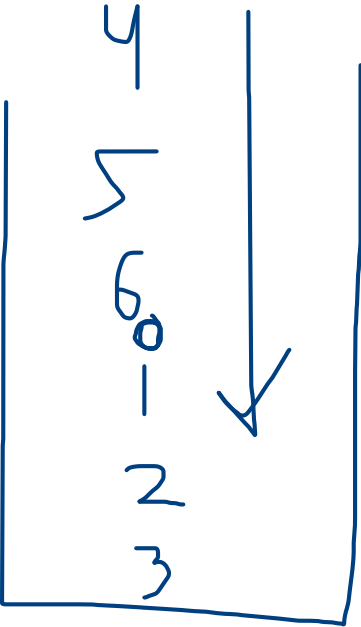
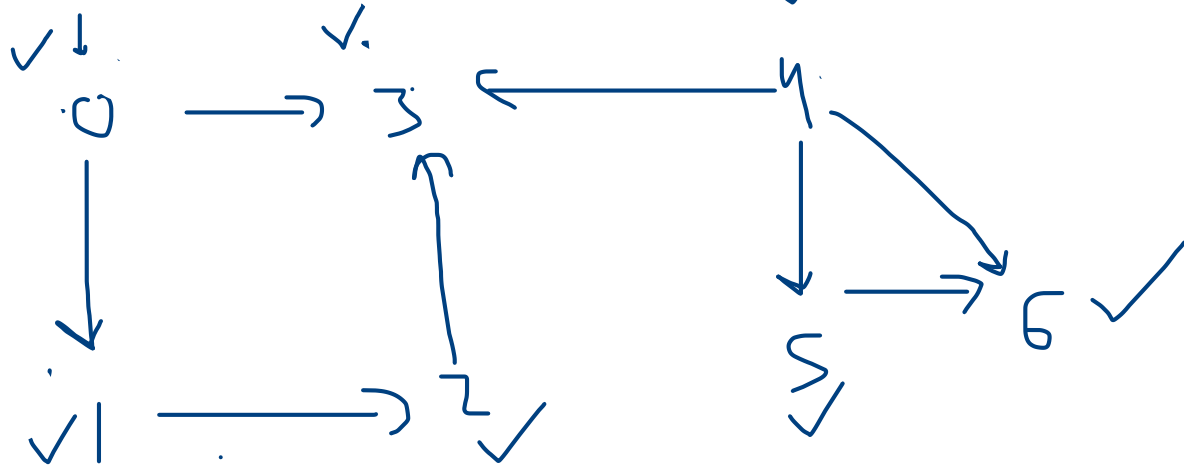


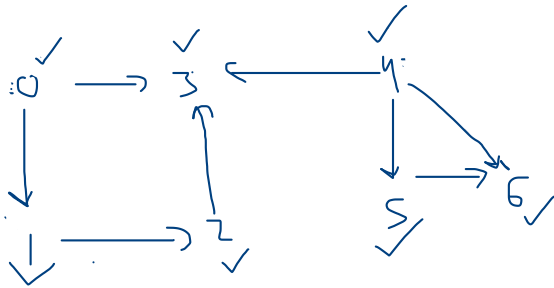
→ 0 1 2 4 5 6 3 ✓

Topological Sort (DFS)

PAGE



4 5 6 0 1 2 } ✓



5
4
3
2
1
0

```

public static void dfs(ArrayList<Edge>[] graph, int v) {
    boolean[] visited = new boolean[v];
    Stack<Integer> topologicalOrder = new Stack<>();
    for (int i = 0; i < v; i++) {
        if (!visited[i]) {
            dfsUtil(i, graph, visited, topologicalOrder);
        }
    }
}

```

```

while(! topologicalOrder.isEmpty()) {
    System.out.print(topologicalOrder.pop() + " ");
}
}

```

```

public static void dfsUtil(int src, ArrayList<Edge>[] graph, boolean[] visited, Stack<Integer> topologicalOrder) {

    visited[src] = true;

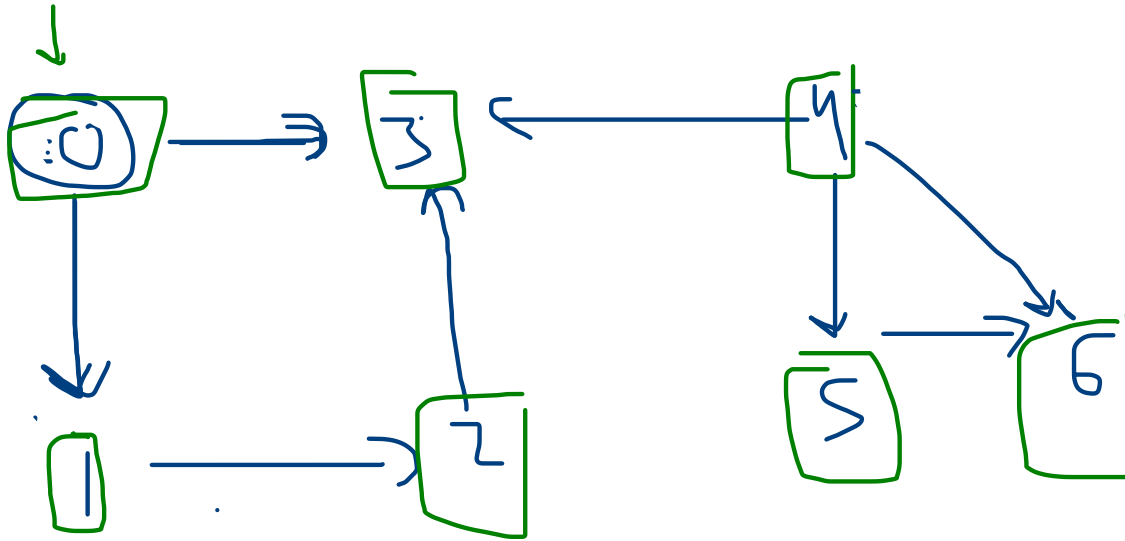
    for (Edge e: graph[src]) {
        if (!visited[e.nbr]) {
            dfsUtil(e.nbr, graph, visited, topologicalOrder);
        }
    }

    topologicalOrder.add(src);
}

