Humuhumunukunukuapua'a UFMG

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1 Estruturas

1.1 BIT

```
// BIT de soma 1-based. v 0-based
// Para mudar o valor da posicao p para x,
// faca: poe(x - query(p, p), p)
// l_bound(x) retorna o menor p tal que
// \text{ query}(1, p+1) > x 	 (0 based!)
// Complexidades:
// build - O(n)
// poe - O(\log(n))
// query - 0(log(n))
// l_bound - O(log(n))
// d432a4
1a8 int n:
7f4 int bit[MAX];
b69 int v[MAX];
0a8 void build() {
       bit[0] = 0:
b91
33 c
       for (int i = 1; i <= n; i++) bit[i] = v[i - 1];
78a
       for (int i = 1: i <= n: i++) {
           int j = i + (i \& -i);
edf
           if (j <= n) bit[j] += bit[i];</pre>
b8a
b39
       }
5dc }
// soma x na posicao p
235 void poe(int x, int p) {
        for (; p <= n; p += p & -p) bit[p] += x;
ec3 }
// soma [1, p]
Obf int pref(int p) {
7c9
       int ret = 0:
       for (; p; p -= p & -p) ret += bit[p];
805
edf
       return ret:
601 }
// soma [a, b]
4ea int query(int a, int b) {
70c return pref(b) - pref(a - 1);
```

```
3dd }
e4a int l_bound(ll x) {
        int p = 0;
676
        for (int i = MAX2; i+1; i--) if (p + (1<<i) <= n
729
            and bit [p + (1 << i)] <= x) x -= bit <math>[p += (1 << i)];
74e
        return p:
c89 }
1.2 BIT 2D
// BIT de soma, update incrementa posicao
// Tem que construir com um vetor com todos os pontos
// que vc quer um dia atualizar (os pontos q vc vai chamar update)
// Complexidades:
// construir - O(n log(n))
// update e query - O(log^2(n))
// 6a760a
a6b template < class T = int > struct bit2d {
        vector <T> X:
        vector < vector < T >> Y, t;
a84
709
        int ub(vector<T>& v, T x) {
dde
             return upper_bound(v.begin(), v.end(), x) - v.begin();
9cc
5cb
        bit2d(vector<pair<T, T>> v) {
2e1
            for (auto [x, y] : v) X.push_back(x);
fd4
            sort(X.begin(), X.end());
1ee
            X.erase(unique(X.begin(), X.end()), X.end());
d56
            t.resize(X.size() + 1);
d12
            Y.resize(t.size());
            sort(v.begin(), v.end(), [](auto a, auto b) {
3d0
e8f
                 return a.second < b.second; });</pre>
            for (auto [x, y] : v) for (int i = ub(X, x); i < t.size();
961
   i += i&-i)
                if (!Y[i].size() or Y[i].back() != y)
b75
   Y[i].push_back(y);
7c7
            for (int i = 0; i < t.size(); i++) t[i].resize(Y[i].size()</pre>
   + 1);
8cc }
e78
        void update(T x, T v, T v) {
2a9
            for (int i = ub(X, x); i < t.size(); i += i\&-i)
```

```
for (int j = ub(Y[i], y); j < t[i].size(); j += j\&-j)
   t[i][j] += v;
       }
533
        T query(T x, T y) {
5d2
966
            T ans = 0:
c54
            for (int i = ub(X, x): i: i -= i\&-i)
4fb
                for (int j = ub(Y[i], y); j; j -= j\&-j) ans += t[i][j];
ba7
            return ans;
62d
46d
        T query (T x1, T v1, T x2, T v2) {
            return query(x2, y2)-query(x2, y1-1)-query(x1-1,
   y2) + query(x1-1, y1-1);
232
6a7 };
1.3 BIT com update em range
```

```
// Operacoes O-based
// query(l, r) retorna a soma de v[l..r]
// update(l, r, x) soma x em v[l..r]
// Complexidades:
// build - O(n)
// query - 0(log(n))
// update - 0(log(n))
// f91737
e04 namespace bit {
        11 bit[2][MAX+2];
3ba
1a8
        int n;
        void build(int n2, int* v) {
61c
1e3
            n = n2:
535
            for (int i = 1; i <= n; i++)
edd
                bit [1] [min(n+1, i+(i\&-i))] += bit [1][i] += v[i-1];
db0
637
        11 get(int x, int i) {
b73
            11 \text{ ret} = 0;
360
            for (; i; i -= i&-i) ret += bit[x][i];
edf
            return ret;
99c
        }
        void add(int x, int i, ll val) {
503
            for (; i <= n; i += i&-i) bit[x][i] += val;</pre>
bf6
        }
162
        11 get2(int p) {
с7с
            return get(0, p) * p + get(1, p);
```

```
153
02a
        11 query(int 1, int r) {
ff5
            return get2(r+1) - get2(1);
633
        }
089
        void update(int 1, int r, 11 x) {
            add(0, 1+1, x), add(0, r+2, -x);
e5f
f58
            add(1, 1+1, -x*1), add(1, r+2, x*(r+1));
        }
e5f
f91 };
```

```
1.4 BIT-Sort Tree
// Tipo uma MergeSort Tree usando Bit
// Apesar da complexidade ser pior, fica melhor na pratica.
// query(1, r, k) retorna o numero de elementos menores que k
// no intervalo [1, r]
// Usa O(n log(n)) de memoria
// Complexidades:
// construir - O(n log^2(n))
// query - O(log^2(n))
// 8d0749
6fa template < typename T > struct ms_bit {
1a8
        int n:
b2f
        vector < vector < T >> bit;
899
        ms_bit(vector<T>& v) : n(v.size()), bit(n+1) {
830
            for (int i = 0; i < n; i++)</pre>
d51
                for (int j = i+1; j \le n; j += j\&-j)
dad
                    bit[i].push_back(v[i]);
535
            for (int i = 1; i <= n; i++)</pre>
eec
                 sort(bit[i].begin(), bit[i].end());
        }
b4d
257
        int p_query(int i, T k) {
7c9
            int ret = 0;
be8
            for (i++: i: i -= i&-i)
                ret += lower_bound(bit[i].begin(), bit[i].end(), k) -
1bd
    bit[i].begin();
edf
            return ret;
6f9
690
        int query(int 1, int r, T k) {
83d
            return p_query(r, k) - p_query(l-1, k);
bcc
        }
```

1.5 Convex Hull Trick Dinamico

```
// para double, use LINF = 1/.0, div(a, b) = a/b
// update(x) atualiza o ponto de intersecao da reta x
// overlap(x) verifica se a reta x sobrepoe a proxima
// add(a, b) adiciona reta da forma ax + b
// query(x) computa maximo de ax + b para entre as retas
//
// O(log(n)) amortizado por insercao
// O(log(n)) por query
// 978376
72c struct Line {
073
        mutable 11 a, b, p;
        bool operator<(const Line& o) const { return a < o.a; }</pre>
8e3
abf
        bool operator<(ll x) const { return p < x; }</pre>
469 }:
326 struct dynamic_hull : multiset <Line, less <>> {
        11 div(ll a, ll b) {
33a
a20
             return a / b - ((a ^ b) < 0 and a % b);
        }
a8a
bbb
        void update(iterator x) {
b2a
             if (next(x) == end()) x -> p = LINF:
             else if (x->a == next(x)->a) x->p = x->b >= next(x)->b?
772
   LINF : -LINF;
424
             else x \rightarrow p = div(next(x) \rightarrow b - x \rightarrow b, x \rightarrow a - next(x) \rightarrow a);
0c4
71c
        bool overlap(iterator x) {
f18
             update(x);
cfa
             if (next(x) == end()) return 0;
            if (x->a == next(x)->a) return x->b >= next(x)->b;
a4a
d40
             return x - p >= next(x) - p;
        }
901
176
        void add(ll a, ll b) {
             auto x = insert({a, b, 0});
1c7
4ab
             while (overlap(x)) erase(next(x)), update(x);
            if (x != begin() and !overlap(prev(x))) x = prev(x),
   update(x);
             while (x != begin() and overlap(prev(x)))
0 fc
4d2
                 x = prev(x), erase(next(x)), update(x);
48f
        }
```

1.6 Convex Hull Trick Estatico

```
// adds tem que serem feitos em ordem de slope
// queries tem que ser feitas em ordem de x
// linear
// 30323e
4b5 struct CHT {
942
        int it:
ac1
        vector<11> a, b;
45e
        CHT(): it(0){}
Obb
        ll eval(int i, ll x){
93d
            return a[i]*x + b[i];
b2a
        }
63a
        bool useless(){
a20
            int sz = a.size();
35f
            int r = sz-1, m = sz-2, l = sz-3:
            return (b[1] - b[r])*(a[m] - a[1]) <
d71
413
                 (b[1] - b[m])*(a[r] - a[1]);
a0c
        }
bf4
        void add(ll A, ll B){
7f5
             a.push_back(A); b.push_back(B);
565
             while (!a.empty()){
233
                 if ((a.size() < 3) || !useless()) break;</pre>
ecb
                 a.erase(a.end() - 2);
                b.erase(b.end() - 2);
568
b21
            }
        }
165
81b
        ll get(ll x){
d27
            it = min(it, int(a.size()) - 1);
             while (it+1 < a.size()){</pre>
46a
3c4
                 if (eval(it+1, x) > eval(it, x)) it++;
f97
                 else break;
fe9
420
            return eval(it, x);
88a
        }
303 };
```

1.7 DSU

```
// Une dois conjuntos e acha a qual conjunto um elemento pertence por
   seu id
//
// find e unite: O(a(n)) \sim = O(1) amortizado
// 8e197e
8d3 struct dsu {
825
        vector<int> id, sz;
b33
        dsu(int n) : id(n), sz(n, 1) { iota(id.begin(), id.end(), 0); }
0cf
        int find(int a) { return a == id[a] ? a : id[a] = find(id[a]);
   }
440
        void unite(int a, int b) {
            a = find(a), b = find(b);
605
d54
            if (a == b) return:
956
            if (sz[a] < sz[b]) swap(a, b);</pre>
6d0
            sz[a] += sz[b], id[b] = a;
ea7
       }
8e1 };
// DSU de bipartido
//
// Une dois vertices e acha a qual componente um vertice pertence
// Informa se a componente de um vertice e bipartida
//
// find e unite: O(log(n))
// 118050
8d3 struct dsu {
6f7
        vector < int > id, sz, bip, c;
        dsu(int n) : id(n), sz(n, 1), bip(n, 1), c(n) {
5b4
db8
            iota(id.begin(), id.end(), 0);
f25
        }
        int find(int a) { return a == id[a] ? a : find(id[a]); }
ef0
        int color(int a) { return a == id[a] ? c[a] : c[a] ^
   color(id[a]); }
        void unite(int a, int b) {
440
263
            bool change = color(a) == color(b);
605
            a = find(a), b = find(b);
a89
            if (a == b) {
```

```
4ed
                if (change) bip[a] = 0;
505
                return;
32d
            }
            if (sz[a] < sz[b]) swap(a, b);
956
            if (change) c[b] = 1;
efe
2cd
            sz[a] += sz[b], id[b] = a, bip[a] &= bip[b];
        }
22b
118 };
// DSU Persistente
//
// Persistencia parcial, ou seja, tem que ir
// incrementando o 't' no une
//
// find e unite: O(log(n))
// 6c63a4
8d3 struct dsu {
        vector<int> id, sz, ti;
733
        dsu(int n) : id(n), sz(n, 1), ti(n, -INF) {
db8
            iota(id.begin(), id.end(), 0);
        }
aad
5e6
        int find(int a, int t) {
6ba
            if (id[a] == a or ti[a] > t) return a;
ea5
            return find(id[a], t);
6cb
        }
fa0
        void unite(int a, int b, int t) {
84f
            a = find(a, t), b = find(b, t);
d54
            if (a == b) return;
956
            if (sz[a] < sz[b]) swap(a, b);</pre>
35d
            sz[a] += sz[b], id[b] = a, ti[b] = t;
513
        }
6c6 };
// DSU com rollback
//
// checkpoint(): salva o estado atual de todas as variaveis
// rollback(): retorna para o valor das variaveis para
// o ultimo checkpoint
// Sempre que uma variavel muda de valor, adiciona na stack
//
```

```
// find e unite: O(log(n))
// checkpoint: 0(1)
// rollback: O(m) em que m e o numero de vezes que alguma
// variavel mudou de valor desde o ultimo checkpoint
// c6e923
8d3 struct dsu {
825
        vector < int > id, sz;
27 c
        stack<stack<pair<int&, int>>> st;
98d
        dsu(int n) : id(n), sz(n, 1) {
1cc
            iota(id.begin(), id.end(), 0), st.emplace();
8cd
bdf
        void save(int &x) { st.top().emplace(x, x); }
30d
        void checkpoint() { st.emplace(); }
5cf
        void rollback() {
ba9
            while(st.top().size()) {
6bf
                auto [end, val] = st.top().top(); st.top().pop();
149
                end = val;
            }
f9a
25a
            st.pop();
3c6
        }
ef0
        int find(int a) { return a == id[a] ? a : find(id[a]); }
440
        void unite(int a, int b) {
605
            a = find(a), b = find(b);
d54
            if (a == b) return:
956
            if (sz[a] < sz[b]) swap(a, b);
            save(sz[a]), save(id[b]);
803
            sz[a] += sz[b], id[b] = a;
6d0
1b9
        }
c6e };
1.8 Li-Chao Tree
// Adiciona retas (ax+b), e computa o minimo entre as retas
// em um dado 'x'
// Cuidado com overflow!
// Se tiver overflow, tenta comprimir o 'x' ou usar
// convex hull trick
//
// O(log(MA-MI)), O(n) de memoria
// 59ba68
```

```
5b0 template < 11 MI = 11(-1e9), 11 MA = 11(1e9) > struct lichao {
        struct line {
b3a
12d
             ll a, b;
cef
             array<int, 2> ch;
fdf
             line(ll a_{-} = 0, ll b_{-} = LINF):
423
                 a(a_{-}), b(b_{-}), ch(\{-1, -1\}) \{\}
888
             11 operator ()(11 x) { return a*x + b; }
d1d
        };
17b
        vector<line> ln:
df8
        int ch(int p, int d) {
e85
             if (\ln \lceil p \rceil, \cosh \lceil d \rceil == -1) {
9af
                 ln[p].ch[d] = ln.size();
cdc
                 ln.emplace_back();
bc1
ef2
             return ln[p].ch[d];
86a
        }
021
        lichao() { ln.emplace_back(); }
c33
        void add(line s, ll l=MI, ll r=MA, int p=0) {
3e3
             11 m = (1+r)/2;
911
             bool L = s(1) < ln[p](1);
d37
             bool M = s(m) < ln[p](m);
03b
             bool R = s(r) < ln[p](r);
825
             if (M) swap(ln[p], s), swap(ln[p].ch, s.ch);
cac
             if (s.b == LINF) return;
f6d
             if (L != M) add(s, 1, m-1, ch(p, 0));
898
             else if (R != M) add(s, m+1, r, ch(p, 1));
76e
        }
092
        11 query(int x, 11 1=MI, 11 r=MA, int p=0) {
11b
            11 m = (1+r)/2, ret = ln[p](x);
9db
             if (ret == LINF) return ret;
529
             if (x < m) return min(ret, query(x, 1, m-1, ch(p, 0)));
81a
             return min(ret, query(x, m+1, r, ch(p, 1)));
fba
        }
59b };
1.9 Li-Chao Tree - Lazy
// Sendo N = MA-MI:
// insert(\{a, b\}) minimiza tudo com ax+b - O(\log N)
// insert(\{a, b\}, 1, r) minimiza com ax+b no range [1, r] - 0(\log^2 N)
// shift({a, b}) soma ax+b em tudo - O(1)
// shift({a, b}, 1, r) soma ax+b no range [1, r] - O(log^2 N)
// query(x) retorna o valor da posicao x - O(\log N)
```

```
// No inicio eh tudo LINF, se inserir {0, 0} fica tudo 0
// O(n log N) de memoria ; O(n) de memoria se nao usar as operações de
   range
// 285a00
41c template <int MI = int(-1e9), int MA = int(1e9) > struct lichao {
        struct line {
b3a
12d
            ll a, b;
            11 la, lb; // lazy
158
cef
            array<int, 2> ch;
fdf
            line(ll a_{-} = 0, ll b_{-} = LINF):
b09
                a(a), b(b), la(0), lb(0), ch(\{-1, -1\}) {}
            11 operator ()(11 x) { return a*x + b; }
888
92e
        };
17b
        vector <line > ln;
        int ch(int p, int d) {
df8
            if (ln[p].ch[d] == -1) {
e85
9af
                ln[p].ch[d] = ln.size();
cdc
                ln.emplace_back();
bc1
ef2
            return ln[p].ch[d];
86a
021
        lichao() { ln.emplace_back(); }
ceb
        void prop(int p, int 1, int r) {
ff8
            if (\ln[p]. \ln == 0 \text{ and } \ln[p]. \ln == 0) return:
1d3
            ln[p].a += ln[p].la, ln[p].b += ln[p].lb;
579
            if (1 != r) {
                int pl = ch(p, 0), pr = ch(p, 1);
b9e
0d7
                ln[pl].la += ln[p].la, ln[pl].lb += ln[p].lb;
                ln[pr].la += ln[p].la, ln[pr].lb += ln[p].lb;
fa8
77f
            }
01e
            ln[p].la = ln[p].lb = 0;
        }
89b
c06
        11 query(int x, int p=0, int l=MI, int r=MA) {
6b9
            prop(p, 1, r);
6f3
            ll ret = ln[p](x);
            if (ln[p].ch[0] == -1 and ln[p].ch[1] == -1) return ret;
33b
            int m = 1 + (r-1)/2;
90d
da9
            if (x <= m) return min(ret, query(x, ch(p, 0), 1, m));</pre>
            return min(ret, query(x, ch(p, 1), m+1, r));
c55
953
        }
5df
        void push(line s, int p, int l, int r) {
```

```
6b9
            prop(p, 1, r);
90d
            int m = 1 + (r-1)/2;
911
            bool L = s(1) < ln[p](1);
d37
            bool M = s(m) < ln[p](m);
03ъ
            bool R = s(r) < ln[p](r);
c3f
            if (M) swap(ln[p].a, s.a), swap(ln[p].b, s.b);
cac
            if (s.b == LINF) return;
c49
            if (L != M) push(s, ch(p, 0), 1, m);
29e
            else if (R != M) push(s, ch(p, 1), m+1, r);
        }
ceb
        void insert(line s, int a=MI, int b=MA, int p=0, int l=MI, int
a8e
   r=MA) {
6b9
            prop(p, 1, r):
2d3
            if (a \le 1 \text{ and } r \le b) \text{ return push}(s, p, 1, r);
1dd
            if (b < 1 or r < a) return;
90d
            int m = 1 + (r-1)/2;
            insert(s, a, b, ch(p, 0), 1, m);
f1e
952
            insert(s, a, b, ch(p, 1), m+1, r);
375
        }
97a
        void shift(line s, int a=MI, int b=MA, int p=0, int l=MI, int
   r=MA) {
6b9
            prop(p, 1, r);
90d
            int m = 1 + (r-1)/2;
9a3
            if (a \le 1 \text{ and } r \le b)
ada
                ln[p].la += s.a, ln[p].lb += s.b;
505
                return:
570
            }
1dd
            if (b < l or r < a) return;
fdd
            if (ln[p].b != LINF) {
                 push(ln[p], ch(p, 0), 1, m);
751
ade
                push(ln[p], ch(p, 1), m+1, r);
c2f
                 ln[p].a = 0, ln[p].b = LINF;
199
            }
a04
            shift(s, a, b, ch(p, 0), 1, m);
e7d
            shift(s, a, b, ch(p, 1), m+1, r);
d43
        }
285 };
1.10 MergeSort Tree
// Se for construida sobre um arrav:
//
        count(i, j, a, b) retorna quantos
//
        elementos de v[i..j] pertencem a [a, b]
//
        report(i, j, a, b) retorna os indices dos
```

```
//
        elementos de v[i..j] que pertencem a [a, b]
//
        retorna o vetor ordenado
```

```
// Se for construida sobre pontos (x, y):
        count(x1, x2, y1, x2) retorna quantos pontos
//
//
        pertencem ao retangulo (x1, y1), (x2, y2)
//
        report(x1, x2, y1, y2) retorna os indices dos pontos que
//
        pertencem ao retangulo (x1, y1), (x2, y2)
//
        retorna os pontos ordenados lexicograficamente
//
        (assume x1 \le x2, y1 \le y2)
//
// kth(y1, y2, k) retorna o indice do ponto com k-esimo menor
// x dentre os pontos que possuem y em [y1, y2] (0 based)
// Se quiser usar para achar k-esimo valor em range, construir
// com ms_tree t(v, true), e chamar kth(1, r, k)
// Usa O(n log(n)) de memoria
// Complexidades:
// construir - O(n log(n))
// count - O(log(n))
// report - O(\log(n) + k) para k indices retornados
// kth - O(log(n))
// 1cef03
c6c template <typename T = int> struct ms_tree {
6f7
        vector<tuple<T, T, int>> v;
1a8
5ee
        vector < vector < tuple < T, T, int >>> t; // {y, idx, left}
6ae
        vector <T> vy;
        ms_tree(vector < pair < T, T >> \& vv) : n(vv.size()), t(4*n), vy(n) {
78c
e80
            for (int i = 0; i < n; i++) v.push_back({vv[i].first,</pre>
   vv[i].second, i});
            sort(v.begin(), v.end());
fca
            build(1, 0, n-1);
224
01a
            for (int i = 0; i < n; i++) vy[i] = get < 0 > (t[1][i+1]);
45e
        ms_tree(vector<T>& vv, bool inv = false) { // inv: inverte
dac
   indice e valor
            vector < pair < T, T >> v2;
8e8
            for (int i = 0; i < vv.size(); i++)</pre>
e1e
196
                inv ? v2.push_back({vv[i], i}) : v2.push_back({i,
   vv[i]});
            *this = ms_tree(v2);
cca
f23
2c6
        void build(int p, int 1, int r) {
            t[p].push_back({get<0>(v[1]), get<0>(v[r]), 0}); //
1d2
   {min_x, max_x, 0}
            if (1 == r) return t[p].push_back({get<1>(v[1]),
5c8
```

```
get <2 > (v[1]), 0});
ee4
             int m = (1+r)/2;
             build(2*p, 1, m), build(2*p+1, m+1, r);
bd9
             int L = 0, R = 0;
32d
             while (t[p].size() \le r-l+1) {
a03
68e
                 int left = get<2>(t[p].back());
                 if (L > m-1 \text{ or } (R+m+1 \le r \text{ and } t[2*p+1][1+R] \le
4aa
    t[2*p][1+L])) {
                      t[p].push_back(t[2*p+1][1 + R++]);
8cf
da0
                      get < 2 > (t[p].back()) = left;
5e2
                      continue:
ce0
                 }
249
                 t[p].push_back(t[2*p][1 + L++]);
339
                 get < 2 > (t[p].back()) = left + 1;
208
             }
2eb
        }
         int get_1(T y) { return lower_bound(vy.begin(), vy.end(), y) -
dd3
    vy.begin(); }
         int get_r(T y) { return upper_bound(vy.begin(), vy.end(), y) -
    vy.begin(); }
         int count(T x1, T x2, T y1, T y2) {
f62
902
             function < int(int, int, int) > dfs = [&](int p, int 1, int
   r) {
7 c 6
                 if (1 == r \text{ or } x2 < get < 0 > (t[p][0]) \text{ or } get < 1 > (t[p][0])
    < x1) return 0:
                 if (x1 \le get<0>(t[p][0]) and get<1>(t[p][0]) \le x2)
2bb
    return r-1;
                 int nl = get<2>(t[p][1]), nr = get<2>(t[p][r]);
784
eb6
                 return dfs(2*p, nl, nr) + dfs(2*p+1, l-nl, r-nr);
122
             };
7cb
             return dfs(1, get_l(y1), get_r(y2));
f65
002
        vector<int> report(T x1, T x2, T y1, T y2) {
4b8
             vector < int > ret;
85e
             function < void(int, int, int) > dfs = [&](int p, int 1, int
   r) {
882
                 if (1 == r \text{ or } x2 < get < 0 > (t[p][0]) \text{ or } get < 1 > (t[p][0])
    < x1) return:
                 if (x1 \le get<0>(t[p][0]) and get<1>(t[p][0]) \le x2) {
8da
                      for (int i = 1; i < r; i++)</pre>
e00
   ret.push_back(get<1>(t[p][i+1]));
505
                      return;
                 }
067
784
                 int nl = get<2>(t[p][1]), nr = get<2>(t[p][r]);
```

```
194
                dfs(2*p, nl, nr), dfs(2*p+1, l-nl, r-nr);
12b
            };
8ad
            dfs(1, get_l(y1), get_r(y2));
edf
            return ret;
668
985
        int kth(T y1, T y2, int k) {
902
            function < int (int, int, int) > dfs = [&] (int p, int 1, int
   r) {
150
                if (k >= r-1) {
941
                    k -= r-1:
daa
                    return -1;
b8d
                }
8da
                if (r-l == 1) return get<1>(t[p][l+1]);
784
                int nl = get<2>(t[p][1]), nr = get<2>(t[p][r]);
072
                int left = dfs(2*p, nl, nr);
                if (left != -1) return left;
3b6
04d
                return dfs(2*p+1, 1-n1, r-nr);
a1b
            }:
            return dfs(1, get_l(v1), get_r(v2));
7cb
635
        }
1ce };
1.11 Min queue - deque
// Tudo O(1) amortizado
// c13c57
1dc template < class T> struct minqueue {
2d8
        deque<pair<T, int>> q;
3fc
        void push(T x) {
56e
            int ct = 1:
953
            while (q.size() and x < q.front().first)</pre>
75f
                ct += q.front().second, q.pop_front();
987
            q.emplace_front(x, ct);
e8d
42d
        void pop() {
            if (q.back().second > 1) q.back().second--;
aa2
c51
            else q.pop_back();
5fd
        T min() { return q.back().first; }
ea6
c13 };
1.12 Min queue - stack
```

// Tudo O(1) amortizado

// fe0cad

```
557 template < class T> struct minstack {
        stack<pair<T, T>> s;
81f
3fc
        void push(T x) {
12b
            if (!s.size()) s.push({x, x});
9d9
            else s.emplace(x, std::min(s.top().second, x));
f8d
4f0
        T top() { return s.top().first; }
94a
        T pop() {
1f2
            T ans = s.top().first;
2eb
            s.pop();
ba7
            return ans:
013
        }
614
        int size() { return s.size(); }
13b
        T min() { return s.top().second; }
4c0 };
1dc template < class T> struct mingueue {
        minstack <T> s1, s2;
7cd
        void push(T x) { s1.push(x); }
        void move() {
c96
d4d
            if (s2.size()) return;
d92
            while (s1.size()) {
7ae
                T x = s1.pop();
489
                s2.push(x);
656
            }
        }
ef1
787
        T front() { return move(), s2.top(); }
        T pop() { return move(), s2.pop(); }
23a
7f3
        int size() { return s1.size()+s2.size(); }
19c
        T min() {
cd6
            if (!s1.size()) return s2.min():
58e
            else if (!s2.size()) return s1.min();
31d
            return std::min(s1.min(), s2.min());
9c7
        }
6d3 };
1.13 Order Statistic Set
// Funciona do C++11 pra cima
// 901923
774 #include <ext/pb_ds/assoc_container.hpp>
30f #include <ext/pb_ds/tree_policy.hpp>
```

0d7 using namespace __gnu_pbds;

```
4fc template <class T>
def
        using ord_set = tree<T, null_type, less<T>, rb_tree_tag,
3a1
        tree_order_statistics_node_update>;
// para declarar:
// ord_set < int > s;
// coisas do set normal funcionam:
// for (auto i : s) cout << i << endl;
// cout << s.size() << endl;
// k-esimo maior elemento O(log|s|):
// k=0: menor elemento
// cout << *s.find_by_order(k) << endl;</pre>
// quantos sao menores do que k O(log|s|):
// cout << s.order_of_key(k) << endl;</pre>
// Para fazer um multiset, tem que
// usar ord_set<pair<int, int>> com o
// segundo parametro sendo algo para diferenciar
// os ementos iguais.
// s.order_of_key({k, -INF}) vai retornar o
// numero de elementos < k
```

1.14 Priority Queue DS

```
// Mantem updates aplicados em uma estrutura de dados
// que permita rollback e nao seja amortizada.
// Cada update possui uma prioridade,
// sendo possivel remover o update com maior prioridade.
// Os updates devem ser comutativos, ou seja, o estado
// da estrutura deve ser o mesmo independente da ordem
// que eles sejam aplicados.
//
// Complexidades:
// update - O(log(n) + T(n))
// query - T(n)
// pop - O(log(n) * T(n)) amortizado
// onde T(n) eh a complexidade do update
//
// 54a75e
// assumes all priorities are distinct
945 template < typename DS, typename UPD > struct priority_queue_ds {
df4
a7e
        vector < tuple < UPD, int, int >> upd; // {u, p, idx_in_pos}
866
        set < pair < int , int >> st;
927
        vector < int > pos;
```

```
cf0
        priority_queue_ds(int n) : D(n) {}
6af
        void update(UPD u, int p) {
9ab
            D.update(u);
d07
            st.emplace(p, pos.size());
6ca
            upd.emplace_back(u, p, pos.size());
e3d
            pos.push_back(upd.size() - 1);
6af
        }
427
        int query(int a) {
            return D.find(a);
aa3
2d3
        }
42d
        void pop() {
25f
            int k = 1, min_p; // k = number of pops we will do
43e
            vector<tuple<UPD, int, int>> small, big;
639
            auto it = st.end();
231
            for (int qt = 0; qt++ < (k+1)/2;) {
049
                it--:
3ab
                min_p = it->first;
80f
                int i = pos[it->second];
e82
                if (qt > 1) big.push_back(upd[i]);
84b
                k = max<int>(k, upd.size() - i);
b9a
            }
b3d
            for (int i = 0; i < k; i++) {
a62
                D.rollback():
6d8
                auto [u, p, idx] = upd.rbegin()[i];
86d
                if (p < min_p) small.emplace_back(u, p, idx);</pre>
            }
588
23e
            st.erase(prev(st.end()));
623
            upd.erase(upd.end() - k, upd.end());
a25
            small.insert(small.end(), big.rbegin(), big.rend());
            for (auto [u, p, idx] : small) {
06f
9ab
                D.update(u);
c8e
                upd.emplace_back(u, p, idx);
a7d
                pos[idx] = upd.size() - 1;
            }
ec7
        }
bd1
54a };
1.15 Range color
```

// update(l, r, c) colore o range [l, r] com a cor c,

```
// e retorna os ranges que foram coloridos {1, r, cor}
// query(i) returna a cor da posicao i
//
// Complexidades (para q operacoes):
// update - O(log(q)) amortizado
// query - O(log(q))
// 9e9cab
df6 template < typename T > struct color {
        set < tuple < int , int , T >> se;
071
        vector<tuple<int, int, T>> update(int 1, int r, T val) {
9c4
            auto it = se.upper bound({r. INF. val});
            if (it != se.begin() and get<1>(*prev(it)) > r) {
753
                auto [L, R, V] = *--it;
e91
3f0
                se.erase(it);
bfd
                se.emplace(L, r, V), se.emplace(r+1, R, V);
683
            it = se.lower_bound({1, -INF, val});
d9e
            if (it != se.begin() and get<1>(*prev(it)) >= 1) {
516
                auto [L, R, V] = *--it;
e91
3f0
                se.erase(it):
75a
                se.emplace(L, 1-1, V), it = se.emplace(1, R, V).first;
b65
            }
d7b
            vector<tuple<int, int, T>> ret;
7a1
            for (; it != se.end() and get<0>(*it) <= r; it =</pre>
   se.erase(it))
8c0
                ret.push_back(*it);
            se.emplace(1, r, val);
b4a
edf
            return ret;
b6c
ff9
        T query(int i) {
            auto it = se.upper_bound({i, INF, T()});
c31
8e7
            if (it == se.begin() or get<1>(*--it) < i) return -1; //</pre>
   nao tem
53d
            return get <2>(*it);
daf
        }
9e9 };
1.16 RMQ < O(n), O(1) > - min queue
// O(n) pra buildar, query O(1)
// Se tiver varios minimos, retorna
// o de menor indice
// bab412
1a5 template < typename T > struct rmq {
```

```
517
         vector <T> v;
fcc
         int n; static const int b = 30;
70e
         vector < int > mask, t;
         int op(int x, int y) { return v[x] \leftarrow v[y] ? x : y; }
183
         int msb(int x) { return __builtin_clz(1)-__builtin_clz(x); }
 ee1
         int small(int r. int sz = b) { return
    r-msb(mask[r]&((1<<sz)-1));}
         rmq() {}
6ad
         rmq(const\ vector < T > \&\ v_) : v(v_), n(v.size()), mask(n), t(n) {
43c
2e5
             for (int i = 0, at = 0; i < n; mask[i++] = at |= 1) {
a61
                 at = (at << 1) & ((1 << b) -1);
c00
                 while (at and op(i-msb(at&-at), i) == i) at ^= at&-at;
 c2f
             for (int i = 0; i < n/b; i++) t[i] = small(b*i+b-1);
 ea4
39d
             for (int j = 1; (1<<j) <= n/b; j++) for (int i = 0;
    i+(1<< i) <= n/b; i++)
                 t[n/b*j+i] = op(t[n/b*(j-1)+i],
ba5
    t[n/b*(j-1)+i+(1<<(j-1))]);
41a
        }
e34
         int index_query(int 1, int r) {
27b
             if (r-l+1 \le b) return small(r, r-l+1);
             int x = 1/b+1, y = r/b-1;
e80
fd3
             if (x > y) return op(small(l+b-1), small(r));
 a4e
             int j = msb(y-x+1);
ea3
             int ans = op(small(1+b-1), op(t[n/b*j+x],
    t[n/b*j+y-(1<< j)+1]));
be6
             return op(ans, small(r));
62a
093
         T query(int 1, int r) { return v[index_query(1, r)]; }
bab };
1.17 SegTreap
// Muda uma posicao do plano, e faz query de operacao
// associativa e comutativa em retangulo
// Mudar ZERO e op
// Esparso nas duas coordenadas, inicialmente eh tudo ZERO
//
// Para query com distancia de manhattan <= d, faca
// nx = x+y, ny = x-y
// Update em (nx, ny), query em ((nx-d, ny-d), (nx+d, ny+d))
// Valores no X tem que ser de O ateh NX
// Para q operacoes, usa O(q log(NX)) de memoria, e as
// operacoes custa O(log(q) log(NX))
// 75f2d0
```

```
55b const int ZERO = INF;
560 const int op(int 1, int r) { return min(1, r); }
878 mt19937 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
aa1 template < typename T > struct treap {
3c9
        struct node {
b19
            node *1, *r;
ee1
            int p;
850
            pair<11, 11> idx; // {y, x}
36d
            T val. mi:
bc2
            node(ll x, ll y, T val_) : l(NULL), r(NULL), p(rng()),
1b5
                idx(pair(y, x)), val(val_), mi(val) {}
01e
            void update() {
d6e
                mi = val;
                if (1) mi = op(mi, 1->mi);
182
                if (r) mi = op(mi, r \rightarrow mi);
b68
            }
282
6e1
        };
bb7
        node* root;
84b
        treap() { root = NULL; }
cec
        \simtreap() {
            vector < node *> q = {root};
609
402
            while (q.size()) {
e5d
                node* x = q.back(); q.pop_back();
                if (!x) continue;
ee9
1 c 7
                q.push_back(x->1), q.push_back(x->r);
bf0
                delete x;
653
            }
50e
        treap(treap&& t) : treap() { swap(root, t.root); }
225
bcf
        void join(node* 1, node* r, node*& i) { // assume que 1 < r</pre>
986
            if (!1 or !r) return void(i = 1 ? 1 : r);
80e
            if (1->p > r->p) join(1->r, r, 1->r), i = 1;
fa0
            else join(1, r->1, r->1), i = r;
bda
            i->update();
671
c82
        void split(node* i, node*& 1, node*& r, pair<11, 11> idx) {
            if (!i) return void(r = 1 = NULL);
26a
13c
            if (i->idx < idx) split(i->r, i->r, r, idx), l = i;
            else split(i \rightarrow 1, l, i \rightarrow 1, idx), r = i;
d26
            i->update();
bda
```

```
143
d3b
        void update(ll x, ll y, T v) {
df9
            node *L, *M, *R;
            split(root, M, R, pair(y, x+1)), split(M, L, M, pair(y,
8b2
   x));
            if (M) M \rightarrow val = M \rightarrow mi = v;
1 e 4
            else M = new node(x, y, v);
9e5
69d
            join(L, M, M), join(M, R, root);
58e
        T query(ll ly, ll ry) {
91b
df9
            node *L, *M, *R;
1c0
            split(root, M, R, pair(ry, LINF)), split(M, L, M, pair(ly,
   0)):
0f7
            T ret = M ? M->mi : ZERO;
69d
            join(L, M, M), join(M, R, root);
edf
            return ret;
1ae
        }
bdf };
46a template < typename T > struct segtreap {
c4f
        vector<treap<T>> seg;
6e7
        vector < int > ch[2];
e4e
        11 NX:
253
        segtreap(11 NX_{-}) : seg(1), NX(NX_{-}) \{ ch[0].push_back(-1), \}
   ch[1].push_back(-1); }
a71
        int get_ch(int i, int d){
e51
            if (ch[d][i] == -1) {
2d6
                 ch[d][i] = seg.size();
23e
                 seg.emplace_back();
842
                 ch[0].push_back(-1), ch[1].push_back(-1);
3e1
            }
968
            return ch[d][i]:
        }
bb6
        T query(11 lx, 11 rx, 11 ly, 11 ry, int p, 11 l, 11 r) {
10c
003
            if (rx < 1 or r < 1x) return ZERO;</pre>
fOf
            if (lx <= 1 and r <= rx) return seg[p].query(ly, ry);</pre>
            11 m = 1 + (r-1)/2:
e6a
            return op(query(lx, rx, ly, ry, get_ch(p, 0), 1, m),
354
060
                     query(lx, rx, ly, ry, get_ch(p, 1), m+1, r));
a5e
f48
        T query(ll lx, ll rx, ll ly, ll ry) { return query(lx, rx, ly,
   ry, 0, 0, NX); }
```

```
249
        void update(ll x, ll y, T val, int p, ll l, ll r) {
                                                                            ee4
                                                                                        int m = (1+r)/2:
            if (1 == r) return seg[p].update(x, y, val);
73c
                                                                            b1f
                                                                                        return query(a, b, 2*p, 1, m) + query(a, b, 2*p+1, m+1, r);
e6a
            11 m = 1 + (r-1)/2;
                                                                            4c5
cc5
            if (x <= m) update(x, y, val, get_ch(p, 0), 1, m);</pre>
                                                                            cfb
                                                                                    11 update(int a, int b, int x, int p=1, int l=0, int r=n-1) {
5a2
            else update(x, y, val, get_ch(p, 1), m+1, r);
                                                                            6b9
                                                                                        prop(p, 1, r);
            seg[p].update(x, y, val);
                                                                            9a3
                                                                                        if (a <= 1 and r <= b) {</pre>
980
cc2
        }
                                                                            b94
                                                                                            lazy[p] += x;
517
        void update(11 x, 11 y, T val) { update(x, y, val, 0, 0, NX); }
                                                                           6b9
                                                                                            prop(p, 1, r);
40a };
                                                                            534
                                                                                            return seg[p];
                                                                            821
                                                                                        }
                                                                            e9f
                                                                                        if (b < 1 or r < a) return seg[p];</pre>
1.18 SegTree
                                                                            ee4
                                                                                        int m = (1+r)/2;
                                                                            fdb
                                                                                        return seg[p] = update(a, b, x, 2*p, 1, m) +
// Recursiva com Lazy Propagation
                                                                            7fd
                                                                                            update(a, b, x, 2*p+1, m+1, r);
// Query: soma do range [a, b]
                                                                            75 c
                                                                                   }
// Update: soma x em cada elemento do range [a, b]
                                                                           Oaf }:
// Pode usar a seguinte funcao para indexar os nohs:
// f(1, r) = (1+r) | (1!=r), usando 2N de memoria
//
                                                                           // Se tiver uma seg de max, da pra descobrir em O(log(n))
// Complexidades:
                                                                            // o primeiro e ultimo elemento >= val numa range:
// build - O(n)
// query - O(log(n))
                                                                           // primeira posicao >= val em [a, b] (ou -1 se nao tem)
// update - 0(log(n))
                                                                           // 68c3e5
                                                                           119 int get_left(int a, int b, int val, int p=1, int l=0, int r=n-1) {
// Oafec1
                                                                            6b9
                                                                                    prop(p, 1, r);
aa4 namespace seg {
                                                                                    if (b < l or r < a or seg[p] < val) return -1;</pre>
                                                                            f38
        11 \text{ seg}[4*MAX], lazy[4*MAX];
005
                                                                                    if (r == 1) return 1;
                                                                            205
052
        int n. *v:
                                                                            ee4
                                                                                    int m = (1+r)/2:
                                                                                    int x = get_left(a, b, val, 2*p, 1, m);
                                                                            753
d22
        11 build(int p=1, int l=0, int r=n-1) {
                                                                            50e
                                                                                    if (x != -1) return x;
3 c 7
            lazy[p] = 0;
                                                                                    return get_left(a, b, val, 2*p+1, m+1, r);
                                                                            c3c
6cd
            if (1 == r) return seg[p] = v[1];
                                                                            68c }
ee4
            int m = (1+r)/2:
193
            return seg[p] = build(2*p, 1, m) + build(2*p+1, m+1, r);
                                                                            // ultima posicao >= val em [a. b] (ou -1 se nao tem)
c71
                                                                           // 1b71df
        void build(int n2, int* v2) {
0d8
                                                                            992 int get_right(int a, int b, int val, int p=1, int l=0, int r=n-1) {
680
            n = n2, v = v2;
                                                                            6b9
                                                                                    prop(p, 1, r);
6f2
            build();
                                                                            f38
                                                                                    if (b < l \text{ or } r < a \text{ or } seg[p] < val) return -1;
acb
        }
                                                                            205
                                                                                    if (r == 1) return 1:
ceb
        void prop(int p, int l, int r) {
                                                                            ee4
                                                                                    int m = (1+r)/2;
            seg[p] += lazy[p]*(r-l+1);
cdf
                                                                                    int x = get_right(a, b, val, 2*p+1, m+1, r);
                                                                            1b1
            if (1 != r) lazy[2*p] += lazy[p], lazy[2*p+1] += lazy[p];
2c9
                                                                            50e
                                                                                    if (x != -1) return x;
3c7
            lazv[p] = 0:
                                                                            6a7
                                                                                    return get_right(a, b, val, 2*p, 1, m);
c10
                                                                            1b7 }
2c3
        ll query(int a, int b, int p=1, int l=0, int r=n-1) {
6b9
            prop(p, 1, r);
                                                                           // Se tiver uma seg de soma sobre um array nao negativo v, da pra
            if (a <= 1 and r <= b) return seg[p];</pre>
527
                                                                            // descobrir em O(\log(n)) o maior j tal que v[i]+v[i+1]+...+v[j-1] <
            if (b < 1 or r < a) return 0;</pre>
786
```

```
val
// 2b8ea7
6a9 int lower_bound(int i, ll& val, int p, int l, int r) {
6b9
        prop(p, 1, r);
        if (r < i) return n:</pre>
6e8
        if (i <= 1 and seg[p] < val) {</pre>
b5d
bff
            val -= seg[p];
041
            return n;
634
        if (1 == r) return 1;
3ce
        int m = (1+r)/2;
ee4
514
        int x = lower_bound(i, val, 2*p, 1, m);
ee0
        if (x != n) return x:
8b9
        return lower_bound(i, val, 2*p+1, m+1, r);
2b8 }
```

1.19 SegTree 2D Iterativa

```
// Consultas 0-based
// Um valor inicial em (x, y) deve ser colocado em seg[x+n][y+n]
// Query: soma do retangulo ((x1, y1), (x2, y2))
// Update: muda o valor da posicao (x, y) para val
// Nao pergunte como que essa coisa funciona
// Para query com distancia de manhattan <= d, faca
// nx = x+y, ny = x-y
// Update em (nx, ny), query em ((nx-d, ny-d), (nx+d, ny+d))
// Se for de min/max, pode tirar os if's da 'query', e fazer
// sempre as 4 operacoes. Fica mais rapido
// Complexidades:
// build - O(n^2)
// \text{ query - } O(\log^2(n))
// update - O(log^2(n))
// 67b9e5
731 int seg[2*MAX][2*MAX], n;
0a8 void build() {
        for (int x = 2*n; x; x--) for (int y = 2*n; y; y--) {
919
c81
            if (x < n) seg[x][y] = seg[2*x][y] + seg[2*x+1][y];
            if (y < n) seg[x][y] = seg[x][2*y] + seg[x][2*y+1];
fe9
d51
499 }
251 int query(int x1, int y1, int x2, int y2) {
```

```
827
         int ret = 0, y3 = y1 + n, y4 = y2 + n;
83e
         for (x1 += n, x2 += n; x1 <= x2; ++x1 /= 2, --x2 /= 2)
0f2
              for (y1 = y3, y2 = y4; y1 \le y2; ++y1 /= 2, --y2 /= 2) {
554
                  if (x1\%2 == 1 \text{ and } y1\%2 == 1) \text{ ret } += \text{seg}[x1][y1];
                  if (x1\%2 == 1 \text{ and } y2\%2 == 0) \text{ ret } += \text{seg}[x1][y2];
6b0
                  if (x2\%2 == 0 \text{ and } y1\%2 == 1) \text{ ret } += \text{seg}[x2][y1];
c01
5d4
                  if (x2\%2 == 0 \text{ and } y2\%2 == 0) \text{ ret } += \text{seg}[x2][y2];
2d0
             }
edf
         return ret;
ff1 }
767 void update(int x, int y, int val) {
66a
         int y2 = y += n;
192
         for (x += n; x; x /= 2, y = y2) {
970
              if (x \ge n) seg[x][y] = val;
              else seg[x][y] = seg[2*x][y] + seg[2*x+1][y];
ba9
3b1
              while (y /= 2) \operatorname{seg}[x][y] = \operatorname{seg}[x][2*y] + \operatorname{seg}[x][2*y+1];
         }
b8b
62e }
1.20 SegTree Beats
// \text{ query}(a, b) - \{\{\min(v[a..b]), \max(v[a..b])\}, \sup(v[a..b])\}
// updatemin(a, b, x) faz com que v[i] <- min(v[i], x),</pre>
// para i em [a. b]
// updatemax faz o mesmo com max, e updatesum soma x
// em todo mundo do intervalo [a, b]
//
// Complexidades:
// build - O(n)
// query - O(log(n))
// update - O(log^2 (n)) amortizado
// (se nao usar updatesum, fica log(n) amortizado)
// 41672b
7c6 #define f first
Oab #define s second
f39 namespace beats {
3c9
         struct node {
526
              int tam:
125
              ll sum, lazy; // lazy pra soma
4f3
              ll mi1, mi2, mi; // mi = #mi1
c61
              ll ma1, ma2, ma; // ma = #ma1
```

```
426
            node(11 x = 0) {
ba6
                sum = mi1 = ma1 = x:
b29
                mi2 = LINF, ma2 = -LINF;
                mi = ma = tam = 1:
62c
c60
                lazv = 0:
b00
770
            node(const node& 1. const node& r) {
a95
                sum = 1.sum + r.sum, tam = 1.tam + r.tam;
c60
                lazv = 0;
                if (1.mi1 > r.mi1) {
797
230
                     mi1 = r.mi1, mi = r.mi;
ea2
                     mi2 = min(1.mi1, r.mi2):
dcd
                } else if (1.mi1 < r.mi1) {</pre>
                     mi1 = 1.mi1. mi = 1.mi:
e34
4b3
                     mi2 = min(r.mi1, l.mi2);
9d9
                } else {
a39
                     mi1 = 1.mi1, mi = 1.mi+r.mi;
                     mi2 = min(1.mi2, r.mi2):
834
5b7
                if (1.ma1 < r.ma1) {</pre>
cd0
                     ma1 = r.ma1, ma = r.ma;
6a0
96d
                     ma2 = max(1.ma1, r.ma2);
                } else if (1.ma1 > r.ma1) {
5f0
                     ma1 = l.ma1, ma = l.ma;
ae0
                     ma2 = max(r.ma1, l.ma2);
2ca
949
                } else {
                     ma1 = l.ma1, ma = l.ma+r.ma;
db2
c05
                     ma2 = max(1.ma2. r.ma2):
                }
710
1ba
            }
4b4
            void setmin(11 x) {
55e
                if (x >= ma1) return;
                sum += (x - ma1)*ma:
463
be5
                if (mi1 == ma1) mi1 = x:
                if (mi2 == ma1) mi2 = x:
0a0
b81
                ma1 = x:
0 c 3
            }
6cb
            void setmax(ll x) {
e25
                if (x <= mi1) return;</pre>
7e8
                sum += (x - mi1)*mi;
                if (ma1 == mi1) ma1 = x:
0bb
                if (ma2 == mi1) ma2 = x;
c32
1ff
                mi1 = x;
            }
a86
4cf
            void setsum(ll x) {
fe8
                mi1 += x, mi2 += x, ma1 += x, ma2 += x;
620
                sum += x*tam:
```

```
c46
                 lazv += x;
b53
47f
        };
        node seg[4*MAX];
62b
        int n, *v;
052
93ъ
        node build(int p=1, int l=0, int r=n-1) {
d84
             if (1 == r) return seg[p] = {v[1]};
             int m = (1+r)/2:
ee4
3d6
             return seg[p] = \{build(2*p, 1, m), build(2*p+1, m+1, r)\};
444
        }
0d8
        void build(int n2. int* v2) {
680
             n = n2. v = v2:
6f2
             build();
        }
acb
        void prop(int p, int 1, int r) {
ceb
8ce
             if (1 == r) return;
abd
             for (int k = 0; k < 2; k++) {
                 if (seg[p].lazy) seg[2*p+k].setsum(seg[p].lazy);
d07
                 seg[2*p+k].setmin(seg[p].ma1);
843
f79
                 seg[2*p+k].setmax(seg[p].mi1);
585
431
             seg[p].lazy = 0;
7ee
055
        pair < pair < ll. 11 > . 11 > guery (int a. int b. int p=1. int l=0.
   int r=n-1) {
e07
             if (b < 1 \text{ or } r < a) \text{ return } \{\{LINF, -LINF\}, 0\}:
             if (a \le 1 \text{ and } r \le b) \text{ return } \{\{seg[p].mi1, seg[p].ma1\},
9be
   seg[p].sum};
6b9
             prop(p, 1, r);
             int m = (1+r)/2;
ee4
             auto L = query(a, b, 2*p, 1, m), R = query(a, b, 2*p+1.
e6f
   m+1. r):
             return {{min(L.f.f, R.f.f), max(L.f.s, R.f.s)}, L.s+R.s};
96d
e9d
2c8
        node updatemin(int a, int b, ll x, int p=1, int l=0, int
   r=n-1) {
744
             if (b < l or r < a or seg[p].ma1 <= x) return seg[p];</pre>
309
             if (a \le 1 \text{ and } r \le b \text{ and } seg[p].ma2 < x) {
                 seg[p].setmin(x);
ccd
534
                 return seg[p];
bbf
6b9
             prop(p, 1, r);
ee4
             int m = (1+r)/2;
96a
             return seg[p] = \{updatemin(a, b, x, 2*p, 1, m),
4db
                              updatemin(a, b, x, 2*p+1, m+1, r)};
```

```
aad
044
        node updatemax(int a, int b, ll x, int p=1, int l=0, int
   r=n-1) {
b59
            if (b < 1 or r < a or seg[p].mi1 >= x) return seg[p];
a9e
             if (a \le 1 \text{ and } r \le b \text{ and } seg[p].mi2 > x) {
                 seg[p].setmax(x);
e8a
534
                 return seg[p];
            }
e9b
6b9
             prop(p, 1, r);
ee4
             int m = (1+r)/2;
             return seg[p] = \{updatemax(a, b, x, 2*p, 1, m),
ee3
98b
                              updatemax(a, b, x, 2*p+1, m+1, r)};
323
        }
        node updatesum(int a, int b, ll x, int p=1, int l=0, int
   r=n-1) {
e9f
            if (b < 1 or r < a) return seg[p];</pre>
             if (a \le 1 \text{ and } r \le b) \{
9a3
8f4
                 seg[p].setsum(x);
534
                 return seg[p];
            }
596
6b9
            prop(p, 1, r);
ee4
            int m = (1+r)/2;
7b6
             return seg[p] = \{updatesum(a, b, x, 2*p, 1, m),
483
                              updatesum(a, b, x, 2*p+1, m+1, r)};
111
        }
0d2 }:
```

