Last time - Sylow Meanens.

One more example, to show that things can get more complicated quitely.

Let's by to classify groups of order 12. If G1=12 then Sylon gives

- * a subgroup HCG, |H|=4; the number of these is $S_2 \in \{1,3\}$ $(S_2|3)$, $S_2 \equiv 1 \mod 2$
- a subgroup $K \subset G$, |K| = 3; the number is $s_3 \in \{1,4\}$ ($s_3 \mid 4$, $s_3 \equiv 1 \mod 3$)
- * At least one of these is normal: indeed, if s3=4 then the nontrivial elements of $k_1,...,k_k$ all have order 3, and $k_i \cap k_j = \{e\}$ (order divide 3, <3), so ce have 8 elements of order 3. So there are at most 4 elements of order & {1,2,4}, hence s=1 and H is normal.
- * If both H and K are normal them $G \cong H \times K$ (using |G| = |H|.|K|, $H \cap K = \{e\}$) and so G is abelian, one of $\mathbb{Z}/4 \times \mathbb{Z}/3 \cong \mathbb{Z}/12$ see last time (Z/2×Z/2)×Z/3 = Z/2×Z/6.
- 4 If H is normal but K isn't, consider the action of 6 on {k1, k2, k3, k4} by canjugation. Conjugation by a nonhival element of K, maps K, to trell but doesn't his any of the 3 other: inteed recall the stabilizer of K_i is $\{g \in G \mid gK_i g' = k_i\} = N(K_i)$, and by orbit stabilizer, $|N(k_i)| = \frac{|G|}{s_0} = \frac{12}{4} = 3$, so $N(k_i) = k_i$. So, a norbinal element of k_1 acts on {k1, k2, k3, k4} by a 3-cycle peneting {k2, k3, k4}, and similarly for others. Here the ation of G on { Kir. Kir} gives a homom. 4:G -> 54
 y: -> 3-4/les

This implies In(4) > A4, here = A4, and G=A4.

- + If K is normal but H isn't, then the are 2 subcars H= 7/4 or 7/2×7/2!
 - if H= Z/4, let x EH generator, let K= {e,y,y²}, then G= KxH is determined by the conjugation action of H on K, ie. need to know xyx' EK. Can't have xyx'=e (=) y=e) or xyx'=y (=) x and y commute, G≃H×K abellan).

So instead $xyx^{2} = y^{2} (=y^{2})$.

Then G is generally by x,y, with $x^{4} = y^{3} = e$ and $xy = y^{2}x$. This group is unfamiliar to us - semidical product 2/3 × 2/4, where 2/4 acts on the normal subgroup 2/3 by 2/4 -> Aut(2/3)= { +id}

-> if $H = \frac{7}{2} \times \frac{7}{2}$, then look at conjugation action $H = A \times (k) \approx \frac{7}{2}$, recens $k \cdot (q) = \frac{7}{2} \times 2$, denote by z its generator, $x \in H$ sto z, z greate H, y generator of k, then G is generator, $x \in H$ sto z, z greate H, y generator of k, then G is generator, $y \in H$ such $y \in \mathbb{R}^2 = 2^2 = y^3 = e$.

(An check this is actually G = De.

(The subgroup generator by y and z is $z \in \mathbb{R}^2 = 2^2 \times (2^2 \times 2^2)$.

Anomal in G, take $y = n \cdot t$ and $y \in \mathbb{R}^2 = 2^2 \times (2^2 \times 2^2)$. $x = n \cdot t$ any reflection.

Thus there are 5 isome days of grays of order 12: $(2/12, 2/2 \times 2/6, A_4, 2/3 \times 2/4, D_6)$.

Generators, proentations, and Cayley graph. Recall:

* The free group For on a generators a,, ..., an.

Elenes are all reduced words $a_{i1}^{m_1} ... a_{ik}^{m_k}$ $k \ge 0$ (enphy und is e)

(non reduced words: reduce by:

oif $i_j = i_{j+1}$, combine $a_i^m a_i^{m'} \rightarrow a_i^{m+m'}$ oif an exponent is zero, remove a_i^0).