```

DAG

Kahn's Algorithm (Topological Sort using BFS)

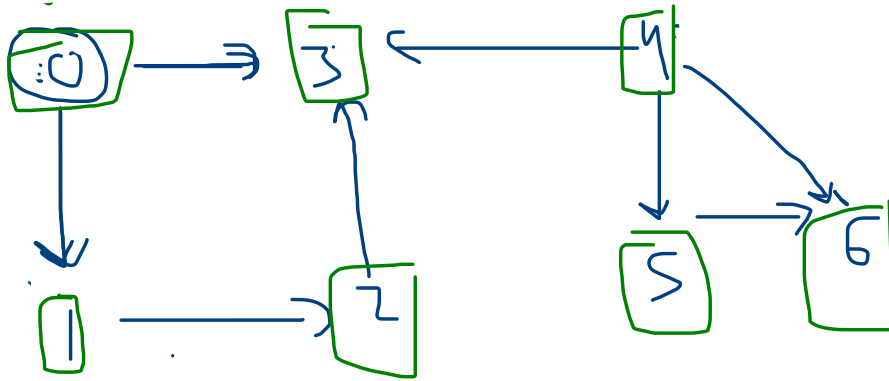


in degree =

0	1	2	3	4	5	6
0	1	1	3	0	1	2

out degree =

0	1	2	3	4	5	6
2	1	1	0	3	1	0



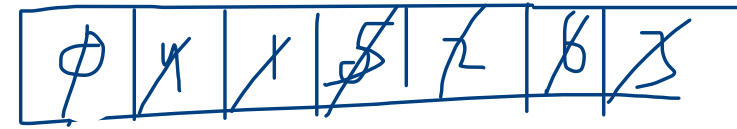
remove que
print
loop in nbr and decrement indegree
add vertex with indeg = 0 in que

✓ indegree =

0	1	2	3	4	5	6
0	1	1	3	0	1	2
	0	0	1 + 0		0	+ 0

ind = 0

que

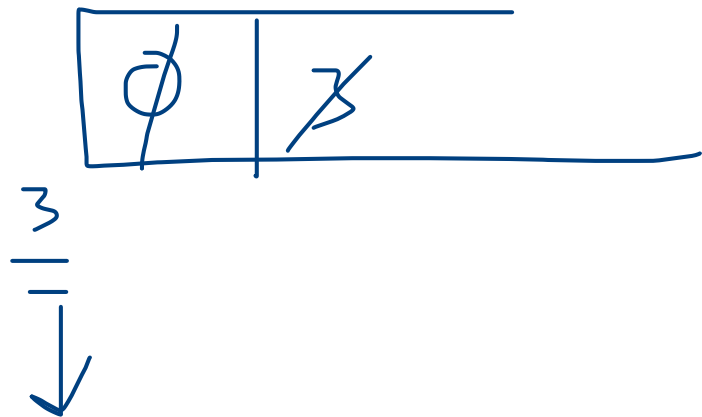
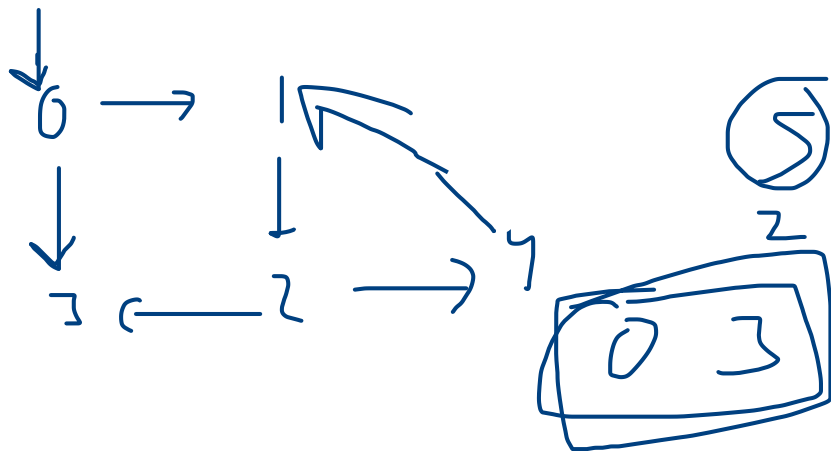


