# Rábalabaxúrias [UFMG]

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```
// ex : str_hash < 31, 1e9 + 7 > h(s);
        11 \text{ val} = h(10, 20);
// Complexidades:
// build - O(|s|)
// get_hash - 0(1)
typedef long long 11;
template < int P, int MOD > struct str_hash {
    int n;
    string s;
    vector<ll> h, power;
    str_hash(string s_): n(s_.size()), s(s_), h(n), power(n){
        power[0] = 1;
        for (int i = 1; i < n; i++) power[i] = power[i-1]*P</pre>
           % MOD:
        h[0] = s[0];
        for (int i = 1; i < n; i++) h[i] = (h[i-1]*P + s[i])
           % MOD;
    }
    11 operator()(int i, int j){
        if (!i) return h[j];
        return (h[j] - h[i-1]*power[j-i+1] % MOD + MOD) %
           MOD;
    }
};
1.2 KMP
// Primeiro chama a funcao process com o padrao
// Depois chama match com (texto, padrao)
// Vai retornar o numero de ocorrencias do padrao
// p eh 1-based
// Complexidades:
// process - O(m)
// match - 0(n + m)
// n = |texto| e m = |padrao|
int p[MAX];
```

```
void process(string& s) {
    int i = 0, j = -1;
    p[0] = -1;
    while (i < s.size()) {</pre>
        while (j \ge 0 \text{ and } s[i] != s[j]) j = p[j];
        i++, j++;
        p[i] = j;
    }
}
int match(string& s, string& t) {
    process(t);
    int i = 0, j = 0, ans = 0;
    while (i < s.size()) {</pre>
        while (j \ge 0 \text{ and } s[i] != t[j]) j = p[j];
        i++, j++;
        if (j == t.size()) j = p[j], ans++;
    }
    return ans;
}
1.3 Z
// Complexidades:
// z - O(|s|)
// \text{ match - } O(|s| + |p|)
vector<int> get_z(string s) {
    int n = s.size();
    vector < int > z(n, 0);
    // intervalo da ultima substring valida
    int 1 = 0, r = 0;
    for (int i = 1; i < n; i++) {</pre>
        // estimativa pra z[i]
        if (i \le r) z[i] = min(r - i + 1, z[i - 1]);
        // calcula valor correto
        while (i + z[i] < n \text{ and } s[z[i]] == s[i + z[i]])
            z[i]++;
        // atualiza [l, r]
        if (i + z[i] - 1 > r) l = i, r = i + z[i] - 1;
    }
```

```
return z;
}
// quantas vezes p aparece em s
int match(string s, string p) {
    int n = s.size(), m = p.size();
    vector < int > z = get_z(p + s);
    int ret = 0;
    for (int i = m; i < n + m; i++)</pre>
        if (z[i] >= m) ret++;
    return ret;
}
1.4 Suffix Array
// kasai recebe o suffix array e calcula lcp[i],
// o lcp entre s[sa[i],...,n-1] e s[sa[i+1],..,n-1]
//
// Complexidades:
// suffix_array - O(n log(n))
// kasai - O(n)
vector<int> suffix_array(string s) {
    s += "$":
    int n = s.size(), N = max(n, 260);
    vector < int > sa(n), ra(n);
    for(int i = 0; i < n; i++) sa[i] = i, ra[i] = s[i];</pre>
    for(int k = 0; k < n; k ? k *= 2 : k++) {
        vector < int > nsa(sa), nra(n), cnt(N);
```

for (int i = 0; i < n; i++) nsa[i] = (nsa[i]-k+n)%n,

for(int i = n-1; i+1; i--) sa[--cnt[ra[nsa[i]]]] =

for(int i = 1, r = 0; i < n; i++) nra[sa[i]] = r +=</pre>

for(int i = 1; i < N; i++) cnt[i] += cnt[i-1];</pre>

cnt[ra[i]]++;

ra[sa[i]] !=

nsa[i];

```
ra[sa[i-1]] or ra[(sa[i]+k)%n] !=
                ra[(sa[i-1]+k)%n];
        ra = nra;
    return vector < int > (sa.begin() + 1, sa.end());
}
vector<int> kasai(string s, vector<int> sa) {
    int n = s.size(), k = 0;
    vector < int > ra(n), lcp(n);
    for (int i = 0; i < n; i++) ra[sa[i]] = i;
    for (int i = 0; i < n; i++, k -= !!k) {
        if (ra[i] == n-1) { k = 0; continue; }
        int j = sa[ra[i]+1];
        while (i+k < n \text{ and } j+k < n \text{ and } s[i+k] == s[j+k]) k++;
        lcp[ra[i]] = k;
    }
    return lcp;
}
1.5 Suffix Array Rafael
// O(n log^2(n))
struct suffix_array{
    string &s;
    int n;
    vector<int> p, r, aux, lcp;
    seg_tree<int, min_el> st;
    suffix_array(string &s):
        s(s), n(s.size()), p(n), r(n), aux(n), lcp(n){
            for (int i = 0; i < n; i++){</pre>
                 p[i] = i;
                r[i] = s[i];
            auto rank = [&](int i){
                 if (i >= n) return -i;
                 return r[i];
            for (int d = 1; d < n; d *= 2) {
                 auto t = [&](int i){
```

```
return make_pair(rank(i), rank(i+d));
                 };
                 sort(p.begin(), p.end(),
                         [&](int &i, int &j){
                         return t(i) < t(j);
                     );
                 aux[p[0]] = 0;
                 for (int i = 1; i < n; i++)</pre>
                     aux[p[i]] = aux[p[i-1]] + (t(p[i]) >
                        t(p[i-1]));
                 for (int j = 0; j < n; j++) r[j] = aux[j];
                 if (aux[p[n-1]] == n-1) break;
            }
            int h = 0;
            for (int i = 0; i < n; i++){</pre>
                 if (r[i] == n-1){
                     lcp[r[i]] = 0;
                     continue;
                 }
                 int j = p[r[i] + 1];
                 while (i + h < n \&\& j + h < n \&\& s[i+h] ==
                    s[j+h]) h++;
                 lcp[r[i]] = h;
                 h = max(0, h-1);
             st = seg_tree < int, min_el > (&lcp);
    int query(int 1, int r){
        return st.query(1, r);
    11 distinct_substrings(){
        11 \text{ ans} = p[0] + 1;
        for (int i = 1; i < n; i++)</pre>
             ans += p[i] - lcp[i-1] + 1;
        return ans;
    }
};
```

## 2 Estruturas

## 2.1 BIT 2D

```
// BIT 1-based
// Para mudar o valor da posicao (x, y) para k,
// faca: poe(x, y, k - sum(x, y, x, y))
//
// Complexidades:
// poe - O(log^2(n))
// \text{ query - O(log^2(n))}
int n;
int bit[MAX][MAX];
void poe(int x, int y, int k) {
    for (int y2 = y; x \le n; x += x & -x)
        for (y = y2; y \le n; y += y \& -y)
            bit[x][y] += k;
}
int sum(int x, int y) {
    int ret = 0;
    for (int y2 = y; x; x -= x & -x)
        for (y = y2; y; y -= y \& -y)
            ret += bit[x][y];
    return ret;
}
int query(int x, int y, int z, int w) {
    return sum(z, w) - sum(x-1, w)
        - sum(z, y-1) + sum(x-1, y-1);
}
     Wavelet Tree
2.2
// Usa O(sigma + n log(sigma)) de memoria,
// onde sigma = MAXN - MINN
// Depois do build, o v fica ordenado
```

// count(i, j, x, y) retorna o numero de elementos de

```
// v[i, j) que pertencem a [x, y]
// kth(i, j, k) retorna o elemento que estaria
// na poscicao k-1 de v[i, j), se ele fosse ordenado
// sum(i, j, x, y) retorna a soma dos elementos de
// v[i, j) que pertencem a [x, y]
// sumk(i, j, k) retorna a soma dos k-esimos menores
// elementos de v[i, j) (sum(i, j, 1) retorna o menor)
//
// Complexidades:
// build - O(n log(sigma))
// count - O(log(sigma))
// kth - O(log(sigma))
// sum - O(log(sigma))
// sumk - O(log(sigma))
int n, v[MAX];
vector < int > > esq(4*(MAXN-MINN)), pref(4*(MAXN-MINN));
void build(int b = 0, int e = n, int p = 1, int l = MINN,
   int r = MAXN) {
    int m = (1+r)/2; esq[p].push_back(0);
       pref[p].push_back(0);
    for (int i = b; i < e; i++) {</pre>
        esq[p].push_back(esq[p].back()+(v[i]<=m));</pre>
        pref[p].push_back(pref[p].back()+v[i]);
    if (l == r) return;
    int m2 = stable_partition(v+b, v+e, [=](int i){return i
       <= m;}) - v;
    build(b, m2, 2*p, 1, m), build(m2, e, 2*p+1, m+1, r);
}
int count(int i, int j, int x, int y, int p = 1, int l =
   MINN, int r = MAXN) {
    if (y < 1 \text{ or } r < x) \text{ return } 0;
    if (x \le 1 \text{ and } r \le y) \text{ return } j-i;
    int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
    return count(ei, ej, x, y, 2*p, 1, m)+count(i-ei, j-ej,
       x, y, 2*p+1, m+1, r);
}
int kth(int i, int j, int k, int p=1, int l = MINN, int r =
```

```
if (1 == r) return 1;
    int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
    if (k <= ej-ei) return kth(ei, ej, k, 2*p, 1, m);</pre>
    return kth(i-ei, j-ej, k-(ej-ei), 2*p+1, m+1, r);
}
int sum(int i, int j, int x, int y, int p = 1, int l = MINN,
   int r = MAXN) {
    if (y < 1 \text{ or } r < x) \text{ return } 0;
    if (x <= 1 and r <= y) return pref[p][j]-pref[p][i];</pre>
    int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
    return sum(ei, ej, x, y, 2*p, 1, m) + sum(i-ei, j-ej, x,
       y, 2*p+1, m+1, r);
}
int sumk(int i, int j, int k, int p = 1, int l = MINN, int r
   = MAXN)
    if (1 == r) return 1*k;
    int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
    if (k <= ej-ei) return sumk(ei, ej, k, 2*p, 1, m);</pre>
    return pref[2*p][ej]-pref[2*p][ei]+sumk(i-ei, j-ej,
       k-(ej-ei), 2*p+1, m+1, r);
}
2.3 MergeSort Tree
// query(a, b, val) retorna numero de
// elementos em [a, b] <= val</pre>
// Usa O(n log(n)) de memoria
//
// Complexidades:
// build - O(n log(n))
// query - O(log^2(n))
#define ALL(x) x.begin(),x.end()
int v[MAX], n;
vector < vector < int > > tree(4*MAX);
void build(int p, int l, int r) {
    if (1 == r) return tree[p].push_back(cr[1]);
```

MAXN) {

## 2.4 Order Statistic Set

