Humuhumunukunukuapua'a UFMG

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Índice			1.17 SegTreap	
		_	1.18 SegTree	15
1	Estruturas	5	1.19 SegTree 2D Iterativa	16
	1.1 BIT	5	1.20 SegTree Beats	17
	1.2 BIT 2D	5	1.21 SegTree Colorida	
	1.3 BIT com update em range	6	1.22 SegTree Esparsa - Lazy	
	1.4 BIT-Sort Tree	6	1.23 SegTree Esparsa - O(q) memoria	
	1.5 Convex Hull Trick Dinamico	6		
	1.6 Convex Hull Trick Estatico	7	1.24 SegTree Iterativa	
	1.7 DSU	7	1.25 SegTree Iterativa com Lazy Propagation	
	1.8 Li-Chao Tree	9	1.26 SegTree PA	23
	1.9 Li-Chao Tree - Lazy	9	1.27 SegTree Persistente	24
	1.10 MergeSort Tree	10	1.28 SegTree Persistente com Lazy	24
	1.11 Min queue - deque	12	1.29 Sparse Table	25
	1.12 Min queue - stack	12	1.30 Sparse Table Disjunta	25
	1.13 Order Statistic Set	12	1.31 Splay Tree	26
	1.14 Priority Queue DS	13	1.32 Splay Tree Implicita	27
	1.15 Range color	13	1.33 Split-Merge Set	29
	1.16 RMQ $<$ O(n), O(1) $>$ - min queue		1.34 SQRT Tree	31

	1.35 Treap	31		2.21 Kosaraju	49
	1.36 Treap Implicita	32		2.22 Kruskal	49
	1.37 Treap Persistent Implicita	33		2.23 Kuhn	50
	1.38 Wavelet Tree	34		2.24 LCA com binary lifting	50
2	Grafos	35		2.25 LCA com HLD	51
4	2.1 AGM Direcionada	35		2.26 LCA com RMQ	52
				2.27 Line Tree	52
		36		2.28 Link-cut Tree	53
	2.3 Bellman-Ford			2.29 Link-cut Tree - aresta	54
	2.4 Block-Cut Tree			2.30 Link-cut Tree - vertice	55
		37		2.31 Max flow com lower bound	56
	2.6 Centro de arvore	38		2.32 MinCostMaxFlow	57
	2.7 Centroid			2.33 Prufer code	58
	2.8 Centroid decomposition	39		2.34 Sack (DSU em arvores)	59
	2.9 Centroid Tree	40		2.35 Stable Marriage	59
	2.10 Dijkstra	40		2.36 Tarjan para SCC	
	2.11 Dinitz	40		2.37 Topological Sort	
	2.12 Dominator Tree	41		2.38 Vertex cover	
	2.13 Euler Path / Euler Cycle	42		2.39 Virtual Tree	
	2.14 Euler Tour Tree	43		2.66 Virtual 11cc	O1
	2.15 Floyd-Warshall	45	3	Problemas	61
	2.16 Functional Graph	45		3.1 Angle Range Intersection	61
	2.17 HLD - aresta	46		3.2 Area da Uniao de Retangulos	62
	2.18 HLD - vertice	47		3.3 Area Maxima de Histograma	63
	2.19 HLD sem Update	48		3.4 Binomial modular	63
	2.20 Isomorfismo de arvores	48		3.5 Closest pair of points	64

3.6	Coloracao de Grafo de Intervalo	64	3.32 Sequencia de de Brujin	79
3.7	Conectividade Dinamica DC	65	3.33 Shortest Addition Chain	79
3.8	Conectividade Dinamica LCT	65	3.34 Simple Polygon	79
3.9	Conj. Indep. Maximo com Peso em Grafo de Intervalo	67	3.35 Sweep Direction	80
3.10	Convex Hull Dinamico	67	3.36 Triangulacao de Delaunay	80
3.11	Distancia maxima entre dois pontos	68	3.37 Triangulos em Grafos	82
3.12	Distinct Range Query	68		
3.13	Distinct Range Query com Update	69	Matematica	82
3.14	Dominator Points	70	4.1 2-SAT	82
3.15	DP de Dominacao 3D	70	4.2 Avaliacao de Interpolacao	83
3.16	Gray Code	71	4.3 Berlekamp-Massey	83
3.17	Half-plane intersection	71	4.4 Binomial Distribution	84
3.18	Heap Sort	71	4.5 Convolucao de GCD / LCM	84
3.19	Hungaro	72	4.6 Coprime Basis	85
3.20	Inversion Count	72	4.7 Crivo de Eratosthenes	85
3.21	LIS - recupera	72	4.8 Deteccao de ciclo - Tortoise and Hare	87
3.22	LIS - tamanho	73	4.9 Division Trick	87
3.23	Minimum Enclosing Circle	73	4.10 Equação Diofantina Linear	87
3.24	Minkowski Sum	74	4.11 Euclides estendido	88
3.25	MO	74	4.12 Exponenciacao rapida	88
3.26	MO - DSU	75	4.13 Fast Walsh Hadamard Transform	88
3.27	MO em Arvores	76	4.14 FFT	88
3.28	Palindromic Factorization	77	4.15 Gauss	90
3.29	Parsing de Expressao	77	4.16 Gauss - Z2	90
3.30	RMQ com Divide and Conquer	78	4.17 Integracao Numerica	91
3.31	Segment Intersection	78	4.18 Inverso Modular	91

	4.19	Karatsuba	91		6.11	Suffix Automaton	. 105
	4.20	Logaritmo Discreto	92		6.12	Trie	. 105
	4.21	Miller-Rabin	92		6.13	Z	. 106
	4.22	Pollard's Rho Alg	92				
	4.23	Produto de dois long long mod m	93	7	Prir	nitivas	106
	4.24	Simplex	93		7.1	Aritmetica Modular	. 106
	4.25	Teorema Chines do Resto	94		7.2	Big Integer	. 107
	4.26	Totiente	94		7.3	Fracao	. 110
5	DP		95		7.4	Geometria	. 110
	5.1	Divide and Conquer DP	95		7.5	Geometria - inteiro	. 114
	5.2	Longest Common Subsequence	95		7.6	Geometria 3D	. 117
	5.3	Mochila	96			Matriz	
	5.4	SOS DP	96				
	5.5	Subset sum	96		1.8	Matroid	. 120
6	Stri	ngs	97	8	Ext	ra	123
	6.1	Aho-corasick	97		8.1	fastIO.cpp	. 123
	6.2	eertree	97		8.2	vimre	. 123
	6.3	KMP	98		8.3	timer.cpp	. 123
	6.4	Manacher	98		8.4	rand.cpp	. 123
	6.5	$Min/max\ suffix/cyclic\ shift . \ . \ . \ . \ . \ . \ . \ . \ . \ .$	99		8.5	template.cpp	123
	6.6	String Hashing	99			• • • • • • • • • • • • • • • • • • • •	
	6.7	String Hashing - modulo $2^61 - 1 \dots \dots \dots \dots \dots$	99			debug.cpp	
	6.8	Suffix Array - O(n log n)	100		8.7	stress.sh	. 124
	6.9	Suffix Array - $O(n)$	100		8.8	makefile	. 124
	0.10	Suffix Array Dinamico	102		8.0	hash.sh	194

1 Estruturas

1.1 BIT

```
// BIT de soma 1-based, v 0-based
// Para mudar o valor da posicao p para x,
// faca: poe(x - query(p, p), p)
// l_bound(x) retorna o menor p tal que
// \text{ query}(1, p+1) > x 	 (0 based!)
// Complexidades:
// build - O(n)
// poe - O(\log(n))
// query - 0(log(n))
// l_bound - O(log(n))
1a8 int n;
7f4 int bit[MAX]:
b69 int v[MAX];
0a8 void build() {
b91
       bit[0] = 0;
       for (int i = 1; i <= n; i++) bit[i] = v[i - 1];</pre>
       for (int i = 1; i <= n; i++) {
78a
edf
            int i = i + (i \& -i):
            if (j <= n) bit[j] += bit[i];</pre>
b8a
        }
b39
5dc }
// soma x na posicao p
235 void poe(int x, int p) {
9c7
        for (; p <= n; p += p & -p) bit[p] += x;</pre>
ec3 }
// soma [1, p]
Obf int pref(int p) {
7c9 int ret = 0;
       for (; p; p -= p & -p) ret += bit[p];
805
        return ret;
601 }
// soma [a, b]
4ea int query(int a, int b) {
70c
       return pref(b) - pref(a - 1);
3dd }
```

```
e4a int l_bound(ll x) {
        int p = 0;
1ba
676
        for (int i = MAX2; i+1; i--) if (p + (1 << i) <= n
729
            and bit [p + (1 << i)] <= x) x -= bit <math>[p += (1 << i)];
74e
        return p;
c89 }
1.2 BIT 2D
// BIT de soma, update incrementa posicao
// Tem que construir com um vetor com todos os pontos
// que vc quer um dia atualizar (os pontos q vc vai chamar update)
// Complexidades:
// construir - O(n log(n))
// update e query - O(\log^2(n))
a6b template < class T = int > struct bit2d {
        vector <T> X;
a84
        vector < vector < T >> Y. t:
709
        int ub(vector<T>& v, T x) {
            return upper_bound(v.begin(), v.end(), x) - v.begin();
dde
9cc
5cb
        bit2d(vector<pair<T, T>> v) {
2e1
            for (auto [x, y] : v) X.push_back(x);
fd4
            sort(X.begin(), X.end());
1ee
            X.erase(unique(X.begin(), X.end()), X.end());
d56
            t.resize(X.size() + 1);
d12
            Y.resize(t.size()):
3d0
            sort(v.begin(), v.end(), [](auto a, auto b) {
e8f
                return a.second < b.second; });</pre>
            for (auto [x, y]: v) for (int i = ub(X, x); i < t.size();
   i += i\&-i)
b75
                if (!Y[i].size() or Y[i].back() != y)
   Y[i].push_back(y);
            for (int i = 0; i < t.size(); i++) t[i].resize(Y[i].size()</pre>
7c7
   + 1);
8cc }
        void update(T x, T v, T v) {
e78
2a9
            for (int i = ub(X, x); i < t.size(); i += i\&-i)
cd2
                for (int j = ub(Y[i], y); j < t[i].size(); j += j&-j)
   t[i][j] += v;
```

```
533
        }
        T query(T x, T y) {
5d2
            T ans = 0;
966
            for (int i = ub(X, x); i; i -= i&-i)
c54
                 for (int j = ub(Y[i], y); j; j -= j&-j) ans += t[i][j];
4fb
ba7
            return ans:
62d
46d
        T query (T x1, T y1, T x2, T y2) {
            return query (x2, y2) -query (x2, y1-1) -query (x1-1, y2)
fcf
   v2) + querv(x1-1, v1-1);
232
6a7 }:
```

1.3 BIT com update em range

```
// Operacoes 0-based
// query(l, r) retorna a soma de v[l..r]
// update(l, r, x) soma x em v[l..r]
// Complexidades:
// build - O(n)
// query - 0(log(n))
// update - 0(log(n))
e04 namespace bit {
        11 bit[2][MAX+2];
1a8
        int n;
        void build(int n2, int* v) {
61 c
1e3
            n = n2;
535
            for (int i = 1; i <= n; i++)
edd
                bit [1] [min(n+1, i+(i\&-i))] += bit [1][i] += v[i-1];
db0
637
        ll get(int x, int i) {
b73
            11 ret = 0:
360
            for (; i; i -= i&-i) ret += bit[x][i];
edf
            return ret;
99c
        }
        void add(int x, int i, ll val) {
20 c
            for (; i <= n; i += i&-i) bit[x][i] += val;</pre>
503
bf6
        }
        11 get2(int p) {
c7c
            return get(0, p) * p + get(1, p);
153
02a
        11 query(int 1, int r) {
ff5
            return get2(r+1) - get2(1);
```

```
633
089
        void update(int 1, int r, ll x) {
            add(0, 1+1, x), add(0, r+2, -x);
e5f
f58
            add(1, 1+1, -x*1), add(1, r+2, x*(r+1));
        }
e5f
f91 };
1.4 BIT-Sort Tree
// Tipo uma MergeSort Tree usando Bit
// Apesar da complexidade ser pior, fica melhor na pratica.
//
// query(1, r, k) retorna o numero de elementos menores que k
// no intervalo [1, r]
//
// Usa O(n log(n)) de memoria
// Complexidades:
// construir - O(n log^2(n))
// query - O(log^2(n))
6fa template < typename T> struct ms_bit {
1a8
        int n;
b2f
        vector < vector < T >> bit;
        ms\_bit(vector < T > \& v) : n(v.size()), bit(n+1) {
899
830
            for (int i = 0: i < n: i++)
d51
                for (int j = i+1; j \le n; j += j\&-j)
                    bit[j].push_back(v[i]);
dad
535
            for (int i = 1; i <= n; i++)</pre>
                 sort(bit[i].begin(), bit[i].end());
eec
        }
b4d
257
        int p_query(int i, T k) {
7 c.9
            int ret = 0:
be8
            for (i++; i; i -= i&-i)
1bd
                ret += lower_bound(bit[i].begin(), bit[i].end(), k) -
   bit[i].begin();
edf
            return ret;
6f9
        }
690
        int query(int 1, int r, T k) {
83d
            return p_query(r, k) - p_query(l-1, k);
bcc
        }
8d0 };
```

1.5 Convex Hull Trick Dinamico

```
// para double, use LINF = 1/.0, div(a, b) = a/b
// update(x) atualiza o ponto de intersecao da reta x
// overlap(x) verifica se a reta x sobrepoe a proxima
// add(a, b) adiciona reta da forma ax + b
// query(x) computa maximo de ax + b para entre as retas
// O(log(n)) amortizado por insercao
// O(log(n)) por query
72c struct Line {
073
        mutable ll a, b, p;
        bool operator<(const Line& o) const { return a < o.a; }</pre>
abf
        bool operator<(ll x) const { return p < x; }</pre>
469 };
326 struct dynamic_hull : multiset <Line, less <>> {
        11 div(ll a, ll b) {
33a
            return a / b - ((a ^ b) < 0 and a % b);
a20
a8a
bbb
        void update(iterator x) {
            if (next(x) == end()) x->p = LINF;
b2a
            else if (x->a == next(x)->a) x->p = x->b >= next(x)->b?
772
   LINF : -LINF;
            else x -> p = div(next(x) -> b - x -> b, x -> a - next(x) -> a);
424
0c4
71c
        bool overlap(iterator x) {
f18
            update(x);
cfa
            if (next(x) == end()) return 0;
            if (x->a == next(x)->a) return x->b >= next(x)->b;
a4a
d40
            return x - p >= next(x) - p;
901
        }
176
        void add(ll a, ll b) {
1c7
            auto x = insert({a, b, 0});
4ab
            while (overlap(x)) erase(next(x)), update(x);
            if (x != begin() and !overlap(prev(x))) x = prev(x),
   update(x);
0fc
            while (x != begin() and overlap(prev(x)))
4d2
                x = prev(x), erase(next(x)), update(x);
48f
        }
4ad
        ll query(ll x) {
229
            assert(!empty());
7d1
            auto 1 = *lower_bound(x);
            return 1.a * x + 1.b:
aba
```

```
3f5
8f2 };
1.6 Convex Hull Trick Estatico
// adds tem que serem feitos em ordem de slope
// queries tem que ser feitas em ordem de x
//
// linear
4b5 struct CHT {
942
         int it;
         vector<11> a, b;
ac1
45e
         CHT():it(0){}
         ll eval(int i, ll x){
0bb
93d
             return a[i]*x + b[i];
b2a
        }
63a
        bool useless(){
 a 20
             int sz = a.size();
35f
             int r = sz-1, m = sz-2, 1 = sz-3;
d71
             return (b[1] - b[r])*(a[m] - a[1]) <
413
                 (b[1] - b[m])*(a[r] - a[1]);
a0c
bf4
         void add(ll A, ll B){
7f5
             a.push_back(A); b.push_back(B);
565
             while (!a.empty()){
233
                 if ((a.size() < 3) || !useless()) break;</pre>
                 a.erase(a.end() - 2):
 ecb
568
                 b.erase(b.end() - 2);
b21
            }
165
         }
81b
        ll get(ll x){
             it = min(it, int(a.size()) - 1);
d27
46a
             while (it+1 < a.size()){
3c4
                 if (eval(it+1, x) > eval(it, x)) it++;
f97
                 else break;
             }
fe9
420
             return eval(it, x);
88a
        }
303 };
1.7 DSU
// Une dois conjuntos e acha a qual conjunto um elemento pertence por
    sen id
//
// find e unite: O(a(n)) \sim = O(1) amortizado
```

```
8d3 struct dsu {
825
        vector < int > id, sz;
        dsu(int n) : id(n), sz(n, 1) { iota(id.begin(), id.end(), 0); }
b33
0cf
        int find(int a) { return a == id[a] ? a : id[a] = find(id[a]):
  }
       void unite(int a, int b) {
440
605
            a = find(a), b = find(b);
d54
            if (a == b) return;
956
            if (sz[a] < sz[b]) swap(a, b);
6d0
            sz[a] += sz[b], id[b] = a;
ea7
       }
8e1 };
// DSU de bipartido
// Une dois vertices e acha a qual componente um vertice pertence
// Informa se a componente de um vertice e bipartida
// find e unite: O(log(n))
8d3 struct dsu {
6f7
        vector < int > id, sz, bip, c;
5b4
        dsu(int n) : id(n), sz(n, 1), bip(n, 1), c(n) {
db8
            iota(id.begin(), id.end(), 0);
f25
       }
ef0
        int find(int a) { return a == id[a] ? a : find(id[a]); }
        int color(int a) { return a == id[a] ? c[a] : c[a] ^
   color(id[a]): }
440
        void unite(int a, int b) {
263
            bool change = color(a) == color(b);
605
            a = find(a), b = find(b);
a89
            if (a == b) {
4ed
                if (change) bip[a] = 0;
505
                return;
            }
32d
956
            if (sz[a] < sz[b]) swap(a, b);
            if (change) c[b] = 1;
efe
            sz[a] += sz[b], id[b] = a, bip[a] &= bip[b];
2cd
22b
        }
```

```
118 };
// DSU Persistente
// Persistencia parcial, ou seja, tem que ir
// incrementando o 't' no une
//
// find e unite: O(log(n))
8d3 struct dsu {
        vector<int> id, sz, ti;
733
        dsu(int n) : id(n), sz(n, 1), ti(n, -INF) {
db8
            iota(id.begin(), id.end(), 0);
        }
aad
5e6
        int find(int a, int t) {
            if (id[a] == a or ti[a] > t) return a;
6ba
ea5
            return find(id[a], t);
        }
6cb
fa0
        void unite(int a, int b, int t) {
84f
            a = find(a, t), b = find(b, t);
d54
            if (a == b) return;
956
            if (sz[a] < sz[b]) swap(a, b);
35d
            sz[a] += sz[b], id[b] = a, ti[b] = t;
513
       }
6c6 };
// DSU com rollback
// checkpoint(): salva o estado atual de todas as variaveis
// rollback(): retorna para o valor das variaveis para
// o ultimo checkpoint
//
// Sempre que uma variavel muda de valor, adiciona na stack
// find e unite: O(log(n))
// checkpoint: 0(1)
// rollback: O(m) em que m e o numero de vezes que alguma
// variavel mudou de valor desde o ultimo checkpoint
8d3 struct dsu {
825
        vector < int > id, sz;
27 c
        stack<stack<pair<int&, int>>> st;
```

```
98d
        dsu(int n) : id(n), sz(n, 1) {
            iota(id.begin(), id.end(), 0), st.emplace();
1cc
8cd
bdf
        void save(int &x) { st.top().emplace(x, x); }
30d
        void checkpoint() { st.emplace(); }
5cf
        void rollback() {
            while(st.top().size()) {
ba9
6bf
                auto [end, val] = st.top().top(); st.top().pop();
149
                end = val:
f9a
           }
25a
            st.pop();
       }
3c6
ef0
        int find(int a) { return a == id[a] ? a : find(id[a]); }
440
        void unite(int a, int b) {
605
            a = find(a), b = find(b);
d54
            if (a == b) return:
956
            if (sz[a] < sz[b]) swap(a, b);
            save(sz[a]), save(id[b]);
803
6d0
            sz[a] += sz[b], id[b] = a;
1b9
       }
c6e };
```

1.8 Li-Chao Tree

```
// Adiciona retas (ax+b), e computa o minimo entre as retas
// em um dado 'x'
// Cuidado com overflow!
// Se tiver overflow, tenta comprimir o 'x' ou usar
// convex hull trick
//
// O(log(MA-MI)), O(n) de memoria
5b0 template<11 MI = 11(-1e9), 11 MA = 11(1e9) > struct lichao {
        struct line {
b3a
12d
            ll a. b:
            array < int, 2 > ch;
cef
fdf
            line(ll a_{-} = 0, ll b_{-} = LINF):
423
                a(a_{-}), b(b_{-}), ch(\{-1, -1\})  {}
888
            11 operator ()(11 x) { return a*x + b; }
d1d
17b
        vector < line > ln;
```

```
df8
        int ch(int p, int d) {
e85
            if (ln[p].ch[d] == -1) {
9af
                ln[p].ch[d] = ln.size();
cdc
                ln.emplace_back();
bc1
ef2
            return ln[p].ch[d];
86a
        }
021
        lichao() { ln.emplace_back(); }
c33
        void add(line s, ll l=MI, ll r=MA, int p=0) {
3e3
            11 m = (1+r)/2;
911
            bool L = s(1) < ln[p](1);
d37
            bool M = s(m) < ln[p](m):
03b
            bool R = s(r) < ln[p](r);
            if (M) swap(ln[p], s), swap(ln[p].ch, s.ch);
825
            if (s.b == LINF) return;
cac
f6d
            if (L != M) add(s, 1, m-1, ch(p, 0));
898
            else if (R != M) add(s, m+1, r, ch(p, 1));
76e
       11 query(int x, 11 1=MI, 11 r=MA, int p=0) {
092
11b
            11 m = (1+r)/2, ret = ln[p](x);
9db
            if (ret == LINF) return ret;
529
            if (x < m) return min(ret, query(x, 1, m-1, ch(p, 0)));</pre>
81a
            return min(ret, query(x, m+1, r, ch(p, 1)));
fba
       }
59b }:
```

1.9 Li-Chao Tree - Lazy

```
// Sendo N = MA-MI:
// insert(\{a, b\}) minimiza tudo com ax+b - 0(\log N)
// insert(\{a, b\}, 1, r) minimiza com ax+b no range [1, r] - 0(\log^2 N)
// shift({a, b}) soma ax+b em tudo - O(1)
// shift({a, b}, 1, r) soma ax+b no range [1, r] - O(log^2 N)
// query(x) retorna o valor da posicao x - O(\log N)
//
// No inicio eh tudo LINF, se inserir {0, 0} fica tudo 0
// O(n log N) de memoria ; O(n) de memoria se nao usar as operacoes de
   range
41c template <int MI = int(-1e9), int MA = int(1e9) > struct lichao {
        struct line {
b3a
12d
            ll a, b;
158
            11 la, lb; // lazy
cef
            array<int, 2> ch;
fdf
            line(ll a_= 0, ll b_= LINF):
```

```
b09
                a(a_{-}), b(b_{-}), la(0), lb(0), ch(\{-1, -1\})  {}
            11 operator ()(11 x) { return a*x + b; }
888
92e
        };
17b
        vector < line > ln;
        int ch(int p, int d) {
df8
e85
            if (\ln \lceil p \rceil, \cosh \lceil d \rceil == -1) {
9af
                ln[p].ch[d] = ln.size();
cdc
                ln.emplace_back();
bc1
ef2
            return ln[p].ch[d];
86a
021
        lichao() { ln.emplace back(); }
ceb
        void prop(int p, int 1, int r) {
ff8
            if (ln[p].la == 0 and ln[p].lb == 0) return;
1d3
            ln[p].a += ln[p].la, ln[p].b += ln[p].lb;
579
            if (1 != r) {
                int pl = ch(p, 0), pr = ch(p, 1);
b9e
                ln[pl].la += ln[p].la, ln[pl].lb += ln[p].lb;
0d7
                ln[pr].la += ln[p].la, ln[pr].lb += ln[p].lb;
fa8
77f
01e
            ln[p].la = ln[p].lb = 0;
89b
        }
c06
        ll query(int x, int p=0, int l=MI, int r=MA) {
6b9
            prop(p, 1, r);
6f3
            ll ret = ln[p](x):
            if (ln[p].ch[0] == -1 and ln[p].ch[1] == -1) return ret;
33b
90d
            int m = 1 + (r-1)/2;
            if (x <= m) return min(ret, query(x, ch(p, 0), 1, m));</pre>
da9
c55
            return min(ret, query(x, ch(p, 1), m+1, r));
        }
953
        void push(line s, int p, int l, int r) {
5df
6b9
            prop(p, 1, r);
90d
            int m = 1 + (r-1)/2;
911
            bool L = s(1) < ln[p](1);
d37
            bool M = s(m) < ln[p](m);
03b
            bool R = s(r) < ln[p](r);
            if (M) swap(ln[p].a, s.a), swap(ln[p].b, s.b);
c3f
            if (s.b == LINF) return;
cac
            if (L != M) push(s, ch(p, 0), 1, m);
c49
            else if (R != M) push(s, ch(p, 1), m+1, r);
29 e
ceb
        void insert(line s, int a=MI, int b=MA, int p=0, int l=MI, int
   r = MA) {
```

```
prop(p, 1, r);
6b9
2d3
            if (a \le 1 \text{ and } r \le b) \text{ return } push(s, p, l, r);
1dd
            if (b < 1 or r < a) return;
90d
            int m = 1 + (r-1)/2;
f1e
            insert(s, a, b, ch(p, 0), 1, m);
            insert(s, a, b, ch(p, 1), m+1, r);
952
375
        }
97a
        void shift(line s, int a=MI, int b=MA, int p=0, int l=MI, int
   r=MA) {
6b9
            prop(p, 1, r);
90d
            int m = 1 + (r-1)/2:
9a3
            if (a \le 1 \text{ and } r \le b)
ada
                 ln[p].la += s.a, ln[p].lb += s.b;
505
                 return;
570
            }
1dd
            if (b < 1 or r < a) return;
            if (ln[p].b != LINF) {
fdd
751
                 push(ln[p], ch(p, 0), 1, m);
ade
                 push(ln[p], ch(p, 1), m+1, r);
                 ln[p].a = 0, ln[p].b = LINF;
c2f
199
a04
            shift(s, a, b, ch(p, 0), 1, m);
e7d
             shift(s, a, b, ch(p, 1), m+1, r);
d43
        }
285 };
```

1.10 MergeSort Tree

```
// Se for construida sobre um array:
//
        count(i, j, a, b) retorna quantos
//
         elementos de v[i..j] pertencem a [a, b]
//
         report(i, j, a, b) retorna os indices dos
         elementos de v[i..j] que pertencem a [a, b]
//
//
         retorna o vetor ordenado
// Se for construida sobre pontos (x, y):
//
         count(x1, x2, y1, x2) retorna quantos pontos
//
         pertencem ao retangulo (x1, y1), (x2, y2)
//
         report(x1, x2, y1, y2) retorna os indices dos pontos que
//
         pertencem ao retangulo (x1, y1), (x2, y2)
//
        retorna os pontos ordenados lexicograficamente
//
         (assume x1 \le x2, y1 \le y2)
// kth(y1, y2, k) retorna o indice do ponto com k-esimo menor
// x dentre os pontos que possuem y em [y1, y2] (0 based)
// Se quiser usar para achar k-esimo valor em range, construir
// com ms_tree t(v, true), e chamar kth(l, r, k)
```

```
// Usa O(n log(n)) de memoria
//
// Complexidades:
// construir - O(n log(n))
// count - 0(log(n))
// report - O(\log(n) + k) para k indices retornados
// kth - O(log(n))
c6c template <typename T = int> struct ms_tree {
6f7
        vector < tuple < T, T, int >> v;
1a8
5ee
        vector < vector < tuple < T, T, int >>> t; // {y, idx, left}
6ae
        vector <T> vy;
        ms_tree(vector < pair < T, T >> \& vv) : n(vv.size()), t(4*n), vy(n) {
78c
e80
             for (int i = 0; i < n; i++) v.push_back({vv[i].first,</pre>
   vv[i].second, i});
             sort(v.begin(), v.end());
fca
             build(1, 0, n-1);
224
01a
             for (int i = 0; i < n; i++) vy[i] = get < 0 > (t[1][i+1]);
45e
        ms_tree(vector<T>& vv, bool inv = false) { // inv: inverte
   indice e valor
             vector < pair < T, T >> v2;
8e8
e1e
            for (int i = 0; i < vv.size(); i++)</pre>
                 inv ? v2.push_back({vv[i], i}) : v2.push_back({i,
196
   vv[i]}):
             *this = ms_tree(v2);
cca
f23
        void build(int p, int 1, int r) {
2.06
            t[p].push_back({get<0>(v[1]), get<0>(v[r]), 0}); //
1d2
   {min_x, max_x, 0}
5c8
             if (1 == r) return t[p].push_back({get<1>(v[1]),
   get <2>(v[1]), 0});
            int m = (1+r)/2;
ee4
             build(2*p, 1, m), build(2*p+1, m+1, r);
bd9
32d
            int L = 0, R = 0:
             while (t[p].size() \le r-l+1) {
a03
                 int left = get <2>(t[p].back());
68e
                 if (L > m-1 \text{ or } (R+m+1 \le r \text{ and } t[2*p+1][1+R] \le
   t[2*p][1+L])) {
                     t[p].push_back(t[2*p+1][1 + R++]);
8cf
                     get < 2 > (t[p].back()) = left;
da0
                     continue:
5e2
                 }
ce0
```

```
249
                 t[p].push_back(t[2*p][1 + L++]);
339
                 get < 2 > (t[p].back()) = left + 1;
208
             }
        }
2eb
         int get_1(T y) { return lower_bound(vy.begin(), vy.end(), y) -
    vy.begin(); }
         int get_r(T y) { return upper_bound(vy.begin(), vy.end(), y) -
    vv.begin(); }
f62
         int count(T x1, T x2, T y1, T y2) {
902
             function < int(int, int, int) > dfs = [&](int p, int 1, int
   r) {
7 c 6
                 if (1 == r \text{ or } x2 < get < 0 > (t[p][0]) \text{ or } get < 1 > (t[p][0])
    < x1) return 0;
                 if (x1 \le get<0>(t[p][0]) and get<1>(t[p][0]) \le x2)
2bb
    return r-1;
                 int nl = get<2>(t[p][1]), nr = get<2>(t[p][r]);
784
                 return dfs(2*p, nl, nr) + dfs(2*p+1, l-nl, r-nr);
eb6
122
             }:
7cb
             return dfs(1, get_l(y1), get_r(y2));
f65
002
        vector<int> report(T x1, T x2, T y1, T y2) {
4b8
             vector < int > ret;
85 e
             function < void(int, int, int) > dfs = [&](int p, int 1, int
   r) {
882
                 if (1 == r \text{ or } x2 < get < 0 > (t[p][0]) \text{ or } get < 1 > (t[p][0])
    < x1) return:
                 if (x1 \le get<0>(t[p][0]) and get<1>(t[p][0]) \le x2) {
8da
                      for (int i = 1; i < r; i++)</pre>
e00
    ret.push_back(get<1>(t[p][i+1]));
505
                      return;
                 }
067
784
                 int nl = get<2>(t[p][1]), nr = get<2>(t[p][r]);
                 dfs(2*p, nl, nr), dfs(2*p+1, l-nl, r-nr);
194
12b
             dfs(1, get_l(y1), get_r(y2));
8ad
             return ret;
edf
668
985
         int kth(T v1, T v2, int k) {
902
             function < int (int, int, int) > dfs = [&] (int p, int 1, int
   r) {
150
                 if (k >= r-1) {
941
                      k \rightarrow r-1:
daa
                      return -1;
b8d
                 }
8da
                 if (r-l == 1) return get<1>(t[p][l+1]);
```

```
784
                int nl = get<2>(t[p][1]), nr = get<2>(t[p][r]);
072
                int left = dfs(2*p, nl, nr);
3b6
                if (left != -1) return left;
                return dfs(2*p+1, l-nl, r-nr);
04d
a1b
            }:
7cb
            return dfs(1, get_l(y1), get_r(y2));
635
1ce }:
1.11 Min queue - deque
// Tudo O(1) amortizado
1dc template < class T> struct minqueue {
        deque<pair<T, int>> q;
2d8
3fc
        void push(T x) {
56e
            int ct = 1:
953
            while (q.size() and x < q.front().first)</pre>
75f
                ct += q.front().second, q.pop_front();
987
            q.emplace_front(x, ct);
        }
e8d
42d
        void pop() {
            if (q.back().second > 1) q.back().second--;
aa2
c51
            else q.pop_back();
5fd
ea6
        T min() { return q.back().first; }
c13 }:
1.12 Min queue - stack
// Tudo O(1) amortizado
557 template < class T> struct minstack {
81f
        stack<pair<T, T>> s;
3fc
        void push(T x) {
12b
            if (!s.size()) s.push({x, x});
9d9
            else s.emplace(x, std::min(s.top().second, x));
f8d
        T top() { return s.top().first; }
4f0
```

94a

1f2 2eb

ba7

013 614 } () gog T

s.pop();

return ans;

T ans = s.top().first;

int size() { return s.size(); }

```
13b
         T min() { return s.top().second; }
 4c0 };
 1dc template < class T> struct minqueue {
         minstack <T> s1, s2;
7cd
         void push(T x) { s1.push(x); }
 c96
         void move() {
 d4d
             if (s2.size()) return;
 d92
             while (s1.size()) {
 7ae
                 T x = s1.pop();
 489
                 s2.push(x);
 656
             }
 ef1
         }
 787
         T front() { return move(), s2.top(); }
         T pop() { return move(), s2.pop(); }
 23a
7f3
         int size() { return s1.size()+s2.size(); }
        T min() {
 19c
 cd6
             if (!s1.size()) return s2.min();
 58e
             else if (!s2.size()) return s1.min();
 31d
             return std::min(s1.min(), s2.min());
9c7
        }
6d3 };
1.13 Order Statistic Set
// Funciona do C++11 pra cima
774 #include <ext/pb_ds/assoc_container.hpp>
 30f #include <ext/pb_ds/tree_policy.hpp>
 0d7 using namespace __gnu_pbds;
 4fc template <class T>
 def
         using ord_set = tree<T, null_type, less<T>, rb_tree_tag,
 3a1
         tree_order_statistics_node_update >;
// para declarar:
// ord set <int> s:
// coisas do set normal funcionam:
// for (auto i : s) cout << i << endl;
// cout << s.size() << endl:
// k-esimo maior elemento O(log|s|):
// k=0: menor elemento
// cout << *s.find_by_order(k) << endl;</pre>
// quantos sao menores do que k O(log|s|):
// cout << s.order_of_key(k) << endl;</pre>
// Para fazer um multiset, tem que
```

```
// usar ord_set<pair<int, int>> com o
// segundo parametro sendo algo para diferenciar
// os ementos iguais.
// s.order_of_key({k, -INF}) vai retornar o
// numero de elementos < k</pre>
```

1.14 Priority Queue DS

```
// Mantem updates aplicados em uma estrutura de dados
// que permita rollback e nao seja amortizada.
// Cada update possui uma prioridade,
// sendo possivel remover o update com maior prioridade.
// Os updates devem ser comutativos, ou seja, o estado
// da estrutura deve ser o mesmo independente da ordem
// que eles sejam aplicados.
//
// Complexidades:
// update - O(log(n) + T(n))
// query - T(n)
// pop - O(log(n) * T(n)) amortizado
// onde T(n) eh a complexidade do update
// assumes all priorities are distinct
945 template < typename DS, typename UPD > struct priority_queue_ds {
df4
        DS D;
a7e
        vector < tuple < UPD, int, int >> upd; // {u, p, idx_in_pos}
866
        set < pair < int , int >> st;
927
        vector < int > pos;
        priority_queue_ds(int n) : D(n) {}
cf0
6af
        void update(UPD u, int p) {
9ab
            D.update(u);
d07
            st.emplace(p, pos.size());
6ca
            upd.emplace_back(u, p, pos.size());
e3d
            pos.push_back(upd.size() - 1);
        }
6af
427
        int query(int a) {
            return D.find(a);
aa3
2d3
        }
42d
        void pop() {
25 f
            int k = 1, min_p; // k = number of pops we will do
43e
            vector < tuple < UPD, int, int >> small, big;
639
            auto it = st.end():
```

```
231
            for (int qt = 0; qt++ < (k+1)/2;) {
049
                it--:
3ab
                min_p = it->first;
80f
                int i = pos[it->second];
e82
                if (qt > 1) big.push_back(upd[i]);
84ъ
                k = max < int > (k, upd.size() - i);
b9a
            }
b3d
            for (int i = 0; i < k; i++) {</pre>
a62
                D.rollback():
6d8
                auto [u, p, idx] = upd.rbegin()[i];
86d
                 if (p < min_p) small.emplace_back(u, p, idx);</pre>
588
            }
23e
            st.erase(prev(st.end()));
            upd.erase(upd.end() - k, upd.end());
623
            small.insert(small.end(), big.rbegin(), big.rend());
a25
06f
            for (auto [u, p, idx] : small) {
9ab
                D.update(u);
c8e
                upd.emplace_back(u, p, idx);
a7d
                pos[idx] = upd.size() - 1;
            }
ec7
bd1
        }
54a }:
1.15 Range color
// update(l, r, c) colore o range [l, r] com a cor c,
// e retorna os ranges que foram coloridos {1, r, cor}
// query(i) returna a cor da posicao i
//
// Complexidades (para q operacoes):
// update - O(log(q)) amortizado
// query - O(log(q))
df6 template < typename T> struct color {
f0c
        set < tuple < int , int , T >> se;
071
        vector<tuple<int, int, T>> update(int 1, int r, T val) {
```

auto it = se.upper_bound({r, INF, val});

auto [L, R, V] = *--it;

it = se.lower_bound({1, -INF, val});

se.erase(it);

}

if (it != se.begin() and get<1>(*prev(it)) > r) {

se.emplace(L, r, V), se.emplace(r+1, R, V);

9c4

753

e91

3f0

bfd

683

d9e

```
516
            if (it != se.begin() and get<1>(*prev(it)) >= 1) {
                                                                            e34
                                                                                    int index_query(int 1, int r) {
                auto [L, R, V] = *--it;
                                                                            27b
e91
                                                                                        if (r-l+1 \le b) return small(r, r-l+1);
3f0
                se.erase(it);
                                                                            e80
                                                                                        int x = 1/b+1, y = r/b-1;
                                                                                        if (x > y) return op(small(1+b-1), small(r));
75a
                se.emplace(L, 1-1, V), it = se.emplace(1, R, V).first;
                                                                            fd3
b65
                                                                            a4e
                                                                                        int j = msb(y-x+1);
            vector<tuple<int, int, T>> ret;
                                                                                        int ans = op(small(1+b-1), op(t[n/b*j+x],
d7b
                                                                            ea3
7a1
            for (; it != se.end() and get<0>(*it) <= r; it =</pre>
                                                                               t[n/b*j+y-(1<<j)+1]));
                                                                                        return op(ans, small(r));
   se.erase(it))
                                                                            be6
8c0
                ret.push_back(*it);
                                                                            62a
            se.emplace(1, r, val);
                                                                                    T query(int 1, int r) { return v[index_query(1, r)]; }
b4a
                                                                            093
edf
            return ret;
                                                                            bab };
b6c
ff9
        T querv(int i) {
                                                                           1.17 SegTreap
c31
            auto it = se.upper_bound({i, INF, T()});
8e7
            if (it == se.begin() or get<1>(*--it) < i) return -1; //</pre>
                                                                           // Muda uma posicao do plano, e faz query de operacao
   nao tem
                                                                           // associativa e comutativa em retangulo
            return get <2>(*it);
53d
                                                                           // Mudar ZERO e op
daf
                                                                           // Esparso nas duas coordenadas, inicialmente eh tudo ZERO
9e9 };
                                                                           // Para query com distancia de manhattan <= d, faca
1.16 RMQ < O(n), O(1) > - min queue
                                                                           // nx = x+y, ny = x-y
                                                                           // Update em (nx, ny), query em ((nx-d, ny-d), (nx+d, ny+d))
// O(n) pra buildar, query O(1)
                                                                           //
// Se tiver varios minimos, retorna
                                                                           // Valores no X tem que ser de O ateh NX
// o de menor indice
                                                                           // Para q operacoes, usa O(q log(NX)) de memoria, e as
                                                                           // operacoes custa O(log(q) log(NX))
1a5 template < typename T > struct rmq {
517
        vector <T> v;
                                                                            55b const int ZERO = INF;
fcc
        int n; static const int b = 30;
                                                                            560 const int op(int 1, int r) { return min(1, r); }
70e
        vector < int > mask, t;
                                                                            878 mt19937 rng((int)
183
        int op(int x, int y) { return v[x] \leftarrow v[y] ? x : y; }
                                                                               chrono::steady_clock::now().time_since_epoch().count());
ee1
        int msb(int x) { return __builtin_clz(1)-__builtin_clz(x); }
c92
        int small(int r, int sz = b) { return
                                                                            aa1 template < typename T > struct treap {
   r-msb(mask[r]&((1<<sz)-1));}
                                                                            3c9
                                                                                    struct node {
                                                                            b19
6ad
        rmq() {}
                                                                                        node *1, *r;
43c
        rmq(const \ vector < T > \& \ v_) : v(v_), n(v.size()), mask(n), t(n) {
                                                                            ee1
                                                                                        int p;
2e5
            for (int i = 0, at = 0; i < n; mask[i++] = at |= 1) {
                                                                            850
                                                                                        pair<11, 11> idx; // {y, x}
                at = (at << 1) &((1 << b) -1);
a61
                                                                            36d
                                                                                        T val, mi;
                while (at and op(i-msb(at&-at), i) == i) at ^= at&-at;
                                                                                        node(ll x, ll y, T val_) : l(NULL), r(NULL), p(rng()),
c00
                                                                            bc2
                                                                                            idx(pair(y, x)), val(val_), mi(val) {}
c2f
                                                                            1 b 5
ea4
            for (int i = 0; i < n/b; i++) t[i] = small(b*i+b-1);
                                                                           01e
                                                                                        void update() {
            for (int j = 1; (1<<j) <= n/b; j++) for (int i = 0;
                                                                            d6e
                                                                                            mi = val:
   i+(1<< j) <= n/b; i++)
                                                                            182
                                                                                            if (1) mi = op(mi, 1->mi);
                t[n/b*j+i] = op(t[n/b*(j-1)+i],
                                                                            b68
                                                                                            if (r) mi = op(mi, r->mi);
   t[n/b*(j-1)+i+(1<<(j-1))]);
                                                                           282
                                                                                        }
```

41a

6e1

};

```
bb7
        node* root;
        treap() { root = NULL; }
84b
        \simtreap() {
cec
            vector < node *> q = {root};
609
402
            while (q.size()) {
e5d
                node* x = q.back(); q.pop_back();
ee9
                if (!x) continue;
                q.push_back(x->1), q.push_back(x->r);
1 c 7
bf0
                delete x;
653
            }
50e
225
        treap(treap&& t) : treap() { swap(root, t.root); }
bcf
        void join(node* 1, node* r, node*& i) { // assume que 1 < r</pre>
986
            if (!1 or !r) return void(i = 1 ? 1 : r);
            if (1->p > r->p) join(1->r, r, 1->r), i = 1;
80e
            else join(1, r->1, r->1), i = r;
fa0
            i->update():
bda
671
        }
c82
        void split(node* i, node*& 1, node*& r, pair<11, 11> idx) {
            if (!i) return void(r = l = NULL);
26a
13c
            if (i->idx < idx) split(i->r, i->r, r, idx), l = i;
d26
            else split(i \rightarrow 1, l, i \rightarrow 1, idx), r = i;
bda
            i->update();
143
d3b
        void update(ll x, ll y, T v) {
df9
            node *L, *M, *R;
8b2
            split(root, M, R, pair(y, x+1)), split(M, L, M, pair(y,
   x));
1e4
            if (M) M->val = M->mi = v;
9e5
            else M = new node(x, y, v);
69d
            join(L, M, M), join(M, R, root);
58e
91b
        T query(ll ly, ll ry) {
df9
            node *L, *M, *R;
1c0
            split(root, M, R, pair(ry, LINF)), split(M, L, M, pair(ly,
   0));
0f7
            T ret = M ? M->mi : ZERO;
            join(L, M, M), join(M, R, root);
69d
            return ret;
edf
        }
1ae
bdf };
46a template < typename T > struct segtreap {
c4f
        vector < treap < T >> seg;
```

```
6e7
        vector < int > ch[2]:
e4e
        ll NX:
        segtreap(11 NX_{-}) : seg(1), NX(NX_{-}) \{ ch[0].push_back(-1), 
253
    ch[1].push_back(-1); }
a71
        int get_ch(int i, int d){
            if (ch[d][i] == -1) {
e51
2d6
                ch[d][i] = seg.size();
23e
                seg.emplace_back();
842
                ch[0].push_back(-1), ch[1].push_back(-1);
3e1
968
            return ch[d][i]:
bb6
        }
        T query(11 lx, 11 rx, 11 ly, 11 ry, int p, 11 l, 11 r) {
10c
003
            if (rx < 1 or r < 1x) return ZERO;</pre>
fOf
            if (lx <= 1 and r <= rx) return seg[p].query(ly, ry);</pre>
            11 m = 1 + (r-1)/2:
e6a
354
            return op(query(lx, rx, ly, ry, get_ch(p, 0), 1, m),
060
                     query(lx, rx, ly, ry, get_ch(p, 1), m+1, r));
a5e
        }
f48
        T query(ll lx, ll rx, ll ly, ll ry) { return query(lx, rx, ly,
   ry, 0, 0, NX); }
        void update(ll x, ll y, T val, int p, ll l, ll r) {
249
73c
            if (1 == r) return seg[p].update(x, y, val);
e6a
            11 m = 1 + (r-1)/2;
сс5
            if (x <= m) update(x, y, val, get_ch(p, 0), 1, m);</pre>
5a2
            else update(x, y, val, get_ch(p, 1), m+1, r);
980
            seg[p].update(x, y, val);
cc2
        }
517
        void update(11 x, 11 y, T val) { update(x, y, val, 0, 0, NX); }
40a };
1.18 SegTree
// Recursiva com Lazy Propagation
// Query: soma do range [a, b]
// Update: soma x em cada elemento do range [a, b]
// Pode usar a seguinte funcao para indexar os nohs:
// f(1, r) = (1+r) | (1!=r), usando 2N de memoria
```

// Complexidades:

// query - O(log(n))

// build - O(n)

```
// update - O(log(n))
aa4 namespace seg {
        ll seg [4*MAX], lazy [4*MAX];
005
052
        int n, *v;
d22
        ll build(int p=1, int l=0, int r=n-1) {
3c7
            lazv[p] = 0;
6cd
            if (1 == r) return seg[p] = v[1];
ee4
            int m = (1+r)/2;
193
            return seg[p] = build(2*p, 1, m) + build(2*p+1, m+1, r);
c71
0d8
        void build(int n2. int* v2) {
680
            n = n2, v = v2;
6f2
            build();
acb
        }
ceb
        void prop(int p, int l, int r) {
            seg[p] += lazy[p]*(r-l+1);
cdf
2c9
            if (1 != r) lazy[2*p] += lazy[p], lazy[2*p+1] += lazy[p];
3c7
            lazy[p] = 0;
c10
        }
2c3
        ll query(int a, int b, int p=1, int l=0, int r=n-1) {
6b9
            prop(p, 1, r);
527
            if (a \le 1 \text{ and } r \le b) \text{ return seg}[p];
786
            if (b < 1 or r < a) return 0;</pre>
ee4
            int m = (1+r)/2:
b1f
            return query(a, b, 2*p, 1, m) + query(a, b, 2*p+1, m+1, r);
4c5
cfb
        11 update(int a, int b, int x, int p=1, int l=0, int r=n-1) {
6b9
            prop(p, 1, r);
9a3
            if (a <= 1 and r <= b) {</pre>
b94
                lazv[p] += x;
6b9
                prop(p, 1, r);
534
                return seg[p];
821
e9f
            if (b < 1 or r < a) return seg[p];</pre>
            int m = (1+r)/2;
ee4
fdb
            return seg[p] = update(a, b, x, 2*p, 1, m) +
                update(a, b, x, 2*p+1, m+1, r);
7fd
75c
        }
0af };
// Se tiver uma seg de max, da pra descobrir em O(log(n))
// o primeiro e ultimo elemento >= val numa range:
// primeira posicao >= val em [a, b] (ou -1 se nao tem)
                                                                          // Para query com distancia de manhattan <= d, faca
```

```
119 int get_left(int a, int b, int val, int p=1, int l=0, int r=n-1) {
6b9
        prop(p, 1, r);
f38
        if (b < l \text{ or } r < a \text{ or } seg[p] < val) return -1;
205
        if (r == 1) return 1;
        int m = (1+r)/2:
ee4
753
        int x = get_left(a, b, val, 2*p, 1, m);
50e
        if (x != -1) return x:
        return get_left(a, b, val, 2*p+1, m+1, r);
сЗс
68c }
// ultima posicao >= val em [a, b] (ou -1 se nao tem)
992 int get_right(int a, int b, int val, int p=1, int l=0, int r=n-1) {
6b9
        prop(p, 1, r):
f38
        if (b < l or r < a or seg[p] < val) return -1;</pre>
205
        if (r == 1) return 1;
        int m = (1+r)/2;
ee4
        int x = get_right(a, b, val, 2*p+1, m+1, r);
1b1
        if (x != -1) return x;
50e
6a7
        return get_right(a, b, val, 2*p, 1, m);
1b7 }
// Se tiver uma seg de soma sobre um array nao negativo v, da pra
// descobrir em O(\log(n)) o maior j tal que v[i]+v[i+1]+...+v[j-1] <
   val
6a9 int lower_bound(int i, ll& val, int p, int l, int r) {
6b9
        prop(p. 1. r):
        if (r < i) return n;</pre>
6e8
b5d
        if (i <= l and seg[p] < val) {</pre>
bff
            val -= seg[p];
041
            return n;
634
        }
Зсе
        if (1 == r) return 1;
        int m = (1+r)/2;
514
        int x = lower_bound(i, val, 2*p, 1, m);
        if (x != n) return x;
ee0
8b9
        return lower_bound(i, val, 2*p+1, m+1, r);
2b8 }
1.19 SegTree 2D Iterativa
// Consultas O-based
// Um valor inicial em (x, y) deve ser colocado em seg[x+n][y+n]
// Query: soma do retangulo ((x1, y1), (x2, y2))
// Update: muda o valor da posicao (x, y) para val
// Nao pergunte como que essa coisa funciona
//
```

```
// nx = x+y, ny = x-y
// Update em (nx, ny), query em ((nx-d, ny-d), (nx+d, ny+d))
// Se for de min/max, pode tirar os if's da 'query', e fazer
// sempre as 4 operacoes. Fica mais rapido
// Complexidades:
// build - O(n^2)
// \text{ query - } O(\log^2(n))
// update - 0(log^2(n))
731 int seg[2*MAX][2*MAX], n;
0a8 void build() {
919
        for (int x = 2*n; x; x--) for (int y = 2*n; y; y--) {
             if (x < n) seg[x][y] = seg[2*x][y] + seg[2*x+1][y];
c81
fe9
             if (y < n) seg[x][y] = seg[x][2*y] + seg[x][2*y+1];
d51
499 }
251 int query(int x1, int y1, int x2, int y2) {
827
         int ret = 0, v3 = v1 + n, v4 = v2 + n;
83e
         for (x1 += n, x2 += n; x1 <= x2; ++x1 /= 2, --x2 /= 2)
0f2
             for (y1 = y3, y2 = y4; y1 \le y2; ++y1 /= 2, --y2 /= 2) {
554
                  if (x1\%2 == 1 \text{ and } y1\%2 == 1) \text{ ret } += \text{seg}[x1][y1];
6b0
                  if (x1\%2 == 1 \text{ and } y2\%2 == 0) \text{ ret } += \text{seg}[x1][y2];
c01
                  if (x2\%2 == 0 \text{ and } y1\%2 == 1) \text{ ret } += \text{seg}[x2][y1];
5d4
                 if (x2\%2 == 0 \text{ and } v2\%2 == 0) \text{ ret } += \text{seg}[x2][v2]:
2d0
             }
edf
         return ret;
ff1 }
767 void update(int x, int y, int val) {
         int v2 = v += n;
66a
192
        for (x += n; x; x /= 2, y = y2) {
970
             if (x \ge n) seg[x][y] = val;
ba9
             else seg[x][y] = seg[2*x][y] + seg[2*x+1][y];
3b1
             while (y /= 2) \text{ seg}[x][y] = \text{seg}[x][2*y] + \text{seg}[x][2*y+1];
484
        }
62e }
      SegTree Beats
1.20
// \text{ query}(a, b) - \{\{\min(v[a..b]), \max(v[a..b])\}, \sup(v[a..b])\}
// updatemin(a, b, x) faz com que v[i] <- min(v[i], x),</pre>
```

```
// para i em [a, b]
// updatemax faz o mesmo com max, e updatesum soma x
// em todo mundo do intervalo [a, b]
//
// Complexidades:
// build - O(n)
// query - O(log(n))
// update - O(log^2 (n)) amortizado
// (se nao usar updatesum, fica log(n) amortizado)
7c6 #define f first
Oab #define s second
f39 namespace beats {
3c9
        struct node {
526
            int tam:
125
            ll sum, lazy; // lazy pra soma
4f3
            ll mi1, mi2, mi; // mi = #mi1
            ll ma1, ma2, ma; // ma = #ma1
c61
426
            node(11 x = 0) {
                sum = mi1 = ma1 = x:
ba6
b29
                mi2 = LINF, ma2 = -LINF;
62c
                mi = ma = tam = 1;
c60
                lazv = 0:
b00
770
            node(const node& 1, const node& r) {
a95
                sum = 1.sum + r.sum, tam = 1.tam + r.tam;
c60
                lazv = 0;
797
                if (1.mi1 > r.mi1) {
230
                    mi1 = r.mi1, mi = r.mi;
ea2
                    mi2 = min(1.mi1, r.mi2);
                } else if (1.mi1 < r.mi1) {</pre>
dcd
e34
                    mi1 = 1.mi1. mi = 1.mi:
                    mi2 = min(r.mi1, 1.mi2);
4b3
9d9
                } else {
                    mi1 = 1.mi1, mi = 1.mi+r.mi;
a39
83d
                    mi2 = min(1.mi2, r.mi2);
5b7
cd0
                if (1.ma1 < r.ma1) {</pre>
6a0
                    ma1 = r.ma1, ma = r.ma;
                    ma2 = max(1.ma1, r.ma2);
96d
5f0
                } else if (1.ma1 > r.ma1) {
ae0
                    ma1 = 1.ma1. ma = 1.ma:
2ca
                    ma2 = max(r.ma1, 1.ma2);
949
                } else {
db2
                    ma1 = 1.ma1, ma = 1.ma+r.ma;
```

```
c05
                     ma2 = max(1.ma2, r.ma2);
                }
710
1ba
            }
4b4
            void setmin(ll x) {
                if (x >= ma1) return;
55e
                sum += (x - ma1)*ma;
463
be5
                if (mi1 == ma1) mi1 = x:
                if (mi2 == ma1) mi2 = x;
0a0
b81
                ma1 = x;
            }
0c3
6cb
            void setmax(ll x) {
e25
                if (x <= mi1) return;</pre>
7e8
                sum += (x - mi1)*mi:
                if (ma1 == mi1) ma1 = x:
0bb
c32
                if (ma2 == mi1) ma2 = x;
1ff
                mi1 = x:
a86
            }
            void setsum(ll x) {
4cf
                mi1 += x, mi2 += x, ma1 += x, ma2 += x;
fe8
620
                sum += x*tam:
c46
                lazy += x;
b53
            }
47f
        };
62b
        node seg[4*MAX];
052
        int n. *v:
93ъ
        node build(int p=1, int l=0, int r=n-1) {
            if (1 == r) return seg[p] = {v[1]};
d84
ee4
            int m = (1+r)/2;
            return seg[p] = \{build(2*p, 1, m), build(2*p+1, m+1, r)\};
3d6
444
        void build(int n2. int* v2) {
0d8
680
            n = n2, v = v2:
6f2
            build():
acb
        }
        void prop(int p, int 1, int r) {
ceb
8ce
            if (1 == r) return;
abd
            for (int k = 0; k < 2; k++) {
d07
                if (seg[p].lazy) seg[2*p+k].setsum(seg[p].lazy);
                seg[2*p+k].setmin(seg[p].ma1);
843
f79
                seg[2*p+k].setmax(seg[p].mi1);
585
431
            seg[p].lazy = 0;
7ee
055
        pair < pair < 11 , 11 > , 11 > query (int a, int b, int p=1, int 1=0,
   int r=n-1) {
```

```
e07
             if (b < l or r < a) return {{LINF, -LINF}, 0};</pre>
             if (a \le 1 \text{ and } r \le b) \text{ return } \{\{seg[p].mi1, seg[p].ma1\},
9be
   seg[p].sum};
6b9
             prop(p, 1, r);
             int m = (1+r)/2:
ee4
             auto L = query(a, b, 2*p, 1, m), R = query(a, b, 2*p+1,
e6f
   m+1. r):
             return {{min(L.f.f, R.f.f), max(L.f.s, R.f.s)}, L.s+R.s};
96d
e9d
2c8
        node updatemin(int a, int b, ll x, int p=1, int l=0, int
   r=n-1) {
744
             if (b < 1 or r < a or seg[p].ma1 <= x) return seg[p];</pre>
309
             if (a \le 1 \text{ and } r \le b \text{ and } seg[p].ma2 < x) {
                 seg[p].setmin(x);
ccd
534
                 return seg[p];
bbf
6b9
             prop(p, 1, r);
             int m = (1+r)/2:
ee4
96a
             return seg[p] = \{updatemin(a, b, x, 2*p, 1, m),
                              updatemin(a, b, x, 2*p+1, m+1, r)};
4db
aad
        }
044
        node updatemax(int a, int b, ll x, int p=1, int l=0, int
   r=n-1) {
b59
             if (b < 1 or r < a or seg[p].mi1 >= x) return seg[p];
a9e
             if (a \le 1 \text{ and } r \le b \text{ and } seg[p].mi2 > x) {
e8a
                 seg[p].setmax(x);
534
                 return seg[p];
e9b
             }
6b9
             prop(p, 1, r);
             int m = (1+r)/2;
ee4
             return seg[p] = \{updatemax(a, b, x, 2*p, 1, m),
ee3
98ъ
                              updatemax(a, b, x, 2*p+1, m+1, r)};
323
        }
        node updatesum(int a, int b, ll x, int p=1, int l=0, int
aee
   r=n-1)
             if (b < l or r < a) return seg[p];</pre>
e9f
9a3
             if (a \le 1 \text{ and } r \le b)
8f4
                 seg[p].setsum(x);
534
                 return seg[p];
596
             }
6b9
             prop(p, 1, r);
             int m = (1+r)/2;
ee4
7b6
             return seg[p] = \{updatesum(a, b, x, 2*p, 1, m),
483
                              updatesum(a, b, x, 2*p+1, m+1, r)};
111
        }
0d2 }:
```

