## Prolog Programming for Artificial Intelligence

Three edition 2001

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Addision Wesley

## **GNU Prolog**

- This Prolog compiler complies with the ISO standard for Prolog (with useful extensions like global variables, ability to interface with the operating system, etc) and produces a native binary that can be run standalone. It is smart enough to avoid linking unused built-in predicates. It also has an interactive interpreter and a Prolog debugger as well as a low-level WAM debugger. You can interface with C code (both ways). Platforms supported include Linux (i86), SunOS (sparc) and Solaris (sparc).
  - <a href="http://www.thefreecountry.com/compilers/prolog.shtml">http://www.thefreecountry.com/compilers/prolog.shtml</a> (Free Prolog Compilers and Interpreters )
  - <a href="http://www.gprolog.org/">http://www.gprolog.org/</a> (The GNU Prolog web site)
  - <a href="http://www.thefreecountry.com/documentation/onlineprolog.sht">http://www.thefreecountry.com/documentation/onlineprolog.sht</a>
     <a href="mailto:ml">ml</a> (Online Prolog Tutorials)

## Some Examples

```
\circ L= [1,2,3,4], member(X, L).
```

- L= [a,b,c], length(L, X).
- $\circ$  X is 3 + 4.
- blue\_box.
   red\_box.
   green\_circle.
   blue\_circle.
   orange\_triangle.

#### **Tutorial**

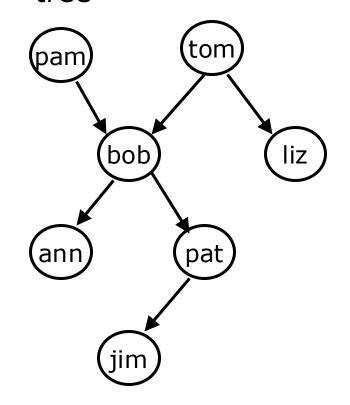
#### A Short Tutorial on Prolog

- As its name indicates, this is a short tutorial on Prolog programming.
- You can learn it by yourself.

#### Part 1 The Prolog Language

# Chapter 1 Introduction to Prolog

Given a whole family tree



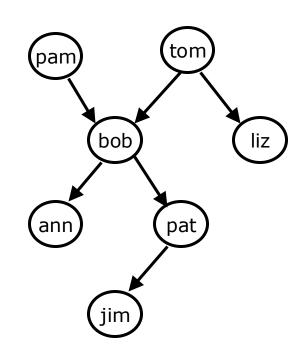
 The tree defined by the Prolog program:

```
parent( pam, bob).
  % Pam is a parent of
  Bob

parent( tom, bob).
parent( tom, liz).
parent( bob, ann).
parent( bob, pat).
parent( pat, jim).
```

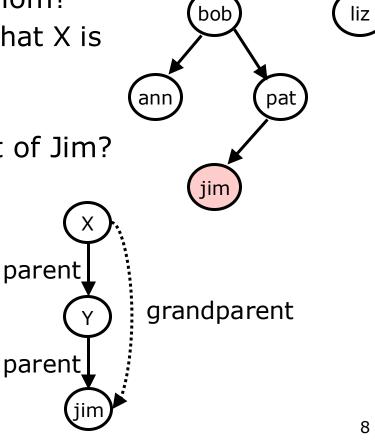
#### Ouestions:

- Is Bob a parent of Pat?
  - o ?- parent( bob, pat).
  - o ?- parent( liz, pat).
  - o ?- parent( tom, ben).
- Who is Liz's parent?
  - o ?- parent( X, liz).
- Who are Bob's children?
  - ?- parent( bob, X).



#### Ouestions:

- Who is a parent of whom?
  - Find X and Y such that X is a parent of Y.
  - ?- parent( X, Y).
- Who is a grandparent of Jim?
  - ?- parent( Y, jim), parent( X, Y).

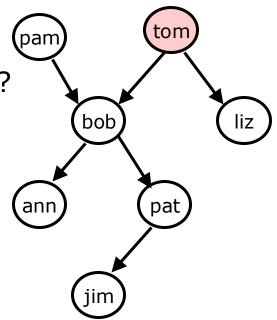


tom

#### Ouestions:

• Who are Tom's grandchildren?

- ?- parent( tom, X),parent( X, Y).
- Do Ann and Pat have a common parent?
  - ?- parent( X, ann),parent( X, pat).



