; i >= ini; i -= 9) {
883
                int at = 0;
05е
                for (int j = max(ini, i - 8); j <= i; j++)</pre>
5 b 1
   at = 10*at + (s[j]-'0');
1fd
                v.push_back(at);
cbb
df8
            if (!v.size()) neg = 0;
cbb
2ff
        string to_string() const {
            if (!v.size()) return "0";
8be
793
            string ret;
73e
            if (neg) ret += '-';
3 e 9
            for (int i = v.size()-1; i >= 0; i--) {
                string at = ::to_string(v[i]);
582
                int add = 9 - at.size();
ced
                if (i+1 < v.size()) for (int j = 0; j < add;
75e
```

```
i++) ret += '0':
f9f
                ret += at;
cbb
edf
            return ret;
cbb
d2f
        friend istream& operator>>(istream& in, bint& val) {
eb6
            string s; in >> s;
            val = s;
966
091
            return in;
cbb
99 d
        friend ostream& operator << (ostream& out, const bint&
   val) {
8b9
            string s = val.to_string();
396
            out << s;
fe8
            return out:
cbb
        }
        // operators
d41
60a
        friend bint abs(bint val) {
            val.neg = 0;
c5f
            return val;
d94
cbb
bee
        friend bint operator-(bint val) {
815
            if (val != 0) val.neg ^= 1;
            return val;
d94
cbb
        bint& operator=(const bint& val) { v = val.v, neg =
41f
   val.neg; return *this; }
        bint& operator=(long long val) {
249
0a6
            v.clear(), neg = 0;
3a6
            if (val < 0) neg = 1, val *= -1;
            for (; val; val /= BASE) v.push_back(val % BASE);
fdc
357
            return *this;
cbb
3bd
        int cmp(const bint& r) const { // menor: -1 | igual:
  0 | maior: 1
b14
            if (neg != r.neg) return neg ? -1 : 1;
0bb
            if (v.size() != r.v.size()) {
                int ret = v.size() < r.v.size() ? -1 : 1;</pre>
ff7
                return neg ? -ret : ret;
91b
cbb
            for (int i = int(v.size())-1; i >= 0; i--) {
478
```

```
405
                if (v[i] != r.v[i]) {
2e5
                    int ret = v[i] < r.v[i] ? -1 : 1;</pre>
91b
                    return neg ? -ret : ret;
cbb
                }
            }
cbb
bb3
            return 0;
cbb
152
        friend bool operator < (const bint& 1, const bint& r)
   { return l.cmp(r) == -1; }
        friend bool operator > (const bint& 1, const bint& r)
   { return 1.cmp(r) == 1; }
        friend bool operator <= (const bint& 1, const bint& r)</pre>
   { return 1.cmp(r) <= 0; }
        friend bool operator >= (const bint& 1, const bint& r)
   { return 1.cmp(r) >= 0; }
        friend bool operator == (const bint % 1, const bint % r)
   { return 1.cmp(r) == 0; }
        friend bool operator!=(const bint& 1, const bint& r)
   { return 1.cmp(r) != 0; }
38e
        bint& operator +=(const bint& r) {
6bf
            if (!r.v.size()) return *this;
            if (neg != r.neg) return *this -= -r;
a93
            for (int i = 0, c = 0; i < r.v.size() or c; i++)</pre>
256
  {
                if (i == v.size()) v.push back(0);
e28
08f
                v[i] += c + (i < r.v.size() ? r.v[i] : 0);
                if ((c = v[i] >= BASE)) v[i] -= BASE;
baa
cbb
            }
357
            return *this;
cbb
54c
        friend bint operator+(bint a, const bint& b) {
   return a += b; }
9 c8
        bint& operator -=(const bint& r) {
6bf
            if (!r.v.size()) return *this;
524
            if (neg != r.neg) return *this += -r;
            if ((!neg and *this < r) or (neg and r < *this))</pre>
358
  {
b10
                *this = r - *this;
                neg ^= 1;
a10
357
                return *this;
cbb
            }
```

```
256
            for (int i = 0, c = 0; i < r.v.size() or c; i++)
  {
9ef
                v[i] = c + (i < r.v.size() ? r.v[i] : 0);
c8c
                if ((c = v[i] < 0)) v[i] += BASE;
            }
cbb
0eb
            trim();
357
            return *this;
cbb
f44
        friend bint operator-(bint a, const bint& b) {
   return a -= b; }
        // operators de * / %
d41
        bint& operator *=(int val) {
6b0
            if (val < 0) val *= -1, neg ^= 1;</pre>
bca
            for (int i = 0, c = 0; i < v.size() or c; i++) {
566
e28
                if (i == v.size()) v.push_back(0);
352
                long long at = (long long) v[i] * val + c;
6a3
                v[i] = at % BASE;
                c = at / BASE;
b3d
            }
cbb
0eb
            trim();
357
            return *this;
cbb
480
        friend bint operator *(bint a, int b) { return a *=
  b; }
        friend bint operator *(int a, bint b) { return b *=
d5c
a; }
13b
        using cplx = complex <double >;
        void fft(vector < cplx > & a, bool f, int N,
bfb
   vector < int > & rev) const {
            for (int i = 0; i < N; i++) if (i < rev[i])
bc7
   swap(a[i], a[rev[i]]);
            vector < cplx > roots(N);
bad
            for (int n = 2: n <= N: n *= 2) {
192
                const static double PI = acos(-1);
4e9
71a
                for (int i = 0; i < n/2; i++) {
                     double alpha = (2*PI*i)/n;
40 d
1a1
                    if (f) alpha = -alpha;
3f6
                     roots[i] = cplx(cos(alpha), sin(alpha));
cbb
3e9
                for (int pos = 0; pos \langle N; pos += n \rangle
                    for (int l = pos, r = pos+n/2, m = 0; m
898
```

```
< n/2; 1++, r++, m++) 
297
                         auto t = roots[m]*a[r];
                         a[r] = a[1] - t;
254
                         a[1] = a[1] + t;
b8f
cbb
                     }
cbb
3f1
            if (!f) return;
            auto invN = cplx(1)/cplx(N);
08b
873
            for (int i = 0; i < N; i++) a[i] *= invN;
cbb
0e0
        vector <long long > convolution(const vector <int > & a,
   const vector<int>& b) const {
            vector < cplx > l(a.begin(), a.end()), r(b.begin(),
ff9
   b.end());
996
            int ln = l.size(), rn = r.size(), N = ln+rn+1, n
   = 1, log_n = 0;
821
            while (n \le N) n \le 1, \log_n + +;
            vector < int > rev(n);
808
            for (int i = 0; i < n; i++) {</pre>
603
434
                rev[i] = 0;
f44
                for (int j = 0; j < log_n; j++) if (i >> j & 1)
4ff
                     rev[i] = 1 << (log n-1-j);
cbb
            }
230
            l.resize(n), r.resize(n);
a89
            fft(l, false, n, rev), fft(r, false, n, rev);
            for (int i = 0; i < n; i++) l[i] *= r[i];
917
            fft(1, true, n, rev);
88b
7ae
            vector < long long > ret;
            for (auto& i : 1) ret.push_back(round(i.real()));
c14
edf
            return ret;
cbb
633
        vector < int > convert_base(const vector < int > & a, int
   from, int to) const {
498
            static vector <long long > pot(10, 1);
671
            if (pot[1] == 1) for (int i = 1; i < 10; i++)
   pot[i] = 10*pot[i-1];
            vector < int > ret;
4b8
156
            long long at = 0;
            int digits = 0;
608
941
            for (int i : a) {
                at += i * pot[digits];
412
035
                digits += from;
```