1.21 SegTree Colorida

```
// Cada posicao tem um valor e uma cor
// O construtor receve um vector de {valor, cor}
// e o numero de cores (as cores devem estar em [0, c-1])
// querv(c, a, b) retorna a soma dos valores
// de todo mundo em [a, b] que tem cor c
// update(c, a, b, x) soma x em todo mundo em
// [a, b] que tem cor c
// paint(c1, c2, a, b) faz com que todo mundo
// em [a, b] que tem cor c1 passe a ter cor c2
// Complexidades:
// construir - O(n log(n)) espaco e tempo
// query - 0(log(n))
// update - O(log(n))
// paint - O(log(n)) amortizado
04f struct seg_color {
        struct node {
b19
            node *1. *r:
0f9
            int cnt;
9ca
            ll val, lazv;
277
            node(): 1(NULL), r(NULL), cnt(0), val(0), lazy(0) {}
01e
            void update() {
d0a
                cnt = 0, val = 0;
bc4
                for (auto i : {1, r}) if (i) {
c89
                    i->prop();
281
                    cnt += i->cnt, val += i->val;
                }
684
554
            }
a9c
            void prop() {
2dd
                if (!lazy) return;
3f7
                val += lazv*(ll)cnt;
b64
                for (auto i : {1, r}) if (i) i->lazy += lazy;
                lazv = 0;
c60
            }
e24
       };
514
1a8
        int n:
9b0
        vector < node *> seg;
        seg_color(vector<pair<int, int>>& v, int c) : n(v.size()),
   seg(c, NULL) {
830
            for (int i = 0; i < n; i++)</pre>
                seg[v[i].second] = insert(seg[v[i].second], i,
   v[i].first, 0, n-1);
```

```
94a
        }
3c7
        \simseg_color() {
dde
            queue < node *> q;
3a6
            for (auto i : seg) q.push(i);
402
            while (q.size()) {
20b
                auto i = q.front(); q.pop();
dab
                if (!i) continue;
7c7
                q.push(i->1), q.push(i->r);
5се
                delete i;
            }
c60
139
        }
40b
        node* insert(node* at, int idx, int val, int l, int r) {
1a4
            if (!at) at = new node();
232
            if (l == r) return at->cnt = 1, at->val = val, at;
ee4
            int m = (1+r)/2;
137
            if (idx \le m) at->1 = insert(at->1, idx, val, 1, m);
            else at->r = insert(at->r, idx, val, m+1, r);
3e6
cff
            return at->update(), at;
d6e
        }
870
        11 query(node* at, int a, int b, int l, int r) {
61b
            if (!at or b < 1 or r < a) return 0;
d9f
            at->prop();
cb2
            if (a <= l and r <= b) return at->val;
ee4
            int m = (1+r)/2;
4c4
            return query(at->1, a, b, 1, m) + query(at->r, a, b, m+1,
   r);
8c3
e54
        11 query(int c, int a, int b) { return query(seg[c], a, b, 0,
   n-1); }
91 c
        void update(node* at, int a, int b, int x, int l, int r) {
fba
            if (!at or b < l or r < a) return;</pre>
d9f
            at->prop();
9a3
            if (a \le 1 \text{ and } r \le b)
e9a
                at->lazv += x;
cb2
                return void(at->prop());
051
            }
            int m = (1+r)/2;
ee4
0b0
            update(at->1, a, b, x, 1, m), update(at->r, a, b, x, m+1,
   r);
7b4
            at->update();
9fd
        void update(int c, int a, int b, int x) { update(seg[c], a, b,
a40
   x, 0, n-1); }
        void paint(node*& from, node*& to, int a, int b, int 1, int r)
70c
   {
10f
            if (to == from or !from or b < l or r < a) return;
```

```
e85
            from ->prop();
889
            if (to) to->prop();
9a3
            if (a <= 1 and r <= b) {</pre>
                if (!to) {
24d
38f
                    to = from:
140
                    from = NULL;
505
                    return:
                }
e5f
                int m = (1+r)/2;
ee4
                paint(from->1, to->1, a, b, 1, m), paint(from->r,
1cb
   to->r, a, b, m+1, r);
                to->update();
72d
270
                delete from:
140
                from = NULL;
505
                return;
            }
a0e
019
            if (!to) to = new node();
            int m = (1+r)/2;
ee4
            paint(from->1, to->1, a, b, 1, m), paint(from->r, to->r,
1cb
   a, b, m+1, r);
45a
            from ->update(), to ->update();
4aa
471
        void paint(int c1, int c2, int a, int b) { paint(seg[c1],
   seg[c2], a, b, 0, n-1); }
293 }:
1.22 SegTree Esparsa - Lazy
// Query: soma do range [a, b]
// Update: flipa os valores de [a, b]
// O MAX tem q ser Q log N para Q updates
//
// Complexidades:
// build - O(1)
// query - O(log(n))
// update - 0(log(n))
aa4 namespace seg {
6de
        int seg[MAX], lazy[MAX], R[MAX], L[MAX], ptr;
e9a
        int get_l(int i){
            if (L[i] == 0) L[i] = ptr++;
3db
a96
            return L[i]:
b6e
        }
943
        int get_r(int i){
71b
            if (R[i] == 0) R[i] = ptr++;
283
            return R[i];
```

43a

}

```
e71
        void build() { ptr = 2; }
ceb
        void prop(int p, int l, int r) {
b77
             if (!lazy[p]) return;
76c
             seg[p] = r-l+1 - seg[p];
213
             if (1 != r) lazy[get_l(p)]^=lazy[p],
    lazy[get_r(p)]^=lazy[p];
3c7
             lazv[p] = 0;
        }
20b
158
        int query(int a, int b, int p=1, int 1=0, int r=N-1) {
6b9
             prop(p, 1, r):
786
             if (b < 1 \text{ or } r < a) \text{ return } 0;
527
             if (a <= 1 and r <= b) return seg[p];</pre>
             int m = (1+r)/2;
ee4
818
             return query(a, b, get_l(p), l, m)+query(a, b, get_r(p),
   m+1, r);
        }
0d9
51f
        int update(int a, int b, int p=1, int l=0, int r=N-1) {
6b9
             prop(p, 1, r);
e9f
             if (b < 1 or r < a) return seg[p];</pre>
9a3
             if (a <= 1 and r <= b) {</pre>
ab6
                 lazy[p] ^= 1;
6b9
                 prop(p, 1, r);
534
                 return seg[p];
8e4
             int m = (1+r)/2;
ee4
             return seg[p] = update(a, b, get_l(p), l, m)+update(a, b,
    get_r(p), m+1, r);
1dc
dc3 }:
1.23 SegTree Esparsa - O(q) memoria
// Query: min do range [a, b]
// Update: troca o valor de uma posicao
// Usa O(q) de memoria para q updates
// Complexidades:
// query - O(log(n))
// update - O(log(n))
13d template < typename T > struct seg {
3c9
        struct node {
```

```
d53
            node* ch[2];
970
            char d;
            T v;
ca0
            T mi:
c4e
d4e
            node(int d_, T v_, T val) : d(d_), v(v_) {
                ch[0] = ch[1] = NULL;
e71
d6e
                mi = val;
            }
065
b32
            node(node* x) : d(x->d), v(x->v), mi(x->mi) {
c99
                ch[0] = x -> ch[0], ch[1] = x -> ch[1];
cb7
            }
01e
            void update() {
909
                mi = numeric_limits <T>::max();
                for (int i = 0; i < 2; i++) if (ch[i])
151
b5a
                    mi = min(mi, ch[i]->mi);
            }
fe3
530
        };
bb7
        node* root;
9c5
        char n;
ba7
        seg() : root(NULL), n(0) {}
512
        \simseg() {
4c0
            std::vector<node*> q = {root};
402
            while (q.size()) {
e5d
                node* x = q.back(); q.pop_back();
                if (!x) continue;
ee9
73f
                q.push_back(x->ch[0]), q.push_back(x->ch[1]);
bf0
                delete x;
d3e
            }
        }
d8c
        char msb(T v, char l, char r) { // msb in range (l, r]
1a6
8e4
            for (char i = r; i > 1; i--) if (v>>i&1) return i;
            return -1;
daa
688
430
        void cut(node* at, T v, char i) {
677
            char d = msb(v ^ at -> v, at -> d, i);
23b
            if (d == -1) return; // no need to split
ebf
            node* nxt = new node(at);
d43
            at -> ch[v>> d&1] = NULL:
34f
            at -> ch[!(v>>d&1)] = nxt;
150
            at -> d = d;
        }
0b3
```

```
6e5
         node* update(node* at, T idx, T val, char i) {
             if (!at) return new node(-1, idx, val);
 c8c
d67
             cut(at, idx, i);
             if (at->d == -1) { // leaf
1a2
792
                 at->mi = val:
 ce6
                 return at;
a6f
             }
b29
             bool dir = idx>>at->d&1;
c8f
             at - ch[dir] = update(at - ch[dir], idx, val, at - cd - 1);
7b4
             at->update();
ce6
             return at;
76d
        }
85c
         void update(T idx, T val) {
8f4
             while (idx >> n) n++;
61e
             root = update(root, idx, val, n-1);
79d
        }
         T query(node* at, T a, T b, T l, T r, char i) {
9d8
df0
             if (!at or b < l or r < a) return numeric_limits<T>::max();
fd3
             if (a <= 1 and r <= b) return at->mi;
841
             T m = 1 + (r-1)/2;
c85
             if (at->d < i) {</pre>
c59
                 if ((at->v>>i\&1) == 0) return query(at, a, b, 1, m,
   i-1);
                 else return query(at, a, b, m+1, r, i-1);
ca4
934
373
             return min(query(at->ch[0], a, b, 1, m, i-1),
    query(at->ch[1], a, b, m+1, r, i-1);
2db
         T query (T 1, T r) \{ return query (root, 1, r, 0, (T(1) << n) -1,
034
    n-1): }
d7f };
1.24 SegTree Iterativa
// Consultas 0-based
// Valores iniciais devem estar em (seg[n], ..., seg[2*n-1])
// Query: soma do range [a, b]
// Update: muda o valor da posicao p para x
//
// Complexidades:
// build - O(n)
// query - O(log(n))
// update - O(log(n))
6a4 int seg[2 * MAX];
la8 int n;
```

```
0a8 void build() {
d15
        for (int i = n - 1; i; i--) seg[i] = seg[2*i] + seg[2*i+1];
9a8 }
4ea int query(int a, int b) {
7c9
        int ret = 0:
728
        for (a += n, b += n; a <= b; ++a /= 2, --b /= 2) {
4ea
            if (a % 2 == 1) ret += seg[a];
            if (b \% 2 == 0) ret += seg[b];
244
ac0
edf
        return ret;
24a }
ff3 void update(int p, int x) {
        seg[p += n] = x;
37 d
c8c
        while (p /= 2) seg[p] = seg[2*p] + seg[2*p+1];
02d }
```

1.25 SegTree Iterativa com Lazy Propagation

```
// Query: soma do range [a, b]
// Update: soma x em cada elemento do range [a, b]
// Para mudar, mudar as funcoes junta, poe e query
// LOG = ceil(log2(MAX))
//
// Complexidades:
// build - O(n)
// query - O(log(n))
// update - 0(log(n))
aa4 namespace seg {
6db
        11 seg[2*MAX], lazy[2*MAX];
1a8
        int n;
9b3
        ll junta(ll a, ll b) {
534
            return a+b:
e26
        }
        // soma x na posicao p de tamanho tam
        void poe(int p, ll x, int tam, bool prop=1) {
1b4
517
            seg[p] += x*tam;
6ae
            if (prop and p < n) lazy[p] += x;</pre>
8bc
        // atualiza todos os pais da folha p
b1e
        void sobe(int p) {
```

```
d5a
            for (int tam = 2; p /= 2; tam *= 2) {
4ca
                seg[p] = junta(seg[2*p], seg[2*p+1]);
388
                poe(p, lazy[p], tam, 0);
            }
acd
b76
        }
        // propaga o caminho da raiz ate a folha p
        void prop(int p) {
a0a
076
            int tam = 1 << (LOG-1);</pre>
            for (int s = LOG; s; s--, tam /= 2) {
0a8
4b1
                int i = p >> s;
27 c
                if (lazy[i]) {
860
                    poe(2*i, lazy[i], tam);
e38
                    poe(2*i+1, lazy[i], tam);
b97
                    lazy[i] = 0;
                }
de8
3ed
            }
        }
e29
61c
        void build(int n2, int* v) {
1e3
            n = n2:
95f
            for (int i = 0; i < n; i++) seg[n+i] = v[i];</pre>
            for (int i = n-1; i; i--) seg[i] = junta(seg[2*i],
c41
   seg[2*i+1]);
f4c
            for (int i = 0; i < 2*n; i++) lazy[i] = 0;</pre>
8bb
4f3
        11 query(int a, int b) {
b73
            ll ret = 0;
b48
            for (prop(a+=n), prop(b+=n); a \le b; ++a/=2, --b/=2) {
                if (a%2 == 1) ret = junta(ret, seg[a]);
a8e
c58
                if (b%2 == 0) ret = junta(ret, seg[b]);
510
            }
edf
            return ret:
38b
        }
        void update(int a, int b, int x) {
a28
c2d
            int a2 = a += n, b2 = b += n, tam = 1;
Off
            for (; a <= b; ++a/=2, --b/=2, tam *= 2) {
32a
                if (a\%2 == 1) poe(a, x, tam);
                if (b\%2 == 0) poe(b, x, tam);
9da
9bc
0f7
            sobe(a2), sobe(b2);
        }
adc
6dc };
```

1.26 SegTree PA

```
// Segtree de PA
// update_set(1, r, A, R) seta [1, r] para PA(A, R),
// update_add soma PA(A, R) em [1, r]
// query(l, r) retorna a soma de [l, r]
// PA(A, R) eh a PA: [A+R, A+2R, A+3R, ...]
// Complexidades:
// construir - O(n)
// update_set, update_add, query - O(log(n))
dc7 struct seg_pa {
350
        struct Data {
8f5
            ll sum:
662
            11 set_a, set_r, add_a, add_r;
9b7
            Data(): sum(0), set_a(LINF), set_r(0), add_a(0), add_r(0)
   {}
eb6
        };
        vector < Data > seg;
16a
1a8
        int n;
d45
        seg_pa(int n_) {
e95
            n = n_{-};
            seg = vector < Data > (4*n);
fc3
ce0
       }
ceb
        void prop(int p, int l, int r) {
d5a
            int tam = r-1+1:
c3f
            11 &sum = seg[p].sum, &set_a = seg[p].set_a, &set_r =
   seg[p].set_r,
a1b
                \&add_a = seg[p].add_a, \&add_r = seg[p].add_r;
c02
            if (set a != LINF) {
660
                set_a += add_a, set_r += add_r;
                sum = set_a*tam + set_r*tam*(tam+1)/2;
06e
579
                if (1 != r) {
                    int m = (1+r)/2;
ee4
                    seg[2*p].set_a = set_a;
886
358
                    seg[2*p].set_r = set_r;
                    seg[2*p].add_a = seg[2*p].add_r = 0;
ed6
f0c
                    seg[2*p+1].set_a = set_a + set_r * (m-l+1);
471
                    seg[2*p+1].set_r = set_r;
d48
                    seg[2*p+1].add_a = seg[2*p+1].add_r = 0;
```

```
a97
                }
823
                set_a = LINF, set_r = 0;
953
                 add_a = add_r = 0;
            } else if (add_a or add_r) {
105
18b
                sum += add_a*tam + add_r*tam*(tam+1)/2;
579
                if (1 != r) {
ee4
                     int m = (1+r)/2:
ffO
                     seg[2*p].add_a += add_a;
                     seg[2*p].add_r += add_r;
ec0
06c
                     seg[2*p+1].add_a += add_a + add_r * (m-l+1);
a6d
                     seg[2*p+1].add r += add r:
8af
953
                add_a = add_r = 0;
551
            }
07f
        }
        int inter(pair<int, int> a, pair<int, int> b) {
0b7
            if (a.first > b.first) swap(a, b);
98c
eef
            return max(0, min(a.second, b.second) - b.first + 1);
628
        11 set(int a, int b, ll aa, ll rr, int p, int l, int r) {
be1
6b9
            prop(p, 1, r);
457
            if (b < 1 or r < a) return seg[p].sum;</pre>
9a3
            if (a \le 1 \text{ and } r \le b)
91c
                seg[p].set_a = aa;
774
                seg[p].set_r = rr;
6b9
                prop(p, 1, r);
254
                return seg[p].sum;
            }
8ee
ee4
            int m = (1+r)/2;
963
            int tam_1 = inter({1, m}, {a, b});
c34
            return seg[p].sum = set(a, b, aa, rr, 2*p, 1, m) +
                 set(a, b, aa + rr * tam_l, rr, 2*p+1, m+1, r);
365
8e2
        }
f55
        void update_set(int 1, int r, 11 aa, 11 rr) {
6f7
            set(1, r, aa, rr, 1, 0, n-1);
913
5f6
        11 add(int a, int b, ll aa, ll rr, int p, int l, int r) {
6b9
            prop(p, 1, r):
            if (b < 1 or r < a) return seg[p].sum;</pre>
457
9a3
            if (a \le 1 \text{ and } r \le b) \{
359
                 seg[p].add_a += aa;
1ee
                 seg[p].add_r += rr;
6b9
                prop(p, 1, r);
254
                return seg[p].sum;
```

```
d19
ee4
            int m = (1+r)/2:
963
            int tam_l = inter({1, m}, {a, b});
            return seg[p].sum = add(a, b, aa, rr, 2*p, 1, m) +
586
695
                add(a, b, aa + rr * tam_l, rr, 2*p+1, m+1, r);
904
848
        void update_add(int 1, int r, 11 aa, 11 rr) {
            add(1, r, aa, rr, 1, 0, n-1);
afa
81e
f45
        ll query(int a, int b, int p, int l, int r) {
6b9
            prop(p, 1, r);
786
            if (b < 1 or r < a) return 0;</pre>
e9a
            if (a <= 1 and r <= b) return seg[p].sum:
            int m = (1+r)/2:
ee4
b1f
            return query(a, b, 2*p, 1, m) + query(a, b, 2*p+1, m+1, r);
f6e
bfc
        11 query(int 1, int r) { return query(1, r, 1, 0, n-1); }
bc4 }:
```

1.27 SegTree Persistente

```
// SegTree de soma, update de somar numa posicao
//
// query(a, b, t) retorna a query de [a, b] na versao t
// update(a, x, t) faz um update v[a]+=x a partir da
// versao de t. criando uma nova versao e retornando seu id
// Por default, faz o update a partir da ultima versao
// build - O(n)
// query - 0(log(n))
// update - O(log(n))
54a const int MAX = 1e5+10, UPD = 1e5+10, LOG = 18;
6de const int MAXS = 2*MAX+UPD*LOG;
f6e namespace perseg {
bd6
        11 seg[MAXS];
f4e
        int rt[UPD], L[MAXS], R[MAXS], cnt, t;
052
        int n, *v;
        11 build(int p, int l, int r) {
3 c 4
6cd
            if (1 == r) return seg[p] = v[1];
855
            L[p] = cnt++, R[p] = cnt++;
ee4
            int m = (1+r)/2;
275
            return seg[p] = build(L[p], 1, m) + build(R[p], m+1, r);
39d
0d8
        void build(int n2, int* v2) {
```

```
680
            n = n2, v = v2;
856
            rt[0] = cnt++:
c50
            build(0, 0, n-1);
a2e
        }
f45
        11 query(int a, int b, int p, int 1, int r) {
            if (b < 1 or r < a) return 0;
786
527
            if (a <= 1 and r <= b) return seg[p]:
            int m = (1+r)/2;
ee4
            return query(a, b, L[p], 1, m) + query(a, b, R[p], m+1, r);
1ed
4d2
182
        11 query(int a, int b, int tt) {
c13
            return query(a, b, rt[tt], 0, n-1);
726
bb3
        ll update(int a, int x, int lp, int p, int l, int r) {
747
            if (l == r) return seg[p] = seg[lp]+x;
            int m = (1+r)/2;
ee4
ab8
            if (a <= m)
                return seg[p] = update(a, x, L[lp], L[p]=cnt++, 1, m)
b48
   + seg[R[p]=R[lp]];
            return seg[p] = seg[L[p]=L[lp]] + update(a, x, R[lp],
8a9
   R[p] = cnt + + , m+1, r);
788
        int update(int a, int x, int tt=t) {
6f6
            update(a, x, rt[tt], rt[++t]=cnt++, 0, n-1);
ab3
e0d
            return t:
d63
        }
26f }:
```

1.28 SegTree Persistente com Lazy

```
// Nao propaga, meio estranho de mexer, mas da
//
// query(a, b, t) retorna a query de [a, b] na versao t
// update(a, b, x, t) faz um update v[a..b]+=x a partir da
// versao de t, criando uma nova versao e retornando seu id
// Por default, faz o update a partir da ultima versao
//
// build - O(n)
// query - O(log(n))
// update - O(log(n))

54a const int MAX = 1e5+10, UPD = 1e5+10, LOG = 18;
ab3 const int MAXS = 2*MAX + 4*UPD*LOG;

f6e namespace perseg {
9eb    int seg[MAXS];
f1e    int rt[UPD], L[MAXS], R[MAXS], cnt, t;
```

```
052
        int n, *v;
                                                                            // MAX2 = log(MAX)
        int build(int p, int 1, int r) {
                                                                            // Complexidades:
adf
            if (1 == r) return seg[p] = v[1];
                                                                            // build - O(n log(n))
6cd
855
            L[p] = cnt++, R[p] = cnt++;
                                                                            // query - 0(1)
            int m = (1+r)/2;
ee4
0.1d
            return seg[p] = max(build(L[p], 1, m), build(R[p], m+1,
                                                                            cca namespace sparse {
                                                                            710
                                                                                     int m[MAX2][MAX], n;
   r));
                                                                            61c
                                                                                     void build(int n2, int* v) {
ffd
        void build(int n2. int *v2) {
850
                                                                            1 e 3
                                                                                         n = n2:
680
            n = n2, v = v2;
                                                                            78e
                                                                                         for (int i = 0; i < n; i++) m[0][i] = v[i];</pre>
856
            rt[0] = cnt++:
                                                                                         for (int j = 1; (1<<j) <= n; j++) for (int i = 0; i+(1<<<math>j)
                                                                            a1c
c50
            build(0, 0, n-1):
                                                                                <= n; i++)
a2e
        }
                                                                            5d5
                                                                                             m[j][i] = min(m[j-1][i], m[j-1][i+(1<<(j-1))]);
976
        int query(int a, int b, int p, int l, int r) {
                                                                            cae
                                                                                    }
27b
            if (b < 1 or r < a) return -INF;</pre>
                                                                                     int query(int a, int b) {
                                                                            4ea
                                                                                         int j = __builtin_clz(1) - __builtin_clz(b-a+1);
793
            if (a <= 1 and r <= b) return lazy[p] + seg[p];</pre>
                                                                            ee5
                                                                                         return min(m[j][a], m[j][b-(1<<j)+1]);</pre>
ee4
            int m = (1+r)/2:
                                                                            dc3
            int ret = lazy[p] + max(query(a, b, L[p], 1, m), query(a,
                                                                            fba
                                                                                    }
7a2
   b, R[p], m+1, r));
                                                                            7aa }
            return ret;
edf
9a7
                                                                            1.30 Sparse Table Disjunta
442
        int query(int a, int b, int tt) {
c13
            return query(a, b, rt[tt], 0, n-1);
                                                                            // Resolve qualquer operacao associativa
a 0.5
                                                                            // MAX2 = log(MAX)
bc1
        int update(int a, int b, int x, int lp, int p, int l, int r) {
                                                                            //
3f6
            tie(seg[p], lazy[p], L[p], R[p]) = {seg[lp], lazy[lp],
                                                                            // Complexidades:
   L[lp], R[lp]};
                                                                            // build - O(n log(n))
            if (b < l or r < a) return seg[p] + lazy[p];</pre>
847
                                                                            // query - 0(1)
32a
            if (a <= 1 and r <= b) return seg[p] + (lazy[p] += x);
                                                                            cca namespace sparse {
            int m = (1+r)/2;
ee4
                                                                            9bf
                                                                                     int m[MAX2][2*MAX], n, v[2*MAX];
            seg[p] = max(update(a, b, x, L[lp], L[p] = cnt++, l, m),
24a
                                                                            5f7
                                                                                     int op(int a, int b) { return min(a, b); }
                          update(a, b, x, R[lp], R[p] = cnt++, m+1, r)):
bdb
                                                                            0d8
                                                                                     void build(int n2, int* v2) {
            lazv[p] = lazv[lp];
1ed
                                                                            1e3
                                                                                         n = n2:
1b7
            return seg[p] + lazy[p];
                                                                            df4
                                                                                         for (int i = 0; i < n; i++) v[i] = v2[i];
877
        }
                                                                            a84
                                                                                         while (n&(n-1)) n++;
cbf
        int update(int a, int b, int x, int tt=t) {
                                                                            3d2
                                                                                         for (int j = 0; (1<<j) < n; j++) {
aa8
            assert(tt <= t):
                                                                                             int len = 1<<j;</pre>
                                                                            1 c 0
661
            update(a, b, x, rt[tt], rt[++t]=cnt++, 0, n-1);
                                                                            d9b
                                                                                             for (int c = len: c < n: c += 2*len) {
e0d
            return t:
                                                                            332
                                                                                                 m[j][c] = v[c], m[j][c-1] = v[c-1];
aad
                                                                            668
                                                                                                 for (int i = c+1; i < c+len; i++) m[j][i] =</pre>
f27 };
                                                                                op(m[j][i-1], v[i]);
                                                                            432
                                                                                                 for (int i = c-2; i >= c-len; i--) m[j][i] =
     Sparse Table
                                                                                op(v[i], m[j][i+1]);
                                                                            eda
                                                                                            }
// Resolve RMQ
                                                                            f4d
                                                                                         }
```

```
ce3
9e3
        int query(int 1, int r) {
f13
            if (1 == r) return v[1];
            int j = __builtin_clz(1) - __builtin_clz(1^r);
e6d
d67
            return op(m[j][1], m[j][r]);
a7b
fd8 }
1.31 Splay Tree
// SEMPRE QUE DESCER NA ARVORE, DAR SPLAY NO
// NODE MAIS PROFUNDO VISITADO
// Todas as operacoes sao O(\log(n)) amortizado
// Se quiser colocar mais informacao no node,
// mudar em 'update'
538 template < typename T > struct splaytree {
3c9
        struct node {
183
            node *ch[2], *p;
e4d
            int sz;
f48
            T val:
da0
            node(T v) {
696
                ch[0] = ch[1] = p = NULL;
a 26
                sz = 1:
250
                val = v;
2d0
            }
01e
            void update() {
a26
                sz = 1;
c7c
                for (int i = 0; i < 2; i++) if (ch[i]) {
d5f
                     sz += ch[i]->sz;
486
                }
f45
aa3
        };
bb7
        node* root;
        splaytree() { root = NULL; }
fbc
214
        splaytree(const splaytree& t) {
cbf
            throw logic_error("Nao copiar a splaytree!");
1f1
891
        \simsplaytree() {
609
            vector < node *> q = {root};
402
            while (q.size()) {
e5d
                node* x = q.back(); q.pop_back();
                if (!x) continue;
ee9
73f
                q.push_back(x->ch[0]), q.push_back(x->ch[1]);
bf0
                delete x;
```

```
d3e
            }
837
        }
94f
        void rotate(node* x) { // x vai ficar em cima
d9b
            node *p = x->p, *pp = p->p;
            if (pp) pp - ch[pp - ch[1] == p] = x;
ecf
286
            bool d = p \rightarrow ch[0] == x;
            p->ch[!d] = x->ch[d], x->ch[d] = p;
d63
bad
            if (p->ch[!d]) p->ch[!d]->p = p;
fc2
            x->p = pp, p->p = x;
1ea
            p->update(), x->update();
007
        }
3fa
        node* splay(node* x) {
a39
            if (!x) return x;
4ea
            root = x;
3cf
            while (x->p) {
d9b
                 node *p = x->p, *pp = p->p;
359
                 if (!pp) return rotate(x), x; // zig
                 if ((pp->ch[0] == p)^(p->ch[0] == x))
еЗс
a2b
                     rotate(x), rotate(x); // zigzag
4b2
                 else rotate(p), rotate(x); // zigzig
028
ea5
            return x;
379
        }
313
        node* insert(T v, bool lb=0) {
b64
            if (!root) return lb ? NULL : root = new node(v);
            node *x = root, *last = NULL;;
002
31e
            while (1) {
5d7
                 bool d = x -> val < v;
Ofd
                if (!d) last = x;
c2e
                 if (x->val == v) break;
c16
                 if (x->ch[d]) x = x->ch[d];
4e6
                 else {
dea
                     if (lb) break:
                     x \rightarrow ch[d] = new node(v);
055
99c
                     x - ch[d] - p = x;
30e
                     x = x - ch[d];
c2b
                     break;
68a
                }
1ab
            }
0b6
            splay(x);
            return lb ? splay(last) : x;
61c
622
сОс
        int size() { return root ? root->sz : 0; }
2ca
        int count(T v) { return insert(v, 1) and root->val == v; }
111
        node* lower_bound(T v) { return insert(v, 1); }
26b
        void erase(T v) {
```

```
446
            if (!count(v)) return;
                                                                             3c9
                                                                                     struct node {
            node *x = root, *1 = x -> ch[0];
                                                                                         node *ch[2], *p;
bce
                                                                             183
268
            if (!1) {
                                                                             e4d
                                                                                         int sz;
                                                                             875
8b1
                root = x->ch[1];
                                                                                         T val, sub, lazy;
                if (root) root->p = NULL;
32e
                                                                             aa6
                                                                                         bool rev;
                                                                             da0
8f3
                return delete x;
                                                                                         node(T v) {
a86
            }
                                                                             696
                                                                                             ch[0] = ch[1] = p = NULL;
5e7
                                                                             a26
            root = 1, 1->p = NULL;
                                                                                             sz = 1;
902
            while (1->ch[1]) 1 = 1->ch[1];
                                                                             1e4
                                                                                             sub = val = v;
bab
                                                                                             lazy = 0;
            splay(1);
                                                                             c60
f0e
            1 - ch[1] = x - ch[1];
                                                                             b67
                                                                                             rev = false;
7d9
            if (1->ch[1]) 1->ch[1]->p = 1;
                                                                             48f
                                                                                         }
bf0
            delete x:
                                                                             a9c
                                                                                         void prop() {
62a
            1->update();
                                                                             0ec
                                                                                             if (lazy) {
007
        }
                                                                             924
                                                                                                 val += lazy, sub += lazy*sz;
24a
                                                                             091
        int order_of_key(T v) {
                                                                                                 if (ch[0]) ch[0]->lazy += lazy;
            if (!lower_bound(v)) return root ? root->sz : 0;
62b
                                                                             1a8
                                                                                                 if (ch[1]) ch[1]->lazy += lazy;
            return root -> ch [0] ? root -> ch [0] -> sz : 0;
                                                                                             }
1cc
                                                                             a98
                                                                             1bb
                                                                                             if (rev) {
b00
                                                                                                 swap(ch[0], ch[1]);
db6
        node* find_by_order(int k) {
                                                                             80a
084
            if (k >= size()) return NULL;
                                                                             628
                                                                                                 if (ch[0]) ch[0]->rev ^= 1;
52f
            node* x = root;
                                                                                                 if (ch[1]) ch[1]->rev ^= 1;
                                                                             adc
            while (1) {
                                                                             30a
                                                                                             }
31e
20f
                if (x->ch[0] \text{ and } x->ch[0]->sz >= k+1) x = x->ch[0];
                                                                             a32
                                                                                             lazy = 0, rev = 0;
4e6
                else {
                                                                             6bf
                                                                                         }
a1c
                     if (x->ch[0]) k -= x->ch[0]->sz;
                                                                             01e
                                                                                         void update() {
                                                                                             sz = 1, sub = val;
1dc
                    if (!k) return splay(x);
                                                                             0 c 3
eb8
                    k--, x = x->ch[1];
                                                                             c7c
                                                                                             for (int i = 0; i < 2; i++) if (ch[i]) {
                }
                                                                             05f
aca
                                                                                                 ch[i]->prop();
e01
            }
                                                                             d5f
                                                                                                 sz += ch[i] -> sz;
        }
                                                                                                 sub += ch[i]->sub;
Ode
                                                                             4 a 1
19c
        T min() {
                                                                             6c1
                                                                                             }
                                                                             e98
                                                                                         }
52f
            node* x = root:
6f6
            while (x->ch[0]) x = x->ch[0]; // max -> ch[1]
                                                                             b4a
                                                                                     }:
3e9
            return splay(x)->val;
70e
        }
                                                                             bb7
                                                                                     node* root;
4ff };
                                                                             5d9
                                                                                     splay() { root = NULL; }
                                                                             9b1
                                                                                     splay(node* x) {
      Splay Tree Implicita
                                                                             4ea
                                                                                         root = x;
                                                                             32e
                                                                                         if (root) root->p = NULL;
// vector da NASA
                                                                             371
// Um pouco mais rapido q a treap
                                                                             1b7
                                                                                     splay(vector < T > v) { // O(n)}
// O construtor a partir do vector
                                                                             950
                                                                                         root = NULL:
// eh linear, todas as outras operacoes
                                                                            806
                                                                                         for (T i : v) {
// custam O(log(n)) amortizado
                                                                            2a0
                                                                                             node* x = new node(i);
                                                                                             x -  ch[0] = root:
                                                                            bd1
081 template < typename T > struct splay {
```

```
37a
                if (root) root->p = x;
4ea
                root = x;
a0a
                root ->update();
17c
            }
c6b
a9e
        splay(const splay& t) {
e62
            throw logic_error("Nao copiar a splay!");
d4d
        }
5ab
       \simsplay() {
609
            vector < node *> q = {root};
402
            while (q.size()) {
e5d
                node* x = q.back(); q.pop_back();
ee9
                if (!x) continue:
                q.push_back(x->ch[0]), q.push_back(x->ch[1]);
73f
bf0
                delete x;
d3e
            }
a1c
        }
        int size(node* x) { return x ? x->sz : 0; }
73c
94f
        void rotate(node* x) { // x vai ficar em cima
d9b
            node *p = x->p, *pp = p->p;
ecf
            if (pp) pp - ch[pp - ch[1] == p] = x;
            bool d = p - ch[0] = x;
286
d63
            p - ch[!d] = x - ch[d], x - ch[d] = p;
bad
            if (p->ch[!d]) p->ch[!d]->p = p;
fc2
            x->p = pp, p->p = x;
1ea
            p->update(), x->update();
007
       }
6a0
        node* splaya(node* x) {
a39
            if (!x) return x;
            root = x, x->update();
be6
3cf
            while (x->p) {
                node *p = x->p, *pp = p->p;
d9b
359
                if (!pp) return rotate(x), x: // zig
                if ((pp->ch[0] == p)^(p->ch[0] == x))
e3c
a2b
                    rotate(x), rotate(x); // zigzag
4b2
                else rotate(p), rotate(x); // zigzig
028
            }
ea5
            return x;
21a
a7f
        node* find(int v) {
a2e
            if (!root) return NULL;
            node *x = root;
52f
6cd
            int key = 0;
31e
            while (1) {
857
                x->prop();
                bool d = key + size(x->ch[0]) < v;
ba1
```

```
877
                  if (\text{kev} + \text{size}(x->\text{ch}[0]) != v \text{ and } x->\text{ch}[d]) {
15e
                      if (d) key += size(x->ch[0])+1;
30e
                      x = x - ch[d];
a30
                 } else break;
3c3
152
             return splaya(x);
f19
        }
         int size() { return root ? root->sz : 0; }
c0c
c26
         void join(splay<T>& 1) { // assume gue 1 < *this</pre>
690
             if (!size()) swap(root, 1.root);
579
             if (!size() or !l.size()) return;
             node* x = 1.root:
bee
31e
             while (1) {
857
                 x->prop();
34d
                 if (!x->ch[1]) break;
bd8
                 x = x -> ch[1];
fa3
             }
             1.splaya(x), root->prop(), root->update();
147
42b
             x - ch[1] = root, x - ch[1] - p = x;
             root = 1.root, 1.root = NULL;
0aa
a0a
             root ->update();
7e6
         node* split(int v) { // retorna os elementos < v</pre>
5ed
398
             if (v <= 0) return NULL;</pre>
060
             if (v >= size()) {
f87
                 node* ret = root:
950
                  root = NULL;
8c9
                  ret ->update();
edf
                  return ret;
d0f
             }
             find(v);
adc
a59
             node* 1 = root -> ch[0];
             root -> ch [0] = NULL;
4df
5a3
             if (1) 1 \rightarrow p = NULL:
             root ->update();
a0a
792
             return 1;
826
        }
511
         T& operator [](int i) {
9d4
             find(i):
ae0
             return root ->val;
        }
829
231
         void push_back(T v) { // 0(1)
             node* r = new node(v);
a01
             r \rightarrow ch[0] = root;
0de
b11
             if (root) root->p = r;
             root = r, root->update();
b13
315
        }
```

```
b7a
        T query(int 1, int r) {
95f
            splay <T > M(split(r+1));
                                                                            bb7
                                                                                     node* root;
5ff
            splay <T> L(M.split(1));
                                                                            fd0
                                                                                     T N;
d1c
            T ans = M.root->sub;
49c
            M.join(L), join(M);
                                                                                     sms() : root(NULL), N(0) {}
                                                                            f34
                                                                                     sms(T v) : sms() { while (v >= N) N = 2*N+1; }
ba7
            return ans;
                                                                            83ъ
ca3
        }
                                                                            5e1
                                                                                     sms(const sms& t) : root(NULL), N(t.N) {
41f
                                                                            3af
                                                                                         for (SIZE_T i = 0; i < t.size(); i++) {</pre>
        void update(int 1, int r, T s) {
95f
            splay <T> M(split(r+1));
                                                                            a0f
                                                                                             T at = t[i];
5ff
                                                                                             SIZE_T qt = t.count(at);
            splay <T > L(M.split(1));
                                                                            e6d
996
            M.root->lazy += s;
                                                                            a43
                                                                                             insert(at, qt);
49c
            M.join(L), join(M);
                                                                            f42
                                                                                             i += qt-1;
9e9
                                                                            1e9
                                                                                         }
8c1
        void reverse(int 1, int r) {
                                                                            ea8
95f
            splay <T > M(split(r+1));
                                                                            a96
                                                                                     sms(initializer_list<T> v) : sms() { for (T i : v) insert(i); }
5ff
            splay <T> L(M.split(1));
                                                                            2dd
                                                                                     \simsms() {
945
            M.root->rev ^= 1;
                                                                            609
                                                                                         vector < node *> q = {root};
                                                                                         while (q.size()) {
49c
            M.join(L), join(M);
                                                                            402
                                                                                             node* x = q.back(); q.pop_back();
c1a
                                                                            e5d
2fb
        void erase(int 1, int r) {
                                                                                             if (!x) continue;
                                                                            ee9
            splay <T> M(split(r+1));
95f
                                                                            1 c 7
                                                                                             q.push_back(x->1), q.push_back(x->r);
5ff
            splay <T> L(M.split(1));
                                                                            bf0
                                                                                             delete x;
                                                                            653
                                                                                         }
dcc
            join(L);
68e
        }
                                                                            f0d
                                                                                     }
a35 };
                                                                            fdc
                                                                                     friend void swap(sms& a, sms& b) {
                                                                            49e
                                                                                         swap(a.root, b.root), swap(a.N, b.N);
1.33 Split-Merge Set
                                                                            984
                                                                                     }
                                                                            83e
                                                                                     sms& operator =(const sms& v) {
// Representa um conjunto de inteiros nao negativos
                                                                            768
                                                                                         sms tmp = v;
// Todas as operacoes custam O(log(N)),
                                                                            420
                                                                                         swap(tmp, *this);
// em que N = maior elemento do set,
                                                                            357
                                                                                         return *this;
// exceto o merge, que custa O(log(N)) amortizado
                                                                                     }
                                                                            e9b
// Usa O(\min(N, n \log(N))) de memoria, sendo 'n' o
                                                                            d06
                                                                                     SIZE T size() const { return root ? root->cnt : 0: }
// numero de elementos distintos no set
                                                                                     SIZE_T count(node* x) const { return x ? x->cnt : 0; }
                                                                            17f
                                                                            75a
                                                                                     void clear() {
2dc template < typename T, bool MULTI = false, typename SIZE_T = int > struct
                                                                            0a0
                                                                                         sms tmp;
   sms {
                                                                            4ac
                                                                                         swap(*this, tmp);
3c9
        struct node {
                                                                            fcb
b19
            node *1, *r;
                                                                            a06
                                                                                     void expand(T v) {
15f
            SIZE T cnt:
                                                                                         for (; N < v; N = 2*N+1) if (root) {
                                                                            bc3
            node() : 1(NULL), r(NULL), cnt(0) {}
658
                                                                            63 c
                                                                                             node* nroot = new node();
01e
            void update() {
                                                                            956
                                                                                             nroot ->1 = root;
a 0 1
                 cnt = 0:
                                                                            897
                                                                                             root = nroot;
d8a
                if (1) cnt += 1->cnt;
                                                                            a0a
                                                                                             root ->update();
e49
                if (r) cnt += r->cnt:
                                                                            449
                                                                                         }
74d
            }
                                                                            9f0
                                                                                     }
```

84f

};

```
node* insert(node* at, T idx, SIZE_T qt, T 1, T r) {
b14
                                                                             84a
1a4
            if (!at) at = new node();
                                                                              r);
893
            if (1 == r) {
                                                                             4e6
                                                                                     }
435
                at->cnt += qt;
                                                                             0a9
beb
                if (!MULTI) at->cnt = 1;
                                                                             ffc
ce6
                return at:
                                                                             df2
            }
a53
841
            T m = 1 + (r-1)/2;
                                                                             e68
            if (idx \le m) at->1 = insert(at->1, idx, qt, 1, m);
                                                                             809
a02
8d9
            else at->r = insert(at->r, idx, qt, m+1, r);
                                                                             c43
cff
            return at->update(), at;
                                                                             4a5
83b
                                                                             40c
cf7
        void insert(T v, SIZE_T qt=1) { // insere 'qt' ocorrencias de
                                                                             841
   , <sub>V</sub> ,
                                                                             5c2
882
            if (qt <= 0) return erase(v, -qt);</pre>
                                                                             4e6
72b
            assert(v >= 0);
                                                                             b4a
f52
            expand(v):
                                                                             ded
5e9
            root = insert(root, v, qt, 0, N);
                                                                             fa6
                                                                                          }
f62
        }
                                                                             41a
                                                                             792
f06
        node* erase(node* at, T idx, SIZE_T qt, T 1, T r) {
                                                                             67f
                                                                                      }
28c
            if (!at) return at:
54b
            if (1 == r) at->cnt = at->cnt < qt ? 0 : at->cnt - qt;
                                                                             78c
4e6
            else {
                                                                             347
                T m = 1 + (r-1)/2;
841
                                                                             504
                if (idx <= m) at->1 = erase(at->1, idx, qt, 1, m);
                                                                             599
281
                else at->r = erase(at->r, idx, qt, m+1, r);
ba1
                                                                             55d
7b4
                 at->update();
                                                                             792
d3d
            }
                                                                             92c
            if (!at->cnt) delete at, at = NULL;
135
                                                                             f58
ce6
            return at;
                                                                             f4f
                                                                             792
e1f
43d
        void erase(T v, SIZE_T qt=1) { // remove 'qt' ocorrencias de
                                                                             06a
                                                                                     }
   , v ,
                                                                             f59
9c3
            if (v < 0 \text{ or } v > N \text{ or } !qt) \text{ return};
                                                                             068
9dc
            if (qt < 0) insert(v, -qt);</pre>
                                                                             785
b1d
            root = erase(root, v, qt, 0, N);
                                                                             938
b32
                                                                             ee2
        void erase_all(T v) { // remove todos os 'v'
8d6
                                                                             2f6
                                                                                     }
347
            if (v < 0 \text{ or } v > N) return:
9f2
            root = erase(root, v, numeric_limits < SIZE_T >:: max(), 0, N);
                                                                             dc6
569
        }
                                                                             7ca
                                                                             640
0fe
        SIZE_T count(node* at, T a, T b, T 1, T r) const {
                                                                             386
61b
            if (!at or b < 1 or r < a) return 0:
                                                                             4e6
0fe
            if (a <= 1 and r <= b) return at->cnt;
                                                                             85 e
```

```
841
            T m = 1 + (r-1)/2:
            return count(at->1, a, b, 1, m) + count(at->r, a, b, m+1,
        SIZE_T count(T v) const { return count(root, v, v, 0, N); }
        SIZE_T order_of_key(T v) { return count(root, 0, v-1, 0, N); }
        SIZE_T lower_bound(T v) { return order_of_key(v); }
        const T operator [](SIZE_T i) const { // i-esimo menor elemento
            assert(i >= 0 and i < size());</pre>
            node* at = root;
            T 1 = 0, r = N;
            while (1 < r) {
                T m = 1 + (r-1)/2;
                if (count(at->1) > i) at = at->1, r = m;
                    i -= count(at->1);
                    at = at->r; 1 = m+1;
                }
            return 1;
        node* merge(node* 1, node* r) {
            if (!1 or !r) return 1 ? 1 : r;
            if (!1->1 and !1->r) { // folha
                if (MULTI) 1->cnt += r->cnt;
                delete r:
                return 1;
            1->1 = merge(1->1, r->1), 1->r = merge(1->r, r->r);
            1->update(), delete r;
            return 1:
        void merge(sms& s) { // mergeia dois sets
            if (N > s.N) swap(*this, s);
            expand(s.N);
            root = merge(root, s.root);
            s.root = NULL:
        node* split(node*& x, SIZE_T k) {
            if (k <= 0 or !x) return NULL;</pre>
            node* ret = new node():
            if (!x->1 \text{ and } !x->r) x->cnt -= k, ret->cnt += k;
                if (k \le count(x->1)) ret->1 = split(x->1, k);
```

```
4e6
                else {
06f
                    ret->r = split(x->r, k - count(x->1));
                    swap(x->1, ret->1);
cfd
63b
                }
674
                ret->update(), x->update();
379
d5b
            if (!x->cnt) delete x, x = NULL;
edf
            return ret;
f18
02b
        void split(SIZE_T k, sms& s) { // pega os 'k' menores
e63
            s.clear();
6e5
            s.root = split(root, min(k, size()));
e3c
9a6
        }
        // pega os menores que 'k'
        void split_val(T k, sms& s) { split(order_of_key(k), s); }
131
2d2 };
```

1.34 SQRT Tree

```
// RMQ em O(log log n) com O(n log log n) pra buildar
// Funciona com qualquer operacao associativa
// Tao rapido quanto a sparse table, mas usa menos memoria
// (log log (1e9) < 5, entao a query eh praticamente O(1))
// build - O(n log log n)
// query - O(log log n)
97a namespace sqrtTree {
052
        int n. *v:
        int pref[4][MAX], sulf[4][MAX], get1[4][MAX], entre[4][MAX],
ec7
   sz[4];
5f7
        int op(int a, int b) { return min(a, b); }
c72
        inline int getblk(int p, int i) { return (i-getl[p][i])/sz[p];
   }
2c6
        void build(int p, int 1, int r) {
bc8
            if (1+1 >= r) return;
            for (int i = 1; i <= r; i++) getl[p][i] = 1;</pre>
368
            for (int L = 1; L <= r; L += sz[p]) {</pre>
f16
                int R = min(L+sz[p]-1, r);
191
89 c
                pref[p][L] = v[L], sulf[p][R] = v[R];
                for (int i = L+1; i <= R; i++) pref[p][i] =</pre>
   op(pref[p][i-1], v[i]);
                for (int i = R-1; i >= L; i--) sulf[p][i] = op(v[i],
d9a
   sulf[p][i+1]);
221
                build(p+1, L, R);
```

```
c7b
            for (int i = 0; i <= sz[p]; i++) {</pre>
695
                int at = entre[p][l+i*sz[p]+i] = sulf[p][l+i*sz[p]];
ca5
759
                for (int j = i+1; j \le sz[p]; j++)
   entre[p][l+i*sz[p]+j] = at =
23a
                         op(at, sulf[p][l+j*sz[p]]);
c51
            }
        }
861
0d8
        void build(int n2, int* v2) {
680
            n = n2, v = v2:
44c
            for (int p = 0; p < 4; p++) sz[p] = n2 = sqrt(n2);
c50
            build(0, 0, n-1);
940
        }
9e3
        int query(int 1, int r) {
792
            if (1+1 >= r) return 1 == r ? v[1] : op(v[1], v[r]);
1ba
            int p = 0;
4ba
            while (getblk(p, 1) == getblk(p, r)) p++;
            int ans = sulf[p][1], a = getblk(p, 1)+1, b = getblk(p,
9e4
   r)-1;
8bf
            if (a <= b) ans = op(ans, entre[p][get1[p][1]+a*sz[p]+b]);</pre>
dea
            return op(ans, pref[p][r]);
589
        }
8ff }
1.35 Treap
```

```
// Todas as operacoes custam
// O(log(n)) com alta probabilidade, exceto meld
// meld custa O(log^2 n) amortizado com alta prob.,
// e permite unir duas treaps sem restricao adicional
// Na pratica, esse meld tem constante muito boa e
// o pior caso eh meio estranho de acontecer
878 mt19937 rng((int)
    chrono::steady_clock::now().time_since_epoch().count());
aa1 template < typename T > struct treap {
3c9
        struct node {
b19
            node *1, *r;
284
            int p, sz;
36d
            T val, mi;
4c7
            node(T v) : 1(NULL), r(NULL), p(rng()), sz(1), val(v),
   mi(v) {}
01e
            void update() {
                sz = 1:
a26
d6e
                mi = val;
bd7
                if (1) sz += 1->sz, mi = min(mi, 1->mi);
```

```
a54
                if (r) sz += r->sz, mi = min(mi, r->mi);
            }
660
        };
c1b
bb7
        node* root:
84b
        treap() { root = NULL: }
2d8
        treap(const treap& t) {
465
            throw logic_error("Nao copiar a treap!");
1e9
cec
        \simtreap() {
609
            vector < node *> q = {root};
402
            while (a.size()) {
e5d
                 node* x = q.back(); q.pop_back();
ee9
                if (!x) continue;
                q.push_back(x->1), q.push_back(x->r);
1c7
bf0
                 delete x;
653
            }
        }
50e
        int size(node* x) { return x ? x->sz : 0; }
73c
b2b
        int size() { return size(root); }
        void join(node* 1, node* r, node*& i) { // assume que 1 < r</pre>
bcf
986
            if (!1 or !r) return void(i = 1 ? 1 : r);
80e
            if (1->p > r->p) join(1->r, r, 1->r), i = 1;
fa0
            else join(1, r->1, r->1), i = r;
bda
            i->update();
671
        }
        void split(node* i, node*& 1, node*& r, T v) {
ece
26a
            if (!i) return void(r = 1 = NULL);
            if (i\rightarrow val < v) split(i\rightarrow r, i\rightarrow r, r, v), l = i;
f05
807
            else split(i \rightarrow 1, l, i \rightarrow 1, v), r = i;
            i->update():
bda
2cd
        }
3fc
        void split_leq(node* i, node*& 1, node*& r, T v) {
26a
            if (!i) return void(r = 1 = NULL);
            if (i->val <= v) split_leq(i->r, i->r, r, v), l = i;
181
58f
            else split_leg(i->1, 1, i->1, v), r = i;
bda
            i->update();
70f
        }
e13
        int count(node* i, T v) {
6b4
            if (!i) return 0;
352
            if (i->val == v) return 1;
            if (v < i->val) return count(i->1, v);
8d0
4d0
            return count(i->r, v);
5e6
        }
26d
        void index_split(node* i, node*& 1, node*& r, int v, int key =
```

```
) (O
26a
            if (!i) return void(r = l = NULL);
            if (key + size(i->1) < v) index_split(i->r, i->r, r, v,
c10
   key+size(i->1)+1), l = i;
            else index_split(i \rightarrow 1, 1, i \rightarrow 1, v, key), r = i;
e5a
bda
            i->update();
ccf
        }
a1f
        int count(T v) {
e06
            return count(root, v);
980
        }
c27
        void insert(T v) {
980
            if (count(v)) return;
0.31
            node *L. *R:
            split(root, L, R, v);
d42
585
            node* at = new node(v);
59f
            join(L, at, L);
a28
            join(L, R, root);
37c
        }
26b
        void erase(T v) {
df9
            node *L. *M. *R:
            split_leq(root, M, R, v), split(M, L, M, v);
b6b
f17
            if (M) delete M;
f38
            M = NULL:
a28
            join(L, R, root);
b92
e77
        void meld(treap& t) { // segmented merge
4a6
            node *L = root, *R = t.root;
950
            root = NULL:
6b1
            while (L or R) {
fe2
                 if (!L or (L and R and L->mi > R->mi)) std::swap(L, R);
                 if (!R) join(root, L, root), L = NULL;
5 e 1
                 else if (L->mi == R->mi) {
3c9
a76
                     node* LL:
439
                     split(L, LL, L, R->mi+1);
359
                     delete LL:
9d9
                } else {
a76
                     node* LL:
537
                     split(L, LL, L, R->mi);
dbb
                     join(root, LL, root);
fc4
                }
576
            }
689
            t.root = NULL;
8e7
651 };
```

