

浙江大学 2019 - 2020 学年秋冬学期
《计算机组成与设计》课程期末考试试卷

课程号: 67190020 , 开课学院: 信息与工程学院

考试试卷: ☒ A 卷、B 卷 (请在选定项上打 \checkmark)

考试形式: ☒ 闭、开卷 (请在选定项上打 \checkmark),

允许带 1 张 A4 大小的手写资料和计算器入场

考试日期: 2020 年 1 月 12 日, 考试时间: 120 分钟

诚信考试, 沉着应考, 杜绝违纪。

考生姓名: _____ 学号: _____ 所属院系 (专业): _____

题序	一	二	三	四	五	六	七	八	总分
得分									
评卷人									

I. CHOICE (Only one correct answer, 60 points)

1. Compute the average CPI (Clock cycle Per Instruction) of the following instructions. (_____)

Instruction type	CPI	Frequency
ALU	1	50%
Branch	2	20%
Load	2	20%
Store	2	10%

A. 1.0

B. 1.5

C. 2.0

D. 2.5

2. What is the content of stack pointer (SP)? (_____)

A. address of the current instruction

B. address of the next instruction

C. address of the top element of the stack

D. size of the stack

3. What is the function of the compiler? (_____)

A. Translates assembly language into binary instructions.

B. Translates source code into intermediates and immediately executes it.

C. Combines independent programs and resolves labels into an executable file.

D. Translate a high-level language into assembly language.

4. The bit used to signify that the cache location is updated is (_____).
- A. Dirty bit
 - B. Update bit
 - C. Reference bit
 - D. Flag bit
5. Interrupts can be generated in response to (_____).
- A. detected program errors such as arithmetic overflow or division by zero
 - B. detected hardware faults
 - C. input/output activities
 - D. All of the above.
6. Write Through technique is used in which memory for updating the data? (_____)
- A. Virtual memory
 - B. Main memory
 - C. Auxiliary memory
 - D. Cache memory
7. The execution of the following two instructions may have the (_____).
- ```
lw x5,0(x6)
lw x7,0(x5)
```
- A. RAW (Read after Write)
  - B. WAW (Write after Write)
  - C. WAR (Write after Read)
  - D. No data dependency.
8. Forwarding is a technique used in a pipeline to reduce the number of stall cycles caused by hazards. Each sequence of instructions shown below causes a hazard for the version of the RISC-V pipeline that we studied in class. The pipeline bubble that would be caused by some of these hazards can be avoided through the use of forwarding. Others cannot. Mark the sequences for which the bubble cannot be avoided through forwarding. (\_\_\_\_\_)
- A. lw x5, 0(x6)  
sll x9, x10, x11  
add x7, x8, x5
  - B. add x7, x8, x9  
beq x7, x0, L2
  - C. sub x5, x6, x7  
add x8, x9, x5
  - D. lw x5, 0(x6)  
add x7, x8, x5

9. The RISC-V addressing mode of “jal x1, 100” is (\_\_\_\_\_).
- A. Immediate addressing
  - B. Base addressing
  - C. PC-relative addressing
  - D. None of the above.
10. Which type of parallel computing architecture is no longer commonly encountered in machines today? (\_\_\_\_\_)
- A. MIMD (Multiple Instruction/Multiple Data Stream)
  - B. MISD (Multiple Instruction/Single Data Stream)
  - C. SIMD (Single Instruction/Multiple Data Stream)
  - D. SISD (Single Instruction/Single Data Stream)
11. Calculate AMAT (Average Memory Access Time) for a machine with the following specs: L1 hits take 3 cycles, L1 local miss rate is 25%. L2 hits take 10 cycles, L2 local miss rate is 60%. L3 hits take 100 cycles, L3 global miss rate is 9%. Main memory accesses take 1000 cycles and all data is available in memory. (\_\_\_\_\_)
- A. 105.5 cycles
  - B. 110.5 cycles
  - C. 24.5 cycles
  - D. None of the above.
12. Suppose you have a cache with capacity of  $2^{15}$  bytes, with 32-byte blocks. Assume the cache is two-way set associative. How many bits are used to select the set? (\_\_\_\_\_)
- A. 9 bits
  - B. 8 bits
  - C. 7 bits
  - D. None of the above.
13. You have a cache with 8B blocks, and a total size of 64B. The cache is two-way set associative. Use a least-recently used replacement policy. Given this sequence of accesses on word-addressed memory (4B), what is the hit rate? (\_\_\_\_\_)
- 0, 1, 2, 15, 17, 0, 32, 1
- A. 12.5%
  - B. 25.0%
  - C. 37.5%
  - D. None of the above.
14. For a computer with an 8-bit address space, a 64-byte, 2-way set-associative cache has 4B cache line size. How many bits are there in such cache's tag? (\_\_\_\_\_)
- A. 2
  - B. 3
  - C. 5
  - D. None of the above.

