# Rábalabaxúrias [UFMG]

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	7.5 Manacher	112	// O(n log(n))
	7.6 eertree	113	int n;
	7.7 String hashing - modulo 2^61 - 1	113	<pre>vector &lt; int &gt; g[MAX]; int subsize[MAX];</pre>
	7.8 Max Suffix	114	<pre>int rem[MAX]; int pai[MAX];</pre>
	7.9 KMP	114	<pre>void dfs(int k, int last) {</pre>
	7.10 Suffix Array - O(n log n)	115	<pre>subsize[k] = 1; for (int i = 0; i &lt; (int) g[k].size(); i++)</pre>
	7.11 Ahocorasick	115	<pre>if (g[k][i] != last and !rem[g[k][i]]) {     dfs(g[k][i], k);</pre>
	7.12 Trie	116	<pre>subsize[k] += subsize[g[k][i]]; }</pre>
			}
3	Extra	118	<pre>int centroid(int k, int last, int size) {</pre>
	8.1 makefile	118	<pre>for (int i = 0; i &lt; (int) g[k].size(); i++) {   int u = g[k][i];</pre>
	8.2 debug.cpp	118	<pre>if (rem[u] or u == last) continue;</pre>

```
if (subsize[u] > size / 2)
            return centroid(u, k, size);
    // k eh o centroid
    return k:
}
//vector<int> dist[MAX];
//void dfs_dist(int k, int last, int d=0) {
      dist[k].push_back(d);
      for (int j : g[k]) if (j != last and !rem[j])
//
          dfs_dist(j, k, d+1);
//}
void decomp(int k, int last = -1) {
    dfs(k, k);
    // acha e tira o centroid
    int c = centroid(k, k, subsize[k]);
    rem[c] = 1;
    pai[c] = last;
    //dfs_dist(c, c);
    // decompoe as sub-arvores
    for (int i = 0; i < (int) g[c].size(); i++)</pre>
        if (!rem[g[c][i]]) decomp(g[c][i], c);
}
void build() {
    memset(rem, 0, sizeof rem);
    decomp(0);
    //for (int i = 0; i < n; i++) reverse(dist[i].begin(),
       dist[i].end());
}
     Tarjan para Pontes
// Computa pontos de articulação
// e pontes
```

dfs\_art(1, -1, d);

```
Prufer code
```

// O(n+m)

}

}

int d = 0;

}

int in[MAX];

int low[MAX];

int parent[MAX];

vector < int > g[MAX];

parent[v] = p;

 $if (p == -1){$ 

int k = 0;

memset(in, -1, sizeof in);

low[v] = in[v] = d++;

is\_art[v] = false; for (int j : g[v]){

void dfs\_art(int v, int p, int &d){

if (j == p) continue;

dfs\_art(j, v, d);

if (low[j] >= in[v]) is\_art[v] = true;

low[v] = min(low[v], low[j]);

else low[v] = min(low[v], in[j]);

k += (parent[j] == v); if (k > 1) is\_art[v] = true;

//if (low[j] > in[v]) this edge is a bridge

 $if (in[j] == -1){$ 

is\_art[v] = false;

for (int j : g[v])

bool is\_art[MAX];

```
// Traduz de lista de arestas para prufer code
// e vice-versa
// Os vertices tem label de 0 a n-1
// Todo array com n-2 posicoes e valores de
// O a n-1 sao prufer codes validos
//
// O(n)
vector < int > to_prufer(vector < pair < int , int >> tree) {
    int n = tree.size()+1;
    vector < int > d(n, 0);
    vector < vector < int >> g(n);
    for (auto [a, b] : tree) d[a]++, d[b]++,
        g[a].push_back(b), g[b].push_back(a);
    vector < int > pai(n, -1);
    queue < int > q; q.push(n-1);
    while (q.size()) {
        int u = q.front(); q.pop();
        for (int v : g[u]) if (v != pai[u])
            pai[v] = u, q.push(v);
    }
    int idx, x;
    idx = x = find(d.begin(), d.end(), 1) - d.begin();
    vector < int > ret;
    for (int i = 0; i < n-2; i++) {</pre>
        int y = pai[x];
        ret.push_back(y);
        if (--d[y] == 1 \text{ and } y < idx) x = y;
        else idx = x = find(d.begin()+idx+1, d.end(), 1) -
            d.begin();
    return ret;
}
vector<pair<int, int>> from_prufer(vector<int> p) {
    int n = p.size()+2;
    vector < int > d(n, 1);
    for (int i : p) d[i]++;
    p.push_back(n-1);
    int idx, x;
    idx = x = find(d.begin(), d.end(), 1) - d.begin();
    vector<pair<int, int>> ret;
```

```
for (int y : p) {
        ret.push_back({x, y});
        if (--d[y] == 1 \text{ and } y < idx) x = y;
        else idx = x = find(d.begin()+idx+1, d.end(), 1) -
           d.begin();
    }
    return ret;
}
1.4 Dinic
// O(min(m * max_flow, n^2 m))
// Grafo com capacidades 1 -> O(sqrt(n)*m)
struct dinic {
    const bool scaling = false; // com scaling -> 0(nm
       log(MAXCAP)),
    int lim;
                                 // com constante alta
    struct edge {
        int to, cap, rev, flow; // para, capacidade, id da
           reversa, fluxo
        bool res; // se a aresta eh residual
        edge(int to_, int cap_, int rev_, bool res_)
            : to(to_), cap(cap_), rev(rev_), flow(0),
               res(res ) {}
    };
    vector < vector < edge >> g;
    vector < int > lev, beg;
    dinic(int n): g(n) {}
    void add(int a, int b, int c) { // de a pra b com cap. c
        g[a].push_back(edge(b, c, g[b].size(), false));
        g[b].push_back(edge(a, 0, g[a].size()-1, true));
    }
    bool bfs(int s, int t) {
        lev = vector<int>(g.size(), -1); lev[s] = 0;
        beg = vector<int>(g.size(), 0);
        queue < int > q; q.push(s);
        while (q.size()) {
```

```
int u = q.front(); q.pop();
        for (auto& i : g[u]) {
            if (lev[i.to] != -1 or (i.flow == i.cap))
               continue:
            if (scaling and i.cap - i.flow < lim)</pre>
               continue;
            lev[i.to] = lev[u] + 1;
            q.push(i.to);
        }
    }
    return lev[t] != -1;
int dfs(int v, int s, int f = INF){
    if (!f or v == s) return f;
    for (int& i = beg[v]; i < g[v].size(); i++) {</pre>
        auto& e = g[v][i];
        if (lev[e.to] != lev[v] + 1) continue;
       int foi = dfs(e.to, s, min(f, e.cap - e.flow));
        if (!foi) continue;
        e.flow += foi, g[e.to][e.rev].flow -= foi;
        return foi;
    }
    return 0;
11 max_flow(int s, int t) {
   11 f = 0;
    for (lim = scaling ? (1 << 30) : 1; lim; lim /= 2)
        while (bfs(s, t)) while (int ff = dfs(s, t)) f
    return f;
vector<pair<int, int> > get_cut(int s, int t) {
    max_flow(s, t);
    vector<pair<int, int> > cut;
    vector<int> vis(g.size(), 0), st = {s};
    vis[s] = 1:
    while (st.size()) {
        int u = st.back(); st.pop_back();
        for (auto e : g[u]) if (!vis[e.to] and e.flow <</pre>
           e.cap)
            vis[e.to] = 1, st.push_back(e.to);
    }
```

#### 1.5 Dominator Tree - Kawakami

```
// Se vira pra usar ai
//
// build - O(n)
// dominates - O(1)
int n;
namespace DTree {
    vector < int > g[MAX];
    // The dominator tree
    vector<int> tree[MAX];
    int dfs_l[MAX], dfs_r[MAX];
    // Auxiliary data
    vector < int > rg[MAX], bucket[MAX];
    int idom[MAX], sdom[MAX], prv[MAX], pre[MAX];
    int ancestor[MAX], label[MAX];
    vector<int> preorder;
    void dfs(int v) {
        static int t = 0;
        pre[v] = ++t;
        sdom[v] = label[v] = v;
        preorder.push_back(v);
        for (int nxt: g[v]) {
            if (sdom[nxt] == -1) {
                prv[nxt] = v;
                dfs(nxt);
```

```
rg[nxt].push_back(v);
    }
}
int eval(int v) {
    if (ancestor[v] == -1) return v;
    if (ancestor[ancestor[v]] == -1) return label[v];
    int u = eval(ancestor[v]);
    if (pre[sdom[u]] < pre[sdom[label[v]]]) label[v] = u;</pre>
    ancestor[v] = ancestor[u];
    return label[v];
}
void dfs2(int v) {
    static int t = 0;
    dfs_1[v] = t++;
    for (int nxt: tree[v]) dfs2(nxt);
    dfs_r[v] = t++;
}
void build(int s) {
    for (int i = 0; i < n; i++) {</pre>
        sdom[i] = pre[i] = ancestor[i] = -1;
        rg[i].clear();
        tree[i].clear();
        bucket[i].clear();
    preorder.clear();
    dfs(s);
    if (preorder.size() == 1) return;
    for (int i = int(preorder.size()) - 1; i >= 1; i--) {
        int w = preorder[i];
        for (int v: rg[w]) {
            int u = eval(v);
            if (pre[sdom[u]] < pre[sdom[w]]) sdom[w] =</pre>
               sdom[u];
        }
        bucket[sdom[w]].push_back(w);
        ancestor[w] = prv[w];
        for (int v: bucket[prv[w]]) {
            int u = eval(v);
            idom[v] = (u == v) ? sdom[v] : u;
        bucket[prv[w]].clear();
    }
```

```
for (int i = 1; i < preorder.size(); i++) {
    int w = preorder[i];
    if (idom[w] != sdom[w]) idom[w] = idom[idom[w]];
        tree[idom[w]].push_back(w);
    }
    idom[s] = sdom[s] = -1;
    dfs2(s);
}

// Whether every path from s to v passes through u
bool dominates(int u, int v) {
    if (pre[v] == -1) return 1; // vacuously true
    return dfs_l[u] <= dfs_l[v] && dfs_r[v] <= dfs_r[u];
}
};</pre>
```

# 1.6 Kosaraju

```
// O(n + m)
int n;
vector < int > g[MAX];
vector<int> gi[MAX]; // grafo invertido
int vis[MAX];
stack<int> S;
int comp[MAX]; // componente conexo de cada vertice
void dfs(int k) {
    vis[k] = 1:
    for (int i = 0; i < (int) g[k].size(); i++)</pre>
        if (!vis[g[k][i]]) dfs(g[k][i]);
    S.push(k);
}
void scc(int k, int c) {
    vis[k] = 1;
    comp[k] = c;
    for (int i = 0; i < (int) gi[k].size(); i++)</pre>
        if (!vis[gi[k][i]]) scc(gi[k][i], c);
```

```
void kosaraju() {
    for (int i = 0; i < n; i++) vis[i] = 0;
    for (int i = 0; i < n; i++) if (!vis[i]) dfs(i);

for (int i = 0; i < n; i++) vis[i] = 0;
    while (S.size()) {
        int u = S.top();
        S.pop();
        if (!vis[u]) scc(u, u);
    }
}</pre>
```

# 1.7 Dijkstra

```
// encontra menor distancia de x
// para todos os vertices
// se ao final do algoritmo d[i] = INF,
// entao x nao alcanca i
//
// O(m log(n))
int d[MAX]:
vector<pair<int,int>> g[MAX]; // {vizinho, peso}
int n;
void dijkstra(int x) {
    for(int i=0; i < n; i++) d[i] = INF;</pre>
    d[x] = 0;
    priority_queue < pair < int , int >> pq;
    pq.push({0,x});
    while(pq.size()) {
        auto [ndist,u] = pq.top(); pq.pop();
        if(-ndist > d[u]) continue;
        for (auto [idx,w]: g[u]) if (d[idx] > d[u] + w) {
            d[idx] = d[u] + w;
```

```
pq.push({-d[idx], idx});
}
}
```

#### 1.8 Heavy-Light Decomposition sem Update

```
// query de min do caminho
//
// Complexidades:
// build - O(n)
// query_path - O(log(n))
#define f first
#define s second
namespace hld {
    vector<pair<int, int> > g[MAX];
    int pos[MAX], sz[MAX];
    int sobe[MAX], pai[MAX];
    int h[MAX], v[MAX], t;
    int men[MAX], seg[2*MAX];
    void build_hld(int k, int p = -1, int f = 1) {
        v[pos[k] = t++] = sobe[k]; sz[k] = 1;
        for (auto& i : g[k]) if (i.f != p) {
            sobe[i.f] = i.s; pai[i.f] = k;
            h[i.f] = (i == g[k][0] ? h[k] : i.f);
            men[i.f] = (i == g[k][0] ? min(men[k], i.s) :
               i.s):
            build_hld(i.f, k, f); sz[k] += sz[i.f];
            if (sz[i.f] > sz[g[k][0].f] or g[k][0].f == p)
               swap(i, g[k][0]);
        if (p*f == -1) build_hld(h[k] = k, -1, t = 0);
    }
    void build(int root = 0) {
        t = 0;
        build_hld(root);
```

# 1.9 LCA com binary lifting

```
// Assume que um vertice eh ancestral dele mesmo, ou seja,
// se a eh ancestral de b, lca(a, b) = a
// MAX2 = ceil(log(MAX))
// Complexidades:
// build - O(n log(n))
// lca - O(log(n))
vector < vector < int > > g(MAX);
int n, p;
int pai[MAX2][MAX];
int in[MAX], out[MAX];
void dfs(int k) {
    in[k] = p++;
    for (int i = 0; i < (int) g[k].size(); i++)</pre>
        if (in[g[k][i]] == -1) {
            pai[0][g[k][i]] = k;
            dfs(g[k][i]);
        }
    out[k] = p++;
```

```
}
void build(int raiz) {
    for (int i = 0; i < n; i++) pai[0][i] = i;</pre>
    p = 0, memset(in, -1, sizeof in);
    dfs(raiz);
    // pd dos pais
    for (int k = 1; k < MAX2; k++) for (int i = 0; i < n;
        pai[k][i] = pai[k - 1][pai[k - 1][i]];
}
bool anc(int a, int b) { // se a eh ancestral de b
    return in[a] <= in[b] and out[a] >= out[b];
}
int lca(int a, int b) {
    if (anc(a, b)) return a;
    if (anc(b, a)) return b;
    // sobe a
    for (int k = MAX2 - 1; k >= 0; k--)
        if (!anc(pai[k][a], b)) a = pai[k][a];
    return pai[0][a];
}
// Alternativamente:
// 'binary lifting' gastando O(n) de memoria
// Da pra add folhas e fazer queries online
// 3 vezes o tempo do binary lifting normal
//
// build - O(n)
// kth, lca, dist - O(log(n))
int d[MAX], p[MAX], pp[MAX];
void set_root(int i) { p[i] = pp[i] = i, d[i] = 0; }
void add_leaf(int i, int u) {
    p[i] = u, d[i] = d[u]+1;
```

```
pp[i] = 2*d[pp[u]] == d[pp[pp[u]]]+d[u] ? pp[pp[u]] : u;
}
int kth(int i, int k) {
    int dd = max(0, d[i]-k);
    while (d[i] > dd) i = d[pp[i]] >= dd ? pp[i] : p[i];
    return i;
}
int lca(int a, int b) {
    if (d[a] < d[b]) swap(a, b);</pre>
    while (d[a] > d[b]) a = d[pp[a]] >= d[b] ? pp[a] : p[a];
    while (a != b) {
        if (pp[a] != pp[b]) a = pp[a], b = pp[b];
        else a = p[a], b = p[b];
    }
    return a;
}
int dist(int a, int b) { return d[a]+d[b]-2*d[lca(a,b)]; }
vector < int > g[MAX];
void build(int i, int pai=-1) {
    if (pai == -1) set_root(i);
    for (int j : g[i]) if (j != pai) {
        add_leaf(j, i);
        build(j, i);
}
1.10 Heavy-Light Decomposition - vertice
// SegTree de soma
// query / update de soma dos vertices
```

```
// SegTree de soma
// query / update de soma dos vertices
//
// Complexidades:
// build - O(n)
// query_path - O(log^2 (n))
// update_path - O(log^2 (n))
```

```
// query_subtree - O(log(n))
// update_subtree - O(log(n))
namespace seg { ... }
namespace hld {
    vector < int > g[MAX];
    int pos[MAX], sz[MAX];
    int peso[MAX], pai[MAX];
    int h[MAX], v[MAX], t;
    void build_hld(int k, int p = -1, int f = 1) {
        v[pos[k] = t++] = peso[k]; sz[k] = 1;
        for (auto& i : g[k]) if (i != p) {
            pai[i] = k;
            h[i] = (i == g[k][0] ? h[k] : i);
            build_hld(i, k, f); sz[k] += sz[i];
            if (sz[i] > sz[g[k][0]] or g[k][0] == p) swap(i,
               g[k][0]);
        }
        if (p*f == -1) build_hld(h[k] = k, -1, t = 0);
    }
    void build(int root = 0) {
        t = 0;
        build_hld(root);
        seg::build(t, v);
    }
    ll query_path(int a, int b) {
        if (pos[a] < pos[b]) swap(a, b);</pre>
        if (h[a] == h[b]) return seg::query(pos[b], pos[a]);
        return seg::query(pos[h[a]], pos[a]) +
           query_path(pai[h[a]], b);
    }
    void update_path(int a, int b, int x) {
        if (pos[a] < pos[b]) swap(a, b);
        if (h[a] == h[b]) return (void)seg::update(pos[b],
           pos[a], x);
        seg::update(pos[h[a]], pos[a], x);
           update_path(pai[h[a]], b, x);
```

```
}
ll query_subtree(int a) {
    return seg::query(pos[a], pos[a]+sz[a]-1);
}
void update_subtree(int a, int x) {
    seg::update(pos[a], pos[a]+sz[a]-1, x);
}
int lca(int a, int b) {
    if (pos[a] < pos[b]) swap(a, b);
    return h[a] == h[b] ? b : lca(pai[h[a]], b);
}</pre>
```

#### 1.11 LCA com HLD

```
// Assume que um vertice eh ancestral dele mesmo, ou seja,
// se a eh ancestral de b, lca(a, b) = a
// Para buildar pasta chamar build(root)
// anc(a, b) responde se 'a' eh ancestral de 'b'
//
// Complexidades:
// build - O(n)
// lca - O(log(n))
// anc - 0(1)
vector < int > g[MAX];
int pos[MAX], h[MAX], sz[MAX];
int pai[MAX], t;
void build(int k, int p = -1, int f = 1) {
    pos[k] = t++; sz[k] = 1;
    for (int& i : g[k]) if (i != p) {
        pai[i] = k;
        h[i] = (i == g[k][0] ? h[k] : i);
        build(i, k, f); sz[k] += sz[i];
        if (sz[i] > sz[g[k][0]] or g[k][0] == p) swap(i,
           g[k][0]);
    if (p*f == -1) t = 0, h[k] = k, build(k, -1, 0);
```

```
}
int lca(int a, int b) {
    if (pos[a] < pos[b]) swap(a, b);</pre>
    return h[a] == h[b] ? b : lca(pai[h[a]], b);
}
bool anc(int a, int b) {
    return pos[a] <= pos[b] and pos[b] <= pos[a]+sz[a]-1;</pre>
}
1.12 LCA com RMQ
// Assume que um vertice eh ancestral dele mesmo, ou seja,
// se a eh ancestral de b, lca(a, b) = a
// dist(a, b) retorna a distancia entre a e b
//
// Complexidades:
// build - O(n)
// lca - 0(1)
// dist - O(1)
template < typename T> struct rmq {
    vector <T> v;
    int n; static const int b = 30;
    vector < int > mask, t;
    int op(int x, int y) { return v[x] < v[y] ? x : y; }
    int msb(int x) { return
        __builtin_clz(1)-__builtin_clz(x); }
    rmq() {}
    rmq(const \ vector < T > \& v_) : v(v_), n(v.size()), mask(n),
        for (int i = 0, at = 0; i < n; mask[i++] = at |= 1) {
            at = (at << 1) &((1 << b) -1);
            while (at and op(i, i-msb(at&-at)) == i) at ^=
                at&-at;
        for (int i = 0; i < n/b; i++) t[i] =</pre>
           b*i+b-1-msb(mask[b*i+b-1]);
```

```
for (int j = 1; (1<<j) <= n/b; j++) for (int i = 0;
           i+(1<< j) <= n/b; i++)
            t[n/b*j+i] = op(t[n/b*(j-1)+i],
               t[n/b*(j-1)+i+(1<<(j-1))]);
    }
    int small(int r, int sz = b) { return
       r-msb(mask[r]&((1<<sz)-1));}
    T query(int 1, int r) {
        if (r-l+1 <= b) return small(r, r-l+1);</pre>
        int ans = op(small(1+b-1), small(r));
        int x = 1/b+1, y = r/b-1;
        if (x <= y) {
            int j = msb(y-x+1);
            ans = op(ans, op(t[n/b*j+x],
               t[n/b*j+y-(1<<j)+1]));
        }
        return ans;
    }
};
namespace lca {
    vector < int > g[MAX];
    int v[2*MAX], pos[MAX], dep[2*MAX];
    int t;
    rmq<int> RMQ;
    void dfs(int i, int d = 0, int p = -1) {
        v[t] = i, pos[i] = t, dep[t++] = d;
        for (int j : g[i]) if (j != p) {
            dfs(j, d+1, i);
            v[t] = i, dep[t++] = d;
        }
    }
    void build(int n, int root) {
        t = 0;
        dfs(root);
        RMQ = rmq < int > (vector < int > (dep, dep + 2*n - 1));
    int lca(int a, int b) {
        a = pos[a], b = pos[b];
        return v[RMQ.query(min(a, b), max(a, b))];
    }
```

## 1.13 Heavy-Light Decomposition - aresta

```
// SegTree de soma
// query / update de soma das arestas
// Complexidades:
// build - O(n)
// \text{ query_path - } O(\log^2 (n))
// update_path - O(log^2 (n))
// query_subtree - O(log(n))
// update_subtree - O(log(n))
#define f first
#define s second
namespace seg { ... }
namespace hld {
    vector<pair<int, int> > g[MAX];
    int pos[MAX], sz[MAX];
    int sobe[MAX], pai[MAX];
    int h[MAX], v[MAX], t;
    void build_hld(int k, int p = -1, int f = 1) {
        v[pos[k] = t++] = sobe[k]; sz[k] = 1;
        for (auto& i : g[k]) if (i.f != p) {
            sobe[i.f] = i.s; pai[i.f] = k;
            h[i.f] = (i == g[k][0] ? h[k] : i.f);
            build_hld(i.f, k, f); sz[k] += sz[i.f];
            if (sz[i.f] > sz[g[k][0].f] or g[k][0].f == p)
               swap(i, g[k][0]);
        if (p*f == -1) build_hld(h[k] = k, -1, t = 0);
```

```
}
void build(int root = 0) {
    t = 0;
    build_hld(root);
    seg::build(t, v);
ll query_path(int a, int b) {
    if (a == b) return 0;
    if (pos[a] < pos[b]) swap(a, b);</pre>
    if (h[a] == h[b]) return seg::query(pos[b]+1,
       pos[a]):
    return seg::query(pos[h[a]], pos[a]) +
       query_path(pai[h[a]], b);
void update_path(int a, int b, int x) {
    if (a == b) return:
    if (pos[a] < pos[b]) swap(a, b);
    if (h[a] == h[b]) return (void)seg::update(pos[b]+1,
       pos[a], x);
    seg::update(pos[h[a]], pos[a], x);
       update_path(pai[h[a]], b, x);
11 query_subtree(int a) {
    if (sz[a] == 1) return 0;
    return seg::query(pos[a]+1, pos[a]+sz[a]-1);
void update_subtree(int a, int x) {
    if (sz[a] == 1) return;
    seg::update(pos[a]+1, pos[a]+sz[a]-1, x);
int lca(int a, int b) {
    if (pos[a] < pos[b]) swap(a, b);
    return h[a] == h[b] ? b : lca(pai[h[a]], b);
}
```

#### 1.14 Bellman-Ford

}

```
// Calcula a menor distancia
// entre a e todos os vertices e
// detecta ciclo negativo
// Retorna 1 se ha ciclo negativo
// Nao precisa representar o grafo,
// soh armazenar as arestas
//
// O(nm)
int n, m;
int d[MAX];
vector<pair<int, int> > ar; // vetor de arestas
vector < int > w;
                             // peso das arestas
bool bellman_ford(int a) {
    for (int i = 0; i < n; i++) d[i] = INF;</pre>
    d[a] = 0:
    for (int i = 0; i <= n; i++)
        for (int j = 0; j < m; j++) {</pre>
            if (d[ar[j].second] > d[ar[j].first] + w[j]) {
                if (i == n) return 1;
                d[ar[j].second] = d[ar[j].first] + w[j];
            }
        }
    return 0;
}
```

# | 1.15 | Tarjan para SCC

```
// O(n + m)
int n;
vector<int> g[MAX];
stack<int> s;
int vis[MAX], comp[MAX];
int id[MAX], p;
```

```
// se quiser comprimir ciclo em grafo nao direcionado,
// colocar um if na dfs para nao voltar pro vertice que veio
int dfs(int k) {
    int lo = id[k] = p++;
    s.push(k);
    vis[k] = 2; // ta na pilha
    // calcula o menor cara q ele alcanca
    // que ainda nao esta em um scc
    for (int i = 0; i < g[k].size(); i++) {</pre>
        if (!vis[g[k][i]])
            lo = min(lo, dfs(g[k][i]));
        else if (vis[g[k][i]] == 2)
            lo = min(lo, id[g[k][i]]);
    }
    // nao alcanca ninguem menor -> comeca scc
    if (lo == id[k]) while (1) {
        int u = s.top();
        s.pop(); vis[u] = 1;
        comp[u] = k;
        if (u == k) break;
    return lo;
}
void tarjan() {
    memset(vis, 0, sizeof(vis));
    p = 0;
    for (int i = 0; i < n; i++) if (!vis[i]) dfs(i);</pre>
}
1.16 Centro da Arvore
// Centro eh o vertice que minimiza
// a maior distancia dele pra alguem
// O centro fica no meio do diametro
// A funcao center retorna um par com
```

```
// o diametro e o centro
// O(n+m)
vector < int > g[MAX];
int n, vis[MAX];
int d[2][MAX];
// retorna ultimo vertice visitado
int bfs(int k, int x) {
        queue < int > q; q.push(k);
    memset(vis, 0, sizeof(vis));
    vis[k] = 1;
    d[x][k] = 0;
    int last = k;
    while (q.size()) {
        int u = q.front(); q.pop();
        last = u;
        for (int i : g[u]) if (!vis[i]) {
            vis[i] = 1;
            q.push(i);
            d[x][i] = d[x][u] + 1;
        }
    }
    return last;
pair<int, int> center() {
    int a = bfs(0, 0);
    int b = bfs(a, 1);
    bfs(b, 0);
    int c, mi = INF;
   for (int i = 0; i < n; i++) if (max(d[0][i], d[1][i]) <</pre>
        mi = max(d[0][i], d[1][i]), c = i;
   return {d[0][a], c};
```

# 1.17 Sack (DSU em arvores)

```
// Responde queries de todas as sub-arvores
// offline
// O(n log(n))
int sz[MAX], cor[MAX], cnt[MAX];
vector < int > g[MAX];
void build(int k, int d=0) {
    sz[k] = 1;
    for (auto& i : g[k]) {
        build(i, d+1); sz[k] += sz[i];
        if (sz[i] > sz[g[k][0]]) swap(i, g[k][0]);
}
void compute(int k, int x, bool dont=1) {
    cnt[cor[k]] += x;
    for (int i = dont; i < g[k].size(); i++)</pre>
        compute(g[k][i], x, 0);
}
void solve(int k, bool keep=0) {
    for (int i = int(g[k].size())-1; i >= 0; i--)
        solve(g[k][i], !i);
    compute(k, 1);
        // agora cnt[i] tem quantas vezes a cor
        // i aparece na sub-arvore do k
    if (!keep) compute(k, -1, 0);
}
```

## 1.18 AGM Direcionada

```
// Fala o menor custo para selecionar arestas tal que
// o vertice 'r' alcance todos
// Se nao tem como, retorna LINF
```

```
// O(m log(n))
struct node {
    pair<ll, int> val;
    ll lazy;
    node *1, *r;
    node() {}
    node(ii v) : val(v), lazy(0), l(NULL), r(NULL) {}
    void prop() {
        val.f += lazy;
        if (1) 1->lazy += lazy;
        if (r) r->lazy += lazy;
        lazy = 0;
   }
};
void merge(node*& a, node* b) {
    if (!a) swap(a, b);
    if (!b) return;
   a->prop(), b->prop();
    if (a->val > b->val) swap(a, b);
    merge(rand()%2 ? a->1 : a->r, b);
pair<11, int> pop(node*& R) {
    R->prop();
    auto ret = R->val;
    node* tmp = R;
    merge(R->1, R->r);
    R = R -> 1;
    if (R) R->lazy -= ret.f;
    delete tmp;
    return ret;
void apaga(node* R) { if (R) apaga(R->1), apaga(R->r),
   delete R; }
ll dmst(int n, int r, vector<pair<ii, int>>& ar) {
    vector < int > p(n); iota(p.begin(), p.end(), 0);
    function < int(int) > find = [&](int k) { return
       p[k] == k?k:p[k] = find(p[k]); };
    vector < node *> h(n);
```

```
for (auto e : ar) merge(h[e.f.s], new node({e.s,
       e.f.f}));
    vector < int > pai(n, -1), path(n);
    pai[r] = r;
    11 \text{ ans} = 0;
    for (int i = 0; i < n; i++) { // vai conectando todo
       mundo
        int u = i, at = 0;
        while (pai[u] == -1) {
            if (!h[u]) { // nao tem
                for (auto i : h) apaga(i);
                return LINF;
            path[at++] = u, pai[u] = i;
            auto [mi, v] = pop(h[u]);
            ans += mi;
            if (pai[u = find(v)] == i) { // ciclo}
                while (find(v = path[--at]) != u)
                     merge(h[u], h[v]), h[v] = NULL,
                        p[find(v)] = u;
                pai[u] = -1;
            }
        }
    for (auto i : h) apaga(i);
    return ans;
}
1.19 Centroid
// Computa os 2 centroids da arvore
//
// O(n)
int n, subsize[MAX];
vector < int > g[MAX];
```

void dfs(int k, int p=-1) {

```
subsize[k] = 1:
   for (int i : g[k]) if (i != p) {
        dfs(i, k);
        subsize[k] += subsize[i];
    }
}
int centroid(int k, int p=-1, int size=-1) {
    if (size == -1) size = subsize[k];
   for (int i : g[k]) if (i != p) if (subsize[i] > size/2)
        return centroid(i, k, size, t);
    return k;
}
pair<int, int> centroids(int k=0) {
    dfs(k);
   int i = centroid(k), i2 = i;
    for (int j : g[i]) if (2*subsize[j] == subsize[k]) i2 =
       j;
    return {i, i2};
}
1.20 Topological Sort
// Retorna uma ordenacaoo topologica de g
// Se g nao for DAG retorna um vetor vazio
//
// O(n + m)
vector < int > g[MAX];
vector < int > topo_sort(int n) {
    vector < int > ret(n,-1), vis(n,0);
    int pos = n-1, dag = 1;
    function < void(int) > dfs = [&] (int v) {
        vis[v] = 1;
        for(auto u : g[v]) {
            if(vis[u] == 1) dag = 0;
            else if(!vis[u]) dfs(u);
```

```
}
    ret[pos--] = v, vis[v] = 2;
};

for(int i=0; i<n; i++) if(!vis[i]) dfs(i);
    if(!dag) ret.clear();
    return ret;
}</pre>
```

#### 1.21 Kruskal

```
// Gera e retorna uma AGM e seu custo total a partir do
   vetor de arestas (edg)
// do grafo
//
// O(m log(m) + m a(m))
typedef tuple <int, int, int> t3;
vector < t3 > edg; // {peso,[x,y]}
// DSU em O(a(n))
void dsu_build();
int find(int a);
void unite(int a, int b);
pair<11, vector <t3>> kruskal(int n) {
    dsu_build(n);
    sort(edg.begin(), edg.end());
    11 cost = 0;
    vector < t3 > mst;
    for(auto [w,x,y] : edg) if(find(x) != find(y)) {
        mst.push_back({w,x,y});
        cost += w;
        unite(x,y);
    return {cost,mst};
}
```

# 1.22 Blossom - matching maximo em grafo geral

```
// O(n^3)
// Se for bipartido, nao precisa da funcao
// 'contract', e roda em O(nm)
vector < int > g[MAX];
int match[MAX]; // match[i] = com quem i esta matchzado ou -1
int n, pai[MAX], base[MAX], vis[MAX];
queue < int > q;
void contract(int u, int v, bool first = 1) {
    static vector < bool > bloss;
    static int 1;
    if (first) {
        bloss = vector < bool > (n, 0);
        vector < bool > teve(n, 0);
        int k = u; l = v;
        while (1) {
            teve[k = base[k]] = 1;
            if (match[k] == -1) break;
            k = pai[match[k]];
        while (!teve[l = base[l]]) l = pai[match[l]];
    }
    while (base[u] != 1) {
        bloss[base[u]] = bloss[base[match[u]]] = 1;
        pai[u] = v;
        v = match[u];
        u = pai[match[u]];
    }
    if (!first) return;
    contract(v, u, 0);
   for (int i = 0; i < n; i++) if (bloss[base[i]]) {</pre>
        base[i] = 1;
        if (!vis[i]) q.push(i);
        vis[i] = 1;
    }
}
int getpath(int s) {
    for (int i = 0; i < n; i++) base[i] = i, pai[i] = -1,
```

```
vis[i] = 0:
    vis[s] = 1; q = queue < int > (); q.push(s);
    while (q.size()) {
        int u = q.front(); q.pop();
        for (int i : g[u]) {
            if (base[i] == base[u] or match[u] == i)
               continue:
            if (i == s or (match[i] != -1 and pai[match[i]]
                ! = -1))
                 contract(u, i);
            else if (pai[i] == -1) {
                pai[i] = u;
                if (match[i] == -1) return i;
                i = match[i];
                vis[i] = 1; q.push(i);
            }
        }
    return -1;
}
int blossom() {
    int ans = 0;
    memset(match, -1, sizeof(match));
    for (int i = 0; i < n; i++) if (match[i] == -1)</pre>
        for (int j : g[i]) if (match[j] == -1) {
            match[i] = j;
            match[j] = i;
            ans++;
            break;
        }
    for (int i = 0; i < n; i++) if (match[i] == -1) {</pre>
        int j = getpath(i);
        if (j == -1) continue;
        ans++:
        while (j != -1) {
            int p = pai[j], pp = match[p];
            match[p] = j;
            match[j] = p;
            j = pp;
        }
    }
```

```
return ans;
}
```

