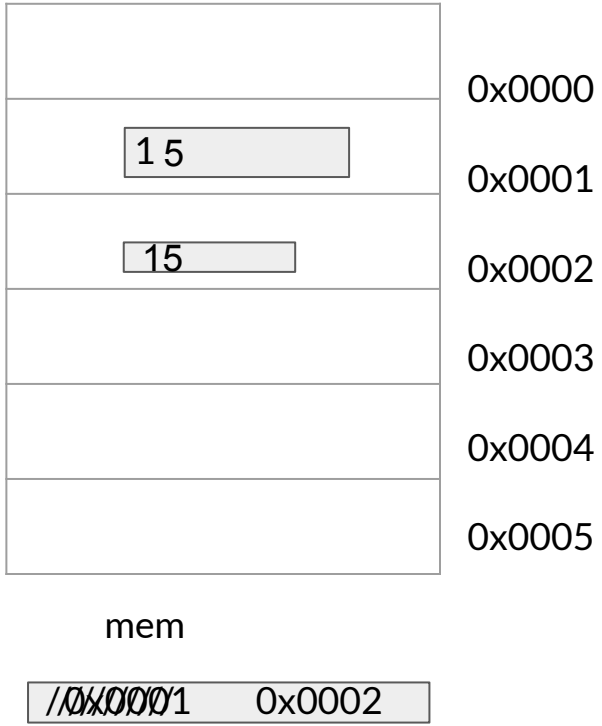
Section 5

Sheet 4

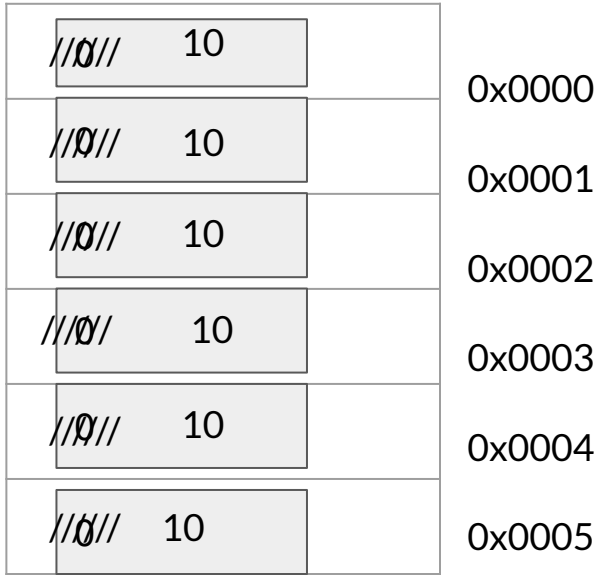
pointer

```
Void main ()
{
    Char x=5;
    Char * pr;
    ptr= &x;
    *ptr=15;
    ptr++;
    *ptr=15;
}
```



pointer

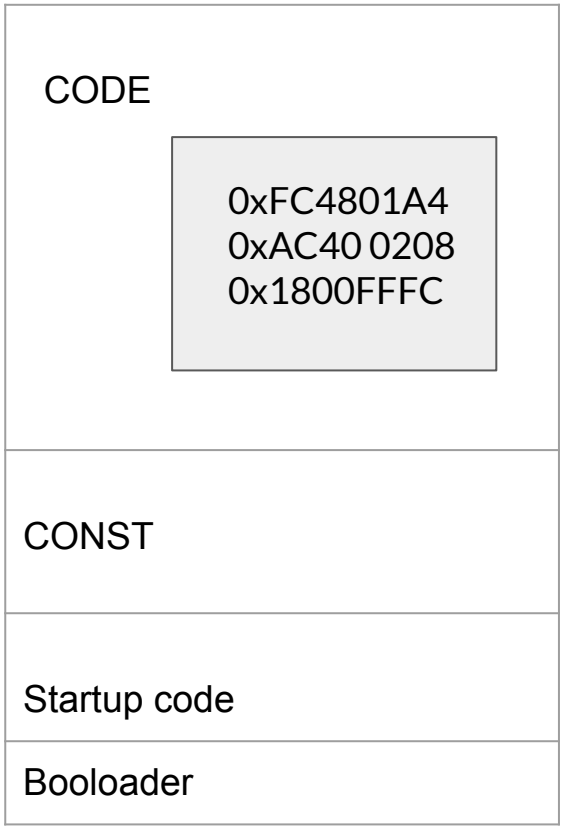
```
Void main ()
{
   Char arr[6]={0,0,0,0,0,0};
   Char *ptr=arr;
   for(i=0;i<6;i++)
     *(ptr+i)=10;

}
```



