

Chapter 7.1

The Disjoint Set ADT

Equivalence Class

1) Definition of Equivalence Class:

Suppose we have a set $U=\{1,2,\dots,n\}$ of n elements and a set $R=\{(i_1,j_1), (i_2,j_2),\dots, (i_r,j_r)\}$ of r relations. The relation R is an equivalence relation iff the following conditions are true(symbol ' \equiv ' represent the equivalence relation on sets, x,y,z are elements in set) :

- Reflexive $x \equiv x$.
- Symmetric $x \equiv y, y \equiv x$
- Transitive $x \equiv y$ and $y \equiv z$, then $x \equiv z$

Equivalence Class

例如:

判别 3 个数 a, b, c 能否构成三角形的三条边?

能构成三角形的等价类:

$\{ (3,4,5), (4,5,6), (5,6,7), \dots \}$

不能构成三角形的等价类:

$\{ (1,2,3), (2,3,5), \dots \}$

Equivalence Class

2) Example:

set $s = \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11\}$

pairs of equivalence:

$(0\ 4), (3\ 1), (6\ 10), (8\ 9), (7\ 4), (6\ 8), (3\ 5), (2\ 11),$
 $(11\ 0)$

Initial: $\{0\}, \{1\}, \{2\}, \{3\}, \{4\}, \{5\}, \{6\}, \{7\}, \{8\},$
 $\{9\}, \{10\}, \{11\}$

$0 \equiv 4$ $\{0, 4\}, \{1\}, \{2\}, \{3\}, \{5\}, \{6\}, \{7\}, \{8\}, \{9\},$
 $\{10\}, \{11\}$

$3 \equiv 1$ $\{0, 4\}, \{1, 3\}, \{2\}, \{5\}, \{6\}, \{7\}, \{8\}, \{9\},$
 $\{10\}, \{11\}$

Equivalence Class

$6 \equiv 10$ $\{0, 4\}, \{1, 3\}, \{2\}, \{5\}, \{6, 10\}, \{7\}, \{8\},$
 $\{9\}, \{11\}$

$8 \equiv 9$ $\{0, 4\}, \{1, 3\}, \{2\}, \{5\}, \{6, 10\}, \{7\},$
 $\{8, 9\}, \{11\}$

$7 \equiv 4$ $\{0, 4, 7\}, \{1, 3\}, \{2\}, \{5\}, \{6, 10\}, \{8, 9\}, \{11\}$

$6 \equiv 8$ $\{0, 4, 7\}, \{1, 3\}, \{2\}, \{5\}, \{6, 8, 9, 10\}, \{11\}$

$3 \equiv 5$ $\{0, 4, 7\}, \{1, 3, 5\}, \{2\}, \{6, 8, 9, 10\}, \{11\}$

$2 \equiv 11$ $\{0, 4, 7\}, \{1, 3, 5\}, \{2, 11\}, \{6, 8, 9, 10\}$

$11 \equiv 0$ $\{0, 4, 7, 2, 11\}, \{1, 3, 5\}, \{6, 8, 9, 10\}$

Equivalence Class

3) Online equivalence class operation

- **Combine(a,b) : combine the equivalence classes that contains elements a and b into a single class**
- **Find(e) : determine the class that currently contains element e.**

