

# Overview

## Relational Database

## Hands-On SQL

BMI701 Introduction of Biomedical Informatics  
Lab Session 1

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Wei-Hung Weng

September 7, 2016

HMS DBMI — MGH LCS



- Wei-Hung Weng, MD
- Research fellow in MGH Lab of Comp Sci
- NLP, medical ontology, machine learning, databases
- Every Thursday 12:30pm, about 1 hour
- Flexible, not mandatory
- Website
  - [Github repository](#) (codes and slides)
  - [ckbjimmy@gmail.com](mailto:ckbjimmy@gmail.com)

# Schedule

Lab	Date	Topic
1	Sep 08	Overview / relational database / SQL
2	Sep 15	Database design / Database normalization / R basic
3	Sep 22	Ontology / database using R
4	Sep 26	<b>(Monday before &amp; after Adam's class)</b> NoSQL / Project discussion
	Oct 06	<b>No lab (Presentation week)</b>
	Oct 13	<b>No lab (Columbus day)</b>
5	Oct 20	NLP using R or py / regular expression

# Schedule

Lab	Date	Topic
6	Oct 27	MetaMap / cTAKES demonstration
7	Nov 03	Data visualization using ggplot2
8	Nov 10	Bioinformatics tools / GWAS
	Nov 17	<b>No lab (AMIA)</b>
9	Nov 24	ML using Weka, R or python (or unsupervised learning / feature engineering)
10	Dec 01	Deep learning on image classification (py)
	Dec 08	<b>No lab (Final week)</b>

- Background? (clinicians / scientists and engineers)
- CS courses? (introduction / data structure / algorithm)
- Math courses?
- Programming language?
  - C/C++, Java, Python, R, Matlab/Octave, Perl, ...

# Today

- Database
- Relational database
- Simple SQL syntax
- MySQL

# Why Database?

- I believe that you've learned a lot about database from Adam's class
- RAM is expensive and small  $\rightarrow$  HD is cheaper and large
- You can't put all data into RAM (unless you are \$\$\$\$)
- But it's also not easy to manage data on your HD
- Therefore we need database solution
  - $O(1) < O(\log(n)) < O(n) < O(n \log(n)) < O(n^m) < O(2^n)$

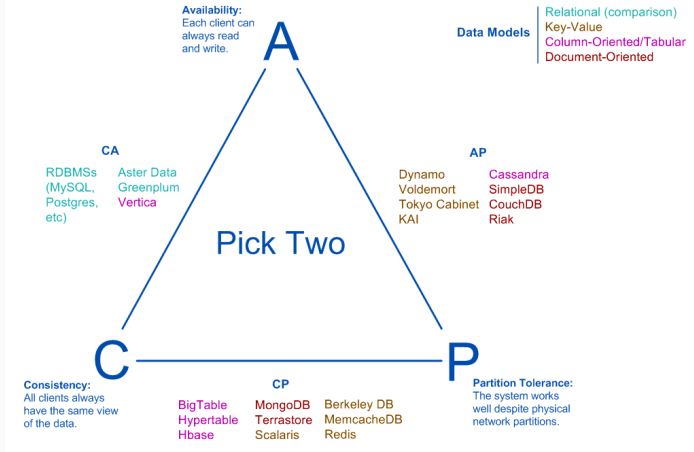
# Database CAP

- Consistency
  - All nodes should see the same data at the same time (transaction)
- Availability/Speed
  - Every request receives a response about whether the transaction is passed or failed
- Partition tolerance
  - Able to continue operation even if there are partially loss/failure of the system
- Eric Brewer 2000
- You may only pick up two of three



- Structured Query Language (higher C and A)
  - RDB
  - Can't deal with extremely large data
- NoSQL for big data (higher P)
- Big data?
  - Google/facebook/twitter-scale (10TB to 10PB)
  - AWS: No need to use NoSQL if your user < 10M
- So we will talk about NoSQL later

## Visual Guide to NoSQL Systems



# Popular RDBMS

- MS SQL Server: you will probably use this in your lab/company
- Oracle: awesome if you are \$\$\$\$
- **MySQL** / MariaDB
  - Open-sourced, fast, easy to scale-up, community, BUT no JSON
- **PostgreSQL**
  - Open-sourced, fast enough, can scale-up, can use JSON, array, ...

