INT 307 Multimedia Security System

Image Representation and Compression Sichen.Liu@xjtlu.edu.cn







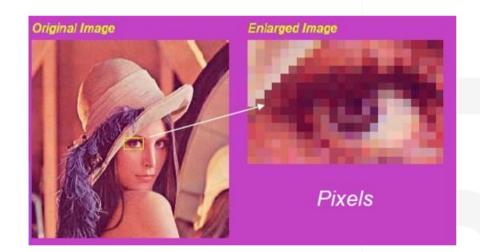
Aims

- Understand how people perceive image and video
- Master how images are presented and compressed in computer systems
- Understand the technology about information encoding



Image / Video Representation

- Images are composited by pixels
 - Each pixel has a color
 - The density of pixel is known as ppi
 - Human visual system has a resolution of approximate 300 ppi
- Video is composited by frames
 - Each frame presents the image at a moment
 - The rate of frame known as fps
 - Human visual system approximates 60 fps





Color Space

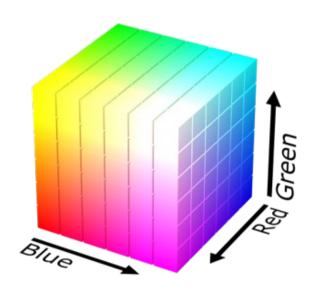
Usually, a vector is used to represent the color of a pixel, where each element in the vector represents a component

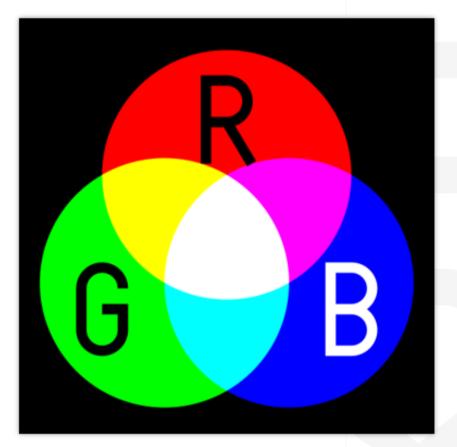
- RGB → Screen display (Red, Green and Blue)
- CMY(K) → Printer: Cyan, Magenta and Yellow (with Black)
- YUV → MPEG / JPEG system: separate light and color



RGB System

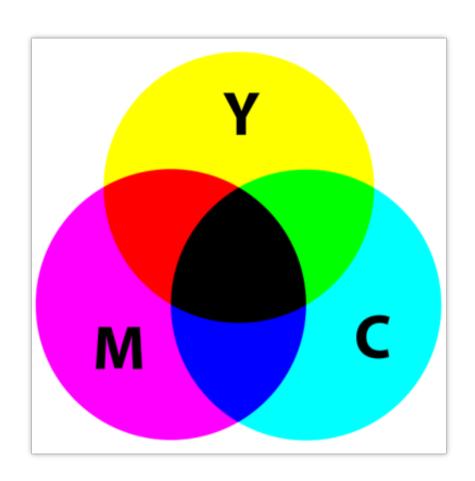
Additive Colour System







CMY(K) System



Subtractive color system

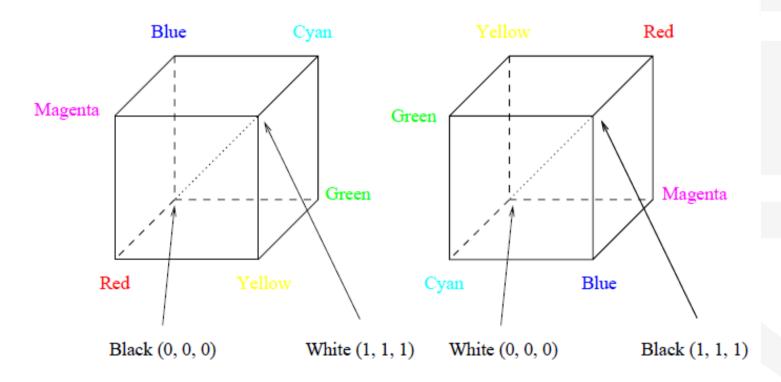
$$\begin{bmatrix} C \\ M \\ Y \end{bmatrix} = \begin{bmatrix} 1 \\ 1 \\ 1 \end{bmatrix} - \begin{bmatrix} R \\ G \\ B \end{bmatrix}$$