- It is easy in Prolog to define a relation.
- The user can easily query the Prolog system about relations defined in the program.
- A Prolog program consists of clauses. Each clause terminates with a full stop.
- The arguments of relations can be
  - Atoms: concrete objects or constants
  - Variables: general objects such as X and Y
- Questions to the system consist of one or more goals.
- An answer to a question can be either positive (succeeded) or negative (failed).
- If several answers satisfy the question then Prolog will find as many of them as desired by the user.

#### o Facts:

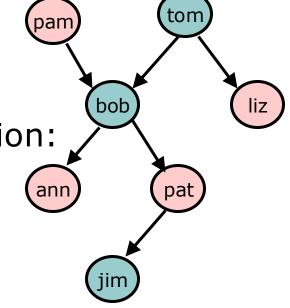
- female( pam).
- male( tom).
- male( bob).
- female( liz).
- female( ann).
- female( pat).
- male( jim).

Define the "offspring" relation:

- Fact: offspring(liz, tom).
- Rule: offspring(Y, X):parent(X, Y).
  - For all X and Y,Y is an offspring of X ifX is a parent of Y.

% Pam is female

% Tom is male



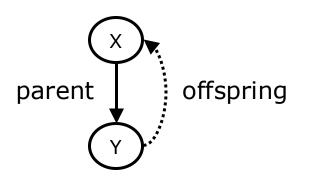
#### o Rules have:

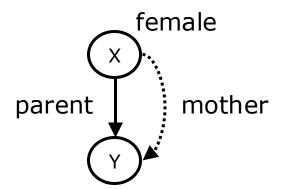
- A condition part (body)
  - o the right-hand side of the rule
- A conclusion part (head)
  - the left-hand side of the rule

#### • Example:

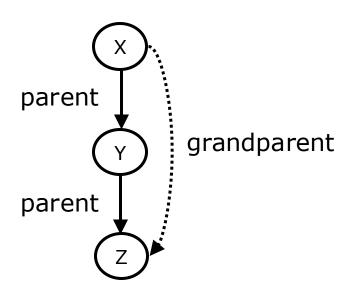
- offspring(Y, X):- parent(X, Y).
- The rule is general in the sense that it is applicable to any objects X and Y.
- A special case of the general rule:
  - offspring( liz, tom) :- parent( tom, liz).
- ?- offspring(liz, tom).
- ?- offspring( X, Y).

- Openion of the Define the Defi
  - mother( X, Y) :- parent( X, Y), female( X).
  - For all X and Y,
    - X is the mother of Y if
    - X is a parent of Y and
    - X is a female.

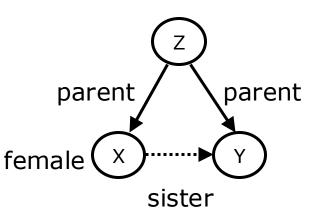




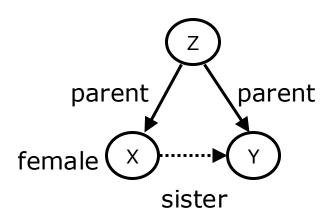
- Define the "grandparent" relation:
  - grandparent( X, Z) :parent( X, Y), parent( Y, Z).



- Openion of the Define the Sister relation:
  - sister( X, Y):parent( Z, X), parent( Z, Y), female(X).
  - For any X and Y,
     X is a sister of Y if
     (1) both X and Y have the same parent, and
     (2) X is female.
  - ?- sister( ann, pat).
  - ?- sister( X, pat).
  - ?- sister( pat, pat).
    - o Pat is a sister to herself?!



- To correct the "sister" relation:
  - sister( X, Y): parent( Z, X), parent( Z, Y), female(X),
     different( X, Y).
  - different (X, Y) is satisfied if and only if X and Y are not equal. (only assumption here)



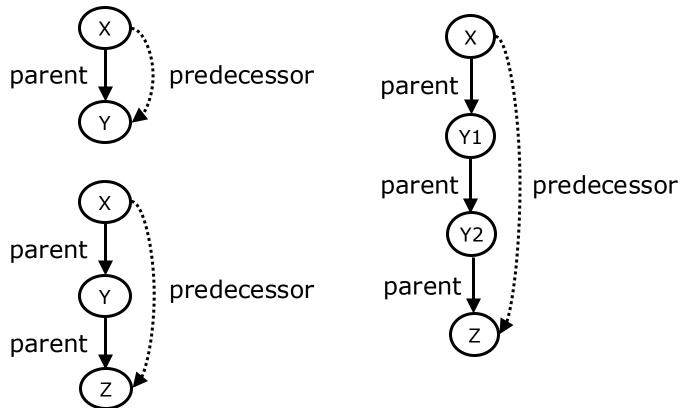
- Prolog clauses consist of
  - Head
  - Body: a list of goal separated by commas (,)
- Prolog clauses are of three types:
  - Facts:
    - declare things that are always true
    - facts are clauses that have a head and the empty body
  - Rules:
    - declare things that are true depending on a given condition
    - rules have the head and the (non-empty) body
  - Questions:
    - the user can ask the program what things are true
    - questions only have the body

