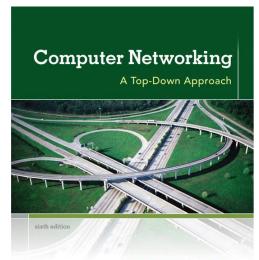
# Chapter 7 Multimedia Networking



KUROSE ROSS

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# Computer Networking: A Top Down Approach 6<sup>th</sup> edition Jim Kurose, Keith Ross Addison-Wesley

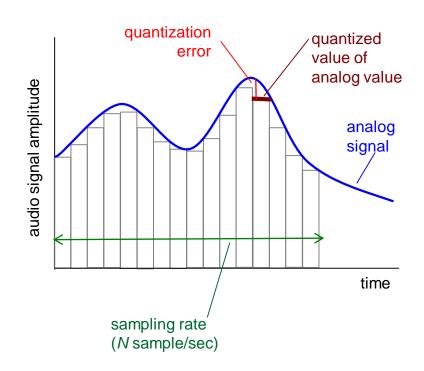
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#### Multimedia networking: outline

- 7. I multimedia networking applications
- 7.2 streaming stored video
- 7.3 voice-over-IP
- 7.4 protocols for real-time conversational applications
- 7.5 network support for multimedia

#### Multimedia: audio

- analog audio signal sampled at constant rate
  - telephone: 8,000 samples/sec
  - CD music: 44,100 samples/sec
- each sample quantized, i.e., rounded
  - e.g., 2<sup>8</sup>=256 possible quantized values
  - each quantized value represented by bits, e.g., 8 bits for 256 values



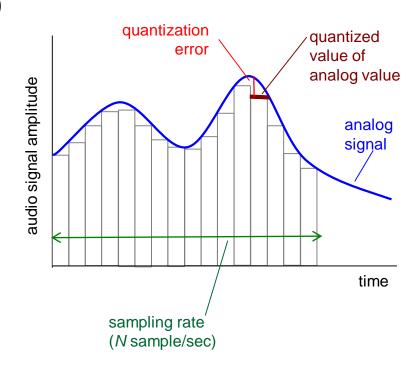
#### Multimedia: audio

- example: 8,000 samples/sec,256 quantized values: 64,000 bps
- receiver converts bits back to analog signal:
  - some quality reduction

#### example rates

CD: 1.411 Mbps

MP3: 96, 128, 160 kbps



#### Multimedia: video

- video: sequence of images displayed at constant rate
  - e.g. 24 images/sec
- digital image: array of pixels
  - each pixel represented by bits
- coding: use redundancy within and between images to decrease # bits used to encode image
  - spatial (within image)
  - temporal (from one image to next)

spatial coding example: instead of sending N values of same color (all purple), send only two values: color value (purple) and number of repeated values (N)



frame i

temporal coding example:\
instead of sending
complete frame at i+1,
send only differences from
frame i



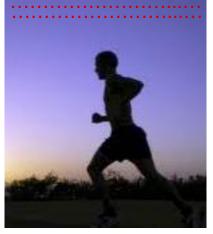
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Multmedia Networking

#### Multimedia: video

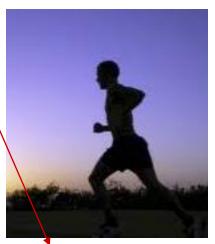
- CBR: (constant bit rate): video encoding rate fixed
- VBR: (variable bit rate): video encoding rate changes as amount of spatial, temporal coding changes
- examples:
  - MPEG I (CD-ROM) 1.5 Mbps
  - MPEG2 (DVD) 3-6 Mbps
  - MPEG4 (often used in Internet, < I Mbps)</li>

spatial coding example: instead of sending N values of same color (all purple), send only two values: color value (purple) and number of repeated values (N)



frame i

temporal coding example: instead of sending complete frame at i+1, send only differences from frame i



frame *i*+1

Multmedia Networking

#### Multimedia networking: 3 application types

- streaming, stored audio, video
  - streaming: can begin playout before downloading entire file
  - stored (at server): can transmit faster than audio/video will be rendered (implies storing/buffering at client)
  - e.g., YouTube, Netflix, Hulu
- conversational voice/video over IP
  - interactive nature of human-to-human conversation limits delay tolerance
  - e.g., Skype, WhatsApp
- streaming live audio, video
  - e.g., live sporting event (football)

