

Relational Algebra Operators

Relational Data Model

- Structure
 - TABLE or RELATION is the only element
- Value constraints
 - Unique or keys
 - NULLs
- Operations
 - Relational algebra or algebra for tables

TABLE Or Relation

- Schema or table header
 - Attributes or columns
 - Type or domain
 - Primitive: int, float, char[], string or varchar[]
 - Containers not allowed
- A table is seen as a collection (or multiset) of tuples
 - Cannot index in the table

Relational Algebra

- Set of operations or functions on tables
 - Input schema(s) \rightarrow Output schema
 - Input tuples \rightarrow Output tuples
- Single table operations
 - Select column, select tuple (row), aggregate, grouping
- Multiple table operations
 - Product and Join, Union, Intersection, Difference

Projection π

- Input table

- $T(A,B,C)$

- **A B C**

1 2 3

3 4 6

8 5 4

7 4 3

- $T' = \pi_{A, (A+B+C) \text{ AS } S'}(T)$

- Output table: T'

- Schema

- $T'(A,S')$

- Same number of tuples as T

- No duplicate elimination

- **A S'**

1 6

3 13

8 17

7 14

Selection σ

- Input table
 - T(A,B,C)

• **A B C**

1 2 3

3 4 6

8 5 4

7 4 3

- **$T' = \sigma_{A>1 \text{ AND } B+C>A}(T)$**
- Output table: T'
 - Schema
 - T'(A,B,C)
 - Same schema as T
 - Only tuples satisfying predicate

• **A B C**

3 4 6

8 5 4

Duplicate Elimination δ

- Input table
- $T' = \delta(T)$

T(A,B)

0 1

2 3

0 1

2 4

3 4

- Output table: T'

- Schema

- T'(A,B)

- Same schema as T

- Only distinct tuples

- At most the same number of tuples from T

T'(A,B)

0 1

2 3

2 4

3 4

Sorting τ

- Input table

$T(A,B)$

0 1

2 3

0 1

2 4

3 4

- $T' = \tau_{B [DESC]}(T)$

- Output table: T'

- Schema

- $T'(A,B)$

- Same schema as T

- Same tuples sorted

$T'(A,B)$

2 4

3 4

2 3

0 1

0 1

Aggregations

SUM, AVG, COUNT, MIN, MAX

- Input table

T(A,B)

0 1

2 3

0 1

2 4

3 4

- $T' = \text{SUM}_A(T)$

- $T'' = \text{MAX}_{A+B}(T)$

- Output table: T'

- Schema

- T'(X)

- Single tuple with aggregate result

T'(X)

7

T''(X)

7

GroupBy Aggregations γ

- Input table

T(A,B)

0 1

2 3

0 1

2 4

3 4

- $T' = \gamma_{A, \text{MIN}(B) \text{ AS } MB}(T)$

- Output table: T'

- Schema

- T'(A, MB)

- Arguments of γ

- Tuples have distinct values for A and group aggregate value for other attributes

T'(A,MB)

0 1

2 3

3 4

Set Operations \cup , \cap , $-$

- Input tables

R(A,B) S(A,B)

1 1 1 2

1 2 4 3

3 4

- Schema of R, S, and result table T' is the same (A,B)

- Union: $T' = R \cup S$

1 1

1 2

3 4

4 3

- Difference: $T' = R - S$

1 1

3 4

- Difference: $T' = S - R$

4 3

- Intersection: $T' = R \cap S$

1 2

Cartesian Product \times

- $R(A) = \{1,1,2,3\}$
- $S(B) = \{1,3,4\}$
- $T = R \times S(A,B) = \{$
 $(1,1),(1,3),(1,4),$
 $(1,1),(1,3),(1,4),$
 $(2,1),(2,3),(2,4),$
 $(3,1),(3,3),(3,4)\}$
- The result consists of pairs of one element from R and one from S
- Every element from R is paired with every element from S
- The number of elements in $R \times S$ is $|R| \times |S|$, i.e., the size of R multiplied by the size of S

- The schema of the result is the **union** of the R schema and the S schema
 - $R(A)$
 - $S(B)$
 - $T(A,B) = A \cup B$

Join ⋈

- $R(A) = \{1,1,2,3\}$
- $S(B) = \{1,3,4\}$
- $T = R \bowtie_{A=B} S = \{$
 $(\color{red}{1},\color{green}{1}),(\color{red}{1},\color{green}{3}),(\color{red}{1},\color{green}{4}),$
 $(\color{red}{1},\color{green}{1}),(\color{red}{1},\color{green}{3}),(\color{red}{1},\color{green}{4}),$
 $(\color{red}{2},\color{green}{1}),(\color{red}{2},\color{green}{3}),(\color{red}{2},\color{green}{4}),$
 $(\color{red}{3},\color{green}{1}),(\color{red}{3},\color{green}{3}),(\color{red}{3},\color{green}{4})\} = \{(\color{red}{1},\color{green}{1}),(\color{red}{1},\color{green}{1}),(\color{red}{3},\color{green}{3})\}$
- Join condition between attributes from the two tables
- Only those tuples from the Cartesian product that satisfy the join condition are included in the result
- The schema of the result is the **union** of the R schema and the S schema
 - $R(A)$
 - $S(B)$
 - $T(A,B) = A \cup B$
- $R \bowtie_{A=B} S = \sigma_{A=B}(R \times S)$

Outer Joins

$R(A,B)$	$S(B,C)$	$R \bowtie S$ [natural join] (A,B,C)	$R \bowtie_o S$ [full outer join] (A,B,C)
0 1	0 1		2 3 4
2 3	2 4		2 3 4
0 1	2 5	2 3 4	0 1 -
2 4	3 4	2 3 4	0 1 -
3 4	0 2		2 4 -
	3 4		3 4 -
			- 0 1
			- 2 4
			- 2 5
			- 0 2

Left (Right) Outer Joins

$R \bowtie_o S$ [full outer join]

$R(A,B)$ $S(B,C)$

0	1	0	1	2	3	4
2	3	2	4	2	3	4
0	1	2	5	0	1	-
2	4	3	4	2	4	-
3	4	0	2	3	4	-
		-	0	-	0	1
		-	2	-	2	4
		-	3	-	2	5
		-	4	-	0	2

$R \bowtie_L S$

[left outer join]

(A,B,C)

2	3	4
2	3	4
0	1	-
2	3	4
0	1	-
0	1	-
2	4	-
3	4	-

$R \bowtie_R S$

[right outer join]

(A,B,C)

2	3	4
2	3	4
-	0	1
-	2	4
-	2	5
-	0	2

Relational Algebra \leftrightarrow SQL

- SELECT \leftrightarrow Projection π
- FROM \leftrightarrow Input tables
- WHERE \leftrightarrow Selection σ , Join predicates
- DISTINCT \leftrightarrow Duplicate elimination δ
- ORDER BY \leftrightarrow Sorting τ
- GROUP BY \leftrightarrow GroupBy aggregations γ
- UNION, INTERSECT, EXCEPT \leftrightarrow Set operations $\cup, \cap, -$
- JOIN \leftrightarrow Join