Chapter-One

Introduction to compiler Design

Preliminaries Required

- Basic knowledge of FSA and CFG.
- Basic knowledge of programming languages.
 - ✓ Knowledge of a high programming language for the programming assignments.

Textbook:

Alfred V. Aho, Monica S. Lam, Ravi Sethi, and Jeffrey D. Ullman, "Compilers: Principles, Techniques, and Tools" Second Edition, Addison-Wesley, 2007.

Objective of the course

- To learn techniques of a modern compiler
- Be able to build a compiler for a simplified (programming) language
- Know how to use compiler construction tools, such as *generators of scanners and parsers etc.*
- Be familiar with *compiler analysis* and *optimization techniques*.
- ... learn how to work on a larger software project!

Topics will be covered in Chapter-1

- *History of compiler*
- Why study compiler?
- What is a compiler?
- Cousins of the compiler
- Structure of compiler
- Phases of compiler
- The Grouping of Phases into Passes

- Compiler Construction Tools
- The Evolution of Programming Languages
- Application of compiler technology
- Impacts on compiler
- Programming language basics

History of compiler

- No high-level languages were available, so all programming was done in *machine language* and assembly language.
- As expensive as these early computers were, most of the money companies spent was for software development, due to the complexities of assembly.
- In 1953, **John Backus** came up with the idea of "speed coding", and developed the first interpreter.
- Unfortunately, this was *10-20 times* slower than programs written in assembly.
- *He was sure he could do better.*

John Backus

■ In 1954, Backus and his team released a research paper titled "Preliminary Report, Specifications for the IBM Mathematical FORmula TRANslating System, FORTRAN."

History of compiler (cont'd)

- The initial release of **FORTRAN** was in 1956, totaling 25,000 lines of assembly code.
- Compiled programs run almost as fast as handwritten assembly!
- Projects that had taken two weeks to write now took only 2 hours.
- By 1958 more than half of all software was written in FORTRAN.
- Programming languages are notations for describing computations to people and to machines.
- The world as we know it depends on programming languages, because all the software running on all the computers was written in some *programming language*.
- But, before a program can be run, it first must be translated into a form in which it can be executed by a computer.
- The software systems that do this translation are called *compilers*.

Programming Language (PL)

- What's a Programming language?
 - o Formal language designed to communicate instructions to a computer.
- There are two major types of PL: *Low-level and high level language*
- i. Low-level language (LLL)
 - Very close to hardware components
 - Language that are easier for HW to understand
 - Machine oriented and require extensive knowledge of computer HW and its configuration.
 - Two categories of LLL:
 - a. Machine Language & b. Assembly Language

i. Low-level language (LLL) cont'd

a. Machine Language

- is only language that is directly understood by the computer
- It does not need to be translated
- Understood by CPU.
- Example: binary notation (0 & 1)- 00011101

b. Assembly Language

- The first step to improve the *program instructions* make machine language is more readable by humans.
- Set of symbol and letters
- 2nd Generation language. Example: MOVE, ADD, SUB, END

ii. High-Level programming (HLL)

- Is programming language that uses English and mathematical symbol in its instructions like +, -, % and any information
- HLL- that's what you usually use? E.g. *Java*, *C*++, *python*, *etc*
- HLL is most close to the logic of human language
- E.g. consider ATM machine, some one want withdrawal \$100.
- Instructions in high level in computer language would looks something like this:

```
if balance <X:
    print 'Insufficient balance'
else
    Print 'please take your money'</pre>
```

Why study compilers?

- We study compiler construction for the following reasons:
 - Writing a compiler gives a student experience with large-scale applications development
 - To understand performance issues for programming
 - Increases understanding of language semantics
 - Seeing the machine code generated for language constructs helps understand performance issues for languages.

What is a compiler?

• A **compiler** is a program that translates source program written in one language into the target program of another language.

