RSA Chosen Ciphertext Attack

RSA CCA

- ➤ Eve intercepts Ciphertext C sent from Alice to Bob as Eve is interested in the corresponding plaintext M.
 - We have $C = M^e \mod n$
- Now Eve **chooses** another ciphertext C' to send to Bob such that Bob decrypts it and sends the message back to Eve.
 - $C' = C \times r^e \mod n$
 - r is a random number chosen such that gcd(n,r)=1
 - Bob decrypts C' as $Y = (C')^d \mod n$
- Eve gets Y and then determine the plaintext M because:
 - Y= $(C \times r^e)^d \mod n$
 - Y= $(M^e \times r^e)^d \mod n$
 - Y= $M \times r \mod n$ Since ed are inverses
 - Y, r and n are known and required to know M
 - M= $Y \times r^{-1} \mod n$
 - Use Extended Euclidean algorithm to find r^{-1}

Extended Euclidean Algorithm

For given integers a and b, the extended Euclidean algorithm not only calculate the greatest common divisor but also two additional integers x and y that satisfy the following equation:

$$ax + by = d = \gcd(a, b)$$

- It should be clear that x and y will have opposite signs.
- ➤ Useful for later crypto computations
- \triangleright follow sequence of divisions for GCD but assume at each step i, can find x &y: r = ax + by at end find GCD value and also x & y
- More generally, the extended Euclidean algorithm can be used to find a multiplicative inverse in Z_n for any n. If we apply the extended Euclidean algorithm to the equation nx + by = d, and the algorithm yields d = 1, then $y = b^{-1}$ in Z_n .

Extended Euclidean Algorithm

Extended Euclidean Algorithm					
Calculate	Which satisfies	Calculate	Which satisfies		
$r_{-1} = a$		$x_{-1} = 1; y_{-1} = 0$	$a = ax_{-1} + by_{-1}$		
$r_0 = b$	1 1 1 1 1	$x_0 = 0; y_0 = 1$	$b = ax_0 + by_0$		
$r_1 = a \mod b$	$a = q_1 b + r_1$	$x_1 = x_{-1} - q_1 x_0 = 1$	$r_1 = ax_1 + by_1$		
$q_1 = \lfloor a/b \rfloor$		$y_1 = y_{-1} - q_1 y_0 = -q_1$			
$r_2 = b \mod r_1$	$b = q_2 r_1 + r_2$	$x_2 = x_0 - q_2 x_1$	$r_2 = ax_2 + by_2$		
$q_2 = \lfloor b/r_1 \rfloor$		$y_2 = y_0 - q_2 y_1$			
$r_3 = r_1 \bmod r_2$	$r_1 = q_3 r_2 + r_3$	$x_3 = x_1 - q_3 x_2$	$r_3 = ax_3 + by_3$		
$q_3 = \lfloor r_1/r_2 \rfloor$	100	$y_3 = y_1 - q_3 y_2$			
•	•	•	•		
•	•	•	•		
•	•	•	•		
$r_n = r_{n-2} \bmod r_{n-1}$	$r_{n-2} = q_n r_{n-1} + r_n$	$x_n = x_{n-2} - q_n x_{n-1}$	$r_n = ax_n + by_n$		
$q_n = \lfloor r_{n-2}/r_{n-1} \rfloor$		$y_n = y_{n-2} - q_n y_{n-1}$			
$r_{n+1} = r_{n-1} \bmod r_n = 0$	$r_{n-1} = q_{n+1}r_n + 0$	11.00	$d = \gcd(a, b) = r_n$		
$q_{n+1} = \lfloor r_{n-1}/r_n \rfloor$			$x = x_n; y = y_n$		

Extended Euclidean Algorithm: Example

a = 1759 b = 550

Table 4.4 Extended Euclidean Algorithm Example

i	ri	q_i	x_i	yi
-1	1759		1	0
0	550		0	1
1	109	3	1	-3
2	5	5	-5	16
3	4	21	106	-339
4	1	1	-111	355
5	0	4		

Result: d = 1; x = -111; y = 355

Thus, $b^{-1} = 355$.

Thank You