15. The drawback of building a large memory with DRAM (Dynamic Random Access Memory) is (\_\_\_\_\_).

- A. The large cost factor
- B. The inefficient memory organization
- C. The Slow speed of operation
- D. All of the mentioned

16. What is the decimal of the binary real number  $10.01 \times 2^{-1}$ ? (\_\_\_\_\_)

- A. 1.125
- B. 2.75
- C. 2.625
- D. 1.375

17. What is the RISC-V assembly code for the binary? (\_\_\_\_\_)

00100100011000111010110000100011

- A. sw t1, 600(t2)
- B. sw t1, 1200(t2)
- C. sw t2, 600(t1)
- D. sw t2, 1200(t1)

18. The effectiveness of the cache memory is based on the property of (\_\_\_\_\_).

- A. Locality of reference
- B. Memory localization
- C. Memory size
- D. None of the mentioned

19. An L2 cache, out of 100 total accesses to the cache system, missed 20 times. If L1 cache had a miss rate of 50%, what is the local miss rate of L2 cache? (\_\_\_\_\_)

- A. 20%
- B. 30%
- C. 35%
- D. 40%

20. Virtual memory uses a page table to track the mapping of virtual addresses to physical addresses. Assume 4 KiB pages, 32-bit virtual address. A cache designer wants to design a virtually indexed, physically tagged cache. To make an 8 KiB direct-mapped cache, assuming 4 words per block. What is the associativity? (\_\_\_\_\_)

- A. The cache is direct-mapped.
- B. The cache is two-way set associative.
- C. The cache is four-way set associative.
- D. None of the above.

## II. TRUE OR FALSE (10 points)

1. \_\_\_\_ Multiple levels of page tables will increase the total amount of page table storage.
2. \_\_\_\_ Static RAM is typically used to implement Cache.
3. \_\_\_\_ First-level caches are more concerned about miss rate, and second-level caches are more concerned about hit time.
4. \_\_\_\_ Two's complement in 8 bits for -128 is 1000 0000.
5. \_\_\_\_ When meeting cache misses, “no write allocate” means only writing to main memory.
6. \_\_\_\_ Pipelining a processor implementation probably will increase throughput.
7. \_\_\_\_ If hit rates are well below 0.9, then they're called as speedy computers.
8. \_\_\_\_ Assume two cache designs  $C_A$  and  $C_B$  have the same block size.  $C_A$  is a 32 KiB 2-way set associative cache and  $C_B$  is a 16 KiB direct-mapped cache. The length of the tag, measured in the number of bits, is the same in  $C_A$  and in  $C_B$ .
9. \_\_\_\_ Each stage in pipelining should be completed within 5 cycles.
10. \_\_\_\_ If a data cache does not contain a dirty bit, then it must be using a write-through policy.