1.36 Treap Implicita

```
// Todas as operacoes custam
// O(log(n)) com alta probabilidade
878 mt19937 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
aa1 template < typename T > struct treap {
        struct node {
3c9
b19
             node *1, *r;
284
             int p, sz;
875
            T val, sub, lazy;
             bool rev;
aa6
8dc
             node(T v) : 1(NULL), r(NULL), p(rng()), sz(1), val(v),
   sub(v), lazy(0), rev(0) {}
a9c
             void prop() {
0ec
                 if (lazy) {
                     val += lazy, sub += lazy*sz;
924
b87
                     if (1) 1->lazy += lazy;
d3b
                     if (r) r->lazv += lazv;
                 }
cea
1bb
                 if (rev) {
e4f
                     swap(1, r);
dc8
                     if (1) 1->rev ^= 1;
                     if (r) r->rev ^= 1;
f2f
3e5
a32
                 lazy = 0, rev = 0;
ca6
             }
01e
             void update() {
0 c 3
                 sz = 1, sub = val;
a09
                 if (1) 1 - prop(), sz += 1 - prop(), sz += 1 - prop();
095
                 if (r) r \rightarrow prop(), sz += r \rightarrow sz, sub += r \rightarrow sub;
360
             }
d37
        };
bb7
        node* root;
        treap() { root = NULL; }
84b
2d8
        treap(const treap& t) {
465
             throw logic_error("Nao copiar a treap!");
1e9
        }
cec
        \simtreap() {
             vector < node *> q = {root};
609
402
             while (q.size()) {
e5d
                 node* x = q.back(); q.pop_back();
ee9
                 if (!x) continue;
1 c 7
                 q.push_back(x->1), q.push_back(x->r);
bf0
                 delete x:
```

```
653
            }
        }
50e
73c
        int size(node* x) { return x ? x->sz : 0; }
        int size() { return size(root); }
b2b
        void join(node* 1, node* r, node*& i) { // assume que 1 < r</pre>
bcf
986
            if (!1 or !r) return void(i = 1 ? 1 : r);
            1->prop(), r->prop();
161
80e
            if (1->p > r->p) join(1->r, r, 1->r), i = 1;
            else join(1, r->1, r->1), i = r;
fa0
bda
            i->update();
b57
a20
        void split(node* i, node*& 1, node*& r, int v, int key = 0) {
26a
            if (!i) return void(r = 1 = NULL);
c89
            i->prop();
            if (key + size(i->1) < v) split(i->r, i->r, r, v,
5bd
   key+size(i->1)+1), l = i;
219
             else split(i \rightarrow 1, l, i \rightarrow 1, v, key), r = i;
bda
            i->update();
        }
d37
231
        void push_back(T v) {
2e0
            node* i = new node(v);
7ab
             join(root, i, root);
46b
        }
b7a
        T query(int 1, int r) {
df9
            node *L. *M. *R:
dca
             split(root, M, R, r+1), split(M, L, M, 1);
d43
            T ans = M->sub:
69d
            join(L, M, M), join(M, R, root);
ba7
            return ans;
1 f 7
41f
        void update(int 1, int r, T s) {
df9
             node *L. *M. *R:
dca
             split(root, M, R, r+1), split(M, L, M, 1);
8f6
            M \rightarrow lazv += s;
69d
             join(L, M, M), join(M, R, root);
29f
8c1
        void reverse(int 1, int r) {
df9
            node *L, *M, *R;
dca
             split(root, M, R, r+1), split(M, L, M, 1);
66a
            M \rightarrow rev = 1:
69d
             join(L, M, M), join(M, R, root);
ea8
        }
139 };
```

1.37 Treap Persistent Implicita

```
// Todas as operacoes custam
// O(log(n)) com alta probabilidade
6cf mt19937_64 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
3c9 struct node {
        node *1, *r;
b19
f14
       ll sz, val, sub;
        node(ll\ v): l(NULL), r(NULL), sz(1), val(v), sub(v) {}
        node(node* x) : 1(x->1), r(x->r), sz(x->sz), val(x->val),
   sub(x->sub) {}
01e
       void update() {
0 c 3
            sz = 1, sub = val;
77e
           if (1) sz += 1->sz, sub += 1->sub;
           if (r) sz += r->sz, sub += r->sub;
124
            sub %= MOD;
472
       }
95f };
bc9 ll size(node* x) { return x ? x->sz : 0; }
761 void update(node* x) { if (x) x->update(); }
828 node* copy(node* x) { return x ? new node(x) : NULL; }
b02 node* join(node* 1, node* r) {
e1f
        if (!1 or !r) return 1 ? copy(1) : copy(r);
48b
        node* ret:
49f
       if (rng() % (size(1) + size(r)) < size(1)) {</pre>
7eb
           ret = copy(1);
cc1
            ret -> r = join(ret -> r, r);
949
       } else {
4c5
            ret = copv(r);
            ret->1 = join(1, ret->1);
551
348
74f
        return update(ret), ret;
2cc }
723 void split(node* x, node*& 1, node*& r, 11 v, 11 key = 0) {
421
        if (!x) return void(1 = r = NULL):
b4b
        if (kev + size(x->1) < v) {
72f
           1 = copy(x);
d70
            split(1->r, 1->r, r, v, key+size(1->1)+1);
9d9
        } else {
303
            r = copy(x);
417
            split(r->1, 1, r->1, v, key);
474
da2
        update(1), update(r);
```

```
666 }
f9e vector < node *> treap;
139 void init(const vector<ll>& v) {
        treap = {NULL};
bbd
969
        for (auto i : v) treap[0] = join(treap[0], new node(i));
286 }
1.38 Wavelet Tree
// Usa O(sigma + n log(sigma)) de memoria,
// onde sigma = MAXN - MINN
// Depois do build, o v fica ordenado
// count(i, j, x, y) retorna o numero de elementos de
// v[i, j) que pertencem a [x, y]
// kth(i, j, k) retorna o elemento que estaria
// na poscicao k-1 de v[i, j), se ele fosse ordenado
// sum(i, j, x, y) retorna a soma dos elementos de
// v[i, j) que pertencem a [x, y]
// sumk(i, j, k) retorna a soma dos k-esimos menores
// elementos de v[i, j) (sum(i, j, 1) retorna o menor)
//
// Complexidades:
// build - O(n log(sigma))
// count - O(log(sigma))
// kth - 0(log(sigma))
// sum - O(log(sigma))
// sumk - O(log(sigma))
597 int n, v[MAX];
578 vector < int > esq[4*(MAXN-MINN)], pref[4*(MAXN-MINN)];
f8d void build(int b = 0, int e = n, int p = 1, int l = MINN, int r =
    MAXN) {
58f
        int m = (1+r)/2; esq[p].push_back(0); pref[p].push_back(0);
        for (int i = b; i < e; i++) {</pre>
f2f
            esq[p].push_back(esq[p].back()+(v[i]<=m));</pre>
6b9
26f
            pref[p].push_back(pref[p].back()+v[i]);
206
8ce
        if (1 == r) return;
3a7
        int m2 = stable_partition(v+b, v+e, [=](int i){return i <=</pre>
   m:}) - v:
347
        build(b, m2, 2*p, 1, m), build(m2, e, 2*p+1, m+1, r);
Ofb }
540 int count(int i, int j, int x, int y, int p = 1, int l = MINN, int
```

```
r = MAXN)  {
        if (y < 1 \text{ or } r < x) \text{ return } 0;
        if (x \le 1 \text{ and } r \le y) \text{ return } j-i;
4db
       int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
ddc
        return count(ei, ej, x, y, 2*p, l, m)+count(i-ei, j-ej, x, y,
0a5
   2*p+1, m+1, r);
3cf }
f62 int kth(int i, int j, int k, int p=1, int l = MINN, int r = MAXN) {
        if (1 == r) return 1:
ddc
        int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
585
        if (k <= ej-ei) return kth(ei, ej, k, 2*p, 1, m);</pre>
28b
        return kth(i-ei, j-ej, k-(ej-ei), 2*p+1, m+1, r);
8b6 }
f2c int sum(int i, int j, int x, int y, int p = 1, int l = MINN, int r
   = MAXN)
        if (v < 1 \text{ or } r < x) return 0:
2ad
        if (x <= 1 and r <= y) return pref[p][i]-pref[p][i];</pre>
2a9
       int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
ddc
        return sum(ei, ej, x, y, 2*p, 1, m) + sum(i-ei, j-ej, x, y,
   2*p+1, m+1, r);
b6d }
b84 int sumk(int i, int j, int k, int p = 1, int l = MINN, int r =
   MAXN) {
8a1
       if (1 == r) return 1*k;
ddc
       int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
       if (k <= ej-ei) return sumk(ei, ej, k, 2*p, 1, m);</pre>
        return pref[2*p][ei]-pref[2*p][ei]+sumk(i-ei, j-ej, k-(ej-ei),
   2*p+1, m+1, r);
940 }
2 Grafos
```

2.1 AGM Directionada

```
// Fala o menor custo para selecionar arestas tal que
// o vertice 'r' alcance todos
// Se nao tem como, retorna LINF
//
// O(m log(n))

3c9 struct node {
f31    pair<11, int> val;
4e4    ll lazy;
```

```
b19
        node *1, *r;
f93
        node() {}
        node(pair < int, int > v) : val(v), lazy(0), l(NULL), r(NULL) {}
c53
        void prop() {
a9c
768
            val.first += lazy;
b87
            if (1) 1->lazy += lazy;
d3b
            if (r) r->lazy += lazy;
            lazv = 0;
c60
        }
296 };
de5 void merge(node*& a, node* b) {
        if (!a) swap(a, b);
802
        if (!b) return;
626
        a->prop(), b->prop();
        if (a->val > b->val) swap(a, b);
d04
4b0
        merge(rand()%2 ? a->1 : a->r, b);
b82 }
d01 pair<ll, int> pop(node*& R) {
e8f
        R->prop();
        auto ret = R->val;
22e
        node* tmp = R;
af0
3f3
        merge(R->1, R->r);
6c9
        R = R - > 1;
3e4
        if (R) R->lazy -= ret.first;
7c3
        delete tmp:
edf
        return ret;
c4e }
6f6 void apaga(node* R) { if (R) apaga(R->1), apaga(R->r), delete R; }
f13 ll dmst(int n, int r, vector<pair<int, int>, int>>& ar) {
94e
        vector < int > p(n); iota(p.begin(), p.end(), 0);
        function < int(int) > find = [&](int k) { return
a23
   p[k]==k?k:p[k]=find(p[k]): }:
2d7
        vector < node *> h(n);
56f
        for (auto e : ar) merge(h[e.first.second], new node({e.second,
   e.first.first}));
fd1
        vector < int > pai(n, -1), path(n);
66e
        pai[r] = r;
04b
        11 \text{ ans} = 0;
603
        for (int i = 0; i < n; i++) { // vai conectando todo mundo
2a3
            int u = i. at = 0:
cae
            while (pai[u] == -1) {
daa
                if (!h[u]) { // nao tem
947
                    for (auto i : h) apaga(i);
77c
                    return LINF:
```

```
dd1
                path[at++] = u, pai[u] = i;
167
55e
                auto [mi, v] = pop(h[u]);
64 c
                ans += mi;
5e2
                if (pai[u = find(v)] == i) { // ciclo
86f
                    while (find(v = path[--at]) != u)
                        merge(h[u], h[v]), h[v] = NULL, p[find(v)] = u;
621
57a
                    pai[u] = -1;
                }
850
            }
ce8
5df
947
        for (auto i : h) apaga(i);
ba7
        return ans;
e02 }
2.2 Articulation Points
```

```
// Computa os pontos de articulação (vertices criticos) de um grafo
// art[i] armazena o numero de novas componentes criadas ao deletar
   vertice i
// se art[i] >= 1, entao vertice i eh ponto de articulacao
//
// O(n+m)
1a8 int n:
789 vector < vector < int >> g;
4ce stack < int > s;
b66 vector < int > id, art;
3e1 int dfs_art(int i, int& t, int p = -1) {
cf0
        int lo = id[i] = t++;
18e
        s.push(i);
        for (int j : g[i]) if (j != p) {
            if (id[i] == -1) {
9a3
206
                int val = dfs_art(j, t, i);
                lo = min(lo, val);
0 c 3
                if (val >= id[i]) {
588
                    art[i]++:
66a
bd9
                    while (s.top() != j) s.pop();
2eb
                    s.pop();
1f3
                // if (val > id[i]) aresta i-j eh ponte
238
            }
328
            else lo = min(lo, id[j]);
```

```
762
3bd
        if (p == -1 and art[i]) art[i]--;
253
        return lo;
8e1 }
d79 void compute_art_points() {
        id = vector<int>(n, -1);
597
        art = vector<int>(n, 0);
a62
6bb
        int t = 0;
        for (int i = 0; i < n; i++) if (id[i] == -1)
d41
625
            dfs_art(i, t, -1);
379 }
```

2.3 Bellman-Ford

```
// Calcula a menor distancia
// entre a e todos os vertices e
// detecta ciclo negativo
// Retorna 1 se ha ciclo negativo
// Nao precisa representar o grafo,
// soh armazenar as arestas
//
// O(nm)
14e int n, m;
248 int d[MAX];
e93 vector <pair <int. int >> ar: // vetor de arestas
9e2 vector < int > w;
                                // peso das arestas
6be bool bellman ford(int a) {
        for (int i = 0; i < n; i++) d[i] = INF;</pre>
8a8
        d[a] = 0:
        for (int i = 0: i <= n: i++)
4e3
891
            for (int j = 0; j < m; j++) {
6e4
                if (d[ar[j].second] > d[ar[j].first] + w[j]) {
705
                    if (i == n) return 1;
e93
                    d[ar[j].second] = d[ar[j].first] + w[j];
84b
                }
a82
            }
bb3
        return 0;
6eb }
```

2.4 Block-Cut Tree

```
// Cria a block-cut tree, uma arvore com os blocos
// e os pontos de articulação
// Blocos sao componentes 2-vertice-conexos maximais
// Uma 2-coloração da arvore eh tal que uma cor são
// os blocos, e a outra cor sao os pontos de art.
// Funciona para grafo nao conexo
//
// art[i] responde o numero de novas componentes conexas
// criadas apos a remocao de i do grafo g
// Se art[i] >= 1, i eh ponto de articulação
//
// Para todo i <= blocks.size()</pre>
// blocks[i] eh uma componente 2-vertce-conexa maximal
// edgblocks[i] sao as arestas do bloco i
// tree[i] eh um vertice da arvore que corresponde ao bloco i
// pos[i] responde a qual vertice da arvore vertice i pertence
// Arvore tem no maximo 2n vertices
//
// O(n+m)
d10 struct block_cut_tree {
        vector < vector < int >> g, blocks, tree;
43b
        vector < vector < pair < int , int >>> edgblocks;
4ce
        stack<int> s:
6c0
        stack<pair<int, int>> s2;
2bb
        vector < int > id, art, pos;
763
        block_cut_tree(vector<vector<int>> g_) : g(g_) {
af1
            int n = g.size();
37a
            id.resize(n, -1), art.resize(n), pos.resize(n);
6f2
            build();
       }
6bd
        int dfs(int i, int& t, int p = -1) {
df6
cf0
            int lo = id[i] = t++;
18e
            s.push(i);
827
            if (p != -1) s2.emplace(i, p);
53f
            for (int j : g[i]) if (j != p and id[j] != -1)
   s2.emplace(i, j);
            for (int j : g[i]) if (j != p) {
cac
                if (id[j] == -1) {
9a3
121
                    int val = dfs(j, t, i);
                    lo = min(lo, val);
0c3
```

```
588
                     if (val >= id[i]) {
66a
                         art[i]++:
483
                         blocks.emplace_back(1, i);
                         while (blocks.back().back() != j)
110
                             blocks.back().push_back(s.top()), s.pop();
138
128
                         edgblocks.emplace_back(1, s2.top()), s2.pop();
47e
                         while (edgblocks.back().back() != pair(j, i))
                             edgblocks.back().push_back(s2.top()),
bce
    s2.pop();
870
                     // if (val > id[i]) aresta i-j eh ponte
85c
328
                 else lo = min(lo, id[j]);
             }
344
3bd
             if (p == -1 and art[i]) art[i]--;
253
             return lo:
726
        }
0a8
        void build() {
6bb
             int t = 0:
             for (int i = 0; i < g.size(); i++) if (id[i] == -1) dfs(i,</pre>
abf
    t, -1);
56c
             tree.resize(blocks.size()):
             for (int i = 0; i < g.size(); i++) if (art[i])</pre>
f7d
                 pos[i] = tree.size(), tree.emplace_back();
965
973
             for (int i = 0; i < blocks.size(); i++) for (int j :</pre>
    blocks[i]) {
403
                 if (!art[i]) pos[i] = i;
                 else tree[i].push_back(pos[j]),
101
    tree[pos[j]].push_back(i);
3df
        }
c03
056 };
2.5 Blossom
// Matching maximo em grafo geral
//
// O(n^3)
// Se for bipartido, nao precisa da funcao
// 'contract', e roda em O(nm)
042 vector < int > g[MAX];
```

```
128 int match[MAX]; // match[i] = com quem i esta matchzado ou -1
1f1 int n, pai[MAX], base[MAX], vis[MAX];
26a queue < int > q;
107 void contract(int u. int v. bool first = 1) {
        static vector < bool > bloss:
165
fbe
        static int 1:
418
        if (first) {
a47
            bloss = vector < bool > (n, 0);
042
            vector < bool > teve(n. 0);
ddf
            int k = u; l = v;
31e
            while (1) {
297
                teve[k = base[k]] = 1:
116
                if (match[k] == -1) break;
dfa
                k = pai[match[k]];
68b
            }
d31
            while (!teve[l = base[l]]) l = pai[match[l]];
5d6
2e9
        while (base[u] != 1) {
            bloss[base[u]] = bloss[base[match[u]]] = 1:
e29
            pai[u] = v;
8fa
            v = match[u];
0b0
            u = pai[match[u]];
a51
58e
        }
71c
        if (!first) return;
95e
        contract(v. u. 0);
6ee
        for (int i = 0; i < n; i++) if (bloss[base[i]]) {</pre>
594
            base[i] = 1:
            if (!vis[i]) q.push(i);
ca7
29a
            vis[i] = 1;
857
       }
e35 }
f10 int getpath(int s) {
        for (int i = 0; i < n; i++) base[i] = i, pai[i] = -1, vis[i] =
88f
   0;
        vis[s] = 1; q = queue < int > (); q.push(s);
ded
        while (q.size()) {
402
be1
            int u = q.front(); q.pop();
bdc
            for (int i : g[u]) {
                if (base[i] == base[u] or match[u] == i) continue;
7a2
                if (i == s or (match[i] != -1 and pai[match[i]] != -1))
e35
                    contract(u. i):
4f2
                else if (pai[i] == -1) {
e2e
545
                    pai[i] = u;
                    if (match[i] == -1) return i;
f6a
                    i = match[i]:
818
```

```
29d
                     vis[i] = 1; q.push(i);
90e
                }
0b5
            }
634
        }
daa
        return -1;
a16 }
83f int blossom() {
        int ans = 0;
1a4
315
        memset(match, -1, sizeof(match));
2e3
        for (int i = 0; i < n; i++) if (match[i] == -1)</pre>
f76
            for (int j : g[i]) if (match[j] == -1) {
1bc
                 match[i] = i:
f1d
                 match[i] = i;
Odf
                 ans++;
c2b
                 break;
723
            }
        for (int i = 0; i < n; i++) if (match[i] == -1) {</pre>
da8
7e3
            int j = getpath(i);
5f2
            if (j == -1) continue;
Odf
            ans++;
3a0
            while (j != -1) {
ef0
                 int p = pai[j], pp = match[p];
348
                 match[p] = j;
fe9
                match[j] = p;
55d
                j = pp;
797
            }
f70
        }
ba7
        return ans;
fcd }
2.6 Centro de arvore
   diametro
```

```
// Retorna o diametro e o(s) centro(s) da arvore
// Uma arvore tem sempre um ou dois centros e estes estao no meio do
//
// O(n)
042 vector < int > g[MAX];
df1 int d[MAX], par[MAX];
544 pair<int, vector<int>> center() {
a95
        int f, df;
36d
        function < void(int) > dfs = [&] (int v) {
d47
            if (d[v] > df) f = v, df = d[v];
e68
            for (int u : g[v]) if (u != par[v])
```

```
1a5
                d[u] = d[v] + 1, par[u] = v, dfs(u);
90d
        f = df = par[0] = -1, d[0] = 0;
41e
        dfs(0):
        int root = f;
c2d
        f = df = par[root] = -1, d[root] = 0:
0f6
        dfs(root):
14e
761
        vector<int> c:
87e
        while (f != -1) {
999
            if (d[f] == df/2 \text{ or } d[f] == (df+1)/2) \text{ c.push back}(f):
19c
            f = par[f]:
3bf
        }
00f
        return {df, c};
9c7 }
2.7 Centroid
```

```
// Computa os 2 centroids da arvore
// O(n)
97a int n, subsize[MAX];
042 vector < int > g[MAX];
98f void dfs(int k, int p=-1) {
        subsize[k] = 1;
bd2
       for (int i : g[k]) if (i != p) {
6e5
801
            dfs(i, k);
2e3
            subsize[k] += subsize[i]:
1b2
       }
5a5 }
2e8 int centroid(int k, int p=-1, int size=-1) {
       if (size == -1) size = subsize[k];
e73
8df
       for (int i : g[k]) if (i != p) if (subsize[i] > size/2)
bab
            return centroid(i, k, size);
839
       return k:
b6a }
f20 pair <int, int > centroids(int k=0) {
051
       dfs(k);
909
        int i = centroid(k), i2 = i;
8dd
       for (int j : g[i]) if (2*subsize[j] == subsize[k]) i2 = j;
0cb
       return {i, i2};
```

```
cf4 }
```

2.8 Centroid decomposition

```
// decomp(0, k) computa numero de caminhos com 'k' arestas
// Mudar depois do comentario
//
// O(n log(n))
042 vector < int > g[MAX];
ba8 int sz[MAX], rem[MAX];
747 void dfs(vector<int>& path, int i, int l=-1, int d=0) {
        path.push_back(d);
        for (int j : g[i]) if (j != 1 and !rem[j]) dfs(path, j, i,
75f
   d+1):
3e9 }
071 int dfs sz(int i. int l=-1) {
        sz[i] = 1:
e5c
        for (int j : g[i]) if (j != l and !rem[j]) sz[i] += dfs_sz(j,
   i):
191
        return sz[i];
86b }
85a int centroid(int i, int 1, int size) {
        for (int i : g[i]) if (i != 1 and !rem[i] and sz[i] > size / 2)
735
            return centroid(j, i, size);
d9a
        return i;
96e }
d79 ll decomp(int i, int k) {
106
        int c = centroid(i, i, dfs_sz(i));
a67
        rem[c] = 1;
        // gasta O(n) aqui - dfs sem ir pros caras removidos
04b
        11 \text{ ans} = 0:
020
        vector < int > cnt(sz[i]);
878
        cnt[0] = 1;
        for (int j : g[c]) if (!rem[j]) {
0a8
5b4
            vector < int > path;
            dfs(path, j);
baf
1a1
            for (int d: path) if (0 \le k-d-1 \text{ and } k-d-1 \le sz[i])
285
                ans += cnt[k-d-1];
e8b
            for (int d : path) cnt[d+1]++;
fa2
        }
```

```
1c1     for (int j : g[c]) if (!rem[j]) ans += decomp(j, k);
3f1     rem[c] = 0;
ba7     return ans;
193 }
```

2.9 Centroid Tree

```
// Constroi a centroid tree
// p[i] eh o pai de i na centroid-tree
// dist[i][k] = distancia na arvore original entre i
// e o k-esimo ancestral na arvore da centroid
// O(n log(n)) de tempo e memoria
845 vector < int > g[MAX], dist[MAX];
c1e int sz[MAX], rem[MAX], p[MAX];
071 int dfs_sz(int i, int l=-1) {
02c
        sz[i] = 1:
e5c
        for (int j : g[i]) if (j != l and !rem[j]) sz[i] += dfs_sz(j,
   i):
191
        return sz[i];
86b }
85a int centroid(int i, int 1, int size) {
        for (int j : g[i]) if (j != 1 and !rem[j] and sz[j] > size / 2)
994
735
            return centroid(i, i, size):
d9a
        return i;
96e }
324 void dfs_dist(int i, int 1, int d=0) {
541
        dist[i].push_back(d);
5a1
        for (int j : g[i]) if (j != l and !rem[j])
82a
            dfs_dist(j, i, d+1);
645 }
27e void decomp(int i, int l = -1) {
        int c = centroid(i, i, dfs_sz(i));
106
        rem[c] = 1, p[c] = 1;
1b9
534
        dfs dist(c. c):
        for (int j : g[c]) if (!rem[j]) decomp(j, c);
a2a
ebd }
76c void build(int n) {
235
        for (int i = 0; i < n; i++) rem[i] = 0, dist[i].clear();</pre>
867
        decomp(0);
96b
        for (int i = 0; i < n; i++) reverse(dist[i].begin(),</pre>
```

```
dist[i].end());
a78 }
2.10 Diikstra
// encontra menor distancia de x
// para todos os vertices
// se ao final do algoritmo d[i] = LINF,
// entao x nao alcanca i
// O(m log(n))
eff ll d[MAX]:
c0d vector<pair<int, int>> g[MAX]; // {vizinho, peso}
1a8 int n;
abc void dijkstra(int v) {
        for (int i = 0; i < n; i++) d[i] = LINF;</pre>
        d[v] = 0:
a7f
88c
        priority_queue < pair < ll, int >> pq;
        pq.emplace(0, v);
b32
265
         while (pq.size()) {
             auto [ndist, u] = pq.top(); pq.pop();
a25
             if (-ndist > d[u]) continue;
953
             for (auto [idx, w] : g[u]) if (d[idx] > d[u] + w) {
cda
331
                 d[idx] = d[u] + w;
a84
                 pq.emplace(-d[idx], idx);
c56
             }
        }
e5c
fec }
2.11 Dinitz
// O(min(m * max_flow, n^2 m))
// Grafo com capacidades 1: O(\min(m \text{ sqrt}(m), m * n^{2/3}))
// Todo vertice tem grau de entrada ou saida 1: O(m sqrt(n))
472 struct dinitz {
         const bool scaling = false; // com scaling -> 0(nm
   log(MAXCAP)),
206
        int lim;
                                     // com constante alta
670
         struct edge {
358
             int to, cap, rev, flow;
7f9
             bool res;
```

```
d36
            edge(int to_, int cap_, int rev_, bool res_)
                : to(to_), cap(cap_), rev(rev_), flow(0), res(res_) {}
a94
f70
        };
002
        vector < vector < edge >> g;
216
        vector < int > lev, beg;
a71
        11 F:
190
        dinitz(int n) : g(n), F(0) {}
        void add(int a, int b, int c) {
087
bae
            g[a].emplace_back(b, c, g[b].size(), false);
4c6
            g[b].emplace_back(a, 0, g[a].size()-1, true);
5c2
123
        bool bfs(int s, int t) {
            lev = vector \langle int \rangle (g.size(), -1); lev[s] = 0;
90f
64 c
            beg = vector<int>(g.size(), 0);
8b2
            queue < int > q; q.push(s);
            while (q.size()) {
402
                int u = q.front(); q.pop();
be1
                for (auto& i : g[u]) {
bd9
                    if (lev[i.to] != -1 or (i.flow == i.cap)) continue;
dbc
b4f
                    if (scaling and i.cap - i.flow < lim) continue;</pre>
                    lev[i.to] = lev[u] + 1;
185
                    q.push(i.to);
8ca
f97
                }
e87
            }
            return lev[t] != -1;
0de
742
        }
dfb
        int dfs(int v, int s, int f = INF) {
50b
            if (!f or v == s) return f;
            for (int& i = beg[v]; i < g[v].size(); i++) {</pre>
88f
027
                auto& e = g[v][i];
                if (lev[e.to] != lev[v] + 1) continue;
206
                int foi = dfs(e.to, s, min(f, e.cap - e.flow));
ee0
                if (!foi) continue:
749
3c5
                e.flow += foi, g[e.to][e.rev].flow -= foi;
                return foi;
45c
618
            }
bb3
            return 0;
4b1
ff6
        11 max flow(int s. int t) {
            for (lim = scaling ? (1 << 30) : 1; lim; lim /= 2)
a86
9d1
                while (bfs(s, t)) while (int ff = dfs(s, t)) F += ff;
4ff
            return F:
8b9
        }
86f };
```

```
// Recupera as arestas do corte s-t
dbd vector<pair<int, int>> get_cut(dinitz& g, int s, int t) {
        g.max_flow(s, t);
f07
        vector < pair < int , int >> cut;
68 c
        vector < int > vis(g.g.size(), 0), st = {s};
1b0
321
        vis[s] = 1;
3c6
        while (st.size()) {
b17
            int u = st.back(); st.pop_back();
322
            for (auto e : g.g[u]) if (!vis[e.to] and e.flow < e.cap)</pre>
                vis[e.to] = 1, st.push_back(e.to);
c17
d14
        }
481
        for (int i = 0; i < g.g.size(); i++) for (auto e : g.g[i])</pre>
9d2
            if (vis[i] and !vis[e.to] and !e.res) cut.emplace back(i.
   e.to):
d1b
        return cut;
1e8 }
2.12 Dominator Tree
// Codigo do Kawakami. Se vira pra usar ai
//
// build - O(m log(n))
// dominates - 0(1)
1a8 int n;
bbf namespace d tree {
        vector < int > g[MAX];
        // The dominator tree
b39
        vector<int> tree[MAX];
```

int dfs_l[MAX], dfs_r[MAX];

vector < int > rg[MAX], bucket[MAX];

int ancestor[MAX], label[MAX];

sdom[v] = label[v] = v;

if (sdom[nxt] == -1) {

prv[nxt] = v;

preorder.push_back(v);

for (int nxt: g[v]) {

int idom[MAX], sdom[MAX], prv[MAX], pre[MAX];

// Auxiliary data

void dfs(int v) {

vector<int> preorder;

pre[v] = ++t:

static int t = 0;

5af

a2e

3ef

44b

563

76a

6a1

db6

767

a3d

806

56c

eed

```
900
                    dfs(nxt);
f48
2b5
                rg[nxt].push_back(v);
            }
5a1
        }
d6a
62e
        int eval(int v) {
c93
            if (ancestor[v] == -1) return v:
            if (ancestor[ancestor[v]] == -1) return label[v];
a75
f33
            int u = eval(ancestor[v]);
            if (pre[sdom[u]] < pre[sdom[label[v]]]) label[v] = u;</pre>
b49
66e
            ancestor[v] = ancestor[u];
c24
            return label[v]:
0b9
       }
4b2
        void dfs2(int v) {
6a1
            static int t = 0;
330
            dfs_1[v] = t++;
5e0
            for (int nxt: tree[v]) dfs2(nxt);
            dfs_r[v] = t++;
8e2
cfa
        void build(int s) {
c2c
            for (int i = 0: i < n: i++) {</pre>
603
e6f
                sdom[i] = pre[i] = ancestor[i] = -1;
                rg[i].clear();
2e1
50a
                tree[i].clear();
666
                bucket[i].clear();
3ba
772
            preorder.clear();
c6c
            dfs(s):
            if (preorder.size() == 1) return;
12b
3c7
            for (int i = int(preorder.size()) - 1; i >= 1; i--) {
                int w = preorder[i];
6.6
a52
                for (int v: rg[w]) {
                    int u = eval(v):
5c1
a17
                    if (pre[sdom[u]] < pre[sdom[w]]) sdom[w] = sdom[u];</pre>
018
680
                bucket[sdom[w]].push_back(w);
                ancestor[w] = prv[w];
ea7
b99
                for (int v: bucket[prv[w]]) {
5c1
                    int u = eval(v):
977
                    idom[v] = (u == v) ? sdom[v] : u;
aff
                bucket[prv[w]].clear();
2cc
0a3
            for (int i = 1; i < preorder.size(); i++) {</pre>
d0c
6c6
                int w = preorder[i];
                if (idom[w] != sdom[w]) idom[w] = idom[idom[w]];
14b
                tree[idom[w]].push_back(w);
32f
```

```
c58
            idom[s] = sdom[s] = -1;
8ac
1 b 6
            dfs2(s);
d09
        }
        // Whether every path from s to v passes through u
490
        bool dominates(int u. int v) {
c75
            if (pre[v] == -1) return 1; // vacuously true
2ea
            return dfs_l[u] <= dfs_l[v] && dfs_r[v] <= dfs_r[u];</pre>
332
        }
ce9 };
```

2.13 Euler Path / Euler Cycle

```
// Para declarar: 'euler <true > E(n); ' se guiser
// direcionado e com 'n' vertices
// As funcoes retornam um par com um booleano
// indicando se possui o cycle/path que voce pediu,
// e um vector de {vertice, id da aresta para chegar no vertice}
// Se for get_path, na primeira posicao o id vai ser -1
// get path(src) tenta achar um caminho ou ciclo euleriano
// comecando no vertice 'src'.
// Se achar um ciclo, o primeiro e ultimo vertice serao 'src'.
// Se for um P3, um possiveo retorno seria [0, 1, 2, 0]
// get_cycle() acha um ciclo euleriano se o grafo for euleriano.
// Se for um P3, um possivel retorno seria [0, 1, 2]
// (vertie inicial nao repete)
//
// O(n+m)
63f template <bool directed=false > struct euler {
1a8
4c0
        vector < vector < pair < int , int >>> g;
d63
        vector<int> used;
30f
        euler(int n_) : n(n_), g(n) {}
50f
        void add(int a. int b) {
4cd
            int at = used.size();
c51
            used.push_back(0);
74 e
            g[a].emplace_back(b, at);
fab
            if (!directed) g[b].emplace_back(a, at);
411
        }
d41 #warning chamar para o src certo!
        pair < bool, vector < pair < int, int >>> get_path(int src) {
eed
baf
            if (!used.size()) return {true, {}};
b25
            vector < int > beg(n, 0);
4ec
            for (int& i : used) i = 0;
```

```
// {{vertice, anterior}, label}
363
            vector<pair<pair<int, int>, int>> ret, st = {{{src, -1}},
   -1}};
3c6
            while (st.size()) {
8ff
                int at = st.back().first.first;
                int& it = beg[at];
002
8 a 1
                while (it < g[at].size() and used[g[at][it].second])</pre>
   it++;
                if (it == g[at].size()) {
8e4
944
                    if (ret.size() and ret.back().first.second != at)
b82
                         return {false, {}};
420
                    ret.push_back(st.back()), st.pop_back();
9d9
                } else {
                    st.push_back({{g[at][it].first, at},
   g[at][it].second});
eb8
                    used[g[at][it].second] = 1;
                }
d14
b3a
a19
            if (ret.size() != used.size()+1) return {false, {}};
f77
            vector < pair < int , int >> ans;
            for (auto i : ret) ans.emplace_back(i.first.first,
fdf
   i.second):
459
            reverse(ans.begin(), ans.end());
997
            return {true, ans};
844
9b6
        pair < bool, vector < pair < int, int >>> get_cycle() {
baf
            if (!used.size()) return {true, {}};
ad1
            int src = 0:
34b
            while (!g[src].size()) src++;
687
            auto ans = get_path(src);
            if (!ans.first or ans.second[0].first !=
   ans.second.back().first)
                return {false, {}};
b82
350
            ans.second[0].second = ans.second.back().second:
8b8
            ans.second.pop_back();
ba7
            return ans;
48f
        }
711 };
2.14 Euler Tour Tree
// Mantem uma floresta enraizada dinamicamente
// e permite queries/updates em sub-arvore
// Chamar ETT E(n, v), passando n = numero de vertices
// e v = vector com os valores de cada vertice (se for vazio,
```

// constroi tudo com 0

```
//
// link(v, u) cria uma aresta de v pra u, de forma que u se torna
// o pai de v (eh preciso que v seja raiz anteriormente)
// cut(v) corta a resta de v para o pai
// query(v) retorna a soma dos valores da sub-arvore de v
// update(v, val) soma val em todos os vertices da sub-arvore de v
// update_v(v, val) muda o valor do vertice v para val
// is_in_subtree(v, u) responde se o vertice u esta na sub-arvore de v
// Tudo O(log(n)) com alta probabilidade
878 mt19937 rng((int)
    chrono::steady_clock::now().time_since_epoch().count());
9f9 template < typename T > struct ETT {
        // treap
3c9
        struct node {
ed1
            node *1, *r, *p;
fa4
            int pr, sz;
875
            T val, sub, lazy;
53e
            int id;
ffd
            bool f; // se eh o 'first'
5ef
            int qt_f; // numero de firsts na subarvore
            node(int id_, T v, bool f_ = 0) : 1(NULL), r(NULL),
7a8
   p(NULL), pr(rng()),
62b
                 sz(1), val(v), sub(v), lazy(), id(id_), f(f_),
   qt_f(f_) {}
a9c
            void prop() {
d09
                 if (lazy != T()) {
021
                     if (f) val += lazy;
971
                     sub += lazy*sz;
b87
                     if (1) 1->lazy += lazy;
d3b
                     if (r) r->lazy += lazy;
30c
                }
                 lazy = T();
bfd
0 f c
            }
01e
            void update() {
8da
                 sz = 1, sub = val, qt_f = f;
                 if (1) 1 - \text{prop}(), sz += 1 - \text{sz}, sub += 1 - \text{sub}, qt_f +=
171
   1->qt_f;
117
                 if (r) r - prop(), sz += r - sz, sub += r - sub, qt_f +=
   r->qt_f;
ccb
bff
        };
bb7
        node* root;
```

```
73c
        int size(node* x) { return x ? x->sz : 0; }
                                                                           8ac
                                                                                       }
        void join(node* 1, node* r, node*& i) { // assume que 1 < r</pre>
bcf
                                                                           ec3
986
            if (!1 or !r) return void(i = 1 ? 1 : r);
                                                                           83f
                                                                                   ETT(const ETT& t) { throw logic_error("Nao copiar a ETT!"); }
            1->prop(), r->prop();
                                                                           c09
                                                                                   \simETT() {
161
ff5
            if (1->pr > r->pr) join(1->r, r, 1->r), 1->r->p = i = 1;
                                                                           609
                                                                                       vector < node *> q = {root};
            else join(1, r->1, r->1), r->1->p = i = r;
982
                                                                           402
                                                                                       while (q.size()) {
bda
            i->update();
                                                                           e5d
                                                                                           node* x = q.back(); q.pop_back();
84d
        }
                                                                                           if (!x) continue;
                                                                           ee9
a20
        void split(node* i, node*& 1, node*& r, int v, int key = 0) {
                                                                           1c7
                                                                                           q.push_back(x->1), q.push_back(x->r);
            if (!i) return void(r = 1 = NULL);
26a
                                                                           bf0
                                                                                           delete x:
c89
            i->prop();
                                                                           653
                                                                                       }
d9e
            if (key + size(i->1) < v) {
                                                                           672
                                                                                   }
                split(i->r, i->r, r, v, key+size(i->l)+1), l = i;
448
                if (r) r -> p = NULL;
a21
                                                                           153
                                                                                   pair<int, int> get_range(int i) {
6e8
                if (i->r) i->r->p = i;
                                                                           670
                                                                                       return {get_idx(first[i]), get_idx(last[i])};
9d9
            } else {
                                                                           ada
98d
                split(i->1, 1, i->1, v, key), r = i;
                                                                           7af
                                                                                   void link(int v, int u) { // 'v' tem que ser raiz
                                                                                       auto [lv, rv] = get_range(v);
5a3
                if (1) 1->p = NULL;
                                                                           890
                if (i->1) i->1->p = i;
                                                                           f13
                                                                                       int ru = get_idx(last[u]);
899
            }
a3f
bda
            i->update();
                                                                           4b4
                                                                                       node* V;
134
                                                                           df9
                                                                                       node *L, *M, *R;
                                                                                       split(root, M, R, rv+1), split(M, L, M, lv);
ac7
        int get_idx(node* i) {
                                                                           117
6cf
            int ret = size(i->1);
                                                                           f1e
                                                                                       V = M;
482
            for (; i->p; i = i->p) {
                                                                           a28
                                                                                       join(L, R, root);
fbf
                node* pai = i->p:
8a6
                if (i != pai->1) ret += size(pai->1) + 1;
                                                                           e66
                                                                                       split(root, L, R, ru+1);
            }
e22
                                                                           367
                                                                                       join(L, V, L);
edf
                                                                           7e8
                                                                                       join(L, last[u] = new node(u, T() /* elemento neutro */),
            return ret;
479
                                                                              L);
        node* get_min(node* i) {
                                                                                       join(L, R, root);
048
                                                                           a28
433
            if (!i) return NULL;
                                                                           8d9
                                                                                   }
f8e
            return i->1 ? get_min(i->1) : i;
                                                                           4e6
                                                                                   void cut(int v) {
0de
        }
                                                                           892
                                                                                       auto [1, r] = get_range(v);
        node* get_max(node* i) {
f03
433
            if (!i) return NULL;
                                                                           df9
                                                                                       node *L, *M, *R;
424
            return i->r ? get_max(i->r) : i;
                                                                                       split(root, M, R, r+1), split(M, L, M, 1);
                                                                           dca
e92
        }
                                                                           de6
                                                                                       node *LL = get_max(L), *RR = get_min(R);
        // fim da treap
                                                                           710
                                                                                       if (LL and RR and LL->id == RR->id) { // remove duplicata
                                                                           e8b
                                                                                            if (last[RR->id] == RR) last[RR->id] = LL;
                                                                           992
                                                                                            node *A. *B:
4fb
        vector < node *> first, last;
                                                                                            split(R, A, B, 1);
                                                                           6b3
f82
        ETT(int n, vector<T> v = {}) : root(NULL), first(n), last(n) {
                                                                                            delete A;
                                                                           10c
            if (!v.size()) v = vector<T>(n);
                                                                           9d5
с5е
                                                                                            R = B:
603
            for (int i = 0; i < n; i++) {</pre>
                                                                           7c0
                first[i] = last[i] = new node(i, v[i], 1);
a00
                                                                           a 28
                                                                                       join(L, R, root);
469
                join(root, first[i], root);
                                                                                       join(root, M, root);
                                                                           a0d
```

```
6ff
        T query(int v) {
808
892
            auto [1, r] = get_range(v);
df9
            node *L, *M, *R;
            split(root, M, R, r+1), split(M, L, M, 1);
dca
d43
            T ans = M->sub;
69d
            join(L, M, M), join(M, R, root);
ba7
            return ans;
ede
        void update(int v, T val) { // soma val em todo mundo da
   subarvore
            auto [1, r] = get_range(v);
892
df9
            node *L. *M. *R:
dca
            split(root, M, R, r+1), split(M, L, M, 1);
409
            M->lazy += val;
            join(L, M, M), join(M, R, root);
69d
61c
        }
        void update_v(int v, T val) { // muda o valor de v pra val
129
            int l = get_idx(first[v]);
ac1
df9
            node *L, *M, *R;
d0c
            split(root, M, R, 1+1), split(M, L, M, 1);
            M \rightarrow val = M \rightarrow sub = val;
25 e
69d
            join(L, M, M), join(M, R, root);
630
        }
934
        bool is_in_subtree(int v, int u) { // se u ta na subtree de v
890
            auto [lv, rv] = get_range(v);
            auto [lu, ru] = get_range(u);
6ec
732
            return lv <= lu and ru <= rv:
        }
a21
355
        void print(node* i) {
            if (!i) return;
eae
            print(i->1);
a1e
            cout << i->id+1 << " ":
743
            print(i->r);
f15
59f
065
        void print() { print(root); cout << endl; }</pre>
045 };
2.15 Floyd-Warshall
// encontra o menor caminho entre todo
// par de vertices e detecta ciclo negativo
// returna 1 sse ha ciclo negativo
// d[i][i] deve ser 0
// para i != j, d[i][j] deve ser w se ha uma aresta
```

// (i, j) de peso w, INF caso contrario

2.16 Functional Graph

return 0;

// O(n^3)

1a8 int n;

830

f90

0ab

830

753

bb3

192 }

ae5 int d[MAX][MAX];

73c bool floyd_warshall() {

for (int k = 0; k < n; k++)

for (int i = 0; i < n; i++)</pre>

for (int j = 0; j < n; j++)

for (int i = 0; i < n; i++)

if (d[i][i] < 0) return 1;</pre>

```
2.16 Functional Graph
// rt[i] fala o ID da raiz associada ao vertice i
// d[i] fala a profundidade (0 sse ta no ciclo)
// pos[i] fala a posicao de i no array que eh a concat. dos ciclos
// build(f, val) recebe a funcao f e o custo de ir de
// i para f[i] (por default, val = f)
// f_k(i, k) fala onde i vai parar se seguir k arestas
// path(i, k) fala o custo (soma) seguir k arestas a partir de i
// Se quiser outra operacao, da pra alterar facil o codigo
// Codigo um pouco louco, tenho que admitir
//
// build - O(n)
// f_k - O(log(min(n, k)))
// path - O(\log(\min(n, k)))
6ef namespace func_graph {
1a8
        int n;
ce2
        int f[MAX], vis[MAX], d[MAX];
        int p[MAX], pp[MAX], rt[MAX], pos[MAX];
f82
ebd
        int sz[MAX], comp;
        vector < vector < int >> ciclo:
6a9
        11 val[MAX], jmp[MAX], seg[2*MAX];
97 c
        11 op(ll a, ll b) { return a+b; }; // mudar a operacao aqui
27b
        void dfs(int i, int t = 2) {
9 c 9
            vis[i] = t:
f09
            if (vis[f[i]] \ge 2) \{ // comeca ciclo - f[i] eh o rep.
e0a
                d[i] = 0, rt[i] = comp;
```

d[i][j] = min(d[i][j], d[i][k] + d[k][j]);

```
74c
                sz[comp] = t - vis[f[i]] + 1;
                p[i] = pp[i] = i, jmp[i] = val[i];
97b
15c
                ciclo.emplace_back();
                ciclo.back().push_back(i);
bfb
9d9
            } else {
                if (!vis[f[i]]) dfs(f[i], t+1);
c16
8c0
                rt[i] = rt[f[i]]:
                if (sz[comp]+1) { // to no ciclo
195
d0f
                    d[i] = 0;
97b
                    p[i] = pp[i] = i, jmp[i] = val[i];
bfb
                    ciclo.back().push_back(i);
9d9
                } else { // nao to no ciclo
00d
                    d[i] = d[f[i]]+1, p[i] = f[i]:
511
                    pp[i] = 2*d[pp[f[i]]] == d[pp[pp[f[i]]]]+d[f[i]] ?
   pp[pp[f[i]]] : f[i];
                    jmp[i] = pp[i] == f[i] ? val[i] : op(val[i],
114
   op(jmp[f[i]], jmp[pp[f[i]]]));
bb3
                }
d08
            if (f[ciclo[rt[i]][0]] == i) comp++; // fim do ciclo
e4a
29a
            vis[i] = 1;
0ba
        void build(vector<int> f_, vector<int> val_ = {}) {
1da
bcb
            n = f_size(), comp = 0;
527
            if (!val_.size()) val_ = f_;
830
            for (int i = 0: i < n: i++)</pre>
                f[i] = f_[i], val[i] = val_[i], vis[i] = 0, sz[i] = -1;
998
            ciclo.clear();
e74
158
            for (int i = 0; i < n; i++) if (!vis[i]) dfs(i);</pre>
6bb
            int t = 0:
daa
            for (auto& c : ciclo) {
336
                reverse(c.begin(), c.end());
ea5
                for (int i : c) {
                    pos[j] = t;
85b
948
                    seg[n+t] = val[j];
c82
                    t++:
25e
                }
cbc
            for (int i = n-1; i; i--) seg[i] = op(seg[2*i],
   seg[2*i+1]);
90b
       }
        int f_k(int i, ll k) {
283
1b1
            while (d[i] and k) {
77b
                int big = d[i] - d[pp[i]];
                if (big <= k) k -= big, i = pp[i];</pre>
ded
```

```
584
                 else k--, i = p[i];
09 c
            }
77e
            if (!k) return i;
            return ciclo[rt[i]][(pos[i] - pos[ciclo[rt[i]][0]] + k) %
    sz[rt[i]]]:
        }
f34
047
        ll path(int i, ll k) {
            auto query = [&](int 1, int r) {
3cf
3e4
                 11 q = 0;
                 for (1 += n, r += n; 1 <= r; ++1/=2, --r/=2) {
47a
27 e
                     if (1\%2 == 1) q = op(q, seg[1]);
1f2
                     if (r\%2 == 0) q = op(q, seg[r]);
598
                 }
bef
                 return q;
6e1
            };
b73
            11 \text{ ret} = 0;
1b1
             while (d[i] and k) {
77b
                 int big = d[i] - d[pp[i]];
327
                 if (big <= k) k -= big, ret = op(ret, jmp[i]), i =</pre>
   pp[i];
f9e
                 else k--, ret = op(ret, val[i]), i = p[i];
7e3
            if (!k) return ret;
e3c
a9e
            int first = pos[ciclo[rt[i]][0]], last =
    pos[ciclo[rt[i]].back()];
            // k/sz[rt[i]] voltas completas
430
            if (k/sz[rt[i]]) ret = op(ret, k/sz[rt[i]] * query(first,
   last)):
            k %= sz[rt[i]]:
9af
еЗс
            if (!k) return ret;
            int l = pos[i], r = first + (pos[i] - first + k - 1) %
8ea
    sz[rt[i]]:
            if (1 <= r) return op(ret, query(1, r));</pre>
982
687
            return op(ret, op(query(1, last), query(first, r)));
380
        }
51f }
2.17 HLD - aresta
// SegTree de soma
// query / update de soma das arestas
// Complexidades:
// build - O(n)
```

// query_path - 0(log^2 (n))

```
// update_path - O(log^2 (n))
// query_subtree - O(log(n))
// update_subtree - O(log(n))
// namespace seg { ... }
826 namespace hld {
        vector < pair < int , int > > g[MAX];
c0d
e65
        int pos[MAX], sz[MAX];
7c0
        int sobe[MAX], pai[MAX];
096
        int h[MAX], v[MAX], t;
0ce
        void build hld(int k, int p = -1, int f = 1) {
            v[pos[k] = t++] = sobe[k]; sz[k] = 1;
180
            for (auto& i : g[k]) if (i.first != p) {
418
dd2
                auto [u, w] = i;
a76
                sobe[u] = w; pai[u] = k;
                h[u] = (i == g[k][0] ? h[k] : u);
0 c 1
                build_hld(u, k, f); sz[k] += sz[u];
da7
                if (sz[u] > sz[g[k][0].first] or g[k][0].first == p)
865
9a3
                    swap(i, g[k][0]);
            }
804
667
            if (p*f == -1) build_hld(h[k] = k, -1, t = 0);
4dd
1f8
        void build(int root = 0) {
a34
            t = 0:
295
            build hld(root):
c83
            seg::build(t, v);
ea2
        11 query_path(int a, int b) {
3fc
2d5
            if (a == b) return 0;
            if (pos[a] < pos[b]) swap(a, b);</pre>
aa1
29b
            if (h[a] == h[b]) return seg::query(pos[b]+1, pos[a]);
fca
            return seg::query(pos[h[a]], pos[a]) +
   query_path(pai[h[a]], b);
87f
920
        void update_path(int a, int b, int x) {
            if (a == b) return;
d54
            if (pos[a] < pos[b]) swap(a, b);</pre>
aa1
            if (h[a] == h[b]) return (void)seg::update(pos[b]+1,
   pos[a], x);
            seg::update(pos[h[a]], pos[a], x); update_path(pai[h[a]],
701
   b, x);
dbf
```

```
d0a
        11 query_subtree(int a) {
b9f
            if (sz[a] == 1) return 0;
            return seg::query(pos[a]+1, pos[a]+sz[a]-1);
2f6
77f
        void update subtree(int a. int x) {
acc
            if (sz[a] == 1) return;
a5a
            seg::update(pos[a]+1, pos[a]+sz[a]-1, x);
9cd
a46
        }
7be
        int lca(int a, int b) {
            if (pos[a] < pos[b]) swap(a, b);</pre>
aa1
ca5
            return h[a] == h[b] ? b : lca(pai[h[a]], b);
219
        }
599 }
2.18 HLD - vertice
// SegTree de soma
// query / update de soma dos vertices
//
// Complexidades:
// build - O(n)
// query_path - 0(log^2 (n))
// update_path - O(log^2 (n))
// query_subtree - O(log(n))
// update_subtree - O(log(n))
// namespace seg { ... }
826 namespace hld {
        vector < int > g[MAX];
042
e65
        int pos[MAX], sz[MAX];
bd4
        int peso[MAX], pai[MAX];
096
        int h[MAX], v[MAX], t;
Осе
        void build_hld(int k, int p = -1, int f = 1) {
b18
            v[pos[k] = t++] = peso[k]; sz[k] = 1;
b94
            for (auto& i : g[k]) if (i != p) {
78d
                pai[i] = k;
                h[i] = (i == g[k][0] ? h[k] : i);
26e
193
                build_hld(i, k, f); sz[k] += sz[i];
cd1
                if (sz[i] > sz[g[k][0]] or g[k][0] == p) swap(i,
    g[k][0]);
d94
667
            if (p*f == -1) build_hld(h[k] = k, -1, t = 0);
a83
        }
1f8
        void build(int root = 0) {
```

```
a34
            t = 0:
            build hld(root):
295
c83
            seg::build(t, v);
ea2
        }
        11 query_path(int a, int b) {
3fc
            if (pos[a] < pos[b]) swap(a, b);</pre>
aa1
4bf
            if (h[a] == h[b]) return seg::query(pos[b], pos[a]);
            return seg::query(pos[h[a]], pos[a]) +
   query_path(pai[h[a]], b);
c17
920
        void update_path(int a, int b, int x) {
            if (pos[a] < pos[b]) swap(a, b);</pre>
aa1
            if (h[a] == h[b]) return (void)seg::update(pos[b], pos[a],
198
   x);
            seg::update(pos[h[a]], pos[a], x); update_path(pai[h[a]],
701
   b, x);
421
        11 query_subtree(int a) {
d0a
b3e
            return seg::query(pos[a], pos[a]+sz[a]-1);
ba2
acc
        void update_subtree(int a, int x) {
a22
            seg::update(pos[a], pos[a]+sz[a]-1, x);
480
7be
        int lca(int a. int b) {
            if (pos[a] < pos[b]) swap(a, b);</pre>
aa1
            return h[a] == h[b] ? b : lca(pai[h[a]], b);
ca5
219
de3 }
      HLD sem Update
// query de min do caminho
// Complexidades:
// build - O(n)
```

```
// query de min do caminho
//
// Complexidades:
// build - O(n)
// query_path - O(log(n))

826 namespace hld {
cod vector<pair<int, int> > g[MAX];
e65 int pos[MAX], sz[MAX];
7c0 int sobe[MAX], pai[MAX];
096 int h[MAX], v[MAX], t;
ea2 int men[MAX], seg[2*MAX];
Oce void build_hld(int k, int p = -1, int f = 1) {
```

```
180
             v[pos[k] = t++] = sobe[k]; sz[k] = 1;
418
             for (auto& i : g[k]) if (i.first != p) {
1f5
                 sobe[i.first] = i.second; pai[i.first] = k;
                 h[i.first] = (i == g[k][0] ? h[k] : i.first);
6fa
87b
                 men[i.first] = (i == g[k][0] ? min(men[k], i.second) :
    i.second);
4b2
                 build hld(i.first, k, f): sz[k] += sz[i.first]:
bc3
                 if (sz[i.first] > sz[g[k][0].first] or g[k][0].first
    == p)
9a3
                     swap(i, g[k][0]);
ea4
667
             if (p*f == -1) build hld(h\lceil k \rceil = k, -1, t = 0):
8ec
1f8
        void build(int root = 0) {
a34
            t = 0:
295
            build_hld(root);
             for (int i = 0; i < t; i++) seg[i+t] = v[i];</pre>
3ae
             for (int i = t-1; i; i--) seg[i] = min(seg[2*i],
8db
    seg[2*i+1]);
        }
ea5
f04
        int query_path(int a, int b) {
             if (a == b) return INF;
490
aa1
             if (pos[a] < pos[b]) swap(a, b);
98f
             if (h[a] != h[b]) return min(men[a], query_path(pai[h[a]],
    b));
46b
             int ans = INF, x = pos[b]+1+t, y = pos[a]+t;
             for (; x \le y; ++x/=2, --y/=2) ans = min({ans, seg[x],
    seg[v]});
ba7
             return ans:
3a9
        }
ee6 }:
2.20 Isomorfismo de arvores
// thash() retorna o hash da arvore (usando centroids como vertices
    especiais).
// Duas arvores sao isomorfas sse seu hash eh o mesmo
// O(|V|.log(|V|))
91f map < vector < int >, int > mphash;
df6 struct tree {
1a8
        int n;
789
        vector < vector < int >> g;
```

```
347
        vector < int > sz, cs;
        tree(int n_-): n(n_-), g(n_-), sz(n_-) {}
1b5
        void dfs_centroid(int v, int p) {
76b
588
            sz[v] = 1:
fa7
            bool cent = true:
            for (int u : g[v]) if (u != p) {
18e
365
                dfs_centroid(u, v), sz[v] += sz[u];
                if(sz[u] > n/2) cent = false;
e90
            }
ece
1f6
            if (cent and n - sz[v] <= n/2) cs.push_back(v);</pre>
368
784
        int fhash(int v, int p) {
544
            vector < int > h;
332
            for (int u : g[v]) if (u != p) h.push_back(fhash(u, v));
1c9
            sort(h.begin(), h.end());
            if (!mphash.count(h)) mphash[h] = mphash.size();
3ac
bbc
            return mphash[h];
748
        }
        11 thash() {
38f
            cs.clear():
23a
3a5
            dfs_centroid(0, -1);
16d
            if (cs.size() == 1) return fhash(cs[0], -1);
772
            11 h1 = fhash(cs[0], cs[1]), h2 = fhash(cs[1], cs[0]);
            return (min(h1, h2) << 30) + max(h1, h2);</pre>
fae
        }
138
4dd }:
```