```
// Funciona do C++11 pra cima
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
template <class T>
    using ord_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
// para declarar:
ord_set < int > s;
// coisas do set normal funcionam:
for (auto i : s) cout << i << endl;</pre>
cout << s.size() << endl;</pre>
// k-esimo maior elemento O(log|s|):
// k=0: menor elemento
cout << *s.find_by_order(k) << endl;</pre>
// quantos sao menores do que k O(log|s|):
cout << s.order_of_key(k) << endl;</pre>
// Para fazer um multiset, tem que
// usar ord_set<pair<int, int> > com o
// segundo parametro sendo algo para diferenciar
```

```
// os ementos iguais.
// s.order_of_key({k, -INF}) vai retornar o
// numero de elementos < k</pre>
```

## 2.5 SQRT decomposition

```
// O-indexed
// MAX2 = sqrt(MAX)
// O bloco da posicao x eh
// sempre x/q
//
// Complexidades:
// build - O(n)
// query - 0(sqrt(n))
int n, q;
int v[MAX];
int bl[MAX2];
void build() {
    q = (int) sqrt(n);
    // computa cada bloco
   for (int i = 0; i <= q; i++) {</pre>
        bl[i] = INF;
        for (int j = 0; j < q and q * i + j < n; <math>j++)
            bl[i] = min(bl[i], v[q * i + j]);
    }
}
int query(int a, int b) {
    int ret = INF;
    // linear no bloco de a
    for (; a <= b and a % q; a++) ret = min(ret, v[a]);
    // bloco por bloco
    for (; a + q <= b; a += q) ret = min(ret, bl[a / q]);</pre>
    // linear no bloco de b
    for (; a <= b; a++) ret = min(ret, v[a]);</pre>
```

```
return ret;
}
     Treap
// insert - O(log(n))
// erase - O(log(n))
// query - O(log(n))
mt19937 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
template < typename T > struct treap {
    struct node {
        int p;
        int 1, r;
        T v;
        int sz;
        T min_s;
        node(){}
        node(T \ v):p(rng()), 1(-1), r(-1), v(v){}
    } t[MAX];
    int size(int i){
        if (i == -1) return 0;
        return t[i].sz;
    }
    void update(int i){
        if (i == -1) return;
        t[i].min_s = t[i].v;
        int 1 = t[i].1;
        int r = t[i].r;
        t[i].sz = 1 + size(1) + size(r);
        if (1 != -1)
            t[i].min_s = min(t[i].min_s, t[l].min_s);
        if (r != -1)
            t[i].min_s = min(t[i].min_s, t[r].min_s);
    }
```

void split(int i, int k, int &l, int &r){ //key

if (i == -1){

```
1 = -1: r = -1:
        return;
    }
   if (t[i].v < k){</pre>
        split(t[i].r, k, 1, r);
        t[i].r = 1;
        1 = i;
    }
    else{
        split(t[i].1, k, 1, r);
        t[i].1 = r;
        r = i;
    }
    update(i);
}
void split_implicit(int i, int k, int &l, int &r, int sz
   = 0){}
    if (i == -1){
        1 = -1; r = -1;
        return;
    }
    int inc = size(t[i].1); //quantidade elementos menor
       que k
    if (sz+inc < k){
        split_implicit(t[i].r, k, l, r, sz+inc+1);
        t[i].r = 1;
        1 = i;
    }
    else{
        split_implicit(t[i].1, k, 1, r, sz);
        t[i].l = r;
        r = i;
    }
    update(i);
}
int merge(int 1, int r){ //priority
    if (1 == -1) {
        update(r);
        return r;
    if (r == -1) {
        update(1);
```

```
return 1:
    }
    if (t[1].p > t[r].p){
       t[1].r = merge(t[1].r, r);
        update(1);
        return 1;
    }
    else{
        t[r].l = merge(l, t[r].l);
        update(r);
        return r;
    }
}
int it = 0;
void insert(int &root, T v){
    int M = it++;
    t[M] = node(v);
    if (root == -1) {
        root = M;
        return;
    }
    root = merge(root, M);
}
T query(int &root, int L, int R){
    int 1, m, r;
    split_implicit(root, R+1, m, r);
    split_implicit(m, L, 1, m);
    T ans = t[m].min_s;
    1 = merge(1, m);
    1 = merge(1, r);
    root = 1;
    return ans;
void erase(int &root, int pos){
    int 1, m, r;
    split_implicit(root, pos+1, m, r);
    split_implicit(m, pos, 1, m);
    1 = merge(1, r);
    root = 1;
```

```
}
};
2.7 SparseTable
// MAX2 = log(MAX)
// Complexidades:
// build - O(n log(n))
// query - O(1)
int n;
int v[MAX];
int m[MAX][MAX2]; // m[i][j] : posicao do minimo
                   // em [v[i], v[i + 2^j - 1]]
void build() {
    for (int i = 0; i < n; i++) m[i][0] = i;</pre>
    for (int j = 1; 1 << j <= n; j++) {
        int tam = 1 << j;</pre>
        for (int i = 0; i + tam <= n; i++) {</pre>
             if (v[m[i][j - 1]] < v[m[i + tam/2][j - 1]])</pre>
                 m[i][j] = m[i][j - 1];
             else m[i][j] = m[i + tam/2][j - 1];
        }
    }
}
int query(int a, int b) {
    int j = (int) \log 2(b - a + 1);
    return min(v[m[a][j]], v[m[b - (1 << j) + 1][j]]);</pre>
}
     DSU Persistente
// Complexidades:
// build - O(n)
// find - O(log(n))
// une - O(log(n))
```

```
int n, p[MAX], sz[MAX], ti[MAX];
                                                                         seg[2*i+1];
                                                                 }
void build() {
    for (int i = 0; i < n; i++) {</pre>
                                                                 int query(int a, int b) {
        p[i] = i;
                                                                     int ret = 0;
        sz[i] = 1;
        ti[i] = -INF;
   }
}
                                                                      }
                                                                     return ret;
int find(int k, int t) {
                                                                 }
    if (p[k] == k or ti[k] > t) return k;
    return find(p[k], t);
                                                                 void update(int p, int x) {
}
                                                                      seg[p += n] = x;
void une(int a, int b, int t) {
                                                                 }
    a = find(a); b = find(b);
    if (a == b) return;
    if (sz[a] > sz[b]) swap(a, b);
                                                                 // SegTree 1-based
    sz[b] += sz[a];
    p[a] = b;
    ti[a] = t;
}
                                                                 // Complexidades:
     SegTree Iterativa
                                                                 // build - O(n)
                                                                 // query - O(log(n))
                                                                 // update - O(log(n))
// Consultas 0-based
// Valores iniciais devem estar em (seg[n], ..., seg[2*n-1])
// Query: soma do range [a, b]
                                                                 int seg[2*MAX];
// Update: muda o valor da posicao p para x
                                                                 int lazy[2*MAX];
                                                                 int n:
// Complexidades:
// build - O(n)
                                                                 void build() {
// query - O(log(n))
// update - 0(log(n))
                                                                         seg[2*i+1];
int seg[2 * MAX];
int n;
```

void build() {

```
for (int i = n - 1; i; i--) seg[i] = seg[2*i] +
   for(a += n, b += n; a <= b; ++a /= 2, --b /= 2) {
       if (a % 2 == 1) ret += seg[a];
       if (b \% 2 == 0) ret += seg[b];
    while (p /= 2) seg[p] = seg[2*p] + seg[2*p+1];
2.10 SegTree Iterativa com Lazy
// Valores iniciais devem estar em (seg[n], ..., seg[2*n-1])
// Query: soma do range [a, b], 0-based
// Update: soma x em cada elemento do range [a, b], 0-based
    for (int i = n - 1; i; i--) seg[i] = seg[2*i] +
    memset(lazy, 0, sizeof(lazy));
// soma x na posicao p de tamanho tam
void poe(int p, int x, int tam) {
```

```
seg[p] += x * tam:
    if (p < n) lazy[p] += x;
}
// atualiza todos os pais da folha p
void sobe(int p) {
    for (int tam = 2; p /= 2; tam *= 2)
        seg[p] = seg[2*p] + seg[2*p+1] + lazy[p] * tam;
}
// propaga o caminho da raiz ate a folha p
void prop(int p) {
    int tam = 1 << 29;</pre>
    for (int s = 30; s; s--, tam /= 2) {
        int i = p \gg s;
        if (lazy[i]) {
            poe(2*i, lazy[i], tam);
            poe(2*i+1, lazy[i], tam);
            lazy[i] = 0;
       }
}
int query(int a, int b) {
    prop(a += n), prop(b += n);
    int ret = 0;
    for(; a <= b; a /= 2, b /= 2) {
        if (a \% 2 == 1) ret += seg[a++];
        if (b \% 2 == 0) ret += seg[b--];
    }
    return ret;
}
void update(int a, int b, int x) {
    int a2 = a += n, b2 = b += n, tam = 1;
    for (; a <= b; a /= 2, b /= 2, tam *= 2) {
        if (a \% 2 == 1) poe(a++, x, tam);
        if (b \% 2 == 0) poe(b--, x, tam);
    sobe(a2), sobe(b2);
}
```

## 2.11 SegTree Espaca

```
// Query: soma do range [a, b]
// Update: flipa os valores de [a, b]
// Complexidades:
// build - O(n)
// query - O(log^2(n))
// update - 0(log^2(n))
namespace seg {
    unordered_map < int , int > t;
    unordered_map < int , int > lazy;
    void build() { t.clear(), lazy.clear(); }
    void prop(int p, int l, int r) {
        if (!lazy[p]) return;
        t[p] = r-l+1-t[p];
        if (1 != r) lazy[2*p]^=lazy[p], lazy[2*p+1]^=lazy[p];
        lazy[p] = 0;
    }
    int query(int a, int b, int p=1, int l=0, int r=N-1) {
        prop(p, l, r);
        if (b < 1 \text{ or } r < a) \text{ return } 0;
        if (a <= 1 and r <= b) return t[p];</pre>
        int m = 1+r >> 1;
        return query (a, b, 2*p, 1, m) + query (a, b, 2*p+1,
            m+1, r);
    }
    int update(int a, int b, int p=1, int 1=0, int r=N-1) {
        prop(p, 1, r);
        if (b < 1 or r < a) return t[p];</pre>
        if (a <= 1 and r <= b) {</pre>
             lazy[p] ^= 1;
            prop(p, 1, r);
            return t[p];
        }
        int m = 1+r >> 1;
```

