IUTORIAL SHEET - 6

aue the applications of MST?

Minimum Spanning Tues on MST de a minimum weight spanning tree which is a subset of the edges of a connected, edge-weighted undirected graph that E with the minimum possible total edge weight.

- Applications

 i) Consider n stations are to be linked using a communication network and trying of communication link between any two stations involves a cost. The ideal solution would be to extend a subgraph truined as minimum cost spanning
- Spanning several cities then we can use the concept of minimum spanning trees.
- iii) Designing LAN
 - and consumer markets.
- v) Suppose you want to apply a set of houses with electric power water Telephone Unes Sewage line.
- Quez) Analyse the time and space complexity of lawn, - Keuiskal, Dijkstuu and Bellmon ford algorithm.

Prime Complexity Space Complexity

Prime's Algorithm O(IEI log IVI)

Kerne Rai's Algorithm O(IEI rog IEI)

Dij'kstra's Algorithm O(V2)

Bellman ford's Algorithm O(VE)

O(E)

Que3) Apply Kauskal & Paim's Algorithm on graph given on sight

Side to compute MST and its weight.

Keuskal's Algorithm

Killiskai S Ingolitini							
Source(v)	Destination(v)	Weight (w)					
6	7	10	and the contraction of				
520	8	2 × * 4 × 4 ×	1 12 3				
2	5	4~	8 \4 \q				
6	8	6 *					
6 2	3	オレ \	8 9 9 9				
7	8 4 (1)(1)	7 x	(7-6) (5)				
0	7	8~					
J	2	8 4	Weight = 1+2+2+4+4+7+8+9 = 37				
3000	4 4 10000	9400	\mathcal{D}				
40000	5 1	10(8))	The Kunnala of Dypolar and				
l l	7	11 *					
3	5	14 x					

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Que 4) Given a directed weighted graph. You are also given the Shoutest path from a source vertex '6' to a destination vertex 't'. Does the shoutest path dermain same in the modified graph in following cases?

The Shoutest path may change. The season is there may be different number of edges in different paths of from 's' to 't'. Example - Let Shoutest path be of weight 15 and has edge 5. Let there be another path with 2 edges and total weight 25. The

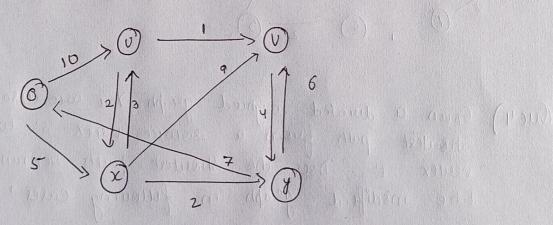
weight of the showlest path is incheased by 5. 10 8 becomes 15+50. Weight of the other pair is incheased by 2+10 & becomes 25+20. So, the showlest path changes to the other path with weight as 45.

(i) If weight of every edge is multiplied by 10 units

By we multiply all edges weight by 10, the shootest path doesn't change.

The reason is simple, usigns of all path from 's' to 2' get multiplied by some amount. The number of edges on a path doesn't matter. It is like changing units of weights.

Que 5) Apply Dijkstra and Bellman Algorithm on graph given on suight side to compute shortest path to all nodes from node s.



Dijketera's Algouithm

Node	Shoutest	distance	ferom	Source	Node
U	8		U		
20	5				
V	9				
y	7				

B ellmon	found a	gouithm.			
15←	°	20	8	∞ ≈ ₂	~ F
2rd		10	Ø 11	(A)	
3 and	° S	8)*/ q	<u>S</u>	\$ 7
Y ⁺ h		8	9	2	7
graph	doesn't	hove	-ue	cycle	