2.21 Kosaraju

```
// O(n + m)
1a8 int n;
042 vector < int > g[MAX];
58d vector < int > gi[MAX]; // grafo invertido
c5a int vis[MAX]:
ee6 stack < int > S;
a52 int comp[MAX]; // componente conexo de cada vertice
1ca void dfs(int k) {
59a
        vis[k] = 1:
54f
        for (int i = 0; i < (int) g[k].size(); i++)</pre>
8d5
            if (!vis[g[k][i]]) dfs(g[k][i]);
58f
        S.push(k);
89c }
```

```
436 void scc(int k, int c) {
59a
        vis[k] = 1;
52c
        comp[k] = c;
        for (int i = 0; i < (int) gi[k].size(); i++)</pre>
ff0
            if (!vis[gi[k][i]]) scc(gi[k][i], c);
bf6
088 }
db8 void kosaraju() {
        for (int i = 0; i < n; i++) vis[i] = 0;</pre>
158
        for (int i = 0; i < n; i++) if (!vis[i]) dfs(i);
991
        for (int i = 0; i < n; i++) vis[i] = 0;
d32
        while (S.size()) {
70b
            int u = S.top();
7de
            S.pop();
f43
            if (!vis[u]) scc(u, u);
207
        }
e21 }
2.22 Kruskal
// Gera e retorna uma AGM e seu custo total a partir do vetor de
   arestas (edg)
// do grafo
//
// O(m log(m) + m a(m))
// 864875
1b9 vector<tuple<int, int, int>> edg; // {peso,[x,y]}
// DSU em O(a(n))
4a6 void dsu_build();
d78 int find(int a);
369 void unite(int a, int b);
c67 pair<ll, vector<tuple<int, int, int>>> kruskal(int n) {
8d2
        dsu_build(n);
        sort(edg.begin(), edg.end());
e31
        11 cost = 0:
854
979
        vector<tuple<int, int, int>> mst;
        for (auto [w,x,y] : edg) if (find(x) != find(y)) {
fea
9de
            mst.emplace_back(w, x, y);
45f
            cost += w:
05a
            unite(x,y);
```

ca2

}

```
5df
        return {cost, mst};
b6a }
2.23 Kuhn
// Computa matching maximo em grafo bipartido
// 'n' e 'm' sao quantos vertices tem em cada particao
// chamar add(i, j) para add aresta entre o cara i
// da particao A, e o cara j da particao B
// (entao i < n, j < m)
// Para recuperar o matching, basta olhar 'ma' e 'mb'
// 'recover' recupera o min vertex cover como um par de
// {caras da particao A, caras da particao B}
// O(|V| * |E|)
// Na pratica, parece rodar tao rapido quanto o Dinic
878 mt19937 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
6c6 struct kuhn {
14e
        int n. m:
        vector < vector < int >> g;
789
        vector < int > vis, ma, mb;
d3f
40e
        kuhn(int n_, int m_) : n(n_), m(m_), g(n),
8af
            vis(n+m), ma(n, -1), mb(m, -1) {}
ba6
        void add(int a, int b) { g[a].push_back(b); }
        bool dfs(int i) {
caf
29a
            vis[i] = 1:
            for (int j : g[i]) if (!vis[n+j]) {
29b
8c9
                vis[n+j] = 1;
2cf
                if (mb[j] == -1 or dfs(mb[j])) {
                    ma[i] = j, mb[j] = i;
bfe
8a6
                    return true;
                }
b17
82a
            }
d1f
            return false;
4ef
bf7
        int matching() {
1ae
            int ret = 0, aum = 1;
5a8
            for (auto& i : g) shuffle(i.begin(), i.end(), rng);
392
            while (aum) {
618
                for (int j = 0; j < m; j++) vis[n+j] = 0;
```

c5d

aum = 0;

```
830
                 for (int i = 0; i < n; i++)</pre>
01f
                     if (ma[i] == -1 and dfs(i)) ret++, aum = 1;
085
            }
edf
             return ret;
        }
2ee
b0d };
ebf pair < vector < int >, vector < int >> recover(kuhn& K) {
        K.matching();
e80
        int n = K.n. m = K.m;
50 c
9d0
        for (int i = 0; i < n+m; i++) K.vis[i] = 0;</pre>
        for (int i = 0; i < n; i++) if (K.ma[i] == -1) K.dfs(i);</pre>
bde
8ad
        vector<int> ca. cb:
576
        for (int i = 0; i < n; i++) if (!K.vis[i]) ca.push_back(i);</pre>
f24
        for (int i = 0; i < m; i++) if (K.vis[n+i]) cb.push_back(i);
        return {ca, cb};
aad
55f }
2.24 LCA com binary lifting
// Assume que um vertice eh ancestral dele mesmo, ou seja,
// se a eh ancestral de b, lca(a, b) = a
// MAX2 = ceil(log(MAX))
// Complexidades:
// build - O(n log(n))
// lca - O(log(n))
677 vector < vector < int > > g(MAX);
41c int n, p;
e75 int pai[MAX2][MAX];
999 int in[MAX], out[MAX];
1ca void dfs(int k) {
fdf
        in[k] = p++;
54f
        for (int i = 0; i < (int) g[k].size(); i++)</pre>
9b7
             if (in[g[k][i]] == -1) {
                 pai[0][g[k][i]] = k;
ba6
c38
                 dfs(g[k][i]);
e2d
26f
        out[k] = p++;
691 }
c11 void build(int raiz) {
        for (int i = 0; i < n; i++) pai[0][i] = i;</pre>
a 67
c63
        p = 0, memset(in, -1, sizeof in);
ecb
        dfs(raiz);
```

```
// pd dos pais
        for (int k = 1; k < MAX2; k++) for (int i = 0; i < n; i++)
511
d38
            pai[k][i] = pai[k - 1][pai[k - 1][i]];
530 }
OOf bool anc(int a. int b) { // se a eh ancestral de b
        return in[a] <= in[b] and out[a] >= out[b];
2d6 }
7be int lca(int a, int b) {
       if (anc(a, b)) return a;
e52
        if (anc(b, a)) return b:
        // sobe a
        for (int k = MAX2 - 1; k >= 0; k--)
f70
acf
            if (!anc(pai[k][a], b)) a = pai[k][a];
847
        return pai[0][a];
5c4 }
// Alternativamente:
// 'binary lifting' gastando O(n) de memoria
// Da pra add folhas e fazer queries online
// 3 vezes o tempo do binary lifting normal
// build - O(n)
// kth, lca, dist - O(log(n))
9c6 int d[MAX], p[MAX], pp[MAX];
d40 void set_root(int i) { p[i] = pp[i] = i, d[i] = 0; }
e9d void add leaf(int i. int u) {
        p[i] = u, d[i] = d[u]+1;
e0b
b15
        pp[i] = 2*d[pp[u]] == d[pp[pp[u]]]+d[u] ? pp[pp[u]] : u;
33f }
c37 int kth(int i, int k) {
4e3
        int dd = max(0, d[i]-k);
        while (d[i] > dd) i = d[pp[i]] >= dd ? pp[i] : p[i];
935
d9a
        return i;
f3c }
7be int lca(int a, int b) {
        if (d[a] < d[b]) swap(a, b);</pre>
a69
        while (d[a] > d[b]) a = d[pp[a]] >= d[b] ? pp[a] : p[a];
6cd
```

```
984
        while (a != b) {
932
            if (pp[a] != pp[b]) a = pp[a], b = pp[b];
e7c
            else a = p[a], b = p[b];
        }
4ea
3f5
        return a;
21d }
4fe int dist(int a, int b) { return d[a]+d[b]-2*d[lca(a,b)]; }
042 vector <int> g[MAX];
3ab void build(int i, int pai=-1) {
        if (pai == -1) set root(i):
15f
        for (int j : g[i]) if (j != pai) {
d31
            add_leaf(j, i);
b21
            build(j, i);
43b
        }
74a }
2.25 LCA com HLD
// Assume que um vertice eh ancestral dele mesmo, ou seja,
// se a eh ancestral de b, lca(a, b) = a
// Para buildar pasta chamar build(root)
// anc(a, b) responde se 'a' eh ancestral de 'b'
//
// Complexidades:
// build - O(n)
// lca - O(log(n))
// anc - O(1)
042 vector < int > g[MAX];
713 int pos[MAX], h[MAX], sz[MAX];
ff1 int pai[MAX], t;
8bf void build(int k, int p = -1, int f = 1) {
        pos[k] = t++; sz[k] = 1;
bce
e26
        for (int& i : g[k]) if (i != p) {
78d
            pai[i] = k;
            h[i] = (i == g[k][0] ? h[k] : i);
26e
            build(i, k, f); sz[k] += sz[i];
cb8
cd1
            if (sz[i] > sz[g[k][0]] or g[k][0] == p) swap(i, g[k][0]);
917
3da
        if (p*f == -1) t = 0, h[k] = k, build(k, -1, 0);
1b9 }
```

```
7be int lca(int a, int b) {
                                                                            a4e
                                                                                             int j = msb(y-x+1);
        if (pos[a] < pos[b]) swap(a, b);</pre>
                                                                            002
                                                                                             ans = op(ans, op(t[n/b*j+x], t[n/b*j+y-(1<<j)+1]));
        return h[a] == h[b] ? b : lca(pai[h[a]], b);
                                                                            4b6
ca5
219 }
                                                                            ba7
                                                                                        return ans;
                                                                                    }
                                                                            6bf
00f bool anc(int a, int b) {
                                                                            021 };
        return pos[a] \le pos[b] and pos[b] \le pos[a] + sz[a] - 1;
272 }
                                                                            065 namespace lca {
                                                                            042
                                                                                    vector < int > g[MAX];
                                                                            8ec
                                                                                    int v[2*MAX], pos[MAX], dep[2*MAX];
2.26 LCA com RMQ
                                                                            8bd
                                                                                    int t;
                                                                                    rmq<int> RMQ;
                                                                            2de
// Assume que um vertice eh ancestral dele mesmo, ou seja,
// se a eh ancestral de b, lca(a, b) = a
                                                                            4cf
                                                                                    void dfs(int i, int d = 0, int p = -1) {
// dist(a, b) retorna a distancia entre a e b
                                                                                        v[t] = i, pos[i] = t, dep[t++] = d;
                                                                            c97
                                                                                        for (int j : g[i]) if (j != p) {
                                                                            cac
// Complexidades:
                                                                            8ec
                                                                                             dfs(j, d+1, i);
// build - O(n)
                                                                                             v[t] = i, dep[t++] = d;
                                                                            cf2
// lca - 0(1)
                                                                            843
// dist - 0(1)
                                                                                    }
                                                                            d6a
                                                                            789
                                                                                    void build(int n, int root) {
1a5 template < typename T > struct rmq {
                                                                                        t = 0:
                                                                            a34
        vector <T> v:
517
                                                                            14e
                                                                                        dfs(root):
fcc
        int n; static const int b = 30;
                                                                            3f4
                                                                                         RMQ = rmq < int > (vector < int > (dep, dep + 2*n-1));
70e
        vector < int > mask, t;
                                                                            657
                                                                                    }
                                                                            7be
                                                                                    int lca(int a, int b) {
        int op(int x, int y) { return v[x] < v[y] ? x : y; }
18e
                                                                                        a = pos[a], b = pos[b];
                                                                            ab7
        int msb(int x) { return __builtin_clz(1)-__builtin_clz(x); }
ee1
                                                                            9c0
                                                                                        return v[RMQ.query(min(a, b), max(a, b))];
        rmq() {}
6ad
                                                                            5db
        rmq(const \ vector < T > \& \ v_) : v(v_), n(v.size()), mask(n), t(n) 
43c
                                                                            b5d
                                                                                    int dist(int a, int b) {
            for (int i = 0, at = 0; i < n; \max\{i++\} = at |= 1) {
2e5
                                                                            670
                                                                                        return dep[pos[a]] + dep[pos[b]] - 2*dep[pos[lca(a, b)]];
a61
                at = (at << 1) &((1 << b) -1);
                                                                            5b7
                                                                                    }
76a
                while (at and op(i, i-msb(at&-at)) == i) at ^= at&-at;
                                                                            645 }
53c
            }
243
            for (int i = 0; i < n/b; i++) t[i] =</pre>
                                                                            2.27 Line Tree
   b*i+b-1-msb(mask[b*i+b-1]);
            for (int j = 1; (1<<j) <= n/b; j++) for (int i = 0;
39d
   i+(1<<j) <= n/b; i++)
                                                                            // Reduz min-query em arvore para RMQ
                t[n/b*j+i] = op(t[n/b*(j-1)+i],
                                                                            // Se o grafo nao for uma arvore, as queries
   t[n/b*(j-1)+i+(1<<(j-1))]);
                                                                            // sao sobre a arvore geradora maxima
                                                                            // Queries de minimo
2d3
        int small(int r, int sz = b) { return
   r-msb(mask[r]&((1<<sz)-1)): }
                                                                            // build - O(n log(n))
b7a
        T query(int 1, int r) {
                                                                            // query - O(log(n))
27b
            if (r-l+1 <= b) return small(r, r-l+1);</pre>
7bf
            int ans = op(small(l+b-1), small(r));
                                                                            1a8 int n:
e80
            int x = 1/b+1, y = r/b-1;
e25
            if (x \le y) {
                                                                           | 3ae namespace linetree {
```

```
int id[MAX], seg[2*MAX], pos[MAX];
f37
        vector < int > v[MAX], val[MAX];
43f
430
        vector<pair<int, pair<int, int> > ar;
        void add(int a, int b, int p) { ar.push_back({p, {a, b}}); }
dc6
0a8
        void build() {
b09
            sort(ar.rbegin(), ar.rend());
            for (int i = 0; i < n; i++) id[i] = i, v[i] = {i},</pre>
   val[i].clear();
            for (auto i : ar) {
8bb
c91
                int a = id[i.second.first], b = id[i.second.second];
f6f
                if (a == b) continue;
c58
                if (v[a].size() < v[b].size()) swap(a, b):
                for (auto j : v[b]) id[j] = a, v[a].push_back(j);
fb8
482
                val[a].push_back(i.first);
78b
                for (auto j : val[b]) val[a].push_back(j);
e39
                v[b].clear(), val[b].clear();
            }
012
8e8
            vector<int> vv;
            for (int i = 0; i < n; i++) for (int j = 0; j <
2ce
   v[i].size(); j++) {
                pos[v[i][j]] = vv.size();
e52
                if (j + 1 < v[i].size()) vv.push_back(val[i][j]);</pre>
941
1cb
                else vv.push_back(0);
475
            }
bb4
            for (int i = n; i < 2*n; i++) seg[i] = vv[i-n];</pre>
            for (int i = n-1; i; i--) seg[i] = min(seg[2*i],
   seg[2*i+1]);
9fe
        int query(int a, int b) {
4ea
            if (id[a] != id[b]) return 0; // nao estao conectados
596
ab7
            a = pos[a], b = pos[b];
            if (a > b) swap(a, b);
d11
199
            b--:
            int ans = INF;
38a
513
            for (a += n, b += n; a \le b; ++a/=2, --b/=2) ans =
   min({ans, seg[a], seg[b]});
            return ans;
ba7
952
00f };
2.28 Link-cut Tree
// Link-cut tree padrao
//
// Todas as operacoes sao O(log(n)) amortizado
```

```
1ef namespace lct {
        struct node {
3c9
19f
            int p, ch[2];
062
            node() \{ p = ch[0] = ch[1] = -1; \}
f43
        }:
5f3
        node t[MAX];
971
        bool is_root(int x) {
657
            return t[x].p == -1 or (t[t[x].p].ch[0] != x and
   t[t[x].p].ch[1] != x);
cf1
ed6
        void rotate(int x) {
            int p = t[x].p, pp = t[p].p;
497
fc4
            if (!is_root(p)) t[pp].ch[t[pp].ch[1] == p] = x;
251
            bool d = t[p].ch[0] == x;
461
            t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
            if (t[p].ch[!d]+1) t[t[p].ch[!d]].p = p;
a76
8fa
            t[x].p = pp, t[p].p = x;
49b
       }
07c
        void splay(int x) {
18c
            while (!is_root(x)) {
497
                int p = t[x].p, pp = t[p].p;
0 c 5
                if (!is\_root(p)) rotate((t[pp].ch[0] == p)^(t[p].ch[0]
   == x) ? x : p);
64f
                rotate(x):
d8d
            }
4fa
        }
f16
        int access(int v) {
0eb
            int last = -1;
            for (int w = v; w+1; last = w, splay(v), w = t[v].p)
01a
024
                splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
            return last;
3d3
0a4
        }
        int find_root(int v) {
e89
5e3
            access(v);
3de
            while (t[v].ch[0]+1) v = t[v].ch[0];
f05
            return splay(v), v;
ee7
        void link(int v, int w) { // v deve ser raiz
142
5e3
            access(v):
            t[v].p = w;
10d
c56
        void cut(int v) { // remove aresta de v pro pai
4e6
5e3
            access(v);
264
            t[v].ch[0] = t[t[v].ch[0]].p = -1;
5f5
        }
```

```
bbb
        int lca(int v, int w) {
            return access(v), access(w);
948
b6d
e4e }
     Link-cut Tree - aresta
// Valores nas arestas
// rootify(v) torna v a raiz de sua arvore
// query(v, w) retorna a soma do caminho v--w
// update(v, w, x) soma x nas arestas do caminho v--w
//
// Todas as operacoes sao O(log(n)) amortizado
1ef namespace lct {
3c9
        struct node {
19f
            int p, ch[2];
810
            ll val, sub;
aa6
            bool rev:
04a
            int sz, ar;
4e4
           ll lazv:
            node() {}
f93
            node(int v, int ar_) :
7a8
546
            p(-1), val(v), sub(v), rev(0), sz(ar_{-}), ar(ar_{-}), lazy(0) {
b07
                ch[0] = ch[1] = -1;
53b
            }
6e0
        }:
        node t[2*MAX]; // MAXN + MAXQ
c53
```

```
99e
       map<pair<int, int>, int> aresta;
e4d
       int sz;
95a
       void prop(int x) {
dc1
            if (t[x].lazy) {
25e
                if (t[x].ar) t[x].val += t[x].lazy;
2ab
               t[x].sub += t[x].lazy*t[x].sz;
                if (t[x].ch[0]+1) t[t[x].ch[0]].lazy += t[x].lazy;
edc
942
                if (t[x].ch[1]+1) t[t[x].ch[1]].lazy += t[x].lazy;
           }
1ba
            if (t[x].rev) {
aa2
                swap(t[x].ch[0], t[x].ch[1]);
f95
379
                if (t[x].ch[0]+1) t[t[x].ch[0]].rev ^= 1;
                if (t[x].ch[1]+1) t[t[x].ch[1]].rev ^= 1;
c3d
50e
230
            t[x].lazy = 0, t[x].rev = 0;
f9d
564
        void update(int x) {
```

```
1a3
            t[x].sz = t[x].ar, t[x].sub = t[x].val;
            for (int i = 0; i < 2; i++) if (t[x].ch[i]+1) {
8ca
621
                prop(t[x].ch[i]);
c4f
                t[x].sz += t[t[x].ch[i]].sz;
269
                t[x].sub += t[t[x].ch[i]].sub:
400
       }
28b
971
        bool is_root(int x) {
657
            return t[x].p == -1 or (t[t[x].p].ch[0] != x and
   t[t[x].p].ch[1] != x);
cf1
       }
ed6
        void rotate(int x) {
497
            int p = t[x].p. pp = t[p].p:
            if (!is_root(p)) t[pp].ch[t[pp].ch[1] == p] = x;
fc4
            bool d = t[p].ch[0] == x;
251
461
            t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
a76
            if (t[p].ch[!d]+1) t[t[p].ch[!d]].p = p;
            t[x].p = pp, t[p].p = x;
8fa
444
            update(p), update(x);
f31
       }
238
        int splay(int x) {
18c
            while (!is_root(x)) {
497
                int p = t[x].p, pp = t[p].p;
77b
                if (!is_root(p)) prop(pp);
be5
                prop(p), prop(x);
0c5
                if (!is\_root(p)) rotate((t[pp].ch[0] == p)^(t[p].ch[0]
   == x) ? x : p);
64f
                rotate(x):
72c
            return prop(x), x;
aab
       }
08f
f16
        int access(int v) {
0eb
            int last = -1:
d9f
            for (int w = v; w+1; update(last = w), splay(v), w =
   t[v].p)
                splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
024
3d3
            return last;
294
9f1
        void make_tree(int v, int w=0, int ar=0) { t[v] = node(w, ar);
}
        int find_root(int v) {
e89
            access(v), prop(v);
13f
9f0
            while (t[v].ch[0]+1) v = t[v].ch[0], prop(v);
637
            return splay(v);
16a
82f
        bool conn(int v, int w) {
2cf
            access(v). access(w):
```

```
b9b
            return v == w ? true : t[v].p != -1;
ec0
277
        void rootify(int v) {
5e3
            access(v);
a02
            t[v].rev ^= 1;
a05
971
        11 query(int v, int w) {
            rootify(w), access(v);
b54
249
            return t[v].sub;
652
3fa
        void update(int v, int w, int x) {
b54
            rootify(w), access(v);
12c
            t[v].lazv += x:
74f
        void link_(int v, int w) {
204
821
            rootify(w);
389
            t[w].p = v;
523
6b8
        void link(int v, int w, int x) { // v--w com peso x
379
            int id = MAX + sz++;
            aresta[make_pair(v, w)] = id;
110
a88
            make_tree(id, x, 1);
            link_(v, id), link_(id, w);
c88
58c
        }
e63
        void cut_(int v, int w) {
b54
            rootifv(w). access(v):
            t[v].ch[0] = t[t[v].ch[0]].p = -1;
264
7cd
031
        void cut(int v, int w) {
b0f
            int id = aresta[make_pair(v, w)];
            cut_(v, id), cut_(id, w);
a4a
840
        int lca(int v, int w) {
bbb
5e3
            access(v):
a8b
            return access(w);
524
       }
9ce }
     Link-cut Tree - vertice
// Valores nos vertices
// make tree(v, w) cria uma nova arvore com um
// vertice soh com valor 'w'
// rootify(v) torna v a raiz de sua arvore
// query(v, w) retorna a soma do caminho v--w
// update(v, w, x) soma x nos vertices do caminho v--w
```

//

```
// Todas as operacoes sao O(log(n)) amortizado
1ef namespace lct {
3c9
        struct node {
19f
            int p, ch[2];
810
            ll val, sub;
aa6
            bool rev:
            int sz;
e4d
4e4
            ll lazv;
f93
            node() {}
aa0
            node(int v) : p(-1), val(v), sub(v), rev(0), sz(1),
   lazv(0) {
b07
                ch[0] = ch[1] = -1:
с4е
2b7
        };
5f3
        node t[MAX];
95a
        void prop(int x) {
            if (t[x].lazy) {
dc1
9f7
                t[x].val += t[x].lazy, t[x].sub += t[x].lazy*t[x].sz;
                if (t[x].ch[0]+1) t[t[x].ch[0]].lazy += t[x].lazy;
edc
942
                if (t[x].ch[1]+1) t[t[x].ch[1]].lazy += t[x].lazy;
e26
            }
aa2
            if (t[x].rev) {
f95
                swap(t[x].ch[0], t[x].ch[1]);
379
                if (t[x].ch[0]+1) t[t[x].ch[0]].rev ^= 1;
c3d
                if (t[x].ch[1]+1) t[t[x].ch[1]].rev ^= 1;
50e
230
            t[x].lazy = 0, t[x].rev = 0;
        }
c62
564
        void update(int x) {
            t[x].sz = 1, t[x].sub = t[x].val;
ec2
            for (int i = 0; i < 2; i++) if (t[x].ch[i]+1) {
8ca
                prop(t[x].ch[i]);
621
                t[x].sz += t[t[x].ch[i]].sz;
c4f
269
                t[x].sub += t[t[x].ch[i]].sub;
400
            }
da7
        }
971
        bool is_root(int x) {
            return t[x].p == -1 or (t[t[x].p].ch[0] != x and
   t[t[x].p].ch[1] != x);
cf1
       }
        void rotate(int x) {
ed6
            int p = t[x].p, pp = t[p].p;
497
fc4
            if (!is_root(p)) t[pp].ch[t[pp].ch[1] == p] = x;
251
            bool d = t[p].ch[0] == x;
```

```
461
            t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
            if (t[p].ch[!d]+1) t[t[p].ch[!d]].p = p;
a76
8fa
            t[x].p = pp, t[p].p = x;
            update(p), update(x);
444
f31
       }
238
        int splay(int x) {
18c
            while (!is root(x)) {
497
                int p = t[x].p, pp = t[p].p;
77b
                if (!is_root(p)) prop(pp);
be5
                prop(p), prop(x);
0c5
                if (!is_root(p)) rotate((t[pp].ch[0] == p)^(t[p].ch[0]
   == x) ? x : p);
64f
                rotate(x):
72c
aab
            return prop(x), x;
       }
08f
f16
        int access(int v) {
            int last = -1:
0eb
            for (int w = v; w+1; update(last = w), splay(v), w =
d9f
   t[v].p)
                splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
024
3d3
            return last;
294
f17
        void make_tree(int v, int w) { t[v] = node(w); }
e89
        int find_root(int v) {
13f
            access(v), prop(v);
            while (t[v].ch[0]+1) v = t[v].ch[0], prop(v);
9f0
637
            return splay(v);
       }
16a
f94
        bool connected(int v, int w) {
2cf
            access(v), access(w);
b9b
            return v == w ? true : t[v].p != -1;
ec6
277
        void rootifv(int v) {
            access(v):
5e3
a02
            t[v].rev ^= 1;
       }
a05
971
        11 query(int v, int w) {
b54
            rootify(w), access(v);
249
            return t[v].sub;
652
       }
        void update(int v, int w, int x) {
3fa
b54
            rootify(w), access(v);
12c
            t[v].lazy += x;
74f
142
        void link(int v, int w) {
821
            rootify(w);
```

```
389
            t[w].p = v;
8a8
031
        void cut(int v, int w) {
b54
            rootify(w), access(v);
264
            t[v].ch[0] = t[t[v].ch[0]].p = -1;
d9a
bbb
        int lca(int v, int w) {
5e3
            access(v);
a8b
            return access(w);
524
        }
f9f }
```

2.31 Max flow com lower bound

```
// add(a, b, l, r):
// adiciona aresta de a pra b, onde precisa passar f de fluxo, l <= f
   <= r
// add(a, b, c):
// adiciona aresta de a pra b com capacidade c
//
// Mesma complexidade do Dinitz
cd5 struct lb_max_flow : dinitz {
        vector < int > d:
5ce
d8c
        lb_max_flow(int n) : dinitz(n + 2), d(n, 0) {}
        void add(int a, int b, int 1, int r) {
b12
c97
            d[a] -= 1:
            d[b] += 1;
f1b
4c0
            dinitz::add(a, b, r - 1);
ed4
        }
087
        void add(int a, int b, int c) {
0f3
            dinitz::add(a, b, c);
039
        }
7a1
        bool has_circulation() {
50c
            int n = d.size();
854
            11 cost = 0:
603
            for (int i = 0; i < n; i++) {</pre>
c69
                if (d[i] > 0) {
f56
                     cost += d[i]:
57a
                     dinitz::add(n, i, d[i]);
9c7
                } else if (d[i] < 0) {</pre>
                     dinitz::add(i, n+1, -d[i]);
b76
dc7
                }
676
            }
067
            return (dinitz::max_flow(n, n+1) == cost);
```

```
110
        bool has_flow(int src, int snk) {
7bd
387
            dinitz::add(snk, src, INF);
            return has_circulation();
e40
cc1
        11 max_flow(int src, int snk) {
4eb
ee8
            if (!has flow(src. snk)) return -1;
            dinitz::F = 0;
4ad
fe5
            return dinitz::max_flow(src, snk);
619
e8f };
```

2.32 MinCostMaxFlow

```
// min_cost_flow(s, t, f) computa o par (fluxo, custo)
// com max(fluxo) <= f que tenha min(custo)</pre>
// min_cost_flow(s, t) -> Fluxo maximo de custo minimo de s pra t
// Se for um dag, da pra substituir o SPFA por uma DP pra nao
// pagar O(nm) no comeco
// Se nao tiver aresta com custo negativo, nao precisa do SPFA
//
// O(nm + f * m log n)
123 template < typename T > struct mcmf {
670
        struct edge {
            int to, rev, flow, cap; // para, id da reversa, fluxo,
b75
   capacidade
            bool res; // se eh reversa
7f9
            T cost; // custo da unidade de fluxo
635
            edge(): to(0), rev(0), flow(0), cap(0), cost(0),
   res(false) {}
1d7
            edge(int to_, int rev_, int flow_, int cap_, T cost_, bool
   res_)
f8d
                : to(to_), rev(rev_), flow(flow_), cap(cap_),
   res(res_), cost(cost_) {}
723
       };
        vector < vector < edge >> g;
002
168
        vector < int > par_idx, par;
f1e
        T inf:
a03
        vector<T> dist;
        mcmf(int n) : g(n), par_idx(n), par(n),
   inf(numeric_limits <T>::max()/3) {}
91c
        void add(int u, int v, int w, T cost) { // de u pra v com cap
   w e custo cost
```

```
2fc
            edge a = edge(v, g[v].size(), 0, w, cost, false);
234
            edge b = edge(u, g[u].size(), 0, 0, -cost, true);
b24
            g[u].push_back(a);
            g[v].push_back(b);
c12
0ed
        }
        vector<T> spfa(int s) { // nao precisa se nao tiver custo
8bc
   negativo
871
            deque < int > q;
3d1
            vector < bool > is_inside(g.size(), 0);
577
            dist = vector<T>(g.size(), inf);
a93
            dist[s] = 0;
a30
            q.push_back(s);
            is_inside[s] = true;
ecb
14d
            while (!q.empty()) {
b1e
                 int v = q.front();
ced
                 q.pop_front();
48d
                 is_inside[v] = false;
76e
                 for (int i = 0; i < g[v].size(); i++) {</pre>
9d4
                     auto [to, rev, flow, cap, res, cost] = g[v][i];
e61
                     if (flow < cap and dist[v] + cost < dist[to]) {</pre>
943
                         dist[to] = dist[v] + cost:
ed6
                         if (is_inside[to]) continue;
020
                         if (!q.empty() and dist[to] > dist[q.front()])
   q.push_back(to);
b33
                         else q.push_front(to);
b52
                         is_inside[to] = true;
2d1
                     }
8cd
                }
f2c
8d7
            return dist;
96c
        }
2a2
        bool dijkstra(int s, int t, vector<T>& pot) {
489
            priority_queue < pair < T, int > , vector < pair < T, int >> ,
   greater<>> q;
577
            dist = vector<T>(g.size(), inf);
            dist[s] = 0;
a93
            q.emplace(0, s);
115
402
            while (q.size()) {
91b
                 auto [d, v] = q.top();
833
                q.pop();
68b
                 if (dist[v] < d) continue;</pre>
```

```
76e
                for (int i = 0; i < g[v].size(); i++) {</pre>
                     auto [to, rev, flow, cap, res, cost] = g[v][i];
9d4
                     cost += pot[v] - pot[to];
e8c
                     if (flow < cap and dist[v] + cost < dist[to]) {</pre>
e61
                         dist[to] = dist[v] + cost;
943
                         q.emplace(dist[to], to);
441
88b
                         par_idx[to] = i, par[to] = v;
873
                     }
de3
                }
            }
9d4
1d4
            return dist[t] < inf;</pre>
c68
        }
3d2
        pair < int , T> min_cost_flow(int s, int t, int flow = INF) {
3dd
            vector <T> pot(g.size(), 0);
            pot = spfa(s); // mudar algoritmo de caminho minimo aqui
9e4
d22
            int f = 0:
ce8
            T ret = 0;
4a0
            while (f < flow and dijkstra(s, t, pot)) {</pre>
bda
                for (int i = 0; i < g.size(); i++)</pre>
                     if (dist[i] < inf) pot[i] += dist[i];</pre>
d2a
71b
                int mn_flow = flow - f, u = t;
045
                while (u != s){
90f
                     mn_flow = min(mn_flow,
                         g[par[u]][par_idx[u]].cap -
   g[par[u]][par_idx[u]].flow);
3d1
                     u = par[u];
935
                }
                ret += pot[t] * mn_flow;
1f2
476
                u = t:
045
                while (u != s) {
e09
                     g[par[u]][par_idx[u]].flow += mn_flow;
                     g[u][g[par[u]][par_idx[u]].rev].flow -= mn_flow;
d98
3d1
                     u = par[u];
                }
bcc
04d
                f += mn flow:
36d
15b
            return make_pair(f, ret);
ссЗ
        }
        // Opcional: retorna as arestas originais por onde passa flow
```

```
= cap
182
        vector<pair<int,int>> recover() {
24a
             vector < pair < int , int >> used;
             for (int i = 0; i < g.size(); i++) for (edge e : g[i])</pre>
2a4
587
                 if(e.flow == e.cap && !e.res) used.push_back({i,
    e.to}):
             return used;
f6b
390
697 };
2.33 Prufer code
// Traduz de lista de arestas para prufer code
// e vice-versa
// Os vertices tem label de O a n-1
// Todo array com n-2 posicoes e valores de
// O a n-1 sao prufer codes validos
//
// O(n)
47d vector < int > to_prufer (vector < pair < int , int >> tree) {
1fa
        int n = tree.size()+1;
2cf
        vector < int > d(n, 0);
        vector < vector < int >> g(n);
4aa
f87
        for (auto [a, b] : tree) d[a]++, d[b]++,
             g[a].push_back(b), g[b].push_back(a);
f60
c5a
        vector < int > pai(n, -1);
        queue < int > q; q.push(n-1);
260
402
        while (q.size()) {
be1
             int u = q.front(); q.pop();
34c
             for (int v : g[u]) if (v != pai[u])
9c9
                 pai[v] = u, q.push(v);
70d
        }
399
        int idx, x;
897
        idx = x = find(d.begin(), d.end(), 1) - d.begin();
        vector<int> ret;
4b8
b28
        for (int i = 0; i < n-2; i++) {
d4b
             int y = pai[x];
e81
             ret.push_back(v);
666
             if (-d[y] == 1 \text{ and } y < idx) x = y;
             else idx = x = find(d.begin()+idx+1, d.end(), 1) -
    d.begin();
5f9
edf
        return ret;
d3b }
4d8 vector<pair<int, int>> from_prufer(vector<int> p) {
```

```
455
        int n = p.size()+2;
126
        vector < int > d(n, 1);
650
        for (int i : p) d[i]++;
85b
        p.push_back(n-1);
399
        int idx, x:
897
        idx = x = find(d.begin(), d.end(), 1) - d.begin();
1df
        vector < pair < int , int >> ret;
b06
        for (int y : p) {
dab
            ret.push_back({x, y});
666
            if (--d[y] == 1 \text{ and } y < idx) x = y;
367
            else idx = x = find(d.begin()+idx+1, d.end(), 1) -
   d.begin();
c3b
edf
        return ret;
765 }
2.34 Sack (DSU em arvores)
// Responde queries de todas as sub-arvores
// offline
//
// O(n log(n))
6bf int sz[MAX], cor[MAX], cnt[MAX];
042 vector < int > g[MAX];
6df void build(int k, int d=0) {
        sz[k] = 1;
e8f
01a
        for (auto& i : g[k]) {
            build(i, d+1); sz[k] += sz[i];
30f
925
            if (sz[i] > sz[g[k][0]]) swap(i, g[k][0]);
011
        }
189 }
74f void compute(int k, int x, bool dont=1) {
de9
        cnt[cor[k]] += x;
        for (int i = dont; i < g[k].size(); i++)</pre>
828
b5c
            compute(g[k][i], x, 0);
896 }
dc4 void solve(int k, bool keep=0) {
32a
        for (int i = int(g[k].size())-1; i >= 0; i--)
b4c
            solve(g[k][i], !i);
4a0
        compute(k, 1);
        // agora cnt[i] tem quantas vezes a cor
        // i aparece na sub-arvore do k
```

```
830
         if (!keep) compute(k, -1, 0);
8bc }
2.35 Stable Marriage
// Emparelha todos os elementos de A com elementos de B
// de forma que nao exista um par x \in A, y \in B
// e x nao pareado com y tal que x prefira parear com y
// e y prefira parear com x.
// a[i] contem os elementos de B ordenados por preferencia de i
// b[j] contem os elementos de A ordenados por preferencia de j
// |A| <= |B|
//
// Retorna um vetor v de tamanho |A| onde v[i] guarda o match de i.
// O(|A| * |B|)
380 vector<int> stable_marriage(vector<vector<int>> &a,
    vector < vector < int >> &b) {
         int n = a.size(), m = b.size();
652
         assert(a[0].size() == m and b[0].size() == n and n <= m);
83e
        vector < int > match(m, -1), it(n, 0);
017
e6f
         vector inv_b(m, vector < int > (n));
         for (int i = 0; i < m; i++) for (int j = 0; j < n; j++)
a34
9f2
             inv_b[i][b[i][j]] = j;
26a
         queue < int > q;
         for (int i = 0; i < n; i++) q.push(i);</pre>
5af
402
         while (q.size()) {
379
             int i = q.front(); q.pop();
4b8
             int j = a[i][it[i]];
57c
             if (match[j] == -1) match[j] = i;
02d
             else if (inv_b[j][i] < inv_b[j][match[j]]) {</pre>
5d1
                 q.emplace(match[j]);
e7d
                 it[match[j]]++;
f1d
                 match[i] = i;
             } else q.emplace(i), it[i]++;
bc4
258
        }
825
        vector<int> ret(n);
d72
         for (int i = 0; i < m; i++) if (match[i] != -1) ret[match[i]]</pre>
   = i:
edf
        return ret;
Off }
```

2.36 Tarjan para SCC

```
// O(n + m)
042 vector < int > g[MAX];
4ce stack<int> s;
a42 int vis[MAX], comp[MAX];
3fd int id[MAX];
// se quiser comprimir ciclo ou achar ponte em grafo nao direcionado,
// colocar um if na dfs para nao voltar pro pai da DFS tree
f32 int dfs(int i, int& t) {
        int lo = id[i] = t++:
cf0
18e
       s.push(i);
       vis[i] = 2;
0c2
        for (int j : g[i]) {
48e
740
            if (!vis[j]) lo = min(lo, dfs(j, t));
994
            else if (vis[j] == 2) lo = min(lo, id[j]);
d64
        // aresta de i pro pai eh uma ponte (no caso nao direcionado)
        if (lo == id[i]) while (1) {
3de
            int u = s.top(); s.pop();
3c3
            vis[u] = 1, comp[u] = i;
9c5
2ef
            if (u == i) break;
266
       }
253
        return lo;
38a }
f93 void tarjan(int n) {
6bb
        int t = 0;
991
        for (int i = 0; i < n; i++) vis[i] = 0;</pre>
3be
        for (int i = 0; i < n; i++) if (!vis[i]) dfs(i, t);</pre>
ea1 }
     Topological Sort
// Retorna uma ordenacaoo topologica de g
// Se g nao for DAG retorna um vetor vazio
//
// O(n + m)
042 vector < int > g[MAX];
```

```
b6a vector<int> topo_sort(int n) {
        vector < int > ret(n,-1), vis(n,0);
f51
        int pos = n-1, dag = 1;
36d
        function < void(int) > dfs = [&](int v) {
             vis[v] = 1;
cca
440
             for (auto u : g[v]) {
152
                 if (vis[u] == 1) dag = 0;
532
                 else if (!vis[u]) dfs(u);
e37
d44
             ret[pos--] = v, vis[v] = 2;
57e
        };
158
        for (int i = 0; i < n; i++) if (!vis[i]) dfs(i);</pre>
d8f
        if (!dag) ret.clear();
        return ret;
edf
d6b }
2.38 Vertex cover
// Encontra o tamanho do vertex cover minimo
// Da pra alterar facil pra achar os vertices
// Parece rodar com < 2 s pra N = 90
//
// O(n * 1.38^n)
76a namespace cover {
5a4
        const int MAX = 96;
042
        vector < int > g[MAX];
823
        bitset < MAX > bs [MAX];
1a8
        int n;
697
        void add(int i, int j) {
bd0
            if (i == j) return;
            n = \max(\{n, i+1, j+1\});
78c
200
             bs[i][j] = bs[j][i] = 1;
203
        }
6c0
        int rec(bitset < MAX > m) {
             int ans = 0;
1a4
25b
             for (int x = 0; x < n; x++) if (m[x]) {
002
                 bitset < MAX > comp;
4bf
                 function < void(int) > dfs = [&](int i) {
b96
                     comp[i] = 1, m[i] = 0;
0c3
                     for (int j : g[i]) if (m[j]) dfs(j);
815
                 };
```

```
963
                dfs(x);
d34
                int ma, deg = -1, cyc = 1;
                for (int i = 0; i < n; i++) if (comp[i]) {</pre>
417
                     int d = (bs[i]&comp).count();
d0b
                     if (d <= 1) cvc = 0;
18a
c1f
                    if (d > deg) deg = d, ma = i;
                }
d8e
269
                if (deg <= 2) { // caminho ou ciclo</pre>
                     ans += (comp.count() + cyc) / 2;
340
5e2
                     continue;
702
                }
3f9
                comp[ma] = 0;
                // ou ta no cover, ou nao ta no cover
                ans += min(1 + rec(comp), deg + rec(comp & ~bs[ma]));
1dd
            }
6e6
ba7
            return ans;
2ec
f5c
        int solve() {
3c5
            bitset < MAX > m:
603
            for (int i = 0; i < n; i++) {</pre>
939
                m[i] = 1;
                for (int j = 0; j < n; j++)
f90
741
                     if (bs[i][j]) g[i].push_back(j);
13e
            }
4f9
            return rec(m);
708
        }
9c5 }
```

2.39 Virtual Tree

```
// Comprime uma arvore dado um conjunto S de vertices, de forma que
// o conjunto de vertices da arvore comprimida contenha S e seja
// minimal e fechado sobre a operacao de LCA
// Se |S| = k, a arvore comprimida tem menos que 2k vertices
// As arestas de virt possuem a distancia do vertice ate o vizinho
// Retorna a raiz da virtual tree
//
// lca::pos deve ser a ordem de visitacao no dfs
// voce pode usar o LCAcomHLD, por exemplo
//
// O(k log(k))
b36 vector<pair<int, int>> virt[MAX];
d41 #warning lembrar de buildar o LCA antes
```

```
c14 int build_virt(vector<int> v) {
        auto cmp = [&](int i, int j) { return lca::pos[i] <</pre>
   lca::pos[j]; };
074
        sort(v.begin(), v.end(), cmp);
        for (int i = v.size()-1; i; i--) v.push_back(lca::lca(v[i],
e85
   v[i-1]));
074
        sort(v.begin(), v.end(), cmp);
        v.erase(unique(v.begin(), v.end()), v.end());
d76
        for (int i = 0; i < v.size(); i++) virt[v[i]].clear();</pre>
37 c
        for (int i = 1; i < v.size(); i++) virt[lca::lca(v[i-1],</pre>
   v[i])].clear();
        for (int i = 1; i < v.size(); i++) {</pre>
ad7
51b
            int parent = lca::lca(v[i-1], v[i]);
290
            int d = lca::dist(parent, v[i]);
d41 #warning soh to colocando aresta descendo
            virt[parent].emplace_back(v[i], d);
4d0
fe5
        }
832
        return v[0];
142 }
```

3 Problemas

3.1 Angle Range Intersection

```
// Computa intersecao de angulos
// Os angulos (arcos) precisam ter comprimeiro < pi
// (caso contrario a intersecao eh estranha)
//
// Tudo 0(1)
32a struct angle_range {
75e
        static constexpr ld ALL = 1e9, NIL = -1e9;
395
        ld 1, r;
        angle_range() : 1(ALL), r(ALL) {}
c77
        angle_range(1d 1_, 1d r_) : 1(1_), r(r_) { fix(1), fix(r); }
894
        void fix(ld& theta) {
4ee
da7
             if (theta == ALL or theta == NIL) return;
323
             if (theta > 2*pi) theta -= 2*pi;
868
            if (theta < 0) theta += 2*pi;</pre>
625
        }
2ee
        bool empty() { return 1 == NIL; }
931
        bool contains(ld q) {
40f
            fix(q);
4d7
            if (1 == ALL) return true;
fec
            if (1 == NIL) return false;
```

```
if (1 < r) return 1 < q and q < r;
6a6
075
            return q > 1 or q < r;</pre>
800
9c7
        friend angle_range operator &(angle_range p, angle_range q) {
743
            if (p.l == ALL or q.l == NIL) return q;
            if (q.1 == ALL or p.1 == NIL) return p;
20f
7d5
            if (p.1 > p.r \text{ and } q.1 > q.r) \text{ return } \{\max(p.1, q.1),
   min(p.r, q.r)};
            if (q.1 > q.r) swap(p.1, q.1), swap(p.r, q.r);
aa6
            if (p.1 > p.r) {
868
249
                if (q.r > p.1) return \{max(q.1, p.1), q.r\};
6f7
                 else if (q.1 < p.r) return {q.1, min(q.r, p.r)};</pre>
270
                return {NIL, NIL};
337
            }
5a8
            if (max(p.1, q.1) > min(p.r, q.r)) return {NIL, NIL};
bcb
            return {max(p.1, q.1), min(p.r, q.r)};
142
        }
5e1 };
```

3.2 Area da Uniao de Retangulos

```
// O(n log(n))
// 5d8d2f
aa4 namespace seg {
6b3
        pair < int , 11 > seg[4*MAX];
b1b
        ll lazy[4*MAX], *v;
1a8
        int n;
e01
        pair < int , ll > merge(pair < int , ll > 1, pair < int , ll > r) {
719
            if (1.second == r.second) return {1.first+r.first,
   1.second:
            else if (1.second < r.second) return 1;</pre>
53b
aa0
            else return r;
d82
        }
6fc
        pair < int, ll > build(int p=1, int l=0, int r=n-1) {
3c7
            lazv[p] = 0;
bf8
            if (1 == r) return seg[p] = {1, v[1]};
            int m = (1+r)/2;
ee4
            return seg[p] = merge(build(2*p, 1, m), build(2*p+1, m+1,
432
   r));
f94
d9e
        void build(int n2, 11* v2) {
680
            n = n2, v = v2:
6f2
            build();
f8a
        }
```

```
ceb
        void prop(int p, int 1, int r) {
208
            seg[p].second += lazy[p];
2c9
            if (1 != r) lazy[2*p] += lazy[p], lazy[2*p+1] += lazy[p];
3c7
            lazv[p] = 0;
bf2
693
        pair < int, ll > query (int a, int b, int p=1, int l=0, int r=n-1)
{
6b9
            prop(p, 1, r);
527
            if (a <= 1 and r <= b) return seg[p];</pre>
            if (b < 1 or r < a) return {0, LINF};</pre>
9b7
ee4
            int m = (1+r)/2;
            return merge (query (a, b, 2*p, 1, m), query (a, b, 2*p+1,
eeb
   m+1. r)):
786
07c
        pair < int, 11 > update(int a, int b, int x, int p=1, int 1=0,
   int r=n-1) {
6b9
            prop(p, 1, r);
            if (a \le 1 \text{ and } r \le b) \{
9a3
                lazy[p] += x;
b94
6b9
                prop(p, 1, r);
534
                return seg[p];
821
e9f
            if (b < 1 or r < a) return seg[p];</pre>
ee4
            int m = (1+r)/2;
086
            return seg[p] = merge(update(a, b, x, 2*p, 1, m),
579
                     update(a, b, x, 2*p+1, m+1, r));
c65
        }
043 }:
eb5 ll seg_vec[MAX];
8be 11 area_sq(vector<pair<int, int>, pair<int, int>>> &sq){
28 c
        vector<pair<int, int>, pair<int, int>>> up;
60a
        for (auto it : sq){
619
            int x1, y1, x2, y2;
ae0
            tie(x1, y1) = it.first;
68e
            tie(x2, y2) = it.second;
80f
            up.push_back({{x1+1, 1}, {v1, v2}});
            up.push_back({{x2+1, -1}, {y1, y2}});
aee
6c3
        }
092
        sort(up.begin(), up.end());
049
        memset(seg_vec, 0, sizeof seg_vec);
6fe
        11 H_MAX = MAX;
156
        seg::build(H_MAX-1, seg_vec);
7ba
        auto it = up.begin();
04b
        11 \text{ ans} = 0;
f14
        while (it != up.end()){
```

```
07f
            11 L = (*it).first.first;
            while (it != up.end() && (*it).first.first == L){
718
127
                int x, inc, y1, y2;
d35
                tie(x, inc) = it->first;
d3d
                tie(y1, y2) = it -> second;
5d1
                seg::update(y1+1, y2, inc);
40d
               it++:
9b1
            }
852
            if (it == up.end()) break;
            11 R = (*it).first.first;
d8a
f59
            11 W = R-L:
efd
            auto jt = seg::query(0, H_MAX-1);
91a
           11 H = H_MAX - 1;
            if (jt.second == 0) H -= jt.first;
e8a
8df
            ans += W*H;
5c8
       }
ba7
        return ans;
385 }
```

3.3 Area Maxima de Histograma

```
// Assume que todas as barras tem largura 1,
// e altura dada no vetor v
//
// O(n)
15e ll area(vector<int> v) {
b73
       11 \text{ ret} = 0;
4ce
        stack<int> s:
       // valores iniciais pra dar tudo certo
447
       v.insert(v.begin(), -1);
d56
        v.insert(v.end(), -1);
1f8
        s.push(0);
0be
        for(int i = 0; i < (int) v.size(); i++) {</pre>
            while (v[s.top()] > v[i]) {
78e
265
                11 h = v[s.top()]; s.pop();
                ret = max(ret, h * (i - s.top() - 1));
de1
            }
40a
18e
            s.push(i);
020
edf
        return ret;
e43 }
```