```
return t[p] = update(a, b, 2*p, 1, m)+update(a, b,
           2*p+1, m+1, r);
    }
};
2.12 SegTree
// Query: soma do range [a, b]
// Update: soma x em cada elemento do range [a, b]
// Complexidades:
// build - O(n)
// query - O(log(n))
// update - O(log(n))
namespace seg {
    ll seg[4*MAX], lazy[4*MAX];
    int n, *v;
    ll build(int p=1, int l=0, int r=n-1) {
        lazy[p] = 0;
        if (1 == r) return seg[p] = v[1];
        int m = (1+r)/2;
        return seg[p] = build(2*p, 1, m) + build(2*p+1, m+1,
           r);
    void build(int n2, int* v2) {
        n = n2, v = v2;
        build();
    void prop(int p, int 1, int r) {
        seg[p] += lazy[p]*(r-l+1);
        if (1 != r) lazy[2*p] += lazy[p], lazy[2*p+1] +=
           lazy[p];
        lazy[p] = 0;
    }
    11 query(int a, int b, int p=1, int l=0, int r=n-1) {
        prop(p, 1, r);
        if (a <= l and r <= b) return seg[p];</pre>
        if (b < 1 or r < a) return 0;
        int m = (1+r)/2;
        return query (a, b, 2*p, 1, m) + query (a, b, 2*p+1, m)
```

```
m+1, r);
    }
   ll update(int a, int b, int x, int p=1, int l=0, int
       r=n-1) {
        prop(p, 1, r);
        if (a <= 1 and r <= b) {</pre>
            lazy[p] += x;
            prop(p, 1, r);
            return seg[p];
        }
        if (b < 1 or r < a) return seg[p];</pre>
        int m = (1+r)/2;
        return seg[p] = update(a, b, x, 2*p, 1, m) +
            update(a, b, x, 2*p+1, m+1, r);
   }
};
2.13 SegTree 2D
// Consultas 0-based
// Um valor inicial em (x, y) deve ser colocado em
   seg[x+n][y+n]
// Query: soma do retangulo ((x1, y1), (x2, y2))
// Update: muda o valor da posicao (x, y) para val
// Nao pergunte como que essa coisa funciona
//
// Para query com distancia de manhattan <= d, faca
// nx = x+y, ny = x-y
// Update em (nx, ny), query em ((nx-d, ny-d), (nx+d, ny+d))
// Se for de min/max, pode tirar os if's da 'query', e fazer
// sempre as 4 operacoes. Fica mais rapido
//
// Complexidades:
// build - O(n^2)
// query - O(log^2(n))
// update - 0(log^2(n))
int seg[2*MAX][2*MAX], n;
void build() {
```

for (int x = 2\*n; x; x--) for (int y = 2\*n; y; y--) {

```
if (x < n) seg[x][y] = seg[2*x][y] + seg[2*x+1][y];
         if (y < n) seg[x][y] = seg[x][2*y] + seg[x][2*y+1];
}
int query(int x1, int y1, int x2, int y2) {
    int ret = 0, y3 = y1 + n, y4 = y2 + n;
    for (x1 += n, x2 += n; x1 \le x2; ++x1 /= 2, --x2 /= 2)
         for (y1 = y3, y2 = y4; y1 \le y2; ++y1 /= 2, --y2 /=
             2) {
             if (x1\%2 == 1 \text{ and } y1\%2 == 1) \text{ ret } += \text{seg}[x1][y1];
             if (x1\%2 == 1 \text{ and } y2\%2 == 0) \text{ ret } += \text{seg}[x1][y2];
             if (x2\%2 == 0 \text{ and } y1\%2 == 1) \text{ ret } += \text{seg}[x2][y1];
             if (x2\%2 == 0 \text{ and } y2\%2 == 0) \text{ ret } += \text{seg}[x2][y2];
         }
    return ret;
}
void update(int x, int y, int val) {
    int y2 = y += n;
    for (x += n; x; x /= 2, y = y2) {
         if (x \ge n) seg[x][y] = val;
         else seg[x][y] = seg[2*x][y] + seg[2*x+1][y];
         while (y /= 2) seg[x][y] = seg[x][2*y] +
            seg[x][2*y+1];
}
2.14 BIT
// BIT 1-based, v 0-based
// Para mudar o valor da posicao p para x,
// faca: poe(x - query(p, p), p)
// l_bound(x) retorna o menor p tal que
// query(1, p+1) > x (0 based!)
// Complexidades:
// build - O(n)
// poe - O(\log(n))
```

// query - O(log(n))

```
// l_bound - O(log(n))
int n;
int bit[MAX];
int v[MAX];
void build() {
    bit[0] = 0;
    for (int i = 1; i <= n; i++) bit[i] = v[i - 1];</pre>
    for (int i = 1; i <= n; i++) {
        int j = i + (i \& -i);
        if (j <= n) bit[j] += bit[i];</pre>
    }
}
// soma x na posicao p
void poe(int x, int p) {
    for (; p \le n; p += p \& -p) bit [p] += x;
// soma [1, p]
int pref(int p) {
   int ret = 0;
   for (; p; p -= p & -p) ret += bit[p];
    return ret;
}
// soma [a, b]
int query(int a, int b) {
    return pref(b) - pref(a - 1);
}
int l_bound(ll x) {
   int p = 0;
    for (int i = MAX2; i+1; i--) if (p + (1<<i) <= n
        and bit [p + (1 << i)] <= x) x -= bit <math>[p += (1 << i)];
    return p;
}
```

## 2.15 Trie

```
// N deve ser maior ou igual ao numero de nos da trie
// fim indica se alguma palavra acaba nesse no
// Complexidade:
// Inserir e conferir string S -> O(|S|)
// usar static trie T
// T.insert(s) para inserir
// T.find(s) para ver se ta
// T.prefix(s) printa as strings
// que tem s como prefixo
struct trie{
    map < char , int > t[MAX+5];
    int p;
    trie(){
        p = 1;
    void insert(string s){
        s += '$';
        int i = 0;
        for (char c : s){
            auto it = t[i].find(c);
            if (it == t[i].end())
                i = t[i][c] = p++;
            else
                i = it->second;
        }
    }
    bool find(string s){
        s += '$';
        int i = 0;
        for (char c : s){
            auto it = t[i].find(c);
            if (it == t[i].end()) return false;
            i = it->second;
        }
        return true;
    void prefix(string &l, int i){
        if (t[i].find('$') != t[i].end())
            cout << " " << 1 << endl;
```

```
for (auto p : t[i]){
            1 += p.first;
            prefix(l, p.second, k);
            l.pop_back();
        }
    }
    void prefix(string s){
        int i = 0;
        for (char c : s){
            auto it = t[i].find(c);
            if (it == t[i].end()) return;
            i = it->second;
        }
        int k = 0;
        prefix(s, i, k);
   }
};
```

## 3 Matematica

## 3.1 Crivo

```
// "O" crivo
//
// Encontra maior divisor primo
// Um numero eh primo sse div[x] == x
// fact fatora um numero <= lim
// A fatoracao sai ordenada
//
// crivo - O(n log(log(n)))
// fact - O(log(n))

int divi[MAX];

void crivo(int lim) {
   for (int i = 1; i <= lim; i++) divi[i] = 1;

   for (int i = 2; i <= lim; i++) if (divi[i] == 1)
        for (int j = i; j <= lim; j += i) divi[j] = i;
}</pre>
```

```
void fact(vector<int>& v, int n) {
    if (n != divi[n]) fact(v, n/divi[n]);
    v.push_back(divi[n]);
}
// Crivo de divisores
// Encontra numero de divisores
// ou soma dos divisores
// O(n log(n))
int divi[MAX];
void crivo(int lim) {
    for (int i = 1; i <= lim; i++) divi[i] = 1;</pre>
    for (int i = 2; i <= lim; i++)</pre>
        for (int j = i; j <= lim; j += i) {</pre>
            // para numero de divisores
            divi[j]++;
            // para soma dos divisores
             divi[j] += i;
        }
}
// Crivo de totiente
// Encontra o valor da funcao
// totiente de Euler
// O(n log(log(n)))
int tot[MAX];
void crivo(int lim) {
    for (int i = 1; i <= lim; i++) tot[i] = i;</pre>
    for (int i = 2; i <= lim; i++) if (tot[i] == i)</pre>
        for (int j = i; j <= lim; j += i)</pre>
            tot[j] -= tot[j] / i;
```

```
}
    MDC Extendido
// acha x e y tal que ax + by = mdc(a, b)
// O(log(min(a, b)))
int mdce(int a, int b, int *x, int *y){
   if(!a){
        *x = 0:
       *y = 1;
        return b;
   int X, Y;
   int mdc = mdce(b % a, a, &X, &Y);
   *x = Y - (b / a) * X;
    *y = X;
   return mdc;
}
    Ordem Grupo
// O grupo Zn eh ciclico sse n =
// 1, 2, 4, p^k ou 2 p^k, p primo impar
// Retorna -1 se nao achar
// O(sqrt(n) log(n))
int tot(int n); // totiente em O(sqrt(n))
int expo(int a, int b, int m); // (a^b) %m em O(log(b))
// acha todos os divisores ordenados em O(sqrt(n))
vector < int > div(int n) {
    vector<int> ret1, ret2;
   for (int i = 1; i*i <= n; i++) if (n % i == 0) {
        ret1.pb(i);
       if (i*i != n) ret2.pb(n/i);
```

}

```
for (int i = ret2.size()-1; i+1; i--) ret1.pb(ret2[i]);
    return ret1;
}
int ordem(int a, int n) {
    vector < int > v = div(tot(n));
    for (int i : v) if (expo(a, i, n) == 1) return i;
    return -1;
}
3.4 Mod Inverse
// Computa o inverso de a modulo b
// Se b eh primo, basta fazer
// a^{(b-2)}
long long inv(long long a, long long b){
    return 1<a ? b - inv(b%a,a)*b/a : 1;
}
     Miller-Rabin
3.5
// Testa se n eh primo, n <= 3 * 10^18
// O(log(n)), considerando multiplicacao
// e exponenciacao constantes
// multiplicacao modular
ll mul(ll x, ll y, ll m); // x*y mod m
ll exp(ll x, ll y, ll m); // x^y mod m;
bool prime(ll n) {
    if (n < 2) return 0;
    if (n <= 3) return 1;
    if (n \% 2 == 0) return 0;
```

11 d = n - 1;

while (d % 2 == 0) r++, d /= 2;

int r = 0;

