

Logarithmic Geometry and Hodge Theory

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X
 \downarrow smooth projective
 S

$S \xrightarrow{\text{period map}} \text{moduli of PHS}$

Ex. $\Delta^* = \{q \in \mathbb{C} : q \neq 0, |q| < 1\}$

$$E_q = \mathbb{C}^x / q\mathbb{Z}$$

E
 \downarrow family of elliptic curve
 Δ^*

$$\Delta^* \longrightarrow \mathrm{PSL}_2(\mathbb{Z}) \backslash \mathcal{H}$$

$$q \longmapsto \frac{1}{2\pi i} \log(q)$$

As $q \rightarrow 0$, E_q becomes a nodal curve.



Q. In what sense can a singular variety carry a PHS?

Can we extend period maps?

A. Kato-Utsui. Yes! Something called "logarithmic PHS".

"Log smoothness"

Ex. $xy = t$

$$\begin{array}{ccc} \mathbb{A}_{x,y}^2 & & (x,y) \\ \downarrow t & & \downarrow \\ \mathbb{A}_t^1 & & t = xy \end{array}$$

$$\Omega^1_f = \langle dx, dy : xdy + ydx = 0 \rangle$$

$$dt \mapsto d(xy) = xdy + ydx$$

At (0,0), relation is trivial!

Def'n If X is a variety, and D a reduced effective divisor, we can define

$$\Omega^1_X(\log D) = \text{"1-forms on } X \text{ w/ logarithmic poles along } D"$$

$$\Omega^1_{\mathbb{A}^1_0}(\log (t=0)) = \frac{dt}{t} \mathbb{C}[t]$$

$$\downarrow$$

$$d(\log t)$$

$$\Omega^1_{\mathbb{A}^2}(\log \begin{matrix} x=0 \\ y=0 \end{matrix}) = \langle \frac{dx}{x}, \frac{dy}{y} \rangle$$

$$\Omega^1_f(\log) = \langle \frac{dx}{x}, \frac{dy}{y} : \frac{dx}{x} + \frac{dy}{y} = 0 \rangle$$

$$\frac{dt}{t} \mapsto \frac{d(xy)}{xy} = \frac{dx}{x} + \frac{dy}{y}$$

Def'n A log structure on a scheme X is a sheaf of ^{comm.} monoids \mathcal{M}_X on $X_{\text{ét}}$ and a map $d: \mathcal{M}_X \rightarrow \mathcal{O}_X$ such that $d^{-1}(\mathcal{O}_X^\times) \xrightarrow{d} \mathcal{O}_X^\times$ is an isom. of sheaves of monoids.

Monoid: A set M w/ an assoc., unital binary operation.

Ex $(\mathbb{N}, +)$, (k^\times, \cdot) , (k, \cdot)

If M is a monoid, there's a group M^{gp} assoc. to it. Eg. $(\mathbb{N}, +) \rightsquigarrow (\mathbb{Z}, +)$

Ex \mathbb{A}^1_t w/ log pole along $t=0$.

$$\mathcal{M}(U) = \{f \in \mathcal{O}_X(U) : f \text{ only vanishes along } t=0\}$$

Ex. If X a var. and D a reduced effective divisor,

$$\mathcal{M}(U) = \{ f \in \mathcal{O}_X(U) : f \text{ only vanishes along } D \}$$

Def'n A morphism of log schemes $(X, \mathcal{M}_X) \rightarrow (Y, \mathcal{M}_Y)$

= the data of a map $f: X \rightarrow Y$ of schemes, and a map $f^* \mathcal{M}_Y \rightarrow \mathcal{M}_X$

of sheaves of monoids, s.t.

