

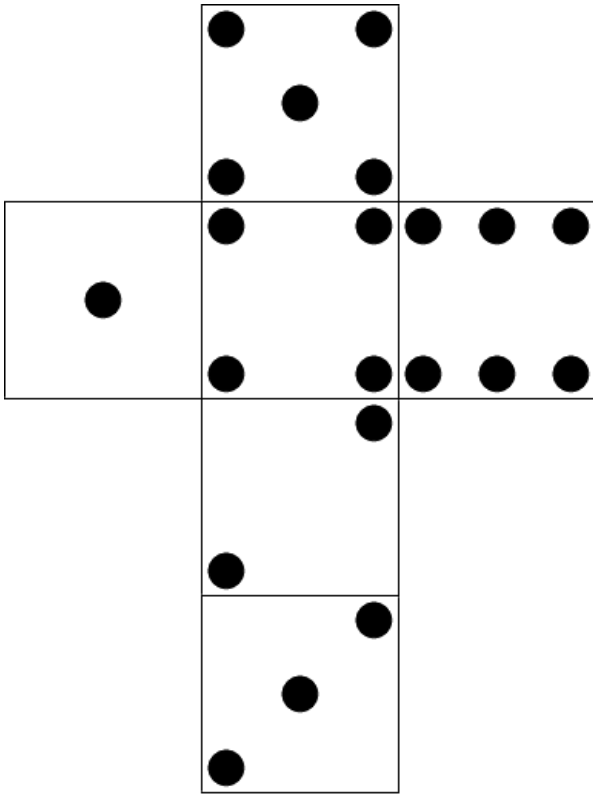
output: standard output

In the sixth example, note that  $\begin{pmatrix} 1 \\ 0 \end{pmatrix}$  is a valid solution.

## B. Dice Tower

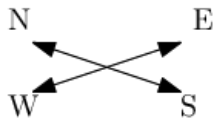
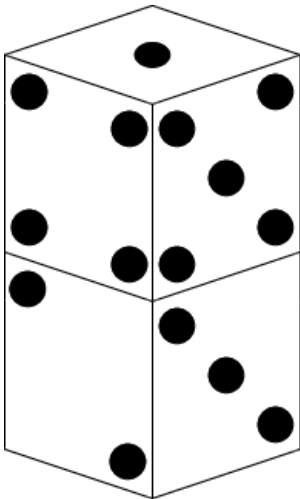
time limit per test: 1 second  
memory limit per test: 256 megabytes  
input: standard input  
output: standard output

Bob is playing with  $n$ -sided dice. A net of such standard cube is shown below.



He has an unlimited supply of these dice and wants to build a tower by stacking multiple dice on top of each other, while choosing the orientation of each dice. Then he counts the number of visible pips on the faces of the dice.

For example, the number of visible pips on the tower below is 18 — the number visible on the top is 1, from the south 2 and 3, from the west 2 and 3, from the north 2 and 3 and from the east 2 and 3.



The one at the bottom and the two sixes by which the dice are touching are not visible, so they are not counted towards total.

Bob also has  $k$  favourite integers  $a_1, a_2, \dots, a_k$ , and for every such integer his goal is to build such a tower that the number of visible pips is exactly  $a_i$ . For each of Bob's favourite integers determine whether it is possible to build a tower that has exactly that many visible pips.

### Input

The first line contains a single integer  $k$  ( $1 \leq k \leq 10^5$ ) — the number of favourite integers of Bob.

The second line contains  $k$  space-separated integers  $a_1, a_2, \dots, a_k$  ( $1 \leq a_i \leq 10^9$ ) — Bob's favourite integers.

### Output

For each of Bob's favourite integers, output "YES" if it is possible to build the tower, or "NO" otherwise (quotes for clarity).

### Example

input
4 29 34 19 38
output
YES YES YES NO

**Note**

The first example is mentioned in the problem statement.

In the second example, one can build the tower by flipping the top dice from the previous tower.

In the third example, one can use a single die that has 4 on top.

The fourth example is impossible.

### C. Diverse Matrix

time limit per test: 1 second  
memory limit per test: 256 megabytes  
input: standard input  
output: standard output

Let  $a$  be a matrix of size  $n \times m$  containing positive integers, not necessarily distinct. Rows of the matrix are numbered from 1 to  $n$ , columns are numbered from 1 to  $m$ . We can construct an array  $d$  consisting of  $n$  integers as follows: for each  $i$ , let  $d_i$  be the greatest common divisor of integers in the  $i$ -th row, and for each  $j$  let  $d_j$  be the greatest common divisor of integers in the  $j$ -th column.