```
// com esses primos, o teste funciona garantido para n
       <= 3*10^18
    // funciona para n <= 3*10^24 com os primos ate 41
    int a[9] = {2, 3, 5, 7, 11, 13, 17, 19, 23};
    // outra opcao para n <= 2^64:
    // int a[7] = {2, 325, 9375, 28178, 450775, 9780504,
       1795265022};
    for (int i = 0; i < 9; i++) {
        if (a[i] >= n) break;
        ll x = exp(a[i], d, n);
        if (x == 1 \text{ or } x == n - 1) \text{ continue};
        bool deu = 1;
        for (int j = 0; j < r - 1; j++) {
            x = mul(x, x, n);
            if (x == n - 1) {
                deu = 0;
                break;
            }
        }
        if (deu) return 0;
    }
    return 1;
3.6 Pollard's Rho
// Usa o algoritmo de deteccao de ciclo de Brent
// A fatoracao nao sai necessariamente ordenada
// O algoritmo rho encontra um fator de n,
// e funciona muito bem quando n possui um fator pequeno
// Eh recomendado chamar srand(time(NULL)) na main
// Complexidades (considerando mul constante):
// rho - esperado O(n^{(1/4)}) no pior caso
// fact - esperado menos que O(n^{(1/4)} \log(n)) no pior caso
11 mdc(11 a, 11 b) { return !b ? a : mdc(b, a % b); }
11 mul(11 x, 11 y, 11 m) {
    if (!y) return 0;
```

```
ll ret = mul(x, y >> 1, m);
    ret = (ret + ret) % m;
    if (y & 1) ret = (ret + x) % m;
    return ret;
}
ll exp(ll x, ll y, ll m) {
    if (!y) return 1;
    ll ret = exp(x, y >> 1, m);
    ret = mul(ret, ret, m);
    if (y & 1) ret = mul(ret, x, m);
    return ret;
}
bool prime(ll n) {
    if (n < 2) return 0;
    if (n <= 3) return 1;
    if (n % 2 == 0) return 0;
    11 d = n - 1;
    int r = 0;
    while (d \% 2 == 0) \{
        r++;
        d /= 2;
    }
    int a[9] = {2, 3, 5, 7, 11, 13, 17, 19, 23};
    for (int i = 0; i < 9; i++) {
        if (a[i] >= n) break;
        ll x = exp(a[i], d, n);
        if (x == 1 \text{ or } x == n - 1) \text{ continue};
        bool deu = 1;
        for (int j = 0; j < r - 1; j++) {
            x = mul(x, x, n);
            if (x == n - 1) {
                 deu = 0;
                 break:
            }
        }
```

```
if (deu) return 0:
    }
    return 1;
}
11 rho(11 n) {
    if (n == 1 or prime(n)) return n;
    if (n % 2 == 0) return 2;
    while (1) {
        11 x = 2, y = 2;
        11 \text{ ciclo} = 2, i = 0;
        ll c = (rand() / (double) RAND_MAX) * (n - 1) + 1;
        11 d = 1:
        while (d == 1) {
            if (++i == ciclo) ciclo *= 2, y = x;
            x = (mul(x, x, n) + c) \% n;
            if (x == y) break;
            d = mdc(abs(x - y), n);
        }
        if (x != y) return d;
    }
}
void fact(ll n, vector<ll>& v) {
    if (n == 1) return;
    if (prime(n)) v.pb(n);
    else {
        ll d = rho(n);
        fact(d, v);
        fact(n / d, v);
    }
}
     Fast Pow
// (x^v mod m) em O(log(v))
```

```
typedef long long int 11;
ll pow(ll x, ll y, ll m) { // iterativo
    ll ret = 1;
    while (y) {
        if (y & 1) ret = (ret * x) % m;
        v >>= 1;
        x = (x * x) % m;
    return ret;
}
ll pow(ll x, ll y, ll m) \{ // \text{ recursivo} \}
    if (y == 0) return 1;
    ll ret = pow(x, y / 2, m);
    ret = (ret * ret) % m;
    if (y & 1) ret = (ret * x) % m;
    return ret;
}
     Totiente
3.8
// O(sqrt(n))
int tot(int n){
    int ret = n;
    for (int i = 2; i*i <= n; i++) if (n % i == 0) {
        while (n \% i == 0) n /= i;
        ret -= ret / i;
    if (n > 1) ret -= ret / n;
    return ret;
}
     Div. Pol.
// Divide p1 por p2
```

```
// Retorna um par com o quociente e o resto
// Os coeficientes devem estar em ordem
// decrescente pelo grau. Ex:
// 3x^2 + 2x - 1 \rightarrow [3, 2, -1]
// O(nm), onde n e m sao os tamanhos dos
// polinomios
typedef vector<int> vi;
pair < vi , vi > div(vi p1, vi p2) {
    vi quoc, resto;
    int a = p1.size(), b = p2.size();
    for (int i = 0; i <= a - b; i++) {</pre>
        int k = p1[i] / p2[0];
        quoc.pb(k);
        for (int j = i; j < i + b; j++)</pre>
            p1[j] = k * p2[j - i];
    }
    for (int i = a - b + 1; i < a; i++)
        resto.pb(p1[i]);
    return mp(quoc, resto);
}
    Problemas
4.1 MO
// O(n sqrt(n) + q)
void add(int pos){
    occ[a[pos]]++;
    counter += (occ[a[pos]] == 1);
}
void remove(int pos){
    occ[a[pos]]--;
    counter -= (occ[a[pos]] == 0);
```

```
}
vector<pii> query(q);
vector<pair<pii, int>> s(q);
for (int i = 0; i < q; i++){</pre>
    int 1, r;
    scanf("%d%d", &1, &r);
    1--; r--;
    query[i] = pii(1, r);
    s[i] = \{\{1/SQ, r\}, i\};
sort(s.begin(), s.end()); //sort queries
for (int i = 0; i < q; i++){</pre>
    int iq = s[i].second;
    pii q = query[iq];
    while (L < q.first){</pre>
        remove(L);
        L++;
    while (L > q.first){
        L--;
        add(L);
    while (R < q.second){</pre>
        R++;
        add(R);
    while (R > q.second){
        remove(R);
        R--;
    ans[iq] = counter;
}
     MergeSort
// Melhor do Brasil, segundo o autor
// O(n log(n))
long long merge_sort(int 1, int r, vector<int> &t){
    if (1 >= r) return 0;
```

```
int m = (1+r)/2;
    auto ans = merge_sort(1, m, t) + merge_sort(m+1, r, t);
    static vector<int> aux; if (aux.size() != t.size())
       aux.resize(t.size());
    for (int i = 1; i <= r; i++) aux[i] = t[i];</pre>
    int i_l = l, i_r = m+1, i = l;
    auto move_1 = [&](){
        t[i++] = aux[i_1++];
    };
    auto move_r = [\&](){
        t[i++] = aux[i_r++];
    };
    while (i \le r){
        if (i_l > m) move_r();
        else if (i_r > r) move_l();
        else{
            if (aux[i_1] <= aux[i_r]) move_1();</pre>
             else{
                 move_r();
                 ans += m - i_1 + 1;
            }
        }
    }
    return ans;
//inversions to turn r into l
template < typename T > 11 inv_count(vector < T > &1, vector < T >
   &r){
    int n = 1.size();
    map < T , int > occ;
    map<pair<T, int>, int> rk;
    for (int i = 0; i < n; i++)</pre>
        rk[make_pair(1[i], occ[1[i]]++)] = i;
    occ.clear();
    vector < int > v(n);
    for (int i = 0; i < n; i++)</pre>
        v[i] = rk[make_pair(r[i], occ[r[i]]++)];
    return merge_sort(0, n-1, v);
```

```
}
```

## 4.3 Min. Circ. Vasek

```
// O(n) com alta probabilidade
const long double EPS = 1e-12;
struct pt {
    long double x, y;
    pt() {}
    pt(long double x, long double y) : x(x), y(y) {}
    pt(const pt& p) : x(p.x), y(p.y) {}
    pt operator + (const pt& p) const { return pt(x+p.x,
       y+p.y); }
    pt operator - (const pt& p) const { return pt(x-p.x,
       y-p.y); }
    pt operator * (long double c) const { return pt(x*c, y*c
       ); }
    pt operator / (long double c) const { return pt(x/c, y/c
       ); }
};
long double dot(pt p, pt q) { return p.x*q.x+p.y*q.y; }
long double dist2(pt p, pt q) { return dot(p-q, p-q); }
long double cross(pt p, pt q) { return p.x*q.y-p.y*q.x; }
pt rotate90(pt p) { return pt(p.y, -p.x); }
pt interline(pt a, pt b, pt c, pt d) {
    b = b-a: d = c-d: c = c-a:
    return a+b*cross(c, d)/cross(b, d);
}
pt center(pt a, pt b, pt c) {
    b = (a+b)/2;
    c = (a+c)/2;
    return interline(b, b+rotate90(a-b), c, c+rotate90(a-c));
}
struct circle {
    pt cen;
```

```
long double r;
    circle() {}
    circle(pt cen, long double r) : cen(cen), r(r) {}
};
bool inside(circle& c, pt& p) {
    return c.r*c.r+1e-9 > dist2(p, c.cen);
pt bestof3(pt a, pt b, pt c) {
    if (dot(b-a, c-a) < 1e-9) return (b+c)/2;
    if (dot(a-b, c-b) < 1e-9) return (a+c)/2;
    if (dot(a-c, b-c) < 1e-9) return (a+b)/2;
    return center(a, b, c);
}
circle minCirc(vector<pt> v) {
    int n = v.size();
    random_shuffle(v.begin(), v.end());
    pt p = pt(0, 0);
    circle ret = circle(p, 0);
    for (int i = 0; i < n; i++) if (!inside(ret, v[i])) {</pre>
        ret = circle(v[i], 0);
        for (int j = 0; j < i; j++) if (!inside(ret, v[j])) {</pre>
            ret = circle((v[i]+v[j])/2, sqrt(dist2(v[i],
                v[i]))/2);
            for (int k = 0; k < j; k++) if (!inside(ret,</pre>
                v[k])) {
                p = bestof3(v[i], v[j], v[k]);
                ret = circle(p, sqrt(dist2(p, v[i])));
            }
        }
    }
    return ret:
}
4.4 nim
// Calcula movimento otimo do jogo classico de Nim
// Assume que o estado atual eh perdedor
// Funcao move retorna um par com a pilha (0 indexed)
// e quanto deve ser tirado dela
```