```
684
                while (digits >= to) {
0c8
                    ret.push_back(at % pot[to]);
cf9
                    at /= pot[to];
                    digits -= to;
fd4
                }
cbb
            }
cbb
944
            ret.push_back(at);
            while (ret.size() and ret.back() == 0)
384
   ret.pop_back();
edf
            return ret;
cbb
        bint operator*(const bint& r) const { // O(n log(n))
edb
            bint ret:
2af
            ret.neg = neg ^ r.neg;
968
d5d
            auto conv = convolution(convert_base(v, 9, 4),
   convert_base(r.v, 9, 4));
           long long c = 0;
a0e
            for (auto i : conv) {
a 74
               long long at = i+c;
f6d
               ret.v.push_back(at % 10000);
4cb
               c = at / 10000;
a25
cbb
            }
3cb
            for (; c; c /= 10000) ret.v.push_back(c%10000);
0e2
            ret.v = convert_base(ret.v, 4, 9);
25 c
           if (!ret.v.size()) ret.neg = 0;
edf
            return ret;
cbb
359
        bint& operator*=(const bint& r) { return *this =
   *this * r; };
        bint& operator/=(int val) {
9a3
           if (val < 0) neg ^{-} 1, val *= -1;
d9a
           for (int i = int(v.size())-1, c = 0; i >= 0;
f18
  i--) {
                long long at = v[i] + c * (long long) BASE;
2a7
e02
               v[i] = at / val:
fb1
                c = at % val:
           }
cbb
0eb
            trim();
357
            return *this;
cbb
e74
        friend bint operator/(bint a, int b) { return a /=
   b; }
```

```
4a9
        int operator %=(int val) {
23b
            if (val < 0) val *= -1;
156
            long long at = 0;
            for (int i = int(v.size())-1; i >= 0; i--)
f31
1 b3
                at = (BASE * at + v[i]) \% val;
d22
            if (neg) at *=-1;
ce6
            return at;
cbb
2fb
        friend int operator % (bint a, int b) { return a % = b;
 }
13b
        friend pair <bint, bint > divmod(const bint & a_, const
   bint& b_) { // O(n^2)
            if (a_ == 0) return {0, 0};
611
            int norm = BASE / (b_.v.back() + 1);
d8a
            bint a = abs(a_) * norm;
b4e
            bint b = abs(b<sub>_</sub>) * norm;
027
14d
            bint q, r;
c91
            for (int i = a.v.size() - 1; i >= 0; i--) {
                r *= BASE, r += a.v[i];
b71
                long long upper = b.v.size() < r.v.size() ?</pre>
4ff
   r.v[b.v.size()] : 0;
86d
                int lower = b.v.size() - 1 < r.v.size() ?</pre>
   r.v[b.v.size() - 1] : 0;
                int d = (upper * BASE + lower) / b.v.back();
431
5 d4
                r \rightarrow b*d;
                while (r < 0) r += b, d--; // roda O(1) vezes
30f
738
                q.v.push_back(d);
cbb
a48
            reverse(q.v.begin(), q.v.end());
ae2
            q.neg = a_.neg ^ b_.neg;
            r.neg = a_.neg;
88b
8 e 5
            q.trim(), r.trim();
0ef
            return {q, r / norm};
cbb
        bint operator/(const bint& val) { return
1 d8
   divmod(*this, val).first; }
        bint& operator/=(const bint& val) { return *this =
7 f 9
   *this / val: }
        bint operator%(const bint& val) { return
1f9
   divmod(*this, val).second; }
        bint& operator%=(const bint& val) { return *this =
df5
   *this % val; }
```

#### 7.3 Matroid

```
// Matroids de Grafo e Particao
// De modo geral, toda Matroid contem um build() linear
// e uma funcao constante oracle()
// oracle(i) responde se o conjunto continua independente
// apos adicao do elemento i
// oracle(i, j) responde se o conjunto continua indepente
// apos trocar o elemento i pelo elemento j
//
// Intersecao sem peso O(r^2 n)
// em que n eh o tamanho do conjunto e r eh o tamanho da
   resposta
// Matroid Grafica
// Matroid das florestas de um grafo
// Um conjunto de arestas eh independente se formam uma
   floresta
//
// build() : O(n)
// oracle() : O(1)
// 691847
fda struct graphic_matroid {
        int n, m, t;
5da
        vector<array<int, 2>> edges;
32 c
        vector < vector < int >> g;
789
        vector<int> comp, in, out;
62e
513
        graphic_matroid(int n_, vector<array<int, 2>> edges_)
a1f
            : n(n_{-}), m(edges_{-}.size()), edges(edges_{-}), g(n),
   comp(n), in(n), out(n) {}
        void dfs(int u) {
315
            in[u] = t++;
ab8
            for (auto v : g[u]) if (in[v] == -1)
17 d
                comp[v] = comp[u], dfs(v);
863
677
            out[u] = t;
cbb
945
        void build(vector<int> I) {
a34
            t = 0;
741
            for (int u = 0; u < n; u++) g[u].clear(), in[u]
```

```
= -1:
667
            for (int e : I) {
                auto [u, v] = edges[e];
d00
                g[u].push_back(v), g[v].push_back(u);
125
cbb
809
            for (int u = 0; u < n; u++) if (in[u] == -1)
                comp[u] = u, dfs(u);
a7d
cbb
f31
        bool is_ancestor(int u, int v) {
a68
            return in[u] <= in[v] and in[v] < out[u];</pre>
cbb
e6b
        bool oracle(int e) {
            return comp[edges[e][0]] != comp[edges[e][1]];
453
cbb
f75
        bool oracle(int e, int f) {
574
            if (oracle(f)) return true;
622
            int u = edges[e][in[edges[e][0]] <</pre>
   in[edges[e][1]]];
            return is_ancestor(u, edges[f][0]) !=
ff2
   is_ancestor(u, edges[f][1]);
cbb
214 };
    // Matroid de particao ou cores
   // Um conjunto eh independente se a quantidade de
       elementos
    // de cada cor nao excede a capacidade da cor
    // Quando todas as capacidades sao 1, um conjunto eh
       independente
    // se todas as suas cores sao distintas
    //
    // build() : O(n)
    // oracle() : O(1)
    // caa72a
994 struct partition_matroid {
        vector < int > cap, color, d;
501
608
        partition_matroid(vector<int> cap_, vector<int>
   color )
04d
            : cap(cap_), color(color_), d(cap.size()) {}
        void build(vector<int> I) {
945
def
            fill(d.begin(), d.end(), 0);
```

```
e9d
            for (int u : I) d[color[u]]++;
cbb
514
        bool oracle(int u) {
            return d[color[u]] < cap[color[u]];</pre>
0a1
cbb
        }
f7f
        bool oracle(int u, int v) {
            return color[u] == color[v] or oracle(v);
2f7
cbb
214 };
    // Intersecao de matroid sem pesos
    // Dadas duas matroids M1 e M2 definidas sobre o mesmo
    // conjunto I, retorna o maior subconjunto de I
    // que eh independente tanto para M1 quanto para M2
    // O(r^2*n)
    // 899f94
    // Matroid "pesada" deve ser a M2
132 template < typename Matroid1, typename Matroid2 >
801 vector <int> matroid_intersection(int n, Matroid1 M1,
   Matroid2 M2) {
f5b
        vector < bool > b(n);
        vector < int > I[2];
a64
        bool converged = false;
a8b
        while (!converged) {
0 c 1
742
            I[0].clear(), I[1].clear();
            for (int u = 0; u < n; u++) I[b[u]].push_back(u);</pre>
99 d
09 d
            M1.build(I[1]), M2.build(I[1]);
            vector < bool > target(n), pushed(n);
289
26a
            queue < int > q;
            for (int u : I[0]) {
5 c 5
2b2
                 target[u] = M2.oracle(u);
c1b
                if (M1.oracle(u)) pushed[u] = true,
   q.push(u);
cbb
3fe
            vector < int > p(n, -1);
            converged = true;
07a
            while (q.size()) {
402
                int u = q.front(); q.pop();
be1
5c6
                if (target[u]) {
```