#### 1.23 Max flow com lower bound

```
// Manda passar pelo menos 'lb' de fluxo
// em cada aresta
// O(dinic)
struct lb_max_flow : dinic {
    vector < int > d;
    vector < int > e;
    lb_max_flow(int n):dinic(n + 2), d(n, 0){}
    void add(int a, int b, int c, int lb = 0){
        c -= lb:
        d[a] -= lb;
        d[b] += lb;
        dinic::add(a, b, c);
    }
    bool check_flow(int src, int snk, int F){
        int n = d.size();
        d[src] += F;
        d[snk] -= F;
        for (int i = 0; i < n; i++){</pre>
             if (d[i] > 0){
                 dinic::add(n, i, d[i]);
             } else if (d[i] < 0){</pre>
                 dinic::add(i, n+1, -d[i]);
        }
        int f = \max_{f \in \mathcal{F}} \{n, n+1\};
        return (f == F);
    }
};
```

#### 1.24 Isomorfismo de Arvores

```
// Duas arvores T1 e T2 sao isomorfas
// sse T1.getHash() = T2.getHash()
// O(n log(n))
map < vector < int > , int > mapp;
struct tree {
    int n:
    vector < vector < int > > g;
    vector < int > subsize;
    tree(int n) {
        g.resize(n);
        subsize.resize(n);
    void dfs(int k, int p=-1) {
        subsize[k] = 1;
        for (int i : g[k]) if (i != p) {
            dfs(i, k);
            subsize[k] += subsize[i];
        }
    }
    int centroid(int k, int p=-1, int size=-1) {
        if (size == -1) size = subsize[k];
        for (int i : g[k]) if (i != p)
            if (subsize[i] > size/2)
                return centroid(i, k, size);
        return k:
    pair < int , int > centroids(int k=0) {
        dfs(k):
        int i = centroid(k), i2 = i;
        for (int j : g[i]) if (2*subsize[j] == subsize[k])
           i2 = j;
        return {i, i2};
    int hashh(int k, int p=-1) {
        vector < int > v;
        for (int i : g[k]) if (i != p) v.push_back(hashh(i,
```

```
k));
sort(v.begin(), v.end());
if (!mapp.count(v)) mapp[v] = int(mapp.size());
return mapp[v];
}
ll getHash(int k=0) {
  pair<int, int> c = centroids(k);
  ll a = hashh(c.first), b = hashh(c.second);
  if (a > b) swap(a, b);
  return (a<<30)+b;
};</pre>
```

# 1.25 Algoritmo de Kuhn

```
// Computa matching maximo em grafo bipartido
// 'n' e 'm' sao quantos vertices tem em cada particao
// chamar add(i, j) para add aresta entre o cara i
// da particao A, e o cara j da particao B
// (entao i < n, j < m)
// Para recuperar o matching, basta olhar 'ma' e 'mb'
// recover() recupera o min vertex cover como um par de
// {caras da particao A, caras da particao B}
//
// O(|V| * |E|)
// Na pratica, parece rodar tao rapido quanto o Dinic
mt19937 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
struct kuhn {
    int n, m;
    vector < vector < int >> g;
    vector < int > vis, ma, mb;
    kuhn(int n_, int m_) : n(n_), m(m_), g(n),
        vis(n+m), ma(n, -1), mb(m, -1) {}
    void add(int a, int b) { g[a].pb(b); }
```

```
bool dfs(int i) {
    vis[i] = 1;
    for (int j : g[i]) if (!vis[n+j]) {
        vis[n+j] = 1;
        if (mb[j] == -1 or dfs(mb[j])) {
            ma[i] = j, mb[j] = i;
            return true;
        }
    }
    return false;
int matching() {
    int ret = 0, aum = 1;
    for (auto& i : g) shuffle(i.begin(), i.end(), rng);
    while (aum) {
        for (int j = 0; j < m; j++) vis[n+j] = 0;
        aum = 0:
        for (int i = 0; i < n; i++)</pre>
            if (ma[i] == -1 and dfs(i)) ret++, aum = 1;
    }
    return ret;
pair < vector < int >, vector < int >> recover() {
    int M = matching();
    for (int i = 0; i < n+m; i++) vis[i] = 0;</pre>
    for (int i = 0; i < n; i++) if (ma[i] == -1)
       assert(!dfs(i));
    vector<int> ca, cb;
    for (int i = 0; i < n; i++) if (!vis[i])</pre>
       ca.push_back(i);
    for (int i = 0; i < m; i++) if (vis[n+i])</pre>
       cb.push_back(i);
    //assert(ca.size() + cb.size() == M);
    return {ca. cb}:
}
  Virtual Tree
```

#### 1.26

};

// Comprime uma arvore dado um conjunto S de vertices, de

```
forma que
// o conjunto de vertices da arvore comprimida contenha S e
// minimal e fechado sobre a operacao de LCA
// Se |S| = k, a arvore comprimida tem O(k) vertices
//
// O(k log(k))
template < typename T > struct rmq {
    vector <T> v;
    int n; static const int b = 30;
    vector < int > mask, t;
    int op(int x, int y) { return v[x] < v[y] ? x : y; }
    int msb(int x) { return
       __builtin_clz(1)-__builtin_clz(x); }
    rmq() {}
    rmq(const vector < T > \& v_) : v(v_), n(v.size()), mask(n),
       t(n) {
        for (int i = 0, at = 0; i < n; mask[i++] = at |= 1) {
            at = (at << 1) & ((1 << b) -1);
            while (at and op(i, i-msb(at&-at)) == i) at ^=
               at&-at;
        for (int i = 0; i < n/b; i++) t[i] =</pre>
           b*i+b-1-msb(mask[b*i+b-1]);
        for (int j = 1; (1<<j) <= n/b; j++) for (int i = 0;
           i+(1<< j) <= n/b; i++)
            t[n/b*j+i] = op(t[n/b*(j-1)+i],
               t[n/b*(j-1)+i+(1<<(j-1))]);
    }
    int small(int r, int sz = b) { return
       r-msb(mask[r]&((1<<sz)-1));}
    T query(int 1, int r) {
        if (r-l+1 \le b) return small(r, r-l+1);
        int ans = op(small(l+b-1), small(r));
        int x = 1/b+1, y = r/b-1;
        if (x <= y) {
            int j = msb(y-x+1);
            ans = op(ans, op(t[n/b*j+x],
               t[n/b*j+y-(1<<j)+1]));
        }
```

```
return ans:
    }
};
namespace lca {
    vector < int > g[MAX];
    int v[2*MAX], pos[MAX], dep[2*MAX];
    int t;
    rmq < int > RMQ;
    void dfs(int i, int d = 0, int p = -1) {
        v[t] = i, pos[i] = t, dep[t++] = d;
        for (int j : g[i]) if (j != p) {
            dfs(j, d+1, i);
            v[t] = i, dep[t++] = d;
        }
    }
    void build(int n, int root) {
        t = 0;
        dfs(root);
        RMQ = rmq < int > (vector < int > (dep, dep + 2*n - 1));
    int lca(int a, int b) {
        a = pos[a], b = pos[b];
        return v[RMQ.query(min(a, b), max(a, b))];
    int dist(int a, int b) {
        return dep[pos[a]] + dep[pos[b]] - 2*dep[pos[lca(a,
            b)]];
    }
}
vector < int > virt[MAX];
#warning lembrar de buildar o LCA antes
int build_virt(vector<int> v) {
    auto cmp = [&](int i, int j) { return lca::pos[i] <</pre>
       lca::pos[j]; };
    sort(v.begin(), v.end(), cmp);
    for (int i = v.size()-1; i; i--)
       v.push_back(lca::lca(v[i], v[i-1]));
    sort(v.begin(), v.end(), cmp);
```

# 1.27 Euler Path / Euler Cycle

```
// Para declarar: 'euler < true > E(n); ' se quiser
// direcionado e com 'n' vertices
// As funcoes retornam um par com um booleano
// indicando se possui o cycle/path que voce pediu,
// e um vector de {vertice, id da aresta para chegar no
   vertice}
// Se for get_path, na primeira posicao o id vai ser -1
// get_path(src) tenta achar um caminho ou ciclo euleriano
// comecando no vertice 'src'.
// Se achar um ciclo, o primeiro e ultimo vertice serao
   'src'.
// Se for um P3, um possiveo retorno seria [0, 1, 2, 0]
// get_cycle() acha um ciclo euleriano se o grafo for
   euleriano.
// Se for um P3, um possivel retorno seria [0, 1, 2]
// (vertie inicial nao repete)
//
// O(n+m)
template < bool directed = false > struct euler {
    vector < vector < ii >> g;
    vector<int> used;
    euler(int n_) : n(n_), g(n) {}
    void add(int a, int b) {
        int at = used.size();
        used.push_back(0);
        g[a].push_back({b, at});
```

```
if (!directed) g[b].push_back({a, at});
#warning chamar para o src certo!
    pair < bool, vector < ii >> get_path(int src) {
        if (!used.size()) return {true, {}};
        vector < int > beg(n, 0);
        for (int& i : used) i = 0;
        // {{vertice, anterior}, label}
        vector<pair<ii, int>> ret, st = {{src, -1}, -1}};
        while (st.size()) {
            int at = st.back().f.f;
            int& it = beg[at];
            while (it < g[at].size() and used[g[at][it].s])</pre>
               it++:
            if (it == g[at].size()) {
                if (ret.size() and ret.back().f.s != at)
                    return {false, {}};
                ret.push_back(st.back()), st.pop_back();
            } else {
                st.push_back({{g[at][it].f, at},
                   g[at][it].s});
                used[g[at][it].s] = 1;
            }
        }
        if (ret.size() != used.size()+1) return {false, {}};
        vector<ii> ans;
        for (auto i : ret) ans.pb({i.f.f, i.s});
        reverse(ans.begin(), ans.end());
        return {true, ans};
    pair < bool, vector < ii >> get_cycle() {
        if (!used.size()) return {true, {}};
        int src = 0;
        while (!g[src].size()) src++;
        auto ans = get_path(src);
        if (!ans.f or ans.s[0].f != ans.s.back().f) return
           {false, {}};
        ans.s[0].s = ans.s.back().s;
        ans.s.pop_back();
        return ans;
};
```

# 1.28 MinCostMaxFlow Papa

```
// min_cost_flow(s, t, f) computa o par (fluxo, custo)
// com max(fluxo) <= f que tenha min(custo)</pre>
// min_cost_flow(s, t) -> Fluxo maximo de custo minimo de s
// Se tomar TLE, aleatorizar a ordem dos vertices no SPFA
template < typename T > struct mcmf {
    struct edge {
        int to, rev, flow, cap; // para, id da reversa,
           fluxo, capacidade
        bool res; // se eh reversa
        T cost; // custo da unidade de fluxo
        edge(): to(0), rev(0), flow(0), cap(0), cost(0),
           res(false){}
        edge(int to_, int rev_, int flow_, int cap_, T
           cost_, bool res_)
            : to(to_), rev(rev_), flow(flow_), cap(cap_),
               res(res_), cost(cost_){}
    };
    vector < vector < edge >> g;
    vector<int> par_idx, par;
    mcmf(int n) : g(n), par_idx(n), par(n) {}
    void add(int u, int v, int w, T cost) { // de u pra v
       com cap w e custo cost
        edge a = edge(v, g[v].size(), 0, w, cost, false);
        edge b = edge(u, g[u].size(), 0, 0, -cost, true);
        g[u].push_back(a);
        g[v].push_back(b);
    }
    bool spfa(int s, int t) {
        deque < int > q;
        vector < bool > is_inside(g.size(), 0);
        T inf = numeric_limits <T>::max() / 3;
        vector<T> dist(g.size(), inf);
```

```
dist[s] = 0;
    is_inside[s] = true;
    q.push_back(s);
    while (!q.empty()) {
        int u = q.front();
        is_inside[u] = false;
        q.pop_front();
        for (int i = 0; i < (int)g[u].size(); i++)</pre>
            if (g[u][i].cap > g[u][i].flow && dist[u] +
               g[u][i].cost < dist[g[u][i].to]) {</pre>
                dist[g[u][i].to] = dist[u] +
                    g[u][i].cost;
                par_idx[g[u][i].to] = i;
                par[g[u][i].to] = u;
                if (is_inside[g[u][i].to]) continue;
                if (!q.empty() && dist[g[u][i].to] >
                    dist[q.front()])
                    q.push_back(g[u][i].to);
                 else q.push_front(g[u][i].to);
                 is_inside[g[u][i].to] = true;
            }
    }
    return dist[t] != inf;
}
pair < int , T > min_cost_flow(int s, int t, int flow = INF)
   ₹
    int f = 0:
    T ret = 0:
    while (f <= flow && spfa(s, t)) {</pre>
        int mn_flow = flow - f, u = t;
        while (u != s){
            mn_flow = min(mn_flow,
                g[par[u]][par_idx[u]].cap -
               g[par[u]][par_idx[u]].flow);
            u = par[u];
```

```
}
            u = t;
            while (u != s) {
                g[par[u]][par_idx[u]].flow += mn_flow;
                g[u][g[par[u]][par_idx[u]].rev].flow -=
                    mn_flow;
                ret += g[par[u]][par_idx[u]].cost * mn_flow;
                u = par[u];
            }
            f += mn_flow;
        }
        return make_pair(f, ret);
    }
    // Opicional: Retorna todas as arestas originais por
       onde passa fluxo = capacidade.
    vector < pair < int , int >> recover() {
        vector < pair < int , int >> used;
        for (int i = 0; i < g.size(); i++) for (edge e :</pre>
           g[i])
            if(e.flow == e.cap && !e.res) used.push_back({i,
                e.to});
        return used;
   }
};
      Link-cut Tree - aresta
1.29
// Valores nas arestas
// rootify(v) torna v a raiz de sua arvore
// query(v, w) retorna a soma do caminho v--w
// update(v, w, x) soma x nas arestas do caminho v--w
```

// Todas as operacoes sao O(log(n)) amortizado

24

namespace lct {

struct node {

```
int p, ch[2];
    ll val, sub;
    bool rev;
    int sz, ar;
    ll lazy;
    node() {}
    node(int v, int ar_) :
    p(-1), val(v), sub(v), rev(0), sz(ar_), ar(ar_),
       lazv(0) {
       ch[0] = ch[1] = -1;
    }
};
node t[2*MAX]; // MAXN + MAXQ
map<ii, int> aresta;
int sz;
void prop(int x) {
    if (t[x].lazy) {
       if (t[x].ar) t[x].val += t[x].lazy;
       t[x].sub += t[x].lazy*t[x].sz;
       if (t[x].ch[0]+1) t[t[x].ch[0]].lazy +=
           t[x].lazy;
        if (t[x].ch[1]+1) t[t[x].ch[1]].lazy +=
           t[x].lazy;
    }
    if (t[x].rev) {
        swap(t[x].ch[0], t[x].ch[1]);
        if (t[x].ch[0]+1) t[t[x].ch[0]].rev ^= 1;
        if (t[x].ch[1]+1) t[t[x].ch[1]].rev ^= 1;
    t[x].lazy = 0, t[x].rev = 0;
}
void update(int x) {
    t[x].sz = t[x].ar, t[x].sub = t[x].val;
    for (int i = 0; i < 2; i++) if (t[x].ch[i]+1) {</pre>
        prop(t[x].ch[i]);
       t[x].sz += t[t[x].ch[i]].sz;
       t[x].sub += t[t[x].ch[i]].sub;
    }
}
bool is_root(int x) {
```

```
return t[x].p == -1 or (t[t[x].p].ch[0] != x and
       t[t[x].p].ch[1] != x);
}
void rotate(int x) {
    int p = t[x].p, pp = t[p].p;
    if (!is_root(p)) t[pp].ch[t[pp].ch[1] == p] = x;
    bool d = t[p].ch[0] == x;
    t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
    if (t[p].ch[!d]+1) t[t[p].ch[!d]].p = p;
    t[x].p = pp, t[p].p = x;
    update(p), update(x);
}
int splay(int x) {
    while (!is_root(x)) {
        int p = t[x].p, pp = t[p].p;
        if (!is_root(p)) prop(pp);
        prop(p), prop(x);
        if (!is_root(p)) rotate((t[pp].ch[0] ==
           p)^(t[p].ch[0] == x) ? x : p);
        rotate(x);
    }
    return prop(x), x;
}
int access(int v) {
    int last = -1;
    for (int w = v; w+1; update(last = w), splay(v), w =
       t[v].p)
        splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
    return last;
}
void make_tree(int v, int w=0, int ar=0) { t[v] =
   node(w, ar); }
int find_root(int v) {
    access(v), prop(v);
    while (t[v].ch[0]+1) v = t[v].ch[0], prop(v);
    return splay(v);
}
bool conn(int v, int w) {
    access(v), access(w);
    return v == w ? true : t[v].p != -1;
}
void rootify(int v) {
```

```
access(v):
        t[v].rev ^= 1;
    11 query(int v, int w) {
        rootify(w), access(v);
        return t[v].sub;
    void update(int v, int w, int x) {
        rootify(w), access(v);
        t[v].lazy += x;
    void link_(int v, int w) {
        rootify(w);
       t[w].p = v;
    void link(int v, int w, int x) { // v--w com peso x
        int id = MAX + sz++;
        aresta[make_pair(v, w)] = id;
        make_tree(id, x, 1);
        link_(v, id), link_(id, w);
    }
    void cut_(int v, int w) {
        rootify(w), access(v);
        t[v].ch[0] = t[t[v].ch[0]].p = -1;
    }
    void cut(int v, int w) {
        int id = aresta[make_pair(v, w)];
        cut_(v, id), cut_(id, w);
    }
    int lca(int v, int w) {
        access(v);
        return access(w);
   }
1.30 Link-cut Tree - vertice
```

}

```
// Valores nos vertices
// make_tree(v, w) cria uma nova arvore com um
// vertice soh com valor 'w'
```

```
// rootify(v) torna v a raiz de sua arvore
// query(v, w) retorna a soma do caminho v--w
// update(v, w, x) soma x nos vertices do caminho v--w
//
// Todas as operacoes sao O(log(n)) amortizado
namespace lct {
    struct node {
        int p, ch[2];
        ll val, sub;
        bool rev;
        int sz;
        ll lazy;
        node() {}
        node(int v) : p(-1), val(v), sub(v), rev(0), sz(1),
           lazy(0) {
            ch[0] = ch[1] = -1;
        }
    };
    node t[MAX];
    void prop(int x) {
        if (t[x].lazy) {
            t[x].val += t[x].lazy, t[x].sub +=
               t[x].lazy*t[x].sz;
            if (t[x].ch[0]+1) t[t[x].ch[0]].lazy +=
               t[x].lazv;
            if (t[x].ch[1]+1) t[t[x].ch[1]].lazy +=
               t[x].lazy;
        }
        if (t[x].rev) {
            swap(t[x].ch[0], t[x].ch[1]);
            if (t[x].ch[0]+1) t[t[x].ch[0]].rev ^= 1;
            if (t[x].ch[1]+1) t[t[x].ch[1]].rev ^= 1;
        t[x].lazy = 0, t[x].rev = 0;
    }
    void update(int x) {
        t[x].sz = 1, t[x].sub = t[x].val;
        for (int i = 0; i < 2; i++) if (t[x].ch[i]+1) {</pre>
            prop(t[x].ch[i]);
```

```
t[x].sz += t[t[x].ch[i]].sz;
        t[x].sub += t[t[x].ch[i]].sub;
    }
}
bool is_root(int x) {
    return t[x].p == -1 or (t[t[x].p].ch[0] != x and
       t[t[x].p].ch[1] != x);
void rotate(int x) {
    int p = t[x].p, pp = t[p].p;
    if (!is_root(p)) t[pp].ch[t[pp].ch[1] == p] = x;
    bool d = t[p].ch[0] == x;
    t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
    if (t[p].ch[!d]+1) t[t[p].ch[!d]].p = p;
    t[x].p = pp, t[p].p = x;
    update(p), update(x);
}
int splay(int x) {
    while (!is_root(x)) {
        int p = t[x].p, pp = t[p].p;
        if (!is_root(p)) prop(pp);
        prop(p), prop(x);
        if (!is_root(p)) rotate((t[pp].ch[0] ==
           p)^{(t[p].ch[0] == x)} ? x : p);
        rotate(x);
    }
    return prop(x), x;
int access(int v) {
    int last = -1;
    for (int w = v; w+1; update(last = w), splay(v), w =
        splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
    return last:
}
void make_tree(int v, int w) { t[v] = node(w); }
int find_root(int v) {
    access(v), prop(v);
    while (t[v].ch[0]+1) v = t[v].ch[0], prop(v);
    return splay(v);
}
bool connected(int v, int w) {
```

```
access(v), access(w);
        return v == w ? true : t[v].p != -1;
    }
    void rootify(int v) {
        access(v);
        t[v].rev ^= 1;
    }
    11 query(int v, int w) {
        rootify(w), access(v);
        return t[v].sub;
    }
    void update(int v, int w, int x) {
        rootify(w), access(v);
        t[v].lazy += x;
    }
    void link(int v, int w) {
        rootify(w);
        t[w].p = v;
    }
    void cut(int v, int w) {
        rootify(w), access(v);
        t[v].ch[0] = t[t[v].ch[0]].p = -1;
    }
    int lca(int v, int w) {
        access(v);
        return access(w);
   }
}
```

## 1.31 Link-cut Tree

```
// Link-cut tree padrao
//
// Todas as operacoes sao O(log(n)) amortizado
namespace lct {
    struct node {
        int p, ch[2];
        node() { p = ch[0] = ch[1] = -1; }
};
```

```
node t[MAX];
bool is_root(int x) {
    return t[x].p == -1 or (t[t[x].p].ch[0] != x and
       t[t[x].p].ch[1] != x);
void rotate(int x) {
    int p = t[x].p, pp = t[p].p;
    if (!is_root(p)) t[pp].ch[t[pp].ch[1] == p] = x;
    bool d = t[p].ch[0] == x;
    t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
    if (t[p].ch[!d]+1) t[t[p].ch[!d]].p = p;
    t[x].p = pp, t[p].p = x;
}
void splay(int x) {
    while (!is_root(x)) {
        int p = t[x].p, pp = t[p].p;
       if (!is_root(p)) rotate((t[pp].ch[0] ==
           p)^{(t[p].ch[0]} == x) ? x : p);
        rotate(x);
    }
}
int access(int v) {
    int last = -1;
    for (int w = v; w+1; last = w, splay(v), w = t[v].p)
        splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
    return last;
}
int find_root(int v) {
    access(v);
    while (t[v].ch[0]+1) v = t[v].ch[0];
    return splay(v), v;
}
void link(int v, int w) { // v deve ser raiz
    access(v):
    t[v].p = w;
void cut(int v) { // remove aresta de v pro pai
    access(v);
    t[v].ch[0] = t[t[v].ch[0]].p = -1;
}
```

```
int lca(int v, int w) {
    return access(v), access(w);
}
```

#### 1.32 Line Tree

```
// Reduz min-query em arvore para RMQ
// Se o grafo nao for uma arvore, as queries
// sao sobre a arvore geradora maxima
// Queries de minimo
//
// build - O(n log(n))
// query - O(log(n))
int n;
namespace linetree {
    int id[MAX], seg[2*MAX], pos[MAX];
    vector < int > v[MAX], val[MAX];
    vector<pair<int, pair<int, int> > ar;
    void add(int a, int b, int p) { ar.pb({p, {a, b}}); }
    void build() {
        sort(ar.rbegin(), ar.rend());
        for (int i = 0; i < n; i++) id[i] = i, v[i] = {i},
           val[i].clear();
        for (auto i : ar) {
            int a = id[i.second.first], b =
               id[i.second.second]:
            if (a == b) continue:
            if (v[a].size() < v[b].size()) swap(a, b);</pre>
            for (auto j : v[b]) id[j] = a, v[a].push_back(j);
            val[a].push_back(i.first);
            for (auto j : val[b]) val[a].push_back(j);
            v[b].clear(), val[b].clear();
        }
        vector < int > vv;
        for (int i = 0; i < n; i++) for (int j = 0; j <
           v[i].size(); j++) {
```

```
pos[v[i][j]] = vv.size();
            if (j + 1 < v[i].size()) vv.push_back(val[i][j]);</pre>
            else vv.push_back(0);
        }
        for (int i = n; i < 2*n; i++) seg[i] = vv[i-n];</pre>
        for (int i = n-1; i; i--) seg[i] = min(seg[2*i],
           seg[2*i+1]);
    int query(int a, int b) {
        if (id[a] != id[b]) return 0; // nao estao conectados
        a = pos[a], b = pos[b];
        if (a > b) swap(a, b);
        b--;
        int ans = INF;
        for (a += n, b += n; a \le b; ++a/=2, --b/=2) ans =
           min({ans, seg[a], seg[b]});
        return ans;
    }
};
```

# 1.33 Floyd-Warshall

```
// encontra o menor caminho entre todo
// par de vertices e detecta ciclo negativo
// returna 1 sse ha ciclo negativo
// d[i][i] deve ser 0
// para i != j, d[i][j] deve ser w se ha uma aresta
// (i, j) de peso w, INF caso contrario
//
// O(n^3)

int n;
int d[MAX][MAX];

bool floyd_warshall() {
   for (int k = 0; k < n; k++)
   for (int i = 0; i < n; i++)
   for (int j = 0; j < n; j++)
        d[i][j] = min(d[i][j], d[i][k] + d[k][j]);</pre>
```

```
for (int i = 0; i < n; i++)
        if (d[i][i] < 0) return 1;

return 0;
}</pre>
```

# 2 Matematica

# 2.1 Ordem de elemento do grupo

```
// Calcula a ordem de a em Z_n
// O grupo Zn eh ciclico sse n =
// 1, 2, 4, p^k ou 2 p^k, p primo impar
// Retorna -1 se nao achar
//
// O(sqrt(n) log(n))
int tot(int n); // totiente em O(sqrt(n))
int expo(int a, int b, int m); // (a^b) %m em O(log(b))
// acha todos os divisores ordenados em O(sqrt(n))
vector < int > div(int n) {
    vector<int> ret1, ret2;
    for (int i = 1; i*i <= n; i++) if (n % i == 0) {
        ret1.pb(i);
        if (i*i != n) ret2.pb(n/i);
    }
    for (int i = ret2.size()-1; i+1; i--) ret1.pb(ret2[i]);
    return ret1;
}
int ordem(int a, int n) {
    vector < int > v = div(tot(n));
    for (int i : v) if (expo(a, i, n) == 1) return i;
    return -1;
}
```

#### 2.2 Teorema Chines do Resto

```
// Combina equacoes modulares lineares: x = a (mod m)
// O m final eh o lcm dos m's, e a resposta eh unica mod o
   lcm
// Os m nao precisam ser coprimos
// Se nao tiver solucao, o 'a' vai ser -1
tuple < 11, 11, 11 > ext_gcd(11 a, 11 b) {
    if (!a) return {b, 0, 1};
    auto [g, x, y] = ext_gcd(b%a, a);
    return \{g, y - b/a*x, x\};
}
struct crt {
    ll a, m;
    crt() : a(0), m(1) {}
    crt(ll a_, ll m_) : a(a_), m(m_) {}
    crt operator * (crt C) {
        auto [g, x, y] = ext_gcd(m, C.m);
        if ((a - C.a) \% g) a = -1;
        if (a == -1 \text{ or } C.a == -1) \text{ return } crt(-1, 0);
        11 lcm = m/g*C.m;
        11 ans = a + (x*(C.a-a)/g \% (C.m/g))*m;
        return crt((ans % lcm + lcm) % lcm, lcm);
    }
};
```

# 2.3 Pollard's Rho Alg

```
// Usa o algoritmo de deteccao de ciclo de Brent
// A fatoracao nao sai necessariamente ordenada
// O algoritmo rho encontra um fator de n,
// e funciona muito bem quando n possui um fator pequeno
//
// Complexidades (considerando mul constante):
// rho - esperado O(n^(1/4)) no pior caso
// fact - esperado menos que O(n^(1/4) log(n)) no pior caso
```

```
mt19937 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
ll uniform(ll 1, ll r){
    uniform_int_distribution < ll > uid(l, r);
    return uid(rng);
}
11 mul(ll a, ll b, ll m) {
   ll ret = a*b - ll(a*(long double)b/m+0.5)*m;
   return ret < 0 ? ret+m : ret;</pre>
}
ll expo(ll a, ll b, ll m) {
   if (!b) return 1;
    ll ans = expo(mul(a, a, m), b/2, m);
   return b%2 ? mul(a, ans, m) : ans;
}
bool prime(ll n) {
   if (n < 2) return 0;
   if (n <= 3) return 1;
    if (n % 2 == 0) return 0;
    11 d = n - 1;
    int r = 0;
    while (d % 2 == 0) {
        r++;
        d /= 2;
    }
    for (int i : {2, 325, 9375, 28178, 450775, 9780504,
       795265022}) {
       if (i >= n) break;
        ll x = expo(i, d, n);
        if (x == 1 or x == n - 1) continue;
        bool deu = 1;
        for (int j = 0; j < r - 1; j++) {
            x = mul(x, x, n);
           if (x == n - 1) {
                deu = 0:
```

```
break:
            }
        if (deu) return 0;
    return 1;
}
11 rho(11 n) {
    if (n == 1 or prime(n)) return n;
    if (n % 2 == 0) return 2;
    while (1) {
        11 x = 2, y = 2, ciclo = 2, i = 0, d = 1;
        ll c = uniform(1, n-1):
        while (d == 1) {
            if (++i == ciclo) ciclo *= 2, y = x;
            x = (mul(x, x, n) + c) \% n;
            if (x == y) break;
            d = \_gcd(abs(x-y), n);
        if (x != y) return d;
}
void fact(ll n, vector<ll>& v) {
    if (n == 1) return;
    if (prime(n)) v.pb(n);
    else {
        11 d = rho(n):
       fact(d, v);
        fact(n / d, v);
   }
}
```

#### 2.4 Algoritmo de Euclides extendido

```
// acha x e y tal que ax + by = mdc(a, b) (nao eh unico)
//
// O(log(min(a, b)))
tuple < 11, 11, 11 > ext_gcd(11 a, 11 b) {
    if (!a) return {b, 0, 1};
    auto [g, x, y] = ext_gcd(b\%a, a);
   return \{g, y - b/a*x, x\};
}
     Eliminação Gaussiana de XOR.
// insert(mask) insere uma mask no espaco vetorial
// get(X) retorna outra uma mask com os caras da base
// cujo xor da X, ou -1 se n tem como
//
// O(log(MAXN))
int basis[LOG]; // basis[i] = mask do cara com bit mais
   significativo i
int rk: // tamanho da base
void insert(int mask) {
    for (int i = LOG - 1; i >= 0; i--) if (mask>>i&1) {
        if (!basis[i]) {
            basis[i] = mask, rk++;
            return:
        }
        mask ^= basis[i];
   }
}
int get(int mask) {
   int ret = 0;
   for (int i = LOG - 1; i >= 0; i--) if (mask>>i&1) {
        if (!basis[i]) return -1;
        mask ^= basis[i], ret |= (1<<i);
    return ret;
```

}

# 2.6 Algoritmo de Euclides

```
// O(log(min(a, b)))
int mdc(int a, int b) {
    return !b ? a : mdc(b, a % b);
}
```

#### 2.7 Binomial Distribution

```
// binom(n, k, p) retorna a probabilidade de k sucessos
// numa binomial(n, p)
mt19937 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
double logfact[MAX];
void calc(){
    logfact[0] = 0;
    for (int i = 1; i < MAX; i++)</pre>
        logfact[i] = logfact[i-1] + log(i);
}
double binom(int n, int k, double p){
    return exp(logfact[n] - logfact[k] - logfact[n-k] + k *
       log(p) + (n-k) * log(1 - p));
}
int main(){//if you want to sample from a bin(n, p)
    calc():
    int n; double p;
    cin >> n >> p;
    binomial_distribution < int > distribution (n, p);
    int IT = 1e5;
    vector<int> freq(n+1, 0);
    for (int i = 0; i < IT; i++){</pre>
        int v = distribution(rng);
        //P(v == k) = (n \text{ choose } k)p^k (1-p)^(n-k) = binom(n,
           k, p)
        freq[v]++;
```

#### 2.8 Divisão de Polinomios

```
// Divide p1 por p2
// Retorna um par com o quociente e o resto
// Os coeficientes devem estar em ordem
// decrescente pelo grau. Ex:
// 3x^2 + 2x - 1 \rightarrow [3, 2, -1]
// O(nm), onde n e m sao os tamanhos dos
// polinomios
typedef vector<int> vi;
pair < vi , vi > div(vi p1, vi p2) {
    vi quoc, resto;
    int a = p1.size(), b = p2.size();
    for (int i = 0; i <= a - b; i++) {
        int k = p1[i] / p2[0];
        quoc.pb(k);
        for (int j = i; j < i + b; j++)
            p1[j] = k * p2[j - i];
    }
    for (int i = a - b + 1; i < a; i++)
        resto.pb(p1[i]);
    return mp(quoc, resto);
}
```

