

Algoritmo DFS, Depth-First Search (búsqueda en profundidad).
Búsqueda de caminos en profundidad.

Algorithm 2.3: Graph depth-first search with a stack.

StackDFS($G, node$) $\rightarrow visited$

Input: $G = (V, E)$, a graph

$node$, the starting vertex in G

Output: $visited$, an array of size $|V|$ such that $visited[i]$ is TRUE if we have visited node i , FALSE otherwise

```
1   $S \leftarrow \text{CreateStack}()$ 
2   $visited \leftarrow \text{CreateArray}(|V|)$ 
3  for  $i \leftarrow 0$  to  $|V|$  do
4       $visited[i] \leftarrow \text{FALSE}$ 
5   $\text{Push}(S, node)$ 
6  while not  $\text{IsStackEmpty}(S)$  do
7       $c \leftarrow \text{Pop}(s)$ 
8       $visited[c] \leftarrow \text{TRUE}$ 
9      foreach  $v$  in  $\text{AdjacencyList}(G, c)$  do
10         if not  $visited[v]$  then
11              $\text{Push}(S, v)$ 
12 return  $visited$ 
```

Algoritmo BFS, Breadth-First Search (búsqueda en amplitud).
Camino más cortos en grafos con aristas sin peso

Algorithm 2.5: Graph breadth-first search.

$\text{BFS}(G, \text{node}) \rightarrow \text{visited}$

Input: $G = (V, E)$, a graph

node , the starting vertex in G

Output: visited , an array of size $|V|$ such that $\text{visited}[i]$ is TRUE if we have visited node i , FALSE otherwise

```
1   $Q \leftarrow \text{CreateQueue}()$ 
2   $\text{visited} \leftarrow \text{CreateArray}(|V|)$ 
3   $\text{inqueue} \leftarrow \text{CreateArray}(|V|)$ 
4  for  $i \leftarrow 0$  to  $|V|$  do
5       $\text{visited}[i] \leftarrow \text{FALSE}$ 
6       $\text{inqueue}[i] \leftarrow \text{FALSE}$ 
7   $\text{Enqueue}(Q, \text{node})$ 
8   $\text{inqueue}[\text{node}] \leftarrow \text{TRUE}$ 
9  while not  $\text{IsQueueEmpty}(Q)$  do
10      $c \leftarrow \text{Dequeue}(Q)$ 
11      $\text{inqueue}[c] \leftarrow \text{FALSE}$ 
12      $\text{visited}[c] \leftarrow \text{TRUE}$ 
13     foreach  $v$  in  $\text{AdjacencyList}(G, c)$  do
14         if not  $\text{visited}[v]$  and not  $\text{inqueue}[v]$  then
15              $\text{Enqueue}(Q, v)$ 
16              $\text{inqueue}[v] \leftarrow \text{TRUE}$ 
17 return  $\text{visited}$ 
```

Algoritmo de Dijkstra.

Camino más cortos en grafo con aristas con pesos positivos.

Algorithm 7.1: Dijkstra's algorithm.

$\text{Dijkstra}(G, s) \rightarrow (pred, dist)$

Input: $G = (V, E)$, a graph

s , the starting node

Output: $pred$, an array of size $|V|$ such that $pred[i]$ is the predecessor of node i in the shortest path from s

$dist$, an array of size $|V|$ such that $dist[i]$ is the length of the shortest path calculated from node s to i

```
1   $pred \leftarrow \text{CreateArray}(|V|)$ 
2   $dist \leftarrow \text{CreateArray}(|V|)$ 
3   $pq \leftarrow \text{CreatePQ}()$ 
4  foreach  $v$  in  $V$  do
5       $pred[v] \leftarrow -1$ 
6      if  $v \neq s$  then
7           $dist[v] \leftarrow \infty$ 
8      else
9           $dist[v] \leftarrow 0$ 
10      $\text{InsertInPQ}(pq, v, dist[v])$ 
11 while  $\text{SizePQ}(pq) \neq 0$  do
12      $u \leftarrow \text{ExtractMinFromPQ}(pq)$ 
13     foreach  $v$  in  $\text{AdjacencyList}(G, u)$  do
14         if  $dist[v] > dist[u] + \text{Weight}(G, u, v)$  then
15              $dist[v] \leftarrow dist[u] + \text{Weight}(G, u, v)$ 
16              $pred[v] \leftarrow u$ 
17              $\text{UpdatePQ}(pq, v, dist[v])$ 
18 return  $(pred, dist)$ 
```

Algoritmo de Bellman-Ford.

Camino más cortos en grafo con aristas con pesos (puede haber pesos negativos).

Algorithm 8.1: Bellman-Ford.

$\text{BellmanFord}(G, s) \rightarrow (pred, dist)$

Input: $G = (V, E)$, a graph

s , the starting node

Output: $pred$, an array of size $|V|$ such that $pred[i]$ is the predecessor of node i in the shortest path from s

$dist$, an array of size $|V|$ such that $dist[i]$ is the length of the shortest path calculated from node s to i

```
1   $pred \leftarrow \text{CreateArray}(|V|)$ 
2   $dist \leftarrow \text{CreateArray}(|V|)$ 
3  foreach  $v$  in  $V$  do
4       $pred[v] \leftarrow -1$ 
5      if  $v \neq s$  then
6           $dist[v] \leftarrow \infty$ 
7      else
8           $dist[v] \leftarrow 0$ 
9  for  $i \leftarrow 0$  to  $|V|$  do
10     foreach  $(u, v)$  in  $E$  do
11         if  $dist[v] > dist[u] + \text{Weight}(G, u, v)$  then
12              $dist[v] \leftarrow dist[u] + \text{Weight}(G, u, v)$ 
13              $pred[v] \leftarrow u$ 
14  return  $(pred, dist)$ 
```
