

# Exploring different merging algorithms for balanced trees and their time complexity optimization.

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July 23, 2025

word count: ???

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# 1 Introduction

A *data structure* is a way to store and organize data in order to facilitate access and modifications (Cormen, Leiserson, Rivest, & Stein, 2022). Designing and choosing more efficient data structures has always been a great pursuit for computer scientists, for optimal data structures can save huge amount of computing resources, especially in face of large amount of data. Basic data structures include ordered data structures like arrays, linked lists and binary search trees and unordered data structures like hashtables.

For ordered data structures, merging two or more instances of them while maintaining its ordered property may be frequently used in practice. For example, to investigate the factors affecting the school grade, data from different schools may be grouped and merged according to various factors. The efficiency of combination varies significantly based on the data structure itself and the algorithm used in the process.

This essay will focus on investigating the theoretical time complexity (need definitions aa) and actual performance of merging algorithms of different data structures, namely arrays, and BSTs, which are the most commonly used data structure in real life.

**Research question:** How does different algorithm affect the efficiency of merging two instances of ordered data structures?

## 2 Theory

### 2.1 Data structure terminology

When a homogeneous relation (a binary relation between two elements)  $\leq$  on a set of elements  $X$  satisfies:

1. Antisymmetry:  $\forall u, v \in X, (u \leq v \wedge v \leq u) \Leftrightarrow u = v.$
2. Totality:  $\forall u, v \in X, u \leq v \vee v \leq u.$
3. Transitivity:  $\forall u, v, w \in X, (u \leq v \wedge v \leq w) \Rightarrow u \leq w.$

We say  $P = (X, \leq)$  is a total order. For example  $P = (\mathbb{R}, \leq)$ , where  $\leq$  is numerical comparison, is a total order. But  $P = (\{S : S \subset \mathbb{R}\}, \subset)$  is not a total order.

Ordered data structures can store elements that satisfies a total order while maintaining their order.

## 2.2 Optimality

When merging two instances of size  $n$  and  $m$  respectively, there are in total  $\binom{n+m}{n}$  possible outcomes. According to the decision tree theory, each of them corresponds to a decision tree leaf node. Since the merging algorithm is comparison based, the decision tree has to be a binary tree (i.e. Each node has at most two children). The height of the decision tree is therefore no lower than  $O(\log_2(\binom{n+m}{n}))$ .

According to Sterling's approximation,

$$n! \approx \sqrt{2\pi n} \left(\frac{n}{e}\right)^n \quad (1)$$

which means

$$O(\log(n!)) = O\left(\frac{1}{2} \log(2\pi n) + n \log n - n \log e\right) = O(n \log n) \quad (2)$$

Using the definition of combination number,

$$\binom{n+m}{n} = \frac{(n+m)!}{n!m!} \quad (3)$$

which is approximately  $O(n \log_2(\frac{n}{m}))$  (provided  $n \leq m$ ).

## Appendix

Listing 1: AvlSet

```

1 #include <vector>
2 #include <algorithm>
3 #include <memory>
4 #include <list>
5 #include <functional>
6 #include <stack>
7
8 template <typename T, typename Compare = std::less<T>>
```

```

9  class AVLSet {
10 private:
11     struct Node {
12         T key;
13         Node* left;
14         Node* right;
15         int height;
16
17         template <typename... Args>
18         Node(Args&&... args)
19             : key(std::forward<Args>(args)...),
20               left(nullptr),
21               right(nullptr),
22               height(1) {}
23     };
24
25     Node* root;
26     size_t size;
27     Compare comp;
28
29     // Memory pool management
30     std::list<Node> node_storage;
31     std::vector<Node*> free_nodes;
32
33     // Helper functions for merging
34     int height(Node* node) const {
35         return node ? node->height : 0;
36     }
37
38     // Memory pool operations
39     Node* create_node(const T& key) {
40         if (!free_nodes.empty()) {
41             Node* node = free_nodes.back();
42             free_nodes.pop_back();
43             *node = Node(key);
44             return node;
45         }
46         node_storage.emplace_back(key);
47         return &node_storage.back();
48     }