## 2.9 Deteccao de ciclo - Tortoise and Hare

```
// Linear no tanto que tem que andar pra ciclar,
// O(1) de memoria
// Retorna um par com o tanto que tem que andar
// do f0 ate o inicio do ciclo e o tam do ciclo
pair<11, 11> find_cycle() {
    11 \text{ tort} = f(f0);
    ll hare = f(f(f0));
    11 t = 0;
    while (tort != hare) {
        tort = f(tort);
        hare = f(f(hare));
        t++;
    }
    11 st = 0;
    tort = f0;
    while (tort != hare) {
        tort = f(tort);
        hare = f(hare);
        st++;
    }
    11 len = 1;
    hare = f(tort);
    while (tort != hare) {
        hare = f(hare);
        len ++;
    return {st, len};
}
2.10 FFT
// Exemplos na main
// Soma O(n) & Multiplicacao O(nlogn)
template < typename T > void fft(vector < T > &a, bool f, int N,
   vector<int> &rev){
    for (int i = 0; i < N; i++)</pre>
```

```
if (i < rev[i])</pre>
             swap(a[i], a[rev[i]]);
    int 1, r, m;
    vector <T> roots(N);
    for (int n = 2; n <= N; n *= 2){</pre>
        T \text{ root} = T :: rt(f, n, N);
        roots[0] = 1;
        for (int i = 1; i < n/2; i++)</pre>
             roots[i] = roots[i-1]*root;
        for (int pos = 0; pos < N; pos += n){
            1 = pos+0, r = pos+n/2, m = 0;
             while (m < n/2) {
                 auto t = roots[m]*a[r]:
                 a[r] = a[1] - t;
                 a[1] = a[1] + t:
                 1++; r++; m++;
            }
        }
    }
    if (f) {
        auto invN = T(1)/N;
        for(int i = 0; i < N; i++) a[i] = a[i]*invN;</pre>
    }
}
template < typename T> struct poly : vector < T> {
    poly(const vector<int> &coef):vector<T>(coef.size()){
        for (int i = 0; i < coef.size(); i++) this->at(i) =
            coef[i];
    poly(const vector<T> &coef):vector<T>(coef){}
    poly(unsigned size, T val = 0):vector<T>(size, val){}
    poly(){}
    T operator()(T x){
        T ans = 0, curr_x(1);
        for (auto c : *this) {
             ans += c*curr_x;
             curr_x *= x;
        return ans;
    }
```

```
poly<T> operator+(const poly<T> &r){
    polv < T > 1 = *this;
    int sz = max(l.size(), r.size());
    l.resize(sz);
    for (int i = 0; i < r.size(); i++)</pre>
        l[i] += r[i];
    return 1:
poly<T> operator - (poly<T> &r){
    for (auto &it : r) it = -it;
    return (*this)+r;
poly<T> operator*(poly<T> r){
    poly < T > 1 = *this;
    int ln = 1.size(), rn = r.size();
    int N = ln+rn+1;
    int log_n = T::fft_len(N);
    int n = 1 << log_n;</pre>
    vector < int > rev(n);
    for (int i = 0; i < n; ++i){</pre>
        rev[i] = 0;
        for (int j = 0; j < log_n; ++j)</pre>
            if (i & (1<<j))</pre>
                 rev[i] = 1 << (log_n-1-j);
    }
    if (N > n) throw logic_error("resulting poly to
       big");
    l.resize(n);
    r.resize(n);
    fft(1, false, n, rev);
    fft(r, false, n, rev);
    for (int i = 0; i < n; i++)</pre>
        l[i] *= r[i];
    fft(1, true, n, rev);
    return 1;
friend ostream& operator << (ostream &out, const poly <T>
   &p){
    if (p.empty()) return out;
    out << p.at(0);
    for (int i = 1; i < p.size(); i++)</pre>
        out << " + " << p.at(i) << "x^" << i;
```

```
out << endl:
        return out;
    }
};
mt19937 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
const int MOD = 998244353;
using mint = mod_int < MOD >;
int main(){
    uniform_int_distribution < int > uid(0, MOD-1);
    int n = (1 << mint::fft_len()/2);</pre>
    auto rand_vec = [&](){
        vector < int > rd(n);
        for (int &i : rd) i = uid(rng);
        return rd;
    };
    poly<mint> p = rand_vec();
    poly < mint > q = rand_vec();
    poly < mint > sum = p+q;
    poly<mint> mult = p*q;
    for (int i = 1; i <= 5000; i++){</pre>
        int x = uid(rng);
        auto P = p(x), Q = q(x), M = mult(x);
        if (P*Q != M) throw logic_error("bad implementation
           :(");
    cout << "sucesso!" << endl;</pre>
    exit(0);
    exit(0):
}
2.11 Baby step Giant step
```

// Resolve Logaritmo Discreto a^x = b mod m, m primo em

```
0(sqrt(n)*hash(n))
// Meet In The Middle, decompondo x = i * ceil(sqrt(n)) - j,
   i,j<=ceil(sqrt(n))
int babyStep(int a, int b, int m)
{
    unordered_map < int , int > mapp;
    int sq=sqrt(m)+1;
    11 asq=1;
    for(int i=0; i<sq; i++)</pre>
         asq=(asq*a)%m;
    11 curr=asq;
    for(int i=1; i<=sq; i++)</pre>
         if(!mapp.count(curr))
             mapp[curr]=i;
         curr=(curr*asq)%m;
    int ret=INF;
    curr=b;
    for(int j=0; j<=sq; j++)</pre>
         if (mapp.count(curr))
             ret=min(ret,(int)(mapp[curr]*sq-j));
         curr=(curr*a)%m;
    if(ret<INF) return ret;</pre>
    return -1;
}
int main()
{
    int a,b,m;
    while(cin>>a>>b>>m,a or b or m)
         int x=babyStep(a,m,b);
         if(x!=-1)
             cout << x << endl;</pre>
         else
             cout << "No Solution" << endl;</pre>
    return 0;
```

#### 2.12 Miller-Rabin

```
// Testa se n eh primo, n <= 3 * 10^18
//
// O(log(n)), considerando multiplicacao
// e exponenciacao constantes
// multiplicacao modular
11 mul(ll a, ll b, ll m) {
    return (a*b-ll(a*(long double)b/m+0.5)*m+m)%m;
}
ll expo(ll a, ll b, ll m) {
    if (!b) return 1;
    11 ans = expo(mul(a, a, m), b/2, m);
   return b%2 ? mul(a, ans, m) : ans;
}
bool prime(ll n) {
   if (n < 2) return 0;
   if (n <= 3) return 1;
   if (n % 2 == 0) return 0;
    11 d = n - 1;
    int r = 0;
    while (d % 2 == 0) {
        r++;
        d /= 2;
    }
    // com esses primos, o teste funciona garantido para n
       <= 2^64
    // funciona para n <= 3*10^24 com os primos ate 41
    for (int i: {2, 325, 9375, 28178, 450775, 9780504,
       1795265022}) {
        if (i >= n) break;
        ll x = expo(i, d, n);
        if (x == 1 \text{ or } x == n - 1) \text{ continue};
        bool deu = 1;
        for (int j = 0; j < r - 1; j++) {
```

```
x = mul(x, x, n);
if (x == n - 1) {
    deu = 0;
    break;
}
if (deu) return 0;
}
return 1;
}
```

#### 2.13 Inverso Modular

```
// Computa o inverso de a modulo b
// Se b eh primo, basta fazer
// a^(b-2)

11 inv(11 a, 11 b) {
    return a > 1 ? b - inv(b%a, a)*b/a : 1;
}

// computa o inverso modular de 1..MAX-1 modulo um primo
11 inv[MAX]:
inv[1] = 1;
for (int i = 2; i < MAX; i++) inv[i] = MOD -
    MOD/i*inv[MOD%i]%MOD;</pre>
```

## 2.14 Karatsuba

```
// Os pragmas realmente deixam muito mais rapido
// Para n ~ 2e5, roda em < 1 s
//
// O(n^1.58)

#pragma GCC optimize("Ofast")
#pragma GCC target ("avx,avx2")
template < typename T > void kar(T* a, T* b, int n, T* r, T* tmp) {
```

```
if (n <= 64) {
        for (int i = 0; i < n; i++) for (int j = 0; j < n;
            r[i+j] += a[i] * b[j];
        return;
    }
    int mid = n/2;
    T * atmp = tmp, *btmp = tmp+mid, *E = tmp+n;
    memset(E, 0, sizeof(E[0])*n);
    for (int i = 0; i < mid; i++) {</pre>
        atmp[i] = a[i] + a[i+mid];
        btmp[i] = b[i] + b[i+mid];
    }
    kar(atmp, btmp, mid, E, tmp+2*n);
    kar(a, b, mid, r, tmp+2*n);
    kar(a+mid, b+mid, mid, r+n, tmp+2*n);
    for (int i = 0; i < mid; i++) {</pre>
        T \text{ temp} = r[i+mid];
        r[i+mid] += E[i] - r[i] - r[i+2*mid];
        r[i+2*mid] += E[i+mid] - temp - r[i+3*mid];
    }
}
template < typename T> vector < T> karatsuba (vector < T> a,
   vector <T> b) {
    int n = max(a.size(), b.size());
    while (n&(n-1)) n++;
    a.resize(n), b.resize(n);
    vector \langle T \rangle ret(2*n), tmp(4*n);
    kar(&a[0], &b[0], n, &ret[0], &tmp[0]);
    return ret;
}
2.15 Variações do crivo de Eratosthenes
```

```
// "0" crivo
//
// Encontra maior divisor primo
// Um numero eh primo sse div[x] == x
// fact fatora um numero <= lim</pre>
```

```
// A fatoracao sai ordenada
// crivo - O(n log(log(n)))
// fact - O(log(n))
int divi[MAX];
void crivo(int lim) {
    for (int i = 1; i <= lim; i++) divi[i] = 1;</pre>
    for (int i = 2; i <= lim; i++) if (divi[i] == 1)</pre>
        for (int j = i; j <= lim; j += i) divi[j] = i;</pre>
}
void fact(vector<int>& v, int n) {
    if (n != divi[n]) fact(v, n/divi[n]);
    v.push_back(divi[n]);
}
// Crivo de divisores
// Encontra numero de divisores
// ou soma dos divisores
//
// O(n log(n))
int divi[MAX];
void crivo(int lim) {
    for (int i = 1; i <= lim; i++) divi[i] = 1;</pre>
    for (int i = 2; i <= lim; i++)</pre>
        for (int j = i; j <= lim; j += i) {</pre>
            // para numero de divisores
            divi[j]++;
            // para soma dos divisores
             divi[j] += i;
}
// Crivo de totiente
```

```
// Encontra o valor da funcao
// totiente de Euler
// O(n log(log(n)))
int tot[MAX];
void crivo(int lim) {
    for (int i = 1; i <= lim; i++) tot[i] = i;</pre>
    for (int i = 2; i <= lim; i++) if (tot[i] == i)</pre>
        for (int j = i; j <= lim; j += i)</pre>
            tot[j] -= tot[j] / i;
}
// Crivo de funcao de mobius
// O(n log(log(n))
char meb[MAX];
void crivo(int lim) {
    for (int i = 2; i <= lim; i++) meb[i] = 2;</pre>
    meb[1] = 1;
    for (int i = 2; i <= lim; i++) if (meb[i] == 2)</pre>
        for (int j = i; j <= lim; j += i) if (meb[j]) {</pre>
            if (meb[j] == 2) meb[j] = 1;
            meb[j] *= j/i\%i ? -1 : 0;
2.16 Totiente
// O(sqrt(n))
int tot(int n){
    int ret = n;
    for (int i = 2; i*i <= n; i++) if (n % i == 0) {
        while (n \% i == 0) n /= i;
```

```
ret -= ret / i:
    if (n > 1) ret -= ret / n;
    return ret;
}
2.17 2-SAT
// Retorna se eh possivel atribuir valores
// ans[i] fala se a variavel 'i' eh verdadeira
// Grafo tem que caber 2n vertices
// add(x, y) adiciona implicacao x -> y
// Para adicionar uma clausula (x ou y)
// chamar add(nao(x), y)
// Se x tem que ser verdadeiro, chamar add(nao(x), x)
// O(|V|+|E|)
vector < int > g[MAX];
int n, vis[MAX], comp[MAX];
stack<int> s;
int id[MAX], p;
vector < int > ord;
int ans[MAX];
int dfs(int k) {
    int lo = id[k] = p++;
    s.push(k);
    vis[k] = 2;
    for (int i = 0; i < g[k].size(); i++) {</pre>
        if (!vis[g[k][i]])
            lo = min(lo, dfs(g[k][i]));
        else if (vis[g[k][i]] == 2)
            lo = min(lo, id[g[k][i]]);
    if (lo == id[k]) while (1) {
```

int u = s.top();

comp[u] = k;

s.pop(); vis[u] = 1;

```
ord.pb(u):
        if (u == k) break;
    return lo;
void tarjan() {
    memset(vis, 0, sizeof(vis));
    p = 0;
    for (int i = 0; i < 2*n; i++) if (!vis[i]) dfs(i);</pre>
}
int nao(int x) { return (x + n) \% (2*n); }
// x \rightarrow y = !x ou y
void add(int x, int y){
    g[x].pb(y);
    // contraposicao
    g[nao(y)].pb(nao(x));
}
bool doisSAT(){
    tarjan();
    for (int i = 0; i < n; i++)</pre>
        if (comp[i] == comp[nao(i)]) return 0;
    memset(ans, -1, sizeof(ans));
    for(auto at: ord) {
        if (ans[at] != -1) continue;
        ans[at] = 1;
        ans[nao(at)] = 0;
    }
    return 1;
}
```

#### 2.18 Exponenciacao rapida

### 2.19 Produto de dois long long mod m

# 3 Primitivas

# 3.1 Complex

```
struct cplx{
   double r, i;
```

```
cplx(complex < double > c):r(c.real()), i(c.imag()){}
cplx(){}
cplx(double r_{-}, double i_{-} = 0):r(r_{-}), i(i_{-})
double abs(){ return hypot(r, i); }
double abs2(){ return r*r + i*i; }
cplx inv() { return cplx(r/abs2(), i/abs2()); }
cplx& operator+=(cplx a){
    r += a.r; i += a.i;
    return *this;
}
cplx& operator -=(cplx a){
    r -= a.r; i -= a.i;
    return *this;
}
cplx& operator*=(cplx a){
    double r_{-} = r*a.r - i*a.i;
    double i_ = r*a.i + i*a.r;
    r = r_{.};
    i = i_;
    return *this;
}
cplx conj(){
    return cplx(r, -i);
cplx& operator/=(cplx a){
    auto a_ = a.inv();
    return (*this)*=a_;
cplx operator-(){ return cplx(-r, -i); }
cplx& operator = (double e){
    return *this = pow(complex < double > (r, i), e);
friend ostream &operator << (ostream &out, cplx a) {</pre>
    return out << a.r << " + " << a.i << "i":
friend cplx operator+(cplx a, cplx b){ return a+=b; }
friend cplx operator-(cplx a, cplx b){ return a-=b; }
friend cplx operator*(cplx a, cplx b){ return a*=b; }
friend cplx operator/(cplx a, cplx b){ return a/=b; }
friend cplx operator^(cplx a, double e) { return a^=e; }
//fft
```

```
static int fft_len(int N){
   int n = 1, log_n = 0;
   while (n <= N) { n <<= 1; log_n++; }
   return log_n;
}
static cplx rt(bool f, int n, int N){
   const static double PI = acos(-1);
   double alpha = (2*PI)/n;
   if (f) alpha = -alpha;
   return cplx(cos(alpha), sin(alpha));
}
};</pre>
```

#### 3.2 Aritmetica Modular

```
// O mod tem q ser primo
template < int p> struct mod_int {
    ll pow(ll b, ll e) {
        if (e == 0) return 1;
        ll r = pow(b*b%p, e/2);
        if (e\%2 == 1) r = (r*b)\%p;
        return r;
    11 inv(11 b) { return pow(b, p-2); }
    using m = mod_int;
    int v;
    mod int() {}
    mod int(ll v ) {
        if (v_ >= p || v_ <= -p) v_ %= p;
        if (v_{-} < 0) v_{-} += p;
        v = v_{-};
    }
    m& operator+=(const m &a) {
        v += a.v;
        if (v >= p) v -= p;
        return *this;
    }
    m& operator -= (const m &a) {
```

```
v -= a.v:
    if (v < 0) v += p;
    return *this;
}
m& operator*=(const m &a) {
    v = (v*ll(a.v))\%p;
    return *this;
}
m& operator/=(const m &a) {
    v = (v*inv(a.v))%p;
    return *this;
}
m operator-(){ return m(-v); }
m& operator^=(11 e) {
    if (e < 0){
        v = inv(v):
        e = -e;
    v = pow(v, e\%(p-1));
    return *this;
}
bool operator == (const m &a) { return v == a.v; }
bool operator!=(const m &a) { return v != a.v; }
friend istream &operator>>(istream &in, m& a) {
    11 val; in >> val;
    a = m(val);
    return in;
friend ostream &operator << (ostream &out, m a) {</pre>
    return out << a.v;</pre>
friend m operator+(m a, m b) { return a+=b; }
friend m operator-(m a, m b) { return a-=b; }
friend m operator*(m a, m b) { return a*=b; }
friend m operator/(m a, m b) { return a/=b; }
friend m operator^(m a, ll e) { return a^=e; }
static int fft_len(int n = -1){
    // max k such that 2^k | p-1
    if (p == 998244353) return 20;
    throw logic_error("find an order");
    return -1;
```

```
static m rt(bool f, int n, int N){
    // an element of order fft_len
    if (p == 998244353){
        m r(695449733);
        if (f) r = r^(-1);
        return r^(N/n);
    }

    throw logic_error("find a root");
    return -1; // return x so that x^(2^k) != x*x^(2^k)
        = 1
}

};

typedef mod_int<(int)1e9+7> mint;
```

#### 3.3 Primitivas de Polinomios

```
#include <bits/stdc++.h>
using namespace std;
namespace algebra {
    const int inf = 1e9;
    const int magic = 500; // threshold for sizes to run the
       naive algo
    namespace fft {
        const int maxn = 1 << 18;</pre>
        typedef double ftype;
        typedef complex <ftype > point;
        point w[maxn];
        const ftype pi = acos(-1);
        bool initiated = 0;
        void init() {
            if(!initiated) {
                 for(int i = 1; i < maxn; i *= 2) {</pre>
                     for(int j = 0; j < i; j++) {</pre>
```

```
w[i + j] = polar(ftype(1), pi * j /
                    i);
            }
        }
        initiated = 1;
    }
}
template < typename T>
    void fft(T *in, point *out, int n, int k = 1) {
        if(n == 1) {
            *out = *in;
        } else {
            n /= 2:
            fft(in, out, n, 2 * k);
            fft(in + k, out + n, n, 2 * k);
            for(int i = 0; i < n; i++) {</pre>
                 auto t = out[i + n] * w[i + n];
                 out[i + n] = out[i] - t;
                 out[i] += t;
            }
        }
    }
template < typename T>
    void mul_slow(vector<T> &a, const vector<T> &b) {
        vector <T> res(a.size() + b.size() - 1);
        for(size_t i = 0; i < a.size(); i++) {</pre>
            for(size_t j = 0; j < b.size(); j++) {</pre>
                 res[i + j] += a[i] * b[j];
            }
        }
        a = res;
    }
template < typename T>
    void mul(vector<T> &a, const vector<T> &b) {
        if(min(a.size(), b.size()) < magic) {</pre>
            mul_slow(a, b);
            return;
        init();
```

```
shift) - 1;
            size_t n = a.size() + b.size() - 1;
            while(__builtin_popcount(n) != 1) {
                n++;
            }
            a.resize(n);
            static point A[maxn], B[maxn];
            static point C[maxn], D[maxn];
            for(size_t i = 0; i < n; i++) {</pre>
                A[i] = point(a[i] & mask, a[i] >> shift);
                if(i < b.size()) {</pre>
                    B[i] = point(b[i] & mask, b[i] >>
                        shift):
                } else {
                    B[i] = 0:
                }
            }
            fft(A, C, n); fft(B, D, n);
            for(size_t i = 0; i < n; i++) {</pre>
                point c0 = C[i] + conj(C[(n - i) \% n]);
                point c1 = C[i] - conj(C[(n - i) \% n]);
                point d0 = D[i] + conj(D[(n - i) \% n]);
                point d1 = D[i] - conj(D[(n - i) \% n]);
                A[i] = c0 * d0 - point(0, 1) * c1 * d1;
                B[i] = c0 * d1 + d0 * c1;
            }
            fft(A, C, n); fft(B, D, n);
            reverse(C + 1, C + n);
            reverse(D + 1, D + n);
            int t = 4 * n;
            for(size_t i = 0; i < n; i++) {</pre>
                int64_t A0 = llround(real(C[i]) / t);
                T A1 = llround(imag(D[i]) / t);
                T A2 = llround(imag(C[i]) / t);
                a[i] = A0 + (A1 << shift) + (A2 << 2 *
                    shift):
            }
            return;
template < typename T>
```

static const int shift = 15, mask = (1 <<</pre>

```
T bpow(T x, size_t n) {
        return n ? n % 2 ? x * bpow(x, n - 1) : bpow(x *
           x, n / 2) : T(1);
template < typename T>
   T bpow(T x, size_t n, T m) {
        return n ? n % 2 ? x * bpow(x, n - 1, m) % m :
           bpow(x * x \% m, n / 2, m) : T(1);
template < typename T>
   T gcd(const T &a, const T &b) {
        return b == T(0) ? a : gcd(b, a % b);
template < typename T>
    T nCr(T n, int r) { // runs in O(r)}
        T res(1):
        for(int i = 0; i < r; i++) {</pre>
            res *= (n - T(i));
            res /= (i + 1);
        return res;
   }
template < int m>
    struct modular {
        int64_t r;
        modular() : r(0) {}
        modular(int64_t rr) : r(rr) {if(abs(r) >= m) r}
           %= m; if(r < 0) r += m;}
        modular inv() const {return bpow(*this, m - 2);}
        modular operator * (const modular &t) const
           {return (r * t.r) % m;}
        modular operator / (const modular &t) const
           {return *this * t.inv();}
        modular operator += (const modular &t) {r +=
           t.r; if(r >= m) r -= m; return *this;}
        modular operator -= (const modular &t) {r -=
           t.r; if(r < 0) r += m; return *this;}
        modular operator + (const modular &t) const
           {return modular(*this) += t;}
        modular operator - (const modular &t) const
           {return modular(*this) -= t;}
```

```
modular operator *= (const modular &t) {return
           *this = *this * t;}
        modular operator /= (const modular &t) {return
           *this = *this / t;}
        bool operator == (const modular &t) const
           {return r == t.r;}
        bool operator != (const modular &t) const
           {return r != t.r;}
        operator int64_t() const {return r;}
   };
template < int T>
    istream& operator >> (istream &in, modular<T> &x) {
       return in >> x.r:
   }
template < typename T>
   struct poly {
        vector <T> a;
        void normalize() { // get rid of leading zeroes
            while(!a.empty() && a.back() == T(0)) {
                a.pop_back();
           }
        }
        polv(){}
        poly(T a0) : a{a0}{normalize();}
        poly(vector <T> t) : a(t) { normalize();}
        poly operator += (const poly &t) {
            a.resize(max(a.size(), t.a.size()));
            for(size_t i = 0; i < t.a.size(); i++) {</pre>
                a[i] += t.a[i];
            normalize();
            return *this;
        poly operator -= (const poly &t) {
            a.resize(max(a.size(), t.a.size()));
```

```
for(size_t i = 0; i < t.a.size(); i++) {</pre>
        a[i] -= t.a[i];
    normalize();
    return *this;
poly operator + (const poly &t) const {return
   polv(*this) += t;}
poly operator - (const poly &t) const {return
   poly(*this) -= t;}
poly mod_xk(size_t k) const { // get same
   polynomial mod x^k
   k = min(k, a.size());
    return vector <T > (begin(a), begin(a) + k);
poly mul_xk(size_t k) const { // multiply by x^k
   poly res(*this);
   res.a.insert(begin(res.a), k, 0);
    return res;
}
poly div_xk(size_t k) const { // divide by x^k,
   dropping coefficients
   k = min(k, a.size());
   return vector<T>(begin(a) + k, end(a));
poly substr(size_t l, size_t r) const { //
   return mod_xk(r).div_xk(l)
   1 = min(1, a.size());
   r = min(r, a.size());
    return vector<T>(begin(a) + 1, begin(a) + r);
poly inv(size_t n) const { // get inverse series
   mod x^n
    assert(!is_zero());
   poly ans = a[0].inv();
    size_t a = 1;
    while (a < n) {
        poly C = (ans * mod_xk(2 * a)).substr(a,
           2 * a):
        ans -= (ans * C).mod_xk(a).mul_xk(a);
        a *= 2:
```

```
}
    return ans.mod_xk(n);
}
poly operator *= (const poly &t) {fft::mul(a,
   t.a); normalize(); return *this;}
polv operator * (const polv &t) const {return
   poly(*this) *= t;}
poly reverse(size_t n, bool rev = 0) const { //
   reverses and leaves only n terms
    poly res(*this);
    if(rev) { // If rev = 1 then tail goes to
        res.a.resize(max(n. res.a.size())):
    std::reverse(res.a.begin(), res.a.end());
    return res.mod_xk(n);
}
pair < poly , poly > divmod_slow(const poly &b)
   const { // when divisor or quotient is small
    vector<T> A(a);
    vector<T> res:
    while(A.size() >= b.a.size()) {
        res.push_back(A.back() / b.a.back());
        if(res.back() != T(0)) {
            for(size_t i = 0; i < b.a.size();</pre>
               i++) {
                A[A.size() - i - 1] -=
                   res.back() * b.a[b.a.size() -
                   i - 1]:
            }
        A.pop_back();
    std::reverse(begin(res), end(res));
    return {res, A};
}
pair < poly , poly > divmod(const poly &b) const {
   // returns quotiend and remainder of a mod b
```

```
if(deg() < b.deg()) {</pre>
        return {poly{0}, *this};
    int d = deg() - b.deg();
    if(min(d, b.deg()) < magic) {</pre>
        return divmod_slow(b);
    poly D = (reverse(d + 1) * b.reverse(d +
       1).inv(d + 1)).mod xk(d + 1).reverse(d +
       1. 1):
    return {D, *this - D * b};
}
poly operator / (const poly &t) const {return
   divmod(t).first:}
poly operator % (const poly &t) const {return
   divmod(t).second;}
poly operator /= (const poly &t) {return *this =
   divmod(t).first;}
poly operator %= (const poly &t) {return *this =
   divmod(t).second;}
poly operator *= (const T &x) {
    for(auto &it: a) {
        it *= x;
    normalize();
    return *this;
poly operator /= (const T &x) {
    for(auto &it: a) {
        it /= x;
    normalize();
    return *this:
poly operator * (const T &x) const {return
   poly(*this) *= x;}
poly operator / (const T &x) const {return
   poly(*this) /= x;}
void print() const {
    for(auto it: a) {
```

```
cout << it << ' ':
    cout << endl;</pre>
T eval(T x) const { // evaluates in single point
    T res(0);
    for(int i = int(a.size()) - 1; i >= 0; i--) {
        res *= x;
        res += a[i];
    }
    return res;
}
T& lead() { // leading coefficient
    return a.back();
}
int deg() const { // degree
    return a.empty() ? -inf : a.size() - 1;
bool is_zero() const { // is polynomial zero
    return a.empty();
T operator [](int idx) const {
    return idx >= (int)a.size() || idx < 0 ?</pre>
       T(0) : a[idx];
}
T& coef(size_t idx) { // mutable reference at
   coefficient
   return a[idx];
bool operator == (const poly &t) const {return a
   == t.a:}
bool operator != (const poly &t) const {return a
   != t.a:}
poly deriv() { // calculate derivative
    vector <T> res;
    for(int i = 1; i <= deg(); i++) {</pre>
        res.push_back(T(i) * a[i]);
    }
```

```
return res;
poly integr() { // calculate integral with C = 0
    vector < T > res = {0};
   for(int i = 0; i <= deg(); i++) {</pre>
        res.push_back(a[i] / T(i + 1));
    return res;
size_t leading_xk() const { // Let p(x) = x^k *
   t(x), return k
   if(is_zero()) {
        return inf;
    int res = 0;
    while (a[res] == T(0)) {
        res++;
    return res;
poly log(size_t n) { // calculate log p(x) mod
   x^n
    assert(a[0] == T(1));
    return (deriv().mod_xk(n) *
       inv(n)).integr().mod_xk(n);
poly exp(size_t n) { // calculate exp p(x) mod
   x^n
   if(is_zero()) {
        return T(1);
    assert(a[0] == T(0));
    poly ans = T(1);
    size_t a = 1;
    while(a < n) {</pre>
        poly C = ans.log(2 * a).div_xk(a) -
           substr(a, 2 * a);
        ans -= (ans * C).mod_xk(a).mul_xk(a);
        a *= 2;
    return ans.mod_xk(n);
```

```
}
poly pow_slow(size_t k, size_t n) { // if k is
    return k ? k % 2 ? (*this * pow_slow(k - 1,
       n)).mod_xk(n) : (*this *
       *this).mod_xk(n).pow_slow(k / 2, n):
       T(1);
}
poly pow(size_t k, size_t n) { // calculate
   p^k(n) mod x^n
    if(is_zero()) {
        return *this;
    }
    if(k < magic) {</pre>
        return pow_slow(k, n);
    int i = leading_xk();
    T j = a[i];
    poly t = div_xk(i) / j;
    return bpow(j, k) * (t.log(n) *
       T(k) .exp(n).mul_xk(i * k).mod_xk(n);
poly mulx(T x) { // component-wise
   multiplication with x^k
    T cur = 1;
    polv res(*this);
    for(int i = 0; i <= deg(); i++) {
        res.coef(i) *= cur;
        cur *= x;
    }
    return res;
poly mulx_sq(T x) { // component-wise
   multiplication with x^{k^2}
    T cur = x;
    T \text{ total} = 1;
    T xx = x * x;
    poly res(*this);
    for(int i = 0; i <= deg(); i++) {</pre>
        res.coef(i) *= total;
        total *= cur;
        cur *= xx;
```

```
}
    return res;
vector<T> chirpz_even(T z, int n) { // P(1),
   P(z^2), P(z^4), ..., P(z^2(n-1))
   int m = deg();
    if(is_zero()) {
        return vector <T>(n, 0);
    }
    vector < T > vv(m + n);
    T zi = z.inv();
   T zz = zi * zi;
    T cur = zi;
    T \text{ total} = 1;
    for (int i = 0; i \le max(n - 1, m); i++) {
        if(i <= m) {vv[m - i] = total;}</pre>
        if(i < n) {vv[m + i] = total;}</pre>
        total *= cur;
        cur *= zz;
    poly w = (mulx_sq(z) * vv).substr(m, m +
       n).mulx_sq(z);
    vector <T> res(n);
    for(int i = 0; i < n; i++) {</pre>
        res[i] = w[i];
    return res;
vector<T> chirpz(T z, int n) { // P(1), P(z),
   P(z^2), ..., P(z^{(n-1)})
    auto even = chirpz_even(z, (n + 1) / 2);
    auto odd = mulx(z).chirpz_even(z, n / 2);
    vector <T> ans(n);
    for(int i = 0; i < n / 2; i++) {</pre>
        ans [2 * i] = even[i];
        ans[2 * i + 1] = odd[i];
    if(n % 2 == 1) {
        ans[n - 1] = even.back();
    return ans;
```

```
template < typename iter >
    vector<T> eval(vector<poly> &tree, int v,
       iter 1, iter r) { // auxiliary evaluation
       function
        if(r - 1 == 1) {
            return {eval(*1)};
        } else {
            auto m = 1 + (r - 1) / 2;
            auto A = (*this \% tree[2 *
               v]).eval(tree, 2 * v, 1, m);
            auto B = (*this \% tree[2 * v +
               1]).eval(tree, 2 * v + 1, m, r);
            A.insert(end(A), begin(B), end(B));
            return A;
vector<T> eval(vector<T> x) { // evaluate
   polynomial in (x1, ..., xn)
   int n = x.size();
    if(is_zero()) {
        return vector <T>(n, T(0));
    vector<poly> tree(4 * n);
    build(tree, 1, begin(x), end(x));
    return eval(tree, 1, begin(x), end(x));
template < typename iter >
    poly inter(vector<poly> &tree, int v, iter
       1, iter r, iter ly, iter ry) { //
       auxiliary interpolation function
        if(r - 1 == 1) {
            return {*ly / a[0]};
        } else {
            auto m = 1 + (r - 1) / 2:
            auto my = ly + (ry - ly) / 2;
            auto A = (*this \% tree[2 *
               v]).inter(tree, 2 * v, 1, m, ly,
               my);
            auto B = (*this \% tree[2 * v +
               1]).inter(tree, 2 * v + 1, m, r,
               mv, rv);
            return A * tree[2 * v + 1] + B *
```

```
tree[2 * v]:
                }
   };
template < typename T>
    poly<T> operator * (const T& a, const poly<T>& b) {
        return b * a;
template < typename T>
    poly<T> xk(int k) { // return x^k
        return poly<T>{1}.mul_xk(k);
   }
template < typename T>
    T resultant(poly<T> a, poly<T> b) { // computes
       resultant of a and b
        if(b.is zero()) {
            return 0;
        } else if(b.deg() == 0) {
            return bpow(b.lead(), a.deg());
        } else {
            int pw = a.deg();
            a \%= b;
            pw -= a.deg();
            T \text{ mul} = bpow(b.lead(), pw) * T((b.deg() &
                a.deg() & 1) ? -1 : 1);
            T ans = resultant(b, a);
            return ans * mul;
        }
template < typename iter >
    poly<typename iter::value_type> kmul(iter L, iter R)
       \{ // \text{ computes } (x-a1)(x-a2)...(x-an) \text{ without } 
       building tree
        if(R - L == 1) {
            return vector < typename
                iter::value_type>{-*L, 1};
        } else {
            iter M = L + (R - L) / 2;
            return kmul(L, M) * kmul(M, R);
```

```
}
    template < typename T, typename iter >
        poly<T> build(vector<poly<T>> &res, int v, iter L,
           iter R) { // builds evaluation tree for
           (x-a1)(x-a2)...(x-an)
            if(R - L == 1) {
                return res[v] = vector<T>{-*L, 1};
                iter M = L + (R - L) / 2;
                return res[v] = build(res, 2 * v, L, M) *
                    build(res, 2 * v + 1, M, R);
            }
        }
    template < typename T>
        poly<T> inter(vector<T> x, vector<T> y) { //
           interpolates minimum polynomial from (xi, yi)
           pairs
            int n = x.size();
            vector<poly<T>> tree(4 * n);
            return build(tree, 1, begin(x),
               end(x)).deriv().inter(tree, 1, begin(x),
               end(x), begin(y), end(y));
        }
};
using namespace algebra;
const int mod = 1e9 + 7;
typedef modular < mod > base;
typedef poly<base> polyn;
using namespace algebra;
signed main() {
    ios::sync_with_stdio(0);
    cin.tie(0);
    int n = 100000;
    polyn a;
    vector < base > x;
    for(int i = 0; i <= n; i++) {</pre>
        a.a.push_back(1 + rand() % 100);
        x.push_back(1 + rand() \% (2 * n));
```

```
    sort(begin(x), end(x));
    x.erase(unique(begin(x), end(x)), end(x));
    auto b = a.eval(x);
    cout << clock() / double(CLOCKS_PER_SEC) << endl;
    auto c = inter(x, b);
    polyn md = kmul(begin(x), end(x));
    cout << clock() / double(CLOCKS_PER_SEC) << endl;
    assert(c == a % md);
    return 0;
}
</pre>
```