1.21 SegTree Colorida

```
// Cada posicao tem um valor e uma cor
// O construtor receve um vector de {valor, cor}
// e o numero de cores (as cores devem estar em [0, c-1])
// query(c, a, b) retorna a soma dos valores
// de todo mundo em [a, b] que tem cor c
// update(c, a, b, x) soma x em todo mundo em
// [a, b] que tem cor c
// paint(c1, c2, a, b) faz com que todo mundo
// em [a, b] que tem cor c1 passe a ter cor c2
//
// Complexidades:
// construir - O(n log(n)) espaco e tempo
// query - O(log(n))
// update - O(log(n))
// paint - O(log(n)) amortizado
// 2938e8
04f struct seg_color {
```

```
3c9
        struct node {
b19
            node *1, *r;
0f9
            int cnt;
9ca
            ll val, lazv;
277
            node(): 1(NULL), r(NULL), cnt(0), val(0), lazy(0) {}
01e
            void update() {
d0a
                cnt = 0, val = 0;
bc4
                for (auto i : {1, r}) if (i) {
c89
                     i->prop();
281
                     cnt += i->cnt, val += i->val;
68d
                }
554
            }
a9c
            void prop() {
2dd
                if (!lazy) return;
3f7
                val += lazy*(ll)cnt;
b64
                for (auto i : {1, r}) if (i) i->lazy += lazy;
c60
                lazv = 0;
e24
            }
514
        };
1a8
        int n;
9b0
        vector < node *> seg;
        seg_color(vector<pair<int, int>>& v, int c) : n(v.size()),
6e0
   seg(c, NULL) {
830
            for (int i = 0; i < n; i++)</pre>
                seg[v[i].second] = insert(seg[v[i].second], i,
   v[i].first, 0, n-1);
94a
        }
3c7
        \simseg_color() {
dde
            queue < node *> q;
3a6
            for (auto i : seg) q.push(i);
            while (q.size()) {
402
20b
                auto i = q.front(); q.pop();
dab
                if (!i) continue;
7c7
                q.push(i->1), q.push(i->r);
5се
                delete i;
c60
            }
139
        }
40b
        node* insert(node* at, int idx, int val, int l, int r) {
1a4
            if (!at) at = new node();
232
            if (1 == r) return at->cnt = 1, at->val = val, at;
ee4
            int m = (1+r)/2:
137
            if (idx \le m) at->1 = insert(at->1, idx, val, 1, m);
3e6
            else at->r = insert(at->r, idx, val, m+1, r);
cff
            return at->update(), at;
```

```
d6e
870
        11 query(node* at, int a, int b, int l, int r) {
61b
            if (!at or b < l or r < a) return 0;</pre>
d9f
            at->prop();
ch2
            if (a <= l and r <= b) return at->val:
ee4
            int m = (1+r)/2;
4c4
            return query(at->1, a, b, 1, m) + query(at->r, a, b, m+1,
   r);
8c3
        11 query(int c, int a, int b) { return query(seg[c], a, b, 0,
   n-1); }
91c
        void update(node* at, int a, int b, int x, int 1, int r) {
fba
            if (!at or b < l or r < a) return:</pre>
            at->prop():
d9f
            if (a <= 1 and r <= b) {</pre>
9a3
                at - > lazv += x;
e9a
cb2
                return void(at->prop());
051
            int m = (1+r)/2;
ee4
            update(at->1, a, b, x, 1, m), update(at->r, a, b, x, m+1,
0b0
   r);
7b4
            at ->update();
9fd
        void update(int c, int a, int b, int x) { update(seg[c], a, b,
a40
   x. 0. n-1): 
70c
        void paint(node*& from. node*& to. int a. int b. int l. int r)
   {
10f
            if (to == from or !from or b < l or r < a) return:
e85
            from ->prop();
889
            if (to) to->prop();
            if (a \le 1 \text{ and } r \le b)
9a3
24d
                if (!to) {
38f
                    to = from:
140
                    from = NULL:
505
                     return;
                }
e5f
ee4
                int m = (1+r)/2;
                paint(from -> 1, to -> 1, a, b, 1, m), paint(from -> r,
   to->r, a, b, m+1, r);
72d
                to->update();
270
                delete from:
                from = NULL;
140
505
                return;
            }
a0e
            if (!to) to = new node();
019
            int m = (1+r)/2;
ee4
            paint(from ->1, to ->1, a, b, 1, m), paint(from ->r, to ->r,
1cb
```

1.22 SegTree Esparsa - Lazy

```
// Query: soma do range [a, b]
// Update: flipa os valores de [a, b]
// O MAX tem q ser Q log N para Q updates
//
// Complexidades:
// build - O(1)
// query - O(log(n))
// update - 0(log(n))
// dc37e6
aa4 namespace seg {
6de
        int seg[MAX], lazy[MAX], R[MAX], L[MAX], ptr;
e9a
        int get_l(int i){
3db
            if (L[i] == 0) L[i] = ptr++;
a 96
            return L[i]:
b6e
        }
943
        int get_r(int i){
71b
            if (R[i] == 0) R[i] = ptr++;
283
            return R[i];
        }
43a
e71
        void build() { ptr = 2; }
ceb
        void prop(int p, int 1, int r) {
            if (!lazv[p]) return;
b77
76c
            seg[p] = r-l+1 - seg[p];
            if (1 != r) lazy[get_l(p)]^=lazy[p],
213
   lazy[get_r(p)]^=lazy[p];
3c7
            lazv[p] = 0;
20b
        }
        int query(int a, int b, int p=1, int l=0, int r=N-1) {
158
6b9
            prop(p, 1, r):
            if (b < 1 \text{ or } r < a) \text{ return } 0:
786
527
            if (a <= 1 and r <= b) return seg[p];</pre>
ee4
            int m = (1+r)/2;
818
            return query(a, b, get_l(p), l, m)+query(a, b, get_r(p),
```

```
m+1, r);
      }
0d9
        int update(int a, int b, int p=1, int l=0, int r=N-1) {
6b9
            prop(p, 1, r);
e9f
            if (b < l or r < a) return seg[p];</pre>
9a3
            if (a \le 1 \text{ and } r \le b)
                lazy[p] ^= 1;
ab6
6b9
                prop(p, 1, r);
534
                return seg[p];
8e4
            }
            int m = (1+r)/2;
43a
            return seg[p] = update(a, b, get_l(p), l, m)+update(a, b,
   get_r(p), m+1, r);
1dc
dc3 }:
      SegTree Esparsa - O(q) memoria
```

```
// Query: min do range [a, b]
// Update: troca o valor de uma posicao
// Usa O(q) de memoria para q updates
//
// Complexidades:
// query - 0(log(n))
// update - 0(log(n))
// 072a21
13d template < typename T > struct seg {
3c9
        struct node {
d53
            node* ch[2];
970
            char d;
ca0
            T v;
            T mi;
            node(int d_, T v_, T val) : d(d_), v(v_) {
d4e
                ch[0] = ch[1] = NULL;
e71
d6e
                mi = val;
065
            node(node* x) : d(x->d), v(x->v), mi(x->mi) {
b32
c99
                ch[0] = x -> ch[0], ch[1] = x -> ch[1];
cb7
            }
01e
            void update() {
909
                mi = numeric_limits <T>::max();
151
                for (int i = 0; i < 2; i++) if (ch[i])
b5a
                     mi = min(mi, ch[i]->mi);
```

```
fe3
            }
530
        };
        node* root;
bb7
9c5
        char n:
ba7
        seg() : root(NULL), n(0) {}
512
        \simseg() {
4c0
            std::vector<node*> q = {root};
402
            while (q.size()) {
e5d
                 node* x = q.back(); q.pop_back();
                 if (!x) continue;
ee9
73f
                 q.push_back(x->ch[0]), q.push_back(x->ch[1]);
bf0
                 delete x;
d3e
            }
        }
d8c
        char msb(T v, char l, char r) { // msb in range (l, r]
1a6
            for (char i = r; i > 1; i--) if (v>>i&1) return i;
8e4
daa
            return -1:
688
        }
430
        void cut(node* at, T v, char i) {
677
            char d = msb(v ^a at -> v, at -> d, i);
23b
            if (d == -1) return; // no need to split
ebf
            node* nxt = new node(at):
d43
            at -> ch \lceil v >> d&1 \rceil = NULL:
34f
            at -> ch[!(v>>d&1)] = nxt;
150
            at -> d = d:
        }
0b3
6e5
        node* update(node* at, T idx, T val, char i) {
c8c
            if (!at) return new node(-1, idx, val);
d67
            cut(at. idx. i):
1a2
            if (at->d == -1) { // leaf
792
                 at->mi = val;
ce6
                 return at;
            }
a6f
b29
            bool dir = idx>>at->d&1;
c8f
            at->ch[dir] = update(at->ch[dir], idx, val, at->d-1);
7b4
            at->update();
ce6
            return at;
76d
85 c
        void update(T idx, T val) {
8f4
            while (idx >> n) n++;
61e
            root = update(root, idx, val, n-1);
79d
        }
```

```
9d8
        T query(node* at, T a, T b, T l, T r, char i) {
            if (!at or b < l or r < a) return numeric_limits<T>::max();
df0
fd3
            if (a <= 1 and r <= b) return at->mi;
            T m = 1 + (r-1)/2;
841
            if (at->d < i) {</pre>
c85
c59
                if ((at->v>>i&1) == 0) return query(at, a, b, 1, m,
   i-1):
                else return query(at, a, b, m+1, r, i-1);
ca4
934
373
            return min(query(at->ch[0], a, b, 1, m, i-1),
   query(at->ch[1], a, b, m+1, r, i-1));
2db
6f6
        T query (T 1, T r) { return query (root, 1, r, 0, (1 \le n) - 1,
   n-1); }
072 };
```

1.24 SegTree Iterativa

```
// Consultas 0-based
// Valores iniciais devem estar em (seg[n], ..., seg[2*n-1])
// Query: soma do range [a, b]
// Update: muda o valor da posicao p para x
//
// Complexidades:
// build - O(n)
// query - O(log(n))
// update - O(log(n))
// 779519
6a4 int seg[2 * MAX];
1a8 int n;
0a8 void build() {
d15
        for (int i = n - 1; i; i--) seg[i] = seg[2*i] + seg[2*i+1];
9a8 }
4ea int query(int a, int b) {
7c9
        int ret = 0;
        for (a += n, b += n; a <= b; ++a /= 2, --b /= 2) {
728
4ea
            if (a % 2 == 1) ret += seg[a];
            if (b % 2 == 0) ret += seg[b];
244
ac0
        }
edf
        return ret;
24a }
ff3 void update(int p, int x) {
37d
        seg[p += n] = x;
```

```
c8c while (p /= 2) seg[p] = seg[2*p] + seg[2*p+1];
02d }
```

1.25 SegTree Iterativa com Lazy Propagation

```
// Query: soma do range [a, b]
// Update: soma x em cada elemento do range [a, b]
// Para mudar, mudar as funcoes junta, poe e query
// LOG = ceil(log2(MAX))
//
// Complexidades:
// build - O(n)
// query - O(log(n))
// update - 0(log(n))
// 6dc475
aa4 namespace seg {
        11 seg[2*MAX], lazy[2*MAX];
1a8
        int n:
9b3
        ll junta(ll a, ll b) {
534
            return a+b;
e26
        }
        // soma x na posicao p de tamanho tam
1 b 4
        void poe(int p, ll x, int tam, bool prop=1) {
517
            seg[p] += x*tam:
6ae
            if (prop and p < n) lazy[p] += x;</pre>
8bc
        }
        // atualiza todos os pais da folha p
b1e
        void sobe(int p) {
d5a
            for (int tam = 2; p /= 2; tam *= 2) {
4ca
                 seg[p] = junta(seg[2*p], seg[2*p+1]);
388
                poe(p, lazy[p], tam, 0);
            }
acd
b76
        }
        // propaga o caminho da raiz ate a folha p
a0a
        void prop(int p) {
076
            int tam = 1 << (LOG-1);</pre>
0a8
            for (int s = LOG; s; s--, tam /= 2) {
4b1
                int i = p >> s;
27 c
                if (lazv[i]) {
860
                     poe(2*i, lazy[i], tam);
e38
                     poe(2*i+1, lazy[i], tam);
b97
                     lazy[i] = 0;
```

```
de8
3ed
e29
        }
61c
        void build(int n2, int* v) {
            n = n2;
1e3
95f
            for (int i = 0; i < n; i++) seg[n+i] = v[i];
            for (int i = n-1; i; i--) seg[i] = junta(seg[2*i],
   seg[2*i+1]);
            for (int i = 0; i < 2*n; i++) lazy[i] = 0;
        }
8bb
4f3
        11 query(int a, int b) {
b73
            ll ret = 0;
b48
            for (prop(a+=n), prop(b+=n); a \le b; ++a/=2, --b/=2) {
                if (a%2 == 1) ret = junta(ret, seg[a]);
a8e
c58
                if (b%2 == 0) ret = junta(ret, seg[b]);
            }
510
edf
            return ret;
38b
        }
        void update(int a, int b, int x) {
a28
            int a2 = a += n, b2 = b += n, tam = 1;
c2d
Off
            for (; a <= b; ++a/=2, --b/=2, tam *= 2) {
32a
                if (a\%2 == 1) poe(a, x, tam);
9da
                if (b\%2 == 0) poe(b, x, tam);
9bc
            sobe(a2), sobe(b2);
0f7
adc
6dc };
1.26 SegTree PA
// Segtree de PA
// update_set(1, r, A, R) seta [1, r] para PA(A, R),
// update_add soma PA(A, R) em [1, r]
// query(l, r) retorna a soma de [l, r]
// PA(A, R) eh a PA: [A+R, A+2R, A+3R, ...]
//
// Complexidades:
// construir - O(n)
// update_set, update_add, query - O(log(n))
// bc4746
```

dc7 struct seg_pa {

struct Data {

350

```
8f5
            11 sum:
662
            11 set_a, set_r, add_a, add_r;
9b7
            Data(): sum(0), set_a(LINF), set_r(0), add_a(0), add_r(0)
   {}
eb6
        };
16a
        vector < Data > seg;
1a8
        int n:
        seg_pa(int n_) {
d45
e95
            n = n :
fc3
            seg = vector < Data > (4*n);
        }
ce0
ceb
        void prop(int p, int l, int r) {
d5a
            int tam = r-1+1:
c3f
            11 &sum = seg[p].sum, &set_a = seg[p].set_a, &set_r =
   seg[p].set_r,
a1b
                &add_a = seg[p].add_a, &add_r = seg[p].add_r;
c02
            if (set a != LINF) {
660
                set_a += add_a, set_r += add_r;
06e
                sum = set_a*tam + set_r*tam*(tam+1)/2;
579
                if (1 != r) {
                    int m = (1+r)/2;
ee4
886
                     seg[2*p].set_a = set_a;
358
                     seg[2*p].set_r = set_r;
ed6
                     seg[2*p].add_a = seg[2*p].add_r = 0;
f0c
                     seg[2*p+1].set_a = set_a + set_r * (m-l+1);
471
                     seg[2*p+1].set_r = set_r;
d48
                     seg[2*p+1].add_a = seg[2*p+1].add_r = 0;
                }
a97
823
                set a = LINF, set r = 0:
                add_a = add_r = 0;
953
105
            } else if (add_a or add_r) {
                sum += add_a*tam + add_r*tam*(tam+1)/2;
18b
579
                if (1 != r) {
                    int m = (1+r)/2;
ee4
ff0
                     seg[2*p].add_a += add_a;
ec0
                     seg[2*p].add_r += add_r;
06c
                     seg[2*p+1].add_a += add_a + add_r * (m-l+1);
a6d
                     seg[2*p+1].add_r += add_r;
8af
                }
953
                add a = add r = 0:
```

```
551
07f
        int inter(pair<int, int> a, pair<int, int> b) {
0b7
98c
            if (a.first > b.first) swap(a, b);
            return max(0, min(a.second, b.second) - b.first + 1);
eef
628
        }
        11 set(int a, int b, ll aa, ll rr, int p, int l, int r) {
be1
6b9
            prop(p, 1, r);
457
            if (b < l or r < a) return seg[p].sum;</pre>
9a3
            if (a <= 1 and r <= b) {</pre>
91c
                seg[p].set_a = aa;
774
                seg[p].set r = rr:
                prop(p, 1, r);
6b9
254
                return seg[p].sum;
            }
8ee
ee4
            int m = (1+r)/2;
            int tam_1 = inter({1, m}, {a, b});
963
            return seg[p].sum = set(a, b, aa, rr, 2*p, 1, m) +
c34
365
                set(a. b. aa + rr * tam 1. rr. 2*p+1. m+1. r):
8e2
        }
f55
        void update_set(int 1, int r, 11 aa, 11 rr) {
            set(1, r, aa, rr, 1, 0, n-1);
6f7
913
        }
5f6
        11 add(int a, int b, ll aa, ll rr, int p, int l, int r) {
6b9
            prop(p, 1, r):
            if (b < 1 or r < a) return seg[p].sum;</pre>
457
            if (a <= 1 and r <= b) {</pre>
9a3
                seg[p].add_a += aa;
359
1ee
                seg[p].add_r += rr;
                prop(p, 1, r);
6b9
254
                return seg[p].sum;
d19
            }
ee4
            int m = (1+r)/2:
963
            int tam_1 = inter({1, m}, {a, b});
586
            return seg[p].sum = add(a, b, aa, rr, 2*p, 1, m) +
695
                add(a, b, aa + rr * tam_1, rr, 2*p+1, m+1, r);
904
848
        void update_add(int 1, int r, 11 aa, 11 rr) {
afa
            add(1, r, aa, rr, 1, 0, n-1);
81e
f45
        11 query(int a, int b, int p, int l, int r) {
6b9
            prop(p, 1, r);
            if (b < 1 \text{ or } r < a) \text{ return } 0:
786
e9a
            if (a <= 1 and r <= b) return seg[p].sum;</pre>
            int m = (1+r)/2:
ee4
            return query(a, b, 2*p, 1, m) + query(a, b, 2*p+1, m+1, r);
b1f
```

```
f6e
bfc
        11 query(int 1, int r) { return query(1, r, 1, 0, n-1); }
bc4 };
1.27 SegTree Persistente
// SegTree de soma, update de somar numa posicao
// query(a, b, t) retorna a query de [a, b] na versao t
// update(a, x, t) faz um update v[a]+=x a partir da
// versao de t. criando uma nova versao e retornando seu id
// Por default, faz o update a partir da ultima versao
//
// build - O(n)
// query - 0(log(n))
// update - 0(log(n))
// 50ab73
54a const int MAX = 1e5+10. UPD = 1e5+10. LOG = 18:
6de const int MAXS = 2*MAX+UPD*LOG;
f6e namespace perseg {
bd6
        11 seg[MAXS];
f4e
        int rt[UPD], L[MAXS], R[MAXS], cnt, t;
052
        int n, *v;
3c4
        11 build(int p, int 1, int r) {
            if (1 == r) return seg[p] = v[1];
6cd
855
            L[p] = cnt++, R[p] = cnt++;
            int m = (1+r)/2:
ee4
275
            return seg[p] = build(L[p], 1, m) + build(R[p], m+1, r);
39d
0d8
        void build(int n2, int* v2) {
680
            n = n2, v = v2;
            rt[0] = cnt++:
856
            build(0, 0, n-1);
c50
a2e
f45
        11 query(int a, int b, int p, int l, int r) {
            if (b < 1 or r < a) return 0;
786
527
            if (a <= l and r <= b) return seg[p];</pre>
            int m = (1+r)/2;
ee4
1ed
            return query(a, b, L[p], 1, m) + query(a, b, R[p], m+1, r);
4d2
        }
182
        11 query(int a, int b, int tt) {
c13
            return query(a, b, rt[tt], 0, n-1);
726
        }
```

11 update(int a, int x, int lp, int p, int l, int r) {

bb3

```
if (1 == r) return seg[p] = seg[lp]+x;
747
ee4
            int m = (1+r)/2:
ab8
            if (a <= m)
                return seg[p] = update(a, x, L[lp], L[p]=cnt++, 1, m)
   + seg[R[p]=R[lp]];
            return seg[p] = seg[L[p]=L[lp]] + update(a, x, R[lp],
   R[p]=cnt++, m+1, r):
788
6f6
        int update(int a, int x, int tt=t) {
            update(a, x, rt[tt], rt[++t]=cnt++, 0, n-1);
ab3
e0d
            return t;
d63
       }
26f }:
```

1.28 SegTree Persistente com Lazy

```
// Nao propaga, meio estranho de mexer, mas da
// query(a, b, t) retorna a query de [a, b] na versao t
// update(a, b, x, t) faz um update v[a..b]+=x a partir da
// versao de t. criando uma nova versao e retornando seu id
// Por default, faz o update a partir da ultima versao
//
// build - O(n)
// query - 0(log(n))
// update - 0(log(n))
// 7447e3
54a const int MAX = 1e5+10, UPD = 1e5+10, LOG = 18;
ab3 const int MAXS = 2*MAX + 4*UPD*LOG:
f6e namespace perseg {
9eb
        int seg[MAXS];
f4e
       int rt[UPD], L[MAXS], R[MAXS], cnt, t;
052
        int n, *v;
adf
        int build(int p, int l, int r) {
6cd
            if (1 == r) return seg[p] = v[1];
855
           L[p] = cnt++, R[p] = cnt++;
            int m = (1+r)/2;
ee4
            return seg[p] = max(build(L[p], 1, m), build(R[p], m+1,
01d
   r));
ffd
0d8
        void build(int n2, int *v2) {
680
            n = n2, v = v2;
856
            rt[0] = cnt++;
c50
            build(0, 0, n-1);
```

```
a2e
976
        int query(int a, int b, int p, int l, int r) {
27b
             if (b < 1 or r < a) return -INF;</pre>
793
             if (a <= l and r <= b) return lazy[p] + seg[p];</pre>
             int m = (1+r)/2:
ee4
7a2
             int ret = lazy[p] + max(query(a, b, L[p], 1, m), query(a,
   b. R[p]. m+1. r)):
             return ret;
edf
9a7
442
        int query(int a, int b, int tt) {
c13
             return query(a, b, rt[tt], 0, n-1);
a05
bc1
        int update(int a, int b, int x, int lp, int p, int l, int r) {
3f6
             tie(seg[p], lazy[p], L[p], R[p]) = {seg[lp], lazy[lp],
   L[lp], R[lp]};
            if (b < l or r < a) return seg[p] + lazy[p];</pre>
847
             if (a \le 1 \text{ and } r \le b) \text{ return seg}[p] + (lazy[p] += x);
32a
             int m = (1+r)/2;
ee4
             seg[p] = max(update(a, b, x, L[lp], L[p] = cnt++, l, m),
24a
bdb
                          update(a, b, x, R[lp], R[p] = cnt++, m+1, r);
1ed
             lazv[p] = lazv[lp];
             return seg[p] + lazy[p];
1b7
877
        }
cbf
        int update(int a, int b, int x, int tt=t) {
             assert(tt <= t);</pre>
aa8
661
             update(a, b, x, rt[tt], rt[++t]=cnt++, 0, n-1);
e0d
             return t:
        }
aad
f27 };
1.29 Sparse Table
// Resolve RMQ
// MAX2 = log(MAX)
//
// Complexidades:
// build - O(n log(n))
// query - 0(1)
// 7aa4c9
cca namespace sparse {
710
        int m[MAX2][MAX], n;
```

for (int i = 0; i < n; i++) m[0][i] = v[i];

for (int j = 1; (1<<j) <= n; j++) for (int i = 0; i+(1<<<math>j)

61c

1e3

78e

a1c

void build(int n2, int* v) {

n = n2:

```
<= n; i++)
                m[j][i] = min(m[j-1][i], m[j-1][i+(1<<(j-1))]);
5d5
cae
4ea
        int query(int a, int b) {
            int j = __builtin_clz(1) - __builtin_clz(b-a+1);
ee5
            return min(m[j][a], m[j][b-(1<<j)+1]);</pre>
dc3
fba
        }
7aa }
      Sparse Table Disjunta
// Resolve qualquer operacao associativa
// MAX2 = log(MAX)
//
// Complexidades:
// build - O(n log(n))
// query - 0(1)
// fd81ae
cca namespace sparse {
9bf
        int m[MAX2][2*MAX], n, v[2*MAX];
        int op(int a, int b) { return min(a, b); }
5f7
850
        void build(int n2, int* v2) {
1e3
            n = n2:
df4
            for (int i = 0; i < n; i++) v[i] = v2[i];
a84
            while (n&(n-1)) n++;
3d2
            for (int j = 0; (1<<j) < n; j++) {
1c0
                int len = 1<<j;</pre>
d9b
                for (int c = len; c < n; c += 2*len) {
332
                    m[j][c] = v[c], m[j][c-1] = v[c-1];
                    for (int i = c+1; i < c+len; i++) m[j][i] =</pre>
   op(m[j][i-1], v[i]);
                    for (int i = c-2; i >= c-len; i--) m[j][i] =
432
   op(v[i], m[j][i+1]);
                }
eda
            }
f4d
ce3
        int query(int 1, int r) {
9e3
```

int j = __builtin_clz(1) - __builtin_clz(1^r);

1.31 Splay Tree

}

f13

e6d

d67

a7b

fd8 }

// SEMPRE QUE DESCER NA ARVORE, DAR SPLAY NO

if (1 == r) return v[1];

return op(m[j][1], m[j][r]);

```
// NODE MAIS PROFUNDO VISITADO
// Todas as operacoes sao O(log(n)) amortizado
// Se guiser colocar mais informação no node,
// mudar em 'update'
// 4ff2b3
538 template < typename T > struct splaytree {
3c9
        struct node {
183
            node *ch[2], *p;
e4d
            int sz:
f48
            T val;
da0
            node(T v) {
696
                 ch[0] = ch[1] = p = NULL;
a26
                 sz = 1;
250
                 val = v;
2d0
            }
01e
            void update() {
a26
                 sz = 1:
с7с
                 for (int i = 0; i < 2; i++) if (ch[i]) {
d5f
                     sz += ch[i]->sz:
486
                 }
f45
            }
aa3
        };
bb7
        node* root;
fbc
        splaytree() { root = NULL; }
214
        splaytree(const splaytree& t) {
             throw logic_error("Nao copiar a splaytree!");
cbf
1f1
        }
891
        \simsplaytree() {
609
            vector < node *> q = {root};
402
            while (q.size()) {
e5d
                 node* x = q.back(); q.pop_back();
ee9
                 if (!x) continue;
73f
                 q.push_back(x->ch[0]), q.push_back(x->ch[1]);
bf0
                 delete x;
d3e
            }
837
        }
94f
        void rotate(node* x) { // x vai ficar em cima
d9b
            node *p = x->p, *pp = p->p;
ecf
            if (pp) pp - ch[pp - ch[1] == p] = x;
286
            bool d = p \rightarrow ch[0] == x;
d63
            p - ch[!d] = x - ch[d], x - ch[d] = p;
            if (p->ch[!d]) p->ch[!d]->p = p;
bad
fc2
            x->p = pp, p->p = x;
```

```
1ea
            p->update(), x->update();
007
3fa
        node* splay(node* x) {
a39
            if (!x) return x;
4ea
            root = x:
3cf
            while (x->p) {
d9b
                node *p = x->p, *pp = p->p;
359
                if (!pp) return rotate(x), x; // zig
e3c
                if ((pp->ch[0] == p)^(p->ch[0] == x))
a2b
                    rotate(x), rotate(x); // zigzag
4b2
                else rotate(p), rotate(x); // zigzig
028
            }
ea5
            return x;
379
313
        node* insert(T v, bool lb=0) {
b64
            if (!root) return lb ? NULL : root = new node(v);
            node *x = root, *last = NULL;;
002
31e
            while (1) {
5d7
                bool d = x - > val < v;
                if (!d) last = x:
0fd
                if (x->val == v) break:
c2e
c16
                if (x->ch[d]) x = x->ch[d];
                else {
4e6
                    if (lb) break;
dea
055
                    x - ch[d] = new node(v);
99c
                    x \rightarrow ch[d] \rightarrow p = x;
30e
                    x = x -> ch[d];
c2b
                    break:
                }
68a
1ab
            }
0b6
            splay(x);
61c
            return lb ? splay(last) : x;
622
сОс
        int size() { return root ? root->sz : 0: }
        int count(T v) { return insert(v, 1) and root->val == v; }
2ca
111
        node* lower_bound(T v) { return insert(v, 1); }
26b
        void erase(T v) {
446
            if (!count(v)) return;
bce
            node *x = root, *1 = x -> ch[0];
268
            if (!1) {
                root = x->ch[1];
8b1
                if (root) root->p = NULL;
32e
8f3
                return delete x;
            }
a86
5e7
            root = 1, 1 - p = NULL;
            while (1->ch[1]) 1 = 1->ch[1];
902
bab
            splay(1);
```

```
f0e
            1 - ch[1] = x - ch[1];
7d9
            if (1->ch[1]) 1->ch[1]->p = 1;
bf0
            delete x;
62a
            1->update();
007
        }
24a
        int order_of_key(T v) {
62b
            if (!lower bound(v)) return root ? root->sz : 0:
            return root -> ch[0] ? root -> ch[0] -> sz : 0;
1cc
b00
        }
db6
        node* find_by_order(int k) {
084
            if (k >= size()) return NULL;
52f
            node* x = root;
31e
            while (1) {
20f
                if (x->ch[0] \text{ and } x->ch[0]->sz >= k+1) x = x->ch[0];
4e6
                else {
                     if (x->ch[0]) k -= x->ch[0]->sz;
a1c
                     if (!k) return splay(x);
1dc
                     k--, x = x->ch[1];
eb8
aca
                }
            }
e01
        }
0de
19c
        T min() {
52f
            node* x = root;
6f6
            while (x->ch[0]) x = x->ch[0]; // max -> ch[1]
3e9
            return splay(x)->val;
70e
        }
4ff };
1.32 Splay Tree Implicita
```

```
// vector da NASA
// Um pouco mais rapido q a treap
// O construtor a partir do vector
// eh linear, todas as outras operacoes
// custam O(log(n)) amortizado
// a3575a
081 template < typename T > struct splay {
         struct node {
3c9
183
             node *ch[2], *p;
e4d
             int sz;
875
             T val, sub, lazy;
aa6
             bool rev;
da0
             node(T v) {
696
                 ch[0] = ch[1] = p = NULL;
a26
                 sz = 1;
1 e 4
                 sub = val = v;
```

```
c60
                lazv = 0;
b67
                rev = false;
48f
            }
a9c
            void prop() {
0ec
                if (lazv) {
                     val += lazy, sub += lazy*sz;
924
091
                     if (ch[0]) ch[0]->lazy += lazy;
1a8
                    if (ch[1]) ch[1]->lazy += lazy;
a98
                }
                if (rev) {
1bb
80a
                     swap(ch[0], ch[1]);
628
                     if (ch[0]) ch[0]->rev ^= 1;
adc
                    if (ch[1]) ch[1]->rev ~= 1:
30a
a32
                lazy = 0, rev = 0;
            }
6bf
01e
            void update() {
0 c 3
                sz = 1, sub = val;
                for (int i = 0; i < 2; i++) if (ch[i]) {
c7c
05f
                     ch[i]->prop();
d5f
                     sz += ch[i] -> sz;
                     sub += ch[i]->sub;
4a1
                }
6 c 1
e98
            }
        };
b4a
bb7
        node* root;
        splay() { root = NULL; }
5d9
9b1
        splay(node* x) {
4ea
            root = x:
32e
            if (root) root->p = NULL;
371
1b7
        splay(vector < T > v) { // O(n)}
950
            root = NULL;
806
            for (T i : v) {
2a0
                node* x = new node(i);
bd1
                x - ch[0] = root;
37a
                if (root) root->p = x;
4ea
                root = x;
a0a
                root ->update();
            }
17 c
c6b
a9e
        splay(const splay& t) {
e62
            throw logic_error("Nao copiar a splay!");
d4d
        \simsplay() {
5ab
```

```
609
             vector < node *> q = {root};
402
             while (q.size()) {
e5d
                 node* x = q.back(); q.pop_back();
ee9
                 if (!x) continue;
73f
                 q.push_back(x->ch[0]), q.push_back(x->ch[1]);
bf0
                 delete x;
d3e
            }
        }
a1c
         int size(node* x) { return x ? x->sz : 0; }
73c
94f
         void rotate(node* x) { // x vai ficar em cima
d9b
             node *p = x->p, *pp = p->p;
ecf
             if (pp) pp -> ch[pp -> ch[1] == p] = x:
286
             bool d = p \rightarrow ch[0] == x;
d63
             p - ch[!d] = x - ch[d], x - ch[d] = p;
             if (p->ch[!d]) p->ch[!d]->p = p;
bad
fc2
             x->p = pp, p->p = x;
1ea
             p->update(), x->update();
007
        }
6a0
         node* splaya(node* x) {
a39
             if (!x) return x;
be6
             root = x, x->update();
3cf
             while (x->p) {
d9b
                 node *p = x->p, *pp = p->p;
359
                 if (!pp) return rotate(x), x; // zig
еЗс
                 if ((pp->ch[0] == p)^(p->ch[0] == x))
a2b
                     rotate(x), rotate(x); // zigzag
4b2
                 else rotate(p), rotate(x); // zigzig
028
             }
ea5
             return x;
        }
21a
a7f
         node* find(int v) {
a2e
             if (!root) return NULL;
52f
             node *x = root;
6cd
             int key = 0;
31e
             while (1) {
857
                 x->prop();
ba1
                 bool d = kev + size(x->ch[0]) < v;
877
                 if (\text{key} + \text{size}(x->\text{ch}[0]) != v \text{ and } x->\text{ch}[d]) {
15e
                     if (d) key += size(x->ch[0])+1;
                     x = x - > ch[d]:
30e
a30
                 } else break;
3c3
152
             return splaya(x);
f19
        }
сОс
         int size() { return root ? root->sz : 0; }
         void join(splay<T>& 1) { // assume que l < *this</pre>
c26
```

```
690
            if (!size()) swap(root, l.root);
579
            if (!size() or !l.size()) return;
bee
            node* x = 1.root;
            while (1) {
31e
857
                x->prop();
                if (!x->ch[1]) break;
34d
bd8
                x = x -  ch[1]:
fa3
147
            1.splaya(x), root->prop(), root->update();
            x - ch[1] = root, x - ch[1] - p = x;
42b
0aa
            root = 1.root, 1.root = NULL;
a0a
            root ->update();
7e6
5ed
        node* split(int v) { // retorna os elementos < v</pre>
            if (v <= 0) return NULL;</pre>
398
060
            if (v >= size()) {
f87
                node* ret = root;
                root = NULL:
950
8c9
                ret->update();
edf
                return ret:
            }
d0f
adc
            find(v);
            node* 1 = root -> ch[0];
a59
4df
            root -> ch [0] = NULL;
5a3
            if (1) 1->p = NULL;
a0a
            root ->update();
792
            return 1;
826
511
        T& operator [](int i) {
9d4
            find(i);
ae0
            return root ->val;
829
231
        void push_back(T v) { // 0(1)
            node* r = new node(v):
a01
            r \rightarrow ch[0] = root:
0de
b11
            if (root) root->p = r;
b13
            root = r, root->update();
315
b7a
        T query(int 1, int r) {
95f
            splay <T > M(split(r+1));
5ff
            splay <T> L(M.split(1));
d1c
            T ans = M.root->sub;
49c
            M.join(L), join(M);
ba7
            return ans;
ca3
41f
        void update(int 1, int r, T s) {
95f
            splay <T> M(split(r+1));
```

```
5ff
            splav <T> L(M.split(1));
996
            M.root->lazy += s;
49c
            M.join(L), join(M);
9e9
        }
8c1
        void reverse(int 1. int r) {
95f
            splay <T> M(split(r+1));
5ff
            splay<T> L(M.split(1));
945
            M.root->rev ^= 1;
49c
            M. join(L), join(M);
        }
c1a
2fb
        void erase(int 1, int r) {
95f
            splay <T> M(split(r+1));
5ff
            splay <T> L(M.split(1));
dcc
            join(L);
        }
68e
a35 };
1.33 Split-Merge Set
// Representa um conjunto de inteiros nao negativos
// Todas as operacoes custam O(log(N)),
// em que N = maior elemento do set,
// exceto o merge, que custa O(log(N)) amortizado
// Usa O(min(N, n log(N))) de memoria, sendo 'n' o
// numero de elementos distintos no set
// 2d2d8a
2dc template < typename T, bool MULTI = false, typename SIZE_T = int > struct
   sms {
3 c 9
        struct node {
b19
            node *1, *r;
15f
            SIZE_T cnt;
658
            node() : l(NULL), r(NULL), cnt(0) {}
01e
            void update() {
a01
                cnt = 0:
d8a
                if (1) cnt += 1->cnt;
e49
                if (r) cnt += r->cnt;
74d
            }
        };
84f
        node* root;
bb7
fd0
        T N:
f34
        sms() : root(NULL), N(0) {}
83b
        sms(T v) : sms() { while (v >= N) N = 2*N+1; }
5e1
        sms(const sms& t) : root(NULL), N(t.N) {
3af
            for (SIZE_T i = 0; i < t.size(); i++) {</pre>
```

```
a0f
                T at = t[i]:
                                                                          841
                SIZE_T qt = t.count(at);
e6d
                                                                          a02
a43
               insert(at, qt);
                                                                          8d9
f42
               i += qt-1;
                                                                          cff
1e9
            }
                                                                          83b
                                                                           cf7
ea8
        sms(initializer list<T> v) : sms() { for (T i : v) insert(i): }
a96
2dd
        \simsms() {
                                                                          882
609
            vector < node *> q = {root};
                                                                          72b
402
            while (q.size()) {
                                                                          f52
e5d
                node* x = q.back(); q.pop_back();
                                                                          5e9
ee9
                if (!x) continue:
1c7
                q.push_back(x->1), q.push_back(x->r);
bf0
                delete x;
            }
653
f0d
        }
                                                                          54b
                                                                          4e6
        friend void swap(sms& a, sms& b) {
                                                                          841
fdc
49e
            swap(a.root, b.root), swap(a.N, b.N);
                                                                          281
984
        }
                                                                          ba1
83e
        sms& operator =(const sms& v) {
                                                                          7b4
768
            sms tmp = v;
                                                                          d3d
            swap(tmp, *this);
420
                                                                          135
357
            return *this;
                                                                          ce6
e9b
                                                                          e1f
d06
        SIZE T size() const { return root ? root->cnt : 0: }
        SIZE_T count(node* x) const { return x ? x->cnt : 0; }
17f
75a
        void clear() {
                                                                          9c3
0a0
            sms tmp:
                                                                          9dc
4ac
            swap(*this, tmp);
                                                                          b1d
fcb
        }
a06
        void expand(T v) {
bc3
            for (: N < v: N = 2*N+1) if (root) {
                                                                          347
                node* nroot = new node():
63c
                                                                          9f2
956
                nroot ->1 = root:
897
               root = nroot;
a0a
               root->update();
dd9
            }
9f0
       }
        node* insert(node* at, T idx, SIZE_T qt, T 1, T r) {
b14
1a4
            if (!at) at = new node();
            if (1 == r) {
893
435
                at->cnt += qt;
                                                                          0a9
beb
               if (!MULTI) at->cnt = 1;
ce6
               return at:
a53
            }
```

```
T m = 1 + (r-1)/2:
            if (idx \le m) at->1 = insert(at->1, idx, qt, 1, m);
            else at->r = insert(at->r, idx, qt, m+1, r);
            return at ->update(), at;
        void insert(T v, SIZE_T qt=1) { // insere 'qt' ocorrencias de
            if (qt <= 0) return erase(v, -qt);</pre>
            assert(v >= 0);
            expand(v):
            root = insert(root, v, qt, 0, N);
f62
        }
f06
        node* erase(node* at, T idx, SIZE_T qt, T 1, T r) {
28 c
            if (!at) return at;
            if (1 == r) at->cnt = at->cnt < qt ? 0 : at->cnt - qt;
            else {
                T m = 1 + (r-1)/2:
                if (idx \le m) at->1 = erase(at->1, idx, qt, 1, m);
                else at->r = erase(at->r, idx, qt, m+1, r);
                at->update();
            if (!at->cnt) delete at, at = NULL;
            return at;
43d
        void erase(T v, SIZE_T qt=1) { // remove 'qt' ocorrencias de
, <sub>17</sub> ,
            if (v < 0 \text{ or } v > N \text{ or } !at) return:
            if (qt < 0) insert(v, -qt);</pre>
            root = erase(root, v, qt, 0, N);
b32
        }
8d6
        void erase_all(T v) { // remove todos os 'v'
            if (v < 0 \text{ or } v > N) return:
            root = erase(root, v, numeric limits < SIZE T >:: max(), 0, N);
569
        }
0fe
        SIZE_T count(node* at, T a, T b, T 1, T r) const {
61b
            if (!at or b < l or r < a) return 0;</pre>
0fe
            if (a <= 1 and r <= b) return at->cnt;
841
            T m = 1 + (r-1)/2;
            return count(at->1, a, b, 1, m) + count(at->r, a, b, m+1,
84a
r);
4e6
        SIZE_T count(T v) const { return count(root, v, v, 0, N); }
ffc
        SIZE_T order_of_kev(T v) { return count(root, 0, v-1, 0, N); }
df2
        SIZE_T lower_bound(T v) { return order_of_key(v); }
```

```
e68
        const T operator [](SIZE_T i) const { // i-esimo menor elemento
809
            assert(i >= 0 and i < size());</pre>
c43
            node* at = root;
4a5
            T 1 = 0, r = N;
40c
            while (1 < r) {
841
                T m = 1 + (r-1)/2;
5c2
                if (count(at->1) > i) at = at->1, r = m:
4e6
                else {
b4a
                    i -= count(at->1);
                    at = at->r; 1 = m+1;
ded
fa6
                }
41a
            }
792
            return 1:
67f
        }
78c
        node* merge(node* 1, node* r) {
347
            if (!1 or !r) return 1 ? 1 : r;
504
            if (!1->1 and !1->r) { // folha
599
                if (MULTI) 1->cnt += r->cnt;
55d
                delete r:
792
                return 1;
92c
            1->1 = merge(1->1, r->1), 1->r = merge(1->r, r->r);
f58
f4f
            1->update(), delete r;
792
            return 1:
06a
        void merge(sms& s) { // mergeia dois sets
f59
068
            if (N > s.N) swap(*this, s);
785
            expand(s.N);
938
            root = merge(root, s.root);
            s.root = NULL:
ee2
2f6
        }
dc6
        node* split(node*& x. SIZE T k) {
            if (k <= 0 or !x) return NULL:
7ca
            node* ret = new node();
6d0
386
            if (!x->1 \text{ and } !x->r) x->cnt -= k, ret->cnt += k;
4e6
85e
                if (k \le count(x->1)) ret->1 = split(x->1, k);
4e6
                else {
                    ret->r = split(x->r, k - count(x->1));
06f
                    swap(x->1, ret->1);
cfd
63b
                ret ->update(), x ->update();
674
379
d5b
            if (!x->cnt) delete x, x = NULL;
edf
            return ret:
```

```
f18
02b
        void split(SIZE_T k, sms& s) { // pega os 'k' menores
e63
            s.clear();
6e5
            s.root = split(root, min(k, size()));
            s.N = N:
e3c
9a6
        }
        // pega os menores que 'k'
        void split_val(T k, sms& s) { split(order_of_key(k), s); }
131
2d2 };
1.34 SQRT Tree
// RMQ em O(log log n) com O(n log log n) pra buildar
// Funciona com qualquer operacao associativa
// Tao rapido quanto a sparse table, mas usa menos memoria
// (log log (1e9) < 5, entao a query eh praticamente O(1))
//
// build - O(n log log n)
// query - O(log log n)
// 8ff986
97a namespace sqrtTree {
052
        int n, *v;
        int pref[4][MAX], sulf[4][MAX], getl[4][MAX], entre[4][MAX],
ec7
   sz[4];
5f7
        int op(int a, int b) { return min(a, b); }
        inline int getblk(int p, int i) { return (i-getl[p][i])/sz[p];
c72
 }
        void build(int p, int l, int r) {
2c6
bc8
            if (1+1 >= r) return;
368
            for (int i = 1; i <= r; i++) getl[p][i] = 1;</pre>
f16
            for (int L = 1; L <= r; L += sz[p]) {</pre>
                int R = min(L+sz[p]-1, r);
191
                pref[p][L] = v[L], sulf[p][R] = v[R];
89 c
                for (int i = L+1; i <= R; i++) pref[p][i] =</pre>
59f
   op(pref[p][i-1], v[i]);
d9a
                for (int i = R-1; i >= L; i--) sulf[p][i] = op(v[i],
    sulf[p][i+1]);
221
                build(p+1, L, R);
c7b
695
            for (int i = 0: i \le sz[p]: i++) {
                int at = entre[p][1+i*sz[p]+i] = sulf[p][1+i*sz[p]];
ca5
759
                for (int j = i+1; j <= sz[p]; j++)</pre>
   entre[p][l+i*sz[p]+j] = at =
23a
                        op(at, sulf[p][l+j*sz[p]]);
c51
            }
```

```
861
        }
0d8
        void build(int n2, int* v2) {
680
            n = n2, v = v2;
            for (int p = 0; p < 4; p++) sz[p] = n2 = sqrt(n2);
44c
c50
            build(0, 0, n-1);
940
9e3
        int query(int 1, int r) {
792
            if (l+1 >= r) return l == r ? v[l] : op(v[l], v[r]);
1ba
            int p = 0;
4ba
            while (getblk(p, 1) == getblk(p, r)) p++;
9e4
            int ans = sulf[p][1], a = getblk(p, 1)+1, b = getblk(p, 1)+1
   r)-1;
8bf
            if (a \le b) ans = op(ans, entre[p][get1[p][1]+a*sz[p]+b]);
dea
            return op(ans, pref[p][r]);
589
        }
8ff }
1.35 Treap
// Todas as operacoes custam
```

```
// O(log(n)) com alta probabilidade, exceto meld
// meld custa O(log^2 n) amortizado com alta prob.,
// e permite unir duas treaps sem restricao adicional
// Na pratica, esse meld tem constante muito boa e
// o pior caso eh meio estranho de acontecer
// bd93e2
878 mt19937 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
aa1 template < typename T> struct treap {
3c9
        struct node {
b19
            node *1, *r;
284
            int p, sz;
36d
            T val, mi;
4c7
            node(T v) : 1(NULL), r(NULL), p(rng()), sz(1), val(v),
   mi(v) {}
            void update() {
01e
                sz = 1;
a26
d6e
                mi = val:
                if (1) sz += 1->sz, mi = min(mi, 1->mi);
bd7
a54
                if (r) sz += r->sz, mi = min(mi, r->mi):
            }
660
        };
c1b
bb7
        node* root;
```

```
84b
         treap() { root = NULL; }
2d8
         treap(const treap& t) {
465
             throw logic_error("Nao copiar a treap!");
        }
1e9
        \simtreap() {
cec
609
             vector < node *> q = {root};
402
             while (q.size()) {
e5d
                 node* x = q.back(); q.pop_back();
ee9
                 if (!x) continue;
1 c 7
                 q.push_back(x->1), q.push_back(x->r);
bf0
                 delete x;
653
             }
50e
        }
73c
         int size(node* x) { return x ? x->sz : 0; }
        int size() { return size(root); }
b2b
bcf
         void join(node* 1, node* r, node*& i) { // assume que 1 < r</pre>
             if (!l or !r) return void(i = 1 ? 1 : r);
986
             if (1->p > r->p) join(1->r, r, 1->r), i = 1;
80e
fa0
             else join(1, r->1, r->1), i = r;
bda
             i->update();
671
ece
        void split(node* i, node*& 1, node*& r, T v) {
26a
             if (!i) return void(r = 1 = NULL);
f 0.5
             if (i\rightarrow val < v) split(i\rightarrow r, i\rightarrow r, r, v), l = i;
807
             else split(i \rightarrow 1, l, i \rightarrow 1, v), r = i;
bda
             i->update();
2cd
        }
3fc
         void split_leq(node* i, node*& 1, node*& r, T v) {
26a
             if (!i) return void(r = 1 = NULL);
             if (i\rightarrow val \leftarrow v) split_leq(i\rightarrow r, i\rightarrow r, r, v), l = i;
181
58f
             else split_leg(i\rightarrow 1, 1, i\rightarrow 1, v), r = i;
             i->update();
bda
70f
        }
e13
        int count(node* i, T v) {
6b4
             if (!i) return 0;
352
             if (i->val == v) return 1;
8d0
             if (v < i->val) return count(i->1, v);
4d0
             return count(i->r, v);
5e6
26d
        void index_split(node* i, node*& 1, node*& r, int v, int key =
   0) {
26a
             if (!i) return void(r = 1 = NULL);
c10
             if (\text{key} + \text{size}(i->1) < v) index_split(i->r, i->r, r, v,
    kev+size(i->1)+1), l = i;
             else index_split(i->1, l, i->1, v, key), r = i;
e5a
bda
             i->update();
```

```
ccf
        }
a1f
        int count(T v) {
                                                                             aa1 template < typename T> struct treap {
e06
            return count(root, v);
                                                                             3c9
                                                                                      struct node {
                                                                             b19
980
        }
                                                                                          node *1, *r;
c27
        void insert(T v) {
                                                                             284
                                                                                          int p, sz;
980
            if (count(v)) return;
                                                                             875
                                                                                          T val, sub, lazy;
031
            node *L, *R;
                                                                             aa6
                                                                                          bool rev:
d42
            split(root, L, R, v);
                                                                             8dc
                                                                                          node(T v) : 1(NULL), r(NULL), p(rng()), sz(1), val(v),
585
            node* at = new node(v);
                                                                                 sub(v), lazv(0), rev(0) {}
59f
            join(L, at, L);
                                                                                          void prop() {
                                                                             a9c
a28
            join(L, R, root);
                                                                             0ec
                                                                                              if (lazy) {
37 c
                                                                             924
                                                                                                  val += lazy, sub += lazy*sz;
26b
        void erase(T v) {
                                                                             b87
                                                                                                  if (1) 1->lazy += lazy;
df9
            node *L, *M, *R;
                                                                             d3b
                                                                                                  if (r) r->lazy += lazy;
                                                                                              }
b6b
            split_leq(root, M, R, v), split(M, L, M, v);
                                                                             cea
                                                                                              if (rev) {
f17
            if (M) delete M;
                                                                             1bb
f38
            M = NULL;
                                                                             e4f
                                                                                                  swap(1, r);
a28
            join(L, R, root);
                                                                             dc8
                                                                                                  if (1) 1->rev ^= 1;
b92
                                                                             f2f
                                                                                                  if (r) r->rev ^= 1;
e77
        void meld(treap& t) { // segmented merge
                                                                             3e5
                                                                                              }
4a6
            node *L = root, *R = t.root;
                                                                             a32
                                                                                              lazy = 0, rev = 0;
950
            root = NULL:
                                                                             ca6
6b1
            while (L or R) {
                                                                             01e
                                                                                          void update() {
fe2
                if (!L or (L and R and L->mi > R->mi)) std::swap(L, R);
                                                                             0c3
                                                                                              sz = 1, sub = val;
5e1
                if (!R) join(root, L, root), L = NULL;
                                                                             a09
                                                                                              if (1) 1->prop(), sz += 1->sz, sub += 1->sub;
3c9
                 else if (L->mi == R->mi) {
                                                                             095
                                                                                              if (r) r \rightarrow prop(), sz += r \rightarrow sz, sub += r \rightarrow sub;
a76
                     node* LL;
                                                                             360
                                                                                          }
439
                     split(L, LL, L, R->mi+1);
                                                                             d37
                                                                                      };
359
                     delete LL;
9d9
                } else {
                                                                             bb7
                                                                                      node* root;
a76
                     node* LL:
537
                     split(L, LL, L, R->mi);
                                                                             84b
                                                                                      treap() { root = NULL; }
                                                                             2d8
                                                                                      treap(const treap& t) {
dbb
                     join(root, LL, root);
fc4
                }
                                                                             465
                                                                                          throw logic_error("Nao copiar a treap!");
576
                                                                             1e9
                                                                                     }
                                                                                     \simtreap() {
689
            t.root = NULL;
                                                                             cec
8e7
                                                                             609
                                                                                          vector < node *> q = {root};
        }
651 };
                                                                             402
                                                                                          while (q.size()) {
                                                                             e5d
                                                                                              node* x = q.back(); q.pop_back();
                                                                             ee9
                                                                                              if (!x) continue;
      Treap Implicita
                                                                             1c7
                                                                                              q.push_back(x->1), q.push_back(x->r);
                                                                             bf0
                                                                                              delete x;
// Todas as operacoes custam
                                                                             653
                                                                                          }
// O(log(n)) com alta probabilidade
                                                                                     }
                                                                             50e
// 63ba4d
                                                                             73c
                                                                                      int size(node* x) { return x ? x->sz : 0; }
878 mt19937 rng((int)
                                                                             b2b
                                                                                      int size() { return size(root); }
   chrono::steady_clock::now().time_since_epoch().count());
```

```
bcf
        void join(node* 1, node* r, node*& i) { // assume que 1 < r</pre>
             if (!1 or !r) return void(i = 1 ? 1 : r);
986
161
             1->prop(), r->prop();
             if (1->p > r->p) join(1->r, r, 1->r), i = 1;
80e
             else join(1, r->1, r->1), i = r;
fa0
             i->update();
bda
b57
        }
a20
        void split(node* i, node*& 1, node*& r, int v, int key = 0) {
26a
             if (!i) return void(r = 1 = NULL);
c89
             i->prop();
5bd
             if (\text{key} + \text{size}(i\rightarrow 1) < v) split(i\rightarrow r, i\rightarrow r, r, v, v)
   key+size(i->1)+1), l = i;
219
             else split(i \rightarrow 1, l, i \rightarrow 1, v, key), r = i;
             i->update();
bda
d37
        }
231
        void push_back(T v) {
2e0
             node* i = new node(v);
7ab
             join(root, i, root);
46b
b7a
        T query(int 1, int r) {
df9
             node *L, *M, *R;
             split(root, M, R, r+1), split(M, L, M, 1);
dca
d43
             T ans = M->sub;
69d
             join(L, M, M), join(M, R, root);
ba7
             return ans;
1f7
41f
        void update(int 1, int r, T s) {
df9
             node *L, *M, *R;
             split(root, M, R, r+1), split(M, L, M, 1);
dca
8f6
             M \rightarrow lazy += s;
             join(L, M, M), join(M, R, root);
69d
29f
8c1
        void reverse(int 1, int r) {
df9
             node *L. *M. *R:
             split(root, M, R, r+1), split(M, L, M, 1);
dca
66a
             M \rightarrow rev = 1:
69d
             join(L, M, M), join(M, R, root);
ea8
        }
139 };
      Treap Persistent Implicita
// Todas as operacoes custam
// O(log(n)) com alta probabilidade
// fb8013
6cf mt19937_64 rng((int)
```

```
chrono::steady_clock::now().time_since_epoch().count());
3c9 struct node {
b19
        node *1, *r;
f14
        ll sz, val, sub;
304
        node(ll\ v): l(NULL), r(NULL), sz(1), val(v), sub(v) {}
c12
        node(node* x) : l(x->l), r(x->r), sz(x->sz), val(x->val),
   sub(x->sub) {}
01e
        void update() {
0c3
            sz = 1, sub = val;
77e
            if (1) sz += 1->sz, sub += 1->sub;
d6e
            if (r) sz += r->sz, sub += r->sub;
124
            sub %= MOD:
472
        }
95f };
bc9 ll size(node* x) { return x ? x->sz : 0; }
761 void update(node* x) { if (x) x->update(); }
828 node* copy(node* x) { return x ? new node(x) : NULL; }
b02 node* join(node* 1, node* r) {
        if (!1 or !r) return 1 ? copy(1) : copy(r);
e1f
48b
        node* ret;
49f
        if (rng() % (size(1) + size(r)) < size(1)) {</pre>
7eb
            ret = copy(1);
cc1
            ret -> r = join(ret -> r, r);
9d9
        } else {
4c5
            ret = copy(r);
551
            ret->1 = join(1, ret->1);
348
74 f
        return update(ret), ret;
2cc }
723 void split(node* x, node*& 1, node*& r, 11 v, 11 key = 0) {
        if (!x) return void(l = r = NULL);
421
b4b
        if (key + size(x->1) < v) {
72f
            1 = copy(x);
d70
            split(1->r, 1->r, r, v, kev+size(1->1)+1);
9d9
        } else {
303
            r = copv(x);
417
            split(r->1, l, r->l, v, key);
474
        update(1), update(r);
da2
666 }
f9e vector < node *> treap;
```

```
139 void init(const vector<11>& v) {
bbd         treap = {NULL};
969         for (auto i : v) treap[0] = join(treap[0], new node(i));
286 }

1.38 Wavelet Tree

// Usa O(sigma + n log(sigma)) de memoria,
```

```
// onde sigma = MAXN - MINN
// Depois do build, o v fica ordenado
// count(i, j, x, y) retorna o numero de elementos de
// v[i, j) que pertencem a [x, y]
// kth(i, j, k) retorna o elemento que estaria
// na poscicao k-1 de v[i, j), se ele fosse ordenado
// sum(i, j, x, y) retorna a soma dos elementos de
// v[i, j) que pertencem a [x, y]
// sumk(i, j, k) retorna a soma dos k-esimos menores
// elementos de v[i, j) (sum(i, j, 1) retorna o menor)
//
// Complexidades:
// build - O(n log(sigma))
// count - O(log(sigma))
// kth - 0(log(sigma))
// sum - O(log(sigma))
// sumk - O(log(sigma))
// 782344
597 int n, v[MAX];
578 vector < int > esq[4*(MAXN-MINN)], pref[4*(MAXN-MINN)];
f8d void build(int b = 0, int e = n, int p = 1, int l = MINN, int r =
   MAXN) {
58f
        int m = (1+r)/2; esq[p].push_back(0); pref[p].push_back(0);
f2f
        for (int i = b; i < e; i++) {</pre>
6b9
            esq[p].push_back(esq[p].back()+(v[i] <= m));
            pref[p].push_back(pref[p].back()+v[i]);
26f
206
        if (1 == r) return;
8ce
3a7
        int m2 = stable_partition(v+b, v+e, [=](int i){return i <=
347
        build(b, m2, 2*p, 1, m), build(m2, e, 2*p+1, m+1, r);
Ofb }
540 int count(int i, int j, int x, int y, int p = 1, int l = MINN, int
   r = MAXN)
2ad
       if (y < 1 \text{ or } r < x) \text{ return } 0;
4db
        if (x <= 1 and r <= y) return j-i;</pre>
```

```
ddc
        int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
        return count(ei, ej, x, y, 2*p, 1, m)+count(i-ei, j-ej, x, y,
   2*p+1, m+1, r);
3cf }
f62 int kth(int i, int j, int k, int p=1, int l = MINN, int r = MAXN) {
Зсе
        if (1 == r) return 1:
        int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
585
        if (k <= ej-ei) return kth(ei, ej, k, 2*p, 1, m);</pre>
28b
        return kth(i-ei, j-ej, k-(ej-ei), 2*p+1, m+1, r);
8b6 }
f2c int sum(int i, int j, int x, int y, int p = 1, int l = MINN, int r
   = MAXN)
2ad
        if (y < 1 \text{ or } r < x) \text{ return } 0;
        if (x <= 1 and r <= y) return pref[p][j]-pref[p][i];</pre>
2a9
ddc
        int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
        return sum(ei, ej, x, y, 2*p, 1, m) + sum(i-ei, j-ej, x, y,
   2*p+1, m+1, r);
b6d }
b84 int sumk(int i, int j, int k, int p = 1, int l = MINN, int r =
   MAXN) {
       if (1 == r) return l*k;
8a1
ddc
        int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
50c
        if (k <= ej-ei) return sumk(ei, ej, k, 2*p, 1, m);</pre>
        return pref[2*p][ej]-pref[2*p][ei]+sumk(i-ei, j-ej, k-(ej-ei),
   2*p+1. m+1. r):
940 }
```