```
101
                    converged = false:
                    for (int v = u; v != -1; v = p[v]) b[v]
c32
   = !b[v];
c2b
                    break;
cbb
e78
                for (int v : I[!b[u]]) if (!pushed[v]) {
                    if ((b[u] and M1.oracle(u, v)) or (b[v]
34d
   and M2.oracle(v, u)))
                         p[v] = u, pushed[v] = true,
bae
   q.push(v);
cbb
                }
            }
cbb
cbb
b68
        return I[1];
cbb }
    // Intersecao de matroid com pesos
    // Dadas duas matroids M1 e M2 e uma funcao de pesos w,
       todas definidas sobre
    // um conjunto I retorna o maior subconjunto de I
       (desempatado pelo menor peso)
    // que eh independente tanto para M1 quanto para M2
   // A resposta eh construida incrementando o tamanho
       conjunto I de 1 em 1
    // Se nao tiver custo negativo, nao precisa de SPFA
    //
    // O(r^3*n) com SPFA
    // O(r^2*n*log(n)) com Dijkstra e potencial
    // 3a09d1
42a template < typename T, typename Matroid1, typename
   Matroid2>
2b5 vector <int > weighted_matroid_intersection(int n,
   vector <T> w, Matroid1 M1, Matroid2 M2) {
        vector <bool> b(n), target(n), is_inside(n);
6 c 9
563
        vector < int > I[2], from(n);
        vector < pair < T, int >> d(n);
e35
169
        auto check_edge = [&](int u, int v) {
            return (b[u] and M1.oracle(u, v)) or (b[v] and
   M2.oracle(v, u));
214
        };
667
        while (true) {
```

```
742
            I[0].clear(), I[1].clear();
            for (int u = 0; u < n; u++) I[b[u]].push_back(u);</pre>
99 d
            // I[1] contem o conjunto de tamanho I[1].size()
d41
   de menor peso
            M1.build(I[1]), M2.build(I[1]);
09 d
687
            for (int u = 0; u < n; u++) {
                 target[u] = false, is_inside[u] = false,
ea5
   from[u] = -1;
                 d[u] = {numeric_limits < T >:: max(), INF};
961
cbb
8d3
            deque <T> q;
476
            sort(I[0].begin(), I[0].end(), [&](int i, int
   j) { return w[i] < w[j]; });</pre>
            for (int u : I[0]) {
5 c 5
                 target[u] = M2.oracle(u);
2b2
5a7
                 if (M1.oracle(u)) {
4ef
                     if (is_inside[u]) continue;
7cc
                     d[u] = \{w[u], 0\};
427
                     if (!q.empty() and d[u] > d[q.front()])
   q.push_back(u);
655
                     else q.push_front(u);
4ae
                     is inside[u] = true;
                }
cbb
cbb
402
            while (q.size()) {
97a
                 int u = q.front(); q.pop_front();
6f3
                 is_inside[u] = false;
57a
                 for (int v : I[!b[u]]) if (check_edge(u, v))
   {
                     pair < T, int > nd(d[u].first + w[v],
9de
   d[u].second + 1);
61b
                     if (nd < d[v]) {</pre>
                         from[v] = u, d[v] = nd;
6ac
                         if (is_inside[v]) continue;
bd7
                         if (q.size() and d[v] >
eec
   d[q.front()]) q.push_back(v);
275
                         else q.push_front(v);
                         is_inside[v] = true;
587
                     }
cbb
                 }
cbb
cbb
cc8
            pair < T , int > mini =
```

```
pair(numeric_limits <T>::max(), INF);
489
            int targ = -1;
259
            for (int u : I[0]) if (target[u] and d[u] < mini)</pre>
                mini = d[u], targ = u;
2b9
            if (targ != -1) for (int u = targ; u != -1; u =
e14
   from[u])
d89
                b[u] = !b[u], w[u] *= -1;
f97
            else break:
cbb
b68
        return I[1];
cbb }
```

#### 7.4 Primitivas de fração

```
// Funciona com o Big Int
// cdb445
a4e template < typename T = int > struct frac {
a40
        T num, den;
        template < class U, class V>
e3f
61d
        frac(U num_ = 0, V den_ = 1) : num(num_), den(den_) {
            assert(den != 0);
bad
            if (den < 0) num *= -1, den *= -1;
583
a51
            T g = gcd(abs(num), den);
572
            num \neq g, den \neq g;
cbb
        friend bool operator <(const frac& 1, const frac& r) {</pre>
51f
            return l.num * r.den < r.num * l.den;</pre>
fa0
cbb
4 b 5
        friend frac operator+(const frac& 1, const frac& r) {
b61
            return {1.num*r.den + 1.den*r.num, 1.den*r.den};
cbb
74d
        friend frac operator - (const frac& 1, const frac& r) {
            return {1.num*r.den - 1.den*r.num, 1.den*r.den};
2cd
cbb
        friend frac operator*(const frac& 1, const frac& r) {
c80
510
            return {1.num*r.num, 1.den*r.den};
cbb
        friend frac operator/(const frac& 1, const frac& r) {
a1b
8f3
            return {1.num*r.den, 1.den*r.num};
cbb
```

# 7.5 Primitivas de matriz - exponenciacao

```
// d05c24
945 #define MODULAR false
5ed template < typename T > struct matrix : vector < vector < T >> {
       int n, m;
14e
30f
       void print() {
            for (int i = 0; i < n; i++) {
603
70f
                for (int j = 0; j < m; j++) cout <<
   (*this)[i][j] << " ";
1fb
                 cout << endl;</pre>
            }
cbb
       }
cbb
        matrix(int n_, int m_, bool ident = false) :
aa3
                 vector < vector < T > (n , vector < T > (m , 0)),
b14
   n(n), m(m)
94e
            if (ident) {
               assert(n == m);
df7
                for (int i = 0; i < n; i++) (*this)[i][i] =
            }
cbb
cbb
        matrix(const vector < vector < T >> & c) :
b83
   vector < vector < T >> (c),
            n(c.size()), m(c[0].size()) {}
a3d
        matrix(const initializer_list < initializer_list < T >> &
efc
   c) {
f7e
            vector < vector < T >> val;
            for (auto& i : c) val.push_back(i);
212
303
            *this = matrix(val);
cbb
388
        matrix<T> operator*(matrix<T>& r) {
```

```
1 e2
            assert(m == r.n);
82c
            matrix <T> M(n, r.m);
d69
            for (int i = 0; i < n; i++) for (int k = 0; k < 0
   m; k++)
df4
                for (int j = 0; j < r.m; j++) {
                     T \text{ add} = (*this)[i][k] * r[k][j];
e34
f98 #if MODULAR
d41 #warning Usar matrix<11> e soh colocar valores em [0,
   MOD) na matriz!
8 b 6
                     M[i][j] += add%MOD;
                     if (M[i][j] >= MOD) M[i][j] -= MOD;
983
8c1 #else
7 bb
                     M[i][j] += add;
f2e #endif
                 }
cbb
474
            return M;
cbb
528
        matrix <T> operator^(ll e) {
            matrix <T> M(n, n, true), at = *this;
f 10
c87
            while (e) {
2e2
               if (e\&1) M = M*at;
cc2
                e >>= 1;
c80
                at = at*at;
cbb
474
            return M;
cbb
582
        void apply_transform(matrix M, ll e){
1 c3
            auto& v = *this;
c87
            while (e) {
9ba
                if (e\&1) v = M*v;
cc2
                e >>= 1;
419
                M = M * M;
            }
cbb
        }
cbb
214 };
```

### 7.6 Primitivas Geometricas

```
c83 typedef double ld;
e3b const ld DINF = 1e18;
43a const ld pi = acos(-1.0);
107 const ld eps = 1e-9;
```