Combine(a,b) is equivalent to

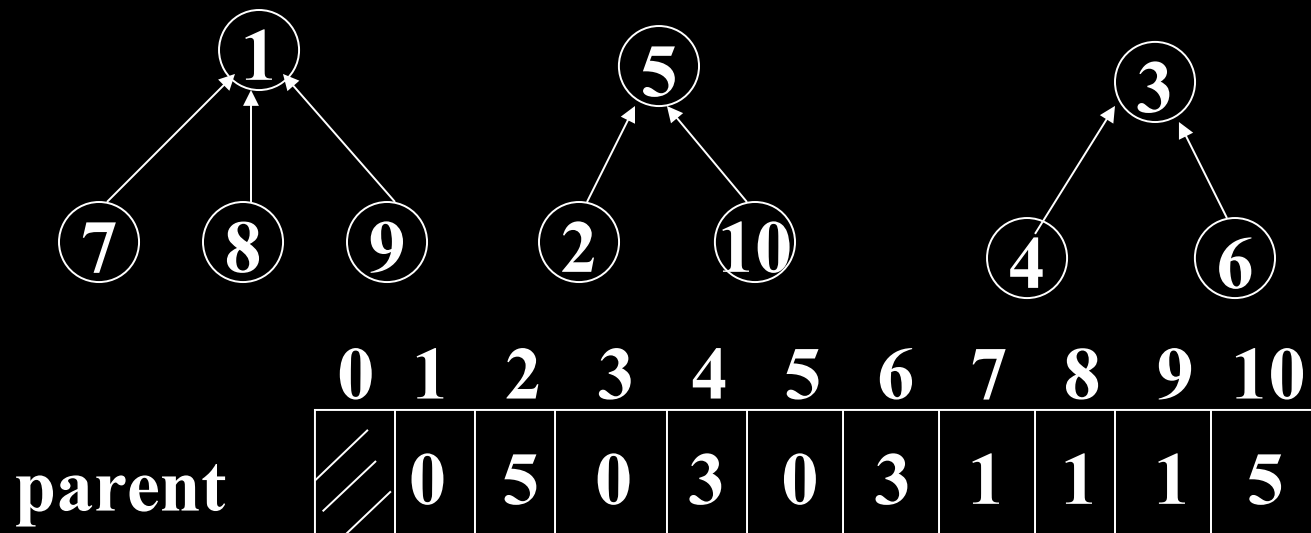
i=Find(a); j=Find(b); if(i!=j) Union(i,j);

Equivalence Class

4) Tree Representation(Union-Find sets)

Example:

$S_1 = \{1, 7, 8, 9\}$, $S_2 = \{5, 2, 10\}$, $S_3 = \{3, 4, 6\}$, they all belong to $S = \{1, 2, 3, \dots, 10\}$



Equivalence Class

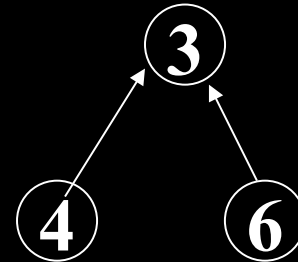
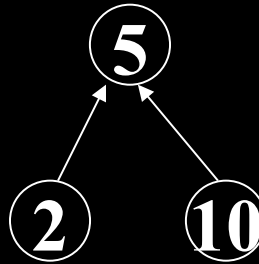
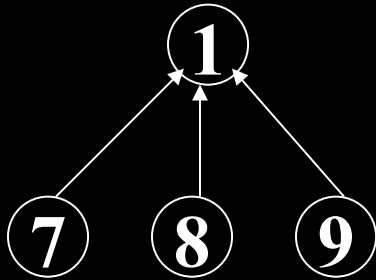
simple tree solution to union-find problem

```
void Initialize(int n)
{
    parent=new int[n+1];
    for(int e=1;e<=n;e++)
        parent[e]=0;
}
```

```
int Find(int e)
{
    while(parent[e])
        e=parent[e];
    return e;
}
```

```
void Union(int i, int j)
{
    parent[j]=i;
}
```

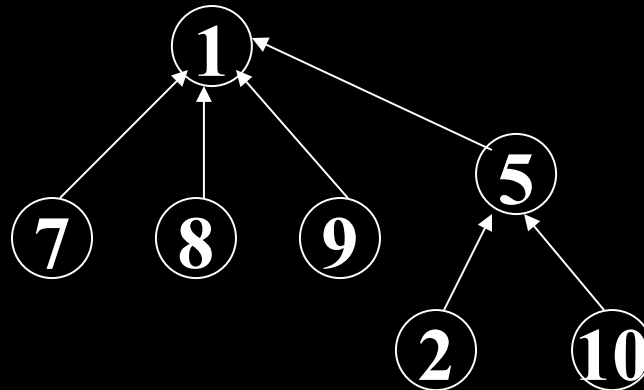

Equivalence Class



parent

0	1	2	3	4	5	6	7	8	9	10
	0	5	0	3	10	3	1	1	1	5

Union (1,5)