0 4 1 5 2

verify

ind = 0

$D \subset G$
 \equiv



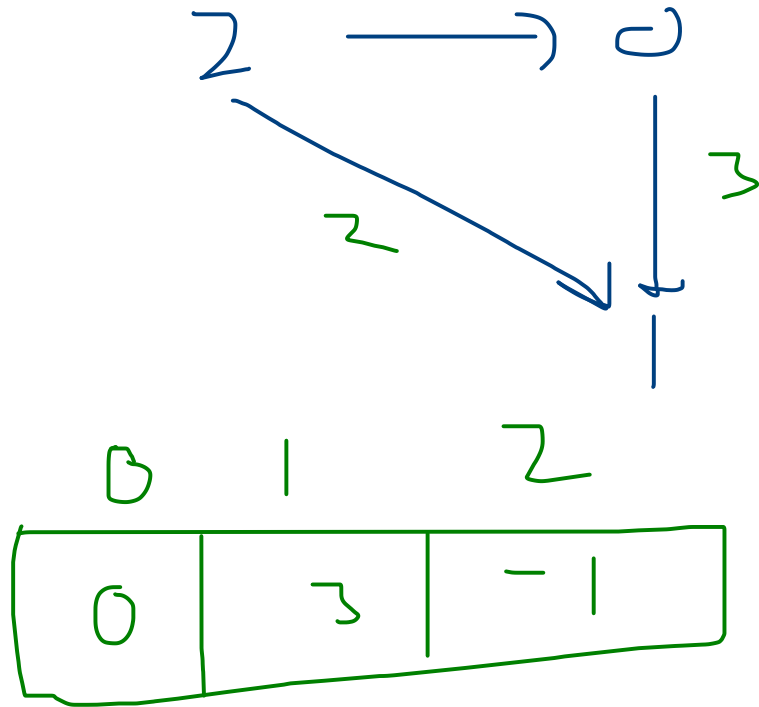
ind=

0	1	2	3	4
0	2	1	1	1
	1		0	

Course Schedule

Shortest Path in DAG

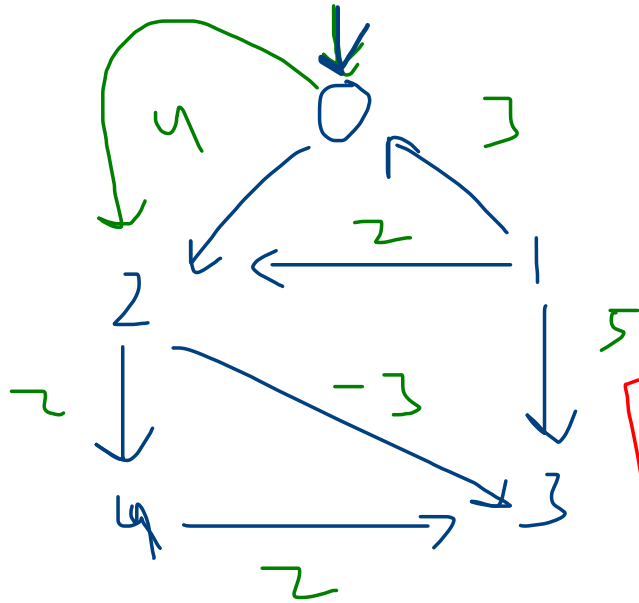
$g_u = 0$



src: 0

==

BFS



\emptyset	2	3	4
-------------	---	---	---

$$if (wt[u,v] + dist[u] < dist[v])$$

$$\{$$

$$dist[v]$$

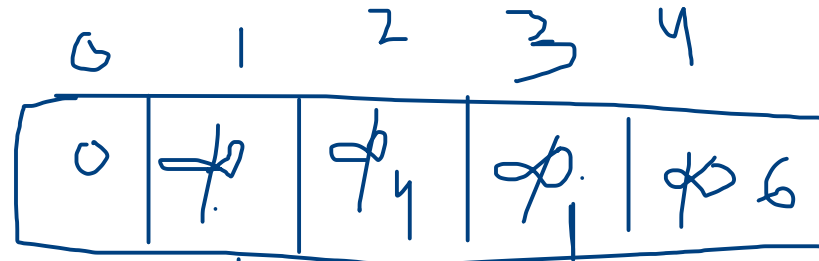
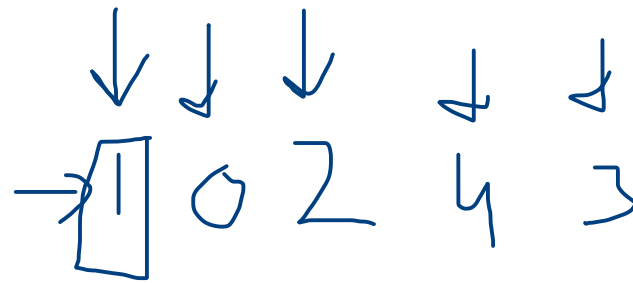
$$= wt[u,v] + dist[u]$$

