Data Structures & Algorithms

Problem statement:

Minesweeper is a single-player puzzle video game. The objective of the game is to clear a rectangular board containing hidden "mines" without detonating any of them, with help from clues about the number of neighboring mines in each field.

The game will be played through a console-based interface. The user will choose a size N for the board (>=5). The board will be randomly generated, with the density of the elements roughly equal to 25%. The NxN board will be presented with the unexplored spots marked with '-' and a number of flags equal to the number of mines. The user can explore a spot or mark it with a flag. Whenever the user chooses to explore a spot he will either: lose the game if the spot is a mine or otherwise expand it such that he gains new information on the surrounding mines. The game ends when a mine is hit or when the there are no empty spots left on the board.

Justification:

The game is played on a rectangular board where the majority of the elements are 0 (the empty spots), thus the only relevant elements remain the spots occupied by a mine. This enables us to use a sparse matrix, since, by ignoring the empty spots, the data requires significantly less storage.

ADT Domain specification & representation:

SparseMatrix = $\{ m \mid m \text{ is a container that represents a two-dimensional array, where each element has a unique position determined by 2 indexes (line & column) and a non-zero value. The elements are stored under the form triplet = (line, column, value), with <math>\{ \text{ line } \in \text{Integer, column } \in \text{Integer, value } \}$.

The hashing will be done using the division method, while the collisions will be resolved by separate chaining.

For this particular problem perfect hashing could be achieved with a hash function that uses the line and column indexes; an example would be: lineIndex + columnIndex * noColumns.

TElem:

line: Integer

column: Integer

value: TValue

Node:

next: \tangle Node

element: TElem

```
SparseMatrix:
```

matrix: \tangle Node[]
noLines: Integer
noColumns: Integer
hash: TFunction

size: Integer

ADT Interface specification:

```
create(m, nol, noc)
  Pre: nol, noc \in Integer, nol > 0, noc > 0
  Post: m ∈ SparseMatrix with nol lines and noc columns
modify(m, l, c, v)
  Pre: m \in SparseMatrix; l, c \in Integer, v \in TValue
  Post: m1 \in SparseMatrix; m1 = m where the TValue with key k on (1, c) = v
  Throws: InvalidPositionException if l > m.noLines or c > m.noColumns
lineNr(m)
  Pre: m \in SparseMatrix
  Post: lineNr \in Integer, lineNr <- the number of lines in the matrix
columnNr(m)
  Pre: m \in SparseMatrix
  Post: columnNr \in Integer, columnNr < - the number of columns in the matrix
getElement(m, l, c)
  Pre: m \in SparseMatrix; l,c \in Integer
  Post: e \in TValue, e \le the value at (l, c) mapped with the key k
  Throws: InvalidPositionException if l > m.noLines or c > m.noColumns
```

ADT Implementation:

```
Subalgorithm create(m, nol, noc):

for i <- 0, i < size, 1 do

m.matrix[i] <- NULL
```

```
m.noLines <- nol
m.noColumns <- noc
```

End-Subalgorithm

Complexity:

BC =AC = WC = $\Theta(n)$, n = m.size - whenever a new matrix is created all the elements are always initialized with NULL