# 3.4 Primitivas de matriz - exponenciacao

```
#define MODULAR false
template < typename T> struct matrix : vector < T>> {
    int n, m;
    void print() {
        for (int i = 0; i < n; i++) {
             for (int j = 0; j < m; j++) cout <<</pre>
                (*this)[i][j] << " ";
             cout << endl;</pre>
        }
    }
    matrix(int n_, int m_, bool ident = false) :
             vector < vector < T >> (n_, vector < T > (m_, 0)), n(n_),
                m(m) {
        if (ident) {
             assert(n == m);
             for (int i = 0; i < n; i++) (*this)[i][i] = 1;</pre>
    }
    matrix(const vector < vector < T >> & c) :
       vector < vector < T >> (c),
        n(c.size()), m(c[0].size()) {}
    matrix<T> operator*(matrix<T>& r) {
        assert(m == r.n);
        matrix<T> M(n, r.m);
```

```
for (int i = 0; i < n; i++) for (int k = 0; k < m;
           k++)
            for (int j = 0; j < r.m; j++) {
                T \text{ add} = (*this)[i][k] * r[k][j];
#if MODULAR
                M[i][j] += add%MOD;
                if (M[i][j] >= MOD) M[i][j] -= MOD;
#else
                M[i][j] += add;
#endif
            }
        return M;
    }
    matrix<T> operator^(ll e){
        matrix <T> M(n, n, true), at = *this;
        while (e) {
            if (e\&1) M = M*at;
           e >>= 1;
            at = at*at;
        }
        return M;
    void apply_transform(matrix M, ll e){
        auto& v = *this;
        while (e) {
            if (e\&1) v = M*v;
            e >>= 1;
            M = M * M;
        }
    }
};
3.5 Primitivas Geometricas
typedef double ld;
const ld DINF = 1e18;
```

```
typedef double ld;
const ld DINF = 1e18;
const ld pi = acos(-1.0);
const ld eps = 1e-9;
#define sq(x) ((x)*(x))
```

```
bool eq(ld a, ld b) {
    return abs(a - b) <= eps;</pre>
struct pt { // ponto
    ld x, y;
    pt() {}
    pt(ld x_{-}, ld y_{-}) : x(x_{-}), y(y_{-}) \{ \}
    bool operator < (const pt p) const {</pre>
        if (!eq(x, p.x)) return x < p.x;
        if (!eq(y, p.y)) return y < p.y;
        return 0:
    }
    bool operator == (const pt p) const {
        return eq(x, p.x) and eq(y, p.y);
    }
    pt operator + (const pt p) const { return pt(x+p.x,
       y+p.y); }
    pt operator - (const pt p) const { return pt(x-p.x,
       y-p.y); }
    pt operator * (const ld c) const { return pt(x*c , y*c
       ); }
    pt operator / (const ld c) const { return pt(x/c , y/c
       ); }
    ld operator * (const pt p) const { return x*p.x + y*p.y;
    ld operator ^ (const pt p) const { return x*p.y - y*p.x;
    friend istream& operator >> (istream& in, pt& p) {
        return in >> p.x >> p.y;
    }
};
struct line { // reta
    pt p, q;
    line() {}
    line(pt p_, pt q_) : p(p_), q(q_) {}
    friend istream& operator >> (istream& in, line& r) {
        return in >> r.p >> r.q;
    }
};
```

```
// PONTO & VETOR
ld dist(pt p, pt q) { // distancia
    return hypot(p.y - q.y, p.x - q.x);
}
ld dist2(pt p, pt q) { // quadrado da distancia
    return sq(p.x - q.x) + sq(p.y - q.y);
}
ld norm(pt v) { // norma do vetor
    return dist(pt(0, 0), v);
}
ld angle(pt v) { // angulo do vetor com o eixo x
    ld ang = atan2(v.y, v.x);
    if (ang < 0) ang += 2*pi;</pre>
    return ang;
}
ld sarea(pt p, pt q, pt r) { // area com sinal
    return ((q-p)^(r-q))/2;
}
bool col(pt p, pt q, pt r) { // se p, q e r sao colin.
    return eq(sarea(p, q, r), 0);
}
int paral(pt u, pt v) { // se u e v sao paralelos
    if (!eq(u^v, 0)) return 0;
    if ((u.x > eps) == (v.x > eps) and (u.y > eps) == (v.y >
       eps))
        return 1:
    return -1:
}
bool ccw(pt p, pt q, pt r) { // se p, q, r sao ccw
    return sarea(p, q, r) > eps;
}
pt rotate(pt p, ld th) { // rotaciona o ponto th radianos
```

```
return pt(p.x * cos(th) - p.y * sin(th),
            p.x * sin(th) + p.y * cos(th));
}
pt rotate90(pt p) { // rotaciona 90 graus
   return pt(-p.y, p.x);
// RETA
bool isvert(line r) { // se r eh vertical
    return eq(r.p.x, r.q.x);
}
ld getm(line r) { // coef. ang. de r
   if (isvert(r)) return DINF;
   return (r.p.y - r.q.y) / (r.p.x - r.q.x);
}
ld getn(line r) { // coef. lin. de r
    if (isvert(r)) return DINF;
    return r.p.y - getm(r) * r.p.x;
}
bool paraline(line r, line s) { // se r e s sao paralelas
    return paral(r.p - r.q, s.p - s.q);
}
bool isinseg(pt p, line r) { // se p pertence ao seg de r
    if (p == r.p or p == r.q) return 1;
   return paral(p - r.p, p - r.q) == -1;
}
ld get_t(pt v, line r) { // retorna t tal que t*v pertence a
   reta r
    return (r.p^r.q) / ((r.p-r.q)^v);
}
pt proj(pt p, line r) { // projecao do ponto p na reta r
   if (r.p == r.q) return r.p;
   r.q = r.q - r.p; p = p - r.p;
    pt proj = r.q * ((p*r.q) / (r.q*r.q));
```

```
return proj + r.p;
}
pt inter(line r, line s) { // r inter s
    if (paraline(r, s)) return pt(DINF, DINF);
    if (isvert(r)) return pt(r.p.x, getm(s) * r.p.x +
       getn(s));
    if (isvert(s)) return pt(s.p.x, getm(r) * s.p.x +
       getn(r));
    1d x = (getn(s) - getn(r)) / (getm(r) - getm(s));
    return pt(x, getm(r) * x + getn(r));
}
bool interseg(line r, line s) { // se o seg de r intersecta
   o seg de s
    if (isinseg(r.p, s) or isinseg(r.q, s)
        or isinseg(s.p, r) or isinseg(s.q, r)) return 1;
    return ccw(r.p, r.q, s.p) != ccw(r.p, r.q, s.q) and
            ccw(s.p, s.q, r.p) != ccw(s.p, s.q, r.q);
}
ld disttoline(pt p, line r) { // distancia do ponto a reta
    return 2 * abs(sarea(p, r.p, r.q)) / dist(r.p, r.q);
}
ld disttoseg(pt p, line r) { // distancia do ponto ao seg
    if ((r.q - r.p)*(p - r.p) < 0) return dist(r.p, p);
    if ((r.p - r.q)*(p - r.q) < 0) return dist(r.q, p);
    return disttoline(p, r);
}
ld distseg(line a, line b) { // distancia entre seg
    if (interseg(a, b)) return 0;
    ld ret = DINF;
    ret = min(ret, disttoseg(a.p, b));
    ret = min(ret, disttoseg(a.q, b));
    ret = min(ret, disttoseg(b.p, a));
    ret = min(ret, disttoseg(b.q, a));
```

```
return ret;
}
// POLIGONO
// distancia entre os retangulos a e b (lados paralelos aos
// assume que ta representado (inferior esquerdo, superior
   direito)
ld dist_rect(pair<pt, pt> a, pair<pt, pt> b) {
    1d hor = 0, vert = 0;
    if (a.s.x < b.f.x) hor = b.f.x - a.s.x;
    else if (b.s.x < a.f.x) hor = a.f.x - b.s.x;
    if (a.s.y < b.f.y) vert = b.f.y - a.s.y;
    else if (b.s.y < a.f.y) vert = a.f.y - b.s.y;</pre>
    return dist(pt(0, 0), pt(hor, vert));
}
ld polarea(vector<pt> v) { // area do poligono
    ld ret = 0;
    for (int i = 0; i < v.size(); i++)</pre>
        ret += sarea(pt(0, 0), v[i], v[(i + 1) % v.size()]);
    return abs(ret);
}
// se o ponto ta dentro do poligono: retorna O se ta fora,
// 1 se ta no interior e 2 se ta na borda
int inpol(vector < pt > & v, pt p) { // O(n)
    int qt = 0;
    for (int i = 0; i < v.size(); i++) {</pre>
        if (v[i] == p) return 2;
        int j = (i+1)%v.size();
        bool igual = eq(v[i].y, p.y) and eq(v[j].y, p.y),
           baixo = v[i].y+eps < p.y;</pre>
        if (!igual and baixo == (v[j].y+eps < p.y)) continue;</pre>
        auto t = (p-v[i])^(v[j]-v[i]);
        if (eq(t, 0)) return 2;
        if (!igual and baixo == (t > eps)) qt += baixo ? 1 :
            -1;
    }
    return qt != 0;
```

```
}
bool interpol(vector<pt> v1, vector<pt> v2) { // se dois
   poligonos se intersectam - O(n*m)
    int n = v1.size(), m = v2.size();
    for (int i = 0; i < n; i++) if (inpol(v2, v1[i])) return</pre>
    for (int i = 0; i < n; i++) if (inpol(v1, v2[i])) return</pre>
       1:
    for (int i = 0; i < n; i++) for (int j = 0; j < m; j++)
        if (interseg(line(v1[i], v1[(i+1)%n]), line(v2[j],
           v2[(j+1)%m]))) return 1;
    return 0:
}
ld distpol(vector<pt> v1, vector<pt> v2) { // distancia
   entre poligonos
    if (interpol(v1, v2)) return 0;
    ld ret = DINF;
    for (int i = 0; i < v1.size(); i++) for (int j = 0; j <</pre>
       v2.size(); j++)
        ret = min(ret, distseg(line(v1[i], v1[(i + 1) %
           v1.size()]),
                     line(v2[j], v2[(j + 1) % v2.size()])));
    return ret;
}
vector<pt> convex_hull(vector<pt> v) { // convex hull - O(n
   log(n))
    if (v.size() <= 1) return v;</pre>
    vector <pt> 1, u;
    sort(v.begin(), v.end());
    for (int i = 0; i < v.size(); i++) {</pre>
        while (1.size() > 1 \text{ and } !ccw(l[1.size()-2],
           1.back(), v[i]))
            1.pop_back();
        1.pb(v[i]);
    for (int i = v.size() - 1; i >= 0; i--) {
        while (u.size() > 1 and !ccw(u[u.size()-2],
```

```
u.back(), v[i]))
            u.pop_back();
        u.pb(v[i]);
    }
    1.pop_back(); u.pop_back();
   for (pt i : u) l.pb(i);
    return 1;
}
struct convex_pol {
    vector < pt > pol;
    convex_pol(vector<pt> v) : pol(convex_hull(v)) {}
    bool is_inside(pt p) { // se o ponto ta dentro do hull -
       O(log(n))
        if (pol.size() == 1) return p == pol[0];
        int 1 = 1, r = pol.size();
        while (1 < r) {
            int m = (1+r)/2;
            if (ccw(p, pol[0], pol[m])) l = m+1;
            else r = m;
        if (1 == 1) return isinseg(p, line(pol[0], pol[1]));
        if (1 == pol.size()) return false;
        return !ccw(p, pol[1], pol[1-1]);
   }
};
// os segmentos precisam ser ter o p < q</pre>
bool operator < (const line& a, const line& b) { //</pre>
   comparador pro sweepline
   if (a.p == b.p) return ccw(a.p, a.q, b.q);
   if (!eq(a.p.x, a.q.x) and (eq(b.p.x, b.q.x) or a.p.x+eps
       < b.p.x))
        return ccw(a.p, a.q, b.p);
    return ccw(a.p, b.q, b.p);
}
// CIRCUNFERENCIA
pt getcenter(pt a, pt b, pt c) { // centro da circunf dado 3
   pontos
```

```
b = (a + b) / 2:
    c = (a + c) / 2;
    return inter(line(b, b + rotate90(a - b)),
            line(c, c + rotate90(a - c)));
}
vector<pt> circ_line_inter(pt a, pt b, pt c, ld r) { //
   intersecao da circunf (c, r) e reta ab
    vector<pt> ret;
    b = b-a, a = a-c;
    1d A = b*b;
    1d B = a*b:
    1d C = a*a - r*r;
    1d D = B*B - A*C;
    if (D < -eps) return ret;</pre>
    ret.push_back(c+a+b*(-B+sqrt(D+eps))/A);
    if (D > eps) ret.push_back(c+a+b*(-B-sqrt(D))/A);
    return ret;
}
vector<pt> circ_inter(pt a, pt b, ld r, ld R) { //
   intersecao da circunf (a, r) e (b, R)
    vector<pt> ret;
    1d d = dist(a, b);
    if (d > r+R or d+min(r, R) < max(r, R)) return ret;</pre>
    1d x = (d*d-R*R+r*r)/(2*d);
    1d v = sqrt(r*r-x*x);
    pt v = (b-a)/d;
    ret.push_back(a+v*x + rotate90(v)*y);
    if (y > 0) ret.push_back(a+v*x - rotate90(v)*y);
    return ret;
}
// comparador pro set para fazer sweep angle com segmentos
double ang;
struct cmp {
    bool operator () (const line& a, const line& b) const {
        line r = line(pt(0, 0), rotate(pt(1, 0), ang));
        return norm(inter(r, a)) < norm(inter(r, b));</pre>
   }
};
```

#### 3.6 Primitivas Geometricas Inteiras

```
#define sq(x) ((x)*(11)(x))
struct pt { // ponto
    int x, y;
    pt() {}
    pt(int x_{-}, int y_{-}) : x(x_{-}), y(y_{-}) \{ \}
    bool operator < (const pt p) const {</pre>
        if (x != p.x) return x < p.x;
        return y < p.y;</pre>
    }
    bool operator == (const pt p) const {
        return x == p.x and y == p.y;
    }
    pt operator + (const pt p) const { return pt(x+p.x,
       (v.q+v); }
    pt operator - (const pt p) const { return pt(x-p.x,
       y-p.y); }
    pt operator * (const int c) const { return pt(x*c, y*c);
    11 operator * (const pt p) const { return x*(11)p.x +
       v*(ll)p.v; }
    ll operator ^ (const pt p) const { return x*(ll)p.y -
       v*(ll)p.x; }
    friend istream& operator >> (istream& in, pt& p) {
        return in >> p.x >> p.y;
    }
};
struct line { // reta
    pt p, q;
    line() {}
    line(pt p_, pt q_) : p(p_), q(q_) {}
    friend istream& operator >> (istream& in, line& r) {
        return in >> r.p >> r.q;
    }
};
// PONTO & VETOR
11 dist2(pt p, pt q) { // quadrado da distancia
```

```
return sq(p.x - q.x) + sq(p.y - q.y);
}
ll sarea2(pt p, pt q, pt r) { // 2 * area com sinal
    return (q-p)^(r-q);
}
bool col(pt p, pt q, pt r) { // se p, q e r sao colin.
    return sarea2(p, q, r) == 0;
}
int paral(pt u, pt v) { // se u e v sao paralelos
    if (u^v) return 0;
    if ((u.x > 0) == (v.x > 0) and (u.y > 0) == (v.y > 0))
       return 1:
    return -1:
}
bool ccw(pt p, pt q, pt r) { // se p, q, r sao ccw
    return sarea2(p, q, r) > 0;
}
int quad(pt p) { // quadrante de um ponto
  return (p.x<0)^3*(p.y<0);
}
bool compare_angle(pt p, pt q) { // retorna se ang(p) <</pre>
   ang(q)
    if (quad(p) != quad(q)) return quad(p) < quad(q);</pre>
    return ccw(q, pt(0, 0), p);
}
pt rotate90(pt p) { // rotaciona 90 graus
    return pt(-p.y, p.x);
}
// RETA
bool paraline(line r, line s) { // se r e s sao paralelas
    return paral(r.p - r.q, s.p - s.q);
}
```

```
bool isinseg(pt p, line r) { // se p pertence ao seg de r
   if (p == r.p or p == r.q) return 1;
   return paral(p - r.p, p - r.q) == -1;
}
bool interseg(line r, line s) { // se o seg de r intersecta
   o seg de s
   if (isinseg(r.p, s) or isinseg(r.q, s)
        or isinseg(s.p, r) or isinseg(s.q, r)) return 1;
   return ccw(r.p, r.q, s.p) != ccw(r.p, r.q, s.q) and
            ccw(s.p, s.q, r.p) != ccw(s.p, s.q, r.q);
}
int segpoints(line r) { // numero de pontos inteiros no
    return 1 + \_gcd(abs(r.p.x - r.q.x), abs(r.p.y - r.q.y));
}
double get_t(pt v, line r) { // retorna t tal que t*v
   pertence a reta r
   return (r.p^r.q) / (double) ((r.p-r.q)^v);
}
// POLIGONO
// quadrado da distancia entre os retangulos a e b (lados
   paralelos aos eixos)
// assume que ta representado (inferior esquerdo, superior
   direito)
11 dist2_rect(pair<pt, pt> a, pair<pt, pt> b) {
   int hor = 0, vert = 0;
   if (a.s.x < b.f.x) hor = b.f.x - a.s.x;
    else if (b.s.x < a.f.x) hor = a.f.x - b.s.x:
   if (a.s.y < b.f.y) vert = b.f.y - a.s.y;</pre>
    else if (b.s.y < a.f.y) vert = a.f.y - b.s.y;</pre>
    return sq(hor) + sq(vert);
}
11 polarea2(vector<pt> v) { // 2 * area do poligono
   ll ret = 0;
   for (int i = 0; i < v.size(); i++)</pre>
```

```
ret += sarea2(pt(0, 0), v[i], v[(i + 1) % v.size()]);
    return abs(ret);
}
// se o ponto ta dentro do poligono: retorna O se ta fora,
// 1 se ta no interior e 2 se ta na borda
int inpol(vector<pt>& v, pt p) { // O(n)
    int qt = 0;
    for (int i = 0; i < v.size(); i++) {</pre>
        if (v[i] == p) return 2;
        int j = (i+1)%v.size();
        bool igual = v[i].y == p.y and v[j].y == p.y, baixo
           = v[i].y < p.y;
        if (!igual and baixo == (v[j].y < p.y)) continue;</pre>
        auto t = (p-v[i])^(v[j]-v[i]);
        if (!t) return 2;
        if (!igual and baixo == (t > 0)) qt += baixo ? 1 :
            -1;
    return qt != 0;
}
vector<pt> convex_hull(vector<pt> v) { // convex hull - O(n
   log(n))
    if (v.size() <= 1) return v;</pre>
    vector <pt> 1, u;
    sort(v.begin(), v.end());
    for (int i = 0; i < v.size(); i++) {</pre>
        while (1.size() > 1 and !ccw(l[1.size()-2],
           1.back(), v[i]))
            l.pop_back();
        1.pb(v[i]);
    }
    for (int i = v.size() - 1; i >= 0; i--) {
        while (u.size() > 1 and !ccw(u[u.size()-2],
           u.back(), v[i]))
            u.pop_back();
        u.pb(v[i]);
    }
    1.pop_back(); u.pop_back();
    for (pt i : u) 1.pb(i);
    return 1:
```

```
11 interior_points(vector<pt> v) { // pontos inteiros dentro
   de um poligono simples
   11 b = 0;
   for (int i = 0; i < v.size(); i++)</pre>
        b += segpoints(line(v[i], v[(i+1)\%v.size()])) - 1;
    return (polarea2(v) - b) / 2 + 1;
}
struct convex_pol {
    vector<pt> pol;
    convex_pol(vector<pt> v) : pol(convex_hull(v)) {}
    bool is_inside(pt p) { // se o ponto ta dentro do hull -
       O(\log(n))
        if (pol.size() == 1) return p == pol[0];
        int 1 = 1, r = pol.size();
        while (1 < r) {
            int m = (1+r)/2;
            if (ccw(p, pol[0], pol[m])) l = m+1;
            else r = m;
        if (l == 1) return isinseg(p, line(pol[0], pol[1]));
        if (1 == pol.size()) return false;
        return !ccw(p, pol[1], pol[1-1]);
   }
};
// os segmentos precisam ser ter o p < q</pre>
bool operator < (const line& a, const line& b) { //</pre>
   comparador pro sweepline
   if (a.p == b.p) return ccw(a.p, a.q, b.q);
   if (a.p.x != a.q.x and (b.p.x == b.q.x or a.p.x < b.p.x))
        return ccw(a.p, a.q, b.p);
    return ccw(a.p, b.q, b.p);
}
// comparador pro set pra fazer sweep angle com segmentos
pt dir;
struct cmp {
```

```
bool operator () (const line& a, const line& b) const {
        return get_t(dir, a) < get_t(dir, b);</pre>
};
```

#### Primitivas Geometricas 3D

```
typedef double ld;
const ld DINF = 1e18:
const ld pi = acos(-1.0);
const ld eps = 1e-9;
#define sq(x)((x)*(x))
bool eq(ld a, ld b) {
    return abs(a - b) <= eps;</pre>
}
struct pt { // ponto
    ld x, y, z;
    pt() {}
    pt(ld x_{-}, ld y_{-}, ld z_{-}) : x(x_{-}), y(y_{-}), z(z_{-}) {}
    bool operator < (const pt p) const {</pre>
        if (!eq(x, p.x)) return x < p.x;
        if (!eq(y, p.y)) return y < p.y;</pre>
        if (!eq(z, p.z)) return z < p.z;
        return 0;
    bool operator == (const pt p) const {
        return eq(x, p.x) and eq(y, p.y) and eq(z, p.z);
    pt operator + (const pt p) const { return pt(x+p.x,
       y+p.y, z+p.z); }
    pt operator - (const pt p) const { return pt(x-p.x,
       y-p.y, z-p.z); }
    pt operator * (const ld c) const { return pt(x*c , y*c
       , z*c ); }
    pt operator / (const ld c) const { return pt(x/c , y/c
       , z/c ); }
    ld operator * (const pt p) const { return x*p.x + y*p.y
```

```
+ z*p.z; }
    void rotate_x(ld a) {
        1d ny = y*cos(a) - z*sin(a);
        1d nz = y*sin(a) + z*cos(a);
        y = ny;
        z = nz;
    }
    void rotate_y(ld a) {
        1d nx = x*cos(a) + z*sin(a);
        1d nz = -x*sin(a) + z*cos(a);
        x = nx;
        z = nz;
    }
    void rotate_z(ld a) {
        1d nx = x*cos(a) - y*sin(a);
        ld ny = x*sin(a) + y*cos(a);
        x = nx;
        y = ny;
    }
};
// converte de coordenadas polares para cartesianas
// (angulos devem estar em radianos)
pt convert(ld rho, ld th, ld phi) {
    return pt(sin(phi) * cos(th), sin(phi) * sin(th),
       cos(phi)) * rho;
}
// distancia
ld dist(pt a, pt b) {
    return sqrt(sq(a.x-b.x) + sq(a.y-b.y) + sq(a.z-b.z));
}
```

#### Primitivas de matriz

```
11 mod(ll v) { return (v + MOD) % MOD; }
11 sum(11 1, 11 r) { return mod(1+r); }
11 mult(ll 1, ll r) { return mod(l*r); }
11 inverse(ll 1){ return inv(1, MOD); }
```

```
bool equal(ll\ l,\ ll\ r) { return mod(l-r) == 0; }
template < typename T> struct matrix {
    vector < vector < T >> in;
    int row, col;
    void print(){//
        for (int i = 0; i < row; i++){</pre>
             for (int j = 0; j < col; j++)</pre>
                 cout << in[i][j] << " ";</pre>
             cout << endl:</pre>
        }
    }
    matrix(int row, int col, int op = 0):row(row), col(col),
       in(row, vector<T>(col, 0)){
        if (op) for (int i = 0; i < row; i++) in[i][i] = 1;</pre>
    matrix(initializer_list<initializer_list<T>> c):
        row(c.size()), col((*c.begin()).size()){
             in = vector < vector < T >> (row, vector < T > (col, 0));
            int i, j;
            i = 0;
            for (auto &it : c){
                 i = 0;
                 for (auto &jt : it){
                     in[i][j] = jt;
                     j++;
                 }
                 i++;
            }
    T &operator()(int i, int j){ return in[i][j]; }
    //in case of a transposed matrix, swap i and j
    matrix<T>& operator*=(T t){
        matrix < T > &l = *this;
        for (int i = 0; i < row; i++)</pre>
            for (int j = 0; j < col; j++)</pre>
                 l(i, j) = mult(l(i, j), t); //% MOD) % MOD;
        return 1;
    matrix<T> operator+(matrix<T> &r){
```

```
matrix<T> &1 = *this:
    matrix<T> m(row, col, 0);
    for (int i = 0; i < row; i++)</pre>
        for (int j = 0; j < col; j++)
             m(i, j) = sum(l(i, j), r(i, j)); //% MOD) %
                MOD;
    return m;
}
matrix <T> operator*(matrix <T> &r){
    matrix <T> &1 = *this;
    int row = l.row;
    int col = r.col;
    int K = 1.col;
    matrix<T> m(row, col, 0);
    for (int i = 0; i < row; i++)</pre>
        for (int j = 0; j < col; j++)</pre>
             for (int k = 0; k < K; k++)
                 m(i, j) = sum(m(i, j), mult(l(i, k),
                    r(k, j)));
    return m;
}
matrix<T> operator^(long long e){
    matrix < T > &m = (*this);
    if (e == 0) return matrix(m.row, m.row, 1);
    if (e == 1) return m;
    if (e == 2) return m*m;
    auto m_{-} = m^{(e/2)}; m_{-} = m_{*}m_{-};
    if (e\%2 == 1) m_{-} = m_{-} * m;
    return m_;
}
void multiply_r(int i, T k){
    matrix < T > &m = (*this);
    for (int j = 0; j < col; j++)</pre>
        m(i, j) = mult(m(i, j), k);
void multiply_c(int j, T k){
    matrix < T > &m = (*this);
    for (int i = 0; i < row; j++)</pre>
        m(i, j) = mult(m(i, j), k);
void sum_r(int i1, int i2, T k){
    matrix < T > &m = (*this);
```

```
for (int j = 0; j < col; j++)</pre>
             m(i1, j) = sum(m(i1, j), mult(k, m(i2, j)));
    }
    bool gaussian(int I, int J){
        matrix < T > &m = (*this);
        T \text{ tmp} = m(I, J);
        if (equal(tmp, 0)) return false;
        multiply_r(I, inverse(tmp));
        for (int i = 0; i < row; i++)</pre>
             if (i != I) sum_r(i, I, mult(-1, m(i, J)));
        multiply_r(I, tmp);
        return true;
    T determinant(){
        matrix < T > m = (*this);
        for (int i = 0; i < row; i++)</pre>
             if (!m.gaussian(i, i)) return 0;
        T ans = 1;
        for (int i = 0; i < row; i++)</pre>
             ans = mult(ans, m(i, i));
        return ans;
};
```

### 4 Estruturas

# 4.1 BIT com update em range

```
// Operacoes 0-based
// query(l, r) retorna a soma de v[l..r]
// update(l, r, x) soma x em v[l..r]
//
// Complexidades:
// build - O(n)
// query - O(log(n))
// update - O(log(n))
```

```
namespace bit {
    11 bit[2][MAX+2];
    int n;
    void build(int n2, int* v) {
        n = n2;
        for (int i = 1; i <= n; i++)
            bit [1] [min(n+1, i+(i\&-i))] += bit [1][i] +=
                v[i-1];
    }
    ll get(int x, int i) {
        11 \text{ ret} = 0;
        for (; i; i -= i&-i) ret += bit[x][i];
        return ret;
    }
    void add(int x, int i, ll val) {
        for (; i <= n; i += i&-i) bit[x][i] += val;</pre>
    }
    11 get2(int p) {
        return get(0, p) * p + get(1, p);
    }
    11 query(int 1, int r) {
        return get2(r+1) - get2(1);
    void update(int 1, int r, 11 x) {
        add(0, 1+1, x), add(0, r+2, -x);
        add(1, 1+1, -x*1), add(1, r+2, x*(r+1));
    }
};
```

### 4.2 Sparse Table

```
// Resolve RMQ
// MAX2 = log(MAX)
//
// Complexidades:
// build - O(n log(n))
// query - O(1)
namespace sparse {
```

## 4.3 Min queue - stack

```
// Tudo O(1) amortizado
template < class T> struct minstack {
    stack<pair<T, T> > s;
    void push(T x) {
        if (!s.size()) s.push({x, x});
        else s.push({x, std::min(s.top().second, x)});
    }
    T top() { return s.top().first; }
    T pop() {
        T ans = s.top().first;
        s.pop();
        return ans;
   }
    T size() { return s.size(); }
    T min() { return s.top().second; }
};
template < class T> struct minqueue {
    minstack <T> s1, s2;
    void push(T x) { s1.push(x); }
    void move() {
```

```
if (s2.size()) return;
while (s1.size()) {
    T x = s1.pop();
    s2.push(x);
}

T front() { return move(), s2.top(); }
T pop() { return move(), s2.pop(); }
T size() { return s1.size()+s2.size(); }
T min() {
    if (!s1.size()) return s2.min();
    else if (!s2.size()) return s1.min();
    return std::min(s1.min(), s2.min());
}
};
```

### 4.4 Min queue - deque

```
// Tudo 0(1) amortizado

template < class T > struct minqueue {
    deque < pair < T, int > > q;

    void push(T x) {
        int ct = 1;
        while (q.size() and x < q.front().f)
            ct += q.front().s, q.pop_front();
        q.push_front({x, ct});
    }

    void pop() {
        if (q.back().s > 1) q.back().s--;
        else q.pop_back();
    }
    T min() { return q.back().f; }
};
```

### 4.5 Splay Tree

```
// SEMPRE QUE DESCER NA ARVORE, DAR SPLAY NO
// NODE MAIS PROFUNDO VISITADO
// Todas as operacoes sao O(log(n)) amortizado
// Se quiser colocar mais informacao no node,
// mudar em 'update'
template < typename T > struct splaytree {
    struct node {
        node *ch[2], *p;
        int sz;
        T val;
        node(T v) {
            ch[0] = ch[1] = p = NULL;
            sz = 1;
            val = v;
        void update() {
            sz = 1;
            for (int i = 0; i < 2; i++) if (ch[i]) {
                sz += ch[i]->sz;
            }
        }
    };
    node* root;
    splaytree() { root = NULL; }
    splaytree(const splaytree& t) {
        throw logic_error("Nao copiar a splaytree!");
    \simsplaytree() {
        vector < node *> q = {root};
        while (q.size()) {
            node* x = q.back(); q.pop_back();
            if (!x) continue;
            q.push_back(x->ch[0]), q.push_back(x->ch[1]);
            delete x;
        }
    }
    void rotate(node* x) { // x vai ficar em cima
        node *p = x->p, *pp = p->p;
```

```
if (pp) pp -> ch[pp -> ch[1] == p] = x;
    bool d = p -> ch[0] == x;
    p - ch[!d] = x - ch[d], x - ch[d] = p;
    if (p->ch[!d]) p->ch[!d]->p = p;
    x->p = pp, p->p = x;
    p->update(), x->update();
node* splay(node* x) {
    if (!x) return x;
    root = x;
    while (x->p) {
        node *p = x->p, *pp = p->p;
        if (!pp) return rotate(x), x; // zig
        if ((pp->ch[0] == p)^(p->ch[0] == x))
            rotate(x), rotate(x); // zigzag
        else rotate(p), rotate(x); // zigzig
    }
    return x;
node* insert(T v, bool lb=0) {
    if (!root) return lb ? NULL : root = new node(v);
    node *x = root, *last = NULL;;
    while (1) {
        bool d = x -> val < v;
        if (!d) last = x;
        if (x->val == v) break;
        if (x->ch[d]) x = x->ch[d];
        else {
            if (lb) break;
            x \rightarrow ch[d] = new node(v);
            x - ch[d] - p = x;
            x = x -  ch[d];
            break;
        }
    }
    splay(x);
    return lb ? splay(last) : x;
}
int size() { return root ? root->sz : 0; }
int count(T v) { return insert(v, 1) and root->val == v;
node* lower_bound(T v) { return insert(v, 1); }
```

```
void erase(T v) {
        if (!count(v)) return;
        node *x = root, *1 = x -> ch[0];
        if (!1) {
            root = x -> ch[1];
            if (root) root->p = NULL;
            return delete x;
        }
        root = 1, 1->p = NULL;
        while (1->ch[1]) 1 = 1->ch[1];
        splay(1);
        1 - ch[1] = x - ch[1];
        if (1->ch[1]) 1->ch[1]->p = 1;
        delete x;
        1->update();
    int order_of_key(T v) {
        if (!lower_bound(v)) return root ? root->sz : 0;
        return root->ch[0] ? root->ch[0]->sz : 0;
    node* find_by_order(int k) {
        node* x = root;
        while (1) {
            if (x->ch[0] \text{ and } x->ch[0]->sz >= k+1) x =
               x - > ch[0];
            else {
                if (x->ch[0]) k -= x->ch[0]->sz;
                if (!k) return splay(x);
                if (!x->ch[1]) {
                     splay(x);
                    return NULL;
                }
                k--, x = x->ch[1];
            }
        }
    T min() {
        node* x = root;
        while (x->ch[0]) x = x->ch[0]; // max -> ch[1]
        return splay(x)->val;
    }
};
```

# 4.6 SQRT-decomposition

```
// Resolve RMQ
// 0-indexed
// MAX2 = sqrt(MAX)
// O bloco da posicao x eh
// sempre x/q
//
// Complexidades:
// build - O(n)
// query - 0(sqrt(n))
int n, q;
int v[MAX];
int bl[MAX2];
void build() {
    q = (int) sqrt(n);
    // computa cada bloco
    for (int i = 0; i <= q; i++) {
        bl[i] = INF;
        for (int j = 0; j < q and q * i + j < n; j++)
            bl[i] = min(bl[i], v[q * i + j]);
    }
}
int query(int a, int b) {
    int ret = INF;
    // linear no bloco de a
    for (; a <= b and a % q; a++) ret = min(ret, v[a]);
    // bloco por bloco
    for (; a + q <= b; a += q) ret = min(ret, bl[a / q]);</pre>
    // linear no bloco de b
    for (; a <= b; a++) ret = min(ret, v[a]);</pre>
    return ret:
}
```

#### 4.7 BIT 2D

```
// BIT de soma 1-based
// Para mudar o valor da posicao (x, y) para k,
// faca: poe(x, y, k - sum(x, y, x, y))
// Complexidades:
// poe - 0(\log^2(n))
// \text{ query - } O(\log^2(n))
int n;
int bit[MAX][MAX];
void poe(int x, int y, int k) {
    for (int y2 = y; x \le n; x += x & -x)
        for (y = y2; y \le n; y += y \& -y)
            bit[x][y] += k;
}
int sum(int x, int y) {
    int ret = 0;
    for (int y2 = y; x; x -= x & -x)
        for (y = y2; y; y -= y & -y)
            ret += bit[x][y];
    return ret;
}
int query(int x, int y, int z, int w) {
    return sum(z, w) - sum(x-1, w)
        - sum(z, y-1) + sum(x-1, y-1);
}
4.8 DSU
// Une dois conjuntos e acha a qual conjunto um elemento
   pertence por seu id
//
// dsu_build: 0(n)
// find e unite: O(a(n)) \sim = O(1) amortizado
```

```
int id[MAX], sz[MAX];
void dsu_build(int n) { for(int i=0; i<n; i++) sz[i] = 1,</pre>
   id[i] = i; }
int find(int a) { return id[a] = a == id[a] ? a :
   find(id[a]); }
void unite(int a, int b) {
    a = find(a), b = find(b);
   if(a == b) return;
   if(sz[a] < sz[b]) swap(a,b);
    sz[a] += sz[b];
    id[b] = a;
}
4.9 Treap Implicita
// Todas as operacoes custam
// O(log(n)) com alta probabilidade
mt19937 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
template < typename T> struct treap {
    struct node {
        node *1, *r;
        int p, sz;
        T val, sub, lazy;
        bool rev;
        node(T \ v) : l(NULL), r(NULL), p(rng()), sz(1),
           val(v), sub(v), lazy(0), rev(0) {}
        void prop() {
            if (lazy) {
                val += lazy, sub += lazy*sz;
                if (1) 1->lazy += lazy;
                if (r) r->lazy += lazy;
```

```
if (rev) {
             swap(1, r);
             if (1) 1->rev ^= 1;
             if (r) r->rev ^= 1;
        }
        lazy = 0, rev = 0;
    }
    void update() {
        sz = 1, sub = val;
        if (1) 1->prop(), sz += 1->sz, sub += 1->sub;
        if (r) r\rightarrow prop(), sz += r\rightarrow sz, sub += r\rightarrow sub;
    }
};
node* root;
treap() { root = NULL; }
treap(const treap& t) {
    throw logic_error("Nao copiar a treap!");
\simtreap() {
    vector < node *> q = {root};
    while (q.size()) {
        node* x = q.back(); q.pop_back();
        if (!x) continue;
        q.push_back(x->1), q.push_back(x->r);
        delete x;
    }
}
int size(node* x) { return x ? x->sz : 0; }
int size() { return size(root); }
void join(node* 1, node* r, node*& i) { // assume que 1
   < r
    if (!1 or !r) return void(i = 1 ? 1 : r);
    1->prop(), r->prop();
    if (1->p > r->p) join(1->r, r, 1->r), i = 1;
    else join(1, r \rightarrow 1, r \rightarrow 1), i = r;
    i->update();
void split(node* i, node*& 1, node*& r, int v, int key =
   0) {
```

```
if (!i) return void(r = 1 = NULL);
        i->prop();
        if (key + size(i->1) < v) split(i->r, i->r, r, v,
            key+size(i->1)+1), 1 = i;
        else split(i \rightarrow 1, l, i \rightarrow 1, v, key), r = i;
        i->update();
    }
    void push_back(T v) {
        node* i = new node(v);
        join(root, i, root);
    }
    T query(int 1, int r) {
        node *L, *M, *R;
        split(root, M, R, r+1), split(M, L, M, 1);
        T ans = M->sub;
        join(L, M, M), join(M, R, root);
        return ans;
    }
    void update(int 1, int r, T s) {
        node *L, *M, *R;
        split(root, M, R, r+1), split(M, L, M, 1);
        M->lazy += s;
        join(L, M, M), join(M, R, root);
    }
    void reverse(int 1, int r) {
        node *L, *M, *R;
        split(root, M, R, r+1), split(M, L, M, 1);
        M \rightarrow rev = 1;
        join(L, M, M), join(M, R, root);
    }
};
```