$$\}$$

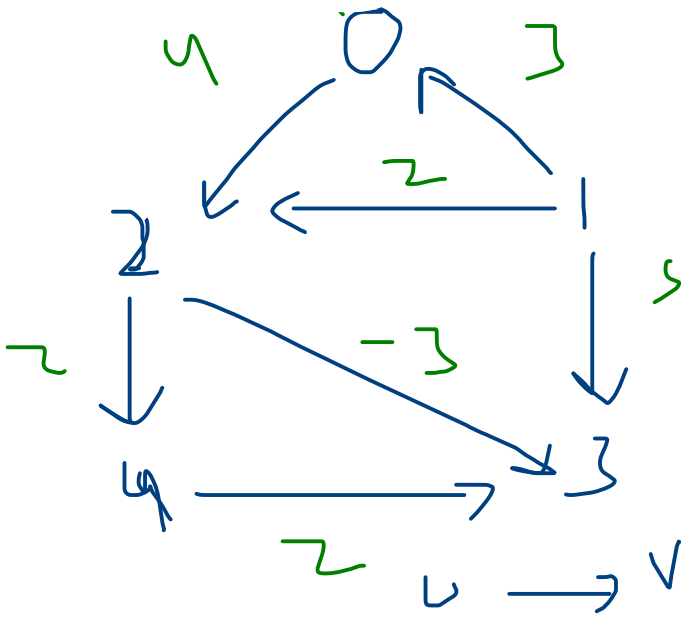
dist

0	1	2	3	4
0	1	4	1	6

6 → v

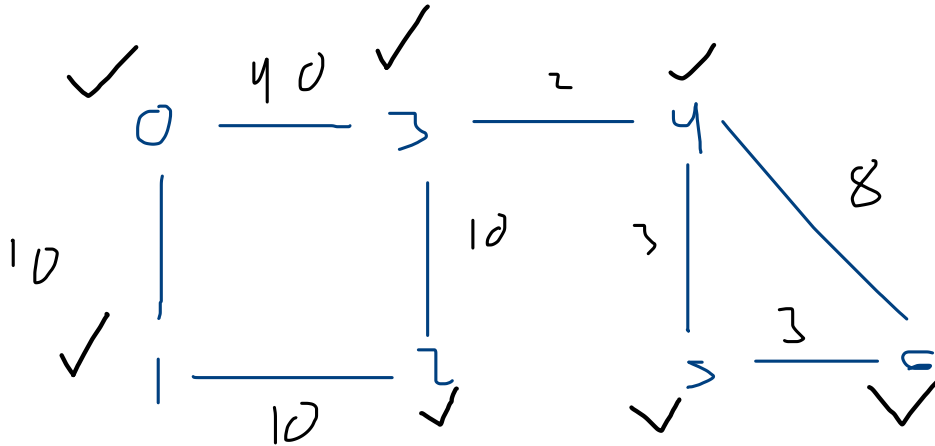


if (d[v] > d[u] + w) {
 d[v] = d[u] + w
}



PC

Dijkstra's shortest path



$$V + E \log V$$

O	PSF	W SF
0	0	0

1	0 - 1	10
2	0 - 1	20
3	0 - 1	30
4	0 - 1	40

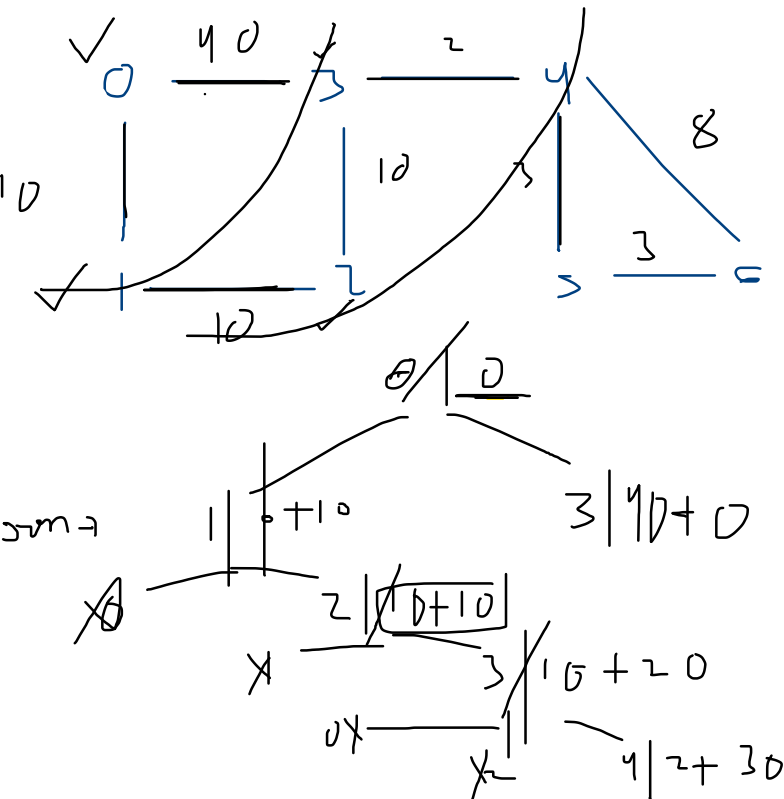
2	0 - 1	20
3	0 - 1	30
4	0 - 1	40
5	0 - 1	50

3	0 - 1	30
4	0 - 1	40
5	0 - 1	50
6	0 - 1	60

4	0 - 1	40
5	0 - 1	50
6	0 - 1	60
7	0 - 1	70

5	0 - 1	50
6	0 - 1	60
7	0 - 1	70
8	0 - 1	80

6	0 - 1	60
7	0 - 1	70
8	0 - 1	80
9	0 - 1	90



```
pq.add(new Triplet(src, 0));
```

```
boolean[] visited = new boolean[v];
```

```
int[] ans = new int[v];
```

```
while(pq.size() > 0) {
```

```
    Triplet rem = pq.remove();
```

```
    if (visited[rem.vtx]) {
```

```
        continue;
```

```
    }
```

```
    visited[rem.vtx] = true;
```

```
    ans[rem.vtx] = rem.cost;
```

```
    for (Edge e: graph[rem.vtx]) {
```

```
        if (!visited[e.nbr]) {
```

```
            Triplet temp = new Triplet(e.nbr, e.wt + rem.cost);
```

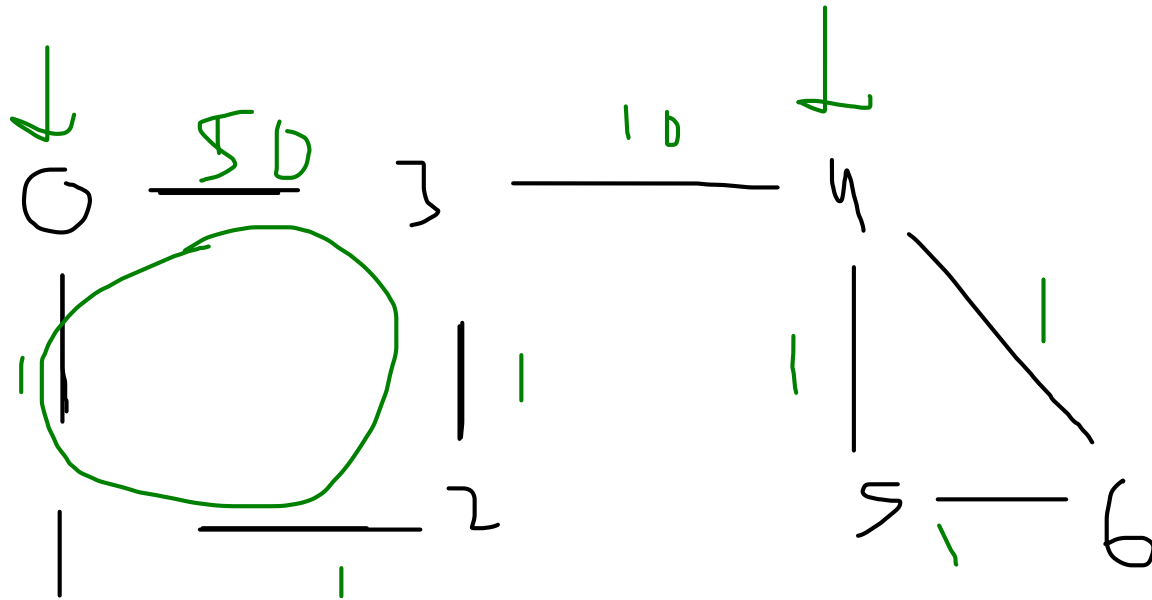
```
            pq.add(temp);
```

```
        }
```

```
    }
```