```

```

49
50 Node* create_node(T&& key) {
51     if (!free_nodes.empty()) {
52         Node* node = free_nodes.back();
53         free_nodes.pop_back();
54         *node = Node(std::move(key));
55         return node;
56     }
57     node_storage.emplace_back(std::move(key));
58     return &node_storage.back();
59 }
60
61 // Generates a balanced subtree in O(n) time out from a
62 // ordered sequence.
63 // Returns the root node pointer.
64 // *Preconditions
65 // keys have to be ordered
66 // _RandAccIt is the random access iterator
67 // (*bg) and (*ed) should be of type T
68
69 template<typename _RandAccIt>
70 Node* build(const _RandAccIt& bg, const _RandAccIt& ed){
71     if(bg == ed) return nullptr;
72     auto it = bg;
73     if(++it == ed) return create_node(*bg);
74     it = bg + (ed - bg) / 2; // The same as (bg + ed)/2 but
75     // avoids overflow problems
76     auto cur = create_node(*it);
77     cur->left = build(bg, it);
78     cur->right = build(it + 1, ed);
79     update_height(cur);
80     return cur;
81 }
82
83 void recycle_node(Node* node) {
84     free_nodes.push_back(node);
85 }
86
87 void update_height(Node* node) {

```

```

86     node->height = 1 + std::max(height(node->left), height(
      node->right));
87 }
88
89 Node* rotate_right(Node* y) {
90     Node* x = y->left;
91     Node* T2 = x->right;
92
93     x->right = y;
94     y->left = T2;
95
96     update_height(y);
97     update_height(x);
98
99     return x;
100 }
101
102 Node* rotate_left(Node* x) {
103     Node* y = x->right;
104     Node* T2 = y->left;
105
106     y->left = x;
107     x->right = T2;
108
109     update_height(x);
110     update_height(y);
111
112     return y;
113 }
114
115 int balance_factor(Node* node) const {
116     return node ? height(node->left) - height(node->right) :
      0;
117 }
118
119 Node* balance(Node* node) {
120     update_height(node);
121     int bf = balance_factor(node);
122
123     // Left Heavy

```

```

124     if (bf > 1) {
125         if (balance_factor(node->left) < 0)
126             node->left = rotate_left(node->left);
127         return rotate_right(node);
128     }
129     // Right Heavy
130     if (bf < -1) {
131         if (balance_factor(node->right) > 0)
132             node->right = rotate_right(node->right);
133         return rotate_left(node);
134     }
135     return node;
136 }
137
138 template <typename Func>
139 void traverse_in_order(Node* node, Func f) const {
140     if (!node) return;
141     std::stack<Node*> stack;
142     Node* current = node;
143     while (current || !stack.empty()) {
144         while (current) {
145             stack.push(current);
146             current = current->left;
147         }
148         current = stack.top();
149         stack.pop();
150         f(current->key);
151         current = current->right;
152     }
153 }
154
155 void recycle_tree(Node* node) {
156     if (!node) return;
157     std::stack<Node*> stack;
158     Node* current = node;
159     Node* last_visited = nullptr;
160
161     while (current || !stack.empty()) {
162         if (current) {
163             stack.push(current);

```



```

164         current = current->left;
165     } else {
166         Node* top = stack.top();
167         if (top->right && top->right != last_visited) {
168             current = top->right;
169         } else {
170             recycle_node(top);
171             last_visited = top;
172             stack.pop();
173         }
174     }
175 }
176 }
177
178 void insert_with_path(std::vector<Node*>& path, const T& key)
179 {
180 }
181
182 public:
183     AVLSet() : root(nullptr), size(0), comp(Compare()) {}
184
185     // Move operations
186     AVLSet(AVLSet&& other) noexcept
187         : root(other.root),
188           size(other.size),
189           node_storage(std::move(other.node_storage)),
190           free_nodes(std::move(other.free_nodes)) {
191         other.root = nullptr;
192         other.size = 0;
193     }
194
195     AVLSet& operator=(AVLSet&& other) noexcept {
196         if (this != &other) {
197             clear();
198             root = other.root;
199             size = other.size;
200             node_storage = std::move(other.node_storage);
201             free_nodes = std::move(other.free_nodes);
202             other.root = nullptr;