```
b32 #define sq(x) ((x)*(x))
d97 bool eq(ld a, ld b) {
        return abs(a - b) <= eps;</pre>
cbb }
   // a8b7d6
b2a struct pt { // ponto
c1e
       ld x, y;
       pt(ld x_{-} = 0, ld y_{-} = 0) : x(x_{-}), y(y_{-}) {}
3dd
       bool operator < (const pt p) const {</pre>
5bc
           if (!eq(x, p.x)) return x < p.x;
059
           if (!eq(y, p.y)) return y < p.y;
f98
bb3
           return 0:
cbb
        }
a83
        bool operator == (const pt p) const {
            return eq(x, p.x) and eq(y, p.y);
ed0
cbb
        pt operator + (const pt p) const { return pt(x+p.x,
cb9
  y+p.y); }
        pt operator - (const pt p) const { return pt(x-p.x,
a24
  y-p.y); }
      pt operator * (const ld c) const { return pt(x*c ,
   y*c ); }
    pt operator / (const ld c) const { return pt(x/c ,
   y/c ); }
3b6 ld operator * (const pt p) const { return x*p.x +
   y*p.y; }
6df
       ld operator ^ (const pt p) const { return x*p.y -
  y*p.x; }
5ed
        friend istream& operator >> (istream& in, pt& p) {
           return in >> p.x >> p.y;
e37
        }
cbb
214 };
    // 7ab617
b3a struct line { // reta
730
       pt p, q;
       line() {}
0d6
       line(pt p_, pt q_) : p(p_), q(q_) {}
4b8
       friend istream& operator >> (istream& in, line& r) {
8d7
```

```
4cb
          return in >> r.p >> r.q;
cbb }
214 };
   // PONTO & VETOR
   // c684fb
364 ld dist(pt p, pt q) { // distancia
cbb }
  // 80f2b6
9d7 ld dist2(pt p, pt q) { // quadrado da distancia
       return sq(p.x - q.x) + sq(p.y - q.y);
cbb }
  // cf7f33
483 ld norm(pt v) { // norma do vetor
       return dist(pt(0, 0), v);
490
cbb }
   // 404df7
589 ld angle(pt v) { // angulo do vetor com o eixo x
1d ang = atan2(v.y, v.x);
    if (ang < 0) ang += 2*pi;</pre>
6f8
19 c
       return ang;
cbb }
   // 1b1d4a
298 ld sarea(pt p, pt q, pt r) { // area com sinal
       return ((q-p)^{(r-q)})/2;
606
cbb }
   // 98c42f
e32 bool col(pt p, pt q, pt r) { // se p, q e r sao colin.
e7d
       return eq(sarea(p, q, r), 0);
cbb }
   // 85d09d
Ocd bool ccw(pt p, pt q, pt r) { // se p, q, r sao ccw
fa7 return sarea(p, q, r) > eps;
cbb }
```

```
// 41a7b4
1ef pt rotate(pt p, ld th) { // rotaciona o ponto th radianos
       return pt(p.x * cos(th) - p.y * sin(th),
               p.x * sin(th) + p.y * cos(th));
ff1
cbb }
    // e4ad5e
ab1 pt rotate90(pt p) { // rotaciona 90 graus
       return pt(-p.y, p.x);
a0d
cbb }
    // RETA
    // 0fb984
edc bool isvert(line r) { // se r eh vertical
       return eq(r.p.x, r.q.x);
cbb }
   // 726d68
099 bool isinseg(pt p, line r) { // se p pertence ao seg de r
       pt a = r.p - p, b = r.q - p;
f65
       return eq((a \hat{b}), 0) and (a * b) < eps;
b04
cbb }
    // a0a30b
98d ld get_t(pt v, line r) { // retorna t tal que t*v
   pertence a reta r
       return (r.p^r.q) / ((r.p-r.q)^v);
6ee
cbb }
    // 2329fe
256 pt proj(pt p, line r) { // projecao do ponto p na reta r
bea if (r.p == r.q) return r.p;
      r.q = r.q - r.p; p = p - r.p;
97a
       pt proj = r.q * ((p*r.q) / (r.q*r.q));
9f8
2 \, c \, d
       return proj + r.p;
cbb }
    // 111fd2
d5c pt inter(line r, line s) { // r inter s
if (eq((r.p - r.q) ^ (s.p - s.q), 0)) return
```

```
pt(DINF, DINF);
       r.q = r.q - r.p, s.p = s.p - r.p, s.q = s.q - r.p;
205
       return r.q * get_t(r.q, s) + r.p;
543
cbb }
   // 35998c
676 bool interseg(line r, line s) { // se o seg de r
   intersecta o seg de s
       if (isinseg(r.p, s) or isinseg(r.q, s)
19b
            or isinseg(s.p, r) or isinseg(s.q, r)) return 1;
c21
9fa
        return ccw(r.p, r.q, s.p) != ccw(r.p, r.q, s.q) and
                ccw(s.p, s.q, r.p) != ccw(s.p, s.q, r.q);
413
cbb }
   // 1b72e1
fcb ld disttoline(pt p, line r) { // distancia do ponto a
   reta
        return 2 * abs(sarea(p, r.p, r.q)) / dist(r.p, r.q);
89a
cbb }
   // 3679c0
bcc ld disttoseg(pt p, line r) { // distancia do ponto ao seg
       if ((r.q - r.p)*(p - r.p) < 0) return dist(r.p, p);
       if ((r.p - r.q)*(p - r.q) < 0) return dist(r.q, p);
951
       return disttoline(p, r);
a19
cbb }
    // 222358
11d ld distseg(line a, line b) { // distancia entre seg
4df
        if (interseg(a, b)) return 0;
349
        ld ret = DINF;
       ret = min(ret, disttoseg(a.p, b));
341
       ret = min(ret, disttoseg(a.q, b));
ceb
        ret = min(ret, disttoseg(b.p, a));
093
        ret = min(ret, disttoseg(b.q, a));
448
edf
        return ret;
cbb }
    // POLIGONO
```

```
// corta poligono com a reta r deixando os pontos p tal
    // ccw(r.p, r.q, p)
    // 2538f9
1a9 vector<pt> cut_polygon(vector<pt> v, line r) { // O(n)
        vector<pt> ret;
8af
        for (int j = 0; j < v.size(); j++) {</pre>
8a4
            if (ccw(r.p, r.q, v[j])) ret.push_back(v[j]);
dac
            if (v.size() == 1) continue;
dce
            line s(v[j], v[(j+1)\%v.size()]);
030
            pt p = inter(r, s);
ae3
            if (isinseg(p, s)) ret.push_back(p);
a3d
cbb
8a1
        ret.erase(unique(ret.begin(), ret.end()), ret.end());
        if (ret.size() > 1 and ret.back() == ret[0])
24 d
   ret.pop_back();
        return ret;
edf
cbb }
    // distancia entre os retangulos a e b (lados paralelos
       aos eixos)
    // assume que ta representado (inferior esquerdo,
       superior direito)
    // 630253
5f5 ld dist_rect(pair<pt, pt> a, pair<pt, pt> b) {
080
        1d hor = 0, vert = 0;
34b
        if (a.second.x < b.first.x) hor = b.first.x -</pre>
   a.second.x;
        else if (b.second.x < a.first.x) hor = a.first.x -</pre>
f5f
   b.second.x:
4fd
        if (a.second.y < b.first.y) vert = b.first.y -</pre>
   a.second.y;
        else if (b.second.y < a.first.y) vert = a.first.y -</pre>
80a
   b.second.y;
        return dist(pt(0, 0), pt(hor, vert));
96f
cbb }
    // 5df9cf
13d ld polarea(vector<pt> v) { // area do poligono
        ld ret = 0;
9 c 5
c6e
        for (int i = 0; i < v.size(); i++)</pre>
```