```
// XOR deve estar armazenado em x
// Para mudar um valor, faca insere(novo_valor),
// atualize o XOR e mude o valor em v
//
// MAX2 = teto do log do maior elemento
// possivel nas pilhas
// O(log(n)) amortizado
int v[MAX], n, x;
stack<int> pi[MAX2];
void insere(int p) {
    for (int i = 0; i < MAX2; i++) if (v[p] & (1 << i))
       pi[i].push(p);
}
pair < int , int > move() {
    int bit = 0; while (x >> bit) bit++; bit--;
    // tira os caras invalidos
    while ((v[pi[bit].top()] & (1 << bit)) == 0)</pre>
       pi[bit].pop();
    int cara = pi[bit].top();
    int tirei = v[cara] - (x^v[cara]);
    v[cara] -= tirei;
    insere(cara);
    return make_pair(cara, tirei);
}
// Acha o movimento otimo baseado
// em v apenas
//
// O(n)
pair < int , int > move() {
    int x = 0;
    for (int i = 0; i < n; i++) x ^= v[i];</pre>
```

```
for (int i = 0; i < n; i++) if ((v[i]^x) < v[i])</pre>
        return make_pair(i, v[i] - (v[i]^x));
}
4.5 LIS
// Calcula uma LIS
// Para ter o tamanho basta fazer lis().size()
// Implementacao do algotitmo descrito em:
// https://goo.gl/HiFkn2
//
// O(n log(n))
const int INF = 0x3f3f3f3f;
int n, v[MAX];
vector<int> lis() {
    int I[n + 1], L[n];
    // pra BB funfar bacana
    I[0] = -INF;
    for (int i = 1; i <= n; i++) I[i] = INF;</pre>
    for (int i = 0; i < n; i++) {
        // BB
        int 1 = 0, r = n;
        while (1 < r) {</pre>
            int m = (1 + r) / 2;
            if (I[m] >= v[i]) r = m;
            else 1 = m + 1;
        }
        // ultimo elemento com tamanho l eh v[i]
        I[1] = v[i]:
        // tamanho da LIS terminando com o
        // elemento v[i] eh l
        L[i] = 1;
    }
    // reconstroi LIS
    vector<int> ret;
```

```
int m = -INF, p;
for (int i = 0; i < n; i++) if (L[i] > m) {
    m = L[i];
    p = i;
}
ret.push_back(v[p]);
int last = m;
while (p--) if (L[p] == m - 1) {
    ret.push_back(v[p]);
    m = L[p];
}
reverse(ret.begin(), ret.end());
return ret;
```

#### 4.6 InversionCount

}

```
// O(n log(n))
int n;
int v[MAX];
// bit de soma
void poe(int p);
int query(int p);
// converte valores do array pra
// numeros de 1 a n
void conv() {
    vector<int> a:
    for (int i = 0; i < n; i++) a.push_back(v[i]);</pre>
    sort(a.begin(), a.end());
    for (int i = 0; i < n; i++)</pre>
        v[i] = 1 + (lower_bound(a.begin(), a.end(), v[i]) -
            a.begin());
}
long long inv() {
    conv();
```

```
build();
    long long ret = 0;
    for (int i = n - 1; i >= 0; i--) {
        ret += query(v[i] - 1);
        poe(v[i]);
    }
    return ret;
}
     Min Fixed Range
//ans[i] = min_{0} <= j < k v[i+j]
vector<int> min_k(vector<int> &v, int k){
    int n = v.size();
    deque < int > d;
    auto put = [&](int i){
        while (!d.empty() && v[d.back()] > v[i])
            d.pop_back();
        d.push_back(i);
    for (int i = 0; i < k-1; i++)</pre>
        put(i);
    vector < int > ans (n-k+1);
    for (int i = 0; i < n-k+1; i++) {
        put(i+k-1);
        while (i > d.front()) d.pop_front();
        ans[i] = v[d.front()];
    }
    return ans;
}
4.8 LIS 2
// O(n log(n))
template < typename T > int lis(vector < T > &v) {
    vector <T> ans;
    for (T t : v){
        auto it = upper_bound(ans.begin(), ans.end(), t);
```

```
if (it == ans.end()) ans.push_back(t);
        else *it = t;
    return ans.size()
}
     Min. Circ.
// O(n) com alta probabilidade
const double EPS = 1e-12;
mt19937 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
struct pt {
    double x, y;
    pt(double x_{=} = 0, double y_{=} = 0) : x(x_{=}), y(y_{=}) {}
    pt operator + (const pt& p) const { return pt(x+p.x,
       y+p.y); }
    pt operator - (const pt& p) const { return pt(x-p.x,
       y-p.y); }
    pt operator * (double c) const { return pt(x*c, y*c); }
    pt operator / (double c) const { return pt(x/c, y/c); }
};
double dot(pt p, pt q) { return p.x*q.x+p.y*q.y; }
double cross(pt p, pt q) { return p.x*q.y-p.y*q.x; }
double dist(pt p, pt q) { return sqrt(dot(p-q, p-q)); }
pt center(pt p, pt q, pt r) {
    pt a = p-r, b = q-r;
    pt c = pt(dot(a, p+r)/2, dot(b, q+r)/2);
    return pt(cross(c, pt(a.y, b.y)), cross(pt(a.x, b.x),
       c)) / cross(a, b);
}
struct circle {
    pt cen;
    double r;
    circle(pt cen_, double r_) : cen(cen_), r(r_) {}
    circle(pt a, pt b, pt c) {
        cen = center(a, b, c);
```

```
r = dist(cen. a):
    bool inside(pt p) { return dist(p, cen) < r+EPS; }</pre>
};
circle minCirc(vector<pt> v) {
    shuffle(v.begin(), v.end(), rng);
    circle ret = circle(pt(0, 0), 0);
    for (int i = 0; i < v.size(); i++) if</pre>
       (!ret.inside(v[i])) {
        ret = circle(v[i], 0);
        for (int j = 0; j < i; j++) if (!ret.inside(v[j])) {</pre>
            ret = circle((v[i]+v[j])/2, dist(v[i], v[j])/2);
            for (int k = 0; k < j; k++) if
                (!ret.inside(v[k]))
                ret = circle(v[i], v[j], v[k]);
        }
    }
    return ret;
}
4.10 Area Histograma
// Assume que todas as barras tem largura 1,
// e altura dada no vetor v
//
// O(n)
typedef long long 11;
11 area(vector<int> v) {
    11 \text{ ret} = 0:
    stack<int> s;
    // valores iniciais pra dar tudo certo
    v.insert(v.begin(), -1);
    v.insert(v.end(), -1);
    s.push(0);
    for(int i = 0; i < (int) v.size(); i++) {</pre>
        while (v[s.top()] > v[i]) {
            11 h = v[s.top()]; s.pop();
            ret = \max(\text{ret}, h * (i - s.top() - 1));
```

```
}
        s.push(i);
    return ret;
}
4.11 CHT
// linear
struct CHT {
    int it;
    vector<ll> a, b;
    CHT():it(0){}
    ll eval(int i, ll x){
        return a[i]*x + b[i];
    }
    bool useless(){
        int sz = a.size();
        int r = sz-1, m = sz-2, 1 = sz-3;
        return (b[1] - b[r])*(a[m] - a[1]) <
            (b[1] - b[m])*(a[r] - a[1]);
    void add(ll A, ll B){
        a.push_back(A); b.push_back(B);
        while (!a.empty()){
            if ((a.size() < 3) || !useless()) break;</pre>
            a.erase(a.end() - 2);
            b.erase(b.end() - 2);
        }
    }
    11 get(11 x){
        it = min(it, int(a.size()) - 1);
        while (it+1 < a.size()){</pre>
            if (eval(it+1, x) > eval(it, x)) it++;
            else break;
        }
        return eval(it, x);
    }
};
```

## 5 Papa

## 5.1 LIS Rec Resp

```
#include <bits/stdc++.h>
using namespace std;
#define sc(a) scanf("%d", &a)
typedef long long int 11;
const int INF = 0x3f3f3f3f;
#define MAXN 100100
int aux[MAXN], endLis[MAXN];
//usar upper_bound se puder >=
vector < int > LisRec(vector < int > v) {
    int n=v.size();
    int lis=0;
    for (int i = 0; i < n; i++){
        int it = lower_bound(aux, aux+lis, v[i]) - aux;
        endLis[i] = it+1;
        lis = max(lis, it+1);
        aux[it] = v[i];
    }
    vector < int > resp;
    int prev=INF;
    for(int i=n-1;i>=0;i--){
        if(endLis[i] == lis && v[i] <= prev){</pre>
            lis--;
            prev=v[i];
            resp.push_back(i);
        }
    reverse(resp.begin(),resp.end());
    return resp;
}
int main()
    int n;
    sc(n);
    vector<int> v(n);
```

```
for(int i=0;i<n;i++)</pre>
         sc(v[i]);
     cout << LisRec(v).size() << endl;</pre>
     return 0;
}
      Aho Corasick
```

```
const int N=100010;
const int M=26;
//N= tamanho da trie, M tamanho do alfabeto
int to[N][M], Link[N], fim[N];
int idx = 1;
void add_str(string &s)
{
    int v = 0;
    for (int i = 0; i < s.size(); i++) {</pre>
        if (!to[v][s[i]]) to[v][s[i]] = idx++;
        v = to[v][s[i]];
    fim[v] = 1;
}
void process()
{
    queue < int > fila;
    fila.push(0);
    while (!fila.empty()) {
        int cur = fila.front();
        fila.pop();
        int 1 = Link[cur];
        fim[cur] |= fim[1];
        for (int i = 0; i < M; i++) {</pre>
            if (to[cur][i]) {
                if (cur != 0) {
                     Link[to[cur][i]] = to[1][i];
                }
                 else
                     Link[to[cur][i]] = 0;
                fila.push(to[cur][i]);
            }
            else {
```