2 Grafos

2.1 AGM Direcionada

```
// Fala o menor custo para selecionar arestas tal que
// o vertice 'r' alcance todos
// Se nao tem como, retorna LINF
//
// O(m log(n))
// dc345b

3c9 struct node {
f31    pair<ll, int> val;
4e4    ll lazy;
b19    node *l, *r;
f93    node() {}
```

```
c53
        node(pair<int, int> v) : val(v), lazy(0), l(NULL), r(NULL) {}
        void prop() {
a9c
            val.first += lazy;
768
            if (1) 1->lazy += lazy;
b87
            if (r) r->lazy += lazy;
d3b
c60
            lazv = 0:
05b
        }
296 };
de5 void merge(node*& a, node* b) {
        if (!a) swap(a, b);
c11
802
        if (!b) return;
626
        a->prop(), b->prop();
        if (a->val > b->val) swap(a, b);
d04
4b0
        merge(rand()%2? a->1 : a->r, b);
b82 }
d01 pair <11, int > pop(node *& R) {
e8f
        R->prop();
22e
        auto ret = R->val;
af0
        node* tmp = R;
3f3
        merge(R->1, R->r);
        R = R - > 1;
6c9
3e4
       if (R) R->lazy -= ret.first;
7c3
        delete tmp;
edf
        return ret;
c4e }
6f6 void apaga(node* R) { if (R) apaga(R->1), apaga(R->r), delete R; }
f13 ll dmst(int n, int r, vector<pair<int, int>, int>>& ar) {
94e
        vector < int > p(n); iota(p.begin(), p.end(), 0);
        function < int(int) > find = [&](int k) { return
   p[k] == k?k:p[k] = find(p[k]); };
        vector < node *> h(n);
2d7
        for (auto e : ar) merge(h[e.first.second], new node({e.second,
56f
   e.first.first}));
fd1
        vector < int > pai(n, -1), path(n);
        pai[r] = r;
66e
04b
       11 \text{ ans} = 0;
603
        for (int i = 0; i < n; i++) { // vai conectando todo mundo
            int u = i. at = 0:
2a3
            while (pai[u] == -1) {
cae
                if (!h[u]) { // nao tem
daa
947
                    for (auto i : h) apaga(i);
77 c
                    return LINF;
dd1
                }
167
                path[at++] = u, pai[u] = i;
```

```
55e
                auto [mi, v] = pop(h[u]);
64c
                ans += mi:
5e2
                if (pai[u = find(v)] == i) { // ciclo
                     while (find(v = path[--at]) != u)
86f
621
                         merge(h[u], h[v]), h[v] = NULL, p[find(v)] = u;
                    pai[u] = -1;
57a
                }
8b0
            }
ce8
        }
5df
947
        for (auto i : h) apaga(i);
ba7
        return ans:
e02 }
     Articulation Points
// Computa os pontos de articulação (vertices criticos) de um grafo
// art[i] armazena o numero de novas componentes criadas ao deletar
   vertice i
// se art[i] >= 1. entao vertice i eh ponto de articulação
// O(n+m)
// 0e405b
1a8 int n;
789 vector < vector < int >> g;
4ce stack<int> s;
b66 vector < int > id, art;
```

3e1 int dfs_art(int i, int& t, int p = -1) {

if (id[j] == -1) {

for (int j : g[i]) if (j != p) {

lo = min(lo, val);

if (val >= id[i]) {

art[i]++:

s.pop();

else lo = min(lo, id[j]);

int val = dfs_art(j, t, i);

while (s.top() != j) s.pop();

// if (val > id[i]) aresta i-j eh ponte

int lo = id[i] = t++:

s.push(i);

cf0

18e

cac

9a3

206

0 c 3

588

66a

bd9

2eb

1f3

238

328

762

}

```
if (p == -1 and art[i]) art[i]--;
3bd
253
        return lo:
8e1 }
d79 void compute_art_points() {
        id = vector < int > (n, -1);
597
a62
       art = vector < int > (n, 0):
6bb
       int t = 0;
d41
       for (int i = 0; i < n; i++) if (id[i] == -1)
625
            dfs_art(i, t, -1);
379 }
```

2.3 Bellman-Ford

```
// Calcula a menor distancia
// entre a e todos os vertices e
// detecta ciclo negativo
// Retorna 1 se ha ciclo negativo
// Nao precisa representar o grafo,
// soh armazenar as arestas
//
// O(nm)
// 03059b
14e int n, m;
248 int d[MAX];
e93 vector <pair <int. int >> ar: // vetor de arestas
9e2 vector < int > w;
                               // peso das arestas
6be bool bellman ford(int a) {
        for (int i = 0; i < n; i++) d[i] = INF;</pre>
        d[a] = 0;
8a8
        for (int i = 0: i <= n: i++)
4e3
891
            for (int j = 0; j < m; j++) {</pre>
6e4
                if (d[ar[j].second] > d[ar[j].first] + w[j]) {
705
                    if (i == n) return 1;
                    d[ar[j].second] = d[ar[j].first] + w[j];
e93
84b
                }
            }
a82
bb3
        return 0;
6eb }
```

2.4 Block-Cut Tree

```
// Cria a block-cut tree, uma arvore com os blocos
// e os pontos de articulação
// Blocos sao componentes 2-vertice-conexos maximais
// Uma 2-coloração da arvore eh tal que uma cor são
// os blocos, e a outra cor sao os pontos de art.
// Funciona para grafo nao conexo
//
// art[i] responde o numero de novas componentes conexas
// criadas apos a remocao de i do grafo g
// Se art[i] >= 1, i eh ponto de articulação
//
// Para todo i <= blocks.size()</pre>
// blocks[i] eh uma componente 2-vertce-conexa maximal
// edgblocks[i] sao as arestas do bloco i
// tree[i] eh um vertice da arvore que corresponde ao bloco i
// pos[i] responde a qual vertice da arvore vertice i pertence
// Arvore tem no maximo 2n vertices
// O(n+m)
// 056fa2
d10 struct block_cut_tree {
        vector < vector < int >> g, blocks, tree;
d8e
43b
        vector < vector < pair < int , int >>> edgblocks;
4ce
        stack<int> s:
6c0
        stack<pair<int, int>> s2;
2bb
        vector < int > id, art, pos;
        block_cut_tree(vector<vector<int>> g_) : g(g_) {
763
            int n = g.size();
af1
37a
            id.resize(n, -1), art.resize(n), pos.resize(n);
6f2
            build():
6bd
        }
        int dfs(int i, int& t, int p = -1) {
df6
cf0
            int lo = id[i] = t++;
18e
            s.push(i);
827
            if (p != -1) s2.emplace(i, p);
53f
            for (int j : g[i]) if (j != p and id[j] != -1)
   s2.emplace(i, j);
cac
            for (int j : g[i]) if (j != p) {
9a3
                if (id[i] == -1) {
121
                    int val = dfs(j, t, i);
0 c 3
                    lo = min(lo. val):
```

```
if (val >= id[i]) {
588
                        art[i]++;
66a
                        blocks.emplace_back(1, i);
                        while (blocks.back().back() != j)
110
                             blocks.back().push_back(s.top()), s.pop();
138
                        edgblocks.emplace_back(1, s2.top()), s2.pop();
128
47e
                        while (edgblocks.back().back() != pair(j, i))
                             edgblocks.back().push_back(s2.top()),
bce
   s2.pop();
                    // if (val > id[i]) aresta i-i eh ponte
85 c
328
                else lo = min(lo, id[j]);
            }
344
            if (p == -1 and art[i]) art[i]--;
3bd
253
            return lo;
       }
726
        void build() {
0a8
6bb
            int t = 0:
            for (int i = 0; i < g.size(); i++) if (id[i] == -1) dfs(i,
abf
   t, -1);
56c
            tree.resize(blocks.size());
f7d
            for (int i = 0; i < g.size(); i++) if (art[i])</pre>
                pos[i] = tree.size(), tree.emplace_back();
965
            for (int i = 0; i < blocks.size(); i++) for (int j :</pre>
   blocks[i]) {
403
                if (!art[j]) pos[j] = i;
                else tree[i].push_back(pos[j]),
   tree[pos[j]].push_back(i);
3df
c03
        }
056 };
2.5 Blossom - matching maximo em grafo geral
// O(n^3)
// Se for bipartido, nao precisa da funcao
// 'contract', e roda em O(nm)
// 4426a4
```

042 vector < int > g[MAX];

```
128 int match[MAX]; // match[i] = com quem i esta matchzado ou -1
1f1 int n, pai[MAX], base[MAX], vis[MAX];
26a queue < int > q;
107 void contract(int u, int v, bool first = 1) {
        static vector < bool > bloss;
165
fbe
        static int 1:
418
        if (first) {
a47
            bloss = vector < bool > (n, 0);
042
            vector < bool > teve(n, 0);
ddf
            int k = u; l = v;
31e
            while (1) {
297
                teve[k = base[k]] = 1:
116
                if (match[k] == -1) break;
dfa
                k = pai[match[k]];
68b
d31
            while (!teve[l = base[l]]) l = pai[match[l]];
5d6
        }
2e9
        while (base[u] != 1) {
            bloss[base[u]] = bloss[base[match[u]]] = 1:
e29
8fa
            pai[u] = v;
0b0
            v = match[u];
            u = pai[match[u]];
a51
58e
        }
71 c
        if (!first) return;
95e
        contract(v. u. 0):
        for (int i = 0; i < n; i++) if (bloss[base[i]]) {</pre>
6ee
594
            base[i] = 1:
            if (!vis[i]) q.push(i);
ca7
29a
            vis[i] = 1;
        }
857
e35 }
f10 int getpath(int s) {
        for (int i = 0; i < n; i++) base[i] = i, pai[i] = -1, vis[i] =</pre>
88f
   0;
ded
        vis[s] = 1; q = queue < int > (); q.push(s);
402
        while (q.size()) {
            int u = q.front(); q.pop();
be1
bdc
            for (int i : g[u]) {
                if (base[i] == base[u] or match[u] == i) continue;
7a2
                if (i == s or (match[i] != -1 and pai[match[i]] != -1))
e35
4f2
                     contract(u, i);
e2e
                else if (pai[i] == -1) {
545
                     pai[i] = u;
                     if (match[i] == -1) return i;
f6a
818
                    i = match[i]:
```

```
29d
                    vis[i] = 1; q.push(i);
                }
90e
0b5
            }
634
        }
daa
        return -1;
a16 }
83f int blossom() {
1a4
        int ans = 0;
315
        memset(match, -1, sizeof(match));
2e3
        for (int i = 0; i < n; i++) if (match[i] == -1)</pre>
f76
            for (int j : g[i]) if (match[j] == -1) {
1bc
                match[i] = i:
f1d
                match[j] = i;
Odf
                ans++;
c2b
                break;
723
            }
       for (int i = 0; i < n; i++) if (match[i] == -1) {</pre>
da8
7e3
            int j = getpath(i);
5f2
            if (j == -1) continue;
0df
            ans++;
3a0
            while (j != -1) {
                int p = pai[j], pp = match[p];
ef0
348
                match[p] = j;
fe9
                match[j] = p;
55d
                j = pp;
            }
797
f70
ba7
        return ans;
fcd }
2.6 Centro de arvore
// Retorna o diametro e o(s) centro(s) da arvore
// Uma arvore tem sempre um ou dois centros e estes estao no meio do
   diametro
//
// O(n)
// cladeb
```

042 vector<int> g[MAX];
df1 int d[MAX], par[MAX];

int f. df:

a 95

36d

d47

544 pair <int, vector <int>> center() {

function < void(int) > dfs = [&] (int v) {

if (d[v] > df) f = v, df = d[v];

```
for (int u : g[v]) if (u != par[v])
e68
                d[u] = d[v] + 1, par[u] = v, dfs(u);
1a5
90d
        };
        f = df = par[0] = -1, d[0] = 0;
1b0
41e
        dfs(0):
c2d
        int root = f:
        f = df = par[root] = -1, d[root] = 0;
0f6
14e
        dfs(root);
761
        vector<int> c;
87e
        while (f != -1) {
999
            if (d[f] == df/2 \text{ or } d[f] == (df+1)/2) \text{ c.push_back}(f);
19c
            f = par[f];
        }
3bf
00f
        return {df, c};
9c7 }
2.7 Centroid
// Computa os 2 centroids da arvore
//
// O(n)
// e16075
97a int n. subsize[MAX]:
042 vector < int > g[MAX];
98f void dfs(int k, int p=-1) {
bd2
        subsize[k] = 1;
6e5
        for (int i : g[k]) if (i != p) {
801
            dfs(i, k);
2e3
            subsize[k] += subsize[i];
1b2
        }
5a5 }
2e8 int centroid(int k, int p=-1, int size=-1) {
        if (size == -1) size = subsize[k];
e73
8df
        for (int i : g[k]) if (i != p) if (subsize[i] > size/2)
bab
            return centroid(i, k, size);
839
        return k:
b6a }
f20 pair < int , int > centroids (int k=0) {
051
        dfs(k);
        int i = centroid(k), i2 = i;
909
```

2.8 Centroid decomposition

```
// decomp(0, k) computa numero de caminhos com 'k' arestas
// Mudar depois do comentario
// O(n log(n))
// fe2541
042 vector < int > g[MAX];
ba8 int sz[MAX], rem[MAX];
747 void dfs(vector<int>& path, int i, int l=-1, int d=0) {
        path.push_back(d);
        for (int j : g[i]) if (j != 1 and !rem[j]) dfs(path, j, i,
   d+1):
3e9 }
071 int dfs sz(int i. int l=-1) {
02c
        sz[i] = 1;
        for (int j : g[i]) if (j != l and !rem[j]) sz[i] += dfs_sz(j,
   i);
        return sz[i];
191
86b }
85a int centroid(int i, int 1, int size) {
        for (int j : g[i]) if (j != 1 and !rem[j] and sz[j] > size / 2)
735
            return centroid(j, i, size);
d9a
        return i:
96e }
d79 ll decomp(int i, int k) {
106
       int c = centroid(i, i, dfs_sz(i));
a67
        rem[c] = 1;
        // gasta O(n) aqui - dfs sem ir pros caras removidos
        11 \text{ ans} = 0:
04b
        vector < int > cnt(sz[i]);
020
878
        cnt[0] = 1:
        for (int j : g[c]) if (!rem[j]) {
0a8
5b4
            vector < int > path;
baf
            dfs(path, j);
1a1
            for (int d : path) if (0 \le k-d-1 \text{ and } k-d-1 \le sz[i])
                ans += cnt[k-d-1]:
285
```

```
// Constroi a centroid tree
// p[i] eh o pai de i na centroid-tree
// dist[i][k] = distancia na arvore original entre i
// e o k-esimo ancestral na arvore da centroid
// O(n log(n)) de tempo e memoria
// a0e7c7
845 vector <int > g[MAX], dist[MAX];
c1e int sz[MAX], rem[MAX], p[MAX];
071 int dfs sz(int i, int l=-1) {
02c
        sz[i] = 1;
        for (int j : g[i]) if (j != l and !rem[j]) sz[i] += dfs_sz(j,
e5c
   i);
        return sz[i];
191
86b }
85a int centroid(int i, int 1, int size) {
        for (int j : g[i]) if (j != 1 and !rem[j] and sz[j] > size / 2)
994
735
            return centroid(j, i, size);
d9a
        return i:
96e }
324 void dfs_dist(int i, int 1, int d=0) {
541
        dist[i].push_back(d);
5a1
        for (int j : g[i]) if (j != l and !rem[j])
82a
            dfs_dist(j, i, d+1);
645 }
27e void decomp(int i, int l = -1) {
106
        int c = centroid(i, i, dfs_sz(i));
        rem[c] = 1, p[c] = 1;
1 b 9
534
        dfs_dist(c, c);
a2a
        for (int j : g[c]) if (!rem[j]) decomp(j, c);
ebd }
```

```
76c void build(int n) {
        for (int i = 0; i < n; i++) rem[i] = 0, dist[i].clear();</pre>
867
        for (int i = 0; i < n; i++) reverse(dist[i].begin(),
   dist[i].end());
a78 }
2.10 Dijkstra
// encontra menor distancia de x
// para todos os vertices
// se ao final do algoritmo d[i] = LINF,
// entao x nao alcanca i
// O(m log(n))
// 695ac4
eff ll d[MAX];
c0d vector<pair<int, int>> g[MAX]; // {vizinho, peso}
1a8 int n;
abc void dijkstra(int v) {
        for (int i = 0; i < n; i++) d[i] = LINF;</pre>
22c
a7f
        d[v] = 0;
88 c
        priority_queue < pair < ll, int >> pq;
b32
        pq.emplace(0, v);
265
        while (pq.size()) {
a25
            auto [ndist, u] = pq.top(); pq.pop();
            if (-ndist > d[u]) continue;
953
            for (auto [idx, w] : g[u]) if (d[idx] > d[u] + w) {
cda
                d[idx] = d[u] + w:
331
a84
                pq.emplace(-d[idx], idx);
c56
            }
        }
e5c
fec }
2.11 Dinitz
// O(min(m * max flow, n^2 m))
// Grafo com capacidades 1: O(\min(m \text{ sqrt}(m), m * n^2(2/3)))
// Todo vertice tem grau de entrada ou saida 1: O(m sqrt(n))
// 86fd2c
472 struct dinitz {
```

```
61f
        const bool scaling = false; // com scaling -> 0(nm
   log(MAXCAP)),
        int lim;
                                     // com constante alta
206
670
        struct edge {
358
            int to, cap, rev, flow;
7f9
            bool res;
d36
            edge(int to_, int cap_, int rev_, bool res_)
                : to(to_), cap(cap_), rev(rev_), flow(0), res(res_) {}
a94
f70
        };
002
        vector < vector < edge >> g;
216
        vector<int> lev, beg;
a71
        11 F:
190
        dinitz(int n) : g(n), F(0) {}
        void add(int a, int b, int c) {
087
            g[a].emplace_back(b, c, g[b].size(), false);
bae
4c6
            g[b].emplace_back(a, 0, g[a].size()-1, true);
5c2
        }
123
        bool bfs(int s, int t) {
90f
            lev = vector<int>(g.size(), -1); lev[s] = 0;
            beg = vector<int>(g.size(), 0);
64 c
8b2
            queue < int > q; q.push(s);
402
            while (q.size()) {
be1
                int u = q.front(); q.pop();
bd9
                for (auto& i : g[u]) {
dbc
                    if (lev[i.to] != -1 or (i.flow == i.cap)) continue;
b4f
                    if (scaling and i.cap - i.flow < lim) continue;</pre>
                    lev[i.to] = lev[u] + 1;
185
8ca
                    q.push(i.to);
f97
                }
e87
            }
            return lev[t] != -1;
0de
742
        }
        int dfs(int v, int s, int f = INF) {
dfb
50b
            if (!f or v == s) return f;
88f
            for (int& i = beg[v]; i < g[v].size(); i++) {</pre>
027
                auto& e = g[v][i];
206
                if (lev[e.to] != lev[v] + 1) continue;
ee0
                int foi = dfs(e.to, s, min(f, e.cap - e.flow));
749
                if (!foi) continue:
                e.flow += foi, g[e.to][e.rev].flow -= foi;
3 c 5
                return foi;
45c
            }
618
bb3
            return 0;
4b1
        }
ff6
        ll max_flow(int s, int t) {
```

```
a86
            for (lim = scaling ? (1 << 30) : 1; lim; lim /= 2)
                                                                             76a
                                                                                     void dfs(int v) {
                 while (bfs(s, t)) while (int ff = dfs(s, t)) F += ff;
9d1
                                                                             6a1
                                                                                         static int t = 0:
4ff
            return F;
                                                                             db6
                                                                                         pre[v] = ++t;
                                                                                         sdom[v] = label[v] = v;
8b9
        }
                                                                             767
86f }:
                                                                                         preorder.push_back(v);
                                                                             a3d
                                                                                         for (int nxt: g[v]) {
                                                                             80b
// Recupera as arestas do corte s-t
                                                                             56c
                                                                                             if (sdom[nxt] == -1) {
// 1e889c
                                                                                                 prv[nxt] = v;
                                                                             eed
dbd vector<pair<int, int>> get_cut(dinitz& g, int s, int t) {
                                                                             900
                                                                                                  dfs(nxt);
        g.max_flow(s, t);
                                                                             f48
68 c
        vector < pair < int , int >> cut;
                                                                             2b5
                                                                                             rg[nxt].push_back(v);
1b0
        vector < int > vis(g.g.size(), 0), st = {s};
                                                                             5a1
                                                                                         }
321
        vis[s] = 1:
                                                                             d6a
                                                                                     }
        while (st.size()) {
3c6
                                                                             62e
                                                                                     int eval(int v) {
b17
            int u = st.back(); st.pop_back();
                                                                             c93
                                                                                         if (ancestor[v] == -1) return v;
322
            for (auto e : g.g[u]) if (!vis[e.to] and e.flow < e.cap)</pre>
                                                                             a75
                                                                                         if (ancestor[ancestor[v]] == -1) return label[v];
c17
                 vis[e.to] = 1, st.push_back(e.to);
                                                                             f33
                                                                                         int u = eval(ancestor[v]);
d14
                                                                             b49
                                                                                         if (pre[sdom[u]] < pre[sdom[label[v]]]) label[v] = u;</pre>
        for (int i = 0; i < g.g.size(); i++) for (auto e : g.g[i])
                                                                             66e
                                                                                         ancestor[v] = ancestor[u];
481
            if (vis[i] and !vis[e.to] and !e.res) cut.emplace_back(i,
                                                                                         return label[v]:
                                                                             c24
                                                                                     }
   e.to):
                                                                             0ъ9
                                                                             4b2
                                                                                     void dfs2(int v) {
d1b
        return cut;
                                                                             6a1
                                                                                         static int t = 0:
1e8 }
                                                                             330
                                                                                         dfs_1[v] = t++;
                                                                             5e0
                                                                                         for (int nxt: tree[v]) dfs2(nxt);
2.12 Dominator Tree - Kawakami
                                                                             8e2
                                                                                         dfs r[v] = t++:
                                                                             cfa
                                                                                     }
// Se vira pra usar ai
                                                                             c2c
                                                                                     void build(int s) {
                                                                             603
                                                                                         for (int i = 0; i < n; i++) {</pre>
// build - O(m log(n))
                                                                             e6f
                                                                                             sdom[i] = pre[i] = ancestor[i] = -1;
// dominates - \Omega(1)
                                                                             2e1
                                                                                             rg[i].clear();
// c80920
                                                                             50a
                                                                                             tree[i].clear();
                                                                                             bucket[i].clear();
                                                                             666
1a8 int n;
                                                                             3ba
                                                                                         }
                                                                             772
                                                                                         preorder.clear();
bbf namespace d_tree {
                                                                             c6c
                                                                                         dfs(s);
042
        vector < int > g[MAX];
                                                                             12b
                                                                                         if (preorder.size() == 1) return;
                                                                             3c7
                                                                                         for (int i = int(preorder.size()) - 1; i >= 1; i--) {
        // The dominator tree
                                                                             6c6
                                                                                              int w = preorder[i];
        vector < int > tree[MAX];
b39
                                                                             a52
                                                                                             for (int v: rg[w]) {
        int dfs_1[MAX], dfs_r[MAX];
5af
                                                                                                  int u = eval(v);
                                                                             5 c 1
                                                                                                  if (pre[sdom[u]] < pre[sdom[w]]) sdom[w] = sdom[u];</pre>
                                                                             a17
        // Auxiliary data
                                                                             018
        vector < int > rg[MAX], bucket[MAX];
a2e
                                                                                             bucket[sdom[w]].push_back(w);
                                                                             680
3ef
        int idom[MAX], sdom[MAX], prv[MAX], pre[MAX];
                                                                             ea7
                                                                                             ancestor[w] = prv[w];
44b
        int ancestor[MAX], label[MAX];
                                                                             b99
                                                                                             for (int v: bucket[prv[w]]) {
```

5c1

int u = eval(v);

563

vector<int> preorder;

```
977
                    idom[v] = (u == v) ? sdom[v] : u;
                }
aff
2cc
                bucket[prv[w]].clear();
            }
0a3
            for (int i = 1; i < preorder.size(); i++) {</pre>
d0c
6c6
                int w = preorder[i];
14b
                if (idom[w] != sdom[w]) idom[w] = idom[idom[w]];
                tree[idom[w]].push_back(w);
32f
c58
            idom[s] = sdom[s] = -1;
8ac
1b6
            dfs2(s);
d09
        }
        // Whether every path from s to v passes through u
490
        bool dominates(int u, int v) {
c75
            if (pre[v] == -1) return 1; // vacuously true
2ea
            return dfs_l[u] <= dfs_l[v] && dfs_r[v] <= dfs_r[u];</pre>
332
ce9 };
```

2.13 Euler Path / Euler Cycle

```
// Para declarar: 'euler < true > E(n); ' se guiser
// direcionado e com 'n' vertices
// As funcoes retornam um par com um booleano
// indicando se possui o cycle/path que voce pediu,
// e um vector de {vertice. id da aresta para chegar no vertice}
// Se for get_path, na primeira posicao o id vai ser -1
// get_path(src) tenta achar um caminho ou ciclo euleriano
// comecando no vertice 'src'.
// Se achar um ciclo, o primeiro e ultimo vertice serao 'src'.
// Se for um P3, um possiveo retorno seria [0, 1, 2, 0]
// get_cycle() acha um ciclo euleriano se o grafo for euleriano.
// Se for um P3, um possivel retorno seria [0, 1, 2]
// (vertie inicial nao repete)
//
// O(n+m)
// 7113df
63f template <bool directed=false > struct euler {
1a8
4c0
        vector < vector < pair < int , int >>> g;
        vector < int > used;
d63
30f
        euler(int n_) : n(n_), g(n) {}
50f
        void add(int a, int b) {
4cd
            int at = used.size();
```

```
c51
            used.push_back(0);
74e
            g[a].emplace_back(b, at);
fab
            if (!directed) g[b].emplace_back(a, at);
        }
411
d41 #warning chamar para o src certo!
        pair < bool, vector < pair < int, int >>> get_path(int src) {
eed
baf
            if (!used.size()) return {true, {}};
b25
            vector < int > beg(n, 0);
            for (int& i : used) i = 0;
4ec
            // {{vertice, anterior}, label}
            vector<pair<pair<int, int>, int>> ret, st = {{{src, -1}},
363
   -1}};
3c6
            while (st.size()) {
8ff
                int at = st.back().first.first;
002
                int& it = beg[at];
8a1
                while (it < g[at].size() and used[g[at][it].second])</pre>
   it++;
                if (it == g[at].size()) {
8e4
                     if (ret.size() and ret.back().first.second != at)
9dd
                         return {false, {}}:
b82
420
                     ret.push_back(st.back()), st.pop_back();
9d9
daa
                     st.push_back({{g[at][it].first, at},
   g[at][it].second});
eb8
                     used[g[at][it].second] = 1;
d14
                }
b3a
a 19
            if (ret.size() != used.size()+1) return {false, {}};
f77
            vector<pair<int, int>> ans;
fdf
            for (auto i : ret) ans.emplace_back(i.first.first,
   i.second):
459
            reverse(ans.begin(), ans.end());
997
            return {true, ans};
844
        }
9b6
        pair < bool, vector < pair < int, int >>> get_cycle() {
baf
            if (!used.size()) return {true, {}};
ad1
            int src = 0;
34b
            while (!g[src].size()) src++;
687
            auto ans = get_path(src);
33c
            if (!ans.first or ans.second[0].first !=
   ans.second.back().first)
b82
                return {false, {}};
            ans.second[0].second = ans.second.back().second;
350
            ans.second.pop_back();
868
ba7
            return ans;
48f
        }
711 }:
```

2.14 Euler Tour Tree

```
// Mantem uma floresta enraizada dinamicamente
// e permite queries/updates em sub-arvore
// Chamar ETT E(n, v), passando n = numero de vertices
// e v = vector com os valores de cada vertice (se for vazio,
// constroi tudo com 0
// link(v, u) cria uma aresta de v pra u, de forma que u se torna
// o pai de v (eh preciso que v seja raiz anteriormente)
// cut(v) corta a resta de v para o pai
// query(v) retorna a soma dos valores da sub-arvore de v
// update(v, val) soma val em todos os vertices da sub-arvore de v
// update_v(v, val) muda o valor do vertice v para val
// is_in_subtree(v, u) responde se o vertice u esta na sub-arvore de v
// Tudo O(log(n)) com alta probabilidade
// c97d63
878 mt19937 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
9f9 template < typename T > struct ETT {
       // treap
3c9
        struct node {
ed1
            node *1, *r, *p;
fa4
            int pr, sz;
875
            T val, sub, lazy;
53e
           int id:
ffd
            bool f; // se eh o 'first'
5ef
            int qt_f; // numero de firsts na subarvore
7a8
            node(int id_, T v, bool f_ = 0) : l(NULL), r(NULL),
   p(NULL), pr(rng()),
                sz(1), val(v), sub(v), lazy(), id(id_{-}), f(f_{-}).
   qt_f(f_) {}
a9c
            void prop() {
d09
                if (lazy != T()) {
021
                    if (f) val += lazy;
971
                    sub += lazv*sz:
b87
                    if (1) 1->lazy += lazy;
d3b
                    if (r) r->lazy += lazy;
30 c
bfd
                lazy = T();
0 f c
01e
            void update() {
8da
                sz = 1, sub = val, qt_f = f;
```

```
171
                if (1) 1->prop(), sz += 1->sz, sub += 1->sub, qt_f +=
   1->qt_f;
                if (r) r \rightarrow prop(), sz += r \rightarrow sz, sub += r \rightarrow sub, qt_f +=
117
   r->qt_f;
           }
ccb
bff
        };
bb7
        node* root;
73c
        int size(node* x) { return x ? x->sz : 0; }
bcf
        void join(node* 1, node* r, node*& i) { // assume que 1 < r</pre>
986
            if (!l or !r) return void(i = 1 ? 1 : r);
161
            1->prop(), r->prop();
            if (1->pr > r->pr) join(1->r, r, 1->r), 1->r->p = i = 1;
ff5
            else join(1, r->1, r->1), r->1->p = i = r;
982
bda
            i->update();
84d
        }
        void split(node* i, node*& 1, node*& r, int v, int key = 0) {
a20
26a
            if (!i) return void(r = 1 = NULL);
c89
            i->prop():
d9e
            if (key + size(i->1) < v) {
448
                split(i->r, i->r, r, v, key+size(i->l)+1), l = i;
a21
                if (r) r - p = NULL;
6e8
                if (i->r) i->r->p = i;
949
984
                split(i->1, 1, i->1, v, key), r = i;
                if (1) 1->p = NULL;
5a3
899
                if (i->1) i->1->p = i;
a3f
bda
            i->update();
        }
134
ac7
        int get_idx(node* i) {
            int ret = size(i->1);
6cf
482
            for (: i->p: i = i->p) {
fbf
                node* pai = i->p;
8a6
                if (i != pai->1) ret += size(pai->1) + 1;
e22
            }
edf
            return ret;
479
048
        node* get_min(node* i) {
433
            if (!i) return NULL:
            return i->1 ? get_min(i->1) : i;
f8e
0de
f03
        node* get_max(node* i) {
433
            if (!i) return NULL;
424
            return i->r ? get_max(i->r) : i;
e92
        }
```

```
// fim da treap
        vector < node *> first, last;
4fb
        ETT(int n, vector<T> v = {}) : root(NULL), first(n), last(n) {
f82
            if (!v.size()) v = vector < T > (n);
с5е
603
            for (int i = 0: i < n: i++) {
                first[i] = last[i] = new node(i, v[i], 1);
a00
469
                join(root, first[i], root);
            }
8ac
ec3
83f
        ETT(const ETT& t) { throw logic_error("Nao copiar a ETT!"); }
c09
       \simETT() {
609
            vector < node *> q = {root};
402
            while (q.size()) {
e5d
                node* x = q.back(); q.pop_back();
ee9
                if (!x) continue;
                q.push_back(x->1), q.push_back(x->r);
1c7
bf0
                delete x;
653
            }
672
       }
153
        pair < int , int > get_range(int i) {
670
            return {get_idx(first[i]), get_idx(last[i])};
ada
7af
        void link(int v, int u) { // 'v' tem que ser raiz
            auto [lv. rv] = get_range(v);
890
f13
            int ru = get_idx(last[u]);
4b4
            node* V;
df9
            node *L, *M, *R;
117
            split(root, M, R, rv+1), split(M, L, M, lv);
f1e
            V = M:
a28
            ioin(L. R. root);
e66
            split(root, L, R, ru+1);
367
            join(L, V, L);
7e8
            join(L, last[u] = new node(u, T() /* elemento neutro */),
   L);
a28
            join(L, R, root);
8d9
       }
        void cut(int v) {
4e6
892
            auto [1, r] = get_range(v);
df9
            node *L, *M, *R;
            split(root, M, R, r+1), split(M, L, M, 1);
dca
            node *LL = get_max(L), *RR = get_min(R);
de6
```

```
710
            if (LL and RR and LL->id == RR->id) { // remove duplicata
                  if (last[RR->id] == RR) last[RR->id] = LL;
e8b
992
                  node *A, *B;
6b3
                  split(R, A, B, 1);
                  delete A:
10c
9d5
                  R = B:
7c0
            }
a28
            join(L, R, root);
a0d
            join(root, M, root);
        }
6ff
808
        T query(int v) {
892
            auto [1, r] = get_range(v);
df9
            node *L. *M. *R:
            split(root, M, R, r+1), split(M, L, M, 1);
dca
d43
            T ans = M->sub;
69d
            join(L, M, M), join(M, R, root);
ba7
            return ans;
ede
93b
        void update(int v, T val) { // soma val em todo mundo da
   subarvore
892
            auto [1, r] = get_range(v);
df9
            node *L, *M, *R;
            split(root, M, R, r+1), split(M, L, M, 1);
dca
409
            M->lazy += val;
69d
            join(L, M, M), join(M, R, root);
61c
129
        void update_v(int v, T val) { // muda o valor de v pra val
ac1
            int l = get_idx(first[v]);
df9
            node *L, *M, *R;
d0c
            split(root, M, R, l+1), split(M, L, M, 1);
            M \rightarrow val = M \rightarrow sub = val;
25 e
69d
            join(L, M, M), join(M, R, root);
630
        }
934
        bool is in subtree(int v. int u) { // se u ta na subtree de v
            auto [lv, rv] = get_range(v);
890
6ec
            auto [lu, ru] = get_range(u);
732
            return lv <= lu and ru <= rv;</pre>
a21
        }
355
        void print(node* i) {
            if (!i) return;
eae
a1e
            print(i->1);
743
            cout << i->id+1 << " ";
f15
            print(i->r);
59f
065
        void print() { print(root); cout << endl; }</pre>
045 }:
```

2.15 Floyd-Warshall

```
// encontra o menor caminho entre todo
// par de vertices e detecta ciclo negativo
// returna 1 sse ha ciclo negativo
// d[i][i] deve ser 0
// para i != j, d[i][j] deve ser w se ha uma aresta
// (i, j) de peso w, INF caso contrario
// O(n^3)
// ea05be
1a8 int n:
ae5 int d[MAX][MAX];
73c bool floyd_warshall() {
        for (int k = 0; k < n; k++)
e22
830
        for (int i = 0; i < n; i++)</pre>
f90
        for (int j = 0; j < n; j++)
0ab
            d[i][j] = min(d[i][j], d[i][k] + d[k][j]);
830
       for (int i = 0; i < n; i++)
753
            if (d[i][i] < 0) return 1;</pre>
bb3
        return 0:
192 }
```

2.16 Functional Graph

```
// rt[i] fala o ID da raiz associada ao vertice i
// d[i] fala a profundidade (0 sse ta no ciclo)
// pos[i] fala a posicao de i no array que eh a concat. dos ciclos
// build(f, val) recebe a funcao f e o custo de ir de
// i para f[i] (por default, val = f)
// f_k(i, k) fala onde i vai parar se seguir k arestas
// path(i, k) fala o custo (soma) seguir k arestas a partir de i
// Se quiser outra operacao, da pra alterar facil o codigo
// Codigo um pouco louco, tenho que admitir
//
// build - O(n)
// f_k - O(log(min(n, k)))
// path - O(\log(\min(n, k)))
// 51fabe
6ef namespace func_graph {
1a8
        int n;
ce2
        int f[MAX], vis[MAX], d[MAX];
```

```
f82
        int p[MAX], pp[MAX], rt[MAX], pos[MAX];
ebd
        int sz[MAX], comp;
6a9
        vector < vector < int >> ciclo;
405
        11 val[MAX], jmp[MAX], seg[2*MAX];
97 c
        11 op(11 a, 11 b) { return a+b; }; // mudar a operacao aqui
27b
        void dfs(int i, int t = 2) {
9c9
            vis[i] = t;
f09
            if (vis[f[i]] \ge 2) \{ // comeca ciclo - f[i] eh o rep.
e0a
                d[i] = 0, rt[i] = comp;
74c
                sz[comp] = t - vis[f[i]] + 1;
97ъ
                p[i] = pp[i] = i, jmp[i] = val[i];
15c
                ciclo.emplace_back();
bfb
                ciclo.back().push_back(i);
9d9
            } else {
c16
                if (!vis[f[i]]) dfs(f[i], t+1);
8c0
                rt[i] = rt[f[i]];
195
                if (sz[comp]+1) { // to no ciclo
d0f
                    d[i] = 0;
97b
                    p[i] = pp[i] = i, jmp[i] = val[i];
                    ciclo.back().push_back(i);
bfb
9d9
                } else { // nao to no ciclo
                    d[i] = d[f[i]]+1, p[i] = f[i];
00d
511
                    pp[i] = 2*d[pp[f[i]]] == d[pp[pp[f[i]]]]+d[f[i]]?
   pp[pp[f[i]]] : f[i];
114
                    jmp[i] = pp[i] == f[i] ? val[i] : op(val[i],
   op(jmp[f[i]], jmp[pp[f[i]]]));
bb3
                }
80b
            }
            if (f[ciclo[rt[i]][0]] == i) comp++; // fim do ciclo
e4a
29a
            vis[i] = 1;
0ba
1da
        void build(vector<int> f_, vector<int> val_ = {}) {
bcb
            n = f .size(). comp = 0:
527
            if (!val_.size()) val_ = f_;
830
            for (int i = 0; i < n; i++)</pre>
998
                f[i] = f_[i], val[i] = val_[i], vis[i] = 0, sz[i] = -1;
e74
            ciclo.clear():
            for (int i = 0; i < n; i++) if (!vis[i]) dfs(i);</pre>
158
6bb
            int t = 0:
daa
            for (auto& c : ciclo) {
336
                reverse(c.begin(), c.end());
ea5
                for (int j : c) {
85b
                    pos[i] = t;
948
                    seg[n+t] = val[j];
c82
                    t++:
```

```
25e
cbc
            for (int i = n-1; i; i--) seg[i] = op(seg[2*i],
dc1
   seg[2*i+1]);
90b
       }
283
        int f k(int i, ll k) {
1b1
            while (d[i] and k) {
77b
                int big = d[i] - d[pp[i]];
ded
                if (big <= k) k -= big, i = pp[i];</pre>
584
                else k--, i = p[i];
09c
            }
77e
            if (!k) return i:
            return ciclo[rt[i]][(pos[i] - pos[ciclo[rt[i]][0]] + k) %
   sz[rt[i]]];
f34
       }
047
        ll path(int i, ll k) {
            auto query = [&](int 1, int r) {
3cf
3 e 4
                11 q = 0;
                for (1 += n, r += n; 1 <= r; ++1/=2, --r/=2) {
47a
                    if (1\%2 == 1) q = op(q, seg[1]);
27e
                    if (r\%2 == 0) q = op(q, seg[r]);
1f2
                }
598
bef
                return q;
6e1
            };
b73
            ll ret = 0:
1b1
            while (d[i] and k) {
77b
                int big = d[i] - d[pp[i]];
                if (big <= k) k -= big, ret = op(ret, jmp[i]), i =</pre>
327
   pp[i];
                else k--, ret = op(ret, val[i]), i = p[i];
f9e
7e3
            if (!k) return ret:
            int first = pos[ciclo[rt[i]][0]], last =
   pos[ciclo[rt[i]].back()];
            // k/sz[rt[i]] voltas completas
            if (k/sz[rt[i]]) ret = op(ret, k/sz[rt[i]] * query(first,
430
   last)):
            k %= sz[rt[i]]:
9af
еЗс
            if (!k) return ret;
            int l = pos[i], r = first + (pos[i] - first + k - 1) %
   sz[rt[i]]:
            if (1 <= r) return op(ret, query(1, r));</pre>
982
687
            return op(ret, op(query(l, last), query(first, r)));
380
        }
```

51f }

2.17 Heavy-Light Decomposition - aresta

```
// SegTree de soma
// query / update de soma das arestas
//
// Complexidades:
// build - O(n)
// query_path - 0(log^2 (n))
// update_path - O(log^2 (n))
// query_subtree - O(log(n))
// update_subtree - O(log(n))
// namespace seg { ... }
// 599946
826 namespace hld {
c0d
        vector < pair < int , int > > g[MAX];
e65
        int pos[MAX], sz[MAX];
7c0
        int sobe[MAX], pai[MAX];
096
        int h[MAX], v[MAX], t;
        void build_hld(int k, int p = -1, int f = 1) {
0ce
180
            v[pos[k] = t++] = sobe[k]; sz[k] = 1;
            for (auto& i : g[k]) if (i.first != p) {
418
dd2
                 auto [u. w] = i:
                 sobe[u] = w; pai[u] = k;
a76
                h[u] = (i == g[k][0] ? h[k] : u);
0 c 1
                 build_hld(u, k, f); sz[k] += sz[u];
da7
865
                 if (sz[u] > sz[g[k][0].first] or g[k][0].first == p)
9a3
                     swap(i, g[k][0]);
804
667
             if (p*f == -1) build_hld(h[k] = k, -1, t = 0);
4dd
1f8
        void build(int root = 0) {
a34
            t = 0:
295
            build_hld(root);
c83
            seg::build(t, v);
ea2
3fc
        11 query_path(int a, int b) {
            if (a == b) return 0:
2d5
            if (pos[a] < pos[b]) swap(a, b);
aa1
29b
            if (h[a] == h[b]) return seg::query(pos[b]+1, pos[a]);
fca
            return seg::query(pos[h[a]], pos[a]) +
```

```
query_path(pai[h[a]], b);
87f
920
        void update_path(int a, int b, int x) {
            if (a == b) return;
d54
            if (pos[a] < pos[b]) swap(a, b);</pre>
aa1
881
            if (h[a] == h[b]) return (void)seg::update(pos[b]+1,
   pos[a], x);
            seg::update(pos[h[a]], pos[a], x); update_path(pai[h[a]],
701
   b, x);
dbf
        11 query_subtree(int a) {
d0a
b9f
            if (sz[a] == 1) return 0:
2f6
            return seg::query(pos[a]+1, pos[a]+sz[a]-1);
77f
        void update_subtree(int a, int x) {
acc
a5a
            if (sz[a] == 1) return;
9cd
            seg::update(pos[a]+1, pos[a]+sz[a]-1, x);
a46
7be
        int lca(int a, int b) {
            if (pos[a] < pos[b]) swap(a, b);</pre>
aa1
            return h[a] == h[b] ? b : lca(pai[h[a]], b);
ca5
        }
219
599 }
```

Heavy-Light Decomposition - vertice

```
// SegTree de soma
// query / update de soma dos vertices
// Complexidades:
// build - O(n)
// query_path - 0(log^2 (n))
// update_path - O(log^2 (n))
// query_subtree - O(log(n))
// update_subtree - O(log(n))
// namespace seg { ... }
// de3d84
826 namespace hld {
042
        vector < int > g[MAX];
e65
        int pos[MAX], sz[MAX];
bd4
        int peso[MAX], pai[MAX];
096
        int h[MAX], v[MAX], t;
0ce
        void build_hld(int k, int p = -1, int f = 1) {
```

```
b18
            v[pos[k] = t++] = peso[k]; sz[k] = 1;
            for (auto& i : g[k]) if (i != p) {
b94
78d
                pai[i] = k;
                h[i] = (i == g[k][0] ? h[k] : i);
26e
                build_hld(i, k, f); sz[k] += sz[i];
193
cd1
                if (sz[i] > sz[g[k][0]] or g[k][0] == p) swap(i,
   g[k][0]);
d94
667
            if (p*f == -1) build_hld(h[k] = k, -1, t = 0);
a83
        }
1f8
        void build(int root = 0) {
a34
            t = 0:
295
            build_hld(root);
c83
            seg::build(t, v);
ea2
        }
3fc
        11 query_path(int a, int b) {
            if (pos[a] < pos[b]) swap(a, b);</pre>
aa1
4bf
            if (h[a] == h[b]) return seg::query(pos[b], pos[a]);
fca
            return seg::query(pos[h[a]], pos[a]) +
   query_path(pai[h[a]], b);
        }
c17
        void update_path(int a, int b, int x) {
920
            if (pos[a] < pos[b]) swap(a, b);</pre>
aa1
198
            if (h[a] == h[b]) return (void)seg::update(pos[b], pos[a],
   x):
701
            seg::update(pos[h[a]], pos[a], x); update_path(pai[h[a]],
   b, x);
421
        }
d0a
        11 query_subtree(int a) {
            return seg::query(pos[a], pos[a]+sz[a]-1);
b3e
ba2
        }
        void update_subtree(int a, int x) {
acc
a22
            seg::update(pos[a], pos[a]+sz[a]-1, x);
        }
480
7be
        int lca(int a, int b) {
            if (pos[a] < pos[b]) swap(a, b);
aa1
ca5
            return h[a] == h[b] ? b : lca(pai[h[a]], b);
        }
219
de3 }
      Heavy-Light Decomposition sem Update
// query de min do caminho
```

```
// Complexidades:
// build - O(n)
// query_path - O(log(n))
// ee6991
826 namespace hld {
c0d
        vector < pair < int , int > > g[MAX];
        int pos[MAX], sz[MAX];
e65
7c0
        int sobe[MAX], pai[MAX];
096
        int h[MAX], v[MAX], t;
        int men[MAX], seg[2*MAX];
ea2
        void build hld(int k, int p = -1, int f = 1) {
0ce
            v[pos[k] = t++] = sobe[k]; sz[k] = 1;
180
418
            for (auto& i : g[k]) if (i.first != p) {
1f5
                sobe[i.first] = i.second; pai[i.first] = k;
6fa
                h[i.first] = (i == g[k][0] ? h[k] : i.first);
                men[i.first] = (i == g[k][0] ? min(men[k], i.second) :
87b
   i.second);
                build hld(i.first, k, f): sz[k] += sz[i.first]:
4b2
                if (sz[i.first] > sz[g[k][0].first] or g[k][0].first
bc3
                     swap(i, g[k][0]);
9a3
            }
667
            if (p*f == -1) build hld(h\lceil k \rceil = k, -1, t = 0):
8ec
1f8
        void build(int root = 0) {
            t = 0:
a34
295
            build_hld(root);
            for (int i = 0; i < t; i++) seg[i+t] = v[i];</pre>
3ae
            for (int i = t-1; i; i--) seg[i] = min(seg[2*i],
   seg[2*i+1]);
       }
ea5
        int query_path(int a, int b) {
f04
490
            if (a == b) return INF;
            if (pos[a] < pos[b]) swap(a, b);</pre>
aa1
98f
            if (h[a] != h[b]) return min(men[a], query_path(pai[h[a]],
   b));
            int ans = INF, x = pos[b]+1+t, y = pos[a]+t;
46b
            for (; x \le y; ++x/=2, --y/=2) ans = min({ans, seg[x],
646
   seg[y]});
ba7
            return ans;
3a9
ee6 };
```

2.20 Isomorfismo de arvores

```
// thash() retorna o hash da arvore (usando centroids como vertices
   especiais).
// Duas arvores sao isomorfas sse seu hash eh o mesmo
// O(|V|.log(|V|))
// 8fb6bb
91f map < vector < int >, int > mphash;
df6 struct tree {
1a8
        int n:
789
        vector < vector < int >> g;
347
        vector<int> sz, cs;
        tree(int n_{-}) : n(n_{-}), g(n_{-}), sz(n_{-}) {}
1b5
76b
        void dfs_centroid(int v, int p) {
588
            sz[v] = 1;
fa7
            bool cent = true;
18e
            for (int u : g[v]) if (u != p) {
365
                 dfs_centroid(u, v), sz[v] += sz[u];
                if(sz[u] > n/2) cent = false;
e90
ece
1f6
            if (cent and n - sz[v] \le n/2) cs.push_back(v);
368
784
        int fhash(int v, int p) {
544
            vector < int > h;
332
            for (int u : g[v]) if (u != p) h.push_back(fhash(u, v));
1c9
            sort(h.begin(), h.end());
            if (!mphash.count(h)) mphash[h] = mphash.size();
3ac
bbc
            return mphash[h];
748
        }
38f
        11 thash() {
23a
            cs.clear():
3a5
            dfs_centroid(0, -1);
16d
            if (cs.size() == 1) return fhash(cs[0], -1);
772
            ll h1 = fhash(cs[0], cs[1]), h2 = fhash(cs[1], cs[0]);
fae
            return (min(h1, h2) << 30) + max(h1, h2);
138
        }
4dd };
```