Equivalence Class

Java

```
public class DisjSets
{
    public DisjSets( int numElements )
        public void union( int root1, int root2 )
        public int find( int x )
        private int [ ] s;
}

public DisjSets( int numElements )
{
    s = new int [ numElements ];
    for( int i = 0; i < s.length; i++ )
        s[ i ] = -1; // 一个根结点
}
```

Equivalence Class

```
public void union( int root1, int root2 )
```

```
{  s[ root2 ] = root1;
```

```
}
```

```
public int find( int x )
```

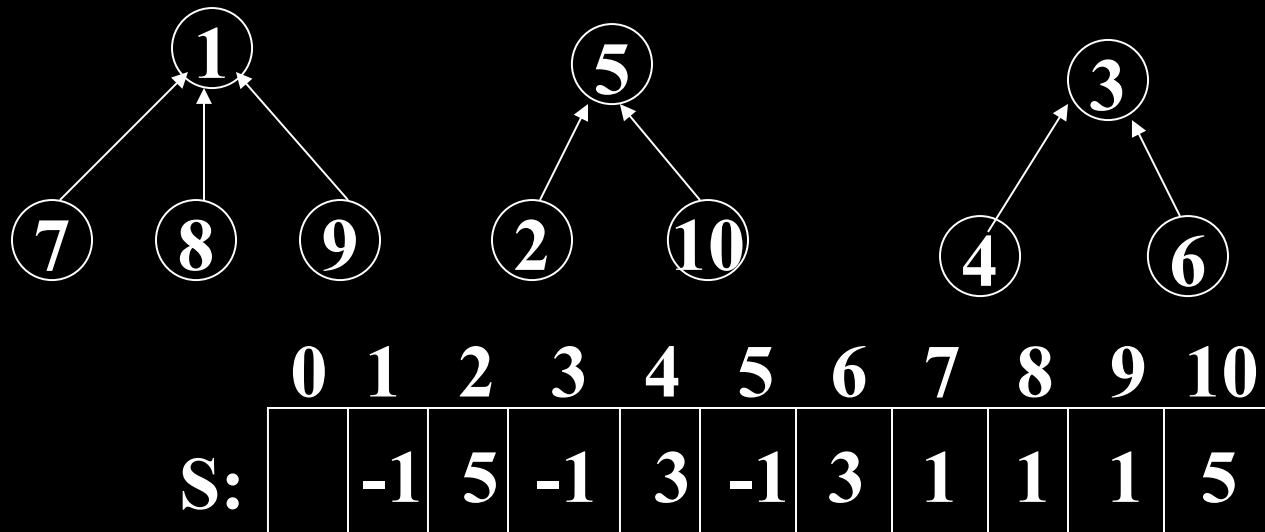
```
{  if( s[x] < 0 )
```

```
    return x;
```

```
    else
```

```
        return find( s[ x ] );
```

```
}
```



Equivalence Class

5) Performance Evaluation

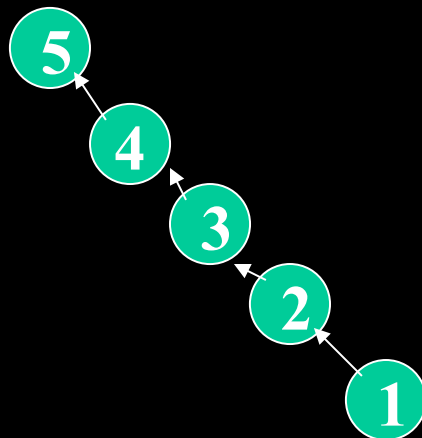
Time complexity: Find-- $O(h)$,

Union-- $\theta(1)$

Assume that u times unions and f times finds are to be performed, $f > u$,

in the worst case a tree with m elements can have a height of m :

Union(2,1), Union(3,2), Union(4,3), Union(5,4)...