CMY(K) System

Subtractive color system

- No perfect Black → K component for black
- Printer system



The RGB Cube

The CMY Cube

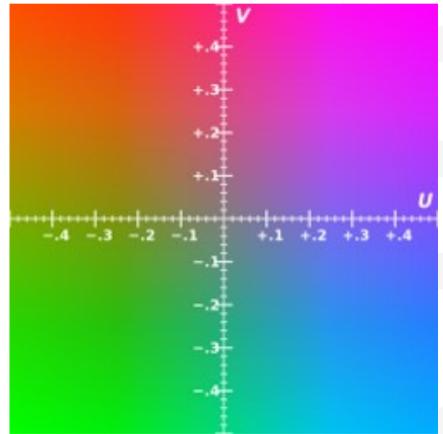


YUV System

- Separate Lightness and Color
- Easier compression as human visual system is more sensitive for light

$$\begin{bmatrix} Y' \\ U \\ V \end{bmatrix} = \begin{bmatrix} 0.299 & 0.587 & 0.114 \\ -0.299 & -0.587 & 0.886 \\ 0.701 & -0.587 & -0.114 \end{bmatrix} \begin{bmatrix} R' \\ G' \\ B' \end{bmatrix}$$

When Y = 0.5, the YUV system seems like





Calculation Question

- For an image RGB system is used with a depth of 8 bits for each component, how many bits needed to represent a picture whose resolution is 1920×1080?
- With the same configuration of the previous question as a frame of a piece of video, how many bits needed to record a one-hour video whose resolution is 1920×1080 with a frame rate of 60 fps?



Facts on Multimedia

- Fact 1: A film frame = 480 x 260 pixels = 374,400 bytes = 2,995,200 bits =599,040 alphabets = 124,025 English words
- Fact 2: A 2 hour 10 mins movie = 187,200 frames = 70,087,680,000 bytes= 23,217,480,000 English words = 7,067 years of words a person can talk
- Fact 3: 128 hours of video = 13,648,457 frames = 31,131,584,804,571 bytes= storage space of 31 1-TB hard disks at my laptop

Why Needs Compression?

- Demand

- VHS (Video Home System): $352 \times 240 \times 30 \times 24 = 60.8 \text{ Mb} / \text{ s}$
- DVD (Digital Versatile Disc): 720 x 480 x 30 x 24 = 248.8 Mb / s
- HDTV (High Definition Television): 1280 x 720 x 60 x 24 = 1327.2
 Mb / s

- Supply

- Telephone: < 56 Kb/s
- ISDN: 64 144 Kb/s
- T1: 1.5 Mb/s
- Ethernet: 10 100 Mb/s
- 802.11b: 1 10 Mb/s
- 802.11g: 50 100 Mb/s
- CD-ROM: 1.2 Mb/s

Compression

As a result, we need to compress the media

■ Image Compression

- Video Compression
 - Intra-frame compression
 - Inter-frame compression

Compression ratio

- If the compression and decompression processes induce no information loss, then the compression scheme is lossless; otherwise, it is lossy.
- Compression ratio:

$$compression \ ratio = \frac{B_0}{B_1}$$

- B_0 number of bits before compression
- B_1 number of bits after compression

Information Measurement

- Information Measure: Consider a symbol x with an occurrence probability p, its information content

$$I = \log \frac{1}{p(x)} = -\log p(x)$$

- The smaller the probability, the more info. the symbol contains
- The occurrence probability somewhat related to the uncertainty of the symbol

Information Entropy

 The Entropy is defined as the average information content per symbol of the source. The Entropy, H, can be expressed as follows

$$H = -\sum_{i=1}^{m} p_i \log_2 p_i \quad \text{bits}$$

- From this definition, the entropy of an information source is a function of occurrence probabilities.
- The entropy reaches the Max. when all symbols in the set are equally probable