```
80f
            ret += sarea(pt(0, 0), v[i], v[(i + 1) %
   v.size()]);
        return abs(ret);
d03
cbb }
    // se o ponto ta dentro do poligono: retorna O se ta
    // 1 se ta no interior e 2 se ta na borda
    // a6423f
8e7 int inpol(vector\langle pt \rangle \& v, pt p) { // O(n)
8de
        int qt = 0;
f 1 4
        for (int i = 0; i < v.size(); i++) {</pre>
             if (p == v[i]) return 2;
bda
6af
            int j = (i+1)%v.size();
            if (eq(p.y, v[i].y) and eq(p.y, v[j].y)) {
e38
                 if ((v[i]-p)*(v[j]-p) < eps) return 2;</pre>
97f
5 e 2
                 continue;
             }
cbb
388
             bool baixo = v[i].y+eps < p.y;</pre>
464
            if (baixo == (v[j].y+eps < p.y)) continue;</pre>
             auto t = (p-v[i])^(v[j]-v[i]);
366
1 b 4
             if (eq(t, 0)) return 2;
            if (baixo == (t > eps)) qt += baixo ? 1 : -1;
839
cbb
        return qt != 0;
b84
cbb }
    // c58350
6ff bool interpol(vector<pt> v1, vector<pt> v2) { // se dois
   poligonos se intersectam - O(n*m)
7 d 1
        int n = v1.size(), m = v2.size();
        for (int i = 0; i < n; i++) if (inpol(v2, v1[i]))</pre>
c36
   return 1;
        for (int i = 0; i < n; i++) if (inpol(v1, v2[i]))</pre>
ab8
   return 1:
        for (int i = 0; i < n; i++) for (int j = 0; j < m;
523
   j++)
0 c8
             if (interseg(line(v1[i], v1[(i+1)%n]),
   line(v2[j], v2[(j+1)%m]))) return 1;
        return 0;
bb3
cbb }
```

```
// 12559f
494 ld distpol(vector<pt> v1, vector<pt> v2) { // distancia
   entre poligonos
        if (interpol(v1, v2)) return 0;
f6b
349
        ld ret = DINF;
        for (int i = 0; i < v1.size(); i++) for (int j = 0;
   j < v2.size(); j++)
            ret = min(ret, distseg(line(v1[i], v1[(i + 1) %
6 c 2
   v1.size()]),
949
                         line(v2[i], v2[(i + 1) %
   v2.size()]));
edf
        return ret;
cbb }
    // 32623c
138 vector <pt> convex_hull(vector <pt> v) { // convex hull -
   O(n log(n))
        if (v.size() <= 1) return v;</pre>
52d
        vector<pt> 1, u;
526
        sort(v.begin(), v.end());
fca
        for (int i = 0; i < v.size(); i++) {</pre>
f14
            while (1.size() > 1 and !ccw(l[1.size()-2],
543
   1.back(), v[i]))
                l.pop_back();
364
c35
            l.push_back(v[i]);
cbb
3e9
        for (int i = v.size() - 1; i >= 0; i--) {
2eb
            while (u.size() > 1 \text{ and } !ccw(u[u.size()-2],
   u.back(), v[i]))
7a8
                u.pop_back();
a95
            u.push_back(v[i]);
cbb
        1.pop_back(); u.pop_back();
cfc
82b
        for (pt i : u) l.push_back(i);
792
        return 1:
cbb }
483 struct convex pol {
f50
        vector<pt> pol;
```

```
d41
        // nao pode ter ponto colinear no convex hull
d98
        convex pol() {}
a04
         convex_pol(vector<pt> v) : pol(convex_hull(v)) {}
        // se o ponto ta dentro do hull - O(\log(n))
d41
d41
        // 800813
8af
        bool is inside(pt p) {
            if (pol.size() == 1) return p == pol[0];
eae
            int 1 = 1, r = pol.size();
67f
40 c
             while (1 < r) {
ee4
                 int m = (1+r)/2;
48f
                 if (ccw(p, pol[0], pol[m])) 1 = m+1;
ef3
                 else r = m:
cbb
00a
            if (1 == 1) return isinseg(p, line(pol[0],
   pol[1]));
9e7
            if (1 == pol.size()) return false;
             return !ccw(p, pol[1], pol[1-1]);
1 c 0
cbb
        // ponto extremo em relacao a cmp(p, q) = p mais
d41
   extremo q
d41
        // (copiado de
   https://github.com/gustavoM32/caderno-zika)
        // 56ccd2
d41
719
        int extreme(const function < bool(pt, pt) > & cmp) {
             int n = pol.size();
b1c
            auto extr = [&](int i, bool& cur_dir) {
4a2
22a
                 \operatorname{cur\_dir} = \operatorname{cmp}(\operatorname{pol}[(i+1)\%n], \operatorname{pol}[i]);
61a
                 return !cur_dir and !cmp(pol[(i+n-1)%n],
   pol[i]);
214
             };
63d
             bool last_dir, cur_dir;
a0d
             if (extr(0, last_dir)) return 0;
993
             int 1 = 0, r = n;
ead
             while (1+1 < r) {
ee4
                 int m = (1+r)/2;
f29
                 if (extr(m, cur_dir)) return m;
44a
                 bool rel_dir = cmp(pol[m], pol[l]);
b18
                if ((!last_dir and cur_dir) or
                          (last dir == cur dir and rel dir ==
261
   cur dir)) {
8a6
                     1 = m;
```

```
1 f 1
                    last_dir = cur_dir;
b6c
                } else r = m;
cbb
792
            return 1;
cbb
        int max_dot(pt v) {
316
            return extreme([&](pt p, pt q) { return p*v >
ec1
   q*v; });
cbb
       }
        pair < int , int > tangents(pt p) {
a54
            auto L = [\&](pt q, pt r) \{ return ccw(p, q, r); \}
08 c
  };
            auto R = [&](pt q, pt r) { return ccw(p, r, q);
422
   };
            return {extreme(L), extreme(R)};
fa8
cbb
214 };
    // CIRCUNFERENCIA
    // a125e4
911 pt getcenter(pt a, pt b, pt c) { // centro da circunf
   dado 3 pontos
       b = (a + b) / 2;
174
        c = (a + c) / 2;
2ae
       return inter(line(b, b + rotate90(a - b)),
98b
3f8
                line(c, c + rotate90(a - c)));
cbb }
    // cd80c0
4b3 vector <pt> circ_line_inter(pt a, pt b, pt c, ld r) { //
   intersecao da circunf (c, r) e reta ab
        vector<pt> ret;
8af
        b = b-a, a = a-c:
f2b
4b1
        1d A = b*b:
20a
        1d B = a*b:
        1d C = a*a - r*r;
2e9
1fa
        1d D = B*B - A*C;
       if (D < -eps) return ret;</pre>
818
        ret.push back(c+a+b*(-B+sqrt(D+eps))/A);
dc5
        if (D > eps) ret.push_back(c+a+b*(-B-sqrt(D))/A);
20 e
edf
        return ret:
```

```
cbb }
    // fb11d8
ad2 vector <pt> circ_inter(pt a, pt b, ld r, ld R) { //
   intersecao da circunf (a, r) e (b, R)
        vector <pt> ret;
8af
        1d d = dist(a, b);
b7e
       if (d > r+R \text{ or } d+min(r, R) < max(r, R)) return ret;
5 ce
       1d x = (d*d-R*R+r*r)/(2*d);
398
       1d y = sqrt(r*r-x*x);
183
325
       pt v = (b-a)/d;
76e
        ret.push_back(a+v*x + rotate90(v)*y);
        if (y > 0) ret.push_back(a+v*x - rotate90(v)*y);
2cb
        return ret:
edf
cbb }
    // 3a44fb
6e0 bool operator <(const line& a, const line& b) { //
   comparador pra reta
       // assume que as retas tem p < q
d41
       pt v1 = a.q - a.p, v2 = b.q - b.p;
a13
       if (!eq(angle(v1), angle(v2))) return angle(v1) <</pre>
f82
   angle(v2);
780
        return ccw(a.p, a.q, b.p); // mesmo angulo
cbb }
b14 bool operator ==(const line& a, const line& b) {
76c
        return !(a < b) and !(b < a);
cbb }
    // comparador pro set pra fazer sweep line com segmentos
   // 36729f
2c4 struct cmp_sweepline {
        bool operator () (const line& a, const line& b)
   const {
d41
           // assume que os segmentos tem p < q
191
            if (a.p == b.p) return ccw(a.p, a.q, b.q);
            if (!eq(a.p.x, a.q.x) and (eq(b.p.x, b.q.x) or
231
   a.p.x+eps < b.p.x)
              return ccw(a.p, a.q, b.p);
780
            return ccw(a.p, b.q, b.p);
dc0
cbb
214 }:
```