Subalgorithm modify(m, l, c, v)

```
if l > m.noLines or c > m.noColumns then
  @Throw InvalidPositionException
end-if
k < -1 + c * m.noColumns
position \leq- m.hash(k)
node <- m.matrix[position]</pre>
if node = NIL and v = NIL then
  return
else if node = NIL and v != NIL then
  @allocate(newNode)
  newNode.next <- NIL
  newNode.element.line <- 1
  newNode.element.column <- c
  newNode.element.value <- v
  node <- newNode
  return;
else if node != NIL and v = NIL then
  aux <- NULL
  if node.element.value = v then
    aux <- node.next
    @de-allocate node
    node <- aux
    m.matrix[hash] <- aux
```

```
return;
     end-if
     while node.next != NULL and node.next.element.value != value do
       node <- node.next
     end-while
    if node.next != NULL and node.next.element.value = value
       aux <- node.next.next</pre>
       @de-allocate node.next
       node.next <- aux
       return;
     end-if
  else
     while node != NULL and node.element.value != value
       node <- node.next
     end-while
    if node != NULL and node.element.value = value do
       node.element.value <- v
     end-if
  end-if
End-Subalgorithm
Complexity:
BC = O(1) – value associated to the key is null
AC = O(
WC = O(1 + \alpha) – the value associated to the key is a list of length \alpha
Function getElement(m, 1, c)
  if l > m.noLines or c > m.noColumns or l < 0 or c < 0 then
     @Throw InvalidPositionException
  end-if
  k < -1 + c * m.noColumns
```

```
position <- m.hash(k)
  node <- m.matrix[position]</pre>
  if node != NIL then
     while node != NIL and node.key != key and node.next != NIL do
       node = node.next
     end-while
     if node == NIL then
       telem.line <- -1
       telem.column <- -1
       telem.value <- 0
       return telem
     end-if
     return node.value
  else
     telem.line <- -1
     telem.column <- -1
     telem.value <- 0
     return telem
  end-if
End-Function
Complexity:
BC = O(1) – the specified element does not exist or it is the head of the list
AC = O(
WC = O(1 + \alpha) – the specified element exists on position \alpha in the list
Function lineNr(m)
  lineNr <- m.noLines
End-Function
Complexity:
BC =AC = WC = \Theta(1) – the number of lines directly accessed at m.noLines
```

Function columnNr(m)

columnNr <- m.noColumns

End-Function

Complexity:

BC = AC = WC = $\Theta(1)$ – the number of columns directly accessed at m.noColumns

ADT Tests:

```
void Tests::testMine()
       Mine mine{ 1, 2 };
       assert(mine.toString() == "Mine={x=1, y=2, marked=0}");
       assert(mine.getX() == 1);
       assert(mine.getY() == 2);
       assert(mine.isMarked() == false);
       mine.setMarked(true);
       mine.setX(10);
       mine.setY(20);
       assert(mine.getX() == 10);
       assert(mine.getY() == 20);
       assert(mine.isMarked() == true);
}
void Tests::testHashTable()
{
       HashTable<int, Mine> table = HashTable<int, Mine>();
       Mine mine1{ 1, 2 };
       Mine mine2{ 3, 4 };
       Mine mine3{ 5, 6 };
       Mine mine4{ 9, 10 };
       table.insert(1, mine1);
       table.insert(24, mine2);
       table.insert(2, mine3);
       assert(table.getSize() == 3);
       table.insert(47, mine4);
       assert(table.getSize() == 4);
       assert(table.getValue(47) == mine4);
       assert(table.getValue(1) == mine1);
       assert(table.removeKeyValue(24, mine2) == true);
       assert(table.removeKeyValue(24, mine2) == false);
```

```
assert(table.removeKeyValue(7, mine2) == false);
       assert(table.removeKeyValue(2, mine3) == true);
       assert(table.getSize() == 2);
       assert(table.update(1, mine1, mine2) == true);
       assert(table.update(5, mine1, mine2) == false);
       assert(table.update(1, mine4, mine2) == true);
}
void Tests::testSparseMatrix()
       SparseMatrix matrix = SparseMatrix(5, 5);
       assert(matrix.columnNr() == 5);
       assert(matrix.lineNr() == 5);
       Mine mine;
       bool getElemTest = false;
       try {
               mine = matrix.getElement(6, 6);
       } catch (InvalidPositionException ex)
              getElemTest = true;
       assert(getElemTest);
       mine = matrix.getElement(3, 2):
       assert(mine.getX() == -1 \&\& mine.getY() == -1 \&\& mine.isMarked() == false);
       mine.setMarked(true);
       mine.setX(1);
       mine.setY(1):
       matrix.modify(1, 1, mine);
       assert(matrix.getElement(1, 1) == mine);
       Mine newMine{ 2, 4 };
       matrix.modify(1, 1, newMine);
       assert(matrix.getElement(1, 1) == newMine);
}
```