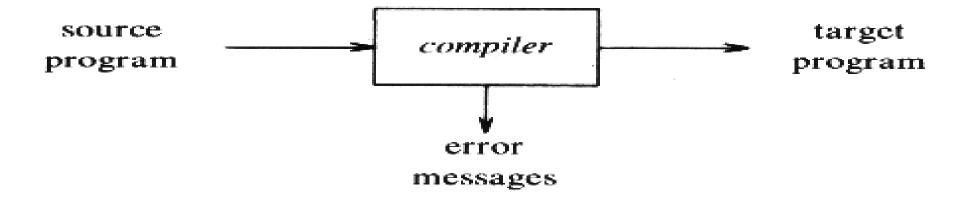


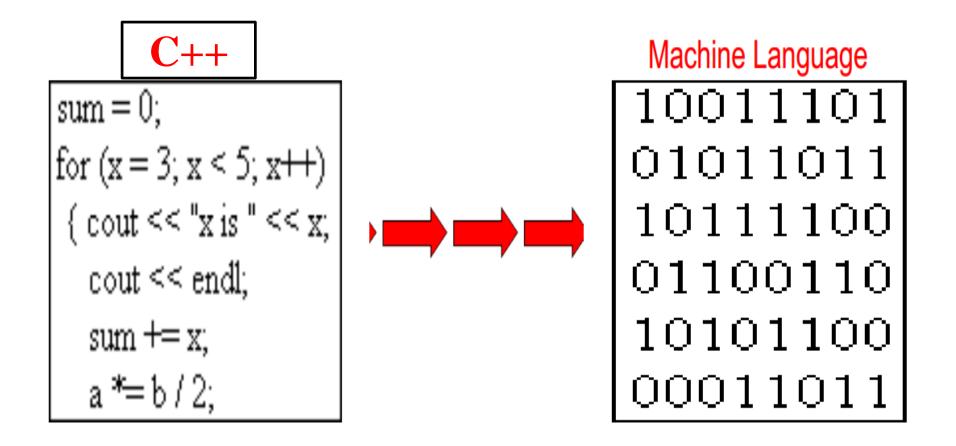
Fig. 1.1. A compiler.

• An important role of the compiler is to report any errors in the source program that it detects during the translation process.

What is a compiler? (cont'd)

- Usually the **source language** (**program**) is a high level language like **Java**, **C**, **C**++, **Python**, etc.
- whereas the **target language** (**program**) is machine code or "code" that a computer's processor understands.
- The source language is optimized for humans.
 - o It is more user-friendly, to some extent platform-independent.
 - o They are easier to read, write, and maintain, and hence it is easy to avoid errors.
- A program written in any language must be translated into a form that is understood by the computer.
 - ✓ This form is typically known as Machine Language (ML) or Machine Code, or Object Code.
- Consists of streams of 0's and 1's

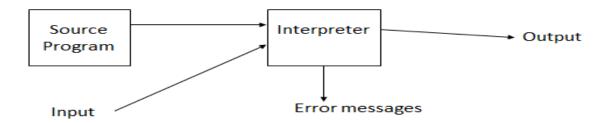
Example: From Source Code into Machine language



• So, the program that carries out the translation activity is called the *compiler*.

Interpreters

- An interpreter is another common kind of language processor.
- Instead of producing a target program as a translation,
- an **interpreter** appears to directly execute the operations specified in the source program on inputs supplied by the user.



- *Interpreter:* processes an internal form of the source program and data at the same time (at run time).
- no object program is generated.

Compiler vs Interpreter

- Compiler takes entire program as input and creates all executable instructions together as compiled code.
- Compilation initially takes time, because it has to create whole **HLL** into machine language at once, and save it on **Hard disk**.
- The file got saved is called *compiled code* or *object code*.
- The compiled code can be used again and again by the **OS** for execution.
- Once the compiled code is done,
 - o the OS do not need the compiler once again
 - Its easy for the OS to schedule the executable instructions to the CPU
 - o If error is occurs then the compiler does not create any single machine instruction

Compiler vs Interpreter

But , An Interpreter:

- o takes single instructions as input and create one or more corresponding machine instructions.
- o does not create a file on the disk.
- Overall delay by an Interpreter is more than a compiler
 - o Because there is no any saved object file on the disk.
- The OS always need to wait for the Interpreter for execution of the code.
- If error occurs then interpreter creates machine instructions for all the line before the error.

