

17/WAKU2-RLN-RELAY:
Privacy-Preserving
Peer-to-Peer
Economic
Spam Protection

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Link to the paper:
https://github.com/vacp2p/research/blob/master/rln-research/Waku_RLN_Relay.pdf

Today, I am going to present waku-rln-relay which is a privacy preserving economic spam protection that suits private and anonymous messaging systems
This is a joint work by me, Oskar Thoren, Barry Whitehat, Wei Jie Koh, Onur Kilic, and Kobi Gurkan
The paper is also available in the link below

Contents

- WAKU2
- WAKU2-RELAY: Privacy-preserving p2p transport protocol
- Spam issue in WAKU2-RELAY
- Privacy-Preservation and Spam protection
- State-of-the-art p2p spam protections
- WAKU2-RLN-RELAY: Privacy-Preserving Peer-to-Peer Economic Spam Protection
- Future work

The agenda is as illustrated

I'll start by discussing WAKU2 and WAKU-RELAY, I will review the spam issue and the related studies, then I'll explain waku-rln-relay architecture and how it excels its counterparts. In the end, I will shed light on the future work.

WAKU2 [1]

- A family of modular, privacy-preserving peer-to-peer (p2p) protocols for private, secure, censorship resistant communication
- Suitable for resource restricted devices e.g., mobile phones
- WAKU2 protocols include:
 - **WAKU2-RELAY: privacy-preserving transport**
 - WAKU2-STORE: historical message storage
 - WAKU2-FILTER: light version of WAKU2-RELAY for bandwidth limited devices
 - **WAKU2-RLN-RELAY: spam-protected version of WAKU2-RELAY**
 - And many more ...
- For the full list of RFCs is available in rfc.vac.dev

[1] <https://rfc.vac.dev/spec/10/>

Waku2 is a stack of peer-to-peer (p2p) protocols that enable anonymous and privacy-preserving communication.

It is designed to be able to run in resource-restricted environments.

The focus of today's talk is WAKU-RELAY and WAKU-RLN-RELAY which are the transport layers of Waku

WAKU2-RELAY [1]

- Publisher-Subscriber Model
- Gossip-based Routing (extension of libp2p GossipSub-v1.1 [2])
- Anonymous and Privacy-Preserving

[1] <https://rfc.vac.dev/spec/11/>

[2] <https://github.com/libp2p/specs/tree/master/pubsub/gossipsub>

Let me start by WAKU2-RELAY transport protocol

It follows publisher-subscriber messaging model

Implements gossip-based routing protocol

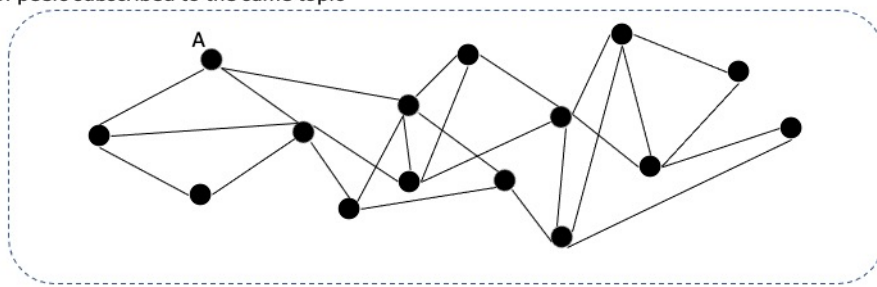
Is a minor extension of libp2p GossipSub protocol

Its end goal is to enable an anonymous and privacy preserving p2p network layer

WAKU2-RELAY

- Peers subscribed to the same topic form a mesh

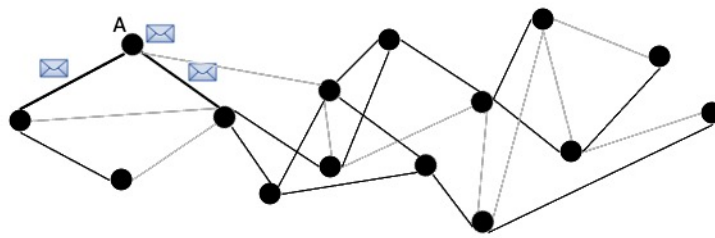
Mesh of peers subscribed to the same topic



Peers in wakurelay congregate around topics they are interested in and can send messages to that topic or receive messages within that topic

WAKU2-RELAY

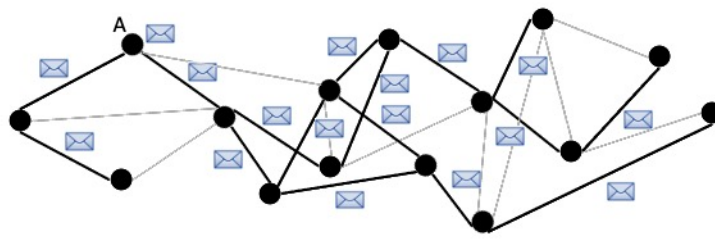
- Peers subscribed to the same topic form a mesh
- Peers route messages by sending them to a subset of their connections



To publish a message, the author forwards its message to a subset of neighbors.

