# CS 332/532 Systems Programming

Lecture 30

Thread / 2

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## **Agenda**

- Thread
- pthread
- data sharing

## pthread3.c

```
#include <stdio.h>
#include <stdlib.h>
#include <pthread.h>
typedef struct foo {
   pthread_t ptid; /* thread id returned by pthread_create */
   int tid;
               /* user managed thread id (0 through nthreads-1) */
   int nthreads;
                    /* total no. of threads created */
} F00;
jvoid *compute(void *args) {
  F00 * info = (F00 *)args;
  printf("Hello, I am thread %d of %d\n", info->tid, info->nthreads);
  return (NULL);
int main(int argc, char **argv) {
  int i, nthreads;
  F00 *info;
  if (argc != 2) {
    printf("Usage: %s <# of threads>\n",argv[0]);
    exit(-1);
```

```
nthreads = atoi(argv[1]); // no. of threads
// allocate structure
info = (F00 *)malloc(sizeof(F00)*nthreads);
// create threads
for ( i = 0; i < nthreads; i++) {</pre>
  info[i].tid = i;
  info[i].nthreads = nthreads;
  pthread_create(&info[i].ptid, NULL, compute, (void *)&info[i]);
// wait for them to complete
for ( i = 0; i < nthreads; i++)</pre>
  pthread_join(info[i].ptid, NULL);
free(info);
printf("Exiting main program\n");
return 0;
```

```
(base) mahmutunan@MacBook-Pro lecture31 % gcc pthread3.c -o exercise6 -lpthread (base) mahmutunan@MacBook-Pro lecture31 % ./exercise6 4
Hello, I am thread 1 of 4
Hello, I am thread 0 of 4
Hello, I am thread 2 of 4
Hello, I am thread 3 of 4
Exiting main program
```

#### Thread Synchronization using Mutexes

- We can use the mutexes provided by the Pthreads library to control access to critical sections of the program and provide synchronization across the threads.
- We will use the example of computing the sum of the elements in a vector to illustrate the use of mutexes.
- We have a vector of N elements, we would like to assign each thread to compute the partial sum of N/P elements (where P is the number of threads), then we will update the shared global variable sum with the partial sums using mutex locks.

 The API for the mutex lock and unlock functions are shown below:

```
#include <pthread.h>
int pthread_mutex_lock(pthread_mutex_t *mutex);
int pthread_mutex_unlock(pthread_mutex_t
*mutex);
```

#### pthread\_sum.c

- The mutex variable is of type pthread\_mutex\_t and can be initially statically by assigning the value PTHREAD\_MUTEX\_INITIALIZER.
- Note that the mutex variable must be declared in global scope since it will be shared among multiple threads.
- A mutex can also be initialized dynamically using the function pthread\_mutex\_init.
- The pthread\_mutex\_destroy function can be used to destroy the mutex that was initialized using pthread\_mutex\_init.

```
#include <stdio.h>
#include <stdlib.h>
#include <pthread.h>
#include <unistd.h>
pthread_mutex_t mutex=PTHREAD_MUTEX_INITIALIZER;
double *a=NULL, sum=0.0;
int
       N, size;
void *compute(void *arg) {
    int myStart, myEnd, myN, i;
    long tid = (long)arg;
    // determine start and end of computation for the current thread
    myN = N/size;
    myStart = tid*myN;
    myEnd = myStart + myN;
    if (tid == (size-1)) myEnd = N;
    // compute partial sum
    double mysum = 0.0;
    for (i=myStart; i<myEnd; i++)</pre>
      mysum += a[i];
    // grab the lock, update global sum, and release lock
    pthread_mutex_lock(&mutex);
    sum += mysum;
    pthread_mutex_unlock(&mutex);
    return (NULL);
```

```
jint main(int argc, char **argv) -{
           long i;
           pthread_t *tid;
           if (argc != 3) {
              printf("Usage: %s <# of elements> <# of threads>\n",argv[0]);
              exit(-1);
           }
           N = atoi(argv[1]); // no. of elements
           size = atoi(argv[2]); // no. of threads
           // allocate vector and initialize
           tid = (pthread_t *)malloc(sizeof(pthread_t)*size);
           a = (double *)malloc(sizeof(double)*N);
           for (i=0; i<N; i++)
             a[i] = (double)(i + 1);
           // create threads
           for ( i = 0; i < size; i++)
             pthread_create(&tid[i], NULL, compute, (void *)i);
           // wait for them to complete
           for ( i = 0; i < size; i++)
             pthread_join(tid[i], NULL);
           printf("The total is %g, it should be equal to %g\n",
                  sum, ((double)N*(N+1))/2);
           return 0;
       Н
64
```