2.21 Kosaraju

```
// O(n + m)
// a4f310
```

```
1a8 int n:
042 vector < int > g[MAX];
58d vector <int> gi[MAX]; // grafo invertido
c5a int vis[MAX]:
ee6 stack<int> S;
a52 int comp[MAX]; // componente conexo de cada vertice
1ca void dfs(int k) {
59a
       vis[k] = 1:
54f
       for (int i = 0; i < (int) g[k].size(); i++)</pre>
            if (!vis[g[k][i]]) dfs(g[k][i]);
58f
       S.push(k);
89c }
436 void scc(int k, int c) {
       vis[k] = 1:
59a
       comp[k] = c;
52c
ff0 for (int i = 0; i < (int) gi[k].size(); i++)
bf6
            if (!vis[gi[k][i]]) scc(gi[k][i], c);
088 }
db8 void kosaraju() {
       for (int i = 0; i < n; i++) vis[i] = 0;</pre>
158
       for (int i = 0; i < n; i++) if (!vis[i]) dfs(i);</pre>
991
       for (int i = 0; i < n; i++) vis[i] = 0;
d32
       while (S.size()) {
70b
           int u = S.top();
7de
           S.pop();
f43
            if (!vis[u]) scc(u, u);
       }
207
e21 }
2.22 Kruskal
// Gera e retorna uma AGM e seu custo total a partir do vetor de
   arestas (edg)
// do grafo
// O(m log(m) + m a(m))
// 864875
1b9 vector<tuple<int, int, int>> edg; // {peso,[x,y]}
```

// DSU em O(a(n))

```
4a6 void dsu_build();
d78 int find(int a);
369 void unite(int a, int b);
c67 pair<11, vector<tuple<int, int, int>>> kruskal(int n) {
        dsu_build(n);
8d2
e31
        sort(edg.begin(), edg.end());
854
        11 cost = 0;
979
        vector < tuple < int , int , int >> mst;
fea
        for (auto [w,x,y] : edg) if (find(x) != find(y)) {
9de
            mst.emplace_back(w, x, y);
45f
            cost += w:
05a
            unite(x,y);
ca2
5df
        return {cost, mst};
b6a }
2.23 Kuhn
// Computa matching maximo em grafo bipartido
// 'n' e 'm' sao quantos vertices tem em cada particao
// chamar add(i, j) para add aresta entre o cara i
// da particao A, e o cara j da particao B
// (entao i < n, j < m)
// Para recuperar o matching, basta olhar 'ma' e 'mb'
// 'recover' recupera o min vertex cover como um par de
// {caras da particao A, caras da particao B}
//
// O(|V| * |E|)
// Na pratica, parece rodar tao rapido quanto o Dinic
878 mt19937 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
// b0dda3
6c6 struct kuhn {
14e
       int n, m;
789
        vector < vector < int >> g;
        vector < int > vis, ma, mb;
d3f
40e
        kuhn(int n_, int m_) : n(n_), m(m_), g(n),
            vis(n+m), ma(n, -1), mb(m, -1) {}
        void add(int a, int b) { g[a].push_back(b); }
ba6
caf
        bool dfs(int i) {
```

```
29a
            vis[i] = 1;
            for (int j : g[i]) if (!vis[n+j]) {
29b
8c9
                vis[n+j] = 1;
                if (mb[i] == -1 or dfs(mb[i])) {
2cf
                    ma[i] = j, mb[j] = i;
bfe
8a6
                    return true;
                }
b17
            }
82a
d1f
            return false;
4ef
bf7
        int matching() {
1ae
            int ret = 0, aum = 1;
            for (auto& i : g) shuffle(i.begin(), i.end(), rng);
5a8
392
            while (aum) {
618
                for (int j = 0; j < m; j++) vis[n+j] = 0;
c5d
830
               for (int i = 0; i < n; i++)
                    if (ma[i] == -1 and dfs(i)) ret++, aum = 1;
01f
            }
085
edf
            return ret:
2ee
        }
b0d };
// 55fb67
ebf pair < vector < int > , vector < int >> recover (kuhn & K) {
       K.matching():
50c
       int n = K.n, m = K.m;
       for (int i = 0; i < n+m; i++) K.vis[i] = 0;
9d0
       for (int i = 0; i < n; i++) if (K.ma[i] == -1) K.dfs(i);
8ad
       vector < int > ca, cb;
576
       for (int i = 0; i < n; i++) if (!K.vis[i]) ca.push_back(i);</pre>
f24
       for (int i = 0; i < m; i++) if (K.vis[n+i]) cb.push_back(i);</pre>
        return {ca. cb}:
aad
55f }
2.24 LCA com binary lifting
// Assume que um vertice eh ancestral dele mesmo, ou seja,
// se a eh ancestral de b, lca(a, b) = a
// MAX2 = ceil(log(MAX))
// Complexidades:
// build - O(n log(n))
// lca - O(log(n))
// b674ca
677 vector < vector < int > > g(MAX);
```

```
41c int n, p;
e75 int pai[MAX2][MAX];
999 int in [MAX], out [MAX];
1ca void dfs(int k) {
fdf
        in[k] = p++;
54f
        for (int i = 0; i < (int) g[k].size(); i++)</pre>
            if (in[g[k][i]] == -1) {
9b7
ba6
                pai[0][g[k][i]] = k;
c38
                dfs(g[k][i]);
e2d
            }
26f
        out[k] = p++;
691 }
c11 void build(int raiz) {
        for (int i = 0; i < n; i++) pai[0][i] = i;
c63
        p = 0, memset(in, -1, sizeof in);
ecb
        dfs(raiz):
        // pd dos pais
511
        for (int k = 1; k < MAX2; k++) for (int i = 0; i < n; i++)
d38
            pai[k][i] = pai[k - 1][pai[k - 1][i]];
530 }
00f bool anc(int a, int b) { // se a eh ancestral de b
        return in[a] <= in[b] and out[a] >= out[b];
2d6 }
7be int lca(int a, int b) {
86d
        if (anc(a, b)) return a;
e52
        if (anc(b, a)) return b;
        // sobe a
f70
        for (int k = MAX2 - 1: k \ge 0: k - -)
            if (!anc(pai[k][a], b)) a = pai[k][a];
acf
        return pai[0][a];
847
5c4 }
// Alternativamente:
// 'binary lifting' gastando O(n) de memoria
// Da pra add folhas e fazer queries online
// 3 vezes o tempo do binary lifting normal
//
// build - O(n)
// kth, lca, dist - O(log(n))
// 89a97a
```

```
9c6 int d[MAX], p[MAX], pp[MAX];
d40 void set_root(int i) { p[i] = pp[i] = i, d[i] = 0; }
e9d void add_leaf(int i, int u) {
e0b
       p[i] = u. d[i] = d[u]+1:
        pp[i] = 2*d[pp[u]] == d[pp[pp[u]]]+d[u] ? pp[pp[u]] : u;
b15
33f }
c37 int kth(int i, int k) {
        int dd = max(0, d[i]-k);
935
        while (d[i] > dd) i = d[pp[i]] >= dd? pp[i] : p[i];
        return i;
d9a
f3c }
7be int lca(int a, int b) {
       if (d[a] < d[b]) swap(a, b);</pre>
        while (d[a] > d[b]) a = d[pp[a]] >= d[b]? pp[a] : p[a];
6cd
       while (a != b) {
984
            if (pp[a] != pp[b]) a = pp[a], b = pp[b];
932
            else a = p[a], b = p[b];
e7c
       }
4ea
3f5
        return a;
21d }
4fe int dist(int a, int b) { return d[a]+d[b]-2*d[lca(a,b)]; }
042 vector < int > g[MAX];
3ab void build(int i, int pai=-1) {
5cf
       if (pai == -1) set_root(i);
       for (int j : g[i]) if (j != pai) {
d31
            add leaf(i, i):
b21
            build(j, i);
43b
       }
74a }
2.25 LCA com HLD
// Assume que um vertice eh ancestral dele mesmo, ou seja,
// se a eh ancestral de b. lca(a, b) = a
// Para buildar pasta chamar build(root)
// anc(a, b) responde se 'a' eh ancestral de 'b'
//
// Complexidades:
// build - O(n)
```

```
// lca - O(log(n))
// anc - O(1)
// fb22c1
042 vector < int > g[MAX];
713 int pos[MAX], h[MAX], sz[MAX];
ff1 int pai[MAX], t;
8bf void build(int k, int p = -1, int f = 1) {
        pos[k] = t++; sz[k] = 1;
e26
        for (int& i : g[k]) if (i != p) {
78d
            pai[i] = k;
            h[i] = (i == g[k][0] ? h[k] : i);
26e
            build(i, k, f); sz[k] += sz[i];
cb8
            if (sz[i] > sz[g[k][0]] or g[k][0] == p) swap(i, g[k][0]);
cd1
917
        }
        if (p*f == -1) t = 0, h[k] = k, build(k, -1, 0);
3da
1b9 }
7be int lca(int a, int b) {
        if (pos[a] < pos[b]) swap(a, b);</pre>
aa1
        return h[a] == h[b] ? b : lca(pai[h[a]], b);
ca5
219 }
00f bool anc(int a. int b) {
        return pos[a] <= pos[b] and pos[b] <= pos[a]+sz[a]-1;</pre>
272 }
2.26 LCA com RMQ
// Assume que um vertice eh ancestral dele mesmo, ou seja,
// se a eh ancestral de b, lca(a, b) = a
// dist(a, b) retorna a distancia entre a e b
// Complexidades:
// build - O(n)
// lca - 0(1)
// dist - O(1)
// 22cde8 - rmg + lca
// 0214e8
1a5 template < typename T> struct rmq {
517
        vector <T> v;
fcc
        int n; static const int b = 30;
70e
        vector < int > mask, t;
```

```
18e
        int op(int x, int y) { return v[x] < v[y] ? x : y; }
        int msb(int x) { return __builtin_clz(1)-__builtin_clz(x); }
ee1
6ad
        rmq() {}
        rmq(const vector < T > \& v_) : v(v_), n(v.size()), mask(n), t(n) {
43c
2e5
            for (int i = 0, at = 0; i < n; \max\{i++\} = at |= 1) {
                at = (at <<1) &((1 << b) -1);
a61
76a
                while (at and op(i, i-msb(at&-at)) == i) at ^= at&-at;
53c
            }
243
            for (int i = 0; i < n/b; i++) t[i] =</pre>
   b*i+b-1-msb(mask[b*i+b-1]):
            for (int j = 1; (1<<j) <= n/b; j++) for (int i = 0;
39d
   i+(1<< i) <= n/b: i++)
                t[n/b*j+i] = op(t[n/b*(j-1)+i],
   t[n/b*(j-1)+i+(1<<(j-1))]);
2d3
        int small(int r, int sz = b) { return
c92
   r-msb(mask[r]&((1<<sz)-1)); }
        T query(int 1, int r) {
b7a
27b
            if (r-l+1 \le b) return small(r, r-l+1);
7bf
            int ans = op(small(l+b-1), small(r));
            int x = 1/b+1, y = r/b-1;
e80
e25
            if (x <= y) {
a4e
                int j = msb(y-x+1);
002
                ans = op(ans, op(t[n/b*j+x], t[n/b*j+y-(1<<j)+1]));
4b6
            }
ba7
            return ans:
6bf
        }
021 }:
// 645120
065 namespace lca {
042
        vector < int > g[MAX];
        int v[2*MAX], pos[MAX], dep[2*MAX];
8ec
8bd
        int t:
        rmq<int> RMQ;
2de
4cf
        void dfs(int i, int d = 0, int p = -1) {
c97
            v[t] = i, pos[i] = t, dep[t++] = d;
cac
            for (int j : g[i]) if (j != p) {
8ec
                dfs(i, d+1, i);
                v[t] = i, dep[t++] = d;
cf2
843
            }
d6a
        void build(int n, int root) {
789
a34
            t = 0;
            dfs(root):
14e
3f4
            RMQ = rmq < int > (vector < int > (dep, dep + 2*n - 1));
```

```
657
7be
        int lca(int a, int b) {
ab7
            a = pos[a], b = pos[b];
            return v[RMQ.query(min(a, b), max(a, b))];
9c0
5db
b5d
        int dist(int a, int b) {
670
            return dep[pos[a]] + dep[pos[b]] - 2*dep[pos[lca(a, b)]];
5b7
        }
645 }
2.27 Line Tree
// Reduz min-query em arvore para RMQ
// Se o grafo nao for uma arvore, as queries
// sao sobre a arvore geradora maxima
// Queries de minimo
//
// build - O(n log(n))
// query - O(log(n))
// b1f418
1a8 int n;
3ae namespace linetree {
f37
        int id[MAX], seg[2*MAX], pos[MAX];
        vector < int > v[MAX], val[MAX];
43f
430
        vector<pair<int, pair<int, int> > ar;
dc6
        void add(int a, int b, int p) { ar.push_back({p, {a, b}}); }
        void build() {
0a8
b09
            sort(ar.rbegin(), ar.rend());
0e3
            for (int i = 0; i < n; i++) id[i] = i, v[i] = \{i\},
   val[i].clear();
8bb
            for (auto i : ar) {
                int a = id[i.second.first], b = id[i.second.second];
c.91
f6f
                if (a == b) continue;
                if (v[a].size() < v[b].size()) swap(a, b);</pre>
c58
                for (auto j : v[b]) id[j] = a, v[a].push_back(j);
fb8
482
                val[a].push_back(i.first);
78b
                for (auto j : val[b]) val[a].push_back(j);
e39
                v[b].clear(), val[b].clear();
012
            }
8e8
            vector < int > vv;
            for (int i = 0; i < n; i++) for (int j = 0; j <
2ce
   v[i].size(); j++) {
e52
                pos[v[i][j]] = vv.size();
941
                if (j + 1 < v[i].size()) vv.push_back(val[i][j]);</pre>
```

```
1cb
                else vv.push_back(0);
            }
475
bb4
            for (int i = n; i < 2*n; i++) seg[i] = vv[i-n];
            for (int i = n-1; i; i--) seg[i] = min(seg[2*i],
   seg[2*i+1]);
9fe
4ea
        int query(int a, int b) {
            if (id[a] != id[b]) return 0; // nao estao conectados
596
ab7
            a = pos[a], b = pos[b];
            if (a > b) swap(a, b);
d11
199
            b--;
38a
            int ans = INF:
513
            for (a += n, b += n; a <= b; ++a/=2, --b/=2) ans =
   min({ans, seg[a], seg[b]});
ba7
            return ans;
952
       }
00f };
2.28 Link-cut Tree
// Link-cut tree padrao
//
// Todas as operacoes sao O(log(n)) amortizado
// e4e663
1ef namespace lct {
3c9
        struct node {
            int p, ch[2];
19f
            node() \{ p = ch[0] = ch[1] = -1; \}
062
f43
       };
5f3
        node t[MAX];
971
        bool is_root(int x) {
657
            return t[x].p == -1 or (t[t[x].p].ch[0] != x and
   t[t[x].p].ch[1] != x);
       }
cf1
```

void rotate(int x) {

void splay(int x) {

int p = t[x].p, pp = t[p].p;

bool d = t[p].ch[0] == x;

t[x].p = pp, t[p].p = x;

while (!is_root(x)) {

if (!is_root(p)) t[pp].ch[t[pp].ch[1] == p] = x;

t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;

if (t[p].ch[!d]+1) t[t[p].ch[!d]].p = p;

ed6

497

fc4

251

461

a76

8fa

49b

07c

18c

```
497
                int p = t[x].p, pp = t[p].p;
0c5
                if (!is_root(p)) rotate((t[pp].ch[0] == p)^(t[p].ch[0]
   == x) ? x : p);
64f
                rotate(x);
            }
d8d
4fa
        }
f16
        int access(int v) {
0eb
            int last = -1;
01a
            for (int w = v; w+1; last = w, splay(v), w = t[v].p)
                splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
024
3d3
            return last;
0a4
        }
e89
        int find root(int v) {
5e3
            access(v):
3de
            while (t[v].ch[0]+1) v = t[v].ch[0];
f05
            return splay(v), v;
ee7
        }
        void link(int v, int w) { // v deve ser raiz
142
5e3
            access(v);
10d
            t[v].p = w;
c56
        void cut(int v) { // remove aresta de v pro pai
4e6
5e3
            access(v):
264
            t[v].ch[0] = t[t[v].ch[0]].p = -1;
5f5
        }
bbb
        int lca(int v. int w) {
948
            return access(v), access(w);
b6d
        }
e4e }
      Link-cut Tree - aresta
// Valores nas arestas
// rootify(v) torna v a raiz de sua arvore
// query(v, w) retorna a soma do caminho v--w
// update(v, w, x) soma x nas arestas do caminho v--w
// Todas as operacoes sao O(log(n)) amortizado
// 9ce48f
1ef namespace lct {
3c9
        struct node {
19f
            int p, ch[2];
810
            ll val, sub;
aa6
            bool rev:
04a
            int sz, ar;
4 e 4
            ll lazy;
```

```
f93
            node() {}
7a8
            node(int v, int ar_) :
546
            p(-1), val(v), sub(v), rev(0), sz(ar_{-}), ar(ar_{-}), lazy(0) {
                ch[0] = ch[1] = -1;
b07
53b
            }
6e0
        };
        node t[2*MAX]; // MAXN + MAXQ
c53
99e
        map<pair<int, int>, int> aresta;
e4d
        int sz:
95a
        void prop(int x) {
dc1
            if (t[x].lazv) {
                if (t[x].ar) t[x].val += t[x].lazy;
25e
                t[x].sub += t[x].lazy*t[x].sz;
2ab
edc
                if (t[x].ch[0]+1) t[t[x].ch[0]].lazy += t[x].lazy;
942
                if (t[x].ch[1]+1) t[t[x].ch[1]].lazy += t[x].lazy;
            }
1ba
            if (t[x].rev) {
aa2
                swap(t[x].ch[0], t[x].ch[1]);
f95
379
                if (t[x].ch[0]+1) t[t[x].ch[0]].rev ^= 1;
c3d
                if (t[x].ch[1]+1) t[t[x].ch[1]].rev ^= 1;
            }
50e
230
            t[x].lazy = 0, t[x].rev = 0;
f9d
        }
564
        void update(int x) {
            t[x].sz = t[x].ar, t[x].sub = t[x].val;
1a3
            for (int i = 0; i < 2; i++) if (t[x].ch[i]+1) {
8ca
621
                prop(t[x].ch[i]);
c4f
                t[x].sz += t[t[x].ch[i]].sz;
269
                t[x].sub += t[t[x].ch[i]].sub;
400
            }
28Ъ
        }
971
        bool is root(int x) {
            return t[x].p == -1 or (t[t[x].p].ch[0] != x and
   t[t[x].p].ch[1] != x);
       }
cf1
ed6
        void rotate(int x) {
497
            int p = t[x].p, pp = t[p].p;
fc4
            if (!is_root(p)) t[pp].ch[t[pp].ch[1] == p] = x;
            bool d = t[p].ch[0] == x;
251
461
            t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
a76
            if (t[p].ch[!d]+1) t[t[p].ch[!d]].p = p;
            t[x].p = pp, t[p].p = x;
8fa
444
            update(p), update(x);
f31
238
        int splay(int x) {
```

```
18c
            while (!is_root(x)) {
497
                int p = t[x].p, pp = t[p].p;
77b
                if (!is_root(p)) prop(pp);
be5
                prop(p), prop(x);
                if (!is\_root(p)) rotate((t[pp].ch[0] == p)^(t[p].ch[0]
0 c 5
   == x) ? x : p);
64f
                rotate(x);
72c
            }
            return prop(x), x;
aab
        }
08f
f16
        int access(int v) {
0eb
            int last = -1:
d9f
            for (int w = v; w+1; update(last = w), splay(v), w =
   t[v].p)
024
                splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
3d3
            return last;
294
        }
        void make_tree(int v, int w=0, int ar=0) { t[v] = node(w, ar);
9f1
}
e89
        int find root(int v) {
13f
            access(v), prop(v);
9f0
            while (t[v].ch[0]+1) v = t[v].ch[0], prop(v);
637
            return splay(v);
16a
        }
82f
        bool conn(int v, int w) {
2cf
            access(v). access(w):
b9b
            return v == w ? true : t[v].p != -1;
ec0
        }
        void rootify(int v) {
277
5e3
            access(v);
            t[v].rev ^= 1;
a02
a05
        }
971
        11 query(int v, int w) {
            rootifv(w). access(v):
b54
249
            return t[v].sub;
652
        }
3fa
        void update(int v, int w, int x) {
b54
            rootify(w), access(v);
12c
            t[v].lazy += x;
74f
        }
        void link_(int v, int w) {
204
821
            rootify(w);
389
            t[w].p = v;
523
6b8
        void link(int v, int w, int x) { // v--w com peso x
379
            int id = MAX + sz++;
110
            aresta[make_pair(v, w)] = id;
```

```
a88
            make_tree(id, x, 1);
c88
            link_(v, id), link_(id, w);
58c
       }
        void cut_(int v, int w) {
e63
b54
            rootify(w), access(v);
            t[v].ch[0] = t[t[v].ch[0]].p = -1;
264
7cd
       }
031
        void cut(int v, int w) {
b0f
            int id = aresta[make_pair(v, w)];
            cut_(v, id), cut_(id, w);
a4a
840
       }
bbb
        int lca(int v, int w) {
5e3
            access(v):
a8b
            return access(w);
       }
524
9ce }
```

2.30 Link-cut Tree - vertice

```
// Valores nos vertices
// make tree(v. w) cria uma nova arvore com um
// vertice soh com valor 'w'
// rootify(v) torna v a raiz de sua arvore
// query(v, w) retorna a soma do caminho v--w
// update(v, w, x) soma x nos vertices do caminho v--w
//
// Todas as operacoes sao O(log(n)) amortizado
// f9f489
1ef namespace lct {
3c9
        struct node {
19f
            int p, ch[2];
810
            ll val, sub;
aa6
            bool rev;
e4d
            int sz;
4 e 4
            ll lazy;
f93
            node() {}
            node(int v) : p(-1), val(v), sub(v), rev(0), sz(1),
aa0
   lazy(0) {
                ch[0] = ch[1] = -1:
b07
            }
c4e
2b7
        };
        node t[MAX];
5f3
95a
        void prop(int x) {
dc1
            if (t[x].lazy) {
```

```
9f7
                t[x].val += t[x].lazy, t[x].sub += t[x].lazy*t[x].sz;
                if (t[x].ch[0]+1) t[t[x].ch[0]].lazy += t[x].lazy;
edc
942
                if (t[x].ch[1]+1) t[t[x].ch[1]].lazy += t[x].lazy;
e26
            }
            if (t[x].rev) {
aa2
f95
                swap(t[x].ch[0], t[x].ch[1]);
379
                if (t[x].ch[0]+1) t[t[x].ch[0]].rev ^= 1;
c3d
                if (t[x].ch[1]+1) t[t[x].ch[1]].rev ^= 1;
50e
230
            t[x].lazy = 0, t[x].rev = 0;
c62
       }
564
        void update(int x) {
ec2
            t[x].sz = 1, t[x].sub = t[x].val:
            for (int i = 0; i < 2; i++) if (t[x].ch[i]+1) {
8ca
621
                prop(t[x].ch[i]);
                t[x].sz += t[t[x].ch[i]].sz;
c4f
269
                t[x].sub += t[t[x].ch[i]].sub;
            }
400
da7
        }
        bool is root(int x) {
971
            return t[x].p == -1 or (t[t[x].p].ch[0] != x and
657
   t[t[x].p].ch[1] != x);
       }
cf1
        void rotate(int x) {
ed6
497
            int p = t[x].p, pp = t[p].p;
fc4
            if (!is_root(p)) t[pp].ch[t[pp].ch[1] == p] = x;
            bool d = t[p].ch[0] == x;
251
461
            t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
            if (t[p].ch[!d]+1) t[t[p].ch[!d]].p = p;
a76
8fa
            t[x].p = pp, t[p].p = x;
444
            update(p), update(x);
f31
       }
238
        int splay(int x) {
            while (!is root(x)) {
18c
497
                int p = t[x].p, pp = t[p].p;
77b
                if (!is_root(p)) prop(pp);
be5
                prop(p), prop(x);
0 c 5
                if (!is_root(p)) rotate((t[pp].ch[0] == p)^(t[p].ch[0]
   == x) ? x : p);
64f
                rotate(x);
72 c
            }
aab
            return prop(x), x;
08f
f16
       int access(int v) {
0eb
            int last = -1;
d9f
            for (int w = v; w+1; update(last = w), splay(v), w =
   t[v].p)
```

```
splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
024
3d3
            return last;
294
f17
        void make_tree(int v, int w) { t[v] = node(w); }
e89
        int find root(int v) {
13f
            access(v), prop(v);
            while (t[v].ch[0]+1) v = t[v].ch[0], prop(v);
9f0
637
            return splay(v);
16a
f94
        bool connected(int v, int w) {
2cf
            access(v), access(w);
b9b
            return v == w ? true : t[v].p != -1;
ec6
277
        void rootify(int v) {
5e3
            access(v);
a02
            t[v].rev ^= 1;
a05
        11 query(int v, int w) {
971
b54
            rootify(w), access(v);
249
            return t[v].sub;
652
        }
3fa
        void update(int v, int w, int x) {
            rootify(w), access(v);
b54
12c
            t[v].lazv += x;
74f
142
        void link(int v, int w) {
821
            rootify(w);
389
            t[w].p = v:
8a8
031
        void cut(int v, int w) {
b54
            rootify(w), access(v);
264
            t[v].ch[0] = t[t[v].ch[0]].p = -1;
d9a
bbb
        int lca(int v. int w) {
            access(v):
5e3
a8b
            return access(w);
524
        }
f9f }
```

2.31 Max flow com lower bound nas arestas

```
// add(a, b, l, r):
// adiciona aresta de a pra b, onde precisa passar f de fluxo, l <= f
    <= r
// add(a, b, c):
// adiciona aresta de a pra b com capacidade c
//</pre>
```

```
// Mesma complexidade do Dinitz
// e8f30e
cd5 struct lb_max_flow : dinitz {
        vector < int > d:
5ce
d8c
        lb_max_flow(int n) : dinitz(n + 2), d(n, 0) {}
b12
        void add(int a, int b, int 1, int r) {
c97
            d[a] -= 1;
f1b
            d[b] += 1;
            dinitz::add(a, b, r - 1);
4c0
ed4
        }
087
        void add(int a, int b, int c) {
0f3
            dinitz::add(a, b, c):
039
        }
        bool has_circulation() {
7a1
50c
            int n = d.size();
            11 cost = 0:
854
603
            for (int i = 0; i < n; i++) {</pre>
c69
                if (d[i] > 0) {
f56
                    cost += d[i];
57a
                    dinitz::add(n, i, d[i]);
9 c 7
                } else if (d[i] < 0) {</pre>
b76
                     dinitz::add(i, n+1, -d[i]);
dc7
                }
676
            }
067
            return (dinitz::max flow(n, n+1) == cost);
110
7bd
        bool has_flow(int src, int snk) {
387
            dinitz::add(snk, src, INF);
e40
            return has_circulation();
cc1
        }
4eb
        11 max flow(int src, int snk) {
            if (!has_flow(src, snk)) return -1;
ee8
4ad
            dinitz::F = 0;
fe5
            return dinitz::max_flow(src, snk);
619
        }
e8f };
```

2.32 MinCostMaxFlow

```
// min_cost_flow(s, t, f) computa o par (fluxo, custo)
// com max(fluxo) <= f que tenha min(custo)
// min_cost_flow(s, t) -> Fluxo maximo de custo minimo de s pra t
// Se for um dag, da pra substituir o SPFA por uma DP pra nao
// pagar O(nm) no comeco
```

```
// Se nao tiver aresta com custo negativo, nao precisa do SPFA
// O(nm + f * m log n)
// 697b4c
123 template < typename T> struct mcmf {
670
        struct edge {
b75
            int to, rev, flow, cap; // para, id da reversa, fluxo,
   capacidade
            bool res: // se eh reversa
7f9
635
            T cost; // custo da unidade de fluxo
            edge(): to(0), rev(0), flow(0), cap(0), cost(0),
   res(false) {}
1d7
            edge(int to_, int rev_, int flow_, int cap_, T cost_, bool
   res_)
f8d
                : to(to_), rev(rev_), flow(flow_), cap(cap_),
   res(res_), cost(cost_) {}
723
        };
002
        vector < vector < edge >> g;
168
        vector < int > par_idx, par;
f1e
        T inf:
        vector <T> dist;
a03
        mcmf(int n) : g(n), par_idx(n), par(n),
   inf(numeric limits <T>::max()/3) {}
91 c
        void add(int u, int v, int w, T cost) { // de u pra v com cap
            edge a = edge(v, g[v].size(), 0, w, cost, false);
2fc
            edge b = edge(u, g[u].size(), 0, 0, -cost, true);
234
            g[u].push_back(a);
b24
c12
            g[v].push_back(b);
        }
0ed
        vector<T> spfa(int s) { // nao precisa se nao tiver custo
8bc
   negativo
871
            deque < int > q;
3d1
            vector < bool > is_inside(g.size(), 0);
            dist = vector <T>(g.size(), inf);
577
            dist[s] = 0;
a93
a30
            q.push_back(s);
            is_inside[s] = true;
ecb
14d
            while (!q.empty()) {
```

```
b1e
                 int v = q.front();
ced
                 q.pop_front();
                 is_inside[v] = false;
48d
                 for (int i = 0; i < g[v].size(); i++) {</pre>
76e
                     auto [to, rev, flow, cap, res, cost] = g[v][i];
9d4
e61
                     if (flow < cap and dist[v] + cost < dist[to]) {</pre>
                         dist[to] = dist[v] + cost;
943
                         if (is_inside[to]) continue;
ed6
020
                          if (!q.empty() and dist[to] > dist[q.front()])
   q.push_back(to);
b33
                          else q.push_front(to);
                          is_inside[to] = true;
b52
2d1
                     }
                 }
8cd
f2c
            }
8d7
            return dist;
96c
2a2
        bool dijkstra(int s, int t, vector<T>& pot) {
             priority_queue < pair < T, int > , vector < pair < T, int >> ,
489
   greater<>> q;
577
             dist = vector <T>(g.size(), inf);
             dist[s] = 0;
a93
            q.emplace(0, s);
115
402
             while (q.size()) {
91b
                 auto [d, v] = q.top();
833
                 q.pop();
68b
                 if (dist[v] < d) continue;</pre>
76e
                 for (int i = 0; i < g[v].size(); i++) {</pre>
                     auto [to, rev, flow, cap, res, cost] = g[v][i];
9d4
                     cost += pot[v] - pot[to];
e8c
                     if (flow < cap and dist[v] + cost < dist[to]) {</pre>
e61
943
                         dist[to] = dist[v] + cost:
                          q.emplace(dist[to], to);
441
88b
                         par_idx[to] = i, par[to] = v;
873
                     }
de3
                 }
9d4
1d4
            return dist[t] < inf;</pre>
        }
c68
        pair < int , T > min_cost_flow(int s, int t, int flow = INF) {
3d2
3dd
            vector <T> pot(g.size(), 0);
            pot = spfa(s); // mudar algoritmo de caminho minimo aqui
9e4
d22
            int f = 0:
```

```
ce8
            T ret = 0;
4a0
            while (f < flow and dijkstra(s, t, pot)) {</pre>
                 for (int i = 0; i < g.size(); i++)</pre>
bda
                     if (dist[i] < inf) pot[i] += dist[i];</pre>
d2a
71b
                 int mn_flow = flow - f, u = t;
045
                 while (u != s){
                     mn_flow = min(mn_flow,
90f
07d
                         g[par[u]][par_idx[u]].cap -
   g[par[u]][par_idx[u]].flow);
                     u = par[u];
3d1
935
                }
1f2
                ret += pot[t] * mn_flow;
476
                u = t:
045
                 while (u != s) {
e09
                     g[par[u]][par_idx[u]].flow += mn_flow;
d98
                     g[u][g[par[u]][par_idx[u]].rev].flow -= mn_flow;
3d1
                     u = par[u];
                }
04d
                f += mn_flow;
36d
            }
15b
            return make_pair(f, ret);
cc3
        }
        // Opcional: retorna as arestas originais por onde passa flow
182
        vector < pair < int , int >> recover() {
24a
            vector < pair < int , int >> used;
            for (int i = 0; i < g.size(); i++) for (edge e : g[i])</pre>
2a4
587
                 if (e.flow == e.cap && !e.res) used.push back({i.
   e.to});
f6b
            return used;
390
697 };
2.33 Prufer code
// Traduz de lista de arestas para prufer code
// e vice-versa
// Os vertices tem label de O a n-1
// Todo array com n-2 posicoes e valores de
// O a n-1 sao prufer codes validos
//
```

```
// O(n)
// d3b324
47d vector <int > to_prufer (vector <pair <int, int >> tree) {
        int n = tree.size()+1;
1fa
2cf
        vector<int> d(n, 0);
4aa
        vector < vector < int >> g(n);
        for (auto [a, b] : tree) d[a]++, d[b]++,
f87
f60
             g[a].push_back(b), g[b].push_back(a);
        vector < int > pai(n, -1);
c5a
260
        queue < int > q; q.push(n-1);
402
        while (q.size()) {
be1
             int u = q.front(); q.pop();
34 c
             for (int v : g[u]) if (v != pai[u])
9c9
                 pai[v] = u, q.push(v);
        }
70d
399
        int idx, x;
897
        idx = x = find(d.begin(), d.end(), 1) - d.begin();
4b8
        vector < int > ret;
b28
        for (int i = 0; i < n-2; i++) {
d4b
             int y = pai[x];
e81
            ret.push_back(y);
666
             if (-d[y] == 1 \text{ and } y < idx) x = y;
             else idx = x = find(d.begin()+idx+1, d.end(), 1) -
367
   d.begin();
5f9
        }
edf
        return ret;
d3b }
// 765413
4d8 vector < pair < int , int >> from prufer (vector < int > p) {
455
        int n = p.size()+2;
126
        vector < int > d(n, 1);
650
        for (int i : p) d[i]++;
85b
        p.push_back(n-1);
399
        int idx, x;
        idx = x = find(d.begin(), d.end(), 1) - d.begin();
897
1df
        vector < pair < int , int >> ret;
b06
        for (int y : p) {
dab
             ret.push_back({x, y});
666
             if (-d[y] == 1 \text{ and } y < idx) x = y;
             else idx = x = find(d.begin()+idx+1, d.end(), 1) -
367
   d.begin();
c3b
edf
        return ret;
765 }
```

2.34 Sack (DSU em arvores)

```
// Responde queries de todas as sub-arvores
// offline
//
// O(n log(n))
// bb361f
6bf int sz[MAX], cor[MAX], cnt[MAX];
042 vector < int > g[MAX];
6df void build(int k, int d=0) {
        sz[k] = 1:
e8f
01a
       for (auto& i : g[k]) {
            build(i, d+1); sz[k] += sz[i];
30f
925
            if (sz[i] > sz[g[k][0]]) swap(i, g[k][0]);
011
       }
189 }
74f void compute(int k, int x, bool dont=1) {
de9
        cnt[cor[k]] += x;
828
        for (int i = dont; i < g[k].size(); i++)</pre>
b5c
            compute(g[k][i], x, 0);
896 }
dc4 void solve(int k, bool keep=0) {
32a
        for (int i = int(g[k].size())-1; i >= 0; i--)
            solve(g[k][i], !i);
b4c
4a0
        compute(k, 1);
        // agora cnt[i] tem quantas vezes a cor
        // i aparece na sub-arvore do k
830
        if (!keep) compute(k, -1, 0);
8bc }
```

2.35 Stable Marriage

```
// Emparelha todos os elementos de A com elementos de B
// de forma que nao exista um par x \in A, y \in B
// e x nao pareado com y tal que x prefira parear com y
// e y prefira parear com x.
//
// a[i] contem os elementos de B ordenados por preferencia de i
// b[j] contem os elementos de A ordenados por preferencia de j
// |A| <= |B|
//</pre>
```

```
// Retorna um vetor v de tamanho |A| onde v[i] guarda o match de i.
// O(|A| * |B|)
// Off8d5
380 vector<int> stable_marriage(vector<vector<int>> &a,
    vector<vector<int>> &b) {
        int n = a.size(), m = b.size();
652
83e
        assert(a[0].size() == m and b[0].size() == n and n <= m);
        vector < int > match(m, -1), it(n, 0);
017
e6f
        vector inv_b(m, vector < int > (n));
a34
        for (int i = 0; i < m; i++) for (int j = 0; j < n; j++)
9f2
            inv_b[i][b[i][j]] = j;
26a
        queue < int > q;
5af
        for (int i = 0; i < n; i++) q.push(i);</pre>
402
        while (q.size()) {
379
            int i = q.front(); q.pop();
            int j = a[i][it[i]];
4b8
57c
            if (match[j] == -1) match[j] = i;
02d
             else if (inv_b[j][i] < inv_b[j][match[j]]) {</pre>
5d1
                 q.emplace(match[j]);
e7d
                it[match[j]]++;
f1d
                 match[j] = i;
bc4
            } else q.emplace(i), it[i]++;
258
        }
        vector<int> ret(n):
825
d72
        for (int i = 0; i < m; i++) if (match[i] != -1) ret[match[i]]</pre>
   = i;
edf
        return ret;
Off }
2.36 Tarjan para SCC
// O(n + m)
// 573bfa
042 vector < int > g[MAX];
4ce stack<int> s;
a42 int vis[MAX], comp[MAX];
3fd int id[MAX]:
// se quiser comprimir ciclo ou achar ponte em grafo nao direcionado,
// colocar um if na dfs para nao voltar pro pai da DFS tree
f32 int dfs(int i, int& t) {
```

```
cf0
        int lo = id[i] = t++;
18e
        s.push(i);
0c2
        vis[i] = 2;
        for (int j : g[i]) {
48e
740
            if (!vis[j]) lo = min(lo, dfs(j, t));
994
            else if (vis[j] == 2) lo = min(lo, id[j]);
d64
        }
        // aresta de i pro pai eh uma ponte (no caso nao direcionado)
3de
        if (lo == id[i]) while (1) {
3c3
            int u = s.top(); s.pop();
9c5
            vis[u] = 1, comp[u] = i;
            if (u == i) break;
2ef
266
        }
253
        return lo;
38a }
f93 void tarjan(int n) {
        int t = 0:
6bb
        for (int i = 0; i < n; i++) vis[i] = 0;</pre>
991
3be
        for (int i = 0; i < n; i++) if (!vis[i]) dfs(i, t);</pre>
ea1 }
2.37 Topological Sort
// Retorna uma ordenacaoo topologica de g
// Se g nao for DAG retorna um vetor vazio
// O(n + m)
// bdc95e
042 vector < int > g[MAX];
b6a vector<int> topo_sort(int n) {
46e
        vector < int > ret(n,-1), vis(n,0);
f51
        int pos = n-1, dag = 1;
        function < void(int) > dfs = [&](int v) {
36d
cca
            vis[v] = 1:
440
            for (auto u : g[v]) {
152
                if (vis[u] == 1) dag = 0;
532
                else if (!vis[u]) dfs(u);
e37
            }
d44
            ret[pos--] = v, vis[v] = 2;
```

```
57e
        };
        for (int i = 0; i < n; i++) if (!vis[i]) dfs(i);</pre>
158
        if (!dag) ret.clear();
d8f
        return ret;
edf
d6b }
2.38 Vertex cover
// Encontra o tamanho do vertex cover minimo
// Da pra alterar facil pra achar os vertices
// Parece rodar com < 2 s pra N = 90
// O(n * 1.38^n)
// 9c5024
76a namespace cover {
5a4
        const int MAX = 96:
        vector < int > g[MAX];
042
823
        bitset < MAX > bs[MAX];
1a8
        int n:
697
        void add(int i, int j) {
bd0
            if (i == j) return;
78c
            n = max({n, i+1, j+1});
200
            bs[i][j] = bs[j][i] = 1;
        }
203
        int rec(bitset < MAX > m) {
6.0
1a4
            int ans = 0;
25b
            for (int x = 0; x < n; x++) if (m[x]) {
002
                 bitset < MAX > comp;
4bf
                 function < void(int) > dfs = [&](int i) {
b96
                     comp[i] = 1, m[i] = 0;
0 c 3
                     for (int j : g[i]) if (m[j]) dfs(j);
815
                 };
963
                 dfs(x);
d34
                 int ma, deg = -1, cyc = 1;
                 for (int i = 0; i < n; i++) if (comp[i]) {</pre>
417
dOb
                     int d = (bs[i]&comp).count();
                     if (d <= 1) cyc = 0;
18a
c1f
                     if (d > deg) deg = d, ma = i;
d8e
269
                 if (deg <= 2) { // caminho ou ciclo</pre>
340
                     ans += (comp.count() + cyc) / 2;
```

```
5e2
                     continue;
                 }
702
3f9
                 comp[ma] = 0;
                 // ou ta no cover, ou nao ta no cover
1dd
                 ans += \min(1 + rec(comp), deg + rec(comp & \sim bs[ma]));
6e6
            }
ba7
            return ans;
2ec
f5c
        int solve() {
3c5
            bitset < MAX > m;
603
            for (int i = 0; i < n; i++) {</pre>
939
                 m[i] = 1:
f90
                 for (int j = 0; j < n; j++)
741
                     if (bs[i][j]) g[i].push_back(j);
13e
            }
4f9
            return rec(m);
708
9c5 }
```

2.39 Virtual Tree

```
// Comprime uma arvore dado um conjunto S de vertices, de forma que
// o conjunto de vertices da arvore comprimida contenha S e seja
// minimal e fechado sobre a operação de LCA
// Se |S| = k, a arvore comprimida tem menos que 2k vertices
// As arestas de virt possuem a distancia do vertice ate o vizinho
// Retorna a raiz da virtual tree
//
// lca::pos deve ser a ordem de visitacao no dfs
// voce pode usar o LCAcomHLD, por exemplo
//
// O(k log(k))
// 42d990
b36 vector <pair <int, int>> virt[MAX];
d41 #warning lembrar de buildar o LCA antes
c14 int build_virt(vector<int> v) {
        auto cmp = [&](int i, int j) { return lca::pos[i] <</pre>
   lca::pos[j]; };
074
        sort(v.begin(), v.end(), cmp);
        for (int i = v.size()-1; i; i--) v.push_back(lca::lca(v[i],
   v[i-1]));
        sort(v.begin(), v.end(), cmp);
074
d76
        v.erase(unique(v.begin(), v.end()), v.end());
37 c
        for (int i = 0; i < v.size(); i++) virt[v[i]].clear();</pre>
```

```
197
        for (int i = 1; i < v.size(); i++) virt[lca::lca(v[i-1],</pre>
   v[i])].clear();
        for (int i = 1; i < v.size(); i++) {</pre>
ad7
            int parent = lca::lca(v[i-1], v[i]);
51b
            int d = lca::dist(parent, v[i]);
290
d41 #warning soh to colocando aresta descendo
4d0
            virt[parent].emplace_back(v[i], d);
fe5
832
        return v[0];
142 }
```

3 Problemas

3.1 Algoritmo Hungaro

```
// Resolve o problema de assignment (matriz n x n)
// Colocar os valores da matriz em 'a' (pode < 0)</pre>
// assignment() retorna um par com o valor do
// assignment minimo, e a coluna escolhida por cada linha
//
// O(n^3)
// 64c53e
a6a template < typename T > struct hungarian {
1a8
         int n;
a08
         vector < vector < T >> a:
f36
         vector<T> u, v;
5ff
         vector < int > p, way;
f1e
         T inf:
c3f
         hungarian(int n_{-}): n(n_{-}), u(n+1), v(n+1), p(n+1), way(n+1) {
b2f
             a = vector < vector < T >> (n, vector < T > (n));
1f3
             inf = numeric_limits <T>::max();
78f
         }
         pair <T, vector <int>> assignment() {
d67
78a
             for (int i = 1; i <= n; i++) {</pre>
8c9
                 p[0] = i;
625
                  int i0 = 0;
ce7
                  vector <T> minv(n+1, inf);
241
                  vector < int > used(n+1, 0);
016
                  do {
472
                      used[j0] = true;
d24
                      int i0 = p[i0], i1 = -1;
7e5
                      T delta = inf;
9ac
                      for (int j = 1; j <= n; j++) if (!used[j]) {</pre>
7bf
                          T cur = a[i0-1][j-1] - u[i0] - v[j];
```

```
9f2
                         if (cur < minv[j]) minv[j] = cur, wav[j] = j0;</pre>
821
                         if (minv[j] < delta) delta = minv[j], j1 = j;</pre>
                     }
4d1
f63
                     for (int j = 0; j \le n; j++)
2c5
                         if (used[j]) u[p[j]] += delta, v[j] -= delta;
                         else minv[j] -= delta;
6ec
6d4
                     j0 = j1;
                } while (p[j0] != 0);
f4f
016
                 do {
4 c 5
                     int j1 = way[j0];
0d7
                     p[j0] = p[j1];
6d4
                     j0 = j1;
886
                 } while (j0);
38d
306
            vector < int > ans(n);
            for (int j = 1; j \le n; j++) ans[p[j]-1] = j-1;
6db
            return make_pair(-v[0], ans);
da3
979
        }
64c };
```

3.2 Algoritmo MO - queries em caminhos de arvore

```
// Problema que resolve: https://www.spoj.com/problems/COT2/
// Complexidade sendo c = O(update) e SQ = sqrt(n):
// O((n + q) * sqrt(n) * c)
// 395329
1bc const int MAX = 40010, SQ = 400;
042 vector < int > g[MAX];
c54 namespace LCA { ... }
249 int in[MAX], out[MAX], vtx[2 * MAX];
81b bool on [MAX];
4c3 int dif, freq[MAX];
9e2 vector < int > w;
d9a void dfs(int v, int p, int &t) {
659
        vtx[t] = v, in[v] = t++;
        for (int u : g[v]) if (u != p) {
            dfs(u, v, t);
c53
e0f
217
        vtx[t] = v, out[v] = t++;
42b }
```

```
e5f void update(int p) { // faca alteracoes aqui
        int v = vtx[p];
bbc
0ec
        if (not on[v]) { // insere vtx v
            dif += (freq[w[v]] == 0);
31c
b20
            freq[w[v]]++;
cf7
        }
4e6
        else { // retira o vertice v
0a9
            dif \rightarrow (freq[w[v]] == 1);
            freq[w[v]]--;
fd3
2c8
        }
73e
        on[v] = not on[v];
ea9 }
a3a vector < tuple < int , int , int >> build_queries (const vector < pair < int ,
   int>>& q) {
        LCA::build(0);
ea6
f77
        vector<tuple<int, int, int>> ret;
        for (auto [1, r] : q){
aa9
d24
            if (in[r] < in[l]) swap(l, r);</pre>
6f9
            int p = LCA::lca(l, r);
826
            int init = (p == 1) ? in[1] : out[1];
07a
            ret.emplace_back(init, in[r], in[p]);
b0e
        }
edf
        return ret;
8e6 }
f31 vector<int> mo_tree(const vector<pair<int, int>>& vq){
        int t = 0;
        dfs(0, -1, t);
dab
        auto q = build_queries(vq);
af1
f48
        vector < int > ord(q.size());
        iota(ord.begin(), ord.end(), 0);
be8
d01
        sort(ord.begin(), ord.end(), [&] (int 1, int r) {
             int bl = get<0>(q[1]) / SQ, br = <math>get<0>(q[r]) / SQ;
d8d
596
            if (bl != br) return bl < br;</pre>
158
            else if (b1 % 2 == 1) return get<1>(q[1]) < get<1>(q[r]);
f1d
             else return get<1>(q[1]) > get<1>(q[r]);
0a8
        });
        memset(freq, 0, sizeof freq);
80e
bf6
        dif = 0:
ff2
        vector<int> ret(q.size());
        int 1 = 0, r = -1:
3d9
```

```
8b0
        for (int i : ord) {
             auto [ql, qr, qp] = q[i];
3c7
af7
             while (r < qr) update(++r);</pre>
             while (1 > q1) update(--1);
d6b
             while (1 < q1) update(1++);</pre>
951
            while (r > qr) update(r--);
6a1
            if (qp < 1 \text{ or } qp > r)  { // se LCA estah entre as pontas
3d8
74b
                 update(qp);
                 ret[i] = dif;
2e1
74b
                 update(qp);
e83
0fe
             else ret[i] = dif:
Ofd
edf
        return ret;
48d }
```

3.3 Angle Range Intersection

```
// Computa intersecao de angulos
// Os angulos (arcos) precisam ter comprimeiro < pi
// (caso contrario a intersecao eh estranha)
// Tudo 0(1)
// 5e1c85
32a struct angle range {
        static constexpr ld ALL = 1e9, NIL = -1e9;
75e
395
        ld 1, r;
        angle_range() : 1(ALL), r(ALL) {}
c77
894
        angle_range(ld l_, ld r_) : l(l_), r(r_) { fix(l), fix(r); }
4ee
        void fix(ld& theta) {
da7
            if (theta == ALL or theta == NIL) return;
323
            if (theta > 2*pi) theta -= 2*pi;
868
            if (theta < 0) theta += 2*pi;</pre>
625
2ee
        bool empty() { return l == NIL; }
931
        bool contains(ld q) {
40f
            fix(a):
4d7
            if (1 == ALL) return true;
fec
            if (1 == NIL) return false;
            if (1 < r) return 1 < q and q < r;
6a6
075
            return q > 1 or q < r;</pre>
800
        friend angle_range operator &(angle_range p, angle_range q) {
9c7
            if (p.l == ALL or q.l == NIL) return q;
743
```

```
20f
            if (q.l == ALL or p.l == NIL) return p;
7d5
            if (p.1 > p.r \text{ and } q.1 > q.r) \text{ return } \{\max(p.1, q.1),
   min(p.r, q.r)};
            if (q.1 > q.r) swap(p.1, q.1), swap(p.r, q.r);
aa6
            if (p.1 > p.r) {
848
249
                 if (q.r > p.l) return {max(q.l, p.l) , q.r};
6f7
                 else if (q.1 < p.r) return \{q.1, \min(q.r, p.r)\};
270
                 return {NIL, NIL};
337
5a8
            if (max(p.1, q.1) > min(p.r, q.r)) return {NIL, NIL};
bcb
            return {max(p.1, q.1), min(p.r, q.r)};
142
        }
5e1 }:
```

3.4 Area da Uniao de Retangulos

```
// O(n log(n))
// 5d8d2f
aa4 namespace seg {
6b3
        pair<int, 11> seg[4*MAX];
        11 lazy[4*MAX], *v;
b<sub>1</sub>b
1a8
        int n;
e01
        pair<int, 11> merge(pair<int, 11> 1, pair<int, 11> r){
719
             if (1.second == r.second) return {1.first+r.first,
   1.second}:
53b
             else if (1.second < r.second) return 1;</pre>
aa0
             else return r;
        }
d82
6fc
        pair < int , ll > build(int p=1, int l=0, int r=n-1) {
3c7
            lazv[p] = 0;
            if (1 == r) return seg[p] = {1, v[1]};
bf8
ee4
            int m = (1+r)/2;
            return seg[p] = merge(build(2*p, 1, m), build(2*p+1, m+1,
432
   r));
f94
        }
        void build(int n2, 11* v2) {
d9e
680
            n = n2, v = v2;
6f2
            build();
f8a
        }
ceb
        void prop(int p, int l, int r) {
208
            seg[p].second += lazy[p];
2c9
            if (1 != r) lazy[2*p] += lazy[p], lazy[2*p+1] += lazy[p];
3c7
            lazv[p] = 0;
bf2
        }
```

```
693
        pair < int, ll > query (int a, int b, int p=1, int l=0, int r=n-1)
   {
6b9
            prop(p, 1, r);
            if (a <= 1 and r <= b) return seg[p];</pre>
527
9ъ7
            if (b < 1 or r < a) return {0, LINF};</pre>
            int m = (1+r)/2;
ee4
            return merge (query (a, b, 2*p, 1, m), query (a, b, 2*p+1),
   m+1, r));
786
07c
        pair < int, 11 > update(int a, int b, int x, int p=1, int 1=0,
   int r=n-1) {
            prop(p, 1, r);
6b9
9a3
            if (a \le 1 \text{ and } r \le b)
b94
                lazv[p] += x;
6b9
                prop(p, 1, r);
                return seg[p];
534
821
            }
            if (b < l or r < a) return seg[p];</pre>
e9f
            int m = (1+r)/2;
ee4
086
            return seg[p] = merge(update(a, b, x, 2*p, 1, m),
                     update(a, b, x, 2*p+1, m+1, r));
579
c65
        }
043 };
eb5 ll seg_vec[MAX];
8be ll area_sq(vector<pair<pair<int, int>, pair<int, int>>> &sq){
28c
        vector<pair<int, int>, pair<int, int>>> up;
60a
        for (auto it : sq){
619
            int x1, y1, x2, y2;
            tie(x1, y1) = it.first;
ae0
68e
            tie(x2, y2) = it.second;
80f
            up.push_back({{x1+1, 1}, {y1, y2}});
            up.push back(\{\{x2+1, -1\}, \{v1, v2\}\}\}):
aee
6c3
092
        sort(up.begin(), up.end());
049
        memset(seg_vec, 0, sizeof seg_vec);
6fe
        11 H_MAX = MAX;
156
        seg::build(H_MAX-1, seg_vec);
7ba
        auto it = up.begin();
04b
        11 \text{ ans} = 0:
        while (it != up.end()){
f14
07f
            11 L = (*it).first.first;
718
            while (it != up.end() && (*it).first.first == L){
127
                int x, inc, v1, v2;
d35
                tie(x, inc) = it->first;
d3d
                tie(y1, y2) = it -> second;
```

```
5d1
                seg::update(v1+1, v2, inc);
40d
9b1
            }
852
            if (it == up.end()) break;
            11 R = (*it).first.first:
d8a
f59
            11 W = R-L:
            auto jt = seg::query(0, H_MAX-1);
efd
91a
            11 H = H_MAX - 1;
e8a
            if (jt.second == 0) H -= jt.first;
8df
            ans += W*H;
5c8
        }
ba7
        return ans:
385 }
```

3.5 Area Maxima de Histograma

```
// Assume que todas as barras tem largura 1,
// e altura dada no vetor v
//
// O(n)
// e43846
15e ll area(vector<int> v) {
b73
        ll ret = 0;
        stack < int > s:
4ce
        // valores iniciais pra dar tudo certo
        v.insert(v.begin(), -1);
447
d56
        v.insert(v.end(), -1);
1f8
        s.push(0);
        for(int i = 0; i < (int) v.size(); i++) {</pre>
0be
78e
            while (v[s.top()] > v[i]) {
265
                11 h = v[s.top()]; s.pop();
                ret = max(ret, h * (i - s.top() - 1));
de1
40a
18e
            s.push(i);
020
        }
edf
        return ret;
e43 }
```

