

Compilers Principals - Chapter3

Zhixin Zhang, 3210106357

Problems: 3.6, 3.9, 3.13, 3.14.

1 Problem 3.6

a. Calculate nullable, FIRST, and FOLLOW for this grammar:

$$S \rightarrow uBDz$$

$$B \rightarrow Bv$$

$$B \rightarrow w$$

$$D \rightarrow EF$$

$$E \rightarrow y$$

$$E \rightarrow$$

$$F \rightarrow x$$

$$F \rightarrow$$

b. Construct the LL(1) parsing table.

c. Give evidence that this grammar is not LL(1).

d. Modify the grammar **as little as possible** to make an LL(1) grammar that accepts the same language.

解:

a.

	nullable	FIRST	FOLLOW
S	no	u	
B	no	w	x, y, z
D	yes	x, y	z
E	yes	y	x, z
F	yes	x	z

表 1 nullable, FIRST and FOLLOW

b.

	u	v	w	x	y	z
S	$S \rightarrow uBDz$					
B			$B \rightarrow w, B \rightarrow Bv$			
D				$D \rightarrow EF$	$D \rightarrow EF$	$D \rightarrow EF$
E				$E \rightarrow$	$E \rightarrow y$	$E \rightarrow$
F				$F \rightarrow x$		$F \rightarrow$

表 2 LL(1) parsing table

c. No, according to the parsing table, we get two possible production when the top of stack of non-terminal is B and the input token is w, $B \rightarrow w, B \rightarrow Bv$.

d. I modify the grammar as follows:

$$S \rightarrow uBEFz$$

$$B \rightarrow wB'$$

$$B' \rightarrow vB'$$

$$B' \rightarrow$$

$$E \rightarrow y$$

$$E \rightarrow$$

$$F \rightarrow x$$

$$F \rightarrow$$

□

2 Problem 3.9

Diagram the LR(0) states for Grammar 3.26, build the SLR parsing table, and identify the conflicts.

解:

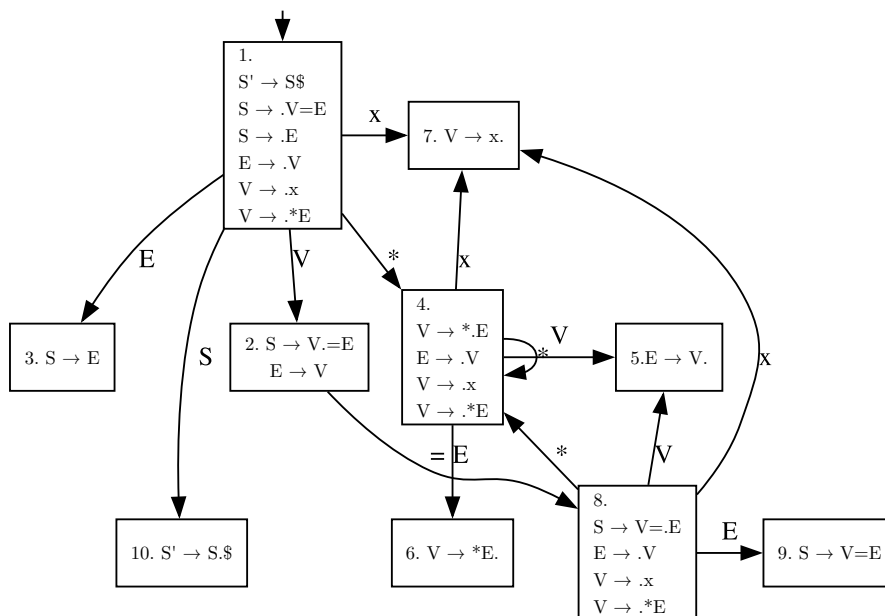


图 1 LR(0) state diagram

	x	*	\$	=	E	V	S
1	s7	s4			g3	g2	g10
2			r3	s8, r3			
3	r10	r10	r10	r10			
4	s7	s4			g6	g5	
5			r6	r6			
6			r5	r5			
7			r2	r2			
8	s7	s4			g9	g5	
9			r10				
10			a				

表 3 SLR parsing table

□

3 Problem 3.13

Show that this grammar is LALR(1) but not SLR:

- | | | | |
|------|-----------------------|------|-----------------------|
| $_0$ | $S \rightarrow X \$$ | $_3$ | $X \rightarrow d c$ |
| $_1$ | $X \rightarrow M a$ | $_4$ | $X \rightarrow b d a$ |
| $_2$ | $X \rightarrow b M c$ | $_5$ | $M \rightarrow d$ |

解:

Follow(x) = { \$ }, Follow(M) = { a, c }

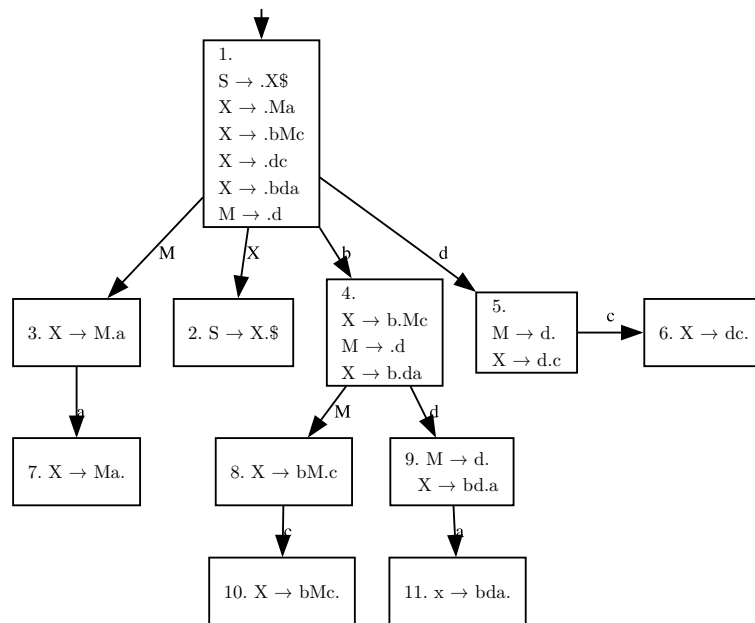


图 2 SLR state diagram

At state 5 and the input is c, we have two possible ways, either to reduce to state 3 or to shift to state 6. So it is not SLR.

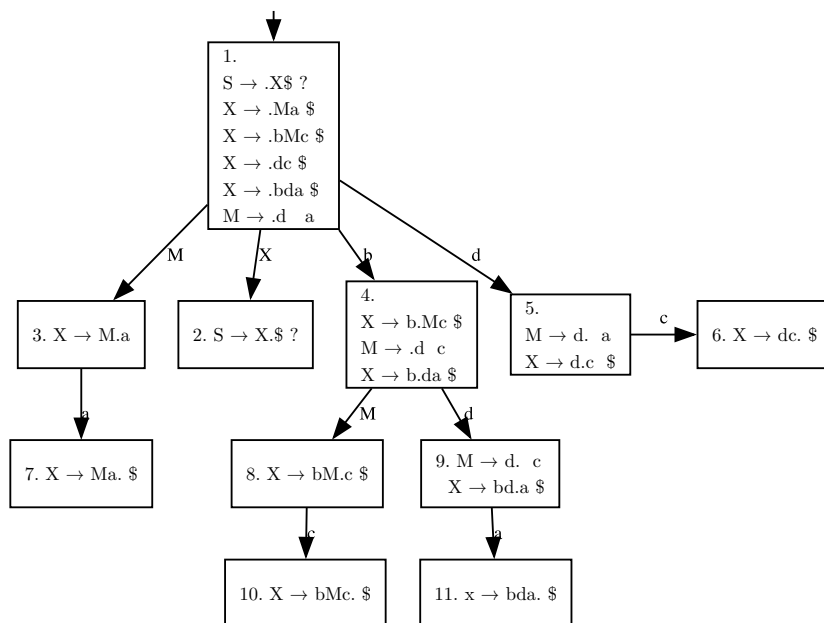


图 3 LR(1) state diagram

The language is LALR(1).

	a	b	c	d	\$	X	M
1		s4		s5		g2	g3
2					a		
3	s6						
4				s8			g7
5	r3			s9			
6					r2		
7			s10				
8	s11		r7				
9					r2		
10					r2		
11					r2		

表 4 LALR(1) parsing table

□

4 Problem 3.14

Show that this grammar is LL(1) but not LALR(1):

- | | | | |
|---|---------------------|---|---------------------|
| 1 | $S \rightarrow (X$ | 5 | $X \rightarrow F]$ |
| 2 | $S \rightarrow E]$ | 6 | $E \rightarrow A$ |
| 3 | $S \rightarrow F)$ | 7 | $F \rightarrow A$ |
| 4 | $X \rightarrow E)$ | 8 | $A \rightarrow$ |

解: The language is LL(1).

	()	[]
X		$X \rightarrow E)$	$X \rightarrow F]$	
S	$S \rightarrow (X$	$S \rightarrow F)$	$S \rightarrow E]$	
E		$E \rightarrow A$	$E \rightarrow A$	
F		$F \rightarrow A$	$F \rightarrow A$	
A		$A \rightarrow$	$A \rightarrow$	

图 4 LL(1) Parsing table

In state 7, there is a reduce-reduce conflict, so it is not LALR(1).

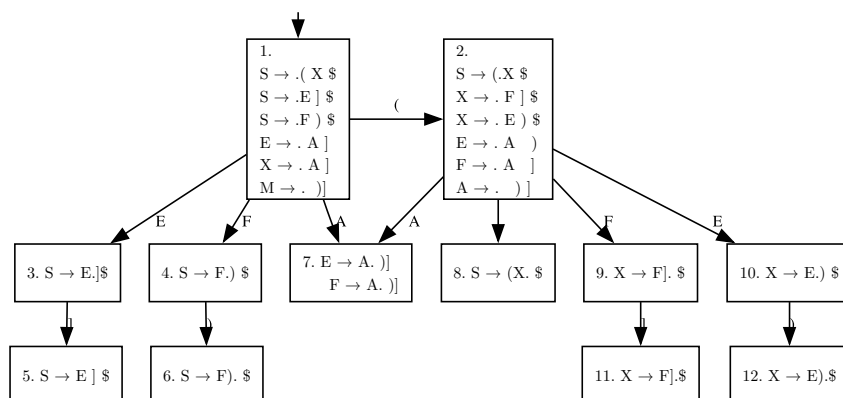


图 5 LALR(1) state diagram

□