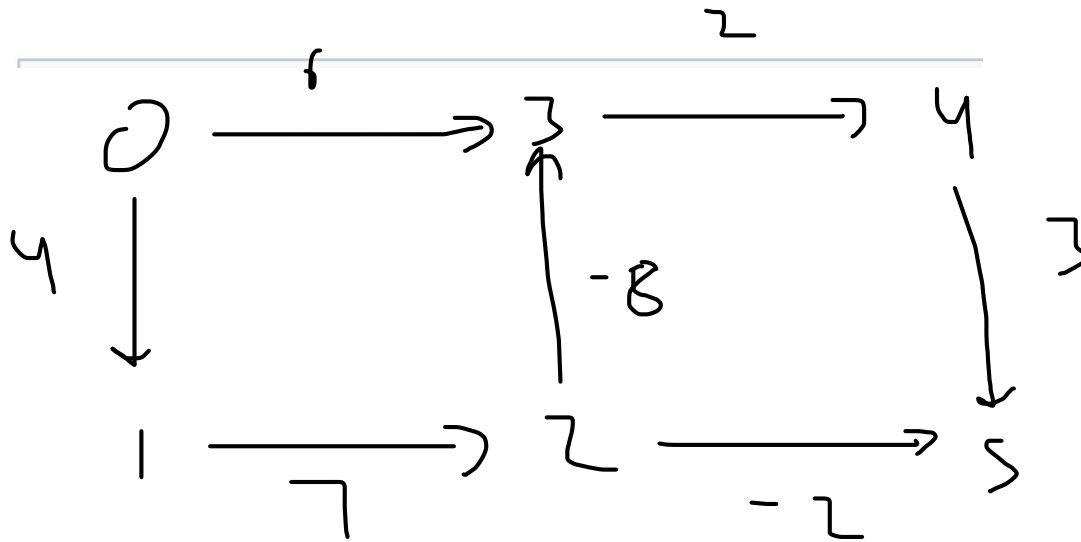
```
}
```

```
return ans;
```

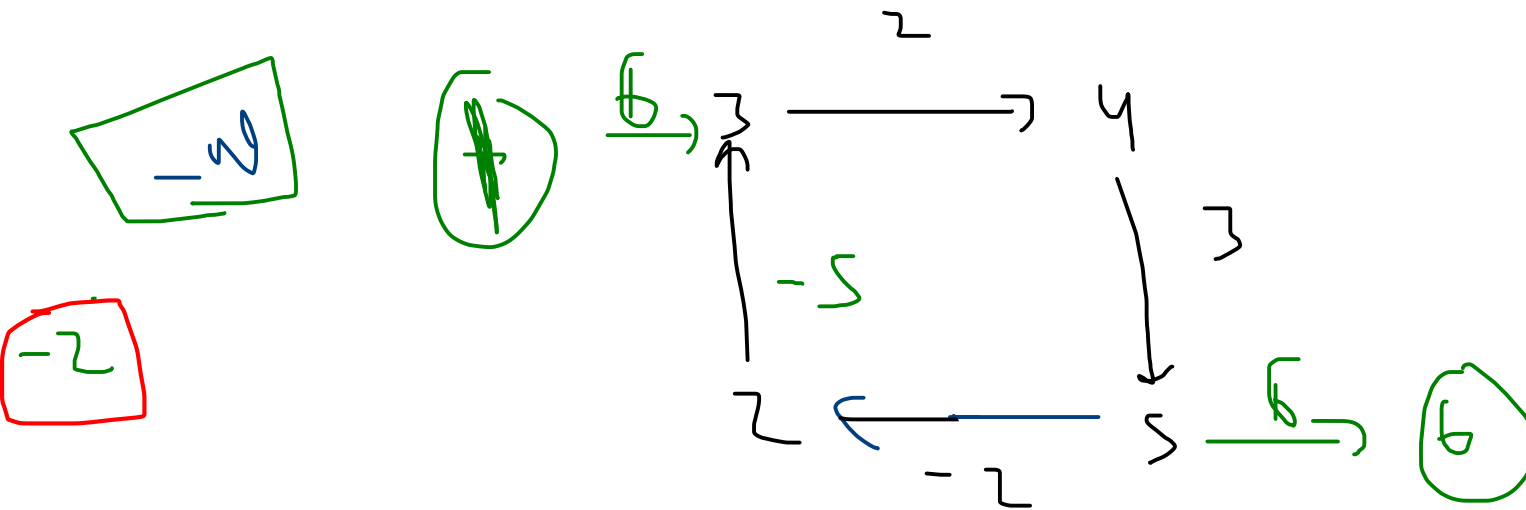


we will never get shorted distance b/w two vertices by looping into cycle

Bellman Ford



negative weight cycle is never possible

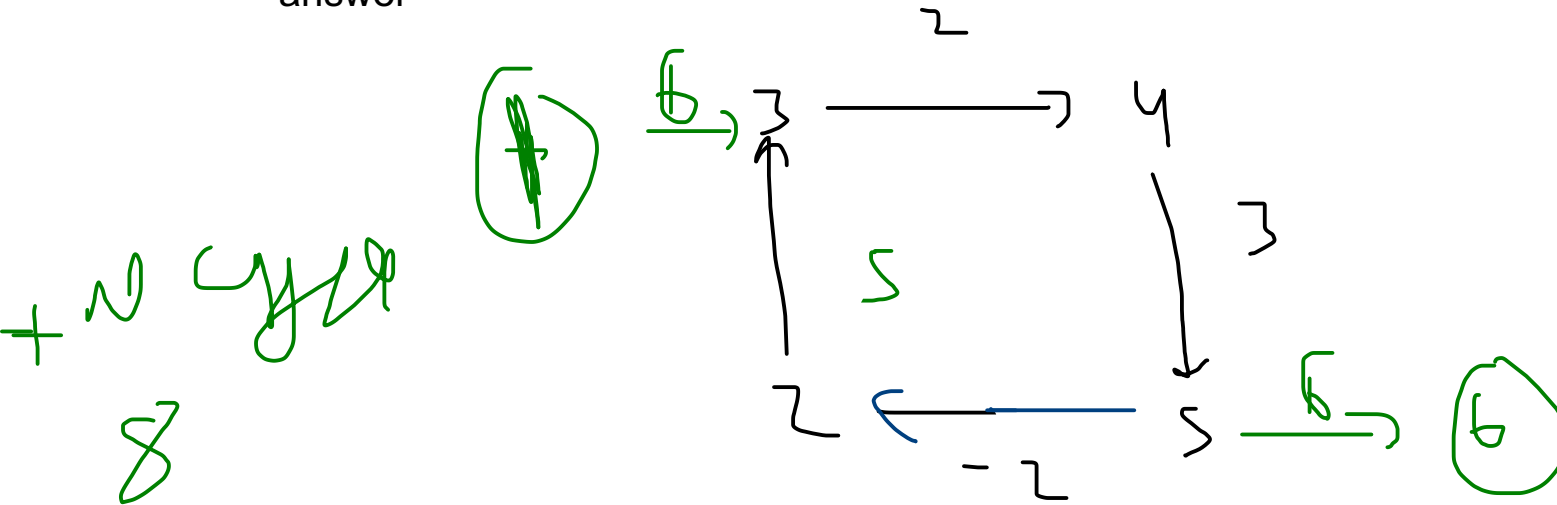


$$6 + 2 + 3 + 6 = 17$$

$$6 + (-2) + 2 + 3 + 6 = 15$$

