Exploring different merging algorithms for balanced trees and their time complexity optimization.

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1 Introduction

A data structure is a way to store and organize data in order to facilitate access and modifications (Cormen, Leiserson, Rivest, & Stein, 2022). Designing and choosing more efficient data structures has always been a great persuit for computer scientists, for optimal data structures can save huge amount of computing resources, especially in face of large amount of data. Basic data structures include ordered data structures like arrays, linked lists and binary search trees and unordered data structures like hashtables.

For ordered data structures, merging two or more instances of them while maintaining its ordered property may be frequently used in practice. For example, to investigate the factors affacting the school grade, data from different schools may be grouped and merged according to various factors. The efficiency of combination varies significantly based on the data structure itself and the algorithm used in the process.

This essay will focus on investigating the theoratical time complexity (need definitions aa) and actual performance of merging algorithms of different data structures, namely arrays, and BSTs, which are the most commonly used data structure in real life.

Research question: How does different algorithm affect the efficiency of merging two instances of ordered data structures?

2 Theory

2.1 Data structure terminology

When a homogeneous relation (a binary relation between two elements) \leq on a set of elements X satisfies:

1. Antisymmetry: $\forall u, v \in X, (u \le v \land v \le u) \Leftrightarrow u = v.$

2. Totality: $\forall u, v \in X, u \leq v \lor v \leq u$.

3. Transitivity: $\forall u, v, w \in X, (u \le v \land v \le w) \Rightarrow u \le w.$

We say $P = (X, \leq)$ is a total order. For example $P = (\mathbb{R}, \leq)$, where \leq is numerical comparison, is a total order. But $P = (\{S : S \subset \mathbb{R}\}, \subset)$ is not a total order.

Ordered data structures can store elements that satisfies a total order while maintaining their order. In C++, arrays, vectors, linked lists and sets can be ordered data structures, but unordered sets are not ordered data structures.

Definition

Ordered data structures are data structures that can store elements that satisfies a total order while maintaining their order.

2.2 Optimality

When merging two instances of size n and m respectively, there are in total $\binom{n+m}{n}$ possible outcomes. According to the decision tree theory, each of them corresponds to a decision tree leaf node. Since the merging algorithm is comparison based, the decision tree has to be a binary tree (i.e. Each node has at most two children). The height of the decision tree is therefore no lower than $O(\log_2(\binom{n+m}{n}))$.

According to Sterling's approximation,

$$n! \approx \sqrt{2\pi n} \left(\frac{n}{e}\right)^n \tag{1}$$

which means

$$O(\log(n!)) = O(\frac{1}{2}\log(2\pi n) + n\log n - n\log e) = O(n\log n)$$
 (2)

Using the definition of combination number,

$$\binom{n+m}{n} = \frac{(n+m)!}{n!m!} \tag{3}$$

which is approximately $O(n \log_2(\frac{n}{m}))$ (provided $n \leq m$).

Theorem

The optimal time complexity of merging two instances of ordered data structures is $O(n \log(\frac{n}{m}))$, provided the moving and comparision cost is O(1).

