

Time Complexity:

$$n = r - p$$

Line	Instruction	# of execution times
	Partition(A,p,r)	
1	x = A[r]	1
2	i = p - 1	1
3	for j = p to r - 1	n+1
4	if A[j] ≤ x	n
5	i = i + 1	n
6	A[i] ↔ A[j]	n
7	A[i + 1] ↔ A[r]	1
8	return i + 1	1

Worst, Average & Best:

$$T(n) = 1 + 1 + (n + 1) + n + n + n + 1 + 1 = 5 + 4n = O(n)$$

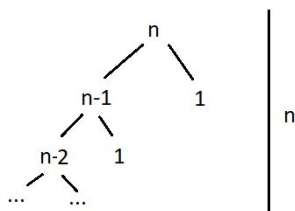
$O(n)$, $\Omega(1)$, $\Theta(n)$

Line	Instruction	# of execution times
	QuickSort(A,p,r)	
1	if p < r	1
2	q = Partition(A,p,r)	O(n)
3	QuickSort(A,p,q-1)	T(n/2)
4	QuickSort(A,q+1,r)	T(n/2)

Worst:

$$T(n) = 1 + O(n) + T(n-1) + T(1) \approx O(n) + T(n-1) + T(1)$$

$$n - i2 = 1 \rightarrow i = 2(n-1) = 2n - 2 \approx n$$



$$\sum_{i=0}^n n = O(n^2)$$

$$O(n^2)$$

Average & Best:

$$T(n) = 1 + O(n) + 2T(n/2) \approx O(n) + 2T(n/2)$$

by the Master method:

$$T(1) \rightarrow \Theta(1), a = 2, b = 2, a = b^1 \text{ is true}$$

$$T(n) = \Theta(n^1 \log n) = \Theta(n \log n)$$

$$\Theta(n \log n), \Omega(n \log n)$$

Line	Instruction	# of execution times
	Rand-Parti (A,p,r)	
1	i = Random(p,r)	1
2	A[r] \leftrightarrow A[i]	1
3	return Partition(A,p,r)	O(n)

Worst, Average & Best:

$$T(n) = 1 + 1 + O(n) = O(n)$$

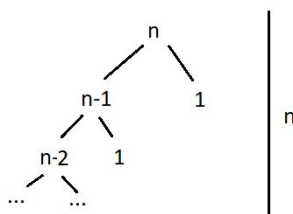
$$O(n), \Omega(n), \Theta(n)$$

Line	Instruction	# of execution times
	Randomized-QS (A,p,r)	
1	if p < r	1
2	q = Rand-Parti(A,p,r)	O(n)
3	Randomized-QS(A,p,q-1)	T(n/2)
4	Randomized-QS(A,q+1,r)	T(n/2)

Worst:

$$T(n) = 1 + O(n) + T(n-1) + T(1) \approx O(n) + T(n-1) + T(1)$$

$$n - i2 = 1 \rightarrow i = 2(n-1) = 2n-2 \approx n$$



$$\sum_{i=0}^n n = O(n^2)$$

$$O(n^2)$$

Average & Best:

$$T(n) = 1 + O(n) + 2T(n/2) \approx O(n) + 2T(n/2)$$

by the Master method:

$$T(1) \rightarrow \Theta(1), a = 2, b = 2, a = b^1 \text{ is true}$$

$$T(n) = \Theta(n^1 \log n) = \Theta(n \log n)$$

$$\Theta(n \log n), \Omega(n \log n)$$

Theoretical treatments:

Prueba	Variante	Estado	Tamaño (n)	Tiempo (ms)
1	Normal	Ascending	10	$O(n \log n)$
2			100	$O(n \log n)$
3			1000	$O(n \log n)$
4			10000	$O(n \log n)$
6		Descending	10	$O(n^2)$
7			100	$O(n^2)$
8			1000	$O(n^2)$
9			10000	$O(n^2)$
11		Random	10	$\Theta(n \log n)$
12			100	$\Theta(n \log n)$
13			1000	$\Theta(n \log n)$
14			10000	$\Theta(n \log n)$
16	Random	Ascending	10	$\Theta(n \log n)$
17			100	$\Theta(n \log n)$
18			1000	$\Theta(n \log n)$
19			10000	$\Theta(n \log n)$
21		Descending	10	$\Theta(n \log n)$
22			100	$\Theta(n \log n)$
23			1000	$\Theta(n \log n)$
24			10000	$\Theta(n \log n)$
26		Random	10	$\Theta(n \log n)$

27			100	$\Theta(n \log n)$
28			1000	$\Theta(n \log n)$
29			10000	$\Theta(n \log n)$

Experimental Unit:

-Quicksort Algorithm

Response Values:

-Execution time of the QuickSort method

Experimental Factors:

- **Studied:**

- Array status.
- Array size.
- Algorithm variant.