- A variable can be substituted by another object.
- Variables are assumed to be universally quantified and are read as "for all".
  - For example:

```
hasachild( X) :- parent( X, Y). can be read in two way
```

- (a) For all X and Y,if X is a parent of Y then X has a child.
- (b) For all X,X has a child if there is some Y such that X is a parent of Y.

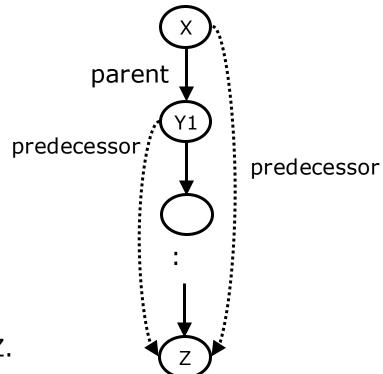
Define the "predecessor" relation



Define the "predecessor" relation

```
predecessor( X, Z):-
  parent( X, Z).
predecessor( X, Z):-
  parent( X, Y),
  predecessor( Y, Z).
```

- For all X and Z,
   X is a predecessor of Z if
   there is a Y such that
   (1) X is a parent of Y and
   (2) Y is a predecessor of Z.
- ?- predecessor( pam, X).



```
% Figure 1.8 The family program.
                                       mother(X,Y):-
                                          parent(X,Y),
parent( pam, bob).
                                         female(X).
parent( tom, bob).
parent( tom, liz).
                                       grandparent(X, Z) :-
parent(bob, ann).
                                          parent(X,Y),
parent(bob, pat).
                                          parent(Y, Z).
parent( pat, jim).
                                       sister(X,Y):-
female(pam).
                                          parent(Z, X),
female(liz).
                                         parent( Z, Y),
                                         female(X),
female(ann).
                                         different(X, Y).
female(pat).
male(tom).
male(bob).
                                       predecessor(X, Z) :- % Rule pr1
male(jim).
                                          parent(X, Z).
offspring(Y, X) :-
                                       predecessor( X, Z) :- % Rule pr2
  parent(X, Y).
                                          parent(X,Y),
                                          predecessor(Y, Z).
```

#### o Procedure:

 In figure 1.8, there are two "predecessor relation" clauses.

```
predecessor( X, Z) :- parent( X, Z).
predecessor( X, Z) :- parent( X, Y), predecessor( Y, Z).
```

Such a set of clauses is called a procedure.

#### o Comments:

```
/* This is a comment */
% This is also a comment
```

#### **Trace and Notrace**

```
! ?- trace.
The debugger will first creep -- showing everything
    (trace)
(15 ms) yes
                                                     X = bob
{trace}
                                                     Z = jim
                                                             1 Redo: predecessor(bob,jim)?
| ?- predecessor( X, Z).
                                                            2 Redo: predecessor(pat,jim) ?
    1 1 Call: predecessor(_16,_17)?
                                                            3 Call: parent(pat, 144)?
    2 2 Call: parent( 16, 17)?
                                                            3 Exit: parent(pat,jim)?
       2 Exit: parent(pam,bob) ?
       1 Exit: predecessor(pam,bob)?
                                                            3 Fail: parent(jim,_17) ?
                                                            3 Call: parent(jim, 144)?
X = pam
                                                            3 Fail: parent(jim,_132) ?
Z = bob ? ;
                                                            2 Fail: predecessor(jim, 17)?
       1 Redo: predecessor(pam,bob)?
                                                             1 Fail: predecessor(16, 17)?
      2 Redo: parent(pam,bob) ?
    2 2 Exit: parent(tom,bob)?
                                                     (266 ms) no
       1 Exit: predecessor(tom,bob)?
                                                     {trace}
X = tom
                                                     ! ?- notrace.
Z = bob ?;
                                                     The debugger is switched off
                                                     yes
```

- To answer a question, Prolog tries to satisfy all the goals.
- To satisfy a goal means to demonstrate that the goal is true, assuming that the relations in the program is true.
- Prolog accepts facts and rules as a set of axioms, and the user's question as a conjectured (推測的) theorem.
- Example:
  - Axioms: All men are fallible (會犯錯的).