### 4.10 Splay Tree Implicita

```
// vector da NASA
// Um pouco mais rapido q a treap
// O construtor a partir do vector
// eh linear, todas as outras operacoes
// custam O(log(n)) amortizado
```

```
template < typename T > struct splay {
    struct node {
        node *ch[2], *p;
        int sz;
        T val, sub, lazy;
        bool rev;
        node(T v) {
            ch[0] = ch[1] = p = NULL;
            sz = 1;
            sub = val = v;
            lazy = 0;
            rev = false;
        }
        void prop() {
            if (lazy) {
                val += lazy, sub += lazy*sz;
                if (ch[0]) ch[0]->lazy += lazy;
                if (ch[1]) ch[1]->lazy += lazy;
            }
            if (rev) {
                swap(ch[0], ch[1]);
                if (ch[0]) ch[0]->rev ^= 1;
                if (ch[1]) ch[1]->rev ^= 1;
            lazy = 0, rev = 0;
        void update() {
            sz = 1, sub = val;
            for (int i = 0; i < 2; i++) if (ch[i]) {</pre>
                ch[i]->prop();
                sz += ch[i]->sz;
                sub += ch[i] -> sub;
            }
        }
    };
    node* root;
    splay() { root = NULL; }
    splay(node* x) {
        root = x;
        if (root) root->p = NULL;
```

```
}
splay(vector < T > v) { // O(n)}
    root = NULL;
    for (T i : v) {
        node* x = new node(i);
        x \rightarrow ch[0] = root;
        if (root) root->p = x;
        root = x;
        root ->update();
    }
}
splay(const splay& t) {
    throw logic_error("Nao copiar a splay!");
\simsplay() {
    vector < node *> q = {root};
    while (q.size()) {
        node* x = q.back(); q.pop_back();
        if (!x) continue;
        q.push_back(x->ch[0]), q.push_back(x->ch[1]);
        delete x;
    }
}
int size(node* x) { return x ? x->sz : 0; }
void rotate(node* x) { // x vai ficar em cima
    node *p = x->p, *pp = p->p;
    if (pp) pp - ch[pp - ch[1] == p] = x;
    bool d = p \rightarrow ch[0] == x;
    p - ch[!d] = x - ch[d], x - ch[d] = p;
    if (p->ch[!d]) p->ch[!d]->p = p;
    x->p = pp, p->p = x;
    p->update(), x->update();
}
node* splaya(node* x) {
    if (!x) return x;
    root = x, x->update();
    while (x->p) {
        node *p = x->p, *pp = p->p;
        if (!pp) return rotate(x), x; // zig
        if ((pp->ch[0] == p)^(p->ch[0] == x))
            rotate(x), rotate(x); // zigzag
```

```
else rotate(p), rotate(x); // zigzig
    }
    return x;
node* find(int v) {
    if (!root) return NULL;
    node *x = root;
    int key = 0;
    while (1) {
        x->prop();
        bool d = key + size(x->ch[0]) < v;
        if (\text{key} + \text{size}(x->\text{ch}[0]) != v \text{ and } x->\text{ch}[d]) {
             if (d) key += size(x->ch[0])+1;
             x = x - ch[d];
        } else break;
    return splaya(x);
}
int size() { return root ? root->sz : 0; }
void join(splay<T>& 1) { // assume que 1 < *this</pre>
    if (!size()) swap(root, 1.root);
    if (!size() or !l.size()) return;
    node* x = 1.root;
    while (1) {
        x->prop();
        if (!x->ch[1]) break;
        x = x -> ch[1];
    1.splaya(x), root->prop(), root->update();
    x - ch[1] = root, x - ch[1] - p = x;
    root = 1.root, 1.root = NULL;
    root ->update();
}
node* split(int v) { // retorna os elementos < v</pre>
    if (v <= 0) return NULL;</pre>
    if (v >= size()) {
        node* ret = root;
        root = NULL;
        ret ->update();
        return ret;
    }
    find(v);
```

```
node*1 = root -> ch[0];
    root -> ch [0] = NULL;
    if (1) 1->p = NULL;
    root ->update();
    return 1;
}
T& operator [](int i) {
    find(i);
    return root -> val;
}
void push_back(T v) { // 0(1)
    node* r = new node(v);
    r - > ch[0] = root;
    if (root) root->p = r;
    root = r, root->update();
}
T query(int 1, int r) {
    splay <T> M(split(r+1));
    splay<T> L(M.split(1));
    T ans = M.root->sub;
    M. join(L), join(M);
    return ans;
}
void update(int 1, int r, T s) {
    splay <T> M(split(r+1));
    splay <T> L(M.split(1));
    M.root->lazy += s;
    M. join(L), join(M);
}
void reverse(int 1, int r) {
    splay <T> M(split(r+1));
    splay<T> L(M.split(1));
    M.root->rev ^= 1;
    M.join(L), join(M);
}
void erase(int 1, int r) {
    splay <T> M(split(r+1));
    splay<T> L(M.split(1));
    join(L);
}
```

};

#### 4.11 MergeSort Tree

//

```
// query(a, b, val) retorna numero de
// elementos em [a, b] <= val</pre>
// Usa O(n log(n)) de memoria
// Complexidades:
// build - O(n log(n))
// query - O(log^2(n))
#define ALL(x) x.begin(),x.end()
int v[MAX], n;
vector < int > tree [4*MAX];
void build(int p, int l, int r) {
    if (1 == r) return tree[p].push_back(v[1]);
    int m = (1+r)/2;
    build(2*p, 1, m), build(2*p+1, m+1, r);
    merge(ALL(tree[2*p]), ALL(tree[2*p+1]),
       back_inserter(tree[p]));
}
int query(int a, int b, int val, int p=1, int l=0, int
   r=n-1) {
    if (b < l or r < a) return 0; // to fora</pre>
    if (a \le 1 \text{ and } r \le b) // to totalmente dentro
        return lower_bound(ALL(tree[p]), val+1) -
            tree[p].begin();
    int m = (1+r)/2;
    return query(a, b, val, 2*p, 1, m) + query(a, b, val,
       2*p+1, m+1, r);
}
4.12 SegTree
// Recursiva com Lazy Propagation
// Query: soma do range [a, b]
// Update: soma x em cada elemento do range [a, b]
```

```
// Complexidades:
// build - O(n)
// query - 0(log(n))
// update - 0(log(n))
namespace seg {
    ll seg[4*MAX], lazy[4*MAX];
    int n, *v;
    ll build(int p=1, int l=0, int r=n-1) {
        lazy[p] = 0;
        if (1 == r) return seg[p] = v[1];
        int m = (1+r)/2;
        return seg[p] = build(2*p, 1, m) + build(2*p+1, m+1,
           r):
    }
    void build(int n2, int* v2) {
        n = n2, v = v2;
        build();
    }
    void prop(int p, int l, int r) {
        seg[p] += lazy[p]*(r-l+1);
        if (1 != r) lazy[2*p] += lazy[p], lazy[2*p+1] +=
           lazy[p];
        lazy[p] = 0;
    }
    ll query(int a, int b, int p=1, int l=0, int r=n-1) {
        prop(p, l, r);
        if (a <= 1 and r <= b) return seg[p];</pre>
        if (b < 1 \text{ or } r < a) \text{ return } 0;
        int m = (1+r)/2;
        return query (a, b, 2*p, 1, m) + query (a, b, 2*p+1,
           m+1, r);
    }
    ll update(int a, int b, int x, int p=1, int l=0, int
       r=n-1) {
        prop(p, 1, r);
        if (a <= 1 and r <= b) {</pre>
            lazy[p] += x;
            prop(p, 1, r);
            return seg[p];
        }
```

```
if (b < l or r < a) return seg[p];</pre>
        int m = (1+r)/2;
        return seg[p] = update(a, b, x, 2*p, 1, m) +
            update(a, b, x, 2*p+1, m+1, r);
};
// Se tiver uma seg de max, da pra descobrir em O(\log(n))
// o primeiro e ultimo elemento >= val numa range:
// primeira posicao >= val em [a, b] (ou -1 se nao tem)
int get_left(int a, int b, int val, int p=1, int l=0, int
   r=n-1) {
   if (b < 1 or r < a or seg[p] < val) return -1;</pre>
    if (r == 1) return 1;
    int m = (1+r)/2;
    int x = get_left(a, b, val, 2*p, 1, m);
    if (x != -1) return x;
    return get_left(a, b, val, 2*p+1, m+1, r);
}
// ultima posicao >= val em [a, b] (ou -1 se nao tem)
int get_right(int a, int b, int val, int p=1, int l=0, int
   r=n-1) {
   if (b < l or r < a or seg[p] < val) return -1;</pre>
    if (r == 1) return 1;
    int m = (1+r)/2;
    int x = get_right(a, b, val, 2*p+1, m+1, r);
    if (x != -1) return x;
    return get_right(a, b, val, 2*p, 1, m);
}
4.13 SegTree Colorida
// Cada posicao tem um valor e uma cor
```

```
// Cada posicao tem um valor e uma cor
// O construtor receve um vector de {valor, cor}
// e o numero de cores (as cores devem estar em [0, c-1])
// query(c, a, b) retorna a soma dos valores
// de todo mundo em [a, b] que tem cor c
// update(c, a, b, x) soma x em todo mundo em
```

```
// [a, b] que tem cor c
// paint(c1, c2, a, b) faz com que todo mundo
// em [a, b] que tem cor c1 passe a ter cor c2
//
// Complexidades:
// construir - O(n log(n)) espaco e tempo
// query - O(log(n))
// update - 0(log(n))
// paint - O(log(n)) amortizado
struct seg_color {
    struct node {
        node *1, *r;
        int cnt;
        ll val, lazy;
        node() : 1(NULL), r(NULL), cnt(0), val(0), lazy(0) {}
        void update() {
            cnt = 0, val = 0;
            for (auto i : {1, r}) if (i) {
                i->prop();
                cnt += i->cnt, val += i->val;
            }
        }
        void prop() {
            if (!lazy) return;
            val += lazy*(ll)cnt;
            for (auto i : {1, r}) if (i) i->lazy += lazy;
            lazy = 0;
        }
    };
    int n;
    vector < node *> seg;
    seg_color(vector<pair<int, int>>& v, int c) :
       n(v.size()), seg(c, NULL) {
        for (int i = 0; i < n; i++)</pre>
            seg[v[i].second] = insert(seg[v[i].second], i,
               v[i].first, 0, n-1);
    \simseg_color() {
        queue < node *> q;
```

```
for (auto i : seg) q.push(i);
    while (q.size()) {
        auto i = q.front(); q.pop();
        if (!i) continue;
        q.push(i \rightarrow l), q.push(i \rightarrow r);
        delete i;
    }
}
node* insert(node* at, int idx, int val, int l, int r) {
    if (!at) at = new node();
    if (l == r) return at->cnt = 1, at->val = val, at;
    int m = (1+r)/2;
    if (idx <= m) at->l = insert(at->l, idx, val, l, m);
    else at->r = insert(at->r, idx, val, m+1, r);
    return at->update(), at;
}
11 query(node* at, int a, int b, int l, int r) {
    if (!at or b < l or r < a) return 0;</pre>
    at ->prop();
    if (a <= l and r <= b) return at->val;
    int m = (1+r)/2;
    return query(at->1, a, b, 1, m) + query(at->r, a, b,
       m+1, r);
}
11 query(int c, int a, int b) { return query(seg[c], a,
   b, 0, n-1); }
void update(node* at, int a, int b, int x, int l, int r)
   {
    if (!at or b < l or r < a) return;</pre>
    at->prop();
    if (a <= 1 and r <= b) {
        at - > lazy += x;
        return void(at->prop());
    }
    int m = (1+r)/2;
    update(at->1, a, b, x, 1, m), update(at->r, a, b, x,
       m+1, r);
    at ->update();
void update(int c, int a, int b, int x) { update(seg[c],
   a, b, x, 0, n-1); }
```

```
void paint(node*& from, node*& to, int a, int b, int 1,
       int r) {
        if (to == from or !from or b < 1 or r < a) return;
        from ->prop();
        if (to) to->prop();
        if (a <= 1 and r <= b) {
            if (!to) {
                to = from;
                from = NULL;
                return;
            int m = (1+r)/2;
            paint(from->1, to->1, a, b, 1, m),
               paint(from->r, to->r, a, b, m+1, r);
            to->update();
            delete from;
            from = NULL;
            return;
        if (!to) to = new node();
        int m = (1+r)/2;
        paint(from->1, to->1, a, b, 1, m), paint(from->r,
           to->r, a, b, m+1, r);
        from ->update(), to ->update();
   }
    void paint(int c1, int c2, int a, int b) {
       paint(seg[c1], seg[c2], a, b, 0, n-1); }
};
```

# 4.14 SegTree Iterativa com Lazy Propagation

```
// Query: soma do range [a, b]
// Update: soma x em cada elemento do range [a, b]
// Para mudar, mudar as funcoes junta, poe e query
// LOG = ceil(log2(MAX))
//
// Complexidades:
// build - O(n)
// query - O(log(n))
// update - O(log(n))
```

```
namespace seg {
    11 seg[2*MAX], lazy[2*MAX];
    int n;
    ll junta(ll a, ll b) {
        return a+b;
    // soma x na posicao p de tamanho tam
    void poe(int p, ll x, int tam, bool prop=1) {
        seg[p] += x*tam;
        if (prop and p < n) lazy[p] += x;</pre>
    }
    // atualiza todos os pais da folha p
    void sobe(int p) {
        for (int tam = 2; p /= 2; tam *= 2) {
            seg[p] = junta(seg[2*p], seg[2*p+1]);
            poe(p, lazy[p], tam, 0);
        }
    }
    // propaga o caminho da raiz ate a folha p
    void prop(int p) {
        int tam = 1 << (LOG-1);
        for (int s = LOG; s; s--, tam /= 2) {
            int i = p >> s;
            if (lazy[i]) {
                poe(2*i, lazy[i], tam);
                poe(2*i+1, lazy[i], tam);
                lazy[i] = 0;
            }
        }
   }
    void build(int n2, int* v) {
        n = n2:
        for (int i = 0; i < n; i++) seg[n+i] = v[i];</pre>
        for (int i = n-1; i; i--) seg[i] = junta(seg[2*i],
           seg[2*i+1]);
        for (int i = 0; i < 2*n; i++) lazy[i] = 0;
```

```
}
    11 query(int a, int b) {
        11 \text{ ret} = 0;
        for (prop(a+=n), prop(b+=n); a \le b; ++a/=2, --b/=2)
            if (a%2 == 1) ret = junta(ret, seg[a]);
            if (b%2 == 0) ret = junta(ret, seg[b]);
        }
        return ret;
    }
    void update(int a, int b, int x) {
        int a2 = a += n, b2 = b += n, tam = 1;
        for (; a <= b; ++a/=2, --b/=2, tam *= 2) {
            if (a\%2 == 1) poe(a, x, tam);
            if (b\%2 == 0) poe(b, x, tam);
        sobe(a2), sobe(b2);
    }
};
4.15 SegTree Beats
// \text{ query}(a, b) - \{\{\min(v[a..b]), \max(v[a..b])\}, \sup(v[a..b])\}
// updatemin(a, b, x) faz com que v[i] <- min(v[i], x),</pre>
// para i em [a, b]
// updatemax faz o mesmo com max, e updatesum soma x
// em todo mundo do intervalo [a. b]
//
// Complexidades:
// build - O(n)
// query - O(log(n))
// update - O(log^2 (n)) amortizado
// (se nao usar updatesum, fica log(n) amortizado)
#define f first
```

#define s second

typedef long long 11;

const 11 LINF = 0x3f3f3f3f3f3f3f3f11;

```
namespace beats {
    struct node {
        int tam;
        ll sum, lazy; // lazy pra soma
        ll mi1, mi2, mi; // mi = #mi1
        ll ma1, ma2, ma; // ma = #ma1
        node(11 x = 0) {
            sum = mi1 = ma1 = x:
            mi2 = LINF, ma2 = -LINF;
            mi = ma = tam = 1;
            lazy = 0;
        }
        node(const node& 1, const node& r) {
            sum = 1.sum + r.sum, tam = 1.tam + r.tam;
            lazy = 0;
            if (1.mi1 > r.mi1) {
                mi1 = r.mi1, mi = r.mi;
                mi2 = min(1.mi1, r.mi2);
            } else if (l.mi1 < r.mi1) {</pre>
                mi1 = 1.mi1, mi = 1.mi;
                mi2 = min(r.mi1, 1.mi2);
            } else {
                mi1 = 1.mi1, mi = 1.mi+r.mi;
                mi2 = min(1.mi2, r.mi2);
            if (1.ma1 < r.ma1) {</pre>
                ma1 = r.ma1, ma = r.ma;
                ma2 = max(1.ma1, r.ma2);
            } else if (l.ma1 > r.ma1) {
                ma1 = l.ma1, ma = l.ma;
                ma2 = max(r.ma1, l.ma2);
            } else {
                ma1 = l.ma1, ma = l.ma+r.ma;
                ma2 = max(1.ma2, r.ma2);
            }
        }
        void setmin(ll x) {
            if (x >= ma1) return;
            sum += (x - ma1)*ma;
            if (mi1 == ma1) mi1 = x;
```

```
if (mi2 == ma1) mi2 = x;
        ma1 = x;
    void setmax(ll x) {
        if (x <= mi1) return;</pre>
        sum += (x - mi1)*mi;
        if (ma1 == mi1) ma1 = x;
        if (ma2 == mi1) ma2 = x;
        mi1 = x:
    void setsum(ll x) {
        mi1 += x, mi2 += x, ma1 += x, ma2 += x;
        sum += x*tam;
        lazy += x;
   }
};
node seg[4*MAX];
int n, *v;
node build(int p=1, int l=0, int r=n-1) {
   if (1 == r) return seg[p] = {v[1]};
    int m = (1+r)/2;
    return seg[p] = \{build(2*p, 1, m), build(2*p+1, m+1,
       r)}:
}
void build(int n2, int* v2) {
    n = n2, v = v2;
    build();
}
void prop(int p, int l, int r) {
    if (1 == r) return;
    for (int k = 0; k < 2; k++) {
        if (seg[p].lazy) seg[2*p+k].setsum(seg[p].lazy);
        seg[2*p+k].setmin(seg[p].ma1);
        seg[2*p+k].setmax(seg[p].mi1);
    }
    seg[p].lazy = 0;
pair < pair < 11, 11 > , 11 > query(int a, int b, int p=1, int
   1=0, int r=n-1) {
   if (b < 1 or r < a) return {{LINF, -LINF}, 0};</pre>
```

```
if (a \le 1 \text{ and } r \le b) \text{ return } \{\{seg[p].mi1,
        seg[p].ma1}, seg[p].sum};
    prop(p, 1, r);
    int m = (1+r)/2;
    auto L = query(a, b, 2*p, 1, m), R = query(a, b,
        2*p+1, m+1, r);
    return {{min(L.f.f, R.f.f), max(L.f.s, R.f.s)},
        L.s+R.s};
}
node updatemin(int a, int b, 11 x, int p=1, int 1=0, int
   r=n-1) {
    if (b < 1 or r < a or seg[p].ma1 <= x) return seg[p];</pre>
    if (a \le 1 \text{ and } r \le b \text{ and } seg[p].ma2 < x) {
        seg[p].setmin(x);
        return seg[p];
    prop(p, 1, r);
    int m = (1+r)/2;
    return seg[p] = {updatemin(a, b, x, 2*p, 1, m),
                     updatemin(a, b, x, 2*p+1, m+1, r)};
}
node updatemax(int a, int b, ll x, int p=1, int l=0, int
   r=n-1) {
    if (b < l or r < a or seg[p].mi1 >= x) return seg[p];
    if (a \le 1 \text{ and } r \le b \text{ and } seg[p].mi2 > x) {
        seg[p].setmax(x);
        return seg[p];
    prop(p, 1, r);
    int m = (1+r)/2;
    return seg[p] = \{updatemax(a, b, x, 2*p, 1, m),
                      updatemax(a, b, x, 2*p+1, m+1, r)};
}
node updatesum(int a, int b, ll x, int p=1, int l=0, int
   r=n-1) {
    if (b < 1 or r < a) return seg[p];</pre>
    if (a <= 1 and r <= b) {</pre>
        seg[p].setsum(x);
        return seg[p];
    prop(p, 1, r);
    int m = (1+r)/2;
```

#### 4.16 SegTree Esparca

```
// Query: soma do range [a, b]
// Update: flipa os valores de [a, b]
// O MAX tem q ser Q log N para Q updates
//
// Complexidades:
// build - O(1)
// query - 0(log(n))
// update - O(log(n))
namespace seg {
    int seg[MAX], lazy[MAX], R[MAX], L[MAX], ptr;
    int get_l(int i){
        if (L[i] == 0) L[i] = ptr++;
        return L[i];
    }
    int get_r(int i){
        if (R[i] == 0) R[i] = ptr++;
        return R[i];
    }
    void build() { ptr = 2; }
    void prop(int p, int 1, int r) {
        if (!lazy[p]) return;
        seg[p] = r-l+1 - seg[p];
        if (1 != r) lazy[get_l(p)]^=lazy[p],
           lazy[get_r(p)]^=lazy[p];
        lazy[p] = 0;
    }
    int query(int a, int b, int p=1, int l=0, int r=N-1) {
        prop(p, 1, r);
        if (b < 1 or r < a) return 0;
```

```
if (a <= l and r <= b) return seg[p];</pre>
        int m = (1+r)/2;
        return query(a, b, get_l(p), l, m)+query(a, b,
           get_r(p), m+1, r);
    }
    int update(int a, int b, int p=1, int l=0, int r=N-1) {
        prop(p, 1, r);
        if (b < l or r < a) return seg[p];</pre>
        if (a <= 1 and r <= b) {
            lazy[p] ^= 1;
            prop(p, 1, r);
            return seg[p];
        int m = (1+r)/2;
        return seg[p] = update(a, b, get_l(p), l,
           m)+update(a, b, get_r(p), m+1, r);
};
```

# 4.17 SegTree Iterativa

```
// Consultas O-based
// Valores iniciais devem estar em (seg[n], ..., seg[2*n-1])
// Query: soma do range [a, b]
// Update: muda o valor da posicao p para x
//
// Complexidades:
// build - O(n)
// query - O(log(n))
// update - O(log(n))
int seg[2 * MAX];
int n;

void build() {
   for (int i = n - 1; i; i--) seg[i] = seg[2*i] +
        seg[2*i+1];
}
```

```
int query(int a, int b) {
    int ret = 0;
    for(a += n, b += n; a <= b; ++a /= 2, --b /= 2) {
        if (a % 2 == 1) ret += seg[a];
        if (b % 2 == 0) ret += seg[b];
    }
    return ret;
}

void update(int p, int x) {
    seg[p += n] = x;
    while (p /= 2) seg[p] = seg[2*p] + seg[2*p+1];
}</pre>
```

# 4.18 SegTree Persistente

```
// SegTree de soma, update de somar numa posicao
//
// query(a, b, t) retorna a query de [a, b] na versao t
// update(a, x, t) faz um update v[a]+=x a partir da
// versao de t, criando uma nova versao e retornando seu id
// Por default, faz o update a partir da ultima versao
//
// build - O(n)
// query - O(log(n))
// update - O(log(n))
const int MAX = 3e4+10, UPD = 2e5+10, LOG = 20;
const int MAXS = 4*MAX+UPD*LOG;
namespace perseg {
    11 seg[MAXS];
   int rt[UPD], L[MAXS], R[MAXS], cnt, t;
    int n, *v;
    ll build(int p, int l, int r) {
        if (1 == r) return seg[p] = v[1];
        L[p] = cnt++, R[p] = cnt++;
        int m = (1+r)/2;
```

```
return seg[p] = build(L[p], 1, m) + build(R[p], m+1,
           r);
    void build(int n2, int* v2) {
        n = n2, v = v2;
        rt[0] = cnt++;
        build(0, 0, n-1);
    }
    ll query(int a, int b, int p, int l, int r) {
        if (b < 1 or r < a) return 0;
        if (a <= 1 and r <= b) return seg[p];</pre>
        int m = (1+r)/2;
        return query(a, b, L[p], 1, m) + query(a, b, R[p],
           m+1, r);
    }
    ll query(int a, int b, int tt) {
        return query(a, b, rt[tt], 0, n-1);
    ll update(int a, int x, int lp, int p, int l, int r) {
        if (l == r) return seg[p] = seg[lp]+x;
        int m = (1+r)/2;
        if (a <= m)
            return seg[p] = update(a, x, L[lp], L[p]=cnt++,
               1, m) + seg[R[p]=R[lp]];
        return seg[p] = seg[L[p]=L[lp]] + update(a, x,
           R[lp], R[p] = cnt ++, m+1, r);
    int update(int a, int x, int tt=t) {
        update(a, x, rt[tt], rt[++t]=cnt++, 0, n-1);
        return t;
   }
};
```

#### 4.19 SegTree 2D Iterativa

```
// Consultas 0-based
// Um valor inicial em (x, y) deve ser colocado em
   seg[x+n][y+n]
// Query: soma do retangulo ((x1, y1), (x2, y2))
// Update: muda o valor da posicao (x, y) para val
```

```
// Nao pergunte como que essa coisa funciona
//
// Para query com distancia de manhattan <= d, faca
// nx = x+y, ny = x-y
// Update em (nx, ny), query em ((nx-d, ny-d), (nx+d, ny+d))
//
// Se for de min/max, pode tirar os if's da 'query', e fazer
// sempre as 4 operacoes. Fica mais rapido
//
// Complexidades:
// build - O(n^2)
// \text{ query - O(log^2(n))}
// update - 0(log^2(n))
int seg[2*MAX][2*MAX], n;
void build() {
    for (int x = 2*n; x; x--) for (int y = 2*n; y; y--) {
        if (x < n) seg[x][y] = seg[2*x][y] + seg[2*x+1][y];
        if (y < n) seg[x][y] = seg[x][2*y] + seg[x][2*y+1];
    }
}
int query(int x1, int y1, int x2, int y2) {
    int ret = 0, y3 = y1 + n, y4 = y2 + n;
    for (x1 += n, x2 += n; x1 <= x2; ++x1 /= 2, --x2 /= 2)
        for (y1 = y3, y2 = y4; y1 \le y2; ++y1 /= 2, --y2 /=
            2) {
             if (x1\%2 == 1 \text{ and } y1\%2 == 1) \text{ ret } += \text{seg}[x1][y1];
             if (x1\%2 == 1 \text{ and } y2\%2 == 0) \text{ ret } += \text{seg}[x1][y2];
             if (x2\%2 == 0 \text{ and } y1\%2 == 1) \text{ ret } += \text{seg}[x2][y1];
             if (x2\%2 == 0 \text{ and } y2\%2 == 0) \text{ ret } += \text{seg}[x2][y2];
        }
    return ret;
}
void update(int x, int y, int val) {
    int y2 = y += n;
    for (x += n; x; x /= 2, y = y2) {
        if (x \ge n) seg[x][y] = val;
        else seg[x][y] = seg[2*x][y] + seg[2*x+1][y];
```

```
while (y /= 2) seg[x][y] = seg[x][2*y] +
         seg[x][2*y+1];
}
```

## 4.20 Split-Merge Set

```
// Representa um conjunto de inteiros nao negativos
// Todas as operacoes custam O(log(N)),
// em que N = maior elemento do set,
// exceto o merge, que custa O(log(N)) amortizado
// Usa O(min(N, n log(N))) de memoria, sendo 'n' o
// numero de elementos distintos no set
template < typename T, bool MULTI = false, typename SIZE_T = int >
   struct sms {
    struct node {
        node *1, *r;
        SIZE_T cnt;
        node() : l(NULL), r(NULL), cnt(0) {}
        void update() {
            cnt = 0;
            if (1) cnt += 1->cnt;
            if (r) cnt += r->cnt;
        }
    };
    node* root;
    T N;
    sms() : root(NULL), N(0) {}
    sms(T v) : sms() { while (v >= N) N = 2*N+1; }
    sms(const sms& t) : root(NULL), N(t.N) {
        for (SIZE_T i = 0; i < t.size(); i++) {</pre>
            T at = t[i];
            SIZE_T qt = t.count(at);
            insert(at, qt);
            i += qt-1;
        }
```

```
}
sms(initializer_list<T> v) : sms() { for (T i : v)
   insert(i); }
\simsms() {
    vector < node *> q = {root};
    while (q.size()) {
        node* x = q.back(); q.pop_back();
        if (!x) continue;
        q.push_back(x->1), q.push_back(x->r);
        delete x;
    }
}
friend void swap(sms& a, sms& b) {
    swap(a.root, b.root), swap(a.N, b.N);
}
SIZE_T size() const { return root ? root->cnt : 0; }
SIZE_T count(node* x) const { return x ? x->cnt : 0; }
void clear() {
    sms tmp;
    swap(*this, tmp);
}
void expand(T v) {
    for (; N < v; N = 2*N+1) if (root) {
        node* nroot = new node();
        nroot ->1 = root;
        root = nroot;
        root ->update();
   }
}
node* insert(node* at, T idx, SIZE_T qt, T 1, T r) {
    if (!at) at = new node();
    if (1 == r) {
        at->cnt += qt;
        if (!MULTI) at->cnt = 1;
        return at;
    T m = 1 + (r-1)/2;
    if (idx \le m) at->1 = insert(at->1, idx, qt, 1, m);
    else at->r = insert(at->r, idx, qt, m+1, r);
    return at->update(), at;
```

```
}
void insert(T v, SIZE_T qt=1) { // insere 'qt'
   ocorrencias de 'v'
    if (qt <= 0) return erase(v, -qt);</pre>
    assert(v >= 0);
    expand(v);
    root = insert(root, v, qt, 0, N);
}
node* erase(node* at, T idx, SIZE_T qt, T 1, T r) {
    if (!at) return at;
    if (1 == r) at->cnt = at->cnt < qt ? 0 : at->cnt -
    else {
        T m = 1 + (r-1)/2;
        if (idx \le m) at->1 = erase(at->1, idx, qt, 1,
            m):
        else at->r = erase(at->r, idx, qt, m+1, r);
        at ->update();
    if (!at->cnt) delete at, at = NULL;
    return at;
void erase(T v, SIZE_T qt=1) { // remove 'qt'
   ocorrencias de 'v'
    if (v < 0 \text{ or } v > N \text{ or } !qt) \text{ return};
    if (qt < 0) insert(v, -qt);</pre>
    root = erase(root, v, qt, 0, N);
}
void erase_all(T v) { // remove todos os 'v'
    if (v < 0 \text{ or } v > N) return;
    root = erase(root, v, numeric_limits < SIZE_T >:: max(),
        0, N);
}
SIZE_T count(node* at, T a, T b, T l, T r) const {
    if (!at or b < 1 or r < a) return 0;</pre>
    if (a <= 1 and r <= b) return at->cnt;
    T m = 1 + (r-1)/2;
    return count(at->1, a, b, 1, m) + count(at->r, a, b,
       m+1, r);
}
```

```
SIZE_T count(T v) const { return count(root, v, v, 0,
SIZE_T order_of_key(T v) { return count(root, 0, v-1, 0,
SIZE_T lower_bound(T v) { return order_of_key(v); }
const T operator [](SIZE_T i) const { // i-esimo menor
   elemento
    assert(i >= 0 and i < size());</pre>
    node* at = root;
    T 1 = 0, r = N;
    while (1 < r) {
        T m = 1 + (r-1)/2;
        if (count(at->1) > i) at = at->1, r = m;
            i -= count(at->1):
            at = at->r; l = m+1;
        }
    }
    return 1;
}
node* merge(node* 1, node* r) {
    if (!1 or !r) return 1 ? 1 : r;
    if (!1->1 \text{ and } !1->r) \{ // \text{ folha} \}
        if (MULTI) 1->cnt += r->cnt;
        delete r;
        return 1;
    1 - > 1 = merge(1 - > 1, r - > 1), 1 - > r = merge(1 - > r, r - > r);
    1->update(), delete r;
    return 1;
}
void merge(sms& s) { // mergeia dois sets
    if (N > s.N) swap(*this, s);
    expand(s.N);
    root = merge(root, s.root);
    s.root = NULL;
}
node* split(node*& x, SIZE_T k) {
    if (k <= 0 or !x) return NULL;
```

```
node* ret = new node();
        if (!x->1 \text{ and } !x->r) x->cnt -= k, ret->cnt += k;
            if (k \le count(x->1)) ret->1 = split(x->1, k);
            else {
                ret->r = split(x->r, k - count(x->1));
                 swap(x->1, ret->1);
            ret ->update(), x->update();
        }
        if (!x->cnt) delete x, x = NULL;
        return ret;
    }
    void split(SIZE_T k, sms& s) { // pega os 'k' menores
        s.clear();
        s.root = split(root, min(k, size()));
        s.N = N;
    }
    // pega os menores que 'k'
    void split_val(T k, sms& s) { split(order_of_key(k), s);
};
```

