### Round-Robin Problem

#### **Problem Description**

设有n个运动员要进行网球循环赛。设计一个满足下列条件的比赛日程表:

- 每个选手必须与其他 n-1 个选手各赛一次;
- 每个选手一天只能赛一次;
- 当 n 是偶数时,循环赛进行 n-1 天。
- 当 n 是奇数时,循环赛进行 n 天。

### **Problem Analysis**

我们在课堂上讨论过这个问题的一个特殊的情况:  $n=2^k$ 。

这个特殊的情况保证我们每次都能将选手等分,这样一来我们就可以利用**分治**的策略解决问题:每次把参赛选手平均分成两组,先递归地让他们进行内部的循环赛(并且这两组的情况完全对称);内部赛结束后,两组选手再与彼方小组的成员比赛。这时候两组人数相同,只需依次进行比赛即可。如下图所示。

为了是自动风风的英。是自己不得起了然间的了,但即以不是自己发展了。 第十日///5/6							
1	2	3	4	5	6	7	8
2	1	4	3	6	5	8	7
3	4	1	2	7	8	5	6
4	3	2	1	8	7	6	5
5	6	7	8	1	2	3	4
6	5	8	7	2	1	4	3
7	8	5	6	3	4	1	2
8	7	6	5	4	3	2	1

对于一般的 n, 我们并不能保证每次都能**等分**。事实上,只要 n 不是 n 的幂次,在划分过程中一定会出现奇数的情况。为了模拟上面的情况,在出现奇数时,我们采用**引入一个虚拟选手**的做法,使选手人数重新变为偶数,从而可以等分。

对于一个规模为 n 的问题,安排的日程表中出现虚拟选手意味着与他交手的选手在当轮轮空。而当我们要从规模为  $\lfloor \frac{n}{2} \rfloor$  的问题得到 n 个选手的循环赛日程表时,由于两组的对称性,**虚拟选手出现的位置将完全对称**,那么我们可以给同时轮空的两个选手安排一场比赛。之后再使两组的选手依次与对组选手交手。这样我们实现了比赛时间的最小化。

#### **Data structure**

我们考虑用一个矩阵来存储这个日程表。其中每一行代表一个选手,每一列代表一天,矩阵中的元素代表对手。例如,矩阵第i行,第j列上的元素,就是选手i在第j天比赛的对手。

	day 1	day 2	day 3	day 4	day 5
1					
2					
3					
4					
5					
6					

## Algorithm example

这里我们以6个选手为例,详述我们的算法。我们将仅描述从 3 个选手的子问题合并到 6 个选手的原问题的步骤。首先我们对 3 个选手的情况求解,结果里出现了虚拟选手; 1-3 号与 4-6 号情况完全一致。如下图:

	Day 1	Day 2	Day  3	Day 4	Day 5	Day 6
1	2	3	_			
2	1	_	3			
3	_	1	2			
4	5	6	_			
5	4	_	6			
6	_	4	5			

由于对称性,我们注意到i号和i+3号总是同时轮空,因此当他们轮空时,我们可以给他们安排比赛。如下图:

	Day 1	Day  2	Day 3	Day 4	Day 5	Day 6
1	2	3	4			
2	1	5	3			
3	6	1	2			
4	5	6	1			
5	4	2	6			
6	3	4	5			

接下来要进行的是组间的比赛。对于第一个组中的 i, 由于 i 和 i+3 可能已经比赛过,因此,第 i 个选手的第 4 天的对手从 i+4 开始,依次增加,如果超过了 6, 我们再填上 i+3。每行皆如此。如下图:

	Day 1	Day  2	Day  3	Day 4	$\mathrm{Day}\ 5$	$\mathrm{Day}\ 6$
1	2	3	4	5		
2	1	5	3	6		
3	6	1	2	4		
4	5	6	1	3		
5	4	2	6	1		
6	3	4	5	2		

这样,我们通过这个例子完成了对算法的描述。

# **Implementation**

```
def roundRobin(n):
    '''
    Scheduling for the round robin tournament.

Parameters:
    n -- the number of players. Could be any positive integers.

Returns:

schedule -- the schedule for the round robin tournament.
    Each row represent the opponents of a player, listed in time order.
    When n is even, schedule should be an n* (n-1) matrix;
    When it is odd, the size should be (n+1)*n with a virtual player.
    The virtual player's number exceeds n.
    '''
    if n & 1:
        n += 1

# base case
    if n == 2:
        return array([[2],[1]]) # 2>n, so it represent a virtual player
```

