

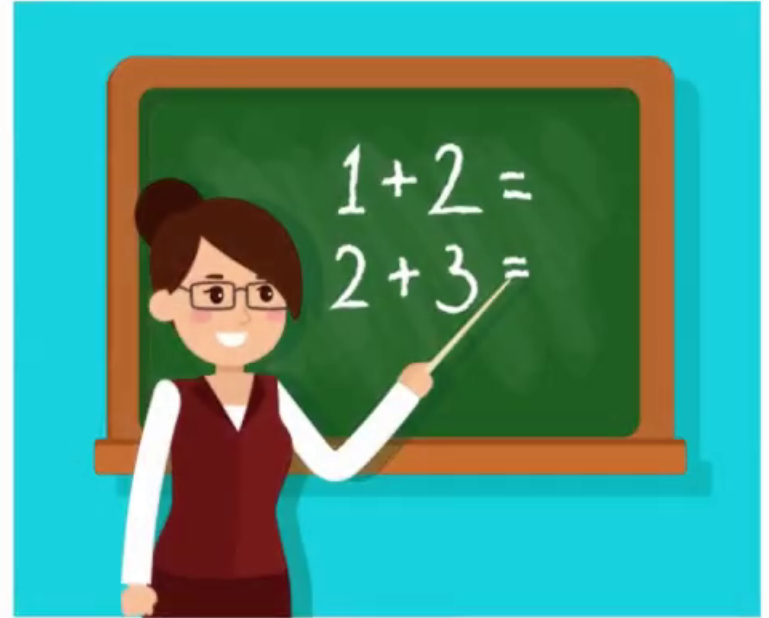
RANGED SUM

ANIMATED INTUITION

$ARR[L] += 1$

$ARR[R] -= 1$

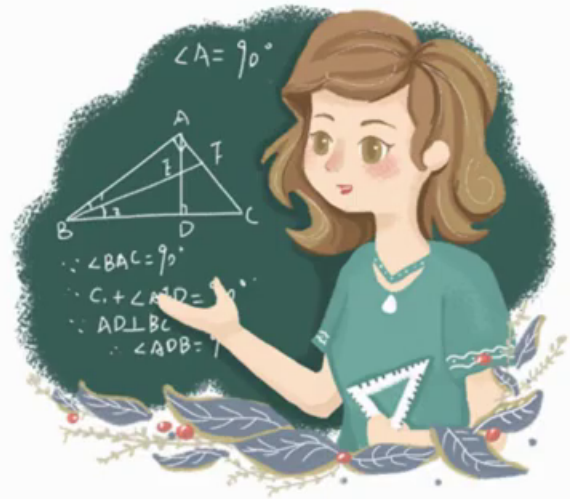
ADD FROM L TO R



Ranged Sum

problem description :

given an array of zeros
of size n and t queries
in each query you are given
two numbers l r add one to
the numbers from l to r
print the final array



Ranged Sum

problem :

input :

5

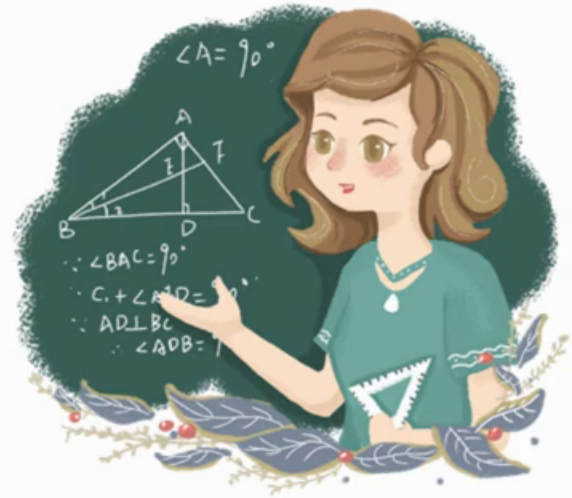
2

2 4

1 4

output :