Combining Both Interpreter and Compiler

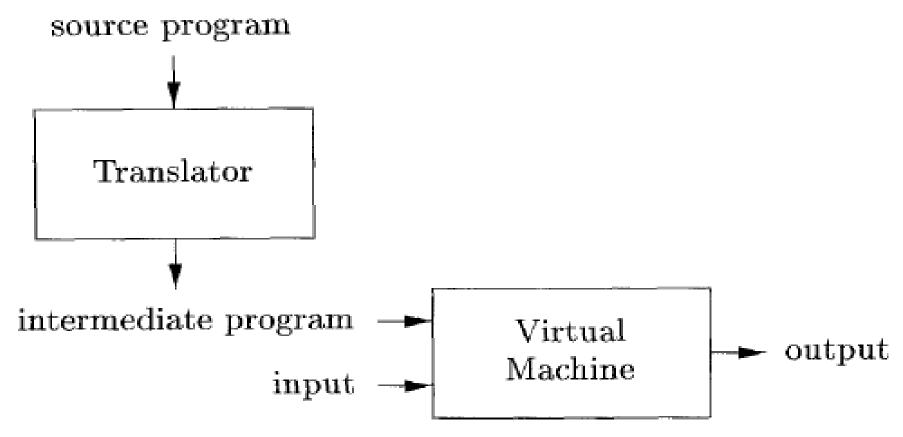


Fig: A hybrid compiler

Example: How its work

- Java language processors combine compilation and interpretation, as shown in figure.
- A Java source program may first be compiled into an intermediate form called bytecodes.
- The bytecodes are then interpreted by a virtual machine.
- In order to achieve faster processing of inputs to outputs, some *Java compilers*,
 called *just-in-time* compilers,
- translate the **bytecodes** into machine language immediately before they run the intermediate program to process the input.

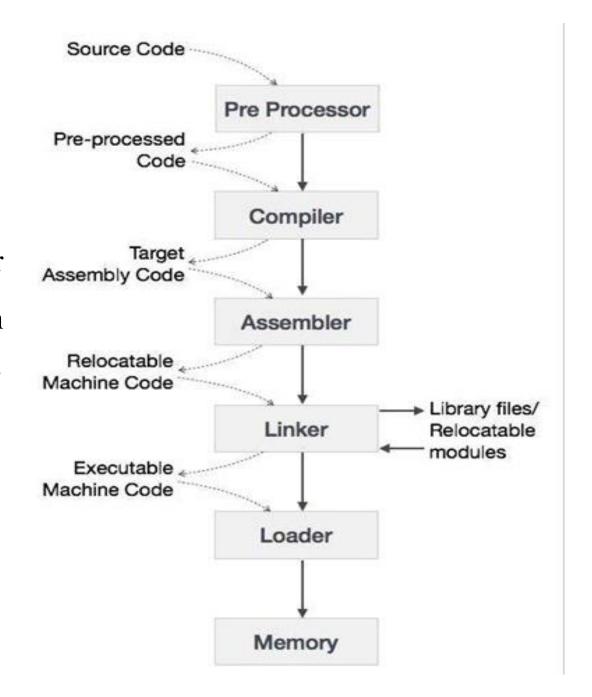
Example #1

Programming language use compiler

- \circ C++
- Programming language use Interpreter
 - o Python, perl
- Programming language use both
 - o Java

Cousins of a compiler

■ In addition to a *compiler*, several other programs may be required to create an *executable target program*, as shown in Fig



a. Pre-processor:

- is a program that processes its input data to produce output that is used as input to another program.
- The output is said to be a preprocessed form of the input data, which is often used by some subsequent programs like *compilers*.
- It includes all the header files and also evaluates if any macro is included.
- It is the optional because if any language which does not support #include and macro preprocessor is not required.

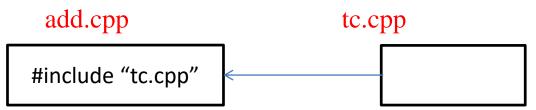
They may perform the following functions:

i. Macro processing

- o A **macro** is a rule or pattern that specifies how a certain input sequence should be mapped to an output sequence according to a defined procedure.
- o E.g C program

```
#define Cube(x);
Void main(){
Cout<<cube(3);</pre>
```

- ii. File Inclusion: Preprocessor includes header files into the program text.
 - When the preprocessor finds an **#include** directive it replaces it by the entire content of the specified file.



b. Compiler:

• It takes pure high level language as a input and convert into assembly code.

c. Assembler

• It takes the assembly code that has been generated by the compiler and convert it into *relocatable machine code*.

