clickhouse执行计划创析

前置条件

版本 21.3 本文主要探究命中索引

创建表

插入数据

```
insert into default.test(a,b,c) values(1,1,1);
insert into default.test(a,b,c) values(5,2,2),(5,3,3);
insert into default.test(a,b,c) values(3,10,4),(3,9,5),(3,8,6),(3,7,7),(3,6,8),
(3,5,9),(3,4,10);
```

查询语句与生成的执行计划

```
select * from test WHERE b = 3 limit 2
```

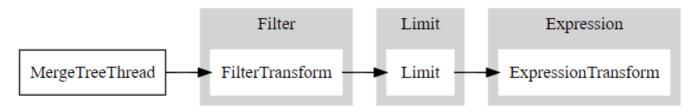
注意:此时需要将optimize_move_to_prewhere关闭,否则filterstep会被优化在prewhere逻辑,表现出来的pipeline不一样没有下面的直观。

pipeline

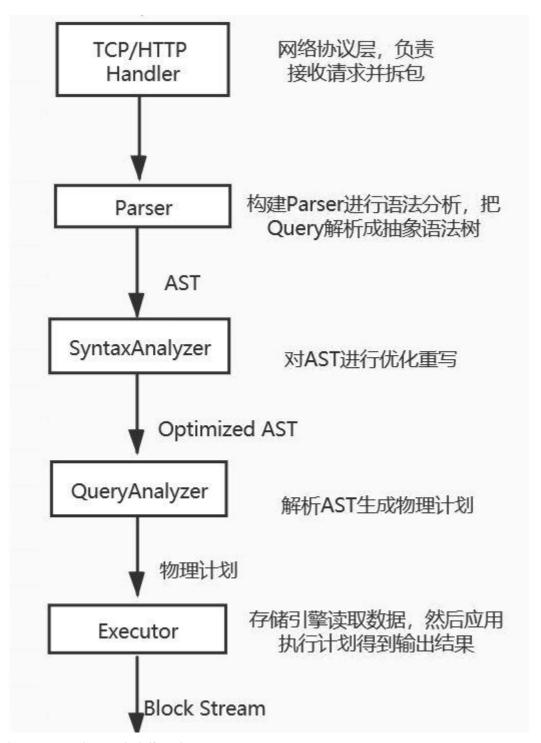
```
cxplain
(Expression)
ExpressionTransform
(Limit)
Limit
(Filter)
FilterTransform
(SettingQuotaAndLimits)
(ReadFromStorage)
```

MergeTreeThread 0 → 1

(ReadFromMergeTree)就是 QueryPlan 中的一个算子·MergeTreeThread则代表实现这个算子向 Pipeline 中添加的 Processor, Processor 后面的数字0 -> 1表示InputPort数量为 0,OutputPort数量为 1.如果 Processor 后面没有数字则默认InputPort和OutputPort数量都是 1。下面是由上面查询计划画出的dag图



怎么生成的执行计划呢?



接下来跟着代码一步步往下走

语法解析

```
bool ParserQueryWithOutput::parseImpl(Pos & pos, ASTPtr & node, Expected & expected)
{
    ParserShowTablesQuery show_tables_p;
    ParserSelectWithUnionQuery select_p;
    ...
    //逐个去解析·有一个解析成功就返回
    bool parsed =
        explain_p.parse(pos, query, expected)
    || select_p.parse(pos, query, expected)
```

AST重写优化

```
InterpreterSelectQuery::InterpreterSelectQuery(
)
{
    ...
    //优化和重写在这个表达式中
    auto analyze = [&] (bool try_move_to_prewhere)
    {
        ...
        //主要逻辑在TreeRewriter中
        syntax_analyzer_result = TreeRewriter(*context).analyzeSelect(
    }
}
```

解释器构造query plan

创建pipeline

```
BlockIO InterpreterSelectQuery::execute()
    BlockIO res;
    QueryPlan query_plan;
    buildQueryPlan(query_plan);
    res.pipeline =
std::move(*query_plan.buildQueryPipeline(QueryPlanOptimizationSettings(context-
>getSettingsRef())));
    return res;
//下面是创建pipeline的
QueryPipelinePtr QueryPlan::buildQueryPipeline(const QueryPlanOptimizationSettings
& optimization_settings)
{
    checkInitialized();
    optimize(optimization_settings);
    struct Frame
        Node * node;
        QueryPipelines pipelines = {};
    };
    QueryPipelinePtr last_pipeline;
    std::stack<Frame> stack;
    stack.push(Frame{.node = root});
    while (!stack.empty())
        auto & frame = stack.top();
        if (last_pipeline)
            frame.pipelines.emplace_back(std::move(last_pipeline));
            last_pipeline = nullptr;
```

```
size_t next_child = frame.pipelines.size();
       if (next_child == frame.node->children.size())
           bool limit max threads = frame.pipelines.empty();
           // 将当前算子对应的 Transformer 添加到 Pipeline.
            last_pipeline = frame.node->step-
>updatePipeline(std::move(frame.pipelines));
           if (limit_max_threads && max_threads)
                last_pipeline->limitMaxThreads(max_threads);
            stack.pop();
       }
       else
            stack.push(Frame{.node = frame.node->children[next_child]});
   }
   for (auto & context : interpreter_context)
        last_pipeline->addInterpreterContext(std::move(context));
   return last_pipeline;
}
```

buildQueryPipeline通过栈实现了 DFS 算法,算子不断入栈,终止条件是:当某个算子的所有 children 算子都完成 updatePipeline。算子的 updatePipeline 逻辑在IQueryPlanStep::updatePipeline()方法,这是一个虚函数。

对于举例的 SQL 来说·updatePipeline的顺序是:ReadFromMergeTree(ISourceStep) -> Filter(FilterStep) -> Limit(LimitStep) -> Expression(ITransformingStep) 。第一个算子ReadFromMergeTree会进入 ISourceStep::updatePipeline()方法,等等。

queryPlan和queryPipeline的关系