```
to[cur][i] = to[1][i];
        }
    }
}
int resolve(string &s)
    int v = 0, r = 0;
    for (int i = 0; i < s.size(); i++) {</pre>
        v = to[v][s[i]];
        if (fim[v]) r++, v = 0;
    }
    return r;
}
    Baby-step Giant-step
//Resolve Logaritmo Discreto a^x = b mod m, m primo em
   0(sqrt(n)*hash(n))
//Meet In The Middle, decompondo x = i * ceil(sqrt(n)) - j,
   i,j<=ceil(sqrt(n))
int babyStep(int a,int b,int m)
    unordered_map < int , int > mapp;
    int sq=sqrt(m)+1;
    ll asq=1;
    for(int i=0; i<sq; i++)</pre>
        asq=(asq*a)%m;
    11 curr=asq;
    for(int i=1; i <= sq; i++)</pre>
        if(!mapp.count(curr))
             mapp[curr]=i;
        curr = (curr * asq) %m;
    }
    int ret=INF;
    curr=b;
    for(int j=0; j<=sq; j++)</pre>
        if (mapp.count(curr))
             ret=min(ret,(int)(mapp[curr]*sq-j));
```

```
curr = (curr * a) %m;
    if(ret<INF) return ret;</pre>
    return -1;
}
int main()
    int a,b,m;
    while(cin>>a>>b>>m,a or b or m)
         int x=babyStep(a,m,b);
         if(x!=-1)
             cout << x << endl;</pre>
         else
             cout << "No Solution" << endl;</pre>
    }
    return 0;
}
5.4 BIT Persistent
#include < bits / stdc++.h>
using namespace std;
typedef long long int 11;
const 11 LINF = 0x3f3f3f3f3f3f3f3f3f11;
#define MAXN 100010
vector < pair < int , ll > > FT [MAXN];
int n:
void clear()
    for(int i=1;i<=n;i++)</pre>
         FT[i].clear();
         FT[i].push_back({-1,0});
    }
}
void add(int i,int v,int time)
{
```

for (; i <= n; i += i & (-i))

```
{
        11 last=FT[i].back().second;
        FT[i].push_back({time,last+v});
    }
}
ll get(int i,int time)
    ll ret=0;
    for(;i>0;i-=i&(-i))
           pos=upper_bound(FT[i].begin(),FT[i].end(),make_pair(time)
        ret+=FT[i][pos].second;
    }
    return ret;
}
11 getRange(int a,int b,int time)
{
    return get(b, time) - get(a-1, time);
}
    Grafos
6.1 LCA com HLD
// Assume que um vertice eh ancestral dele mesmo, ou seja,
// se a eh ancestral de b, lca(a, b) = a
// Para buildar pasta chamar build(root)
//
// Complexidades:
// build - O(n)
// lca - O(log(n))
vector < vector < int > > g(MAX);
int in[MAX], h[MAX], sz[MAX];
int pai[MAX], t;
void build(int k, int p = -1, int f = 1) {
    in[k] = t++; sz[k] = 1;
    for (int& i : g[k]) if (i != p) {
```

```
pai[i] = k;
        h[i] = (i == g[k][0] ? h[k] : i);
        build(i, k, f); sz[k] += sz[i];
        if (sz[i] > sz[g[k][0]]) swap(i, g[k][0]);
    if (p*f == -1) t = 0, h[k] = k, build(k, -1, 0);
}
int lca(int a, int b) {
    if (in[a] < in[b]) swap(a, b);</pre>
    return h[a] == h[b] ? b : lca(pai[h[a]], b);
}
6.2 HLD-aresta
// SegTree de soma
// query / update de soma das arestas
//
// Complexidades:
// build - O(n)
// \text{ query_path - } O(\log^2 (n))
// update_path - O(log^2 (n))
// query_subtree - O(log(n))
// update_subtree - O(log(n))
#define f first
#define s second
namespace seg {
    11 seg[4*MAX], lazy[4*MAX];
    int n, *v;
    ll build(int p=1, int l=0, int r=n-1) {
        lazy[p] = 0;
        if (1 == r) return seg[p] = v[1];
        int m = (1+r)/2;
        return seg[p] = build(2*p, 1, m) + build(2*p+1, m+1,
    }
    void build(int n2, int* v2) {
        n = n2, v = v2;
```

```
build();
    }
    void prop(int p, int l, int r) {
        seg[p] += lazy[p]*(r-l+1);
        if (1 != r) lazy[2*p] += lazy[p], lazy[2*p+1] +=
           lazv[p];
        lazv[p] = 0;
    }
    11 query(int a, int b, int p=1, int l=0, int r=n-1) {
        prop(p, 1, r);
        if (a <= l and r <= b) return seg[p];</pre>
        if (b < 1 or r < a) return 0;
        int m = (1+r)/2;
        return query (a, b, 2*p, 1, m) + query (a, b, 2*p+1, m)
           m+1, r);
   }
    ll update(int a, int b, int x, int p=1, int l=0, int
       r=n-1) {
        prop(p, l, r);
        if (a <= 1 and r <= b) {
            lazy[p] += x;
            prop(p, 1, r);
            return seg[p];
        if (b < l or r < a) return seg[p];</pre>
        int m = (1+r)/2;
        return seg[p] = update(a, b, x, 2*p, 1, m) +
            update(a, b, x, 2*p+1, m+1, r);
   }
};
namespace hld {
    vector < pair < int , int > > g[MAX];
    int in[MAX], out[MAX], sz[MAX];
   int sobe[MAX], pai[MAX];
    int h[MAX], v[MAX], t;
    void build_hld(int k, int p = -1, int f = 1) {
        v[in[k] = t++] = sobe[k]; sz[k] = 1;
        for (auto& i : g[k]) if (i.f != p) {
            sobe[i.f] = i.s; pai[i.f] = k;
            h[i.f] = (i == g[k][0] ? h[k] : i.f);
```

```
build_hld(i.f, k, f); sz[k] += sz[i.f];
            if (sz[i.f] > sz[g[k][0].f]) swap(i, g[k][0]);
        }
        out[k] = t;
        if (p*f == -1) build_hld(h[k] = k, -1, t = 0);
    void build(int root = 0) {
        t = 0:
        build_hld(root);
        seg::build(t, v);
    11 query_path(int a, int b) {
        if (a == b) return 0;
        if (in[a] < in[b]) swap(a, b);</pre>
        if (h[a] == h[b]) return seg::query(in[b]+1, in[a]);
        return seg::query(in[h[a]], in[a]) +
           query_path(pai[h[a]], b);
    void update_path(int a, int b, int x) {
        if (a == b) return;
        if (in[a] < in[b]) swap(a, b);</pre>
        if (h[a] == h[b]) return (void)seg::update(in[b]+1,
           in[a], x);
        seg::update(in[h[a]], in[a], x);
           update_path(pai[h[a]], b, x);
    }
    11 query_subtree(int a) {
        if (in[a] == out[a]-1) return 0;
        return seg::query(in[a]+1, out[a]-1);
    }
    void update_subtree(int a, int x) {
        if (in[a] == out[a]-1) return;
        seg::update(in[a]+1, out[a]-1, x);
    }
    int lca(int a, int b) {
        if (in[a] < in[b]) swap(a, b);</pre>
        return h[a] == h[b] ? b : lca(pai[h[a]], b);
};
```

## 6.3 HLD-vertice

```
// SegTree de soma
// query / update de soma dos vertices
// Complexidades:
// build - O(n)
// \text{ query_path - } O(\log^2 (n))
// update_path - O(log^2 (n))
// query_subtree - O(log(n))
// update_subtree - O(log(n))
namespace seg {
    11 \text{ seg}[4*MAX], lazy[4*MAX];
    int n, *v;
    ll build(int p=1, int l=0, int r=n-1) {
        lazv[p] = 0;
        if (1 == r) return seg[p] = v[1];
        int m = (1+r)/2;
        return seg[p] = build(2*p, 1, m) + build(2*p+1, m+1,
           r);
    }
    void build(int n2, int* v2) {
        n = n2, v = v2;
        build();
    }
    void prop(int p, int l, int r) {
        seg[p] += lazy[p]*(r-l+1);
        if (1 != r) lazy[2*p] += lazy[p], lazy[2*p+1] +=
           lazy[p];
        lazy[p] = 0;
    }
    ll query(int a, int b, int p=1, int l=0, int r=n-1) {
        prop(p, 1, r);
        if (a <= l and r <= b) return seg[p];</pre>
        if (b < 1 or r < a) return 0;
        int m = (1+r)/2;
        return query(a, b, 2*p, 1, m) + query(a, b, 2*p+1,
           m+1, r);
    }
    ll update(int a, int b, int x, int p=1, int l=0, int
```

```
r=n-1) {
        prop(p, 1, r);
        if (a <= 1 and r <= b) {
            lazy[p] += x;
            prop(p, 1, r);
            return seg[p];
        }
        if (b < l or r < a) return seg[p];</pre>
        int m = (1+r)/2;
        return seg[p] = update(a, b, x, 2*p, 1, m) +
            update(a, b, x, 2*p+1, m+1, r);
   }
};
namespace hld {
    vector < int > g[MAX];
    int in[MAX], out[MAX], sz[MAX];
    int peso[MAX], pai[MAX];
    int h[MAX], v[MAX], t;
    void build_hld(int k, int p = -1, int f = 1) {
        v[in[k] = t++] = peso[k]; sz[k] = 1;
        for (auto& i : g[k]) if (i != p) {
            pai[i] = k;
            h[i] = (i == g[k][0] ? h[k] : i);
            build_hld(i, k, f); sz[k] += sz[i];
            if (sz[i] > sz[g[k][0]]) swap(i, g[k][0]);
        }
        out[k] = t:
        if (p*f == -1) build_hld(h[k] = k, -1, t = 0);
    void build(int root = 0) {
        t = 0:
        build_hld(root);
        seg::build(t, v);
    }
    11 query_path(int a, int b) {
        if (a == b) return seg::query(in[a], in[a]);
        if (in[a] < in[b]) swap(a, b);</pre>
        if (h[a] == h[b]) return seg::query(in[b], in[a]);
```

```
return seg::query(in[h[a]], in[a]) +
           query_path(pai[h[a]], b);
   }
    void update_path(int a, int b, int x) {
       if (a == b) return (void)seg::update(in[a], in[a],
           x);
       if (in[a] < in[b]) swap(a, b);</pre>
       if (h[a] == h[b]) return (void)seg::update(in[b],
           in[a], x);
       seg::update(in[h[a]], in[a], x);
           update_path(pai[h[a]], b, x);
   }
    11 query_subtree(int a) {
       if (in[a] == out[a]-1) return seg::query(in[a],
       return seg::query(in[a], out[a]-1);
   }
    void update_subtree(int a, int x) {
       if (in[a] == out[a]-1) return
           (void) seg::update(in[a], in[a], x);
        seg::update(in[a], out[a]-1, x);
   }
   int lca(int a, int b) {
       if (in[a] < in[b]) swap(a, b);</pre>
       return h[a] == h[b] ? b : lca(pai[h[a]], b);
   }
};
6.4 LCA com RMQ
// Assume que um vertice eh ancestral dele mesmo, ou seja,
// se a eh ancestral de b, lca(a, b) = a
//
// Complexidades:
// build - O(n) + build_RMQ
// lca - RMQ
int n;
vector < vector < int > > g(MAX);
aparicao
```