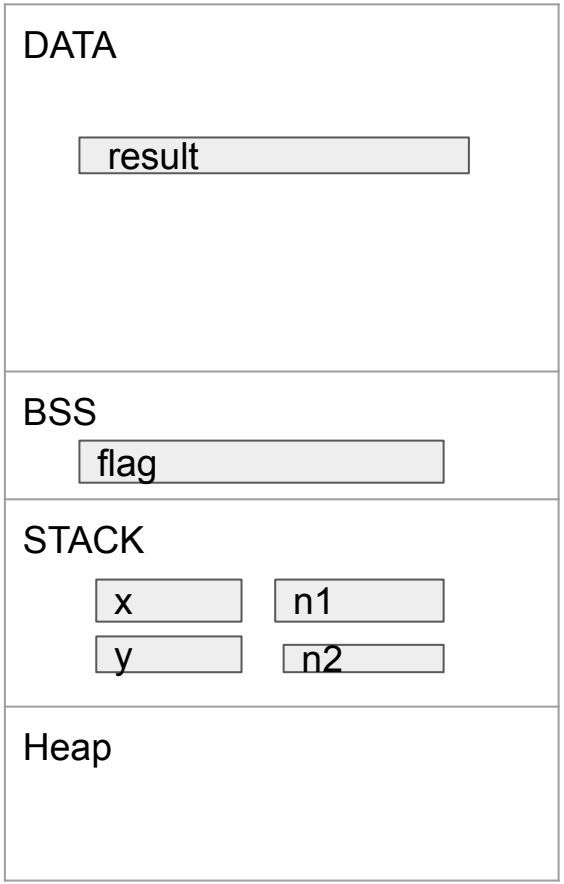
mem

0x000 0x0001 0x0002 and so on

stack




```
int result =1;
char flag;
int add(int x,inty)
{
   result=x+y;
   return result;
}
int sub(int n1,n2)
{
  if(n1>n2){flag=1;}
   result=n1-n2;
   return result;
}

Void main()
  result=sub(5,4);
 if(flag==1)
{
  printf(result);
}
```