3.4 Binomial modular

```
// Computa C(n, k) mod m em O(m + log(m) log(n))
// = O(rapido)
97c ll divi[MAX]:
398 ll expo(ll a, ll b, ll m) {
        if (!b) return 1:
1 c 1
399
        ll ans = expo(a*a\%m, b/2, m);
        if (b\%2) ans *= a:
751
2e9
        return ans%m;
754 }
f0a ll inv(ll a, ll b){
        return 1 < a ? b - inv(b%a,a)*b/a : 1;
bca
041 }
153 template < typename T > tuple < T, T, T > ext_gcd(T a, T b) {
3bd
        if (!a) return {b, 0, 1};
550
        auto [g, x, y] = ext_gcd(b\%a, a);
        return \{g, y - b/a*x, x\};
c59
537 }
bfe template < typename T = 11> struct crt {
627
        Ta, m;
        crt() : a(0), m(1) {}
5f3
        crt(T a_, T m_) : a(a_), m(m_) {}
7eb
911
        crt operator * (crt C) {
238
            auto [g, x, y] = ext_gcd(m, C.m);
dc0
            if ((a - C.a) \% g) a = -1;
            if (a == -1 or C.a == -1) return crt(-1, 0);
4f9
            T lcm = m/g*C.m;
d09
            T ans = a + (x*(C.a-a)/g \% (C.m/g))*m;
eb2
d8d
            return crt((ans % lcm + lcm) % lcm, lcm);
1f2
        }
0d9 };
6f2 pair<11, 11> divide_show(11 n, int p, int k, int pak) {
        if (n == 0) return {0, 1};
d02
        11 blocos = n/pak, falta = n%pak;
        ll periodo = divi[pak], resto = divi[falta];
2ce
616
        ll r = expo(periodo, blocos, pak)*resto%pak;
445
        auto rec = divide_show(n/p, p, k, pak);
a51
        11 y = n/p + rec.first;
```

```
bb9
        r = r*rec.second % pak;
90f
        return {y, r};
533 }
6e6 ll solve_pak(ll n, ll x, int p, int k, int pak) {
d34
        divi[0] = 1:
        for (int i = 1; i <= pak; i++) {</pre>
f2b
            divi[i] = divi[i-1];
901
840
            if (i%p) divi[i] = divi[i] * i % pak;
51a
        }
        auto dn = divide_show(n, p, k, pak), dx = divide_show(x, p, k,
   pak),
162
             dnx = divide\_show(n-x, p, k, pak);
        11 y = dn.first-dx.first-dnx.first, r =
768
            (dn.second*inv(dx.second, pak)%pak)*inv(dnx.second,
b64
   pak)%pak;
035
        return expo(p, y, pak) * r % pak;
d78 }
9dd ll solve(ll n, ll x, int mod) {
        vector < pair < int , int >> f;
490
c3b
        int mod2 = mod;
7b4
        for (int i = 2; i*i <= mod2; i++) if (mod2%i==0) {</pre>
aff
            int c = 0:
75b
            while (mod2\%i==0) mod2 /= i, c++;
2a1
            f.push_back({i, c});
fe7
Off
        if (mod2 > 1) f.push_back({mod2, 1});
e96
        crt ans(0, 1);
a13
        for (int i = 0; i < f.size(); i++) {</pre>
702
            int pak = 1;
7e4
            for (int j = 0; j < f[i].second; j++) pak *= f[i].first;</pre>
            ans = ans * crt(solve_pak(n, x, f[i].first, f[i].second,
   pak), pak);
7fd
5fb
        return ans.a;
689 }
3.5 Closest pair of points
// O(nlogn)
915 pair <pt, pt > closest_pair_of_points(vector <pt > v) {
3d2
        int n = v.size();
fca
        sort(v.begin(), v.end());
```

```
31 c
        for (int i = 1; i < n; i++) if (v[i] == v[i-1]) return</pre>
   {v[i-1], v[i]};
        auto cmp_y = [&](const pt &1, const pt &r) {
c20
b53
            if (1.y != r.y) return 1.y < r.y;
            return 1.x < r.x;</pre>
920
55a
        };
62e
        set < pt, decltype(cmp_y) > s(cmp_y);
3d9
        int 1 = 0, r = -1;
6a2
        11 d2_min = numeric_limits<11>::max();
4d5
        pt pl, pr;
bd1
        const int magic = 5;
a55
        while (r+1 < n) {
7f1
            auto it = s.insert(v[++r]).first:
c92
            int cnt = magic/2;
773
            while (cnt-- and it != s.begin()) it--;
a01
d68
            while (cnt++ < magic and it != s.end()) {</pre>
f19
                 if (!((*it) == v[r])) {
67e
                    11 d2 = dist2(*it, v[r]);
74e
                     if (d2_min > d2) {
                         d2_min = d2;
229
841
                         pl = *it;
4f2
                         pr = v[r];
7d9
                     }
10a
                 }
40d
                it++;
801
eb0
            while (1 < r \text{ and } sq(v[1].x-v[r].x) > d2_min)
   s.erase(v[1++]);
de6
c74
        return {pl, pr};
f90 }
3.6 Coloração de Grafo de Intervalo
// Colore os intervalos com o numero minimo
// de cores de tal forma que dois intervalos
// que se interceptam tem cores diferentes
// As cores vao de 1 ate n
// O(n log(n))
615 vector<int> coloring(vector<pair<int, int>>& v) {
```

3d2

c08

603

150

int n = v.size();

vector<pair<int, pair<int, int>>> ev;

ev.push_back({v[i].first, {1, i}});

for (int i = 0; i < n; i++) {</pre>

```
cda
            ev.push_back({v[i].second, {0, i}});
6a4
49e
        sort(ev.begin(), ev.end());
        vector < int > ans(n), avl(n);
360
265
        for (int i = 0; i < n; i++) avl.push_back(n-i);</pre>
4bf
        for (auto i : ev) {
cbe
            if (i.second.first == 1) {
021
                ans[i.second.second] = avl.back();
a00
                avl.pop_back();
            } else avl.push_back(ans[i.second.second]);
e98
3a6
ba7
        return ans:
83a }
```

3.7 Conectividade Dinamica DC

```
// Offline com Divide and Conquer e
// DSU com rollback
// O(n log^2(n))
8f2 typedef pair <int, int > T;
1cd namespace data {
553
        int n, ans;
573
        int p[MAX], sz[MAX];
        stack<int> S;
ee6
        void build(int n2) {
e5b
1e3
            n = n2:
            for (int i = 0; i < n; i++) p[i] = i, sz[i] = 1;
8a6
0b2
            ans = n;
        }
cba
1b1
        int find(int k) {
006
            while (p[k] != k) k = p[k];
839
            return k;
        }
c1e
072
        void add(T x) {
700
            int a = x.first, b = x.second;
605
            a = find(a), b = find(b);
           if (a == b) return S.push(-1);
843
e7d
            ans - -;
3c6
            if (sz[a] > sz[b]) swap(a, b);
4c2
            S.push(a);
582
            sz[b] += sz[a];
84b
            p[a] = b;
e1a
        }
5eb
        int query() {
```

```
ba7
            return ans;
        }
35c
5cf
        void rollback() {
465
            int u = S.top(); S.pop();
61c
            if (u == -1) return:
270
            sz[p[u]] -= sz[u];
546
            p[u] = u:
Odf
            ans++;
456
        }
568 };
357 int ponta[MAX]; // outra ponta do intervalo ou -1 se for query
4f0 int ans[MAX], n, q;
487 T qu[MAX];
47b void solve(int l = 0, int r = q-1) {
0b1
        if (1 >= r) {
            ans[1] = data::query(); // agora a estrutura ta certa
8c0
505
            return;
        }
f77
962
        int m = (1+r)/2, qnt = 1;
fc7
        for (int i = m+1; i <= r; i++) if (ponta[i]+1 and ponta[i] < 1)</pre>
37d
            data::add(qu[i]), qnt++;
221
        solve(1, m);
593
        while (--qnt) data::rollback();
        for (int i = 1; i <= m; i++) if (ponta[i]+1 and ponta[i] > r)
a2c
37d
            data::add(qu[i]), qnt++;
37b
        solve(m+1, r):
        while (qnt--) data::rollback();
281
0d4 }
3.8 Conectividade Dinamica LCT
// Offline com link-cut trees
// O(n log(n))
1ef namespace lct {
3c9
        struct node {
            int p, ch[2];
19f
a2a
            int val, sub;
aa6
            bool rev;
f93
            node() {}
            node(int v) : p(-1), val(v), sub(v), rev(0) { ch[0] = }
   ch[1] = -1; }
        };
cac
c53
        node t[2*MAX]; // MAXN + MAXQ
```

```
99e
        map<pair<int, int>, int> aresta;
e4d
        int sz:
95a
        void prop(int x) {
aa2
            if (t[x].rev) {
f95
                swap(t[x].ch[0], t[x].ch[1]);
379
                if (t[x].ch[0]+1) t[t[x].ch[0]].rev ^= 1;
                if (t[x].ch[1]+1) t[t[x].ch[1]].rev ^= 1;
c3d
50e
            t[x].rev = 0;
693
750
564
        void update(int x) {
e8d
            t[x].sub = t[x].val:
            for (int i = 0; i < 2; i++) if (t[x].ch[i]+1) {</pre>
8ca
621
                prop(t[x].ch[i]);
78d
                t[x].sub = min(t[x].sub, t[t[x].ch[i]].sub);
3e4
            }
       }
9bf
971
        bool is_root(int x) {
            return t[x].p == -1 or (t[t[x].p].ch[0] != x and
   t[t[x].p].ch[1] != x);
cf1
        void rotate(int x) {
ed6
497
            int p = t[x].p, pp = t[p].p;
fc4
            if (!is_root(p)) t[pp].ch[t[pp].ch[1] == p] = x;
251
            bool d = t[p].ch[0] == x;
            t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
461
a76
            if (t[p].ch[!d]+1) t[t[p].ch[!d]].p = p;
            t[x].p = pp, t[p].p = x;
8fa
444
            update(p), update(x);
       }
f31
238
        int splay(int x) {
            while (!is_root(x)) {
18c
497
                int p = t[x].p. pp = t[p].p:
77b
                if (!is_root(p)) prop(pp);
be5
                prop(p), prop(x);
                if (!is\_root(p)) rotate((t[pp].ch[0] == p)^(t[p].ch[0]
0 c 5
   == x) ? x : p);
64f
                rotate(x):
72c
aab
            return prop(x), x;
08f
        int access(int v) {
f16
0eb
            int last = -1:
            for (int w = v; w+1; update(last = w), splay(v), w =
d9f
   t[v].p)
                splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
024
```

```
3d3
            return last;
        }
294
952
        void make_tree(int v, int w=INF) { t[v] = node(w); }
82f
        bool conn(int v, int w) {
2cf
            access(v). access(w):
b9b
            return v == w ? true : t[v].p != -1;
ec0
        }
277
        void rootify(int v) {
5e3
            access(v);
            t[v].rev ^= 1;
a02
a05
        }
        int query(int v, int w) {
a1d
b54
            rootifv(w). access(v):
249
            return t[v].sub;
c28
        }
204
        void link_(int v, int w) {
821
            rootify(w);
389
            t[w].p = v;
523
        void link(int v, int w, int x) { // v--w com peso x
6b8
379
            int id = MAX + sz++;
110
            aresta[make_pair(v, w)] = id;
ab6
            make_tree(id, x);
c88
            link_(v, id), link_(id, w);
984
        }
e63
        void cut_(int v, int w) {
            rootify(w), access(v);
b54
264
            t[v].ch[0] = t[t[v].ch[0]].p = -1;
7cd
031
        void cut(int v, int w) {
            int id = aresta[make_pair(v, w)];
b0f
a4a
            cut_(v, id), cut_(id, w);
        }
840
0d7 }
893 void dyn_conn() {
        int n, q; cin >> n >> q;
c5f
d6e
        vector <int > p(2*q, -1); // outra ponta do intervalo
b4f
        for (int i = 0; i < n; i++) lct::make_tree(i);</pre>
fbf
        vector<pair<int, int>> qu(q);
139
        map<pair<int, int>, int> m;
        for (int i = 0; i < q; i++) {</pre>
abf
3c2
            char c; cin >> c;
ef6
            if (c == '?') continue;
602
            int a, b; cin >> a >> b; a--, b--;
d11
            if (a > b) swap(a, b);
8a1
            qu[i] = {a, b};
```

```
8d7
            if (c == '+') {
                p[i] = i+q, p[i+q] = i;
94b
906
                m[make_pair(a, b)] = i;
9d9
            } else {
                int j = m[make_pair(a, b)];
412
                p[i] = j, p[j] = i;
ac2
            }
850
9e5
        }
447
        int ans = n;
        for (int i = 0; i < q; i++) {</pre>
abf
87d
            if (p[i] == -1) {
886
                cout << ans << endl: // numero de comp conexos</pre>
5e2
            }
b35
            int a = qu[i].first, b = qu[i].second;
69d
            if (p[i] > i) { // +
c4d
ac5
                if (lct::conn(a, b)) {
                    int mi = lct::query(a, b);
18f
                    if (p[i] < mi) {</pre>
993
dd3
                         p[p[i]] = p[i];
5e2
                         continue;
474
                    lct::cut(qu[p[mi]].first, qu[p[mi]].second), ans++;
6f7
                    p[mi] = mi;
6ea
9a9
d1d
                lct::link(a, b, p[i]), ans--;
            } else if (p[i] != i) lct::cut(a, b), ans++; // -
9d0
c03
       }
56a }
```

3.9 Conj. Indep. Maximo com Peso em Grafo de Intervalo

```
// Retorna os indices ordenados dos intervalos selecionados
// Se tiver empate, retorna o que minimiza o comprimento total
// O(n log(n))
31e vector<int> ind_set(vector<tuple<int, int, int>>& v) {
        vector<tuple<int, int, int>> w;
b27
f14
        for (int i = 0: i < v.size(): i++) {</pre>
            w.push_back(tuple(get<0>(v[i]), 0, i));
e85
6f0
            w.push_back(tuple(get<1>(v[i]), 1, i));
17f
d1d
        sort(w.begin(), w.end());
844
        vector < int > nxt(v.size());
c22
        vector<pair<ll, int>> dp(v.size());
```

```
0eb
        int last = -1;
723
        for (auto [fim, t, i] : w) {
25a
            if (t == 0) {
                nxt[i] = last;
4ca
5e2
                continue;
5fd
78b
            dp[i] = \{0, 0\}:
            if (last != -1) dp[i] = max(dp[i], dp[last]);
cb8
911
            pair<11, int> pega = {get<2>(v[i]), -(get<1>(v[i]) -
   get<0>(v[i]) + 1)};
5d3
            if (nxt[i] != -1) pega.first += dp[nxt[i]].first,
   pega.second += dp[nxt[i]].second;
b08
            if (pega > dp[i]) dp[i] = pega;
7cb
            else nxt[i] = last;
381
            last = i;
        }
b7c
977
        pair<11, int > ans = \{0, 0\};
919
        int idx = -1:
        for (int i = 0; i < v.size(); i++) if (dp[i] > ans) ans =
ceb
   dp[i], idx = i;
4b8
        vector<int> ret;
        while (idx != -1) {
fdd
d69
            if (get < 2 > (v[idx]) > 0 and
                (nxt[idx] == -1 or get<1>(v[nxt[idx]]) <</pre>
a05
   get<0>(v[idx]))) ret.push_back(idx);
e4f
            idx = nxt[idx]:
042
        }
0ea
        sort(ret.begin(), ret.end());
        return ret;
edf
c4d }
```

3.10 Convex Hull Dinamico

```
// insert - O(log n) amortizado
// is_inside - O(log n)
0b9 struct upper {
af8
        set <pt> se;
80b
        set <pt>::iterator it;
25 c
        int is_under(pt p) { // 1 -> inside ; 2 -> border
fe0
            it = se.lower_bound(p);
            if (it == se.end()) return 0:
633
a94
            if (it == se.begin()) return p == *it ? 2 : 0;
ca0
            if (ccw(p, *it, *prev(it))) return 1;
402
            return ccw(p, *prev(it), *it) ? 0 : 2;
dba
        }
```

```
void insert(pt p) {
eaa
712
            if (is_under(p)) return;
            if (it != se.end()) while (next(it) != se.end() and
   !ccw(*next(it), *it, p))
                it = se.erase(it);
316
be3
            if (it != se.begin()) while (--it != se.begin() and
   !ccw(p, *it, *prev(it)))
                it = se.erase(it);
316
0c8
            se.insert(p);
5da
        }
750 }:
06f struct dyn_hull {
d93
        upper U, L;
333
        int is_inside(pt p) {
632
            int u = U.is\_under(p), l = L.is\_under({-p.x, -p.y});
4 c.c
            if (!u or !1) return 0;
fc0
            return max(u, 1);
478
        void insert(pt p) {
86c
            U.insert(p);
925
            L.insert({-p.x, -p.y});
64b
285
        int size() {
7c2
            int ans = U.se.size() + L.se.size();
1c9
            return ans <= 2 ? ans/2 : ans-2;</pre>
ad5
        }
65e }:
3.11 Distancia maxima entre dois pontos
```

```
// \max_{dist2(v)} - O(n \log(n))
// max_dist_manhattan - O(n)
// Quadrado da Distancia Euclidiana (precisa copiar convex_hull, ccw e
   pt)
859 ll max_dist2(vector<pt> v) {
        v = convex_hull(v);
a14
       if (v.size() \le 2) return dist2(v[0], v[1%v.size()]):
       11 ans = 0:
04b
323
       int n = v.size(), j = 0;
603
        for (int i = 0: i < n: i++) {</pre>
057
            while (!ccw(v[(i+1)%n]-v[i], pt(0, 0), v[(i+1)%n]-v[i])) j
   = (j+1) \%n;
```

```
e7a
            ans = \max(\{ans, dist2(v[i], v[j]), dist2(v[(i+1)%n],
   v[i])});
        }
1f6
ba7
        return ans;
bda }
// Distancia de Manhattan
c51 template < typename T> T max_dist_manhattan(vector < pair < T, T>> v) {
8eb
        T min_sum, max_sum, min_dif, max_dif;
4f5
        min_sum = max_sum = v[0].first + v[0].second;
271
        min_dif = max_dif = v[0].first - v[0].second;
c25
        for (auto [x, y] : v) {
1cb
            min sum = min(min sum, x+v):
683
            max_sum = max(max_sum, x+y);
782
            min_dif = min(min_dif, x-y);
            max_dif = max(max_dif, x-y);
af7
e3a
9f0
        return max(max_sum - min_sum, max_dif - min_dif);
4e9 }
     Distinct Range Query
// build - O(n (log n + log(sigma)))
// query - O(log(sigma))
789 namespace perseg { };
53d int qt[MAX];
edc void build(vector<int>& v) {
        int n = v.size();
16b
        perseg::build(n);
663
        map<int, int> last;
05e
        int at = 0;
603
        for (int i = 0; i < n; i++) {</pre>
817
            if (last.count(v[i])) {
a58
                perseg::update(last[v[i]], -1);
69a
                at++;
d1f
            }
4f2
            perseg::update(i, 1);
460
            qt[i] = ++at;
efe
            last[v[i]] = i:
d6f
        }
0f4 }
9e3 int query(int 1, int r) {
080
        return perseg::query(1, r, qt[r]);
```

3.13 Distinct Range Query com Update

```
// build - O(n log(n))
// query - O(log^2(n))
// update - O(log^2(n))
774 #include <ext/pb_ds/assoc_container.hpp>
30f #include <ext/pb_ds/tree_policy.hpp>
0d7 using namespace __gnu_pbds;
4fc template <class T>
        using ord_set = tree<T, null_type, less<T>, rb_tree_tag,
def
        tree_order_statistics_node_update>;
042 int v[MAX], n, nxt[MAX], prv[MAX];
f60 map<int, set<int> > ocor;
e04 namespace bit {
        ord_set<pair<int, int>> bit[MAX];
0a8
        void build() {
            for (int i = 1; i <= n; i++) bit[i].insert({nxt[i-1],</pre>
3e1
   i-1});
78a
            for (int i = 1; i <= n; i++) {</pre>
                int j = i + (i\&-i);
edf
                if (j <= n) for (auto x : bit[i]) bit[j].insert(x);</pre>
d03
            }
5cb
af6
        }
d3f
        int pref(int p, int x) {
7c9
            int ret = 0;
bbf
            for (; p; p -= p\&-p) ret += bit[p].order_of_key({x, -INF});
edf
            return ret;
0e1
        }
d50
        int query(int 1, int r, int x) {
e55
            return pref(r+1, x) - pref(1, x);
9b4
ff3
        void update(int p, int x) {
f17
            int p2 = p;
            for (p++; p \le n; p += p\&-p) {
5ed
                bit[p].erase({nxt[p2], p2});
ca8
f6b
                bit[p].insert({x, p2});
            }
3df
151
        }
c63 }
0a8 void build() {
```

```
383
        for (int i = 0; i < n; i++) nxt[i] = INF;</pre>
        for (int i = 0; i < n; i++) prv[i] = -INF;</pre>
7b3
d07
        vector < pair < int , int >> t;
        for (int i = 0; i < n; i++) t.push_back({v[i], i});</pre>
348
        sort(t.begin(), t.end());
3fd
        for (int i = 0; i < n; i++) {</pre>
603
b40
            if (i and t[i].first == t[i-1].first)
565
                 prv[t[i].second] = t[i-1].second;
            if (i+1 < n and t[i].first == t[i+1].first)</pre>
a8b
                 nxt[t[i].second] = t[i+1].second;
12f
48d
        }
a23
        for (int i = 0: i < n: i++) ocor[v[i]].insert(i):</pre>
1d7
        bit::build();
d44 }
aae void muda(int p, int x) {
f92
        bit::update(p, x);
c3d
        nxt[p] = x:
97c }
4ea int query(int a, int b) {
a0a
        return b-a+1 - bit::query(a, b, b+1);
511 }
ff3 void update(int p, int x) { // mudar valor na pos. p para x
c0b
        if (prv[p] > -INF) muda(prv[p], nxt[p]);
        if (nxt[p] < INF) prv[nxt[p]] = prv[p];</pre>
4ae
5bf
        ocor[v[p]].erase(p);
4b4
        if (!ocor[x].size()) {
19d
            muda(p, INF);
8d4
            prv[p] = -INF:
        } else if (*ocor[x].rbegin() < p) {</pre>
a69
5b5
            int i = *ocor[x].rbegin();
f64
            prv[p] = i;
19d
            muda(p, INF);
5f2
            muda(i, p);
9d9
        } else {
            int i = *ocor[x].lower_bound(p);
d46
33f
            if (prv[i] > -INF) {
f17
                 muda(prv[i], p);
8f9
                 prv[p] = prv[i];
bc4
            } else prv[p] = -INF;
523
            prv[i] = p;
597
            muda(p, i);
```

```
5d7     }
c96     v[p] = x; ocor[x].insert(p);
38e }
```

3.14 Dominator Points

```
// Se um ponto A tem ambas as coordenadas >= B, dizemos
// que A domina B
// is_dominated(p) fala se existe algum ponto no conjunto
// que domina p
// insert(p) insere p no conjunto
// (se p for dominado por alguem, nao vai inserir)
// o multiset 'quina' guarda informacao sobre os pontos
// nao dominados por um elemento do conjunto que nao dominam
// outro ponto nao dominado por um elemento do conjunto
// No caso, armazena os valores de x+y esses pontos
//
// Complexidades:
// is_dominated - O(log(n))
// insert - O(log(n)) amortizado
// query - O(1)
e2a struct dominator_points {
        set < pair < int , int >> se;
baf
        multiset < int > quina;
4dd
        bool is_dominated(pair<int, int> p) {
a85
            auto it = se.lower_bound(p);
80f
            if (it == se.end()) return 0;
633
            return it->second >= p.second;
ab4
28f
99Ъ
        void mid(pair<int, int> a, pair<int, int> b, bool rem) {
29a
            pair < int , int > m = {a.first+1, b.second+1};
b19
            int val = m.first + m.second;
638
            if (!rem) quina.insert(val);
731
            else quina.erase(quina.find(val));
241
7c4
        bool insert(pair<int, int> p) {
            if (is_dominated(p)) return 0;
fb4
            auto it = se.lower_bound(p);
80f
            if (it != se.begin() and it != se.end())
ca9
d4a
                mid(*prev(it), *it, 1);
            while (it != se.begin()) {
1fa
049
                it--;
23 c
                if (it->second > p.second) break;
b86
                if (it != se.begin()) mid(*prev(it), *it, 1);
316
                it = se.erase(it):
```

```
acd
433
            it = se.insert(p).first;
69e
            if (it != se.begin()) mid(*prev(it), *it, 0);
96d
             if (next(it) != se.end()) mid(*it, *next(it), 0);
6a5
            return 1:
688
        }
5eb
        int query() {
956
            if (!quina.size()) return INF;
add
             return *quina.begin();
h8b
        }
09f };
3.15 DP de Dominação 3D
// Computa para todo ponto i,
// dp[i] = 1 + max_{i} dominado por i dp[i]
// em que ser dominado eh ter as 3 coordenadas menores
// Da pra adaptar facil para outras dps
//
// O(n log^2 n), O(n) de memoria
c53 void lis2d(vector<vector<tuple<int, int, int>>>& v, vector<int>&
    dp, int 1, int r) {
893
        if (1 == r) {
56f
            for (int i = 0; i < v[1].size(); i++) {</pre>
                 int ii = get<2>(v[1][i]);
8b5
1ce
                 dp[ii] = max(dp[ii], 1);
4b0
            }
505
            return;
3e4
        }
        int m = (1+r)/2;
ee4
62b
        lis2d(v, dp, 1, m);
325
        vector<tuple<int, int, int>> vv[2];
d44
        vector < int > Z;
871
        for (int i = 1; i <= r; i++) for (auto it : v[i]) {</pre>
2ef
             vv[i > m].push_back(it);
042
            Z.push_back(get<1>(it));
        }
0d1
        sort(vv[0].begin(), vv[0].end());
e9f
9b5
        sort(vv[1].begin(), vv[1].end());
0d1
        sort(Z.begin(), Z.end());
        auto get_z = [&](int z) { return lower_bound(Z.begin(),
    Z.end(), z) - Z.begin(); };
        vector < int > bit(Z.size());
c51
181
        int i = 0;
```

```
e9a
        for (auto [y, z, id] : vv[1]) {
            while (i < vv[0].size() and get<0>(vv[0][i]) < y) {</pre>
6bd
397
                 auto [y2, z2, id2] = vv[0][i++];
                for (int p = get_z(z2)+1; p <= Z.size(); p += p\&-p)
ea0
                     bit[p-1] = max(bit[p-1], dp[id2]);
300
            }
82c
d3b
            int q = 0;
            for (int p = get_z(z); p; p -= p&-p) q = max(q, bit[p-1]);
fd9
614
            dp[id] = max(dp[id], q + 1);
acc
c25
        lis2d(v, dp, m+1, r);
4d6 }
4de vector<int> solve(vector<tuple<int, int, int>> v) {
3d2
        int n = v.size();
        vector<tuple<int, int, int, int>> vv;
cd4
603
        for (int i = 0; i < n; i++) {</pre>
            auto [x, y, z] = v[i];
9be
            vv.emplace_back(x, y, z, i);
5bb
64c
        sort(vv.begin(), vv.end());
bd3
e11
        vector < vector < tuple < int , int , int >>> V;
603
        for (int i = 0; i < n; i++) {</pre>
a5b
            int j = i;
808
            V.emplace_back();
c01
            while (j < n and get <0>(vv[j]) == get <0>(vv[i])) {
ba6
                 auto [x, y, z, id] = vv[j++];
                V.back().emplace_back(y, z, id);
cbb
8bd
            }
452
            i = j-1;
ac4
388
        vector < int > dp(n);
        lis2d(V. dp. 0. V.size()-1):
839
898
        return dp;
bOa }
3.16 Gray Code
// Gera uma permutacao de 0 a 2^n-1, de forma que
// duas posicoes adjacentes diferem em exatamente 1 bit
//
// O(2<sup>n</sup>)
df6 vector<int> gray_code(int n) {
73f
        vector < int > ret(1 << n);</pre>
        for (int i = 0; i < (1<<n); i++) ret[i] = i^(i>>1);
```

f29

```
edf
                         return ret;
 840 }
 3.17 Half-plane intersection
 // Cada half-plane eh identificado por uma reta e a regiao ccw a ela
 // O(n log n)
  f4f vector <pt> hp_intersection(vector <line> &v) {
                         deque < pt > dq = \{\{INF, INF\}, \{-INF, INF\}, \{-INF, -INF\}, \{INF, -INF
            -INF}};
 d41 #warning considerar trocar por compare_angle
  de3
                         sort(v.begin(), v.end(), [&](line r, line s) { return
           angle(r.q-r.p) < angle(s.q-s.p); \});
                         for(int i = 0; i < v.size() and dq.size() > 1; i++) {
 5e9
 c69
                                    pt p1 = dq.front(), p2 = dq.back();
 6c6
                                    while (dq.size() and !ccw(v[i].p, v[i].q, dq.back()))
 47b
                                               p1 = dq.back(), dq.pop_back();
 0a2
                                    while (dq.size() and !ccw(v[i].p, v[i].q, dq.front()))
 7 cf
                                               p2 = dq.front(), dq.pop_front();
 4d9
                                    if (!dq.size()) break;
 606
                                    if (p1 == dq.front() and p2 == dq.back()) continue;
                                    dq.push_back(inter(v[i], line(dq.back(), p1)));
  c9b
 65 c
                                    dq.push_front(inter(v[i], line(dq.front(), p2)));
 fdd
                                    if (dq.size() > 1 and dq.back() == dq.front())
           dq.pop_back();
 4d8
                         return vector < pt > (dq.begin(), dq.end());
 b2b
 f56 }
3.18 Heap Sort
 // O(n log n)
 f18 void down(vector<int>& v, int n, int i) {
                         while ((i = 2*i+1) < n) {
 583
                                    if (i+1 < n and v[i] < v[i+1]) i++;</pre>
                                    if (v[i] < v[(i-1)/2]) break;
 b27
 322
                                    swap(v[i], v[(i-1)/2]);
 170
                        }
 724 }
eb6 void heap_sort(vector<int>& v) {
```

```
int j1 = way[j0];
3d2
        int n = v.size();
                                                                              4c5
        for (int i = n/2-1; i \ge 0; i--) down(v, n, i);
61d
                                                                              0d7
                                                                                                  p[j0] = p[j1];
917
        for (int i = n-1; i > 0; i--)
                                                                             6d4
                                                                                                   j0 = j1;
37f
            swap(v[0], v[i]), down(v, i, 0);
                                                                              886
                                                                                              } while (j0);
b33 }
                                                                              38d
                                                                              306
                                                                                          vector < int > ans(n);
                                                                             6db
                                                                                          for (int j = 1; j \le n; j++) ans [p[j]-1] = j-1;
3.19 Hungaro
                                                                              da3
                                                                                          return make_pair(-v[0], ans);
                                                                              979
                                                                                      }
// Resolve o problema de assignment (matriz n x n)
                                                                              64c };
// Colocar os valores da matriz em 'a' (pode < 0)</pre>
// assignment() retorna um par com o valor do
                                                                             3.20 Inversion Count
// assignment minimo, e a coluna escolhida por cada linha
//
                                                                             // Computa o numero de inversoes para transformar
// O(n^3)
                                                                             // l em r (se nao tem como, retorna -1)
                                                                             //
a6a template < typename T > struct hungarian {
                                                                             // O(n log(n))
1a8
        int n:
        vector < vector < T >> a;
a08
                                                                              37b template < typename T > 11 inv_count(vector < T > 1, vector < T > r = {}) {
f36
        vector <T> u. v:
                                                                                      if (!r.size()) {
                                                                              bb6
5ff
        vector < int > p, way;
                                                                              796
                                                                                          r = 1:
f1e
        T inf:
                                                                             1bc
                                                                                          sort(r.begin(), r.end());
                                                                              dfb
                                                                                      }
c3f
        hungarian(int n_): n(n_{-}), u(n+1), v(n+1), p(n+1), way(n+1) {
                                                                                      int n = 1.size();
                                                                              874
b2f
            a = vector < vector < T >> (n, vector < T > (n));
                                                                              8c0
                                                                                      vector < int > v(n), bit(n):
1f3
            inf = numeric_limits <T>::max();
                                                                              4e9
                                                                                      vector < pair < T, int >> w;
78f
        }
                                                                              61 c
                                                                                      for (int i = 0; i < n; i++) w.push_back({r[i], i+1});</pre>
d67
        pair <T. vector <int >> assignment() {
                                                                              d1d
                                                                                      sort(w.begin(), w.end());
            for (int i = 1; i <= n; i++) {</pre>
78a
                                                                                      for (int i = 0; i < n; i++) {</pre>
                                                                              603
8c9
                 p[0] = i;
                                                                                          auto it = lower_bound(w.begin(), w.end(), make_pair(l[i],
                                                                              bf3
                 int j0 = 0;
625
                                                                                 0));
ce7
                 vector <T> minv(n+1, inf);
                                                                                          if (it == w.end() or it->first != l[i]) return -1; // nao
                                                                             1bf
241
                 vector < int > used(n+1, 0);
                                                                                 da
016
                 do {
                                                                              962
                                                                                          v[i] = it->second;
472
                     used[j0] = true;
                                                                              6c0
                                                                                          it->second = -1;
d24
                     int i0 = p[j0], j1 = -1;
                                                                              964
                                                                                      }
7e5
                     T delta = inf;
9ac
                     for (int j = 1; j <= n; j++) if (!used[j]) {</pre>
                                                                              04b
                                                                                      11 \text{ ans} = 0:
                         T cur = a[i0-1][j-1] - u[i0] - v[j];
7bf
                                                                                      for (int i = n-1; i >= 0; i--) {
                                                                              45b
9f2
                         if (cur < minv[j]) minv[j] = cur, wav[j] = j0;</pre>
                                                                                          for (int j = v[i]-1; j; j = j\&-j) ans += bit[j];
                                                                              2d9
                         if (minv[j] < delta) delta = minv[j], j1 = j;</pre>
821
                                                                                          for (int j = v[i]; j < n; j += j\&-j) bit[j]++;
                                                                              3a1
                     }
4d1
                                                                              ebe
                                                                                      }
f63
                     for (int j = 0; j <= n; j++)
                                                                              ba7
                                                                                      return ans;
2c5
                         if (used[j]) u[p[j]] += delta, v[j] -= delta;
                                                                              eef }
                         else minv[j] -= delta;
6ec
6d4
                     j0 = j1;
f4f
                 } while (p[i0] != 0);
                                                                             3.21 LIS - recupera
016
                 do {
```

```
// Calcula e retorna uma LIS
// O(n.log(n))
121 template < typename T> vector < T> lis(vector < T>& v) {
        int n = v.size(), m = -1;
f0c
        vector <T> d(n+1. INF):
        vector < int > l(n);
aec
       d[0] = -INF;
007
603
        for (int i = 0; i < n; i++) {</pre>
            // Para non-decreasing use upper_bound()
4fd
            int t = lower_bound(d.begin(), d.end(), v[i]) - d.begin();
            d[t] = v[i], l[i] = t, m = max(m, t);
3ad
89 c
4ff
        int p = n;
5a9
        vector <T> ret:
        while (p--) if (1[p] == m) {
cdf
883
            ret.push_back(v[p]);
76b
            m - - ;
f83
        reverse(ret.begin(),ret.end());
969
edf
        return ret:
474 }
3.22 LIS - tamanho
// Calcula o tamanho da LIS
//
// O(n log(n))
84b template < typename T > int lis(vector < T > &v) {
        vector <T> ans;
2da
5e0
        for (T t : v){
            // Para non-decreasing use upper_bound()
            auto it = lower_bound(ans.begin(), ans.end(), t);
fe6
d7f
            if (it == ans.end()) ans.push_back(t);
b94
            else *it = t;
1f5
1eb
        return ans.size();
402 }
      Minimum Enclosing Circle
```

// O(n) com alta probabilidade

```
22c const double EPS = 1e-12;
878 mt19937 rng((int)
    chrono::steady_clock::now().time_since_epoch().count());
b2a struct pt {
662
        double x, y;
        pt(double x_{=} = 0, double y_{=} = 0) : x(x_{=}), y(y_{=}) {}
be7
7af
        pt operator + (const pt& p) const { return pt(x+p.x, y+p.y); }
        pt operator - (const pt& p) const { return pt(x-p.x, y-p.y); }
b23
254
        pt operator * (double c) const { return pt(x*c, y*c); }
701
        pt operator / (double c) const { return pt(x/c, y/c); }
54d }:
2f9 double dot(pt p, pt q) { return p.x*q.x+p.y*q.y; }
dd5 double cross(pt p, pt q) { return p.x*q.y-p.y*q.x; }
e7c double dist(pt p, pt q) { return sqrt(dot(p-q, p-q)); }
3f4 pt center(pt p, pt q, pt r) {
5d9
        pt a = p-r, b = q-r;
        pt c = pt(dot(a, p+r)/2, dot(b, q+r)/2);
        return pt(cross(c, pt(a.y, b.y)), cross(pt(a.x, b.x), c)) /
   cross(a, b):
fc8 }
aa8 struct circle {
f41
        pt cen;
c12
        double r:
898
        circle(pt cen_, double r_) : cen(cen_), r(r_) {}
83 c
        circle(pt a, pt b, pt c) {
            cen = center(a, b, c);
13d
1f1
            r = dist(cen, a);
bc1
        }
cd5
        bool inside(pt p) { return dist(p, cen) < r+EPS; }</pre>
2a6 };
806 circle minCirc(vector<pt> v) {
f21
        shuffle(v.begin(), v.end(), rng);
ae0
        circle ret = circle(pt(0, 0), 0);
618
        for (int i = 0; i < v.size(); i++) if (!ret.inside(v[i])) {</pre>
            ret = circle(v[i], 0);
16a
f11
            for (int j = 0; j < i; j++) if (!ret.inside(v[j])) {</pre>
881
                ret = circle((v[i]+v[j])/2, dist(v[i], v[j])/2);
                for (int k = 0; k < j; k++) if (!ret.inside(v[k]))
b8c
43f
                    ret = circle(v[i], v[j], v[k]);
5f8
            }
        }
6a1
```

```
return ret;
eba }
3.24 Minkowski Sum
// Computa A+B = \{a+b : a \setminus A, b \setminus B\}, em que
// A e B sao poligonos convexos
// A+B eh um poligono convexo com no max |A|+|B| pontos
//
// O(|A|+|B|)
539 vector<pt> minkowski(vector<pt> p, vector<pt> q) {
        auto fix = [](vector < pt > & P) {
            rotate(P.begin(), min_element(P.begin(), P.end()),
515
   P.end());
018
            P.push_back(P[0]), P.push_back(P[1]);
f24
        };
        fix(p), fix(q);
889
        vector < pt > ret;
8af
692
        int i = 0, j = 0;
2ee
        while (i < p.size()-2 \text{ or } j < q.size()-2) {
898
            ret.push_back(p[i] + q[j]);
            auto c = ((p[i+1] - p[i]) ^ (q[i+1] - q[i]));
732
           if (c \ge 0) i = min<int>(i+1, p.size()-2);
ebc
81e
            if (c \le 0) j = min \le int > (j+1, q.size()-2);
9ff
edf
        return ret;
d7c }
c3e ld dist_convex(vector<pt> p, vector<pt> q) {
dc2
        for (pt& i : p) i = i * -1;
        auto s = minkowski(p, q);
44c
        if (inpol(s, pt(0, 0))) return 0;
95d
       return 1:
6a5
       ld ans = DINF;
921
073
        for (int i = 0; i < s.size(); i++) ans = min(ans,
f04
                disttoseg(pt(0, 0), line(s[(i+1)%s.size()], s[i])));
ba7
        return ans;
2f5 }
3.25 MO
// Para ter o bound abaixo, escolher
// SQ = n / sqrt(q)
//
// O(n * sqrt(q))
```

edf

```
0d2 const int MAX = 1e5+10;
6ff const int SQ = sqrt(MAX);
b69 int v[MAX];
b65 int ans, freq[MAX];
9da inline void insert(int p) {
        int o = v[p];
        freq[o]++;
591
        ans += (freq[o] == 1);
21d }
a25 inline void erase(int p) {
        int o = v[p];
ae0
        ans -= (freq[o] == 1);
7ee
        freq[o]--;
ba2
dc7 }
e51 inline ll hilbert(int x, int y) {
        static int N = 1 << (__builtin_clz(0) - __builtin_clz(MAX));</pre>
100
        int rx, ry, s;
b72
        11 d = 0:
        for (s = N/2; s > 0; s /= 2) {
43b
c95
            rx = (x \& s) > 0, ry = (y \& s) > 0;
e3e
            d += s * 11(s) * ((3 * rx) ^ ry);
d2e
            if (rv == 0) {
5aa
                if (rx == 1) x = N-1 - x, y = N-1 - y;
9dd
                swap(x, y);
e2d
            }
888
        }
        return d;
be2
7fa }
bac #define HILBERT true
617 vector<int> MO(vector<pair<int, int>> &q) {
c3b
        ans = 0:
        int m = q.size();
c23
3f8
        vector < int > ord(m);
be8
        iota(ord.begin(), ord.end(), 0);
6a6 #if HILBERT
        vector<ll> h(m):
        for (int i = 0; i < m; i++) h[i] = hilbert(q[i].first,</pre>
74c
   q[i].second):
        sort(ord.begin(), ord.end(), [&](int 1, int r) { return h[1] <</pre>
   h[r]; });
8c1 #else
        sort(ord.begin(), ord.end(), [&](int 1, int r) {
```

```
9c9
            if (q[1].first / SQ != q[r].first / SQ) return q[1].first
                                                                             e7d
   < q[r].first;
                                                                             3c6
            if ((q[1].first / SQ) % 2) return q[1].second >
                                                                             4c2
   q[r].second:
                                                                             582
            return q[1].second < q[r].second;</pre>
a66
                                                                             84b
        }):
                                                                             720
bec
f2e #endif
                                                                             35 c
435
                                                                             5cf
        vector < int > ret(m):
3d9
        int 1 = 0, r = -1;
                                                                             465
                                                                             61 c
8b0
        for (int i : ord) {
                                                                            270
6c6
            int ql, qr;
                                                                             546
4f5
            tie(al. ar) = a[i]:
                                                                            Odf
            while (r < qr) insert(++r);</pre>
026
                                                                             456
232
            while (1 > q1) insert(--1);
75e
            while (1 < q1) erase(1++);</pre>
fe8
            while (r > qr) erase(r--);
            ret[i] = ans;
381
116
edf
        return ret;
fb7 }
                                                                             547
                                                                             c23
3.26 MO - DSU
                                                                            3f8
                                                                             be8
// Dado uma lista de arestas de um grafo, responde
                                                                             d01
// para cada query(1, r), quantos componentes conexos
                                                                             9c9
// o grafo tem se soh considerar as arestas l. l+1, ..., r
// Da pra adaptar pra usar MO com qualquer estrutura rollbackavel
                                                                             a66
//
                                                                             b90
// O(m sqrt(q) log(n))
                                                                             435
8d3 struct dsu {
                                                                             3bd
553
        int n, ans;
                                                                             dd5
2e3
        vector < int > p, sz;
                                                                             5ec
ee6
        stack<int> S;
                                                                             acc
                                                                             00c
4b8
        dsu(int n_{-}) : n(n_{-}), ans(n), p(n), sz(n) {
                                                                             b99
8a6
            for (int i = 0; i < n; i++) p[i] = i, sz[i] = 1;
                                                                             64e
aae
        }
                                                                             259
        int find(int k) {
1b1
                                                                             6b0
006
            while (p[k] != k) k = p[k];
839
            return k:
                                                                             dd5
c1e
        }
                                                                            176
553
        void add(pair<int, int> x) {
                                                                            ae9
700
            int a = x.first, b = x.second;
                                                                             e25
605
            a = find(a), b = find(b);
                                                                             ebc
            if (a == b) return S.push(-1);
843
```

```
ans --;
            if (sz[a] > sz[b]) swap(a, b);
            S.push(a);
            sz[b] += sz[a];
            p[a] = b;
        int query() { return ans; }
        void rollback() {
            int u = S.top(); S.pop();
            if (u == -1) return:
            sz[p[u]] -= sz[u];
            p[u] = u;
            ans++:
        }
9c1 };
1a8 int n;
e93 vector<pair<int, int>> ar;
// 9d242b
617 vector<int> MO(vector<pair<int, int>> &q) {
        int SQ = sqrt(q.size()) + 1;
        int m = q.size();
        vector<int> ord(m);
        iota(ord.begin(), ord.end(), 0);
        sort(ord.begin(), ord.end(), [&](int 1, int r) {
                 if (q[1].first / SQ != q[r].first / SQ) return
   q[1].first < q[r].first;</pre>
                return q[1].second < q[r].second;</pre>
                });
        vector < int > ret(m);
        dsu small(n):
        for (int i = 0; i < m; i++) {</pre>
            auto [1, r] = q[ord[i]];
            if (1 / SQ == r / SQ) {
                for (int k = 1; k <= r; k++) small.add(ar[k]);</pre>
                ret[ord[i]] = small.query();
                for (int k = 1; k <= r; k++) small.rollback();</pre>
            }
        }
        for (int i = 0; i < m; i++) {</pre>
            dsu D(n):
            int fim = q[ord[i]].first/SQ*SQ + SQ - 1;
            int last_r = fim;
            int j = i-1;
```

```
00c
            while (j+1 < m and q[ord[j+1]].first / SQ ==</pre>
   q[ord[i]].first / SQ) {
                auto [1, r] = q[ord[++j]];
                if (1 / SQ == r / SQ) continue;
f58
59b
                 while (last_r < r) D.add(ar[++last_r]);</pre>
                for (int k = 1; k <= fim; k++) D.add(ar[k]);</pre>
2cf
9b2
                ret[ord[j]] = D.query();
572
                for (int k = 1; k <= fim; k++) D.rollback();</pre>
9c8
            }
bdf
            i = j;
e99
        return ret;
edf
9d2 }
3.27 MO em Arvores
```

```
// Problema que resolve: https://www.spoj.com/problems/COT2/
//
// Complexidade sendo c = O(update) e SQ = sgrt(n):
// O((n + q) * sqrt(n) * c)
1bc const int MAX = 40010, SQ = 400;
042 vector < int > g[MAX];
c54 namespace LCA { ... }
249 int in[MAX], out[MAX], vtx[2 * MAX];
81b bool on [MAX];
4c3 int dif, freq[MAX];
9e2 vector < int > w;
d9a void dfs(int v, int p, int &t) {
        vtx[t] = v, in[v] = t++;
659
        for (int u : g[v]) if (u != p) {
18e
            dfs(u, v, t);
c53
e0f
       }
        vtx[t] = v, out[v] = t++;
217
42b }
e5f void update(int p) { // faca alteracoes aqui
bbc
        int v = vtx[p];
```

```
0ec
        if (not on[v]) { // insere vtx v
31c
            dif += (freq[w[v]] == 0);
b20
            frea[w[v]]++;
        }
cf7
        else { // retira o vertice v
4e6
0a9
            dif \rightarrow (freq[w[v]] == 1);
fd3
            freq[w[v]]--;
2c8
        }
73e
        on[v] = not on[v];
ea9 }
a3a vector < tuple < int , int , int >> build_queries (const vector < pair < int ,
   int>>& a) {
ea6
        LCA::build(0);
        vector<tuple<int, int, int>> ret;
f77
        for (auto [1, r] : q){
aa9
d24
            if (in[r] < in[l]) swap(l, r);
6f9
            int p = LCA::lca(1, r);
826
            int init = (p == 1) ? in[1] : out[1];
07a
            ret.emplace_back(init, in[r], in[p]);
        }
b0e
edf
        return ret;
8e6 }
f31 vector<int> mo_tree(const vector<pair<int, int>>& vq){
6bb
        int t = 0;
dab
        dfs(0, -1, t);
af1
        auto q = build_queries(vq);
f48
        vector<int> ord(q.size());
be8
        iota(ord.begin(), ord.end(), 0);
        sort(ord.begin(), ord.end(), [&] (int 1, int r) {
d01
            int bl = get<0>(q[1]) / SQ, br = <math>get<0>(q[r]) / SQ;
d8d
596
            if (bl != br) return bl < br;</pre>
158
            else if (b1 % 2 == 1) return get<1>(q[1]) < get<1>(q[r]);
f1d
            else return get<1>(q[1]) > get<1>(q[r]);
0a8
        });
80e
        memset(freq, 0, sizeof freq);
bf6
        dif = 0:
        vector<int> ret(q.size());
ff2
        int 1 = 0, r = -1;
3d9
8b0
        for (int i : ord) {
            auto [ql, qr, qp] = q[i];
3 c 7
            while (r < qr) update(++r);</pre>
af7
```

```
d6b
             while (1 > q1) update(--1);
             while (1 < q1) update(1++);</pre>
951
             while (r > qr) update(r--);
6a1
            if (qp < 1 \text{ or } qp > r)  { // se LCA estah entre as pontas
3d8
74b
                 update(qp);
2e1
                 ret[i] = dif;
74b
                 update(qp);
e83
0fe
             else ret[i] = dif;
Ofd
edf
        return ret;
48d }
```

3.28 Palindromic Factorization

```
// Precisa da eertree
// Computa o numero de formas de particionar cada
// prefixo da string em strings palindromicas
// O(n log n), considerando alfabeto O(1)
070 struct eertree { ... };
0e7 ll factorization(string s) {
        int n = s.size(), sz = 2;
b19
580
        eertree PT(n):
147
        vector \langle int \rangle diff (n+2), slink (n+2), sans (n+2), dp (n+1);
0ec
        dp[0] = 1;
        for (int i = 1; i <= n; i++) {</pre>
78a
c58
            PT.add(s[i-1]);
a7c
            if (PT.size()+2 > sz) {
6c4
                diff[sz] = PT.len[sz] - PT.len[PT.link[sz]];
241
                if (diff[sz] == diff[PT.link[sz]])
d6f
                     slink[sz] = slink[PT.link[sz]];
f53
                else slink[sz] = PT.link[sz];
eb9
                sz++;
            }
f6a
            for (int v = PT.last; PT.len[v] > 0; v = slink[v]) {
911
                sans[v] = dp[i - (PT.len[slink[v]] + diff[v])];
297
                if (diff[v] == diff[PT.link[v]])
85d
f20
                    sans[v] = (sans[v] + sans[PT.link[v]]) % MOD;
071
                dp[i] = (dp[i] + sans[v]) % MOD;
e5e
            }
fc0
5f0
        return dp[n];
3a7 }
```

3.29 Parsing de Expressao

```
// Operacoes associativas a esquerda por default
// Para mudar isso, colocar em r_assoc
// Operacoes com maior prioridade sao feitas primeiro
cc1 bool blank(char c) {
f34
        return c == ' ';
ec3 }
8e4 bool is_unary(char c) {
        return c == '+' or c == '-';
b6b }
76d bool is_op(char c) {
010
        if (is_unary(c)) return true;
        return c == '*' or c == '/' or c == '+' or c == '-':
31c
4e4 }
fa3 bool r_assoc(char op) {
        // operator unario - deve ser assoc. a direita
cf0
        return op < 0;</pre>
c5c }
79d int priority(char op) {
        // operator unario - deve ter precedencia maior
103
        if (op < 0) return INF;</pre>
        if (op == '*' or op == '/') return 2;
727
        if (op == '+' or op == '-') return 1;
439
daa
        return -1;
966 }
c15 void process_op(stack<int>& st, stack<int>& op) {
        char o = op.top(); op.pop();
88 c
91c
        if (o < 0) {
4e6
            o *= -1:
1e2
            int 1 = st.top(); st.pop();
Off
            if (o == '+') st.push(1);
            if (o == '-') st.push(-1);
7e9
9d9
        } else {
14c
            int r = st.top(); st.pop();
1e2
            int 1 = st.top(); st.pop();
1e4
            if (o == '*') st.push(1 * r);
f55
            if (o == '/') st.push(1 / r);
605
            if (o == '+') st.push(1 + r);
            if (o == '-') st.push(1 - r);
c40
```

```
abe
        }
2b2 }
439 int eval(string& s) {
212
        stack<int> st, op;
        bool un = true;
d0c
1 c f
       for (int i = 0: i < s.size(): i++) {
            if (blank(s[i])) continue;
68d
            if (s[i] == '(') {
139
367
                op.push('(');
99d
                un = true:
130
            } else if (s[i] == ')') {
709
                while (op.top() != '(') process_op(st, op);
75e
                op.pop();
                un = false;
ce2
            } else if (is_op(s[i])) {
146
                char o = s[i]:
4d0
                if (un and is_unary(o)) o *= -1;
37 c
                while (op.size() and (
ae3
                             (!r_assoc(o) and priority(op.top()) >=
   priority(o)) or
c41
                             (r_assoc(o) and priority(op.top()) >
   priority(o))))
                    process_op(st, op);
c47
c00
                op.push(o);
99d
                un = true;
949
            } else {
                int val = 0:
c2b
                while (i < s.size() and isalnum(s[i]))</pre>
                    val = val * 10 + s[i++] - '0':
169
                i--;
25d
                st.push(val);
ce2
                un = false:
            }
с7с
b19
        }
7f6
        while (op.size()) process_op(st, op);
123
        return st.top();
05c }
3.30 RMQ com Divide and Conquer
// Responde todas as queries em
// O(n log(n))
f74 typedef pair <pair <int, int>, int> iii;
```

```
7c6 #define f first
Oab #define s second
87d int n, q, v[MAX];
e3f iii qu[MAX];
aeb int ans[MAX], pref[MAX], sulf[MAX];
0e3 void solve(int l=0, int r=n-1, int ql=0, int qr=q-1) {
8a3
        if (1 > r or q1 > qr) return;
        int m = (1+r)/2:
ee4
        int qL = partition(qu+ql, qu+qr+1, [=](iii x){return x.f.s <</pre>
1b1
    m;}) - qu;
        int qR = partition(qu+qL, qu+qr+1, [=](iii x){return x.f.f
eb0
    <=m;}) - qu;
        pref[m] = sulf[m] = v[m]:
3cd
9f9
        for (int i = m-1; i >= 1; i--) pref[i] = min(v[i], pref[i+1]);
        for (int i = m+1; i <= r; i++) sulf[i] = min(v[i], sulf[i-1]);</pre>
ea8
b2a
        for (int i = qL; i < qR; i++)
             ans[qu[i].s] = min(pref[qu[i].f.f], sulf[qu[i].f.s]);
f3a
364
        solve(l, m-1, ql, qL-1), solve(m+1, r, qR, qr);
13e }
3.31 Segment Intersection
// Verifica, dado n segmentos, se existe algum par de segmentos
// que se intersecta
//
// O(n log n)
6e0 bool operator < (const line& a, const line& b) { // comparador pro
191
        if (a.p == b.p) return ccw(a.p, a.q, b.q);
        if (!eq(a.p.x, a.q.x)) and (eq(b.p.x, b.q.x)) or a.p.x+eps <
    b.p.x))
780
            return ccw(a.p, a.q, b.p);
        return ccw(a.p, b.q, b.p);
dc0
e36 }
8e2 bool has intersection(vector<line> v) {
        auto intersects = [&](pair<line, int> a, pair<line, int> b) {
576
a08
            return interseg(a.first, b.first);
3e6
        }:
e1b
        vector<pair<pt, pair<int, int>>> w;
f 14
        for (int i = 0; i < v.size(); i++) {</pre>
```

```
876
            if (v[i].q < v[i].p) swap(v[i].p, v[i].q);</pre>
            w.push_back({v[i].p, {0, i}});
e1d
034
            w.push_back({v[i].q, {1, i}});
220
        sort(w.begin(), w.end());
d1d
7f2
        set < pair < line , int >> se;
e58
       for (auto i : w) {
bfd
            line at = v[i.second.second];
292
            if (i.second.first == 0) {
145
                auto nxt = se.lower_bound({at, i.second.second});
                if (nxt != se.end() and intersects(*nxt, {at,
   i.second.second})) return 1;
257
                if (nxt != se.begin() and intersects(*(--nxt), {at,
   i.second.second})) return 1;
78f
                se.insert({at, i.second.second});
9d9
            } else {
                auto nxt = se.upper_bound({at, i.second.second}), cur
884
   = nxt, prev = --cur;
                if (nxt != se.end() and prev != se.begin()
b64
                    and intersects(*nxt, *(--prev))) return 1;
4fb
                se.erase(cur):
cca
            }
925
a00
       }
bb3
        return 0;
196 }
```

3.32 Sequencia de de Brujin

```
// Se passar sem o terceiro parametro, gera um vetor com valores
// em [0, k) de tamanho k^n de forma que todos os subarrays ciclicos
// de tamanho n ocorrem exatamente uma vez
// Se passar com um limite lim, gera o menor vetor com valores
// em [0, k) que possui lim subarrays de tamanho n distintos
// (assume que lim <= k^n)</pre>
//
// Linear no tamanho da resposta
860 vector<int> de_brujin(int n, int k, int lim = INF) {
        if (k == 1) return vector < int > (lim == INF ? 1 : n, 0);
b55
        vector<int> 1 = {0}, ret; // 1 eh lyndon word
5f6
        while (true) {
667
c86
            if (1.size() == 0) {
1 b 9
                if (lim == INF) break;
daf
                1.push_back(0);
            }
bae
686
            if (n % 1.size() == 0) for (int i : 1) {
728
                ret.push_back(i);
```

```
c99
                if (ret.size() == n+lim-1) return ret;
            }
56e
630
            int p = 1.size();
905
            while (1.size() < n) 1.push_back(1[1.size()%p]);</pre>
            while (l.size() and l.back() == k-1) l.pop_back();
e7f
88a
            if (1.size()) 1.back()++;
2ef
        }
edf
        return ret;
197 }
```