We call the matrix **diverse** if all  $2m$  numbers  $(d_1, d_2, \dots, d_n, d_{n+1}, d_{n+2}, \dots, d_{n+m})$  are pairwise distinct.

The **magnitude** of a matrix equals to the maximum of  $d$ .

For example, suppose we have the following matrix:

We construct the array  $d$ :

1.  $d_1$  is the greatest common divisor of 12, 15, and 20, that is 1;
2.  $d_2$  is the greatest common divisor of 15, 12, and 20, that is 1;
3.  $d_3$  is the greatest common divisor of 12 and 15, that is 3;
4.  $d_4$  is the greatest common divisor of 12 and 20, that is 4;
5.  $d_5$  is the greatest common divisor of 15 and 20, that is 5.

So  $d = [1, 1, 3, 4, 5]$ . All values in this array are distinct, so the matrix is diverse. The magnitude is equal to 5.

For a given  $n$  and  $m$ , find a diverse matrix that minimises the magnitude. If there are multiple solutions, you may output any of them. If there are no solutions, output a single integer  $-1$ .

**Input**

The only line in the input contains two space separated integers  $n$  and  $m$  ( $1 \leq n, m \leq 10^5$ ) — the number of rows and the number of columns of the matrix to be found.

**Output**

If there is no solution, output a single integer  $-1$ .

Otherwise, output  $n$  rows. The  $i$ -th of them should contain  $m$  space-separated integers, the  $j$ -th of which is  $a_{ij}$  — the positive integer in the  $i$ -th row and  $j$ -th column of a diverse matrix minimizing the magnitude.

Furthermore, it must hold that  $a_{ij} \leq 2 \cdot \max(d_i, d_j)$ . It can be shown that if a solution exists, there is also a solution with this additional constraint (still having minimum possible magnitude).

#### Examples

input
2 2
output
4 12 2 9

input
1 1

output
0

**Note**  
 In the first example, the GCDs of rows are                    and                   , and the GCDs of columns are                    and                   . All GCDs are pairwise distinct and the maximum of them is                   . Since the GCDs have to be distinct and at least                   , it is clear that there are no diverse matrices of size                    with magnitude smaller than                   .

In the second example, no matter what                    is,                    will always hold, so there are no diverse matrices.

## D. Decreasing Debts

time limit per test: 2 seconds  
 memory limit per test: 256 megabytes  
 input: standard input  
 output: standard output

There are                    people in this world, conveniently numbered                    through                   . They are using burles to buy goods and services. Occasionally, a person might not have enough currency to buy what he wants or needs, so he borrows money from someone else, with the idea that he will repay the loan later with interest. Let                    denote the debt of                    towards                   , or                    if there is no such debt.

Sometimes, this becomes very complex, as the person lending money can run into financial troubles before his debtor is able to repay his debt, and finds himself in the need of borrowing money.

When this process runs for a long enough time, it might happen that there are so many debts that they can be consolidated. There are two ways this can be done:

- Let                    and                    such that                    or                   . We can decrease the                    and                    by                    and increase                    and                    by                   , where                   .
- Let                   . We can set                    to                   .

The total debt is defined as the sum of all debts:

Your goal is to use the above rules in any order any number of times, to make the total debt as small as possible. Note that you don't have to minimise the **number** of non-zero debts, only the **total debt**.

**Input**  
 The first line contains two space separated integers                    (                   ) and                    (                   ), representing the number of people and the number of debts, respectively.

                  lines follow, each of which contains three space separated integers                   ,                    (                   ),                    (                   ), meaning that the person                    borrowed                    burles from person                   .

**Output**  
 On the first line print an integer                    (                   ), representing the number of debts after the consolidation. It can be shown that an answer always exists with this additional constraint.

After that print                    lines, -th of which contains three space separated integers                   , meaning that the person                    owes the person                    exactly                    burles. The output must satisfy                   ,                    and                   .

For each pair                   , it should hold that                    or                   . In other words, each pair of people can be included at most once in the output.