stack

PUSH {r0}

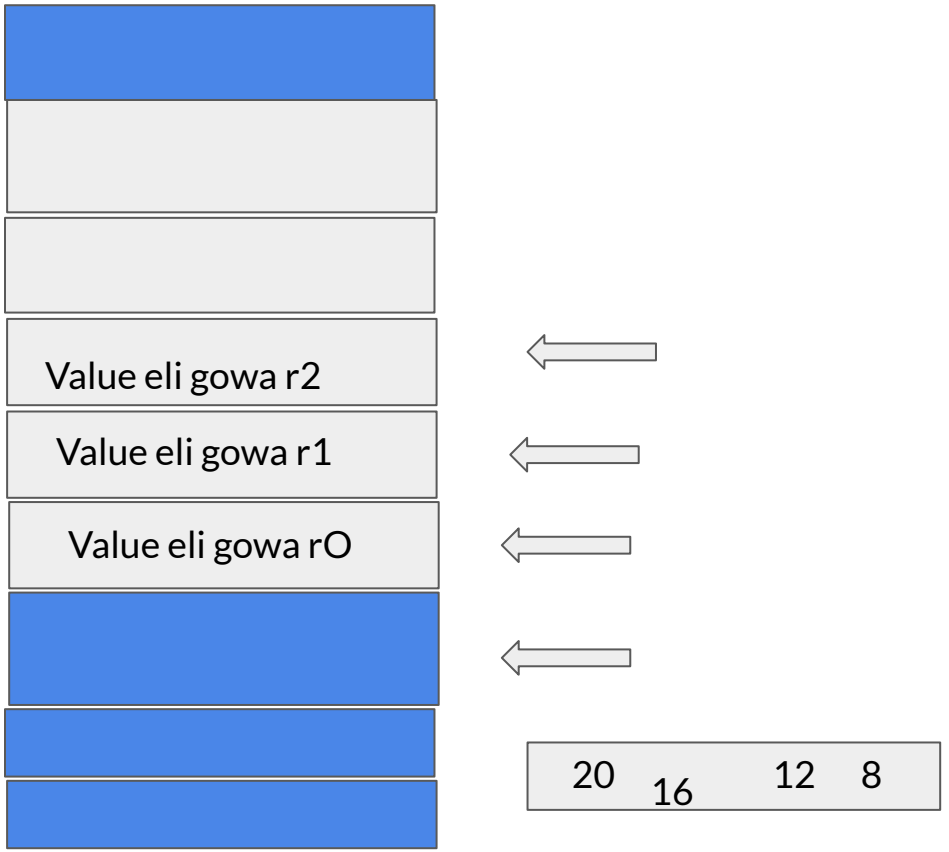
PUSH {r1}

PUSH {r2}

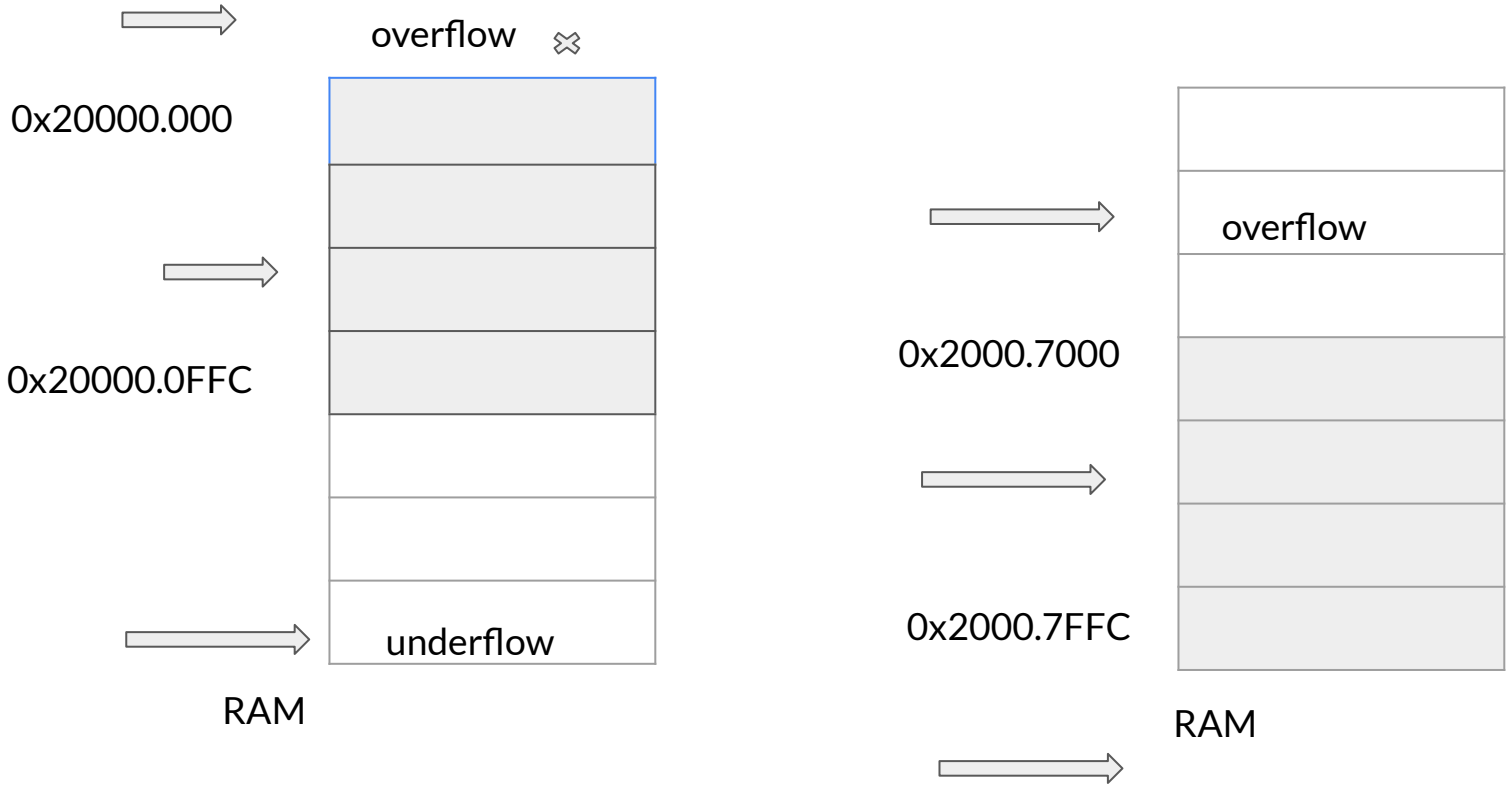
POP {r2}

POP{r1}

POP {r0}



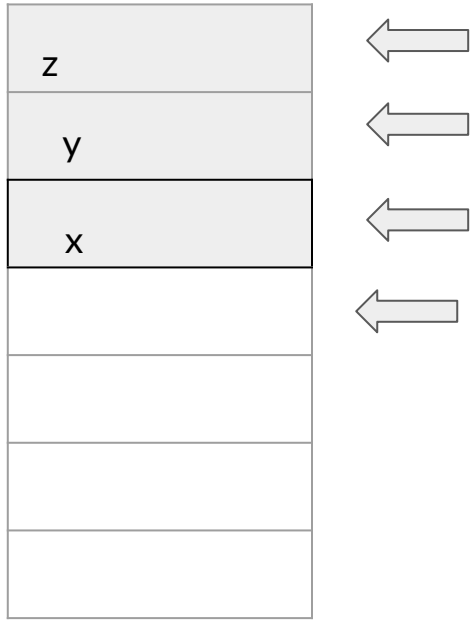
Value @add16	Value @add12	Value @add8
r0	r1	r2


overflow
0x20000.000
0x20000.0FFC
underflow
RAM
overflow
0x2000.7000
0x2000.7FFC
RAM

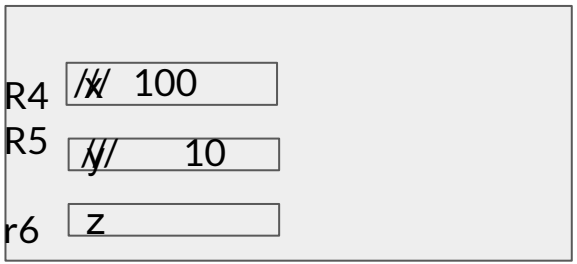
stack

```
void main
{
  r4=x,r5=y,6=z
  f1()
  r4?

}
```

```
f1()
{
        Push {r6,r5,r4}

     Move r4 #100
     add r5 r5 r6
     Pop{r6}
     Pop{r5}
     pop{4}
}
```

