

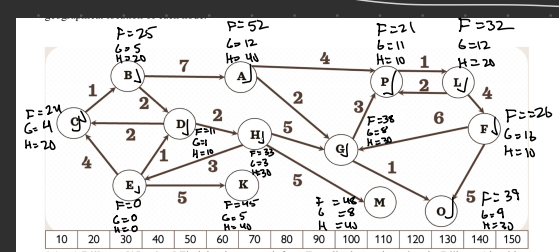


$C: E, D, C = 3$
 $D: E, D = 1$
 $H: E, D, H = 3$
 $B: E, D, C, B = 4$
 $K: E, K = 5$
 $M: E, D, H, M = 8$
 $G: E, D, H, G = 8$
 $O: E, D, H, G, O = 9$
 $A: E, D, C, B, A = 11$
 $P: E, D, H, G, P = 11$
 $L: E, D, H, G, P, L = 12$
 $F: E, D, H, G, P, L, F = 16$

Curr
 E
 D
 H
 C
 B
 E
 M
 G
 O
 A
 P
 L
 F

 Prev
 null
 E
 D
 C
 K
 H
 G
 B
 A
 P
 L

Visited: E, D, H, C, B, K, M, G, O, A, P, L, F cost



Visited: E, D, C, B, H, G, P, L, F, G, M, A

3. Both algorithms found H in 4 iterations therefore they are the same.

74. The time complexity of this code is $O(n^2)$ because we use 2 nested for loops, one to read the $n \times n$ matrix, and the other to check symmetry of matrix which is n^2 in worst case. So, $O(n^2) + O(n^2) = O(n^2)$ where n is the size of the matrix.

The space complexity for this code is n^2 where n is size of the matrix because the $n \times n$ matrix takes up $O(n^2)$ memory.

75. Time complexity is $O(n)$ where n is the number of vertices, because I use for loops for this code and none of them are nested so its all $O(n)$.

The space complexity of this code is also $O(n)$ where n is the # of vertices because I store $O(n)$ elements in the graph I am printing out.

76. The time complexity is $O(V+E)$ where V is the vertex, and E is the edge because in the worst case we go through every single vertex and edge in the graph using dfs.

The space complexity is also $O(V+E)$ where V is the vertex and E is the edge because we store each of the edges at least once which $= O(V+E)$.