WAKU2-RELAY

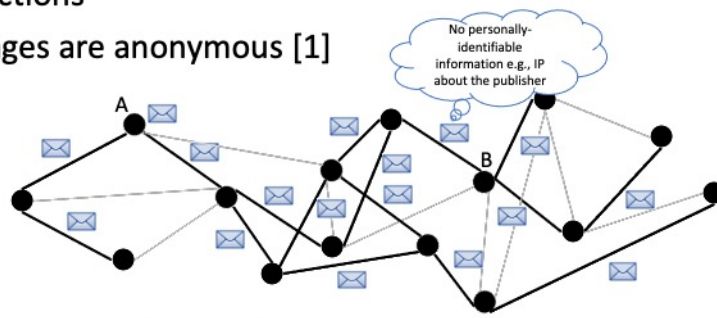
- Peers subscribed to the same topic form a mesh
- Peers route messages by sending them to a subset of their connections



The neighbors proceed similarly till the message gets propagated in the network of the subscribed peers.

WAKU2-RELAY

- Peers subscribed to the same topic form a mesh
- Peers route messages by sending them to a subset of their connections
- Messages are anonymous [1]

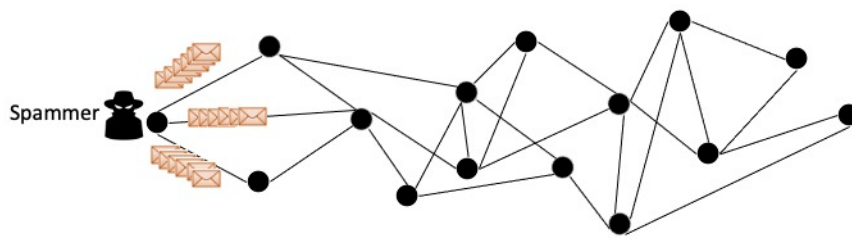


[1] <https://rfc.vac.dev/spec/11/#security-analysis>

Messages are anonymous, i.e., there is no personally identifiable information like IP address attached to them, thus it is not feasible to identify the message origin. For more detailed security analysis see the link below

Spam issue in WAKU2-RELAY

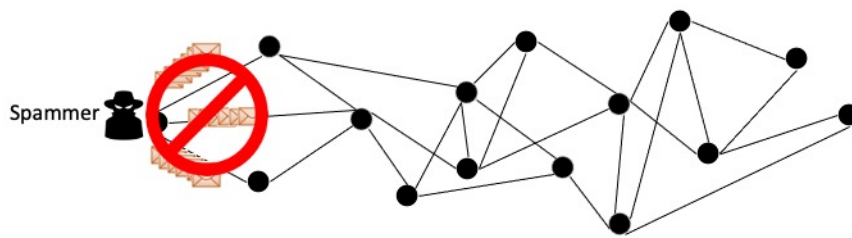
- We define spammers as entities that publish a large number of messages in a short amount of time, and cause denial-of-service



Waku Relay, as an open messaging network, is prone to spam messages and spammers;
we define spammers as entities that publish a large number of messages in a short amount of time, and cause Denial of Service attack.

Spam issue in WAKU2-RELAY

- We define spammers as entities that publish a large number of messages in a short amount of time, and cause denial-of-service
- Spam Protection = Controlled Messaging Rate

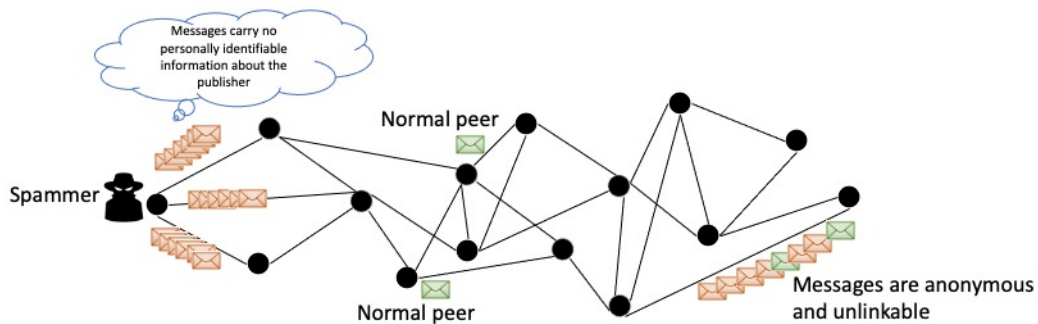


Waku Relay, as an open messaging network, is prone to spam messages and spammers;
we define spammers as entities that publish a large number of messages in a short amount of time, and cause Denial of Service attack.

With that definition, Spammers can be controlled if we can control their messaging rate

Privacy-Preservation and Spam protection

- Messages are anonymous: No Personally Identifiable information is available

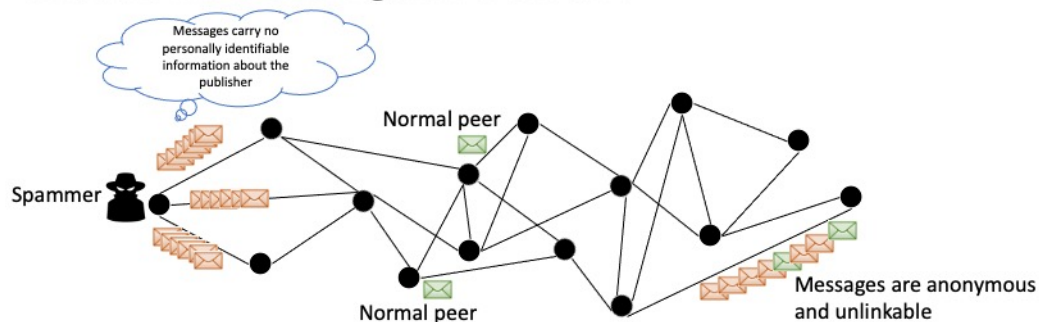


But there is an issue, messages are anonymous (no information like IP is available) and not linkable to their authors