### 4.21 BIT

```
// BIT de soma 1-based, v 0-based
// Para mudar o valor da posicao p para x,
// faca: poe(x - query(p, p), p)
// l_bound(x) retorna o menor p tal que
// query(1, p+1) > x
                      (0 based!)
// Complexidades:
// build - O(n)
// poe - O(log(n))
// query - O(log(n))
// l_bound - O(log(n))
int n;
int bit[MAX];
int v[MAX];
```

```
void build() {
    bit[0] = 0;
    for (int i = 1; i <= n; i++) bit[i] = v[i - 1];</pre>
    for (int i = 1; i <= n; i++) {
        int j = i + (i \& -i);
        if (j <= n) bit[j] += bit[i];</pre>
    }
}
// soma x na posicao p
void poe(int x, int p) {
    for (; p <= n; p += p & -p) bit[p] += x;</pre>
}
// soma [1, p]
int pref(int p) {
    int ret = 0;
    for (; p; p -= p & -p) ret += bit[p];
    return ret;
}
// soma [a, b]
int query(int a, int b) {
    return pref(b) - pref(a - 1);
}
int l_bound(ll x) {
    int p = 0;
   for (int i = MAX2; i+1; i--) if (p + (1<<i) <= n
        and bit [p + (1 << i)] <= x) x -= bit <math>[p += (1 << i)];
    return p;
}
      Order Statistic Set
4.22
```

```
// Funciona do C++11 pra cima
#include <ext/pb_ds/assoc_container.hpp>
```

```
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
template <class T>
    using ord_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
// para declarar:
ord_set < int > s;
// coisas do set normal funcionam:
for (auto i : s) cout << i << endl;</pre>
cout << s.size() << endl;</pre>
// k-esimo maior elemento O(\log|s|):
// k=0: menor elemento
cout << *s.find_by_order(k) << endl;</pre>
// quantos sao menores do que k O(log|s|):
cout << s.order_of_key(k) << endl;</pre>
// Para fazer um multiset, tem que
// usar ord_set<pair<int, int> > com o
// segundo parametro sendo algo para diferenciar
// os ementos iguais.
// s.order_of_key({k, -INF}) vai retornar o
// numero de elementos < k
4.23 Treap
// Todas as operacoes custam
// O(log(n)) com alta probabilidade
mt19937 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
```

node(T v) : l(NULL), r(NULL), p(rng()), sz(1),

template < typename T> struct treap {

struct node {

T val;

node \*1, \*r;

val(v) {}

void update() {

int p, sz;

```
if (1) sz += 1->sz;
        if (r) sz += r->sz;
   }
};
node* root;
treap() { root = NULL; }
treap(const treap& t) {
    throw logic_error("Nao copiar a treap!");
\simtreap() {
    vector < node *> q = {root};
    while (q.size()) {
        node* x = q.back(); q.pop_back();
        if (!x) continue;
        q.push_back(x->1), q.push_back(x->r);
        delete x;
    }
}
int size(node* x) { return x ? x->sz : 0; }
int size() { return size(root); }
void join(node* 1, node* r, node*& i) { // assume que 1
   < r
    if (!l or !r) return void(i = 1 ? l : r);
    if (1->p > r->p) join(1->r, r, 1->r), i = 1;
    else join(1, r->1, r->1), i = r;
    i->update();
}
void split(node* i, node*& l, node*& r, T v) {
    if (!i) return void(r = 1 = NULL);
    if (i->val < v) split(i->r, i->r, r, v), l = i;
    else split(i \rightarrow 1, l, i \rightarrow 1, v), r = i;
    i->update();
}
int count(node* i, T v) {
    if (!i) return 0;
    if (i->val == v) return 1;
    if (v < i->val) return count(i->l, v);
    return count(i->r, v);
```

```
}
    void index_split(node* i, node*& 1, node*& r, int v, int
       kev = 0) {
        if (!i) return void(r = 1 = NULL);
        if (key + size(i->1) < v) index_split(i->r, i->r, r,
           v, key+size(i->1)+1), l = i;
        else index_split(i->1, 1, i->1, v, key), r = i;
        i->update();
    }
    int count(T v) {
        return count(root, v);
    void insert(T v) {
        if (count(v)) return;
        node *L, *R;
        split(root, L, R, v);
        node* at = new node(v);
        join(L, at, L);
        join(L, R, root);
    void erase(T v) {
        node *L, *M, *R;
        split(root, M, R, v+1), split(M, L, M, v);
        if (M) delete M;
        M = NULL;
        join(L, R, root);
    }
};
4.24 Split-Merge Set - Lazy
// Representa um conjunto de inteiros nao negativos
// Todas as operacoes custam O(log(N)),
// em que N = maior elemento do set,
// exceto o merge e o insert_range, que custa O(log(N))
   amortizado
// Usa O(min(N, n log(N))) de memoria, sendo 'n' o
```

// numero de elementos distintos no set

template < typename T > struct sms {

```
struct node {
    node *1, *r;
    int cnt;
    bool flip;
    node(): 1(NULL), r(NULL), cnt(0), flip(0) {}
    void update() {
        cnt = 0;
        if (1) cnt += 1->cnt;
        if (r) cnt += r->cnt;
    }
};
void prop(node* x, int size) {
    if (!x or !x->flip) return;
    x \rightarrow flip = 0;
    x \rightarrow cnt = size - x \rightarrow cnt;
    if (size > 1) {
        if (!x->1) x->1 = new node();
        if (!x->r) x->r = new node();
        x - > 1 - > flip ^= 1;
        x->r->flip ^= 1;
    }
}
node* root;
T N;
sms() : root(NULL), N(0) {}
sms(T v) : sms() { while (v >= N) N = 2*N+1; }
sms(sms& t) : root(NULL), N(t.N) {
    for (int i = 0; i < t.size(); i++) insert(t[i]);</pre>
sms(initializer_list<T> v) : sms() { for (T i : v)
   insert(i): }
void destroy(node* r) {
    vector < node *> q = {r};
    while (q.size()) {
        node* x = q.back(); q.pop_back();
        if (!x) continue;
        q.push_back(x->1), q.push_back(x->r);
        delete x;
    }
```

```
\simsms() { destroy(root); }
friend void swap(sms& a, sms& b) {
    swap(a.root, b.root), swap(a.N, b.N);
int count(node* x, T size) {
    if (!x) return 0;
    prop(x, size);
    return x->cnt;
int size() { return count(root, N+1); }
void clear() {
    sms tmp;
    swap(*this, tmp);
void expand(T v) {
    for (; N < v; N = 2*N+1) if (root) {
        prop(root, N+1);
        node* nroot = new node();
        nroot ->1 = root;
        root = nroot;
        root ->update();
    }
}
node* insert(node* at, T idx, T l, T r) {
    if (!at) at = new node();
    else prop(at, r-l+1);
    if (1 == r) {
        at->cnt = 1;
        return at;
   }
    T m = 1 + (r-1)/2:
    if (idx \le m) at->1 = insert(at->1, idx, 1, m);
    else at->r = insert(at->r, idx, m+1, r);
    return at->update(), at;
}
void insert(T v) {
    assert(v >= 0);
    expand(v);
    root = insert(root, v, 0, N);
```

```
}
node* erase(node* at, T idx, T 1, T r) {
    if (!at) return at;
    prop(at, r-l+1);
    if (1 == r) at -> cnt = 0;
    else {
        T m = 1 + (r-1)/2;
        if (idx \le m) at > 1 = erase(at > 1, idx, 1, m);
        else at->r = erase(at->r, idx, m+1, r);
        at->update();
    return at;
}
void erase(T v) {
    if (v < 0 \text{ or } v > N) return:
    root = erase(root, v, 0, N);
}
int count(node* at, T a, T b, T l, T r) {
    if (!at or b < l or r < a) return 0;
    prop(at, r-l+1);
    if (a <= 1 and r <= b) return at->cnt;
    T m = 1 + (r-1)/2;
    return count(at->1, a, b, 1, m) + count(at->r, a, b,
       m+1, r);
}
int count(T v) { return count(root, v, v, 0, N); }
int order_of_key(T v) { return count(root, 0, v-1, 0,
   N): }
int lower_bound(T v) { return order_of_key(v); }
const T operator [](int i) { // i-esimo menor elemento
    assert(i >= 0 and i < size()):
    node* at = root:
    T 1 = 0, r = N;
    while (1 < r) {
        prop(at, r-l+1);
        T m = 1 + (r-1)/2;
        if (count(at->1, m-1+1) > i) at = at->1, r = m;
            i \rightarrow count(at \rightarrow 1, r-m);
```

```
at = at -> r : l = m+1 :
        }
    }
    return 1;
}
node* merge(node* a, node* b, T tam) {
    if (!a or !b) return a ? a : b;
    prop(a, tam), prop(b, tam);
    if (b \rightarrow cnt == tam) swap(a, b);
    if (tam == 1 \text{ or } a \rightarrow cnt == tam) {
        destroy(b);
        return a;
    }
    a - 1 = merge(a - 1, b - 1, tam > 1), a - r = merge(a - r, tam > 1)
       b->r, tam>>1);
    a->update(), delete b;
    return a;
void merge(sms& s) { // mergeia dois sets
    if (N > s.N) swap(*this, s);
    expand(s.N);
    root = merge(root, s.root, N+1);
    s.root = NULL;
}
node* split(node*& x, int k, T tam) {
    if (k <= 0 or !x) return NULL;</pre>
    prop(x, tam);
    node* ret = new node();
    if (tam == 1) x->cnt = 0, ret->cnt = 1;
    else {
        if (k \le count(x->1, tam>>1)) ret->1 =
            split(x->1, k, tam>>1);
        else {
             ret -> r = split(x -> r, k - count(x -> l,
                tam>>1), tam>>1);
             swap(x->1, ret->1);
        ret->update(), x->update();
    }
    return ret;
```

```
}
void split(int k, sms& s) { // pega os 'k' menores
    s.clear();
    s.root = split(root, min(k, size()), N+1);
    s.N = N;
}
// pega os menores que 'k'
void split_val(T k, sms& s) { split(order_of_key(k), s);
   }
void flip(node*& at, T a, T b, T l, T r) {
    if (!at) at = new node();
    else prop(at, r-l+1);
    if (a <= 1 and r <= b) {</pre>
        at ->flip ^= 1;
        prop(at, r-l+1);
        return:
    if (r < a or b < 1) return;
    T m = 1 + (r-1)/2;
    flip(at->1, a, b, 1, m), flip(at->r, a, b, m+1, r);
    at ->update();
}
void flip(T 1, T r) { // flipa os valores em [1, r]
    assert(1 >= 0 and 1 <= r);
    expand(r);
    flip(root, 1, r, 0, N);
}
// complemento considerando que o universo eh [0, lim]
void complement(T lim) {
    assert(lim >= 0);
    if (lim > N) expand(lim);
    flip(root, 0, lim, 0, N);
    sms tmp:
    split_val(lim+1, tmp);
    swap(*this, tmp);
}
void insert_range(T 1, T r) { // insere todo os valores
   em [1, r]
    sms tmp;
    tmp.flip(l, r);
    merge(tmp);
```

```
// trocar v[...] para ... na query
template < typename T > struct rmq {
    vector <T> v;
    int n, b;
    vector < int > id, st;
    vector < vector < int >> table;
    vector < vector < int >>> entre;
    int op(int x, int y) { return v[x] < v[y] ? x : y; }
    rmq(vector<T>& v_) {
        v = v_{n}, n = v.size();
        b = (\_builtin\_clz(1) - \_builtin\_clz(n) + 1)/4 + 1;
        id.resize(n);
        table.assign(4*b, vector<int>((n+b-1)/b));
        entre.assign(1<<b<<b, vector<vector<int>>(b,
            vector < int > (b, -1));
        for (int i = 0; i < n; i += b) {
            int at = 0, 1 = min(n, i+b);
            st.clear();
            for (int j = i; j < l; j++) {
                 while (st.size() and op(st.back(), j) == j)
                    st.pop_back(), at *= 2;
                 st.push_back(j), at = 2*at+1;
            for (int j = i; j < 1; j++) id[j] = at;</pre>
            if (entre[at][0][0] == -1) for (int x = 0; x <</pre>
               1-i; x++) {
                 entre[at][x][x] = x;
                for (int y = x+1; y < 1-i; y++)
                     entre[at][x][y] =
                        op(i+entre[at][x][y-1], i+y) - i;
            }
```

```
table [0][i/b] = i+entre[at][0][1-i-1];
        for (int j = 1; (1<<j) <= (n+b-1)/b; j++)
            for (int i = 0; i+(1 << j) <= (n+b-1)/b; i++)
                table[j][i] = op(table[j-1][i],
                   table [j-1][i+(1<<(j-1))]);
    T query(int i, int j) {
        if (i/b == j/b) return
           v[i/b*b+entre[id[i]][i%b][j%b]];
        int x = i/b+1, y = j/b-1, ans = i;
        if (x <= y) {
            int t = __builtin_clz(1) - __builtin_clz(y-x+1);
            ans = op(ans, op(table[t][x],
               table [t] [y-(1<<t)+1]);
        ans = op(ans, op(i/b*b+entre[id[i]][i\%b][b-1],
           j/b*b+entre[id[j]][0][j%b]));
        return v[ans];
   }
};
```

#### 4.26 DSU Persistente

```
// Persistencia parcial, ou seja, tem que ir
// incrementando o 't' no une
//
// Complexidades:
// build - O(n)
// find - O(log(n))
// une - O(log(n))

int n, p[MAX], sz[MAX], ti[MAX];

void build() {
   for (int i = 0; i < n; i++) {
      p[i] = i;
      sz[i] = 1;
      ti[i] = -INF;
}</pre>
```

```
int find(int k, int t) {
    if (p[k] == k or ti[k] > t) return k;
    return find(p[k], t);
}

void une(int a, int b, int t) {
    a = find(a, t); b = find(b, t);
    if (a == b) return;
    if (sz[a] > sz[b]) swap(a, b);

    sz[b] += sz[a];
    p[a] = b;
    ti[a] = t;
}
```

#### 4.27 Treap Persistent Implicita

```
// Todas as operacoes custam
// O(log(n)) com alta probabilidade
mt19937_64 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
struct node {
    node *1, *r;
    ll sz, val, sub;
    node(11 v): 1(NULL), r(NULL), sz(1), val(v), sub(v) {}
    node(node* x) : l(x->l), r(x->r), sz(x->sz),
       val(x->val), sub(x->sub) {}
    void update() {
        sz = 1, sub = val;
        if (1) sz += 1->sz, sub += 1->sub;
        if (r) sz += r->sz, sub += r->sub;
        sub %= MOD;
};
ll size(node* x) { return x ? x->sz : 0; }
```

```
void update(node* x) { if (x) x->update(); }
 node* copy(node* x) { return x ? new node(x) : NULL; }
 node* join(node* 1, node* r) {
    if (!1 or !r) return 1 ? copy(1) : copy(r);
     node* ret;
     if (rng() % (size(l) + size(r)) < size(l)) {</pre>
        ret = copy(1);
        ret->r = join(ret->r, r);
     } else {
         ret = copy(r);
         ret - > 1 = join(1, ret - > 1);
     }
     return update(ret), ret;
}
void split(node* x, node*& 1, node*& r, ll v, ll key = 0) {
     if (!x) return void(1 = r = NULL);
     if (key + size(x->1) < v) {
         1 = copv(x);
         split(1->r, 1->r, r, v, key+size(1->1)+1);
     } else {
         r = copy(x);
         split(r->1, 1, r->1, v, key);
     }
     update(1), update(r);
 vector < node *> treap;
 void init(const vector<ll>& v) {
     treap = {NULL};
    for (auto i : v) treap[0] = join(treap[0], new node(i));
}
      RMQ < O(n), O(1) > - min queue
// O(n) pra buildar, query O(1)
// Para retornar o indice, basta
// trocar v[...] para ... na query
```

```
template < typename T > struct rmq {
    vector <T> v;
    int n; static const int b = 30;
    vector < int > mask, t;
    int op(int x, int y) { return v[x] < v[y] ? x : y; }</pre>
    int msb(int x) { return
       __builtin_clz(1) - __builtin_clz(x); }
    int small(int r, int sz = b) { return
       r-msb(mask[r]&((1<<sz)-1));}
    rmg(const \ vector < T > \& \ v_) : v(v_), n(v.size()), mask(n),
       t(n) {
        for (int i = 0, at = 0; i < n; mask[i++] = at |= 1) {
            at = (at << 1) & ((1 << b) -1);
            while (at and op(i, i-msb(at&-at)) == i) at ^=
                at&-at;
        }
        for (int i = 0; i < n/b; i++) t[i] = small(b*i+b-1);
        for (int j = 1; (1<<j) <= n/b; j++) for (int i = 0;
           i+(1<< i) <= n/b; i++)
            t[n/b*j+i] = op(t[n/b*(j-1)+i],
               t[n/b*(j-1)+i+(1<<(j-1))]);
    T query(int 1, int r) {
        if (r-l+1 <= b) return v[small(r, r-l+1)];</pre>
        int ans = op(small(l+b-1), small(r));
        int x = 1/b+1, y = r/b-1;
        if (x <= y) {
            int j = msb(y-x+1);
            ans = op(ans, op(t[n/b*j+x],
                t[n/b*j+y-(1<<j)+1]));
        }
        return v[ans];
    }
};
4.29
      SQRT Tree
// RMQ em O(log log n) com O(n log log n) pra buildar
```

```
// Funciona com qualquer operacao associativa
// Tao rapido quanto a sparse table, mas usa menos memoria
// (log log (1e9) < 5, entao a query eh praticamente O(1))
//
// build - O(n log log n)
// query - O(log log n)
namespace sqrtTree {
    int n, *v;
    int pref[4][MAX], sulf[4][MAX], getl[4][MAX],
       entre[4][MAX], sz[4];
    int op(int a, int b) { return min(a, b); }
    inline int getblk(int p, int i) { return
       (i-getl[p][i])/sz[p]; }
    void build(int p, int l, int r) {
        if (1+1 >= r) return;
        for (int i = 1; i <= r; i++) getl[p][i] = 1;
        for (int L = 1; L <= r; L += sz[p]) {
            int R = min(L+sz[p]-1, r);
            pref[p][L] = v[L], sulf[p][R] = v[R];
            for (int i = L+1; i <= R; i++) pref[p][i] =</pre>
                op(pref[p][i-1], v[i]);
            for (int i = R-1; i >= L; i--) sulf[p][i] =
                op(v[i], sulf[p][i+1]);
            build(p+1, L, R);
        for (int i = 0; i <= sz[p]; i++) {
            int at = entre[p][1+i*sz[p]+i] =
               sulf[p][l+i*sz[p]];
            for (int j = i+1; j <= sz[p]; j++)</pre>
               entre[p][l+i*sz[p]+j] = at =
                    op(at, sulf[p][1+j*sz[p]]);
        }
    void build(int n2, int* v2) {
        n = n2, v = v2;
        for (int p = 0; p < 4; p++) sz[p] = n2 = sqrt(n2);
        build(0, 0, n-1);
    }
    int query(int 1, int r) {
        if (1+1 >= r) return 1 == r ? v[1] : op(v[1], v[r]);
```

#### 4.30 Wavelet Tree

```
// Usa O(sigma + n log(sigma)) de memoria,
// onde sigma = MAXN - MINN
// Depois do build, o v fica ordenado
// count(i, j, x, y) retorna o numero de elementos de
// v[i, j) que pertencem a [x, y]
// kth(i, j, k) retorna o elemento que estaria
// na poscicao k-1 de v[i, j), se ele fosse ordenado
// sum(i, j, x, y) retorna a soma dos elementos de
// v[i, j) que pertencem a [x, y]
// sumk(i, j, k) retorna a soma dos k-esimos menores
// elementos de v[i, j) (sum(i, j, 1) retorna o menor)
//
// Complexidades:
// build - O(n log(sigma))
// count - O(log(sigma))
// kth - 0(log(sigma))
// sum - O(log(sigma))
// sumk - O(log(sigma))
int n, v[MAX];
vector < int > esq[4*(MAXN-MINN)], pref[4*(MAXN-MINN)];
void build(int b = 0, int e = n, int p = 1, int l = MINN,
   int r = MAXN) {
    int m = (1+r)/2; esq[p].push_back(0);
       pref[p].push_back(0);
    for (int i = b; i < e; i++) {</pre>
        esq[p].push_back(esq[p].back()+(v[i]<=m));</pre>
```

```
pref[p].push_back(pref[p].back()+v[i]);
    }
    if (1 == r) return;
    int m2 = stable_partition(v+b, v+e, [=](int i){return i
       <= m;}) - v;
    build(b, m2, 2*p, 1, m), build(m2, e, 2*p+1, m+1, r);
}
int count(int i, int j, int x, int y, int p = 1, int l =
   MINN, int r = MAXN) {
   if (y < 1 or r < x) return 0;</pre>
   if (x \le 1 \text{ and } r \le y) \text{ return } j-i;
   int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
    return count(ei, ej, x, y, 2*p, 1, m)+count(i-ei, j-ej,
       x, y, 2*p+1, m+1, r);
}
int kth(int i, int j, int k, int p=1, int 1 = MINN, int r =
   MAXN) {
   if (1 == r) return 1;
   int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
    if (k <= ej-ei) return kth(ei, ej, k, 2*p, 1, m);</pre>
    return kth(i-ei, j-ej, k-(ej-ei), 2*p+1, m+1, r);
}
int sum(int i, int j, int x, int y, int p = 1, int l = MINN,
   int r = MAXN) {
    if (y < 1 or r < x) return 0;
   if (x <= l and r <= y) return pref[p][j]-pref[p][i];</pre>
    int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
    return sum(ei, ej, x, y, 2*p, 1, m) + sum(i-ei, j-ej, x,
       y, 2*p+1, m+1, r);
}
int sumk(int i, int j, int k, int p = 1, int l = MINN, int r
   = MAXN)
    if (1 == r) return l*k;
    int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
    if (k <= ej-ei) return sumk(ei, ej, k, 2*p, 1, m);</pre>
    return pref[2*p][ej]-pref[2*p][ei]+sumk(i-ei, j-ej,
       k-(ej-ei), 2*p+1, m+1, r);
}
```

#### 5 DP

#### 5.1 SOS DP

```
// O(n 2^n)
//iterative version
for(int mask = 0; mask < (1<<N); ++mask){</pre>
    dp[mask][-1] = A[mask]; //handle base case separately
        (leaf states)
    for(int i = 0; i < N; ++i){</pre>
        if (mask & (1<<i))</pre>
             dp[mask][i] = dp[mask][i-1] +
                dp[mask^(1<<i)][i-1];</pre>
        else dp[mask][i] = dp[mask][i-1];
    F[mask] = dp[mask][N-1];
}
//memory optimized, super easy to code.
for(int i = 0; i < (1 << N); ++i) F[i] = A[i];
for(int i = 0; i < N; ++i) for(int mask = 0; mask < (1<<N);
   ++mask){
    if (mask & (1<<i))</pre>
        F[mask] += F[mask^(1<<i)];
}
```

## 5.2 Convex Hull Trick (Rafael)

```
// linear

struct CHT {
    int it;
    vector<ll> a, b;
    CHT():it(0){}
    ll eval(int i, ll x){
        return a[i]*x + b[i];
    }
    bool useless(){
```

```
int sz = a.size();
        int r = sz-1, m = sz-2, 1 = sz-3;
        return (b[1] - b[r])*(a[m] - a[1]) <
            (b[1] - b[m])*(a[r] - a[1]);
    }
    void add(ll A, ll B){
        a.push_back(A); b.push_back(B);
        while (!a.empty()){
            if ((a.size() < 3) || !useless()) break;</pre>
            a.erase(a.end() - 2);
            b.erase(b.end() - 2);
        }
    }
    ll get(ll x){
        it = min(it, int(a.size()) - 1);
        while (it+1 < a.size()){</pre>
            if (eval(it+1, x) > eval(it, x)) it++;
            else break;
        return eval(it, x);
    }
};
```

## 5.3 Mochila

```
void solve(vector<int>& ans, int 1, int r, int cap) {
    if (1 == r) {
        if (w[1] <= cap) ans.push_back(1);</pre>
        return;
    int m = (1+r)/2;
    get_dp(0, 1, m, cap), get_dp(1, m+1, r, cap);
    int left_cap = -1, opt = -INF;
    for (int j = 0; j \le cap; j++)
        if (int at = dp[0][j] + dp[1][cap - j]; at > opt)
            opt = at, left_cap = j;
    solve(ans, 1, m, left_cap), solve(ans, m+1, r, cap -
       left_cap);
}
vector < int > knapsack(int n, int cap) {
    vector < int > ans;
    solve(ans, 0, n-1, cap);
    return ans;
}
```

## 5.4 Divide and Conquer DP

```
// Tudo 1-based!!!
// Particiona o array em k subarrays
// maximizando o somatorio das queries
//
// O(k n log n), assumindo quer query(l, r) eh O(1)

typedef long long ll;

ll dp[MAX][2];

void solve(int k, int l, int r, int lk, int rk) {
   if (l > r) return;
   int m = (l+r)/2, p = -1;
   ll& ans = dp[m][k&1] = -LINF;
   // ans = dp[m][~k&1], p = m+1; // para intervalos vazios
   for (int i = lk; i <= min(rk, m); i++) {</pre>
```

```
ll at = dp[i-1][~k&1] + query(i, m);
    if (at > ans) ans = at, p = i;
}
solve(k, l, m-1, lk, p), solve(k, m+1, r, p, rk);
}

ll DC(int n, int k) {
    dp[0][0] = dp[0][1] = 0;
    // garante que todo mundo pertence a algum intervalo
    for (int i = 1; i <= n; i++) dp[i][0] = -LINF;
    // se puder usar inervalos vazios, usar solve(i, 1, n, 1, n)
    for (int i = 1; i <= k; i++) solve(i, i, n, i, n);
    return dp[n][k&1];
}</pre>
```

#### 5.5 Longest Common Subsequence

```
// Computa a LCS entre dois arrays usando
// o algoritmo de Hirschberg para recuperar
//
// O(n*m), O(n+m) de memoria
int lcs_s[MAX], lcs_t[MAX];
int dp[2][MAX];
// dp[0][j] = max lcs(s[li...ri], t[lj, lj+j])
void dp_top(int li, int ri, int lj, int rj) {
    memset(dp[0], 0, (rj-lj+1)*sizeof(dp[0][0]));
    for (int i = li; i <= ri; i++) {</pre>
        for (int j = rj; j >= lj; j--)
            dp[0][i - 1i] = max(dp[0][i - 1i],
            (lcs_s[i] == lcs_t[j]) + (j > lj ? dp[0][j-1 -
               li] : 0));
        for (int j = lj+1; j <= rj; j++)</pre>
            dp[0][j-1j] = max(dp[0][j-1j], dp[0][j-1
               -lj]);
    }
}
```

```
// dp[1][j] = max lcs(s[li...ri], t[lj+j, rj])
void dp_bottom(int li, int ri, int lj, int rj) {
    memset(dp[1], 0, (rj-lj+1)*sizeof(dp[1][0]));
    for (int i = ri; i >= li; i--) {
        for (int j = lj; j <= rj; j++)</pre>
            dp[1][i - 1i] = max(dp[1][i - 1i],
            (lcs_s[i] == lcs_t[j]) + (j < rj ? dp[1][j+1 -
               li]: 0));
        for (int j = rj-1; j >= lj; j--)
            dp[1][j-1j] = max(dp[1][j-1j], dp[1][j+1-
               li]);
   }
}
void solve(vector<int>& ans, int li, int ri, int lj, int rj)
    if (li == ri){
        for (int j = lj; j <= rj; j++)</pre>
            if (lcs_s[li] == lcs_t[j]){
                ans.push_back(lcs_t[j]);
                break:
        return;
    if (lj == rj){
        for (int i = li; i <= ri; i++){</pre>
            if (lcs_s[i] == lcs_t[li]){
                ans.push_back(lcs_s[i]);
                break:
            }
        }
        return;
    }
    int mi = (li+ri)/2;
    dp_top(li, mi, lj, rj), dp_bottom(mi+1, ri, lj, rj);
    int j_{-} = 0, mx = -1;
    for (int j = lj-1; j <= rj; j++) {
        int val = 0;
        if (j >= lj) val += dp[0][j - lj];
        if (j < rj) val += dp[1][j+1 - lj];
```

#### 6 Problemas

#### 6.1 Sweep Direction

```
// Passa por todas as ordenacoes dos pontos definitas por
   "direcoes"
// Assume que nao existem pontos coincidentes
// O(n^2 \log n)
void sweep_direction(vector<pt> v) {
    int n = v.size();
    sort(v.begin(), v.end(), [](pt a, pt b) {
        if (a.x != b.x) return a.x < b.x;
        return a.y > b.y;
    });
    vector < int > at(n);
    iota(at.begin(), at.end(), 0);
    vector < ii > swapp;
    for (int i = 0; i < n; i++) for (int j = i+1; j < n; j++)
        swapp.push_back({i, j}), swapp.push_back({j, i});
    sort(swapp.begin(), swapp.end(), [&](ii a, ii b) {
```

```
pt A = rotate90(v[a.f] - v[a.s]);
        pt B = rotate90(v[b.f] - v[b.s]);
        if (quad(A) == quad(B) and !sarea2(pt(0, 0), A, B))
           return a < b;</pre>
        return compare_angle(A, B);
    });
    for (auto par : swapp) {
        assert(abs(at[par.f] - at[par.s]) == 1);
        int l = min(at[par.f], at[par.s]), r = n-1
           max(at[par.f], at[par.s]);
        // l e r sao quantos caras tem de cada lado do par
           de pontos
        // (cada par eh visitado duas vezes)
        swap(v[at[par.f]], v[at[par.s]]);
        swap(at[par.f], at[par.s]);
}
```

#### 6.2 Inversion Count

```
// Computa o numero de inversoes para transformar
// l em r (se nao tem como, retorna -1)
// O(n log(n))
template < typename T > 11 inv_count(vector < T > 1, vector < T > r =
   {}) {
    if (!r.size()) {
        r = 1;
        sort(r.begin(), r.end());
    }
    int n = 1.size();
    vector < int > v(n), bit(n);
    vector<pair<T, int>> w;
    for (int i = 0; i < n; i++) w.push_back({r[i], i+1});</pre>
    sort(w.begin(), w.end());
    for (int i = 0; i < n; i++) {</pre>
        auto it = lower_bound(w.begin(), w.end(),
            make_pair(l[i], 0));
        if (it == w.end() or it->first != l[i]) return -1;
```

```
// nao da
v[i] = it->second;
it->second = -1;
}

ll ans = 0;
for (int i = n-1; i >= 0; i--) {
   for (int j = v[i]-1; j; j -= j&-j) ans += bit[j];
   for (int j = v[i]; j < n; j += j&-j) bit[j]++;
}
return ans;
}</pre>
```

## 6.3 Gray Code

```
// Gera uma permutacao de 0 a 2^n-1, de forma que
// duas posicoes adjacentes diferem em exatamente 1 bit
//
// O(2^n)

vector<int> gray_code(int n) {
    vector<int> ret(1<<n);
    for (int i = 0; i < (1<<n); i++) ret[i] = i^(i>>1);
    return ret;
}
```

## 6.4 Triangulos em Grafos

```
// get_triangles(i) encontra todos os triangulos ijk no grafo
// Custo nas arestas
// retorna {custo do triangulo, {j, k}}
//
// O(m sqrt(m) log(n)) se chamar para todos os vertices
vector<ii> g[MAX]; // {para, peso}

#warning o 'g' deve estar ordenado
vector<pair<int, ii>> get_triangles(int i) {
```

## 6.5 RMQ com Divide and Conquer

```
// Responde todas as queries em
// O(n log(n))
typedef pair<pair<int, int>, int> iii;
#define f first
#define s second
int n, q, v[MAX];
iii qu[MAX];
int ans[MAX], pref[MAX], sulf[MAX];
void solve(int l=0, int r=n-1, int ql=0, int qr=q-1) {
    if (1 > r or q1 > qr) return;
    int m = (1+r)/2;
    int qL = partition(qu+ql, qu+qr+1, [=](iii x){return
       x.f.s < m;) - qu;
    int qR = partition(qu+qL, qu+qr+1, [=](iii x){return
       x.f.f <= m; }) - qu;
    pref(m) = sulf(m) = v(m);
    for (int i = m-1; i >= 1; i--) pref[i] = min(v[i],
       pref[i+1]);
    for (int i = m+1; i <= r; i++) sulf[i] = min(v[i],</pre>
```

```
sulf[i-1]);
    for (int i = qL; i < qR; i++)</pre>
        ans[qu[i].s] = min(pref[qu[i].f.f], sulf[qu[i].f.s]);
    solve(1, m-1, ql, qL-1), solve(m+1, r, qR, qr);
}
6.6 MO - DSU
// Dado uma lista de arestas de um grafo, desponde
// para cada query(1, r), quantos componentes conexos
// o grafo tem se soh considerar as arestas l, l+1, ..., r
// Da pra adaptar pra usar MO com qualquer estrutura
   rollbackavel
//
// O(m \ sqrt(m) \log(n))
struct dsu {
    int n, ans;
    vector < int > p, sz;
    stack<int> S;
    dsu(int n_{-}) : n(n_{-}), ans(n), p(n), sz(n) {
        for (int i = 0; i < n; i++) p[i] = i, sz[i] = 1;</pre>
    }
    int find(int k) {
        while (p[k] != k) k = p[k];
        return k;
    }
    void add(pair<int, int> x) {
        int a = x.first, b = x.second;
        a = find(a), b = find(b);
        if (a == b) return S.push(-1);
        ans --;
        if (sz[a] > sz[b]) swap(a, b);
        S.push(a);
        sz[b] += sz[a];
        p[a] = b;
```

```
int query() { return ans; }
    void rollback() {
        int u = S.top(); S.pop();
        if (u == -1) return;
        sz[p[u]] -= sz[u];
        p[u] = u;
        ans++;
};
int n;
vector<ii> ar; // vetor com as arestas
vector<int> MO(vector<ii> &q) {
    int SQ = sqrt(q.size()) + 1;
    int m = q.size();
    vector < int > ord(m);
    iota(ord.begin(), ord.end(), 0);
    sort(ord.begin(), ord.end(), [&](int 1, int r) {
        if (q[1].first / SQ != q[r].first / SQ) return
            q[l].first < q[r].first;
        return q[1].second < q[r].second;</pre>
    });
    vector < int > ret(m);
    for (int i = 0; i < m; i++) {</pre>
        dsu D(n);
        int fim = q[ord[i]].first/SQ*SQ + SQ - 1;
        int last_r = fim;
        int j = i-1;
        while (j+1 < m and q[ord[j+1]].first / SQ ==</pre>
            q[ord[i]].first / SQ) {
            auto [1, r] = q[ord[++j]];
            if (1 / SQ == r / SQ) {
                 dsu D2(n):
                for (int k = 1; k <= r; k++) D2.add(ar[k]);</pre>
                ret[ord[j]] = D2.query();
                 continue;
            }
             while (last_r < r) D.add(ar[++last_r]);</pre>
```

```
for (int k = 1; k <= fim; k++) D.add(ar[k]);</pre>
        ret[ord[i]] = D.query();
        for (int k = 1; k <= fim; k++) D.rollback();</pre>
    }
    i = j;
return ret;
```