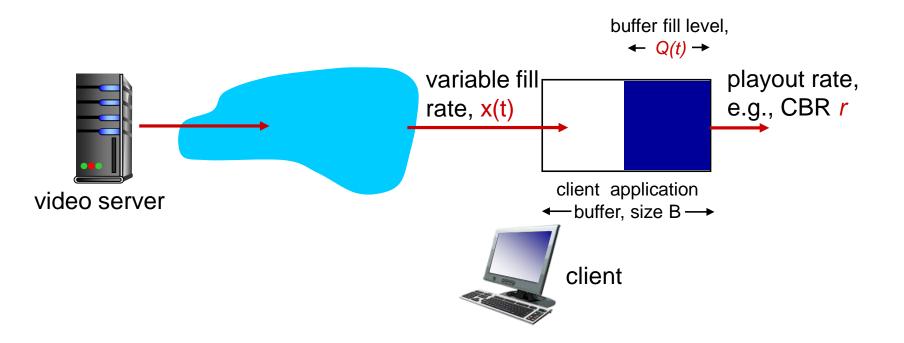
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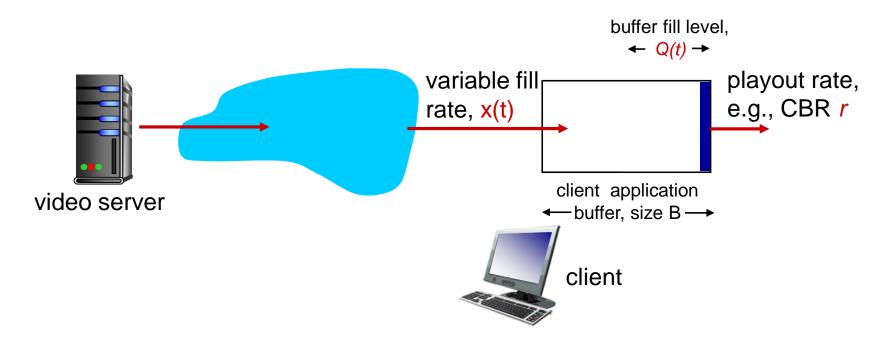
# Streaming stored video: challenges

- continuous playout constraint: once client playout begins, playback must match original timing
  - ... but network delays are variable (jitter), so will need client-side buffer to match playout requirements
- other challenges:
  - client interactivity: pause, fast-forward, rewind, jump through video
  - video packets may be lost, retransmitted

# Client-side buffering, playout



#### Client-side buffering, playout



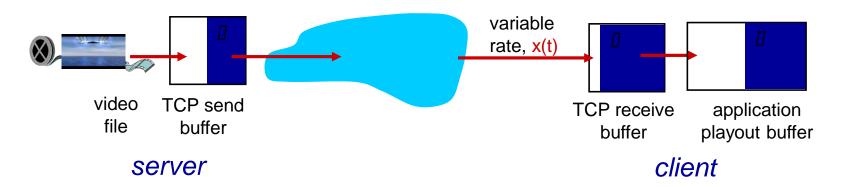
- 1. Initial fill of buffer until playout begins at  $t_p$
- 2. playout begins at t<sub>D.</sub>
- 3. buffer fill level varies over time as fill rate x(t) varies and playout rate r is constant

#### Streaming multimedia: UDP

- server sends at rate appropriate for client
  - often: send rate = encoding rate = constant rate
  - transmission rate can be oblivious to congestion levels
- short playout delay (2-5 seconds) to remove network jitter
- error recovery: application-level, timeipermitting
- RTP [RFC 2326]: multimedia payload types
- UDP may not go through firewalls

#### Streaming multimedia: HTTP

- multimedia file retrieved via HTTP GET
- send at maximum possible rate under TCP



- fill rate fluctuates due to TCP congestion control, retransmissions (in-order delivery)
- larger playout delay: smooth TCP delivery rate
- HTTP/TCP passes more easily through firewalls

#### Streaming multimedia: DASH

- DASH: Dynamic, Adaptive Streaming over HTTP
- \* server:
  - divides video file into multiple chunks
  - each chunk stored, encoded at different rates
  - manifest file: provides URLs for different chunks

#### client:

- periodically measures server-to-client bandwidth
- consulting manifest, requests one chunk at a time
  - chooses maximum coding rate sustainable given current bandwidth
  - can choose different coding rates at different points in time (depending on available bandwidth at time)

#### Streaming multimedia: DASH

- DASH: Dynamic, Adaptive Streaming over HTTP
- "intelligence" at client: client determines
  - when to request chunk (so that buffer starvation, or overflow does not occur)
  - what encoding rate to request (higher quality when more bandwidth available)
  - where to request chunk (can request from URL server that is "close" to client or has high available bandwidth)