- ACID
  - Atomicity (all or none / commit or rollback)
  - Consistency (transaction between two accounts)
  - Isolation (locking)
  - Durability (won't rollback once the transaction is done)
- PK, FK, index
- ER diagram
- Normalization/denormalization (next lab)

- A table has fixed columns
- A column has a datatype
  - Numeric, character, datetime, boolean, geometric, text, JSON, array (postgresql), ...
- Rows can grow (python tuple  $\rightarrow$  list)

## Primary Key (PK)

- Unique identifier in the table
- MUST be unique
- CAN'T be NULL
- Index is created

## Foreign Key (FK)

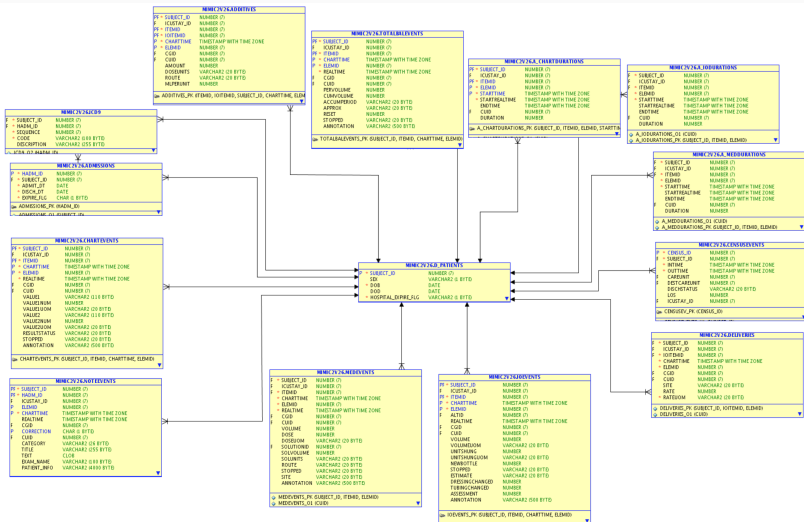
- For referential integrity
- The value **MUST** exist in other tables
- May be NULL or duplicated
- No index is created

- Fast but takes space
- Slow insert/update but fast read/delete
- B+ tree

Task	No index	Index
Create	$O(n)$	$O(n)$
Read	$O(n)$	$O(\log(n))$
Update	$O(1)$	$O(\log(n))$
Delete	$O(n)$	$O(\log(n))$



# Entity-Relationship Diagram (ER Diagram)



# SQL's CRUD

- CRUD: create, read, update, delete
- Insert
- Select
  - **Filter**: where, distinct
  - **Aggregate**: group by
  - **Join**
  - **Subquery**
  - Partition (splitting the data into manageable size)
- Update
- Delete

- Install MySQL Server
- Install MySQL Workbench
- Small databases
- [github.com/ckbjimmy/bmi701lab/blob/master/lab01.sql](https://github.com/ckbjimmy/bmi701lab/blob/master/lab01.sql)

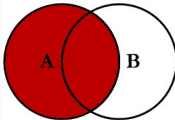
## Filter

- `SELECT * FROM gwas LIMIT 5`
- `SELECT gene FROM gwas WHERE chr_id = 22`
- `SELECT distinct gene, trait, p_value FROM gwas  
WHERE allele like 'A' ORDER BY p_value`
- `SELECT gene, trait FROM gwas WHERE trait like  
'%diabetes_'`
- `SELECT gene, trait FROM gwas WHERE trait like  
'%diabetes%'`
- `SELECT gene, trait FROM gwas WHERE trait like  
'diabetes'`
- `SELECT GROUP_CONCAT(DISTINCT gene SEPARATOR ',')`  
`FROM gwas WHERE chr_id = 3`

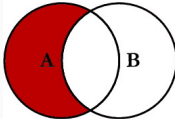
# Aggregation

- `SELECT count(*) FROM sitka`
- `SELECT avg(size) FROM sitka`
- `SELECT max(size) FROM sitka`
- `SELECT min(size) FROM sitka`
- `SELECT stddev(size) FROM sitka`
- `SELECT tree, avg(size) FROM sitka GROUP BY tree`

