

(1)

Ans. to Q No-1

- 1) The equation ~~will~~ for finding remainder will be:

$$A - \left(B * \underset{\text{floor}}{\cancel{ceil}}(A/B) \right)$$

$$[\because \text{Remainder} = \text{Dividend} - (\text{Divisor} \times \text{quotient})]$$

2)

Operator	Count
-	1
*	1
/	1

AmtoQNo - 2

⇒ First, we find the lower bound of the given integer in the array and then find the upper bound. After that, (index of upper bound - index of lower bound) will give us the results.

⇒ We can also do it using,
Binary search first:

```
int idn = binary-search(a, a+n, p) - a;
```

```
int count = 1;
```

Code for counting elements on left side:

```
int l = idn - 1;
```

```
while (l >= 0 && a[l] == p) {
```

```
    ++count;
```

```
    --l;
```

```
}
```

For counting right side:

```
int r = idn + 1;
```

```
while (r < n && a[r] == p) {
```

```
    ++count; ++r;
```

count
is the
answer /
no of
occurrences

③

Ans. to Q No-3

1. Array
2. Linked list
3. Double ended Queue
4. Binary Search Tree

Ans. to Q No-4

We will take 2 balls each time and weigh them. When we will find an imbalance we will take a ball from that imbalance pair and weigh it against another ball.

This will generate the solution. In

this approach, we will need to use the

scale maximum 5 times.

Ans to Q No - 5

We may use a Stack to solve in $O(N)$.

We iterate through the array and ~~if~~

insert the current number in the stack

only if the stack is empty or ~~s.top()~~

is a number $>$ current number.

But when ~~if~~ current number is greater

than $s.top()$, we pop all the elements

currently present in the stack and print

them because the current number is

greater than all those numbers. Only after

that, we put the current number in the

stack and repeat the process. When iteration

is done, we print the numbers left in

the stack as 'number \rightarrow none'.