如果 n 为偶数,上述函数返回的就是最终的结果矩阵;如果 n 为奇数,这个函数返回的矩阵的前 n 行(去掉虚拟选手的赛程)是我们结果矩阵,且如果矩阵中出现了大于 n 的数字,表明这一轮该选手轮空。为了方便观察,我们写如下函数:

```
We simply print the result in a beautified way.
'''

for i in range(n):
    x = list(array(mat[i,:], dtype = int))
    virtual = [i for i in range(len(x)) if x[i] > n]

    if len(virtual) > 0:
        x[squeeze(virtual)] = None # None implies no match that day

    print('The schedule for player', i + 1, 'is:', x)
```

#### 简单看一下结果, 我们挑选了 n = 4k, 4k + 1, 4k + 2, 4k + 3 的多种样例:

```
The schedule for player 1 is: [None]
The schedule for player 1 is: [2, None, 3]
The schedule for player 2 is: [1, 3, None]
The schedule for player 3 is: [None, 2, 1]
The schedule for player 1 is: [2, 4, 3, 6, 7, 8, 5]
The schedule for player 2 is: [1, 3, 4, 7, 8, 5, 6]
The schedule for player 3 is: [4, 2, 1, 8, 5, 6, 7]
The schedule for player 4 is: [3, 1, 2, 5, 6, 7, 8]
The schedule for player 6 is: [5, 7, 8, 1, 4, 3, 2]
The schedule for player 7 is: [8, 6, 5, 2, 1, 4, 3]
The schedule for player 8 is: [7, 5, 6, 3, 2, 1, 4]
The schedule for player 1 is: [2, 4, 3, 6, 7, 8, 5, 10, 11, 12, 13, 14, 15, None, 9]
The schedule for player 2 is: [1, 3, 4, 7, 8, 5, 6, 11, 12, 13, 14, 15, None, 9, 10]
The schedule for player 5 is: [6, 8, 7, 4, 3, 2, 1, 14, 15, None, 9, 10, 11, 12, 13]
The schedule for player 6 is: [5, 7, 8, 1, 4, 3, 2, 15, None, 9, 10, 11, 12, 13, 14]
The schedule for player 7 is: [8, 6, 5, 2, 1, 4, 3, None, 9, 10, 11, 12, 13, 14, 15]
The schedule for player 8 is: [7, 5, 6, 3, 2, 1, 4, 9, 10, 11, 12, 13, 14, 15, None]
The schedule for player 9 is: [10, 12, 11, 14, 15, None, 13, 8, 7, 6, 5, 4, 3, 2, 1]
The schedule for player 10 is: [9, 11, 12, 15, None, 13, 14, 1, 8, 7, 6, 5, 4, 3, 2]
The schedule for player 11 is: [12, 10, 9, None, 13, 14, 15, 2, 1, 8, 7, 6, 5, 4, 3]
The schedule for player 13 is: [14, None, 15, 12, 11, 10, 9, 4, 3, 2, 1, 8, 7, 6, 5]
```

```
The schedule for player 15 is: [None, 14, 13, 10, 9, 12, 11, 6, 5, 4, 3, 2, 1, 8, 7]

The schedule for player 1 is: [2, 4, 3, 6, 7, 8, 5, 9, 10, 11, 12, 13, 14]

The schedule for player 2 is: [1, 3, 4, 7, 9, 5, 6, 10, 11, 12, 13, 14, 8]

The schedule for player 3 is: [4, 2, 1, 10, 5, 6, 7, 11, 12, 13, 14, 8, 9]

The schedule for player 4 is: [3, 1, 2, 5, 6, 7, 11, 12, 13, 14, 8, 9, 10]

The schedule for player 5 is: [6, 12, 7, 4, 3, 2, 1, 13, 14, 8, 9, 10, 11]

The schedule for player 6 is: [5, 7, 13, 1, 4, 3, 2, 14, 8, 9, 10, 11, 12]

The schedule for player 7 is: [14, 6, 5, 2, 1, 4, 3, 8, 9, 10, 11, 12, 13]

The schedule for player 8 is: [9, 11, 10, 13, 14, 1, 12, 7, 6, 5, 4, 3, 2]

The schedule for player 9 is: [8, 10, 11, 14, 2, 12, 13, 1, 7, 6, 5, 4, 3]

The schedule for player 10 is: [11, 9, 8, 3, 12, 13, 14, 2, 1, 7, 6, 5, 4]

The schedule for player 12 is: [13, 5, 14, 11, 10, 9, 8, 4, 3, 2, 1, 7, 6]

The schedule for player 13 is: [12, 14, 6, 8, 11, 10, 9, 5, 4, 3, 2, 1]
```