#### 6.7 LIS2 - Longest Increasing Subsequence

```
// Calcula o tamanho da LIS
// O(n.log(n))
template < typename T > int lis(vector < T > &v) {
    vector <T> ans;
    for (T t : v){
        // Para non-decreasing use upper_bound()
        auto it = lower_bound(ans.begin(), ans.end(), t);
        if (it == ans.end()) ans.push_back(t);
        else *it = t;
    }
    return ans.size();
}
```

#### 6.8 Nim

```
// Calcula movimento otimo do jogo classico de Nim
// Assume que o estado atual eh perdedor
// Funcao move retorna um par com a pilha (O indexed)
// e quanto deve ser tirado dela
// XOR deve estar armazenado em x
// Para mudar um valor, faca insere(novo_valor),
// atualize o XOR e mude o valor em v
```

```
// MAX2 = teto do log do maior elemento
// possivel nas pilhas
// O(log(n)) amortizado
int v[MAX], n, x;
stack<int> pi[MAX2];
void insere(int p) {
    for (int i = 0; i < MAX2; i++) if (v[p] & (1 << i))</pre>
       pi[i].push(p);
}
pair < int , int > move() {
    int bit = 0; while (x >> bit) bit++; bit--;
    // tira os caras invalidos
    while ((v[pi[bit].top()] & (1 << bit)) == 0)</pre>
       pi[bit].pop();
    int cara = pi[bit].top();
    int tirei = v[cara] - (x^v[cara]);
    v[cara] -= tirei;
    insere(cara);
    return make_pair(cara, tirei);
}
// Acha o movimento otimo baseado
// em v apenas
//
// O(n)
pair < int , int > move() {
    int x = 0;
    for (int i = 0; i < n; i++) x ^= v[i];</pre>
    for (int i = 0; i < n; i++) if ((v[i]^x) < v[i])
        return make_pair(i, v[i] - (v[i]^x));
}
```

#### 6.9 Arpa's Trick

```
// Responde RMQ em O((n+q)\log(n)) offline
// Adicionar as queries usando arpa::add(a, b)
// A resposta vai ta em ans[], na ordem que foram colocadas
int n, v[MAX], ans[MAX];
namespace arpa {
    int p[MAX], cnt;
    stack<int> s;
    vector<pair<int, int> > 1[MAX];
   int find(int k) { return p[k] == k ? k : p[k] =
       find(p[k]); }
   void add(int a, int b) { l[b].push_back({a, cnt++}); }
    void solve() {
       for (int i = 0; (p[i]=i) < n; s.push(i++)) {
            while (s.size() and v[s.top()] >= v[i])
               p[s.top()] = i, s.pop();
            for (auto q : l[i]) ans[q.second] =
               v[find(q.first)];
       }
   }
}
6.10 Simple Polygon
// Verifica se um poligono com n pontos eh simples
//
// O(n log n)
bool operator < (const line& a, const line& b) { //
   comparador pro sweepline
   if (a.p == b.p) return ccw(a.p, a.q, b.q);
   if (!eq(a.p.x, a.q.x) and (eq(b.p.x, b.q.x) or a.p.x+eps
       < b.p.x))
       return ccw(a.p, a.q, b.p);
   return ccw(a.p, b.q, b.p);
```

```
bool simple(vector<pt> v) {
    auto intersects = [&](pair<line, int> a, pair<line, int>
        if ((a.s+1)%v.size() == b.s or (b.s+1)%v.size() ==
           a.s) return false;
        return interseg(a.f, b.f);
    };
    vector<line> seg;
    vector<pair<pt, pair<int, int>>> w;
    for (int i = 0; i < v.size(); i++) {</pre>
        pt at = v[i], nxt = v[(i+1)%v.size()];
        if (nxt < at) swap(at, nxt);</pre>
        seg.push_back(line(at, nxt));
        w.push_back({at, {0, i}});
        w.push_back({nxt, {1, i}});
        // casos degenerados estranhos
        if (isinseg(v[(i+2)%v.size()], line(at, nxt)))
           return 0:
        if (isinseg(v[(i+v.size()-1)%v.size()], line(at,
           nxt))) return 0;
    sort(w.begin(), w.end());
    set < pair < line , int >> se;
    for (auto i : w) {
        line at = seg[i.s.s];
        if (i.s.f == 0) {
            auto nxt = se.lower_bound({at, i.s.s});
            if (nxt != se.end() and intersects(*nxt, {at,
               i.s.s})) return 0:
            if (nxt != se.begin() and intersects(*(--nxt),
               {at, i.s.s})) return 0;
            se.insert({at, i.s.s});
        } else {
            auto nxt = se.upper_bound({at, i.s.s}), cur =
               nxt, prev = --cur;
            if (nxt != se.end() and prev != se.begin()
                and intersects(*nxt, *(--prev))) return 0;
            se.erase(cur);
    }
    return 1;
```

## 6.11 Segment Intersection

}

```
// Verifica, dado n segmentos, se existe algum par de
   segmentos
// que se intersecta
//
// O(n log n)
bool operator < (const line& a, const line& b) { //</pre>
   comparador pro sweepline
    if (a.p == b.p) return ccw(a.p, a.q, b.q);
    if (!eq(a.p.x, a.q.x) and (eq(b.p.x, b.q.x) or a.p.x+eps
       < b.p.x))
        return ccw(a.p, a.q, b.p);
    return ccw(a.p, b.q, b.p);
}
bool has_intersection(vector<line> v) {
    auto intersects = [&](pair<line, int> a, pair<line, int>
       b) {
        return interseg(a.f, b.f);
    vector<pair<pt, pair<int, int>>> w;
    for (int i = 0; i < v.size(); i++) {</pre>
        if (v[i].q < v[i].p) swap(v[i].p, v[i].q);</pre>
        w.push_back({v[i].p, {0, i}});
        w.push_back({v[i].q, {1, i}});
    }
    sort(w.begin(), w.end());
    set < pair < line, int >> se;
    for (auto i : w) {
        line at = v[i.s.s];
        if (i.s.f == 0) {
            auto nxt = se.lower_bound({at, i.s.s});
            if (nxt != se.end() and intersects(*nxt, {at,
               i.s.s})) return 1;
            if (nxt != se.begin() and intersects(*(--nxt),
                {at, i.s.s})) return 1;
```

```
se.insert({at, i.s.s});
} else {
    auto nxt = se.upper_bound({at, i.s.s}), cur =
        nxt, prev = --cur;
    if (nxt != se.end() and prev != se.begin()
        and intersects(*nxt, *(--prev))) return 1;
    se.erase(cur);
}
return 0;
}
```

#### 6.12 Conectividade Dinamica

```
// Offline com Divide and Conquer e
// DSU com rollback
// O(n log^2(n))
typedef pair<int, int> T;
namespace data {
    int n, ans;
    int p[MAX], sz[MAX];
    stack<int> S;
    void build(int n2) {
        n = n2;
        for (int i = 0; i < n; i++) p[i] = i, sz[i] = 1;
        ans = n;
    }
    int find(int k) {
        while (p[k] != k) k = p[k];
        return k;
    }
    void add(T x) {
        int a = x.first, b = x.second;
        a = find(a), b = find(b);
        if (a == b) return S.push(-1);
        ans --;
        if (sz[a] > sz[b]) swap(a, b);
```

```
S.push(a);
        sz[b] += sz[a];
        p[a] = b;
    }
    int query() {
        return ans;
    }
    void rollback() {
        int u = S.top(); S.pop();
        if (u == -1) return;
        sz[p[u]] -= sz[u];
        p[u] = u;
        ans++;
   }
};
int ponta[MAX]; // outra ponta do intervalo ou -1 se for
   query
int ans[MAX], n, q;
T qu[MAX];
void solve(int l = 0, int r = q-1) {
    if (1 >= r) {
        ans[1] = data::query(); // agora a estrutura ta certa
        return;
    }
    int m = (1+r)/2, qnt = 1;
    for (int i = m+1; i <= r; i++) if (ponta[i]+1 and
       ponta[i] < 1)</pre>
       data::add(qu[i]), qnt++;
    solve(1, m);
    while (--qnt) data::rollback();
   for (int i = 1; i <= m; i++) if (ponta[i]+1 and ponta[i]
       > r)
        data::add(qu[i]), qnt++;
    solve(m+1, r);
    while (qnt--) data::rollback();
}
```

#### 6.13 Distinct Range Query com Update

```
// build - O(n log(n))
// \text{ query - O(log^2(n))}
// update - O(log^2(n))
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
template <class T>
    using ord_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
int v[MAX], n, nxt[MAX], prv[MAX];
map < int , set < int > > ocor;
namespace bit {
    ord_set < ii > bit [MAX];
    void build() {
        for (int i = 1; i <= n; i++)</pre>
            bit[i].insert({nxt[i-1], i-1});
        for (int i = 1; i <= n; i++) {</pre>
            int j = i + (i\&-i);
            if (j <= n) for (auto x : bit[i])</pre>
                bit[j].insert(x);
        }
    int pref(int p, int x) {
        int ret = 0;
        for (; p; p -= p\&-p) ret += bit[p].order_of_key({x,}
            -INF }):
        return ret;
    int query(int 1, int r, int x) {
        return pref(r+1, x) - pref(1, x);
    void update(int p, int x) {
        int p2 = p;
        for (p++; p <= n; p += p&-p) {</pre>
            bit[p].erase({nxt[p2], p2});
             bit[p].insert({x, p2});
        }
    }
```

```
}
void build() {
    for (int i = 0; i < n; i++) nxt[i] = INF;</pre>
    for (int i = 0; i < n; i++) prv[i] = -INF;</pre>
    vector<ii> t;
    for (int i = 0; i < n; i++) t.push_back({v[i], i});</pre>
    sort(t.begin(), t.end());
    for (int i = 0; i < n; i++) {</pre>
        if (i and t[i].f == t[i-1].f) prv[t[i].s] = t[i-1].s;
        if (i+1 < n \text{ and } t[i].f == t[i+1].f) nxt[t[i].s] =
            t[i+1].s:
    }
    for (int i = 0; i < n; i++) ocor[v[i]].insert(i);</pre>
    bit::build();
}
void muda(int p, int x) {
    bit::update(p, x);
    nxt[p] = x;
}
int query(int a, int b) {
    return b-a+1 - bit::query(a, b, b+1);
}
void update(int p, int x) { // mudar valor na pos. p para x
    if (prv[p] > -INF) muda(prv[p], nxt[p]);
    if (nxt[p] < INF) prv[nxt[p]] = prv[p];</pre>
    ocor[v[p]].erase(p);
    if (!ocor[x].size()) {
        muda(p, INF);
        prv[p] = -INF;
    } else if (*ocor[x].rbegin() < p) {</pre>
        int i = *ocor[x].rbegin();
        prv[p] = i;
        muda(p, INF);
        muda(i, p);
    } else {
```

```
int i = *ocor[x].lower_bound(p);
if (prv[i] > -INF) {
         muda(prv[i], p);
         prv[p] = prv[i];
     } else prv[p] = -INF;
     prv[i] = p;
     muda(p, i);
}
v[p] = x; ocor[x].insert(p);
}
```

#### 6.14 LIS - Longest Increasing Subsequence

```
// Calcula e retorna uma LIS
// O(n.log(n))
template < typename T > vector <T > lis(vector <T > & v) {
    int n = v.size(), m = -1;
    vector <T> d(n+1, INF);
    vector < int > 1(n);
    d[0] = -INF;
    for(int i=0; i<n; i++) {</pre>
        // Para non-decreasing use upper_bound()
        int t = lower_bound(d.begin(), d.end(), v[i]) -
            d.begin();
        d[t] = v[i], l[i] = t, m = max(m, t);
    }
    int p = n;
    vector <T> ret;
    while (p--) if (l[p] == m) {
        ret.push_back(v[p]);
        m - - ;
    reverse(ret.begin(),ret.end());
    return ret;
}
```

### 6.15 Algoritmo Hungaro

```
// Resolve o problema de assignment (matriz n x n)
// Colocar os valores da matriz em 'a' (pode < 0)</pre>
// assignment() retorna um par com o valor do
// assignment minimo, e a coluna escolhida por cada linha
// O(n^3)
template < typename T > struct hungarian {
    int n:
    vector < vector < T >> a;
    vector <T> u, v;
    vector < int > p, way;
    T inf;
    hungarian(int n_{-}): n(n_{-}), u(n+1), v(n+1), p(n+1),
        wav(n+1) {
        a = vector < vector < T >> (n, vector < T > (n));
        inf = numeric_limits <T>::max();
    }
    pair <T, vector <int >> assignment() {
        for (int i = 1; i <= n; i++) {
             p[0] = i;
             int j0 = 0;
             vector <T> minv(n+1, inf);
             vector < int > used(n+1, 0);
             do {
                 used[j0] = true;
                 int i0 = p[j0], j1 = -1;
                 T delta = inf;
                 for (int j = 1; j <= n; j++) if (!used[j]) {</pre>
                     T cur = a[i0-1][j-1] - u[i0] - v[j];
                     if (cur < minv[j]) minv[j] = cur, way[j]</pre>
                     if (minv[j] < delta) delta = minv[j], j1</pre>
                         = j;
                 for (int j = 0; j <= n; j++)
                     if (used[j]) u[p[j]] += delta, v[j] -=
                         delta;
                     else minv[j] -= delta;
```

```
j0 = j1;
} while (p[j0] != 0);
do {
    int j1 = way[j0];
    p[j0] = p[j1];
    j0 = j1;
} while (j0);
}
vector<int> ans(n);
for (int j = 1; j <= n; j++) ans[p[j]-1] = j-1;
return make_pair(-v[0], ans);
}
};</pre>
```

#### 6.16 Mo algorithm - distinct values

```
// O(s*n*f + q*(n/s)*f) optimize over s, insert/erase = O(f)
// for s = sqrt(n), O((n+q)*sqrt(n)*f)
const int MAX = 3e4+10;
const int SQ = sqrt(MAX);
int v[MAX];
int ans, freq[MAX];
inline void insert(int p) {
    int o = v[p];
    freq[o]++;
    ans += (freq[o] == 1);
}
inline void erase(int p) {
    int o = v[p];
    ans -= (freq[o] == 1);
    freq[o]--;
}
inline 11 hilbert(int x, int y) {
    static int N = (1 << 20);
    int rx, ry, s;
```

```
11 d = 0:
    for (s = N/2; s>0; s /= 2) {
        rx = (x \& s) > 0;
        ry = (y \& s) > 0;
        d += s * 11(s) * ((3 * rx) ^ ry);
        if (ry == 0) {
            if (rx == 1) {
                x = N-1 - x;
                y = N-1 - y;
            swap(x, y);
        }
    }
    return d;
}
#define HILBERT true
vector<int> MO(vector<ii> &q) {
    ans = 0;
    int m = q.size();
    vector < int > ord(m);
    iota(ord.begin(), ord.end(), 0);
#if HILBERT
    vector < ll > h(m);
    for (int i = 0; i < m; i++) h[i] = hilbert(q[i].first,</pre>
       q[i].second);
    sort(ord.begin(), ord.end(), [&](int 1, int r) { return
       h[1] < h[r]; \});
#else
    sort(ord.begin(), ord.end(), [&](int 1, int r) {
        if (q[1].first / SQ != q[r].first / SQ) return
           q[l].first < q[r].first;
        if ((q[1].first / SQ) % 2) return q[1].second >
           q[r].second;
        return q[1].second < q[r].second;</pre>
    });
#endif
    vector < int > ret(m);
    int 1 = 0, r = -1;
    for (int i : ord) {
        int ql, qr;
```

```
tie(q1, qr) = q[i];
while (r < qr) insert(++r);
while (l > ql) insert(--l);
while (l < ql) erase(l++);
while (r > qr) erase(r--);
ret[i] = ans;
}
return ret;
}
```

#### 6.17 Binomial modular

```
// Computa C(n, k) mod m em O(m + log(m) log(n))
// = O(rapido)
11 divi[MAX];
ll expo(ll a, ll b, ll m) {
    if (!b) return 1;
    ll ans = expo(a*a\%m, b/2, m);
    if (b\%2) ans *= a;
    return ans%m;
}
11 inv(ll a, ll b){
    return 1<a ? b - inv(b%a,a)*b/a : 1;
}
ll gcde(ll a, ll b, ll& x, ll& y) {
    if (!a) {
        x = 0:
        y = 1;
        return b;
    }
    11 X, Y;
    ll g = gcde(b \% a, a, X, Y);
    x = Y - (b / a) * X;
    y = X;
```

```
return g:
}
struct crt {
    ll a, m;
    crt(ll a_, ll m_) : a(a_), m(m_) {}
    crt operator * (crt C) {
        11 x, y;
        ll g = gcde(m, C.m, x, y);
        if ((a - C.a) \% g) a = -1;
        if (a == -1 \text{ or } C.a == -1) \text{ return } crt(-1, 0);
        11 lcm = m/g*C.m;
        ll ans = a + (x*(C.a-a)/g \% (C.m/g))*m;
        return crt((ans % lcm + lcm) % lcm, lcm);
    }
};
pair < 11, 11 > divide_show(11 n, int p, int k, int pak) {
    if (n == 0) return {0, 1};
    11 blocos = n/pak, falta = n%pak;
    ll periodo = divi[pak], resto = divi[falta];
    ll r = expo(periodo, blocos, pak)*resto%pak;
    auto rec = divide_show(n/p, p, k, pak);
    ll v = n/p + rec.f;
    r = r*rec.s \% pak;
    return {y, r};
}
11 solve_pak(ll n, ll x, int p, int k, int pak) {
    divi[0] = 1;
    for (int i = 1; i <= pak; i++) {</pre>
        divi[i] = divi[i-1];
        if (i%p) divi[i] = divi[i] * i % pak;
    }
    auto dn = divide_show(n, p, k, pak), dx = divide_show(x,
       p, k, pak),
                             dnx = divide\_show(n-x, p, k,
                                 pak);
```

```
11 y = dn.f-dx.f-dnx.f, r = (dn.s*inv(dx.s,
       pak)%pak)*inv(dnx.s, pak)%pak;
    return expo(p, y, pak) * r % pak;
}
ll solve(ll n, ll x, int mod) {
    vector<ii> f;
    int mod2 = mod;
    for (int i = 2; i*i <= mod2; i++) if (mod2%i==0) {</pre>
        int c = 0;
        while (mod2\%i==0) mod2 /= i, c++;
        f.pb({i, c});
    }
    if (mod2 > 1) f.pb(\{mod2, 1\});
    crt ans(0, 1);
    for (int i = 0; i < f.size(); i++) {</pre>
        int pak = 1;
        for (int j = 0; j < f[i].s; j++) pak *= f[i].f;
        ans = ans * crt(solve_pak(n, x, f[i].f, f[i].s,
           pak), pak);
    }
    return ans.a;
}
```

#### 6.18 Colocacao de Grafo de Intervalo

```
// Colore os intervalos com o numero minimo
// de cores de tal forma que dois intervalos
// que se interceptam tem cores diferentes
// As cores vao de 1 ate n
//
// O(n log(n))

vector<int> coloring(vector<pair<int, int>>& v) {
   int n = v.size();
   vector<pair<int, pair<int, int>>> ev;
   for (int i = 0; i < n; i++) {
       ev.push_back({v[i].first, {1, i}});
       ev.push_back({v[i].second, {0, i}});
   }</pre>
```

```
sort(ev.begin(), ev.end());
    vector < int > ans(n), avl(n);
   for (int i = 0; i < n; i++) avl.push_back(n-i);</pre>
    for (auto i : ev) {
        if (i.second.first == 1) {
            ans[i.second.second] = avl.back();
            avl.pop_back();
        } else avl.push_back(ans[i.second.second]);
    }
    return ans;
}
6.19 Distinct Range Query - Wavelet
// build - O(n (log n + log(sigma)))
// query - O(log(sigma))
int v[MAX], n, nxt[MAX];
namespace wav {
    vector < int > esq[4*(1+MAXN-MINN)];
    void build(int b = 0, int e = n, int p = 1, int l =
       MINN, int r = MAXN) {
        if (1 == r) return;
        int m = (1+r)/2; esq[p].push_back(0);
        for (int i = b; i < e; i++)</pre>
           esq[p].push_back(esq[p].back()+(nxt[i]<=m));</pre>
        int m2 = stable_partition(nxt+b, nxt+e, [=](int
           i) {return i <= m;}) - nxt;
        build(b, m2, 2*p, 1, m), build(m2, e, 2*p+1, m+1, r);
    }
    int count(int i, int j, int x, int y, int p = 1, int l =
       MINN, int r = MAXN) {
        if (y < 1 or r < x) return 0;
        if (x <= l and r <= y) return j-i;</pre>
        int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
        return count(ei, ej, x, y, 2*p, 1, m)+count(i-ei,
           j-ej, x, y, 2*p+1, m+1, r);
```

```
}

void build() {
    for (int i = 0; i < n; i++) nxt[i] = MAXN+1;
    vector < ii > t;
    for (int i = 0; i < n; i++) t.push_back({v[i], i});
    sort(t.begin(), t.end());
    for (int i = 0; i < n-1; i++) if (t[i].f == t[i+1].f)
        nxt[t[i].s] = t[i+1].s;

wav::build();
}

int query(int a, int b) {
    return wav::count(a, b+1, b+1, MAXN+1);
}
</pre>
```

## 6.20 Area da Uniao de Retangulos

```
// O(n log(n))
const int MAX = 1e5+10;
namespace seg {
    pair < int , 11 > seg [4*MAX];
    ll lazy[4*MAX], *v;
    int n;
    pair<int, ll> merge(pair<int, ll> l, pair<int, ll> r){
        if (1.second == r.second) return {1.first+r.first,
           1.second}:
        else if (l.second < r.second) return l;</pre>
        else return r;
    }
    pair < int, 11 > build(int p=1, int 1=0, int r=n-1) {
        lazy[p] = 0;
        if (1 == r) return seg[p] = {1, v[1]};
        int m = (1+r)/2;
        return seg[p] = merge(build(2*p, 1, m), build(2*p+1,
```

```
m+1, r)):
    }
    void build(int n2, l1* v2) {
        n = n2, v = v2;
        build();
    }
    void prop(int p, int l, int r) {
        seg[p].second += lazy[p];
        if (1 != r) lazy[2*p] += lazy[p], lazy[2*p+1] +=
           lazy[p];
        lazy[p] = 0;
    }
    pair < int, ll > query (int a, int b, int p=1, int l=0, int
       r=n-1) {
        prop(p, 1, r);
        if (a <= 1 and r <= b) return seg[p];</pre>
        if (b < 1 or r < a) return {0, LINF};</pre>
        int m = (1+r)/2;
        return merge(query(a, b, 2*p, 1, m), query(a, b,
           2*p+1, m+1, r);
    }
    pair<int, 11> update(int a, int b, int x, int p=1, int
       1=0, int r=n-1) {
        prop(p, 1, r);
        if (a <= 1 and r <= b) {
            lazv[p] += x;
            prop(p, 1, r);
            return seg[p];
        if (b < 1 or r < a) return seg[p];</pre>
        int m = (1+r)/2;
        return seg[p] = merge(update(a, b, x, 2*p, 1, m),
                update(a, b, x, 2*p+1, m+1, r));
    }
};
11 seg_vec[MAX];
ll area_sq(vector<pair<ii, ii>> &sq){
    vector<pair<ii, ii>> up;
    for (auto it : sq){
        int x1, y1, x2, y2;
```

```
tie(x1, y1) = it.first;
    tie(x2, y2) = it.second;
    up.push_back({{x1+1, 1}, {y1, y2}});
    up.push_back({{x2+1, -1}, {y1, y2}});
sort(up.begin(), up.end());
memset(seg_vec, 0, sizeof seg_vec);
11 H_MAX = MAX;
seg::build(H_MAX-1, seg_vec);
auto it = up.begin();
11 \text{ ans} = 0;
while (it != up.end()){
    11 L = (*it).first.first;
    while (it != up.end() && (*it).first.first == L){
        int x, inc, y1, y2;
        tie(x, inc) = it->first;
        tie(y1, y2) = it -> second;
        seg::update(y1+1, y2, inc);
        it++;
    }
    if (it == up.end()) break;
    11 R = (*it).first.first;
   11 W = R-L;
    auto jt = seg::query(0, H_MAX-1);
    11 H = H_MAX - 1;
   if (jt.second == 0) H -= jt.first;
    ans += W*H;
}
return ans;
```

#### 6.21 Closest pair of points

}

```
// O(nlogn)
pair < pt, pt > closest_pair_of_points(vector < pt > &v) {
    #warning changes v order
    int n = v.size();
    sort(v.begin(), v.end());
```

```
for (int i = 1; i < n; i++){</pre>
        if (v[i] == v[i-1]){
            return make_pair(v[i-1], v[i]);
    }
    auto cmp_y = [&](const pt &l, const pt &r){
        if (1.y != r.y) return 1.y < r.y;</pre>
        return l.x < r.x;</pre>
    };
    set < pt, decltype(cmp_y) > s(cmp_y);
    int 1 = 0, r = -1;
    11 d2_min = numeric_limits < ll >:: max();
    pt pl, pr;
    const int magic = 5;
    while (r+1 < n) {
        auto it = s.insert(v[++r]).first;
        int cnt = magic/2;
        while (cnt-- && it != s.begin())
             it--;
        cnt = 0;
        while (cnt++ < magic && it != s.end()){</pre>
             if (!((*it) == v[r])){
                 11 d2 = dist2(*it, v[r]);
                 if (d2_min > d2){
                     d2_min = d2;
                     pl = *it;
                     pr = v[r];
             }
             it++;
        while (1 < r \&\& sq(v[1].x-v[r].x) > d2_min)
            s.erase(v[1++]);
    }
    return make_pair(pl, pr);
}
```

## 6.22 Area Maxima de Histograma

// Assume que todas as barras tem largura 1,

```
// e altura dada no vetor v
// O(n)
typedef long long 11;
11 area(vector<int> v) {
    11 \text{ ret} = 0;
    stack<int> s;
    // valores iniciais pra dar tudo certo
    v.insert(v.begin(), -1);
    v.insert(v.end(), -1);
    s.push(0);
    for(int i = 0; i < (int) v.size(); i++) {</pre>
        while (v[s.top()] > v[i]) {
            ll h = v[s.top()]; s.pop();
            ret = max(ret, h * (i - s.top() - 1));
        }
        s.push(i);
    }
    return ret;
}
```

## Conj. Indep. Maximo com Peso em Grafo de Intervalo

```
// Retorna os indices ordenados dos
// intervalos selecionados
// Se tiver empate, retorna o que minimiza o comprimento
   total
// O(n log(n))
vector < int > ind_set(vector < tuple < int, int, int >>& v) {
    vector<tuple<int, int, int>> w;
    for (int i = 0; i < v.size(); i++) {</pre>
        w.push_back(tuple(get<0>(v[i]), 0, i));
        w.push_back(tuple(get<1>(v[i]), 1, i));
```

```
}
    sort(w.begin(), w.end());
    vector < int > nxt(v.size());
    vector < pair < ll, int >> dp(v.size());
    int last = -1;
    for (auto [fim, t, i] : w) {
        if (t == 0) {
            nxt[i] = last;
            continue;
        }
        dp[i] = \{0, 0\};
        if (last != -1) dp[i] = max(dp[i], dp[last]);
        pair<ll, int> pega = {get<2>(v[i]), -(get<1>(v[i]) -
           get < 0 > (v[i]) + 1);
        if (nxt[i] != -1) pega.first += dp[nxt[i]].first,
           pega.second += dp[nxt[i]].second;
        if (pega > dp[i]) dp[i] = pega;
        else nxt[i] = last;
        last = i;
    }
    pair<11, int> ans = {0, 0};
    int idx = -1;
    for (int i = 0; i < v.size(); i++) if (dp[i] > ans) ans
       = dp[i], idx = i;
    vector<int> ret;
    while (idx != -1) {
        if (get < 2 > (v[idx]) > 0 and
            (nxt[idx] == -1 or get<1>(v[nxt[idx]]) <</pre>
               get<0>(v[idx]))) ret.push_back(idx);
        idx = nxt[idx];
    sort(ret.begin(), ret.end());
    return ret:
}
6.24 Distinct Range Query - Persistent Segtree
```

```
// build - O(n (log n + log(sigma)))
// query - O(log(sigma))
```

```
const int MAX = 3e4+10, LOG = 20;
const int MAXS = 4*MAX+MAX*LOG;
namespace perseg {
    11 seg[MAXS];
    int rt[MAX], L[MAXS], R[MAXS], cnt, t;
    int n, *v;
    ll build(int p, int l, int r) {
        if (1 == r) return seg[p] = 0;
        L[p] = cnt++, R[p] = cnt++;
        int m = (1+r)/2;
        return seg[p] = build(L[p], 1, m) + build(R[p], m+1,
           r);
    }
    void build(int n2) {
        n = n2;
        rt[0] = cnt++;
        build(0, 0, n-1);
    }
    ll query(int a, int b, int p, int l, int r) {
        if (b < l or r < a) return 0;
        if (a <= 1 and r <= b) return seg[p];</pre>
        int m = (1+r)/2;
        return query(a, b, L[p], 1, m) + query(a, b, R[p],
           m+1, r);
    ll query(int a, int b, int tt) {
        return query(a, b, rt[tt], 0, n-1);
   }
    11 update(int a, int x, int lp, int p, int l, int r) {
        if (1 == r) return seg[p] = seg[lp] + x;
        int m = (1+r)/2:
        if (a \le m)
            return seg[p] = update(a, x, L[lp], L[p]=cnt++,
               1, m) + seg[R[p]=R[lp]];
        return seg[p] = seg[L[p]=L[lp]] + update(a, x,
           R[lp], R[p] = cnt++, m+1, r);
    }
    void update(int a, int x, int tt=t) {
```

```
update(a, x, rt[tt], rt[++t]=cnt++, 0, n-1);
   }
};
int qt[MAX];
void build(vector<int>& v) {
    int n = v.size();
    perseg::build(n);
    map<int, int> last;
    int at = 0;
    for (int i = 0; i < n; i++) {
        if (last.count(v[i])) {
            perseg::update(last[v[i]], -1);
            at++:
        perseg::update(i, 1);
        qt[i] = ++at;
        last[v[i]] = i;
   }
}
int query(int 1, int r) {
    return perseg::query(1, r, qt[r]);
}
6.25 Dominator Points
// Se um ponto A tem ambas as coordenadas >= B, dizemos
// que A domina B
// is_dominated(p) fala se existe algum ponto no conjunto
// que domina p
// insert(p) insere p no conjunto
// (se p for dominado por alguem, nao vai inserir)
//
// Complexidades:
// is_dominated - O(log(n))
// insert - O(log(n)) amortizado
```

struct dominator\_points {

```
set < pair < int , int >> se;
    dominator_points() {}
    bool is_dominated(pair<int, int> p) {
        auto it = se.lower_bound(p);
        if (it == se.end()) return 0;
        return it->second >= p.second;
    bool insert(pair<int, int> p) {
        if (is_dominated(p)) return 0;
        auto it = se.lower_bound(p);
        while (it != se.begin()) {
            it--;
            if (it->second > p.second) break;
            it = se.erase(it):
        se.insert(p);
        return 1;
};
```

## 6.26 Mo algorithm - DQUERY path on trees

```
int ans, freq[MAX], freqv[MAX];
void dfs(int i, int p, int &t){
    v[t] = i;
    st[i] = t++;
   for (int j : g[i]){
        if (j == p) continue;
        dfs(j, i, t);
    }
    v[t] = i;
    en[i] = t++;
}
void update(int o){//only change this function
    if (freqv[o] == 1){//insert w[o]
        ans += (freq[w[o]] == 0);
        freq[w[o]]++;
    }
    if (freqv[o] != 1){//erase w[o]
        ans -= (freq[w[o]] == 1);
        freq[w[o]]--;
    }
}
void insert(int p){
    int o = v[p];
    freqv[o]++;
    update(o);
}
void erase(int p){
    int o = v[p];
   freqv[o]--;
    update(o);
}
vector<tuple<int, int, int>> make_queries(vector<ii>> &q_){
    vector<tuple<int, int, int>> q;
    for (auto &it : q_){
        int 1, r;
```

```
tie(1, r) = it:
        if (st[r] < st[l]) swap(l, r);</pre>
        int p = LCA::lca(1, r);
        int init = (p == 1) ? st[1] : en[1];
        q.push_back({init, st[r], st[p]});
    return q;
}
vector<int> MO(vector<ii> &q_){
    LCA::build(0);//any LCA alg works
    int t = 0:
    dfs(0, -1, t);
    auto q = make_queries(q_);
    ans = 0:
    memset(freq, 0, sizeof freq);
    memset(freqv, 0, sizeof freqv);
    int m = q.size();
    vector < int > ord(m), ret(m);
    iota(ord.begin(), ord.end(), 0);
    sort(ord.begin(), ord.end(), [&](int 1, int r){
        int sl = get < 0 > (q[1])/SQ;
        int sr = get<0>(q[r])/SQ;
        if (sl != sr) return sl < sr;</pre>
        return get<1>(q[1]) < get<1>(q[r]);
    });
    int 1 = 0, r = 0;
    insert(0):
    for (int i : ord){
        int ql, qr, qp;
        tie(ql, qr, qp) = q[i];
        while (r < qr) insert(++r);</pre>
        while (1 > q1) insert(--1);
        while (1 < q1) erase(1++);</pre>
        while (r > qr) erase(r--);
        if (qp < 1 || qp > r){
            //lca out of range
            insert(qp);
```

## 6.27 Mininum Enclosing Circle

```
// O(n) com alta probabilidade
const double EPS = 1e-12;
mt19937 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
struct pt {
    double x, y;
    pt(double x_{-} = 0, double y_{-} = 0) : x(x_{-}), y(y_{-}) {}
    pt operator + (const pt& p) const { return pt(x+p.x,
       (v.q+v); }
   pt operator - (const pt& p) const { return pt(x-p.x,
       { ; (v.q-v
    pt operator * (double c) const { return pt(x*c, y*c); }
   pt operator / (double c) const { return pt(x/c, y/c); }
};
double dot(pt p, pt q) { return p.x*q.x+p.y*q.y; }
double cross(pt p, pt q) { return p.x*q.y-p.y*q.x; }
double dist(pt p, pt q) { return sqrt(dot(p-q, p-q)); }
pt center(pt p, pt q, pt r) {
    pt a = p-r, b = q-r;
   pt c = pt(dot(a, p+r)/2, dot(b, q+r)/2);
   return pt(cross(c, pt(a.y, b.y)), cross(pt(a.x, b.x),
       c)) / cross(a, b);
}
struct circle {
    pt cen;
```

```
double r;
    circle(pt cen_, double r_) : cen(cen_), r(r_) {}
    circle(pt a, pt b, pt c) {
        cen = center(a, b, c);
        r = dist(cen, a);
    bool inside(pt p) { return dist(p, cen) < r+EPS; }</pre>
};
circle minCirc(vector<pt> v) {
    shuffle(v.begin(), v.end(), rng);
    circle ret = circle(pt(0, 0), 0);
    for (int i = 0; i < v.size(); i++) if</pre>
       (!ret.inside(v[i])) {
        ret = circle(v[i], 0);
        for (int j = 0; j < i; j++) if (!ret.inside(v[j])) {</pre>
            ret = circle((v[i]+v[j])/2, dist(v[i], v[j])/2);
            for (int k = 0; k < j; k++) if
                (!ret.inside(v[k]))
                ret = circle(v[i], v[j], v[k]);
        }
    return ret;
}
```