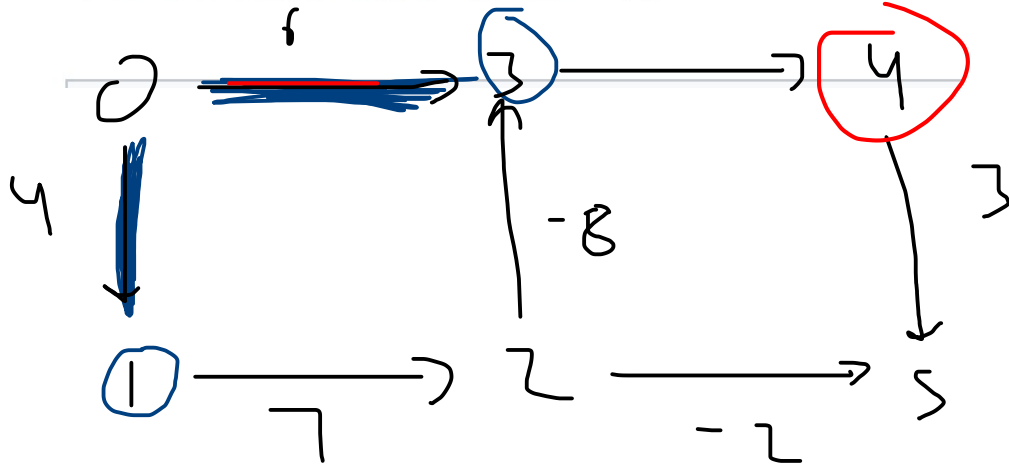
$$6 + (-2) + 2 + 2 + 3 + 6 = 17$$

for positive wt cycle you will never going to complete one cycle to get the answer



st

Bellman Ford



Random

$u \rightarrow v$	w
$4 \rightarrow 5$	3
$3 \rightarrow 4$	2
$2 \rightarrow 3$	-8
$1 \rightarrow 2$	7
$0 \rightarrow 3$	6
$0 \rightarrow 1$	4