3.6 Binomial modular

```
// Computa C(n, k) mod m em O(m + log(m) log(n))
// = O(rapido)
// ed4344
```

```
97c ll divi[MAX];
398 ll expo(ll a, ll b, ll m) {
       if (!b) return 1:
1c1
399
        ll ans = expo(a*a\%m, b/2, m);
751
       if (b\%2) ans *= a:
2e9
        return ans%m;
754 }
f0a ll inv(ll a, ll b){
        return 1<a ? b - inv(b%a,a)*b/a : 1;</pre>
041 }
153 template < typename T > tuple < T, T, T > ext_gcd(T a, T b) {
        if (!a) return {b, 0, 1};
3bd
550
        auto [g, x, y] = ext_gcd(b\%a, a);
        return \{g, y - b/a*x, x\};
c59
537 }
bfe template < typename T = 11 > struct crt {
627
        Ta, m;
5f3
        crt(): a(0), m(1) {}
7eb
        crt(T a_{,} T m_{,}) : a(a_{,}), m(m_{,}) \{ \}
911
        crt operator * (crt C) {
            auto [g, x, y] = ext_gcd(m, C.m);
238
dc0
            if ((a - C.a) \% g) a = -1;
4f9
            if (a == -1 or C.a == -1) return crt(-1, 0);
d09
           T lcm = m/g*C.m;
            T ans = a + (x*(C.a-a)/g \% (C.m/g))*m;
eb2
d8d
            return crt((ans % lcm + lcm) % lcm, lcm);
        }
1f2
0d9 }:
6f2 pair<11, 11> divide_show(11 n, int p, int k, int pak) {
        if (n == 0) return {0, 1};
4f7
d02
        11 blocos = n/pak, falta = n%pak;
2ce
        ll periodo = divi[pak], resto = divi[falta];
616
        11 r = expo(periodo, blocos, pak)*resto%pak;
445
        auto rec = divide_show(n/p, p, k, pak);
        ll v = n/p + rec.first;
a51
bb9
        r = r*rec.second % pak;
90f
        return {y, r};
533 }
```

```
6e6 ll solve_pak(ll n, ll x, int p, int k, int pak) {
        divi[0] = 1;
d34
f2b
        for (int i = 1; i <= pak; i++) {</pre>
            divi[i] = divi[i-1];
901
840
            if (i%p) divi[i] = divi[i] * i % pak;
51a
        }
        auto dn = divide_show(n, p, k, pak), dx = divide_show(x, p, k,
4ac
   pak),
162
              dnx = divide_show(n-x, p, k, pak);
768
        11 y = dn.first-dx.first-dnx.first, r =
b64
            (dn.second*inv(dx.second, pak)%pak)*inv(dnx.second,
    pak)%pak;
035
        return expo(p, y, pak) * r % pak;
d78 }
9dd ll solve(ll n, ll x, int mod) {
        vector<pair<int, int>> f;
490
c3b
        int mod2 = mod;
        for (int i = 2; i*i <= mod2; i++) if (mod2%i==0) {</pre>
7b4
            int c = 0:
aff
75b
            while (mod2\%i==0) mod2 /= i, c++;
2a1
            f.push_back({i, c});
fe7
        }
Off
        if (mod2 > 1) f.push_back({mod2, 1});
        crt ans(0, 1):
e96
a13
        for (int i = 0; i < f.size(); i++) {</pre>
702
            int pak = 1;
7e4
            for (int j = 0; j < f[i].second; j++) pak *= f[i].first;</pre>
            ans = ans * crt(solve_pak(n, x, f[i].first, f[i].second,
    pak), pak);
7fd
        }
5fb
        return ans.a:
689 }
3.7 Closest pair of points
// O(nlogn)
// f90265
915 pair < pt , pt > closest_pair_of_points(vector < pt > v) {
        int n = v.size();
3d2
        sort(v.begin(), v.end());
fca
        for (int i = 1; i < n; i++) if (v[i] == v[i-1]) return</pre>
   {v[i-1], v[i]};
c20
        auto cmp_y = [&](const pt &1, const pt &r) {
```

```
b53
            if (1.y != r.y) return 1.y < r.y;</pre>
920
            return l.x < r.x;</pre>
55a
        };
62e
        set < pt, decltype(cmp_y) > s(cmp_y);
3d9
        int 1 = 0, r = -1:
6a2
        11 d2_min = numeric_limits < ll >:: max();
4d5
        pt pl, pr;
bd1
        const int magic = 5;
a55
        while (r+1 < n) {
            auto it = s.insert(v[++r]).first;
7 f 1
c92
            int cnt = magic/2;
773
            while (cnt-- and it != s.begin()) it--;
a01
            cnt = 0:
d68
            while (cnt++ < magic and it != s.end()) {</pre>
f19
                 if (!((*it) == v[r])) {
                     11 d2 = dist2(*it, v[r]);
67e
74e
                     if (d2_min > d2) {
229
                         d2_min = d2;
                         pl = *it;
841
4f2
                         pr = v[r];
                     }
7d9
                 }
10a
40d
                 it++;
801
            while (1 < r \text{ and } sq(v[1].x-v[r].x) > d2_min)
   s.erase(v[1++]):
de6
c74
        return {pl, pr};
f90 }
```

3.8 Coloração de Grafo de Intervalo

```
// Colore os intervalos com o numero minimo
// de cores de tal forma que dois intervalos
// que se interceptam tem cores diferentes
// As cores vao de 1 ate n
//
// O(n log(n))
// 83a32d
615 vector<int> coloring(vector<pair<int, int>>& v) {
3d2
        int n = v.size();
c08
        vector<pair<int, pair<int, int>>> ev;
603
       for (int i = 0; i < n; i++) {</pre>
150
            ev.push_back({v[i].first, {1, i}});
cda
            ev.push_back({v[i].second, {0, i}});
6a4
        }
```

```
49e
        sort(ev.begin(), ev.end());
360
        vector < int > ans(n), avl(n);
265
        for (int i = 0; i < n; i++) avl.push_back(n-i);
        for (auto i : ev) {
4bf
            if (i.second.first == 1) {
cbe
021
                ans[i.second.second] = avl.back();
a00
                avl.pop_back();
e98
            } else avl.push_back(ans[i.second.second]);
3a6
ba7
        return ans;
83a }
3.9 Conectividade Dinamica
// Offline com Divide and Conquer e
// DSU com rollback
// O(n log^2(n))
// 043d93
8f2 typedef pair <int, int> T;
1cd namespace data {
553
        int n, ans;
573
        int p[MAX], sz[MAX];
        stack<int> S;
ee6
e5b
        void build(int n2) {
1e3
            n = n2:
8a6
            for (int i = 0; i < n; i++) p[i] = i, sz[i] = 1;
0b2
            ans = n:
cba
        }
1b1
        int find(int k) {
006
            while (p[k] != k) k = p[k];
839
            return k;
        }
c1e
072
        void add(T x) {
700
            int a = x.first, b = x.second;
605
            a = find(a), b = find(b);
843
            if (a == b) return S.push(-1);
e7d
            ans --:
3c6
            if (sz[a] > sz[b]) swap(a, b);
4c2
            S.push(a);
582
            sz[b] += sz[a];
84b
            p[a] = b;
e1a
        }
5eb
        int query() {
ba7
            return ans;
```

```
35 c
        }
5cf
        void rollback() {
465
            int u = S.top(); S.pop();
            if (u == -1) return;
61 c
270
            sz[p[u]] -= sz[u]:
            p[u] = u;
546
Odf
            ans++:
456
       }
568 };
357 int ponta[MAX]; // outra ponta do intervalo ou -1 se for query
4f0 int ans[MAX], n, q;
487 T qu[MAX];
47b void solve(int l = 0, int r = q-1) {
        if (1 >= r) {
0b1
8c0
            ans[1] = data::query(); // agora a estrutura ta certa
505
            return:
f77
962
        int m = (1+r)/2, qnt = 1;
        for (int i = m+1; i <= r; i++) if (ponta[i]+1 and ponta[i] < 1)</pre>
fc7
37d
            data::add(qu[i]), qnt++;
        solve(1. m):
221
593
        while (--qnt) data::rollback();
a2c
        for (int i = 1; i <= m; i++) if (ponta[i]+1 and ponta[i] > r)
37d
            data::add(qu[i]), qnt++;
37b
        solve(m+1, r);
281
        while (qnt--) data::rollback();
0d4 }
      Conectividade Dinamica 2
```

```
// Offline com link-cut trees
// O(n log(n))
// d38e4e
1ef namespace lct {
3c9
        struct node {
19f
            int p, ch[2];
a2a
            int val, sub;
aa6
            bool rev;
f93
            node() {}
            node(int v) : p(-1), val(v), sub(v), rev(0) { ch[0] = }
   ch[1] = -1; }
        };
cac
c53
        node t[2*MAX]; // MAXN + MAXQ
```

```
99e
        map<pair<int, int>, int> aresta;
e4d
        int sz:
95a
        void prop(int x) {
            if (t[x].rev) {
aa2
f95
                swap(t[x].ch[0], t[x].ch[1]);
379
                if (t[x].ch[0]+1) t[t[x].ch[0]].rev ^= 1;
c3d
                if (t[x].ch[1]+1) t[t[x].ch[1]].rev ^= 1;
50e
693
            t[x].rev = 0;
750
        }
564
        void update(int x) {
e8d
            t[x].sub = t[x].val:
8ca
            for (int i = 0; i < 2; i++) if (t[x].ch[i]+1) {</pre>
621
                prop(t[x].ch[i]);
78d
                t[x].sub = min(t[x].sub, t[t[x].ch[i]].sub);
3e4
        }
9bf
971
        bool is_root(int x) {
            return t[x].p == -1 or (t[t[x].p].ch[0] != x and
657
   t[t[x].p].ch[1] != x);
cf1
ed6
        void rotate(int x) {
497
            int p = t[x].p, pp = t[p].p;
fc4
            if (!is_root(p)) t[pp].ch[t[pp].ch[1] == p] = x;
251
            bool d = t[p].ch[0] == x;
            t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
461
a76
            if (t[p].ch[!d]+1) t[t[p].ch[!d]].p = p;
            t[x].p = pp, t[p].p = x;
8fa
444
            update(p), update(x);
f31
        }
238
        int splay(int x) {
            while (!is root(x)) {
18c
497
                int p = t[x].p, pp = t[p].p;
77b
                if (!is_root(p)) prop(pp);
be5
                prop(p), prop(x);
                if (!is\_root(p)) rotate((t[pp].ch[0] == p)^(t[p].ch[0]
0 c 5
   == x) ? x : p);
64f
                rotate(x);
72c
            }
aab
            return prop(x), x;
08f
f16
        int access(int v) {
0eb
            int last = -1:
d9f
            for (int w = v; w+1; update(last = w), splay(v), w =
   t[v].p)
024
                splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
```

```
3d3
            return last;
294
952
        void make_tree(int v, int w=INF) { t[v] = node(w); }
        bool conn(int v, int w) {
82f
2cf
            access(v). access(w):
            return v == w ? true : t[v].p != -1;
b9b
ec0
        }
277
        void rootify(int v) {
5e3
            access(v);
            t[v].rev ^= 1;
a02
a05
        }
a1d
        int query(int v, int w) {
b54
            rootifv(w). access(v):
            return t[v].sub;
249
c28
204
        void link_(int v, int w) {
821
            rootify(w);
389
            t[w].p = v;
523
6b8
        void link(int v, int w, int x) { // v--w com peso x
379
            int id = MAX + sz++:
110
            aresta[make_pair(v, w)] = id;
            make_tree(id, x);
ab6
c88
            link_(v, id), link_(id, w);
984
e63
        void cut_(int v, int w) {
            rootify(w), access(v);
b54
264
            t[v].ch[0] = t[t[v].ch[0]].p = -1;
7cd
031
        void cut(int v, int w) {
b0f
            int id = aresta[make_pair(v, w)];
a4a
            cut_(v, id), cut_(id, w);
840
        }
0d7 }
893 void dyn_conn() {
c5f
        int n, q; cin >> n >> q;
d6e
        vector < int > p(2*g, -1); // outra ponta do intervalo
b4f
        for (int i = 0; i < n; i++) lct::make_tree(i);</pre>
fbf
        vector < pair < int , int >> qu(q);
139
        map<pair<int, int>, int> m;
abf
        for (int i = 0; i < q; i++) {
3c2
            char c; cin >> c;
            if (c == '?') continue;
ef6
602
           int a, b; cin >> a >> b; a--, b--;
d11
            if (a > b) swap(a, b);
8a1
            qu[i] = {a, b};
```

```
8d7
            if (c == '+') {
94b
                p[i] = i+q, p[i+q] = i;
                m[make_pair(a, b)] = i;
906
9d9
            } else {
412
                int j = m[make_pair(a, b)];
ac2
                p[i] = j, p[j] = i;
            }
850
        }
9e5
447
        int ans = n;
        for (int i = 0; i < q; i++) {</pre>
abf
87d
            if (p[i] == -1) {
                cout << ans << endl; // numero de comp conexos</pre>
886
5e2
                continue:
b35
            }
            int a = qu[i].first, b = qu[i].second;
69d
c4d
            if (p[i] > i) { // +
ac5
                if (lct::conn(a, b)) {
                     int mi = lct::query(a, b);
18f
993
                     if (p[i] < mi) {</pre>
                         p[p[i]] = p[i];
dd3
5e2
                         continue;
474
6f7
                     lct::cut(qu[p[mi]].first, qu[p[mi]].second), ans++;
6ea
                     p[mi] = mi;
9a9
d1d
                lct::link(a, b, p[i]), ans--;
9d0
            } else if (p[i] != i) lct::cut(a, b), ans++; // -
        }
c03
56a }
3.11 Conj. Indep. Maximo com Peso em Grafo de Intervalo
// Retorna os indices ordenados dos intervalos selecionados
// Se tiver empate, retorna o que minimiza o comprimento total
// O(n log(n))
// c4dbe2
31e vector < int > ind_set(vector < tuple < int, int, int >> & v) {
b27
        vector<tuple<int, int, int>> w;
        for (int i = 0; i < v.size(); i++) {</pre>
f14
            w.push_back(tuple(get<0>(v[i]), 0, i));
e85
            w.push_back(tuple(get<1>(v[i]), 1, i));
6f0
17f
d1d
        sort(w.begin(), w.end());
844
        vector < int > nxt(v.size());
```

```
c22
        vector < pair < 11, int >> dp(v.size());
0eb
        int last = -1:
723
        for (auto [fim, t, i] : w) {
25a
            if (t == 0) {
4ca
                nxt[i] = last;
5e2
                continue;
5fd
            }
78b
            dp[i] = \{0, 0\};
cb8
            if (last != -1) dp[i] = max(dp[i], dp[last]);
            pair<11, int> pega = {get<2>(v[i]), -(get<1>(v[i]) -
911
   get < 0 > (v[i]) + 1);
            if (nxt[i] != -1) pega.first += dp[nxt[i]].first,
   pega.second += dp[nxt[i]].second;
b08
            if (pega > dp[i]) dp[i] = pega;
            else nxt[i] = last;
7cb
            last = i;
381
       }
b7c
        pair <11, int > ans = {0, 0};
977
919
        int idx = -1;
        for (int i = 0; i < v.size(); i++) if (dp[i] > ans) ans =
   dp[i], idx = i;
        vector<int> ret:
4b8
fdd
        while (idx != -1) {
d69
            if (get < 2 > (v[idx]) > 0 and
                (nxt[idx] == -1 or get<1>(v[nxt[idx]]) <</pre>
   get <0>(v[idx]))) ret.push_back(idx);
e4f
            idx = nxt[idx];
042
0ea
        sort(ret.begin(), ret.end());
edf
        return ret;
c4d }
```

3.12 Convex Hull Dinamico

```
// insert - O(log n) amortizado
// is_inside - O(log n)
// 35c4e8
Ob9 struct upper {
af8
        set <pt> se;
80b
        set <pt>::iterator it;
25 c
        int is_under(pt p) { // 1 -> inside ; 2 -> border
fe0
            it = se.lower_bound(p);
633
            if (it == se.end()) return 0;
a94
            if (it == se.begin()) return p == *it ? 2 : 0;
ca0
            if (ccw(p, *it, *prev(it))) return 1;
```

```
}
dba
        void insert(pt p) {
eaa
            if (is_under(p)) return;
712
a86
            if (it != se.end()) while (next(it) != se.end() and
    !ccw(*next(it), *it, p))
                it = se.erase(it);
316
            if (it != se.begin()) while (--it != se.begin() and
be3
    !ccw(p, *it, *prev(it)))
316
                it = se.erase(it);
0c8
            se.insert(p):
5da
        }
750 };
06f struct dyn_hull {
        upper U, L;
d93
333
        int is_inside(pt p) {
632
            int u = U.is_under(p), l = L.is_under({-p.x, -p.y});
            if (!u or !1) return 0:
4 c c
            return max(u, 1);
fc0
478
        }
        void insert(pt p) {
eaa
86c
            U.insert(p):
925
            L.insert({-p.x, -p.y});
64b
        }
285
        int size() {
7c2
            int ans = U.se.size() + L.se.size();
            return ans <= 2 ? ans/2 : ans-2;</pre>
1 c 9
ad5
        }
65e }:
3.13 Distancia maxima entre dois pontos
// \max_{dist2(v)} - O(n \log(n))
// max_dist_manhattan - O(n)
// Quadrado da Distancia Euclidiana (precisa copiar convex_hull, ccw e
   pt)
// bdace4
859 ll max_dist2(vector<pt> v) {
221
        v = convex_hull(v);
a 14
        if (v.size() <= 2) return dist2(v[0], v[1%v.size()]);</pre>
04b
        11 \text{ ans} = 0;
        int n = v.size(), j = 0;
323
```

return ccw(p, *prev(it), *it) ? 0 : 2;

402

```
603
        for (int i = 0; i < n; i++) {</pre>
            while (!ccw(v[(i+1)%n]-v[i], pt(0, 0), v[(j+1)%n]-v[j])) j
057
   = (j+1) \%n;
e7a
            ans = \max(\{ans, dist2(v[i], v[i]), dist2(v[(i+1)%n],
   v[i])});
        }
1f6
ba7
        return ans:
bda }
// Distancia de Manhattan
// 4e96f0
c51 template < typename T> T max_dist_manhattan(vector < pair < T, T>> v) {
        T min sum. max sum. min dif. max dif:
        min_sum = max_sum = v[0].first + v[0].second;
4f5
        min_dif = max_dif = v[0].first - v[0].second;
271
c25
        for (auto [x, y] : v) {
            min_sum = min(min_sum, x+y);
1cb
            max_sum = max(max_sum, x+y);
683
            min_dif = min(min_dif, x-y);
782
af7
            max_dif = max(max_dif, x-y);
e3a
9f0
        return max(max_sum - min_sum, max_dif - min_dif);
4e9 }
3.14 Distinct Range Query
// build - O(n (log n + log(sigma)))
// query - O(log(sigma))
// 5c7aa1
789 namespace perseg { };
53d int qt[MAX];
edc void build(vector<int>& v) {
3d2
        int n = v.size();
16b
        perseg::build(n);
        map < int , int > last;
663
        int at = 0;
05e
603
        for (int i = 0: i < n: i++) {</pre>
817
            if (last.count(v[i])) {
a58
                perseg::update(last[v[i]], -1);
69a
                at++:
d1f
4f2
            perseg::update(i, 1);
460
            qt[i] = ++at;
```

efe

last[v[i]] = i;

```
d6f  }
0f4 }

9e3 int query(int 1, int r) {
080    return perseg::query(1, r, qt[r]);
215 }
```

3.15 Distinct Range Query com Update

```
// build - O(n log(n))
// query - O(log^2(n))
// update - 0(log^2(n))
// 2306f3
774 #include <ext/pb_ds/assoc_container.hpp>
30f #include <ext/pb_ds/tree_policy.hpp>
0d7 using namespace __gnu_pbds;
4fc template <class T>
def
        using ord_set = tree<T, null_type, less<T>, rb_tree_tag,
        tree_order_statistics_node_update >;
042 int v[MAX], n, nxt[MAX], prv[MAX];
f60 map<int, set<int> > ocor;
e04 namespace bit {
686
        ord_set <pair <int, int >> bit [MAX];
0a8
        void build() {
3e1
             for (int i = 1; i <= n; i++) bit[i].insert({nxt[i-1],</pre>
   i-1}):
78a
             for (int i = 1; i <= n; i++) {</pre>
edf
                 int j = i + (i\&-i);
d03
                 if (j <= n) for (auto x : bit[i]) bit[j].insert(x);</pre>
5cb
            }
        }
af6
        int pref(int p, int x) {
d3f
7 c.9
             int ret = 0:
bbf
             for (; p; p -= p&-p) ret += bit[p].order_of_key({x, -INF});
edf
             return ret;
        }
0e1
d50
        int query(int 1, int r, int x) {
e55
             return pref(r+1, x) - pref(l, x);
9b4
        }
ff3
        void update(int p, int x) {
f17
             int p2 = p;
5ed
             for (p++; p <= n; p += p&-p) {</pre>
ca8
                 bit[p].erase({nxt[p2], p2});
```

```
f6b
                 bit[p].insert({x, p2});
            }
3df
151
        }
c63 }
0a8 void build() {
383
        for (int i = 0; i < n; i++) nxt[i] = INF:
        for (int i = 0; i < n; i++) prv[i] = -INF;</pre>
7b3
d07
        vector < pair < int , int >> t;
348
        for (int i = 0; i < n; i++) t.push_back({v[i], i});</pre>
3fd
        sort(t.begin(), t.end());
603
        for (int i = 0; i < n; i++) {</pre>
b40
            if (i and t[i].first == t[i-1].first)
                 prv[t[i].second] = t[i-1].second;
565
            if (i+1 < n and t[i].first == t[i+1].first)</pre>
a8b
12f
                 nxt[t[i].second] = t[i+1].second;
48d
        }
a23
        for (int i = 0; i < n; i++) ocor[v[i]].insert(i);</pre>
1d7
        bit::build();
d44 }
aae void muda(int p, int x) {
f92
        bit::update(p, x);
c3d
        nxt[p] = x;
97c }
4ea int query(int a, int b) {
a0a
        return b-a+1 - bit::query(a, b, b+1);
511 }
ff3 void update(int p, int x) { // mudar valor na pos. p para x
        if (prv[p] > -INF) muda(prv[p], nxt[p]);
c0b
        if (nxt[p] < INF) prv[nxt[p]] = prv[p];</pre>
4ae
5bf
        ocor[v[p]].erase(p);
4b4
        if (!ocor[x].size()) {
19d
            muda(p, INF);
8d4
            prv[p] = -INF;
        } else if (*ocor[x].rbegin() < p) {</pre>
a69
5b5
            int i = *ocor[x].rbegin();
f64
            prv[p] = i;
            muda(p, INF);
19d
5f2
            muda(i, p);
9d9
        } else {
d46
            int i = *ocor[x].lower_bound(p);
```

```
33f
            if (prv[i] > -INF) {
f17
                muda(prv[i], p);
8f9
                prv[p] = prv[i];
bc4
            } else prv[p] = -INF;
523
            prv[i] = p;
597
            muda(p, i);
5d7
       }
        v[p] = x; ocor[x].insert(p);
c96
38e }
```

3.16 Dominator Points

```
// Se um ponto A tem ambas as coordenadas >= B, dizemos
// que A domina B
// is_dominated(p) fala se existe algum ponto no conjunto
// que domina p
// insert(p) insere p no conjunto
// (se p for dominado por alguem, nao vai inserir)
// o multiset 'quina' guarda informacao sobre os pontos
// nao dominados por um elemento do conjunto que nao dominam
// outro ponto nao dominado por um elemento do conjunto
// No caso, armazena os valores de x+y esses pontos
//
// Complexidades:
// is_dominated - O(log(n))
// insert - O(log(n)) amortizado
// auerv - 0(1)
// 09ffdc
e2a struct dominator_points {
baf
        set < pair < int , int >> se;
4dd
        multiset < int > quina;
a85
        bool is_dominated(pair<int, int> p) {
80f
            auto it = se.lower_bound(p);
633
            if (it == se.end()) return 0;
            return it->second >= p.second;
ab4
28f
99b
        void mid(pair<int, int> a, pair<int, int> b, bool rem) {
29a
            pair < int , int > m = {a.first+1, b.second+1};
b19
            int val = m.first + m.second;
638
            if (!rem) quina.insert(val);
731
            else quina.erase(quina.find(val));
241
7 c.4
        bool insert(pair<int, int> p) {
fb4
            if (is_dominated(p)) return 0;
80f
            auto it = se.lower_bound(p);
```

```
ca9
            if (it != se.begin() and it != se.end())
d4a
                mid(*prev(it), *it, 1);
1fa
            while (it != se.begin()) {
049
                it--;
                if (it->second > p.second) break;
23 c
                if (it != se.begin()) mid(*prev(it), *it, 1);
b86
316
                it = se.erase(it):
            }
acd
433
            it = se.insert(p).first;
            if (it != se.begin()) mid(*prev(it), *it, 0);
69e
96d
            if (next(it) != se.end()) mid(*it, *next(it), 0);
6a5
            return 1:
688
       }
5eb
        int query() {
956
            if (!quina.size()) return INF;
            return *quina.begin();
add
b8b
       }
09f };
```

3.17 DP de Dominacao 3D

```
// Computa para todo ponto i,
// dp[i] = 1 + max_{i} dominado por i dp[i]
// em que ser dominado eh ter as 3 coordenadas menores
// Da pra adaptar facil para outras dps
//
// O(n log<sup>2</sup> n). O(n) de memoria
// 7c8896
c53 void lis2d(vector<vector<tuple<int, int, int>>>& v, vector<int>&
   dp, int 1, int r) {
893
        if (1 == r) {
56f
            for (int i = 0; i < v[1].size(); i++) {</pre>
8b5
                int ii = get<2>(v[1][i]);
1ce
                 dp[ii] = max(dp[ii], 1);
            }
4b0
505
            return;
3e4
ee4
        int m = (1+r)/2;
62b
        lis2d(v, dp, 1, m);
325
        vector < tuple < int , int , int >> vv[2];
d44
        vector < int > Z:
871
        for (int i = 1; i <= r; i++) for (auto it : v[i]) {</pre>
2ef
            vv[i > m].push_back(it);
042
            Z.push_back(get<1>(it));
0d1
        }
```

```
e9f
        sort(vv[0].begin(), vv[0].end());
9b5
        sort(vv[1].begin(), vv[1].end());
0d1
        sort(Z.begin(), Z.end());
573
        auto get_z = [&](int z) { return lower_bound(Z.begin(),
   Z.end(), z) - Z.begin(); };
        vector < int > bit(Z.size());
c51
181
        int i = 0;
e9a
        for (auto [v, z, id] : vv[1]) {
             while (i < vv[0].size() and get<0>(vv[0][i]) < y) {</pre>
6bd
397
                 auto [v2, z2, id2] = vv[0][i++];
                 for (int p = get_z(z2)+1; p <= Z.size(); p += p&-p)</pre>
ea0
                     bit[p-1] = max(bit[p-1], dp[id2]);
300
82c
            }
d3b
            int q = 0;
fd9
            for (int p = get_z(z); p; p -= p\&-p) q = max(q, bit[p-1]);
614
            dp[id] = max(dp[id], q + 1);
        }
acc
c25
        lis2d(v, dp, m+1, r);
4d6 }
4de vector < int > solve (vector < tuple < int , int , int >> v) {
        int n = v.size();
3d2
cd4
        vector<tuple<int, int, int, int>> vv;
603
        for (int i = 0; i < n; i++) {</pre>
9be
            auto [x, y, z] = v[i];
5bb
            vv.emplace_back(x, y, z, i);
64c
        }
        sort(vv.begin(), vv.end());
bd3
        vector < vector < tuple < int , int , int >>> V;
e11
603
        for (int i = 0; i < n; i++) {</pre>
a5b
            int j = i;
808
            V.emplace back():
            while (j < n and get <0>(vv[j]) == get <0>(vv[i])) {
c01
ba6
                 auto [x, y, z, id] = vv[j++];
                 V.back().emplace_back(y, z, id);
cbb
8bd
            }
452
            i = j-1;
ac4
        }
388
        vector < int > dp(n);
        lis2d(V, dp, 0, V.size()-1);
839
898
        return dp;
bOa }
```

3.18 Gray Code

```
// Gera uma permutacao de 0 a 2^n-1, de forma que
// duas posicoes adjacentes diferem em exatamente 1 bit
//
// O(2<sup>n</sup>)
// 840df4
df6 vector<int> gray_code(int n) {
        vector < int > ret(1 << n);</pre>
f29
        for (int i = 0; i < (1 << n); i++) ret[i] = i^{(i>)1};
        return ret:
edf
840 }
3.19 Half-plane intersection
// Cada half-plane eh identificado por uma reta e a regiao ccw a ela
// O(n log n)
// f56e1c
f4f vector <pt> hp_intersection(vector <line> &v) {
        deque<pt> dq = {{INF, INF}, {-INF, INF}, {-INF, -INF}, {INF,
   -INF}};
d41 #warning considerar trocar por compare_angle
        sort(v.begin(), v.end(), [&](line r, line s) { return
de3
   angle(r.q-r.p) < angle(s.q-s.p); \});
        for(int i = 0; i < v.size() and dq.size() > 1; i++) {
5e9
c69
            pt p1 = dq.front(), p2 = dq.back();
            while (dq.size() and !ccw(v[i].p, v[i].q, dq.back()))
6c6
47b
                p1 = dq.back(), dq.pop_back();
            while (dq.size() and !ccw(v[i].p, v[i].q, dq.front()))
0a2
7cf
                p2 = dq.front(), dq.pop_front();
4d9
            if (!dq.size()) break;
606
            if (p1 == dq.front() and p2 == dq.back()) continue;
c9b
            dq.push_back(inter(v[i], line(dq.back(), p1)));
            dq.push_front(inter(v[i], line(dq.front(), p2)));
            if (dq.size() > 1 and dq.back() == dq.front())
   dq.pop_back();
4d8
b2b
        return vector < pt > (dq.begin(), dq.end());
f56 }
```

3.20 Heap Sort

```
// O(n log n)
// 385e91
f18 void down(vector<int>& v, int n, int i) {
        while ((i = 2*i+1) < n) {
e1f
583
             if (i+1 < n and v[i] < v[i+1]) i++;</pre>
b27
             if (v[i] < v[(i-1)/2]) break;
322
             swap(v[i], v[(i-1)/2]);
170
        }
724 }
eb6 void heap_sort(vector<int>& v) {
        int n = v.size();
61d
        for (int i = n/2-1; i \ge 0; i--) down(v, n, i);
917
        for (int i = n-1; i > 0; i--)
37f
             swap(v[0], v[i]), down(v, i, 0);
b33 }
3.21 Inversion Count
// Computa o numero de inversoes para transformar
// l em r (se nao tem como, retorna -1)
//
// O(n log(n))
// eef01f
37b template < typename T > 11 inv_count(vector < T > 1, vector < T > r = {}) {
bb6
        if (!r.size()) {
796
            r = 1:
1bc
             sort(r.begin(), r.end());
        }
dfb
874
        int n = 1.size();
8c0
        vector < int > v(n), bit(n);
4e9
        vector<pair<T, int>> w;
61c
        for (int i = 0; i < n; i++) w.push_back({r[i], i+1});</pre>
d1d
        sort(w.begin(), w.end());
603
        for (int i = 0; i < n; i++) {</pre>
bf3
             auto it = lower_bound(w.begin(), w.end(), make_pair(l[i],
    0));
1bf
             if (it == w.end() or it->first != 1[i]) return -1; // nao
   da
962
             v[i] = it -> second;
6c0
             it->second = -1;
964
        }
04b
        11 \text{ ans} = 0:
45b
        for (int i = n-1; i >= 0; i--) {
2d9
             for (int j = v[i]-1; j; j -= j\&-j) ans += bit[j];
```

3.22 LIS - Longest Increasing Subsequence

```
// Calcula e retorna uma LIS
//
// O(n.log(n))
// 4749e8
121 template < typename T> vector < T> lis(vector < T>& v) {
        int n = v.size(), m = -1;
f0c
        vector <T> d(n+1, INF);
aec
        vector < int > l(n);
007
        d[0] = -INF;
        for (int i = 0; i < n; i++) {</pre>
603
            // Para non-decreasing use upper_bound()
4fd
            int t = lower_bound(d.begin(), d.end(), v[i]) - d.begin();
            d[t] = v[i], l[i] = t, m = max(m, t);
3ad
89 c
4ff
        int p = n;
5a9
        vector <T> ret:
        while (p--) if (l[p] == m) {
cdf
883
            ret.push_back(v[p]);
76b
            m - - ;
f83
969
        reverse (ret.begin(), ret.end());
edf
        return ret;
474 }
```

3.23 LIS2 - Longest Increasing Subsequence

3.24 Minimum Enclosing Circle

```
// O(n) com alta probabilidade
// b0a6ba
22c const double EPS = 1e-12;
878 mt19937 rng((int)
    chrono::steady_clock::now().time_since_epoch().count());
b2a struct pt {
662
         double x, y;
        pt(double x_ = 0, double y_ = 0) : x(x_), y(y_) {}
be7
7af
         pt operator + (const pt& p) const { return pt(x+p.x, y+p.y); }
b23
         pt operator - (const pt& p) const { return pt(x-p.x, y-p.y); }
254
         pt operator * (double c) const { return pt(x*c, y*c); }
         pt operator / (double c) const { return pt(x/c, y/c); }
 701
54d };
2f9 double dot(pt p, pt q) { return p.x*q.x+p.y*q.y; }
dd5 double cross(pt p, pt q) { return p.x*q.y-p.y*q.x; }
 e7c double dist(pt p, pt q) { return sqrt(dot(p-q, p-q)); }
3f4 pt center(pt p, pt q, pt r) {
        pt a = p-r, b = q-r;
5d9
        pt c = pt(dot(a, p+r)/2, dot(b, q+r)/2);
        return pt(cross(c, pt(a.y, b.y)), cross(pt(a.x, b.x), c)) /
    cross(a, b);
fc8 }
aa8 struct circle {
f41
        pt cen;
c12
         double r;
898
         circle(pt cen_, double r_) : cen(cen_), r(r_) {}
83 c
         circle(pt a, pt b, pt c) {
             cen = center(a, b, c);
13d
1f1
            r = dist(cen. a):
cd5
         bool inside(pt p) { return dist(p, cen) < r+EPS; }</pre>
2a6 };
806 circle minCirc(vector<pt> v) {
```

```
f21
        shuffle(v.begin(), v.end(), rng);
ae0
        circle ret = circle(pt(0, 0), 0);
618
        for (int i = 0; i < v.size(); i++) if (!ret.inside(v[i])) {</pre>
16a
            ret = circle(v[i], 0);
f11
            for (int j = 0; j < i; j++) if (!ret.inside(v[j])) {</pre>
                 ret = circle((v[i]+v[j])/2, dist(v[i], v[j])/2);
881
b8c
                 for (int k = 0; k < j; k++) if (!ret.inside(v[k]))
                     ret = circle(v[i], v[i], v[k]);
43f
5f8
            }
6a1
edf
        return ret;
eba }
3.25 Minkowski Sum
// Computa A+B = \{a+b : a \setminus in A, b \setminus in B\}, em que
```

```
// A e B sao poligonos convexos
// A+B eh um poligono convexo com no max |A|+|B| pontos
//
// O(|A|+|B|)
// d7cca8
539 vector<pt> minkowski(vector<pt> p, vector<pt> q) {
        auto fix = [](vector < pt > & P) {
051
515
            rotate(P.begin(), min_element(P.begin(), P.end()),
   P.end());
018
            P.push_back(P[0]), P.push_back(P[1]);
f24
        ን:
889
        fix(p), fix(q);
8af
        vector < pt > ret;
692
       int i = 0, j = 0;
2ee
        while (i < p.size()-2 or j < q.size()-2) {</pre>
898
            ret.push_back(p[i] + q[i]);
            auto c = ((p[i+1] - p[i]) ^ (q[j+1] - q[j]));
732
           if (c >= 0) i = min<int>(i+1, p.size()-2);
ebc
81e
           if (c <= 0) j = min<int>(j+1, q.size()-2);
9ff
edf
        return ret;
d7c }
// 2f5dd2
c3e ld dist_convex(vector<pt> p, vector<pt> q) {
        for (pt& i : p) i = i * -1;
dc2
44c
        auto s = minkowski(p, q);
95d
        if (inpol(s, pt(0, 0))) return 0;
6a5
       return 1;
921
        ld ans = DINF;
```

```
073
        for (int i = 0; i < s.size(); i++) ans = min(ans,
f04
                disttoseg(pt(0, 0), line(s[(i+1)%s.size()], s[i])));
ba7
        return ans;
2f5 }
3.26 MO - DSU
// Dado uma lista de arestas de um grafo, responde
// para cada query(1, r), quantos componentes conexos
// o grafo tem se soh considerar as arestas 1, 1+1, ..., r
// Da pra adaptar pra usar MO com qualquer estrutura rollbackavel
//
// O(m sqrt(q) log(n))
// 704722
8d3 struct dsu {
553
        int n. ans:
        vector<int> p, sz;
2e3
ee6
        stack<int> S;
4b8
        dsu(int n_{-}) : n(n_{-}), ans(n), p(n), sz(n) 
            for (int i = 0; i < n; i++) p[i] = i, sz[i] = 1;
8a6
        }
aae
        int find(int k) {
1 b 1
006
            while (p[k] != k) k = p[k];
839
            return k;
c1e
        }
553
        void add(pair<int, int> x) {
700
            int a = x.first, b = x.second;
605
            a = find(a), b = find(b);
843
            if (a == b) return S.push(-1);
e7d
            ans --:
3c6
            if (sz[a] > sz[b]) swap(a, b);
4c2
            S.push(a);
582
            sz[b] += sz[a];
84b
            p[a] = b;
720
35 c
        int query() { return ans; }
5cf
        void rollback() {
465
            int u = S.top(); S.pop();
61c
            if (u == -1) return;
270
            sz[p[u]] -= sz[u]:
            p[u] = u;
546
Odf
            ans++;
456
        }
9c1 };
```

```
1a8 int n:
e93 vector <pair <int, int >> ar;
// 9d242b
617 vector < int > MO(vector < pair < int , int >> &q) {
        int SQ = sqrt(q.size()) + 1;
c23
       int m = q.size();
3f8
        vector < int > ord(m);
be8
        iota(ord.begin(), ord.end(), 0);
        sort(ord.begin(), ord.end(), [&](int 1, int r) {
d01
9c9
                 if (q[1].first / SQ != q[r].first / SQ) return
   q[l].first < q[r].first;
                return a[1].second < a[r].second:
a66
b90
                 }):
435
        vector < int > ret(m);
3bd
        dsu small(n);
        for (int i = 0; i < m; i++) {</pre>
dd5
5ec
            auto [1, r] = q[ord[i]];
            if (1 / SQ == r / SQ) {
acc
00c
                 for (int k = 1; k <= r; k++) small.add(ar[k]);</pre>
b99
                 ret[ord[i]] = small.query();
                for (int k = 1; k <= r; k++) small.rollback();</pre>
64e
259
            }
6b0
        }
dd5
        for (int i = 0; i < m; i++) {</pre>
            dsu D(n);
176
ae9
            int fim = q[ord[i]].first/SQ*SQ + SQ - 1;
e25
            int last_r = fim;
            int j = i-1;
ebc
            while (j+1 < m and g[ord[j+1]].first / SQ ==</pre>
   q[ord[i]].first / SQ) {
                 auto [1, r] = q[ord[++j]];
a0e
f58
                if (1 / SQ == r / SQ) continue;
                 while (last_r < r) D.add(ar[++last_r]);</pre>
59b
                 for (int k = 1; k <= fim; k++) D.add(ar[k]);</pre>
2cf
9b2
                ret[ord[j]] = D.query();
572
                 for (int k = 1; k <= fim; k++) D.rollback();</pre>
            }
9 c 8
bdf
            i = i;
e99
edf
        return ret;
```

```
9d2 }
```

3.27 Mo - numero de distintos em range

```
// Para ter o bound abaixo, escolher
// SQ = n / sqrt(q)
//
// O(n * sqrt(q))
// e94f60
0d2 const int MAX = 1e5+10;
6ff const int SQ = sqrt(MAX);
b69 int v[MAX];
b65 int ans, freq[MAX];
9da inline void insert(int p) {
        int o = v[p];
591
        freq[o]++;
992
        ans += (freq[o] == 1);
21d }
a25 inline void erase(int p) {
        int o = v[p]:
7ee
        ans -= (freq[o] == 1);
        freq[o]--;
ba2
dc7 }
e51 inline ll hilbert(int x, int y) {
        static int N = 1 << (__builtin_clz(0) - __builtin_clz(MAX));</pre>
71e
100
        int rx, ry, s;
b72
        11 d = 0:
43b
        for (s = N/2; s > 0; s /= 2) {
c95
            rx = (x \& s) > 0, ry = (y \& s) > 0;
e3e
            d += s * 11(s) * ((3 * rx) ^ ry);
            if (ry == 0) {
d2e
5aa
                if (rx == 1) x = N-1 - x, y = N-1 - y;
9dd
                swap(x, y);
            }
e2d
888
be2
        return d;
7fa }
bac #define HILBERT true
617 vector <int > MO(vector <pair <int, int >> &q) {
c3b
        ans = 0;
c23
        int m = q.size();
```

```
3f8
        vector < int > ord(m);
be8
        iota(ord.begin(), ord.end(), 0);
6a6 #if HILBERT
8c4
       vector<ll> h(m);
       for (int i = 0; i < m; i++) h[i] = hilbert(q[i].first,</pre>
        sort(ord.begin(), ord.end(), [&](int 1, int r) { return h[1] <</pre>
   h[r]; });
8c1 #else
        sort(ord.begin(), ord.end(), [&](int 1, int r) {
            if (q[1].first / SQ != q[r].first / SQ) return q[1].first
   < q[r].first;
            if ((q[1].first / SQ) % 2) return q[1].second >
   q[r].second;
            return q[1].second < q[r].second;</pre>
a66
bec
        }):
f2e #endif
435
        vector < int > ret(m):
        int 1 = 0, r = -1;
3d9
8b0
       for (int i : ord) {
6c6
            int al, ar;
4f5
            tie(ql, qr) = q[i];
026
            while (r < qr) insert(++r);</pre>
232
            while (1 > q1) insert(--1);
75e
            while (1 < q1) erase(1++);</pre>
           while (r > qr) erase(r--);
fe8
381
           ret[i] = ans:
        }
116
edf
        return ret;
fb7 }
3.28 Palindromic Factorization
// Precisa da eertree
// Computa o numero de formas de particionar cada
// prefixo da string em strings palindromicas
// O(n log n), considerando alfabeto O(1)
// 9e6e22
070 struct eertree { ... };
```

vector $\langle int \rangle$ diff (n+2), slink (n+2), sans (n+2), dp (n+1);

0e7 ll factorization(string s) {

eertree PT(n);

int n = s.size(), sz = 2;

b19

580

147

```
0ec
        dp[0] = 1;
78a
        for (int i = 1; i <= n; i++) {</pre>
c58
            PT.add(s[i-1]);
a7c
            if (PT.size()+2 > sz) {
6c4
                diff[sz] = PT.len[sz] - PT.len[PT.link[sz]];
                if (diff[sz] == diff[PT.link[sz]])
241
d6f
                    slink[sz] = slink[PT.link[sz]]:
f53
                else slink[sz] = PT.link[sz];
eb9
                sz++;
            }
f6a
911
            for (int v = PT.last; PT.len[v] > 0; v = slink[v]) {
297
                sans[v] = dp[i - (PT.len[slink[v]] + diff[v])];
85d
                if (diff[v] == diff[PT.link[v]])
                    sans[v] = (sans[v] + sans[PT.link[v]]) % MOD:
f20
071
                dp[i] = (dp[i] + sans[v]) % MOD;
            }
e5e
fc0
        }
5f0
        return dp[n];
3a7 }
3.29 Parsing de Expressao
// Operacoes associativas a esquerda por default
// Para mudar isso. colocar em r assoc
// Operacoes com maior prioridade sao feitas primeiro
//
// 9ad15a
cc1 bool blank(char c) {
        return c == ' ':
f34
ec3 }
8e4 bool is_unary(char c) {
        return c == '+' or c == '-';
f9c
b6b }
76d bool is_op(char c) {
        if (is_unary(c)) return true;
        return c == '*' or c == '/' or c == '+' or c == '-';
31c
4e4 }
fa3 bool r_assoc(char op) {
        // operator unario - deve ser assoc. a direita
cf0
        return op < 0;</pre>
c5c }
```

| 79d int priority(char op) {

```
// operator unario - deve ter precedencia maior
103
        if (op < 0) return INF;</pre>
        if (op == '*' or op == '/') return 2;
727
        if (op == '+' or op == '-') return 1;
439
        return -1;
daa
966 }
c15 void process_op(stack<int>& st, stack<int>& op) {
        char o = op.top(); op.pop();
88c
91c
        if (o < 0) {
4e6
           o *= -1;
1e2
            int 1 = st.top(); st.pop();
           if (o == '+') st.push(1);
Off
7e9
           if (o == '-') st.push(-1);
9d9
       } else {
14c
            int r = st.top(); st.pop();
           int 1 = st.top(); st.pop();
1e2
           if (o == '*') st.push(1 * r);
1e4
f55
           if (o == '/') st.push(1 / r);
           if (o == '+') st.push(1 + r);
605
c40
            if (o == '-') st.push(1 - r);
       }
abe
2b2 }
439 int eval(string& s) {
212
        stack<int> st, op;
d0c
        bool un = true:
1cf
        for (int i = 0; i < s.size(); i++) {</pre>
68d
            if (blank(s[i])) continue;
139
            if (s[i] == '(') {
                op.push('(');
367
99d
                un = true:
            } else if (s[i] == ')') {
130
709
                while (op.top() != '(') process_op(st, op);
75e
                op.pop();
ce2
                un = false;
146
            } else if (is_op(s[i])) {
4d0
                char o = s[i];
                if (un and is_unary(o)) o *= -1;
37c
                while (op.size() and (
ae3
                            (!r_assoc(o) and priority(op.top()) >=
cd6
   priority(o)) or
                            (r_assoc(o) and priority(op.top()) >
c41
   priority(o))))
c47
                    process_op(st, op);
```

```
c00
                op.push(o);
994
                un = true:
9d9
            } else {
                int val = 0;
da8
c2b
                while (i < s.size() and isalnum(s[i]))</pre>
8a3
                     val = val * 10 + s[i++] - '0';
169
                i--:
25d
                st.push(val);
ce2
                un = false;
            }
c7c
b19
        }
7f6
        while (op.size()) process_op(st, op);
        return st.top();
123
05c }
3.30 RMQ com Divide and Conquer
// Responde todas as queries em
// O(n log(n))
// 5a6ebd
f74 typedef pair <pair <int, int>, int> iii;
7c6 #define f first
Oab #define s second
87d int n, q, v[MAX];
e3f iii qu[MAX];
aeb int ans[MAX], pref[MAX], sulf[MAX];
0e3 void solve(int l=0, int r=n-1, int ql=0, int qr=q-1) {
8a3
        if (1 > r or q1 > qr) return;
ee4
        int m = (1+r)/2;
1b1
        int qL = partition(qu+ql, qu+qr+1, [=](iii x){return x.f.s <</pre>
   m;}) - qu;
        int qR = partition(qu+qL, qu+qr+1, [=](iii x){return x.f.f
eb0
    <=m;}) - qu;
        pref[m] = sulf[m] = v[m];
3cd
9f9
        for (int i = m-1; i >= 1; i--) pref[i] = min(v[i], pref[i+1]);
        for (int i = m+1; i <= r; i++) sulf[i] = min(v[i], sulf[i-1]);</pre>
ea8
b2a
        for (int i = qL; i < qR; i++)
f3a
            ans[qu[i].s] = min(pref[qu[i].f.f], sulf[qu[i].f.s]);
364
        solve(1, m-1, ql, qL-1), solve(m+1, r, qR, qr);
13e }
```

3.31 Segment Intersection

```
// Verifica, dado n segmentos, se existe algum par de segmentos
// que se intersecta
// O(n log n)
// 3957d8
6e0 bool operator < (const line& a, const line& b) { // comparador pro
   sweepline
        if (a.p == b.p) return ccw(a.p, a.q, b.q);
191
        if (!eq(a.p.x, a.q.x)) and (eq(b.p.x, b.q.x)) or a.p.x+eps <
   b.p.x))
780
            return ccw(a.p, a.q, b.p);
dc0
        return ccw(a.p, b.q, b.p);
e36 }
8e2 bool has_intersection(vector<line> v) {
576
        auto intersects = [&](pair<line, int> a, pair<line, int> b) {
a08
            return interseg(a.first, b.first);
3e6
e1b
        vector < pair < pt , pair < int , int >>> w;
f14
        for (int i = 0; i < v.size(); i++) {</pre>
            if (v[i].q < v[i].p) swap(v[i].p, v[i].q);</pre>
876
e1d
            w.push_back({v[i].p, {0, i}});
034
            w.push_back({v[i].q, {1, i}});
220
        sort(w.begin(), w.end());
d1d
7f2
        set < pair < line , int >> se;
e58
        for (auto i : w) {
bfd
            line at = v[i.second.second];
292
            if (i.second.first == 0) {
145
                auto nxt = se.lower_bound({at, i.second.second});
d1e
                if (nxt != se.end() and intersects(*nxt, {at,
   i.second.second})) return 1;
                if (nxt != se.begin() and intersects(*(--nxt), {at,
257
   i.second.second})) return 1;
                se.insert({at, i.second.second});
78f
9d9
                auto nxt = se.upper_bound({at, i.second.second}), cur
   = nxt, prev = --cur;
                if (nxt != se.end() and prev != se.begin()
b64
                     and intersects(*nxt, *(--prev))) return 1;
                se.erase(cur);
cca
925
            }
a00
bb3
        return 0;
```

196 }

3.32 Sequencia de de Brujin

```
// Se passar sem o terceiro parametro, gera um vetor com valores
// em [0, k) de tamanho k^n de forma que todos os subarrays ciclicos
// de tamanho n ocorrem exatamente uma vez
// Se passar com um limite lim, gera o menor vetor com valores
// em [0, k) que possui lim subarrays de tamanho n distintos
// (assume que lim <= k^n)</pre>
//
// Linear no tamanho da resposta
// 19720c
860 vector <int > de_brujin(int n, int k, int lim = INF) {
b55
        if (k == 1) return vector<int>(lim == INF ? 1 : n, 0);
5f6
        vector < int > 1 = \{0\}, ret; // 1 eh lyndon word
667
        while (true) {
c86
            if (1.size() == 0) {
1b9
                if (lim == INF) break;
daf
                1.push_back(0);
bae
686
            if (n % 1.size() == 0) for (int i : 1) {
728
                ret.push_back(i);
c99
                if (ret.size() == n+lim-1) return ret;
56e
            }
630
            int p = 1.size();
905
            while (1.size() < n) 1.push_back(1[1.size()%p]);</pre>
e7f
            while (1.size() and 1.back() == k-1) 1.pop_back();
88a
            if (1.size()) 1.back()++;
2ef
edf
        return ret:
197 }
```

3.33 Shortest Addition Chain

```
// Computa o menor numero de adicoes para construir
// cada valor, comecando com 1 (e podendo salvar variaveis)
// Retorna um par com a dp e o pai na arvore
// A arvore eh tao que o taminho da raiz (1) ate x
// contem os valores que devem ser criados para gerar x
// A profundidade de x na arvore eh dp[x]
// DP funciona para ateh 300, mas a arvore soh funciona
// para ateh 148
//
// 84fcff
```

```
// recuperacao certa soh ateh 148 (erra para 149, 233, 298)
3de pair < vector < int > , vector < int >> addition_chain() {
        int MAX = 301;
16f
875
        vector < int > dp(MAX), p(MAX);
        for (int n = 2: n < MAX: n++) {
1ab
7c0
            pair<int, int> val = {INF, -1};
212
            for (int i = 1; i < n; i++) for (int j = i; j; j = p[j])
94a
                if (j == n-i) val = min(val, pair(dp[i]+1, i));
eb3
            tie(dp[n], p[n]) = val;
            if (n == 9) p[n] = 8;
efe
ba1
            if (n == 149 \text{ or } n == 233) \text{ dp}[n] --;
bcd
717
        return {dp, p};
84f }
3.34 Simple Polygon
// Verifica se um poligono com n pontos eh simples
//
// O(n log n)
// c724a4
6e0 bool operator < (const line& a, const line& b) { // comparador pro
   sweepline
        if (a.p == b.p) return ccw(a.p, a.q, b.q);
191
        if (!eq(a.p.x, a.q.x)) and (eq(b.p.x, b.q.x)) or a.p.x+eps <
231
   b.p.x))
780
            return ccw(a.p, a.q, b.p);
dc0
        return ccw(a.p, b.q, b.p);
e36 }
6f3 bool simple(vector<pt> v) {
576
        auto intersects = [&](pair<line, int> a, pair<line, int> b) {
e72
            if ((a.second+1)%v.size() == b.second or
80e
                 (b.second+1)%v.size() == a.second) return false;
            return interseg(a.first, b.first);
a08
1c5
        };
        vector<line> seg;
41a
e1b
        vector<pair<pt, pair<int, int>>> w;
f14
        for (int i = 0: i < v.size(): i++) {</pre>
```

pt at = v[i], nxt = v[(i+1)%v.size()];

if (isinseg(v[(i+2)%v.size()], line(at, nxt))) return 0;

if (nxt < at) swap(at, nxt);</pre>

w.push_back({at, {0, i}});

w.push_back({nxt, {1, i}});

seg.push_back(line(at, nxt));