#### pthread\_sum.c

#### compile & run

```
(base) mahmutunan@MacBook-Pro lecture32 % gcc pthread_sum.c -o exercise7 -lpthread (base) mahmutunan@MacBook-Pro lecture32 % ./exercise7 200 4
The total is 20100, it should be equal to 20100
(base) mahmutunan@MacBook-Pro lecture32 % ./exercise7 1000 4
The total is 500500, it should be equal to 500500
(base) mahmutunan@MacBook-Pro lecture32 % _
```

#### pthread\_sum2.c

```
#include <stdio.h>
#include <stdlib.h>
#include <pthread.h>
#include <unistd.h>
 double *a=NULL, *partialsum;
int
        N, nthreads;
⊨void *compute(void *arg) {
     int myStart, myEnd, myN, i;
     long tid = (long)arg;
     // determine start and end of computation for the current thread
     myN = N/nthreads;
     myStart = tid*myN;
     myEnd = myStart + myN;
     if (tid == (nthreads-1)) myEnd = N;
     // compute partial sum
     double mysum = 0.0;
     for (i=myStart; i<myEnd; i++)</pre>
       mysum += a[i];
     partialsum[tid] = mysum;
     return (NULL);
```

#### pthread\_sum2.c

```
⊨int main(int argc, char **argv) {
     long i;
     pthread_t *tid;
     double sum = 0.0;
     if (argc != 3) {
        printf("Usage: %s <# of elements> <# of threads>\n",argv[0]);
        exit(-1);
     N = atoi(argv[1]); // no. of elements
     nthreads = atoi(argv[2]); // no. of threads
     // allocate vector and initialize
     tid = (pthread_t *)malloc(sizeof(pthread_t)*nthreads);
     a = (double *)malloc(sizeof(double)*N);
     partialsum = (double *)malloc(sizeof(double)*nthreads);
     for (i=0; i<N; i++)
       a[i] = (double)(i + 1);
     // create threads
     for ( i = 0; i < nthreads; i++)</pre>
       pthread_create(&tid[i], NULL, compute, (void *)i);
```

#### pthread\_sum2.c

```
// wait for them to complete
for ( i = 0; i < nthreads; i++)</pre>
  pthread_join(tid[i], NULL);
for ( i = 0; i < nthreads; i++)</pre>
  sum += partialsum[i];
printf("The total is %g, it should be equal to %g\n",
       sum, ((double)N*(N+1))/2);
free(tid);
free(a);
free(partialsum);
return 0;
```

```
(base) mahmutunan@MacBook-Pro lecture32 % gcc pthread_sum2.c -o exercise8 -lpthread (base) mahmutunan@MacBook-Pro lecture32 % ./exercise8 1000 4
The total is 500500, it should be equal to 500500
(base) mahmutunan@MacBook-Pro lecture32 % ./exercise8 200 4
The total is 20100, it should be equal to 20100
```

```
pthread_t tid[2];
int counter;
void* trythis(void* arg)
    unsigned long i = 0;
    counter += 1;
    printf("\n Job %d has started\n", counter);
    for (i = 0; i < (0xFFFFFFFF); i++)</pre>
    printf("\n Job %d has finished\n", counter);
    return NULL;
1}
int main(void)
    int i = 0;
    int error;
    while (i < 2) {
        error = pthread_create(&(tid[i]), NULL, &trythis, NULL);
        if (error != 0)
            printf("\nThread can't be created : [%s]", strerror(error));
        i++;
    pthread_join(tid[0], NULL);
    pthread_join(tid[1], NULL);
    return 0;
```

```
(base) mahmutunan@MacBook-Pro lecture32 % gcc exercise2.c -o exercise2 -lpthread
(base) mahmutunan@MacBook-Pro lecture32 % ./exercise2

Job 1 has started

Job 2 has started

Job 2 has finished

Job 2 has finished
```

```
#include <pthread.h>
#include <stdio.h>
#include <string.h>
pthread_t tid[2];
int counter;
pthread_mutex_t lock;
void* trythis(void* arg)
{
    pthread_mutex_lock(&lock);
    unsigned long i = 0;
    counter += 1;
    printf("\n Job %d has started\n", counter);
    for (i = 0; i < (0xFFFFFFFF); i++)</pre>
    printf("\n Job %d has finished\n", counter);
    pthread_mutex_unlock(&lock);
    return NULL;
```

```
jint main(void)
{
    int i = 0;
    int error;
    if (pthread_mutex_init(&lock, NULL) != 0) {
         printf("\n mutex init has failed\n");
         return 1;
     }
    while (i < 2) {
         error = pthread_create(&(tid[i]),
                                 NULL,
                                &trythis, NULL);
         if (error != 0)
             printf("\nThread can't be created :[%s]",
                    strerror(error));
         i++;
     pthread_join(tid[0], NULL);
     pthread_join(tid[1], NULL);
     pthread_mutex_destroy(&lock);
    return 0;
```

```
[(base) mahmutunan@MacBook-Pro lecture32 % gcc exercise3.c -o exercise3 -lpthread
[(base) mahmutunan@MacBook-Pro lecture32 % ./exercise3

Job 1 has started

Job 1 has finished

Job 2 has started

Job 2 has finished
```