Example: Address binding

Mov a, b add c, 28

d. Linker:

- A **linker or link editor** is a program that takes one or more objects generated by a compiler and combines them into a single executable program.
- Also used to link all parts of program together for execution.
- Library routine or link to library

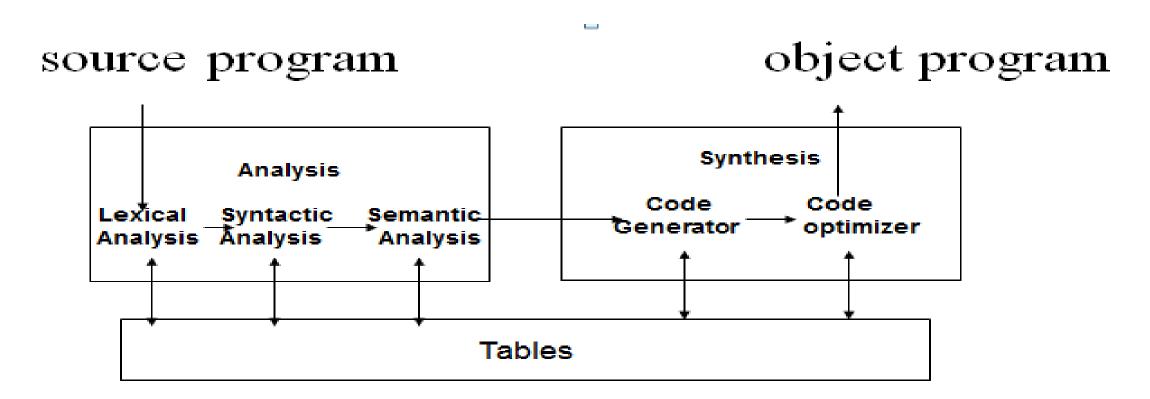
Header file
 printf();
 scanf();

e. Loader:

- A **loader** is the part of an operating system that is responsible for loading programs in memory.
- It loads the executable file into permanent storage.
 - That means: takes the entire program to main memory for execution. Example: exe.file
- **Note:** Here ,we only concentrate on Compiler
 - What exactly the compiler will do?
 - What is internal process of the compiler?

The Structure of a Compiler

Any compiler must perform two major tasks: such as Analysis and Synthesis as shown in fig



The Structure of a Compiler

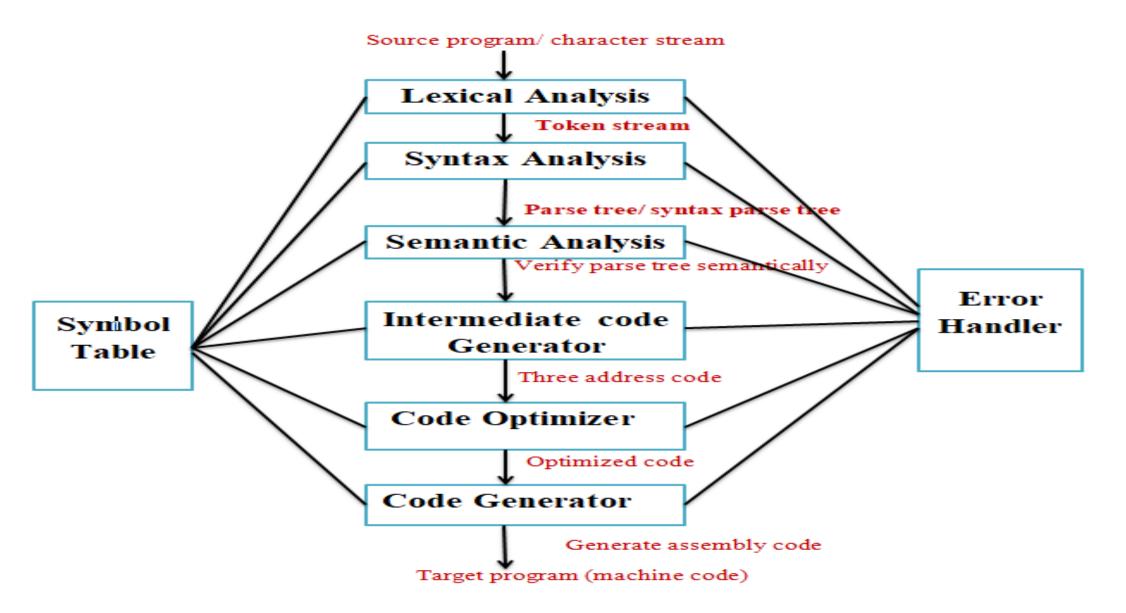
i. Analysis:

- It takes the input source program and breaks it into parts.
- o It then creates an intermediate representation of source program.
- o the *front end* of the compiler and Machine independent
- Lexical Analyzer, Syntax Analyzer and Semantic Analyzer are the parts of this phase.

ii. Synthesis:

- o It takes the intermediate representation as the input and creates the desired target program.
- o the *back end* of the compiler and Machine dependent
- Code Generator, and Code Optimizer are the parts of this phase.

