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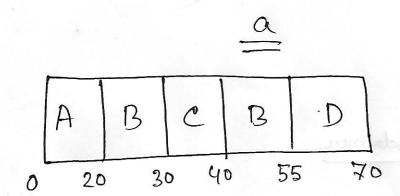
Section: 07

#### Disclaimer

I have not shared any answers with anyone, did not help anyone and did not take help from anyone for this exam. The answers I have given here are done by myself only.

Signature

## Amito & No-4



$$A \rightarrow P_1$$
 $B \rightarrow P_2$ 

$$C \rightarrow P_3$$

Waiting time 
$$\frac{b}{15}$$
 Twinaround time  $P_2 \rightarrow (20-20) + (40-30)$   $P_2 \rightarrow End-Apprival$   $\Rightarrow 55-15$   $\Rightarrow 40$ 

Am. 15

Am. 40

4 content switches are needed here.

#### Ans. to Q NO - 5

$$P_2 \rightarrow 63 + 80 + 76 = 219$$

.. Average waiting time = 
$$\frac{981}{5} = 196.2$$
 (Am.)

## Response time:

$$P_i \rightarrow 0-0=0$$

$$P_2 \rightarrow 313 - 250 = 63$$

Average response time = 
$$\frac{231}{5}$$
 = 46.2 (Ans)

## Am to BNO-3

Criven, cone=4, s= 40%=0.4

Speedup 
$$\leq \frac{4}{0.4 + \frac{1-0.4}{4}}$$

< 1.82 times

Again, if cone = 8 and s=0.4

 $\leq$  2.11 times

: Performance increax =  $\frac{0.29}{2.11}$  × 100%

= 13,74%

: 13.74% < 30°6, & will decide NOT

to buy 8 come processon.

# Am. to QNo-1

The operating system is loaded through the Bootstrap program. The computer abready knows where the Bootstrap program is and that Bootstrap program loads the OS into the main memory for startup.

done ad olide process and

We need 0s for communicating and using the resources from our hardware. We need to install it. Only copying the file is not enough. Because then our computer will not be able to recognize the OS and load it correctly during startup.

1 - 01/ 9 ot met

Enample-1: When a process will need I/o, it will go from ready to waiting state and when I/o request is complete, it will go back to the ready state.

### Enample - 2:

A process can be interrupted to

A process can create a child process and
wait for it's completion. In this state,
the parent process will go into the waiting
state and I when the child process tinishes
enecution, the parent process will go back
to the ready state.

Data parallelism divides the data among different threads and each of those threads perform the same operation on the data.

Task parallelism divides the task itself among Several threads and each thread, perform different operation on the data.

Data parallelism enample: 97 we want to do sum from 1-1000, we can divide this data among 2 different threads. One thread will sum from 1-500 and another thread will sum from 501-1000.

## tank parallelisme enample:

It we want to do, addition and subtraction

on the same number, we can do it using 2 different thread. One will return sum and another will return difference.

Ver, it's possible to do both at same time.

Several Hurrando and each Thread, perform different

Yes, it can prevent starvation. Because in multilevel feedback queue, the processes can switch between different queues. So, if a process is waiting for a long time in a queue, it will simply move to another queue, where it will have higher priority for execution.