- **Not studied:**

- Number of programs being executed
- RAM capacity

Observational Factors:

-Program execution in the computer

Factor Levels:

- **Variant:** Normal, Random
- **Status:** Ascending, Descending, Random.
- **Size:** 10^1 , 10^2 , 10^3 , 10^4

Treatment:

Treatment	Variant	Status	Size(n)	Time (ms)
1	Normal	Ascending	10	
2			100	
3			1000	
4			10000	
6		Descending	10	
7			100	
8			1000	
9			10000	
11		Random	10	

12			100	
13			1000	
14			10000	
16		Ascending	10	
17			100	
18			1000	
19			10000	
21		Descending	10	
22			100	
23			1000	
24			10000	
26			10	
27			100	
28			1000	
29			10000	
	Random	Random		

1000 repetitions per treatment.

1.b. Hitherto the stages of study and experiment design we have completed are planning and realization. The stages that we are missing are analysis, interpretation, and control & final conclusions.

1.c. The program objective is to compare two or more treatments, since in this experiment we created many different treatments to see how the response variables change and to analyze the results and the variance using ANOVA. All of this because we want to understand the behavior of the QuickSort algorithm and we want to draw a valid conclusion.

1.d. Analysis:

Ascending arrays:

In the case of ascending arrays we have the following null hypothesis, alternate hypothesis and alpha value. (Group 1 is the QuickSort algorithm and Group 2 is the Randomized QuickSort algorithm)

$$H_0 : \mu_1 = \mu_2$$

$$H_a : \mu_1 \neq \mu_2$$

$$\alpha = 0.05$$

This is the resulting ANOVA table:

Variation Source	Sum of Squares	Degrees of Freedom	S.S. Average	F	P-Value	F Crit
Between Groups	567842,0792	1	567842,0792	6,07694	0,014538618	3,888374717
Within Groups	18688422,26	200	93442,11129			
Total	19256264,34	201				

Since the P-Value is 0.015, it is lower than alpha. Therefore we must reject the null hypothesis.

Descending arrays:

In the case of descending arrays we have the following null hypothesis, alternate hypothesis and alpha value. (Group 1 is the QuickSort algorithm and Group 2 is the Randomized QuickSort algorithm)

$$H_0 : \mu_1 = \mu_2$$

$$H_a : \mu_1 \neq \mu_2$$

$$\alpha = 0.05$$

This is the resulting ANOVA table:

Variation Source	Sum of Squares	Degrees of Freedom	S.S. Average	F	P-Value	F Crit
Between Groups	113395,0297	1	113395,0297	76,21889	1,00616E-15	3,888374717
Within Groups	297550,9901	200	1487,75495			
Total	410946,0198	201				

Since the P-Value is $1,006 \times 10^{-15}$, it is lower than alpha. Therefore we must reject the null hypothesis.

Random arrays:

In the case of random arrays we have the following null hypothesis, alternate hypothesis and alpha value. (Group 1 is the QuickSort algorithm and Group 2 is the Randomized QuickSort algorithm)

$$H_0 : \mu_1 = \mu_2$$

$$H_a : \mu_1 \neq \mu_2$$

$$\alpha = 0.05$$

This is the resulting ANOVA table:

Variation Source	Sum of Squares	Degrees of Freedom	S.S. Average	F	P-Value	F Crit
Between Groups	227834,9307	1	227834,9307	23,82808	2,14957E-06	3,888374717
Within Groups	1912323,168	200	9561,615842			
Total	2140158,099	201				

Since the P-Value is $2,150 \times 10^{-6}$, it is lower than alpha. Therefore we must reject the null hypothesis.