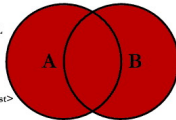
# SQL JOINS



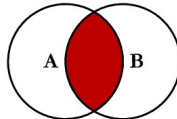
```
SELECT <select_list>
FROM TableA A
LEFT JOIN TableB B
ON A.Key = B.Key
```



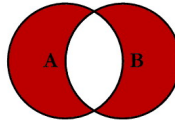
```
SELECT <select_list>
FROM TableA A
LEFT JOIN TableB B
ON A.Key = B.Key
WHERE B.Key IS NULL
```



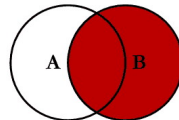
```
SELECT <select_list>
FROM TableA A
FULL OUTER JOIN TableB B
ON A.Key = B.Key
```



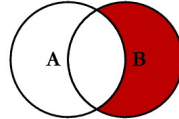
```
SELECT <select_list>
FROM TableA A
INNER JOIN TableB B
ON A.Key = B.Key
```



```
SELECT <select_list>
FROM TableA A
FULL OUTER JOIN TableB B
ON A.Key = B.Key
WHERE A.Key IS NULL
OR B.Key IS NULL
```



```
SELECT <select_list>
FROM TableA A
RIGHT JOIN TableB B
ON A.Key = B.Key
```



```
SELECT <select_list>
FROM TableA A
RIGHT JOIN TableB B
ON A.Key = B.Key
WHERE A.Key IS NULL
```

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- ON col1 = col2
- Inner (default), left, right, full, cross (inner join on TRUE)
- Using PK
- ```
SELECT curr_description_f.conceptid,  
curr_description_f.term,  
curr_textdefinition_f.term FROM curr_description_f  
LEFT JOIN curr_textdefinition_f ON  
curr_description_f.conceptid LIKE  
curr_textdefinition_f.conceptid WHERE  
curr_description_f.conceptid LIKE '%777%' LIMIT 30
```

- `SELECT * FROM sitka WHERE treat LIKE '%ozo%'`
- `SELECT * FROM (SELECT * FROM sitka WHERE treat LIKE '%ozo%') oz WHERE size > 5.0`



# Demonstration

- My database project (MySQL, R, shiny, AWS)

BMI701 | Database Project Find the Experts!

**Instruction**

Entering your query and press  
Submit button

**Author**

Wei-Hung Weng (Jimmy)

Last updated on Oct 4, 2015

### Input your symptom and zip code here

headache2138Submit

Show ▾ entries Search:

| last   | first    | zip      | address                | specialtyId |
|--------|----------|----------|------------------------|-------------|
| EDWARD | WOLPOW   | 21385665 | 300 MOUNT AUBURN ST    | NEUROLOGY   |
| ZONGQI | XIA      | 21382903 | 1563 MASSACHUSETTS AVE | NEUROLOGY   |
| AMY    | KAHN     | 21381053 | 725 CONCORD AVE        | NEUROLOGY   |
| JUSTIN | ROUSSEAU | 21382903 | 1563 MASSACHUSETTS AVE | NEUROLOGY   |
| LINDA  | BUCHWALD | 21385600 | 300 MT AUBURN ST       | NEUROLOGY   |

lastfirstzipaddressspecialtyId

Showing 1 to 5 of 5 entriesPrevious1Next

# Preparation for Next Next Week

- Install R
- Install RStudio
- Register UTS Service
- Download SNOMED (need to wait for UTS permission)

# Take Home Message

- Every Thursday 12:30pm
- `SELECT (DISTINCT) ... FROM ... JOIN ... ON ...  
WHERE ... GROUP BY ... ORDER BY ... LIMIT ...`
- Install R/RStudio, register UTS, download SNOMED
- Contact
  - [Github repository](#)
  - [ckbjimmy@gmail.com](mailto:ckbjimmy@gmail.com)
  - [Linkedin: Wei-Hung Weng](#)