Equivalence Class

①improve Union

two rules:

- **Weight rule:** if the number of nodes in tree i is less than the number in tree j , then make j the parent of i ; otherwise, make i the parent of j .
- **Height rule:** if the height of tree i is less than that of tree j , then make j the parent of i ; otherwise, make i the parent of j .

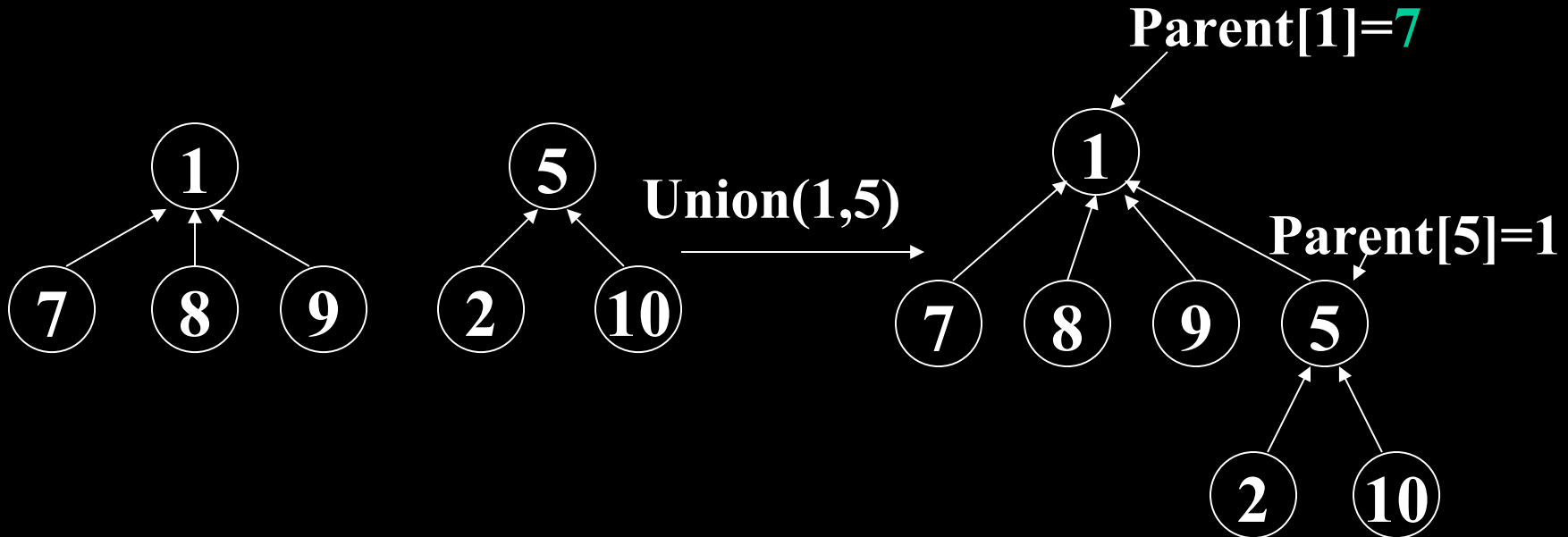
Equivalence Class

6) Performance Enhancement

- ①improve Union in order to decrease the time each find take, so that the height of tree will not increase linearly .
- ②.Improvement of *Find* –path compression

Equivalence Class

Let's discuss the weight rule(C++) :



Besides the *parent* field, each node has a boolean field *root*. The *root* field is true iff the node is presently a root node. The *parent* field of each root node is used to keep a count of the total number of nodes in the tree.