Thread 1	Thread 2	Balance
Read balance: \$1000		\$1000
	Read balance: \$1000	\$1000
	Deposit \$200	\$1000
Deposit \$200		\$1000
Update balance \$1000+\$200		\$1200
	Update balance \$1000+\$200	\$1200

- The fundamental principle is this: Two or more processes can cooperate by means of simple signals, such that a process can be forced to stop at a specified place until it has received a specific signal.
- Any complex coordination requirement can be satisfied by the appropriate structure of signals.
- For signaling, special variables called semaphores are used.
- To transmit a signal via semaphore s, a process executes the primitive semSignal(s).
- To receive a signal via semaphore s, a process executes the primitive semWait(s); if the corresponding signal has not yet been transmitted, the process is suspended until the transmission takes place.

- While mutexes are one solution to ensure synchronization, semaphores provide an alternative and more generalized approach to establish synchronization among multiple threads.
- While mutexes provide a locking mechanism (both lock and unlock are performed by a single thread - cooperative locks), semaphores provide a signaling mechanism (two different threads cooperate with the wait/signal calls) to synchronize access to shared resources.
- Note that it is possible to implement a mutex using a binary semaphore.

A variable that has an integer value upon which only three operations are defined:

 There is no way to inspect or manipulate semaphores other than these three operations

- 1) A semaphore may be initialized to a nonnegative integer value
- 2) The semWait operation decrements the semaphore value
- 3) The semSignal operation increments the semaphore value

- The Pthread library provides the wait and signal calls that are described in the text book. However, since Linux uses the word signal for software interrupts, the wait and signal calls are called sem\_wait and sem\_post.
- The C APIs for sem\_wait and sem\_post are:

```
#include <semaphore.h>
int sem_wait(sem_t *sem);
int sem_post(sem_t *sem);
```

- We have to include the header file semaphore.h to compile and link with the pthread library using lpthread.
- The *sem\_wait* call decrements the semaphore variable *sem*.
- The *sem\_wait* call returns immediately if the value of *sem* is greater than zero after decrementing the semaphore.
- If the current value of the semaphore is zero, then sem\_wait call will block until the semaphore values goes above zero or it is interrupted by a signal handler.
- Similarly, the sem\_post call increments the semaphore variable sem and if the value of sem goes above zero, it will then wake up the corresponding thread or process that was blocked in a sem\_wait call.

- The sem\_init function must be used to initialize a semaphore and the sem\_destroy method to destroy a semaphore.
- The C API for the sem\_init and sem\_destroy are given below:

```
#include <semaphore.h>
int sem_init(sem_t *sem, int pshared, unsigned
int value);
int sem destroy(sem t *sem);
```

 The sem\_init call will assign a semaphore sem with the initial value provided and the second parameter pshared specifies whether the semaphore is shared between threads in with a process or it is shared between processes. • If the value of *pshared* is zero, then the semaphore is shared between threads and the semaphore must be declared in a shared address space so that all threads can access it.

• If the value of *pshared* is nonzero, then the semaphore will be shared between processes and the semaphore must be declared in a shared memory region

#### **Producer/Consumer Problem**

# General Statement:

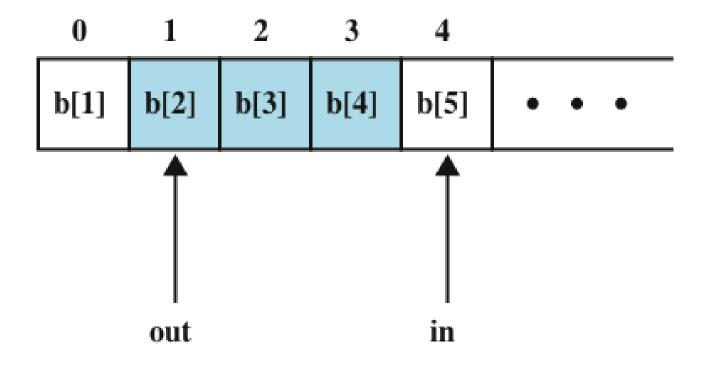
One or more producers are generating data and placing these in a buffer