Information Content

Consider the two blocks of binary data shown below which contains the most information?

$$P_{0} = \frac{63}{64}$$

$$P_{1} = \frac{1}{64}$$

$$P_{1} = \frac{32}{64} = \frac{1}{2}$$

$$P_{1} = \frac{32}{64} = \frac{1}{2}$$

$$P_{1} = \frac{32}{64} = \frac{1}{2}$$

$$H = -\frac{63}{64} \log_{2} \frac{63}{64} - \frac{1}{64} \log_{2} \frac{1}{64}$$

$$= 0.116 \text{ bits/pixel}$$

$$P_{1} = \frac{32}{64} = \frac{1}{2}$$

$$H = -\frac{1}{2} \log_{2} \frac{1}{2} - \frac{1}{2} \log_{2} \frac{1}{2}$$

$$= 1.0 \text{ bits/pixel}$$

$$P_0 = \frac{32}{64} = \frac{1}{2}$$

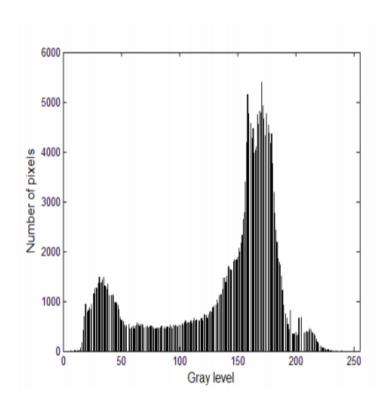
$$P_1 = \frac{32}{64} = \frac{1}{2}$$

$$H = -\frac{1}{2}\log_2\frac{1}{2} - \frac{1}{2}\log_2\frac{1}{2}$$
= 1.0 bits/pixel



Entropy Coding

- Image Histogram: entropy = 7.63 bits/pixel







Entropy Coding

- Fixed length coding

Symbol	Probability	Codeword	Codeword Length
А	0.75	00	2
В	0.125	01	2
С	0.0625	10	2
D	0.0625	11	2

- Average bits/symbol
- Average bits/symbol = 0.75*2 + 0.125*2 + 0.0625*2 + 0.0625*2 = 2.0 bits/pixel

Entropy Coding

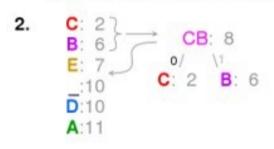
- Variable Length Coding

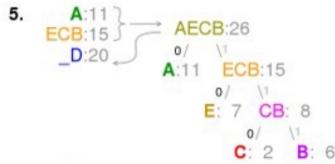
Symbol	Probability	Codeword	Codeword Length
А	0.75	0	1
В	0.125	10	2
С	0.0625	110	3
D	0.0625	111	3

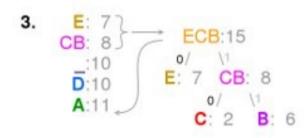
- Average bits/symbol
- Average bits/symbol = 0.75*1 + 0.125*2 + 0.0625*3 + 0.0625*3= 1.375 bits/pixel

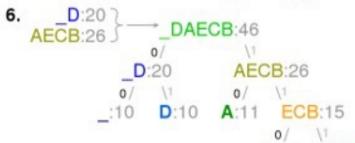
Huffman Codingwords

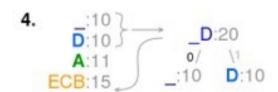
- If the symbol probabilities are known, Huffman codewords can be automatically generated
 - Reorder in decreasing order of probability at each step
 - Merge the two lowest probability symbols at each step
 - 1. "A_DEAD_DAD_CEDED_A_BAD_BABE_A_BEADED_ABACA_BED"