3.33 Shortest Addition Chain

```
// Computa o menor numero de adicoes para construir
// cada valor, comecando com 1 (e podendo salvar variaveis)
// Retorna um par com a dp e o pai na arvore
// A arvore eh tao que o taminho da raiz (1) ate x
// contem os valores que devem ser criados para gerar x
// A profundidade de x na arvore eh dp[x]
// DP funciona para ateh 300, mas a arvore soh funciona
// para ateh 148
// recuperacao certa soh ateh 148 (erra para 149, 233, 298)
3de pair < vector < int > , vector < int >> addition_chain() {
16f
        int MAX = 301:
875
        vector < int > dp(MAX), p(MAX);
1ab
        for (int n = 2; n < MAX; n++) {
7c0
            pair < int , int > val = {INF , -1};
212
            for (int i = 1; i < n; i++) for (int j = i; j; j = p[j])
94a
                if (j == n-i) val = min(val, pair(dp[i]+1, i));
eb3
            tie(dp[n], p[n]) = val;
efe
            if (n == 9) p[n] = 8;
ba1
            if (n == 149 \text{ or } n == 233) dp[n]--;
bcd
717
        return {dp, p};
84f }
```

3.34 Simple Polygon

```
// Verifica se um poligono com n pontos eh simples
//
// O(n log n)

6e0 bool operator < (const line& a, const line& b) { // comparador pro
    sweepline
191    if (a.p == b.p) return ccw(a.p, a.q, b.q);
231    if (!eq(a.p.x, a.q.x) and (eq(b.p.x, b.q.x) or a.p.x+eps <
    b.p.x))</pre>
```

```
780
            return ccw(a.p, a.q, b.p);
        return ccw(a.p, b.q, b.p);
dc0
e36 }
6f3 bool simple(vector<pt> v) {
        auto intersects = [&](pair<line, int> a, pair<line, int> b) {
576
e72
            if ((a.second+1)%v.size() == b.second or
80e
                (b.second+1)%v.size() == a.second) return false;
a08
            return interseg(a.first, b.first);
1 c.5
        }:
41a
        vector<line> seg;
e1b
        vector<pair<pt, pair<int, int>>> w;
f14
        for (int i = 0: i < v.size(): i++) {</pre>
            pt at = v[i], nxt = v[(i+1)%v.size()];
0a8
828
            if (nxt < at) swap(at, nxt);</pre>
937
            seg.push_back(line(at, nxt));
f7e
            w.push_back({at, {0, i}});
            w.push_back({nxt, {1, i}});
69 c
            // casos degenerados estranhos
            if (isinseg(v[(i+2)%v.size()], line(at, nxt))) return 0;
ae8
            if (isinseg(v[(i+v.size()-1)%v.size()], line(at, nxt)))
688
   return 0;
cba
        }
d1d
        sort(w.begin(), w.end());
7f2
        set < pair < line , int >> se;
e58
        for (auto i : w) {
ff8
            line at = seg[i.second.second];
292
            if (i.second.first == 0) {
145
                auto nxt = se.lower_bound({at, i.second.second});
7c4
                if (nxt != se.end() and intersects(*nxt, {at,
   i.second.second})) return 0;
                if (nxt != se.begin() and intersects(*(--nxt), {at,
b34
   i.second.second})) return 0:
78f
                se.insert({at. i.second.second}):
            } else {
9d9
884
                auto nxt = se.upper_bound({at, i.second.second}), cur
   = nxt, prev = --cur;
                if (nxt != se.end() and prev != se.begin()
b64
403
                    and intersects(*nxt, *(--prev))) return 0;
cca
                se.erase(cur);
            }
ad0
        }
d17
6a5
        return 1;
af3 }
```

3.35 Sweep Direction

```
// Passa por todas as ordenacoes dos pontos definitas por "direcoes"
// Assume que nao existem pontos coincidentes
//
// O(n^2 \log n)
4b8 void sweep_direction(vector<pt> v) {
3d2
        int n = v.size():
163
        sort(v.begin(), v.end(), [](pt a, pt b) {
3a5
            if (a.x != b.x) return a.x < b.x;
572
            return a.y > b.y;
79a
        });
b89
        vector < int > at(n);
516
        iota(at.begin(), at.end(), 0);
b79
        vector<pair<int, int>> swapp;
        for (int i = 0; i < n; i++) for (int j = i+1; j < n; j++)
25e
95f
            swapp.push_back({i, j}), swapp.push_back({j, i});
        sort(swapp.begin(), swapp.end(), [&](auto a, auto b) {
269
            pt A = rotate90(v[a.first] - v[a.second]);
134
247
            pt B = rotate90(v[b.first] - v[b.second]);
            if (quad(A) == quad(B) and !sarea2(pt(0, 0), A, B)) return
615
   a < b;
224
            return compare_angle(A, B);
        });
5e7
        for (auto par : swapp) {
4e6
e24
            assert(abs(at[par.first] - at[par.second]) == 1);
            int 1 = min(at[par.first], at[par.second]),
a96
                r = n-1 - max(at[par.first], at[par.second]);
0.43
            // l e r sao quantos caras tem de cada lado do par de
                pontos
            // (cada par eh visitado duas vezes)
9cf
            swap(v[at[par.first]], v[at[par.second]]);
            swap(at[par.first], at[par.second]);
1c0
        }
241
6bb }
```

3.36 Triangulacao de Delaunay

```
// Computa a triangulacao de Delaunay, o dual
// do diagrama de Voronoi (a menos de casos degenerados)
// Retorna um grafo indexado pelos indices dos pontos, e as arestas
// sao as arestas da triangulacao
// As arestas partindo de um vertice ja vem ordenadas por angulo,
// ou seja, se o vertice v nao esta no convex hull, (v, v_i, v_{i+1})
// eh um triangulo da triangulacao, em que v_i eh o i-esimo vizinho
// Usa o alg d&c, precisa representar MAX_COOR^4, por isso __int128
// pra aguentar valores ateh 1e9
```

```
//
// Propriedades:
// 1 - 0 grafo tem no max 3n-6 arestas
// 2 - Para todo triangulo, a circunf. que passa pelos 3 pontos
      nao contem estritamente nenhum ponto
// 3 - A MST euclidiana eh subgrafo desse grafo
// 4 - Cada ponto eh vizinho do ponto mais proximo dele
//
// O(n log n)
2ad typedef struct QuadEdge* Q;
ba5 struct QuadEdge {
53e
        int id:
114
        pt o;
41e
       Q rot, nxt;
       bool used;
3e5
3fc
        QuadEdge(int id_ = -1, pt o_ = pt(INF, INF)) :
            id(id_), o(o_), rot(nullptr), nxt(nullptr), used(false) {}
4ba
        Q rev() const { return rot->rot; }
00f
        Q next() const { return nxt; }
сЗс
        Q prev() const { return rot->next()->rot; }
188
0d4
        pt dest() const { return rev()->o; }
828 }:
91b Q edge(pt from, pt to, int id_from, int id_to) {
сбе
        Q e1 = new QuadEdge(id_from, from);
        Q e2 = new QuadEdge(id_to, to);
61b
8f6
        Q e3 = new QuadEdge;
       Q e4 = new QuadEdge;
5ca
e69
        tie(e1->rot, e2->rot, e3->rot, e4->rot) = \{e3, e4, e2, e1\};
        tie(e1->nxt, e2->nxt, e3->nxt, e4->nxt) = \{e1, e2, e4, e3\};
f22
1ad
        return e1:
c70 }
d8d void splice(Q a, Q b) {
a6f
        swap(a->nxt->rot->nxt, b->nxt->rot->nxt);
da4
        swap(a->nxt, b->nxt);
a58 }
167 void del_edge(Q& e, Q ne) { // delete e and assign e <- ne
cc0
        splice(e, e->prev());
        splice(e->rev(), e->rev()->prev());
eec
7ea
        delete e->rev()->rot, delete e->rev();
        delete e->rot; delete e;
524
6b2
        e = ne:
```

```
18b }
d08 Q conn(Q a, Q b) {
        Q = edge(a->dest(), b->o, a->rev()->id, b->id);
        splice(e, a->rev()->prev());
f2b
d37
        splice(e->rev(), b);
6bf
        return e:
f78 }
d64 bool in_c(pt a, pt b, pt c, pt p) { // p ta na circunf. (a, b, c) ?
268
        \_int128 p2 = p*p, A = a*a - p2, B = b*b - p2, C = c*c - p2;
        return sarea2(p, a, b) * C + sarea2(p, b, c) * A + sarea2(p,
   c. a) * B > 0:
b54 }
540 pair < Q, Q > build_tr(vector < pt > & p, int 1, int r) {
        if (r-1+1 \le 3) {
            Q = edge(p[1], p[1+1], 1, 1+1), b = edge(p[1+1], p[r],
2eb
   1+1, r);
            if (r-1+1 == 2) return {a, a->rev()};
912
0ec
            splice(a->rev(), b);
сЗс
            ll ar = sarea2(p[1], p[1+1], p[r]);
            Q c = ar ? conn(b, a) : 0;
1af
021
            if (ar >= 0) return \{a, b > rev()\};
9db
            return {c->rev(), c};
bce
        }
ee4
        int m = (1+r)/2;
328
        auto [la, ra] = build_tr(p, l, m);
        auto [lb, rb] = build_tr(p, m+1, r);
b93
667
        while (true) {
b99
            if (ccw(lb->o, ra->o, ra->dest())) ra = ra->rev()->prev();
458
            else if (ccw(lb->o, ra->o, lb->dest())) lb =
   lb->rev()->next();
f97
            else break:
24a
ca5
        Q b = conn(lb->rev(), ra);
        auto valid = [&](Q e) { return ccw(e->dest(), b->dest(),
713
   b->o); };
        if (ra->o == la->o) la = b->rev();
ee1
63f
        if (1b->o == rb->o) rb = b;
        while (true) {
667
            Q L = b - > rev() - > next();
71e
            if (valid(L)) while (in_c(b->dest(), b->o, L->dest(),
d11
   L->next()->dest()))
                del_edge(L, L->next());
1c0
            Q R = b - > prev();
c76
2b0
            if (valid(R)) while (in_c(b->dest(), b->o, R->dest(),
```

```
R->prev()->dest()))
541
                del_edge(R, R->prev());
            if (!valid(L) and !valid(R)) break;
a3a
            if (!valid(L) or (valid(R) and in_c(L->dest(), L->o, R->o,
   R->dest())))
36c
                b = conn(R, b \rightarrow rev());
666
            else b = conn(b->rev(), L->rev());
94d
a2b
        return {la, rb};
689 }
b58 vector < vector < int >> delaunay (vector < pt > v) {
3d2
        int n = v.size():
397
        auto tmp = v;
135
        vector < int > idx(n);
        iota(idx.begin(), idx.end(), 0);
295
        sort(idx.begin(), idx.end(), [&](int 1, int r) { return v[1] <</pre>
   v[r]; });
        for (int i = 0; i < n; i++) v[i] = tmp[idx[i]];</pre>
5d8
        assert(unique(v.begin(), v.end()) == v.end());
780
4aa
        vector < vector < int >> g(n);
        bool col = true:
4ec
        for (int i = 2; i < n; i++) if (sarea2(v[i], v[i-1], v[i-2]))
   col = false;
        if (col) {
bf5
aa4
            for (int i = 1; i < n; i++)</pre>
                g[idx[i-1]].push_back(idx[i]),
   g[idx[i]].push_back(idx[i-1]);
96b
            return g;
0ae
d36
        Q e = build_tr(v, 0, n-1).first;
113
        vector <Q> edg = {e};
5d1
        for (int i = 0; i < edg.size(); e = edg[i++]) {</pre>
3ed
            for (Q at = e: !at->used: at = at->next()) {
                at->used = true;
60d
cf8
                g[idx[at->id]].push_back(idx[at->rev()->id]);
                edg.push_back(at->rev());
15d
9f2
            }
d19
        }
96b
        return g;
b43 }
      Triangulos em Grafos
// get_triangles(i) encontra todos os triangulos ijk no grafo
// Custo nas arestas
// retorna {custo do triangulo, {j, k}}
```

```
// O(m sqrt(m) log(n)) se chamar para todos os vertices
c0d vector<pair<int, int>> g[MAX]; // {para, peso}
d41 #warning o 'g' deve estar ordenado
9a5 vector<pair<int, pair<int, int>>> get_triangles(int i) {
771
        vector<pair<int, pair<int, int>>> tri;
b23
        for (pair<int, int> j : g[i]) {
            int a = i, b = j.first;
2b3
6dd
            if (g[a].size() > g[b].size()) swap(a, b);
            for (pair<int, int> c : g[a]) if (c.first != b and c.first
   > i.first) {
525
                auto it = lower_bound(g[b].begin(), g[b].end(),
   make_pair(c.first, -INF));
                if (it == g[b].end() or it->first != c.first) continue;
f55
                tri.push_back({j.second+c.second+it->second, {a == i ?
0aa
   b : a, c.first}});
            }
b5e
7e1
f5e
        return tri;
036 }
```

4 Matematica

4.1 2-SAT

```
// solve() retorna um par, o first fala se eh possivel
// atribuir, o second fala se cada variavel eh verdadeira
// O(|V|+|E|) = O(\#variaveis + \#restricoes)
138 struct sat {
e6c
        int n, tot;
789
        vector < vector < int >> g;
0ca
        vector<int> vis, comp, id, ans;
4ce
        stack<int> s;
141
        sat() {}
        sat(int n_{-}) : n(n_{-}), tot(n), g(2*n) {}
172
f32
        int dfs(int i, int& t) {
cf0
             int lo = id[i] = t++;
efc
             s.push(i), vis[i] = 2;
48e
             for (int j : g[i]) {
740
                 if (!vis[j]) lo = min(lo, dfs(j, t));
```

```
994
                 else if (vis[j] == 2) lo = min(lo, id[j]);
            }
d64
3de
            if (lo == id[i]) while (1) {
                 int u = s.top(); s.pop();
3 c 3
                 vis[u] = 1, comp[u] = i;
9c5
                 if ((u>1) < n \text{ and } ans[u>1] == -1) ans[u>1] = <math>\sim u\&1;
91d
2ef
                 if (u == i) break;
            }
60d
253
             return lo;
dec
74a
        void add_impl(int x, int y) { // x -> y = !x ou y
26a
            x = x >= 0 ? 2*x : -2*x-1:
2b8
            y = y >= 0 ? 2*y : -2*y-1;
a1e
            g[x].push_back(y);
            g[y^1].push_back(x^1);
1e2
ef0
        void add_cl(int x, int y) { // x ou y
e85
0b5
             add_impl(\sim x, v);
254
        }
487
        void add_xor(int x, int y) { // x xor y
0b7
             add_cl(x, y), add_cl(\sim x, \sim y);
        }
9a1
978
        void add_eq(int x, int y) { // x = y
c86
             add_xor(\simx, y);
b91
b10
        void add_true(int x) { // x = T
18b
             add_impl(\sim x, x);
9e2
d14
        void at_most_one(vector<int> v) { // no max um verdadeiro
             g.resize(2*(tot+v.size()));
54d
f14
            for (int i = 0; i < v.size(); i++) {</pre>
                 add_impl(tot+i, ~v[i]);
8c9
a8f
                 if (i) {
                     add_impl(tot+i, tot+i-1);
b6a
3d3
                     add_impl(v[i], tot+i-1);
                 }
0f7
084
            }
258
             tot += v.size();
b00
        }
        pair < bool, vector < int >> solve() {
a8e
27b
             ans = vector < int > (n, -1);
6bb
            int t = 0:
0de
            vis = comp = id = vector \langle int \rangle (2*tot, 0);
            for (int i = 0; i < 2*tot; i++) if (!vis[i]) dfs(i, t);</pre>
53c
            for (int i = 0: i < tot: i++)</pre>
f88
```

```
4c9
                if (comp[2*i] == comp[2*i+1]) return {false, {}};
997
            return {true, ans};
7b3
        }
ef6 };
    Avaliação de Interpolação
// Dado 'n' pontos (i, y[i]), i \in [0, n),
// avalia o polinomio de grau n-1 que passa
// por esses pontos em 'x'
// Tudo modular, precisa do mint
//
// O(n)
ee8 mint evaluate_interpolation(int x, vector<mint> y) {
80e
        int n = y.size();
        vector < mint > sulf(n+1, 1), fat(n, 1), ifat(n);
184
6fa
        for (int i = n-1: i >= 0: i--) sulf[i] = sulf[i+1] * (x - i):
29b
        for (int i = 1; i < n; i++) fat[i] = fat[i-1] * i;</pre>
0da
        ifat[n-1] = 1/fat[n-1];
        for (int i = n-2; i >= 0; i--) ifat[i] = ifat[i+1] * (i + 1);
3db
        mint pref = 1. ans = 0:
ca1
5ea
        for (int i = 0; i < n; pref *= (x - i++)) {</pre>
42f
            mint num = pref * sulf[i+1];
b4e
            mint den = ifat[i] * ifat[n-1 - i];
0 bd
            if ((n-1 - i)\%2) den *= -1;
03f
            ans += y[i] * num * den;
        }
ce6
ba7
        return ans;
4fe }
4.3 Berlekamp-Massey
// guess_kth(s, k) chuta o k-esimo (0-based) termo
// de uma recorrencia linear que gera s
// Para uma rec. lin. de ordem x, se passar 2x termos
// vai gerar a certa
// Usar aritmetica modular
// O(n^2 \log k), em que n = |s|
b7c template < typename T > T evaluate (vector < T > c, vector < T > s, ll k) {
ff2
        int n = c.size();
```

```
9ee
        assert(c.size() <= s.size());
d09
        auto mul = [&](const vector<T> &a, const vector<T> &b) {
564
            vector <T> ret(a.size() + b.size() - 1):
d75
            for (int i = 0; i < a.size(); i++) for (int j = 0; j <</pre>
   b.size(); j++)
cff
                ret[i+j] += a[i] * b[j];
            for (int i = ret.size()-1; i >= n; i--) for (int i = n-1;
   j >= 0; j--)
                ret[i-j-1] += ret[i] * c[j];
            ret.resize(min<int>(ret.size(), n));
16d
edf
            return ret:
3b9
        }:
        vector < T > a = n == 1 ? vector < T > ({c[0]}) : vector < T > ({0, 1}),
1a6
   x = \{1\};
        while (k) {
95f
            if (k\&1) x = mul(x, a):
7f1
            a = mul(a, a), k >>= 1;
b28
8ea
       }
dd6
        x.resize(n);
ce8
       T ret = 0:
e72
        for (int i = 0; i < n; i++) ret += x[i] * s[i];</pre>
edf
        return ret:
7e2 }
192 template < typename T > vector < T > berlekamp_massey(vector < T > s) {
        int n = s.size(), l = 0, m = 1;
222
        vector < T > b(n), c(n);
        T ld = b[0] = c[0] = 1;
46e
620
        for (int i = 0; i < n; i++, m++) {</pre>
793
            T d = s[i]:
            for (int j = 1; j <= 1; j++) d += c[j] * s[i-j];
ab6
            if (d == 0) continue;
5f0
8b4
            vector <T> temp = c;
369
            T coef = d / ld;
ba6
            for (int j = m; j < n; j++) c[j] -= coef * b[j-m];
88f
            if (2 * 1 \le i) 1 = i + 1 - 1, b = temp, 1d = d, m = 0;
76a
        c.resize(1 + 1):
90c
        c.erase(c.begin());
844
0dc
        for (T\& x : c) x = -x;
807
        return c:
4d9 }
2cf template < typename T > T guess_kth(const vector < T > & s, ll k) {
```

```
cc3
         auto c = berlekamp_massev(s);
         return evaluate(c, s, k);
96a
697 }
4.4 Binomial Distribution
// binom(n, k, p) retorna a probabilidade de k sucessos
// numa binomial(n, p)
361 double logfact[MAX];
9e4 void calc() {
7a0
        logfact[0] = 0;
        for (int i = 1; i < MAX; i++) logfact[i] = logfact[i-1] +</pre>
   log(i);
67a }
94c double binom(int n, int k, double p) {
         return exp(logfact[n] - logfact[k] - logfact[n-k] + k * log(p)
    + (n-k) * log(1 - p));
587 }
4.5 Convolucao de GCD / LCM
// O(n log(n))
// multiple_transform(a)[i] = \sum_d a[d * i]
bbe template < typename T > void multiple_transform (vector < T > & v, bool
    inv = false) {
        vector < int > I(v.size()-1);
64a
847
        iota(I.begin(), I.end(), 1);
674
        if (inv) reverse(I.begin(), I.end());
        for (int i : I) for (int j = 2; i*j < v.size(); j++)
dad
            v[i] += (inv ? -1 : 1) * v[i*j];
a8a
338 }
// gcd_convolution(a, b)[k] = \sum_{gcd(i, j) = k} a_i * b_j
fe2 template < typename T > vector < T > gcd_convolution(vector < T > a,
    vector<T> b) {
        multiple_transform(a), multiple_transform(b);
bdf
799
        for (int i = 0: i < a.size(): i++) a[i] *= b[i]:
        multiple_transform(a, true);
3f5
        return a;
984 }
// divisor_transform(a)[i] = \sum_{d|i} a[i/d]
```

```
be7 template < typename T > void divisor_transform(vector < T > & v, bool inv
                                                                             c91
                                                                                                  if (x == 1) {
   = false) {
                                                                             fac
                                                                                                      j = INF;
        vector < int > I(v.size()-1);
                                                                             5e2
64a
                                                                                                      continue;
        iota(I.begin(), I.end(), 1);
                                                                             6e0
                                                                                                  }
        if (!inv) reverse(I.begin(), I.end());
5ea
                                                                             544
                                                                                                 T \& v = basis[i]:
        for (int i : I) for (int j = 2; i*j < v.size(); j++)
                                                                                                  T g = gcd(x, y);
dad
                                                                             3c9
14f
            v[i*j] += (inv ? -1 : 1) * v[i];
                                                                             e10
                                                                                                  if (g == 1) continue;
aa7 }
                                                                             15b
                                                                                                  y /= g, x /= g;
                                                                             8c6
                                                                                                  basis.push_back(g);
                                                                                             }
// lcm_convolution(a, b)[k] = \sum_{lcm(i, j)} = k} a_i * b_j
                                                                             069
b1b template < typename T > vector < T > lcm_convolution (vector < T > a,
                                                                             422
   vector<T> b) {
                                                                                         basis.erase(remove(basis.begin(), basis.end(), 1),
                                                                             fe8
3af
        divisor transform(a), divisor transform(b);
                                                                                basis.end()):
799
        for (int i = 0; i < a.size(); i++) a[i] *= b[i];</pre>
                                                                             1a5
                                                                                     }
d8f
        divisor_transform(a, true);
3f5
        return a;
                                                                             4ba
                                                                                     vector<int> factor(T x) {
f5a }
                                                                             21d
                                                                                         vector < int > fat(basis.size());
                                                                             6fd
                                                                                         for (int i = 0; i < basis.size(); i++) {</pre>
                                                                             25 c
                                                                                              while (x \% basis[i] == 0) x /= basis[i], fat[i]++;
4.6 Coprime Basis
                                                                             8de
                                                                                         }
                                                                             6a7
                                                                                         return fat;
// Dado um conjunto de elementos A constroi uma base B
                                                                             b5d
                                                                                     }
// de fatores coprimos tal que todo elemento A[i]
                                                                             671 };
// pode ser fatorado como A[i] = \prod B[j]^p_ij
                                                                             4.7 Crivo de Eratosthenes
// Sendo n o numero de inserts, a complexidade esperada fica
// O(n*(n*loglog(MAX) + log(MAX)^2))
                                                                             // "O" crivo
                                                                             //
// No pior caso, podemos trocar n*loglog(MAX) por
// se MAX <= 1e6 fica 8*n
                                                                             // Encontra maior divisor primo
// se MAX <= 1e9 fica 10*n
                                                                             // Um numero eh primo sse divi[x] == x
// se MAX <= 1e18 fica 16*n
                                                                             // fact fatora um numero <= lim
// se MAX <= 1e36 fica 26*n
                                                                             // A fatoracao sai ordenada
ebc template <typename T> struct coprime_basis {
                                                                             // crivo - O(n log(log(n)))
a00
        vector <T> basis;
                                                                             // fact - O(log(n))
60e
        coprime_basis() {}
                                                                             f12 int divi[MAX];
        coprime_basis(vector<T> v) { for (T i : v) insert(i); }
055
                                                                             fb9 void crivo(int lim) {
                                                                                     for (int i = 1; i <= lim; i++) divi[i] = 1;</pre>
845
        void insert(T z) {
                                                                             f53
сЗс
            int n = basis.size();
efe
            basis.push_back(z);
                                                                             d46
                                                                                     for (int i = 2; i <= lim; i++) if (divi[i] == 1)</pre>
                                                                                         for (int j = i; j <= lim; j += i) divi[j] = i;</pre>
43 c
            for (int i = n; i < basis.size(); i++) {</pre>
                                                                             018
                for (int j = (i != n) ? i+1 : 0; j < basis.size();</pre>
                                                                             349 }
21 c
   j++) {
```

470 void fact(vector<int>& v, int n) {

if (n != divi[n]) fact(v, n/divi[n]);

4ce

024

if (i == j) continue;

T &x = basis[i];

```
ab4
        v.push_back(divi[n]);
1db }
// Crivo linear
// Mesma coisa que o de cima, mas tambem
// calcula a lista de primos
//
// O(n)
f12 int divi[MAX];
fd3 vector<int> primes;
fb9 void crivo(int lim) {
d5a
       divi[1] = 1;
f70
       for (int i = 2; i <= lim; i++) {</pre>
3eb
            if (divi[i] == 0) divi[i] = i, primes.push_back(i);
3ba
            for (int j : primes) {
522
                if (j > divi[i] or i*j > lim) break;
00Ъ
                divi[i*j] = j;
            }
491
85a
        }
519 }
// Crivo de divisores
// Encontra numero de divisores
// ou soma dos divisores
// O(n log(n))
f12 int divi[MAX];
fb9 void crivo(int lim) {
        for (int i = 1; i <= lim; i++) divi[i] = 1;</pre>
        for (int i = 2; i <= lim; i++)</pre>
424
594
            for (int j = i; j <= lim; j += i) {</pre>
                // para numero de divisores
9e0
                divi[j]++;
                // para soma dos divisores
278
                divi[j] += i;
c58
            }
fc1 }
// Crivo de totiente
//
```

```
// Encontra o valor da funcao
// totiente de Euler
//
// O(n log(log(n)))
5f4 int tot[MAX];
fb9 void crivo(int lim) {
a27
        for (int i = 1; i <= lim; i++) {</pre>
bc9
            tot[i] += i;
feb
            for (int j = 2*i; j <= lim; j += i)</pre>
837
                tot[i] -= tot[i];
678
        }
212 }
// Crivo de função de mobius
// O(n log(log(n)))
4e1 char meb[MAX];
fb9 void crivo(int lim) {
        for (int i = 2; i <= lim; i++) meb[i] = 2;</pre>
        meb[1] = 1;
ace
842
        for (int i = 2; i <= lim; i++) if (meb[i] == 2)
8d8
            for (int j = i; j <= lim; j += i) if (meb[j]) {</pre>
686
                if (meb[j] == 2) meb[j] = 1;
ae1
                meb[j] *= j/i\%i ? -1 : 0;
97f
            }
9bc }
// Crivo linear de funcao multiplicativa
// Computa f(i) para todo 1 <= i <= n, sendo f
// uma funcao multiplicativa (se gcd(a,b) = 1,
// entao f(a*b) = f(a)*f(b)
// f_prime tem que computar f de um primo, e
// add_prime tem que computar f(p^(k+1)) dado f(p^k) e p
// Se quiser computar f(p^k) dado p \in k, usar os comentarios
//
// O(n)
fd3 vector<int> primes;
623 int f[MAX], pot[MAX];
//int expo[MAX];
5c4 void sieve(int lim) {
```

```
// Funcoes para soma dos divisores:
fc9
        auto f_prime = [](int p) { return p+1; };
31c
        auto add_prime = [](int fpak, int p) { return fpak*p+1; };
        //auto f_pak = [](int p, int k) {};
        f[1] = 1;
02d
f70
        for (int i = 2: i <= lim: i++) {</pre>
            if (!pot[i]) {
e6b
e74
                primes.push_back(i);
f05
                f[i] = f_prime(i), pot[i] = i;
                //\expo[i] = 1;
b71
            }
3b9
            for (int p : primes) {
b9f
                if (i*p > lim) break;
569
                if (i%p == 0) {
                    f[i*p] = f[i / pot[i]] * add_prime(f[pot[i]], p);
b97
                    // se for descomentar, tirar a linha de cima também
                    //f[i*p] = f[i / pot[i]] * f_pak(p, expo[i]+1);
                    //\exp [i*p] = \exp [i]+1;
51f
                    pot[i*p] = pot[i] * p;
c2b
                    break;
                } else {
9d9
9ef
                    f[i*p] = f[i] * f[p];
                    pot[i*p] = p;
638
                    //\expo[i*p] = 1;
e92
                }
f31
            }
1bb
        }
350 }
4.8 Deteccao de ciclo - Tortoise and Hare
// Linear no tanto que tem que andar pra ciclar,
// O(1) de memoria
```

```
// Retorna um par com o tanto que tem que andar
// do f0 ate o inicio do ciclo e o tam do ciclo
58d pair<11, 11> find_cycle() {
        11 \text{ tort} = f(f0);
273
b2b
        ll hare = f(f(f0));
        11 t = 0:
b1b
683
        while (tort != hare) {
            tort = f(tort);
b4d
4b2
            hare = f(f(hare));
c82
            t++;
93d
        }
0e8
        11 st = 0;
```

```
909
        tort = f0;
683
        while (tort != hare) {
b4d
            tort = f(tort);
1a2
            hare = f(hare);
397
            st++:
        }
c91
73d
        ll len = 1;
3cd
        hare = f(tort);
683
        while (tort != hare) {
1a2
            hare = f(hare);
040
            len++:
f1a
        }
ebd
        return {st, len};
899 }
```

4.9 Division Trick

4.10 Equação Diofantina Linear

```
// Encontra o numero de solucoes de a*x + b*y = c,
// em que x \in [lx, rx] e y \in [ly, ry]
// Usar o comentario para recuperar as solucoes
// (note que o b ao final eh b/gcd(a, b))
// Cuidado com overflow! Tem que caber o quadrado dos valores
//
// O(log(min(a, b)))
c5e template < typename T > tuple < ll, T, T > ext_gcd(ll a, ll b) {
3bd
        if (!a) return {b, 0, 1};
c4b
        auto [g, x, y] = ext_gcd < T > (b%a, a);
        return \{g, y - b/a*x, x\};
c59
8a8 }
// numero de solucoes de a*[lx, rx] + b*[ly, ry] = c
14c template < typename T = 11> // usar __int128 se for ate 1e18
2a4 ll diophantine(ll a, ll b, ll c, ll lx, ll rx, ll ly, ll ry) {
c80
        if (lx > rx or ly > ry) return 0;
a98
        if (a == 0 \text{ and } b == 0) \text{ return } c ? 0 : (rx-1x+1)*(ry-1y+1);
```

```
8ce
        auto [g, x, y] = ext_gcd < T > (abs(a), abs(b));
        if (c % g != 0) return 0;
9c3
249
        if (a == 0) return (rx-lx+1)*(ly <= c/b and c/b <= ry);
        if (b == 0) return (ry-ly+1)*(lx <= c/a \text{ and } c/a <= rx);
4ce
        x *= a/abs(a) * c/g, y *= b/abs(b) * c/g, a /= g, b /= g;
fb1
b20
        auto shift = [\&](T qt) \{ x += qt*b, y -= qt*a; \};
        auto test = [&](T& k, ll mi, ll ma, ll coef, int t) {
efa
            shift((mi - k)*t / coef);
866
            if (k < mi) shift(coef > 0 ? t : -t);
79d
74d
            if (k > ma) return pair \langle T, T \rangle (rx+2, rx+1);
41f
            T x1 = x:
633
            shift((ma - k)*t / coef);
            if (k > ma) shift(coef > 0 ? -t : t);
c5b
4a9
           return pair<T, T>(x1, x);
8e1
        };
639
        auto [11, r1] = test(x, 1x, rx, b, 1);
        auto [12, r2] = test(v, lv, rv, a, -1);
38e
c43
        if (12 > r2) swap(12, r2);
        T l = max(11, 12), r = min(r1, r2);
50a
        if (1 > r) return 0;
339
        11 k = (r-1) / abs(b) + 1;
42f
839
        return k; // solucoes: x = 1 + [0, k)*|b|
98e }
4.11 Euclides estendido
// Acha x e y tal que ax + by = mdc(a, b) (nao eh unico)
// Assume a, b >= 0
//
// O(log(min(a, b)))
2be tuple < 11, 11, 11 > ext_gcd(11 a, 11 b) {
       if (!a) return {b, 0, 1};
3bd
550
        auto [g, x, y] = ext_gcd(b\%a, a);
c59
        return \{g, y - b/a*x, x\};
354 }
4.12 Exponenciacao rapida
// (x^y mod m) em O(log(y))
03c ll pow(ll x, ll y, ll m) \{ // \text{ iterativo} \}
        ll ret = 1:
c85
```

1b8

895

while (y) {

if (y & 1) ret = (ret * x) % m;

```
cc5
            x = (x * x) % m;
020
edf
        return ret;
12b }
03c ll pow(ll x, ll y, ll m) { // recursivo
        if (!v) return 1;
13a
426
        11 ans = pow(x*x\%m, y/2, m);
688
        return y%2 ? x*ans%m : ans;
7d4 }
4.13 Fast Walsh Hadamard Transform
// FWHT<'|'>(f) eh SOS DP
// FWHT<'&'>(f) eh soma de superset DP
// Se chamar com ^, usar tamanho potencia de 2!!
// O(n log(n))
382 template < char op, class T > vector < T > FWHT (vector < T > f, bool inv =
   false) {
b75
        int n = f.size():
        for (int k = 0; (n-1) >> k; k++) for (int i = 0; i < n; i++) if
    (i >> k&1) {
29e
            int i = i^{(1 << k)}:
             if (op == '\^') f[j] += f[i], f[i] = f[j] - 2*f[i];
627
             if (op == '|') f[i] += (inv ? -1 : 1) * f[i];
a38
93c
             if (op == '&') f[j] += (inv ? -1 : 1) * f[i];
1bb
578
        if (op == '^' and inv) for (auto& i : f) i /= n;
         return f:
abe
50e }
4.14 FFT
// Chamar convolution com vector < complex < double >> para FFT
// Precisa do mint para NTT
//
// O(n log(n))
// Para FFT
488 void get_roots(bool f, int n, vector<complex<double>>& roots) {
f26
         const static double PI = acosl(-1);
71a
        for (int i = 0; i < n/2; i++) {
b1e
             double alpha = i*((2*PI)/n);
1a1
            if (f) alpha = -alpha;
```

23b

v >>= 1;

```
069
            roots[i] = {cos(alpha), sin(alpha)};
804
       }
de5 }
// Para NTT
9f7 template < int p>
97b void get_roots(bool f, int n, vector<mod_int<p>>& roots) {
        mod_int  r;
de9
       int ord;
       if (p == 998244353) {
57a
9b6
          r = 102292;
81b
           ord = (1 << 23):
1cc
       } else if (p == 754974721) {
43a
           r = 739831874;
f0a
            ord = (1 << 24);
b60
        } else if (p == 167772161) {
a2a
            r = 243;
033
            ord = (1 << 25):
        } else assert(false);
cd7
       if (f) r = r^(p - 1 - ord/n);
547
        else r = r^{(ord/n)}:
ee2
        roots[0] = 1:
be4
078
        for (int i = 1; i < n/2; i++) roots[i] = roots[i-1]*r;
63f }
8a2 template < typename T > void fft(vector < T > &a, bool f, int N,
   vector < int > &rev) {
       for (int i = 0; i < N; i++) if (i < rev[i]) swap(a[i],
   a[rev[i]]);
       int 1, r, m;
12b
cb4
       vector <T> roots(N);
       for (int n = 2: n \le N: n *= 2) {
192
0f4
            get roots(f. n. roots):
5dc
            for (int pos = 0; pos < N; pos += n) {</pre>
432
                1 = pos+0, r = pos+n/2, m = 0;
a88
                while (m < n/2) {
297
                    auto t = roots[m]*a[r]:
254
                    a[r] = a[1] - t;
                    a[1] = a[1] + t:
b8f
925
                    1++; r++; m++;
780
                }
            }
4ba
04a
        }
        if (f) {
235
1c5
            auto invN = T(1)/T(N);
```

```
557
             for (int i = 0; i < N; i++) a[i] = a[i]*invN;</pre>
256
        }
e5e }
bf5 template < typename T > vector < T > convolution (vector < T > &a, vector < T >
   &b) {
279
        vector <T> l(a.begin(), a.end());
f41
        vector <T> r(b.begin(), b.end());
        int ln = l.size(), rn = r.size();
7c6
287
        int N = ln+rn-1;
        int n = 1, log_n = 0;
f03
ac4
        while (n \le N) \{ n \le 1; \log_n + +; \}
808
        vector < int > rev(n);
bae
        for (int i = 0: i < n: ++i) {
434
            rev[i] = 0:
920
            for (int j = 0; j < log_n; ++j)</pre>
836
                 if (i & (1 << j)) rev[i] |= 1 << (log_n-1-j);
c34
        }
        assert(N <= n):
143
fa4
        1.resize(n);
        r.resize(n):
7e4
56e
        fft(l, false, n, rev);
fcf
        fft(r, false, n, rev);
917
        for (int i = 0; i < n; i++) 1[i] *= r[i];</pre>
88b
        fft(1, true, n, rev);
5e1
        l.resize(N):
792
        return 1:
700 }
// NTT
6c8 template < int p, typename T> vector < mod_int < p>> ntt (vector < T>& a,
   vector<T>& b) {
d52
        vector < mod_int < p >> A(a.begin(), a.end()), B(b.begin(),
d29
        return convolution(A. B):
3bf }
// Convolução de inteiro
// Precisa do CRT
// Tabela de valores:
// [0,1] - <int, 1>
// [-1e5, 1e5] - <11, 2>
// [-1e9, 1e9] - <__int128, 3>
b3c template < typename T, int mods >
eec vector <T> int_convolution(vector <int>& a, vector <int>& b) {
        static const int M1 = 998244353, M2 = 754974721, M3 =
fe8
```

```
3d8
   167772161;
                                                                                     }
bf5
        auto c1 = ntt < M1 > (a, b);
                                                                                      vector <T> ans(m, 0);
                                                                             b1d
        auto c2 = (mods \ge 2 ? ntt < M2 > (a, b) : vector < mod_int < M2 >> ());
                                                                                      for (int i = 0; i < m; i++) if (where[i] != -1)</pre>
221
                                                                             e1a
f9b
        auto c3 = (mods \ge 3? ntt < M3 > (a, b) : vector < mod_int < M3 >> ());
                                                                                          ans[i] = a[where[i]][m] / a[where[i]][i];
                                                                             12a
                                                                             603
                                                                                      for (int i = 0; i < n; i++) {</pre>
2da
        vector <T> ans:
                                                                             501
                                                                                          T sum = 0:
        for (int i = 0; i < c1.size(); i++) {</pre>
5 c 5
                                                                             a75
                                                                                          for (int j = 0; j < m; j++)
c09
            crt < T > at (c1[i].v, M1);
                                                                             5a9
                                                                                              sum += ans[i] * a[i][j];
316
            if (mods >= 2) at = at * crt<T>(c2[i].v, M2);
                                                                                          if (abs(sum - a[i][m]) > eps)
                                                                             b1f
987
            if (mods >= 3) at = at * crt<T>(c3[i].v, M3);
                                                                             6cd
                                                                                              return pair(0, vector<T>());
b2b
            ans.push_back(at.a);
                                                                              ec9
                                                                                      }
26d
            if (at.a > at.m/2) ans.back() -= at.m;
b9f
                                                                             12e
                                                                                      for (int i = 0; i < m; i++) if (where[i] == -1)</pre>
ba7
        return ans;
                                                                             018
                                                                                          return pair(INF, ans);
5e8 }
                                                                             280
                                                                                      return pair(1, ans);
                                                                             292 }
4.15 Gauss
                                                                             4.16 Gauss - Z2
// Resolve sistema linear
// Retornar um par com o numero de solucoes
                                                                             // D eh dimensao do espaco vetorial
                                                                             // add(v) - adiciona o vetor v na base (retorna se ele jah pertencia
// e alguma solucao, caso exista
//
                                                                                 ao span da base)
// O(n^2 * m)
                                                                             // coord(v) - retorna as coordenadas (c) de v na base atual (basis^T.c
                                                                             // recover(v) - retorna as coordenadas de v nos vetores na ordem em
67a template < typename T>
728 pair <int, vector <T>> gauss(vector <vector <T>> a, vector <T> b) {
                                                                                 que foram inseridos
                                                                             // coord(v).first e recover(v).first - se v pertence ao span
        const double eps = 1e-6;
        int n = a.size(), m = a[0].size();
                                                                             //
f92
        for (int i = 0; i < n; i++) a[i].push_back(b[i]);</pre>
2f0
                                                                             // Complexidade:
                                                                             // add, coord, recover: O(D^2 / 64)
3cb
        vector < int > where (m, -1);
237
        for (int col = 0, row = 0; col < m and row < n; col++) {
                                                                             2a3 template <int D> struct Gauss_z2 {
f05
            int sel = row:
                                                                             3c1
                                                                                      bitset < D > basis[D], keep[D];
b95
            for (int i=row; i<n; ++i)</pre>
                                                                             b16
                                                                                      int rk, in;
e55
                 if (abs(a[i][col]) > abs(a[sel][col])) sel = i;
                                                                                      vector < int > id;
                                                                             482
2c4
            if (abs(a[sel][col]) < eps) continue;</pre>
            for (int i = col; i <= m; i++)</pre>
                                                                             37f
                                                                                      Gauss_z2 () : rk(0), in(-1), id(D, -1) {};
1ae
                 swap(a[sel][i], a[row][i]);
dd2
            where [col] = row;
                                                                                      bool add(bitset <D> v) {
2c3
                                                                             04e
                                                                             42c
                                                                                          in++:
0 c 0
            for (int i = 0; i < n; i++) if (i != row) {
                                                                             fb0
                                                                                          bitset <D> k:
                T c = a[i][col] / a[row][col]:
96c
                                                                             659
                                                                                          for (int i = D - 1; i >= 0; i--) if (v[i]) {
d5c
                for (int j = col; j <= m; j++)</pre>
                                                                             189
                                                                                              if (basis[i][i]) v ^= basis[i], k ^= keep[i];
                     a[i][j] -= a[row][j] * c;
                                                                             4e6
                                                                                              else {
c8f
490
            }
                                                                             ea6
                                                                                                  k[i] = true, id[i] = in, keep[i] = k;
b70
            row++;
                                                                             6се
                                                                                                  basis[i] = v, rk++;
```

```
// computa o inverso modular de 1..MAX-1 modulo um primo
8a6
                     return true:
                }
                                                                             a88 ll inv[MAX]:
b34
09c
            }
                                                                             0f2 inv[1] = 1;
                                                                             Ofa for (int i = 2; i < MAX; i++) inv[i] = MOD - MOD/i*inv[MOD%i]%MOD;
d1f
            return false;
58b
0f6
        pair < bool, bitset < D >> coord(bitset < D > v) {
                                                                             4.19 Karatsuba
944
            bitset <D> c:
659
            for (int i = D - 1; i \ge 0; i--) if (v[i]) {
a39
                 if (basis[i][i]) v ^= basis[i], c[i] = true;
                                                                            // Os pragmas podem ajudar
                 else return {false, bitset <D>()};
8af
                                                                             // Para n \sim 2e5, roda em < 1 s
a08
            }
5db
            return {true, c};
                                                                             // O(n^1.58)
a08
330
        pair < bool, vector < int >> recover (bitset < D > v) {
                                                                             //#pragma GCC optimize("Ofast")
22e
            auto [span, bc] = coord(v);
                                                                             //#pragma GCC target ("avx,avx2")
af8
            if (not span) return {false, {}};
                                                                             77a template < typename T > void kar(T* a, T* b, int n, T* r, T* tmp) {
f79
            bitset <D> aux;
                                                                                     if (n <= 64) {
                                                                             d4c
            for (int i = D - 1; i >= 0; i--) if (bc[i]) aux ^= keep[i];
5a0
                                                                             510
                                                                                         for (int i = 0; i < n; i++) for (int j = 0; j < n; j++)
ea9
            vector < int > oc;
                                                                             212
                                                                                             r[i+j] += a[i] * b[j];
            for (int i = D - 1; i \ge 0; i--) if (aux[i])
                                                                             505
                                                                                         return;
   oc.push_back(id[i]);
                                                                                     }
                                                                             bb8
            return {true, oc};
001
                                                                             194
                                                                                     int mid = n/2:
b75
                                                                             2d7
                                                                                     T * atmp = tmp, *btmp = tmp+mid, *E = tmp+n;
d0a };
                                                                                     memset(E, 0, sizeof(E[0])*n);
                                                                             4 f 1
                                                                             c65
                                                                                     for (int i = 0; i < mid; i++) {</pre>
4.17 Integracao Numerica
                                                                             c72
                                                                                         atmp[i] = a[i] + a[i+mid];
                                                                             4b9
                                                                                         btmp[i] = b[i] + b[i+mid]:
// Metodo de Simpson 3/8
                                                                             a3f
// Integra f no intervalo [a, b], erro cresce proporcional a (b - a)^5
                                                                             38a
                                                                                     kar(atmp, btmp, mid, E, tmp+2*n);
                                                                                     kar(a, b, mid, r, tmp+2*n);
                                                                             b1e
676 const int N = 3*100; // multiplo de 3
                                                                             229
                                                                                     kar(a+mid, b+mid, mid, r+n, tmp+2*n);
287 ld integrate(ld a, ld b, function<ld(ld)> f) {
                                                                             c65
                                                                                     for (int i = 0; i < mid; i++) {</pre>
       ld s = 0, h = (b - a)/N;
b4d
                                                                             735
                                                                                         T \text{ temp} = r[i+mid];
        for (int i = 1 : i < N: i++) s += f(a + i*h)*(i%3 ? 3 : 2):
067
                                                                             de7
                                                                                         r[i+mid] += E[i] - r[i] - r[i+2*mid];
        return (f(a) + s + f(b))*3*h/8;
0da
                                                                             f1e
                                                                                         r[i+2*mid] += E[i+mid] - temp - r[i+3*mid];
c7e }
                                                                             f72
                                                                                     }
                                                                             28f }
4.18 Inverso Modular
                                                                             e38 template < typename T > vector < T > karatsuba (vector < T > a, vector < T > b)
// Computa o inverso de a modulo b
                                                                                {
// Se b eh primo, basta fazer
                                                                             ba3
                                                                                     int n = max(a.size(), b.size());
// a^{(b-2)}
                                                                             a84
                                                                                     while (n&(n-1)) n++:
                                                                                     a.resize(n), b.resize(n);
                                                                             ca9
f0a ll inv(ll a, ll b) {
                                                                                     vector\langle T \rangle ret(2*n), tmp(4*n);
                                                                             ae0
ae1
        return a > 1? b - inv(b\%a, a)*b/a : 1:
                                                                             644
                                                                                     kar(&a[0], &b[0], n, &ret[0], &tmp[0]);
```

edf

f87 }

return ret;

cf9 }

4.20 Logaritmo Discreto

```
// Resolve logaritmo discreto com o algoritmo baby step giant step
// Encontra o menor x tal que a^x = b (mod m)
// Se nao tem, retorna -1
// O(sqrt(m) * log(sqrt(m))
d41
da8 int dlog(int b, int a, int m) {
        if (a == 0) return b ? -1 : 1; // caso nao definido
d41
a6e
        a \%= m, b \%= m;
        int k = 1, shift = 0;
a10
31e
        while (1) {
6e3
           int g = gcd(a, m);
d47
           if (g == 1) break;
d41
9bc
           if (b == k) return shift;
642
           if (b % g) return -1;
c36
            b \neq g, m \neq g, shift++;
9ab
            k = (11) k * a / g % m;
515
       }
d41
        int sq = sqrt(m)+1, giant = 1;
af7
975
        for (int i = 0; i < sq; i++) giant = (11) giant * a % m;
d41
0b5
        vector < pair < int , int >> baby;
33f
       for (int i = 0, cur = b; i <= sq; i++) {
496
            baby.emplace_back(cur, i);
            cur = (11) cur * a % m;
16c
622
        sort(baby.begin(), baby.end());
eb4
d41
9c9
        for (int j = 1, cur = k; j \le sq; j++) {
ace
            cur = (11) cur * giant % m;
            auto it = lower_bound(baby.begin(), baby.end(), pair(cur,
78b
   INF));
            if (it != baby.begin() and (--it)->first == cur)
d26
ac3
                return sq * j - it->second + shift;
        }
b9d
d41
daa
        return -1;
739 }
```

4.21 Miller-Rabin

```
// Testa se n eh primo, n <= 3 * 10^18
```

```
// O(log(n)), considerando multiplicacao
// e exponenciacao constantes
 d8b ll mul(ll a. ll b. ll m) {
         11 ret = a*b - 11((long double)1/m*a*b+0.5)*m;
 e7a
 074
         return ret < 0 ? ret+m : ret:</pre>
 2f3 }
 03c ll pow(ll x, ll y, ll m) {
13a
        if (!v) return 1;
         11 ans = pow(mul(x, x, m), y/2, m);
7fa
         return y%2 ? mul(x, ans, m) : ans;
 539 }
1a2 bool prime(ll n) {
         if (n < 2) return 0;</pre>
 237
         if (n <= 3) return 1;
        if (n % 2 == 0) return 0;
 9de
        ll r = \_builtin\_ctzll(n - 1), d = n >> r;
f6a
         // com esses primos, o teste funciona garantido para n <= 2^64
         // funciona para n <= 3*10^24 com os primos ate 41
         for (int a: {2, 325, 9375, 28178, 450775, 9780504,
 771
    795265022}) {
 da0
             11 x = pow(a, d, n);
 709
             if (x == 1 \text{ or } x == n - 1 \text{ or a } \% n == 0) continue;
             for (int j = 0; j < r - 1; j++) {
 4a2
 10f
                 x = mul(x, x, n);
 df0
                 if (x == n - 1) break;
1ff
 e1b
             if (x != n - 1) return 0;
 e74
         }
 6a5
         return 1;
 5ba }
4.22 Pollard's Rho Alg
// Usa o algoritmo de deteccao de ciclo de Floyd
// com uma otimizacao na qual o gcd eh acumulado
// A fatoração não sai necessariamente ordenada
// O algoritmo rho encontra um fator de n,
// e funciona muito bem quando n possui um fator pequeno
// Complexidades (considerando mul constante):
// rho - esperado O(n^{(1/4)}) no pior caso
```

```
// fact - esperado menos que O(n^{(1/4)} \log(n)) no pior caso
d8b ll mul(ll a, ll b, ll m) {
        11 \text{ ret} = a*b - 11((long double)1/m*a*b+0.5)*m;
e7a
074
        return ret < 0 ? ret+m : ret;</pre>
2f3 }
03c ll pow(ll x, ll y, ll m) {
        if (!v) return 1;
13a
dbc
        11 ans = pow(mul(x, x, m), y/2, m);
7fa
        return y%2 ? mul(x, ans, m) : ans;
539 }
1a2 bool prime(ll n) {
       if (n < 2) return 0;
       if (n <= 3) return 1;</pre>
237
9de
       if (n % 2 == 0) return 0;
f6a
        ll r = \_builtin\_ctzll(n - 1), d = n >> r;
        for (int a: {2, 325, 9375, 28178, 450775, 9780504,
771
   795265022}) {
            11 x = pow(a, d, n);
da0
709
            if (x == 1 \text{ or } x == n - 1 \text{ or a } \% n == 0) continue;
4a2
            for (int j = 0; j < r - 1; j++) {
10f
                x = mul(x, x, n):
df0
                if (x == n - 1) break;
1ff
e1b
            if (x != n - 1) return 0;
e74
        return 1;
6a5
5ba }
9cf ll rho(ll n) {
        if (n == 1 or prime(n)) return n;
0f9
f7c
        auto f = [n](11 x) \{ return mul(x, x, n) + 1; \};
8a5
        11 x = 0, y = 0, t = 30, prd = 2, x0 = 1, q;
533
        while (t % 40 != 0 or gcd(prd, n) == 1) {
8a0
            if (x==y) x = ++x0, y = f(x);
e13
            q = mul(prd, abs(x-y), n);
            if (q != 0) prd = q;
21f
450
            x = f(x), y = f(f(y)), t++;
379
002
        return gcd(prd, n);
523 }
```

```
5b7 vector<ll> fact(ll n) {
1b9
        if (n == 1) return {};
0ec
        if (prime(n)) return {n};
        11 d = rho(n);
0ed
        vector < 11 > 1 = fact(d), r = fact(n / d);
1de
3af
        l.insert(l.end(), r.begin(), r.end());
792
        return 1:
902 }
4.23 Produto de dois long long mod m
// 0(1)
d8b ll mul(ll a, ll b, ll m) { // a*b % m
        11 \text{ ret} = a*b - 11((long double)1/m*a*b+0.5)*m;
074
        return ret < 0 ? ret+m : ret:</pre>
2f3 }
4.24 Simplex
// Maximiza c^T x s.t. Ax \langle = b, x \rangle = 0
//
// 0(2^n), porem executa em 0(n^3) no caso medio
395 const double eps = 1e-7;
493 namespace Simplex {
69c
        vector < vector < double >> T;
14e
        int n, m;
        vector < int > X, Y;
43e
c51
        void pivot(int x, int y) {
8e6
             swap(X[y], Y[x-1]);
d03
             for (int i = 0; i <= m; i++) if (i != y) T[x][i] /=
   T[x][y];
33c
            T[x][y] = 1/T[x][y];
             for (int i = 0; i \le n; i++) if (i != x and abs(T[i][y]) >
38b
    eps) {
774
                 for (int j = 0; j <= m; j++) if (j != y) T[i][j] -=
   T[i][y] * T[x][j];
3d8
                T[i][y] = -T[i][y] * T[x][y];
            }
a7d
e05
        }
        // Retorna o par (valor maximo, vetor solucao)
6f8
        pair < double , vector < double >> simplex(
```

```
e9d
                vector < vector < double >> A, vector < double >> b,
   vector < double > c) {
            n = b.size(), m = c.size();
5bb
            T = vector(n + 1, vector < double > (m + 1));
002
2d9
            X = vector < int > (m);
            Y = vector<int>(n):
0 c 2
115
            for (int i = 0: i < m: i++) X[i] = i:
            for (int i = 0; i < n; i++) Y[i] = i+m;</pre>
51f
5b5
            for (int i = 0; i < m; i++) T[0][i] = -c[i];
            for (int i = 0; i < n; i++) {</pre>
603
ba6
                for (int j = 0; j < m; j++) T[i+1][j] = A[i][j];
                T[i+1][m] = b[i];
eca
07c
            }
667
            while (true) {
                int x = -1, y = -1;
714
2db
                double mn = -eps;
               for (int i = 1; i <= n; i++) if (T[i][m] < mn) mn =
   T[i][m], x = i;
af2
               if (x < 0) break;
               for (int i = 0; i < m; i++) if (T[x][i] < -eps) { y = }
882
   i; break; }
4a6
                if (y < 0) return {-1e18, {}}; // sem solucao para Ax
   <= b
                pivot(x, y);
7fb
472
            }
667
            while (true) {
                int x = -1, y = -1;
714
2db
                double mn = -eps;
               for (int i = 0; i < m; i++) if (T[0][i] < mn) mn =
   T[0][i], y = i;
9b0
               if (v < 0) break;
034
                mn = 1e200:
                for (int i = 1; i \le n; i++) if (T[i][y] > eps and
   T[i][m] / T[i][y] < mn)
48f
                    mn = T[i][m] / T[i][y], x = i;
                if (x < 0) return {1e18, {}}; // c^T x eh ilimitado
53b
7fb
                pivot(x, y);
            }
81e
            vector < double > r(m);
290
            for(int i = 0; i < n; i++) if (Y[i] < m) r[Y[i]] =
32f
   T[i+1][m]:
e59
            return {T[0][m], r};
7a4
        }
a64 }
```