A routing peer won't be able to tell apart spam messages from non-spam

Privacy-Preservation and Spam protection

- Messages are anonymous: No Personally Identifiable information is available
- Solutions like IP blocking are not effective



But there is an issue, messages are anonymous (no information like IP is available) and not linkable to their authors

A routing peer won't be able to tell apart spam messages from non-spam

Thus solutions like IP blocking are not effective

State-of-the-art p2p spam protections

- Proof-of-work [1] deployed by Whisper [2]
 - Computationally expensive
 - Not suitable for network of heterogeneous peers with limited resources
- Peer Scoring [3] in libp2p
 - Local to each peer
 - No global identification of spammer
 - Subject to inexpensive attacks using bots
 - Prone to censorship

[1] Cynthia Dwork and Moni Naor. Pricing via processing or combatting junk mail. In Annual 456 international cryptology conference. Springer, 1992.

[2] <https://eips.ethereum.org/eips/eip-627>.

[3] <https://github.com/libp2p/specs/blob/master/pubsub/gossipsub/gossipsub-v1.1.md#peerscore>.

The state-of-the-art p2p spam protection techniques are Proof of Work (POW) deployed by Whisper and Peer scoring method adopted by libp2p

The PoW is not computationally efficient and does not fit resource limited devices (limited resources won't be able to participate and benefit from the messaging system)

In peer scoring

The solution is local since each peer monitors and scores its direct connections and drops the connections with low scores.

However, a spammer would be still able to continue its activity by switching its connection from one peer to another as soon its score drops a threshold.

Peer scoring is also subject to inexpensive attacks where the spammer can deploy millions of bots to send bulk messages.

It is also prone to censorship

WAKU2-RLN-RELAY [1]

WAKU2-RLN-RELAY = WAKU2-RELAY + Rate Limiting Nullifiers (RLN)

- P2p solution
- Global spam protection
- Privacy preserving
- Efficient
- Economic incentives

[1] <https://rfc.vac.dev/spec/17/>

In WAKU-RLN-RELAY we cope with the aforementioned issues in a nutshell, WAKU-RLN-RELAY provides a private p2p transport protocol that is spam-protected
it uses waku-relay and combines it with rate limiting nullifier construct to control the messaging rate
The end result has a p2p structure, with no central entity involved
it **allows global identification** and removal of spammers
it is **privacy-preserving** since there is no need to personally identifiable information e.g., email address, IP, etc. about peers to be able to identify and block spammers
It is **efficient** i.e., with no unreasonable computational, storage, memory, and bandwidth requirement, as such, it fits the network of heterogeneous peers with limited resources.
It has **economic-incentives, i.e.**, there is a financial punishment for the spammers and a financial reward for those who catch spammers.

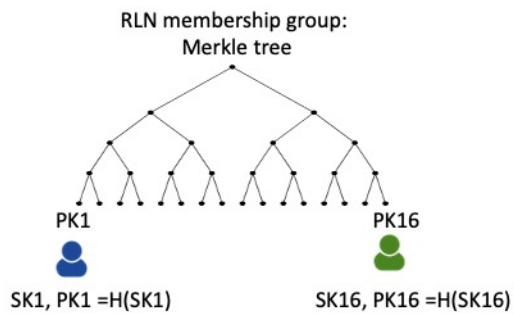
RLN Primitive [1]

- RLN is a zero-knowledge and rate-limited signaling framework
- Each user can only send M messages for each External Nullifier
- External nullifier can be seen as a voting booth where each user can only cast one vote
- M and external nullifier are application dependent
- M=1 for this presentation

[1] <https://ethresear.ch/t/semaphore-rln-rate-limiting-nullifier-for-spam-prevention-in-anonymous-p2p-setting/5009>

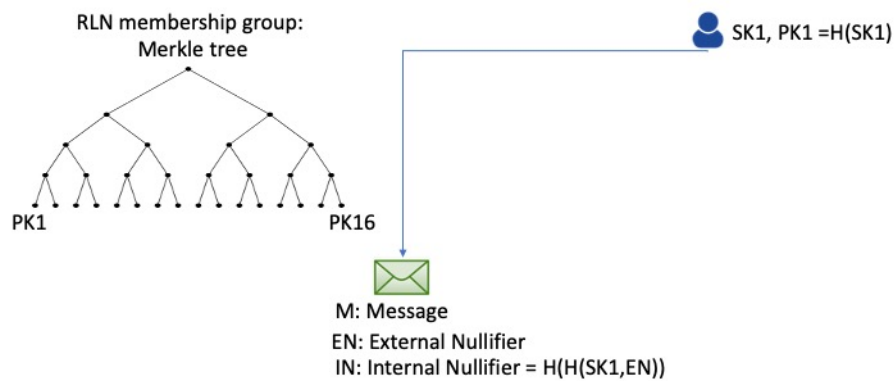
- Lets begin with the RLN construct
- it is a zero-knowledge and rate-limited signaling framework
- It allows a set of users to broadcast arbitrary signals (where signal is any value like a string, vote, etc.) while proving they are among a group of authorized users without disclosing their identities
- The idea is that each user can only send M messages for a specific external nullifier. External nullifier can be seen as a voting booth where each user can only cast one vote
- For the rest of this presentation we consider M as 1