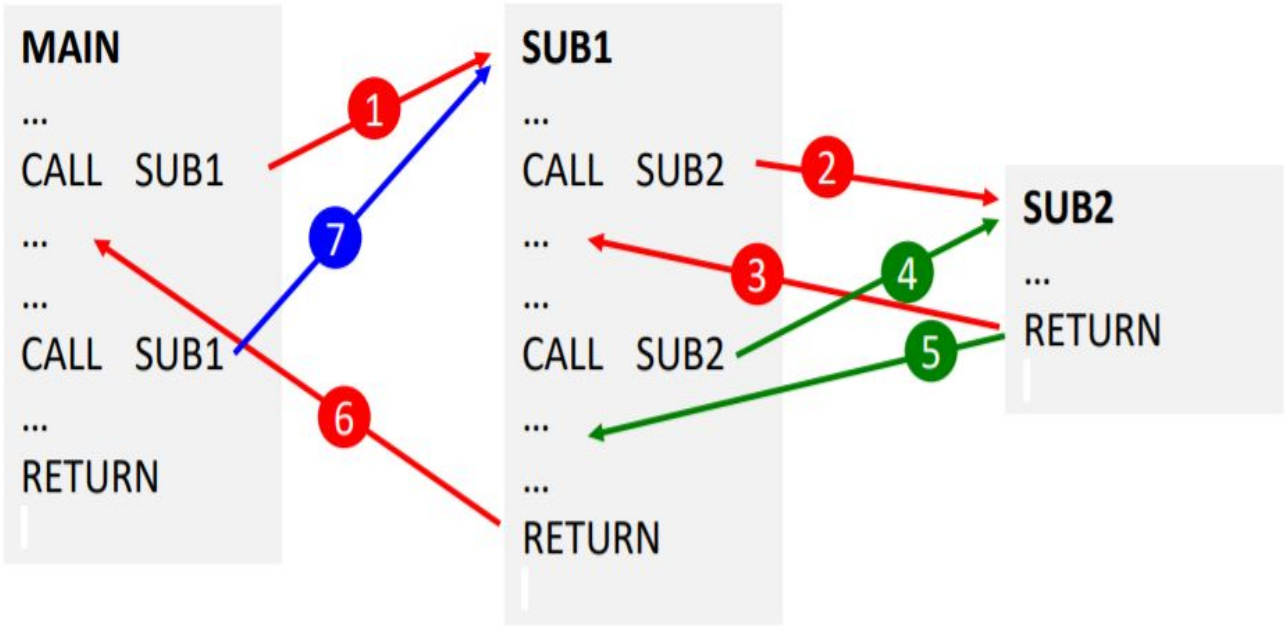


mem



Subroutine

a subroutine is a sequence of instructions that performs a particular task
subroutine called wherever task needs to be performed



Subroutine

- to call a subroutine use the BL (branch and link) instruction
- saves return address in link register (LR = R14)

```
0x00000400    BL      SUB1    ; LR = 0x00000404 (return address)
0x00000404    ...             ;
```

return address = address of next instruction

to return from a subroutine use BX (branch and exchange) specifying the link register (LR)

```
              BX      LR      ; PC (program counter) = LR
```

Subroutine

at the start of every non leaf subroutine, push the contents of LR (link register), which contains the return address, onto the system stack

return from a non leaf subroutine by popping return address from stack and assigning to the PC (program counter)

```
SUB1  PUSH   {LR}     ; push link register onto stack
      ...
      ...
      POP    {PC}     ; return by popping saved return address into PC
```

Stack task

Write by c++ stack:

1- user can create stack with any size ,push,pop and view content.

2- guarantee no stack over flow ,or stack underflow.

3- use copy constructor.

```
 97  }
 98  int main()
 99  {
100   Stack s1(3);
101      s1.push(5);
102      s1.push(14);
103      viewcontent(s1);
104      s1.pop();
105       viewcontent(s1);
106
107      return 0;
108  }
109
```

gs & others

Code::Blocks | Search results | Cccc | Build log | Build messages | CppCheck/Vera++ | CppCheck/Vera++ messages

File	Line	Message
		=== Build: Debug in iti32 (compiler: GNU GCC Compiler) ===
C:\Users\mo...		In member function 'int Stack::stack(Stack&)':
C:\Users\mo...	97	warning: no return statement in function returning non-void [-Wreturn-type]