Equivalence Class

Union with the weight rule

```
void Initialize(int n)
{ root=new bool[n+1];
  parent=new int[n+1];
  for(int e=1;e<=n;e++)
  { parent[e]=1;
    root[e]=true;
  }
}

int Find(int e)
{ while(!root[e])
  e=parent[e];
  return e;
}
```


Equivalence Class

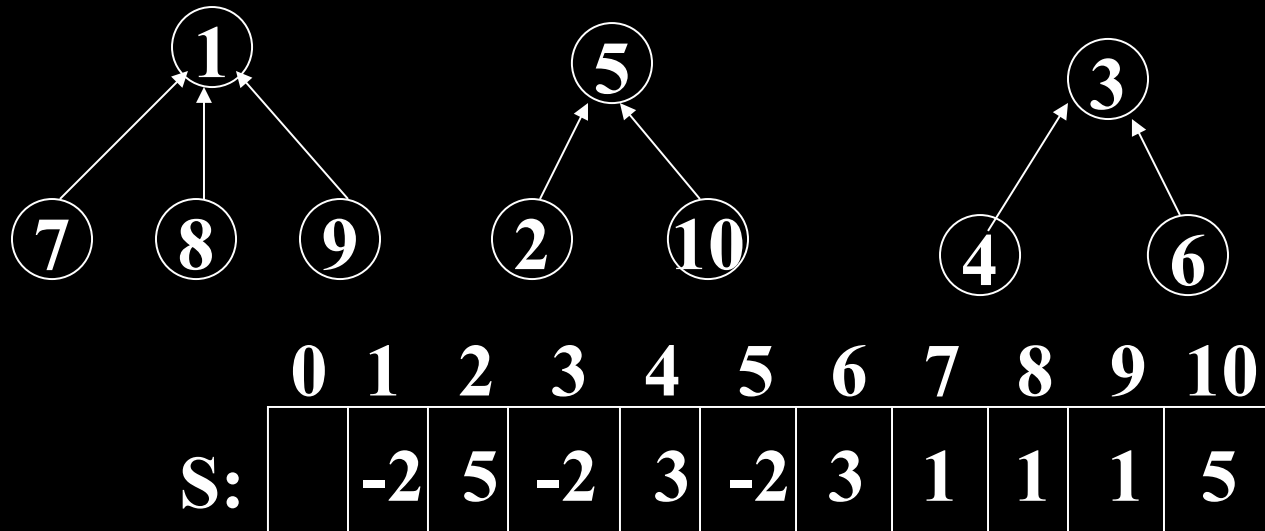
```
void Union(int i, int j)
{ if(parent[i]<parent[j]) // i becomes subtree of j
  { parent[j]=parent[j]+parent[i];
    root[i]=false;
    parent[i]=j;
  }
  else { parent[i]=parent[i]+parent[j];
        root[j]=false;
        parent[j]=i;
      }
}
```

Equivalence Class

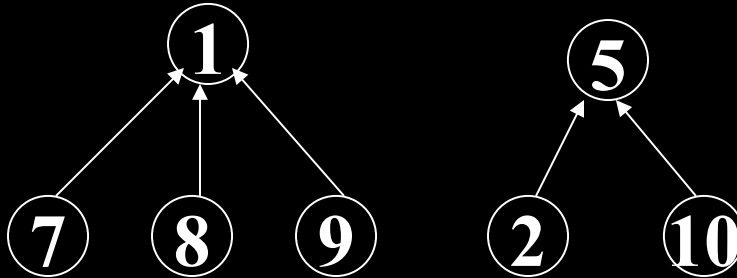
Java(高度规则)

用一个数组来实现，根结点中放负数，而且是代表高度。

```
public void union( int root1, int root2 )  
{  if( s[ root2 ] < s[ root1 ] )  
    s[ root1 ] = root2;  
  else {  if( s[ root1 ] == s[ root2 ] )  
    s[ root1 ]--;  
    s[ root2 ] = root1; } }
```



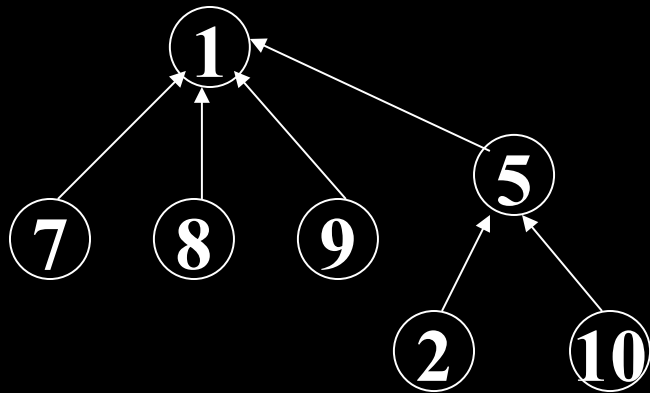
Equivalence Class



union (1, 5)