A single consumer is taking items out of the buffer one at a time

Only one producer or consumer may access the buffer at any one time

# The Problem:

Ensure that the producer won't try to add data into the buffer if its full, and that the consumer won't try to remove data from an empty buffer



Note: shaded area indicates portion of buffer that is occupied

#### Figure 5.11 Infinite Buffer for the Producer/Consumer Problem

- We will implement the bounded-buffer producer-consumer problem using semaphores here. In this exercise we will consider the case of a single producer and single consumer and use threads to create a producer and a consumer.
- We use the pseudo code from the textbook (Figure 5.16 on page 228) and replace semWait and semSignal with sem\_wait and sem\_post.
- This code is based on the examples provided in the classic book - UNIX Networking Programming, Volume 2 by W. Richard Stevens

#### prodcons1.c

```
/* Solution to the single Producer/Consumer problem using semaphores.
  This example uses a circular buffer to put and get the data
   (a bounded-buffer).
   Source: UNIX Network Programming, Volume 2 by W. Richard Stevens
  To compile: gcc -0 -Wall -o <filename> <filename>.c -lpthread
  To run: ./<filename> <#items>
  To enable printing add -DDEBUG to compile:
   gcc -0 -Wall -DDEBUG -o <filename> <filename>.c -lpthread
*/
/* include globals */
#include <stdio.h>
#include <stdlib.h>
#include <pthread.h> /* for POSIX threads */
#include <semaphore.h> /* for POSIX semaphores */
```

```
#define
                       10
        NBUFF
                                                   /* read-only */
int nitems;
struct {
                         /* data shared by producer and consumer */
 int
          buff[NBUFF];
          mutex, nempty, nstored; /* semaphores, not pointers */
 sem_t
} shared;
void *producer(void *), *consumer(void *);
/* end globals */
/* main program */
int main(int argc, char **argv)
 pthread_t tid_producer, tid_consumer;
 if (argc != 2) {
   printf("Usage: %s <#items>\n", argv[0]);
   exit(-1);
```

```
nitems = atoi(argv[1]);
42
         /* initialize three semaphores */
         sem_init(&shared.mutex, 0, 1);
         sem_init(&shared.nempty, 0, NBUFF);
         sem_init(&shared.nstored, 0, 0);
         /* create one producer thread and one consumer thread */
         pthread_create(&tid_producer, NULL, producer, NULL);
         pthread_create(&tid_consumer, NULL, consumer, NULL);
         /* wait for producer and consumer threads */
         pthread_join(tid_producer, NULL);
         pthread_join(tid_consumer, NULL);
         /* remove the semaphores */
         sem_destroy(&shared.mutex);
         sem_destroy(&shared.nempty);
         sem_destroy(&shared.nstored);
         return 0;
       /* end main */
```

```
/* producer function */
void *producer(void *arg)
 int i;
 for (i = 0; i < nitems; i++) {
   sem_wait(&shared.nempty);  /* wait for at least 1 empty slot */
   sem_wait(&shared.mutex);
   shared.buff[i % NBUFF] = i; /* store i into circular buffer */
#ifdef DEBUG
   printf("wrote %d to buffer at location %d\n", i, i % NBUFF);
#endif
   sem_post(&shared.mutex);
   return (NULL);
/* end producer */
```

```
/* consumer function */
void *consumer(void *arg)
 int i;
 for (i = 0; i < nitems; i++) {</pre>
   sem_wait(&shared.nstored);  /* wait for at least 1 stored item */
   sem_wait(&shared.mutex);
   if (shared.buff[i % NBUFF] != i)
     printf("error: buff[%d] = %d\n", i, shared.buff[i % NBUFF]);
#ifdef DEBUG
   printf("read %d from buffer at location %d\n",
           shared.buff[i % NBUFF], i % NBUFF);
#endif
   sem_post(&shared.mutex);
   sem_post(&shared.nempty); /* 1 more empty slot */
 return (NULL);
```