Sheet 4

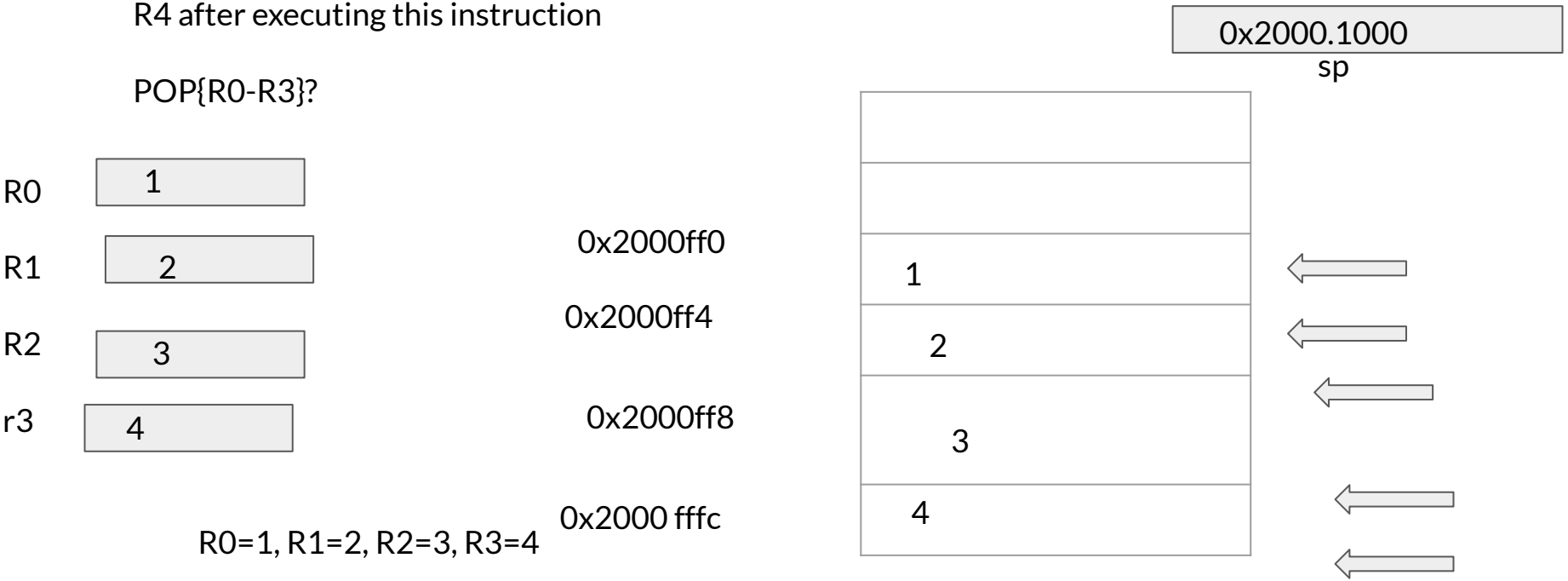
Q1. When does the LR have to be pushed on the stack?

When a function calls another function, the LR is saved in the stack to avoid losing the return address of the
caller function.

Q2. Show the SP value and the content of stack after executing this instruction PUSH{R4, R6-R8} assuming

the SP initially equals 0x2000.1000 and R4=1, R6=2, R7=3, R8=4. What will be the values of the registers R0-
R4 after executing this instruction

POP{R0-R3}?



R0=1, R1=2, R2=3, R3=4

Q3. Explain how does the return from subroutine work in these two functions?

```
Function PUSH {R4,LR}          Function2
    ;stuff                         ;stuff
    POP  {R4,PC}                   BX LR
```

Q4. Write assembly code that pushes registers R1, R3, and R5 onto the stack.

Answer

```
PUSH {R1}
PUSH {R3}
PUSH {R5}
```

Q5. What are the addressing modes used in each of the following instructions?

LDR R0, [R1]	Indexed add
LDR R2, [R1, #4]	Indexed add with immediate offset
MOV R3, #100	Immediate value
PUSH {R0}	Push & pop register addressing mode
MOV R0, #1	Immed add
LDRB R0, [PC, #0x30]	Relative address
LDR R0, =1234567	Relatives address/immediate address

BL func pc relative mode

Q6. Write a complete ARM assembly program that calls the procedure func1 which in turn calls a procedure func2. The procedure func2 is defined in Q1 of Sheet 3.

```
Reset_Handler

    BL func1

DeadLoop B DeadLoop

func1   PUSH {LR}

        BL func2

        POP {LR}

        BX LR
```