// casos degenerados estranhos

0a8 828

937

f7e

69 c

ae8

```
88d
            if (isinseg(v[(i+v.size()-1)%v.size()], line(at, nxt)))
    return 0:
        }
cba
d1d
        sort(w.begin(), w.end());
        set<pair<line, int>> se;
7f2
        for (auto i : w) {
e58
ff8
            line at = seg[i.second.second];
292
            if (i.second.first == 0) {
145
                auto nxt = se.lower_bound({at, i.second.second});
7 c.4
                if (nxt != se.end() and intersects(*nxt, {at,
    i.second.second})) return 0;
                if (nxt != se.begin() and intersects(*(--nxt), {at,
b34
    i.second.second})) return 0:
                 se.insert({at, i.second.second});
78f
949
            } else {
884
                 auto nxt = se.upper_bound({at, i.second.second}), cur
    = nxt, prev = --cur;
                if (nxt != se.end() and prev != se.begin()
b64
                     and intersects(*nxt, *(--prev))) return 0;
403
                se.erase(cur):
cca
            }
ad0
d17
        }
6a5
        return 1;
af3 }
3.35 Sweep Direction
// Passa por todas as ordenacoes dos pontos definitas por "direcoes"
// Assume que nao existem pontos coincidentes
//
// O(n^2 \log n)
// 6bb68d
4b8 void sweep_direction(vector<pt> v) {
3d2
        int n = v.size();
        sort(v.begin(), v.end(), [](pt a, pt b) {
163
3a5
            if (a.x != b.x) return a.x < b.x;
572
            return a.v > b.v;
79a
        });
        vector < int > at(n);
b89
516
        iota(at.begin(), at.end(), 0);
b79
        vector<pair<int, int>> swapp;
        for (int i = 0; i < n; i++) for (int j = i+1; j < n; j++)
25e
95f
            swapp.push_back({i, j}), swapp.push_back({j, i});
269
        sort(swapp.begin(), swapp.end(), [&](auto a, auto b) {
134
            pt A = rotate90(v[a.first] - v[a.second]);
```

```
247
            pt B = rotate90(v[b.first] - v[b.second]);
            if (quad(A) == quad(B) and !sarea2(pt(0, 0), A, B)) return
615
   a < b;
224
            return compare_angle(A, B);
5e7
       }):
        for (auto par : swapp) {
4e6
e24
            assert(abs(at[par.first] - at[par.second]) == 1);
a96
            int 1 = min(at[par.first], at[par.second]),
0d3
                r = n-1 - max(at[par.first], at[par.second]);
            // l e r sao quantos caras tem de cada lado do par de
               pontos
            // (cada par eh visitado duas vezes)
9cf
            swap(v[at[par.first]], v[at[par.second]]);
            swap(at[par.first], at[par.second]);
1c0
241
       }
6bb }
```

3.36 Triangulação de Delaunay

```
// Computa a triangulação de Delaunay, o dual
// do diagrama de Voronoi (a menos de casos degenerados)
// Retorna um grafo indexado pelos indices dos pontos, e as arestas
// sao as arestas da triangulação
// As arestas partindo de um vertice ja vem ordenadas por angulo,
// ou seja, se o vertice v nao esta no convex hull, (v, v_i, v_{i+1})
// eh um triangulo da triangulacao, em que v_i eh o i-esimo vizinho
// Usa o alg d&c. precisa representar MAX COOR^4. por isso int128
// pra aguentar valores ateh 1e9
//
// Propriedades:
// 1 - 0 grafo tem no max 3n-6 arestas
// 2 - Para todo triangulo, a circunf. que passa pelos 3 pontos
// nao contem estritamente nenhum ponto
// 3 - A MST euclidiana eh subgrafo desse grafo
// 4 - Cada ponto eh vizinho do ponto mais proximo dele
//
// O(n log n)
// 362c83
2ad typedef struct QuadEdge* Q;
ba5 struct QuadEdge {
53e
        int id:
       pt o;
       Q rot, nxt;
41e
3e5
       bool used;
3fc
        QuadEdge(int id_ = -1, pt o_ = pt(INF, INF)) :
```

```
4ba
            id(id_), o(o_), rot(nullptr), nxt(nullptr), used(false) {}
        Q rev() const { return rot->rot; }
00f
        Q next() const { return nxt; }
сЗс
188
        Q prev() const { return rot->next()->rot; }
        pt dest() const { return rev()->o; }
0d4
828 }:
91b Q edge(pt from, pt to, int id_from, int id_to) {
        Q e1 = new QuadEdge(id_from, from);
61b
        Q e2 = new QuadEdge(id_to, to);
8f6
        Q e3 = new QuadEdge;
5ca
        Q e4 = new QuadEdge:
        tie(e1->rot, e2->rot, e3->rot, e4->rot) = \{e3, e4, e2, e1\};
e69
f22
        tie(e1->nxt, e2->nxt, e3->nxt, e4->nxt) = \{e1, e2, e4, e3\};
        return e1:
1ad
c70 }
d8d void splice(Q a, Q b) {
        swap(a->nxt->rot->nxt, b->nxt->rot->nxt);
a6f
da4
        swap(a->nxt, b->nxt);
a58 }
167 void del_edge(Q& e, Q ne) { // delete e and assign e <- ne
cc0
        splice(e, e->prev());
eec
        splice(e->rev(). e->rev()->prev()):
        delete e->rev()->rot, delete e->rev();
7ea
524
        delete e->rot: delete e:
6b2
        e = ne;
18b }
d08 Q conn(Q a, Q b) {
        Q = edge(a->dest(), b->o, a->rev()->id, b->id);
f2b
        splice(e. a->rev()->prev()):
        splice(e->rev(), b);
d37
6bf
        return e;
f78 }
d64 bool in_c(pt a, pt b, pt c, pt p) { // p ta na circunf. (a, b, c) ?
268
        \_int128 p2 = p*p, A = a*a - p2, B = b*b - p2, C = c*c - p2;
        return sarea2(p, a, b) * C + sarea2(p, b, c) * A + sarea2(p,
   c, a) * B > 0;
b54 }
540 pair < Q, Q > build_tr(vector < pt > & p, int 1, int r) {
09d
        if (r-1+1 \le 3) {
2eb
            Q = edge(p[1], p[1+1], 1, 1+1), b = edge(p[1+1], p[r],
```

```
1+1, r);
912
            if (r-1+1 == 2) return \{a, a->rev()\};
            splice(a->rev(), b);
0ec
            ll ar = sarea2(p[1], p[1+1], p[r]);
сЗс
            Q c = ar ? conn(b, a) : 0;
1af
            if (ar >= 0) return \{a, b > rev()\};
021
9db
            return {c->rev(), c};
        }
bce
        int m = (1+r)/2;
ee4
328
        auto [la, ra] = build_tr(p, 1, m);
b93
        auto [lb, rb] = build_tr(p, m+1, r);
667
        while (true) {
b99
            if (ccw(lb->o, ra->o, ra->dest())) ra = ra->rev()->prev();
            else if (ccw(lb->o, ra->o, lb->dest())) lb =
   lb->rev()->next();
            else break;
f97
24a
        Q b = conn(lb->rev(), ra);
ca5
        auto valid = [&](Q e) { return ccw(e->dest(), b->dest(),
713
   b -> o): }:
        if (ra->o == la->o) la = b->rev();
        if (1b->o == rb->o) rb = b;
63f
        while (true) {
667
            Q L = b \rightarrow rev() \rightarrow next();
71e
            if (valid(L)) while (in_c(b->dest(), b->o, L->dest(),
   L->next()->dest())
1c0
                 del_edge(L, L->next());
c76
            Q R = b - > prev();
            if (valid(R)) while (in_c(b->dest(), b->o, R->dest(),
   R->prev()->dest()))
541
                 del_edge(R, R->prev());
            if (!valid(L) and !valid(R)) break;
a3a
            if (!valid(L) or (valid(R) and in_c(L->dest(), L->o, R->o,
   R->dest())))
                b = conn(R, b \rightarrow rev());
36c
666
            else b = conn(b->rev(), L->rev());
94d
        }
a2b
        return {la, rb};
689 }
b58 vector < vector < int >> delaunay (vector < pt > v) {
3d2
        int n = v.size();
397
        auto tmp = v;
135
        vector < int > idx(n);
295
        iota(idx.begin(), idx.end(), 0);
        sort(idx.begin(), idx.end(), [&](int 1, int r) { return v[1] <</pre>
   v[r]; });
```

```
5d8
        for (int i = 0; i < n; i++) v[i] = tmp[idx[i]];</pre>
780
        assert(unique(v.begin(), v.end()) == v.end());
        vector < vector < int >> g(n);
4aa
4ec
        bool col = true;
        for (int i = 2; i < n; i++) if (sarea2(v[i], v[i-1], v[i-2]))
   col = false;
bf5
        if (col) {
aa4
            for (int i = 1; i < n; i++)</pre>
839
                g[idx[i-1]].push_back(idx[i]),
   g[idx[i]].push_back(idx[i-1]);
96b
            return g;
        }
0ae
d36
        Q e = build_tr(v, 0, n-1).first;
113
        vector<Q> edg = {e};
        for (int i = 0; i < edg.size(); e = edg[i++]) {</pre>
5d1
            for (Q at = e; !at->used; at = at->next()) {
3ed
60d
                 at->used = true;
                g[idx[at->id]].push_back(idx[at->rev()->id]);
cf8
                 edg.push_back(at->rev());
15d
9f2
            }
        }
d19
96b
        return g;
b43 }
```

3.37 Triangulos em Grafos

```
// get_triangles(i) encontra todos os triangulos ijk no grafo
// Custo nas arestas
// retorna {custo do triangulo, {j, k}}
// O(m sqrt(m) log(n)) se chamar para todos os vertices
// fladbc
c0d vector<pair<int, int>> g[MAX]; // {para, peso}
d41 #warning o 'g' deve estar ordenado
9a5 vector<pair<int, pair<int, int>>> get_triangles(int i) {
771
        vector<pair<int, pair<int, int>>> tri;
        for (pair<int, int> j : g[i]) {
b23
2b3
            int a = i, b = j.first;
            if (g[a].size() > g[b].size()) swap(a, b);
6dd
            for (pair < int, int > c : g[a]) if (c.first != b and c.first
eb0
   > j.first) {
                auto it = lower_bound(g[b].begin(), g[b].end(),
525
   make_pair(c.first, -INF));
                if (it == g[b].end() or it->first != c.first) continue;
f55
0aa
                tri.push_back({j.second+c.second+it->second, {a == i ?
```

```
b : a, c.first}});
b5e       }
7e1    }
f5e    return tri;
036 }
```

4 Matematica

4.1 2-SAT

```
// solve() retorna um par, o first fala se eh possivel
// atribuir, o second fala se cada variavel eh verdadeira
// O(|V|+|E|) = O(\#variaveis + \#restricoes)
// ef6b3b
138 struct sat {
e6c
        int n. tot:
789
        vector < vector < int >> g;
0ca
       vector < int > vis, comp, id, ans;
4ce
        stack<int> s:
141
        sat() {}
172
        sat(int n_{-}) : n(n_{-}), tot(n), g(2*n) {}
f32
        int dfs(int i. int& t) {
            int lo = id[i] = t++;
cf0
efc
            s.push(i), vis[i] = 2;
48e
            for (int j : g[i]) {
740
                if (!vis[j]) lo = min(lo, dfs(j, t));
994
                else if (vis[j] == 2) lo = min(lo, id[j]);
d64
3de
            if (lo == id[i]) while (1) {
3c3
                int u = s.top(); s.pop();
9c5
                vis[u] = 1, comp[u] = i;
                if ((u>1) < n \text{ and } ans[u>1] == -1) ans[u>1] = \sim u\&1:
91d
2ef
                if (u == i) break;
60d
            }
253
            return lo:
dec
74a
        void add_impl(int x, int y) { // x -> y = !x ou y
26a
            x = x >= 0 ? 2*x : -2*x-1;
2b8
            y = y >= 0 ? 2*y : -2*y-1;
a1e
            g[x].push_back(y);
1e2
            g[y^1].push_back(x^1);
```

```
ef0
e85
         void add_cl(int x, int y) { // x ou y
0b5
             add_impl(\sim x, y);
254
        }
         void add_xor(int x, int y) { // x xor y
487
             add_cl(x, y), add_cl(\sim x, \sim y);
0b7
9a1
        }
978
         void add_eq(int x, int y) { // x = y
c86
             add_xor(\simx, y);
        }
b91
b10
         void add_true(int x) { // x = T
18b
             add_impl(\sim x, x);
9e2
        }
d14
         void at_most_one(vector<int> v) { // no max um verdadeiro
54d
             g.resize(2*(tot+v.size()));
f14
             for (int i = 0; i < v.size(); i++) {</pre>
8c9
                 add_impl(tot+i, \sim v[i]);
a8f
                 if (i) {
b6a
                     add_impl(tot+i, tot+i-1);
3d3
                     add_impl(v[i], tot+i-1);
                 }
0f7
084
258
             tot += v.size();
b00
        }
         pair < bool, vector < int >> solve() {
a8e
27b
             ans = vector < int > (n, -1);
             int t = 0;
6bb
0de
             vis = comp = id = vector\langle int \rangle (2*tot, 0);
53c
             for (int i = 0; i < 2*tot; i++) if (!vis[i]) dfs(i, t);</pre>
             for (int i = 0; i < tot; i++)</pre>
f88
4c9
                 if (comp[2*i] == comp[2*i+1]) return {false, {}};
997
             return {true, ans};
7b3
        }
ef6 };
4.2 Algoritmo de Euclides estendido
// Acha x e y tal que ax + by = mdc(a, b) (nao eh unico)
// Assume a, b >= 0
//
// O(log(min(a, b)))
// 35411d
2be tuple < 11, 11, 11 > ext_gcd(11 a, 11 b) {
3bd
        if (!a) return {b, 0, 1};
550
         auto [g, x, y] = ext_gcd(b%a, a);
```

```
c59
        return \{g, y - b/a*x, x\};
                                                                            9ee
354 }
                                                                            d09
                                                                            564
4.3 Avaliação de Interpolação
                                                                            d75
                                                                               b.size(); j++)
// Dado 'n' pontos (i, v[i]), i \in [0, n),
                                                                            cff
// avalia o polinomio de grau n-1 que passa
// por esses pontos em 'x'
                                                                                i >= 0; i--)
// Tudo modular, precisa do mint
                                                                            112
                                                                            16d
// O(n)
// 4fe929
                                                                            3ъ9
                                                                                    }:
ee8 mint evaluate_interpolation(int x, vector<mint> y) {
                                                                            1a6
        int n = v.size();
80e
                                                                               x = \{1\}:
                                                                            95f
                                                                                    while (k) {
184
        vector \leq mint > sulf(n+1, 1), fat(n, 1), ifat(n);
                                                                            7f1
        for (int i = n-1; i \ge 0; i--) sulf[i] = sulf[i+1] * (x - i);
6fa
                                                                            b28
29b
        for (int i = 1: i < n: i++) fat[i] = fat[i-1] * i:
                                                                                    }
                                                                            8ea
        ifat[n-1] = 1/fat[n-1];
0da
                                                                            dd6
                                                                                    x.resize(n);
3db
        for (int i = n-2; i >= 0; i--) ifat[i] = ifat[i+1] * (i + 1);
                                                                                    T ret = 0:
                                                                            ce8
ca1
        mint pref = 1, ans = 0;
                                                                            e72
        for (int i = 0; i < n; pref *= (x - i++)) {</pre>
5ea
                                                                                    return ret:
                                                                            edf
42f
            mint num = pref * sulf[i+1];
                                                                            7e2 }
b4e
            mint den = ifat[i] * ifat[n-1 - i];
0bd
            if ((n-1 - i)\%2) den *= -1:
                                                                            222
03f
            ans += v[i] * num * den;
                                                                            46e
ce6
                                                                            620
ba7
        return ans:
                                                                            793
4fe }
                                                                            ab6
                                                                            5f0
4.4 Berlekamp-Massey
                                                                            8b4
                                                                            369
// guess_kth(s, k) chuta o k-esimo (0-based) termo
                                                                            ba6
// de uma recorrencia linear que gera s
                                                                            88f
// Para uma rec. lin. de ordem x, se passar 2x termos
                                                                            76a
// vai gerar a certa
                                                                            90c
// Usar aritmetica modular
                                                                            844
//
                                                                            0dc
// O(n^2 log k), em que n = |s|
                                                                            807
                                                                                    return c:
// 8644e3
                                                                            4d9 }
b7c template < typename T > T evaluate (vector < T > c, vector < T > s, ll k) {
```

ff2

int n = c.size();

```
assert(c.size() <= s.size());</pre>
        auto mul = [&](const vector<T> &a, const vector<T> &b) {
            vector <T> ret(a.size() + b.size() - 1);
            for (int i = 0; i < a.size(); i++) for (int j = 0; j <</pre>
                ret[i+j] += a[i] * b[j];
            for (int i = ret.size()-1; i >= n; i--) for (int j = n-1;
                ret[i-j-1] += ret[i] * c[j];
            ret.resize(min<int>(ret.size(), n));
            return ret:
        vector < T > a = n == 1 ? vector < T > ({c[0]}) : vector < T > ({0, 1}),
            if (k\&1) x = mul(x, a):
            a = mul(a, a), k >>= 1;
        for (int i = 0; i < n; i++) ret += x[i] * s[i];
192 template < typename T > vector < T > berlekamp_massey(vector < T > s) {
        int n = s.size(), 1 = 0, m = 1;
        vector < T > b(n), c(n);
        T ld = b[0] = c[0] = 1;
        for (int i = 0; i < n; i++, m++) {</pre>
            T d = s[i]:
            for (int j = 1; j <= 1; j++) d += c[j] * s[i-j];
            if (d == 0) continue;
            vector <T> temp = c;
            T coef = d / ld;
            for (int j = m; j < n; j++) c[j] -= coef * b[j-m];
            if (2 * 1 \le i) 1 = i + 1 - 1, b = temp, 1d = d, m = 0;
        c.resize(1 + 1):
        c.erase(c.begin());
        for (T\& x : c) x = -x;
2cf template < typename T > T guess_kth(const vector < T > & s, ll k) {
```

```
cc3     auto c = berlekamp_massey(s);
96a     return evaluate(c, s, k);
697 }
```

4.5 Binomial Distribution

4.6 Convolucao de GCD / LCM

```
// O(n log(n))
// multiple_transform(a)[i] = \sum_d a[d * i]
// 338be8
bbe template < typename T > void multiple_transform (vector < T > & v, bool
   inv = false) {
        vector < int > I(v.size()-1);
64a
       iota(I.begin(), I.end(), 1);
674
       if (inv) reverse(I.begin(), I.end());
        for (int i : I) for (int j = 2; i*j < v.size(); j++)
            v[i] += (inv ? -1 : 1) * v[i*j];
a8a
338 }
// \gcd_{convolution(a, b)[k]} = \sum_{gcd(i, j)} = k} a_i * b_j
fe2 template < typename T > vector < T > gcd_convolution(vector < T > a,
   vector <T> b) {
        multiple_transform(a), multiple_transform(b);
799
        for (int i = 0; i < a.size(); i++) a[i] *= b[i];
dea
        multiple_transform(a, true);
3f5
        return a;
984 }
```

```
// divisor transform(a)[i] = \sum {d|i} a[i/d]
// aa74e5
be7 template < typename T > void divisor_transform (vector < T > & v, bool inv
   = false) {
        vector < int > I(v.size()-1);
64a
847
        iota(I.begin(), I.end(), 1);
        if (!inv) reverse(I.begin(), I.end());
5ea
        for (int i : I) for (int j = 2; i*j < v.size(); j++)</pre>
dad
            v[i*j] += (inv ? -1 : 1) * v[i];
14f
aa7 }
// lcm convolution(a, b)[k] = \sum \{lcm(i, i) = k\} a i * b i
// f5acc1
b1b template < typename T > vector < T > lcm_convolution (vector < T > a,
    vector<T> b) {
3af
        divisor_transform(a), divisor_transform(b);
        for (int i = 0; i < a.size(); i++) a[i] *= b[i];</pre>
799
        divisor_transform(a, true);
d8f
3f5
        return a:
f5a }
4.7 Coprime Basis
// Dado um conjunto de elementos A constroi uma base B
// de fatores coprimos tal que todo elemento A[i]
```

```
// pode ser fatorado como A[i] = \prod B[j]^p_ij
// Sendo n o numero de inserts, a complexidade esperada fica
// O(n*(n*loglog(MAX) + log(MAX)^2))
// No pior caso, podemos trocar n*loglog(MAX) por
// se MAX <= 1e6 fica 8*n
// se MAX <= 1e9 fica 10*n
// se MAX <= 1e18 fica 16*n
// se MAX <= 1e36 fica 26*n
// 6714d3
ebc template <typename T> struct coprime_basis {
        vector<T> basis:
60e
        coprime_basis() {}
        coprime_basis(vector<T> v) { for (T i : v) insert(i); }
055
845
        void insert(T z) {
сЗс
            int n = basis.size();
```

```
efe
            basis.push_back(z);
            for (int i = n; i < basis.size(); i++) {</pre>
43c
                for (int j = (i != n) ? i+1 : 0; j < basis.size();</pre>
21c
   j++) {
4ce
                     if (i == j) continue;
                     T &x = basis[i];
024
c91
                     if (x == 1) {
                         j = INF;
5e2
                         continue;
                     }
6e0
544
                    T \& v = basis[i];
3c9
                    T g = gcd(x, y);
e10
                    if (g == 1) continue;
15b
                     y /= g, x /= g;
8c6
                     basis.push_back(g);
069
                }
422
            }
            basis.erase(remove(basis.begin(), basis.end(), 1),
   basis.end());
       }
1a5
        vector < int > factor(T x) {
4ba
21d
            vector < int > fat(basis.size());
6fd
            for (int i = 0; i < basis.size(); i++) {</pre>
25 c
                 while (x \% basis[i] == 0) x /= basis[i], fat[i]++;
8de
            }
6a7
            return fat;
b5d
        }
671 };
4.8 Crivo de Eratosthenes
```

```
// "0" crivo
//
// Encontra maior divisor primo
// Um numero eh primo sse divi[x] == x
// fact fatora um numero <= lim
// A fatoracao sai ordenada
//
// crivo - O(n log(log(n)))
// fact - O(log(n))

// hash (crivo e fact): def8f3
f12 int divi[MAX];

fb9 void crivo(int lim) {
f53  for (int i = 1; i <= lim; i++) divi[i] = 1;</pre>
```

```
for (int i = 2; i <= lim; i++) if (divi[i] == 1)</pre>
d46
018
            for (int j = i; j <= lim; j += i) divi[j] = i;</pre>
349 }
470 void fact(vector<int>& v, int n) {
        if (n != divi[n]) fact(v, n/divi[n]);
        v.push_back(divi[n]);
ab4
1db }
// Crivo linear
// Mesma coisa que o de cima, mas tambem
// calcula a lista de primos
//
// O(n)
// 792458
f12 int divi[MAX];
fd3 vector<int> primes;
fb9 void crivo(int lim) {
        divi[1] = 1:
d5a
f70
        for (int i = 2; i <= lim; i++) {</pre>
3eb
            if (divi[i] == 0) divi[i] = i, primes.push_back(i);
3ba
            for (int j : primes) {
522
                 if (j > divi[i] or i*j > lim) break;
00b
                 divi[i*j] = j;
491
            }
85a
        }
519 }
// Crivo de divisores
//
// Encontra numero de divisores
// ou soma dos divisores
//
// O(n log(n))
// 9bf7b6
f12 int divi[MAX];
fb9 void crivo(int lim) {
        for (int i = 1; i <= lim; i++) divi[i] = 1;</pre>
424
        for (int i = 2; i <= lim; i++)</pre>
594
            for (int j = i; j <= lim; j += i) {</pre>
```

```
// para numero de divisores
9e0
                divi[j]++;
                // para soma dos divisores
278
                divi[i] += i;
            }
c58
fc1 }
// Crivo de totiente
// Encontra o valor da funcao
// totiente de Euler
// O(n log(log(n)))
// 266461
5f4 int tot[MAX];
fb9 void crivo(int lim) {
        for (int i = 1; i <= lim; i++) {</pre>
a27
bc9
            tot[i] += i:
            for (int j = 2*i; j <= lim; j += i)</pre>
                tot[i] -= tot[i];
837
       }
678
212 }
// Crivo de funcao de mobius
//
// O(n log(log(n)))
// 58d036
4e1 char meb[MAX];
fb9 void crivo(int lim) {
649
        for (int i = 2; i <= lim; i++) meb[i] = 2;</pre>
        meb[1] = 1:
ace
842
        for (int i = 2; i <= lim; i++) if (meb[i] == 2)
            for (int j = i; j <= lim; j += i) if (meb[j]) {</pre>
848
686
                if (meb[j] == 2) meb[j] = 1;
ae1
                meb[j] *= j/i\%i ? -1 : 0;
97f
            }
9bc }
// Crivo linear de funcao multiplicativa
// Computa f(i) para todo 1 <= i <= n, sendo f
// uma funcao multiplicativa (se gcd(a,b) = 1,
// entao f(a*b) = f(a)*f(b)
```

```
// f_prime tem que computar f de um primo, e
// add_prime tem que computar f(p^(k+1)) dado f(p^k) e p
// Se quiser computar f(p^k) dado p e k, usar os comentarios
//
// O(n)
// 66886a
fd3 vector<int> primes;
623 int f[MAX], pot[MAX];
//int expo[MAX];
5c4 void sieve(int lim) {
         // Funcoes para soma dos divisores:
fc9
         auto f_prime = [](int p) { return p+1; };
31c
         auto add_prime = [](int fpak, int p) { return fpak*p+1; };
         //auto f_pak = [](int p, int k) {};
02d
        f[1] = 1;
f70
         for (int i = 2; i <= lim; i++) {</pre>
e6b
             if (!pot[i]) {
e74
                 primes.push_back(i);
f05
                 f[i] = f_prime(i), pot[i] = i;
                 //\exp[i] = 1;
b71
            }
3ъ9
            for (int p : primes) {
                 if (i*p > lim) break;
b9f
569
                 if (i%p == 0) {
b97
                     f[i*p] = f[i / pot[i]] * add_prime(f[pot[i]], p);
                     // se for descomentar, tirar a linha de cima tambem
                     //f[i*p] = f[i / pot[i]] * f_pak(p, expo[i]+1);
                     //\expo[i*p] = \expo[i]+1;
51f
                     pot[i*p] = pot[i] * p;
c2b
                     break:
9d9
                 } else {
9ef
                     f[i*p] = f[i] * f[p];
638
                     pot[i*p] = p;
                     //\expo[i*p] = 1;
e92
                }
f31
             }
         }
1bb
350 }
4.9 Deteccao de ciclo - Tortoise and Hare
// Linear no tanto que tem que andar pra ciclar,
// O(1) de memoria
```

```
// Retorna um par com o tanto que tem que andar
// do fO ate o inicio do ciclo e o tam do ciclo
// 899f20
58d pair<11, 11> find_cycle() {
       11 \text{ tort} = f(f0);
b2b
       ll hare = f(f(f0)):
b1b
       11 t = 0;
683
       while (tort != hare) {
            tort = f(tort);
b4d
4b2
           hare = f(f(hare));
c82
           t++:
93d
       }
0e8
       11 st = 0;
909
       tort = f0;
683
        while (tort != hare) {
b4d
            tort = f(tort);
1a2
           hare = f(hare);
397
            st++;
c.91
       }
       ll len = 1;
73d
3cd
       hare = f(tort):
683
        while (tort != hare) {
1a2
            hare = f(hare);
040
            len++:
f1a
ebd
       return {st, len};
899 }
4.10 Division Trick
// Gera o conjunto n/i, pra todo i, em O(sqrt(n))
// copiei do github do tfg50
// 5bf9bf
79c for(int l = 1, r; l \le n; l = r + 1) {
       r = n / (n / 1);
       // n / i has the same value for 1 <= i <= r
5bf }
4.11 Eliminacao Gaussiana
// Resolve sistema linear
// Retornar um par com o numero de solucoes
// e alguma solucao, caso exista
//
```

```
// O(n^2 * m)
 // 1d10b5
67a template < typename T>
 728 pair<int, vector<T>> gauss(vector<vector<T>> a, vector<T> b) {
         const double eps = 1e-6;
f92
         int n = a.size(), m = a[0].size();
2f0
         for (int i = 0; i < n; i++) a[i].push_back(b[i]);</pre>
 3cb
         vector < int > where (m, -1);
 237
         for (int col = 0, row = 0; col < m and row < n; col++) {
f05
             int sel = row:
b95
             for (int i=row: i<n: ++i)</pre>
 e55
                  if (abs(a[i][col]) > abs(a[sel][col])) sel = i;
2c4
             if (abs(a[sel][col]) < eps) continue;</pre>
             for (int i = col; i <= m; i++)</pre>
1ae
 dd2
                  swap(a[sel][i], a[row][i]);
2c3
             where [col] = row;
             for (int i = 0; i < n; i++) if (i != row) {</pre>
0c0
                 T c = a[i][col] / a[row][col]:
96c
 d5c
                 for (int j = col; j <= m; j++)</pre>
 c8f
                      a[i][j] -= a[row][j] * c;
 490
             }
 b70
             row++;
 3d8
         }
b1d
         vector <T> ans(m. 0):
         for (int i = 0; i < m; i++) if (where[i] != -1)</pre>
 e1a
 12a
             ans[i] = a[where[i]][m] / a[where[i]][i];
         for (int i = 0; i < n; i++) {</pre>
 603
 501
             T sum = 0;
 a75
             for (int j = 0; j < m; j++)
5a9
                  sum += ans[i] * a[i][i]:
 b1f
             if (abs(sum - a[i][m]) > eps)
 6cd
                  return pair(0, vector<T>());
 ec9
         }
12e
         for (int i = 0; i < m; i++) if (where[i] == -1)</pre>
 018
             return pair(INF, ans);
 280
         return pair(1, ans);
 292 }
4.12 Eliminação Gaussiana Z2
// D eh dimensao do espaco vetorial
// add(v) - adiciona o vetor v na base (retorna se ele jah pertencia
```

```
ao span da base)
// coord(v) - retorna as coordenadas (c) de v na base atual (basis^T.c
// recover(v) - retorna as coordenadas de v nos vetores na ordem em
   que foram inseridos
// coord(v).first e recover(v).first - se v pertence ao span
// Complexidade:
// add, coord, recover: O(D^2 / 64)
// d0a4b3
2a3 template <int D> struct Gauss_z2 {
        bitset <D> basis[D]. keep[D];
b16
        int rk, in;
482
        vector < int > id;
37f
        Gauss_z2 () : rk(0), in(-1), id(D, -1) {};
        bool add(bitset < D > v) {
04e
42c
            in++:
fb0
            bitset <D> k;
659
            for (int i = D - 1; i \ge 0; i--) if (v[i]) {
189
                if (basis[i][i]) v ^= basis[i], k ^= keep[i];
                else {
4e6
                    k[i] = true, id[i] = in, keep[i] = k;
ea6
6ce
                    basis[i] = v, rk++;
8a6
                    return true;
                }
b34
            }
09c
d1f
            return false;
58b
0f6
        pair < bool, bitset < D >> coord(bitset < D > v) {
944
            bitset <D> c:
659
            for (int i = D - 1; i \ge 0; i - -) if (v[i]) {
                if (basis[i][i]) v ^= basis[i], c[i] = true;
a39
8af
                else return {false, bitset <D>()};
            }
a08
5db
            return {true, c};
a08
330
        pair < bool, vector < int >> recover(bitset < D > v) {
            auto [span, bc] = coord(v);
22e
            if (not span) return {false, {}};
af8
f79
            bitset <D> aux:
            for (int i = D - 1; i >= 0; i--) if (bc[i]) aux ^= keep[i];
5a0
ea9
            vector < int > oc;
            for (int i = D - 1; i >= 0; i--) if (aux[i])
   oc.push_back(id[i]);
```

```
001          return {true, oc};
b75    }
d0a };
```

4.13 Equação Diofantina Linear

```
// Encontra o numero de solucoes de a*x + b*y = c,
// em que x \in [lx, rx] e y \in [ly, ry]
// Usar o comentario para recuperar as solucoes
// (note que o b ao final eh b/gcd(a, b))
// Cuidado com overflow! Tem que caber o quadrado dos valores
// O(log(min(a, b)))
// 2e8259
c5e template < typename T > tuple < ll, T, T > ext_gcd(ll a, ll b) {
        if (!a) return {b, 0, 1};
3bd
c4b
        auto [g, x, y] = ext_gcd < T > (b%a, a);
c59
        return \{g, y - b/a*x, x\};
8a8 }
// numero de solucoes de a*[lx, rx] + b*[ly, ry] = c
14c template < typename T = 11> // usar __int128 se for ate 1e18
2a4 ll diophantine(ll a, ll b, ll c, ll lx, ll rx, ll ly, ll ry) {
c80
        if (1x > rx \text{ or } 1y > ry) \text{ return } 0;
        if (a == 0 \text{ and } b == 0) \text{ return } c ? 0 : (rx-lx+1)*(ry-ly+1);
a98
8ce
        auto [g, x, y] = ext_gcd < T > (abs(a), abs(b));
        if (c % g != 0) return 0;
9 c 3
249
        if (a == 0) return (rx-lx+1)*(ly <= c/b and c/b <= ry);
        if (b == 0) return (ry-ly+1)*(lx <= c/a and c/a <= rx);
4ce
fb1
        x *= a/abs(a) * c/g, v *= b/abs(b) * c/g, a /= g, b /= g;
b20
        auto shift = [\&](T qt) \{ x += qt*b, y -= qt*a; \};
efa
        auto test = [&](T& k, 11 mi, 11 ma, 11 coef, int t) {
866
            shift((mi - k)*t / coef);
79d
            if (k < mi) shift(coef > 0 ? t : -t);
74d
            if (k > ma) return pair<T, T>(rx+2, rx+1);
            T x1 = x;
41f
            shift((ma - k)*t / coef);
633
            if (k > ma) shift(coef > 0 ? -t : t);
c5b
4a9
            return pair <T, T > (x1, x);
8e1
        }:
639
        auto [11, r1] = test(x, lx, rx, b, 1);
38e
        auto [12, r2] = test(y, ly, ry, a, -1);
c43
        if (12 > r2) swap(12, r2);
        T l = max(11, 12), r = min(r1, r2);
50a
```

```
339 if (1 > r) return 0;

42f ll k = (r-l) / abs(b) + 1;

839 return k; // solucoes: x = 1 + [0, k)*|b|

98e }
```

4.14 Exponenciacao rapida

```
// (x^y mod m) em O(log(y))
// 12b2f8
03c ll pow(ll x, ll y, ll m) \{ // \text{ iterativo} \}
c85
       ll ret = 1;
        while (v) {
           if (y & 1) ret = (ret * x) % m;
895
23b
          y >>= 1;
cc5
            x = (x * x) % m;
020
        }
edf
        return ret;
12b }
// 7d427b
03c ll pow(ll x, ll y, ll m) { // recursivo
        if (!y) return 1;
13a
        ll ans = pow(x*x\%m, y/2, m);
426
88d
        return v%2 ? x*ans%m : ans;
7d4 }
```

4.15 Fast Walsh Hadamard Transform

```
// FWHT<',|'>(f) eh SOS DP
// FWHT<'&'>(f) eh soma de superset DP
// Se chamar com ^, usar tamanho potencia de 2!!
// O(n log(n))
// 50e84f
382 template < char op, class T > vector < T > FWHT (vector < T > f, bool inv =
   false) {
b75
        int n = f.size();
        for (int k = 0; (n-1) >> k; k++) for (int i = 0; i < n; i++) if
29e
            int j = i^(1 << k);
            if (op == '\^') f[i] += f[i], f[i] = f[i] - 2*f[i];
a38
            if (op == ', ') f[i] += (inv ? -1 : 1) * f[i];
93 c
            if (op == '&') f[j] += (inv ? -1 : 1) * f[i];
1bb
578
        if (op == '^' and inv) for (auto& i : f) i /= n;
```

```
abe
        return f;
50e }
4.16 FFT
// Chamar convolution com vector < complex < double >> para FFT
// Precisa do mint para NTT
// O(n log(n))
// Para FFT
// de56b9
488 void get_roots(bool f, int n, vector<complex<double>>& roots) {
        const static double PI = acosl(-1);
71a
        for (int i = 0; i < n/2; i++) {
            double alpha = i*((2*PI)/n);
b1e
1a1
            if (f) alpha = -alpha;
069
            roots[i] = {cos(alpha), sin(alpha)};
804
        }
de5 }
// Para NTT
// 91cd08
9f7 template < int p>
97b void get_roots(bool f, int n, vector<mod_int<p>>& roots) {
1e6
        mod_int  r;
de9
        int ord:
        if (p == 998244353) {
57a
9b6
           r = 102292;
81b
            ord = (1 << 23):
1cc
        } else if (p == 754974721) {
43a
            r = 739831874:
f0a
            ord = (1 << 24);
b60
        } else if (p == 167772161) {
a2a
            r = 243:
033
            ord = (1 << 25);
cd7
        } else assert(false);
        if (f) r = r^(p - 1 - ord/n);
547
        else r = r^(ord/n):
ee2
        roots[0] = 1;
be4
078
        for (int i = 1; i < n/2; i++) roots[i] = roots[i-1]*r;</pre>
63f }
// d5c432
8a2 template < typename T > void fft(vector < T > &a, bool f, int N,
   vector<int> &rev) {
```

```
for (int i = 0; i < N; i++) if (i < rev[i]) swap(a[i],
bc7
   a[rev[i]]):
        int 1, r, m;
12b
        vector <T> roots(N):
192
        for (int n = 2: n \le N: n *= 2) {
            get_roots(f, n, roots);
0f4
            for (int pos = 0; pos < N; pos += n) {
5dc
432
                1 = pos+0, r = pos+n/2, m = 0;
                while (m < n/2) {
a88
297
                    auto t = roots[m]*a[r];
254
                    a[r] = a[1] - t:
b8f
                    a[1] = a[1] + t:
925
                    l++; r++; m++;
                }
780
4ba
            }
04a
        }
       if (f) {
235
            auto invN = T(1)/T(N);
1c5
557
            for (int i = 0: i < N: i++) a[i] = a[i]*invN:
256
        }
e5e }
bf5 template < typename T > vector < T > convolution (vector < T > &a, vector < T >
   &b) {
279
        vector <T> l(a.begin(), a.end());
f41
        vector <T> r(b.begin(), b.end());
        int ln = l.size(), rn = r.size();
7c6
287
       int N = ln+rn-1:
f03
       int n = 1, log_n = 0;
ac4
        while (n \le N) \{ n \le 1; \log_n + +; \}
808
        vector < int > rev(n):
bae
        for (int i = 0; i < n; ++i) {</pre>
434
            rev[i] = 0:
920
            for (int j = 0; j < log_n; ++j)</pre>
                if (i & (1<<j)) rev[i] |= 1 << (log_n-1-j);</pre>
836
c34
143
        assert(N <= n);
fa4
        l.resize(n);
7e4
       r.resize(n):
56e
       fft(1, false, n, rev);
       fft(r, false, n, rev);
fcf
917
       for (int i = 0; i < n; i++) l[i] *= r[i];
88b
        fft(1, true, n, rev);
       l.resize(N):
5e1
792
        return 1;
700 }
```

```
// NTT
// 3bf256
6c8 template < int p, typename T > vector < mod_int < p >> ntt (vector < T > & a,
   vector<T>& b) {
        vector<mod_int<p>>> A(a.begin(), a.end()), B(b.begin(),
d52
d29
        return convolution(A. B);
3bf }
// Convolução de inteiro
//
// Precisa do CRT
// Tabela de valores:
// [0.1] - <int. 1>
// [-1e5, 1e5] - <11, 2>
// [-1e9, 1e9] - <__int128, 3>
//
// 053a7d
b3c template < typename T, int mods >
eec vector <T> int_convolution(vector <int>& a, vector <int>& b) {
        static const int M1 = 998244353, M2 = 754974721, M3 =
   167772161:
bf5
        auto c1 = ntt < M1 > (a, b):
221
        auto c2 = (mods \ge 2 ? ntt < M2 > (a. b) : vector < mod int < M2 >> ()):
        auto c3 = (mods >= 3 ? ntt < M3 > (a, b) : vector < mod_int < M3 >> ());
f9b
2da
        vector <T> ans:
5c5
        for (int i = 0; i < c1.size(); i++) {</pre>
            crt<T> at(c1[i].v. M1);
c09
316
            if (mods \ge 2) at = at * crt<T>(c2[i].v, M2);
987
            if (mods >= 3) at = at * crt<T>(c3[i].v, M3);
b2b
            ans.push back(at.a):
            if (at.a > at.m/2) ans.back() -= at.m:
26d
b9f
        }
ba7
        return ans;
5e8 }
4.17 Integração Numerica - Metodo de Simpson 3/8
// Integra f no intervalo [a, b], erro cresce proporcional a (b - a)^5
// 352415
676 const int N = 3*100: // multiplo de 3
287 ld integrate(ld a, ld b, function < ld(ld) > f) {
```

ld s = 0, h = (b - a)/N;

```
// Se b eh primo, basta fazer
// a^(b-2)

// cf94fe
f0a ll inv(ll a, ll b) {
    ae1      return a > 1 ? b - inv(b%a, a)*b/a : 1;
    cf9 }

// computa o inverso modular de 1..MAX-1 modulo um primo
// 7e4e3
a88 ll inv[MAX]:
0f2 inv[1] = 1;
0fa for (int i = 2; i < MAX; i++) inv[i] = MOD - MOD/i*inv[MOD%i]%MOD;</pre>
```

4.19 Karatsuba

```
// Os pragmas podem ajudar
// Para n \sim 2e5, roda em < 1 s
//
// O(n^1.58)
// 8065d6
//#pragma GCC optimize("Ofast")
//#pragma GCC target ("avx,avx2")
77a template < typename T > void kar(T* a, T* b, int n, T* r, T* tmp) {
d4c
       if (n <= 64) {
510
            for (int i = 0; i < n; i++) for (int j = 0; j < n; j++)
212
                r[i+j] += a[i] * b[j];
505
            return;
bb8
194
       int mid = n/2;
       T *atmp = tmp, *btmp = tmp+mid, *E = tmp+n;
2d7
4 f 1
        memset(E, 0, sizeof(E[0])*n);
c65
        for (int i = 0; i < mid; i++) {</pre>
c72
            atmp[i] = a[i] + a[i+mid];
4b9
            btmp[i] = b[i] + b[i+mid];
a3f
38a
        kar(atmp, btmp, mid, E, tmp+2*n);
b1e
        kar(a, b, mid, r, tmp+2*n);
229
        kar(a+mid, b+mid, mid, r+n, tmp+2*n);
```

```
c65
        for (int i = 0; i < mid; i++) {</pre>
735
            T \text{ temp} = r[i+mid];
de7
            r[i+mid] += E[i] - r[i] - r[i+2*mid];
f1e
            r[i+2*mid] += E[i+mid] - temp - r[i+3*mid];
f72
        }
28f }
e38 template < typename T > vector < T > karatsuba (vector < T > a, vector < T > b)
ba3
        int n = max(a.size(), b.size());
a84
        while (n&(n-1)) n++;
        a.resize(n), b.resize(n);
ae0
        vector < T > ret(2*n), tmp(4*n):
        kar(&a[0], &b[0], n, &ret[0], &tmp[0]);
644
edf
        return ret;
f87 }
4.20 Logaritmo Discreto
// Resolve logaritmo discreto com o algoritmo baby step giant step
// Encontra o menor x tal que a^x = b (mod m)
// Se nao tem. retorna -1
//
// O(sqrt(m) * log(sqrt(m))
// 739fa8
d41
da8 int dlog(int b, int a, int m) {
        if (a == 0) return b ? -1 : 1; // caso nao definido
d41
a6e
        a \%= m, b \%= m;
a10
        int k = 1, shift = 0;
31e
        while (1) {
6e3
            int g = gcd(a, m);
d47
            if (g == 1) break;
d41
9bc
            if (b == k) return shift;
642
            if (b % g) return -1;
c36
            b \neq g, m \neq g, shift++;
            k = (11) k * a / g % m;
9ab
515
        }
d41
af7
        int sq = sqrt(m)+1, giant = 1;
```

for (int i = 0; i < sq; i++) giant = (11) giant * a % m;</pre>

vector < pair < int , int >> baby;

for (int i = 0, cur = b; i <= sq; i++) {

baby.emplace_back(cur, i);

975

d41

0b5

33f

496

```
16c
            cur = (11) cur * a % m;
622
eb4
        sort(baby.begin(), baby.end());
d41
9c9
        for (int j = 1, cur = k; j <= sq; j++) {</pre>
            cur = (11) cur * giant % m;
ace
78b
            auto it = lower_bound(baby.begin(), baby.end(), pair(cur,
   INF)):
d26
            if (it != baby.begin() and (--it)->first == cur)
                return sq * j - it->second + shift;
ac3
b9d
        }
d41
daa
        return -1:
739 }
4.21 Miller-Rabin
// Testa se n eh primo, n \leq 3 * 10^18
//
// O(log(n)), considerando multiplicacao
// e exponenciacao constantes
// 4ebecc
d8b ll mul(ll a, ll b, ll m) {
e7a
        11 \text{ ret} = a*b - 11((long double)1/m*a*b+0.5)*m;
074
        return ret < 0 ? ret+m : ret;</pre>
2f3 }
03c ll pow(ll x, ll y, ll m) {
13a
       if (!y) return 1;
dbc
        ll ans = pow(mul(x, x, m), v/2, m);
7fa
        return y%2 ? mul(x, ans, m) : ans;
539 }
1a2 bool prime(ll n) {
```

if (n < 2) return 0;</pre>

if (n <= 3) return 1;</pre>

if (n % 2 == 0) return 0;

11 x = pow(a, d, n);

 $ll r = _builtin_ctzll(n - 1), d = n >> r;$

// funciona para n \leq 3*10^24 com os primos ate 41

for (int a: {2, 325, 9375, 28178, 450775, 9780504,

if (x == 1 or x == n - 1 or a % n == 0) continue;

// com esses primos, o teste funciona garantido para n <= 2^64

1aa

237

9de

f6a

da0 709

795265022}) {

4.22 Pollard's Rho Alg

```
// Usa o algoritmo de deteccao de ciclo de Floyd
// com uma otimizacao na qual o gcd eh acumulado
// A fatoração não sai necessariamente ordenada
// O algoritmo rho encontra um fator de n,
// e funciona muito bem quando n possui um fator pequeno
//
// Complexidades (considerando mul constante):
// rho - esperado O(n^{(1/4)}) no pior caso
// fact - esperado menos que O(n^{(1/4)} \log(n)) no pior caso
// b00653
d8b ll mul(ll a, ll b, ll m) {
        11 ret = a*b - 11((long double)1/m*a*b+0.5)*m;
074
        return ret < 0 ? ret+m : ret;</pre>
2f3 }
03c ll pow(ll x, ll y, ll m) {
13a
        if (!y) return 1;
dbc
        11 ans = pow(mul(x, x, m), y/2, m);
7fa
        return y%2 ? mul(x, ans, m) : ans;
539 }
1a2 bool prime(ll n) {
        if (n < 2) return 0;
237
        if (n <= 3) return 1;</pre>
        if (n % 2 == 0) return 0;
9de
f6a
        ll r = \_builtin\_ctzll(n - 1), d = n >> r;
        for (int a: {2, 325, 9375, 28178, 450775, 9780504,
771
    795265022}) {
            11 x = pow(a, d, n);
da0
709
            if (x == 1 \text{ or } x == n - 1 \text{ or a } \% n == 0) continue;
4a2
            for (int j = 0; j < r - 1; j++) {
10f
                x = mul(x, x, n);
                if (x == n - 1) break;
df0
```

```
1ff
e1b
            if (x != n - 1) return 0;
e74
        }
6a5
        return 1;
5ba }
9cf ll rho(ll n) {
0f9
        if (n == 1 or prime(n)) return n;
f7c
        auto f = [n](11 x) \{ return mul(x, x, n) + 1; \};
8a5
        11 x = 0, y = 0, t = 30, prd = 2, x0 = 1, q;
533
        while (t % 40 != 0 or gcd(prd, n) == 1) {
8a0
            if (x==y) x = ++x0, y = f(x):
            q = mul(prd, abs(x-y), n);
e13
21f
            if (q != 0) prd = q;
450
            x = f(x), y = f(f(y)), t++;
379
002
        return gcd(prd, n);
523 }
5b7 vector<ll> fact(ll n) {
        if (n == 1) return {};
       if (prime(n)) return {n};
0ec
0ed
       ll d = rho(n);
1de
        vector < 11 > 1 = fact(d), r = fact(n / d);
       1.insert(1.end(), r.begin(), r.end());
3af
        return 1:
792
902 }
4.23 Produto de dois long long mod m
// 0(1)
// 2f3a79
d8b ll mul(ll a, ll b, ll m) { // a*b % m
e7a
        11 \text{ ret} = a*b - 11((long double)1/m*a*b+0.5)*m;
074
        return ret < 0 ? ret+m : ret:</pre>
2f3 }
4.24 Simplex
// Maximiza c^T x s.t. Ax <= b. x >= 0
// O(2^n), porem executa em O(n^3) no caso medio
// 3a08e5
395 const double eps = 1e-7;
```

```
493 namespace Simplex {
69 c
        vector < vector < double >> T;
14e
        int n. m:
        vector < int > X. Y:
43e
c51
        void pivot(int x, int y) {
8e6
             swap(X[y], Y[x-1]);
d03
             for (int i = 0; i <= m; i++) if (i != y) T[x][i] /=
   T[x][y];
33c
            T[x][y] = 1/T[x][y];
38b
             for (int i = 0; i <= n; i++) if (i != x and abs(T[i][y]) >
   eps) {
774
                 for (int j = 0; j <= m; j++) if (j != y) T[i][j] -=
   T[i][y] * T[x][i];
                 T[i][v] = -T[i][v] * T[x][v];
a7d
            }
        }
e05
        // Retorna o par (valor maximo, vetor solucao)
        pair < double , vector < double >> simplex(
6f8
e9d
                 vector < vector < double >> A, vector < double >> b,
   vector < double > c) {
5bb
             n = b.size(), m = c.size();
002
            T = vector(n + 1, vector < double > (m + 1));
2d9
            X = vector < int > (m);
            Y = vector < int > (n):
0c2
             for (int i = 0: i < m: i++) X[i] = i:</pre>
115
             for (int i = 0; i < n; i++) Y[i] = i+m;</pre>
51f
5b5
             for (int i = 0; i < m; i++) T[0][i] = -c[i];
603
             for (int i = 0; i < n; i++) {</pre>
ba6
                 for (int j = 0; j < m; j++) T[i+1][j] = A[i][j];
                 T[i+1][m] = b[i]:
eca
07 c
            }
667
             while (true) {
714
                 int x = -1, y = -1;
2db
                 double mn = -eps;
                 for (int i = 1; i <= n; i++) if (T[i][m] < mn) mn =</pre>
   T[i][m], x = i;
af2
                 if (x < 0) break;
                 for (int i = 0; i < m; i++) if (T[x][i] < -eps) { y = }
882
   i; break; }
4a6
                 if (y < 0) return \{-1e18, \{\}\}; // sem solucao para Ax
<= b
7fb
                 pivot(x, y);
472
            }
```

```
667
            while (true) {
714
                int x = -1, y = -1;
2db
                double mn = -eps;
               for (int i = 0; i < m; i++) if (T[0][i] < mn) mn =
   T[0][i], y = i;
9b0
               if (y < 0) break;
034
                mn = 1e200:
                for (int i = 1; i \le n; i++) if (T[i][y] > eps and
   T[i][m] / T[i][y] < mn
                    mn = T[i][m] / T[i][y], x = i;
                if (x < 0) return {1e18, {}}; // c^T x eh ilimitado
53b
7fb
                pivot(x, y);
81e
            }
290
            vector < double > r(m);
           for (int i = 0; i < n; i++) if (Y[i] < m) r[Y[i]] =
   T[i+1][m];
e59
            return {T[0][m], r};
7a4
       }
a64 }
```

4.25 Teorema Chines do Resto

```
// Combina equações modulares lineares: x = a (mod m)
// O m final eh o lcm dos m's, e a resposta eh unica mod o lcm
// Os m nao precisam ser coprimos
// Se nao tiver solucao, o 'a' vai ser -1
// 7cd7b3
153 template < typename T > tuple < T, T, T > ext_gcd(T a, T b) {
       if (!a) return {b, 0, 1};
550
        auto [g, x, y] = ext_gcd(b\%a, a);
c59
        return \{g, v - b/a*x, x\};
537 }
bfe template < typename T = 11 > struct crt {
627
       T a, m;
5f3
        crt(): a(0), m(1) {}
7eb
        crt(T a_, T m_) : a(a_), m(m_) {}
        crt operator * (crt C) {
911
238
            auto [g, x, y] = ext\_gcd(m, C.m);
dc0
            if ((a - C.a) \% g) a = -1;
4f9
            if (a == -1 or C.a == -1) return crt(-1, 0);
d09
           T lcm = m/g*C.m;
eb2
            T ans = a + (x*(C.a-a)/g \% (C.m/g))*m;
d8d
            return crt((ans % lcm + lcm) % lcm, lcm);
```

```
1f2
0d9 };
4.26 Totiente
// O(sqrt(n))
// faeca3
a7e int tot(int n){
0f6
        int ret = n;
        for (int i = 2; i*i <= n; i++) if (n % i == 0) {
505
            while (n \% i == 0) n /= i:
b0c
125
            ret -= ret / i:
        }
34a
af4
        if (n > 1) ret -= ret / n:
edf
        return ret;
fae }
5 DP
5.1 Divide and Conquer DP
// Particiona o arrav em k subarravs
// minimizando o somatorio das queries
//
// O(k n log n), assumindo quer query(1, r) eh O(1)
// 4efe6b
547 ll dp[MAX][2];
94b void solve(int k, int l, int r, int lk, int rk) {
        if (1 > r) return;
de6
109
        int m = (1+r)/2, p = -1;
        auto& ans = dp[m][k&1] = LINF;
d2b
        for (int i = max(m, lk); i <= rk; i++) {</pre>
6e2
324
            int at = dp[i+1][\sim k\&1] + query(m, i);
```

if (at < ans) ans = at, p = i;</pre>

solve(k, l, m-1, lk, p), solve(k, m+1, r, p, rk);

57d

91f

321

959 }

}

cf1 ll DC(int n, int k) {

dp[n][0] = dp[n][1] = 0;

```
f27
        for (int i = 0; i < n; i++) dp[i][0] = LINF;</pre>
b76
        for (int i = 1; i <= k; i++) solve(i, 0, n-i, 0, n-i);
8e7
        return dp[0][k&1];
5e9 }
```