$$\text{dist}[v] > \text{dist}[u] + \text{wt}(u, v)$$

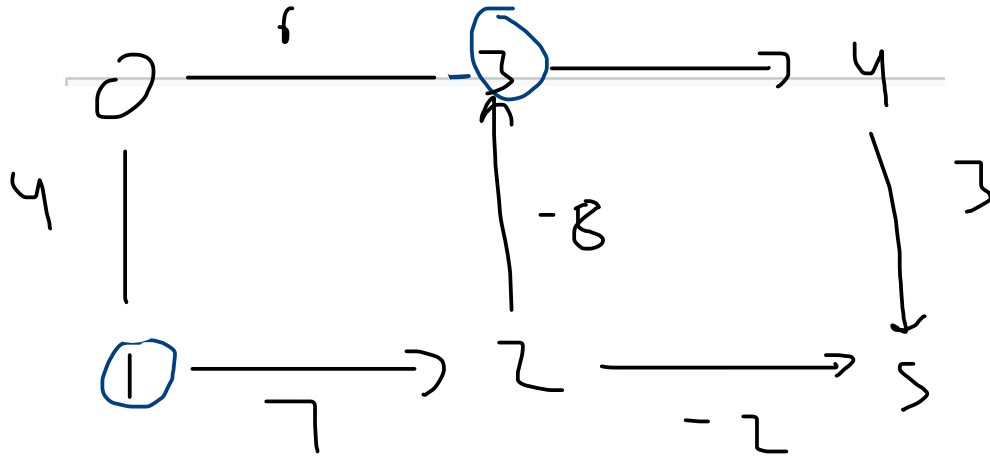
dist

0	1	2	3	4	5
0	4	∞	6	∞	∞

$v - 1$ iterations
 $6 - 1 = 5$

2nd

Bellman Ford



$$\text{dist}[v] > \text{dist}[u] + \text{wt}(u,v)$$

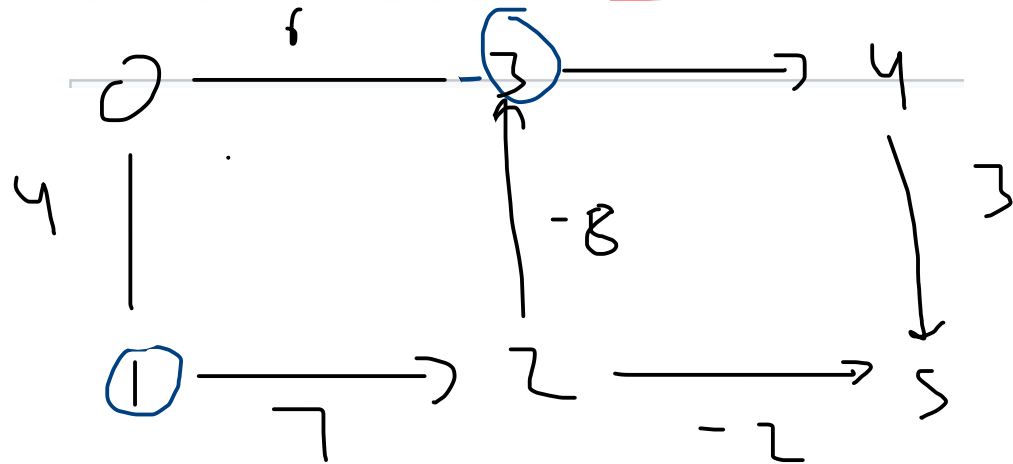
dist

0	1	2	3	4	5
0	4	11	6	8	∞

v - 1 iterations
 $6 - 1 = 5$

$u \rightarrow v$	wt
$\rightarrow 4$	5
$\rightarrow 3$	4
$\rightarrow 2$	3
$\rightarrow 1$	2
$\rightarrow 0$	3
$\rightarrow 0$	1

Bellman Ford



$$\text{dist}[v] > \text{dist}[u] + \text{wt}(u,v)$$

0	→	v	w
4	5	3	
3	4	2	
2	3	-8	
1	2	7	
0	3	6	
0	1	4	

dist

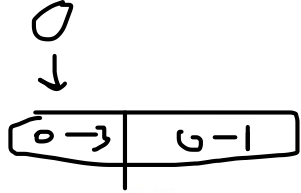
0	1	2	3	4	5
0	4	11	6	8	11
	4	11	3	5	8

v - 1 iterations

$$6 - 1 = 5$$

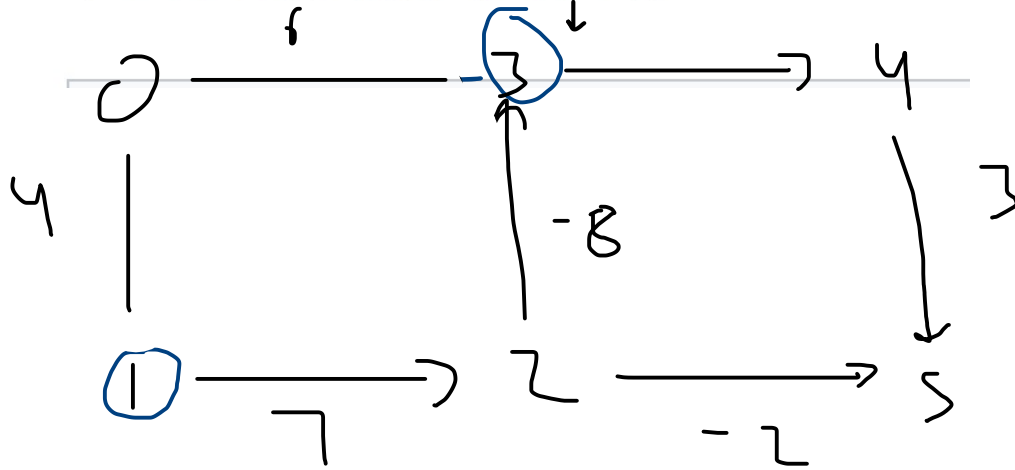
ans of a particular vertice is made in ith iteration whose path is also i edge path away

ith edge path ans ----> in \leq ith iteration



$$\frac{18}{1} + \underline{\underline{2}}$$

Bellman Ford



$$dist[v] > dist[u] + wt(u, v)$$

$u \rightarrow v$	w
4 5	3
3 4	2
\rightarrow 2 3	-8
\rightarrow 1 2	7
\rightarrow 0 3	6
\rightarrow 0 1	4