4.25 Teorema Chines do Resto

```
// Combina equações modulares lineares: x = a (mod m)
// O m final eh o lcm dos m's, e a resposta eh unica mod o lcm
// Os m nao precisam ser coprimos
// Se nao tiver solucao, o 'a' vai ser -1
153 template < typename T > tuple < T, T, T > ext_gcd(T a, T b) {
3bd
        if (!a) return {b, 0, 1};
550
        auto [g, x, y] = ext_gcd(b\%a, a);
        return \{g, y - b/a*x, x\};
c59
537 }
bfe template < typename T = 11 > struct crt {
        Ta, m;
5f3
        crt() : a(0), m(1) {}
7eb
        crt(T a_, T m_) : a(a_), m(m_) {}
911
        crt operator * (crt C) {
238
            auto [g, x, y] = ext_gcd(m, C.m);
            if ((a - C.a) \% g) a = -1;
dc0
4f9
            if (a == -1 or C.a == -1) return crt(-1, 0);
d09
            T lcm = m/g*C.m;
            T ans = a + (x*(C.a-a)/g \% (C.m/g))*m;
eb2
d8d
            return crt((ans % lcm + lcm) % lcm, lcm);
1f2
0d9 }:
```

4.26 Totiente

```
// O(sqrt(n))
a7e int tot(int n){
0f6
        int ret = n;
505
         for (int i = 2; i*i <= n; i++) if (n % i == 0) {</pre>
             while (n % i == 0) n /= i;
b0c
125
             ret -= ret / i:
34a
        }
        if (n > 1) ret -= ret / n;
af4
edf
        return ret;
fae }
```

5 DP

5.1 Divide and Conquer DP

```
// Particiona o arrav em k subarravs
// minimizando o somatorio das queries
// O(k n log n), assumindo quer query(1, r) eh O(1)
547 ll dp[MAX][2];
94b void solve(int k, int l, int r, int lk, int rk) {
de6
        if (1 > r) return;
109
        int m = (1+r)/2, p = -1;
d2b
        auto& ans = dp[m][k&1] = LINF;
6e2
        for (int i = max(m, lk); i <= rk; i++) {</pre>
            int at = dp[i+1][\sim k\&1] + query(m, i);
324
57d
            if (at < ans) ans = at, p = i;</pre>
91f
1ee
        solve(k, l, m-1, lk, p), solve(k, m+1, r, p, rk);
959 }
cf1 ll DC(int n, int k) {
        dp[n][0] = dp[n][1] = 0;
321
        for (int i = 0; i < n; i++) dp[i][0] = LINF;</pre>
f27
b76
        for (int i = 1; i <= k; i++) solve(i, 0, n-i, 0, n-i);
8e7
        return dp[0][k&1];
5e9 }
```

5.2 Longest Common Subsequence

```
// Computa a LCS entre dois arrays usando
// o algoritmo de Hirschberg para recuperar
//
// O(n*m), O(n+m) de memoria
eaf int lcs_s[MAX], lcs_t[MAX];
a6d int dp[2][MAX];
// dp[0][j] = max lcs(s[li...ri], t[lj, lj+j])
d12 void dp_top(int li, int ri, int lj, int rj) {
d13
        memset(dp[0], 0, (rj-lj+1)*sizeof(dp[0][0]));
753
        for (int i = li; i <= ri; i++) {</pre>
9aa
            for (int j = rj; j >= lj; j--)
83b
                dp[0][j-1j] = max(dp[0][j-1j],
```

```
741
                (lcs_s[i] == lcs_t[j]) + (j > 1j ? dp[0][j-1 - 1j] :
   0));
            for (int j = 1j+1; j \le rj; j++)
04c
                dp[0][j-1j] = max(dp[0][j-1j], dp[0][j-1-1j]);
939
        }
09f
58f }
// dp[1][j] = max lcs(s[li...ri], t[lj+j, rj])
ca0 void dp_bottom(int li, int ri, int lj, int rj) {
        memset(dp[1], 0, (rj-lj+1)*sizeof(dp[1][0]));
3a2
        for (int i = ri; i >= li; i--) {
49c
            for (int j = lj; j <= rj; j++)</pre>
dbb
                dp[1][j-lj] = max(dp[1][j-lj],
4da
                (lcs_s[i] == lcs_t[j]) + (j < rj ? dp[1][j+1 - 1j] :
   0));
            for (int j = rj-1; j >= lj; j--)
6ca
769
                dp[1][j-1j] = max(dp[1][j-1j], dp[1][j+1-1j]);
19b
        }
e8a }
93c void solve(vector<int>& ans, int li, int ri, int lj, int rj) {
2ad
        if (li == ri){
49 c
            for (int j = lj; j <= rj; j++)</pre>
f5b
                if (lcs_s[li] == lcs_t[j]){
a66
                     ans.push_back(lcs_t[j]);
c2b
                    break:
840
                }
505
            return:
126
        }
534
        if (li == ri){
753
            for (int i = li: i <= ri: i++){</pre>
88f
                if (lcs_s[i] == lcs_t[li]){
531
                    ans.push_back(lcs_s[i]);
c2b
                    break:
68a
                }
            }
a03
505
            return;
76d
        }
a57
        int mi = (li+ri)/2;
ade
        dp_top(li, mi, lj, rj), dp_bottom(mi+1, ri, lj, rj);
d7a
        int j_{-} = 0, mx = -1;
        for (int j = lj-1; j <= rj; j++) {</pre>
aee
da8
            int val = 0;
2bb
            if (j >= lj) val += dp[0][j - lj];
            if (j < rj) val += dp[1][j+1 - lj];</pre>
b9e
```

```
ba8
            if (val >= mx) mx = val, j_ = j;
14e
6f1
        if (mx == -1) return;
        solve(ans, li, mi, lj, j_), solve(ans, mi+1, ri, j_+1, rj);
c2a
dd5 }
058 vector<int> lcs(const vector<int>& s, const vector<int>& t) {
        for (int i = 0; i < s.size(); i++) lcs_s[i] = s[i];</pre>
577
        for (int i = 0; i < t.size(); i++) lcs t[i] = t[i];
dab
        vector < int > ans;
599
        solve(ans, 0, s.size()-1, 0, t.size()-1):
ba7
        return ans;
17c }
5.3 Mochila
// Resolve mochila, recuperando a resposta
//
// O(n * cap), O(n + cap) de memoria
add int v[MAX], w[MAX]; // valor e peso
582 int dp[2][MAX_CAP];
// DP usando os itens [1, r], com capacidade = cap
0d6 void get_dp(int x, int 1, int r, int cap) {
        memset(dp[x], 0, (cap+1)*sizeof(dp[x][0]));
        for (int i = 1; i \le r; i++) for (int j = cap; j \ge 0; j--)
574
            if (j - w[i] >= 0) dp[x][j] = max(dp[x][j], v[i] + dp[x][j]
   - w[i]]);
b95 }
5ab void solve(vector<int>& ans, int 1, int r, int cap) {
893
        if (1 == r) {
9ff
            if (w[1] <= cap) ans.push_back(1);</pre>
505
            return;
13a
       int m = (1+r)/2;
ee4
283
        get_dp(0, 1, m, cap), get_dp(1, m+1, r, cap);
056
       int left_cap = -1, opt = -INF;
        for (int j = 0; j \le cap; j++)
c94
2f2
            if (int at = dp[0][j] + dp[1][cap - j]; at > opt)
91d
                opt = at, left_cap = j;
da3
        solve(ans, 1, m, left_cap), solve(ans, m+1, r, cap - left_cap);
d75 }
0d7 vector<int> knapsack(int n, int cap) {
```

```
dab
        vector<int> ans:
1e0
        solve(ans, 0, n-1, cap);
        return ans;
ba7
e4d }
5.4 SOS DP
// \Omega(n 2^n)
// soma de sub-conjunto
e03 vector<11> sos dp(vector<11> f) {
        int N = __builtin_ctz(f.size());
6c0
        assert((1<<N) == f.size()):
e59
5a5
        for (int i = 0; i < N; i++) for (int mask = 0; mask < (1<<N);
   mask++)
796
            if (mask>>i&1) f[mask] += f[mask^(1<<ii)];</pre>
abe
        return f:
bec }
// soma de super-conjunto
e03 vector<ll> sos_dp(vector<ll> f) {
        int N = __builtin_ctz(f.size());
        assert((1<<N) == f.size()):
e59
        for (int i = 0; i < N; i++) for (int mask = 0; mask < (1<<N);
5a5
   mask++)
            if (\sim mask >> i\&1) f[mask] += f[mask^(1<<ii)];
a3c
abe
        return f;
dbd }
5.5 Subset sum
// Retorna max(x <= t tal que existe subset de w que soma x)
// O(n * max(w))
// O(max(w)) de memoria
efd int subset_sum(vector<int> w, int t) {
        int pref = 0, k = 0;
bb5
        while (k < w.size() and pref + w[k] <= t) pref += w[k++];</pre>
417
1e7
        if (k == w.size()) return pref;
        int W = *max_element(w.begin(), w.end());
444
        vector \langle int \rangle last, dp(2*W, -1);
44d
d7b
        dp[W - (t-pref)] = k;
54d
        for (int i = k; i < w.size(); i++) {</pre>
288
            last = dp;
```

```
15f
            for (int x = 0; x < W; x++) dp[x+w[i]] = max(dp[x+w[i]],
   last[x]):
           for (int x = 2*W - 1; x > W; x--)
17b
                for (int j = max(0, last[x]); j < dp[x]; j++)
303
595
                    dp[x-w[j]] = max(dp[x-w[j]], j);
        }
867
2fb
        int ans = t:
        while (dp[W - (t-ans)] < 0) ans --;
1 c 1
ba7
        return ans;
488 }
```

6 Strings

6.1 Aho-corasick

```
// query retorna o somatorio do numero de matches de
// todas as stringuinhas na stringona
//
// insert - O(|s| log(SIGMA))
// build - O(N), onde N = somatorio dos tamanhos das strings
// query - O(|s|)
eal namespace aho {
807
        map < char , int > to[MAX];
c87
        int link[MAX], idx, term[MAX], exit[MAX], sobe[MAX];
bfc
        void insert(string& s) {
05e
            int at = 0;
b4f
            for (char c : s) {
b68
                auto it = to[at].find(c);
1c9
                if (it == to[at].end()) at = to[at][c] = ++idx;
361
                else at = it->second;
ff4
142
            term[at]++, sobe[at]++;
6eb
d41 #warning nao esquece de chamar build() depois de inserir
        void build() {
0a8
26a
            queue < int > q;
537
            q.push(0);
            link[0] = exit[0] = -1;
dff
402
            while (q.size()) {
379
                int i = q.front(); q.pop();
3c4
                for (auto [c, j] : to[i]) {
                    int l = link[i]:
5da
102
                    while (1 != -1 and !to[1].count(c)) 1 = link[1];
7a5
                    link[j] = 1 == -1 ? 0 : to[1][c];
```

```
3ab
                     exit[j] = term[link[j]] ? link[j] : exit[link[j]];
                    if (exit[j]+1) sobe[j] += sobe[exit[j]];
6f2
113
                    q.push(j);
f1d
                }
367
            }
768
        }
bc0
        int query(string& s) {
86d
            int at = 0, ans = 0;
b4f
            for (char c : s){
                while (at != -1 and !to[at].count(c)) at = link[at];
1 ca
5b9
                at = at == -1 ? 0 : to[at][c];
2b1
                ans += sobe[at]:
b85
            }
ba7
            return ans;
038
        }
a30 }
6.2 eertree
// Constroi a eertree, caractere a caractere
// Inicializar com a quantidade de caracteres maxima
// size() retorna a quantidade de substrings pal. distintas
// depois de chamar propagate(), cada substring palindromica
// ocorre qt[i] vezes. O propagate() retorna o numero de
// substrings pal. com repeticao
//
// O(n) amortizado, considerando alfabeto O(1)
8eb struct eertree {
7 c.c
        vector < vector < int >> t;
42e
        int n, last, sz;
745
        vector<int> s, len, link, qt;
d36
        eertree(int N) {
ec8
            t = vector(N+2, vector(26, int()));
            s = len = link = qt = vector < int > (N+2);
cee
cd1
288
            link[0] = 1, len[0] = 0, link[1] = 1, len[1] = -1;
            sz = 2, last = 0, n = 1;
688
        }
8dc
244
        void add(char c) {
            s[n++] = c -= 'a';
692
34f
            while (s[n-len[last]-2] != c) last = link[last];
289
            if (!t[last][c]) {
dab
                int prev = link[last];
553
                while (s[n-len[prev]-2] != c) prev = link[prev];
```

```
fb2
                link[sz] = t[prev][c];
                len[sz] = len[last]+2:
3f5
1f8
                t[last][c] = sz++;
            }
f8b
344
            qt[last = t[last][c]]++;
b1d
f17
        int size() { return sz-2; }
2af
        11 propagate() {
b73
            11 \text{ ret} = 0;
            for (int i = n; i > 1; i--) {
ebb
fd3
                qt[link[i]] += qt[i];
db5
                ret += qt[i];
074
            }
edf
            return ret;
ef6
       }
a2e }:
6.3 KMP
// de s em t
```

```
// matching(s, t) retorna os indices das ocorrencias
// autKMP constroi o automato do KMP
// Complexidades:
// pi - O(n)
// match - O(n + m)
// construir o automato - O(|sigma|*n)
// n = |padrao| e m = |texto|
ea8 template < typename T > vector < int > pi(T s) {
019
        vector < int > p(s.size());
725
        for (int i = 1, j = 0; i < s.size(); i++) {</pre>
a51
            while (j and s[j] != s[i]) j = p[j-1];
973
            if (s[i] == s[i]) i++;
f8c
            p[i] = j;
e0a
74e
        return p;
f50 }
c10 template < typename T> vector < int > matching (T& s, T& t) {
        vector < int > p = pi(s), match;
658
a1b
        for (int i = 0, j = 0; i < t.size(); i++) {</pre>
            while (j \text{ and } s[j] != t[i]) j = p[j-1];
c4d
            if (s[i] == t[i]) i++;
310
            if (j == s.size()) match.push_back(i-j+1), j = p[j-1];
028
        }
ed8
        return match;
```

```
c82 }
a2d struct KMPaut : vector<vector<int>> {
47 c
        KMPaut(){}
        KMPaut (string& s) : vector < vector < int >> (26,
6 c 7
    vector < int > (s.size()+1)) {
503
             vector<int> p = pi(s):
04b
             auto& aut = *this;
4fa
             aut[s[0]-'a'][0] = 1;
             for (char c = 0; c < 26; c++)
19a
5d3
                 for (int i = 1; i <= s.size(); i++)</pre>
42b
                     aut[c][i] = s[i] - a' == c ? i+1 : aut[c][p[i-1]]:
4bb
        }
79b };
6.4 Manacher
// manacher recebe um vetor de T e retorna o vetor com tamanho dos
    palindromos
// ret[2*i] = tamanho do maior palindromo centrado em i
// \text{ ret}[2*i+1] = \text{tamanho maior palindromo centrado em i e i+1}
//
// Complexidades:
// manacher - O(n)
// palindrome - <0(n), 0(1)>
// pal_end - O(n)
28a template < typename T > vector < int > manacher (const T& s) {
18f
        int l = 0, r = -1, n = s.size();
fc9
        vector < int > d1(n), d2(n);
603
        for (int i = 0; i < n; i++) {</pre>
821
             int k = i > r ? 1 : min(d1[l+r-i], r-i);
61a
             while (i+k < n \&\& i-k >= 0 \&\& s[i+k] == s[i-k]) k++;
61e
             d1[i] = k--:
9f6
             if (i+k > r) l = i-k, r = i+k;
950
        }
e03
        1 = 0, r = -1;
603
        for (int i = 0; i < n; i++) {</pre>
             int k = i > r ? 0 : min(d2[1+r-i+1], r-i+1); k++;
a64
2c6
             while (i+k \le n \&\& i-k \ge 0 \&\& s[i+k-1] == s[i-k]) k++;
             d2[i] = --k:
eaa
26d
             if (i+k-1 > r) l = i-k, r = i+k-1:
4fe
        }
c41
        vector < int > ret(2*n-1);
e6b
        for (int i = 0; i < n; i++) ret[2*i] = 2*d1[i]-1;</pre>
e1d
        for (int i = 0; i < n-1; i++) ret[2*i+1] = 2*d2[i+1];
edf
        return ret;
```

```
ebb }
// verifica se a string s[i..j] eh palindromo
cac template < typename T> struct palindrome {
f97
        vector < int > man;
b2d
        palindrome(const T& s) : man(manacher(s)) {}
9d7
        bool query(int i, int j) {
            return man[i+j] >= j-i+1;
bad
1e7
60c };
// tamanho do maior palindromo que termina em cada posicao
7cb template < typename T > vector < int > pal_end(const T& s) {
e57
        vector < int > ret(s.size());
fde
        palindrome <T> p(s);
d51
        ret[0] = 1;
       for (int i = 1; i < s.size(); i++) {</pre>
88e
            ret[i] = min(ret[i-1]+2, i+1);
a32
            while (!p.query(i-ret[i]+1, i)) ret[i]--;
6ea
78e
        }
edf
        return ret;
8bd }
    Min/max suffix/cyclic shift
// Computa o indice do menor/maior sufixo/cyclic shift
// da string, lexicograficamente
//
// O(n)
016 template < typename T > int max_suffix(T s, bool mi = false) {
476
        s.push_back(*min_element(s.begin(), s.end())-1);
1a4
        int ans = 0:
88e
        for (int i = 1; i < s.size(); i++) {</pre>
            int j = 0;
eec
            while (ans+j < i and s[i+j] == s[ans+j]) j++;
708
7a2
            if (s[i+j] > s[ans+j]) {
b52
                if (!mi or i != s.size()-2) ans = i;
            } else if (j) i += j-1;
e51
69 c
ba7
        return ans;
f2a }
a1a template < typename T > int min_suffix(T s) {
76b
        for (auto& i : s) i *= -1;
```

s.push_back(*max_element(s.begin(), s.end())+1);

09d

```
925
         return max_suffix(s, true);
ec0 }
97c template < typename T > int max_cyclic_shift(T s) {
         int n = s.size();
163
1ad
         for (int i = 0; i < n; i++) s.push_back(s[i]);</pre>
20a
         return max suffix(s):
d34 }
08a template < typename T > int min_cyclic_shift(T s) {
76b
         for (auto& i : s) i *= -1;
7be
         return max_cyclic_shift(s);
 c7a }
6.6 String Hashing
// Complexidades:
// construtor - O(|s|)
// operator() - 0(1)
878 mt19937 rng((int)
    chrono::steady_clock::now().time_since_epoch().count());
463 int uniform(int 1, int r) {
a7f
         uniform_int_distribution < int > uid(1, r);
f54
         return uid(rng);
d9e }
9e0 template <int MOD> struct str_hash { // 116fcb
         static int P;
dcf
         vector<ll> h, p;
         str_hash(string s) : h(s.size()), p(s.size()) {
 ea8
             p[0] = 1, h[0] = s[0];
7a2
             for (int i = 1; i < s.size(); i++)</pre>
 ad7
                 p[i] = p[i - 1]*P%MOD, h[i] = (h[i - 1]*P + s[i])%MOD;
84 c
1ef
         11 operator()(int 1, int r) { // retorna hash s[1...r]
af7
749
             ll hash = h[r] - (1 ? h[1 - 1]*p[r - 1 + 1]%MOD : 0);
             return hash < 0 ? hash + MOD : hash;</pre>
dfd
3ba
116 }:
217 template <int MOD > int str_hash < MOD >:: P = uniform (256, MOD - 1); //
    1 > |sigma|
6.7 String Hashing - modulo 2<sup>61</sup> - 1
// Quase duas vezes mais lento
```

```
// Complexidades:
// build - O(|s|)
// operator() - 0(1)
9d0 const ll MOD = (111<<61) - 1;
e38 ll mulmod(ll a. ll b) {
        const static ll LOWER = (111 << 30) - 1, GET31 = (111 << 31) - 1;
        11 \ 11 = a\&LOWER, h1 = a>>30, 12 = b\&LOWER, h2 = b>>30;
410
d54
        11 m = 11*h2 + 12*h1, h = h1*h2;
        ll ans = 11*12 + (h>>1) + ((h&1)<<60) + (m>>31) +
   ((m\&GET31) << 30) + 1:
        ans = (ans\&MOD) + (ans>>61), ans = (ans\&MOD) + (ans>>61);
1dd
c0f
        return ans - 1:
f98 }
798 mt19937_64
   rng(chrono::steady_clock::now().time_since_epoch().count());
f89 ll uniform(ll l, ll r) {
        uniform_int_distribution < ll > uid(1, r);
969
f54
        return uid(rng);
cac }
d7d struct str hash {
c20
        static 11 P:
dcf
       vector < 11 > h, p;
ea8
        str_hash(string s) : h(s.size()), p(s.size()) {
7a2
            p[0] = 1, h[0] = s[0];
ad7
            for (int i = 1; i < s.size(); i++)</pre>
                p[i] = mulmod(p[i - 1], P), h[i] = (mulmod(h[i - 1],
   P) + s[i])%MOD;
       }
507
af7
        11 operator()(int 1. int r) { // retorna hash s[1...r]
            ll hash = h[r] - (1 ? mulmod(h[1 - 1], p[r - 1 + 1]) : 0);
538
dfd
            return hash < 0 ? hash + MOD : hash;</pre>
544
        }
148 };
6c5 ll str_hash::P = uniform(256, MOD - 1); // 1 > |sigma|
6.8 Suffix Array - O(n log n)
// kasai recebe o suffix array e calcula lcp[i],
// o lcp entre s[sa[i],...,n-1] e s[sa[i+1],...,n-1]
//
// Complexidades:
// suffix_array - O(n log(n))
```

```
// kasai - O(n)
733 vector <int> suffix_array(string s) {
         s += "$":
         int n = s.size(), N = max(n, 260):
043
2f3
         vector<int> sa(n), ra(n);
29b
         for(int i = 0; i < n; i++) sa[i] = i, ra[i] = s[i];
0a2
         for (int k = 0; k < n; k ? k *= 2 : k++) {
             vector < int > nsa(sa), nra(n), cnt(N);
5ce
fae
             for(int i = 0; i < n; i++) nsa[i] = (nsa[i]-k+n)%n,
    cnt[ra[i]]++:
4c4
             for(int i = 1; i < N; i++) cnt[i] += cnt[i-1];</pre>
368
             for(int i = n-1; i+1; i--) sa[--cnt[ra[nsa[i]]] = nsa[i];
28f
             for(int i = 1, r = 0; i < n; i++) nra[sa[i]] = r +=</pre>
    ra[sa[i]] !=
                 ra[sa[i-1]] or ra[(sa[i]+k)\%n] != ra[(sa[i-1]+k)\%n];
f86
26b
             ra = nra:
d5e
             if (ra[sa[n-1]] == n-1) break;
11e
057
         return vector < int > (sa.begin()+1, sa.end());
ff3 }
481 vector <int > kasai(string s, vector <int > sa) {
 232
         int n = s.size(), k = 0;
408
         vector < int > ra(n). lcp(n):
         for (int i = 0; i < n; i++) ra[sa[i]] = i;</pre>
676
         for (int i = 0; i < n; i++, k -= !!k) {
740
199
             if (ra[i] == n-1) { k = 0; continue; }
             int j = sa[ra[i]+1];
1de
891
             while (i+k < n \text{ and } i+k < n \text{ and } s[i+k] == s[i+k]) k++:
d98
             lcp[ra[i]] = k;
a07
         }
         return lcp;
5ed
fbe }
6.9 Suffix Array - O(n)
// Rapidao
// Computa o suffix array em 'sa', o rank em 'rnk'
// e o lcp em 'lcp'
// query(i, j) retorna o LCP entre s[i..n-1] e s[j..n-1]
//
// Complexidades
```

```
// O(n) para construir
// query - 0(1)
1a5 template < typename T> struct rmq {
517
        vector <T> v:
        int n; static const int b = 30;
fcc
70e
        vector < int > mask, t;
        int op(int x, int y) { return v[x] \le v[y] ? x : y; }
183
        int msb(int x) { return __builtin_clz(1)-__builtin_clz(x); }
ee1
c92
        int small(int r, int sz = b) { return
   r-msb(mask[r]&((1<<sz)-1));}
        rma() {}
6ad
        rmq(const\ vector < T > \&\ v_): v(v_), n(v.size()), mask(n), t(n) {
43c
            for (int i = 0, at = 0; i < n; mask[i++] = at |= 1) {</pre>
2e5
a61
                at = (at << 1) &((1 << b) -1);
c00
                while (at and op(i-msb(at&-at), i) == i) at ^= at&-at;
c2f
            for (int i = 0; i < n/b; i++) t[i] = small(b*i+b-1);
ea4
            for (int j = 1; (1<<j) <= n/b; j++) for (int i = 0;
   i+(1<<j) <= n/b; i++)
                t[n/b*j+i] = op(t[n/b*(j-1)+i],
   t[n/b*(j-1)+i+(1<<(j-1))]);
       }
41a
e34
        int index_query(int 1, int r) {
27b
            if (r-l+1 \le b) return small(r, r-l+1):
            int x = 1/b+1, y = r/b-1;
e80
fd3
            if (x > y) return op(small(1+b-1), small(r));
a4e
            int j = msb(y-x+1);
            int ans = op(small(l+b-1), op(t[n/b*j+x],
ea3
   t[n/b*j+y-(1<<j)+1]));
            return op(ans, small(r));
be6
62a
093
        T query(int 1, int r) { return v[index_query(1, r)]; }
bab };
9d7 struct suffix_array {
        string s;
ac0
1a8
        int n;
5b4
        vector < int > sa, cnt, rnk, lcp;
2de
        rmq<int> RMQ;
        bool cmp(int a1, int b1, int a2, int b2, int a3=0, int b3=0) {
d6e
            return a1 != b1 ? a1 < b1 : (a2 != b2 ? a2 < b2 : a3 < b3):
91d
82d
4a4
        template < typename T > void radix(int * fr, int * to, T * r, int N,
   int k) {
```

```
c17
            cnt = vector < int > (k+1, 0);
            for (int i = 0; i < N; i++) cnt[r[fr[i]]]++;</pre>
bac
703
            for (int i = 1; i <= k; i++) cnt[i] += cnt[i-1];</pre>
            for (int i = N-1; i+1; i--) to[--cnt[r[fr[i]]]] = fr[i];
000
6f3
        }
d66
        void rec(vector<int>& v, int k) {
a76
             auto &tmp = rnk, &m0 = lcp;
3a9
            int N = v.size()-3, sz = (N+2)/3, sz2 = sz+N/3;
7f8
            vector < int > R(sz2+3);
            for (int i = 1, j = 0; j < sz2; i += i%3) R[j++] = i;
74f
b30
            radix(&R[0], &tmp[0], &v[0]+2, sz2, k);
207
            radix(&tmp[0]. &R[0]. &v[0]+1. sz2. k):
             radix(&R[0], &tmp[0], &v[0]+0, sz2, k);
5f1
            int dif = 0;
af5
ed9
            int 10 = -1, 11 = -1, 12 = -1;
d81
            for (int i = 0; i < sz2; i++) {</pre>
8de
                 if (v[tmp[i]] != 10 or v[tmp[i]+1] != 11 or
   v[tmp[i]+2] != 12)
                     10 = v[tmp[i]], 11 = v[tmp[i]+1], 12 =
b43
   v[tmp[i]+2], dif++;
                 if (tmp[i]%3 == 1) R[tmp[i]/3] = dif;
199
                 else R[tmp[i]/3+sz] = dif;
1f5
d18
            }
47f
            if (dif < sz2) {</pre>
146
                 rec(R. dif):
746
                 for (int i = 0; i < sz2; i++) R[sa[i]] = i+1;</pre>
105
            } else for (int i = 0; i < sz2; i++) sa[R[i]-1] = i;</pre>
6f4
            for (int i = 0, j = 0; j < sz2; i++) if (sa[i] < sz)
   tmp[i++] = 3*sa[i]:
7ce
            radix(&tmp[0], &m0[0], &v[0], sz, k):
            for (int i = 0; i < sz2; i++)</pre>
74d
                 sa[i] = sa[i] < sz ? 3*sa[i]+1 : 3*(sa[i]-sz)+2;
с9е
332
            int at = sz2+sz-1, p = sz-1, p2 = sz2-1;
1c9
            while (p \ge 0 \text{ and } p2 \ge 0) {
3b3
                 if ((sa[p2]\%3==1 \text{ and } cmp(v[m0[p]], v[sa[p2]]),
   R[m0[p]/3],
                     R[sa[p2]/3+sz])) or (sa[p2]%3==2 and cmp(v[m0[p]],
Осе
   v[sa[p2]],
af6
                     v[m0[p]+1], v[sa[p2]+1], R[m0[p]/3+sz],
   R[sa[p2]/3+1]))
300
                     sa[at--] = sa[p2--];
cb0
                 else sa[at--] = m0[p--]:
```

```
214
             while (p >= 0) sa[at--] = m0[p--];
f2b
eb6
            if (N\%3==1) for (int i = 0; i < N; i++) sa[i] = sa[i+1];
ee6
        }
        suffix_array(const string& s_) : s(s_), n(s.size()), sa(n+3),
938
e62
                 cnt(n+1), rnk(n), lcp(n-1) {
             vector < int > v(n+3);
9fe
f9b
             for (int i = 0; i < n; i++) v[i] = i;</pre>
             radix(&v[0], &rnk[0], &s[0], n, 256);
eba
e6d
            int dif = 1;
830
            for (int i = 0; i < n; i++)</pre>
419
                 v[rnk[i]] = dif += (i and s[rnk[i]] != s[rnk[i-1]]);
            if (n \ge 2) rec(v, dif);
7cf
fb9
             sa.resize(n);
76f
             for (int i = 0; i < n; i++) rnk[sa[i]] = i;</pre>
892
             for (int i = 0, k = 0; i < n; i++, k -= !!k) {
                 if (rnk[i] == n-1) {
668
                     k = 0:
5a4
5e2
                     continue;
9df
39a
                 int j = sa[rnk[i]+1];
891
                 while (i+k < n \text{ and } j+k < n \text{ and } s[i+k] == s[j+k]) k++;
825
                 lcp[rnk[i]] = k;
a3e
            }
9ff
             RMQ = rmq<int>(lcp);
9a8
        }
588
        int query(int i, int j) {
d97
            if (i == j) return n-i;
223
            i = rnk[i], j = rnk[j];
            return RMQ.query(min(i, j), max(i, j)-1);
c3a
940
        pair<int, int> next(int L, int R, int i, char c) {
71c
024
            int 1 = L, r = R+1;
40c
             while (1 < r) {
ee4
                 int m = (1+r)/2;
                 if (i+sa[m] >= n or s[i+sa[m]] < c) l = m+1;</pre>
e7e
ef3
                 else r = m;
ebe
            }
            if (1 == R+1 \text{ or } s[i+sa[1]] > c) \text{ return } \{-1, -1\};
575
            L = 1;
eb7
9e2
            l = L, r = R+1;
             while (1 < r) {
40c
                 int m = (1+r)/2;
ee4
```

```
1a1
                if (i+sa[m] >= n or s[i+sa[m]] <= c) l = m+1;
ef3
                else r = m;
b5b
            }
56a
            R = 1-1;
            return {L, R};
e13
71b
        // quantas vezes 't' ocorre em 's' - O(|t| log n)
        int count_substr(string& t) {
66d
b2b
            int L = 0, R = n-1;
            for (int i = 0; i < t.size(); i++) {</pre>
c9d
de0
                tie(L, R) = next(L, R, i, t[i]);
4fc
                if (L == -1) return 0;
cff
            }
fbf
            return R-L+1;
aaa
        }
        // exemplo de f que resolve o problema
            https://codeforces.com/edu/course/2/lesson/2/5/practice/contes
        ll f(ll k) \{ return k*(k+1)/2; \}
57e
        11 dfs(int L, int R, int p) { // dfs na suffix tree chamado em
e68
   pre ordem
            int ext = L != R ? RMQ.query(L, R-1) : n - sa[L];
c54
            // Tem 'ext - p' substrings diferentes que ocorrem 'R-L+1'
                vezes
            // O LCP de todas elas eh 'ext'
            ll ans = (ext-p)*f(R-L+1);
f80
            // L eh terminal. e folha sse L == R
63 c
            if (sa[L]+ext == n) L++;
            // se for um SA de varias strings separadas como s#t$u&,
                usar no lugar do if de cima
            // (separadores < 'a', diferentes e inclusive no final)</pre>
            // while (L <= R && (sa[L]+ext == n \mid | s[sa[L]+ext] <
                'a')) {
            // L++;
            // }
            while (L <= R) {</pre>
add
5a8
                int idx = L != R ? RMQ.index_query(L, R-1) : -1;
                if (idx == -1 or lcp[idx] != ext) idx = R;
5ef
478
                ans += dfs(L, idx, ext);
28d
                L = idx+1:
```

```
590     }
ba7     return ans;
e21 }

    // sum over substrings: computa, para toda substring t
        distinta de s,
        // \sum f(# ocorrencias de t em s) - 0 (n)
ca8     ll sos() { return dfs(0, n-1, 0); }
6fa };
```

6.10 Suffix Array Dinamico

```
// Mantem o suffix array, lcp e rank de uma string,
// premitindo push_front e pop_front
// O operador [i] return um par com sa[i] e lcp[i]
// lcp[i] tem o lcp entre sa[i] e sa[i-1] (lcp[0] = 0)
//
// Complexidades:
// Construir sobre uma string de tamanho n: O(n log n)
// push_front e pop_front: O(log n) amortizado
2fe struct dyn_sa {
3c9
        struct node {
1d4
            int sa, lcp;
            node *1, *r, *p;
ed1
f0d
            int sz, mi;
17b
            node(int sa_, int lcp_, node* p_) : sa(sa_), lcp(lcp_),
                1(NULL), r(NULL), p(p_), sz(1), mi(lcp) {}
543
01e
            void update() {
58f
                sz = 1, mi = lcp;
bd7
                if (1) sz += 1->sz, mi = min(mi, 1->mi);
a54
                if (r) sz += r->sz, mi = min(mi, r->mi);
27 c
            }
574
        };
bb7
        node* root;
295
        vector<ll> tag; // tag of a suffix (reversed id)
        string s; // reversed
ac0
        dyn_sa() : root(NULL) {}
cf4
        dyn_sa(string s_) : dyn_sa() {
e45
ae4
            reverse(s_.begin(), s_.end());
            for (char c : s_) push_front(c);
519
2a7
a86
        \sim dyn_sa() {
609
            vector < node *> q = {root};
402
            while (q.size()) {
```

```
e5d
                 node* x = q.back(); q.pop_back();
ee9
                 if (!x) continue;
                 q.push_back(x->1), q.push_back(x->r);
1c7
bf0
                 delete x;
653
            }
        }
8c1
        int size(node* x) { return x ? x->sz : 0; }
73c
        int mirror(int i) { return s.size()-1 - i; }
08e
        bool cmp(int i, int j) {
580
a29
            if (s[i] != s[j]) return s[i] < s[j];</pre>
5b4
            if (i == 0 or j == 0) return i < j;</pre>
988
            return tag[i-1] < tag[j-1];</pre>
9fd
        }
        void fix_path(node* x) { while (x) x->update(), x = x->p; }
919
245
        void flatten(vector < node * > & v, node * x) {
8c8
            if (!x) return;
e96
            flatten(v, x->1);
2a2
            v.push_back(x);
42d
            flatten(v, x->r);
01f
        }
        void build(vector<node*>& v, node*& x, node* p, int L, int R,
964
   11 1. 11 r) {
04c
            if (L > R) return void(x = NULL);
331
            int M = (L+R)/2;
3e3
            11 m = (1+r)/2:
7e5
            x = v[M]:
63e
            x - p = p;
bb3
            tag[x->sa] = m;
            build(v, x->1, x, L, M-1, 1, m-1), build(v, x->r, x, M+1,
ae0
   R, m+1, r);
            x->update();
ca8
a3a
82f
        void fix(node*& x, node* p, ll l, ll r) {
            if (3*max(size(x->1), size(x->r)) \le 2*size(x)) return
7f0
   x->update();
3d1
            vector < node *> v;
Осс
            flatten(v, x);
            build(v, x, p, 0, v.size()-1, 1, r);
ea9
b86
        }
b19
        node* next(node* x) {
728
            if (x->r) {
a91
                x = x - > r:
347
                 while (x->1) x = x->1;
ea5
                 return x;
e7d
            }
402
            while (x->p \text{ and } x->p->r == x) x = x->p;
```

```
137
            return x->p;
48b
b68
        node* prev(node* x) {
            if (x->1) {
e41
a26
                 x = x - > 1:
93c
                 while (x->r) x = x->r;
ea5
                 return x:
            }
9be
6a1
            while (x->p \text{ and } x->p->1 == x) x = x->p;
137
            return x->p;
73e
        }
4f7
        int get lcp(node* x. node* v) {
            if (!x or !y) return 0; // change defaut value here
75a
            if (s[x->sa] != s[y->sa]) return 0;
e51
843
            if (x->sa == 0 \text{ or } y->sa == 0) \text{ return } 1;
4d0
            return 1 + query(mirror(x->sa-1), mirror(y->sa-1));
8d6
ad6
        void add_suf(node*& x, node* p, int id, ll l, ll r) {
            if (!x) {
91e
8e3
                 x = new node(id, 0, p);
8e2
                 node *prv = prev(x), *nxt = next(x);
                 int lcp_cur = get_lcp(prv, x), lcp_nxt = get_lcp(x,
65d
   nxt);
                 if (nxt) nxt->lcp = lcp_nxt, fix_path(nxt);
ca3
71f
                 x \rightarrow 1cp = 1cp cur:
                 tag[id] = (1+r)/2;
7b4
ca8
                 x->update();
                 return;
505
d0e
            }
4a3
            if (cmp(id, x->sa)) add_suf(x->l, x, id, l, tag[x->sa]-1);
c3a
            else add_suf(x \rightarrow r, x, id, tag[x \rightarrow sa]+1, r);
            fix(x, p, l, r);
3db
c98
        }
        void push_front(char c) {
ec2
cc7
            s += c:
493
            tag.push_back(-1);
05e
            add_suf(root, NULL, s.size() - 1, 0, 1e18);
1f2
7f3
        void rem suf(node*& x. int id) {
6cf
            if (x->sa != id) {
864
                 if (tag[id] < tag[x->sa]) return rem_suf(x->1, id);
                 return rem_suf(x->r, id);
e6f
2ae
            }
2cf
            node* nxt = next(x):
09b
            if (nxt) nxt->lcp = min(nxt->lcp, x->lcp), fix_path(nxt);
```

```
b20
             node *p = x - p, *tmp = x;
f3f
             if (!x->1 \text{ or } !x->r) {
2fd
                 x = x->1 ? x->1 : x->r;
753
                 if (x) x-p = p;
9d9
             } else {
7 f 7
                 for (tmp = x->1, p = x; tmp->r; tmp = tmp->r) p = tmp;
f2a
                 x->sa = tmp->sa, x->lcp = tmp->lcp;
482
                 if (tmp->1) tmp->1->p = p;
                 if (p->1 == tmp) p->1 = tmp->1;
14c
a94
                 else p \rightarrow r = tmp \rightarrow 1;
dc4
             }
b5e
             fix_path(p);
7c3
             delete tmp;
510
        }
15b
         void pop_front() {
abe
             if (!s.size()) return;
342
             s.pop_back();
436
             rem_suf(root, s.size());
сбе
             tag.pop_back();
        }
987
         int query(node* x, ll l, ll r, ll a, ll b) {
530
e51
             if (!x \text{ or } tag[x->sa] == -1 \text{ or } r < a \text{ or } b < 1) return
    s.size():
ef5
             if (a <= l and r <= b) return x->mi;
8eb
             int ans = s.size():
e1f
             if (a \le tag[x->sa] \text{ and } tag[x->sa] \le b) ans = min(ans,
   x \rightarrow lcp);
d99
             ans = min(ans, query(x->1, 1, tag[x->sa]-1, a, b));
             ans = min(ans, query(x->r, tag[x->sa]+1, r, a, b));
261
ba7
             return ans;
        }
4c8
588
        int query(int i, int j) { // lcp(s[i..], s[j..])
             if (i == j) return s.size() - i;
209
             ll a = tag[mirror(i)], b = tag[mirror(j)];
29e
710
             int ret = query(root, 0, 1e18, min(a, b)+1, max(a, b));
             return ret;
edf
84e
        // optional: get rank[i], sa[i] and lcp[i]
         int rank(int i) {
044
396
             i = mirror(i):
52f
             node* x = root:
7 c.9
             int ret = 0:
f4c
             while (x) {
33e
                 if (tag[x->sa] < tag[i]) {
f9d
                     ret += size(x->1)+1:
```

```
a91
                      x = x - > r;
6dc
                  } else x = x - > 1;
a19
edf
             return ret;
153
649
         pair<int, int> operator[](int i) {
52f
             node* x = root;
31e
             while (1) {
d4d
                  if (i < size(x->1)) x = x->1;
                  else {
4e6
85f
                      i \rightarrow size(x\rightarrow 1);
                      if (!i) return {mirror(x->sa), x->lcp};
040
                      i--, x = x->r:
b9b
7a2
             }
90c
         }
4c2 };
```

6.11 Suffix Automaton

```
// Automato que aceita os sufixos de uma string
// Todas as funcoes sao lineares
16e namespace sam {
        int cur, sz, len[2*MAX], link[2*MAX], acc[2*MAX];
c1a
        int nxt[2*MAX][26];
0b8
        void add(int c) {
e6a
            int at = cur;
17a
9a6
            len[sz] = len[cur]+1, cur = sz++;
            while (at != -1 and !nxt[at][c]) nxt[at][c] = cur, at =
500
   link[at]:
7ea
            if (at == -1) { link[cur] = 0; return; }
654
            int q = nxt[at][c];
fd9
            if (len[q] == len[at]+1) { link[cur] = q; return; }
            int qq = sz++;
31f
            len[qq] = len[at]+1, link[qq] = link[q];
2c3
9a9
            for (int i = 0; i < 26; i++) nxt[qq][i] = nxt[q][i];
            while (at != -1 and nxt[at][c] == q) nxt[at][c] = qq, at =
   link[at]:
8b8
            link[cur] = link[q] = qq;
61a
       }
94e
        void build(string& s) {
889
            cur = 0, sz = 0, len[0] = 0, link[0] = -1, sz++;
9fe
            for (auto i : s) add(i-'a');
17a
            int at = cur;
            while (at) acc[at] = 1, at = link[at];
121
```

```
0e7
        }
        // coisas que da pra fazer:
28 c
        11 distinct_substrings() {
04b
            11 \text{ ans} = 0:
a1e
            for (int i = 1; i < sz; i++) ans += len[i] - len[link[i]];</pre>
ba7
            return ans:
        }
0d7
a6c
        string longest_common_substring(string& S, string& T) {
419
             build(S):
111
            int at = 0, 1 = 0, ans = 0, pos = -1;
d59
            for (int i = 0; i < T.size(); i++) {</pre>
f2c
                 while (at and !nxt[at][T[i]-'a']) at = link[at]. 1 =
   len[at]:
efa
                 if (nxt[at][T[i]-'a']) at = nxt[at][T[i]-'a'], 1++;
749
                 else at = 0.1 = 0:
                 if (1 > ans) ans = 1, pos = i;
a1a
2b3
20f
            return T.substr(pos-ans+1, ans);
930
        }
46e
        11 dp[2*MAX];
455
        11 paths(int i) {
2a8
            auto& x = dp[i];
dee
            if (x) return x;
483
71c
            for (int j = 0; j < 26; j++) if (nxt[i][j]) x +=
    paths(nxt[i][i]);
ea5
            return x:
d88
        void kth_substring(int k, int at=0) { // k=1 : menor substring
105
   lexicog.
9d2
            for (int i = 0; i < 26; i++) if (k and nxt[at][i]) {</pre>
                 if (paths(nxt[at][i]) >= k) {
d58
d02
                     cout << char('a'+i);
                     kth_substring(k-1, nxt[at][i]);
c43
505
                     return;
69a
                 }
5f4
                 k -= paths(nxt[at][i]);
ef6
            }
a13
        }
c37 };
6.12
     Trie
// trie T() constroi uma trie para o alfabeto das letras minusculas
// trie T(tamanho do alfabeto, menor caracter) também pode ser usado
```

```
// T.insert(s) - O(|s|*sigma)
// T.erase(s) - O(|s|)
// T.find(s) retorna a posicao, O se nao achar - O(|s|)
// T.count_pref(s) numero de strings que possuem s como prefixo -
   0(|s|)
//
// Nao funciona para string vazia
ab5 struct trie {
        vector < vector < int >> to:
e1a
450
        vector < int > end , pref;
       int sigma; char norm;
bb1
        trie(int sigma_=26, char norm_='a') : sigma(sigma_),
   norm(norm) {
58a
            to = {vector < int > (sigma)};
86e
            end = \{0\}, pref = \{0\};
fe1
64e
        void insert(string s) {
            int x = 0;
c67
7e7
            for(auto c : s) {
                int &nxt = to[x][c-norm];
800
dd7
                if(!nxt) {
0aa
                    nxt = to.size();
526
                    to.push_back(vector<int>(sigma));
770
                    end.push_back(0), pref.push_back(0);
933
                }
827
                x = nxt, pref[x]++;
            }
34 c
            end[x]++;
e4e
e6b
6b2
        void erase(string s) {
c67
            int x = 0;
            for(char c : s) {
b4f
800
                int &nxt = to[x][c-norm];
                x = nxt, pref[x]--;
10c
d8e
                if(!pref[x]) nxt = 0;
885
            }
bf0
            end[x]--;
ddd
aee
        int find(string s) {
c67
            int x = 0:
7e7
            for(auto c : s) {
                x = to[x][c-norm];
2ec
a66
                if(!x) return 0;
e12
            }
ea5
            return x;
e77
        }
```

```
839
        int count_pref(string s) {
e2f
            return pref[find(s)];
f40
        }
979 };
6.13 Z
// z[i] = lcp(s, s[i..n))
//
// Complexidades:
//z - 0(|s|)
// \text{ match - } O(|s| + |p|)
a19 vector <int> get_z(string s) {
        int n = s.size();
2b1
        vector < int > z(n, 0);
fae
        int 1 = 0, r = 0;
        for (int i = 1; i < n; i++) {</pre>
6f5
0af
            if (i \le r) z[i] = min(r - i + 1, z[i - 1]);
457
            while (i + z[i] < n \text{ and } s[z[i]] == s[i + z[i]]) z[i]++;
65e
            if (i + z[i] - 1 > r) 1 = i, r = i + z[i] - 1:
        }
5cd
070
        return z;
74a }
   Primitivas
```