```

```

203         other.size = 0;
204     }
205     return *this;
206 }
207
208 // Disable copy operations
209 AVLSet(const AVLSet&) = delete;
210 AVLSet& operator=(const AVLSet&) = delete;
211
212 void clear() {
213     recycle_tree(root);
214     node_storage.clear();
215     free_nodes.clear();
216     root = nullptr;
217     size = 0;
218 }
219
220 bool empty() const {
221     return size == 0;
222 }
223
224 size_t get_size() const {
225     return size;
226 }
227
228 template <typename Func>
229 void traverse_in_order(Func f) const {
230     traverse_in_order(root, f);
231 }
232
233 std::vector<T> items() const {
234     std::vector<T> result;
235     traverse_in_order([&result](const T& key) {
236         result.push_back(key);
237     });
238     return result;
239 }
240
241 /* Merge two sets in O(N+M) time.
242  * Some additional space may be costed.

```

```

243     * But it does not affect the result of the experiment.
244     */
245
246     void linearmerge(AVLSet&& other) {
247         if (other.empty()) return;
248
249         other.free_nodes.clear();
250
251         // Get elements from both trees
252         std::vector<T> q1 = items();
253         std::vector<T> q2 = other.items();
254         std::vector<T> all_elements;
255         all_elements.reserve(q1.size() + q2.size());
256
257         // Merge sorted vectors
258         std::merge(q1.begin(), q1.end(), q2.begin(), q2.end(),
259                 std::back_inserter(all_elements), comp);
260
261         // Recycle both trees
262         recycle_tree(root);
263         recycle_tree(other.root);
264         root = nullptr;
265         other.root = nullptr;
266         size = 0;
267         other.size = 0;
268
269         // Build new tree
270         root = build(all_elements.begin(), all_elements.end());
271         size = all_elements.size();
272     }
273
274     void simplemerge(AVLSet&& other) {
275         if (other.empty()) return;
276
277         if (size < other.size) {
278             // Swap to merge smaller into larger
279             std::swap(root, other.root);
280             std::swap(size, other.size);
281             std::swap(node_storage, other.node_storage);
282             std::swap(free_nodes, other.free_nodes);

```

```

283     }
284
285     // Insert all elements from the smaller tree (now 'other')
        into this
286     other.traverse_in_order([this](const T& key) {
287         this->insert(key);
288     });
289
290     other.clear();
291 }
292
293 private:
294
295 Node* insert(Node* node, const T& key) {
296     if (!node) {
297         size++;
298         return create_node(key);
299     }
300
301     if (comp(key, node->key)) {
302         node->left = insert(node->left, key);
303     } else if (comp(node->key, key)) {
304         node->right = insert(node->right, key);
305     } else {
306         return node;
307     }
308
309     return balance(node);
310 }
311
312 public:
313 void insert(const T& val){
314     root = insert(root, val);
315 }
316 void remove(const T& val){
317     root = remove(root, val);
318 }
319 template<typename RandAccIt>
320 void construct(RandAccIt bg, RandAccIt ed){
321     // Clear the current tree

```

```
322     recycle_tree(root);
323     root = nullptr;
324     // Build new tree
325     root = build(bg, ed);
326     size = static_cast<size_t>(ed - bg);
327 }
328 };
```

## References

Cormen, T. H., Leiserson, C. E., Rivest, R. L., & Stein, C. (2022). *Introduction to algorithms* (Fourth ed.). MIT Press.