### Examples

input
3 2 1 2 10 2 3 5
output
2 1 2 5 1 3 5

input
3 3 1 2 10 2 3 15 3 1 10
output
1

2 3 5
input
4 2 1 2 12 3 4 8
output
2 1 2 12 3 4 8
input
3 4 2 3 1 2 3 2 2 3 4 2 3 8
output
1 2 3 15

**Note**  
In the first example the optimal sequence of operations can be the following:

1. Perform an operation of the first type with 4, 2, 12 and 3. The resulting debts are: 2, 12, 3, all other debts are 0;
2. Perform an operation of the second type with 2. The resulting debts are: 0, 12, 3, all other debts are 0.

In the second example the optimal sequence of operations can be the following:

1. Perform an operation of the first type with 3, 4, 1 and 2. The resulting debts are: 0, 1, 2, all other debts are 0;
2. Perform an operation of the first type with 2, 3, 2 and 1. The resulting debts are: 0, 1, 2, all other debts are 0;
3. Perform an operation of the second type with 1. The resulting debts are: 0, 1, 2, all other debts are 0;
4. Perform an operation of the second type with 2. The resulting debts are: 0, 1, 2, all other debts are 0;
5. Perform an operation of the second type with 1. The resulting debts are: 0, 1, 2, all other debts are 0.

## E. Spaceship Solitaire

time limit per test: 1 second  
memory limit per test: 256 megabytes  
input: standard input  
output: standard output

Bob is playing a game of Spaceship Solitaire. The goal of this game is to build a spaceship. In order to do this, he first needs to accumulate enough resources for the construction. There are  $n$  types of resources, numbered  $1$  through  $n$ . Bob needs at least  $g_i$  pieces of the  $i$ -th resource to build the spaceship. The number  $g_i$  is called **the goal for resource  $i$** .

Each resource takes  $t_i$  turn to produce and in each turn only one resource can be produced. However, there are certain milestones that speed up production. Every milestone is a triple  $(i, j, k)$ , meaning that as soon as Bob has  $k$  units of the resource  $i$ , he receives one unit of the resource  $j$  for free, without him needing to spend a turn. It is possible that getting this free resource allows Bob to claim reward for another milestone. This way, he can obtain a large number of resources in a single turn.

The game is constructed in such a way that there are never two milestones that have the same  $i$  and  $k$ , that is, the award for reaching  $k$  units of resource  $i$  is at most one additional resource.

**A bonus is never awarded for  $i$  of any resource, neither for reaching the goal  $g_i$  nor for going past the goal — formally, for every milestone  $(i, j, k)$ .**

A bonus for reaching certain amount of a resource can be the resource itself, that is,  $j = i$ .

Initially there are no milestones. You are to process  $m$  updates, each of which adds, removes or modifies a milestone. After every update, output the minimum number of turns needed to finish the game, that is, to accumulate at least  $g_i$  of  $i$ -th resource for each  $i$ .

**Input**  
The first line contains a single integer  $n$  ( $1 \leq n \leq 10^5$ ) — the number of types of resources.  
  
The second line contains  $n$  space separated integers  $g_1, g_2, \dots, g_n$  ( $1 \leq g_i \leq 10^5$ ), the  $i$ -th of which is the goal for the  $i$ -th resource.

The third line contains a single integer  $k$  — the number of updates to the game milestones.

Then  $m$  lines follow, the  $i$ -th of which contains three space separated integers  $a_i, b_i, c_i$  ( $0 \leq a_i, b_i, c_i \leq 10^9$ ).

For each triple, perform the following actions:

- First, if there is already a milestone for obtaining  $a_i$  units of resource  $b_i$ , it is removed.
- If  $a_i = 0$ , no new milestone is added.
- If  $a_i > 0$ , add the following milestone: "For reaching  $a_i$  units of resource  $b_i$ , gain one free piece of  $c_i$ ."
- Output the minimum number of turns needed to win the game.

Output

Output  $m$  lines, each consisting of a single integer, the  $i$ -th represents the answer after the  $i$ -th update.

Example

input
2 2 3 5 2 1 1 2 2 1 1 1 1 2 1 2 2 2 0
output
4 3 3 2 3

Note

After the first update, the optimal strategy is as follows. First produce  $b_1$  once, which gives a free resource  $c_1$ . Then, produce  $b_1$  twice and  $b_2$  once, for a total of four turns.