Phases of a compiler



Phases of a compiler (cont'd.)

- Each phase transforms the source program from one representation into another representation.
- Each phase communicate with error handlers and symbol table.
- **Example:** How to compiler translate source program into target program.

The given Source program: Position = initial + rate * 60

• To answer the above example, let us discuss each phase of the compiler as follows:

Phase 1. Lexical Analysis (scanner)

- First phases of the compiler
- It reads the stream of characters making up the source program and groups the characters into meaningful sequences called *lexemes*.
- For each lexeme, the lexical analyzer produces a **token** of the form that it passes on to the subsequent phase, **syntax analysis**.

(token-name, attribute-value)

- Token-name: an abstract symbol is used during syntax analysis, an
- attribute-value: points to an entry in the symbol table for this token.

Example: what will do at lexical phases?

Position = initial + rate * 60

Cont'd

- "position" is a lexeme that is mapped into the *token (id, 1)*, where 1 points to the symbol-table entry for position.
- = is a lexeme that is mapped into the token (=).
- "initial" is a lexeme that is mapped into the *token (id, 2)*, where 2 points to the symbol-table entry for initial.
- + is a lexeme that is mapped into the **token** (+).
- "rate" is a lexeme mapped into the *token (id, 3)*, where 3 points to the symbol-table entry for rate.
- * is a lexeme that is mapped into the **token** (*).
- 60 is a lexeme that is mapped into the token (60)

Cont'd

- Blanks separating the lexemes would be discarded by the lexical analyzer.
- after lexical analysis as the sequence of tokens $\langle \mathbf{i} \mathbf{d}, \mathbf{l} \rangle \langle \mathbf{i} \mathbf{d},$
- In this representation, the token names =, +, and * are abstract symbols for the assignment, addition, and multiplication operators, respectively.

Phase 2. Syntax Analysis (parser tree)

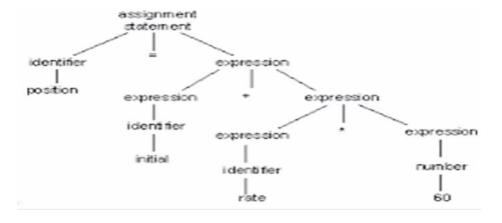
- The second phase of the compiler
- The parser uses the first components of the tokens produced by the lexical analyzer to create a tree
- -like intermediate representation that depicts the grammatical structure of the token stream.
- A typical representation is
 - o a syntax tree in which each interior node represents an operation and
 - o the **children** of the node represent the arguments of the operation.

Cont'd

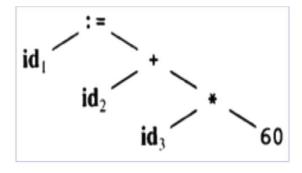
Rule of Syntax analysis:

- •Any identifier is an expression
- •Any number is an expression
- •If E1 and E2 are expression, then E1*E2 and E1+E2 is an expression
- •If id is an identifier and E is an expression then id=E

■This tree shows the order in which the operations in the assignment: Position = initial + rate * 60

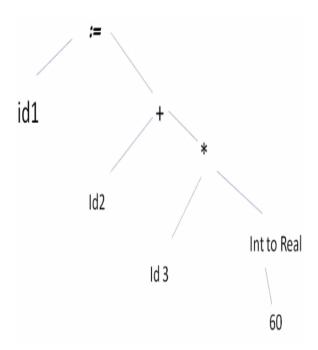


•Sequence of token produced form lexical Analysis: id1=id2+id3*60



Phase 3:- Semantic Analysis

- The semantic analyzer uses the **syntax tree**
- Performing type checking, where the compiler checks that each operator has matching operands.
- For example, many programming language definitions require an array index to be an integer;
- the compiler must report an error if a **floating-point number** is used to index an array.