#### gcc -O -Wall -DDEBUG prodcons1.c -lpthread \$ ./a.out 20

wrote 0 to buffer at location 0

wrote 1 to buffer at location 1

wrote 2 to buffer at location 2

wrote 3 to buffer at location 3

wrote 4 to buffer at location 4

wrote 5 to buffer at location 5

wrote 6 to buffer at location 6

wrote 7 to buffer at location 7

wrote 8 to buffer at location 8

read 0 from buffer at location 0

read 1 from buffer at location 1

read 2 from buffer at location 2

read 3 from buffer at location 3

read 4 from buffer at location 4

read 5 from buffer at location 5

read 6 from buffer at location 6

read 7 from buffer at location 7

read 8 from buffer at location 8

wrote 9 to buffer at location 9

wrote 10 to buffer at location 0

wrote 11 to buffer at location 1

wrote 12 to buffer at location 2

wrote 13 to buffer at location 3

wrote 14 to buffer at location 4

wrote 15 to buffer at location 5

wrote 16 to buffer at location 6

wrote 17 to buffer at location 7

wrote 18 to buffer at location 8

read 9 from buffer at location 9

read 10 from buffer at location 0

read 11 from buffer at location 1

read 12 from buffer at location 2

read 13 from buffer at location 3

read 14 from buffer at location 4

read 15 from buffer at location 5

read 16 from buffer at location 6

read 17 from buffer at location 7

read 18 from buffer at location 8 wrote 19 to buffer at location 9

read 19 from buffer at location 9

- The source code is self-explanatory, we will focus on key sections of the code here.
- First, we define global variables such as the number of items (*nitems*) the producer will produce and the consumer will consume.
- Then we create a shared region, called shared, that will be shared between the producer and consumer threads.
- It contains the buffer shared by the producer and consumer and the three semaphores: one for the mutex lock (*mutex*), one for number of empty slots (*nempty*), and one for the number of slots filled (*nstored*) (these correspond to semaphores *s*, *e*, and *n*, respectively, from the textbook).

 In the main function, we read the number of items to be produced/consumed as a command-line argument and initialize the three semaphores using sem\_init as per the pseudocode from the textbook.

```
nitems = atoi(argv[1]);

/* initialize three semaphores */
sem_init(&shared.mutex, 0, 1);
sem_init(&shared.nempty, 0, NBUFF);
sem_init(&shared.nstored, 0, 0);
```

- Generalization of the semWait and semSignal primitives
  - No other process may access the semaphore until all operations have completed

#### Consists of:

- Current value of the semaphore
- Process ID of the last process to operate on the semaphore
- Number of processes waiting for the semaphore value to be greater than its current value
- Number of processes waiting for the semaphore value to be zero

Solaris provides classic counting semaphores with the following primitives:

- sema\_p() Decrements the semaphore, potentially blocking the thread
- sema\_v() Increments the semaphore, potentially unblocking a waiting thread
- sema\_tryp() Decrements the semaphore if blocking is not required

#### User level:

- Linux provides a semaphore interface corresponding to that in UNIX SVR4
- Internally:
  - Implemented as functions within the kernel and are more efficient than user-visable semaphores
- Three types of kernel semaphores:
  - Binary semaphores
  - Counting semaphores
  - Reader-writer semaphores

Traditional Semaphores		
<pre>void sema init(struct semaphore *sem, int count)</pre>	Initializes the dynamically created semaphore to the given count	
<pre>void init MUTEX(struct semaphore *sem)</pre>	Initializes the dynamically created semaphore with a count of 1 (initially unlocked)	
<pre>void init MUTEX LOCKED(struct semaphore *sem)</pre>	Initializes the dynamically created semaphore with a count of 0 (initially locked)	
void down(struct semaphore *sem)	Attempts to acquire the given semaphore, entering uninterruptible sleep if semaphore is unavailable	
<pre>int down interruptible(struct semaphore *sem)</pre>	Attempts to acquire the given semaphore, entering interruptible sleep if semaphore is unavailable; returns -EINTR value if a signal other than the result of an up operation is received	
<pre>int down trylock(struct semaphore  *sem)</pre>	Attempts to acquire the given semaphore, and returns a nonzero value if semaphore is unavailable	
void up(struct semaphore *sem)	Releases the given semaphore	
Reader-Writer Semaphores		
<pre>void init rwsem(struct rw_semaphore, *rwsem)</pre>	Initializes the dynamically created semaphore with a count of 1	
<pre>void down read(struct rw semaphore, *rwsem)</pre>	Down operation for readers	
<pre>void up read(struct rw semaphore, *rwsem)</pre>	Up operation for readers	
<pre>void down write(struct rw_semaphore, *rwsem)</pre>	Down operation for writers	
<pre>void up write(struct rw semaphore, *rwsem)</pre>	Up operation for writers	

**Table 6.5** 

#### Linux Semaphores

(Table can be found on page 293 in textbook)