RLN Primitive: Membership Tree



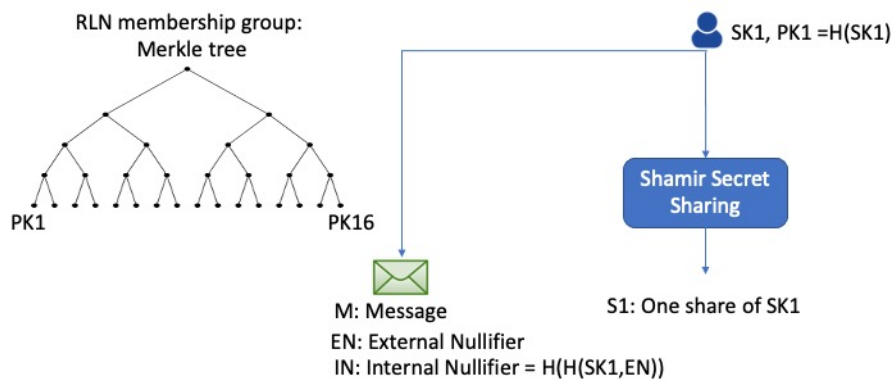
- It consists of A merkle tree that represents a group of authorized users, each user has a pk registered in this tree, the corresponding SK is only known to the user

RLN Primitive: Signaling



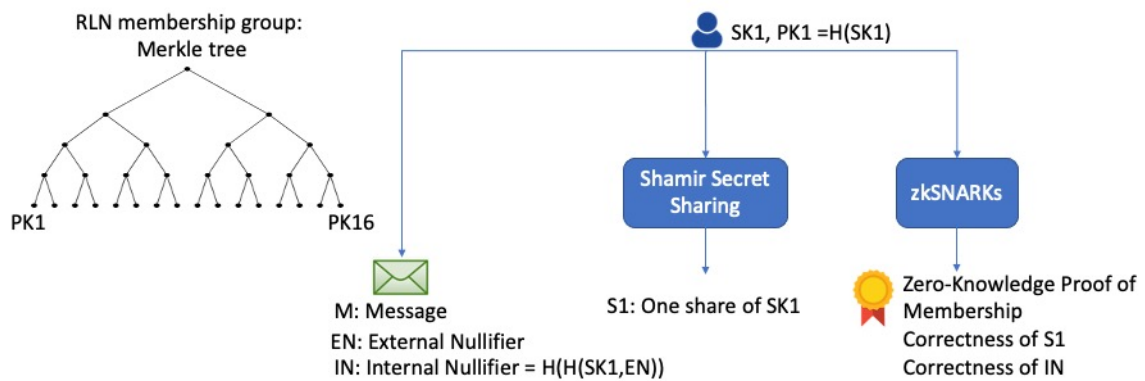
- For the signaling/publishing:
- The user specifies a external nullifier as well as an internal nullifier which is derived from the SK and the external nullifier

RLN Primitive: Signaling



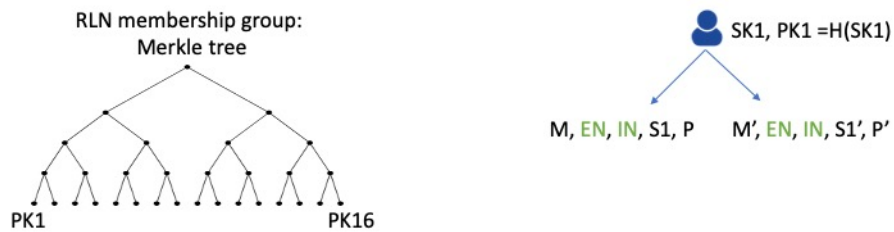
- The user also discloses a share of its secret key using shamir secret sharing scheme
- This share will be used to remove the user from the group in case of double signaling

RLN Primitive: Signaling



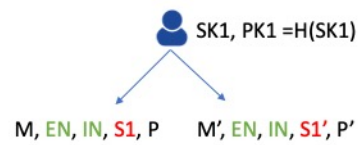
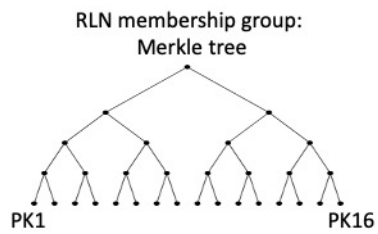
- Finally, the user proves in zero knowledge manner that:
- It is part of the group
- And that has computed the secret share and the Internal nullifier correctly