## III. CACHE (8 points)

The contents of an associative cache are shown below. This cache has four-byte blocks and 32-bit addresses. Tags and sets are shown in binary, and data values in hex. The “...” should be replaced by an appropriate number of 0's.

| Set | V | Tag               | 0    | 1    | 2    | 3    |        |
|-----|---|-------------------|------|------|------|------|--------|
| 0   | 1 | 11110...011110000 | 0x44 | 0x33 | 0x22 | 0x11 | Line A |
|     | 1 | 10000...010110010 | 0xAA | 0xBB | 0xCC | 0xDD | Line B |
| 1   | 1 | 10000...011110000 | 0x01 | 0x02 | 0x03 | 0x04 | Line C |
|     | 1 | 10000...010000000 | 0x12 | 0xAB | 0x34 | 0xEF | Line D |

What are the addresses that are cached in each block? Write the range in hexadecimal using all eight digits, e.g. 0x00000000-0x000000AC.

Line A: \_\_\_\_\_

Line B: \_\_\_\_\_

Line C: \_\_\_\_\_

Line D: \_\_\_\_\_

#### IV. PIPELINE (12 points)

Data hazards occur due to data dependencies among instructions. Forwarding (aka bypassing) can solve many data hazards. We assume that the write is in the first half of the clock cycle and the read is in the second half.

1. Given the RV32I code below, spot the data dependencies in the code below. (2 points)

Figure out how forwarding can resolve data hazards. (2 points)

| Instruction    | C0 | C1 | C2 | C3  | C4  | C5  | C6 |
|----------------|----|----|----|-----|-----|-----|----|
| addi t0, s0,-1 | IF | ID | EX | MEM | WB  |     |    |
| and s2, t0, a0 |    | IF | ID | EX  | MEM | WB  |    |
| sw s0, 100(t0) |    |    | IF | ID  | EX  | MEM | WB |

2. Given the RV32I code below and a pipelined CPU with no forwarding, write out all hazards, using the following form: instruction 1-2. (4 points)

What types are each hazard (data hazard or control hazard)? (2 points)

How many stalls would there need to be in order to fix each data hazard(s)? (2 points)

|   | Instruction       | C0 | C1 | C2 | C3  | C4  | C5  | C6  | C7  | C8 |
|---|-------------------|----|----|----|-----|-----|-----|-----|-----|----|
| 1 | addi t0, s0,-1    | IF | ID | EX | MEM | WB  |     |     |     |    |
| 2 | and s2, t0, a0    |    | IF | ID | EX  | MEM | WB  |     |     |    |
| 3 | sw s2, 100(t0)    |    |    | IF | ID  | EX  | MEM | WB  |     |    |
| 4 | beq s0, s3, label |    |    |    | IF  | ID  | EX  | MEM | WB  |    |
| 5 | addi t2, x0, x0   |    |    |    |     | IF  | ID  | EX  | MEM | WB |

## V. VIRTUAL MEMORY (10 points)

Given the information below:

- 1 MiB of virtual memory
- 256 KiB of physical memory
- 4 KiB page size
- TLB: 8 entries, 2-way set associative

1. How many bits are needed to represent the virtual address space, physical address space and page offset respectively? (3 points)

2. How many bits are in the TLB index and TLB tag respectively? (2 points)

3. Translate virtual address 0x15213 to physical address, given the contents of the TLB and the first 32 entries of the page table below. (5 points)

TLB

| Index | Tag (Hex) | Physical page number (Hex) | Valid |
|-------|-----------|----------------------------|-------|
| 0     | 05        | 13                         | 1     |
|       | 3F        | 15                         | 1     |
| 1     | 10        | 0F                         | 1     |
|       | 0F        | 1E                         | 0     |
| 2     | 1F        | 01                         | 1     |
|       | 11        | 1F                         | 0     |
| 3     | 03        | 2B                         | 1     |
|       | 1D        | 23                         | 0     |