#### 6.28 Min fixed range

```
// https://codeforces.com/contest/1195/problem/E
//
// O(n)
// ans[i] = min_{0 <= j < k} v[i+j]

vector<int> min_k(vector<int> &v, int k){
   int n = v.size();
   deque<int> d;
   auto put = [&](int i){
      while (!d.empty() && v[d.back()] > v[i])
            d.pop_back();
      d.push_back(i);
   };
```

```
for (int i = 0; i < k-1; i++)
        put(i);

vector < int > ans(n-k+1);

for (int i = 0; i < n-k+1; i++) {
        put(i+k-1);
        while (i > d.front()) d.pop_front();
        ans[i] = v[d.front()];
}

return ans;
}
```

#### 6.29 Conectividade Dinamica 2

```
// Offline com link-cut trees
// O(n log(n))
namespace lct {
    struct node {
        int p, ch[2];
        int val, sub;
        bool rev;
        node() {}
        node(int v) : p(-1), val(v), sub(v), rev(0) { ch[0]}
           = ch[1] = -1;
    };
    node t[2*MAX]; // MAXN + MAXQ
    map<ii, int> aresta;
    int sz;
    void prop(int x) {
        if (t[x].rev) {
            swap(t[x].ch[0], t[x].ch[1]);
            if (t[x].ch[0]+1) t[t[x].ch[0]].rev ^= 1;
            if (t[x].ch[1]+1) t[t[x].ch[1]].rev ^= 1;
        t[x].rev = 0;
    }
    void update(int x) {
        t[x].sub = t[x].val;
```

```
for (int i = 0; i < 2; i++) if (t[x].ch[i]+1) {
        prop(t[x].ch[i]);
       t[x].sub = min(t[x].sub, t[t[x].ch[i]].sub);
    }
}
bool is_root(int x) {
    return t[x].p == -1 or (t[t[x].p].ch[0] != x and
       t[t[x].p].ch[1] != x);
}
void rotate(int x) {
    int p = t[x].p, pp = t[p].p;
    if (!is_root(p)) t[pp].ch[t[pp].ch[1] == p] = x;
    bool d = t[p].ch[0] == x;
   t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
    if (t[p].ch[!d]+1) t[t[p].ch[!d]].p = p;
   t[x].p = pp, t[p].p = x;
   update(p), update(x);
int splay(int x) {
    while (!is_root(x)) {
        int p = t[x].p, pp = t[p].p;
       if (!is_root(p)) prop(pp);
        prop(p), prop(x);
        if (!is_root(p)) rotate((t[pp].ch[0] ==
           p)^{(t[p].ch[0] == x)} ? x : p);
        rotate(x);
    }
    return prop(x), x;
}
int access(int v) {
    int last = -1;
    for (int w = v; w+1; update(last = w), splay(v), w =
       t[v].p)
        splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
    return last;
void make_tree(int v, int w=INF) { t[v] = node(w); }
bool conn(int v, int w) {
    access(v), access(w);
    return v == w ? true : t[v].p != -1;
}
void rootify(int v) {
```

```
access(v):
        t[v].rev ^= 1;
   }
    int query(int v, int w) {
        rootify(w), access(v);
        return t[v].sub;
   }
    void link_(int v, int w) {
        rootify(w);
        t[w].p = v;
   }
    void link(int v, int w, int x) { // v--w com peso x
        int id = MAX + sz++;
        aresta[make_pair(v, w)] = id;
        make_tree(id, x);
        link_(v, id), link_(id, w);
   }
    void cut_(int v, int w) {
        rootify(w), access(v);
        t[v].ch[0] = t[t[v].ch[0]].p = -1;
   }
    void cut(int v, int w) {
        int id = aresta[make_pair(v, w)];
        cut_(v, id), cut_(id, w);
   }
void dyn_conn() {
    int n, q; cin >> n >> q;
    vector < int > p(2*q, -1); // outra ponta do intervalo
    for (int i = 0; i < n; i++) lct::make_tree(i);</pre>
    vector<ii> qu(q);
   map<ii, int> m;
    for (int i = 0; i < q; i++) {</pre>
        char c; cin >> c;
        if (c == '?') continue;
        int a, b; cin >> a >> b; a--, b--;
        if (a > b) swap(a, b);
        qu[i] = {a, b};
        if (c == '+') {
            p[i] = i+q, p[i+q] = i;
            m[make_pair(a, b)] = i;
```

}

```
} else {
        int j = m[make_pair(a, b)];
        p[i] = j, p[j] = i;
    }
}
int ans = n;
for (int i = 0; i < q; i++) {</pre>
    if (p[i] == -1) {
        cout << ans << endl; // numero de comp conexos</pre>
        continue;
    }
    int a = qu[i].f, b = qu[i].s;
    if (p[i] > i) { // +
        if (lct::conn(a, b)) {
            int mi = lct::query(a, b);
            if (p[i] < mi) {</pre>
                p[p[i]] = p[i];
                 continue;
            lct::cut(qu[p[mi]].f, qu[p[mi]].s), ans++;
            p[mi] = mi;
        lct::link(a, b, p[i]), ans--;
    } else if (p[i] != i) lct::cut(a, b), ans++; // -
```

## 6.30 Points Inside Polygon

}

```
// Encontra quais pontos estao
// dentro de um poligono simples nao convexo
// o poligono tem lados paralelos aos eixos
// Pontos na borda estao dentro
// Pontos podem estar em ordem horaria ou anti-horaria
//
// O(n log(n))
#define f first
#define s second
#define pb push_back
```

```
typedef long long 11;
typedef pair<int, int> ii;
const 11 N = 1e9+10;
const int MAX = 1e5+10;
int ta[MAX];
namespace seg {
    unordered_map<11, int> seg;
    int query(int a, int b, ll p, ll l, ll r) {
        if (b < l or r < a) return 0;
        if (a <= 1 and r <= b) return seg[p];</pre>
        11 m = (1+r)/2;
        return query(a, b, 2*p, 1, m)+query(a, b, 2*p+1,
           m+1, r);
    }
    int query(ll p) {
        return query (0, p+N, 1, 0, 2*N);
    }
    int update(ll i, int x, ll p, ll l, ll r) {
        if (i < l or r < i) return seg[p];</pre>
        if (1 == r) return seg[p] += x;
        11 m = (1+r)/2;
        return seg[p] = update(i, x, 2*p, 1, m)+update(i, x,
           2*p+1, m+1, r);
    }
    void update(ll a, ll b, int x) {
        if (a > b) return;
        update(a+N, x, 1, 0, 2*N);
        update(b+N+1, -x, 1, 0, 2*N);
   }
};
void pointsInsidePol(vector<ii>& pol, vector<ii>& v) {
    vector<pair<int, pair<int, ii> > ev; // {x, {tipo, {a,
       b}}}
   // -1: poe ; id: query ; 1e9: tira
   for (int i = 0; i < v.size(); i++)</pre>
        ev.pb({v[i].f, {i, {v[i].s, v[i].s}}});
   for (int i = 0; i < pol.size(); i++) {</pre>
        ii u = pol[i], v = pol[(i+1)\%pol.size()];
```

```
if (u.s == v.s) {
        ev.pb(\{min(u.f, v.f), \{-1, \{u.s, u.s\}\}\});
        ev.pb({max(u.f, v.f), {N, {u.s, u.s}}});
        continue;
    }
    int t = N;
    if (u.s > v.s) t = -1;
    ev.pb({u.f, {t, {min(u.s, v.s)+1, max(u.s, v.s)}}});
}
sort(ev.begin(), ev.end());
for (int i = 0; i < v.size(); i++) ta[i] = 0;</pre>
for (auto i : ev) {
    pair<int, ii> j = i.s;
    if (j.f == -1) seg::update(j.s.f, j.s.s, 1);
    else if (j.f == N) seg::update(j.s.f, j.s.s, -1);
    else if (seg::query(j.s.f)) ta[j.f] = 1; // ta dentro
}
```

# 7 Strings

}

## 7.1 Algoritmo Z

```
// Complexidades:
// z - O(|s|)
// match - O(|s| + |p|)

vector<int> get_z(string s) {
   int n = s.size();
   vector<int> z(n, 0);

   // intervalo da ultima substring valida
   int l = 0, r = 0;
   for (int i = 1; i < n; i++) {
        // estimativa pra z[i]
        if (i <= r) z[i] = min(r - i + 1, z[i - 1]);
        // calcula valor correto</pre>
```

```
while (i + z[i] < n and s[z[i]] == s[i + z[i]])
        z[i]++;
        // atualiza [l, r]
        if (i + z[i] - 1 > r) l = i, r = i + z[i] - 1;
}

return z;
}

// quantas vezes p aparece em s
int match(string s, string p) {
    int n = s.size(), m = p.size();
    vector<int> z = get_z(p + s);

int ret = 0;
    for (int i = m; i < n + m; i++)
        if (z[i] >= m) ret++;

return ret;
}
```

## 7.2 String hashing

```
// Para evitar colisao: testar mais de um
// mod; so comparar strings do mesmo tamanho
// ex : str_hash<1e9+7> h(s);
// ll val = h(10, 20);
//
// Complexidades:
// build - O(|s|)
// operator() - O(1)

mt19937 rng((int)
    chrono::steady_clock::now().time_since_epoch().count());
int uniform(int l, int r) {
    uniform_int_distribution<int> uid(l, r);
    return uid(rng);
}
```

```
template < int MOD > struct str_hash {
    static int P;
    int n;
    string s;
    vector<ll> h, power;
    str_hash(string s_) : n(s_.size()), s(s_), h(n),
       power(n) {
        power[0] = 1;
        for (int i = 1; i < n; i++) power[i] = power[i-1]*P</pre>
           % MOD;
        h[0] = s[0];
        for (int i = 1; i < n; i++) h[i] = (h[i-1]*P + s[i])
           % MOD:
    }
    11 operator()(int i, int j) { // retorna hash da
       substring s[i..j]
        if (!i) return h[j];
        ll ret = h[j] - h[i-1]*power[j-i+1] % MOD;
        return ret < 0 ? ret+MOD : ret;</pre>
    }
};
template < int MOD > int str_hash < MOD > :: P = uniform(27, MOD - 1);
// primeiro parametro deve ser maior que o tamanho do
   alfabeto
```

#### 7.3 Automato de Sufixo

```
// Automato que aceita os sufixos de uma string
// Todas as funcoes sao lineares

namespace sam {
   int cur, sz, len[2*MAX], link[2*MAX], acc[2*MAX];
   int nxt[2*MAX][26];

  void add(int c) {
     int at = cur;
     len[sz] = len[cur]+1, cur = sz++;
     while (at != -1 and !nxt[at][c]) nxt[at][c] = cur,
        at = link[at];
   if (at == -1) { link[cur] = 0; return; }
```

```
int q = nxt[at][c];
    if (len[q] == len[at]+1) { link[cur] = q; return; }
    int qq = sz++;
    len[qq] = len[at]+1, link[qq] = link[q];
    for (int i = 0; i < 26; i++) nxt[qq][i] = nxt[q][i];</pre>
    while (at != -1 and nxt[at][c] == q) nxt[at][c] =
       qq, at = link[at];
    link[cur] = link[q] = qq;
}
void build(string& s) {
    cur = 0, sz = 0, len[0] = 0, link[0] = -1, sz++;
    for (auto i : s) add(i-'a');
    int at = cur:
    while (at) acc[at] = 1, at = link[at];
}
// coisas que da pra fazer:
11 distinct_substrings() {
    11 \text{ ans} = 0;
    for (int i = 1; i < sz; i++) ans += len[i] -
       len[link[i]];
    return ans;
}
string longest_common_substring(string& S, string& T) {
    build(S);
    int at = 0, 1 = 0, ans = 0, pos = -1;
    for (int i = 0; i < T.size(); i++) {</pre>
        while (at and !nxt[at][T[i]-'a']) at = link[at],
           1 = len[at];
        if (nxt[at][T[i]-'a']) at = nxt[at][T[i]-'a'],
           1++;
        else at = 0, 1 = 0;
        if (1 > ans) ans = 1, pos = i;
    return T.substr(pos-ans+1, ans);
}
11 dp[2*MAX];
11 paths(int i) {
    auto& x = dp[i];
    if (x) return x;
    x = 1;
    for (int j = 0; j < 26; j++) if (nxt[i][j]) x +=</pre>
```

```
paths(nxt[i][j]);
    return x;
}

void kth_substring(int k, int at=0) { // k=1 : menor
    substring lexicog.
    for (int i = 0; i < 26; i++) if (k and nxt[at][i]) {
        if (paths(nxt[at][i]) >= k) {
            cout << char('a'+i);
            kth_substring(k-1, nxt[at][i]);
            return;
        }
        k -= paths(nxt[at][i]);
    }
};</pre>
```

## 7.4 Suffix Array - O(n)

```
// Rapidao
// Computa o suffix array em 'sa', o rank em 'rnk'
// e o lcp em 'lcp'
// query(i, j) retorna o LCP entre s[i..n-1] e s[j..n-1]
// Complexidades (assumindo rmg <0(n), 0(1)>):
// O(n) para construir
// query - 0(1)
struct suffix_array {
    string s;
    int n;
    vector < int > sa, cnt, rnk, lcp;
    rmq<int> RMQ;
    bool cmp(int a1, int b1, int a2, int b2, int a3=0, int
       b3=0) {
        return a1 != b1 ? a1 < b1 : (a2 != b2 ? a2 < b2 : a3
           < b3);
    template < typename T> void radix(int* fr, int* to, T* r,
       int N, int k) {
```

```
cnt = vector < int > (k+1, 0):
    for (int i = 0; i < N; i++) cnt[r[fr[i]]]++;</pre>
    for (int i = 1; i <= k; i++) cnt[i] += cnt[i-1];</pre>
    for (int i = N-1; i+1; i--) to[--cnt[r[fr[i]]]] =
       fr[i];
}
void rec(vector<int>& v, int k) {
    auto &tmp = rnk, &m0 = lcp;
    int N = v.size()-3, sz = (N+2)/3, sz2 = sz+N/3;
    vector < int > R(sz2+3);
    for (int i = 1, j = 0; j < sz2; i += i%3) R[j++] = i;
    radix(&R[0], &tmp[0], &v[0]+2, sz2, k);
    radix(&tmp[0], &R[0], &v[0]+1, sz2, k);
    radix(&R[0], &tmp[0], &v[0]+0, sz2, k);
    int dif = 0;
    int 10 = -1, 11 = -1, 12 = -1;
    for (int i = 0; i < sz2; i++) {</pre>
        if (v[tmp[i]] != 10 or v[tmp[i]+1] != 11 or
            v[tmp[i]+2] != 12)
            10 = v[tmp[i]], 11 = v[tmp[i]+1], 12 =
                v[tmp[i]+2], dif++;
        if (tmp[i]%3 == 1) R[tmp[i]/3] = dif;
        else R[tmp[i]/3+sz] = dif;
    }
    if (dif < sz2) {</pre>
        rec(R, dif);
        for (int i = 0; i < sz2; i++) R[sa[i]] = i+1;</pre>
    } else for (int i = 0; i < sz2; i++) sa[R[i]-1] = i;</pre>
    for (int i = 0, j = 0; j < sz2; i++) if (sa[i] < sz)
       tmp[j++] = 3*sa[i];
    radix(&tmp[0], &m0[0], &v[0], sz, k);
    for (int i = 0; i < sz2; i++)</pre>
        sa[i] = sa[i] < sz ? 3*sa[i]+1 : 3*(sa[i]-sz)+2;
    int at = sz2+sz-1, p = sz-1, p2 = sz2-1;
    while (p \ge 0 \text{ and } p2 \ge 0) {
        if ((sa[p2]%3==1 and cmp(v[m0[p]], v[sa[p2]],
            R[m0[p]/3],
```

```
R[sa[p2]/3+sz])) or (sa[p2]%3==2 and
                cmp(v[m0[p]], v[sa[p2]],
            v[m0[p]+1], v[sa[p2]+1], R[m0[p]/3+sz],
                R[sa[p2]/3+1]))
            sa[at--] = sa[p2--];
        else sa[at--] = m0[p--];
    }
    while (p >= 0) sa[at--] = m0[p--];
    if (N\%3==1) for (int i = 0; i < N; i++) sa[i] =
       sa[i+1];
}
suffix_array(const string& s_) : s(s_), n(s.size()),
   sa(n+3),
        cnt(n+1), rnk(n), lcp(n-1) {
    vector < int > v(n+3);
    for (int i = 0; i < n; i++) v[i] = i;
    radix(&v[0], &rnk[0], &s[0], n, 256);
    int dif = 1;
    for (int i = 0; i < n; i++)</pre>
        v[rnk[i]] = dif += (i and s[rnk[i]] !=
           s[rnk[i-1]]);
    if (n \ge 2) rec(v, dif);
    sa.resize(n);
    for (int i = 0; i < n; i++) rnk[sa[i]] = i;</pre>
    for (int i = 0, k = 0; i < n; i++, k -= !!k) {
        if (rnk[i] == n-1) {
            k = 0;
            continue;
        }
        int j = sa[rnk[i]+1];
        while (i+k < n \text{ and } j+k < n \text{ and } s[i+k] == s[j+k])
           k++:
        lcp[rnk[i]] = k;
    RMQ = rmq < int > (lcp);
}
int query(int i, int j) {
    if (i == j) return n-i;
    i = rnk[i], j = rnk[j];
```

```
return RMQ.query(min(i, j), max(i, j)-1);
}
pair < int, int > next(int L, int R, int i, char c) {
    int l = L, r = R+1;
    while (1 < r) {
        int m = (1+r)/2;
        if (i+sa[m] >= n or s[i+sa[m]] < c) l = m+1;</pre>
        else r = m;
    }
    if (1 == R+1 \text{ or } s[i+sa[1]] > c) \text{ return } \{-1, -1\};
    L = 1;
    1 = L, r = R+1;
    while (1 < r) {
        int m = (1+r)/2;
        if (i+sa[m] >= n or s[i+sa[m]] <= c) l = m+1;</pre>
        else r = m;
    }
    R = 1-1;
    return {L, R};
}
// quantas vezes 't' ocorre em 's' - O(|t| log n)
int count_substr(string& t) {
    int L = 0, R = n-1;
    for (int i = 0; i < t.size(); i++) {</pre>
        tie(L, R) = next(L, R, i, t[i]);
        if (L == -1) return 0;
    return R-L+1;
}
// exemplo de f que resolve o problema
//
   https://codeforces.com/edu/course/2/lesson/2/5/practice/com
ll f(ll k) { return k*(k+1)/2; }
11 dfs(int L, int R, int p) { // dfs na suffix tree
   chamado em pre ordem
    int ext;
    if (L == R) ext = n - sa[L];
    else ext = RMQ.query(L, R-1);
```

```
// Tem 'ext - p' substrings diferentes que ocorrem
           'R-L+1' vezes
        // O LCP de todas elas eh 'ext'
        ll ans = (ext-p)*f(R-L+1);
        // L eh terminal, e folha sse L == R
        if (sa[L]+ext == n) L++;
        /* se for um SA de varias strings separadas como
           s#t$u&, usar no lugar do if de cima
           (separadores < 'a', diferentes e inclusive no
        while (L \leq R && (sa[L]+ext == n || s[sa[L]+ext] \leq
           'a')) {
           L++:
        } */
        while (L <= R) {
            auto [1, r] = next(L, R, ext, s[sa[L]+ext]);
            ans += dfs(1, r, ext);
           L = r+1;
        }
        return ans;
    }
    // sum over substrings: computa, para toda substring t
       distinta de s,
    // \sum f(# ocorrencias de t em s) - 0 (n log n)
    11 sos() { return dfs(0, n-1, 0); }
}:
```

#### 7.5 Manacher

```
// manacher recebe um vetor de T e retorna o vetor com
   tamanho dos palindromos
// ret[2*i] = tamanho do maior palindromo centrado em i
// ret[2*i+1] = tamanho maior palindromo centrado em i e i+1
//
// Complexidades:
// manacher - O(n)
```

```
// palindrome - <0(n), 0(1)>
// pal_end - O(n)
template < typename T> vector < int > manacher(const vector < T > &
   s) {
    int 1 = 0, r = -1, n = s.size();
    vector < int > d1(n), d2(n);
    for (int i = 0; i < n; i++) {</pre>
        int k = i > r ? 1 : min(d1[l+r-i], r-i);
        while (i+k < n \&\& i-k >= 0 \&\& s[i+k] == s[i-k]) k++;
        d1[i] = k--;
        if (i+k > r) l = i-k, r = i+k;
    }
    1 = 0, r = -1;
    for (int i = 0; i < n; i++) {
        int k = i > r ? 0 : min(d2[1+r-i+1], r-i+1); k++;
        while (i+k \le n \&\& i-k \ge 0 \&\& s[i+k-1] == s[i-k])
           k++;
        d2[i] = --k;
        if (i+k-1 > r) l = i-k, r = i+k-1;
    }
    vector \langle int \rangle ret (2*n-1);
    for (int i = 0; i < n; i++) ret[2*i] = 2*d1[i]-1;</pre>
    for (int i = 0; i < n-1; i++) ret[2*i+1] = 2*d2[i+1];
    return ret;
}
// verifica se a string s[i..j] eh palindromo
template < typename T > struct palindrome {
    vector < int > man;
    palindrome(const vector < T > & s) : man(manacher(s)) {}
    bool query(int i, int j) {
        return man[i+j] >= j-i+1;
    }
};
// tamanho do maior palindromo que termina em cada posicao
template < typename T > vector < int > pal_end(const vector < T > & s)
    vector<int> ret(s.size());
    palindrome <T> p(s);
```

```
ret[0] = 1;
for (int i = 1; i < s.size(); i++) {
    ret[i] = min(ret[i-1]+2, i+1);
    while (!p.query(i-ret[i]+1, i)) ret[i]--;
}
return ret;
}</pre>
```

#### 7.6 eertree

```
// Constroi a eertree, caractere a caractere
// Inicializar com a quantidade de caracteres maxima
// size() retorna a quantidade de substrings pal. distintas
// depois de chamar propagate(), cada substring palindromica
// ocorre qt[i] vezes. O propagate() retorna o numero de
// substrings pal. com repeticao
//
// O(n) amortizado, considerando alfabeto O(1)
struct eertree {
    vector < vector < int >> t;
    int n, last, sz;
    vector<int> s, len, link, qt;
    eertree(int N) {
        t = vector(N+2, vector(26, int()));
        s = len = link = qt = vector < int > (N+2);
        s[0] = -1;
        link[0] = 1, len[0] = 0, link[1] = 1, len[1] = -1;
        sz = 2, last = 0, n = 1;
    }
    void add(char c) {
        s[n++] = c -= 'a';
        while (s[n-len[last]-2] != c) last = link[last];
        if (!t[last][c]) {
            int prev = link[last];
            while (s[n-len[prev]-2] != c) prev = link[prev];
            link[sz] = t[prev][c];
            len[sz] = len[last]+2;
```

```
t[last][c] = sz++;
}
    qt[last = t[last][c]]++;
}
int size() { return sz-2; }
ll propagate() {
    ll ret = 0;
    for (int i = n; i > 1; i--) {
        qt[link[i]] += qt[i];
        ret += qt[i];
    }
    return ret;
}
};
```

#### 7.7 String hashing - modulo 2<sup>61</sup> - 1

```
// Usa modulo 2^61 - 1 \sim 2e18
// Eh quase duas vezes mais lento
// Complexidades:
// build - O(|s|)
// operator() - 0(1)
const 11 MOD = (111<<<61) -1;</pre>
ll mulmod(ll a, ll b) {
    const static ll LOWER = (111<<30)-1, GET31 = (111<<31)-1;</pre>
    11 \ 11 = a\&LOWER, h1 = a>>30, 12 = b\&LOWER, h2 = b>>30;
    11 m = 11*h2 + 12*h1, h = h1*h2;
    ll ans = 11*12 + (h>>1) + ((h&1)<<60) + (m>>31) +
       ((m\&GET31) << 30) + 1;
    ans = (ans\&MOD) + (ans >> 61);
    ans = (ans\&MOD) + (ans>>61);
    return ans-1;
}
mt19937_64
   rng(chrono::steady_clock::now().time_since_epoch().count());
```

```
ll uniform(ll l, ll r) {
    uniform_int_distribution < ll> uid(l, r);
    return uid(rng);
}
struct str_hash {
    static 11 P;
    int n;
    string s;
    vector<ll> h, power;
    str_hash(string s_) : n(s_.size()), s(s_), h(n),
       power(n) {
        power[0] = 1;
        for (int i = 1; i < n; i++) power[i] =</pre>
           mulmod(power[i-1], P);
        h[0] = s[0];
        for (int i = 1; i < n; i++) h[i] = (mulmod(h[i-1],
           P) + s[i]) % MOD;
    11 operator()(int i, int j) { // retorna hash da
       substring s[i..j]
       if (!i) return h[j];
        ll ret = h[j] - mulmod(h[i-1], power[j-i+1]);
        return ret < 0 ? ret+MOD : ret;</pre>
    }
};
11 str_hash::P = uniform(27, MOD-1);
// primeiro parametro deve ser maior que o tamanho do
   alfabeto
     Max Suffix
```

```
// computa o indice do maior sufixo da
// string, lexicograficamente
//
// O(n)
int max_sulf(string s) {
    s += '#';
    int ans = max_element(s.begin(), s.end()) - s.begin();
```

```
for (int i = ans+1, j = 0; i < s.size(); i++) {
        if (ans+j < i and s[i] == s[ans+j]) j++;
        else {
            if (ans+j < i and s[i] > s[ans+j]) ans = i-j;
            j = 0;
        }
    }
    return ans;
}
7.9 KMP
// mathcing(s, t) retorna os indices das ocorrencias
// de s em t
// autKMP constroi o automato do KMP
//
// Complexidades:
// pi - O(n)
// match - 0(n + m)
// construir o automato - O(|sigma|*n)
// n = |padrao| e m = |texto|
vector<int> pi(string s) {
    vector < int > p(s.size());
    for (int i = 1, j = 0; i < s.size(); i++) {</pre>
        while (j \text{ and } s[j] != s[i]) j = p[j-1];
        if (s[j] == s[i]) j++;
        p[i] = j;
    }
    return p;
}
vector<int> matching(string& t, string& s) {
    vector < int > p = pi(s+'$'), match;
    for (int i = 0, j = 0; i < t.size(); i++) {
        while (j \text{ and } s[j] != t[i]) j = p[j-1];
        if (s[j] == t[i]) j++;
        if (j == s.size()) match.push_back(i-j+1);
    }
    return match;
```

## 7.10 Suffix Array - $O(n \log n)$

```
// kasai recebe o suffix array e calcula lcp[i],
// o lcp entre s[sa[i],...,n-1] e s[sa[i+1],..,n-1]
// Complexidades:
// suffix_array - O(n log(n))
// kasai - O(n)
vector<int> suffix_array(string s) {
    s += "$";
    int n = s.size(), N = max(n, 260);
    vector < int > sa(n), ra(n);
    for(int i = 0; i < n; i++) sa[i] = i, ra[i] = s[i];</pre>
    for(int k = 0; k < n; k ? k *= 2 : k++) {
        vector < int > nsa(sa), nra(n), cnt(N);
        for(int i = 0; i < n; i++) nsa[i] = (nsa[i]-k+n)%n,
            cnt[ra[i]]++;
        for(int i = 1; i < N; i++) cnt[i] += cnt[i-1];</pre>
        for(int i = n-1; i+1; i--) sa[--cnt[ra[nsa[i]]]] =
           nsa[i];
```

```
for(int i = 1, r = 0; i < n; i++) nra[sa[i]] = r +=</pre>
           ra[sa[i]] !=
            ra[sa[i-1]] or ra[(sa[i]+k)%n] !=
                ra[(sa[i-1]+k)%n];
        ra = nra;
        if (ra[sa[n-1]] == n-1) break;
    return vector < int > (sa.begin() +1, sa.end());
}
vector<int> kasai(string s, vector<int> sa) {
    int n = s.size(), k = 0;
    vector < int > ra(n), lcp(n);
    for (int i = 0; i < n; i++) ra[sa[i]] = i;
    for (int i = 0; i < n; i++, k -= !!k) {
        if (ra[i] == n-1) { k = 0; continue; }
        int j = sa[ra[i]+1];
        while (i+k < n \text{ and } j+k < n \text{ and } s[i+k] == s[j+k]) k++;
        lcp[ra[i]] = k;
    }
    return lcp;
7.11 Ahocorasick
// Complexidades:
// all linear O(n*sigma)
// example of query returns number of nonoverlapping matches
namespace aho {
    const vector<pair<char, char>> vt = {
        {'a', 'z'},
        {'A', 'Z'},
        {'0', '9'}
    };//example of alphabet
```

void fix(char &c){
 int acc = 0;

```
for (auto p : vt){
            if (p.first <= c && c <= p.second){</pre>
                c = c - p.first + acc;
                return;
            }
            acc += p.second - p.first + 1;
        }
    }
    void unfix(char &c){
        int acc = 0;
        for (auto p : vt){
            int next_acc = acc + p.second - p.first;
            if (acc <= c && c <= next_acc){</pre>
                c = p.first + c - acc;
                return;
            }
            acc = next_acc + 1;
        }
    void fix(string &s){ for (char &c : s) fix(c); }
    void unfix(string &s){ for (char &c : s) unfix(c); }
    const int SIGMA = 70; //fix(vt.back().second) + 1;
    const int MAXN = 1e5+10;
    int to[MAXN][SIGMA];
    int link[MAXN], end[MAXN];
    int idx;
    void init(){
#warning dont forget to init before inserting strings
        memset(to, 0, sizeof to);
        idx = 1;
    void insert(string &s){
        fix(s):
        int v = 0:
        for (char c : s){
            int &w = to[v][c];
            if (!w) w = idx++;
            v = w;
        }
```

```
end[v] = 1:
    }
    void build(){
#warning dont forget to build after inserting strings
         queue < int > q;
         q.push(0);
         while (!q.empty()){
             int cur = q.front(); q.pop();
             int l = link[cur];
             end[cur] |= end[1];
             for (int i = 0; i < SIGMA; i++){</pre>
                 int &w = to[cur][i];
                 if (w){
                     link[w] = ((cur != 0) ? to[1][i] : 0);
                     q.push(w);
                 else w = to[l][i];
             }
         }
    }
    int query(string &s){
        fix(s);
        int v = 0;
         int counter = 0;
         for (char c : s){
             v = to[v][c];
             if (end[v]) {
                 counter++;
                 v = 0; //remove if matches could overlap
             }
         }
         return counter;
    }
}
7.12 Trie
// trie T() constroi uma trie para o alfabeto das letras
    minusculas
// trie T(tamanho do alfabeto, menor caracter) tambem pode
```

```
ser usado
//
// T.insert(s) - 0(|s|*sigma)
// T.erase(s) - O(|s|)
// T.find(s) retorna a posicao, O se nao achar - O(|s|)
// T.count_pref(s) numero de strings que possuem s como
   prefixo - O(|s|)
//
// Nao funciona para string vazia
struct trie {
    vector < vector < int >> to;
    vector<int> end, pref;
    int sigma; char norm;
    trie(int sigma_=26, char norm_='a') : sigma(sigma_),
       norm(norm_) {
        to = {vector < int > (sigma)};
        end = \{0\}, pref = \{0\};
    void insert(string s) {
        int x = 0;
        for(auto c : s) {
            int &nxt = to[x][c-norm];
            if(!nxt) {
                nxt = to.size();
                to.push_back(vector<int>(sigma));
                end.push_back(0), pref.push_back(0);
            x = nxt, pref[x]++;
        end[x]++;
    void erase(string s) {
        int x = 0;
        for(char c : s) {
            int &nxt = to[x][c-norm];
            x = nxt, pref[x] --;
            if(!pref[x]) nxt = 0;
        }
        end[x]--;
    }
    int find(string s) {
```

```
int x = 0;
    for(auto c : s) {
        x = to[x][c-norm];
        if(!x) return 0;
    }
    return x;
}
int count_pref(string s) {
    return pref[find(s)];
}
};
```

## 8 Extra

#### 8.1 makefile

```
CXX = g++
CXXFLAGS = -fsanitize=address,undefined -01
   -fno-omit-frame-pointer -g -Wall -Wshadow -std=c++17
   -Wno-unused-result -Wno-sign-compare -Wno-char-subscripts
#-fuse-ld=gold
```

## 8.2 debug.cpp

```
void debug_out(string s, int line) { cerr << endl; }
template < typename H, typename... T>
void debug_out(string s, int line, H h, T... t) {
   if (s[0] != ',') cerr << "Line(" << line << ") ";
   do { cerr << s[0]; s = s.substr(1);
   } while (s.size() and s[0] != ',');
   cerr << " = " << h;
   debug_out(s, line, t...);
}
#ifdef DEBUG
#define debug(...) debug_out(#__VA_ARGS__, __LINE__,
   __VA_ARGS__)
#else
#define debug(...)
#endif</pre>
```

## 8.3 template.cpp

```
#include <bits/stdc++.h>
using namespace std;

#define _ ios_base::sync_with_stdio(0);cin.tie(0);
#define endl '\n'
#define f first
```

```
#define s second
#define pb push_back

typedef long long ll;
typedef pair < int, int > ii;

const int INF = 0x3f3f3f3f3f;
const ll LINF = 0x3f3f3f3f3f3f3f3f3f3f1l;

int main() { _
    exit(0);
}
```

### 8.4 fastIO.cpp

```
int read_int() {
    bool minus = false;
    int result = 0;
    char ch;
    ch = getchar();
    while (1) {
        if (ch == '-') break;
        if (ch >= '0' && ch <= '9') break;
        ch = getchar();
    }
    if (ch == '-') minus = true;
    else result = ch-'0';
    while (1) {
        ch = getchar();
        if (ch < '0' || ch > '9') break;
        result = result *10 + (ch - '0');
    if (minus) return -result;
    else return result;
}
```

#### 8.5 vimrc

```
set ts=4 si ai sw=4 number mouse=a syntax on
```

#### 8.6 stress.sh

```
make a a2 gen || exit 1
for ((i = 1; ; i++)) do
    ./gen $i > in
    ./a < in > out
    ./a2 < in > out2
    if (! cmp -s out out2) then
        echo "--> entrada:"
        cat in
        echo "--> saida1:"
        cat out
        echo "--> saida2:"
        cat out2
        break;
    fi
    echo $i
done
```

# 8.7 rand.cpp

```
mt19937 rng((int)
    chrono::steady_clock::now().time_since_epoch().count());
int uniform(int 1, int r){
    uniform_int_distribution<int> uid(1, r);
    return uid(rng);
}
```