Repeat until word is reduced.

- This is the "largest" group with a generators, all others are \simeq quotients of F_n . If G is generated by $g_1, g_n \in G$, define a homomorphism $\varphi: F_n \to G$ by retting $a_i \mapsto g_i$ (and so, $\pi_{a_{ij}} \mapsto \pi_{g_{ij}} = \pi_{g_{ij}}$)
- A finitely generated group is said to be finitely proceed if the kernel of φ is the smallest normal subgroup of F_n containing some finite subset $\{r_1,...,r_K\}\subset F_n$, (i.e. the subgroup generated by r_j 's and r_j words in the generators their conjugates $x^{-1}r_jx$).

Write $G \cong \langle a_{1,...}, a_{n} | r_{1,...}, r_{k} \rangle$, then $G \cong F_{n} / \langle canj's of r_{1...}, r_{k} \rangle$ generators relations.

 $\underline{\mathsf{Ex}}; \quad \mathbf{Z}^{\mathsf{n}} \cong \langle a_1,..,a_{\mathsf{n}} \mid a_i;a_j;a_j^{\mathsf{n}} \rangle \forall i,j \rangle.$

 $\underline{E_{x}}$, $S_{3} \cong \langle s_{1}, s_{2} | s_{1}^{2}, s_{2}^{2}, (s_{1}s_{2})^{3} \rangle$

• Concretely, given generators $g_1, ..., g_m \in G$, it is not had to find relations $f_1, ..., f_k$, i.e. words in the free group F_n st. under $f_i : F_n \to G$, $f_i \mapsto e$. If these relations hold in G, then $f_i : f_i : G_i : G_i$

among the gi, ie. when 1... The and their conjugates generate the (4). · How to work w/a group described by generalis and relations? Sometimes we know what G is, but sometimes we don't. Two useful ideas (among many): 1 the Cayley graph 2 named forms. Cayley graph; Given generators $g_1,...,g_n \in G$, the Cayley graph of G has) . vertices = the elements of the groups I . two vertices s, t are connected by an edge labelled g; when t = sg; [her we're doing right nultiplication; one could do left multi-instead] $\cdots \xrightarrow{-2} \xrightarrow{-1} \xrightarrow{0} \xrightarrow{1} \xrightarrow{1} \xrightarrow{1} \xrightarrow{1}$ Ex: Z with its weel generator 1: Ex. S_3 with generators $S_1 = (12)$ $S_2 = (23)$ $(s_i^2 = s_i)$ so edges are undirated $(123) = s_1 s_2 s_1 s_2 s_2 s_2 s_3 = (132)$ $S_1 = (12)$ $S_2 s_3 = (132)$ $S_1 = (12)$ $S_2 s_3 = (132)$ The fait that his closes up shows relation 5,525, = 525,52 (=> (5,52)=e) Since any word in 51,52 with relations $s_1^2 = s_2^2 = e$ can be reduced to (...s, s2 s1...), me can use this to check S3 = (s1, s2 | s1, s2, (s1s2)) Ex: S_{i} with generators $S_{i} = (i i+1), 1 \le i \le 3$: ("permobedion") Faces are \square relation $s_1s_3=s_3s_1$ and $s_1s_2s_1 = s_2s_1s_2$, $s_2s_3s_2 = s_3s_2s_3$ More generally, Sn has bollowing presentation: $S_{n} = \langle s_{i}, ..., s_{n-1} | s_{i}^{2} = 1 \ \forall i , s_{i} s_{i} = s_{i} + 1 \ s_{i} = s_{i+1} s_{i} = s_{i$ · Fait: Gails on its Cayley graph by left multiplication. ·

This is an ison once we have found a complete set of relations

· Word length of an element of G:= shortest distance from e to gEG in the Cayley graph.

This action is transitive on vertices (& on edges of given label gi): graph is very symmetric!

