

Computation Theory (COMP 170), Fall 2020

Recitation 01

[1] Induction Practice

We define the reverse of a string recursively as follows:

$$\begin{aligned} \text{rev}(\varepsilon) &= \varepsilon, \\ \text{rev}(w \cdot \mathbf{a}) &= \mathbf{a} \cdot \text{rev}(w) \end{aligned}$$

for all $\mathbf{a} \in \Sigma$ and $w \in \Sigma^*$. Use this definition and induction on the length of v to prove that

$$\text{for all } u, v \in \Sigma^*, \text{rev}(u \cdot v) = \text{rev}(v) \cdot \text{rev}(u).$$

For best results, use the induction proof paradigm resource to help structure your proof.

[2] DFA Construction

Show that each of the following languages are regular by constructing a DFA that accepts it:

- a. $A = \{x \in \{a, b\}^* \mid \text{the second character of } x \text{ is a } b\}$
- b. $B = \{x \in \{a, b\}^* \mid \text{the second to last character of } x \text{ is a } b\}$

You do not need to prove the correctness of your construction. However, you should give a brief explanation, and make sure any diagrams are laid out in a clear and logical manner (and labeled, when appropriate.)

[3] Twisted

For a language A , let $\text{twisted}(A)$ be the language created by taking the even-length strings in A and exchanging every even-indexed character with the character that follows it. E.g. if

$$A = \{\text{cat}, \text{bear}, \text{monkey}\}$$

then

$$\text{twisted}(A) = \{\text{ebra}, \text{omknye}\}.$$

Show that if A is regular, then $\text{twisted}(A)$ is also regular. In particular, given a general DFA $M = (Q, \Sigma, \delta, s, F)$ accepting A , show how to construct a DFA $M' = (Q', \Sigma, \delta', s', F')$ accepting $\text{twisted}(A)$. You do not need to prove your construction is correct, but should provide an explanation of how your DFA works.