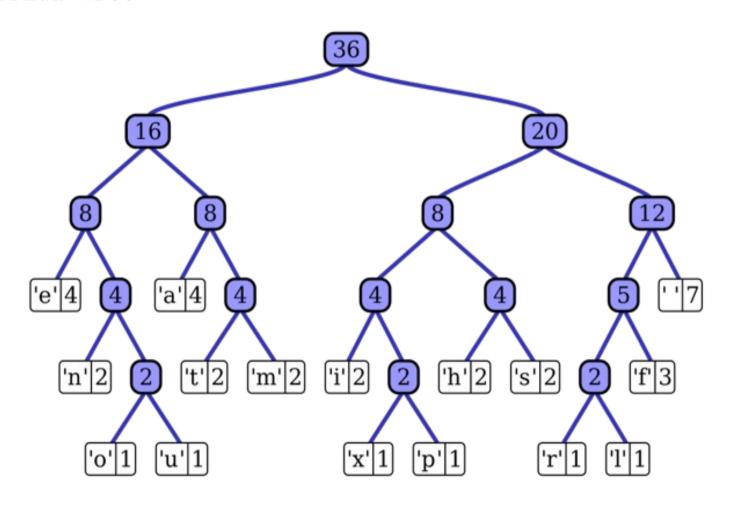


Huffman Codingwords

- Huffman codewords have to be an integer number of bits long
- The symbol probabilities must be known in the decoder size. If not, they must be generated and transmitted to the decoder with the Huffman coded data
- A larger of number of symbols results in a large codebook
- Dynamic Huffman coding scheme exists where the code words are adaptively adjusted during encoding and decoding, but it is complex for implementation

Exercise

Please use Huffman coding to code "this is an example of a huffman tree"





Arithmetic Coding

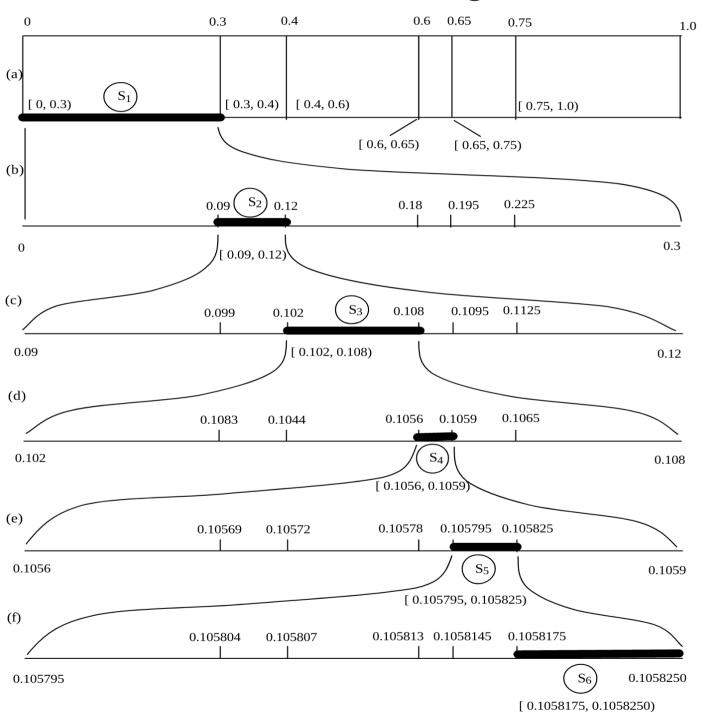
- It overcomes limitation of Huffman coding: non-integer length coding, and probability distribution can be derived in real-time
- It operates by replacing a stream of input symbols with a single floating point output number.
- Consider the following symbols with probabilities (CP is the cumulative probability):

Source symbol	Occurrence probability	Associated subintervals	СР
S_1	0.3	[0, 0.3)	0
S_2	0.1	[0.3, 0.4)	0.3
S_3	0.2	[0.4, 0.6)	0.4
S ₄	0.05	[0.6, 0.65)	0.6
S_5	0.1	[0.65, 0.75)	0.65
S ₆	0.25	[0.75, 1.0)	0.75