Phase 4:- Intermediate Code Generation

- After syntax and semantic analysis of the source program, many compilers generate an explicit low-level or machine
- The considered intermediate form called three-address code, which consists of a sequence of assembly
- ICG-should have two important properties:
 - it should be easy to produce and
 - it should be easy to translate into the target machine.
- Each instructions has at most three operands
- Each operand can act like a register.

```
t1 = int to float (60)
t2 = id3 * t1
t3 = id2 + t2
id1 = t3
```

Phase 5:- Code Optimization

- The code-optimization phase attempts to improve the intermediate code so that generate better target code will result.
- Usually better means: faster, shorter code, or target code that consumes less power.
- The optimizer can deduce that the conversion of 60 from integer to floating point can be done once and for all at compile time,
- so the int to float operation can be eliminated by replacing the integer 60 by the floating-point number 60.0. Moreover, t3 is used only once.

$$t1 = id3 * 60.0$$

 $id1 = id2 + t1$

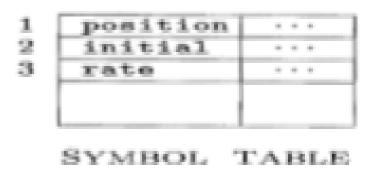
Phase 6:- Code Generation

- takes as input an intermediate representation of the source program and maps it into the target language
- If the target language is machine, code, registers or memory locations are selected for each of the variables used by the program.
- Then, the intermediate instructions are translated into sequences of machine instructions that perform the same task.
- A crucial aspect of code generation is the judicious assignment of registers to hold variables.

```
LDF R2, id3
MULF R2, R2, #60.0
LDF R1, id2
ADDF R1, R1, R2
STF id1, R1
```

7. Symbol-Table Management

- The symbol table is a data structure containing a record for each variable name, with fields for the attributes of the name.
- The data structure should be designed to allow the compiler to find the record for each name quickly and to store or retrieve data from that record quickly.

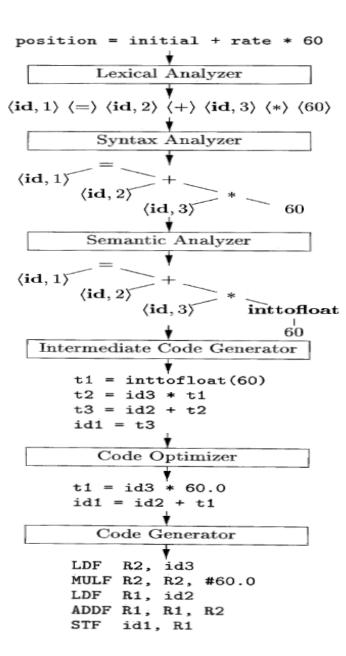


Example:

1	position	
2	initial	
3	rate	

SYMBOL TABLE

How all phases of compiler works together



Grouping of Compiler Phases

Front end

- Consist of those phases that depend on the source language but largely independent of the target machine.
- For example, the front-end phases of lexical analysis, syntax analysis, semantic analysis

Back end

 Consist of those phases that are usually target machine dependent such as code optimization and code generation.

Compiler-Construction Tools (CCTs)

Defining CCTS:

programs or environments that assist in the creation of an entire compiler or its parts.

Uses for CCTs to generate:

- o lexical analyzers, syntax analyzers, semantic analyzers, intermediate code, optimized target code.
- Some commonly used compiler-construction tools include:
 - Scanner generators, Parser generators Syntax-directed translation engines, Codegenerators, Data-flow analysis engines, Compiler-construction toolkits.

Reading Assignment

- 1. How do compiler works?
- 2. How do computers read a code (source program)?
- 3. What are the advantages of (a) a compiler over an interpreter (b) an interpreter over a compiler?
- 4. The Evolution of Programming Languages
- 5. Application of compiler technology
- 6. **Impacts on Compilers**
- 7. Programming language basics

Reading Assignment: Cont'd

- 9. What advantages are there to a language-processing system in which the compiler produces assembly language rather than machine language?
- 10. A compiler that translates a high-level language into another high-level language is called a *source-to-source* translator. What advantages are there to using C as a target language for a compiler?

Reading Assignment: Cont'd

12. Indicate which of the following terms: apply to which of the following languages:

- a) Imperative b) declarative c) von Neumann d) object-oriented e) functional
- f) third-generation g) fourth-generation h) scripting
 - 1) C 2) C++ 3) Cobol
 - 4) Fortran 5) Java 6) Lisp
 - 7) *ML* 8) *Perl*
 - 9) Python 10) VB.

End of Chapter- One