```
// comparador pro set pra fazer sweep angle com segmentos
// f778aa
bef pt dir;
5b0 struct cmp_sweepangle {
d80    bool operator () (const line& a, const line& b)
    const {
522      return get_t(dir, a) + eps < get_t(dir, b);
cbb    }
214 };</pre>
```

#### 7.7 Primitivas Geometricas 3D

```
c83 typedef double ld;
e3b const ld DINF = 1e18;
107 \text{ const} 1d \text{ eps} = 1e-9;
d41
b32 #define sq(x) ((x)*(x))
d41
d97 bool eq(ld a, ld b) {
            return abs(a - b) <= eps;</pre>
ba0
cbb }
d41
b2a struct pt { // ponto
2eb
            ld x, y, z;
a50
            pt(1d x_{-} = 0, 1d y_{-} = 0, 1d z_{-} = 0) : x(x_{-}),
   y(y_{-}), z(z_{-}) \{ \}
             bool operator < (const pt p) const {</pre>
5bc
                      if (!eq(x, p.x)) return x < p.x;
059
                     if (!eq(y, p.y)) return y < p.y;</pre>
f98
                     if (!eq(z, p.z)) return z < p.z;
44 c
                      return 0;
bb3
             }
cbb
a83
             bool operator == (const pt p) const {
                      return eq(x, p.x) and eq(y, p.y) and
41 c
   eq(z, p.z);
cbb
             pt operator + (const pt p) const { return
   pt(x+p.x, y+p.y, z+p.z);}
             pt operator - (const pt p) const { return
392
   pt(x-p.x, y-p.y, z-p.z); }
            pt operator * (const ld c) const { return pt(x*c
```

```
, y*c , z*c ); }
            pt operator / (const ld c) const { return pt(x/c
    , y/c , z/c ); }
a65
            ld operator * (const pt p) const { return x*p.x
   + y*p.y + z*p.z; }
            pt operator ^ (const pt p) const { return
7 f 6
   pt(y*p.z - z*p.y, z*p.x - x*p.z, x*p.y - y*p.x); }
            friend istream& operator >> (istream& in, pt& p)
5ed
   {
9bf
                    return in >> p.x >> p.y >> p.z;
cbb
            }
214 }:
d 4 1
b3a struct line { // reta
730
            pt p, q;
0d6
            line() {}
4 b 8
            line(pt p_{-}, pt q_{-}) : p(p_{-}), q(q_{-}) {}
            friend istream& operator >> (istream& in, line&
8 d 7
 r) {
4cb
                    return in >> r.p >> r.q;
cbb
            }
214 };
d41
79b struct plane { // plano
7 e 1
            array <pt, 3> p; // pontos que definem o plano
29b
            array < ld, 4 > eq; // equacao do plano
bb7
            plane() {}
fb0
            plane(pt p_, pt q_, pt r_) : p(\{p_, q_, r_\}) {
   build(): }
d41
ca9
            friend istream& operator >> (istream& in, plane&
   P) {
2ab
                    return in >> P.p[0] >> P.p[1] >> P.p[2];
70e
                     P.build():
cbb
            }
0a8
            void build() {
da2
                     pt dir = (p[1] - p[0]) ^ (p[2] - p[0]);
7 d5
                    eq = {dir.x, dir.y, dir.z,
   dir*p[0]*(-1)};
cbb
            }
214 };
d41
```

```
// converte de coordenadas polares para cartesianas
   // (angulos devem estar em radianos)
   // phi eh o angulo com o eixo z (cima) theta eh o angulo
      de rotacao ao redor de z
2fb pt convert(ld rho, ld th, ld phi) {
         return pt(sin(phi) * cos(th), sin(phi) *
  sin(th), cos(phi)) * rho;
d41
   // projecao do ponto p na reta r
256 pt proj(pt p, line r) {
           if (r.p == r.q) return r.p;
         r.q = r.q - r.p; p = p - r.p;
97a
           pt proj = r.q * ((p*r.q) / (r.q*r.q));
9f8
2cd
           return proj + r.p;
cbb }
d41
   // projecao do ponto p no plano P
b1a pt proj(pt p, plane P) {
           p = p - P.p[0], P.p[1] = P.p[1] - P.p[0], P.p[2]
  = P.p[2] - P.p[0];
           pt norm = P.p[1] ^ P.p[2];
b69
6ab
           pt proj = p - (norm * (norm * p) / (norm*norm));
           return proj + P.p[0];
467
cbb }
d41
   // distancia
a45 ld dist(pt a, pt b) {
    return sqrt(sq(a.x-b.x) + sq(a.y-b.y) +
   sq(a.z-b.z));
cbb }
d41
   // distancia ponto reta
137 ld distline(pt p, line r) {
       return dist(p, proj(p, r));
ce1
cbb }
d41
   // distancia de ponto para segmento
d43 ld distseg(pt p, line r) {
           if ((r.q - r.p)*(p - r.p) < 0) return dist(r.p,
73 d
  p);
951
           if ((r.p - r.q)*(p - r.q) < 0) return dist(r.q,
```

```
p);
           return distline(p, r);
200
cbb }
d41
   // distancia de ponto a plano com sinal
7cc ld sdist(pt p, plane P) {
           return P.eq[0]*p.x + P.eq[1]*p.y + P.eq[2]*p.z +
    P.eq[3];
cbb }
d41
   // distancia de ponto a plano
768 ld distplane(pt p, plane P) {
       return abs(sdist(p, P));
cbb }
d41
   // se ponto pertence a reta
099 bool isinseg(pt p, line r) {
          return eq(distseg(p, r), 0);
a32
cbb }
d41
   // se ponto pertence ao triangulo definido por P.p
cd2 bool isinpol(pt p, vector<pt> v) {
            assert(v.size() >= 3);
fad
bf4
            pt norm = (v[1]-v[0]) ^ (v[2]-v[1]);
8a4
           bool inside = true;
           int sign = -1;
cec
         for (int i = 0; i < v.size(); i++) {</pre>
f14
834
                    line r(v[(i+1)\%3], v[i]);
2a9
                   if (isinseg(p, r)) return true;
d41
4ef
                    pt ar = v[(i+1)\%3] - v[i];
                   if (sign == -1) sign =
320
   ((ar^(p-v[i]))*norm > 0);
                    else if (((ar^(p-v[i]))*norm > 0) !=
82b
   sign) inside = false;
cbb
aca
            return inside;
cbb }
d41
   // distancia de ponto ate poligono
361 ld distpol(pt p, vector<pt> v) {
            pt p2 = proj(p, plane(v[0], v[1], v[2]));
3 e 7
```

