COS341 Semester Project 2025 : Code-Generation from "SPL"

General assumptions: Scope-Analysis and consistent re-naming of user-defined names have already been carried out, and the Symbol Table has already been filled with the relevant information.

VARIABLES ::= VAR VARIABLES VAR ::= user-defined-name

Translation Advice:

Variable-Declarations do <u>not</u> get translated to target code. The Variable-Declarations were only needed for "filling" the Symbol-Table which the translator (code-generator) is going to consult.

PDEF ::= NAME (PARAM) { BODY }

Translation Advice:

The Sub-Tree "under" tree-node PDEF will later be used for "inlining", as discussed in the Lecture of Thursday the 2nd of October.

FDEF ::= NAME (PARAM) { BODY ; return ATOM }

Translation Advice:

The Sub-Tree "under" tree-node FDEF will later be used for "inlining", as discussed in the Lecture of Thursday the 2nd of October.

BODY ::= local { MAXTHREE } ALGO

PARAM ::= MAXTHREE

MAXTHREE ::=

MAXTHREE ::= VAR

MAXTHREE ::= VAR VAR

MAXTHREE ::= VAR VAR VAR

Translation Advice:

Variable-Declarations do <u>not</u> get translated to target code. The Variable-Declarations were only needed for "filling" the Symbol-Table which the translator (code-generator) is going to consult. **Only the ALGO will get translated**.

MAINPROG ::= var { VARIABLES } ALGO

Translation Advice:

Variable-Declarations do <u>not</u> get translated to target code. The Variable-Declarations were only needed for "filling" the Symbol-Table which the translator (code-generator) is going to consult. **Only the ALGO will get translated**.

ATOM ::= VAR

Translation Advice:

Similar to $\underline{\text{Trans}}(\text{Exp} \to \underline{\text{id}})$ in Fig.6.3 of our blue textbook, however with a normal = (instead of the textbook's :=) in generated target code.

ATOM ::= number

Translation Advice:

Similar to $\underline{\text{Trans}}(\text{Exp} \rightarrow \underline{\text{num}})$ in Fig.6.3 of our blue textbook, however with a normal = (instead of the textbook's :=) in generated target code.

ALGO ::= INSTR; ALGO

Translation Advice:

Similar to $\underline{\text{Trans}}(\text{Stat} \rightarrow \underline{\text{Stat1}}; \underline{\text{Stat2}})$ in Fig.6.5 of our blue textbook.

INSTR ::= halt

Translation Advice:

Trans(halt) is { target_code = " STOP "; return(target_code); }

INSTR ::= print OUTPUT

 $\begin{array}{lll} \text{OUTPUT} & := & \underline{\text{string}} \\ \text{OUTPUT} & ::= & ATOM \\ \text{ATOM} & ::= & \underline{\text{number}} \\ \text{ATOM} & ::= & VAR \\ \end{array}$

Translation Advice:

INSTR ::= ASSIGN ASSIGN ::= VAR = TERM

Translation Advice:

Similar to $\underline{\text{Trans}}(\text{Stat} \rightarrow \underline{\text{id}} := \underline{\text{Exp}})$ in Fig.6.5 of our blue textbook, however with a normal = (instead of the textbook's :=) in generated target code.

INSTR ::= BRANCH

BRANCH ::= if TERM { ALGO } else { ALGO }

Translation Advice:

- Our Target-Language does not have the keyword *ELSE* in its syntax !
- Our Target-Language does not have the keyword LABEL in its syntax!

Therefore, the **translation techniques from Fig.6.5 and Fig.6.6** of our textbook *must get modified* as follows:

- Instead of "LABEL" ++ label ++ we produce: "REM" ++ label ++ "\newline" ++
- Instead of "IF t1 op t2 THEN labelT ELSE labelF" we produce:

```
"IF t1 op t1 THEN labelT ++ \newline ++ code_of_the_else_ALGO ++ \newline ++ GOTO ++ labelExit ++ \newline ++ REM ++ labelT ++ \newline ++ code_of_the_then_ALGO ++ \newline ++ REM ++ labelExit ++ \newline ++ "
```

```
BRANCH ::= if TERM { ALGO }
```

Translation Advice:

For the same reasons as described above, we modify Fig.6.5 and Fig.6.6 to produce:

```
"IF t1 op t1 THEN <u>labelT</u> ++ \newline ++
GOTO ++ <u>labelExit</u> ++ \newline ++
REM ++ <u>labelT</u> ++ \newline ++
code_of_the_then_ALGO ++ \newline ++
REM ++ labelExit ++ \newline ++ "
```

INSTR ::= LOOP

LOOP ::= while TERM { ALGO }
LOOP ::= do { ALGO } until TERM

Translation Advice:

Similar to Figures 6.5 and 6.6 in the textbook though you must keep in mind (as mentioned above) that

- our Target-Language uses the keyword REM instead of the keyword LABEL
- our Target-Language does not have the keyword ELSE

INSTR ::= NAME (INPUT) // procedure call without return
ASSIGN ::= VAR = NAME (INPUT) // function call with return

Translation Advice:

The Expression NAME(INPUT) gets translated as in Textbook Fig.6.3 whereby a CALL command will emerge. (In a *later* phase, the CALL will get *replaced* by way of *inlining*.)

TERM ::= ATOM

TERM ::= (UNOP TERM)

TERM ::= (TERM BINOP TERM)

Translation Advice:

Very similar to Figure 6.3 in the textbook!

However, wherever the textbook generates := we simply generate = (without the colon).

UNOP ::= neg Trans(neg) is simply the minus symbol -

UNOP ::= not

Translation Advice:

As our target language does <u>not</u> have the Boolean negation symbol in its syntax, we must **translate** if (not(TERM)) as we translate if (TERM), however with **swapping** the then-Algo for the else-ALGO (and, vice versa, the else-ALGO for the then-ALGO).

BINOP ::= eq

Translation Advice:

<u>Trans(eq)</u> is simply the equality symbol = (which our target language "overloads" for both value assignments as well as also for the Boolean comparisons).

BINOP ::= > <u>Trans(>)</u> is simply itself: >

BINOP ::= or BINOP ::= and

Translation Advice:

As our target language does <u>not</u> have any Boolean operators in its syntax, we use the "cascading" translation technique from **Fig.6.8** of our textbook whereby, however, we must again keep in mind that

- our Target-Language does not have any keyword ELSE
- our Target-Language uses REM label \newline , instead of LABEL label

BINOP plus Trans(plus) is simply + ::= BINOP Trans(minus) is simply -::= minus BINOP mult Trans(mult) is simply * ::= BINOP div Trans(div) is simply / ::=

The generated target code must be stored in a *.txt (ASCII) output file

<u>IF in doubt</u> → then go to <u>Tutor's Consultation Hour</u> in the Orange Lab

And now: **HAPPY CODING:**)