    Socrates is a man.
  - Theorem: Socrates is fallible.
  - For all X, if X is a man then X is fallible.

```
fallible( X) :- man( X) man( socrates).
```

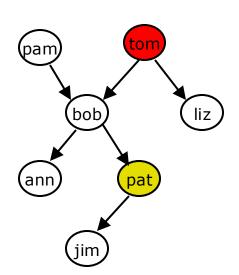
?- fallible( socrates).

?- predecessor( tom, pat).

- (1) parent( bob, pat) → predecessor( bob, pat)
  - Form parent( bob, pat) it follows that predecessor( bob, pat), by rule pr1.
- (2) parent(tom, bob) is fact.
- (3) parent(tom, bob) and parent(bob, pat) → predecessor(tom, pat).
  - Using the fact and the derived fact parent( bob, pat)
     we can conclude predecessor( tom, pat).

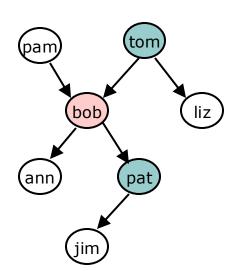
?- predecessor( tom, pat). predecessor( X, Z) :- parent( X, Z). % Rule pr1 predecessor( X, Z) :- parent( X, Y), % Rule pr2 predecessor( Y, Z).

- How dose the Prolog system actually find a proof sequence?
  - Prolog first tries that clause which appears first in the program. (rule pr1)
  - $\circ$  Now, X = tom, Z = pat.
  - The goal predecessor(tom, pat) is then replace by parent(tom, pat). (see Figure 1.9)
  - There is no clause in the program whose head matches the goal parent( tom, pat).
  - Prolog backtracks to the original goal in order to try an alternative way (rule pr2).



?- predecessor( tom, pat).

- Apply rule pr2, X = tom, Z = pat, but Y is not instantiated yet.
- The top goal predecessor( tom, pat) is replaces by two goals: (see Figure 1.10)
  - parent( tom, Y)
  - predecessor( Y, pat)
- The first goal matches one of the facts. (Y = bob)
- The remaining goal has become predecessor( bob, pat)
- Using rule pr1, this goal can be satisfied.
  - predecessor( bob, pat) :- parent( bob, pat)



predecessor(tom, pat) By rule pr1 By rule pr2 parent( tom, Y) parent(tom, pat) predecessor(Y, pat) By fact Y = bobno parent( tom, bob) predecessor( bob, pat) The top goal is satisfied when a path is found from the root node to a leaf node By rule pr1 labeled 'yes'. parent( bob, pat) The execution of Prolog is the searching for such path. ves

#### **Trace**

predecessor(X, Z) :- parent(X, Z). % Rule pr1
predecessor(X, Z) :- parent(X, Y), % Rule pr2
predecessor(Y, Z).

```
predecessor(tom, pat)
                            By rule pr1
                                                              By rule pr2
                                                       parent(tom, Y)
                            parent( tom, pat)
                                                     predecessor( Y, pat)
                                                                By fact
                                                   Y = bob
                                   no
                                                                 parent( tom, bob)
?- predecessor( tom, pat).
       1 Call: predecessor(tom,pat)?
                                                    predecessor(bob, pat)
       2 Call: parent(tom,pat) ?
       2 Fail: parent(tom,pat) ?
                                                                  By rule pr1
       2 Call: parent(tom,_79)?
       2 Exit: parent(tom,bob) ?
                                                      parent(bob, pat)
       2 Call: predecessor(bob,pat)?
       3 Call: parent(bob,pat)?
       3 Exit: parent(bob,pat) ?
                                                             yes
       2 Exit: predecessor(bob,pat)?
       1 Exit: predecessor(tom,pat)?
```

true?

# 1.5 Declarative and procedural meaning of programs

- Two levels of meaning of Prolog programs:
  - The declarative (宣告的) meaning
    - o concerned only with the relations defined by the program
    - determines what will be the output of the program
    - The programmer should concentrate mainly on the declarative meaning and avoid being distracted by the executional details.
  - The procedural (程序的) meaning
    - determines how this output is obtained
    - determines how the relations are actually evaluated by the Prolog system
    - The procedural aspects cannot be completely ignored by the programmer for practical reasons of executional efficiency.

#### Exercise

- Exercise 1.7
  - (a) ?- mother( pam, bob).
  - (b) ?- grandparent(bob, jim).
  - Try to understand how Prolog derives answers to the following questions, using the program of Figure 1.8. (trace)
  - Try to draw the corresponding derivation diagrams in the style of Figures 1.9 to 1.11.
  - Will any backtracking occur at particular questions?