· For infinite groups, we can ask about the growth rate of G.
Given set of generous 9i, how does telements reproculed by words of length & N
grow with N? eop; does it grow polynomially or exponentially?
Even if we change our set of generators to some other of: und length of a given
element changes by a bounded factor only (bound = had lengths of new generators in tens of old ones & vice versa). So the exponential or polynomial nature of the growth is
of old ones & vile bessa). So the exponential or polynomial nature of the growth is
independent of the set of generators. Ex: f.g. abelian groups have polynomial growth / free grays have exp. growth.
For finite groups, the Cayley graph is finite and graph isn't relevant, but question about
word length renain interesting!
$Ex:$ in S_n , $\{s_i = (i \ i + i)\}$: the largest element is $\begin{pmatrix} 1 & 2 & \cdots & n \\ 5 & & J \\ n & n & \cdots & 1 \end{pmatrix}$ of wordlength $\frac{n(n-1)}{2}$, and
the word length of of Sn is #inversions (izj st. o(i)>o(j).
This is best undertood by representing permetations as diagrams
- Composition = stade dayrams
This is best undertood by representing perutations as diagrams - Composition = stack diagrams - Expression in terms of Si comes from decomposing diagram into i it!
layer $\sqrt{\text{single cossing}} = s_i$
· Reservation of Sn < any two dignans for or are related by
(1) $X \sim X \sim$
· Word length w(5) = #inversions is now clear
. We can even lit all the shortest words that can represent a given permutation!
Nomely: $\sigma \in S_n$ has a shortest word ending with S_i \iff $w(\sigma s_i) < w(\sigma)$ Call the set of such i the "ending set" of σ . \iff $\sigma(i+1) < \sigma(i)$. Then for each $i \in ending$ set, upon the process for $\sigma s_i^{-1} = \sigma s_i$.
· For each of Sn we can find a preferred expression of 5 as a word in Sylves Sn.
by choosing at each step the smallest i st. o(i+1) < o(i) to end the word
This gives a normal form for elements of Sn lie-a preferred word representing
each element) and hence a solution to the word problem - when do two words

reprosent the same element? () When does a word reprosent e € 6?).

(5)

For Sn, or other groups where it's well understood how to "calculate elements" (eg. groups of matrices, etc.), we don't need fancy algorithms or normal forms to solve the word problem. In many groups however this is all we have!

Ex: the braid group $B_n = \langle s_1 ... s_{n-1} | s_i s_j = s_j s_i \forall |i-j| \ge 2$, $s_i s_{i+1} s_i = s_{i+1} s_i s_{i+1} \rangle$ $(b.) s_i^2 + 1)$ $s_i = \frac{1}{3} (\frac{1}{3}) \frac{1}{3} \frac$

(important in knot theory etc.)

 $s_{i} = \frac{1}{1} \frac{1}{1}$ $s_{i}^{-1} = \frac{1}{1} \frac{1}{1} \frac{1}{1}$

Algorithmics is similar to Sn, wing:

· penulation braids =

any 2 strands coss at most once, all cossings /

There form a finite set, in bijection with Sn.

. Let Δ = the longer permutation braid. Since its shorter word can start/end with any s_i , $s_i^*(\Delta)$ is still a permutation braid + it conjugates $s_i \leftrightarrow s_{n-i}$.

Thus any elevent of Bn can be written as $g = \Delta^k P_1 \dots P_r$, $P_i = pent braids$ "Moving to the left everything that can be" => can find an expression st.

{ending set of P;} > { starting set of P;+1} \forall j

ie-any way of adding initial letter of a shortest word of Pit, to the end of Pi would cause it to be no large a permutation braid.

· This gives rise to "Garside nand form" and solution to word problem in Bn. (Garside 1969 + Thurston & Elisfai-Morton early 1990s).

Further examples (HW+ next lecture): - semidirect products

- Heisenbez group
- SL2(Z) and PSL2(Z)