

## Assignment 4

The code is available at: <https://github.com/ErlendKK/DAT600-Algorithmic-Theory/tree/main/Assignment%204>

### Problem 1

Let

$x$  = amount of Product X produced

$y$  = amount of Product Y produced

$m$  = heures of machine time

$c$  = heures of craftsman time

#### Constraints:

$$m \leq 40$$

$$c \leq 35$$

$$0,25x + 0,33y = m$$

$$0,33x + 0,50y = c$$

$$0,25x + 0,33y \leq 40;$$

$$0,33x + 0,50y \leq 35;$$

$$x \geq 10;$$

$$y \geq 0;$$

#### Objective function:

$$\text{Revenues} = 200x + 300y$$

$$\text{Costs} = 100m + 20c$$

$$\text{Maximize: } z = \text{Revenues} - \text{Costs}$$

$$= 200x + 300y - 100m - 20c$$

$$= 200x + 300y - 100 \cdot (0,25x + 0,33y) - 20 \cdot (0,33x + 0,50y)$$

$$= 168,4x + 256,7y$$

#### Slackified problem formulation:

$$-168,4x - 256,7y + z = 0$$

$$0,25x + 0,33y + sm = 40$$

$$0,33x + 0,50y + sc = 35$$

$$x \geq 10;$$

$$y \geq 0; sm \geq 0; sc \geq 0;$$

#### Initial simplex tableau:

x	y	sm	sc	Z	C	
0,25	0,33	1	0	0	40	sm
0,33	0,5	0	1	0	35	sc
-168,4	-256,7	0	0	1	0	z

Programmatic solution.

The problem is solved in scipy.

```
from scipy.optimize import linprog

'''
Objective function
Maximize: -168,4x - 256,7y + z = 0;
'''

c = [-168.4, -256.7]

'''
Subject to
0,25x + 0,33y <= 40;
0,33x + 0,50y <= 35;
x >= 10;
y >= 0;
'''

A_ub = [[0.25, 0.33], [0.33, 0.5]]
b_ub = [40, 35]
x_bounds = (10, None)
y_bounds = (0, None)

def solve_problem_one():
    result = linprog(c, A_ub=A_ub, b_ub=b_ub, bounds=[x_bounds, y_bounds], method='highs')
    print(f"Optimal verdi: {-result.fun}")
    print(f"x-verdier: {result.x}")

solve_problem_one()
```

Results:

Optimal verdi: 17958.78

x-verdier: [10. 63.4]

## **Problem 2: Maximum Flow**

### **Problem 2 a)**

let

$G = (V, E)$  be a graph with a set  $V$  of vertices and set  $E$  of edges

$s$  = the source vertex

$t$  = the sink vertex

$u$  and  $v$  = any pair of connected vertices

$f(u,v)$  = the flow on the edge connecting  $u$  and  $v$ .

$C(u, v)$  = the capacity (maximum flow) on the edge connecting  $u$  and  $v$ .

$x(u,v)$  = a boolean representing included/ not included in the minimum cut

Objective function:

Minimize:  $z = \sum(c(u,v) * x(u,v)) - (\text{the sum of capacities for the edges crossing the cut})$

Constraints:

$f(u,v) \leq c(u,v)$  for all  $u,v$  in  $V$

$f(u,v) \geq 0$  for all  $u,v$  in  $V$

The problem is solved with puLP

```
import pulp
import numpy as np

G = np.array([
    [0, 14, 25, 0, 0, 0, 0], # s
    [0, 0, 0, 21, 0, 0, 0], # V1
    [0, 0, 0, 3, 0, 7, 0], # V2
    [0, 6, 13, 0, 10, 5, 0], # V3
    [0, 0, 0, 0, 0, 0, 20], # V4
    [0, 0, 0, 0, 10, 0, 10], # V5
    [0, 0, 0, 0, 0, 0, 0] # t
])

def solve_problem_two_a(G):
    num_nodes = G.shape[0]
    edges = [(i, j) for i in range(num_nodes) for j in range(num_nodes) if G[i, j] > 0]
    prob = pulp.LpProblem("MinimumCut", pulp.LpMinimize)

    # source set (1) or sink set (0)
    in_source_set = pulp.LpVariable.dicts("in_source_set", range(num_nodes), cat=pulp.LpBinary)
    prob += in_source_set[0] == 1
    prob += in_source_set[num_nodes - 1] == 0

    # included/excluded from the cut -> {1,0}
    x = pulp.LpVariable.dicts("x", edges, cat=pulp.LpBinary)

    # Objective Function:
    # Minimize the sum of the capacities of the cut edges

    prob += pulp.lpSum([G[i][j] * x[(i, j)] for (i, j) in edges])

    for (i, j) in edges:
        prob += x[(i, j)] >= in_source_set[i] - in_source_set[j]

    status = prob.solve(pulp.PULP_CBC_CMD(msg=False))
    cut_edges = [(i, j) for (i, j) in edges if pulp.value(x[(i, j)]) == 1]

    print(f"Problem 2a")
    print(f"Minimum Cut Value: {pulp.value(prob.objective)}")
    print(f"Edges in the cut: {cut_edges}\n")

solve_problem_two_a(G)
```

Result:

Minimum Cut Value: 22.0

Edges in the cut: [(2, 5), (3, 4), (3, 5)]

## Problem 2 b)

Based on the textbook (29.25 – 29.28):

Objective function (29.25):

Maximize :  $\sum(f_{sv} - f_{vs})$  over the set of vertices sharing a boundary with s.

Subject to:

$f(u,v) \leq c(u,v)$  for all  $u,v$  in  $V$

$\sum(f(u,v)) = \sum(f(v,u))$ , for all  $u,v$  in  $V - \{s, t\}$

$f(u,v) \geq 0$  for all  $u,v$  in  $V$

According to the textbook this can be rewritten so that it has  $O(V + E)$  constraints. I think the solution is to rewrite the  $V^2$  constraints in terms of edges:

Constraint	Size
$f(u,v) \leq c(u,v)$ for all $(u,v)$ in $E$ :	$E$
$\sum(f(u,v)) = \sum(f(v,u))$ , for all $u,v$ in $V - \{s, t\}$	$V$
$f(u,v) \geq 0$ for all $(u,v)$ in $E$	$E$

$\Rightarrow O(2E + V) = O(V + E)$

The problem is solved in scipy.

```
from scipy.sparse import csr_matrix
from scipy.sparse.csgraph import maximum_flow

def solve_problem_two_b(G):
    graph = csr_matrix(G)
    source = 0
    sink = len(G) - 1
    result = maximum_flow(graph, source, sink)
    print(f"Problem 2b")
    print(f"Maximum flow: {result.flow_value}")

solve_problem_two_b(G)
```

Result: 22