After the second update, the optimal strategy is to produce  $b_1$  three times — the first two times a single unit of resource  $c_1$  is also granted.

After the third update, the game is won as follows.

- First produce  $b_1$  once. This gives a free unit of  $c_1$ . This gives additional bonus of resource  $c_2$ . After the first turn, the number of resources is thus  $(1, 1, 1)$ .
- Next, produce resource  $b_2$  again, which gives another unit of  $c_2$ .
- After this, produce one more unit of  $b_1$ .

The final count of resources is  $(3, 2, 2)$ , and three turns are needed to reach this situation. Notice that we have more of resource  $c_2$  than its goal, which is of no use.

F. Almost Same Distance

time limit per test: 5 seconds  
memory limit per test: 256 megabytes  
input: standard input  
output: standard output

Let  $G$  be a simple graph. Let  $S$  be a non-empty subset of vertices. Then  $S$  is **almost- $k$ -uniform** if for each pair of distinct vertices  $u, v \in S$  the distance between  $u$  and  $v$  is either  $k$  or  $k+1$ .

You are given a tree on  $n$  vertices. For each  $k$  between  $1$  and  $n-1$ , find the maximum size of an almost- $k$ -uniform set.

Input

The first line contains a single integer  $n$  ( $2 \leq n \leq 10^5$ ) – the number of vertices of the tree.

Then  $n-1$  lines follows, the  $i$ -th of which consisting of two space separated integers  $u_i, v_i$  ( $1 \leq u_i, v_i \leq n$ ) meaning that there is an edge between vertices  $u_i$  and  $v_i$ .

It is guaranteed that the given graph is tree.

Output

Output a single line containing  $n$  space separated integers  $ans_1, ans_2, \dots, ans_n$ , where  $ans_k$  is the maximum size of an almost- $k$ -uniform set.

Examples

input
5 1 2 1 3 1 4 4 5

output
4 3 2 1 1

input
6 1 2 1 3 1 4 4 5 4 6
output
4 4 2 1 1 1

**Note**  
Consider the first example.

- The only maximum almost- $k$ -uniform set is  $\{1, 2, 3, 4, 5, 6\}$ .
- One of the maximum almost- $k$ -uniform sets is  $\{1, 2, 3, 4, 5\}$ , another one is  $\{1, 2, 3, 4, 6\}$ .
- A maximum almost- $k$ -uniform set is any pair of vertices on distance  $k$ .
- Any single vertex is an almost- $k$ -uniform set for  $k \leq 6$ .

In the second sample there is an almost- $k$ -uniform set of size  $k$ , and that is  $\{1, 2, 3, 4, 5, 6\}$ .

## G. Permutation Concatenation

time limit per test: 1 second  
memory limit per test: 256 megabytes  
input: standard input  
output: standard output

Let  $n$  be an integer. Consider all permutations on integers  $1$  to  $n$  in lexicographic order, and concatenate them into one big sequence  $S$ . For example, if  $n = 3$ , then  $S = 123132123321$ . The length of this sequence is  $n!$ .

Let  $i$  and  $j$  be a pair of indices. We call the sequence  $S[i..j]$  a **subarray** of  $S$ .

You are given  $i$  and  $j$ . Find the number of distinct subarrays of  $S$  between  $i$  and  $j$ . Since this number may be large, output it modulo  $998244353$  (a prime number).

**Input**  
The only line contains one integer  $n$  ( $1 \leq n \leq 10^5$ ), as described in the problem statement.

**Output**  
Output a single integer — the number of distinct subarrays, modulo  $998244353$ .

input
2
output
8

input
10
output
19210869

**Note**  
In the first example, the sequence  $S$  is  $123132123321$ . It has eight distinct subarrays:  $\{1\}$ ,  $\{2\}$ ,  $\{3\}$ ,  $\{1, 2\}$ ,  $\{2, 3\}$ ,  $\{1, 3\}$ ,  $\{1, 2, 3\}$ , and  $\{1, 2, 3, 1\}$ .

## H. Red-Blue Graph

time limit per test: 4 seconds  
memory limit per test: 512 megabytes  
input: standard input  
output: standard output

There is a directed graph on  $n$  vertices numbered  $1$  through  $n$  where each vertex (except  $1$ ) has two outgoing arcs, red and blue. At any point in time, exactly one of the arcs is *active* for each vertex. Initially, all blue arcs are active and there is a token located at vertex  $1$ . In one second, the vertex with token first switches its active arcs — the inactive arc becomes active and vice versa. Then, the token is moved along the active arc. When the token reaches the vertex  $n$ , it stops. It is guaranteed that  $n$  is reachable via arcs

from every vertex.