Arithmetic Coding

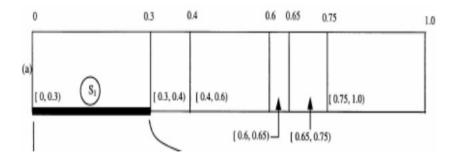
- Suppose we wish to encode the string: S1; S2; S3; S4; S5; S6;
- We start with the interval [L, H) and set to [0, 1) for 6 symbols.
- Since the first symbol is S1, we pick up its subinterval[L, H) = [0.0, 0.3), and any real symbols can be considered as disjoint to be divided in the same way.
- To encode the S2, we use the same procedure as used in above to divide the interval [0, 0.3) into six sub-intervals. We pick up the S2 subinterval [0.09, 0.12)

Arithmetic Coding



Arithmetic Decoding

- For encoding, the input is a source symbol string and the output is a subinterval (called the final subinterval). For our case, it is [0.1058175, 0.1058250)
- Decoding sort of reverses what encoding has done
- The decoder knows the encoding procedure and therefore has the information contained in the following figure

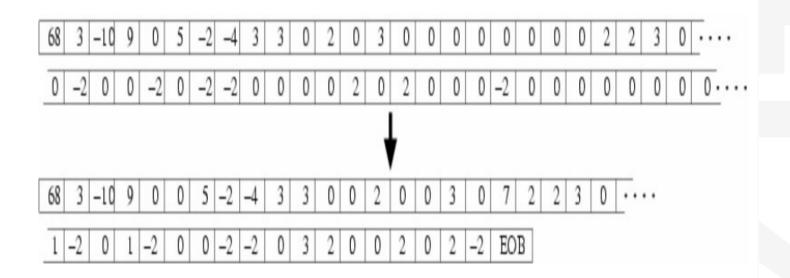


- It compares the lower end point of the final subinterval 0.1058175 with all the end points in the above figure 0 < 0.1058175 < 0.3



Run-length Coding

- In run-length coding, a run of consecutive identical symbols is combined together and represented by a single codeword.
- The various run-lengths are then represented by a variable length (e.g. Huffman) codeword.
- Use an escape code.



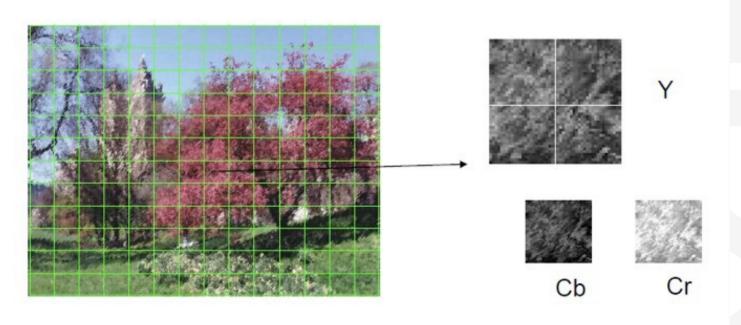


JPEG

- Divide Images into 8×8 (luminance) or 16×16 (chrominance) blocks
- Perform DCT on image blocks
- Apply Quantization
- Perform Zigzag ordering and Run-length encoding
- Entropy Encoding (Huffman Encoding)

Macroblocks

- An image is divided into 8×8 blocks for the luminance components
- An image is divided into 16×16 blocks for the chrominance components
- Chrominance blocks are down-sampled to 8×8 blocks (called 4:2:2 format)
- Zero-Padding on boundary blocks



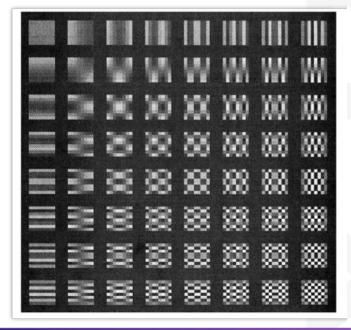
Discrete Frequency-Domain Analysis

$$F(\mu,\nu) = \frac{C(\mu)}{2} \frac{C(\nu)}{2} \sum_{y=0}^{7} \sum_{x=0}^{7} f(x,y) \cos[(2x+1)\mu\pi/16] \cos[(2x+1)\nu\pi/16]$$

$$C(\mu) = egin{cases} rac{1}{\sqrt{2}}, \emph{if} \mu = 0 \ 1, \emph{if} \mu > 0 \end{cases}$$

- Frequency of image
- Rapid changes → High frequency
- Minor changes → Low frequency
- Usually the image is low frequency (redundancy in spatial)

2D-DCT bases





Quantization

- Human eye is good at seeing small difference in brightness over a relative area
- Not good at the brightness variations of high frequency
- Reduce the amount of information in the high frequency component
- Divide by a constant and round to the nearest integer
- Effect: high frequency → zero
- Compression ratio: larger element value → greater compression.