RLN Primitive: Detecting double signaling



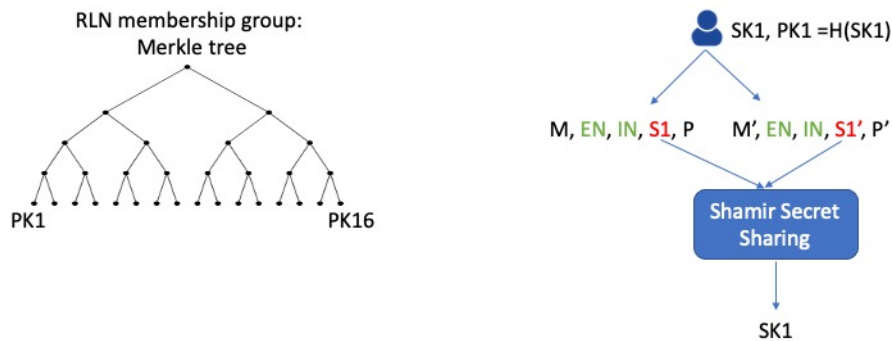
- If a user publishes more than 2 times for the same external nullifier, it will end up having two messages with the same external and internal nullifiers (remember those are deterministic values computed from SK and external nullifier)
- So the double signal can be detected

RLN Primitive: Detecting double signaling



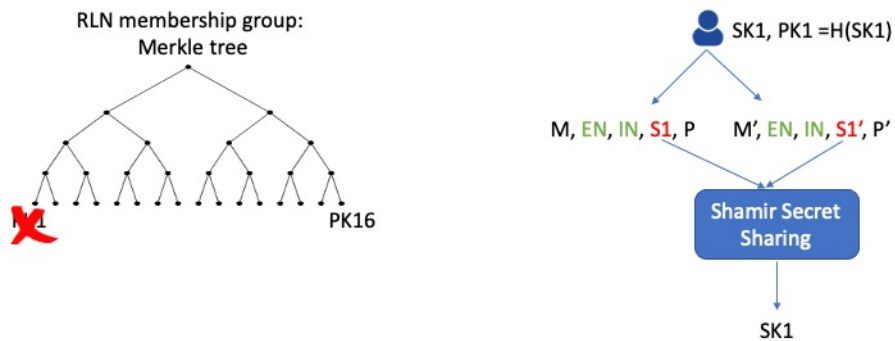
- Also, it will disclose 2 shares of its SK

RLN Primitive: Detecting double signaling



- using which the corresponding SK can be reconstructed (property of shamir secret sharing).

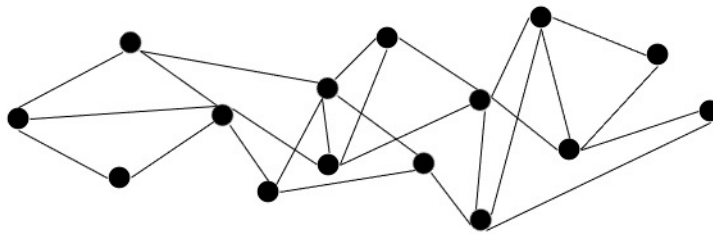
RLN Primitive: Detecting double signaling



- And the user gets removed from the tree. Thus, it Can no longer use that sk for messaging.

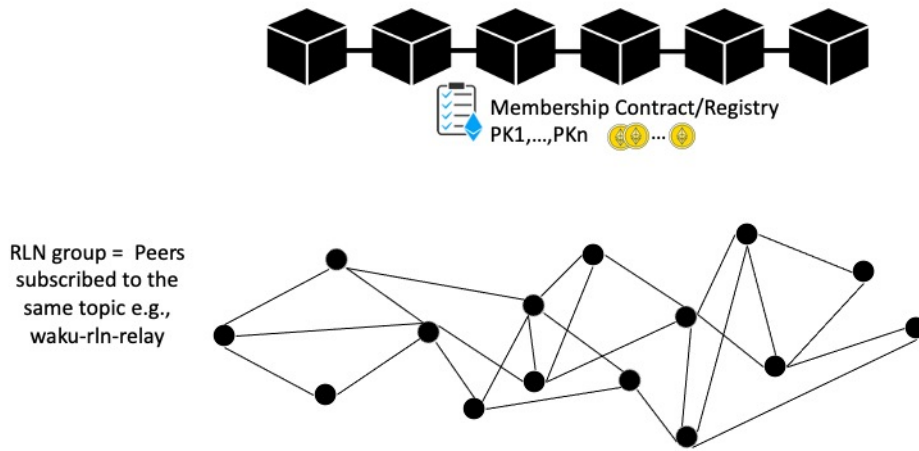
WAKU2-RLN-RELAY

RLN group = Peers
subscribed to the
same topic e.g.,
waku-rln-relay



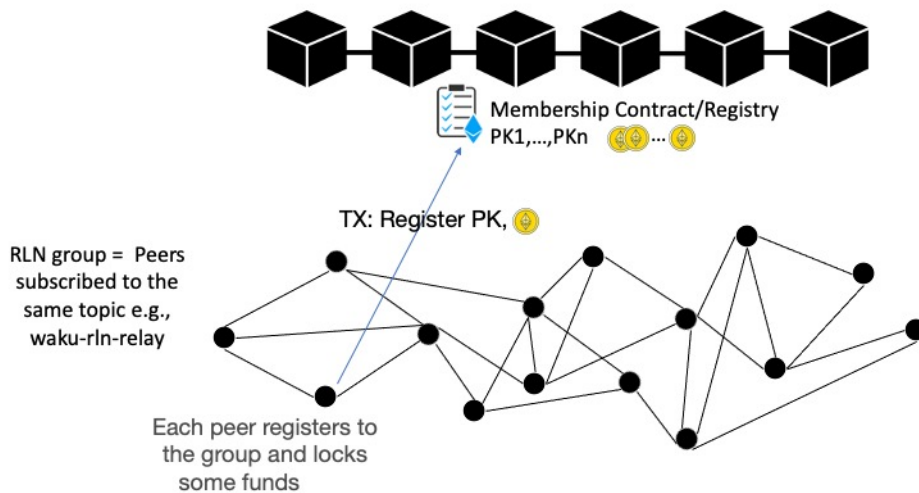
Now lets see the end to end integration of RLN into Waku-rln-relay
In Waku-rln-relay, the rln group consists of peers that belong to the same GossipSub
layer (subscribed to the same topic)

