

```

113 sequence longest_decreasing_powerSet(const sequence& A) {
114     const size_t n = A.size(); // 2 times
115     sequence best; // 1 time
116     std::vector<size_t> stack(n+1, 0); // n times
117     size_t k = 0; // 1 time
118     bool isDecreasing = true; // 1 time
119     while (true) // # loops 2^n
120     {
121         isDecreasing = true;
122         if (stack[k] < n) { // 1 + max(3, 3)
123             stack[k+1] = stack[k] + 1; // 4 times
124             ++k;
125         } else {
126             stack[k-1]++;
127             k--;
128         }
129         if (k == 0) { // 1 + max(3, 0) = 1 time
130             break;
131         }
132         sequence candidate; // 1 time
133         for (size_t i = 1; i <= k; ++i) // 2n times
134         {
135             candidate.push_back(A[stack[i]-1]);
136         }
137         // TODO
138         // write the if statement to test whether candidate determines
139         // a decreasing sequence AND has a size larger than the size
140         // of the current best
141         // if both conditions are satisfied, then stored candidate in best
142         // It is decreasing unless the if statement entered
143         for (size_t i = 1; i < candidate.size(); ++i) { // n times
144             if (candidate[i-1] < candidate[i])
145             {
146                 isDecreasing = false;
147                 break;
148             }
149         }
150         if (isDecreasing && candidate.size() > best.size())
151         {
152             best = candidate;
153         }
154     }
155     return best;
156 }
157 // ----- 6n+6 times

```

generates all power sets
 $|K|: 0, 1, 2, \dots, 2^n$
 each candidate generated in to a sequence

$6 + 2n + n + 2$
 $= 3n + 8$
 $5 + n + 2^n(3n + 8)$
 $S.C. = 3n \cdot 2^n + 8 \cdot 2^n + n + 5$ time units $\in O(n2^n)$

```

65 sequence longest_decreasing_end_to_beginning(const sequence& A) {
66
67     const size_t n = A.size(); 1 time  $1+n$ 
68
69     // populate the array H with 0 values
70     std::vector<size_t> H(n, 0); 1 time  $1+n$ 
71
72     // calculate the values of array H
73     // note that i has to be declared signed, to avoid an infinite loop, since
74     // the loop condition is i >= 0
75     for (signed int i = n-2; i >= 0; i--)  $\sum_{i=0}^{n-2} \sum_{j=i+1}^n 5 = \sum_{i=0}^{n-2} \sum_{j=1}^{n-i} 5 = \sum_{i=0}^{n-2} 5(n-i) = \sum_{i=0}^{n-2} 5n - 5i = 5n(n-1) - 5 \frac{(n-2)(n-1)}{2}$ 
76     {
77         for (size_t j = i+1; j < n; j++)  $= 5n^2 - 25n + 30$ 
78         {
79             if (A[i] > A[j] && H[i] <= H[j])  $3 + \max(2, 0) = 5 \text{ times}$ 
80             {
81                 H[i] = H[j] + 1;
82             }
83         }
84     }  $5n^2 + 5n + 30$ 
85
86     // calculate in max the length of the longest subsequence
87     // by adding 1 to the maximum value in H  $n+2 \text{ times}$ 
88     auto max = *std::max_element(H.begin(), H.end()) + 1; 1 time
89
90     // allocate space for the subsequence R
91     std::vector<int> R(max); 1 time
92
93     // add elements to R by whose H's values are in decreasing order.
94     // starting with max-1
95     // store in index the H values sought  $4 \text{ times}$ 
96
97     size_t index = max-1; 1 time  $2 \text{ times}$ 
98     size_t j = 0; 1 time
99     for (size_t i = 0; i < n; ++i) 1 time  $n \text{ times}$ 
100     {
101         if (H[i] == index) //  $1 + \max(5, 0) = 6 \text{ times}$ 
102         {
103             R[j] = A[i];
104             j++;
105             index--;
106         }
107     }  $6n$ 
108     return sequence(R.begin(), R.begin() + max); 1 time
109 }

```

$S.C. = 5n^2 + 13n + 37 \text{ times}$

$$\underline{h=3}$$

Stack

0	1	2	3
(0)	(1)	(2)	(3)

$$K=0 \quad \text{Stack}[0] \Rightarrow$$

\Rightarrow

0	1	0	0
---	---	---	---

$k=1$ $Stack[1] \Rightarrow$

\Rightarrow

0	1	2	0
---	---	---	---

[7]

$$k=2 \quad \text{stack}[2] =)$$

\Rightarrow

2	1	2	3
---	---	---	---

[3]

$$K=3, \text{ stack}[3] =)$$

\Rightarrow

0	1	3	3
---	---	---	---

[7]

$K=2$ Stack[2] \Rightarrow

d2	3	3
[1]		

$k=1$ Stack[1] = 2

0	2	4	5
---	---	---	---

2

$$k=2 \quad \text{Stack}[2] \Rightarrow$$

0	3	4	3
---	---	---	---

[1]

$$k=1 \text{ Stack}[5]$$

1	3	4	3
---	---	---	---

$k=0$ break!!

loops 8 times | $n=3$

$$2^3 = 8$$

$$z^n = K$$