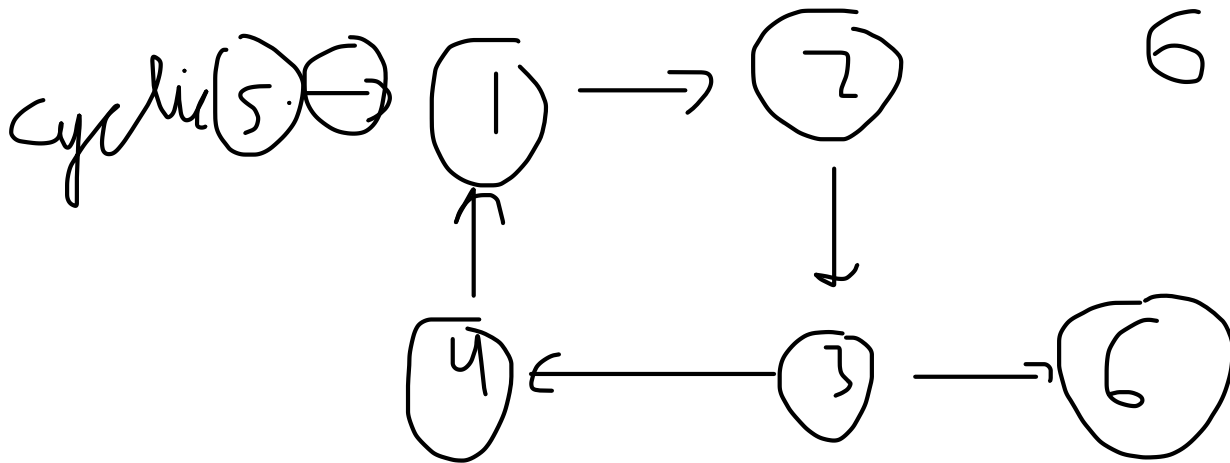
dist

0	1	2	3	4	5
0	7	7	7	7	7
	4	1	6	3	5

v - 1 iterations

$$v - 1$$

Acyclic: $\begin{matrix} \vee & \vee & \vdash r \\ & \vee - 1 & \text{only} \end{matrix}$

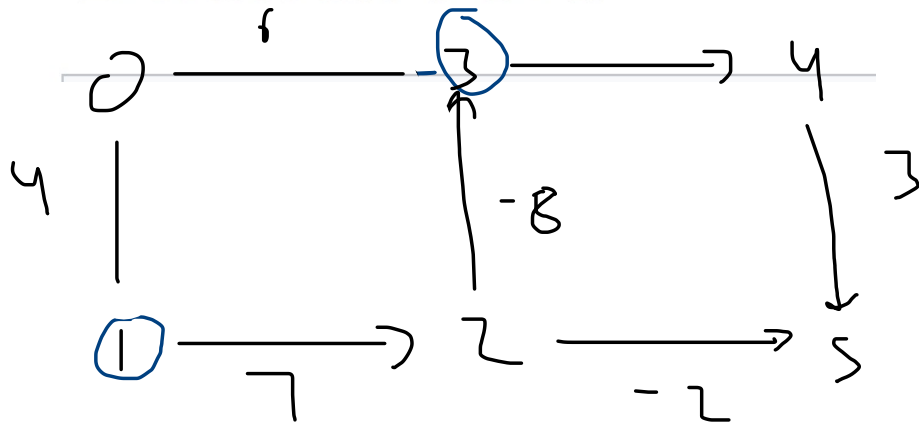


$$6 = v$$

$$= r = \boxed{v - 1}$$

True

Bellman Ford



$$\text{dist}[v] > \text{dist}[u] + \text{wt}(u, v)$$

$$\infty + 3 = \infty$$

dist

0	1	2	3	4	5
0	∞	∞	∞	∞	∞

5

↑

	0	1	3
0	4	5	5
1	3	4	2
2	2	3	-8
3	1	2	7
4	3	3	6

```

int[] dist = new int[vtx + 1];
Arrays.fill(dist, Integer.MAX_VALUE);
dist[src] = 0;

```

```

for (int i = 1; i < vtx; i++) {
    for (int j = 0; j < arr.length; j++) {
        int u = arr[j][0];
        int v = arr[j][1];
        int wt = arr[j][2];

        if (dist[u] == Integer.MAX_VALUE) {
            continue;
        }

        if (dist[v] > dist[u] + wt) {
            dist[v] = dist[u] + wt;
        }
    }
}

```

v - 1 ite

```

for (int i = 1; i < vtx + 1; i++) {
    if (dist[i] == Integer.MAX_VALUE) {
        System.out.print(M + " ");
    } else {
        System.out.print(dist[i] + " ");
    }
}

```

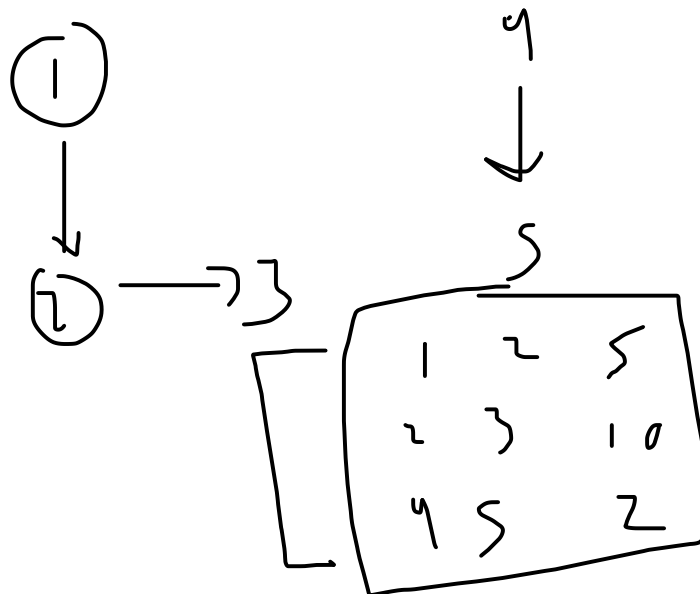
Time Comp of Bellman Ford

```

for (int i = 1; i < V; i++) {
    for (int j = 0; j < arr.length; j++) {
        int u = arr[j][0];
        int v = arr[j][1];
        int wt = arr[j][2];

        if (dist[u] == Integer.MAX_VALUE) {
            continue;
        }

        if (dist[v] > dist[u] + wt) {
            dist[v] = dist[u] + wt;
        }
    }
}
    
```



$$T(n) = \boxed{V \times E} \rightarrow \text{work}$$

\boxed{E}

$$S(n) = O(V)$$

$\boxed{\text{Comments}}$

Shortest Path In A Binary Maze