WAKU2-RLN-RELAY: Registration



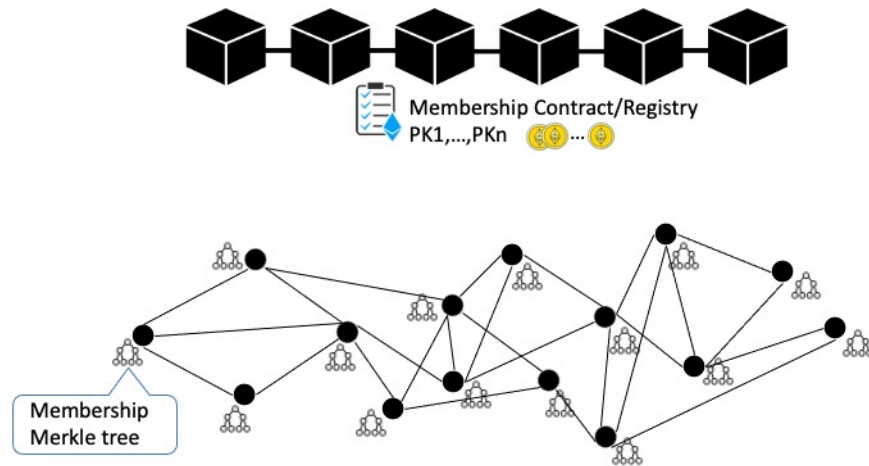
Each peer has a rln pk, and the list of pks is stored in a contract deployed on the Eth blockchain.

WAKU2-RLNR-ELAY: Registration



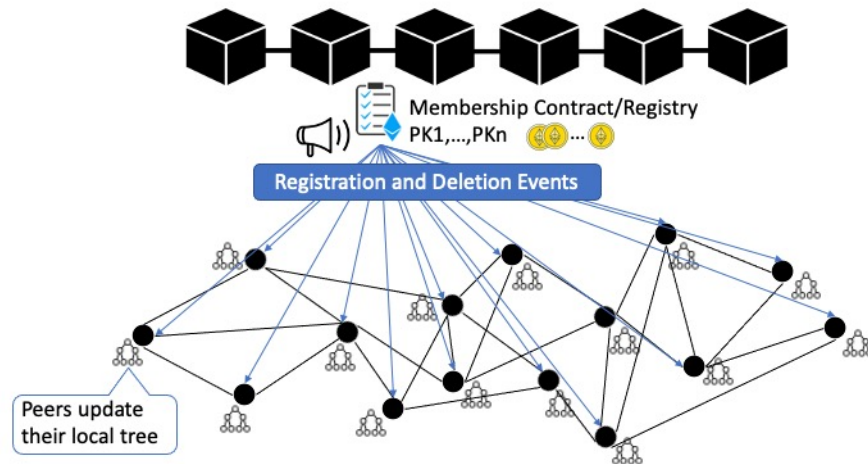
A peer willing to publish a message should register by sending a tx to the contract that contains its rln pk and some amount of Ether. This amount is deposited on the contract to prevent spam activity.

WAKU2-RLN-RELAY: Registration



Peers construct the rln membership Merkle tree locally

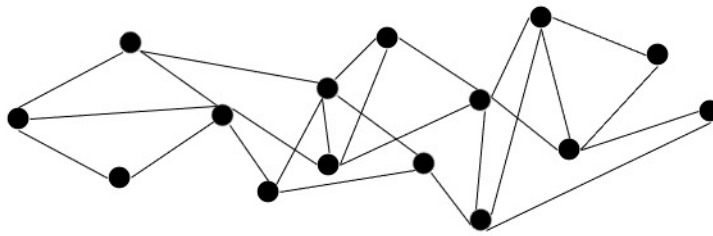
WAKU2-RLN-RELAY: Registration



And listen for the registration and deletion events emitted from the contract to update their trees.

WAKU2-RLN-RELAY: External Nullifier

External Nullifier = Epoch = the number of T seconds that elapsed since the Unix epoch.
Messaging rate is limited to 1 per epoch.

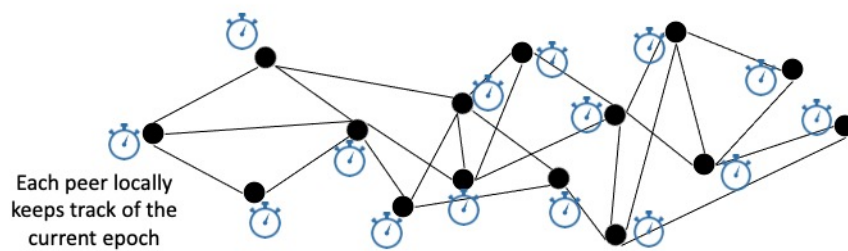


For the external nullifier we define it as the number of T seconds that elapsed since the Unix epoch.