```
int ord[2 * MAX]: // ord[i] : i-esimo vertice na ordem de
   visitacao da dfs
int v[2 * MAX]; // vetor de alturas que eh usado na RMQ
int p;
void dfs(int k, int l) {
    ord[p] = k;
    pos[k] = p;
    v[p++] = 1;
    for (int i = 0; i < (int) g[k].size(); i++)</pre>
        if (pos[g[k][i]] == -1) {
            dfs(g[k][i], l + 1);
            ord[p] = k;
            v[p++] = 1;
        }
}
void build(int root) {
    for (int i = 0; i < n; i++) pos[i] = -1;</pre>
    p = 0;
    dfs(root, 0);
    build_RMQ();
}
int lca(int u, int v) {
    int a = pos[u], b = pos[v];
    if (a > b) swap(a, b);
    return ord[RMQ(a, b)];
}
6.5 HLD sem Update
// query de min do caminho
//
// Complexidades:
// build - O(n)
// query_path - O(log(n))
#define f first
```

#define s second

```
namespace hld {
    vector < pair < int , int > > g[MAX];
    int in[MAX], sz[MAX];
    int sobe[MAX], pai[MAX];
    int h[MAX], v[MAX], t;
    int men[MAX], seg[2*MAX];
    void build_hld(int k, int p = -1, int f = 1) {
        v[in[k] = t++] = sobe[k]; sz[k] = 1;
        for (auto& i : g[k]) if (i.f != p) {
            sobe[i.f] = i.s; pai[i.f] = k;
            h[i.f] = (i == g[k][0] ? h[k] : i.f);
            men[i.f] = (i == g[k][0] ? min(men[k], i.s) :
               i.s):
            build_hld(i.f, k, f); sz[k] += sz[i.f];
            if (sz[i.f] > sz[g[k][0].f]) swap(i, g[k][0]);
        if (p*f == -1) build_hld(h[k] = k, -1, t = 0);
    }
    void build(int root = 0) {
        t = 0;
        build_hld(root);
        for (int i = 0; i < t; i++) seg[i+t] = v[i];</pre>
        for (int i = t-1; i; i--) seg[i] = min(seg[2*i],
           seg[2*i+1]);
    }
    int query_path(int a, int b) {
        if (a == b) return INF;
        if (in[a] < in[b]) swap(a, b);
        if (h[a] != h[b]) return min(men[a],
            query_path(pai[h[a]], b));
        int ans = INF, x = in[b]+1+t, y = in[a]+t;
        for (; x \le y; ++x/=2, --y/=2) ans = min({ans,
           seg[x], seg[y]);
        return ans;
    }
};
```

#### 6.6 LCA

```
// Assume que um vertice eh ancestral dele mesmo, ou seja,
// se a eh ancestral de b, lca(a, b) = a
// MAX2 = ceil(log(MAX))
//
// Complexidades:
// build - O(n log(n))
// lca - O(log(n))
vector < vector < int > > g(MAX);
int n, p;
int pai[MAX2][MAX];
int in[MAX], out[MAX];
void dfs(int k) {
    in[k] = p++;
    for (int i = 0; i < (int) g[k].size(); i++)</pre>
        if (in[g[k][i]] == -1) {
            pai[0][g[k][i]] = k;
            dfs(g[k][i]);
    out[k] = p++;
}
void build(int raiz) {
    for (int i = 0; i < n; i++) pai[0][i] = i;</pre>
    p = 0, memset(in, -1, sizeof in);
    dfs(raiz);
    // pd dos pais
    for (int k = 1; k < MAX2; k++) for (int i = 0; i < n;
       i++)
        pai[k][i] = pai[k - 1][pai[k - 1][i]];
}
bool anc(int a, int b) { // se a eh ancestral de b
    return in[a] <= in[b] and out[a] >= out[b];
}
int lca(int a, int b) {
    if (anc(a, b)) return a;
```

```
if (anc(b, a)) return b;
    // sobe a
    for (int k = MAX2 - 1; k >= 0; k--)
        if (!anc(pai[k][a], b)) a = pai[k][a];
    return pai[0][a];
}
6.7 Blossom
// O(n^3)
// Se for bipartido, nao precisa da funcao
// 'contract', e roda em O(nm)
vector < vector < int > > g(MAX);
int match[MAX]; // match[i] = com quem i esta matchzado ou -1
int n, pai[MAX], base[MAX], vis[MAX];
queue < int > q;
void contract(int u, int v, bool first = 1) {
    static vector < bool > bloss;
    static int 1;
    if (first) {
        bloss = vector < bool > (n, 0);
        vector < bool > teve(n, 0);
        int k = u; l = v;
        while (1) {
            teve[k = base[k]] = 1;
            if (match[k] == -1) break;
            k = pai[match[k]];
        while (!teve[l = base[l]]) l = pai[match[l]];
    }
    while (base[u] != 1) {
        bloss[base[u]] = bloss[base[match[u]]] = 1;
        pai[u] = v;
        v = match[u];
        u = pai[match[u]];
    }
    if (!first) return;
    contract(v, u, 0);
```

```
for (int i = 0; i < n; i++) if (bloss[base[i]]) {</pre>
        base[i] = 1;
        if (!vis[i]) q.push(i);
        vis[i] = 1;
    }
}
int getpath(int s) {
    for (int i = 0; i < n; i++) base[i] = i, pai[i] = -1,</pre>
       vis[i] = 0;
    vis[s] = 1; q = queue < int > (); q.push(s);
    while (q.size()) {
        int u = q.front(); q.pop();
        for (int i : g[u]) {
            if (base[i] == base[u] or match[u] == i)
                continue:
            if (i == s or (match[i] != -1 and pai[match[i]]
                != -1))
                contract(u, i);
            else if (pai[i] == -1) {
                pai[i] = u;
                if (match[i] == -1) return i;
                i = match[i];
                vis[i] = 1; q.push(i);
            }
        }
    }
    return -1;
}
int blossom() {
    int ans = 0;
    memset(match, -1, sizeof(match));
    for (int i = 0; i < n; i++) if (match[i] == -1)</pre>
        for (int j : g[i]) if (match[j] == -1) {
            match[i] = j;
            match[j] = i;
            ans++;
            break;
    for (int i = 0; i < n; i++) if (match[i] == -1) {
        int j = getpath(i);
```

```
if (j == -1) continue;
        ans++;
        while (j != -1) {
            int p = pai[j], pp = match[p];
            match[p] = j;
            match[j] = p;
            j = pp;
        }
    }
    return ans;
}
6.8 Tarjan
// O(n + m)
int n;
vector < vector < int > > g(MAX);
stack<int> s;
int vis[MAX], comp[MAX];
int id[MAX], p;
int dfs(int k) {
    int lo = id[k] = p++;
    s.push(k);
    vis[k] = 2; // ta na pilha
   // calcula o menor cara q ele alcanca
    // que ainda nao esta em um scc
   for (int i = 0; i < g[k].size(); i++) {</pre>
        if (!vis[g[k][i]])
            lo = min(lo, dfs(g[k][i]));
        else if (vis[g[k][i]] == 2)
            lo = min(lo, id[g[k][i]]);
    }
    // nao alcanca ninguem menor -> comeca scc
    if (lo == id[k]) while (1) {
        int u = s.top();
        s.pop(); vis[u] = 1;
        comp[u] = k;
        if (u == k) break;
```

```
return lo;
}

void tarjan() {
    memset(vis, 0, sizeof(vis));

    p = 0;
    for (int i = 0; i < n; i++) if (!vis[i]) dfs(i);
}

6.9 Floyd-Warshall

// encontra o menor caminho entre todo
// par de vertices e detecta ciclo negativo
// returna 1 sse ha ciclo negativo
// d[i][i] deve ser 0
// para i != j, d[i][j] deve ser w se ha uma aresta</pre>
```

```
// par de vertices e detecta ciclo negativo
// returna 1 sse ha ciclo negativo
// d[i][i] deve ser 0
// para i != j, d[i][j] deve ser w se ha uma aresta
// (i, j) de peso w, INF caso contrario
//
// O(n^3)
int n;
int d[MAX][MAX];
bool floyd_warshall() {
    for (int k = 0; k < n; k++)
    for (int i = 0; i < n; i++)</pre>
    for (int j = 0; j < n; j++)
        d[i][j] = min(d[i][j], d[i][k] + d[k][j]);
    for (int i = 0; i < n; i++)</pre>
        if (d[i][i] < 0) return 1;
    return 0;
}
```

## 6.10 Dinic Bruno

```
// O(n^2 m)
```

```
// Grafo bipartido -> O(sqrt(n)*m)
struct edge {
    int p, c, id; // para, capacidade, id
    edge(int p_, int c_, int id_) : p(p_), c(c_), id(id_) {}
};
vector < vector < edge > > g(MAX);
vector < int > lev;
void add(int a, int b, int c) { // de a pra b com cap. c
    g[a].pb(edge(b, c, g[b].size()));
    g[b].pb(edge(a, 0, g[a].size()-1));
}
bool bfs(int s, int t) {
    lev = vector\langle int \rangle(g.size(), -1); lev[s] = 0;
    queue < int > q; q.push(s);
    while (q.size()) {
        int u = q.front(); q.pop();
        for (auto& i : g[u]) {
            if (lev[i.p] != -1 or !i.c) continue;
            lev[i.p] = lev[u] + 1;
            if (i.p == t) return 1;
            q.push(i.p);
        }
    }
    return 0;
}
int dfs(int v, int s, int f = INF){
    if (v == s) return f;
    int tem = f;
    for (auto& i : g[v]) {
        if (lev[i.p] != lev[v] + 1 or !i.c) continue;
        int foi = dfs(i.p, s, min(tem, i.c));
        tem -= foi, i.c -= foi, g[i.p][i.id].c += foi;
    if (f == tem) lev[v] = -1;
    return f - tem;
}
```