You are given  $q$  queries. Each query contains a state of the graph — a pair  $(v, c)$  of the following form:

- $v$  is the vertex where the token is currently located;
- $c$  is a string consisting of  $n$  characters. The  $i$ -th character corresponds to the color of the active edge leading from the  $i$ -th vertex (the character is 'R' if red arc is active, otherwise the character is 'B').

For each query, determine whether the given state is reachable from the initial state and the first time this configuration appears. Note that the two operations (change active arc and traverse it) are atomic — a state is not considered reached if it appears after changing the active arc but before traversing it.

Input

The first line contains a single integer  $n$  (  $1 \leq n \leq 10^5$  ) — the number of vertices.

$n$  lines follow,  $i$ -th of contains two space separated integers  $a_i$  and  $b_i$  (  $1 \leq a_i, b_i \leq n$  ) representing a blue arc  $a_i \rightarrow b_i$  and red arc  $b_i \rightarrow a_i$ , respectively. It is guaranteed that vertex  $1$  is reachable from every vertex.

The next line contains a single integer  $q$  (  $1 \leq q \leq 10^5$  ) — the number of queries.

Then  $q$  lines with queries follow. The  $i$ -th of these lines contains an integer  $v_i$  (  $1 \leq v_i \leq n$  ) and a string  $c_i$  of length  $n$  consiting only of characters 'R' and 'B'. The  $j$ -th of these characters is 'R' if the red arc going from  $j$  is active and 'B' otherwise.

Output

Output  $q$  lines, each containing answer to a single query.

If the state in the  $i$ -th query is unreachable, output the integer  $-1$ . Otherwise, output  $t_i$  — the first time when the state appears (measured in seconds, starting from the initial state of the graph which appears in time  $0$ ).

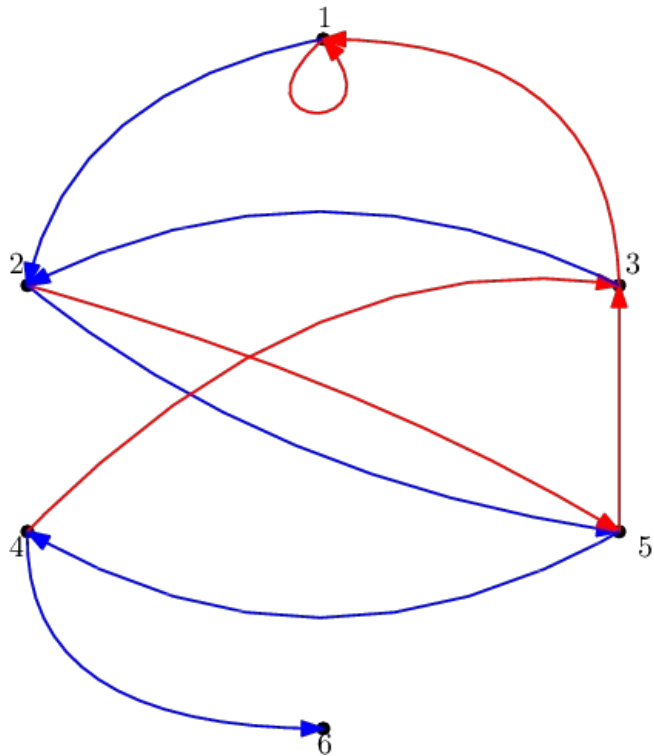
Example

input
6 2 1 5 5 2 1 6 3 4 3 21 1 BBBBB 1 RBBBB 2 BBBBB 5 BRBBB 3 BRBBR 1 BRRBR 1 RRRBR 2 BRRBR 5 BBRBR 4 BBRBB 3 BBRRB 2 BBRRB 5 BRBRB 3 BRBRR 1 BRRRR 1 RRRRR 2 BRRRR 5 BBRRR 4 BBRRB 2 BRBBB 4 BRBBR
output
0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 -1 -1

Note

The graph in the first example is depected in the figure below.





The first 4 queries denote the journey of the token. On the 5-th move the token would reach the vertex 3. The last two queries show states that are unreachable.