Da categories of Chasi wherent sheaces

Dima Arinkin

Q Coh (x) as a 13h cut.

Last time: compartly generated DG cuts.

Det. AcBegenrates B if A = o, and then Ind (A) as B is an equil and then Bc= (Apro-tr) lear

Conversely, if B is compartly generated, any  $A \in B^c$  with  $B^c = (A^{pre-t})^{lcon}$  generates B.

Re (classically) X = schane, Qloh(X) (= Dqc (X1)

Det Qish(x) (ubelian) is defined by gluing.

Sq.(X) is its desired ext.

Better. Consider full subject of D(Ox-mod) with 9. who cohomology.

- QWh(x) & down that enough projecties.

- Are injecties in QGh(X) D'injectiv as sheares?

A) a Dhent.

Plan. Define (Chh(Y) directly by gluing

Det Q(bh(X) = lin D(k) (norks for any prestants X)

for scheme: Spor R C X

it x is classian, R's are chassical

Ruch The lim is over us - cost.

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uinhij = Spec Rij

3) Filuinuinuk ~> Fil...

$$\Sigma g. X = U \cup V$$
,  $Q Coh(X) = Q Coh(U) \times Q Coh(V)$   
 $Q Coh(X) = Q Coh(U) \times Q Coh(V)$ 

Ex. compute Hom's in this language.

## Cordlery. (Q6h(x)) = Perf(X)

Det. F (QCh(X) is perfect if flaper R is perfect for all spec R C) X

(Suffixes to check for an open cover)

[IMPORTANT]. Is QGh(x) compartly generated?

Thum (Thomason - Trobaugh)

Yes, it X is q.r. sep. scheme.

Penf(x) generates QCh(x)

Proof. X= UU U

"portere" affire

D .

O take generators Fd & Pout (4) and extend (how?) to Fac Pert (x).

@ For should generate Pert (X) "modulo" shears supported on Y= X \ U C V

Take enough perfect objects ap in QGh (V) = [ f a QGh (V): Flunu=0] a Con W

to generate it. (it V=V(+1,...,+k) CV, can take C=000 +1 0v)

Extend +. as toPert (x) s.r. as | u = 2

## Operations on Oh Categories

Q Coh (x) = lim D(R) SpecR→x

X geoment separated.

Then 1) Pert  $(x) = Q \cosh(x)^c$ 

lin Pert (R)
2) (Thomason - Trobaugh) Pert (x) generates Q (oh Cr).

Key step in proof

UCX: needed to extend

UVV from Pert (U) to Perf (X). (the actual TT Howsen)

More precisely, evens F F Perf (U) is a direct summer of F on form F & Perf (x

Working with compactly gan. Du cots

 $\mathbb{C}$   $\mathbb{B}$  = Ind(A) =  $\mathbb{D}(A^{*p})$ conp. gen.

RMA B has TI (and all limits) (computed printwise on 4).

B = Funct (AP, D(k))

Prop (Brown Rep) \$ : B'P -> D(k) is representable it it sends \$ to TT.

Proof it) p is do tormined by of AP. QED.

Coollary It B<sub>3</sub> is c. gen., B<sub>2</sub> is cocomplete, then any continuous functor  $F: B_1 \to B_2$  has a right adjoint.

 $\frac{\beta_{ret}}{\beta_{ret}}$   $\forall x \in B_2$ . Hom  $(F(-), x) : B_1^{op} \longrightarrow D(k)$  is representable by  $F^{R}(x)$ . Q(G).

Rule 1.  $F^R$  is continuous ity  $F(B_1^C) \subset B_2^C$ ,

This  $F(B_1^C) \subset B_2^C$ ,

The second to  $B_1^C \cap B_2^C$  is convenient.

2. Assuming this, F(Bi) generates BC iff FR is conservative.

[ F (x) = 0 = > x = 0]

Proposition (DG cats, Q-cats)

B = DG (at

U full subject

A

TFAE: 1) i: A - B admits a right adjoint.