#### Content distribution networks

- challenge: how to stream content (selected from millions of videos) to hundreds of thousands of simultaneous users?
- option 1: single, large "mega-server"
  - single point of failure
  - point of network congestion
  - long path to distant clients
  - multiple copies of video sent over outgoing link
- ....quite simply: this solution doesn't scale

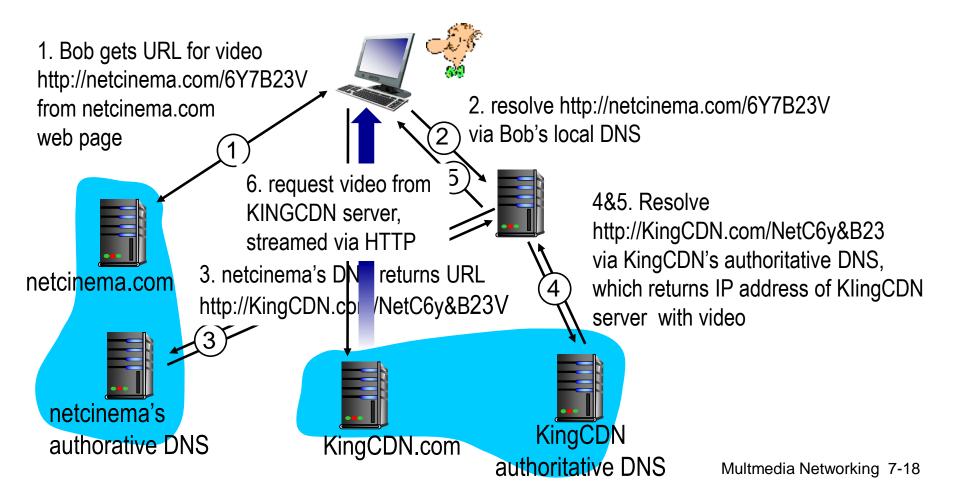
#### Content distribution networks

- challenge: how to stream content (selected from millions of videos) to hundreds of thousands of simultaneous users?
- option 2: store/serve multiple copies of videos at multiple geographically distributed sites (CDN)
  - enter deep: push CDN servers deep into many access networks
    - close to users
    - used by Akamai, 1700 locations

#### CDN: "simple" content access scenario

Bob (client) requests video http://netcinema.com/6Y7B23V

video stored in CDN at http://KingCDN.com/NetC6y&B23V



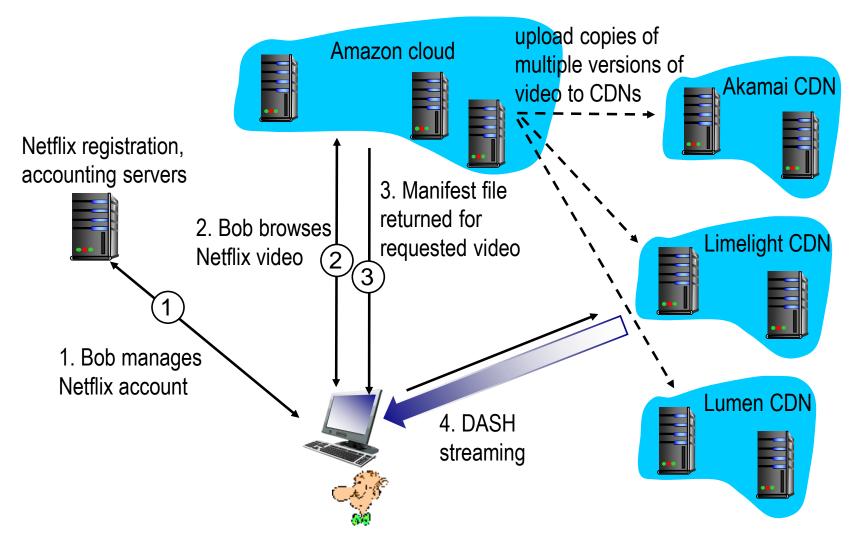
# CDN cluster selection strategy

- challenge: how does CDN DNS select "good"
   CDN node to stream to client
  - pick CDN node geographically closest to client
  - pick CDN node with shortest delay (or min # hops) to client (CDN nodes periodically ping access ISPs, reporting results to CDN DNS)
- alternative: let client decide give client a list of several CDN servers
  - client pings servers, picks "best"
  - Netflix approach