Appendix

Listing 1: AvlSet

```
#include <vector>
  #include <algorithm>
  #include <memory>
4 #include <list>
  #include <functional>
  #include <stack>
  template <typename T, typename Compare = std::less<T>>
  class AVLSet {
  private:
      struct Node {
11
          T key;
12
          Node* left;
          Node* right;
14
          int height;
          template <typename... Args>
          Node(Args&&... args)
              : key(std::forward<Args>(args)...),
                left(nullptr),
                right(nullptr),
21
                height(1) {}
22
      };
23
      Node* root;
      size_t size;
26
      Compare comp;
27
      // Memory pool management
29
      std::list<Node> node_storage;
30
      std::vector<Node*> free_nodes;
      // Helper functions for merging
      int height(Node* node) const {
34
          return node ? node->height : 0;
35
      }
36
```

```
// Memory pool operations
38
      Node* create_node(const T& key) {
39
          if (!free_nodes.empty()) {
40
             Node* node = free_nodes.back();
             free_nodes.pop_back();
              *node = Node(key);
             return node;
          }
          node_storage.emplace_back(key);
46
          return &node_storage.back();
47
      }
      Node* create_node(T&& key) {
          if (!free_nodes.empty()) {
             Node* node = free_nodes.back();
             free_nodes.pop_back();
53
             *node = Node(std::move(key));
             return node;
          }
          node_storage.emplace_back(std::move(key));
          return &node_storage.back();
      }
59
      // Generates a balanced subtree in O(n) time out from a
61
          ordered sequence.
      // Returns the root node pointer.
      // *Preconditions
      // keys have to be ordered
      // _RandAccIt is the random access iterator
      // (*bg) and (*ed) should be of type T
66
67
      template<typename _RandAccIt>
68
      Node* build(const _RandAccIt& bg, const _RandAccIt& ed){
69
          if(bg == ed) return nullptr;
          auto it = bg;
          if(++it == ed) return create_node(*bg);
          it = bg + (ed - bg) / 2; // The same as (bg + ed)/2 but
73
             avoids overflow problems
          auto cur = create_node(*it);
```

```
cur->left = build(bg, it);
75
            cur->right = build(it + 1, ed);
76
            update_height(cur);
            return cur;
78
        }
80
        void recycle_node(Node* node) {
81
            free_nodes.push_back(node);
82
        }
83
84
       void update_height(Node* node) {
85
            node->height = 1 + std::max(height(node->left), height(
                node->right));
        }
87
88
        Node* rotate_right(Node* y) {
89
            Node* x = y \rightarrow left;
90
            Node* T2 = x->right;
91
92
            x->right = y;
            y \rightarrow left = T2;
95
            update_height(y);
96
            update_height(x);
97
98
            return x;
99
        }
100
101
        Node* rotate_left(Node* x) {
102
            Node* y = x->right;
103
            Node* T2 = y->left;
104
            y \rightarrow left = x;
106
            x->right = T2;
107
            update_height(x);
            update_height(y);
110
111
            return y;
112
        }
113
```

```
114
       int balance_factor(Node* node) const {
115
           return node ? height(node->left) - height(node->right) :
116
               0;
       }
117
118
       Node* balance(Node* node) {
119
           update_height(node);
120
           int bf = balance_factor(node);
121
           // Left Heavy
123
           if (bf > 1) {
               if (balance_factor(node->left) < 0)</pre>
125
                   node->left = rotate_left(node->left);
               return rotate_right(node);
127
           }
128
           // Right Heavy
129
           if (bf < -1) {</pre>
130
               if (balance_factor(node->right) > 0)
                   node->right = rotate_right(node->right);
               return rotate_left(node);
           }
134
           return node;
135
       }
136
137
       template <typename Func>
138
       void traverse_in_order(Node* node, Func f) const {
139
           if (!node) return;
           std::stack<Node*> stack;
141
           Node* current = node;
142
           while (current || !stack.empty()) {
143
               while (current) {
144
                   stack.push(current);
145
                   current = current->left;
146
               }
               current = stack.top();
               stack.pop();
149
               f(current->key);
150
               current = current->right;
151
           }
152
```

```
}
154
       void recycle_tree(Node* node) {
155
           if (!node) return;
156
           std::stack<Node*> stack;
           Node* current = node;
158
           Node* last_visited = nullptr;
160
           while (current || !stack.empty()) {
161
               if (current) {
162
                   stack.push(current);
163
                   current = current->left;
               } else {
165
                   Node* top = stack.top();