2) 7 full subject. e CB s.t. 1/1/41e (c-> Hom (1, e)=0)

(6) any 6+B fits into a C

27677

Parx

1) The & in 267 is unique Properties

3) For iR, 1d->iR, i is an ison (because i i, fully faithful)

$$1) \xrightarrow{i^R} B/e . \qquad (also, e \sim B/A)$$

Q (Luotient)

A ZBZ e is a short exact sequence

A is a coloculization of 13]

(A)

Which is why (()) senitree ~ D())

Now suppose B is wwwpate, A is generated by ACCBC.

Then ACB has a right adjoint, S. e=AI gives

A -> B --> e .

Example. I and all of functors presone D.

Application. 0=0 (0) (AC)1=0 (=) B=Ind(A)

Suppose B is optly gen. as well.

Then i) B° -> (e)c (V. adjoint is wate)

2) B° genrates 2 (r.adjoint is conversative)

e = Ind(ec)

ec = ((Im Bc) pre tr) kan = Pert (p(Bc) P)
page 5

$$Q(\omega_{1}(x)) = D(R)$$

$$Q(\omega_{1}(x)) = Q(\omega_{1}(x))$$

$$Q(\omega_{1}(x)) = Q(\omega_{1}(x))$$

$$Q(\omega_{1}(\omega_{2}(x)))$$

$$Q(\omega_{1}(x)) = Q(\omega_{1}(x))$$

$$Q(\omega_{1}(x)) =$$

## Fourier - Mulear transform in Dh cats

Last time austions.

Large world A = 1 B = 1 P

Small world A C -> B - 7 e c

TI The newsures sum

՛.

Recodl\_ Small Dh (ats. Ob(A&B) = Ob(A) × Ob(B)

Hom ANB (a, Nb1, az Nb2) = Hom (a1, az) N Hom (b1, b2)

Has a universal property.

Ruh. For Kuroubian pre-O dy categories, take ((A &B) pre-tr) kan.

Key property.

Ind (A) & Ind (B) Went Ind (A & B)

Large Oh cuts

In general, A&B is defined by universal property using continuous bilinear functors.

functors.

Define ABB as Ind (ABB) / some objects

Cone (\$\phi(a\_{188}b) \rightarrow (\phi a\_i) \otimes b)

Cone (\$\phi(a\_{188}b) \rightarrow (\phi a\_i) \otimes b)

Cone (\$\phi(a\_{188}b) \rightarrow (\phi a\_i) \otimes b)

Dual category.

Small world . A (-) A ">

 $\frac{(\text{Ley property})}{(\text{Ind}(A))^{V}} = \text{Ind}(A^{9P}) = D(A) = \text{funct}(A, D(k))$   $= \text{Funct}(\text{Ind}(A), D(k))^{\text{Cont}}.$ 

Betar (In general). W of () h (ats) complete \_\_\_\_\_ Dual (alegry may on Many 1st exist.

All compartly generated cats one dualizable.

Geometrically:

X = 9 cpt, separated scheme.

Reson (Q (wh (x) = Ind (Pert(x))). Q (wh (x) = Ind (Pert(x) op)

Perf  $(x)^{\circ p} \longrightarrow \text{Perf}(x)$   $\text{Hom}(-, 0_X) \text{ (ordinary clust)}$ 

Ceng. X is affine, QGh(X) = D(R),  $QGh(X)^{V} = D(R^{op})$ .

Therefore. QCh(x) = QCh(x); corresponds to QCh(x)  $\otimes$  QCh(x)  $\rightarrow$  D(k)

Page 7 Perf(x)  $\otimes$  Perf(x)  $\sim$ 

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X, Y y cpt. separated
 QCoh (x) & QCoh (Y) = Ind (penf(x) & penf (Y))
     Quoh (xx') = Ind (Perf (xxy)) FBG.
     answer
 Check that (Part (x) & Part (4)) CQCoh (xxx) banishes.
Corollary. Funct (eigh(x), alcohy)) cont. = ( LL(x) & Q(h(Y) = Q(h(xxy)
          Funct (D(Part (X)OP), D(Part (Y)OP))cont.
             D ( part (x) 00 Part (4) °P)
               D(put (xxy) = Q6h (xxy)
          Qk = Bx (px (-) 8 k)
        Idech(x) & Funct (QGh(x), QGh(x)) cont.
            \beta_{\Lambda} \in QGh(X \times X).
     Ext (Idech(x), Idech(x)) = Ext xxx (Oa,80)
     Hochschild whomseogy.
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