Pseudocode Solution:

Complexity:

```
Subalgorithm create(con, size):

con.board <- SparseMatrix(size, size)

con.view <- SparseMatrix(size, size)

con.numberOfMines <- NULL

initMines()

End-Subalgorithm
```

 $BC = AC = WC = \Theta$ (n+m), n = size, m = numberOfMines – time taken to initialize the matrices and the mines

```
<u>Subalgorithm initMines(con):</u>
  con.numberOfMines <- (con.board.columnNr() * con.board.lineNr()) / 5
  emptyElem.x < --1
  emptyElem.y <- -1
  emptyElem.value <- 0
  for i \le 0, i \le con.numberOfMines, 1 do
    do
       x <- @random % con.board.lineNr()
       y <- @random % con.board.columnNo()
    while con.board.getElement(x, y) = emptyElem
    con.board.modify(x, y, Mine(x, y))
  end-for
End-Subalgorithm
Complexity:
BC = AC = WC = \Theta (n), n – numberOfMines – time taken to initialize n mines
Subalgorithm flag(con, x, y):
  mine <- con.view.getElement(x, y)
  if mine.isMarked() = true then
    mine.setMarked(false)
    con.view.modify(x, y, mine)
    con.board.modify(x, y, mine)
  else
    mine.setMarked(true)
    con.view.modify(x, y, mine)
    con.board.modify(x, y, mine)
```

End-Subalgorithm

end-if

```
Complexity:
BC =
WC =
AC =
Function expand(con, x, y):
  mine.x < --1
  mine.y <- -1
  mine.value = 0
  if con.board.getElement(x, y).isMarked() then
    @throw MarkedSpotException
  end-if
  if con.board.getElement(x, y) != mine then
    expand <- false
  else
    con.fillBoard(x, y)
    expand <- true
End-Function
Complexity:
BC =
AC =
WC =
Subalgorithm fillBoard(con, x, y):
  counter <- 0
  if y + 1 < con.board.columnNr() and con.board.getElement(x, y+1).getX != -1 then
    counter <- counter + 1
  if x + 1 < \text{con.board.lineNr}() and con.board.getElement(x+1, y).getX != -1 then
    counter <- counter + 1
  if y + 1 < con.board.columnNr() and x + 1 < con.board.lineNr() and
```

```
con.board.getElement(x+1, y+1).getX != -1 then
    counter <- counter + 1
  if y - 1 > \text{con.board.columnNr}() and con.board.getElement(x, y-1).getX != -1 then
    counter <- counter + 1
  if x - 1 > \text{con.board.lineNr}() and con.board.getElement(x - 1, y).getX != -1 then
    counter <- counter + 1
  if y - 1 > con.board.columnNr() and x - 1 > con.board.lineNr() and
    con.board.getElement(x-1, y-1).getX != -1 then
    counter <- counter + 1
  if y + 1 < con.board.columnNr() and x - 1 > con.board.lineNr() and
    con.board.getElement(x-1, y+1).getX != -1 then
    counter <- counter + 1
  if y - 1 > con.board.columnNr() and x + 1 > con.board.lineNr() and
    con.board.getElement(x+1, y+1).getX != -1 then
    counter <- counter + 1
  mine \leq- Mine(x, y)
  mine.setProxy(counter)
  con.view.modify(x, y, mine)
End-Subalgorithm
Complexity:
BC =
AC =
WC =
Function getBoardHeight(con):
  getBoardHeight <- con.board.lineNr()</pre>
End-Function
Complexity:
BC = AC = WC = \Theta(1)
```

Function getBoardWidth(con):

getBoardWidth <- con.board.columnNr()</pre>

End-Function

Complexity:

$$BC = AC = WC = \Theta(1)$$