$$f^* \mathcal{M}_Y \rightarrow \mathcal{M}_X$$

$$f^* d_Y \downarrow \quad // \quad \downarrow d_X$$

$$f^* \mathcal{O}_Y \rightarrow \mathcal{O}_X$$

Def'n If $f: (X, \mathcal{M}_X) \rightarrow (Y, \mathcal{M}_Y)$ a morphism of log schemes,

$$\Omega_f^1(\log) = (\Omega_f^1 \oplus (\mathcal{O}_X \otimes_{\mathbb{Z}} \mathcal{M}_X^{\text{gp}})) / \sim$$

$$f \otimes s \approx f d \log(s)$$

$$\sim \left[\begin{array}{l} d(t) \otimes t = dt \\ 1 \otimes s = 0 \quad \text{for } s \in \text{im}(f^* \mathcal{M}_Y \rightarrow \mathcal{M}_X) \end{array} \right]$$

Log smoothness

A morphism $f: (X, \mathcal{M}_X) \rightarrow (Y, \mathcal{M}_Y)$ of log schemes

is log smooth if

1) f is loc. of finite pres.

2) X, Y are fine

3) if $T' \hookrightarrow T$ is a square zero ext'n, N a fine log str. on T ,

and $N' = i^{-1} N$, then in any diagram

$$(T', N') \rightarrow (X, \mathcal{M}_X)$$

$$\downarrow \quad \quad \quad \downarrow$$

$$(T, N) \rightarrow (Y, \mathcal{M}_Y)$$

étale loc. on T ,
he can solve this
lifting problem.

Def A monoid M is
integral if $M \rightarrow M^{\text{gp}}$
is injective.

Def A log scheme (X, \mathcal{M}_X)
is fine if \mathcal{M}_X is étale locally
integral & $\mathcal{M}_X / \mathcal{O}_X^{\times}$ is
finitely gen.
as a monoid.

Thus every nodal curve X/h can be given a log structure M_X so that

$$(X, M_X) \rightarrow (\mathrm{Spec} k, k^\times) \text{ is log smooth.}$$

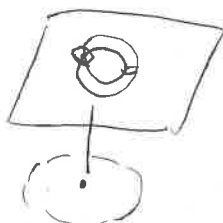
Moreover, any log smooth curve over $(\mathrm{Spec} k, k^\times)$ is nodal.

Def'n let X be a log \mathbb{C} -analytic space.

$$X^{\log} = \left\{ (x, h) : x \in X, h: M_{X,x} \rightarrow S^1 \text{ s.t. } h(\alpha^{-1}(f)) = \frac{f(x)}{|f(x)|} \right\}$$

Example

$$\begin{array}{c} E \\ \downarrow \\ \Delta \end{array}$$



$$\begin{array}{c} E^{\log} \\ \downarrow \\ \Delta^{\log} = S^1 \times [0, 1) \end{array}$$

Let X be a log \mathbb{C} -analytic space. A log VPHS on X is

- a local system $\mathcal{H}_{\mathbb{Z}}$ on X^{\log}

- $\mathcal{Q} = \mathcal{H}_{\mathbb{C}} \otimes \mathcal{H}_{\mathbb{C}} \rightarrow \mathcal{O}_1$

- filtration F^\bullet on $\mathcal{H}_{\mathbb{Z}} \otimes_{\mathbb{Z}} \mathcal{O}_{X^{\log}}$

obeying

1) Griffiths transversality

2) $F^\bullet \mathcal{H}_0$ is pulled back from X

extends $\mathrm{ev}: \mathcal{O}_{X,x} \rightarrow \mathbb{C}$

3) if $x \in X$, and $y = (x, h) \in X^{\log}$, and $s: \mathcal{O}_{X^{\log}, y} \rightarrow \mathbb{C}$

then if $\left| \exp(s(\log f_i)) \right| \ll 1$, then $(\mathcal{H}_{\mathbb{Z}, y}, \mathcal{Q}_y, F^\bullet(s))$ is a PHS

\uparrow
 f_i gen. $M_{X,x} / \mathcal{O}_{X,x}^\times$

$\mathcal{O}_{X,y}^{\otimes \log} (F_y^\bullet)$