166
                   if (top->right && top->right != last_visited) {
167
                       current = top->right;
168
                   } else {
169
                       recycle_node(top);
170
                       last_visited = top;
                       stack.pop();
                   }
               }
174
           }
175
       }
176
177
178
   public:
179
       AVLSet() : root(nullptr), size(0), comp(Compare()) {}
181
       // Move operations
182
       AVLSet(AVLSet&& other) noexcept
183
           : root(other.root),
184
             size(other.size),
185
             node_storage(std::move(other.node_storage)),
186
             free_nodes(std::move(other.free_nodes)) {
           other.root = nullptr;
           other.size = 0;
189
       }
190
191
       AVLSet& operator=(AVLSet&& other) noexcept {
192
```

```
if (this != &other) {
193
               clear();
194
               root = other.root;
195
               size = other.size;
196
               node_storage = std::move(other.node_storage);
               free_nodes = std::move(other.free_nodes);
198
               other.root = nullptr;
199
               other.size = 0;
200
           }
201
           return *this;
202
       }
203
204
       // Disable copy operations
       AVLSet(const AVLSet&) = delete;
206
       AVLSet& operator=(const AVLSet&) = delete;
207
208
       void clear() {
209
           recycle_tree(root);
210
           node_storage.clear();
211
           free_nodes.clear();
212
           root = nullptr;
           size = 0;
214
       }
215
216
       bool empty() const {
217
           return size == 0;
218
219
       size_t get_size() const {
221
222
           return size;
       }
223
224
       template <typename Func>
225
       void traverse_in_order(Func f) const {
226
           traverse_in_order(root, f);
       }
228
       std::vector<T> items() const {
230
           std::vector<T> result;
231
           traverse_in_order([&result](const T& key) {
232
```

```
result.push_back(key);
233
           });
234
           return result;
235
       }
236
       void swap_with(AVLSet& other) {
           std::swap(root, other.root);
           std::swap(size, other.size);
240
           std::swap(node_storage, other.node_storage);
241
           std::swap(free_nodes, other.free_nodes);
242
       }
243
244
       /* Merge two sets in O(N+M) time.
       * Some additional space may be costed.
247
       * But it does not affect the result of the experiment.
248
249
       void linearmerge(AVLSet&& other) {
           if (other.empty()) return;
           other.free_nodes.clear();
255
           // Get elements from both trees
256
           std::vector<T> q1 = items();
257
           std::vector<T> q2 = other.items();
258
           std::vector<T> all_elements;
259
           all_elements.reserve(q1.size() + q2.size());
261
           // Merge sorted vectors
262
           std::merge(q1.begin(), q1.end(), q2.begin(), q2.end(),
263
                     std::back_inserter(all_elements), comp);
264
265
           // Recycle both trees
266
           recycle_tree(root);
           recycle_tree(other.root);
           root = nullptr;
269
           other.root = nullptr;
270
           size = 0;
271
           other.size = 0;
272
```

```
273
           // Build new tree
274
           root = build(all_elements.begin(), all_elements.end());
275
           size = all_elements.size();
276
       }
       void simplemerge(AVLSet&& other) {
279
           if (other.empty()) return;
280
281
           if (size < other.size) { swap_with(other); }</pre>
282
283
           // Insert all elements from the smaller tree (now 'other')
                into this
           other.traverse_in_order([this](const T& key) {
               this->insert(key);
286
           });
287
288
           other.clear();
289
       }
290
291
       private:
292
293
       Node* insert(Node* node, const T& key) {
294
           if (!node) {
295
               size++;
296
               return create_node(key);
297
           }
298
           if (comp(key, node->key)) {
300
               node->left = insert(node->left, key);
301
           } else if (comp(node->key, key)) {
302
               node->right = insert(node->right, key);
303
           } else {
304
               return node;
305
           }
           return balance(node);
308
       }
309
310
       public:
311
```

```
void insert(const T& val){
312
           root = insert(root, val);
313
       }
314
       void remove(const T& val){
315
           root = remove(root, val);
       }
317
       template<typename RandAccIt>
318
       void construct(RandAccIt bg, RandAccIt ed){
319
           // Clear the current tree
320
           recycle_tree(root);
321
           root = nullptr;
322
           // Build new tree
           root = build(bg, ed);
           size = static_cast<size_t>(ed - bg);
325
       }
326
   };
327
```

References

Cormen, T. H., Leiserson, C. E., Rivest, R. L., & Stein, C. (2022). *Introduction to algorithms* (Fourth ed.). MIT Press.