### Case study: Netflix

- 30% downstream US traffic in 2011
- owns very little infrastructure, uses 3<sup>rd</sup> party services:
  - own registration, payment servers
  - Amazon (3<sup>rd</sup> party) cloud services:
    - Netflix uploads studio master to Amazon cloud
    - create multiple version of movie (different endodings) in cloud
    - upload versions from cloud to CDNs
    - Cloud hosts Netflix web pages for user browsing
  - three 3<sup>rd</sup> party CDNs host/stream Netflix content: Akamai, Limelight, Lumen

#### Case study: Netflix



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# Voice-over-IP (VoIP)

- VolP end-end-delay requirement: needed to maintain "conversational" aspect
  - higher delays noticeable, impair interactivity
  - < 150 msec: good</p>
  - > 400 msec bad
  - includes application-level (packetization, playout), network delays
- session initialization: how does callee advertise IP address, port number, encoding algorithms?
- value-added services: call forwarding, screening, recording

#### VoIP characteristics

- speaker's audio: alternating talk spurts, silent periods.
  - 64 kbps during talk spurt
  - 20 msec chunks at 8 Kbytes/sec: 160 bytes of data
- application-layer header added to each chunk
- chunk+header encapsulated into UDP or TCP segment
- application sends segment into socket every 20 msec

### VoIP: packet loss, delay

- network loss: IP datagram lost due to network congestion (router buffer overflow)
- delay loss: IP datagram arrives too late for playout at receiver
  - delays: processing, queueing in network; end-system (sender, receiver) delays
  - typical maximum tolerable delay: 400 ms
- loss tolerance: depending on voice encoding, loss concealment, packet loss rates between 1% and 10% can be tolerated

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#### Real-Time Protocol (RTP)

- RTP specifies packet structure for packets carrying audio, video data
- \* RFC 3550
- RTP packet provides
  - payload type identification
  - packet sequence numbering
  - time stamping

- RTP runs in end systems
- RTP packets encapsulated in UDP segments
- interoperability: if two VoIP applications run RTP, they may be able to work together

#### RTP runs on top of UDP

RTP libraries provide transport-layer interface that extends UDP:

- port numbers, IP addresses
- payload type identification
- packet sequence numbering
- time-stamping

transport RTP
layer UDP
IP
Data Link
Physical

#### RTP and QoS

- \* RTP does *not* provide any mechanism to ensure timely data delivery or other QoS guarantees
- RTP encapsulation only seen at end systems (not by intermediate routers)
  - routers provide best-effort service, making no special effort to ensure that RTP packets arrive at destination in timely matter

#### SIP: Session Initiation Protocol [RFC 3261]

#### long-term vision:

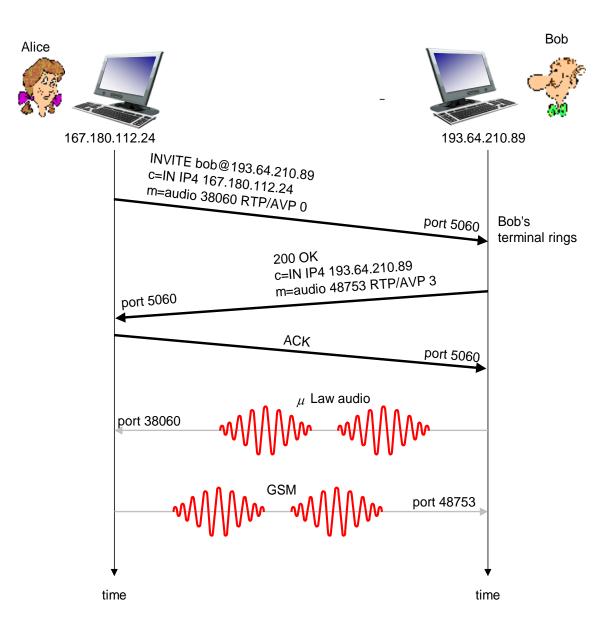
- all telephone calls, video conference calls take place over Internet
- people identified by names or e-mail addresses, rather than by phone numbers
- can reach callee (if callee so desires), no matter where callee roams, no matter what IP device callee is currently using

#### SIP services

- SIP provides mechanisms for call setup:
  - for caller to let callee know she wants to establish a call
  - so caller, callee can agree on media type, encoding
  - to end call

- determine current IP address of callee:
  - maps mnemonic identifier to current IP address
- call management:
  - add new media streams during call
  - change encoding during call
  - invite others
  - transfer, hold calls