Page Table

| VPN(Hex) | PPN(Hex) | Valid | VPN(Hex) | PPN(Hex) | Valid |
|----------|----------|-------|----------|----------|-------|
| 00       | 17       | 1     | 10       | 26       | 0     |
| 01       | 28       | 1     | 11       | 17       | 0     |
| 02       | 14       | 1     | 12       | 0E       | 1     |
| 03       | 0B       | 0     | 13       | 10       | 1     |
| 04       | 26       | 0     | 14       | 13       | 1     |
| 05       | 13       | 0     | 15       | 18       | 1     |
| 06       | 0F       | 1     | 16       | 31       | 1     |
| 07       | 10       | 1     | 17       | 12       | 0     |
| 08       | 1C       | 0     | 18       | 23       | 1     |
| 09       | 25       | 1     | 19       | 04       | 0     |
| 0A       | 31       | 0     | 1A       | 0C       | 1     |
| 0B       | 16       | 1     | 1B       | 2B       | 0     |
| 0C       | 01       | 0     | 1C       | 1E       | 0     |
| 0D       | 15       | 0     | 1D       | 3E       | 1     |
| 0E       | 0C       | 0     | 1E       | 27       | 1     |
| 0F       | 2B       | 1     | 1F       | 15       | 1     |

VPN=\_\_\_\_\_

TLB Tag=\_\_\_\_\_, TLB miss or hit\_\_\_\_\_

Page table miss or hit\_\_\_\_\_

Physical Address=\_\_\_\_\_ (If you don't have enough information to fill the blank, fill the blank with "/")

### Instruction encoding for a reducing RISC-V ISA

| <i>name</i> | <i>format type</i> | <i>instruction format</i>    |            |            |            |                    |                |
|-------------|--------------------|------------------------------|------------|------------|------------|--------------------|----------------|
| <i>beq</i>  | <i>B-type</i>      | <i>imm[12/10:5]</i>          | <i>rs2</i> | <i>rs1</i> | <i>000</i> | <i>imm[4:1/11]</i> | <i>1100011</i> |
| <i>lw</i>   | <i>I-type</i>      | <i>imm[11:0]</i>             |            | <i>rs1</i> | <i>010</i> | <i>rd</i>          | <i>0000011</i> |
| <i>sw</i>   | <i>S-type</i>      | <i>imm[11:5]</i>             | <i>rs2</i> | <i>rs1</i> | <i>010</i> | <i>imm[4:0]</i>    | <i>0100011</i> |
| <i>addi</i> | <i>I-type</i>      | <i>imm[11:0]</i>             |            | <i>rs1</i> | <i>000</i> | <i>rd</i>          | <i>0010011</i> |
| <i>add</i>  | <i>R-type</i>      | <i>0000000</i>               | <i>rs2</i> | <i>rs1</i> | <i>000</i> | <i>rd</i>          | <i>0110011</i> |
| <i>sub</i>  | <i>R-type</i>      | <i>0100000</i>               | <i>rs2</i> | <i>rs1</i> | <i>000</i> | <i>rd</i>          | <i>0110011</i> |
| <i>sll</i>  | <i>R-type</i>      | <i>0000000</i>               | <i>rs2</i> | <i>rs1</i> | <i>001</i> | <i>rd</i>          | <i>0110011</i> |
| <i>srl</i>  | <i>R-type</i>      | <i>0000000</i>               | <i>rs2</i> | <i>rs1</i> | <i>101</i> | <i>rd</i>          | <i>0110011</i> |
| <i>or</i>   | <i>R-type</i>      | <i>0000000</i>               | <i>rs2</i> | <i>rs1</i> | <i>110</i> | <i>rd</i>          | <i>0110011</i> |
| <i>and</i>  | <i>R-type</i>      | <i>0000000</i>               | <i>rs2</i> | <i>rs1</i> | <i>111</i> | <i>rd</i>          | <i>0110011</i> |
| <i>jal</i>  | <i>J-type</i>      | <i>imm[20/10:1/11/19:12]</i> |            |            |            | <i>rd</i>          | <i>1101111</i> |