```
61a
            if (isinpol(p2, v)) return dist(p, p2);
349
           ld ret = DINF;
f14
           for (int i = 0; i < v.size(); i++) {</pre>
                   int j = (i+1)%v.size();
6af
                    ret = min(ret, distseg(p, line(v[i],
5ee
  v[j])));
cbb
edf
            return ret;
cbb }
d41
   // intersecao de plano e segmento
    // BOTH = o segmento esta no plano
   // ONE = um dos pontos do segmento esta no plano
   // PARAL = segmento paralelo ao plano
    // CONCOR = segmento concorrente ao plano
e51 enum RETCODE {BOTH, ONE, PARAL, CONCOR};
26b pair < RETCODE, pt > intersect(plane P, line r) {
        1d d1 = sdist(r.p, P);
fac
        1d d2 = sdist(r.q, P);
f8f
        if (eq(d1, 0) \text{ and } eq(d2, 0))
53a
                   return pair(BOTH, r.p);
504
72 c
        if (eq(d1, 0))
847
                    return pair(ONE, r.p);
485
        if (eq(d2, 0))
168
                    return pair(ONE, r.q);
        if ((d1 > 0 \text{ and } d2 > 0) \text{ or } (d1 < 0 \text{ and } d2 < 0))
3fb
463
            if (eq(d1-d2, 0)) return pair(PARAL, pt());
406
            return pair(CONCOR, pt());
cbb
c84
        1d frac = d1 / (d1 - d2);
        pt res = r.p + ((r.q - r.p) * frac);
3ff
        return pair(ONE, res);
394
cbb }
d41
   // rotaciona p ao redor do eixo u por um angulo a
787 pt rotate(pt p, pt u, ld a) {
           u = u / dist(u, pt());
773
           return u * (u * p) + (u ^ p ^ u) * cos(a) + (u ^
  p) * sin(a);
cbb }
d41
```

### 7.8 Primitivas Geometricas Inteiras

```
2de #define sq(x) ((x)*(11)(x))
    // 840720
b2a struct pt { // ponto
e91
        int x, y;
       pt(int x_{-} = 0, int y_{-} = 0) : x(x_{-}), y(y_{-}) {}
df1
       bool operator < (const pt p) const {</pre>
5bc
            if (x != p.x) return x < p.x;
95a
89 c
          return y < p.y;
cbb
a83
       bool operator == (const pt p) const {
d74
            return x == p.x and y == p.y;
cbb
        pt operator + (const pt p) const { return pt(x+p.x,
cb9
   y+p.y); }
a24
        pt operator - (const pt p) const { return pt(x-p.x,
   y-p.y); }
       pt operator * (const int c) const { return pt(x*c,
   y*c); }
60d ll operator * (const pt p) const { return x*(11)p.x
   + v*(11)p.v; }
       11 operator ^ (const pt p) const { return x*(11)p.y
d86
 - y*(11)p.x; }
5ed
       friend istream& operator >> (istream& in, pt& p) {
e37
            return in >> p.x >> p.y;
cbb
       }
214 };
   // 7ab617
b3a struct line { // reta
730
        pt p, q;
0 d6
       line() {}
       line(pt p_, pt q_) : p(p_), q(q_) {}
4b8
8 d7
       friend istream& operator >> (istream& in, line& r) {
4cb
            return in >> r.p >> r.q;
cbb
214 };
    // PONTO & VETOR
```

```
// 51563e
ea8 ll dist2(pt p, pt q) { // quadrado da distancia
       return sq(p.x - q.x) + sq(p.y - q.y);
cbb }
    // bf431d
5a2 ll sarea2(pt p, pt q, pt r) { // 2 * area com sinal
       return (q-p)^(r-q);
cbb }
    // a082d3
e32 bool col(pt p, pt q, pt r) { // se p, q e r sao colin.
       return sarea2(p, q, r) == 0;
cbb }
    // 42bb09
Ocd bool ccw(pt p, pt q, pt r) { // se p, q, r sao ccw
        return sarea2(p, q, r) > 0;
276
cbb }
    // fcf924
c31 int quad(pt p) { // quadrante de um ponto
       return (p.x<0)^3*(p.y<0);
dbb
cbb }
    // 77187b
2df bool compare_angle(pt p, pt q) { // retorna se ang(p) <
   ang(q)
       if (quad(p) != quad(q)) return quad(p) < quad(q);</pre>
9fc
       return ccw(q, pt(0, 0), p);
ea1
cbb }
    // e4ad5e
ab1 pt rotate90(pt p) { // rotaciona 90 graus
a0d
       return pt(-p.y, p.x);
cbb }
    // RETA
    // c9f07f
099 bool isinseg(pt p, line r) { // se p pertence ao seg de r
f65
       pt a = r.p - p, b = r.q - p;
```

```
2ac
        return (a ^ b) == 0 and (a * b) <= 0:
cbb }
   // 35998c
676 bool interseg(line r, line s) { // se o seg de r
   intersecta o seg de s
       if (isinseg(r.p, s) or isinseg(r.q, s)
19b
            or isinseg(s.p, r) or isinseg(s.q, r)) return 1;
c21
        return ccw(r.p, r.q, s.p) != ccw(r.p, r.q, s.q) and
9fa
                ccw(s.p, s.q, r.p) != ccw(s.p, s.q, r.q);
413
cbb }
   // dd8702
9e0 int segpoints(line r) { // numero de pontos inteiros no
        return 1 + \_gcd(abs(r.p.x - r.q.x), abs(r.p.y -
9ce
   r.q.y));
cbb }
    // d273be
88a double get_t(pt v, line r) { // retorna t tal que t*v
   pertence a reta r
        return (r.p^r.q) / (double) ((r.p-r.q)^v);
1ad
cbb }
   // POLIGONO
    // quadrado da distancia entre os retangulos a e b
       (lados paralelos aos eixos)
   // assume que ta representado (inferior esquerdo,
       superior direito)
    // e13018
485 11 dist2_rect(pair<pt, pt> a, pair<pt, pt> b) {
c59 int hor = 0, vert = 0;
34b
        if (a.second.x < b.first.x) hor = b.first.x -</pre>
   a.second.x:
        else if (b.second.x < a.first.x) hor = a.first.x -</pre>
f5f
   b.second.x;
        if (a.second.y < b.first.y) vert = b.first.y -</pre>
4fd
   a.second.v;
80a
        else if (b.second.y < a.first.y) vert = a.first.y -</pre>
```

```
b.second.y;
        return sq(hor) + sq(vert);
869
cbb }
    // d5f693
9c3 ll polarea2(vector<pt> v) { // 2 * area do poligono
        11 \text{ ret} = 0;
        for (int i = 0; i < v.size(); i++)</pre>
сбе
             ret += sarea2(pt(0, 0), v[i], v[(i + 1) %
532
   v.size()]);
d03
        return abs(ret);
cbb }
    // se o ponto ta dentro do poligono: retorna 0 se ta
    // 1 se ta no interior e 2 se ta na borda
    // afd587
8e7 int inpol(vector\langle pt \rangle \& v, pt p) { // O(n)
        int qt = 0;
8de
        for (int i = 0; i < v.size(); i++) {</pre>
f 14
            if (p == v[i]) return 2;
bda
6af
            int j = (i+1)%v.size();
            if (p.y == v[i].y and p.y == v[j].y) {
cc6
                if ((v[i]-p)*(v[j]-p) <= 0) return 2;</pre>
547
                 continue;
5e2
            }
cbb
78 c
            bool baixo = v[i].y < p.y;</pre>
            if (baixo == (v[j].y < p.y)) continue;</pre>
057
366
             auto t = (p-v[i])^(v[j]-v[i]);
2ad
            if (!t) return 2:
            if (baixo == (t > 0)) qt += baixo ? 1 : -1;
0bb
cbb
        return qt != 0;
b84
cbb }
    // 32623c
138 vector <pt > convex_hull(vector <pt > v) { // convex hull -
   O(n log(n))
        if (v.size() <= 1) return v;</pre>
52d
        vector<pt> 1, u;
526
        sort(v.begin(), v.end());
fca
        for (int i = 0; i < v.size(); i++) {</pre>
f 14
```