Peers are allowed to publish one message per epoch without being slashed

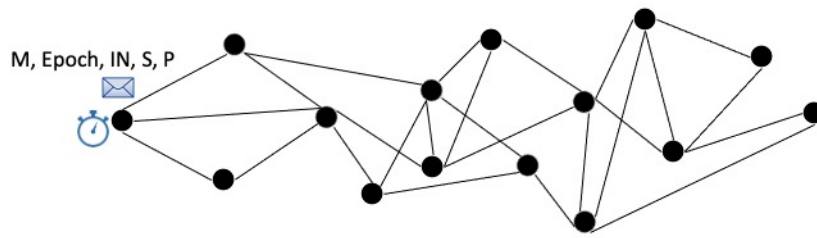
WAKU2-RLN-RELAY: External Nullifier

External Nullifier = Epoch = the number of T seconds that elapsed since the Unix epoch.
Messaging rate is limited to 1 per epoch.



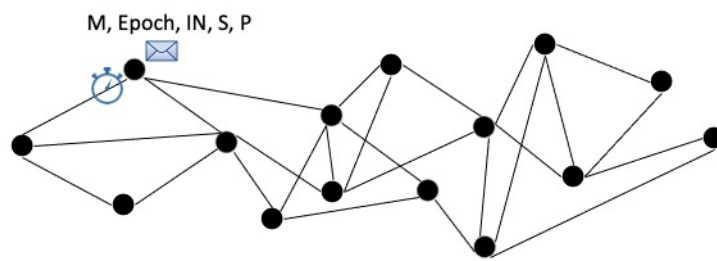
Each peer locally keeps track of the current epoch.

WAKU2-RLN-RELAY: Publishing



Message publishing is the same as RLN framework
The message owner, attaches the external and internal nullifiers, share of its sk, and zero knowledge proof part to the message

WAKU2-RLN-RELAY: Routing



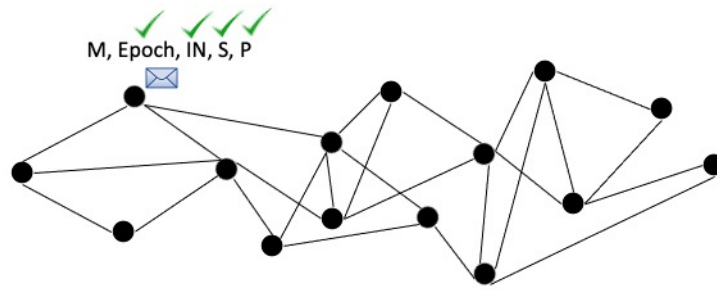
A routing peer follows the regular routing protocol of wakurelay (gossipSub protocol) and in addition checks

That Epoch of the incoming message does not have huge gap with the current epoch

Verifies the membership proof and the correctness of other fields

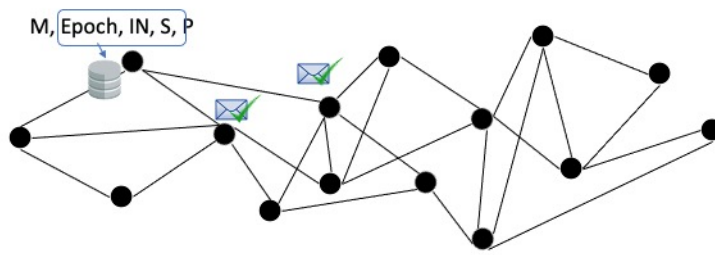
Checks the nullifiers to see whether the rate limit is violated or not (I will explain it a bit more in the next slides)

WAKU2-RLN-RELAY: Routing



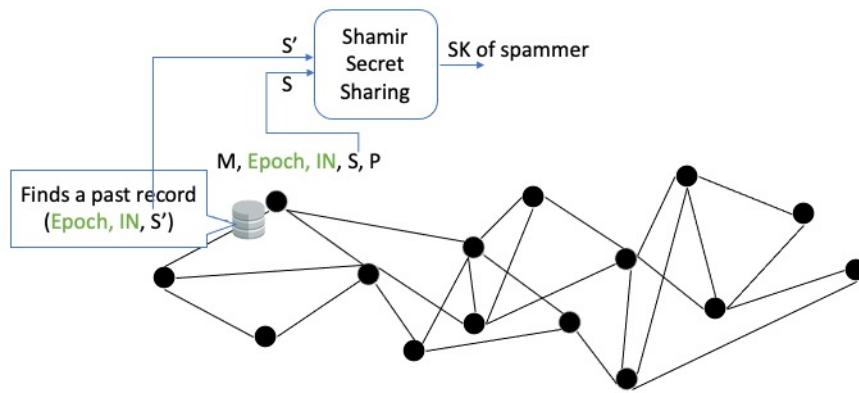
Relays the message if all these conditions are met, discards otherwise