```
int fluxo(int s, int t) {
   int f = 0;
   while (bfs(s, t)) f += dfs(s, t);
   return f;
}
```

## 6.11 Dijkstra

```
// encontra menor distancia de a
// para todos os vertices
// se ao final do algoritmo d[i] = INF,
// entao a nao alcanca i
//
// O(m log(n))
int n;
vector < vector < int > > g(MAX);
vector < vector < int > > w(MAX); // peso das arestas
int d[MAX];
void dijsktra(int a) {
    for (int i = 0; i < n; i++) d[i] = INF;</pre>
    d[a] = 0;
    priority_queue < pair < int , int > > Q;
    Q.push(make_pair(0, a));
    while (Q.size()) {
        int u = Q.top().second, dist = -Q.top().first;
        : () gog. D
        if (dist > d[u]) continue;
        for (int i = 0; i < (int) g[u].size(); i++) {</pre>
            int v = g[u][i];
            if (d[v] > d[u] + w[u][i]) {
                 d[v] = d[u] + w[u][i];
                 Q.push(make_pair(-d[v], v));
            }
        }
    }
```

## 6.12 Art. Point

```
// Computa pontos de articulação
// e pontes
//
// O(n+m)
int in[MAX];
int low[MAX];
int parent[MAX];
vector < int > g[MAX];
bool is_art[MAX];
void dfs_art(int v, int p, int &d){
    parent[v] = p;
    low[v] = in[v] = d++;
    is_art[v] = false;
    for (int j : g[v]){
        if (j == p) continue;
        if (in[j] == -1){
            dfs_art(j, v, d);
            if (low[j] >= in[v]) is_art[v] = true;
            //if (low[j] > in[v]) this edge is a bridge
            low[v] = min(low[v], low[j]);
        else low[v] = min(low[v], in[j]);
    }
    if (p == -1){
        is_art[v] = false;
        int k = 0;
        for (int j : g[v])
            k += (parent[j] == v);
        if (k > 1) is_art[v] = true;
    }
}
int d = 0;
memset(in, -1, sizeof in);
dfs_art(1, -1, d);
```

#### 6.13 Kruskal

```
// Gera AGM a partir do vetor de arestas
//
// O(m log(n))
int n;
vector<pair<int, pair<int, int> > ar; // vetor de arestas
int v[MAX];
// Union-Find em O(log(n))
void build();
int find(int k);
void une(int a, int b);
void kruskal() {
    build():
    sort(ar.begin(), ar.end());
    for (int i = 0; i < (int) ar.size(); i++) {</pre>
        int a = ar[i].s.f, b = ar[i].s.s;
        if (find(a) != find(b)) {
            une(a, b);
            // aresta faz parte da AGM
        }
    }
}
```

## 6.14 Bellman-Ford

```
// Calcula a menor distancia
// entre a e todos os vertices e
// detecta ciclo negativo
// Retorna 1 se ha ciclo negativo
// Nao precisa representar o grafo,
// soh armazenar as arestas
//
// O(nm)
int n, m;
int d[MAX];
vector<pair<int, int> > ar; // vetor de arestas
```

```
vector < int > w:
                       // peso das arestas
bool bellman_ford(int a) {
    for (int i = 0; i < n; i++) d[i] = INF;</pre>
    d[a] = 0;
    for (int i = 0; i <= n; i++)
        for (int j = 0; j < m; j++) {</pre>
            if (d[ar[j].second] > d[ar[j].first] + w[j]) {
                if (i == n) return 1;
                d[ar[j].second] = d[ar[j].first] + w[j];
            }
        }
    return 0;
}
6.15 Centroid
// O(n log(n))
int n;
vector < vector < int > > g(MAX);
int subsize[MAX];
int rem[MAX];
int pai[MAX];
void dfs(int k, int last) {
    subsize[k] = 1;
    for (int i = 0; i < (int) g[k].size(); i++)</pre>
        if (g[k][i] != last and !rem[g[k][i]]) {
            dfs(g[k][i], k);
            subsize[k] += subsize[g[k][i]];
}
int centroid(int k, int last, int size) {
    for (int i = 0; i < (int) g[k].size(); i++) {</pre>
        int u = g[k][i];
        if (rem[u] or u == last) continue;
        if (subsize[u] > size / 2)
```

```
return centroid(u, k, size);
    // k eh o centroid
    return k:
}
void decomp(int k, int last) {
    dfs(k, k);
    // acha e tira o centroid
    int c = centroid(k, k, subsize[k]);
    rem[c] = 1;
    pai[c] = last;
    if (k == last) pai[c] = c;
    // decompoe as sub-arvores
    for (int i = 0; i < (int) g[c].size(); i++)</pre>
        if (!rem[g[c][i]]) decomp(g[c][i], c);
}
void build() {
    memset(rem, 0, sizeof rem);
    decomp(0, 0);
}
6.16 Center
// Centro eh o vertice que minimiza
// a maior distancia dele pra alguem
// O centro fica no meio do diametro
// A funcao center retorna um par com
// o diametro e o centro
//
// O(n+m)
vector < vector < int > > g(MAX);
int n, vis[MAX];
int d[2][MAX];
// retorna ultimo vertice visitado
int bfs(int k, int x) {
        queue < int > q; q.push(k);
```

```
memset(vis, 0, sizeof(vis));
    vis[k] = 1;
    d[x][k] = 0;
    int last = k;
    while (q.size()) {
        int u = q.front(); q.pop();
        last = u;
        for (int i : g[u]) if (!vis[i]) {
            vis[i] = 1;
            q.push(i);
            d[x][i] = d[x][u] + 1;
        }
    }
    return last;
}
pair<int, int> center() {
   int a = bfs(0, 0);
    int b = bfs(a, 1);
    bfs(b, 0);
   int c, mi = INF;
   for (int i = 0; i < n; i++) if (max(d[0][i], d[1][i]) <</pre>
        mi = max(d[0][i], d[1][i]), c = i;
    return {d[0][a], c};
}
6.17 Dinic Dilson
// O(n^2 m)
// Grafo bipartido -> O(sqrt(n)*m)
template <class T> struct dinic {
    struct edge {
        int v, rev;
        edge(int v_, T cap_, int rev_) : v(v_), cap(cap_),
           rev(rev ) {}
    };
    vector < vector < edge >> g;
    vector < int > level;
```

```
queue < int > q;
T flow;
int n;
dinic(int n_{-}) : g(n_{-}), level(n_{-}), n(n_{-}) \{\}
void add_edge(int u, int v, T cap) {
    if (u == v)
        return;
    edge e(v, cap, int(g[v].size()));
    edge r(u, 0, int(g[u].size()));
    g[u].push_back(e);
    g[v].push_back(r);
}
bool build_level_graph(int src, int sink) {
    fill(level.begin(), level.end(), -1);
    while (not q.empty())
        q.pop();
    level[src] = 0;
    q.push(src);
    while (not q.empty()) {
        int u = q.front();
        q.pop();
       for (auto e = g[u].begin(); e != g[u].end();
            if (not e->cap or level[e->v] != -1)
                continue;
            level[e->v] = level[u] + 1;
            if (e->v == sink)
                return true;
            q.push(e->v);
        }
    }
    return false;
}
T blocking_flow(int u, int sink, T f) {
    if (u == sink or not f)
        return f:
    T fu = f;
    for (auto e = g[u].begin(); e != g[u].end(); ++e) {
        if (not e->cap or level[e->v] != level[u] + 1)
            continue:
```

```
T mincap = blocking_flow(e->v, sink, min(fu,
               e->cap));
            if (mincap) {
                g[e->v][e->rev].cap += mincap;
                e->cap -= mincap;
                fu -= mincap;
            }
        }
        if (f == fu)
            level[u] = -1;
        return f - fu;
    }
    T max_flow(int src, int sink) {
        flow = 0;
        while (build_level_graph(src, sink))
            flow += blocking_flow(src, sink,
               numeric_limits <T>::max());
        return flow;
   }
};
6.18 Kosaraju
// O(n + m)
int n:
vector < vector < int > > g(MAX);
vector < vector < int > > gi(MAX); // grafo invertido
int vis[MAX]:
stack<int> S:
int comp[MAX]; // componente conexo de cada vertice
void dfs(int k) {
    vis[k] = 1:
   for (int i = 0; i < (int) g[k].size(); i++)</pre>
        if (!vis[g[k][i]]) dfs(g[k][i]);
    S.push(k);
void scc(int k, int c) {
    vis[k] = 1;
```

```
comp[k] = c;
    for (int i = 0; i < (int) gi[k].size(); i++)</pre>
        if (!vis[gi[k][i]]) scc(gi[k][i], c);
}
void kosaraju() {
    for (int i = 0; i < n; i++) vis[i] = 0;</pre>
    for (int i = 0; i < n; i++) if (!vis[i]) dfs(i);</pre>
    for (int i = 0; i < n; i++) vis[i] = 0;</pre>
    while (S.size()) {
        int u = S.top();
        S.pop();
        if (!vis[u]) scc(u, u);
    }
}
6.19 2 SAT
// Retorna se eh possivel atribuir valores
// Grafo tem que caber 2n vertices
// add(x, y) adiciona implicacao x -> y
// Para adicionar uma clausula (x ou y)
// chamar add(nao(x), y)
// Se x tem que ser verdadeiro, chamar add(nao(x), x)
// O tarjan deve computar o componente conexo
// de cada vertice em comp
//
// O(|V|+|E|)
vector < vector < int > > g(MAX);
int n;
int nao(int x){ return (x + n) \% (2*n); }
// x \rightarrow y = !x ou y
void add(int x, int y){
    g[x].pb(y);
```

// contraposicao

}

g[nao(y)].pb(nao(x));

```
bool doisSAT(){
   tarjan();
   for (int i = 0; i < m; i++)
       if (comp[i] == comp[nao(i)]) return 0;
   return 1;
}
6.20 vimrc
set ts=4 si ai sw=4 number mouse=a
svntax on
6.21 Makefile
CXX = g++
CXXFLAGS = -fsanitize=address -01 -fno-omit-frame-pointer -g
   -Wall -Wshadow -std=c++11 -Wno-unused-result
   -Wno-sign-compare
6.22 Template
#include <bits/stdc++.h>
using namespace std;
#define _ ios_base::sync_with_stdio(0);cin.tie(0);
#define endl '\n'
#define f first
#define s second
#define pb push_back
typedef long long 11;
typedef pair<int, int> ii;
const int INF = 0x3f3f3f3f;
const 11 LINF = 0x3f3f3f3f3f3f3f3f3f11;
int main(){
    exit(0);
}
```