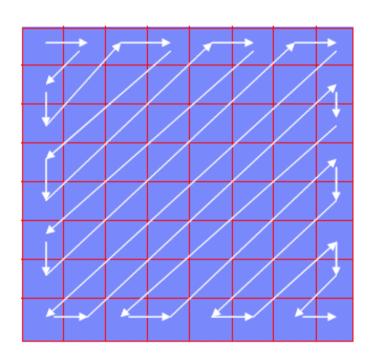
Quantization

- Luminance Quantization Table and Chrominance Quantization Table

16	11	10	16	24	40	51	61	17	18	24	47	99	99	99	99
12	12	14	19	26	58	60	55	18	21	26	66	99	99	99	99
14	13	16	24	40	57	69	56	24	26	56	99	99	99	99	99
14	17	22	29	51	87	80	62	47	66	99	99	99	99	99	99
18	22	37	56	68	109	103	77	99	99	99	99	99	99	99	99
24	35	55	64	81	104	113	92	99	99	99	99	99	99	99	99
49	64	78	87	103	121	120	101	99	99	99	99	99	99	99	99
72	92	95	98	112	100	103	99	99	99	99	99	99	99	99	99
	(a)					(b)									

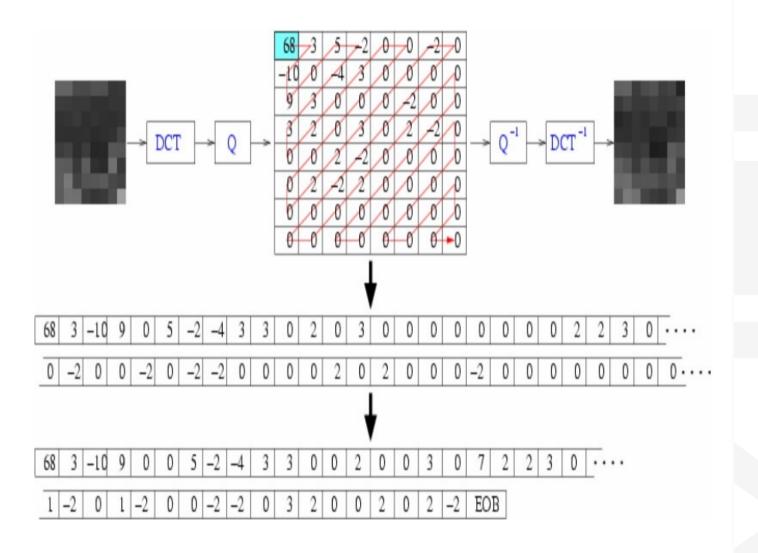
Zig-Zag Order

An order to sort the 2D DCT coefficients to 1D signals



Lossless Compression

Run Length Coding



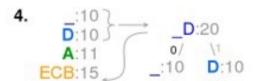


Lossless Compression

Huffman Coding

- 1. "A_DEAD_DAD_CEDED_A_BAD_BABE_A_BEADED_ABACA_BED"
- 2. C: 2 B: 6 E: 7 :10 C: 2 B: 6 D:10 A:11

- 5. A:11 AECB:26
 D:20 A:11 ECB:15
 O/ \1
 E: 7 CB: 8
 O/ \1
 C: 2 B: 6
- 3. E: 7 CB: 8 CB: 15 CB: 10 CB: 7 CB: 8 CB: 7 CB



- D: 01 A: 10 E: 110 C: 1110 B: 1111



JPEG

As an example, the bottom picture shows our compressed image.

This right image is a mere 1.5% of the original size.







THANK YOU