WAKU2-RLN-RELAY: Routing



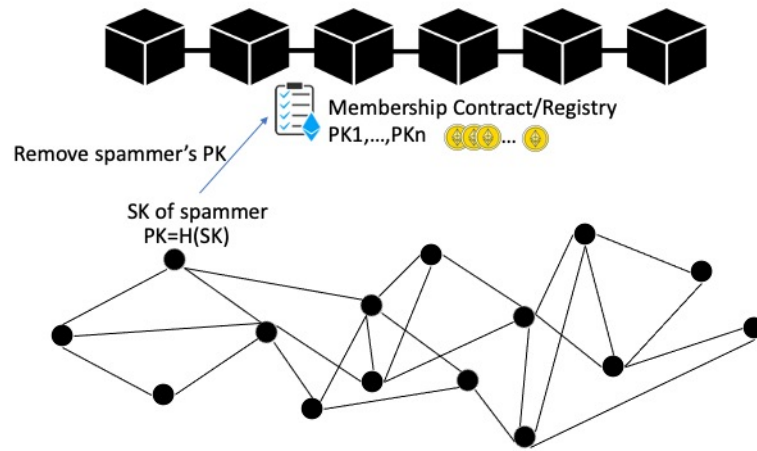
The peer keeps record of the rln metadata to catch spammers and slash them

WAKU2-RLN-RELAY: Slashing



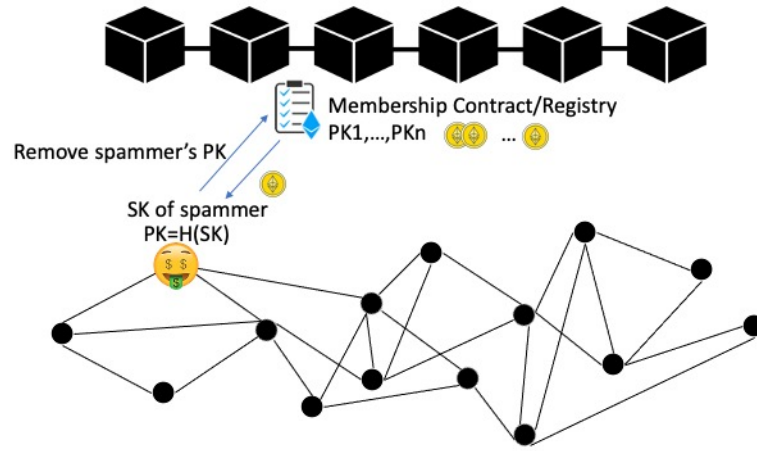
While checking the nullifiers, If the routing peer finds out that the messaging rate is violated, i.e., an old message with the same Epoch and internal nullifier values it reconstructs the sk of the spammer

WAKU2-RLN-RELAY: Slashing



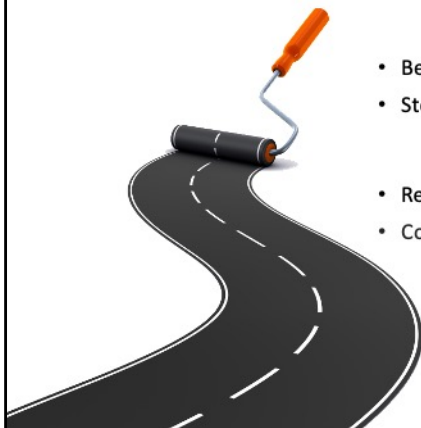
Sends a transaction to the contract and removes the spammer pk from the group

WAKU2-RLN-RELAY: Slashing



and withdraws its fund

Future work



- Benchmarking
- Storage-efficient Merkle tree storage
 - P2p network of full-nodes and light-nodes
 - Partial view of Merkle tree
- Real-time removal of spammers using off-chain/p2p solutions
- Cost-effective way of member insertion and deletion using layer 2 solutions

Now, let's browse our future plan:

Benchmarking is the first on our future work list

The next is to address storage overhead. Currently, peers maintain the entire tree locally which takes up to 67 MB for tree with depth 20 (and almost 274 GB for $d=32$). This overhead won't fit resource-limited devices.

- One solution is to use the light-node and full-node architecture where resource-full nodes retain the entire tree and serve it to the nodes with limited storage.
- Another possible solution is to have a partial view of the tree and yet being able to construct and update the tree root and the authentication path when group state changes

On-chain slashing is subject to delay (the tx has to be mined), so is the removal of spammers. We are seeking an off-chain slashing method, where peers communicate the pk of the slashed user in a p2p manner, to provide real-time spam-protection

The other direction to pursue is to provide a cost-effective way of member insertion and deletion using layer 2 solutions. The reason is that currently these operations cost almost 40 k gas, which translates to 15 USD which might be not affordable by the users, so an alternative solution is worth investigation

References

- Waku-rln-relay specs: <https://rfc.vac.dev/spec/17/>
- Waku-rln-relay paper: https://github.com/vacp2p/research/blob/master/rln-research/Waku_RLN_Relay.pdf
- Vac post on Waku-rln-relay: <https://vac.dev/rln-relay>
- Nim-Waku implementation: <https://github.com/status-im/nim-waku>
- js-Waku implementation: <https://github.com/status-im/js-waku>
- RLN Ethereum research post: <https://ethresear.ch/t/semaphore-rln-rate-limiting-nullifier-for-spam-prevention-in-anonymous-p2p-setting/5009>
- RLN medium post: <https://medium.com/privacy-scaling-explorations/rate-limiting-nullifier-a-spam-protection-mechanism-for-anonymous-environments-bbe4006a57d>
- RLN circuits: <https://github.com/appliedzkp/rln>
- RLN circuits spec: <https://hackmd.io/7GR5Vi28Rz2EpEmLK0E0Aw>
- RLN in Rust: <https://github.com/kilic/rln>

Thank you