#### Example: setting up call to known IP address



- \* Alice's SIP invite message indicates her port number, IP address, encoding she prefers to receive (PCM μlaw)
- Bob's 200 OK message indicates his port number, IP address
- ❖ SIP messages can be sent over TCP or UDP; here sent over RTP/UDP
- default SIP port number is5060

# Setting up a call (more)

- codec negotiation:
  - suppose Bob doesn't have PCM µlaw encoder
  - Bob will instead reply with 606 Not Acceptable Reply, listing his encoders. Alice can then send new INVITE message, advertising different encoder

- rejecting a call
  - Bob can reject with replies "busy," "gone," "payment required," "forbidden"
- media can be sent over RTP or some other protocol

# SIP proxy

- another function of SIP server: proxy
- Alice sends invite message to her proxy server
  - contains address sip:bob@domain.com
  - proxy responsible for routing SIP messages to callee, possibly through multiple proxies
- Bob sends response back through same set of SIP proxies
- proxy returns Bob's SIP response message to Alice
  - contains Bob's IP address
- SIP proxy analogous to local DNS server plus TCP setup

#### Comparison with H.323

- H.323: another signaling protocol for real-time, interactive multimedia
- H.323: complete, vertically integrated suite of protocols for multimedia conferencing: signaling, registration, admission control, transport, codecs
- SIP: single component.
   Works with RTP, but does not mandate it. Can be combined with other protocols, services

- H.323 comes from the ITU (telephony)
- SIP comes from IETF: borrows much of its concepts from HTTP
  - SIP has Web flavor;
     H.323 has telephony flavor

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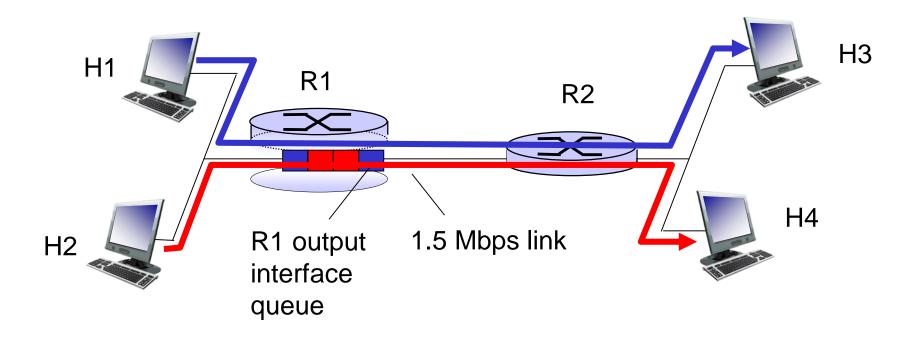
#### Network support for multimedia

Approach	Granularity	Guarantee	Mechanisms	Complex	Deployed?
Making best	All traffic	None or	No network	low	everywhere
of best effort	treated	soft	support (all at		
service	equally		application)		
Differentiated	Traffic	None of	Packet market,	med	some
service	"class"	soft	scheduling,		
			policing.		
Per-	Per-	Soft or hard	Packet market,	high	little to
connection	connection	after flow	scheduling,		none
QoS	flow	admitted	policing, call		
			admission		

#### Dimensioning best effort networks

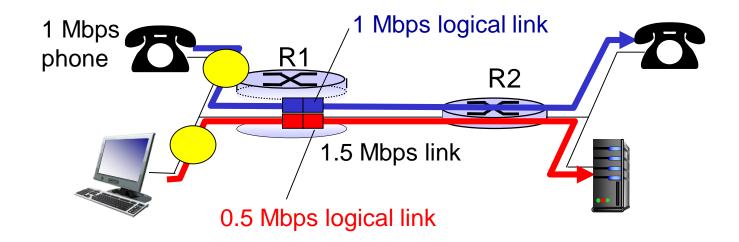
- approach: deploy enough link capacity so that congestion doesn't occur, multimedia traffic flows without delay or loss
  - low complexity of network mechanisms (use current "best effort" network)
  - high bandwidth costs
- challenges:
  - network dimensioning: how much bandwidth is "enough?"
  - estimating network traffic demand: needed to determine how much bandwidth is "enough" (for that much traffic)

#### Multiple classes of service: scenario



#### Principles for QOS guarantees (more)

 allocating fixed (non-sharable) bandwidth to flow: inefficient use of bandwidth if flows doesn't use its allocation



#### Principle 3

while providing isolation, it is desirable to use resources as efficiently as possible

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