1 2 2 2 0



Ranged Sum

arr 0 0 0 0 0

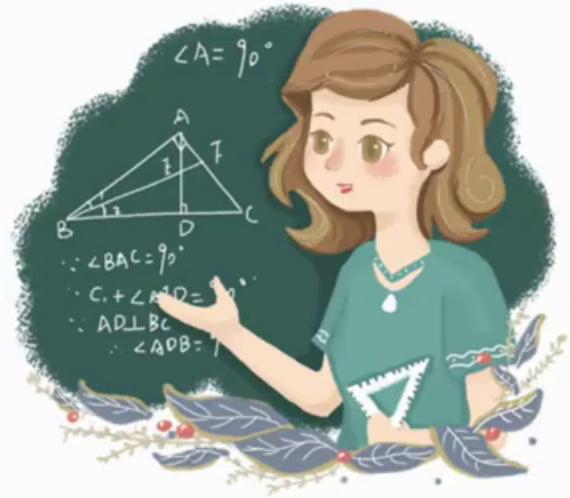
l=2 r=4

l=1 r=4

sum 1 2 2 2 0

if we used two nested
for loops

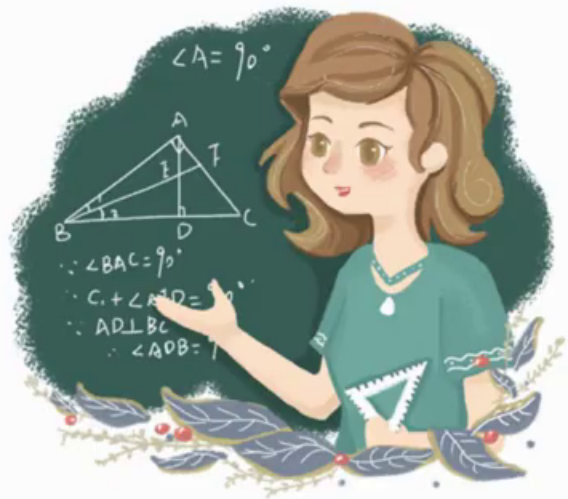
time complexity will be $O(t*n)$



Ranged Sum

okay the idea to solve
this problem efficiently is to

find a way to calculate the final array
by using signs you only know
l and r



Ranged Sum

we only need to know where
to start and where to stop

arr 0 0 0 0 0 0

l=2 r=4

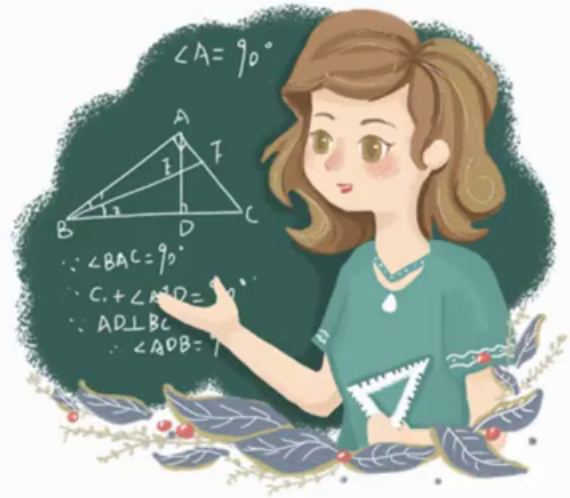
l=1 r=4

arr[l]+=1 arr[r+1]-=1

output : 1 1 0 0 0 -2

if (i){arr[i]+=arr[i-1];}

output : 1 2 2 2 2 0



Ranged Sum

$n=5$ but i am using 6 positions

arr 0 0 0 0 0 0

$l=2$ $r=4$

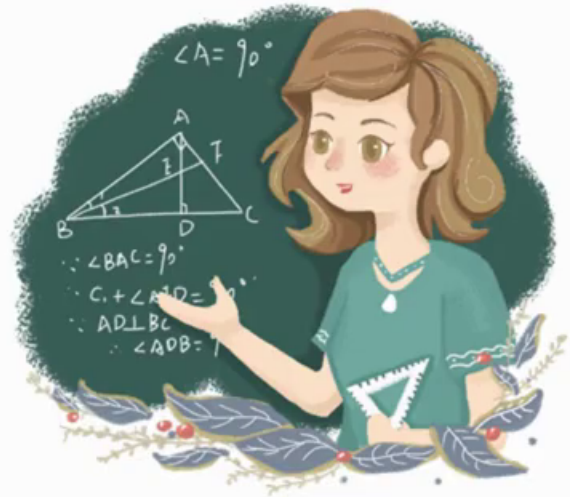
arr[l]+=1 arr[r+1]-=1

arr 0 1 0 0 -1 0

if (i){arr[i]+=arr[i-1];}

arr 0 1 1 1 0 0

**1 is my starting position
and -1 is where i stop**



Ranged Sum

how we improved the complexity ?

inside each query

$\text{arr}[l] += \text{val}$ $\text{arr}[r+1] -= \text{val}$ $O(1)$

after all the queries

$\text{arr}[i] += \text{arr}[i-1]$ in a loop $O(n)$

from start to end

over all take the maximum