s[root1]--

s[root2] = root1

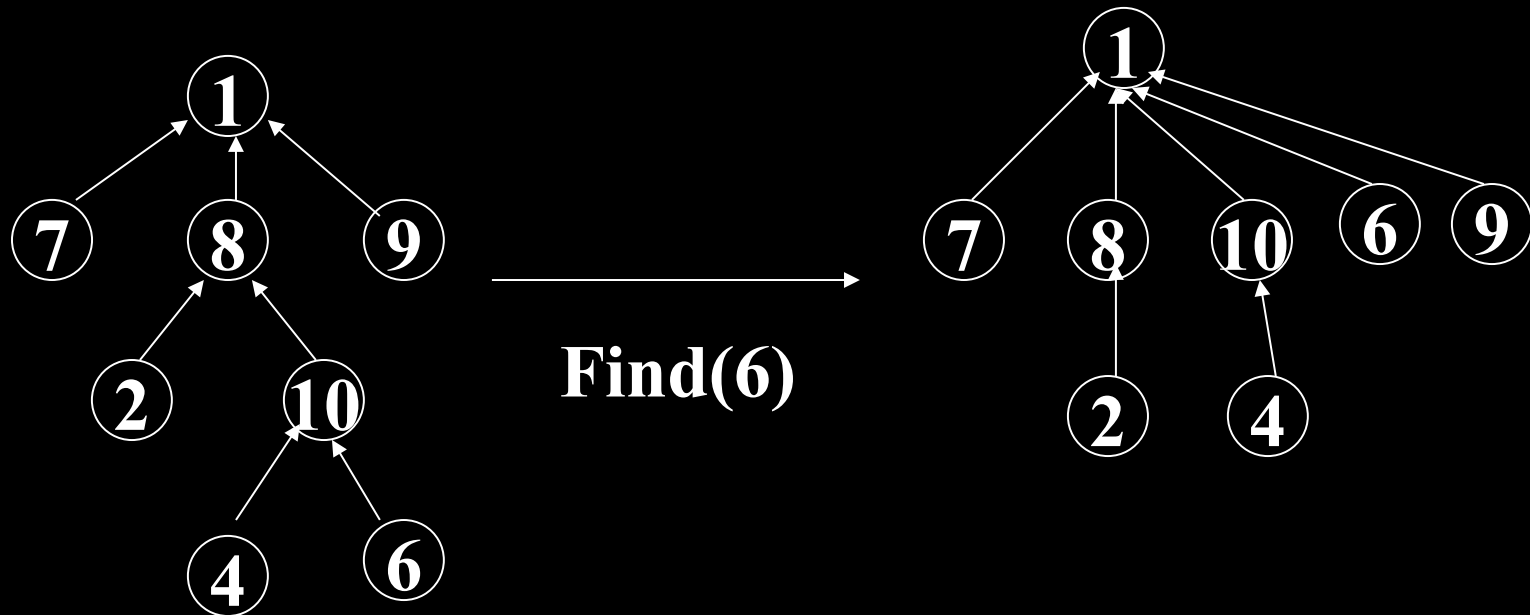


Equivalence Class

②.Improvement of *Find* –path compression

When processing a equivalence pair, we need to operate *Find* twice, *WeightUnion* once.

Example of improvement:



Equivalence Class

```
int Find( int e) { /* C++ */  
{ int j=e;  
  while(!root[j]) j=parent[j];  
  int f=e;  
  while(f!=j)  
    { int pf=parent[f];  parent[f]=j; f=pf; }  
}
```

Equivalence Class

Java

```
public int find( int x )  
{   if( s[ x ] < 0 )  
        return x;  
    else  
        return s[ x ] = find( s[ x ] );  
}
```