#### **Correctness checking**

我们需要验证算法的正确性。 怎样的结果是正确的:

- 每一行是 [1,n] 的不重复数字, 且应包含除行数以外的所有数字。
- 每一列除了轮空时有一个数字大于 n, 其他应是 [1,n] 上的不重复数字,并且应该全部出现。
- 对所有的 *i*. 如果

```
schedule[i,j] = k
```

就有

```
schedule[k-1,j] = i+1
```

```
def check(mat, n):
    ...
    To check the correctness of our algorithm

Parameters:
    mat -- the matrix gained from the function roundRobin.

n -- the parameter for roundRobin. Actually it is needless, since we can determine its value through mat's shape, but with it it is more convenient.

Returns:
    if return True, it means our algorithm works correctly. Else it is incorrect.
    ...
    # check the properties about rows
    for i in range(n):
        b = inld(list(range(l, n+l)), mat[i,:]) # check whether the numbers from l to n is in the i-th row of mat

        assert(squeeze(where(b == False)).shape == ()) # to assure there is only l number missing i.e. the row number
        if squeeze(where(b == False)) != i: # to check if the missing number is the row number
        print(squeeze(where(b == False)))
```

```
return False

if i == n-1 and n&l == 0:
    pass # if n is even there is only n-1 columns
else:
    c = inld(list(range(l, n+1)), mat[:,i]) # check whether the numbers from 1

to n is in the i-th column of mat
    if squeeze(where(c == False)).shape[0] != 0:
        return False

    # check whether there at most 1 larger than n
    d = inld(mat[:,i], list(range(l,n+1)))
    skip = squeeze(where(d == False))
    assert(skip.shape == () or skip.shape == (0,))
    if skip.shape == (0,) and mat[skip, i] <= n:
        return False

# check whether schedule[i,j] = k implies schedule[k-1,h] = i+1

for j in range(mat.shape[1]):
    k = int(mat[i,j])
    if mat[k-1,j] != i+1:
        print(i)
        print(mat[k,j])
        return False

# check the properties about columns

return True</pre>
```

由于循环赛参与者不太可能太多,我们仅检验 n 在 1-200 范围内算法的正确性:

```
result = True
for i in range(1,201):
    result = result and check(roundRobin(i),i)

print(result)
```

```
True
```

这样, 我们认为我们的算法是正确的。

## **Complexity analysis**

根据上面的讨论,我们可以得出  $T(n) = T(\frac{n}{2}) + O(n^2)$ 

根据主定理,我们知道  $T(n) = \Theta(n^2)$  下面我们通过测对不同大小的情况程序的运行时间来进行测试。

```
import time

def comp_test():
    runtime = []
    for i in range(1,1001):
        tic = time.time()
        roundRobin(i)
        toc = time.time()
        runtime.append(toc - tic)
        if i % 50 == 0:
            print('current:', i)
    return runtime
```

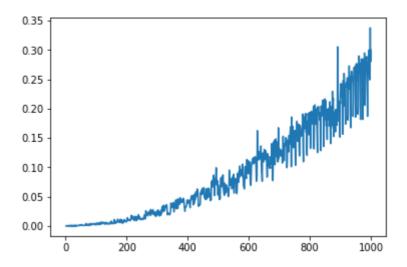
```
runtime = comp_test()
```

```
current: 50
current: 100
current: 150
current: 200
current: 250
current: 300
current: 350
current: 400
current: 450
current: 500
current: 500
current: 550
current: 660
current: 670
current: 750
current: 750
current: 750
current: 800
current: 850
current: 950
current: 950
current: 950
current: 950
```

```
from matplotlib.pyplot import *
% matplotlib inline

plot(list(range(1,1001)), runtime)
```

```
[<matplotlib.lines.Line2D at 0x2ab08a85c50>]
```



从上面的图像我们发现代码的运行时间大致与n的大小成二次方关系,这验证了我们刚才的分析。