5.2 Longest Common Subsequence

```
// Computa a LCS entre dois arrays usando
// o algoritmo de Hirschberg para recuperar
//
// O(n*m), O(n+m) de memoria
// 337bb3
eaf int lcs_s[MAX], lcs_t[MAX];
a6d int dp[2][MAX];
// dp[0][j] = max lcs(s[li...ri], t[lj, lj+j])
d12 void dp_top(int li, int ri, int lj, int rj) {
d13
        memset(dp[0], 0, (rj-lj+1)*sizeof(dp[0][0]));
753
        for (int i = li; i <= ri; i++) {</pre>
9aa
            for (int j = rj; j >= lj; j--)
                dp[0][j-1j] = max(dp[0][j-1j],
83b
                (lcs_s[i] == lcs_t[j]) + (j > 1j ? dp[0][j-1 - 1j] :
741
   0));
            for (int j = lj+1; j <= rj; j++)</pre>
04c
939
                dp[0][j-1j] = max(dp[0][j-1j], dp[0][j-1-1j]);
09f
        }
58f }
// dp[1][j] = max lcs(s[li...ri], t[lj+j, rj])
ca0 void dp_bottom(int li, int ri, int lj, int rj) {
0dd
        memset(dp[1], 0, (rj-lj+1)*sizeof(dp[1][0]));
3a2
        for (int i = ri; i >= li; i--) {
49c
            for (int j = lj; j <= rj; j++)</pre>
dbb
                dp[1][j-1j] = max(dp[1][j-1j],
                (lcs_s[i] == lcs_t[j]) + (j < rj ? dp[1][j+1 - lj] :
4da
   0));
            for (int j = rj-1; j >= lj; j--)
6ca
769
                dp[1][i - 1i] = max(dp[1][i - 1i], dp[1][i+1 - 1i]);
19b
        }
e8a }
93c void solve(vector<int>& ans, int li, int ri, int lj, int rj) {
        if (li == ri){
2ad
49c
            for (int j = lj; j <= rj; j++)</pre>
f5b
                if (lcs_s[li] == lcs_t[j]){
a66
                    ans.push_back(lcs_t[j]);
```

```
c2b
                     break:
                }
840
505
            return:
        }
126
534
        if (lj == rj){
753
            for (int i = li; i <= ri; i++){</pre>
88f
                 if (lcs_s[i] == lcs_t[lj]){
531
                     ans.push_back(lcs_s[i]);
c2b
                     break;
                }
68a
a03
            }
505
            return;
76d
        }
a57
        int mi = (li+ri)/2;
ade
        dp_top(li, mi, lj, rj), dp_bottom(mi+1, ri, lj, rj);
d7a
        int i_{-} = 0, mx = -1;
        for (int j = lj-1; j <= rj; j++) {
aee
da8
            int val = 0:
            if (i >= lj) val += dp[0][j - lj];
2bb
            if (j < rj) val += dp[1][j+1 - lj];
b9e
            if (val >= mx) mx = val, i_ = i;
ba8
14e
        }
6f1
        if (mx == -1) return:
c2a
        solve(ans, li, mi, lj, j_), solve(ans, mi+1, ri, j_+1, rj);
dd5 }
058 vector<int> lcs(const vector<int>& s, const vector<int>& t) {
        for (int i = 0; i < s.size(); i++) lcs_s[i] = s[i];</pre>
953
577
        for (int i = 0; i < t.size(); i++) lcs_t[i] = t[i];</pre>
        vector < int > ans:
dab
599
        solve(ans. 0. s.size()-1. 0. t.size()-1):
ba7
        return ans;
17c }
5.3 Mochila
// Resolve mochila, recuperando a resposta
// O(n * cap), O(n + cap) de memoria
// 400885
```

```
add int v[MAX], w[MAX]; // valor e peso
582 int dp[2][MAX_CAP];
```

```
// DP usando os itens [1, r], com capacidade = cap
0d6 void get_dp(int x, int 1, int r, int cap) {
        memset(dp[x], 0, (cap+1)*sizeof(dp[x][0]));
f8f
        for (int i = 1; i \le r; i++) for (int j = cap; j \ge 0; j--)
574
            if (j - w[i] >= 0) dp[x][j] = max(dp[x][j], v[i] + dp[x][j]
3a9
   - w[i]]);
ъ95 }
5ab void solve(vector<int>& ans, int 1, int r, int cap) {
        if (1 == r) {
893
9ff
            if (w[1] <= cap) ans.push_back(1);</pre>
505
            return:
13a
ee4
        int m = (1+r)/2;
283
        get_dp(0, 1, m, cap), get_dp(1, m+1, r, cap);
       int left_cap = -1, opt = -INF;
056
c94
        for (int j = 0; j <= cap; j++)
2f2
            if (int at = dp[0][j] + dp[1][cap - j]; at > opt)
                opt = at, left_cap = j;
91d
da3
        solve(ans, 1, m, left_cap), solve(ans, m+1, r, cap - left_cap);
d75 }
0d7 vector<int> knapsack(int n, int cap) {
dab
        vector < int > ans;
1e0
        solve(ans, 0, n-1, cap);
ba7
        return ans:
e4d }
5.4 SOS DP
// O(n 2^n)
// soma de sub-conjunto
// bec381
e03 vector<ll> sos_dp(vector<ll> f) {
       int N = __builtin_ctz(f.size());
6c0
        assert((1<<N) == f.size());
e59
        for (int i = 0; i < N; i++) for (int mask = 0; mask < (1<<N);
   mask++)
            if (mask>>i&1) f[mask] += f[mask^(1<<i)];</pre>
796
        return f:
abe
bec }
// soma de super-conjunto
// dbd121
```

e03 vector<ll> sos_dp(vector<ll> f) {

```
6c0
        int N = __builtin_ctz(f.size());
e59
        assert((1<<N) == f.size());
5a5
        for (int i = 0; i < N; i++) for (int mask = 0; mask < (1<<N);
   mask++)
            if (\sim mask >> i\&1) f[mask] += f[mask^(1<<ii)];
a3c
abe
        return f:
dbd }
5.5 Subset sum
// Retorna max(x <= t tal que existe subset de w que soma x)
//
// O(n * max(w))
// O(max(w)) de memoria
// d888b0
efd int subset_sum(vector<int> w, int t) {
        int pref = 0, k = 0;
bb5
417
        while (k < w.size()) and pref + w[k] <= t) pref += w[k++];
        if (k == w.size()) return pref;
1e7
444
        int W = *max_element(w.begin(), w.end());
        vector < int > last, dp(2*W, -1);
44d
d7b
        dp[W - (t-pref)] = k;
54d
        for (int i = k; i < w.size(); i++) {</pre>
288
            last = dp;
            for (int x = 0; x < W; x++) dp[x+w[i]] = max(dp[x+w[i]],
15f
   last[x]);
17b
            for (int x = 2*W - 1; x > W; x--)
                for (int j = max(0, last[x]); j < dp[x]; j++)</pre>
303
595
                     dp[x-w[j]] = max(dp[x-w[j]], j);
        }
867
2fb
        int ans = t;
        while (dp[W - (t-ans)] < 0) ans --;
1 c 1
ba7
        return ans:
d88 }
   Strings
6.1 Aho-corasick
// query retorna o somatorio do numero de matches de
```

```
// todas as stringuinhas na stringona
//
// insert - O(|s| log(SIGMA))
```

```
// build - O(N), onde N = somatorio dos tamanhos das strings
// query - 0(|s|)
// a30d6e
eal namespace aho {
        map < char , int > to[MAX];
807
c87
        int link[MAX], idx, term[MAX], exit[MAX], sobe[MAX];
        void insert(string& s) {
bfc
            int at = 0:
05e
b4f
            for (char c : s) {
b68
                auto it = to[at].find(c);
1c9
                if (it == to[at].end()) at = to[at][c] = ++idx:
361
                else at = it->second:
ff4
            term[at]++, sobe[at]++;
142
6eb
        }
d41 #warning nao esquece de chamar build() depois de inserir
        void build() {
0a8
26a
            queue < int > q;
537
            q.push(0);
            link[0] = exit[0] = -1;
dff
402
            while (q.size()) {
379
                int i = q.front(); q.pop();
3 c 4
                for (auto [c, j] : to[i]) {
5da
                    int l = link[i]:
                    while (1 != -1 and !to[1].count(c)) 1 = link[1];
102
7a5
                    link[i] = 1 == -1 ? 0 : to[1][c]:
                    exit[j] = term[link[j]] ? link[j] : exit[link[j]];
3ab
6f2
                    if (exit[j]+1) sobe[j] += sobe[exit[j]];
113
                    q.push(j);
f1d
                }
            }
367
768
        }
        int query(string& s) {
bc0
86d
            int at = 0, ans = 0;
b4f
            for (char c : s){
                while (at != -1 and !to[at].count(c)) at = link[at];
1ca
5b9
                at = at == -1 ? 0 : to[at][c]:
2b1
                ans += sobe[at];
            }
b85
ba7
            return ans;
038
a30 }
```

6.2 Algoritmo Z

```
// z[i] = lcp(s, s[i..n))
//
// Complexidades:
// z - O(|s|)
// \text{ match - } O(|s| + |p|)
// 74a9e1
a19 vector<int> get_z(string s) {
163
        int n = s.size();
2b1
        vector<int> z(n. 0):
fae
        int 1 = 0, r = 0;
6f5
        for (int i = 1: i < n: i++) {
            if (i \le r) z[i] = min(r - i + 1, z[i - 1]);
0af
            while (i + z[i] < n \text{ and } s[z[i]] == s[i + z[i]]) z[i]++;
457
65e
            if (i + z[i] - 1 > r) l = i, r = i + z[i] - 1;
5cd
        }
070
        return z;
74a }
6.3 Automato de Sufixo
// Automato que aceita os sufixos de uma string
// Todas as funcoes sao lineares
// c37a72
16e namespace sam {
c1a
        int cur, sz, len[2*MAX], link[2*MAX], acc[2*MAX];
0b8
        int nxt[2*MAX][26];
e6a
        void add(int c) {
17a
            int at = cur;
9a6
            len[sz] = len[cur]+1, cur = sz++;
             while (at != -1 and !nxt[at][c]) nxt[at][c] = cur, at =
500
   link[at];
7ea
            if (at == -1) { link[cur] = 0; return; }
654
            int q = nxt[at][c];
fd9
            if (len[q] == len[at]+1) { link[cur] = q; return; }
31f
            int qq = sz++;
2c3
            len[qq] = len[at]+1, link[qq] = link[q];
9a9
            for (int i = 0; i < 26; i++) nxt[qq][i] = nxt[q][i];</pre>
             while (at != -1 and nxt[at][c] == q) nxt[at][c] = qq, at =
   link[at];
868
            link[cur] = link[q] = qq;
61a
        }
94e
        void build(string& s) {
```

```
889
            cur = 0, sz = 0, len[0] = 0, link[0] = -1, sz++;
            for (auto i : s) add(i-'a');
9fe
17a
            int at = cur;
121
            while (at) acc[at] = 1, at = link[at];
0e7
        // coisas que da pra fazer:
        11 distinct_substrings() {
28c
04b
            11 \text{ ans} = 0;
            for (int i = 1; i < sz; i++) ans += len[i] - len[link[i]];</pre>
a1e
ba7
            return ans;
0d7
a6c
        string longest_common_substring(string& S, string& T) {
419
            build(S):
            int at = 0, 1 = 0, ans = 0, pos = -1;
111
d59
            for (int i = 0; i < T.size(); i++) {</pre>
f2c
                while (at and !nxt[at][T[i]-'a']) at = link[at], l =
   len[at]:
                if (nxt[at][T[i]-'a']) at = nxt[at][T[i]-'a'], 1++;
efa
749
                else at = 0.1 = 0:
a1a
                if (1 > ans) ans = 1, pos = i;
2b3
20f
            return T.substr(pos-ans+1, ans);
930
        }
46e
        11 dp[2*MAX];
455
        11 paths(int i) {
2a8
            auto& x = dp[i];
dee
            if (x) return x:
483
            for (int j = 0; j < 26; j++) if (nxt[i][j]) x +=
71c
   paths(nxt[i][j]);
            return x;
ea5
d88
        }
105
        void kth_substring(int k, int at=0) { // k=1 : menor substring
   lexicog.
            for (int i = 0; i < 26; i++) if (k and nxt[at][i]) {</pre>
9d2
d58
                if (paths(nxt[at][i]) >= k) {
d02
                    cout << char('a'+i);</pre>
c43
                    kth_substring(k-1, nxt[at][i]);
505
                    return:
69a
                k -= paths(nxt[at][i]);
5f4
ef6
       }
a13
c37 };
```

6.4 eertree

```
// Constroi a eertree, caractere a caractere
// Inicializar com a quantidade de caracteres maxima
// size() retorna a quantidade de substrings pal. distintas
// depois de chamar propagate(), cada substring palindromica
// ocorre qt[i] vezes. O propagate() retorna o numero de
// substrings pal. com repeticao
//
// O(n) amortizado, considerando alfabeto O(1)
// a2e693
8eb struct eertree {
7 c.c
        vector < vector < int >> t;
        int n, last, sz;
42e
745
         vector<int> s, len, link, qt;
         eertree(int N) {
d36
ec8
             t = vector(N+2, vector(26, int())):
cee
             s = len = link = qt = vector < int > (N+2);
cd1
             s[0] = -1:
             link[0] = 1, len[0] = 0, link[1] = 1, len[1] = -1;
288
688
             sz = 2, last = 0, n = 1;
        }
8dc
         void add(char c) {
244
692
             s[n++] = c -= 'a':
34f
             while (s[n-len[last]-2] != c) last = link[last];
289
             if (!t[last][c]) {
dab
                 int prev = link[last];
553
                 while (s[n-len[prev]-2] != c) prev = link[prev];
fb2
                 link[sz] = t[prev][c];
3f5
                 len[sz] = len[last]+2;
1f8
                 t[last][c] = sz++;
f8b
            }
344
             gt[last = t[last][c]]++;
b1d
f17
        int size() { return sz-2; }
2af
        11 propagate() {
b73
             11 \text{ ret} = 0:
             for (int i = n; i > 1; i--) {
ebb
fd3
                 qt[link[i]] += qt[i];
db5
                 ret += qt[i];
074
edf
             return ret:
ef6
        }
a2e };
```

6.5 KMP

```
// matching(s, t) retorna os indices das ocorrencias
// de s em t
// autKMP constroi o automato do KMP
// Complexidades:
// pi - O(n)
// match - 0(n + m)
// construir o automato - O(|sigma|*n)
// n = |padrao| e m = |texto|
// f50359
ea8 template < typename T > vector < int > pi(T s) {
        vector < int > p(s.size());
        for (int i = 1, j = 0; i < s.size(); i++) {</pre>
725
             while (j \text{ and } s[j] != s[i]) j = p[j-1];
a51
973
            if (s[j] == s[i]) j++;
f8c
            p[i] = j;
e0a
74e
        return p;
f50 }
// c82524
c10 template < typename T> vector < int > matching(T& s, T& t) {
658
        vector < int > p = pi(s), match;
a1b
        for (int i = 0, j = 0; i < t.size(); i++) {</pre>
             while (j \text{ and } s[j] != t[i]) j = p[j-1];
6be
c4d
             if (s[j] == t[i]) j++;
310
             if (j == s.size()) match.push_back(i-j+1), j = p[j-1];
028
ed8
        return match;
c82 }
// 79bd9e
a2d struct KMPaut : vector < vector < int >> {
47c
        KMPaut(){}
        KMPaut (string& s) : vector < vector < int >> (26,
6c7
   vector < int > (s.size()+1)) {
503
             vector < int > p = pi(s);
04b
             auto& aut = *this;
4fa
             aut[s[0]-'a'][0] = 1;
            for (char c = 0; c < 26; c++)
                 for (int i = 1; i <= s.size(); i++)</pre>
5d3
42b
                      aut[c][i] = s[i] - 'a' == c ? i+1 : aut[c][p[i-1]];
4bb
        }
79b };
```

6.6 Manacher

```
// manacher recebe um vetor de T e retorna o vetor com tamanho dos
    palindromos
// ret[2*i] = tamanho do maior palindromo centrado em i
// \text{ ret}[2*i+1] = \text{tamanho maior palindromo centrado em i e i+1}
//
// Complexidades:
// manacher - O(n)
// palindrome - <0(n), 0(1)>
// pal_end - 0(n)
// ebb184
28a template < typename T > vector < int > manacher (const T& s) {
         int 1 = 0, r = -1, n = s.size();
fc9
         vector < int > d1(n), d2(n);
         for (int i = 0; i < n; i++) {</pre>
603
821
             int k = i > r ? 1 : min(d1[l+r-i], r-i);
61a
             while (i+k < n \&\& i-k >= 0 \&\& s[i+k] == s[i-k]) k++:
61e
             d1[i] = k--;
9f6
             if (i+k > r) l = i-k, r = i+k;
950
         }
        1 = 0, r = -1;
e03
603
         for (int i = 0; i < n; i++) {
             int k = i > r ? 0 : min(d2[1+r-i+1], r-i+1); k++;
a64
             while (i+k \le n \&\& i-k \ge 0 \&\& s[i+k-1] == s[i-k]) k++;
2c6
             d2[i] = --k:
eaa
             if (i+k-1 > r) l = i-k, r = i+k-1;
26d
4fe
         }
c41
         vector < int > ret(2*n-1);
e6b
         for (int i = 0; i < n; i++) ret[2*i] = 2*d1[i]-1;
 e1d
         for (int i = 0; i < n-1; i++) ret[2*i+1] = 2*d2[i+1];
edf
         return ret;
ebb }
// 60c6f5
// verifica se a string s[i..j] eh palindromo
cac template < typename T > struct palindrome {
f97
         vector < int > man;
b2d
         palindrome(const T& s) : man(manacher(s)) {}
9d7
         bool query(int i, int j) {
bad
             return man[i+j] >= j-i+1;
1e7
60c }:
// 8bd4d5
```

```
// tamanho do maior palindromo que termina em cada posicao
7cb template < typename T > vector < int > pal_end(const T& s) {
        vector < int > ret(s.size());
e57
fde
        palindrome <T> p(s);
        ret[0] = 1;
d51
88e
        for (int i = 1; i < s.size(); i++) {</pre>
a32
            ret[i] = min(ret[i-1]+2, i+1):
            while (!p.query(i-ret[i]+1, i)) ret[i]--;
6ea
78e
edf
        return ret;
8bd }
6.7 Min/max suffix/cyclic shift
// Computa o indice do menor/maior sufixo/cyclic shift
// da string, lexicograficamente
//
// O(n)
// af0367
016 template < typename T > int max_suffix(T s, bool mi = false) {
        s.push_back(*min_element(s.begin(), s.end())-1);
476
        int ans = 0;
1a4
        for (int i = 1; i < s.size(); i++) {</pre>
88e
            int j = 0;
eec
            while (ans+j < i and s[i+j] == s[ans+j]) j++;
708
7a2
            if (s[i+j] > s[ans+j]) {
                if (!mi or i != s.size()-2) ans = i;
b52
e51
            } else if (j) i += j-1;
69 c
ba7
        return ans;
f2a }
a1a template < typename T > int min_suffix(T s) {
76b
        for (auto& i : s) i *= -1;
        s.push_back(*max_element(s.begin(), s.end())+1);
09d
925
        return max_suffix(s, true);
ec0 }
97c template < typename T > int max_cyclic_shift(T s) {
        int n = s.size();
163
1ad
        for (int i = 0; i < n; i++) s.push_back(s[i]);</pre>
        return max_suffix(s);
20a
d34 }
```

08a template < typename T> int min_cyclic_shift(T s) {

for (auto& i : s) i *= -1;

76b

```
7be
        return max_cyclic_shift(s);
c7a }
6.8 String Hashing
// Complexidades:
// construtor - O(|s|)
// operator() - 0(1)
// 918dfb
878 mt19937 rng((int)
    chrono::steady_clock::now().time_since_epoch().count());
463 int uniform(int 1, int r) {
        uniform_int_distribution < int > uid(1, r);
a7f
f54
        return uid(rng);
d9e }
9e0 template <int MOD> struct str_hash { // 116fcb
        static int P;
c63
dcf
        vector<ll> h, p;
        str_hash(string s) : h(s.size()), p(s.size()) {
ea8
            p[0] = 1, h[0] = s[0];
7a2
            for (int i = 1; i < s.size(); i++)</pre>
ad7
                p[i] = p[i - 1]*P%MOD, h[i] = (h[i - 1]*P + s[i])%MOD;
84 c
1ef
af7
        11 operator()(int 1, int r) { // retorna hash s[1...r]
749
            ll hash = h[r] - (1 ? h[1 - 1]*p[r - 1 + 1]%MOD : 0);
dfd
            return hash < 0 ? hash + MOD : hash;</pre>
3ba
116 }:
217 template < int MOD > int str_hash < MOD > :: P = uniform(256, MOD - 1); //
   1 > |sigma|
6.9 String Hashing - modulo 2<sup>61</sup> - 1
// Quase duas vezes mais lento
// Complexidades:
// build - 0(|s|)
// operator() - 0(1)
//
// d3c0f0
9d0 const 11 MOD = (111<<61) - 1;
e38 ll mulmod(ll a, ll b) {
        const static ll LOWER = (111<<30) - 1, GET31 = (111<<31) - 1;</pre>
```

```
11 \ 11 = a\&LOWER, h1 = a>>30, 12 = b\&LOWER, h2 = b>>30;
410
d54
        11 m = 11*h2 + 12*h1, h = h1*h2;
                                                                            0a2
        ll ans = 11*12 + (h>>1) + ((h&1)<<60) + (m>>31) +
                                                                            5ce
784
   ((m\&GET31) << 30) + 1;
        ans = (ans\&MOD) + (ans>>61), ans = (ans\&MOD) + (ans>>61);
1dd
                                                                            fae
                                                                                cnt[ra[i]]++:
c0f
        return ans - 1;
f98 }
                                                                            4c4
                                                                            368
798 mt19937_64
   rng(chrono::steady_clock::now().time_since_epoch().count());
                                                                            28f
                                                                                ra[sa[i]] !=
f89 ll uniform(ll l, ll r) {
                                                                            f86
969
        uniform int distribution < 11 > uid(1, r):
                                                                            26b
        return uid(rng);
f54
                                                                            d5e
                                                                            11e
cac }
                                                                            057
d7d struct str_hash {
                                                                            ff3 }
        static 11 P;
c20
dcf
        vector<ll> h, p;
        str_hash(string s) : h(s.size()), p(s.size()) {
                                                                            232
ea8
            p[0] = 1, h[0] = s[0];
7a2
                                                                            408
            for (int i = 1; i < s.size(); i++)</pre>
                                                                            676
ad7
                p[i] = mulmod(p[i - 1], P), h[i] = (mulmod(h[i - 1],
632
   P) + s[i])%MOD;
                                                                            740
                                                                            199
507
af7
        11 operator()(int 1, int r) { // retorna hash s[1...r]
                                                                            1de
538
            ll hash = h[r] - (1 ? mulmod(h[1 - 1], p[r - 1 + 1]) : 0);
                                                                            891
dfd
            return hash < 0 ? hash + MOD : hash:</pre>
                                                                            d98
544
        }
                                                                            a07
148 };
                                                                            5ed
                                                                                    return lcp;
6c5 ll str_hash::P = uniform(256, MOD - 1); // l > |sigma|
                                                                            fbe }
6.10 Suffix Array - O(n log n)
// kasai recebe o suffix array e calcula lcp[i],
                                                                            // Rapidao
// o lcp entre s[sa[i],...,n-1] e s[sa[i+1],...,n-1]
//
// Complexidades:
                                                                            //
// suffix_array - O(n log(n))
// kasai - O(n)
                                                                            // Complexidades
// d3a6ce
                                                                            // query - 0(1)
733 vector <int > suffix_array(string s) {
b38
       s += "$";
043
        int n = s.size(), N = max(n, 260);
2f3
       vector < int > sa(n), ra(n);
```

for(int i = 0; i < n; i++) sa[i] = i, ra[i] = s[i];</pre>

29b

```
for(int k = 0; k < n; k ? k *= 2 : k++) {
             vector < int > nsa(sa), nra(n), cnt(N);
             for (int i = 0; i < n; i++) nsa[i] = (nsa[i]-k+n)%n,
             for (int i = 1: i < N: i++) cnt[i] += cnt[i-1]:
             for(int i = n-1; i+1; i--) sa[--cnt[ra[nsa[i]]] = nsa[i];
             for(int i = 1, r = 0; i < n; i++) nra[sa[i]] = r +=</pre>
                 ra[sa[i-1]] or ra[(sa[i]+k)\%n] != ra[(sa[i-1]+k)\%n];
             ra = nra:
             if (ra[sa[n-1]] == n-1) break;
         return vector < int > (sa.begin()+1, sa.end());
481 vector <int > kasai(string s, vector <int > sa) {
         int n = s.size(), k = 0;
         vector < int > ra(n), lcp(n);
         for (int i = 0; i < n; i++) ra[sa[i]] = i;</pre>
         for (int i = 0; i < n; i++, k -= !!k) {</pre>
             if (ra[i] == n-1) { k = 0; continue; }
             int i = sa[ra[i]+1]:
             while (i+k < n \text{ and } j+k < n \text{ and } s[i+k] == s[j+k]) k++;
             lcp[ra[i]] = k;
6.11 Suffix Array - O(n)
// Computa o suffix array em 'sa', o rank em 'rnk'
// e o lcp em 'lcp'
// query(i, j) retorna o LCP entre s[i..n-1] e s[j..n-1]
// O(n) para construir
// hash do arquivo inteiro: fa533e
// bab412
| 1a5 template < typename T> struct rmq {
```

```
517
        vector <T> v;
fcc
        int n; static const int b = 30;
70e
        vector < int > mask, t;
        int op(int x, int y) { return v[x] \leftarrow v[y] ? x : y; }
183
        int msb(int x) { return __builtin_clz(1)-__builtin_clz(x); }
ee1
c92
        int small(int r. int sz = b) { return
   r-msb(mask[r]&((1<<sz)-1)); }
        rmq() {}
6ad
        rmq(const\ vector < T > \&\ v_) : v(v_), n(v.size()), mask(n), t(n) {
43 c
2e5
            for (int i = 0, at = 0; i < n; mask[i++] = at |= 1) {
a61
                at = (at << 1) &((1 << b) -1):
c00
                while (at and op(i-msb(at&-at), i) == i) at ^= at&-at:
c2f
            }
            for (int i = 0; i < n/b; i++) t[i] = small(b*i+b-1);
ea4
            for (int j = 1; (1<<j) <= n/b; j++) for (int i = 0;
   i+(1<< j) <= n/b; i++)
                t[n/b*j+i] = op(t[n/b*(j-1)+i],
ba5
   t[n/b*(j-1)+i+(1<<(j-1))]);
41a
       }
e34
        int index_query(int 1, int r) {
27b
            if (r-l+1 \le b) return small(r, r-l+1);
            int x = 1/b+1, y = r/b-1;
e80
fd3
            if (x > y) return op(small(l+b-1), small(r));
a4e
            int j = msb(y-x+1);
ea3
            int ans = op(small(l+b-1), op(t[n/b*j+x],
   t[n/b*j+y-(1<<j)+1]));
be6
            return op(ans, small(r));
62a
093
        T query(int 1, int r) { return v[index_query(1, r)]; }
bab }:
9d7 struct suffix_array {
ac0
        string s:
        int n:
1a8
5b4
        vector < int > sa, cnt, rnk, lcp;
        rmq<int> RMQ;
2de
        bool cmp(int a1, int b1, int a2, int b2, int a3=0, int b3=0) {
d6e
91d
            return a1 != b1 ? a1 < b1 : (a2 != b2 ? a2 < b2 : a3 < b3);
82d
        template < typename T > void radix(int* fr, int* to, T* r, int N,
4a4
   int k) {
c17
            cnt = vector < int > (k+1, 0):
            for (int i = 0; i < N; i++) cnt[r[fr[i]]]++;</pre>
bac
            for (int i = 1; i <= k; i++) cnt[i] += cnt[i-1];</pre>
703
000
            for (int i = N-1: i+1: i--) to [--cnt[r[fr[i]]]] = fr[i]:
```

```
6f3
d66
        void rec(vector<int>& v, int k) {
a76
            auto &tmp = rnk, &m0 = lcp;
3a9
            int N = v.size()-3, sz = (N+2)/3, sz2 = sz+N/3;
7f8
            vector < int > R(sz2+3):
            for (int i = 1, j = 0; j < sz2; i += i%3) R[j++] = i;
74f
b30
            radix(&R[0], &tmp[0], &v[0]+2, sz2, k);
207
            radix(&tmp[0], &R[0], &v[0]+1, sz2, k);
            radix(&R[0], &tmp[0], &v[0]+0, sz2, k);
5 f 1
af5
            int dif = 0:
ed9
            int 10 = -1, 11 = -1, 12 = -1;
            for (int i = 0; i < sz2; i++) {</pre>
d81
                if (v[tmp[i]] != 10 or v[tmp[i]+1] != 11 or
8de
   v[tmp[i]+2] != 12)
b43
                     10 = v[tmp[i]], 11 = v[tmp[i]+1], 12 =
   v[tmp[i]+2], dif++;
199
                if (tmp[i]%3 == 1) R[tmp[i]/3] = dif;
                else R[tmp[i]/3+sz] = dif;
1f5
            }
d18
47f
            if (dif < sz2) {</pre>
146
                rec(R, dif);
746
                for (int i = 0; i < sz2; i++) R[sa[i]] = i+1;</pre>
105
            } else for (int i = 0; i < sz2; i++) sa[R[i]-1] = i;
6f4
            for (int i = 0, j = 0; j < sz2; i++) if (sa[i] < sz)
   tmp[j++] = 3*sa[i];
7ce
            radix(&tmp[0], &m0[0], &v[0], sz, k);
74d
            for (int i = 0: i < sz2: i++)</pre>
с9е
                 sa[i] = sa[i] < sz ? 3*sa[i]+1 : 3*(sa[i]-sz)+2;
332
            int at = sz2+sz-1, p = sz-1, p2 = sz2-1;
            while (p \ge 0 \text{ and } p2 \ge 0) {
1c9
3b3
                if ((sa[p2]%3==1 and cmp(v[m0[p]], v[sa[p2]],
   R[m0[p]/3],
                     R[sa[p2]/3+sz])) or (sa[p2]%3==2 and cmp(v[m0[p]],
0ce
   v[sa[p2]],
af6
                     v[m0[p]+1], v[sa[p2]+1], R[m0[p]/3+sz],
   R[sa[p2]/3+1]))
300
                     sa[at--] = sa[p2--];
                else sa[at--] = m0[p--];
cb0
            }
214
f2b
            while (p >= 0) sa[at--] = m0[p--];
            if (N\%3==1) for (int i = 0; i < N; i++) sa[i] = sa[i+1];
eb6
ee6
        }
```

```
938
        suffix_array(const string \& s_) : s(s_), n(s.size()), sa(n+3),
e62
                 cnt(n+1), rnk(n), lcp(n-1) {
             vector < int > v(n+3);
9fe
             for (int i = 0; i < n; i++) v[i] = i;</pre>
f9b
             radix(&v[0], &rnk[0], &s[0], n, 256);
eba
e6d
             int dif = 1:
             for (int i = 0; i < n; i++)</pre>
830
419
                 v[rnk[i]] = dif += (i and s[rnk[i]] != s[rnk[i-1]]);
             if (n \ge 2) rec(v, dif);
7 c.f
fb9
             sa.resize(n);
76f
             for (int i = 0; i < n; i++) rnk[sa[i]] = i;
             for (int i = 0, k = 0; i < n; i++, k -= !!k) {
892
668
                 if (rnk[i] == n-1) {
                     k = 0:
5a4
5e2
                     continue;
9df
                 int j = sa[rnk[i]+1];
39a
                 while (i+k < n \text{ and } j+k < n \text{ and } s[i+k] == s[j+k]) k++;
891
                 lcp[rnk[i]] = k;
825
a3e
9ff
             RMQ = rmq<int>(lcp);
9a8
        }
        // hash ateh aqui (sem o RMQ): 1ff700
588
        int query(int i, int j) {
d97
             if (i == j) return n-i;
223
             i = rnk[i], j = rnk[j];
             return RMQ.query(min(i, j), max(i, j)-1);
c3a
940
71c
        pair < int, int > next(int L, int R, int i, char c) {
024
             int 1 = L, r = R+1:
40 c
             while (1 < r) {
                 int m = (1+r)/2;
ee4
e7e
                 if (i+sa[m] >= n or s[i+sa[m]] < c) l = m+1;</pre>
                 else r = m:
ef3
             }
ebe
575
             if (1 == R+1 \text{ or } s[i+sa[1]] > c) \text{ return } \{-1, -1\};
eb7
            L = 1;
            1 = L, r = R+1;
9e2
             while (1 < r) {
40c
ee4
                 int m = (1+r)/2;
                 if (i+sa[m] >= n or s[i+sa[m]] <= c) l = m+1;</pre>
1a1
                 else r = m:
ef3
             }
b5b
```

```
56a
            R = 1-1:
e13
            return {L, R};
71b
        // quantas vezes 't' ocorre em 's' - O(|t| log n)
        int count_substr(string& t) {
66d
b2b
            int L = 0, R = n-1;
c9d
            for (int i = 0: i < t.size(): i++) {</pre>
de0
                tie(L, R) = next(L, R, i, t[i]);
                if (L == -1) return 0:
4fc
cff
fbf
            return R-L+1;
        }
aaa
        // exemplo de f que resolve o problema
           https://codeforces.com/edu/course/2/lesson/2/5/practice/contes
        ll f(ll k) { return k*(k+1)/2; }
57e
        11 dfs(int L, int R, int p) { // dfs na suffix tree chamado em
e68
   pre ordem
c54
            int ext = L != R ? RMQ.query(L, R-1) : n - sa[L];
            // Tem 'ext - p' substrings diferentes que ocorrem 'R-L+1'
                Vezes
            // O LCP de todas elas eh 'ext'
f80
            ll ans = (ext-p)*f(R-L+1);
            // L eh terminal, e folha sse L == R
            if (sa[L]+ext == n) L++;
63 c
            // se for um SA de varias strings separadas como s#t$u&,
                usar no lugar do if de cima
            // (separadores < 'a', diferentes e inclusive no final)</pre>
            // while (L <= R && (sa[L]+ext == n \mid | s[sa[L]+ext] <
                'a')) {
            // L++;
            // }
            while (L <= R) {
add
5a8
                int idx = L != R ? RMQ.index_query(L, R-1) : -1;
                if (idx == -1 or lcp[idx] != ext) idx = R;
5ef
478
                ans += dfs(L, idx, ext);
28d
                L = idx+1:
590
ba7
            return ans;
e21
        }
```

6.12 Suffix Array Dinamico

```
// Mantem o suffix array, lcp e rank de uma string,
// premitindo push_front e pop_front
// O operador [i] return um par com sa[i] e lcp[i]
// lcp[i] tem o lcp entre sa[i] e sa[i-1] (lcp[0] = 0)
// Complexidades:
// Construir sobre uma string de tamanho n: O(n log n)
// push_front e pop_front: O(log n) amortizado
// 4c2a2e
2fe struct dyn_sa {
3c9
        struct node {
1d4
            int sa, lcp;
ed1
            node *1, *r, *p;
f0d
            int sz, mi;
17b
            node(int sa_, int lcp_, node* p_) : sa(sa_), lcp(lcp_),
                1(NULL), r(NULL), p(p_), sz(1), mi(lcp) {}
543
01e
            void update() {
                sz = 1, mi = lcp;
58f
bd7
                if (1) sz += 1->sz, mi = min(mi, 1->mi);
                if (r) sz += r->sz, mi = min(mi, r->mi);
a54
27 c
            }
574
        };
bb7
        node* root;
295
        vector<ll> tag; // tag of a suffix (reversed id)
        string s; // reversed
ac0
        dyn_sa() : root(NULL) {}
cf4
e45
        dyn_sa(string s_) : dyn_sa() {
            reverse(s_.begin(), s_.end());
ae4
            for (char c : s_) push_front(c);
519
2a7
        }
a86
        \sim dyn_sa() {
609
            vector < node *> q = {root};
402
            while (q.size()) {
e5d
                node* x = q.back(); q.pop_back();
ee9
                if (!x) continue;
```

```
1c7
                 q.push_back(x->1), q.push_back(x->r);
bf0
                 delete x:
653
            }
        }
8c1
73c
        int size(node* x) { return x ? x->sz : 0; }
08e
        int mirror(int i) { return s.size()-1 - i; }
580
        bool cmp(int i, int j) {
a29
            if (s[i] != s[j]) return s[i] < s[j];</pre>
            if (i == 0 or j == 0) return i < j;</pre>
5b4
988
            return tag[i-1] < tag[j-1];</pre>
9fd
        }
919
        void fix_path(node* x) { while (x) x->update(), x = x->p; }
245
        void flatten(vector < node * > & v, node * x) {
8c8
            if (!x) return;
e96
            flatten(v, x->1);
2a2
            v.push_back(x);
42d
            flatten(v, x->r);
01f
964
        void build(vector<node*>& v, node*& x, node* p, int L, int R,
   11 1. 11 r) {
04c
            if (L > R) return void(x = NULL);
331
            int M = (L+R)/2:
3e3
            11 m = (1+r)/2;
7e5
            x = v[M]:
63e
            x->p = p;
bb3
            tag[x->sa] = m;
ae0
            build(v, x->1, x, L, M-1, 1, m-1), build(v, x->r, x, M+1,
   R, m+1, r);
ca8
            x->update();
a3a
82f
        void fix(node*& x, node* p, ll l, ll r) {
            if (3*max(size(x->1), size(x->r)) \le 2*size(x)) return
7f0
   x->update();
3d1
            vector < node *> v:
Осс
            flatten(v, x);
            build(v, x, p, 0, v.size()-1, 1, r);
ea9
b86
b19
        node* next(node* x) {
728
            if (x->r) {
a 9 1
                x = x -> r:
347
                 while (x->1) x = x->1;
ea5
                return x;
e7d
402
            while (x->p \text{ and } x->p->r == x) x = x->p;
137
            return x->p;
48b
        }
```

```
b68
        node* prev(node* x) {
            if (x->1) {
e41
a26
                x = x - > 1;
93c
                 while (x->r) x = x->r;
ea5
                return x;
9be
6a1
            while (x->p \text{ and } x->p->l == x) x = x->p;
137
            return x->p;
73e
        }
4f7
        int get_lcp(node* x, node* y) {
75a
            if (!x or !y) return 0; // change defaut value here
e51
            if (s[x->sa] != s[v->sa]) return 0:
            if (x->sa == 0 \text{ or } y->sa == 0) return 1;
843
            return 1 + query(mirror(x->sa-1), mirror(y->sa-1));
4d0
8d6
        }
ad6
        void add_suf(node*& x, node* p, int id, ll l, ll r) {
91e
            if (!x) {
8e3
                x = new node(id, 0, p);
                 node *prv = prev(x), *nxt = next(x);
8e2
65d
                int lcp_cur = get_lcp(prv, x), lcp_nxt = get_lcp(x,
   nxt);
                if (nxt) nxt->lcp = lcp_nxt, fix_path(nxt);
ca3
71f
                x \rightarrow lcp = lcp_cur;
7b4
                tag[id] = (1+r)/2;
ca8
                x->update();
505
                return;
d0e
            }
4a3
            if (cmp(id, x->sa)) add_suf(x->l, x, id, l, tag[x->sa]-1);
c3a
            else add_suf(x \rightarrow r, x, id, tag[x \rightarrow sa]+1, r);
3db
            fix(x, p, 1, r);
c98
ec2
        void push_front(char c) {
cc7
            s += c:
493
            tag.push_back(-1);
05e
            add_suf(root, NULL, s.size() - 1, 0, 1e18);
        }
1f2
7f3
        void rem suf(node*& x. int id) {
6cf
            if (x->sa != id) {
864
                 if (tag[id] < tag[x->sa]) return rem_suf(x->1, id);
e6f
                 return rem_suf(x->r, id);
            }
2ae
2cf
            node* nxt = next(x):
09b
            if (nxt) nxt->lcp = min(nxt->lcp, x->lcp), fix_path(nxt);
b20
            node *p = x - p, *tmp = x;
```

```
f3f
             if (!x->1 \text{ or } !x->r) {
2fd
                 x = x->1 ? x->1 : x->r;
753
                 if (x) x - p = p;
            } else {
9d9
7f7
                 for (tmp = x->1, p = x; tmp->r; tmp = tmp->r) p = tmp;
                 x->sa = tmp->sa, x->lcp = tmp->lcp;
f2a
482
                 if (tmp->1) tmp->1->p = p:
14c
                 if (p->1 == tmp) p->1 = tmp->1;
a94
                 else p \rightarrow r = tmp \rightarrow 1;
             }
dc4
b5e
             fix_path(p);
7c3
             delete tmp;
510
        }
15b
        void pop_front() {
abe
             if (!s.size()) return;
342
             s.pop_back();
436
             rem_suf(root, s.size());
сбе
             tag.pop_back();
987
        }
530
        int query(node* x, ll l, ll r, ll a, ll b) {
             if (!x \text{ or } tag[x->sa] == -1 \text{ or } r < a \text{ or } b < 1) return
e51
   s.size():
             if (a <= l and r <= b) return x->mi;
ef5
8eb
             int ans = s.size();
e1f
             if (a \le tag[x->sa]  and tag[x->sa] \le b) ans = min(ans,
   x \rightarrow lcp);
d99
             ans = min(ans, query(x->1, 1, tag[x->sa]-1, a, b));
             ans = min(ans, query(x->r, tag[x->sa]+1, r, a, b));
261
ba7
             return ans;
4 c 8
        }
588
        int query(int i, int j) { // lcp(s[i..], s[j..])
209
             if (i == j) return s.size() - i;
             11 a = tag[mirror(i)]. b = tag[mirror(i)]:
29e
             int ret = query(root, 0, 1e18, min(a, b)+1, max(a, b));
710
edf
             return ret;
84e
        }
        // optional: get rank[i], sa[i] and lcp[i]
044
        int rank(int i) {
396
             i = mirror(i);
52f
             node* x = root;
7c9
             int ret = 0;
f4c
             while (x) {
33e
                 if (tag[x->sa] < tag[i]) {</pre>
f9d
                     ret += size(x->1)+1;
a 9 1
                     x = x - > r:
6dc
                 else x = x - > 1:
```

```
a19
edf
            return ret;
153
649
        pair < int , int > operator[](int i) {
52f
            node* x = root;
31e
            while (1) {
d4d
                 if (i < size(x->1)) x = x->1:
4e6
                 else {
85f
                     i \rightarrow size(x\rightarrow 1);
                     if (!i) return {mirror(x->sa), x->lcp};
e03
040
                     i--, x = x->r;
b9b
                }
7a2
            }
90c
4c2 };
6.13 Trie
// trie T() constroi uma trie para o alfabeto das letras minusculas
// trie T(tamanho do alfabeto, menor caracter) tambem pode ser usado
//
// T.insert(s) - 0(|s|*sigma)
// T.erase(s) - O(|s|)
// T.find(s) retorna a posicao, O se nao achar - O(|s|)
// T.count_pref(s) numero de strings que possuem s como prefixo -
   0(|s|)
// Nao funciona para string vazia
// 979609
ab5 struct trie {
        vector < vector < int >> to:
e1a
450
        vector<int> end, pref;
af0
        int sigma; char norm;
        trie(int sigma_=26, char norm_='a') : sigma(sigma_),
   norm(norm_) {
58a
            to = {vector < int > (sigma)};
            end = \{0\}, pref = \{0\};
86e
fe1
        }
        void insert(string s) {
64e
c67
            int x = 0:
7e7
            for(auto c : s) {
008
                 int &nxt = to[x][c-norm];
dd7
                 if(!nxt) {
                     nxt = to.size();
Oaa
526
                     to.push_back(vector<int>(sigma));
770
                     end.push_back(0), pref.push_back(0);
```

```
933
                }
827
                x = nxt, pref[x]++;
34c
            }
e4e
            end[x]++;
e6b
        }
6b2
        void erase(string s) {
c67
            int x = 0:
            for(char c : s) {
b4f
800
                int &nxt = to[x][c-norm];
                x = nxt, pref[x]--;
10c
d8e
                if(!pref[x]) nxt = 0;
885
            }
bf0
            end[x]--;
ddd
aee
        int find(string s) {
c67
            int x = 0:
7e7
            for(auto c : s) {
                x = to[x][c-norm]:
2ec
a66
                if(!x) return 0;
            }
e12
ea5
            return x;
e77
839
        int count_pref(string s) {
e2f
            return pref[find(s)];
f40
        }
979 }:
```

7 Primitivas

7.1 Aritmetica Modular

```
// O mod tem q ser primo
// 5a6efb
429 template <int p> struct mod_int {
         ll pow(ll b, ll e) {
02c
a63
             if (e == 0) return 1;
630
             ll r = pow(b*b%p, e/2);
475
             if (e\%2 == 1) r = (r*b)\%p;
4 c 1
             return r;
3ba
        }
ae3
        11 inv(11 b) { return pow(b, p-2); }
4d7
         using m = mod_int;
d93
         int v;
fe0
         mod_int() : v(0) {}
```

```
e12
        mod_int(ll v_) {
019
            if (v_ >= p or v_ <= -p) v_ %= p;
bc6
            if (v_{-} < 0) v_{-} += p;
2e7
            v = v_{-};
7f3
74d
        m& operator+=(const m &a) {
2fd
            v += a.v:
ba5
            if (v >= p) v -= p;
357
            return *this;
c8b
eff
        m& operator -= (const m &a) {
8b4
            v -= a.v:
cc8
            if (v < 0) v += p:
            return *this;
357
f8d
4 c 4
        m& operator*=(const m &a) {
8a5
            v = v * 11(a.v) \% p;
357
            return *this;
d4c
3f9
        m& operator/=(const m &a) {
            v = v* inv(a.v) \% p;
5d6
357
            return *this;
62d
d65
        m operator-(){ return m(-v); }
b3e
        m& operator^=(11 e) {
            if (e < 0){
06d
6e2
                v = inv(v):
00c
                e = -e:
275
ebf
            v = pow(v, e\%(p-1));
357
            return *this;
e84
        bool operator == (const m &a) { return v == a.v; }
423
        bool operator!=(const m &a) { return v != a.v: }
69f
1c6
        friend istream &operator>>(istream &in, m& a) {
d1c
            11 val: in >> val:
d48
            a = m(val);
091
            return in;
870
44f
        friend ostream &operator << (ostream &out, m a) {</pre>
5a0
            return out << a.v;</pre>
214
        friend m operator+(m a, m b) { return a+=b; }
399
f9e
        friend m operator-(m a, m b) { return a-=b; }
9 c 1
        friend m operator*(m a, m b) { return a*=b; }
51b
        friend m operator/(m a, m b) { return a/=b; }
```

```
08f
        friend m operator^(m a, ll e) { return a^=e; }
f94 };
055 typedef mod_int < (int) 1e9+7 > mint;
7.2 Big Integer
// Complexidades: (para n digitos)
// Soma, subtracao, comparacao - O(n)
// Multiplicacao - O(n log(n))
// Divisao, resto - O(n^2)
// 6c3c3a
864 struct bint {
        static const int BASE = 1e9;
990
        vector < int > v;
3bd
        bool neg;
609
        bint() : neg(0) {}
        bint(int val) : bint() { *this = val; }
d53
        bint(long long val) : bint() { *this = val; }
e8f
a0f
        void trim() {
f42
            while (v.size() and v.back() == 0) v.pop_back();
df8
            if (!v.size()) neg = 0;
8e3
        }
        // converter de/para string | cin/cout
        bint(const char* s) : bint() { from_string(string(s)); }
294
548
        bint(const string& s) : bint() { from_string(s); }
4ab
        void from_string(const string& s) {
0a6
            v.clear(), neg = 0;
d72
            int ini = 0;
            while (ini < s.size() and (s[ini] == '-' or s[ini] == '+'
8e2
    or s[ini] == '0'))
71d
                if (s[ini++] == '-') neg = 1;
            for (int i = s.size()-1; i >= ini; i -= 9) {
883
05e
                int at = 0;
                for (int j = max(ini, i - 8); j \le i; j++) at = 10*at
   + (s[i]-'0'):
1fd
                v.push_back(at);
a5a
df8
            if (!v.size()) neg = 0;
e9a
        string to_string() const {
2ff
8be
            if (!v.size()) return "0";
793
            string ret;
```

```
73e
            if (neg) ret += '-';
            for (int i = v.size()-1; i >= 0; i--) {
3e9
582
                string at = ::to_string(v[i]);
                int add = 9 - at.size();
ced
                if (i+1 < v.size()) for (int j = 0; j < add; j++) ret</pre>
75e
   += '0':
f9f
                ret += at:
edf
            return ret;
770
d2f
        friend istream& operator >> (istream& in, bint& val) {
eb6
            string s; in >> s;
966
            val = s:
091
            return in;
328
99d
        friend ostream& operator << (ostream& out, const bint& val) {
8b9
            string s = val.to_string();
396
            out << s;
fe8
            return out;
ce1
       }
        // operators
60a
        friend bint abs(bint val) {
c5f
            val.neg = 0;
d94
            return val:
44b
bee
        friend bint operator-(bint val) {
815
            if (val != 0) val.neg ^= 1;
d94
            return val;
326
        bint& operator=(const bint& val) { v = val.v, neg = val.neg;
   return *this; }
        bint& operator=(long long val) {
249
0a6
            v.clear(). neg = 0:
            if (val < 0) neg = 1, val *= -1;
3a6
            for (; val; val /= BASE) v.push_back(val % BASE);
fdc
357
            return *this;
220
3bd
        int cmp(const bint& r) const { // menor: -1 | igual: 0 |
   maior: 1
            if (neg != r.neg) return neg ? -1 : 1;
b14
            if (v.size() != r.v.size()) {
0bb
ff7
                int ret = v.size() < r.v.size() ? -1 : 1;</pre>
91b
                return neg ? -ret : ret;
1f6
            for (int i = int(v.size())-1; i >= 0; i--) {
478
405
                if (v[i] != r.v[i]) {
```

```
2e5
                     int ret = v[i] < r.v[i] ? -1 : 1;</pre>
91b
                    return neg ? -ret : ret;
9a9
                }
c32
            }
bb3
            return 0;
07d
152
        friend bool operator < (const bint& 1, const bint& r) { return
   1.cmp(r) == -1;}
        friend bool operator > (const bint& 1, const bint& r) { return
c7a
   1.cmp(r) == 1; }
        friend bool operator <= (const bint& 1, const bint& r) { return
   1.cmp(r) <= 0; }
        friend bool operator >= (const bint& 1, const bint& r) { return
   1.cmp(r) >= 0; }
        friend bool operator == (const bint& 1, const bint& r) { return
   1.cmp(r) == 0; }
        friend bool operator!=(const bint& 1, const bint& r) { return
10b
   1.cmp(r) != 0; }
38e
        bint& operator +=(const bint& r) {
6bf
            if (!r.v.size()) return *this;
a93
            if (neg != r.neg) return *this -= -r;
256
            for (int i = 0, c = 0; i < r.v.size() or c; i++) {</pre>
e28
                if (i == v.size()) v.push_back(0);
08f
                v[i] += c + (i < r.v.size() ? r.v[i] : 0);
                if ((c = v[i] >= BASE)) v[i] -= BASE;
baa
8bb
357
            return *this:
ab1
54c
        friend bint operator+(bint a, const bint& b) { return a += b; }
        bint& operator -=(const bint& r) {
9 c 8
6bf
            if (!r.v.size()) return *this;
524
            if (neg != r.neg) return *this += -r;
358
            if ((!neg and *this < r) or (neg and r < *this)) {
                *this = r - *this;
b10
a10
                neg ^= 1;
357
                return *this;
807
256
            for (int i = 0, c = 0; i < r.v.size() or c; i++) {</pre>
9ef
                v[i] = c + (i < r.v.size() ? r.v[i] : 0);
                if ((c = v[i] < 0)) v[i] += BASE:
c8c
687
0eb
            trim();
357
            return *this;
f72
f44
        friend bint operator-(bint a, const bint& b) { return a -= b; }
```

```
// operators de * / %
        bint& operator *=(int val) {
6b0
            if (val < 0) val *= -1, neg ^= 1;</pre>
bca
566
            for (int i = 0, c = 0; i < v.size() or c; i++) {
e28
                if (i == v.size()) v.push_back(0);
352
                long long at = (long long) v[i] * val + c;
6a3
                v[i] = at % BASE:
b3d
                c = at / BASE;
cb1
            }
            trim();
0eb
357
            return *this;
a57
480
        friend bint operator *(bint a. int b) { return a *= b; }
        friend bint operator *(int a, bint b) { return b *= a; }
d5c
13b
        using cplx = complex <double >;
        void fft(vector < cplx > & a, bool f, int N, vector < int > & rev)
bfb
   const {
            for (int i = 0; i < N; i++) if (i < rev[i]) swap(a[i],
bc7
   a[rev[i]]);
            vector < cplx > roots(N);
bad
            for (int n = 2: n <= N: n *= 2) {
192
4e9
                const static double PI = acos(-1);
                for (int i = 0; i < n/2; i++) {
71a
40d
                    double alpha = (2*PI*i)/n;
1a1
                    if (f) alpha = -alpha;
3f6
                    roots[i] = cplx(cos(alpha), sin(alpha));
f16
                }
3e9
                for (int pos = 0: pos < N: pos += n)
                    for (int 1 = pos, r = pos+n/2, m = 0; m < n/2;
   1++, r++, m++) {
                         auto t = roots[m]*a[r]:
297
254
                         a[r] = a[1] - t;
                         a[1] = a[1] + t:
b8f
                    }
b0d
e07
            }
            if (!f) return;
3f1
08b
            auto invN = cplx(1)/cplx(N);
873
            for (int i = 0; i < N; i++) a[i] *= invN;</pre>
c75
0e0
        vector < long long > convolution (const vector < int > & a, const
   vector < int > & b) const {
            vector < cplx > l(a.begin(), a.end()), r(b.begin(), b.end());
ff9
996
            int ln = l.size(), rn = r.size(), N = ln+rn+1, n = 1,
   log_n = 0;
821
            while (n \le N) n \le 1, \log_n + 1;
            vector < int > rev(n):
808
603
            for (int i = 0: i < n: i++) {
```

```
434
                rev[i] = 0;
f44
                for (int j = 0; j < log_n; j++) if (i >> j & 1)
4ff
                    rev[i] = 1 << (log_n-1-j);
256
            }
230
            l.resize(n), r.resize(n);
            fft(1, false, n, rev), fft(r, false, n, rev);
a89
917
            for (int i = 0; i < n; i++) l[i] *= r[i];
88b
            fft(1, true, n, rev);
7ae
            vector < long long > ret;
            for (auto& i : 1) ret.push_back(round(i.real()));
c14
edf
            return ret;
917
633
        vector < int > convert base (const vector < int > & a. int from. int
   to) const {
498
            static vector<long long> pot(10, 1);
671
            if (pot[1] == 1) for (int i = 1; i < 10; i++) pot[i] =
   10*pot[i-1];
4b8
            vector < int > ret:
            long long at = 0;
156
608
            int digits = 0;
941
            for (int i : a) {
412
                at += i * pot[digits];
035
                digits += from;
684
                while (digits >= to) {
0c8
                    ret.push_back(at % pot[to]);
cf9
                    at /= pot[to];
fd4
                    digits -= to;
122
                }
87ъ
            }
944
            ret.push_back(at);
384
            while (ret.size() and ret.back() == 0) ret.pop_back();
edf
            return ret;
090
        }
edb
        bint operator*(const bint& r) const { // O(n log(n))
2af
            bint ret;
968
            ret.neg = neg ^ r.neg;
d5d
            auto conv = convolution(convert_base(v, 9, 4),
   convert_base(r.v, 9, 4));
a0e
            long long c = 0;
a74
            for (auto i : conv) {
f6d
                long long at = i+c;
4cb
                ret.v.push_back(at % 10000);
a25
                c = at / 10000;
773
            }
3cb
            for (; c; c /= 10000) ret.v.push_back(c%10000);
            ret.v = convert_base(ret.v, 4, 9);
0e2
25 c
            if (!ret.v.size()) ret.neg = 0;
```

```
edf
            return ret;
c6b
359
        bint& operator*=(const bint& r) { return *this = *this * r; };
9a3
        bint& operator/=(int val) {
d9a
            if (val < 0) neg ^= 1. val *= -1:</pre>
            for (int i = int(v.size())-1, c = 0; i >= 0; i--) {
f18
2a7
                long long at = v[i] + c * (long long) BASE;
                v[i] = at / val;
e02
fb1
                c = at % val;
            }
fdb
0eb
            trim();
357
            return *this;
db6
e74
        friend bint operator/(bint a, int b) { return a /= b; }
4a9
        int operator %=(int val) {
23b
            if (val < 0) val *= -1;</pre>
156
            long long at = 0;
            for (int i = int(v.size())-1; i >= 0; i--)
f31
                at = (BASE * at + v[i]) \% val;
1b3
d22
            if (neg) at *= -1;
ce6
            return at;
4b4
        friend int operator%(bint a, int b) { return a %= b; }
2fb
        friend pair < bint, bint > divmod(const bint& a_, const bint& b_)
13b
   611
            if (a_ == 0) return {0, 0};
d8a
            int norm = BASE / (b_.v.back() + 1);
b4e
            bint a = abs(a) * norm:
027
            bint b = abs(b_) * norm;
14d
            bint q, r;
            for (int i = a.v.size() - 1; i >= 0; i--) {
c91
b71
                r *= BASE, r += a.v[i];
                long long upper = b.v.size() < r.v.size() ?</pre>
   r.v[b.v.size()] : 0:
86d
                int lower = b.v.size() - 1 < r.v.size() ?</pre>
   r.v[b.v.size() - 1] : 0;
                int d = (upper * BASE + lower) / b.v.back();
431
5d4
30f
                while (r < 0) r += b, d--; // roda O(1) vezes
738
                q.v.push_back(d);
            }
c6a
a48
            reverse(q.v.begin(), q.v.end());
ae2
            q.neg = a_.neg ^ b_.neg;
88b
            r.neg = a_.neg;
8e5
            q.trim(), r.trim();
            return {q, r / norm};
0ef
4fd
        }
```

```
1d8
        bint operator/(const bint& val) { return divmod(*this,
   val).first: }
        bint& operator/=(const bint& val) { return *this = *this /
7f9
   val: }
        bint operator%(const bint& val) { return divmod(*this,
1f9
   val).second; }
        bint& operator%=(const bint& val) { return *this = *this %
   val; }
6c3 };
7.3 Matroid
// Matroids de Grafo e Particao
// De modo geral, toda Matroid contem um build() linear
// e uma funcao constante oracle()
// oracle(i) responde se o conjunto continua independente
// apos adicao do elemento i
// oracle(i, j) responde se o conjunto continua indepente
// apos trocar o elemento i pelo elemento j
//
// Intersecao sem peso O(r^2 n)
// em que n eh o tamanho do conjunto e r eh o tamanho da resposta
// Matroid Grafica
// Matroid das florestas de um grafo
// Um conjunto de arestas eh independente se formam uma floresta
// build() : O(n)
// oracle() : 0(1)
// 691847
fda struct graphic_matroid {
5da
        int n, m, t;
32c
        vector<array<int, 2>> edges;
        vector < vector < int >> g;
789
```

graphic_matroid(int n_, vector<array<int, 2>> edges_)

for (auto v : g[u]) if (in[v] == -1)

comp[v] = comp[u], dfs(v);