7.1 Aritmetica Modular

```
// O mod tem q ser primo
429 template <int p> struct mod_int {
        ll pow(ll b, ll e) {
02c
a63
            if (e == 0) return 1;
630
            ll r = pow(b*b%p, e/2);
475
            if (e\%2 == 1) r = (r*b)\%p;
4 c 1
            return r;
3ba
        }
        11 inv(11 b) { return pow(b, p-2); }
4d7
        using m = mod_int;
d93
        int v;
fe0
        mod_int() : v(0) {}
```

```
e12
        mod_int(ll v_) {
019
            if (v_ >= p or v_ <= -p) v_ %= p;
bc6
            if (v_{-} < 0) v_{-} += p;
2e7
            v = v_{-};
7f3
74d
        m& operator+=(const m &a) {
2fd
            v += a.v:
ba5
            if (v >= p) v -= p;
357
            return *this;
c8b
eff
        m& operator -= (const m &a) {
8b4
            v -= a.v:
            if (v < 0) v += p;
cc8
            return *this;
357
f8d
4c4
        m& operator*=(const m &a) {
8a5
            v = v * 11(a.v) \% p;
357
            return *this;
d4c
3f9
        m& operator/=(const m &a) {
            v = v* inv(a.v) \% p;
5d6
357
            return *this;
62d
d65
        m operator-(){ return m(-v); }
b3e
        m& operator^=(11 e) {
            if (e < 0){
06d
6e2
                v = inv(v):
00c
                e = -e:
275
ebf
            v = pow(v, e\%(p-1));
357
            return *this;
e84
        bool operator == (const m &a) { return v == a.v; }
423
        bool operator!=(const m &a) { return v != a.v: }
69f
1c6
        friend istream &operator>>(istream &in, m& a) {
d1c
            11 val: in >> val:
d48
            a = m(val);
091
            return in;
870
44f
        friend ostream &operator << (ostream &out, m a) {</pre>
5a0
            return out << a.v;</pre>
214
        friend m operator+(m a, m b) { return a+=b; }
399
f9e
        friend m operator-(m a, m b) { return a-=b; }
9c1
        friend m operator*(m a, m b) { return a*=b; }
51b
        friend m operator/(m a, m b) { return a/=b; }
```

```
08f
        friend m operator^(m a, ll e) { return a^=e; }
f94 };
055 typedef mod_int < (int) 1e9+7 > mint;
7.2 Big Integer
// Complexidades: (para n digitos)
// Soma, subtracao, comparacao - O(n)
// Multiplicacao - O(n log(n))
// Divisao, resto - O(n^2)
864 struct bint {
669
        static const int BASE = 1e9;
990
        vector<int> v;
3bd
        bool neg;
        bint() : neg(0) {}
609
d53
        bint(int val) : bint() { *this = val; }
        bint(long long val) : bint() { *this = val; }
a0f
        void trim() {
f42
            while (v.size() and v.back() == 0) v.pop_back();
df8
            if (!v.size()) neg = 0;
8e3
        }
        // converter de/para string | cin/cout
        bint(const char* s) : bint() { from_string(string(s)); }
294
        bint(const string& s) : bint() { from_string(s); }
548
        void from_string(const string& s) {
4ab
0a6
            v.clear(), neg = 0;
d72
            int ini = 0:
            while (ini < s.size() and (s[ini] == '-' or s[ini] == '+'
8e2
   or s[ini] == '0'))
71d
                if (s[ini++] == '-') neg = 1;
            for (int i = s.size()-1; i >= ini; i -= 9) {
883
05e
                int at = 0:
5b1
                for (int j = max(ini, i - 8); j \le i; j++) at = 10*at
   + (s[i]-'0');
1fd
                v.push_back(at);
a5a
df8
            if (!v.size()) neg = 0;
        }
e9a
2ff
        string to_string() const {
8be
            if (!v.size()) return "0";
793
            string ret;
            if (neg) ret += '-';
73e
```

```
3e9
            for (int i = v.size()-1; i >= 0; i--) {
582
                string at = ::to_string(v[i]);
                int add = 9 - at.size();
ced
                if (i+1 < v.size()) for (int j = 0; j < add; j++) ret
   += '0';
f9f
                ret += at;
f64
            }
edf
            return ret;
770
d2f
        friend istream& operator>>(istream& in, bint& val) {
eb6
            string s; in >> s;
966
            val = s:
091
            return in:
328
        friend ostream& operator<<(ostream& out, const bint& val) {</pre>
99d
8b9
            string s = val.to_string();
396
            out << s;
fe8
            return out;
ce1
        // operators
        friend bint abs(bint val) {
60a
c5f
            val.neg = 0;
d94
            return val;
44b
bee
       friend bint operator-(bint val) {
815
            if (val != 0) val.neg ^= 1;
d94
            return val:
326
        bint& operator=(const bint& val) { v = val.v, neg = val.neg;
41f
   return *this: }
249
        bint& operator=(long long val) {
            v.clear(), neg = 0;
0a6
3a6
            if (val < 0) neg = 1. val *= -1:
            for (; val; val /= BASE) v.push_back(val % BASE);
fdc
357
            return *this;
220
        }
        int cmp(const bint& r) const { // menor: -1 | igual: 0 |
   maior: 1
b14
            if (neg != r.neg) return neg ? -1 : 1;
            if (v.size() != r.v.size()) {
0bb
                int ret = v.size() < r.v.size() ? -1 : 1;</pre>
ff7
                return neg ? -ret : ret;
91b
            }
1f6
478
            for (int i = int(v.size())-1; i >= 0; i--) {
                if (v[i] != r.v[i]) {
405
                    int ret = v[i] < r.v[i] ? -1 : 1;</pre>
2e5
```

```
91b
                    return neg ? -ret : ret;
                }
9a9
c32
            }
bb3
            return 0;
07d
        friend bool operator < (const bint& 1, const bint& r) { return
152
   1.cmp(r) == -1: 
        friend bool operator > (const bint& 1, const bint& r) { return
   1.cmp(r) == 1; }
        friend bool operator <= (const bint& 1, const bint& r) { return
   1.cmp(r) <= 0; }
        friend bool operator >= (const bint& 1, const bint& r) { return
   1.cmp(r) >= 0: }
a67
        friend bool operator == (const bint& 1, const bint& r) { return
   1.cmp(r) == 0; }
        friend bool operator!=(const bint& 1, const bint& r) { return
   1.cmp(r) != 0; }
38e
        bint& operator +=(const bint& r) {
            if (!r.v.size()) return *this;
6bf
a 93
            if (neg != r.neg) return *this -= -r;
256
            for (int i = 0, c = 0; i < r.v.size() or c; i++) {
                if (i == v.size()) v.push_back(0);
e28
08f
                v[i] += c + (i < r.v.size() ? r.v[i] : 0);
                if ((c = v[i] >= BASE)) v[i] -= BASE;
baa
8bb
357
            return *this;
ab1
        }
54c
        friend bint operator+(bint a, const bint& b) { return a += b; }
9c8
        bint& operator -=(const bint& r) {
            if (!r.v.size()) return *this;
6bf
524
            if (neg != r.neg) return *this += -r;
            if ((!neg and *this < r) or (neg and r < *this)) {
358
b10
                *this = r - *this:
                neg ^= 1;
a10
357
                return *this;
            }
807
256
            for (int i = 0, c = 0; i < r.v.size() or c; i++) {
9ef
                v[i] = c + (i < r.v.size() ? r.v[i] : 0);
c8c
                if ((c = v[i] < 0)) v[i] += BASE;
            }
687
0eb
            trim();
357
            return *this;
f72
        }
        friend bint operator-(bint a, const bint& b) { return a -= b; }
f44
        // operators de * / %
```

```
6b0
        bint& operator *=(int val) {
bca
            if (val < 0) val *= -1, neg ^= 1;</pre>
566
            for (int i = 0, c = 0; i < v.size() or c; i++) {
                if (i == v.size()) v.push_back(0);
e28
352
                long long at = (long long) v[i] * val + c;
                v[i] = at % BASE;
6a3
b3d
                c = at / BASE:
cb1
            }
0eb
            trim();
357
            return *this;
a57
480
        friend bint operator *(bint a, int b) { return a *= b; }
d5c
        friend bint operator *(int a, bint b) { return b *= a; }
        using cplx = complex < double >;
13b
        void fft(vector<cplx>& a, bool f, int N, vector<int>& rev)
bfb
   const {
            for (int i = 0; i < N; i++) if (i < rev[i]) swap(a[i],
bc7
   a[rev[i]]);
            vector < cplx > roots(N);
bad
            for (int n = 2; n <= N; n *= 2) {</pre>
192
                const static double PI = acos(-1);
4e9
                for (int i = 0; i < n/2; i++) {
71a
                    double alpha = (2*PI*i)/n;
40d
1a1
                    if (f) alpha = -alpha;
3f6
                    roots[i] = cplx(cos(alpha), sin(alpha));
f16
3e9
                for (int pos = 0; pos < N; pos += n)
898
                    for (int 1 = pos, r = pos+n/2, m = 0; m < n/2;
   1++, r++, m++) {
297
                         auto t = roots[m]*a[r];
254
                        a[r] = a[1] - t:
b8f
                         a[1] = a[1] + t;
                    }
b0d
e07
            }
            if (!f) return;
3f1
08b
            auto invN = cplx(1)/cplx(N);
873
            for (int i = 0; i < N; i++) a[i] *= invN;</pre>
c75
0e0
        vector < long long > convolution(const vector < int > & a, const
   vector < int > & b) const {
            vector < cplx > 1(a.begin(), a.end()), r(b.begin(), b.end());
ff9
            int ln = l.size(), rn = r.size(), N = ln+rn+1, n = 1,
996
   log_n = 0;
            while (n \le N) n \le 1, \log_n + 1
821
808
            vector < int > rev(n);
            for (int i = 0; i < n; i++) {</pre>
603
434
                rev[i] = 0:
```

```
f44
                for (int j = 0; j < log_n; j++) if (i >> j & 1)
4ff
                    rev[i] = 1 << (log_n-1-j);
256
            }
230
            l.resize(n), r.resize(n);
a89
            fft(1, false, n, rev), fft(r, false, n, rev);
917
            for (int i = 0; i < n; i++) l[i] *= r[i];</pre>
88b
            fft(l, true, n, rev);
7ae
            vector < long long > ret;
c14
            for (auto& i : 1) ret.push_back(round(i.real()));
edf
            return ret:
917
        }
633
        vector < int > convert base (const vector < int > & a. int from. int
   to) const {
498
            static vector<long long> pot(10, 1);
671
            if (pot[1] == 1) for (int i = 1; i < 10; i++) pot[i] =</pre>
   10*pot[i-1];
4b8
            vector < int > ret;
            long long at = 0;
156
608
            int digits = 0;
            for (int i : a) {
941
                at += i * pot[digits];
412
035
                digits += from;
684
                while (digits >= to) {
0c8
                    ret.push_back(at % pot[to]);
cf9
                    at /= pot[to];
fd4
                    digits -= to;
122
                }
87b
            }
944
            ret.push_back(at);
384
            while (ret.size() and ret.back() == 0) ret.pop_back();
            return ret:
edf
090
        bint operator*(const bint& r) const { // O(n log(n))
edb
2af
            bint ret:
968
            ret.neg = neg ^ r.neg;
d5d
            auto conv = convolution(convert_base(v, 9, 4),
   convert_base(r.v, 9, 4));
            long long c = 0;
a0e
a74
            for (auto i : conv) {
f6d
                long long at = i+c;
                ret.v.push_back(at % 10000);
4cb
a25
                c = at / 10000;
773
            for (; c; c /= 10000) ret.v.push_back(c%10000);
3cb
0e2
            ret.v = convert_base(ret.v, 4, 9);
25 c
            if (!ret.v.size()) ret.neg = 0;
            return ret:
edf
```

```
c6b
        bint& operator*=(const bint& r) { return *this = *this * r; };
359
9a3
        bint& operator/=(int val) {
            if (val < 0) neg ^= 1, val *= -1;</pre>
d9a
            for (int i = int(v.size())-1, c = 0; i >= 0; i--) {
f18
                 long long at = v[i] + c * (long long) BASE;
2a7
e02
                v[i] = at / val:
                c = at % val;
fb1
            }
fdb
            trim();
0eb
357
            return *this;
db6
e74
        friend bint operator/(bint a. int b) { return a /= b: }
        int operator %=(int val) {
4a9
23b
            if (val < 0) val *= -1;</pre>
            long long at = 0;
156
f31
            for (int i = int(v.size())-1; i >= 0; i--)
                 at = (BASE * at + v[i]) \% val:
1b3
d22
            if (neg) at *= -1;
            return at:
ce6
4b4
        }
        friend int operator % (bint a, int b) { return a % = b; }
2fb
        friend pair < bint, bint > divmod(const bint& a_, const bint& b_)
   { // O(n^2)}
            if (a_ == 0) return {0, 0};
611
d8a
            int norm = BASE / (b .v.back() + 1);
            bint a = abs(a<sub>_</sub>) * norm;
b4e
027
            bint b = abs(b) * norm:
14d
            bint q, r;
            for (int i = a.v.size() - 1; i >= 0; i--) {
c91
                r *= BASE. r += a.v[i]:
b71
4ff
                long long upper = b.v.size() < r.v.size() ?</pre>
   r.v[b.v.size()] : 0;
                int lower = b.v.size() - 1 < r.v.size() ?</pre>
86d
   r.v[b.v.size() - 1] : 0;
431
                int d = (upper * BASE + lower) / b.v.back();
5d4
                r \rightarrow b*d:
                 while (r < 0) r += b, d--; // roda O(1) vezes
30f
738
                 q.v.push_back(d);
c6a
            }
a48
            reverse(q.v.begin(), q.v.end());
ae2
            q.neg = a_.neg ^ b_.neg;
            r.neg = a_.neg;
88b
8e5
            q.trim(), r.trim();
0ef
            return {q, r / norm};
4fd
        bint operator/(const bint& val) { return divmod(*this,
1d8
```

```
val).first; }
        bint& operator/=(const bint& val) { return *this = *this /
7f9
   val; }
1f9
        bint operator % (const bint & val) { return divmod(*this,
   val).second: }
        bint& operator%=(const bint& val) { return *this = *this %
df5
   val: }
6c3 };
7.3 Fracao
// Funciona com o Big Int
a4e template < typename T = int > struct frac {
        T num, den;
a40
e3f
        template < class U, class V>
61d
        frac(U num_ = 0, V den_ = 1) : num(num_), den(den_) {
            assert(den != 0):
bad
583
            if (den < 0) num *= -1, den *= -1;
            T g = gcd(abs(num), den);
a51
572
            num \neq g, den \neq g;
        }
fbf
51f
        friend bool operator<(const frac& 1, const frac& r) {</pre>
fa0
            return l.num * r.den < r.num * l.den;</pre>
a4e
4b5
        friend frac operator+(const frac& 1, const frac& r) {
            return {1.num*r.den + 1.den*r.num, 1.den*r.den};
b61
25f
74d
        friend frac operator-(const frac& 1, const frac& r) {
2cd
            return {1.num*r.den - 1.den*r.num, 1.den*r.den};
8a7
        friend frac operator*(const frac& 1, const frac& r) {
c80
            return {1.num*r.num, 1.den*r.den};
510
14b
        friend frac operator/(const frac& 1, const frac& r) {
a1b
            return {1.num*r.den, 1.den*r.num};
8f3
b2c
        }
        friend ostream& operator << (ostream& out, frac f) {</pre>
012
            out << f.num << ',' << f.den;
37a
fe8
            return out;
b49
        }
cdb };
7.4 Geometria
```

c83 typedef double ld;

```
e3b const ld DINF = 1e18;
43a const ld pi = acos(-1.0);
107 const ld eps = 1e-9;
b32 #define sq(x) ((x)*(x))
d97 bool eq(ld a, ld b) {
       return abs(a - b) <= eps;</pre>
ba0
bfc }
b2a struct pt { // ponto
       ld x, y;
3dd
       pt(1d x_{-} = 0, 1d y_{-} = 0) : x(x_{-}), y(y_{-}) {}
       bool operator < (const pt p) const {</pre>
5bc
            if (!eq(x, p.x)) return x < p.x;
059
f98
           if (!eq(y, p.y)) return y < p.y;</pre>
bb3
          return 0;
f61
a83
        bool operator == (const pt p) const {
ed0
           return eq(x, p.x) and eq(y, p.y);
589
        pt operator + (const pt p) const { return pt(x+p.x, y+p.y); }
cb9
       pt operator - (const pt p) const { return pt(x-p.x, y-p.y); }
a24
4a8
       pt operator * (const ld c) const { return pt(x*c , y*c ); }
a60
       pt operator / (const ld c) const { return pt(x/c , y/c ); }
3b6
       ld operator * (const pt p) const { return x*p.x + y*p.y; }
        ld operator ^ (const pt p) const { return x*p.y - y*p.x; }
6df
       friend istream& operator >> (istream& in, pt& p) {
5ed
            return in >> p.x >> p.y;
e37
e45
       }
a8b }:
b3a struct line { // reta
730
       pt p, q;
0d6
       line() {}
4b8 line(pt p_, pt q_) : p(p_), q(q_) {}
8d7
       friend istream& operator >> (istream& in, line& r) {
4cb
            return in >> r.p >> r.q;
858
7ab };
// PONTO & VETOR
364 ld dist(pt p, pt q) { // distancia
       return hypot(p.y - q.y, p.x - q.x);
5f3
c68 }
```

```
9d7 ld dist2(pt p, pt q) { // quadrado da distancia
f24
        return sq(p.x - q.x) + sq(p.y - q.y);
80f }
483 ld norm(pt v) { // norma do vetor
        return dist(pt(0, 0), v);
cf7 }
589 ld angle(pt v) { // angulo do vetor com o eixo x
        ld ang = atan2(v.y, v.x);
6f8
        if (ang < 0) ang += 2*pi;
19c
        return ang:
404 }
298 ld sarea(pt p, pt q, pt r) { // area com sinal
       return ((q-p)^(r-q))/2;
1b1 }
e32 bool col(pt p, pt q, pt r) { // se p, q e r sao colin.
        return eq(sarea(p, q, r), 0);
98c }
Ocd bool ccw(pt p, pt q, pt r) \{ // \text{ se p, q, r sao ccw} \}
fa7
        return sarea(p, q, r) > eps;
85d }
1ef pt rotate(pt p, ld th) { // rotaciona o ponto th radianos
        return pt(p.x * cos(th) - p.y * sin(th),
e5c
ff1
                p.x * sin(th) + p.y * cos(th));
41a }
ab1 pt rotate90(pt p) { // rotaciona 90 graus
        return pt(-p.y, p.x);
e4a }
// RETA
edc bool isvert(line r) { // se r eh vertical
        return eq(r.p.x, r.q.x);
Ofb }
099 bool isinseg(pt p, line r) { // se p pertence ao seg de r
f65
        pt a = r.p - p, b = r.q - p;
        return eq((a \hat{b}), 0) and (a * b) < eps;
b04
726 }
98d ld get_t(pt v, line r) { // retorna t tal que t*v pertence a reta r
```

```
6ee
        return (r.p^r.q) / ((r.p-r.q)^v);
a0a }
256 pt proj(pt p, line r) { // projecao do ponto p na reta r
       if (r.p == r.q) return r.p;
bea
       r.q = r.q - r.p; p = p - r.p;
97a
9f8
       pt proj = r.q * ((p*r.q) / (r.q*r.q));
2cd
        return proj + r.p;
232 }
d5c pt inter(line r, line s) { // r inter s
       if (eq((r.p - r.q) ^ (s.p - s.q), 0)) return pt(DINF, DINF);
205
       r.q = r.q - r.p, s.p = s.p - r.p, s.q = s.q - r.p;
543
        return r.q * get_t(r.q, s) + r.p;
111 }
676 bool interseg(line r, line s) { // se o seg de r intersecta o seg
   de s
        if (isinseg(r.p, s) or isinseg(r.q, s)
19b
            or isinseg(s.p, r) or isinseg(s.q, r)) return 1;
c21
9fa
        return ccw(r.p, r.q, s.p) != ccw(r.p, r.q, s.q) and
413
                ccw(s.p, s.q, r.p) != ccw(s.p, s.q, r.q);
359 }
fcb ld disttoline(pt p, line r) { // distancia do ponto a reta
        return 2 * abs(sarea(p, r.p, r.g)) / dist(r.p, r.g);
89a
1b7 }
bcc ld disttoseg(pt p, line r) { // distancia do ponto ao seg
73d
       if ((r.q - r.p)*(p - r.p) < 0) return dist(r.p, p);
951
        if ((r.p - r.q)*(p - r.q) < 0) return dist(r.q, p);
a19
        return disttoline(p, r);
367 }
11d ld distseg(line a, line b) { // distancia entre seg
4df
        if (interseg(a, b)) return 0;
349
       ld ret = DINF:
       ret = min(ret, disttoseg(a.p, b));
341
       ret = min(ret, disttoseg(a.q, b));
ceb
       ret = min(ret, disttoseg(b.p, a));
093
        ret = min(ret, disttoseg(b.q, a));
448
edf
        return ret;
222 }
```

```
// POLIGONO
// corta poligono com a reta r deixando os pontos p tal que
// ccw(r.p, r.q, p)
1a9 vector<pt> cut_polygon(vector<pt> v, line r) { // O(n)
8af
        vector<pt> ret;
8a4
        for (int j = 0; j < v.size(); j++) {</pre>
            if (ccw(r.p, r.q, v[j])) ret.push_back(v[j]);
dac
            if (v.size() == 1) continue;
dce
030
            line s(v[j], v[(j+1)%v.size()]);
ae3
            pt p = inter(r, s);
a3d
            if (isinseg(p, s)) ret.push_back(p);
d44
8a1
        ret.erase(unique(ret.begin(), ret.end()), ret.end());
24d
        if (ret.size() > 1 and ret.back() == ret[0]) ret.pop_back();
edf
        return ret:
253 }
// distancia entre os retangulos a e b (lados paralelos aos eixos)
// assume que ta representado (inferior esquerdo, superior direito)
5f5 ld dist_rect(pair<pt, pt> a, pair<pt, pt> b) {
        ld hor = 0. vert = 0:
080
34b
        if (a.second.x < b.first.x) hor = b.first.x - a.second.x;</pre>
f5f
        else if (b.second.x < a.first.x) hor = a.first.x - b.second.x;</pre>
4fd
        if (a.second.y < b.first.y) vert = b.first.y - a.second.y;</pre>
80a
        else if (b.second.y < a.first.y) vert = a.first.y - b.second.y;</pre>
96f
        return dist(pt(0, 0), pt(hor, vert));
630 }
13d ld polarea(vector<pt> v) { // area do poligono
9 c 5
        ld ret = 0:
сбе
        for (int i = 0; i < v.size(); i++)</pre>
80f
            ret += sarea(pt(0, 0), v[i], v[(i + 1) % v.size()]);
        return abs(ret):
d03
5df }
// se o ponto ta dentro do poligono: retorna O se ta fora,
// 1 se ta no interior e 2 se ta na borda
8e7 int inpol(vector\phit>& v, pt p) { // O(n)
8de
        int qt = 0;
f14
        for (int i = 0: i < v.size(): i++) {</pre>
            if (p == v[i]) return 2;
bda
            int j = (i+1)%v.size();
6af
e38
            if (eq(p.y, v[i].y) and eq(p.y, v[j].y)) {
97f
                if ((v[i]-p)*(v[j]-p) < eps) return 2;
5e2
                continue:
48b
            }
```

```
388
            bool baixo = v[i].y+eps < p.y;</pre>
            if (baixo == (v[j].y+eps < p.y)) continue;</pre>
464
            auto t = (p-v[i])^(v[j]-v[i]);
366
1b4
            if (eq(t, 0)) return 2;
839
            if (baixo == (t > eps)) qt += baixo ? 1 : -1;
d13
b84
        return qt != 0;
a64 }
6ff bool interpol(vector<pt> v1, vector<pt> v2) { // se dois poligonos
   se intersectam - O(n*m)
        int n = v1.size(), m = v2.size();
7d1
c36
        for (int i = 0: i < n: i++) if (inpol(v2. v1[i])) return 1:
        for (int i = 0; i < n; i++) if (inpol(v1, v2[i])) return 1;</pre>
ab8
        for (int i = 0; i < n; i++) for (int j = 0; j < m; j++)
523
            if (interseg(line(v1[i], v1[(i+1)%n]), line(v2[j],
0c8
   v2[(j+1)%m]))) return 1;
        return 0:
bb3
c58 }
494 ld distpol(vector<pt> v1, vector<pt> v2) { // distancia entre
   poligonos
f6b
        if (interpol(v1, v2)) return 0;
349
        ld ret = DINF;
        for (int i = 0; i < v1.size(); i++) for (int j = 0; j <</pre>
   v2.size(): i++)
            ret = min(ret, distseg(line(v1[i], v1[(i + 1) %
   v1.size()]),
                         line(v2[j], v2[(j + 1) % v2.size()])));
949
edf
        return ret;
125 }
138 vector <pt> convex_hull(vector <pt> v) { // convex hull - O(n log(n))
fca
        sort(v.begin(), v.end());
d76
        v.erase(unique(v.begin(), v.end()), v.end());
52d
        if (v.size() <= 1) return v;</pre>
526
        vector < pt > 1, u;
f14
        for (int i = 0; i < v.size(); i++) {</pre>
            while (1.size() > 1 and !ccw(1.end()[-2], 1.end()[-1],
fb2
   v[i]))
364
                1.pop_back();
c35
            l.push_back(v[i]);
58e
        for (int i = v.size() - 1; i >= 0; i--) {
3e9
f19
            while (u.size() > 1 \text{ and } !ccw(u.end()[-2], u.end()[-1].
```

```
v[i]))
7a8
                 u.pop_back();
            u.push_back(v[i]);
a95
        }
0b8
        1.pop_back(); u.pop_back();
cfc
        for (pt i : u) l.push_back(i);
82b
792
        return 1:
10d }
483 struct convex_pol {
f50
        vector<pt> pol;
        // nao pode ter ponto colinear no convex hull
d98
        convex_pol() {}
        convex_pol(vector < pt > v) : pol(convex_hull(v)) {}
a04
        // se o ponto ta dentro do hull - O(\log(n))
        bool is_inside(pt p) {
8af
b6e
            if (pol.size() == 0) return false;
             if (pol.size() == 1) return p == pol[0];
eae
67f
            int 1 = 1, r = pol.size();
40c
            while (1 < r) {
                 int m = (1+r)/2:
ee4
48f
                 if (ccw(p, pol[0], pol[m])) 1 = m+1;
ef3
                 else r = m:
91c
            }
00a
            if (1 == 1) return isinseg(p, line(pol[0], pol[1]));
            if (1 == pol.size()) return false;
9e7
            return !ccw(p, pol[1], pol[1-1]);
1c0
6b0
        }
        // ponto extremo em relacao a cmp(p, q) = p mais extremo q
        // (copiado de https://github.com/gustavoM32/caderno-zika)
719
        int extreme(const function < bool(pt, pt) > & cmp) {
b1c
            int n = pol.size();
             auto extr = [&](int i, bool& cur_dir) {
4a2
22a
                 \operatorname{cur\_dir} = \operatorname{cmp}(\operatorname{pol}[(i+1)\%n], \operatorname{pol}[i]);
61a
                 return !cur_dir and !cmp(pol[(i+n-1)%n], pol[i]);
364
            };
63d
            bool last_dir, cur_dir;
a0d
            if (extr(0, last_dir)) return 0;
993
            int 1 = 0, r = n:
            while (1+1 < r) {
ead
                 int m = (1+r)/2:
ee4
f29
                 if (extr(m, cur_dir)) return m;
44a
                 bool rel_dir = cmp(pol[m], pol[l]);
b18
                 if ((!last_dir and cur_dir) or
261
                          (last dir == cur dir and rel dir == cur dir)) {
```

```
8a6
                    1 = m:
1 f 1
                    last_dir = cur_dir;
94a
                } else r = m;
606
            }
792
            return 1;
56c
316
        int max dot(pt v) {
ec1
            return extreme([&](pt p, pt q) { return p*v > q*v; });
3b7
a 54
        pair < int , int > tangents(pt p) {
ffb
            auto L = [k](pt q, pt r) \{ return ccw(p, r, q); \};
8fd
            auto R = [\&](pt q, pt r) \{ return ccw(p, q, r); \};
fa8
            return {extreme(L). extreme(R)}:
736
        }
3ec };
// CIRCUNFERENCIA
911 pt getcenter(pt a, pt b, pt c) { // centro da circunf dado 3 pontos
174
        b = (a + b) / 2:
       c = (a + c) / 2;
2ae
98b
       return inter(line(b, b + rotate90(a - b)),
                line(c, c + rotate90(a - c)));
3f8
a12 }
4b3 vector <pt> circ_line_inter(pt a, pt b, pt c, ld r) { // intersecao
   da circunf (c, r) e reta ab
8af
        vector <pt> ret:
        b = b-a, a = a-c;
f2b
       1d A = b*b;
4b1
       1d B = a*b:
2e9
       1d C = a*a - r*r;
       1d D = B*B - A*C:
       if (D < -eps) return ret;</pre>
818
       ret.push_back(c+a+b*(-B+sqrt(D+eps))/A);
dc5
        if (D > eps) ret.push_back(c+a+b*(-B-sqrt(D))/A);
20e
edf
        return ret;
cd8 }
ad2 vector <pt> circ_inter(pt a, pt b, ld r, ld R) { // intersecao da
   circunf (a, r) e (b, R)
       vector <pt> ret;
8af
b7e
       1d d = dist(a, b);
       if (d > r+R \text{ or } d+min(r, R) < max(r, R)) return ret;
398
       1d x = (d*d-R*R+r*r)/(2*d);
       1d y = sqrt(r*r-x*x);
183
325
        pt v = (b-a)/d;
```

```
76e
        ret.push_back(a+v*x + rotate90(v)*y);
        if (y > 0) ret.push_back(a+v*x - rotate90(v)*y);
2cb
        return ret;
edf
fb1 }
6e0 bool operator <(const line& a, const line& b) { // comparador pra
   reta
        // assume que as retas tem p < q</pre>
        pt v1 = a.q - a.p, v2 = b.q - b.p;
a13
f82
        if (!eq(angle(v1), angle(v2))) return angle(v1) < angle(v2);</pre>
780
        return ccw(a.p, a.q, b.p); // mesmo angulo
27e }
b14 bool operator ==(const line& a, const line& b) {
        return !(a < b) and !(b < a):
76c
449 }
// comparador pro set pra fazer sweep line com segmentos
2c4 struct cmp_sweepline {
        bool operator () (const line& a, const line& b) const {
d80
            // assume que os segmentos tem p < q</pre>
            if (a.p == b.p) return ccw(a.p, a.q, b.q);
191
231
            if (!eq(a.p.x, a.q.x)) and (eq(b.p.x, b.q.x)) or a.p.x+eps <
   b.p.x))
780
                 return ccw(a.p, a.q, b.p);
            return ccw(a.p, b.q, b.p);
dc0
243
367 };
// comparador pro set pra fazer sweep angle com segmentos
bef pt dir;
5b0 struct cmp_sweepangle {
d80
        bool operator () (const line& a, const line& b) const {
522
            return get_t(dir, a) + eps < get_t(dir, b);</pre>
653
        }
97f };
7.5 Geometria - inteiro
2de #define sq(x) ((x)*(11)(x))
// 840720
b2a struct pt { // ponto
e91
        int x, y;
df1
        pt(int x_{-} = 0, int y_{-} = 0) : x(x_{-}), y(y_{-}) {}
5bc
        bool operator < (const pt p) const {</pre>
95a
            if (x != p.x) return x < p.x;
89 c
            return y < p.y;</pre>
```

```
dcd
a83
       bool operator == (const pt p) const {
d74
           return x == p.x and y == p.y;
7b4
       }
       pt operator + (const pt p) const { return pt(x+p.x, y+p.y); }
cb9
       pt operator - (const pt p) const { return pt(x-p.x, y-p.y); }
a24
       pt operator * (const int c) const { return pt(x*c, y*c); }
0ef
       11 operator * (const pt p) const { return x*(11)p.x +
   v*(11)p.v: }
       11 operator ^ (const pt p) const { return x*(11)p.y -
   v*(11)p.x; }
     friend istream& operator >> (istream& in, pt& p) {
e37
           return in >> p.x >> p.v;
e45
       }
840 };
// 7ab617
b3a struct line { // reta
730
       pt p, q;
0d6 line() {}
4b8 line(pt p_, pt q_) : p(p_), q(q_) {}
8d7
      friend istream& operator >> (istream& in, line& r) {
4cb
           return in >> r.p >> r.q;
858
       }
7ab }:
// PONTO & VETOR
// 51563e
ea8 ll dist2(pt p, pt q) { // quadrado da distancia
       return sq(p.x - q.x) + sq(p.y - q.y);
515 }
// bf431d
5a2 ll sarea2(pt p, pt q, pt r) \{ // 2 * area com sinal \}
       return (q-p)^(r-q);
bf4 }
// a082d3
e32 bool col(pt p, pt q, pt r) { // se p, q e r sao colin.
       return sarea2(p, q, r) == 0;
a08 }
// 42bb09
Ocd bool ccw(pt p, pt q, pt r) { // se p, q, r sao ccw
       return sarea2(p, q, r) > 0;
42b }
```

```
// fcf924
c31 int quad(pt p) { // quadrante de um ponto
       return (p.x<0)^3*(p.y<0);
fcf }
// 77187b
2df bool compare_angle(pt p, pt q) { // retorna se ang(p) < ang(q)
9fc
        if (quad(p) != quad(q)) return quad(p) < quad(q);</pre>
        return ccw(q, pt(0, 0), p);
ea1
771 }
// e4ad5e
ab1 pt rotate90(pt p) { // rotaciona 90 graus
       return pt(-p.y, p.x);
e4a }
// RETA
// c9f07f
099 bool isinseg(pt p, line r) { // se p pertence ao seg de r
        pt a = r.p - p, b = r.q - p;
f65
        return (a ^ b) == 0 and (a * b) <= 0;
2ac
c9f }
// 35998c
676 bool interseg(line r, line s) { // se o seg de r intersecta o seg
   de s
19b
      if (isinseg(r.p, s) or isinseg(r.q, s)
            or isinseg(s.p, r) or isinseg(s.q, r)) return 1;
c21
9fa
        return ccw(r.p, r.q, s.p) != ccw(r.p, r.q, s.q) and
413
                ccw(s.p, s.q, r.p) != ccw(s.p, s.q, r.q);
359 }
// dd8702
9e0 int segpoints(line r) { // numero de pontos inteiros no segmento
        return 1 + \_gcd(abs(r.p.x - r.q.x), abs(r.p.y - r.q.y));
dd8 }
88a double get_t(pt v, line r) { // retorna t tal que t*v pertence a
       return (r.p^r.q) / (double) ((r.p-r.q)^v);
d27 }
// POLIGONO
```

```
// quadrado da distancia entre os retangulos a e b (lados paralelos
   aos eixos)
// assume que ta representado (inferior esquerdo, superior direito)
// e13018
485 ll dist2_rect(pair<pt, pt> a, pair<pt, pt> b) {
c59
        int hor = 0. vert = 0:
34b
        if (a.second.x < b.first.x) hor = b.first.x - a.second.x;</pre>
f5f
        else if (b.second.x < a.first.x) hor = a.first.x - b.second.x;</pre>
4fd
        if (a.second.y < b.first.y) vert = b.first.y - a.second.y;</pre>
80a
        else if (b.second.y < a.first.y) vert = a.first.y - b.second.y;</pre>
869
        return sq(hor) + sq(vert);
e13 }
// d5f693
9c3 ll polarea2(vector<pt> v) { // 2 * area do poligono
b73
        ll ret = 0;
        for (int i = 0; i < v.size(); i++)</pre>
c6e
532
            ret += sarea2(pt(0, 0), v[i], v[(i + 1) % v.size()]);
d03
        return abs(ret):
d5f }
// se o ponto ta dentro do poligono: retorna O se ta fora,
// 1 se ta no interior e 2 se ta na borda
// afd587
8e7 int inpol(vector\phit>& v, pt p) { // O(n)
8de
        int qt = 0;
f14
        for (int i = 0: i < v.size(): i++) {</pre>
            if (p == v[i]) return 2;
bda
6af
            int j = (i+1)%v.size();
            if (p.y == v[i].y and p.y == v[j].y) {
cc6
547
                if ((v[i]-p)*(v[j]-p) <= 0) return 2;</pre>
5e2
                continue:
b47
            }
            bool baixo = v[i].y < p.y;</pre>
78c
057
            if (baixo == (v[j].y < p.y)) continue;</pre>
366
            auto t = (p-v[i])^(v[j]-v[i]);
2ad
            if (!t) return 2;
0bb
            if (baixo == (t > 0)) qt += baixo ? 1 : -1;
9cf
b84
        return qt != 0;
afd }
// 10d7e0
138 vector<pt> convex_hull(vector<pt> v) { // convex hull - O(n log(n))
fca
        sort(v.begin(), v.end());
d76
        v.erase(unique(v.begin(), v.end()), v.end());
```

```
52d
        if (v.size() <= 1) return v;</pre>
526
        vector<pt> 1, u;
f14
        for (int i = 0; i < v.size(); i++) {</pre>
fb2
            while (1.size() > 1 \text{ and } !ccw(1.end()[-2], 1.end()[-1],
   v[i]))
364
                 1.pop_back();
c35
            1.push_back(v[i]);
        }
58e
3e9
        for (int i = v.size() - 1; i >= 0; i--) {
            while (u.size() > 1 \text{ and } !ccw(u.end()[-2], u.end()[-1],
f19
   v[i]))
7a8
                 u.pop_back();
a95
            u.push_back(v[i]);
0b8
cfc
        1.pop_back(); u.pop_back();
        for (pt i : u) l.push_back(i);
82b
792
        return 1;
10d }
// af2d96
786 ll interior_points(vector<pt> v) { // pontos inteiros dentro de um
   poligono simples
c4e
        11 b = 0:
        for (int i = 0; i < v.size(); i++)</pre>
c6e
0 ce
            b += segpoints(line(v[i], v[(i+1)\%v.size()])) - 1;
a1c
        return (polarea2(v) - b) / 2 + 1;
af2 }
483 struct convex_pol {
f50
        vector<pt> pol;
        // nao pode ter ponto colinear no convex hull
        convex_pol() {}
d98
a04
        convex_pol(vector<pt> v) : pol(convex_hull(v)) {}
        // se o ponto ta dentro do hull - O(log(n))
        // 6b097f
8af
        bool is_inside(pt p) {
b6e
            if (pol.size() == 0) return false;
eae
            if (pol.size() == 1) return p == pol[0];
            int 1 = 1, r = pol.size();
67 f
            while (1 < r) {
40c
                int m = (1+r)/2:
ee4
48f
                if (ccw(p, pol[0], pol[m])) l = m+1;
ef3
                 else r = m;
            }
91 c
00a
            if (1 == 1) return isinseg(p, line(pol[0], pol[1]));
```

```
9e7
            if (1 == pol.size()) return false;
1c0
            return !ccw(p, pol[1], pol[1-1]);
6b0
        // ponto extremo em relacao a cmp(p, q) = p mais extremo q
        // (copiado de https://github.com/gustavoM32/caderno-zika)
        // 56ccd2
719
        int extreme(const function < bool(pt. pt) > & cmp) {
b1c
            int n = pol.size();
4a2
            auto extr = [&](int i, bool& cur_dir) {
22a
                 \operatorname{cur\_dir} = \operatorname{cmp}(\operatorname{pol}[(i+1)\%n], \operatorname{pol}[i]);
61a
                 return !cur_dir and !cmp(pol[(i+n-1)%n], pol[i]);
364
            }:
63d
            bool last dir. cur dir:
            if (extr(0, last_dir)) return 0;
a0d
993
            int 1 = 0, r = n;
            while (1+1 < r) {
ead
ee4
                int m = (1+r)/2;
                if (extr(m, cur_dir)) return m;
f29
                bool rel_dir = cmp(pol[m], pol[1]);
44a
                if ((!last_dir and cur_dir) or
b18
                         (last dir == cur dir and rel dir == cur dir)) {
261
                     1 = m:
8a6
                    last dir = cur dir:
1 f 1
94a
                 } else r = m;
606
            }
792
            return 1;
56c
316
        int max dot(pt v) {
            return extreme([&](pt p, pt q) { return p*v > q*v; });
ec1
3b7
a 54
        pair < int , int > tangents(pt p) {
ffb
            auto L = [\&](pt q, pt r) \{ return ccw(p, r, q); \};
8fd
            auto R = [\&](pt q, pt r) \{ return ccw(p, q, r); \};
            return {extreme(L). extreme(R)}:
fa8
736
3ec };
6e0 bool operator <(const line& a, const line& b) { // comparador pra
   reta
        // assume que as retas tem p < q
        pt v1 = a.q - a.p, v2 = b.q - b.p;
a13
036
        bool b1 = compare_angle(v1, v2), b2 = compare_angle(v2, v1);
        if (b1 or b2) return b1:
73 c
780
        return ccw(a.p, a.q, b.p); // mesmo angulo
b61 }
b14 bool operator ==(const line& a, const line& b) {
```

```
76c
        return !(a < b) and !(b < a);
449 }
// comparador pro set pra fazer sweep line com segmentos
// 6774df
2c4 struct cmp_sweepline {
        bool operator () (const line& a, const line& b) const {
            // assume que os segmentos tem p < q</pre>
191
            if (a.p == b.p) return ccw(a.p, a.q, b.q);
614
            if (a.p.x != a.q.x and (b.p.x == b.q.x or a.p.x < b.p.x))
780
                 return ccw(a.p, a.q, b.p);
dc0
            return ccw(a.p, b.q, b.p);
baf
677 };
// comparador pro set pra fazer sweep angle com segmentos
// 1ee7f5
bef pt dir;
5b0 struct cmp_sweepangle {
        bool operator () (const line& a, const line& b) const {
            return get_t(dir, a) < get_t(dir, b);</pre>
261
dc5
f6d }:
7.6 Geometria 3D
c83 typedef double ld;
e3b const ld DINF = 1e18;
107 const ld eps = 1e-9;
b32 #define sq(x) ((x)*(x))
d97 bool eq(ld a, ld b) {
ba0
            return abs(a - b) <= eps;</pre>
bfc }
b2a struct pt { // ponto
2eb
            ld x, y, z;
            pt(1d x_{-} = 0, 1d y_{-} = 0, 1d z_{-} = 0) : x(x_{-}), y(y_{-}), z(z_{-})
a50
   {}
            bool operator < (const pt p) const {</pre>
5bc
059
                     if (!eq(x, p.x)) return x < p.x;
f98
                     if (!eq(y, p.y)) return y < p.y;
44c
                     if (!eq(z, p.z)) return z < p.z;</pre>
bb3
                     return 0:
6cd
            }
```

bool operator == (const pt p) const {

a83

```
41c
                    return eq(x, p.x) and eq(y, p.y) and eq(z, p.z);
fb5
            pt operator + (const pt p) const { return pt(x+p.x, y+p.y,
44b
   z+p.z); }
            pt operator - (const pt p) const { return pt(x-p.x, y-p.y,
   z-p.z); }
fb7
           pt operator * (const ld c) const { return pt(x*c , y*c ,
   z*c ); }
            pt operator / (const ld c) const { return pt(x/c , y/c ,
7a1
a65
            ld operator * (const pt p) const { return x*p.x + y*p.y +
   z*p.z: }
7f6
            pt operator ^ (const pt p) const { return pt(y*p.z -
   z*p.y, z*p.x - x*p.z, x*p.y - y*p.x); }
5ed
            friend istream& operator >> (istream& in, pt& p) {
9bf
                    return in >> p.x >> p.y >> p.z;
5e8
            }
3ee };
b3a struct line { // reta
730
            pt p, q;
0d6
           line() {}
4b8
            line(pt p_, pt q_) : p(p_), q(q_) {}
8d7
            friend istream& operator >> (istream& in, line& r) {
4cb
                    return in >> r.p >> r.q;
858
            }
7ab }:
79b struct plane { // plano
7e1
            array <pt, 3> p; // pontos que definem o plano
            array <ld, 4> eq; // equacao do plano
29b
bb7
            plane() {}
fb0
            plane(pt p_, pt q_, pt r_) : p({p_, q_, r_}) { build(); }
            friend istream& operator >> (istream& in, plane& P) {
ca9
                    return in >> P.p[0] >> P.p[1] >> P.p[2];
2ab
70e
                    P.build():
544
            }
0a8
            void build() {
da2
                    pt dir = (p[1] - p[0]) ^ (p[2] - p[0]);
7d5
                    eq = \{dir.x, dir.y, dir.z, dir*p[0]*(-1)\};
41a
            }
d5d };
// converte de coordenadas polares para cartesianas
// (angulos devem estar em radianos)
// phi eh o angulo com o eixo z (cima) theta eh o angulo de rotacao ao
```

```
redor de z
2fb pt convert(ld rho, ld th, ld phi) {
           return pt(sin(phi) * cos(th), sin(phi) * sin(th),
   cos(phi)) * rho:
a4f }
// projecao do ponto p na reta r
256 pt proj(pt p, line r) {
bea
            if (r.p == r.q) return r.p;
97a
            r.q = r.q - r.p; p = p - r.p;
9f8
            pt proj = r.q * ((p*r.q) / (r.q*r.q));
2cd
            return proj + r.p;
232 }
// projecao do ponto p no plano P
bla pt proj(pt p, plane P) {
7b6
            p = p - P.p[0], P.p[1] = P.p[1] - P.p[0], P.p[2] = P.p[2]
   - P.p[0];
b69
            pt norm = P.p[1] ^ P.p[2];
            pt proj = p - (norm * (norm * p) / (norm*norm));
6ab
467
            return proj + P.p[0];
4a0 }
// distancia
a45 ld dist(pt a, pt b) {
            return sqrt(sq(a.x-b.x) + sq(a.y-b.y) + sq(a.z-b.z));
2d0 }
// distancia ponto reta
137 ld distline(pt p, line r) {
            return dist(p, proj(p, r));
3c4 }
// distancia de ponto para segmento
d43 ld distseg(pt p, line r) {
            if ((r.q - r.p)*(p - r.p) < 0) return dist(r.p. p);
73d
951
            if ((r.p - r.q)*(p - r.q) < 0) return dist(r.q, p);
200
            return distline(p, r);
42c }
// distancia de ponto a plano com sinal
7cc ld sdist(pt p, plane P) {
150
            return P.eq[0]*p.x + P.eq[1]*p.y + P.eq[2]*p.z + P.eq[3];
d49 }
// distancia de ponto a plano
768 ld distplane(pt p, plane P) {
```

```
сЗе
            return abs(sdist(p, P));
33d }
// se ponto pertence a reta
099 bool isinseg(pt p, line r) {
            return eq(distseg(p, r), 0);
a32
31a }
// se ponto pertence ao triangulo definido por P.p
cd2 bool isinpol(pt p, vector<pt> v) {
fad
            assert(v.size() >= 3);
bf4
            pt norm = (v[1]-v[0]) ^ (v[2]-v[1]);
8a4
            bool inside = true:
cec
           int sign = -1;
           for (int i = 0; i < v.size(); i++) {</pre>
f14
                    line r(v[(i+1)%3], v[i]);
834
                    if (isinseg(p, r)) return true;
2a9
                    pt ar = v[(i+1)\%3] - v[i];
4ef
320
                    if (sign == -1) sign = ((ar^(p-v[i]))*norm > 0);
                    else if (((ar^(p-v[i]))*norm > 0) != sign) inside
   = false;
15e
            }
            return inside;
aca
c81 }
// distancia de ponto ate poligono
361 ld distpol(pt p, vector<pt> v) {
3e7
            pt p2 = proj(p, plane(v[0], v[1], v[2]));
61a
            if (isinpol(p2, v)) return dist(p, p2);
349
            ld ret = DINF:
f14
            for (int i = 0; i < v.size(); i++) {</pre>
                   int j = (i+1)%v.size();
6af
5ee
                    ret = min(ret, distseg(p, line(v[i], v[j])));
7b2
            }
edf
            return ret;
a8d }
// intersecao de plano e segmento
// BOTH = o segmento esta no plano
// ONE = um dos pontos do segmento esta no plano
// PARAL = segmento paralelo ao plano
// CONCOR = segmento concorrente ao plano
e51 enum RETCODE {BOTH, ONE, PARAL, CONCOR};
26b pair < RETCODE, pt > intersect(plane P, line r) {
       1d d1 = sdist(r.p, P);
        1d d2 = sdist(r.q, P);
f8f
```

```
53a
         if (eq(d1, 0) and eq(d2, 0))
504
                     return pair(BOTH, r.p);
72c
         if (eq(d1, 0))
847
                     return pair(ONE, r.p);
         if (eq(d2, 0))
485
                     return pair(ONE, r.q);
168
3fb
         if ((d1 > 0 \text{ and } d2 > 0) \text{ or } (d1 < 0 \text{ and } d2 < 0)) {}
463
             if (eq(d1-d2, 0)) return pair(PARAL, pt());
406
             return pair(CONCOR, pt());
        }
91 c
c84
        1d frac = d1 / (d1 - d2);
3ff
         pt res = r.p + ((r.q - r.p) * frac);
394
         return pair(ONE, res);
b92 }
// rotaciona p ao redor do eixo u por um angulo a
787 pt rotate(pt p, pt u, ld a) {
             u = u / dist(u, pt());
773
             return u * (u * p) + (u ^ p ^ u) * cos(a) + (u ^ p) *
e6f
    sin(a):
7f0 }
7.7 Matriz
945 #define MODULAR false
5ed template < typename T > struct matrix : vector < vector < T >> {
14e
        int n. m:
30f
         void print() {
603
             for (int i = 0; i < n; i++) {</pre>
70f
                 for (int j = 0; j < m; j++) cout << (*this)[i][j] << "
1fb
                 cout << endl;</pre>
d98
             }
        }
101
aa3
         matrix(int n_, int m_, bool ident = false) :
b14
                 vector < vector < T > (n_, vector < T > (m_, 0)), n(n_), m(m_)  {
94e
             if (ident) {
df7
                 assert(n == m);
                 for (int i = 0; i < n; i++) (*this)[i][i] = 1;</pre>
a89
359
             }
527
b83
         matrix(const vector<vector<T>>& c) : vector<vector<T>>(c),
a3d
             n(c.size()), m(c[0].size()) {}
efc
         matrix(const initializer_list<initializer_list<T>>& c) {
f7e
             vector < vector < T >> val;
```

```
212
            for (auto& i : c) val.push_back(i);
            *this = matrix(val);
303
c50
       }
388
        matrix <T> operator*(matrix <T>& r) {
            assert(m == r.n);
1e2
82c
            matrix <T> M(n, r.m);
            for (int i = 0; i < n; i++) for (int k = 0; k < m; k++)
d69
df4
                for (int j = 0; j < r.m; j++) {
                    T \text{ add} = (*this)[i][k] * r[k][j];
e34
f98 #if MODULAR
d41 #warning Usar matrix<ll> e soh colocar valores em [0, MOD) na
   matriz!
                    M[i][j] += add%MOD;
8b6
                    if (M[i][j] >= MOD) M[i][j] -= MOD;
983
8c1 #else
7bb
                    M[i][i] += add;
f2e #endif
                }
620
474
            return M:
394
528
        matrix <T> operator^(ll e){
            matrix <T> M(n, n, true), at = *this;
f10
c87
            while (e) {
2e2
               if (e\&1) M = M*at;
cc2
                e >>= 1:
c80
                at = at*at;
eb6
            }
474
            return M;
ca3
        void apply_transform(matrix M, ll e){
582
1c3
            auto& v = *this;
c87
            while (e) {
9ba
                if (e\&1) v = M*v:
cc2
                e >>= 1:
419
                M = M * M;
d86
            }
4e5
        }
70d };
7.8 Matroid
// Matroids de Grafo e Particao
// De modo geral, toda Matroid contem um build() linear
// e uma funcao constante oracle()
// oracle(i) responde se o conjunto continua independente
// apos adicao do elemento i
```

```
// oracle(i, j) responde se o conjunto continua indepente
// apos trocar o elemento i pelo elemento j
//
// Intersecao sem peso O(r^2 n)
// em que n eh o tamanho do conjunto e r eh o tamanho da resposta
// Matroid Grafica
// Matroid das florestas de um grafo
// Um conjunto de arestas eh independente se formam uma floresta
//
// build() : O(n)
// oracle() : 0(1)
fda struct graphic_matroid {
5da
        int n, m, t;
32c
        vector < array < int , 2>> edges;
789
        vector < vector < int >> g;
62e
        vector<int> comp, in, out;
513
        graphic_matroid(int n_, vector<array<int, 2>> edges_)
            : n(n_), m(edges_.size()), edges(edges_), g(n), comp(n),
    in(n), out(n) {}
315
        void dfs(int u) {
ab8
            in[u] = t++:
17d
            for (auto v : g[u]) if (in[v] == -1)
863
                comp[v] = comp[u], dfs(v);
677
            out[u] = t:
d83
        }
945
        void build(vector<int> I) {
a34
            t = 0:
741
            for (int u = 0; u < n; u++) g[u].clear(), in[u] = -1;
667
            for (int e : I) {
d00
                auto [u, v] = edges[e];
125
                g[u].push_back(v), g[v].push_back(u);
a8a
809
            for (int u = 0; u < n; u++) if (in[u] == -1)
a7d
                comp[u] = u, dfs(u);
207
        }
f31
        bool is_ancestor(int u, int v) {
a68
            return in[u] <= in[v] and in[v] < out[u];</pre>
0c2
        }
        bool oracle(int e) {
e6b
            return comp[edges[e][0]] != comp[edges[e][1]];
453
687
f75
        bool oracle(int e, int f) {
574
            if (oracle(f)) return true;
622
            int u = edges[e][in[edges[e][0]] < in[edges[e][1]]];</pre>
ff2
            return is_ancestor(u, edges[f][0]) != is_ancestor(u,
```

```
edges[f][1]);
8a9
      }
691 };
// Matroid de particao ou cores
// Um conjunto eh independente se a quantidade de elementos
// de cada cor nao excede a capacidade da cor
// Quando todas as capacidades sao 1, um conjunto eh independente
// se todas as suas cores sao distintas
// build() : O(n)
// oracle() : 0(1)
994 struct partition_matroid {
501
        vector < int > cap, color, d;
608
        partition_matroid(vector<int> cap_, vector<int> color_)
04d
            : cap(cap_), color(color_), d(cap.size()) {}
        void build(vector<int> I) {
945
            fill(d.begin(), d.end(), 0);
def
            for (int u : I) d[color[u]]++;
e9d
c58
       }
514
        bool oracle(int u) {
0a1
            return d[color[u]] < cap[color[u]];</pre>
703
       }
f7f
        bool oracle(int u. int v) {
2f7
            return color[u] == color[v] or oracle(v);
        }
4b4
caa }:
// Intersecao de matroid sem pesos
// Dadas duas matroids M1 e M2 definidas sobre o mesmo
// conjunto I, retorna o maior subconjunto de I
// que eh independente tanto para M1 quanto para M2
//
// O(r^2*n)
// Matroid "pesada" deve ser a M2
132 template < typename Matroid1, typename Matroid2 >
801 vector < int > matroid_intersection(int n, Matroid1 M1, Matroid2 M2) {
f5b
        vector < bool > b(n);
a64
        vector < int > I[2]:
a8b
       bool converged = false;
0 c 1
        while (!converged) {
742
            I[0].clear(), I[1].clear();
99d
            for (int u = 0; u < n; u++) I[b[u]].push_back(u);
09d
            M1.build(I[1]), M2.build(I[1]);
```

```
289
            vector < bool > target(n), pushed(n);
26a
            queue < int > q;
5 c 5
            for (int u : I[0]) {
                target[u] = M2.oracle(u);
2b2
                if (M1.oracle(u)) pushed[u] = true, q.push(u);
c1b
0e6
3fe
            vector < int > p(n, -1);
07a
            converged = true;
402
            while (q.size()) {
                int u = q.front(); q.pop();
be1
5c6
                if (target[u]) {
101
                    converged = false;
c32
                    for (int v = u; v != -1; v = p[v]) b[v] = !b[v];
c2b
                    break:
                }
a80
e78
                for (int v : I[!b[u]]) if (!pushed[v]) {
34d
                    if ((b[u] and M1.oracle(u, v)) or (b[v] and
   M2.oracle(v, u)))
bae
                        p[v] = u, pushed[v] = true, q.push(v);
533
                }
            }
1d9
5e7
b68
        return I[1];
381 }
// Intersecao de matroid com pesos
// Dadas duas matroids M1 e M2 e uma funcao de pesos w, todas
   definidas sobre
// um conjunto I retorna o maior subconjunto de I (desempatado pelo
   menor peso)
// que eh independente tanto para M1 quanto para M2
// A resposta eh construida incrementando o tamanho conjunto I de 1 em
// Se nao tiver custo negativo, nao precisa de SPFA
// O(r^3*n) com SPFA
// O(r^2*n*log(n)) com Dijkstra e potencial
42a template < typename T, typename Matroid1, typename Matroid2>
2b5 vector <int > weighted_matroid_intersection(int n, vector <T > w,
   Matroid1 M1, Matroid2 M2) {
        vector < bool > b(n), target(n), is_inside(n);
6c9
        vector < int > I[2], from(n);
563
e35
        vector<pair<T, int>> d(n);
169
        auto check_edge = [&](int u, int v) {
            return (b[u] and M1.oracle(u, v)) or (b[v] and
249
   M2.oracle(v. u)):
```

```
253
        };
                                                                             f97
        while (true) {
                                                                             c7d
                                                                                     }
667
742
            I[0].clear(), I[1].clear();
                                                                             b68
                                                                                     return I[1];
            for (int u = 0; u < n; u++) I[b[u]].push_back(u);
                                                                             8e7 }
99d
            // I[1] contem o conjunto de tamanho I[1].size() de menor
09d
            M1.build(I[1]), M2.build(I[1]):
            for (int u = 0; u < n; u++) {</pre>
687
ea5
                target[u] = false, is_inside[u] = false, from[u] = -1;
                d[u] = {numeric_limits <T>::max(), INF};
961
392
            }
8d3
            deque <T> q;
476
            sort(I[0].begin(), I[0].end(), [&](int i, int j){ return
   w[i] < w[i]; });
            for (int u : I[0]) {
5c5
                target[u] = M2.oracle(u);
2b2
5a7
                if (M1.oracle(u)) {
                    if (is_inside[u]) continue;
4ef
                    d[u] = \{w[u], 0\};
7cc
                     if (!q.empty() and d[u] > d[q.front()])
427
   q.push_back(u);
655
                     else q.push_front(u);
                     is_inside[u] = true;
4ae
764
                }
            }
add
402
            while (q.size()) {
                int u = q.front(); q.pop_front();
97a
6f3
                is inside[u] = false:
                for (int v : I[!b[u]]) if (check_edge(u, v)) {
57a
                     pair <T, int > nd(d[u].first + w[v], d[u].second +
9de
   1);
61b
                    if (nd < d[v]) {</pre>
                         from[v] = u, d[v] = nd;
6ac
                         if (is inside[v]) continue;
bd7
                         if (q.size() and d[v] > d[q.front()])
   q.push_back(v);
                         else q.push_front(v);
275
587
                         is_inside[v] = true;
b3f
                    }
                }
a3b
            }
563
            pair <T, int > mini = pair(numeric_limits <T >:: max(), INF);
cc8
            int targ = -1;
489
            for (int u : I[0]) if (target[u] and d[u] < mini)</pre>
259
2b9
                mini = d[u], targ = u;
            if (targ != -1) for (int u = targ; u != -1; u = from[u])
e14
                b[u] = !b[u], w[u] *= -1:
d89
```

else break;

8 Extra

8.1 fastIO.cpp

```
int read_int() {
    bool minus = false;
    int result = 0;
    char ch;
    ch = getchar();
    while (1) {
        if (ch == '-') break;
        if (ch >= '0' && ch <= '9') break:
        ch = getchar();
   }
    if (ch == '-') minus = true;
    else result = ch-'0';
    while (1) {
        ch = getchar();
        if (ch < '0' || ch > '9') break;
        result = result *10 + (ch - '0');
    if (minus) return -result;
    else return result;
}
8.2 vimrc
set ts=4 sw=4 mouse=a nu ai si undofile
function Hash(1)
    return system("sed '/^\w/d' \| cpp -dD -P -fpreprocessed \| tr -d
       '[:space:]' \| md5sum \| cut -c-6", a:1)
endfunction
function PrintHash() range
    let l = getline(a:firstline, a:lastline)
   let stk = []
    for i in range(len(1))
        call add(stk, i)
        for c in str2list(l[i])
            if c == char2nr('{'})
                call add(stk, i)
            endif
            if c == char2nr(')')
                call remove(stk, -1)
            endif
        endfor
        let ini = stk[-1]
```

```
echo Hash(join(l[ini:i], "\n"))[0:2] 1[i]
        call remove(stk, -1)
    endfor
endfunction
vmap <C-H> :call PrintHash() <CR>
8.3 timer.cpp
// timer T; T() -> retorna o tempo em ms desde que declarou
using namespace chrono;
struct timer : high_resolution_clock {
    const time_point start;
    timer(): start(now()) {}
    int operator()() {
        return duration_cast < milliseconds > (now() - start).count();
};
8.4 rand.cpp
mt19937 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
int uniform(int 1, int r){
    uniform_int_distribution < int > uid(1, r);
    return uid(rng);
}
8.5 template.cpp
#include <bits/stdc++.h>
using namespace std;
#define _ ios_base::sync_with_stdio(0);cin.tie(0);
#define endl '\n'
typedef long long 11;
const int INF = 0x3f3f3f3f;
const 11 LINF = 0x3f3f3f3f3f3f3f3f3f11;
int main() {
    exit(0);
```

debug.cpp

```
void debug_out(string s, int line) { cerr << endl; }</pre>
template < typename H, typename ... T>
void debug_out(string s, int line, H h, T... t) {
    if (s[0] != ',') cerr << "Line(" << line << ") ";</pre>
    do { cerr << s[0]; s = s.substr(1);</pre>
   } while (s.size() and s[0] != ',');
    cerr << " = " << h;
    debug_out(s, line, t...);
}
#ifdef DEBUG
#define debug(...) debug_out(#__VA_ARGS__, __LINE__, __VA_ARGS__)
#else
#define debug(...) 42
#endif
8.7 stress.sh
P=a
```

```
make ${P} ${P}2 gen || exit 1
for ((i = 1; ; i++)) do
    ./gen $i > in
    ./${P} < in > out
    ./${P}2 < in > out2
    if (! cmp -s out out2) then
        echo "--> entrada:"
        cat in
        echo "--> saida1:"
        cat out
        echo "--> saida2:"
        cat out2
        break;
    fi
    echo $i
done
```

makefile

```
CXX = g++
CXXFLAGS = -fsanitize=address, undefined -fno-omit-frame-pointer -g
   -Wall -Wshadow -std=c++17 -Wno-unused-result -Wno-sign-compare
   -Wno-char-subscripts #-fuse-ld=gold
```

8.9 hash.sh

```
# Para usar (hash das linhas [11, 12]):
# bash hash.sh arquivo.cpp 11 12
sed -n 2', 3' p' 1 \mid \text{sed } /^\# \text{w/d'} \mid \text{cpp -dD -P -fpreprocessed} \mid \text{tr}
    -d '[:space:]' | md5sum | cut -c-6
```