

Bellman-Ford Algorithm

1. Definition

Bellman–Ford Algorithm is a single-source shortest path algorithm used to find the minimum distance from a starting node to all other nodes in a weighted graph, even when the graph contains negative edge weights.

Key Characteristics

- Works on **directed** and **undirected** graphs
- Handles **negative weights**
- Detects **negative weight cycles**
- Slower than Dijkstra's algorithm

2. Solution Process (Step-by-Step)

Step 1: Initialization

- Distance to **source** = 0
- Distance to **all other vertices** = ∞

$\text{dist}[\text{source}] = 0$

$\text{dist}[\text{others}] = \infty$

Step 2: Edge Relaxation (Repeat $V - 1$ Times)

For each iteration (1 to $V - 1$):

- Examine all edges ($u \rightarrow v$) with weight w
- Update the distance if a shorter path is found:

if $\text{dist}[u] + w < \text{dist}[v]$:

$\text{dist}[v] = \text{dist}[u] + w$

Reason:

A shortest path in a graph with V vertices must have at most **$V - 1$ edges**.

Step 3: Negative Cycle Check

Perform one more relaxation round:

- If any distance updates again → **Negative cycle exists**

if $\text{dist}[u] + w < \text{dist}[v]$ again:

Negative cycle detected

3. Pseudocode

```
BellmanFord(V, E, source):  
  
    dist = array(V, INF)  
    dist[source] = 0  
  
    for i = 1 to V-1:  
        for each (u, v, w) in edges:  
            if dist[u] + w < dist[v]:  
                dist[v] = dist[u] + w  
  
    for each (u, v, w) in edges:  
        if dist[u] + w < dist[v]:  
            print("Negative cycle")  
            return  
  
    return dist
```

4. Time Complexity

Operation	Complexity
Relaxing edges ($V - 1$ times)	$O(V \times E)$
Checking negative cycle	$O(E)$
Total Time Complexity	$O(VE)$
Space Complexity	$O(V)$