```
543
            while (1.size() > 1 \text{ and } !ccw(l[1.size() - 2],
   1.back(), v[i]))
                1.pop_back();
364
c35
            l.push_back(v[i]);
cbb
        for (int i = v.size() - 1; i >= 0; i--) {
3 e 9
            while (u.size() > 1 \text{ and } !ccw(u[u.size()-2],
2eb
   u.back(), v[i]))
7 a 8
                u.pop_back();
a95
            u.push_back(v[i]);
cbb
cfc
        1.pop_back(); u.pop_back();
82b
        for (pt i : u) l.push_back(i);
792
        return 1:
cbb }
    // af2d96
786 ll interior_points(vector<pt> v) { // pontos inteiros
   dentro de um poligono simples
        11 b = 0;
c4e
        for (int i = 0; i < v.size(); i++)</pre>
c6e
            b += segpoints(line(v[i], v[(i+1)%v.size()])) -
Осе
  1;
        return (polarea2(v) - b) / 2 + 1;
a1c
cbb }
483 struct convex_pol {
f50
        vector <pt> pol;
d41
        // nao pode ter ponto colinear no convex hull
d98
        convex_pol() {}
        convex_pol(vector < pt > v) : pol(convex_hull(v)) {}
a04
        // se o ponto ta dentro do hull - O(\log(n))
d41
d41
        // 800813
8af
        bool is_inside(pt p) {
            if (pol.size() == 1) return p == pol[0];
eae
67f
            int 1 = 1, r = pol.size();
40 c
            while (1 < r) {
ee4
                int m = (1+r)/2;
                if (ccw(p, pol[0], pol[m])) 1 = m+1;
48f
                else r = m:
ef3
```

```
cbb
            if (1 == 1) return isinseg(p, line(pol[0],
00a
   pol[1])):
            if (l == pol.size()) return false;
9e7
             return !ccw(p, pol[1], pol[1-1]);
1 c 0
cbb
d41
        // ponto extremo em relacao a cmp(p, q) = p mais
        // (copiado de
d41
   https://github.com/gustavoM32/caderno-zika)
d41
719
         int extreme(const function < bool(pt, pt) > & cmp) {
             int n = pol.size();
b1c
4a2
             auto extr = [&](int i, bool& cur_dir) {
22a
                  \operatorname{cur}_{\operatorname{dir}} = \operatorname{cmp}(\operatorname{pol}[(i+1)\%n], \operatorname{pol}[i]);
                 return !cur_dir and !cmp(pol[(i+n-1)%n],
61a
   pol[i]);
214
             };
             bool last_dir, cur_dir;
63 d
             if (extr(0, last dir)) return 0;
a0d
993
            int 1 = 0, r = n;
ead
             while (1+1 < r) {
ee4
                 int m = (1+r)/2;
                 if (extr(m, cur_dir)) return m;
f29
                 bool rel_dir = cmp(pol[m], pol[1]);
44a
b18
                 if ((!last dir and cur dir) or
261
                          (last_dir == cur_dir and rel_dir ==
   cur dir)) {
8a6
                      1 = m;
                     last_dir = cur_dir;
1f1
                 } else r = m;
b6c
cbb
792
             return 1;
cbb
316
         int max_dot(pt v) {
             return extreme([&](pt p, pt q) { return p*v >
ec1
   q*v; \});
cbb
         pair<int, int> tangents(pt p) {
a54
08 c
             auto L = [\&](pt q, pt r) \{ return ccw(p, q, r); 
   };
422
             auto R = [\&](pt q, pt r) \{ return ccw(p, r, q); \}
```

```
};
            return {extreme(L), extreme(R)};
fa8
cbb
      }
214 };
    // dca598
6e0 bool operator <(const line& a, const line& b) { //
   comparador pra reta
d41
        // assume que as retas tem p < q</pre>
        pt v1 = a.q - a.p, v2 = b.q - b.p;
a13
        bool b1 = compare_angle(v1, v2), b2 =
   compare_angle(v2, v1);
73c
        if (b1 or b2) return b1;
780
        return ccw(a.p, a.q, b.p); // mesmo angulo
cbb }
b14 bool operator ==(const line& a, const line& b) {
        return !(a < b) and !(b < a);
cbb }
   // comparador pro set pra fazer sweep line com segmentos
   // 6774df
2c4 struct cmp_sweepline {
        bool operator () (const line& a, const line& b)
   const {
            // assume que os segmentos tem p < q</pre>
d41
           if (a.p == b.p) return ccw(a.p, a.q, b.q);
191
614
           if (a.p.x != a.q.x and (b.p.x == b.q.x or a.p.x
   < b.p.x)
780
                return ccw(a.p, a.q, b.p);
            return ccw(a.p, b.q, b.p);
dc0
cbb
       }
214 };
    // comparador pro set pra fazer sweep angle com segmentos
   // 1ee7f5
bef pt dir;
5b0 struct cmp_sweepangle {
        bool operator () (const line& a, const line& b)
   const {
261
            return get_t(dir, a) < get_t(dir, b);</pre>
cbb
214 };
```

# 8 Extra

# 8.1 fastIO.cpp

```
int read_int() {
    bool minus = false;
    int result = 0;
    char ch;
    ch = getchar();
    while (1) {
        if (ch == '-') break;
        if (ch >= '0' && ch <= '9') break;
        ch = getchar();
    if (ch == '-') minus = true;
    else result = ch-'0';
    while (1) {
        ch = getchar();
        if (ch < '0' || ch > '9') break;
        result = result *10 + (ch - '0');
    if (minus) return -result;
    else return result;
```

### 8.2 vimrc

set ts=4 si ai sw=4 nu mouse=a undofile syntax on

# 8.3 timer.cpp

```
};
8.4 rand.cpp
mt19937 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
int uniform(int 1, int r){
    uniform_int_distribution <int> uid(1, r);
    return uid(rng);
}
8.5 template.cpp
#include <bits/stdc++.h>
using namespace std;
#define _ ios_base::sync_with_stdio(0);cin.tie(0);
#define endl '\n'
typedef long long 11;
const int INF = 0x3f3f3f3f;
const 11 LINF = 0x3f3f3f3f3f3f3f3f3f11;
int main() { _
    exit(0);
}
8.6 debug.cpp
void debug_out(string s, int line) { cerr << endl; }</pre>
template < typename H, typename ... T>
void debug_out(string s, int line, H h, T... t) {
   if (s[0] != ',') cerr << "Line(" << line << ") ";</pre>
    do { cerr << s[0]; s = s.substr(1);</pre>
   } while (s.size() and s[0] != ',');
    cerr << " = " << h;
```

```
debug_out(s, line, t...);
}
#ifdef DEBUG
#define debug(...) debug_out(#__VA_ARGS__, __LINE__,
   __VA_ARGS__)
#else
#define debug(...)
#endif
     stress.sh
P = a
make ${P} ${P}2 gen || exit 1
for ((i = 1; ; i++)) do
    ./gen $i > in
    ./${P} < in > out
    ./${P}2 < in > out2
    if (! cmp -s out out2) then
        echo "--> entrada:"
        cat in
        echo "--> saida1:"
        cat out
        echo "--> saida2:"
        cat out2
        break;
    fi
    echo $i
done
     makefile
8.8
CXX = g++
CXXFLAGS = -fsanitize=address, undefined
   -fno-omit-frame-pointer -g -Wall -Wshadow -std=c++17
   -Wno-unused-result -Wno-sign-compare -Wno-char-subscripts
   #-fuse-ld=gold
     hash.sh
8.9
```

# Para usar (hash das linhas [11, 12]):

```
# bash hash.sh arquivo.cpp 11 12
sed -n $2','$3' p' $1 | sed '/^#w/d' | cpp -dD -P
    -fpreprocessed | tr -d '[:space:]' | md5sum | cut -c-6

8.10 linehash.sh

# Para usar:
# bash linehash.sh arquivo.cpp
while read 1; do
    echo $1 > tmp.txt
    h=$(echo $(bash hash.sh tmp.txt 1 1) | cut -c-3)
    echo "$h $1"

done < "$1"</pre>
```