: n(n_), m(edges_.size()), edges(edges_), g(n), comp(n),

vector<int> comp, in, out;

void build(vector<int> I) {

62e

513

a1f

315

ab8

17d

863

677

483

945

a34

in(n), out(n) {}

}

void dfs(int u) {

in[u] = t++:

out[u] = t;

t = 0;

```
for (int u = 0; u < n; u++) g[u].clear(), in[u] = -1;
741
667
            for (int e : I) {
d00
                auto [u, v] = edges[e];
                g[u].push_back(v), g[v].push_back(u);
125
a8a
            for (int u = 0; u < n; u++) if (in[u] == -1)
809
a7d
                comp[u] = u. dfs(u):
207
        }
f31
        bool is_ancestor(int u, int v) {
            return in[u] <= in[v] and in[v] < out[u];</pre>
a68
0c2
        }
e6b
        bool oracle(int e) {
453
            return comp[edges[e][0]] != comp[edges[e][1]];
687
        }
f75
        bool oracle(int e, int f) {
574
            if (oracle(f)) return true;
622
            int u = edges[e][in[edges[e][0]] < in[edges[e][1]]];</pre>
            return is_ancestor(u, edges[f][0]) != is_ancestor(u,
ff2
   edges[f][1]);
8a9
       }
691 }:
// Matroid de particao ou cores
// Um conjunto eh independente se a quantidade de elementos
// de cada cor nao excede a capacidade da cor
// Quando todas as capacidades sao 1, um conjunto eh independente
// se todas as suas cores sao distintas
//
// build() : O(n)
// oracle() : 0(1)
// caa72a
994 struct partition_matroid {
501
        vector < int > cap. color. d:
        partition_matroid(vector<int> cap_, vector<int> color_)
608
04d
            : cap(cap_), color(color_), d(cap.size()) {}
945
        void build(vector<int> I) {
def
            fill(d.begin(), d.end(), 0);
e9d
            for (int u : I) d[color[u]]++;
c58
        }
514
        bool oracle(int u) {
            return d[color[u]] < cap[color[u]];</pre>
0a1
703
f7f
        bool oracle(int u, int v) {
2f7
            return color[u] == color[v] or oracle(v);
        }
4b4
caa }:
```

```
// Intersecao de matroid sem pesos
// Dadas duas matroids M1 e M2 definidas sobre o mesmo
// conjunto I, retorna o maior subconjunto de I
// que eh independente tanto para M1 quanto para M2
// O(r^2*n)
// 899f94
// Matroid "pesada" deve ser a M2
132 template < typename Matroid1, typename Matroid2 >
801 vector < int > matroid_intersection(int n, Matroid1 M1, Matroid2 M2) {
f5b
        vector < bool > b(n):
a64
        vector < int > I[2]:
a8b
        bool converged = false;
        while (!converged) {
0 c 1
742
            I[0].clear(), I[1].clear();
            for (int u = 0; u < n; u++) I[b[u]].push_back(u);
99d
09d
            M1.build(I[1]), M2.build(I[1]);
289
            vector < bool > target(n), pushed(n);
26a
            queue < int > q;
5 c 5
            for (int u : I[0]) {
2b2
                target[u] = M2.oracle(u);
c1b
                if (M1.oracle(u)) pushed[u] = true, q.push(u);
0e6
3fe
            vector < int > p(n, -1);
07a
            converged = true;
402
            while (q.size()) {
be1
                int u = q.front(); q.pop();
5.6
                if (target[u]) {
101
                     converged = false;
                     for (int v = u; v != -1; v = p[v]) b[v] = !b[v];
c32
c2b
                     break:
a80
                for (int v : I[!b[u]]) if (!pushed[v]) {
e78
                     if ((b[u] and M1.oracle(u, v)) or (b[v] and
34d
   M2.oracle(v, u)))
                         p[v] = u, pushed[v] = true, q.push(v);
bae
533
            }
1d9
5e7
        }
        return I[1];
b68
381 }
// Intersecao de matroid com pesos
// Dadas duas matroids M1 e M2 e uma funcao de pesos w, todas
```

```
definidas sobre
// um conjunto I retorna o maior subconjunto de I (desempatado pelo
// que eh independente tanto para M1 quanto para M2
// A resposta eh construida incrementando o tamanho conjunto I de 1 em
// Se nao tiver custo negativo, nao precisa de SPFA
//
// O(r^3*n) com SPFA
// O(r^2*n*log(n)) com Dijkstra e potencial
// 3a09d1
42a template < typename T, typename Matroid1, typename Matroid2>
2b5 vector < int > weighted_matroid_intersection(int n, vector < T > w,
   Matroid1 M1, Matroid2 M2) {
6c9
        vector < bool > b(n), target(n), is_inside(n);
563
        vector < int > I[2], from(n);
e35
        vector < pair < T, int >> d(n);
        auto check_edge = [&](int u, int v) {
169
            return (b[u] and M1.oracle(u, v)) or (b[v] and
   M2.oracle(v, u));
253
        };
        while (true) {
667
742
            I[0].clear(), I[1].clear();
99d
            for (int u = 0; u < n; u++) I[b[u]].push_back(u);
            // I[1] contem o conjunto de tamanho I[1].size() de menor
                peso
09d
            M1.build(I[1]), M2.build(I[1]);
            for (int u = 0; u < n; u++) {
687
                target[u] = false, is_inside[u] = false, from[u] = -1;
ea5
                d[u] = {numeric_limits <T>::max(), INF};
961
392
            }
8d3
            deque <T> q;
476
            sort(I[0].begin(), I[0].end(), [&](int i, int j){ return
   w[i] < w[i]; \});
            for (int u : I[0]) {
5c5
2b2
                target[u] = M2.oracle(u);
5a7
                if (M1.oracle(u)) {
4ef
                    if (is_inside[u]) continue;
7cc
                    d[u] = \{w[u], 0\};
                    if (!q.empty() and d[u] > d[q.front()])
   q.push_back(u);
                    else q.push_front(u);
655
4ae
                    is_inside[u] = true;
764
                }
            }
add
402
            while (q.size()) {
```

```
97a
                 int u = q.front(); q.pop_front();
6f3
                 is_inside[u] = false;
57a
                 for (int v : I[!b[u]]) if (check_edge(u, v)) {
9de
                     pair <T, int > nd(d[u].first + w[v], d[u].second +
   1):
                     if (nd < d[v]) {</pre>
61b
6ac
                         from[v] = u. d[v] = nd:
bd7
                         if (is_inside[v]) continue;
                         if (q.size() and d[v] > d[q.front()])
eec
   q.push_back(v);
275
                         else q.push_front(v);
587
                         is_inside[v] = true;
b3f
                    }
a3b
                }
563
            }
cc8
            pair < T, int > mini = pair (numeric_limits < T >:: max(), INF);
489
            int targ = -1;
259
            for (int u : I[0]) if (target[u] and d[u] < mini)</pre>
2b9
                 mini = d[u], targ = u;
            if (targ != -1) for (int u = targ; u != -1; u = from[u])
e14
489
                b[u] = !b[u], w[u] *= -1;
f97
            else break:
        }
c7d
        return I[1];
b68
8e7 }
```

7.4 Primitivas de fração

```
// Funciona com o Big Int
// cdb445
a4e template < typename T = int > struct frac {
a40
        T num, den;
e3f
        template < class U, class V>
61d
        frac(U num_ = 0, V den_ = 1) : num(num_), den(den_) {
             assert(den != 0);
bad
583
             if (den < 0) num *= -1, den *= -1;
            T g = gcd(abs(num), den);
a51
572
             num \neq g, den \neq g;
        }
fbf
51f
        friend bool operator < (const frac& 1, const frac& r) {</pre>
             return l.num * r.den < r.num * l.den:</pre>
fa0
a4e
4b5
        friend frac operator+(const frac& 1, const frac& r) {
b61
             return {1.num*r.den + 1.den*r.num, 1.den*r.den};
25f
        }
```

```
74d
        friend frac operator-(const frac& 1, const frac& r) {
            return {1.num*r.den - 1.den*r.num, 1.den*r.den};
2cd
8a7
        friend frac operator*(const frac& 1, const frac& r) {
c80
            return {1.num*r.num, 1.den*r.den};
510
14b
a1b
       friend frac operator/(const frac& 1, const frac& r) {
8f3
            return {1.num*r.den, 1.den*r.num};
b2c
       friend ostream& operator << (ostream& out, frac f) {</pre>
012
37a
            out << f.num << ',' << f.den;
            return out:
b49
cdb };
   Primitivas de matriz - exponenciacao
// d05c24
945 #define MODULAR false
```

```
5ed template < typename T > struct matrix : vector < vector < T >> {
        int n. m:
30f
        void print() {
603
            for (int i = 0; i < n; i++) {
                for (int j = 0; j < m; j++) cout << (*this)[i][j] << "
70f
1fb
                cout << endl;</pre>
            }
d98
101
       }
        matrix(int n_, int m_, bool ident = false) :
aa3
                vector < vector < T >> (n_, vector < T > (m_, 0)), n(n_), m(m_)  {
b14
94e
            if (ident) {
df7
                assert(n == m):
                for (int i = 0; i < n; i++) (*this)[i][i] = 1;
a89
            }
359
527
        }
        matrix(const vector<vector<T>>& c) : vector<vector<T>>(c),
b83
            n(c.size()). m(c[0].size()) {}
a3d
        matrix(const initializer_list<initializer_list<T>>& c) {
efc
f7e
            vector < vector < T >> val:
            for (auto& i : c) val.push_back(i);
303
            *this = matrix(val);
        }
c50
388
        matrix<T> operator*(matrix<T>& r) {
```

```
1e2
            assert(m == r.n):
82c
            matrix <T> M(n, r.m):
d69
            for (int i = 0; i < n; i++) for (int k = 0; k < m; k++)
df4
                for (int j = 0; j < r.m; j++) {
                    T \text{ add} = (*this)[i][k] * r[k][j];
e34
f98 #if MODULAR
d41 #warning Usar matrix<ll> e soh colocar valores em [0, MOD] na
   matriz!
8b6
                    M[i][i] += add%MOD;
                     if (M[i][j] >= MOD) M[i][j] -= MOD;
983
8c1 #else
7bb
                    M[i][i] += add:
f2e #endif
620
                }
474
            return M;
394
528
        matrix<T> operator^(ll e){
            matrix<T> M(n, n, true), at = *this;
f10
c87
            while (e) {
2e2
                if (e\&1) M = M*at;
cc2
                e >>= 1;
c80
                at = at*at;
eb6
474
            return M;
ca3
582
        void apply_transform(matrix M, ll e){
            auto& v = *this;
1c3
c87
            while (e) {
                if (e\&1) v = M*v;
9ba
cc2
                e >>= 1;
419
                M = M * M :
d86
            }
4e5
        }
70d }:
```

7.6 Primitivas Geometricas

```
c83 typedef double ld;
e3b const ld DINF = 1e18;
43a const ld pi = acos(-1.0);
107 const ld eps = 1e-9;

b32 #define sq(x) ((x)*(x))

d97 bool eq(ld a, ld b) {
    return abs(a - b) <= eps;
bfc }</pre>
```

```
// a8b7d6
b2a struct pt { // ponto
        ld x, y;
c1e
        pt(1d x_{-} = 0, 1d y_{-} = 0) : x(x_{-}), y(y_{-}) {}
3dd
        bool operator < (const pt p) const {</pre>
5bc
059
            if (!eq(x, p.x)) return x < p.x;
f98
            if (!eq(y, p.y)) return y < p.y;</pre>
bb3
           return 0;
f61
        }
a83
        bool operator == (const pt p) const {
ed0
            return eq(x, p.x) and eq(y, p.y);
589
cb9
        pt operator + (const pt p) const { return pt(x+p.x, y+p.y); }
        pt operator - (const pt p) const { return pt(x-p.x, y-p.y); }
a24
4a8
        pt operator * (const ld c) const { return pt(x*c , y*c ); }
a60
        pt operator / (const ld c) const { return pt(x/c , y/c ); }
3b6
        ld operator * (const pt p) const { return x*p.x + y*p.y; }
        ld operator ^ (const pt p) const { return x*p.y - y*p.x; }
6df
        friend istream& operator >> (istream& in, pt& p) {
5ed
            return in >> p.x >> p.y;
e37
e45
       }
a8b }:
// 7ab617
b3a struct line { // reta
730
        pt p, q;
       line() {}
0d6
4b8 line(pt p_, pt q_) : p(p_), q(q_) {}
8d7
      friend istream& operator >> (istream& in, line& r) {
            return in >> r.p >> r.q;
4cb
858
       }
7ab }:
// PONTO & VETOR
// c684fb
364 ld dist(pt p, pt q) { // distancia
       return hypot(p.y - q.y, p.x - q.x);
c68 }
// 80f2b6
9d7 ld dist2(pt p, pt q) { // quadrado da distancia
       return sq(p.x - q.x) + sq(p.y - q.y);
f24
80f }
// cf7f33
```

```
483 ld norm(pt v) { // norma do vetor
        return dist(pt(0, 0), v);
cf7 }
// 404df7
589 ld angle(pt v) { // angulo do vetor com o eixo x
        ld ang = atan2(v.y, v.x);
6f8
        if (ang < 0) ang += 2*pi;</pre>
19c
        return ang;
404 }
// 1b1d4a
298 ld sarea(pt p, pt q, pt r) { // area com sinal
        return ((q-p)^(r-q))/2;
1b1 }
// 98c42f
e32 bool col(pt p, pt q, pt r) { // se p, q e r sao colin.
       return eq(sarea(p, q, r), 0);
e7d
98c }
// 85d09d
Ocd bool ccw(pt p, pt q, pt r) \{ // \text{ se p, q, r sao ccw} \}
        return sarea(p, q, r) > eps;
85d }
// 41a7b4
lef pt rotate(pt p, ld th) { // rotaciona o ponto th radianos
       return pt(p.x * cos(th) - p.y * sin(th),
ff1
                p.x * sin(th) + p.y * cos(th));
41a }
// e4ad5e
ab1 pt rotate90(pt p) { // rotaciona 90 graus
        return pt(-p.y, p.x);
e4a }
// RETA
// Ofb984
edc bool isvert(line r) { // se r eh vertical
       return eq(r.p.x, r.q.x);
Ofb }
// 726d68
099 bool isinseg(pt p, line r) { // se p pertence ao seg de r
        pt a = r.p - p, b = r.q - p;
f65
```

```
b04
        return eq((a \hat{b}), 0) and (a * b) < eps;
726 }
// a0a30b
98d ld get_t(pt v, line r) { // retorna t tal que t*v pertence a reta r
       return (r.p^r.q) / ((r.p-r.q)^v);
a0a }
// 2329fe
256 pt proj(pt p, line r) { // projecao do ponto p na reta r
bea
       if (r.p == r.q) return r.p;
       r.q = r.q - r.p; p = p - r.p;
9f8
       pt proj = r.q * ((p*r.q) / (r.q*r.q));
2cd
       return proj + r.p;
232 }
// 111fd2
d5c pt inter(line r, line s) { // r inter s
       if (eq((r.p - r.q) ^ (s.p - s.q), 0)) return pt(DINF, DINF);
       r.q = r.q - r.p, s.p = s.p - r.p, s.q = s.q - r.p;
205
       return r.q * get_t(r.q, s) + r.p;
543
111 }
// 35998c
676 bool interseg(line r, line s) { // se o seg de r intersecta o seg
19b
        if (isinseg(r.p, s) or isinseg(r.q, s)
c21
            or isinseg(s.p, r) or isinseg(s.q, r)) return 1;
9fa
       return ccw(r.p, r.q, s.p) != ccw(r.p, r.q, s.q) and
413
                ccw(s.p, s.q, r.p) != ccw(s.p, s.q, r.q);
359 }
// 1b72e1
fcb ld disttoline(pt p, line r) { // distancia do ponto a reta
       return 2 * abs(sarea(p, r.p, r.q)) / dist(r.p, r.q);
1b7 }
// 3679c0
bcc ld disttoseg(pt p, line r) { // distancia do ponto ao seg
       if ((r.q - r.p)*(p - r.p) < 0) return dist(r.p, p);
73d
951
       if ((r.p - r.q)*(p - r.q) < 0) return dist(r.q, p);
a19
       return disttoline(p, r);
367 }
// 222358
11d ld distseg(line a, line b) { // distancia entre seg
```

```
4df
        if (interseg(a, b)) return 0;
        ld ret = DINF;
349
        ret = min(ret, disttoseg(a.p, b));
341
        ret = min(ret, disttoseg(a.q, b));
ceb
        ret = min(ret, disttoseg(b.p, a));
093
        ret = min(ret, disttoseg(b.q, a));
448
edf
        return ret;
222 }
// POLIGONO
// corta poligono com a reta r deixando os pontos p tal que
// ccw(r.p, r.q, p)
// 2538f9
1a9 vector <pt> cut_polygon(vector <pt> v, line r) { // O(n)
        vector < pt > ret;
8af
8a4
        for (int j = 0; j < v.size(); j++) {</pre>
            if (ccw(r.p, r.q, v[j])) ret.push_back(v[j]);
dac
            if (v.size() == 1) continue;
dce
030
            line s(v[j], v[(j+1)\%v.size()]);
            pt p = inter(r, s);
ae3
a3d
            if (isinseg(p, s)) ret.push_back(p);
d44
8a1
        ret.erase(unique(ret.begin(), ret.end()), ret.end());
        if (ret.size() > 1 and ret.back() == ret[0]) ret.pop_back();
24d
edf
        return ret:
253 }
// distancia entre os retangulos a e b (lados paralelos aos eixos)
// assume que ta representado (inferior esquerdo, superior direito)
// 630253
5f5 ld dist_rect(pair<pt, pt> a, pair<pt, pt> b) {
        ld hor = 0. vert = 0:
080
        if (a.second.x < b.first.x) hor = b.first.x - a.second.x;</pre>
34b
f5f
        else if (b.second.x < a.first.x) hor = a.first.x - b.second.x:
4fd
        if (a.second.y < b.first.y) vert = b.first.y - a.second.y;</pre>
80a
        else if (b.second.y < a.first.y) vert = a.first.y - b.second.y;</pre>
96f
        return dist(pt(0, 0), pt(hor, vert));
630 }
// 5df9cf
13d ld polarea(vector<pt> v) { // area do poligono
9c5
        1d ret = 0;
        for (int i = 0; i < v.size(); i++)</pre>
c6e
80f
            ret += sarea(pt(0, 0), v[i], v[(i + 1) % v.size()]);
```

```
d03
        return abs(ret);
5df }
// se o ponto ta dentro do poligono: retorna O se ta fora,
// 1 se ta no interior e 2 se ta na borda
// a6423f
8e7 int inpol(vector\phi) { // O(n)
        int qt = 0;
8de
f14
        for (int i = 0; i < v.size(); i++) {</pre>
            if (p == v[i]) return 2;
bda
6af
            int j = (i+1)%v.size();
e38
            if (eq(p.y, v[i].y) and eq(p.y, v[j].y)) {
97f
                if ((v[i]-p)*(v[j]-p) < eps) return 2;</pre>
5e2
                continue:
            }
48b
388
            bool baixo = v[i].y+eps < p.y;</pre>
464
            if (baixo == (v[j].y+eps < p.y)) continue;</pre>
            auto t = (p-v[i])^(v[j]-v[i]);
366
            if (eq(t, 0)) return 2;
1b4
            if (baixo == (t > eps)) qt += baixo ? 1 : -1;
839
d13
        }
b84
        return qt != 0;
a64 }
// c58350
6ff bool interpol(vector<pt> v1, vector<pt> v2) { // se dois poligonos
   se intersectam - O(n*m)
        int n = v1.size(). m = v2.size();
7d1
       for (int i = 0; i < n; i++) if (inpol(v2, v1[i])) return 1;</pre>
c36
ab8
       for (int i = 0; i < n; i++) if (inpol(v1, v2[i])) return 1;
       for (int i = 0; i < n; i++) for (int j = 0; j < m; j++)
523
0c8
            if (interseg(line(v1[i], v1[(i+1)%n]), line(v2[i],
   v2[(i+1)%m])) return 1:
bb3
        return 0:
c58 }
// 12559f
494 ld distpol(vector<pt> v1, vector<pt> v2) { // distancia entre
f6b
        if (interpol(v1, v2)) return 0;
349
        ld ret = DINF;
       for (int i = 0; i < v1.size(); i++) for (int j = 0; j <</pre>
   v2.size(); j++)
            ret = min(ret, distseg(line(v1[i], v1[(i + 1) %
   v1.size()]).
```

```
9d9
                         line(v2[j], v2[(j + 1) % v2.size()])));
edf
        return ret:
125 }
// 10d7e0
138 vector<pt> convex_hull(vector<pt> v) { // convex hull - O(n log(n))
        sort(v.begin(), v.end());
        v.erase(unique(v.begin(), v.end()), v.end());
d76
52d
        if (v.size() <= 1) return v;</pre>
        vector<pt> 1, u;
526
f14
        for (int i = 0; i < v.size(); i++) {</pre>
fb2
            while (1.size() > 1 \text{ and } !ccw(1.end()[-2], 1.end()[-1],
   v[i]))
364
                1.pop_back();
c35
            1.push_back(v[i]);
58e
3e9
        for (int i = v.size() - 1; i >= 0; i--) {
            while (u.size() > 1 \text{ and } !ccw(u.end()[-2], u.end()[-1],
f19
   v[i]))
7a8
                u.pop_back();
            u.push_back(v[i]);
a95
0b8
cfc
        1.pop_back(); u.pop_back();
82b
        for (pt i : u) l.push_back(i);
792
        return 1:
10d }
483 struct convex_pol {
        vector<pt> pol;
        // nao pode ter ponto colinear no convex hull
d98
        convex_pol() {}
        convex_pol(vector < pt > v) : pol(convex_hull(v)) {}
a04
        // se o ponto ta dentro do hull - O(\log(n))
        // 6b097f
        bool is_inside(pt p) {
8af
b6e
            if (pol.size() == 0) return false;
            if (pol.size() == 1) return p == pol[0];
eae
67f
            int 1 = 1, r = pol.size();
            while (1 < r) {
40 c
                int m = (1+r)/2;
ee4
48f
                if (ccw(p, pol[0], pol[m])) l = m+1;
                else r = m:
ef3
91c
00a
            if (1 == 1) return isinseg(p, line(pol[0], pol[1]));
9e7
            if (1 == pol.size()) return false;
```

```
1 c 0
             return !ccw(p, pol[1], pol[1-1]);
6b0
        // ponto extremo em relacao a cmp(p, q) = p mais extremo q
        // (copiado de https://github.com/gustavoM32/caderno-zika)
        // 56ccd2
        int extreme(const function < bool(pt, pt) > & cmp) {
719
b1c
            int n = pol.size();
4a2
             auto extr = [&](int i, bool& cur_dir) {
22a
                 \operatorname{cur}_{\operatorname{dir}} = \operatorname{cmp}(\operatorname{pol}[(i+1)\%n], \operatorname{pol}[i]);
                 return !cur_dir and !cmp(pol[(i+n-1)%n], pol[i]);
61a
364
            };
63d
            bool last_dir, cur_dir;
a0d
            if (extr(0, last dir)) return 0;
            int 1 = 0, r = n;
993
ead
            while (1+1 < r) {
                 int m = (1+r)/2;
ee4
f29
                 if (extr(m, cur_dir)) return m;
                 bool rel_dir = cmp(pol[m], pol[l]);
44a
                 if ((!last_dir and cur_dir) or
b18
                          (last dir == cur dir and rel dir == cur dir)) {
261
8a6
                     1 = m:
                     last_dir = cur_dir;
1f1
94a
                 } else r = m;
606
            }
792
             return 1;
56c
316
        int max_dot(pt v) {
ec1
             return extreme([&](pt p, pt q) { return p*v > q*v; });
3b7
a54
        pair < int , int > tangents(pt p) {
             auto L = [\&](pt q, pt r) \{ return ccw(p, r, q); \};
ffb
8fd
             auto R = [k](pt q, pt r) \{ return ccw(p, q, r); \};
            return {extreme(L), extreme(R)};
fa8
736
        }
3ec };
// CIRCUNFERENCIA
// a125e4
911 pt getcenter(pt a, pt b, pt c) { // centro da circunf dado 3 pontos
        b = (a + b) / 2:
174
        c = (a + c) / 2;
2ae
       return inter(line(b, b + rotate90(a - b)),
98b
                line(c, c + rotate90(a - c)));
3f8
a12 }
// cd80c0
```

```
4b3 vector <pt> circ_line_inter(pt a, pt b, pt c, ld r) { // intersecao
    da circunf (c, r) e reta ab
        vector<pt> ret;
8af
        b = b-a, a = a-c;
f2b
4b1
        1d A = b*b:
20a
        1d B = a*b:
2e9
        1d C = a*a - r*r:
        1d D = B*B - A*C;
1fa
818
        if (D < -eps) return ret;</pre>
        ret.push_back(c+a+b*(-B+sqrt(D+eps))/A);
dc5
20e
        if (D > eps) ret.push_back(c+a+b*(-B-sqrt(D))/A);
edf
        return ret:
cd8 }
// fb11d8
ad2 vector<pt> circ_inter(pt a, pt b, ld r, ld R) { // intersecao da
    circunf (a, r) e (b, R)
8af
        vector<pt> ret;
b7e
        1d d = dist(a, b);
        if (d > r+R \text{ or } d+min(r, R) < max(r, R)) return ret;
5ce
        1d x = (d*d-R*R+r*r)/(2*d):
398
183
        1d v = sqrt(r*r-x*x);
325
        pt v = (b-a)/d;
76e
        ret.push_back(a+v*x + rotate90(v)*y);
2cb
        if (y > 0) ret.push_back(a+v*x - rotate90(v)*y);
edf
        return ret:
fb1 }
// 3a44fb
6e0 bool operator <(const line& a, const line& b) { // comparador pra
   reta
        // assume que as retas tem p < q</pre>
        pt v1 = a.q - a.p, v2 = b.q - b.p;
a13
f82
        if (!eq(angle(v1), angle(v2))) return angle(v1) < angle(v2);</pre>
780
        return ccw(a.p, a.q, b.p); // mesmo angulo
27e }
b14 bool operator ==(const line& a, const line& b) {
76c
        return !(a < b) and !(b < a);
449 }
// comparador pro set pra fazer sweep line com segmentos
// 36729f
2c4 struct cmp_sweepline {
        bool operator () (const line& a, const line& b) const {
            // assume que os segmentos tem p < q</pre>
191
            if (a.p == b.p) return ccw(a.p, a.q, b.q);
231
            if (!eq(a.p.x, a.q.x)) and (eq(b.p.x, b.q.x)) or a.p.x+eps <
```

```
b.p.x))
780
                return ccw(a.p, a.q, b.p);
dc0
            return ccw(a.p, b.q, b.p);
243
       }
367 }:
// comparador pro set pra fazer sweep angle com segmentos
// f778aa
bef pt dir;
5b0 struct cmp_sweepangle {
d80
        bool operator () (const line& a, const line& b) const {
            return get_t(dir, a) + eps < get_t(dir, b);</pre>
653
97f };
```

7.7 Primitivas Geometricas 3D

```
c83 typedef double ld;
e3b const ld DINF = 1e18:
107 const ld eps = 1e-9;
b32 #define sq(x) ((x)*(x))
d97 bool eq(ld a, ld b) {
ba0
            return abs(a - b) <= eps;</pre>
bfc }
// 3eef01
b2a struct pt { // ponto
2eb
            ld x, y, z;
            pt(1d x_{-} = 0, 1d y_{-} = 0, 1d z_{-} = 0) : x(x_{-}), y(y_{-}), z(z_{-})
a50
  {}
5bc
            bool operator < (const pt p) const {</pre>
059
                    if (!eq(x, p.x)) return x < p.x;
f98
                    if (!eq(y, p.y)) return y < p.y;
                    if (!eq(z, p.z)) return z < p.z;
44c
bb3
                    return 0;
            }
6cd
a83
            bool operator == (const pt p) const {
41 c
                     return eq(x, p.x) and eq(y, p.y) and eq(z, p.z);
fb5
44b
            pt operator + (const pt p) const { return pt(x+p.x, y+p.y,
   z+p.z); }
            pt operator - (const pt p) const { return pt(x-p.x, y-p.y,
   z-p.z); }
fb7
            pt operator * (const ld c) const { return pt(x*c , y*c ,
   z*c ); }
```

```
pt operator / (const ld c) const { return pt(x/c , y/c ,
   z/c ); }
            ld operator * (const pt p) const { return x*p.x + y*p.y +
a65
   z*p.z; }
            pt operator ^ (const pt p) const { return pt(y*p.z -
7f6
   z*p.y, z*p.x - x*p.z, x*p.y - y*p.x); }
5ed
            friend istream& operator >> (istream& in, pt& p) {
                    return in >> p.x >> p.y >> p.z;
9bf
5e8
3ee }:
// 7ab617
b3a struct line { // reta
730
            pt p, q;
0d6
            line() {}
4b8
            line(pt p_, pt q_) : p(p_), q(q_) {}
8d7
            friend istream& operator >> (istream& in, line& r) {
4cb
                    return in >> r.p >> r.q;
858
            }
7ab }:
// d5d580
79b struct plane { // plano
7e1
            array<pt, 3> p; // pontos que definem o plano
29b
            array < ld, 4 > eq; // equacao do plano
bb7
            plane() {}
fb0
            plane(pt p_, pt q_, pt r_) : p({p_, q_, r_}) { build(); }
ca9
            friend istream& operator >> (istream& in, plane& P) {
2ab
                    return in >> P.p[0] >> P.p[1] >> P.p[2];
70e
                    P.build():
544
            void build() {
0a8
da2
                    pt dir = (p[1] - p[0]) ^ (p[2] - p[0]);
7d5
                    eq = \{dir.x, dir.y, dir.z, dir*p[0]*(-1)\};
41a
            }
d5d };
// converte de coordenadas polares para cartesianas
// (angulos devem estar em radianos)
// phi eh o angulo com o eixo z (cima) theta eh o angulo de rotacao ao
   redor de z
// a4f17f
2fb pt convert(ld rho, ld th, ld phi) {
            return pt(sin(phi) * cos(th), sin(phi) * sin(th),
   cos(phi)) * rho;
a4f }
```

```
// projecao do ponto p na reta r
// 2329fe
256 pt proj(pt p, line r) {
            if (r.p == r.q) return r.p;
bea
97a
            r.q = r.q - r.p; p = p - r.p;
9f8
            pt proj = r.q * ((p*r.q) / (r.q*r.q));
2cd
            return proj + r.p;
232 }
// projecao do ponto p no plano P
// 4a0d14
b1a pt proj(pt p, plane P) {
            p = p - P.p[0], P.p[1] = P.p[1] - P.p[0], P.p[2] = P.p[2]
   - P.p[0];
b69
            pt norm = P.p[1] ^ P.p[2];
            pt proj = p - (norm * (norm * p) / (norm*norm));
6ab
467
           return proj + P.p[0];
4a0 }
// distancia
// 2d06b0
a45 ld dist(pt a, pt b) {
fd9
            return sqrt(sq(a.x-b.x) + sq(a.y-b.y) + sq(a.z-b.z));
2d0 }
// distancia ponto reta
// 3c4e1b
137 ld distline(pt p, line r) {
            return dist(p, proj(p, r));
3c4 }
// distancia de ponto para segmento
// 42cbbd
d43 ld distseg(pt p, line r) {
            if ((r.q - r.p)*(p - r.p) < 0) return dist(r.p, p);
951
           if ((r.p - r.q)*(p - r.q) < 0) return dist(r.q, p);
200
           return distline(p, r);
42c }
// distancia de ponto a plano com sinal
// d490d9
7cc ld sdist(pt p, plane P) {
           return P.eq[0]*p.x + P.eq[1]*p.y + P.eq[2]*p.z + P.eq[3];
150
d49 }
// distancia de ponto a plano
```

```
// 33dc8c
768 ld distplane(pt p, plane P) {
            return abs(sdist(p, P));
33d }
// se ponto pertence a reta
// 31a295
099 bool isinseg(pt p, line r) {
a32
            return eq(distseg(p, r), 0);
31a }
// se ponto pertence ao triangulo definido por P.p
// c81f7e
cd2 bool isinpol(pt p, vector<pt> v) {
fad
            assert(v.size() >= 3);
bf4
            pt norm = (v[1]-v[0]) ^ (v[2]-v[1]);
8a4
            bool inside = true;
            int sign = -1;
cec
f14
            for (int i = 0; i < v.size(); i++) {</pre>
834
                    line r(v(i+1)\%3), v(i):
                    if (isinseg(p, r)) return true;
2a9
4ef
                    pt ar = v[(i+1)\%3] - v[i];
320
                    if (sign == -1) sign = ((ar^(p-v[i]))*norm > 0);
82b
                    else if (((ar^(p-v[i]))*norm > 0) != sign) inside
   = false:
15e
aca
            return inside:
c81 }
// distancia de ponto ate poligono
// a8d4c2
361 ld distpol(pt p, vector<pt> v) {
3e7
            pt p2 = proj(p, plane(v[0], v[1], v[2]));
61a
            if (isinpol(p2, v)) return dist(p, p2);
349
            ld ret = DINF:
f14
            for (int i = 0; i < v.size(); i++) {</pre>
6af
                    int j = (i+1)%v.size();
5ee
                    ret = min(ret, distseg(p, line(v[i], v[j])));
7b2
edf
            return ret;
a8d }
// intersecao de plano e segmento
// BOTH = o segmento esta no plano
// ONE = um dos pontos do segmento esta no plano
// PARAL = segmento paralelo ao plano
```

```
// CONCOR = segmento concorrente ao plano
// e2ecac
e51 enum RETCODE {BOTH, ONE, PARAL, CONCOR};
26b pair < RETCODE, pt > intersect(plane P, line r) {
       1d d1 = sdist(r.p. P):
f8f
        1d d2 = sdist(r.q, P);
53a
       if (eq(d1, 0) \text{ and } eq(d2, 0))
                   return pair(BOTH, r.p);
504
72c
       if (eq(d1, 0))
847
                    return pair(ONE, r.p);
485
        if (eq(d2, 0))
168
                   return pair(ONE, r.q);
3fb
        if ((d1 > 0 \text{ and } d2 > 0) \text{ or } (d1 < 0 \text{ and } d2 < 0))
            if (eq(d1-d2, 0)) return pair(PARAL, pt());
463
406
            return pair(CONCOR, pt());
91c
c84
        1d frac = d1 / (d1 - d2);
3ff
        pt res = r.p + ((r.q - r.p) * frac);
394
        return pair(ONE, res);
b92 }
// rotaciona p ao redor do eixo u por um angulo a
// 7f0a40
787 pt rotate(pt p, pt u, ld a) {
773
           u = u / dist(u, pt());
e6f
           return u * (u * p) + (u ^ p ^ u) * cos(a) + (u ^ p) *
   sin(a):
7f0 }
```

7.8 Primitivas Geometricas Inteiras

```
2de #define sq(x) ((x)*(11)(x))
// 840720
b2a struct pt { // ponto
e91
        int x, y;
        pt(int x_{-} = 0, int y_{-} = 0) : x(x_{-}), y(y_{-}) {}
df1
5bc
       bool operator < (const pt p) const {</pre>
95a
          if (x != p.x) return x < p.x;
89 c
           return y < p.y;</pre>
dcd
        bool operator == (const pt p) const {
a83
d74
            return x == p.x and y == p.y;
7b4
cb9
        pt operator + (const pt p) const { return pt(x+p.x, y+p.y); }
a24
        pt operator - (const pt p) const { return pt(x-p.x, y-p.y); }
        pt operator * (const int c) const { return pt(x*c, y*c); }
0ef
```

```
11 operator * (const pt p) const { return x*(11)p.x +
   y*(11)p.y; }
        11 operator ^ (const pt p) const { return x*(11)p.y -
        friend istream& operator >> (istream& in, pt& p) {
5ed
            return in >> p.x >> p.y;
e37
e45
       }
840 }:
// 7ab617
b3a struct line { // reta
        pt p, q;
0d6
       line() {}
        line(pt p_, pt q_) : p(p_), q(q_) {}
4b8
       friend istream& operator >> (istream& in, line& r) {
8d7
4cb
            return in >> r.p >> r.q;
858
       }
7ab };
// PONTO & VETOR
// 51563e
ea8 ll dist2(pt p, pt q) { // quadrado da distancia
       return sq(p.x - q.x) + sq(p.y - q.y);
515 }
// bf431d
5a2 ll sarea2(pt p, pt q, pt r) \{ // 2 * area com sinal \}
       return (q-p)^(r-q);
bf4 }
// a082d3
e32 bool col(pt p, pt q, pt r) \{ // \text{ se p, q e r sao colin.} \}
        return sarea2(p, q, r) == 0:
a08 }
// 42bb09
Ocd bool ccw(pt p, pt q, pt r) { // se p, q, r sao ccw
276 return sarea2(p, q, r) > 0;
42b }
// fcf924
c31 int quad(pt p) { // quadrante de um ponto
        return (p.x<0)^3*(p.y<0);
dbb
fcf }
// 77187b
```

```
2df bool compare_angle(pt p, pt q) { // retorna se ang(p) < ang(q)</pre>
        if (quad(p) != quad(q)) return quad(p) < quad(q);</pre>
9fc
ea1
        return ccw(q, pt(0, 0), p);
771 }
// e4ad5e
ab1 pt rotate90(pt p) { // rotaciona 90 graus
        return pt(-p.y, p.x);
e4a }
// RETA
// c9f07f
099 bool isinseg(pt p, line r) { // se p pertence ao seg de r
f65
        pt a = r.p - p, b = r.q - p;
2ac
        return (a ^ b) == 0 and (a * b) <= 0;
c9f }
// 35998c
676 bool interseg(line r, line s) { // se o seg de r intersecta o seg
   de s
        if (isinseg(r.p, s) or isinseg(r.q, s)
19b
            or isinseg(s.p, r) or isinseg(s.q, r)) return 1;
c21
9fa
        return ccw(r.p, r.q, s.p) != ccw(r.p, r.q, s.q) and
413
                ccw(s.p, s.q, r.p) != ccw(s.p, s.q, r.q);
359 }
// dd8702
9e0 int segpoints(line r) { // numero de pontos inteiros no segmento
        return 1 + \_gcd(abs(r.p.x - r.q.x), abs(r.p.y - r.q.y));
dd8 }
// d273be
88a double get_t(pt v, line r) { // retorna t tal que t*v pertence a
        return (r.p^r.q) / (double) ((r.p-r.q)^v);
1ad
d27 }
// POLIGONO
// quadrado da distancia entre os retangulos a e b (lados paralelos
   aos eixos)
// assume que ta representado (inferior esquerdo, superior direito)
// e13018
485 ll dist2_rect(pair<pt, pt> a, pair<pt, pt> b) {
        int hor = 0, vert = 0;
c59
```

```
34b
        if (a.second.x < b.first.x) hor = b.first.x - a.second.x;</pre>
        else if (b.second.x < a.first.x) hor = a.first.x - b.second.x;</pre>
f5f
4fd
        if (a.second.y < b.first.y) vert = b.first.y - a.second.y;</pre>
80a
        else if (b.second.y < a.first.y) vert = a.first.y - b.second.y;</pre>
869
        return sq(hor) + sq(vert);
e13 }
// d5f693
9c3 ll polarea2(vector<pt> v) { // 2 * area do poligono
        ll ret = 0:
сбе
        for (int i = 0; i < v.size(); i++)</pre>
532
            ret += sarea2(pt(0, 0), v[i], v[(i + 1) % v.size()]);
d03
        return abs(ret):
d5f }
// se o ponto ta dentro do poligono: retorna 0 se ta fora,
// 1 se ta no interior e 2 se ta na borda
// afd587
8e7 int inpol(vector\phi) { // O(n)
        int qt = 0;
8de
f14
        for (int i = 0; i < v.size(); i++) {</pre>
            if (p == v[i]) return 2;
bda
            int j = (i+1)%v.size();
6af
cc6
            if (p.y == v[i].y and p.y == v[j].y) {
547
                 if ((v[i]-p)*(v[j]-p) <= 0) return 2;</pre>
5e2
                 continue:
b47
            }
78c
            bool baixo = v[i].y < p.y;</pre>
057
            if (baixo == (v[j].y < p.y)) continue;</pre>
366
            auto t = (p-v[i])^(v[j]-v[i]);
2ad
            if (!t) return 2:
0bb
            if (baixo == (t > 0)) qt += baixo ? 1 : -1;
9cf
        }
b84
        return qt != 0;
afd }
// 10d7e0
138 vector <pt > convex_hull(vector <pt > v) { // convex hull - O(n log(n))
fca
        sort(v.begin(), v.end());
d76
        v.erase(unique(v.begin(), v.end()), v.end());
        if (v.size() <= 1) return v;</pre>
52d
        vector<pt> 1, u;
526
f14
        for (int i = 0; i < v.size(); i++) {</pre>
             while (l.size() > 1 and !ccw(l.end()[-2], l.end()[-1],
fb2
   v[i]))
364
                 1.pop_back();
            l.push_back(v[i]);
c35
```

```
58e
        for (int i = v.size() - 1; i >= 0; i--) {
3e9
f19
            while (u.size() > 1 \text{ and } !ccw(u.end()[-2], u.end()[-1],
   v[i]))
7a8
                u.pop_back();
            u.push_back(v[i]);
a95
0b8
       }
cfc
       1.pop_back(); u.pop_back();
82b
        for (pt i : u) l.push_back(i);
792
        return 1:
10d }
// af2d96
786 ll interior_points(vector<pt> v) { // pontos inteiros dentro de um
   poligono simples
        11 b = 0:
c4e
c6e
        for (int i = 0; i < v.size(); i++)</pre>
            b += segpoints(line(v[i], v[(i+1)\%v.size()])) - 1;
Осе
        return (polarea2(v) - b) / 2 + 1;
a1c
af2 }
483 struct convex_pol {
f50
        vector<pt> pol;
        // nao pode ter ponto colinear no convex hull
d98
        convex pol() {}
a04
        convex_pol(vector<pt> v) : pol(convex_hull(v)) {}
        // se o ponto ta dentro do hull - O(\log(n))
        // 6b097f
        bool is_inside(pt p) {
8af
b6e
            if (pol.size() == 0) return false;
            if (pol.size() == 1) return p == pol[0];
eae
            int l = 1, r = pol.size();
67f
            while (1 < r) {
40c
                int m = (1+r)/2:
ee4
                if (ccw(p, pol[0], pol[m])) l = m+1;
48f
ef3
                else r = m;
91c
00a
            if (1 == 1) return isinseg(p, line(pol[0], pol[1]));
            if (1 == pol.size()) return false;
9e7
1c0
            return !ccw(p, pol[1], pol[1-1]);
6b0
        // ponto extremo em relacao a cmp(p, q) = p mais extremo q
        // (copiado de https://github.com/gustavoM32/caderno-zika)
        // 56ccd2
719
        int extreme(const function < bool(pt, pt) > & cmp) {
```

```
b1c
             int n = pol.size();
4a2
             auto extr = [&](int i, bool& cur_dir) {
22a
                 \operatorname{cur}_{\operatorname{dir}} = \operatorname{cmp}(\operatorname{pol}[(i+1)\%n], \operatorname{pol}[i]);
61a
                 return !cur_dir and !cmp(pol[(i+n-1)%n], pol[i]);
364
             };
63d
             bool last_dir, cur_dir;
a0d
             if (extr(0, last dir)) return 0:
993
             int 1 = 0, r = n;
             while (1+1 < r) {
ead
                 int m = (1+r)/2:
ee4
f29
                 if (extr(m, cur_dir)) return m;
44a
                 bool rel_dir = cmp(pol[m], pol[l]);
b18
                 if ((!last dir and cur dir) or
261
                          (last_dir == cur_dir and rel_dir == cur_dir)) {
8a6
                     1 = m:
1f1
                     last_dir = cur_dir;
94a
                 } else r = m;
606
792
             return 1;
        }
56c
316
        int max_dot(pt v) {
             return extreme([&](pt p, pt q) { return p*v > q*v; });
ec1
3b7
a54
        pair < int , int > tangents(pt p) {
ffb
             auto L = [\&](pt q, pt r) \{ return ccw(p, r, q); \};
8fd
             auto R = [\&](pt q, pt r) \{ return ccw(p, q, r); \};
fa8
             return {extreme(L), extreme(R)};
736
        }
3ec };
// dca598
6e0 bool operator <(const line& a, const line& b) { // comparador pra
        // assume que as retas tem p < q
        pt v1 = a.q - a.p, v2 = b.q - b.p;
a13
        bool b1 = compare_angle(v1, v2), b2 = compare_angle(v2, v1);
036
73c
        if (b1 or b2) return b1;
780
        return ccw(a.p, a.q, b.p); // mesmo angulo
b61 }
b14 bool operator ==(const line& a, const line& b) {
76c
        return !(a < b) and !(b < a):
449 }
// comparador pro set pra fazer sweep line com segmentos
// 6774df
2c4 struct cmp_sweepline {
        bool operator () (const line& a, const line& b) const {
```

```
// assume que os segmentos tem p < q
191
            if (a.p == b.p) return ccw(a.p, a.q, b.q);
614
            if (a.p.x != a.q.x and (b.p.x == b.q.x or a.p.x < b.p.x))
780
                return ccw(a.p, a.q, b.p);
            return ccw(a.p, b.q, b.p);
dc0
baf
677 };
// comparador pro set pra fazer sweep angle com segmentos
// 1ee7f5
bef pt dir;
5b0 struct cmp_sweepangle {
       bool operator () (const line& a, const line& b) const {
261
            return get_t(dir, a) < get_t(dir, b);</pre>
dc5
f6d };
```

8 Extra

8.1 fastIO.cpp

int read_int() {

bool minus = false;

```
int result = 0;
    char ch;
    ch = getchar();
    while (1) {
       if (ch == '-') break;
        if (ch >= '0' && ch <= '9') break;
        ch = getchar();
    if (ch == '-') minus = true;
    else result = ch-'0';
    while (1) {
        ch = getchar();
        if (ch < '0' || ch > '9') break;
        result = result *10 + (ch - '0');
    if (minus) return -result;
    else return result;
8.2 vimrc
set ts=4 si ai sw=4 nu mouse=a undofile
svntax on
vnoremap <C-H>:w !sed '/^\#w/d' \| cpp -dD -P -fpreprocessed \| tr -d
   '[:space:]' \| md5sum \| cut -c-6<CR>
8.3 timer.cpp
// timer T; T() -> retorna o tempo em ms desde que declarou
using namespace chrono;
struct timer : high_resolution_clock {
    const time_point start;
    timer(): start(now()) {}
    int operator()() {
        return duration_cast < milliseconds > (now() - start).count();
};
```

8.4 rand.cpp

```
mt19937 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
int uniform(int 1, int r){
    uniform_int_distribution < int > uid(1, r);
    return uid(rng);
}
8.5 template.cpp
#include <bits/stdc++.h>
using namespace std;
#define _ ios_base::sync_with_stdio(0);cin.tie(0);
#define endl '\n'
typedef long long 11;
const int INF = 0x3f3f3f3f;
const 11 LINF = 0x3f3f3f3f3f3f3f3f11;
int main() {
    exit(0);
}
    debug.cpp
void debug_out(string s, int line) { cerr << endl; }</pre>
template < typename H, typename ... T>
void debug_out(string s, int line, H h, T... t) {
    if (s[0] != ',') cerr << "Line(" << line << ") ";</pre>
    do { cerr << s[0]; s = s.substr(1);</pre>
    } while (s.size() and s[0] != ',');
    cerr << " = " << h;
    debug_out(s, line, t...);
}
#ifdef DEBUG
#define debug(...) debug_out(#__VA_ARGS__, __LINE__, __VA_ARGS__)
#else
#define debug(...) 42
#endif
    stress.sh
P=a
```

```
make ${P} ${P}2 gen || exit 1
for ((i = 1; ; i++)) do
    ./gen $i > in
    ./${P} < in > out
    ./${P}2 < in > out2
    if (! cmp -s out out2) then
        echo "--> entrada:"
        cat in
        echo "--> saida1:"
        cat out
        echo "--> saida2:"
        cat out2
        break:
   fi
    echo $i
done
8.8 makefile
CXX = g++
CXXFLAGS = -fsanitize=address, undefined -fno-omit-frame-pointer -g
   -Wall -Wshadow -std=c++17 -Wno-unused-result -Wno-sign-compare
   -Wno-char-subscripts #-fuse-ld=gold
8.9 gethash.sh
# Para usar:
# bash gethash.sh arquivo.cpp
echo "" > pref.txt
while IFS= read -r 1; do
    echo "$1" >> pref.txt
    echo "$1" > line.txt
    hp=$(echo $(bash hash.sh pref.txt 1 1000) | cut -c-2)
    hl=$(echo $(bash hash.sh line.txt 1 1000) | cut -c-2)
    echo -e "$hp $hl $1"
done < "$1"
8.10 hash.sh
# Para usar (hash das linhas [11, 12]):
# bash hash.sh arquivo.cpp 11 12
sed -n $2','$3' p' $1 | sed '/^#w/d' | cpp -dD -P -fpreprocessed | tr
   -d '[:space:]' | md5sum | cut -c-6
```